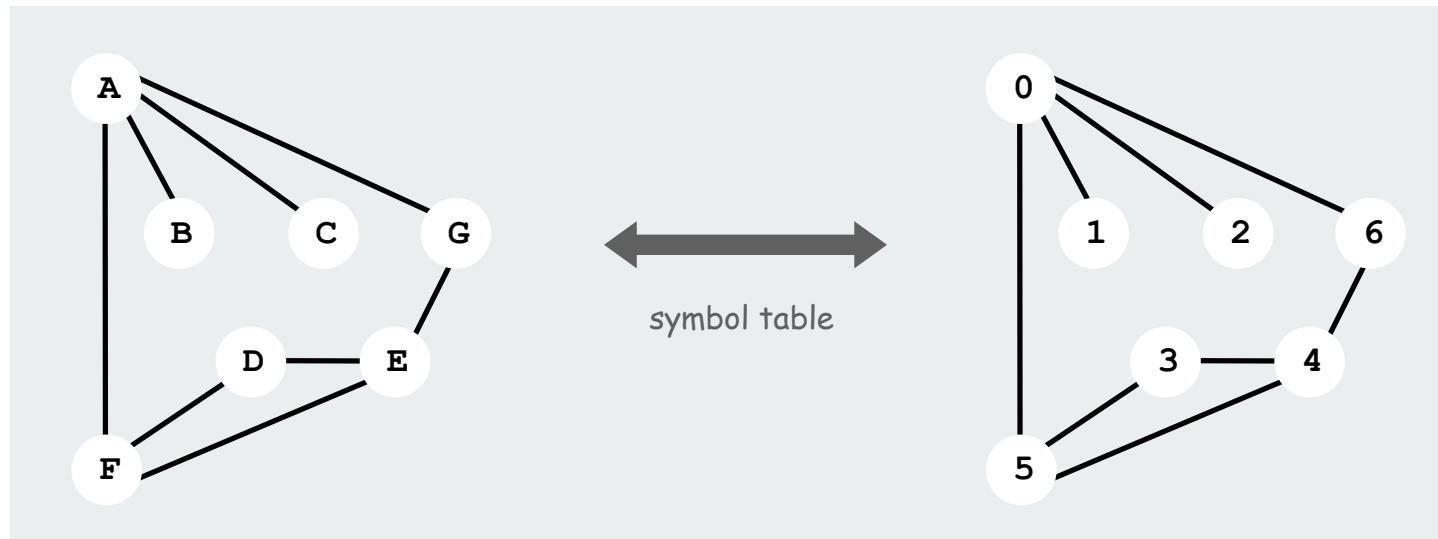


Graph Traversal

M. Elif Karslıgil

Undirected Graph Representation

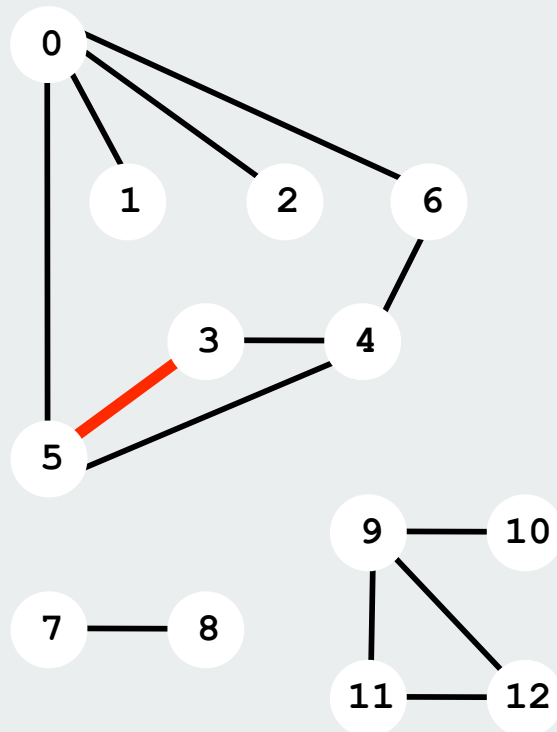
- **Computer** : use integers between 0 and V-1
- **Real world**: convert between names and integers with symbol table



Undirected Graph - Adjacency Matrix

Maintain a two-dimensional $v \times v$ boolean array.

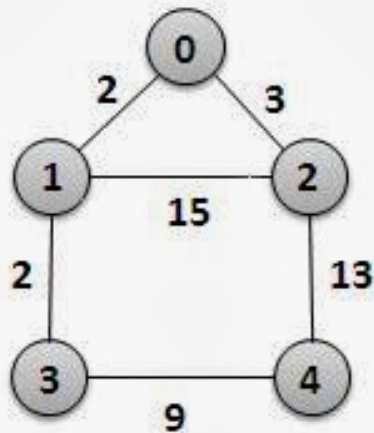
For each edge $v-w$ in graph: $\text{adj}[v][w] = \text{adj}[w][v] = \text{true}$.



two entries for each edge

| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
|----|---|---|---|---|---|---|---|---|---|---|----|----|----|
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 7 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 8 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 9 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
| 10 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 11 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 12 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |

Undirected Weighted Graph - Adjacency Matrix

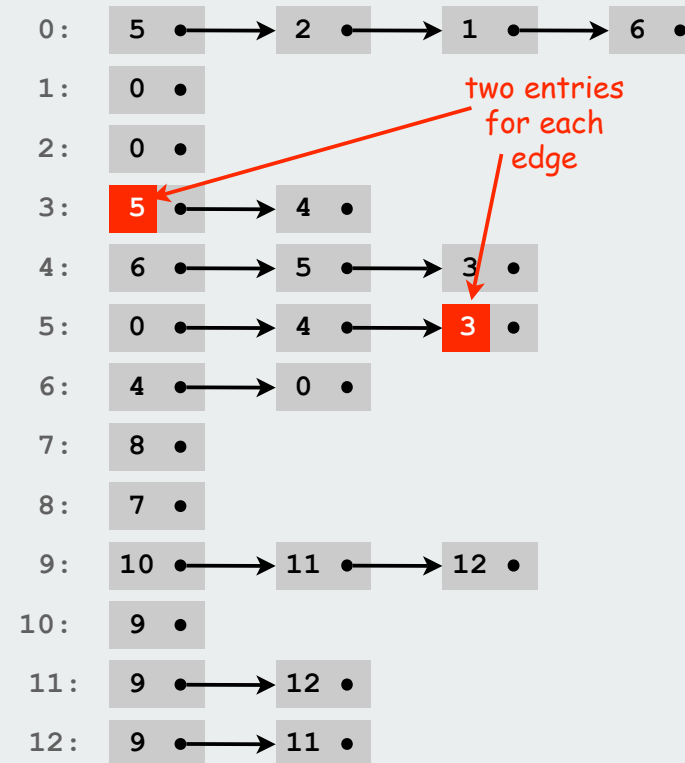
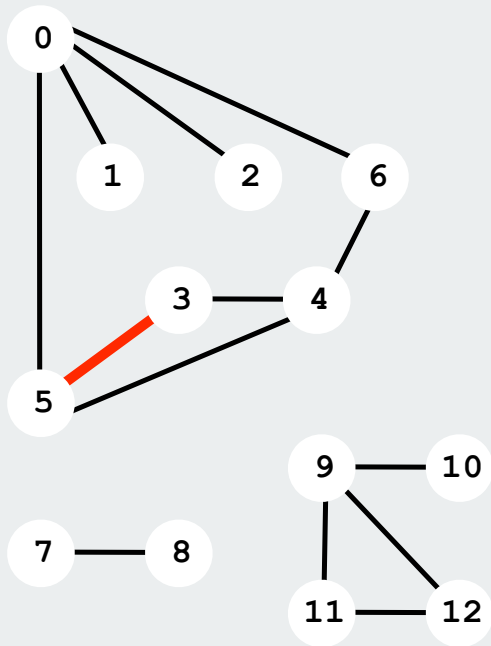


| | 0 | 1 | 2 | 3 | 4 |
|---|---|----|----|---|----|
| 0 | 0 | 2 | 3 | 0 | 0 |
| 1 | 2 | 0 | 15 | 2 | 0 |
| 2 | 3 | 15 | 0 | 0 | 13 |
| 3 | 0 | 2 | 0 | 0 | 9 |
| 4 | 0 | 0 | 13 | 9 | 0 |

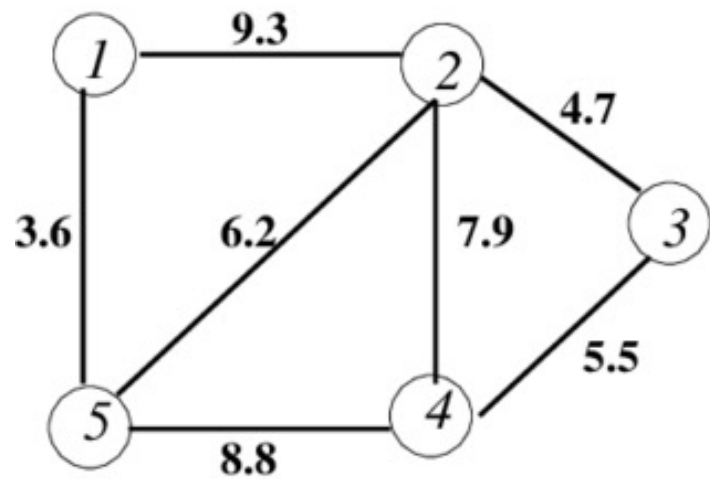
**Adjacency Matrix Representation of
Weighted Graph**

Undirected Graph - Adjacency List

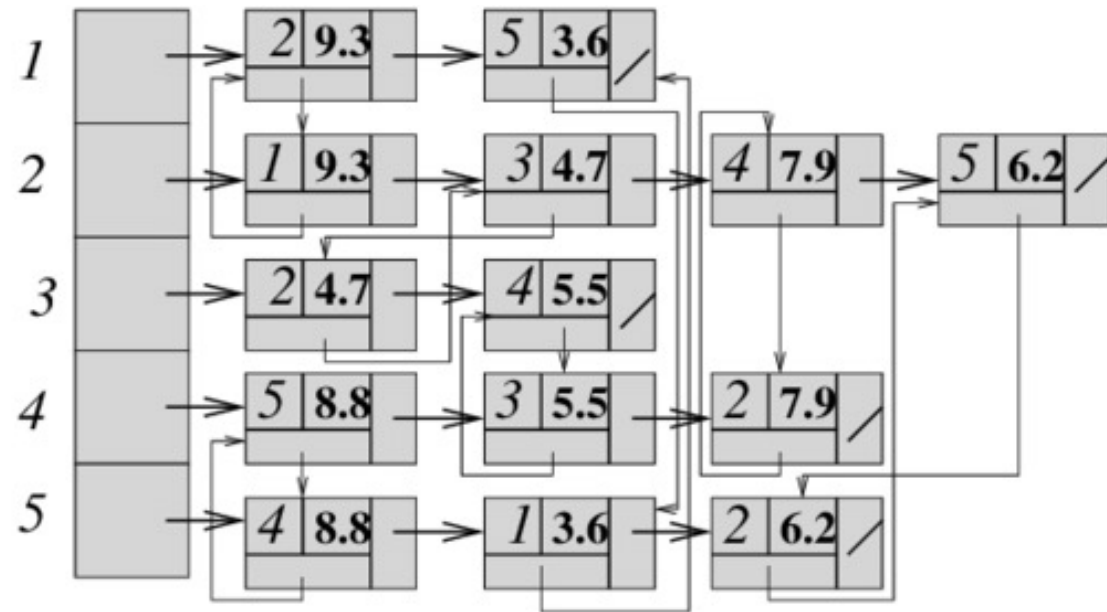
Maintain vertex-indexed array of lists (implementation omitted)



Undirected Weighted Graph - Adjacency List



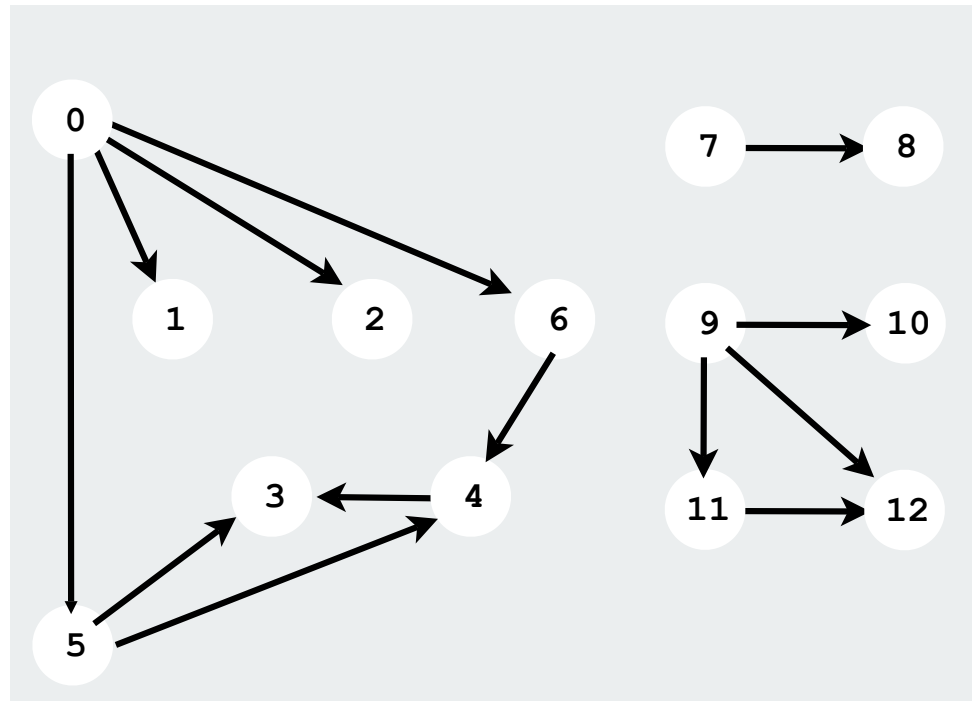
(a)



(b)

Directed Graphs

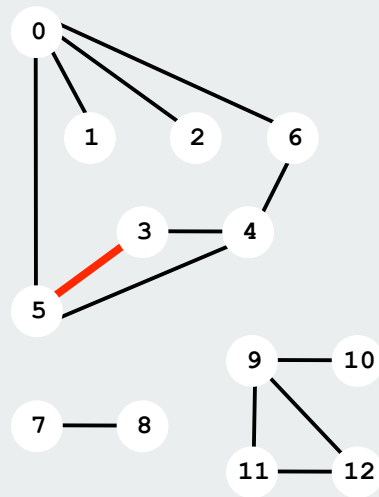
- A **directed graph** is a set of vertices and a collection of directed edges.
- Edges are one way



Graph Traversal

Maintain a two-dimensional $v \times v$ boolean array.

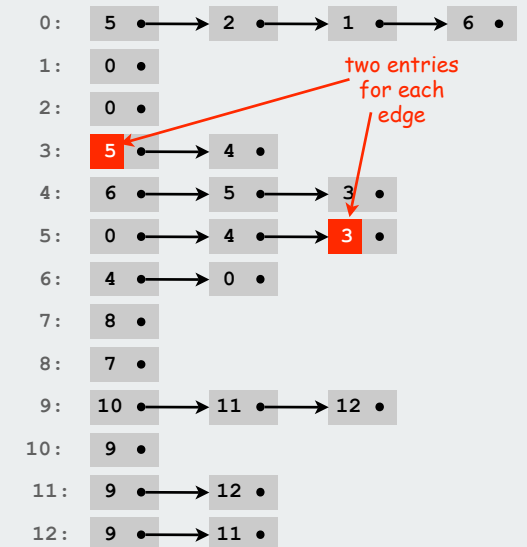
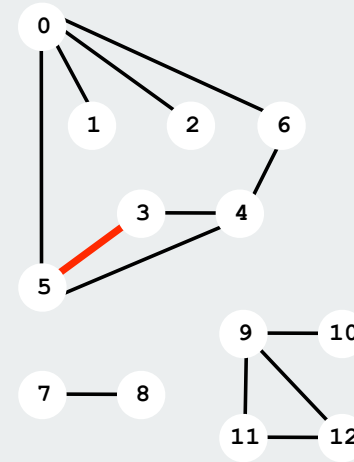
For each edge $v-w$ in graph: $\text{adj}[v][w] = \text{adj}[w][v] = \text{true}$.



two entries
for each
edge

| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
|----|---|---|---|---|---|---|---|---|---|---|----|----|----|
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 2 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 5 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 6 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 7 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 8 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 9 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |
| 10 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 11 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 12 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |

Maintain vertex-indexed array of lists (implementation omitted)



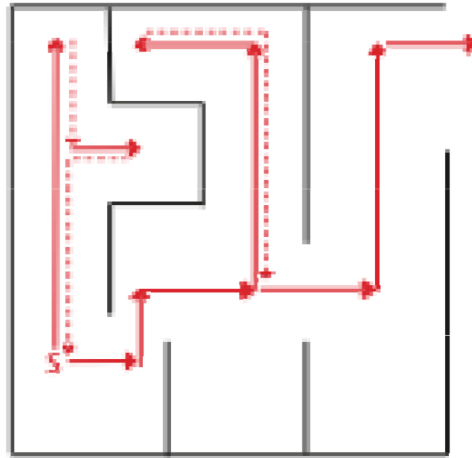
Graph Traversal

- Visit every edge and node in the graph in a systematic way
 - **Depth-first search.** Put unvisited vertices on a **stack**.
 - **Breadth-first search.** Put unvisited vertices on a **queue**.

Graph Traversal – Depth First Search (DFS)

- Like exploring a maze
- From current vertex, move to another
- Until you get stuck
- Then backtrack till you find a new place to explore

- e.g “left-hand” rule



Trémaux Maze Exploration

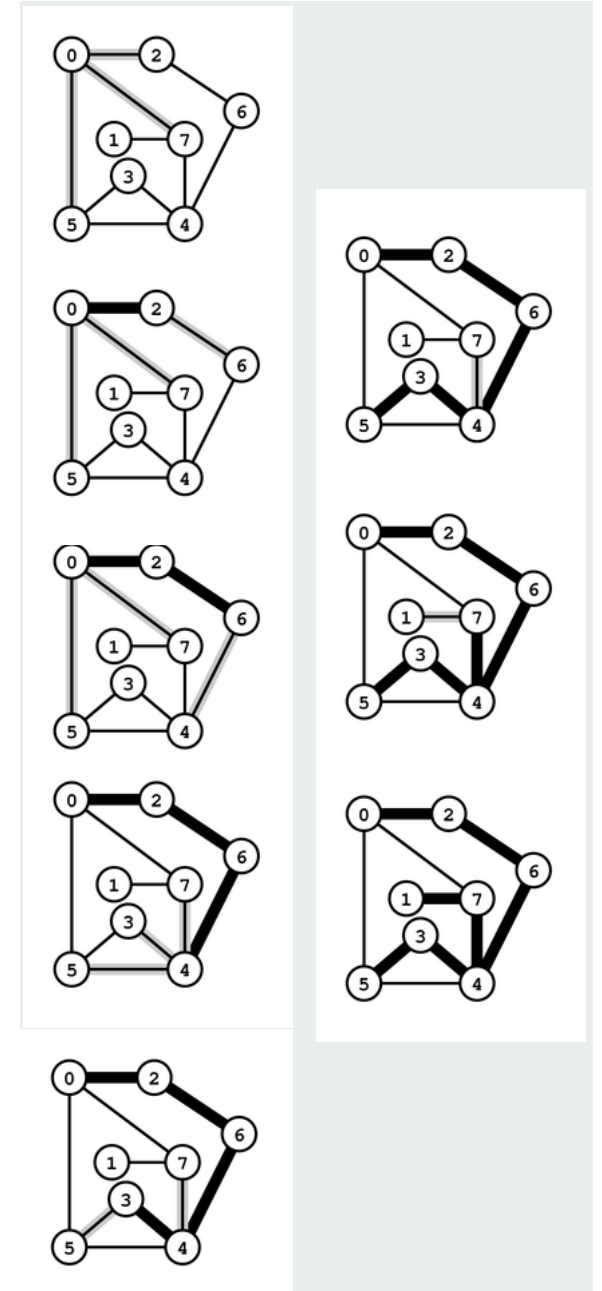
- Unroll a ball of string behind you.
- Mark each visited intersection by turning on a light.
- Mark each visited passage by opening a door

First use? Theseus entered labyrinth to kill the monstrous Minotaur;
Ariadne held ball of string.



Graph Traversal – Depth First Search (DFS)

- **Goal.** Systematically search through a graph.
- **Idea.** Mimic maze exploration.
- **Typical applications.**
 - find all vertices connected to a given s
 - find a path from s to t

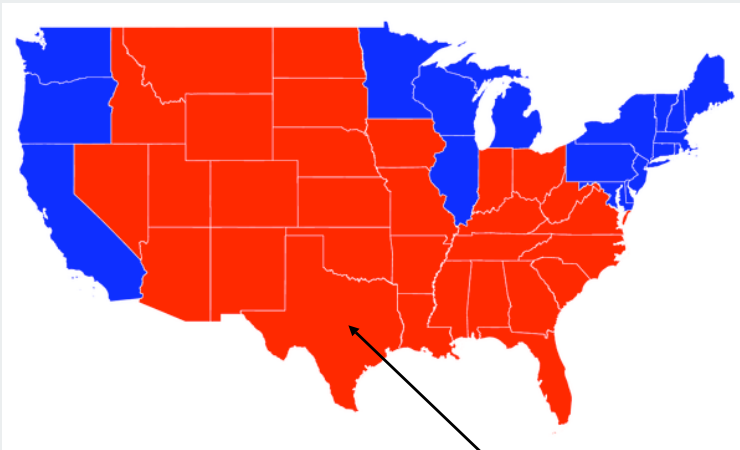


Flood Fill

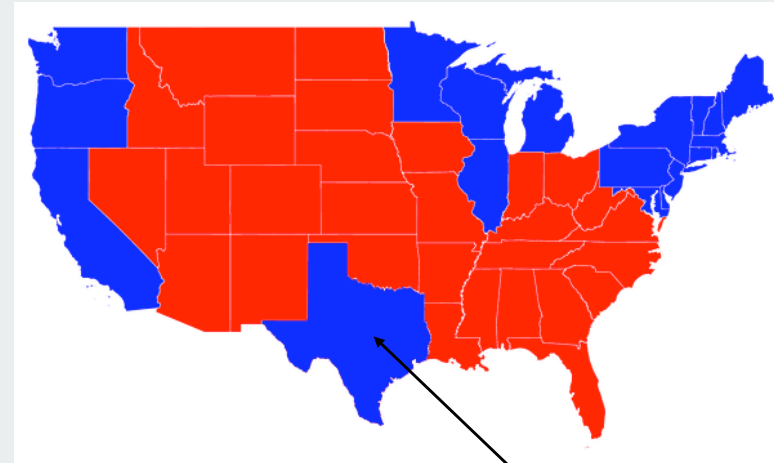
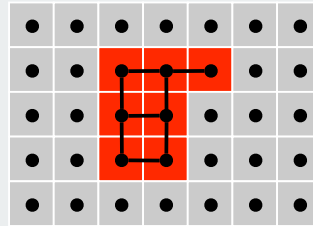
Change color of entire blob of neighboring **red** pixels to **blue**.

Build a **grid graph**

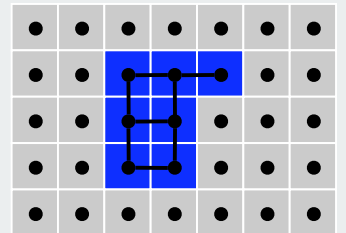
- vertex: pixel.
- edge: between two adjacent lime pixels.
- blob: all pixels connected to given pixel.



recolor red blob to blue



recolor red blob to blue



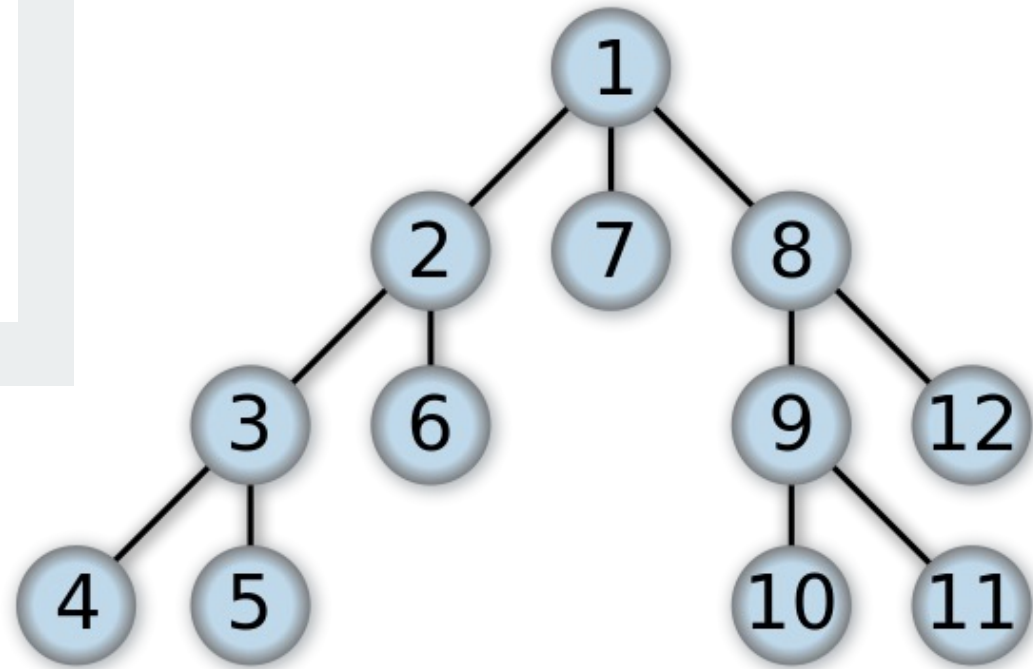
Graph Traversal – Depth First Search (DFS)

DFS (to **visit** a vertex s)

Mark s as visited.

Visit all unmarked vertices v adjacent to s .

↑
recursive

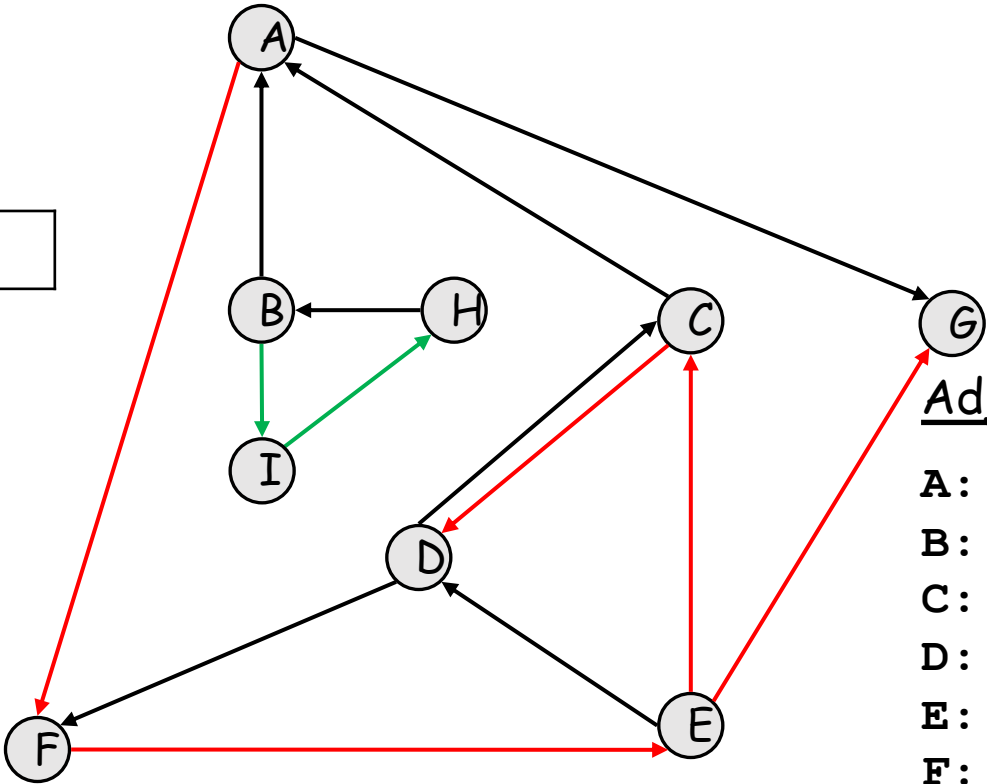


DFS Traversal Example

visited nodes array:

| | | | | | | | | |
|---|---|---|---|---|---|---|---|---|
| A | B | C | D | E | F | G | H | I |
| | | | | | | | | |

How many connected components does the graph have ?
Starting Point is A



Adjacency Lists

- A: F G
- B: A I
- C: A D
- D: C F
- E: C D G
- F: E
- G:
- H: B
- I: H

DFS (to visit a vertex *s*)

Mark *s* as visited.

Visit all unmarked vertices *v* adjacent to *s*.

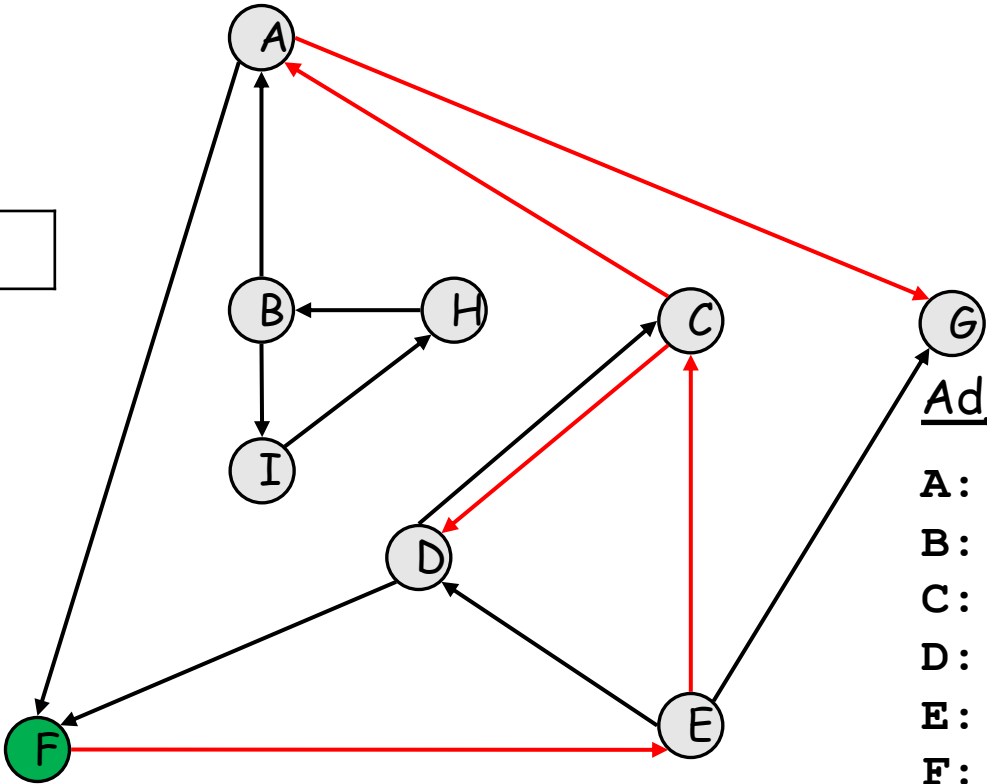
↑
recursive

DFS Traversal Example

visited :

| A | B | C | D | E | F | G | H | I |
|---|---|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 1 | 1 | 1 | 1 | 0 | 0 |

Is there any path from F to I



Adjacency Lists

A: F G
B: A I
C: A D
D: C F
E: C D G
F: E
G:
H: B
I: H

DFS (to visit a vertex *s*)

Mark *s* as visited.

Visit all unmarked vertices *v* adjacent to *s*.

↑
recursive

Graph Traversal – Depth First Search (DFS)

- Recursive DFS:

DFS(G, s)

{

mark s as visited

for all neighbors w of s in Graph G

if w is not visited

DFS(G, w)

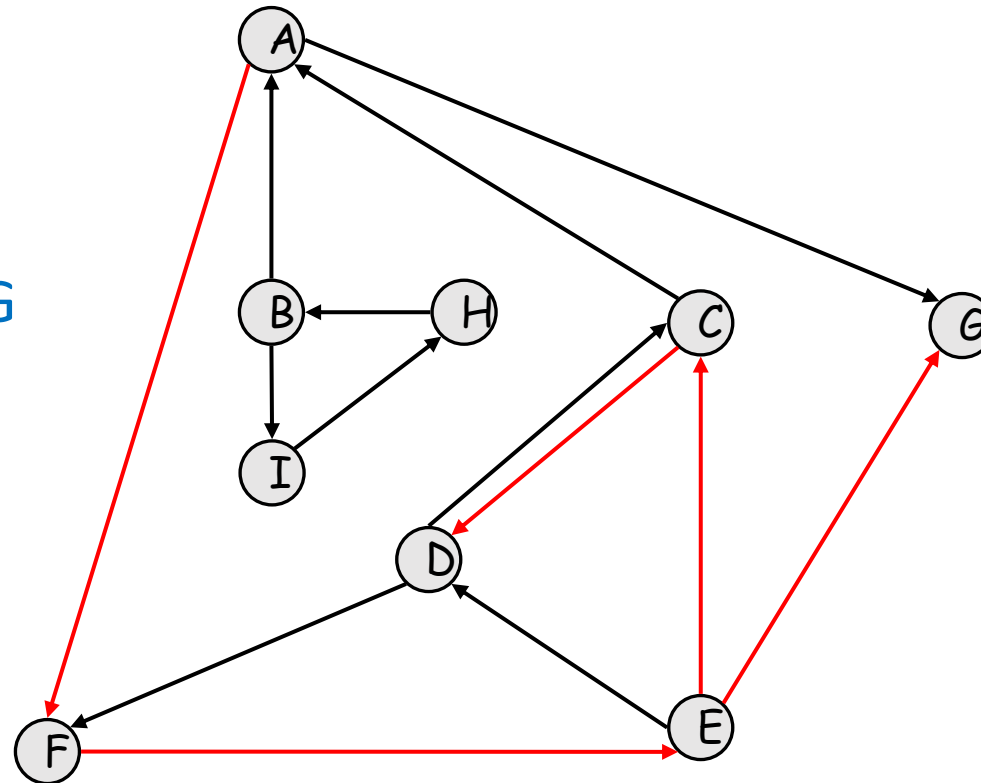
}

Is there a path from A to D ?

A F E C D

Is there a path from A to B ?

A F E C D G



Adjacency Lists

| | |
|----|-------|
| A: | F G |
| B: | A I |
| C: | A D |
| D: | C F |
| E: | C D G |
| F: | E |
| G: | |
| H: | B |
| I: | H |

Depth First Search (DFS)

- To visit a node v :
 - mark it as visited
 - recursively visit all unmarked nodes w adjacent to v
- To traverse a Graph G :
 - initialize all nodes as unmarked
 - visit each unmarked node
- Running time.
 - $O(E)$ since each edge examined at most twice
 - usually less than V to find paths in real graphs

Graph Traversal – Depth First Search (DFS)

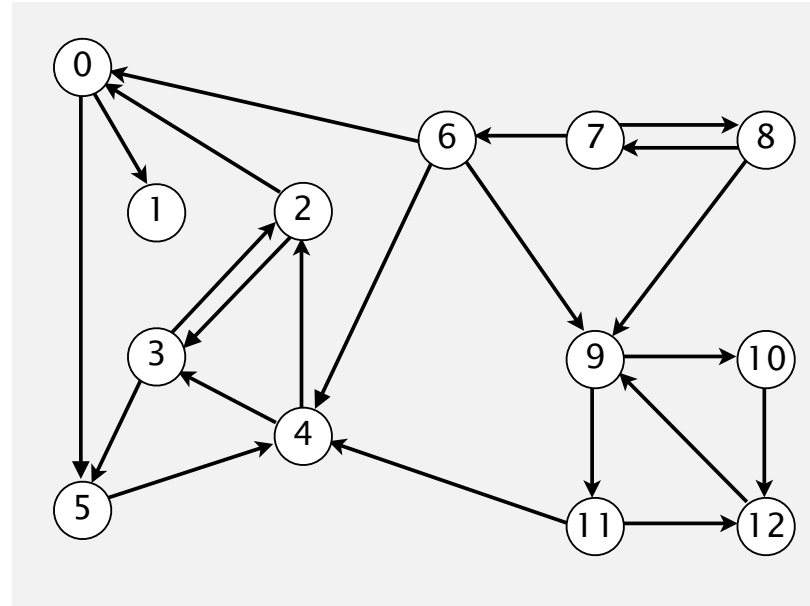
- Non-Recursive DFS: using STACK

```
DFS(G, s)
{
    Stack.push(s)
    mark s as visited
    while(Stack is not empty)
    {
        v = Stack.pop();
        for all neighbors w of v in Graph G
            if w is not visited
            {
                STACK.push(w)
                mark w as visited
            }
    }
}
```

DFS Traversal Example

Adjacency Lists

- 0 : 1 5
- 1 : -
- 2 : 0 3
- 3 : 2 5
- 4 : 2 3
- 5 : 4
- 6 : 0 4 9
- 7 : 6 8
- 8 : 7 9
- 9 : 10 11
- 10 : 12
- 11 : 4 12
- 12 : 9



Is there a path from node 0 to node 9

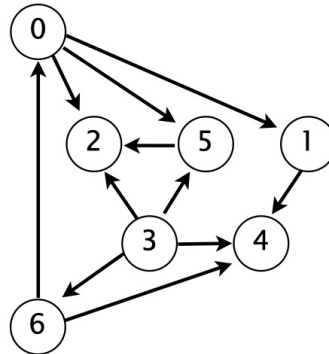
DFS Traversal from node 0 : 0 1 5 4 2 3

Topological Sort

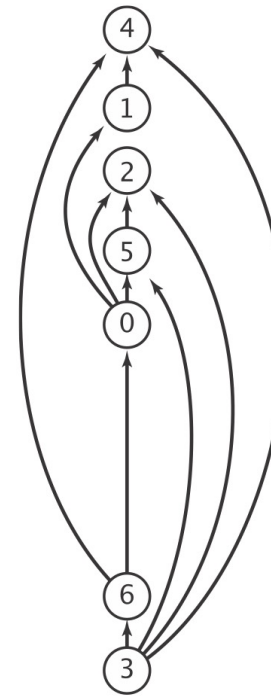
Goal: Given a set of tasks to be completed with precedence constraints, in which order should we schedule the tasks?

- 0. Algorithms
- 1. Complexity Theory
- 2. Artificial Intelligence
- 3. Intro to CS
- 4. Cryptography
- 5. Scientific Computing
- 6. Advanced Programming

tasks



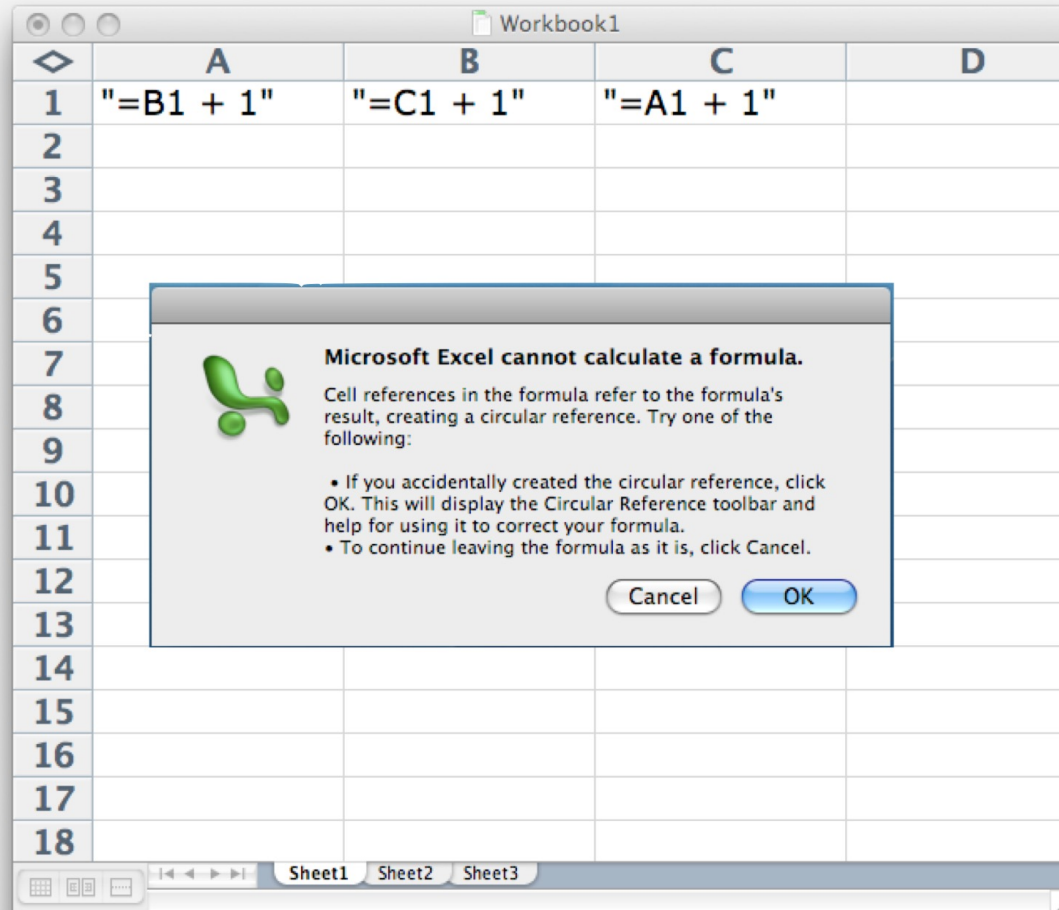
precedence constraint graph



feasible schedule

Directed cycle detection application: spreadsheet recalculation

Microsoft Excel does cycle detection (and has a circular reference toolbar!)



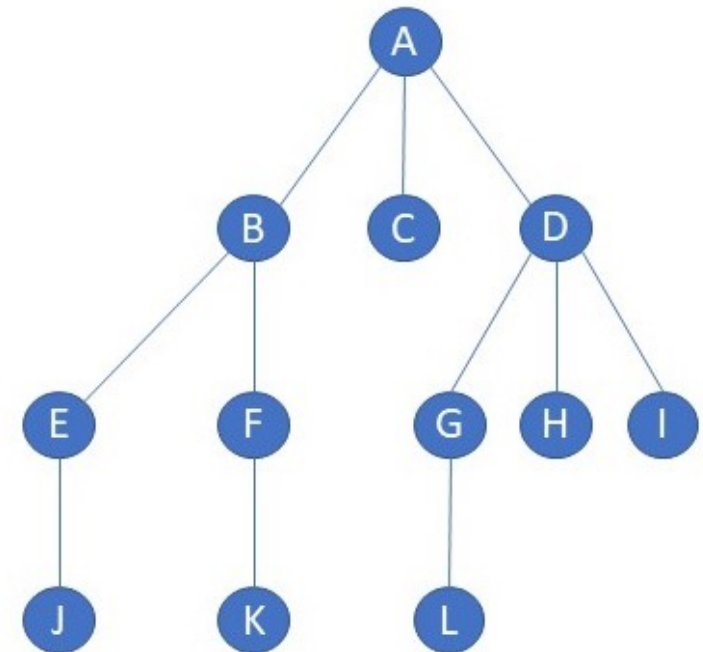
Graph Traversal – Breadth First Search (BFS)

BFS (from source vertex s)

Put s onto a FIFO queue.

Repeat until the queue is empty:

- remove the least recently added vertex v
- add each of v 's unvisited neighbors to the queue, and mark them as visited.



Breadth First Search (BFS)

Graph Traversal – Breadth First Search (BFS)

BFS (from source vertex s)

Put s onto a FIFO queue.

Repeat until the queue is empty:

- remove the least recently added vertex v
- add each of v 's unvisited neighbors to the queue, and mark them as visited.

Visited :

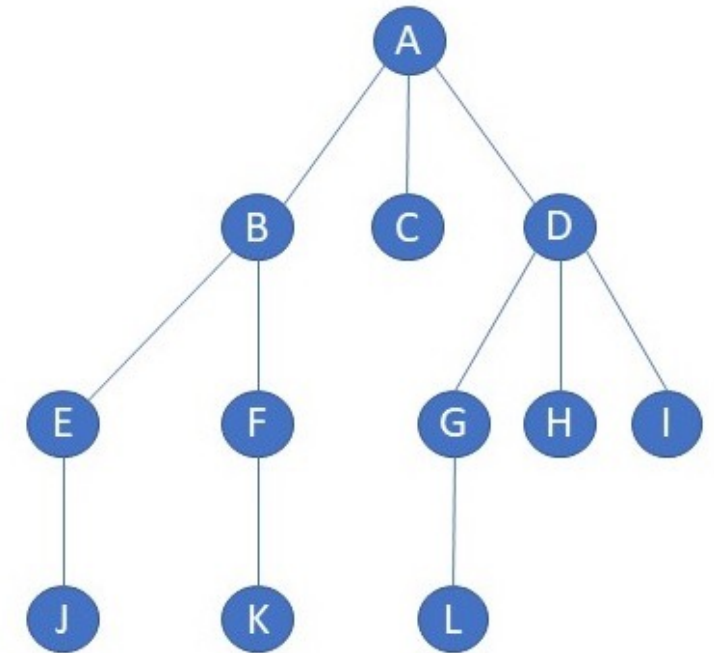
| A | B | C | D | E | F | G | H | I | J | K | L |
|---|---|---|---|---|---|---|---|---|---|---|---|
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 0 |

Path :

| A | B | C | D | E | F | G | H | I | J | K | L |
|---|---|---|---|---|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 1 | 2 | 2 | 2 | 2 | 2 | 3 | 3 | 0 |

Shortest Path from A to K

Queue : A B C D E F G H I J K L



Breadth First Search (BFS)

BFS Traversal Example

visited :

| A | B | C | D | E | F | G | H | I |
|---|---|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 1 | 1 | 1 | 1 | 0 | 0 |

BFS Traversal :

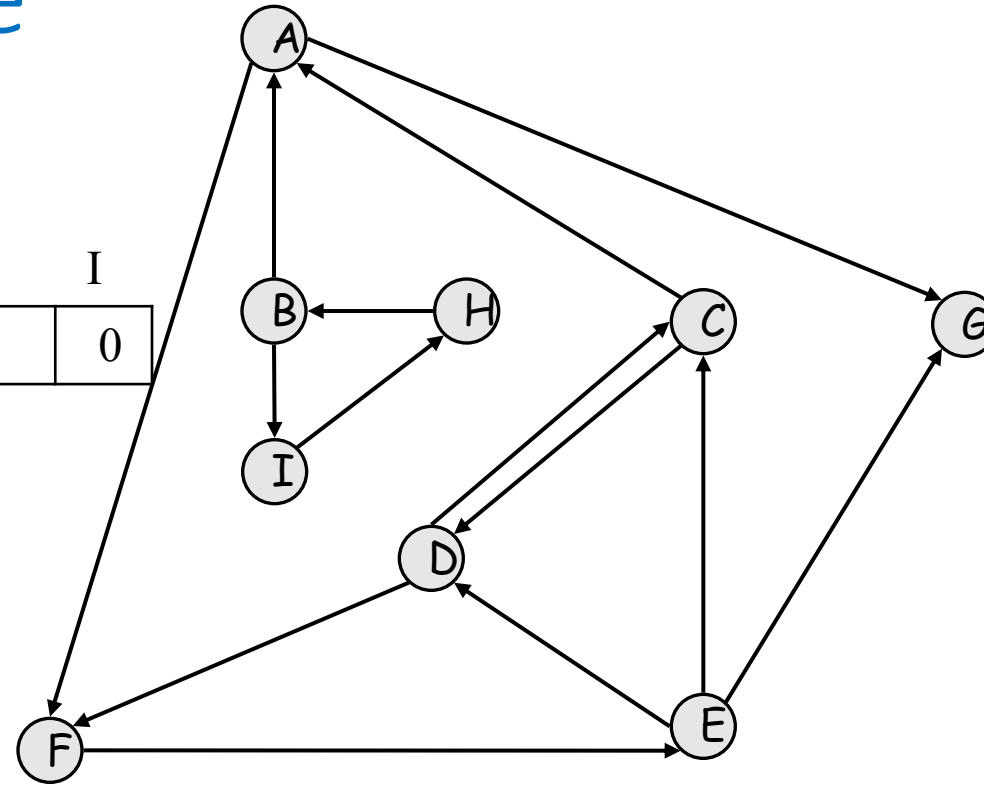
A F G E C D

BFS (from source vertex s)

Put s onto a FIFO queue.

Repeat until the queue is empty:

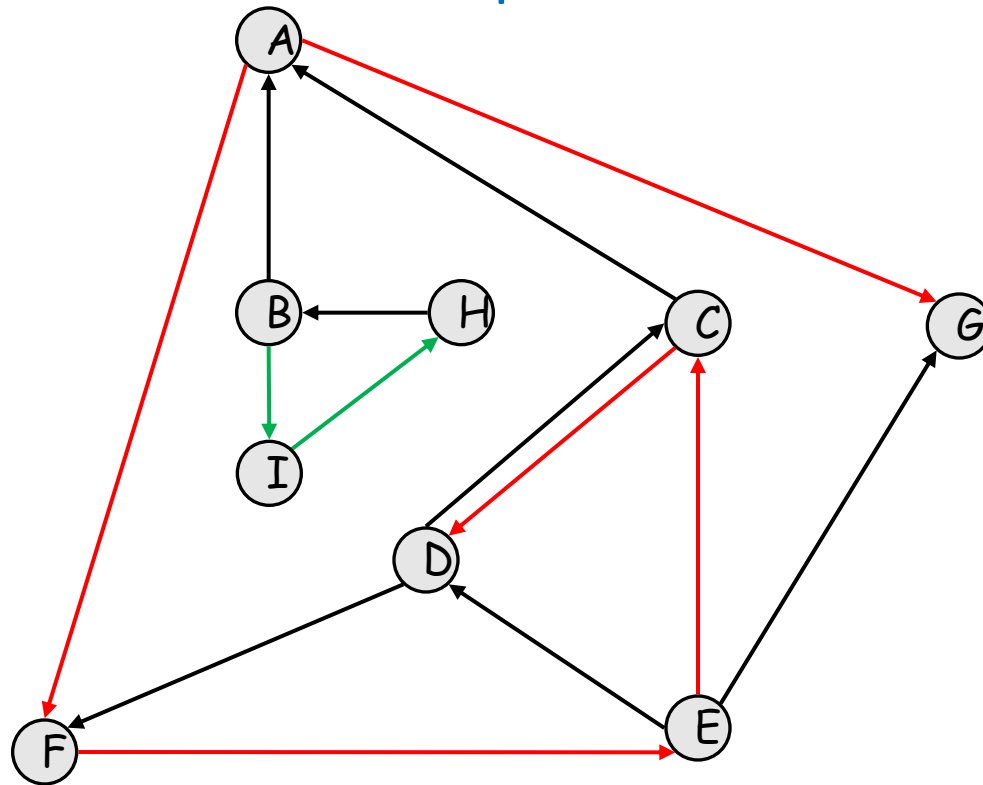
- remove the least recently added vertex v
- add each of v's unvisited neighbors to the queue, and mark them as visited.



Adjacency Lists

A: F G
B: A I
C: A D
D: C F
E: C D G
F: E
G:
H: B
I: H

BFS Traversal Example



Adjacency Lists

A: F G
B: A I
C: A D
D: C F
E: C D G
F: E
G:
H: B
I: H

A F G E C D

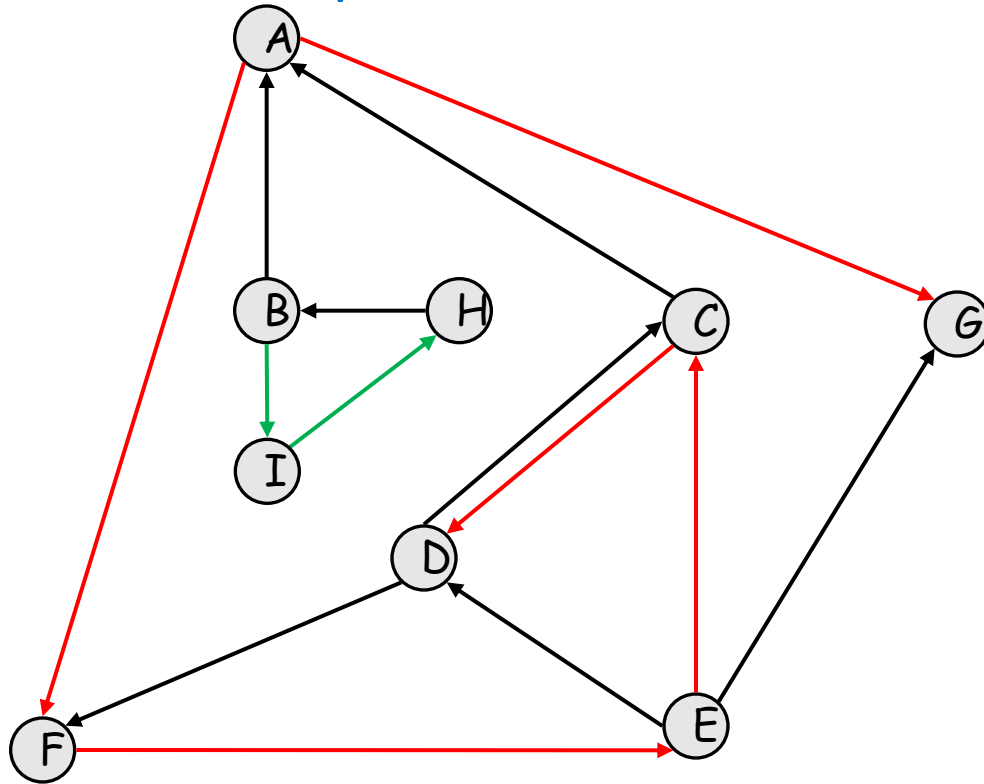
B I H

Graph Traversal – Breadth First Search (BFS)

- BFS : using QUEUE

```
BFS(G, s)
{
    Queue.enqueue(s)
    mark s as visited
    while(Queue is not empty)
    {
        v = Queue.dequeue();
        for all neighbors w of v in Graph G
            if w is not visited
            {
                Queue.enqueue(w)
                mark w as visited
            }
    }
}
```

Traversal Example



Adjacency Lists

A: F G
B: A I
C: A D
D: C F
E: C D G
F: E
G:
H: B
I: H

BFS : A F G E C D

B I H

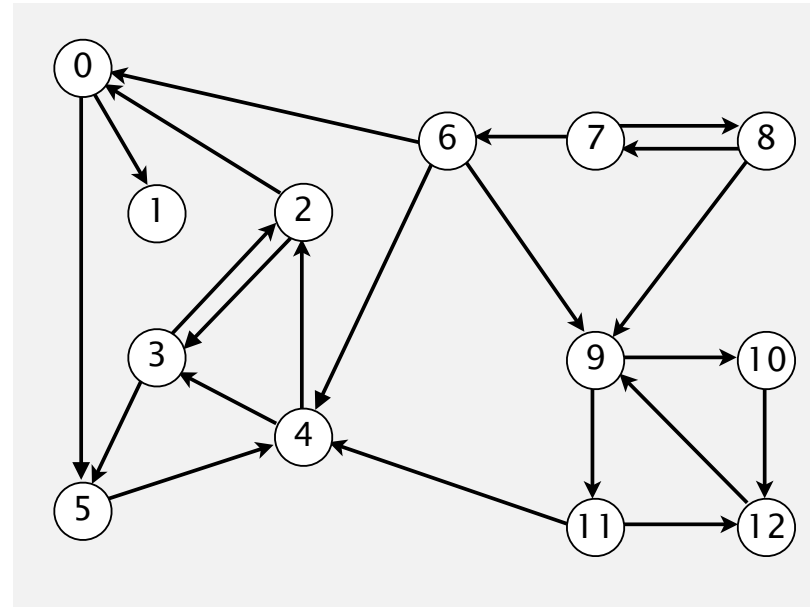
DFS : A F E C D G

B I H

BFS Traversal Example

Adjacency Lists

- 0 : 1 5
- 1 : -
- 2 : 0 3
- 3 : 2 5
- 4 : 2 3
- 5 : 4
- 6 : 0 4 9
- 7 : 6 8
- 8 : 7 9
- 9 : 10 11
- 10 : 12
- 11 : 4 12
- 12 : 9



DFS : **6** 0 1 5 4 2 3 9 10 12 11

BFS : **6 0 4 9 1 5 2 3 10 11 12**

Graph Search Problems

- Is there a path from s to t ?
- Find shortest path (fewest edges) from s to t
- Is there a cycle in the graph ?
- How many connected components exist?