Using AI to Detect Malware in Object Storage

Author: Matthew Battagel, Supervisor: Theodoros Spyridopoulos

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Abstract

Lorem Ipsum

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1. Introduction

Overview

The exponential growth of data generation has made data storage an increasingly important aspect for both individuals and organizations alike. Object storage has emerged as a promising solution due to its ability to store vast amounts of unstructured data in a cost-effective and scalable manner. Unlike traditional storage techniques, object storage stores data as objects with related metadata and unique identifiers, allowing for efficient and cheap storage within buckets.

One of the most widely used object storage platforms is Amazon S3, which provides a highly scalable and reliable solution for storing data. However, an open-source alternative called MinIO has emerged as a promising contender, providing similar features to Amazon S3 while giving customers greater control over their data. MinIO is written in Go and is available for free under the Apache License 3.0 or, for commercial and enterprise purposes, at a reduced cost compared to Amazon S3. (?). MinIO offers a wide range of features, including high performance, data replication, encryption and erasure coding (MinIO, 2023). Most importantly, MinIO is designed to scale out horizontally to ensure that it can handle the demands of large-scale applications.

Scalability is made simple by allowing multiple types of hardware platforms to work together in separate nodes each with their own compute and storage. This is extremely attractive for customers who want to utilises their existing hardware without being tied down to a specific provider. This also applies for customers looking to migrate their data from Amazon S3 to cheaper solution without compromising on the high performance, reliability and scalability of the S3 platform.

While MinIO is a great alternative to Amazon S3, it does not offer any form of malware detection integration. This could put customers off from choosing MinIO as a viable platform to migrate to from Amazon S3 or leave existing users data vulnerable to malware attacks. This project aims to address this issue by integrating a malware detection system into MinIO. An

important goal for the is to negatively impact the scalability or performance as little as possible so that MinIO is still an effective alternative to Amazon S3.

Motivation

Due to the high amount of unstructured data expected to be both written and read to the object store, there are increased risk of encountering malicious files. Therefore malware detection within object storage is crucial in modern cloud storage scenarios. Most popular off-the-shelf object storage platforms, such as AWS, already have integrated third-party antivirus software, such as ClamAV and Sophos (Srinivasan *et al.*, 2022), to mitigate security risks. MinIO on the other hand is vulnerable to malware attacks as it currently does not have any native antivirus integration. This forces customers who require complete virus protection to either not use MinIO or to use potentially costly third-party software. As antivirus scanning is inherently resource intensive, if the software is integrated incorrectly, it could reduce the ability for the storage solution to scale horizontally which negates one of the major benefits of object storage. The purpose of this project is to implement malware detection within MinIO while being mindful to not impact the scalability or performance of the platform.

Project Aims

This project aims to supply an end-to-end solution for detecting malware within the MinIO object storage platform with three main requirements. The solution should be able to detect the latest uploads to the object store and scan them. It should be able to perform this function without significantly impacting the performance of the object store. The solution should also scale alongside MinIO to ensure that it does not bottleneck the object store at high loads. These main three aims can be quantified so that an accurate evaluation of the solution can be made at the end of the project.

- Detection of malware within the object store should match 100% of the malware detection rate as the stand-alone antivirus.
- Performance of the solution to be within 10% of the performance of the stand-alone object store, MinIO.
- Retain the previous metric while both increasing the available resources and changing the platform.

Milestones

There are many key milestones that can be used to measure the progress of the project through the duration of its implementation. It is worth noting that these milestones were created before the implementation of the solution and therefore are subject to change.

- Setup local ClamAV instance
- Setup local MinIO instance
- Setup local Kafka instance
- Detect PUT message on Kafka
- Read JSON of message to find bucket and object path
- Create AV manager service in GoLang

- Request GET on Object using message data
- Scan Object using Clamd
- Act on result of scan Add tags, "scanned", "date_scanned"
- Configuration file using viper
- Add metric collection with Prometheus
- Add structured logging with zap
- Create Database of Audit Logs with postgresql
- Termination of connections to services
- Create tests for each package using mockery
- Prepare for Kubernetes deployment with k3d
- Containerise Aegis with Docker
- Automatically configure and start MinIO, Kafka, Clamd, Postgresql, Prometheus and Aegis with Helm
- Port forward MinIO out of the cluster
- Load balancing / performance analysis for Clamd
- Expose Postgresql and Prometheus for analytics
- Prepare configurations for demo, release etc
- Clean-up version control and write readme.me

These milestones can be broken down even further into smaller tasks. These are best represented in Gantt / burndown charts which can be found in the appendix.

2. Background

Amazon S3 Malware Detection

As MinIO's largest competitor, this project draws a lot of inspiration from Amazon S3s integrated malware detection blog page (Srinivasan *et al.*, 2022). The blog explains Amazons current approach for managing malware detection within their service. Amazon S3 uses a combination of ClamAV and Sophos as their third-party scanning engines due to their out-of-the-box nature. Amazon then gives you the option to use either of these engines or both. The blog goes on to describe the three main interaction mechanisms that Amazon S3 uses to flag files for scanning. Firstly, an API endpoint would be provided to handle all uploads. This forms a queue of uploads which are then scanned before entering the bucket. Next, event-driven scanning is used keep track of all regular file uploads. The antivirus will then scan each file after they have been written to the bucket. Finally, retro-driven scanning is used to scan all existing files within the bucket. The user then has the flexibility to define what types of files should be scanned including defining time windows. This blog has given some useful methodologies of how to keeping track of both incoming and previously scanned files. Creating a system that can match these methods is important for offering a matching level of scalability and security within MinIO.

3. Specification and Design

Specification

The specification for this project is to help guide the project to fufill the aims set out in the previous section. The specification is broken down into three main sections; functional requirements, non-functional requirements and constraints.

Functional Requirements Functional requirements are

Non-Functional Requirements Non-functional requirements are

Constraints

The constraints are the limitations that the solution must adhere to. The main constraint of the project is the strict time limit given to the project. There are a total of 12 weeks to achieve a production ready product which will greatly limit the scope of the project. This means accurately prioritising the features that are most important to the project while also balancing the time spent to implement them. The second constraint is the limited resources available to the project.

Another constraint of the project is that all the external software used must be open source / available for commercial use under license or fee. This is to ensure that the project is legally viable if the solution was to be used commercially.

The final constraint is that the solution must

Architecture

Given the specification above, various potential architectures can be created and evaluated. An optimal design will then be chosen based on which design fits the specification best.

Design 1

This design makes use of the performance benefits of MinIO by allowing PUTs to be initially written to the bucket without being scanned. The design then uses a event queue compatible with MinIO to keep track of all the files that have been uploaded. The queue is then used to trigger a scan of the file once an antivirus is available. The design is shown in figure 1.

This design has many benefits over other potential implementations. Firstly, it uses the storage provided by MinIO to store all incoming files without having to manage a separate storage solution. This removes a lot of complexity from the solution by not having to account for a number of failure conditions that could occur with a high availability, production ready storage solution. For example, the solution would not be responsible for handling partial writes, loss of data, or data corruption. Removing this responsibility allows the solution to focus on the core functionality of the project, the scanning of

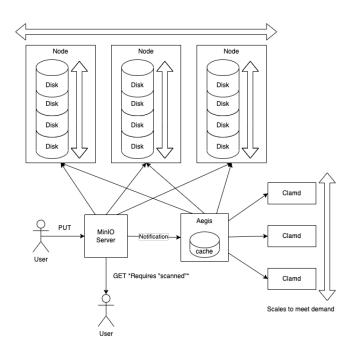


Figure 1. Post-Write Architecture

files, which is essential for keeping the project within the time constraints.

Secondly, the design also makes use of the integrated event queue provided by MinIO. This again removes responsibility from the solution by differing the scalability and reliability requirements of an event queue to MinIO".

Lastly, having Aegis dispatch the files to a scalable number of antivirus scanners allows the solution to scale to meet the demands of the system. This meets a key requirement as the solution is expected to have the capacity for a large number of operations. This method does require the use of a load balancer to effectively distribute the load across the available antivirus scanners.

The design also has a number of drawbacks. Firstly, the design still requires a small about of cache to temporarily store the object when it is being dispatched to the antivirus. Provisioning of this cache has to be large enough to handle the largest file possible to be uploaded to the object store. In reality, this cache would be provisioned even larger to allow for the temporary storage of multiple objects while multiple scans are being performed asynchronously. In addition, the cache needs to be large enough to ensure that the system does not become overwhelmed by the number of objects being scanned as the system scales. This is a minor issue as store capacity is cheap and the provisioning of the cache easy to scale up.

The second drawback is that, for each event, Aegis makes a get request for the object to be scanned. This effectively doubles the number of requests made to the object store. This also means that Aegis must have the ability to get any file expected to be scanned and therefore must have access to the whole storage network. The impact of this drawback is mitigated as the solution is expected to be deployed on the same network as the object store which should reduce the latency of each request made by Aegis. However, this still leaves MinIO to handle twice as many requests with the performance loss being noticed mainly on more distributed storage topologies.

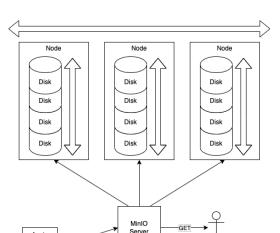
Thirdly, the design only allows for a single Aegis instance to dispatch all incoming objects to available scanners. This is a potential bottleneck for the system as this instance could become overwhelmed by the number of requests it is receiving. This is a minor issue as the dispatching of objects to scanners is not as performance intensive as other areas of the solution, such as the actual scanning, and therefore it is not expected to be a major bottleneck.

Lastly, any object uploaded to the store will have a certain period of time where it remains unchecked. In this time, the user could potentially download an unscanned object or the object could cause harm to the store before it is detected. Although the handling of infected objects is out of scope, in an actual implementation of the solution, the user could be made unable to download unscanned objects until they have been scanned.

Design 2

This design created a wrapper around MinIO that the user interacts with instead of MinIO. This means that all puts go through Aegis before being uploaded to the object store. The design is shown in figure 2.

The main benefit of this design is that the user interacts only with Aegis when uploading files. This means that all incoming files can



be stored within a temporary storage before ever entering the object store. This offers the best protection against malicious files as the user cannot ever download an unscanned or infected file as it is never uploaded to the object store. Infected files can then either be deleted or moved to a separate quarantine store for analysis.

This designs main advantage also comes with a major drawback. This design requires Aegis to handle the full throughput of all the puts to the system. Aegis then has the full responsibility of being available to all puts and, in a failure scenario, to handle the recovery of the system. Additionally, the cache provisioned must be large enough to handle the largest files at maximum throughput with extra room for unexpected delays. This negatively affects the scope of the project by requiring the solution to prioritise features that are already covered by MinIO.

Because MinIO is dependent on Aegis to handle the puts, MinIO must wait to be passed incoming objects sequentially after Aegis has finished processing the previous object. This removes the potential for aggregate performance where

Design 3

Design three is very similar to the second design, however, instead of wrapping outside the MinIO service, it intercepts the writes from the client before objects are written to the object store. With this interception, Aegis can scan the object and decide whether to allow the object to be written to the store or to quarantine the object. The design is shown in figure 3.

This design has similar benefits as the second design. It offers the most protection against malicious files by never allowing either unscanned or infected objects to be stored in the object store. However, it also has similar drawbacks. This is because Aegis is still in sequence with MinIO meaning that for optimal throughput, Aegis would need to match the performance of MinIO.

Similar to the upload queue design, this design also requires Aegis to have a large cache to handle the largest files at maximum throughput. This cache must also be large enough to handle the number of objects being put by MinIO into the store. This issue cannot be mitigated without the risk of compromising performance at increased loads.

However, this design does have an advantage over the second design as there is less responsibility

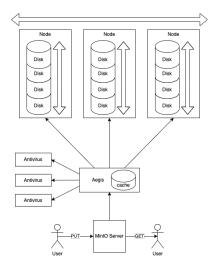


Figure 3. Write Interception Architecture

placed on Aegis to be as failure tolerant. MinIO is still directly responsible for accepting objects into the store and therefore is still responsible for the recovery of the system in a failure scenario. This allows the scope to focus on more related features to malware scanning.

Design 4

The final design distributes Aegis onto each node in the object store. Each Aegis instance relies on. This means that each node has a local instance of Aegis that is responsible for scanning objects before they are written to the store. The design is shown in figure 4.

This design makes use of the distributed nature of MinIO to match the demand when scaling out the system. As more nodes are added, more Aegis instances are added to handle the increased scanning demand.

The major drawback of this design is that it replies on the ability to scan whole files by only using data on a single node. In actual implementations, MinIO makes use of erasure coding to add increased redundancy to the store. Erasure coding splits files into multiple parts, and calculates an additional part that stores the parity. These parts are then distributed among all nodes in the system. This means that the Aegis instance on each node only has access to the part available on their node and

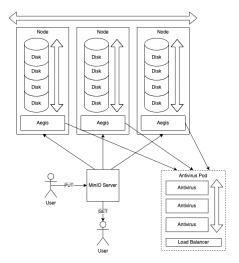


Figure 4. Antivirus per Node Architecture

therefore will not be able to reconstruct the whole file.

Optimal Design

4. Implementation

Event Queue

Packaging

- 5. Results and Evaluation
- 6. Product Issues / Future work
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