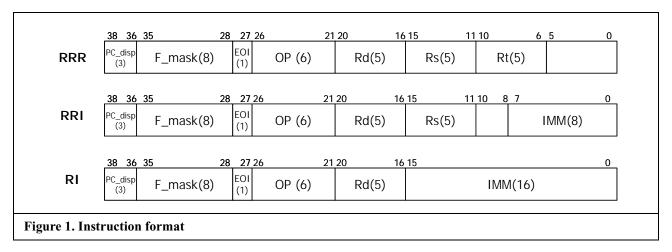
1.0 Micro-ISA

1.1 Instruction word format



1.2 Instruction word fields descriptions

Table 1: Field descriptions

Field	Description								
PC_disp (3)	PC update information. When an instruction w/ non-zero PC_disp is decoded, the content of PC is updated to PC + PC_disp. Usually, PC_disp is the number of z80 instruction bytes, and only the first instruction if the translation group has non-zero values in this field; other instructions have 0 in this field.								
F_mask (8)	Each bit in this field represents a corresponding condition bit as follows:								
	7 6 5 4 3 2 1 0								
	S Z X N X PV N C								
EOI (1)	rent instruction. If the bit is 1, the corresponding field in F register is updated; otherwise, the corresponding bit is not affected. As defined in F register, bit 5 and 3 are unused. End of translation group, wich represents the z80 instruction boundary. This bit controls control transfer. Interrupts can occur only after the instruction								
OP (6)	with this field, or after a taken jump / branch.								
Rd (5)	Instruction opcode Destination register field. In RI format, this field can represent a source register Source register 1 field (left-hand side operand in instruction descriptions)								
Rs (5)									
Rt (5)	Source register 2 field (right-hand side operand in instruction descriptions)								
IMM (8)	8-bit immediate (signed) field. If the instruction requires a 16-bit value from this field, it is sign-extended.								
IMM (16)	16-bit immediate (signed) field.								

1.3 Special-purpose registers

Table 2: Special purpose registers

Register	Description
archPC (16)	Architected 16-bit program counter. Incremented by PC_disp every clock cycle, or updated with the target PC address if a control transfer occurs.
	Note on CTIs (control transfer instructions): When the first instruction in the translation group, archPC is updated by PC_disp, pointing to the next sequential PC. When archPC should be updated by a taken relative branch instruction, you may assume that the offset field of the branch has been pre-calculated so that you don't have to worry about whether archPC is currently pointing to the current PC or the next PC
IFF1 (1)	determines if an incoming maskable interrupt is blocked. The content of this field is updated by EI and DI instructions. See z80 user's manual
IFF2 (1)	contains the previous state of IFF1. The content of this field is updated by EI and DI instructions. See z80 user's manual

1.4 General purpose registers

General purpose registers are accessed through the register file using 5-bit register specifiers. If the register identifier starts with 0, it is in 8/16-bit group, in which all 8-bit values are sign-extended before fed into the ALU. 16-bit group register identifiers always start with 1. In addition to z80 architected registers, this group has 4 temporary registers, which always end with 10. The order of register specifiers were determined by the register indexing scheme used in z80 architecture, in favor of the ease of translation. So don't complain that the order is so messy.

Table 3: General purpose registers

	8-bit regis	ster group	16-bit register group			
register	specifier	description	register	specifier	description	
R0	00000	zero register that always contains 0. Can be used as a 16-bit register	SP	11001	z80 stack pointer	
A	01111	z80 accumulator register	AF	11100	z80 {A,F}	
F	01101	z80 condition code vector	Аг	11100	ZδU {A,Γ}	
В	00001	z80	ВС	10000	z80 {B,C}	
С	00011	z80	БС	10000	Z00 {B,C}	
D	00101	z80	DE	10100	z80 {D,E}	
Е	00111	z80	DE	10100	Z00 {D,E}	
Н	01001	z80	HL	11000	z80 {H,L}	
L	01011	z80	TIL	11000	200 (11,L)	
I	00100	z80	IX	10001	z80	
R	01000	z80	IY	10101	z80	
Т0	00010	8bit temp	T10	10010	{T1,T0}	
T1	00110	8bit temp	110	10010	{11,10}	
T2	01010	8bit temp	T32	11010	{T3,T2}	
Т3	01110	8bit temp	132	11010	{13,12}	
			AF'	11111	z80 shadow AF	
			BC'	10011	z80 shadow BC	
			DE'	10111	z80 shadow DE	

Table 3: General purpose registers

	8-bit regis	ter group	16-bit register group			
register	specifier description		register	specifier	description	
			HL'	11011	z80 shadow HL	

1.5 Instructions

IMPORTANT: Make sure that you are realizing flag bits correctly. In the table below, each instruction has an column "flags generated", in which you can find a note such as "see z80 DAA". Then, look up the table "flag notations" and find the detailed descriptions that you need to implement.

Table 4: Operand notation

notation	description
Rx	8bit register
RRx	16bit register
RRRx	8bit or 16bit register
IMM8	8-bit immediate (in RRI format). You may want to do sign-extention for this field if the actual computation is performed in 16-bit
IMM16	16-bit immediate (in RI format)

Table 5: Flag notations

	Description	S (sign)	Z (zero)	H (half carry / borrow)	P/V (parity / overflow)	N (negate)	C (carry / borrow)
standard (8bit)	standard 8bit add/sub	result (7)	result == 0 ? 1: 0	ADD: carry at bit 3 to 4 SUB: borrow from bit 4	overflow at 8bit computa- tion	ADD: 0 SUB: 1	ADD: carry at bit 7 to 8 SUB: borrow from bit 8
standard(16bit)	standard 16bit add/sub	result (15)	result == 0 ? 1: 0	ADD: carry at bit 11 to 12 SUB: borrow from bit 12	overflow at 16bit compu- tatin	ADD: 0 SUB: 1	ADD: carry at bit 15 to 16 SUB: borrow from bit 8
z80 AND		result (7)*	result(7:0) == 0 ? 1 : 0*	1	overflow at 8bit computa- tion**	0	0
z80 OR		result (7)*	result(7:0) == 0 ? 1: 0*	0	overflow at 8bit computa- tion**	0	0
z80 XOR		result (7)*	result(7:0) == 0 ? 1 : 0*	0	1 when even parity	0	0
z80 shift/rotate for XXX	for z80 RLC, RRC, RL, RR, SLA, SRA, SRL	result (7)*	result(7:0) == 0 ? 1 : 0*	0	1 when even parity	0	consult z80 user's manual for carry flags
z80 shift/rotate for XXXA	for z80 RLCA, RRCA, RLA, RRA	not affected***	not affected***	0	not affected***	0	consult z80 user's manual for carry flags
z80 RRD/RLD		result (7)	result == 0 ? 1 : 0	0	1 when even parity	0	not affected***
z80 bit b		don't care	result == 0 ? 1 : 0	1	don't care	0	not affected***

Table 5: Flag notations

	Description	S (sign)	Z (zero)	H (half carry / borrow)	P/V (parity / overflow)	N (negate)	C (carry / borrow)
z80 in		bit 7 of the data from I/O	the data from I/O == 0 ? 1:0	0	1 when the data from I/O has even par- ity	0	not affected***
z80 DAA***		A (7)	A == 0 ? 1:0	see z80 DAA instruction	1 when A has even parity	not affected***	see z80 DAA instruction

^{*} Instructions may execute with 16-bit registers. However, sign and zero flags are generated by examining only result bits (7:0).

Table 6: Instructions

type	Opcod e	Synopsis	for- mat	Semantic	description	Flags gener- ated	note
	00	add Ra, Rb, Rc	RRR	Ra = Rb + Rc		standard (8bit)	
	01	addi Ra, Rb, IMM8	RRI	Ra = Rb + IMM(signed)		standard (8bit)	
	02	sub Ra, Rb, Rc	RRR	Ra = Rb – Rc		standard (8bit)	
ALU arith- metic	03	add16 RRa, RRb, RRc	RRR	RRa = RRb + RRc		standard (16bit)	
	04	addi16 RRa, RRb, IMM8	RRI	RRa = RRb + IMM(signed)		standard (16bit)	IMM8 is sign extended
	05	sub16 RRa, RRb, RRc	RRR	RRa = RRb – RRc		standard (16bit)	

^{**}In z80 manual, it is unclear that overflow bit for AND and OR operation. Just follow the standard boolean equation for overflow checking.

^{***} it will be taken care of by annotation bits. Therefore, treat it as "don't care"

^{****} follow flag bits in z80 user's manual

Table 6: Instructions

type	Opcod e	Synopsis	for- mat	Semantic	description	Flags gener- ated	note
	06	and RRRa, RRRb, RRRc	RRR	RRRa = bitwise RRRb & RRRc		follow z80 AND	
	07	and RRRa, RRRb, IMM8	RRI	RRRa = bitwise RRRb & IMM8		follow z80 AND	If RRR is a 16bit register, IMM8 is sign extended
	08	or RRRa, RRRb, RRRc	RRR	RRRa = bitwise RRRb RRRc		follow z80 OR	
	09	or RRRa, RRRb, IMM8	RRI	RRRa = bitwise RRRb IMM8		follow z80 OR	if RRR is a 16 bit register, IMM8 is sign extended
	0A	xor RRRa, RRRb, RRRc	RRR	RRRa = bitwise RRRb xor RRRc		follow z80 xor	
	0B	xor RRRa, RRRb, IMM8	RRI	RRRa = bitwise RRRb xor IMM8		follow z80 xor	if RRR is a 16bit register IMM8 is sign extended
	0C	not RRRa, RRRb	RRI	RRRa = bitwise ~RRRb		follow z80 OR	
ALU logi- cal	OD	shiftrotate Ra, Rb, OP	RRI	OP field is defined as follows: RLC: 00 RLCA: 01 RRC: 02 RRCA: 03 RL: 04 RLA: 05 RR: 06 RRA: 07 SLA: 08 SRA: 09 SRL: 0A SLL: 0B	direct imple- mentation of rotate and shift z80 instruc- tions	as defined in z80 shift and rotate instruc- tions	8bit operation note: SLL (shift left logical) operation is undocu- mented and will be added. please find some informa- tion on undocumented z80 instruction
	0E	get4 Ra, Rb, IMM8	RRI	Ra = Rb(7:4) if IMM!=0, Rb(3:0) if IMM==0	get H/L portion 4 bits	none	for RLD operation
	0F	merge44 Ra, Rb, Rc	RRR	Ra = Rb(3:0), Rc(3:0)	merge 2 L portion 4-bit chunks into one 8-bit value	follow RRD / RLD	for RLD operation
	10	mvPC RRa	RRI	RRa = PC	copy the cur- rent content of PC into RRa	none	This is required since PC is a special purpose register
ALU spe- cial	11	DAA		same as z80 DAA	direct imple- mentation of z80 DAA instruction	follow z80 DAA	special ALU support
	12	getbit Ra, Rb, IMM8	RRI	Ra = Rb(bit IMM)	get one bit from Rb, copy the bit into Ra, clears other bits in Ra	see bit b, (??)	8bit op, Ra is either 0 or 1
ALU bit operation	13	ngetbit Ra, Rb, IMM8	RRI	Ra = ~Rb(bit IMM)	get one bit from Rb, invert the bit, copy the bit into Ra, clears other bits in Ra	see bit b, (??)	8bit op, Ra is either 1 or 0
	14	setbit Ra, Rb, IMM8	RRI	Ra(bit IMM) = Rb(0)	set the bit in Ra to Rb(0)	none	8bit op, Rb is either 0 or 1
	15	nsetbit Ra, Rb, IMM8	RRI	Ra(bit IMM) = ~Rb(0)	set the bit in Ra to ~Rb(0)	none	8bit op, Rb is either 1 or 0

Table 6: Instructions

type	Opcod e	Synopsis	for- mat	Semantic	description	Flags gener- ated	note
	20	j IMM16	RI	PC = IMM	jump direct to IMM16	none	always taken
	21	jr RRa	RI	PC = RRa	jump direct to the content of RRa	none	always taken
	22	Jc RRRa, IMM16	RI	if RRRa is not 0, PC = IMM	jump direct to IMM16 if RRRa is true (non zero)	none	conditional, RRRa is either 8 or 16 bits
Control transfer	23	Jrc RRRa, RRb	RRI	if RRRa is not 0, PC = RRb	jump direct to the content of RRRb if RRRa is true	none	conditional, RRRa is either 8 or 16 bits
	24	bne RRRa, RRRb, IMM8	RRI	if RRRa != RRRb, PC = PC + IMM8 (signed)	jump relative if (RRRa == RRRb)	none	relative branch, the offset is dependent on z80 ops, assuming PC has already been incremented at the beginning of the translated instruction group
	25	beq RRRa, RRRb, IMM8	RRI	if RRRa == RRRb, PC = PC + IMM8 (signed)	jump relative if (RRRa != RRRb)	none	relative branch, the offset is dependent on z80 ops, assuming PC has already been incremented at the beginning of the translated instruction group
	30	ld Ra, (RRb + IMM8)	RRI	Ra = [RRb + IMM(signed)]		none	8bit load
load / store	31	st Ra, (RRb + IMM8)	RRI	[RRb + IMM(signed)] = Ra		none	8bit store
Store	16	limm RRRa, IMM16	RI	RRRa = IMM	load 16-bit immediate into the register	none	higher 8bits of IMM are ignored if RRRa is Ra
	17	EI	RRR	set iff to 1	enable inter- rupt	none	interrupt FF (iff1, iff2) con- trol. See z80 EI implemen- tation
Interrupt	18	DI	RRR	set iff to 0	disable inter- rupt	none	interrupt FF (iff1, iff2) control. See z80 DI implementation
	19	IM IMM16	RI	set the interrupt mode to IMM16 (do nothing)		none	will not be implemented since SMS uses only interrupt mode 1.
1/0	32	in Ra, (Rb)	RRI	Ra = I/O port (Rb)	I/O port read	follow z80 In	
I/O	33	out Ra, (Rb)	RRI	I/O port (Rb) = Ra	I/O port write	none	