ACHIEVING HIGH-PERFORMANCE THE Functional WAY

A Functional Pearl on Expressing High-Performance Optimizations as Rewrite Strategies





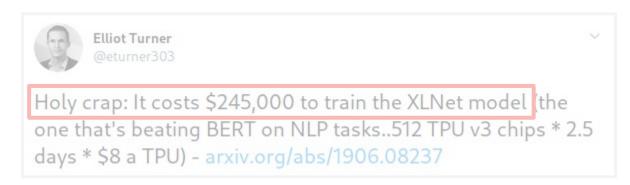


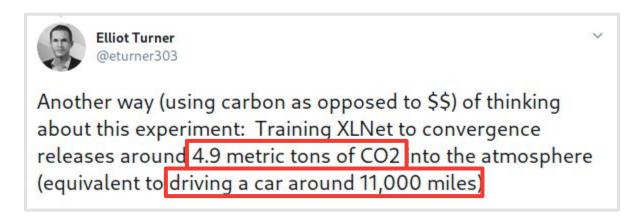
Why do we care?



Holy crap: It costs \$245,000 to train the XLNet model the one that's beating BERT on NLP tasks..512 TPU v3 chips * 2.5 days * \$8 a TPU) - arxiv.org/abs/1906.08237

Why do we care?







H-PERFORMANCE

Why do we care?

Turner mer303

It costs \$245,000 to train the XLNet model the beating BERT on NLP tasks..512 TPU v3 chips * 2.5 a TPU) - arxiv.org/abs/1906.08237

Turner rner303

ray (using carbon as opposed to \$\$) of thinking experiment: Training XLNet to convergence round 4.9 metric tons of CO2 nto the atmosphere it to driving a car around 11,000 miles)

Achieving High-Performance the <u>Functional</u> Way Manual

Naive Matrix Multiplication in



HIGH-PERFO

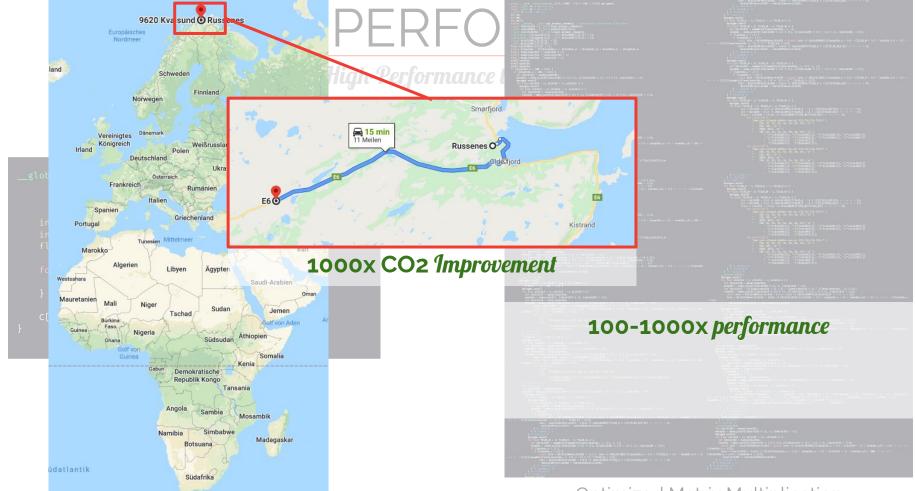
Achieving High-Performance

Naive Matrix Multiplication in

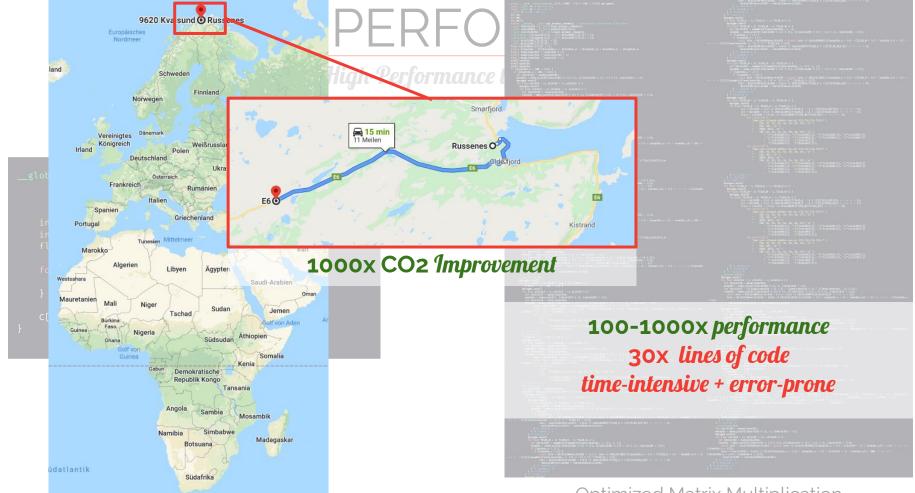


```
100-1000× performance
```

Optimized Matrix Multiplication

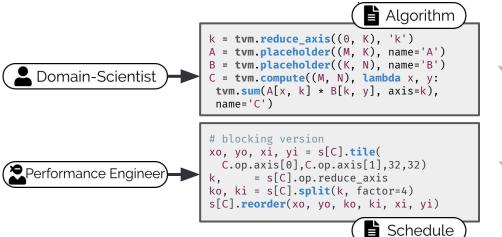


Optimized Matrix Multiplication



Optimized Matrix Multiplication

Achieving High-Performance the <u>Functional</u> Way Decoupled



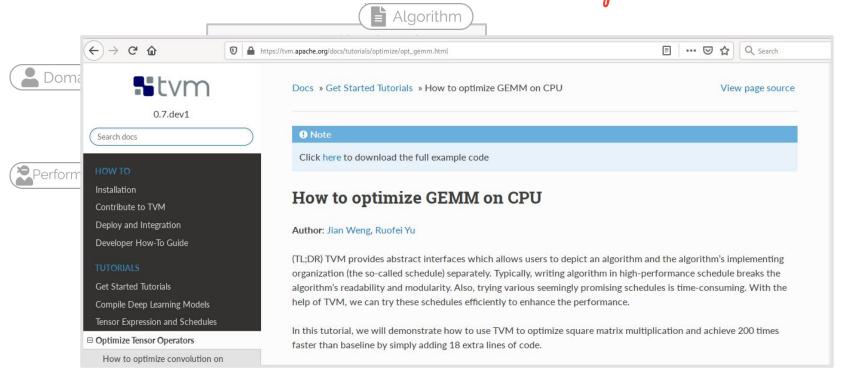
Halide



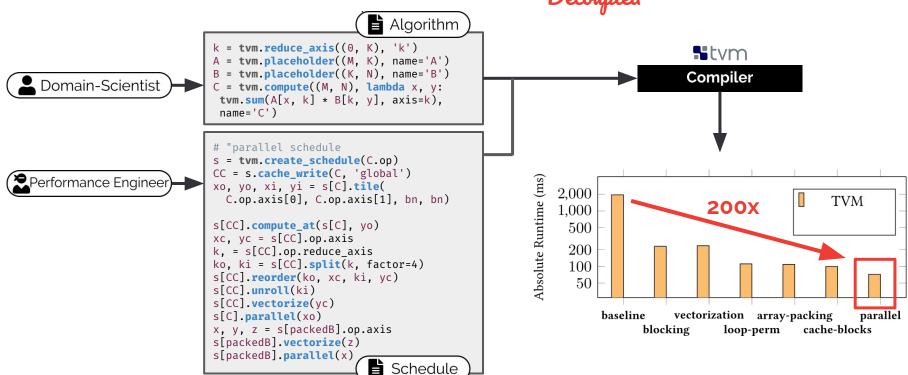
What to compute Tiramisu-Compiler / tiramisu

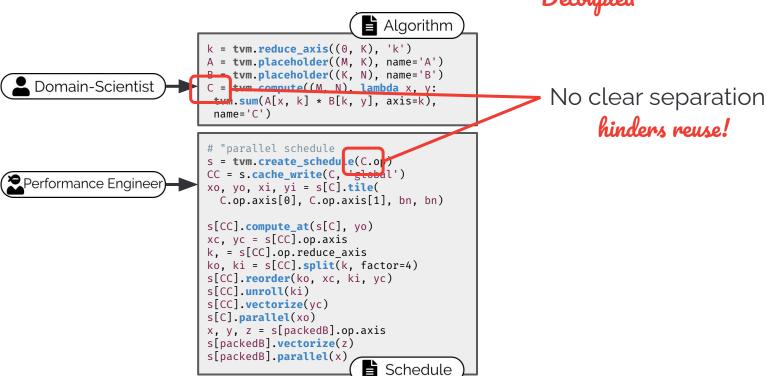


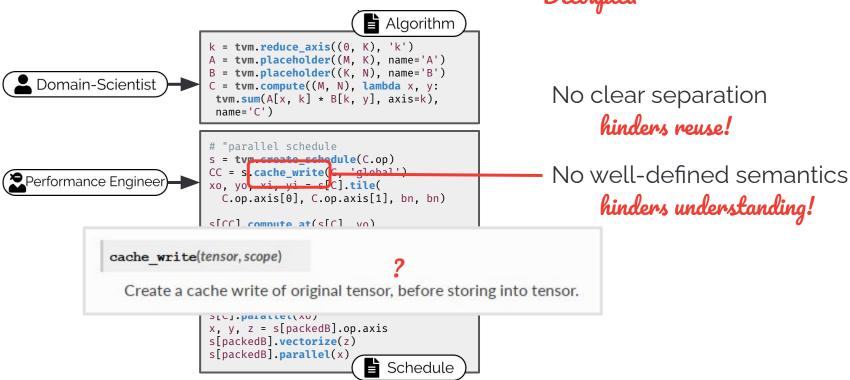
How to optimize



Achieving High-Performance the **Eunctional** Way Decoupled Algorithm k = tvm.reduce_axis((0, K), 'k') **Stvm** A = tvm.placeholder((M, K), name='A') B = tvm.placeholder((K, N), name='B') Compiler Domain-Scientist C = tvm.compute((M, N), lambda x, y: tvm.sum(A[x, k] * B[k, y], axis=k),name='C') # blocking version xo, yo, xi, yi = s[C].tile(C.op.axis[0],C.op.axis[1],32,32) Absolute Runtime (ms) Performance Engineer k, = s[C].op.reduce_axis 2,000 ko, ki = s[C].split(k, factor=4) **TVM** 1,000 s[C].reorder(xo, yo, ko, ki, xi, yi) 500 Schedule 200 100 50 vectorization array-packing baseline parallel blocking loop-perm cache-blocks







Achieving High-Performance the Functional Way Decoupled

```
k = tvm.reduce_axis((0, K), 'k')
A = tvm.placeholder((M, K), name='A')
B = tvm.placeholder((K, N), name='B')
C = tvm.compute((M, N), lambda x, y:
tvm.sum(A[x, k] * B[k, y], axis=k),
name='C')
# "parallel schedule
```

Performance Engineer

```
s = tvm.create_schedule(C.op)
CC = s.cache_write(C__'global')
xo, yo, xi, yi = s[C .tile(
  C.op.axis[0], C.op.axis[1], bn bn)
s[CC].compute_at(s[C], yo)
xc, yc = s[CC].op.axis
k, = s[CC].op.reduce axis
ko, ki = s[CC].split(k, factor=4)
s[CC].reorder(ko, xc, ki, yc)
s[CC].unrott(ki)
s[CC].vectorize(y)
s[C].parallel(x0)
x, y, z = s[packedB].op.axis
s[packedB].vectorize(z)
s[packedB].parallel(x)
                        Schedule
```

No clear separation hinders reuse!

No well-defined semantics hinders understanding!

Optimizations are built-in no extensibility!

Achieving High-Performance the Functional Way Decoupled

```
k = tvm.reduce_axis((0, K), 'k')
A = tvm.placeholder((M, K), name='A')
B = tvm.placeholder((K, N), name='B')
C = tvm.compute((M, N), lambda x, y:
tvm.sum(A[x, k] * B[k, y], axis=k),
name='C')
```

No clear separation

hindons nouse!



We aim for a more principled way to describe and apply optimizations

```
xc, yc = s[CC].op.axis
k, = s[CC].op.reduce_axis
ko, ki = s[CC].split(k, factor=4)
s[CC].reorder(ko, xc, ki, yc)
s[CC].unroit(ki)
s[CC].vectorize(y;)
s[C].paratict(xo)
x, y, z = s[packedB].op.axis
s[packedB].vectorize(z)
s[packedB].parallel(x)
Schedule
```

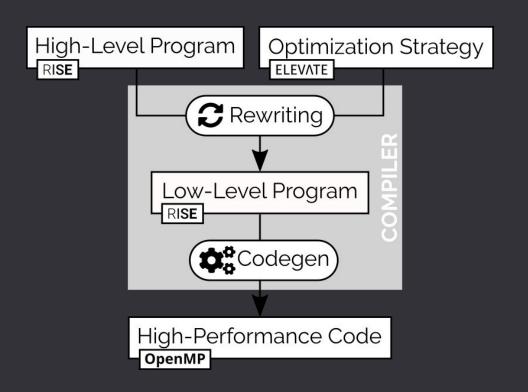
Optimizations are built-in no extensibility!

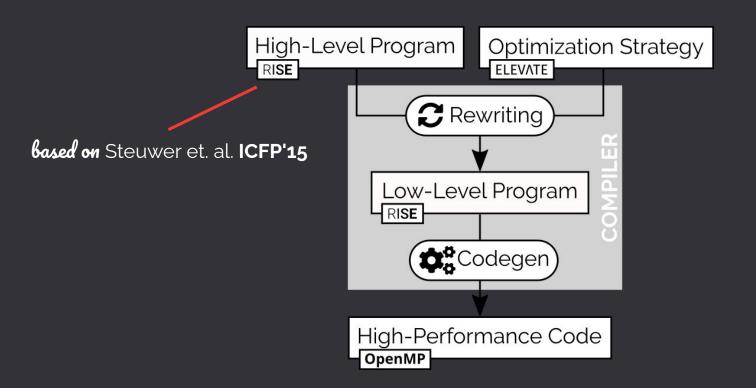
OUR GOALS

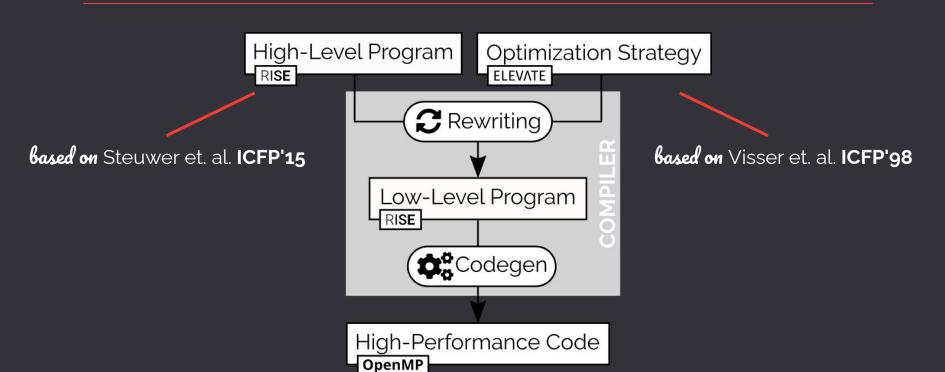
A Principled Way to Separate, Describe, and Apply Optimizations

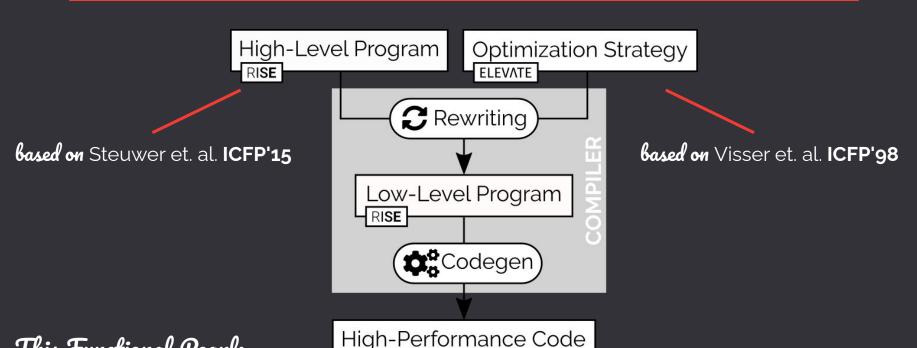
- (1) *Separate concerns*: Computations should be expressed at a high abstraction level only. They should not be changed to express optimizations;
- (2) Facilitate reuse: Optimization strategies should be defined clearly separated from the computational program facilitating reusability of computational programs and strategies;
- (3) *Enable composability*: Computations *and* strategies should be written as compositions of user-defined building blocks (possibly domain-specific ones); *both languages* should facilitate the creation of higher-level abstractions;
- (4) *Allow reasoning*: Computational patterns, but also especially strategies, should have a precise, well-defined semantics allowing reasoning about them;
- (5) *Be explicit*: Implicit default behavior should be avoided to empower users to be in control.

Fundamentally we argue that a more principled high-performance code generation approach should be holistic by considering computation and optimization strategies equally important. As a consequence, a strategy language should be built with the same standards as a language describing computation.









This Functional Pearl:

We apply established functional programming techniques for elegantly expressing high-performance program optimizations as composable rewrite strategies

OpenMP

ELEVATE

A Language for Describing Optimization Strategies

A **Strategy** encodes a program transformation:

```
type Strategy[P] = P => RewriteResult[P]
```

A **RewriteResult** encodes its success or failure:

ELEVATE

A Language for Describing Optimization Strategies

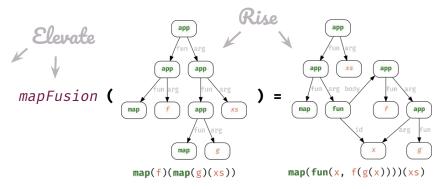
A **Strategy** encodes a program transformation:

```
type Strategy[P] = P => RewriteResult[P]
```

A **RewriteResult** encodes its success or failure:

Rewrite Rules are examples for basic strategies

```
def mapFusion: Strategy[Rise] =
  (p:Rise) => p match {
   case app(app(map, f),
        app(app(map, g), xs)) =>
    Success( map(fun(x => f(g(x))))(xs) )
   case _ => Failure( mapFusion )
}
```



COMBINATORS

How to Build More Powerful Strategies

Sequential Composition (;)

```
def seq[P]: Strategy[P] => Strategy[P] => Strategy[P] =
    fs => ss => p => fs(p) >>= ss
```

Left Choice (<+)

Try

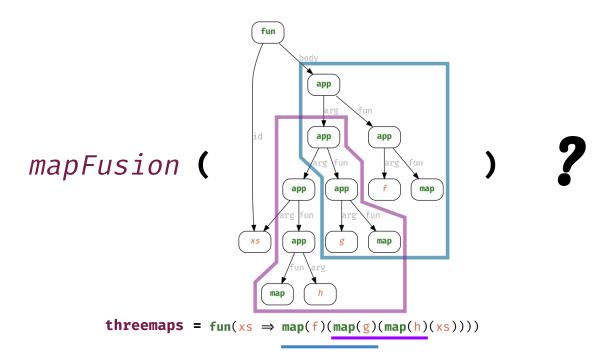
```
def try[P]: Strategy[P] => Strategy[P] =
    s => p => (s <+ id)(p)</pre>
```

Repeat

```
def repeat[P]: Strategy[P] => Strategy[P] =
    s => p => try(s ; repeat(s))(p)
```

TRAVERSALS

Describing Precise Locations



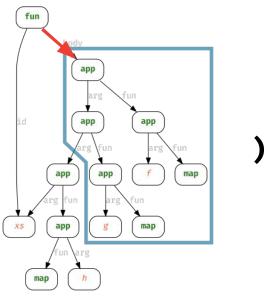
There are two possible locations for successfully applying the rule

TRAVERSALS

Describing Precise Locations

```
def body: Traversal[Rise] = s => p => p match {
  case fun(x,b) => (nb => fun(x,nb) <$> s(b)
  case _ => Failure( body(s) )
  }
apply S at body of function abstraction
```

body(mapFusion) (



```
threemaps = fun(xs, map(f)(map(g)(map(h)(xs))))
```

There are *two possible locations* for successfully applying the rule

TRAVERSALS

Describing Precise Locations

```
def body: Traversal[Rise] = s => p => p match {
      case fun(x,b) \Rightarrow (nb \Rightarrow fun(x,nb) \Leftrightarrow s(b)
      case => Failure( body(s) )
    body(argument(mapFusion)) (
def argument: Traversal[Rise] = s => p => p match {
  case app(f,a) \Rightarrow (na \Rightarrow app(f,na) < s > s(a)
  case _ => Failure( argument(s) )
apply s at argument of function application
                                      threemaps = fun(xs, map(f)(map(g)(map(h)(xs))))
```

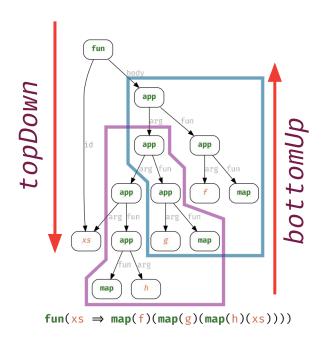
There are two possible locations for successfully applying the rule

NORMALIZATION

More Complex Traversals

Generic Tree Traversals...

```
def topDown: Traversal[Rise] = s => p => (s <+ one(topDown(s)))(p)
def bottomUp: Traversal[Rise] = s => p => (one(topDown(s)) <+ s)(p)
...</pre>
```



NORMALIZATION

More Complex Traversals

Generic Tree Traversals...

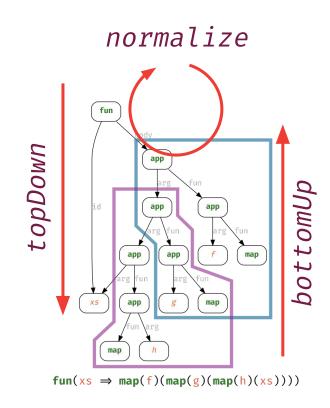
```
def topDown: Traversal[Rise] = s => p => (s <+ one(topDown(s)))(p)
def bottomUp: Traversal[Rise] = s => p => (one(topDown(s)) <+ s)(p)
...</pre>
```

... and a strategy for normalization

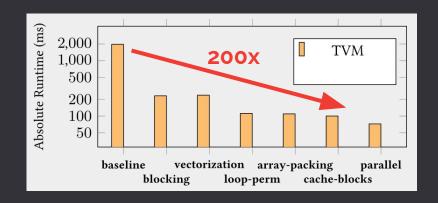
```
def normalize: Traversal[Rise] = s => p => repeat(topDown(s))(p)
```

With these, we define normal-forms like $\beta\eta$ -normal-form

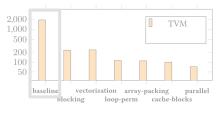
```
def BENF = normalize(betaReduction <+ etaReduction)</pre>
```



Implementing TVM's Scheduling Language



Optimizing Matrix Multiplication - Baseline



RISE

What to compute



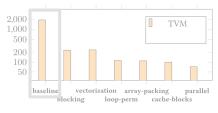
```
// matrix multiplication in RISE
val dot = fun(as, fun(bs, zip(as)(bs) |>
map(fun(ab, mult(fst(ab))(snd(ab)))) |>
reduce(add)(o) ) )
val mm = fun(a, fun(b, a |>
map( fun(arow, transpose(b) |>
map( fun(bcol, dot(arow)(bcol) ))))))

// baseline strategy in ELEVATE
val baseline = ( DFNF ';'
fuseReduceMap 'o' topDown )
(baseline ';' lowerToC)(mm)
```

ELEVATE

How to optimize

Optimizing Matrix Multiplication - Baseline



clear separation

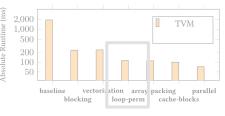
RISE

```
// matrix multiplication in RISE
val dot = fun(as, fun(bs, zip(as)(bs) |>
  map(fun(ab, mult(fst(ab))(snd(ab)))) |>
  reduce(add)(o) )
val mm = fun(a, fun(b, a |>
  map( fun(arow, transpose(b) |>
    map(fun(bcol,
      dot(arow)(bcol) )))) ))
// baseline strategy in ELEVATE
val baseline = ( DFNF ';'
  fuseReduceMap 'a' topDown )
(baseline ';' YowerToC)(mm)
```

Stym



Optimizing Matrix Multiplication - Loop Permutation



facilitate reuse

user-defined vs. built-in

```
val loopPerm = (
tile(32,32) '@' outermost(mapNest(2)) ';;'
fissionReduceMap '@' outermost(appliedReduce) ';;'
split(4) '@' innermost(appliedReduce) ';;'
reorder(Seq(1,2,5,3,6,4)) ';;'
vectorize(32) '@' innermost(isApp(isApp(isMap))))
(loopPerm ';' lowerToC)(mm)
```

ELEVATE



no clear separation of concerns

Optimizing Matrix Multiplication - Array Packing

clear separation of concerns vs. no clear separation

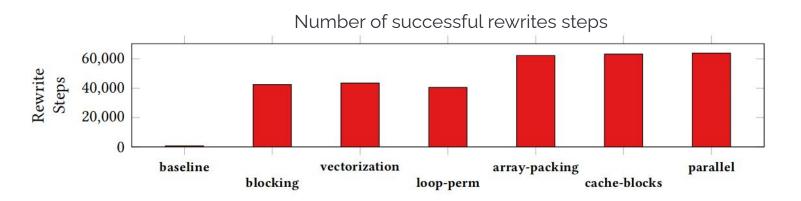
```
val appliedMap = isApp(isApp(isMap))
   val isTransposedB = isApp(isTranspose)
   val packB = storeInMemory(isTransposedB,
    permuteB ';;'
    vectorize(32) 'a' innermost(appliedMap) '::'
                                                     11 # Array packing schedule
    parallel
                   'a' outermost(isMap)
   ) 'a' inLambda
   val par = (
     packB ':: 'loopPerm ':: '
                                                         s[CC].compute_at(s[C], yo)
     (parallel 'a' outermost(isMap))
                                                         xc, yc = s[CC].op.axis
                'a' outermost(isToMem) ';;'
                                                         k, = s[CC].op.reduce axis
13
      unroll 'a' innermost(isReduce))
15
   (par '; 'lowerToC )(mm)
                                                         s[CC].unroll(ki)
                                                         s[CC].vectorize(vc)
```

ELEVATE

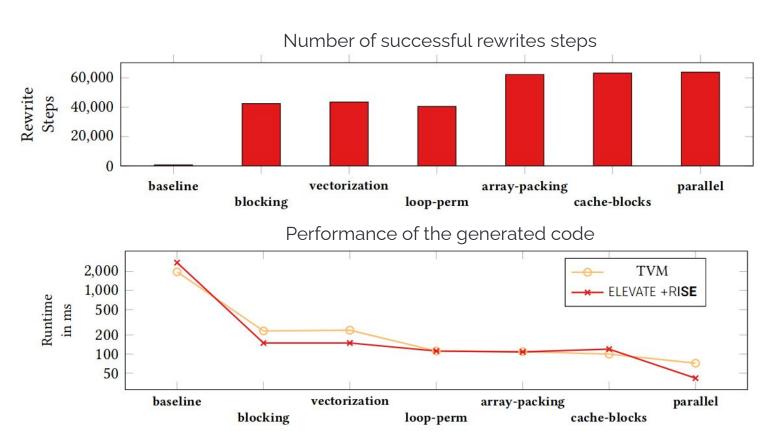
```
1 # Modified algorithm
  bn = 32
    = tvm.reduce_axis((o, K), 'k')
    = tvm.placeholder((M, K), name='A')
    = tvm.placeholder((K, N), name='B')
  pB = tvm.compute((N / bn, K, bn),
     lambda x, y, z: B[y, x * bn + z], name='pB')
  C = tvm.compute((M,N), lambda x,v:
    tvm.sum(A[x,k] * pB[y//bn,k,
    tvm.indexmod(y,bn)], axis=k),name='C')
  s = tvm.create_schedule(C.op)
  CC = s.cache_write(C, 'global')
  xo, yo, xi, yi = s[C].tile(
    C.op.axis[0], C.op.axis[1], bn, bn)
  ko, ki = s[CC].split(k, factor=4)
  s[CC].reorder(ko, xc, ki, yc)
  s[C].parallel(xo)
  x, y, z = s[pB].op.axis
  s[pB].vectorize(z)
  s[pB].parallel(x)
```



Counting Rewrite Steps and Measuring Performance



Counting Rewrite Steps and Measuring Performance



ACHIEVING HIGH-PERFORMANCE THE Functional WAY

...is Open Source!

RISE

rise-lang.org/ github.com/rise-lang

ELEVATE

elevate-lang.org
github.com/elevate-lang







