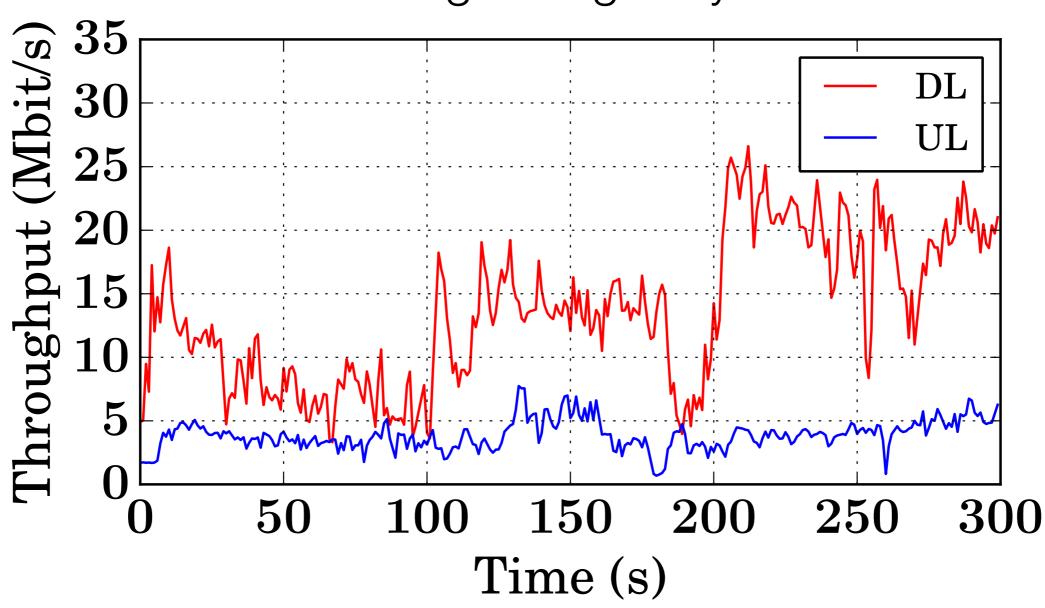
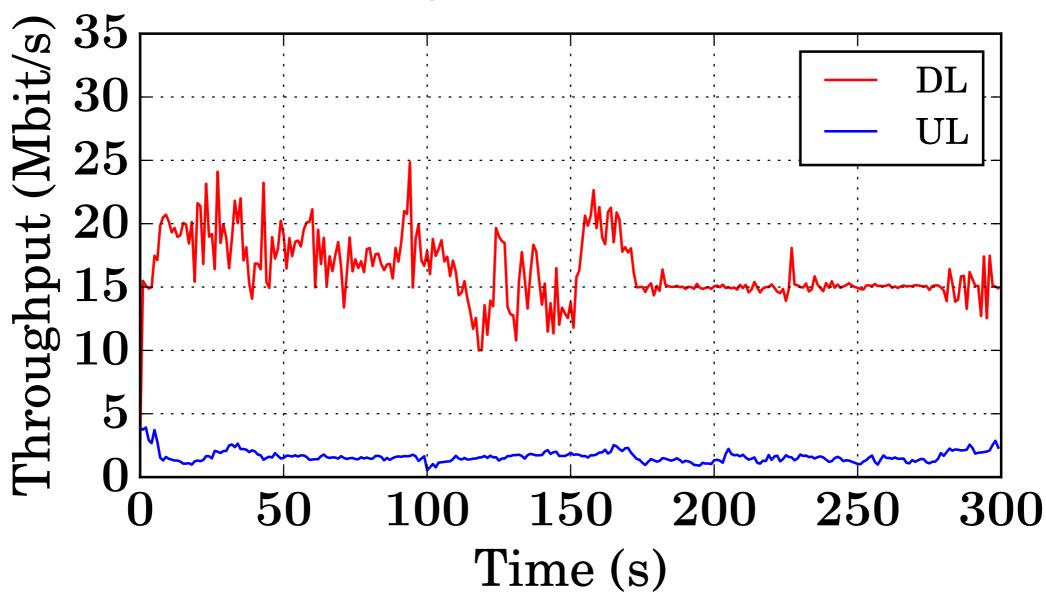
Adaptive Congestion Control for Unpredictable Cellular Networks

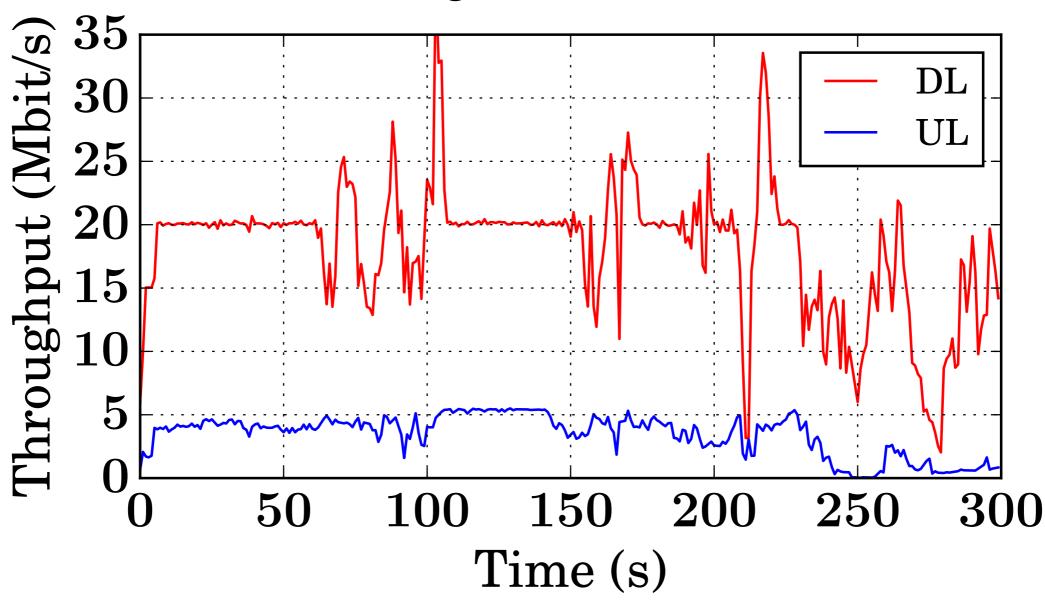




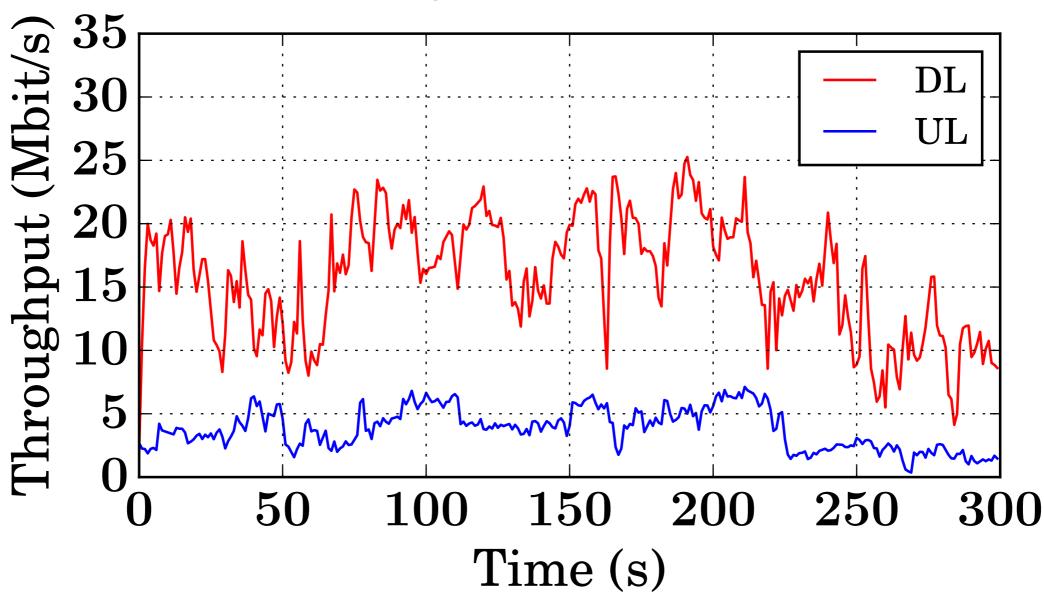






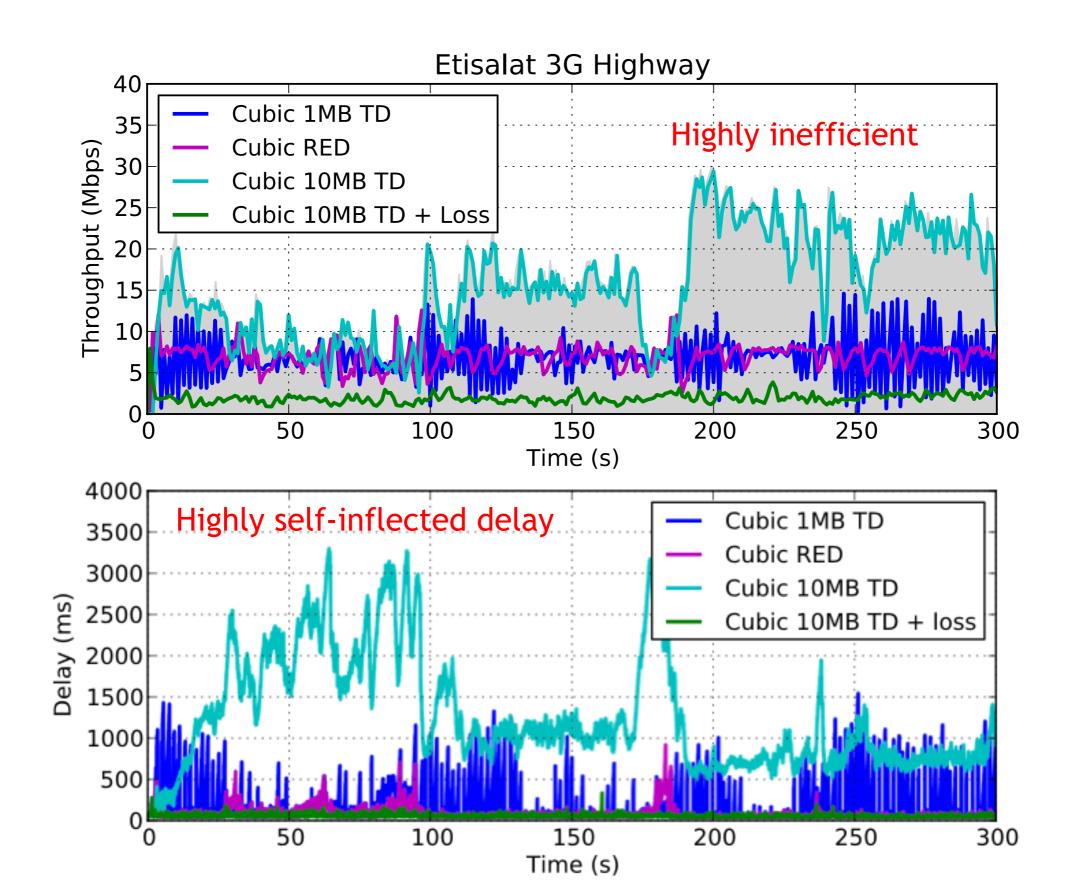






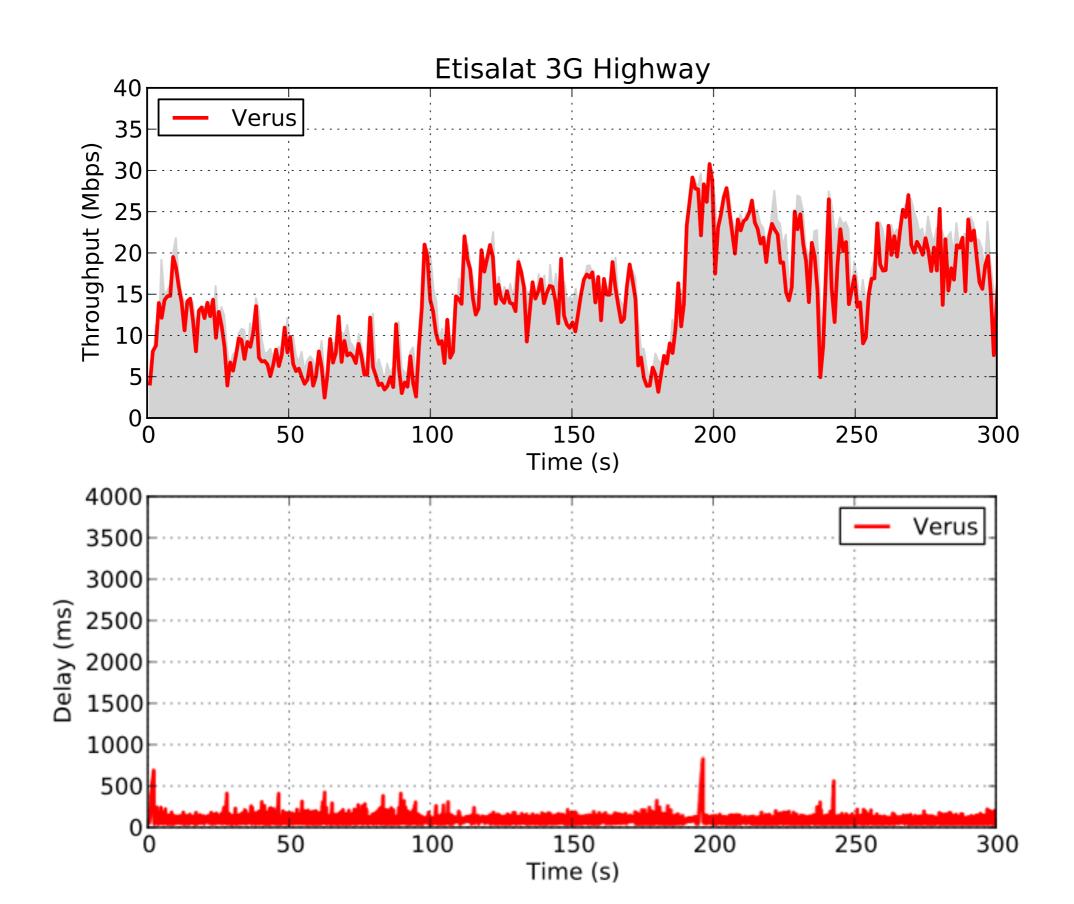
Why do we need a new congestion control protocol for cellular?

TCP cubic over Cellular



Can we do better?

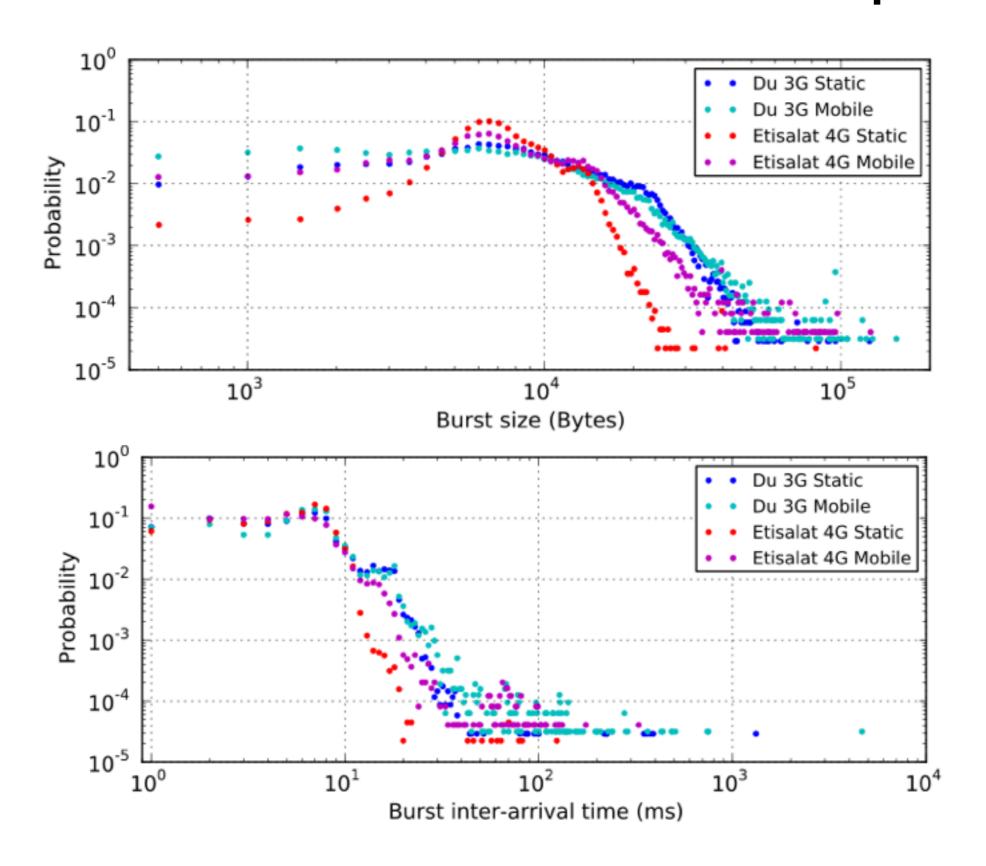
Verus



Why is cellular different?

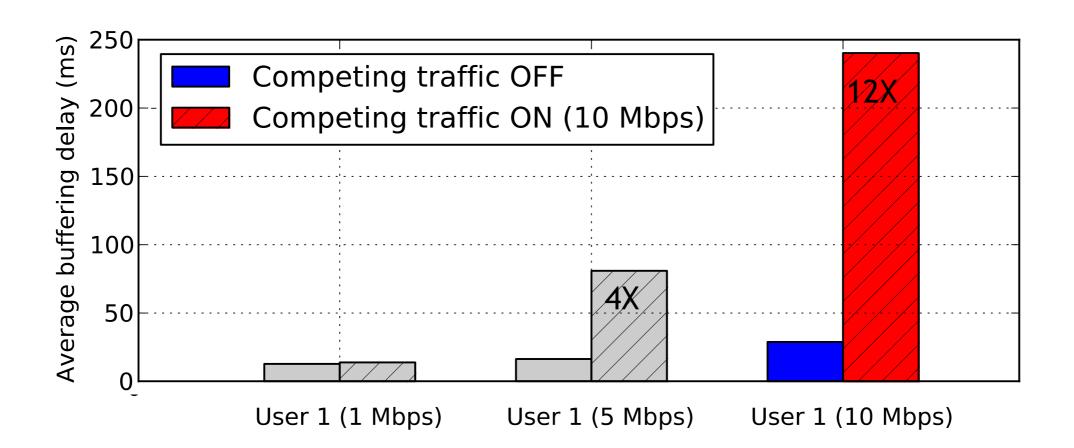
- Sprout* was a big motivation for our work
- Cellular networks experience:
 - Rapidly varying link rates
 - Occasional multi second outages
 - Deep packet queues
- Sprout assumes:
 - Stochastic modeling of the cellular channel
 - Self-interaction is the dominant factor in end-to-end delay rather than competing traffic

Cellular channels are hard to predict



Competing traffic has a major impact

- Two users over real 3G network:
 - User 1 is constantly receiving a downlink stream
 - User 2 with ON/OFF 10 Mbps downlink stream



Verus

 Is an adaptive congestion control protocol for cellular networks

- Design goals:
 - 1. Track fast channel changes
 - 2. Balance throughput and delay
 - 3. Provide fairness between competing flows

Verus design

- Congestion control protocol requires some input:
 - Channel/network state modeling
 - Explicit network cooperation e.g. ECN
 - Loss detection
 - Delay feedback
- We use delay feedback
 - Only change end nodes
 - Proactively avoid congestion
 - Cheap signaling overhead

Verus design

- Don't do channel prediction/modeling
- Build on TCP concepts:
 - Use slow start
 - Use Multiplicative Decrease (MD) on packet loss
 - Replace Additive Increase (AI) with a step based increase/decrease
- Learn the relationship between delay and sending window
 - Delay curve

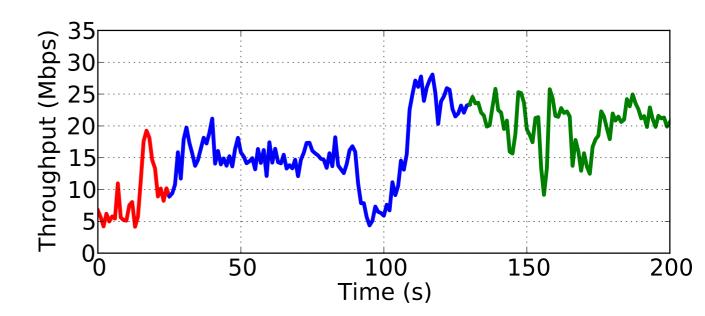
Verus key idea

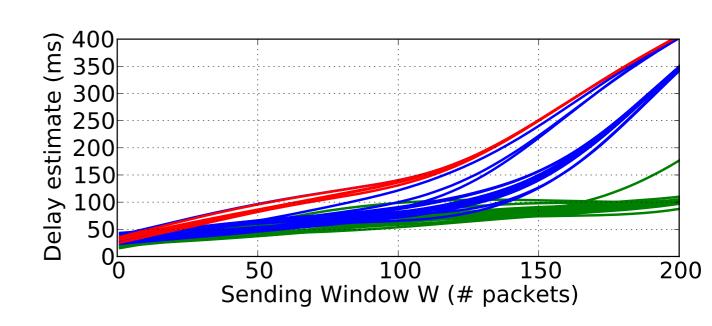
- Every 5 ms epoch we:
 - Estimate network delay
 - Infer the number of packets to be sent (W) to avoid congestion

$$W(t+1) = f(d(t) \pm \delta(t))$$
 Delay curve Estimated delay (D_{est})

Delay curve

- A way to track network changes
- Reflects relationship between sending window and network delay
- Verus dynamically learns the network state
 - Through delay feedback (ACKs)





Verus estimated delay

• The estimated delay ($D_{est,i+1}$) at every epoch is:

$$D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{if } \Delta D_i > 0 \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{if } \Delta D_i > 0 \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i} + \delta_2 & \text{otherwise} \end{cases} & \text{otherwise} \\ D_{est,i} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = \begin{cases} D_{est,i+1} - \delta_1 & \text{otherwise} \\ D_{est,i+1} = S_{est,i+1} & \text{otherwise} \\ D_{est,i+1} & \text{otherwise} \\ D_{est,i+1} = S_{est,i+1} & \text{otherwise} \\ D_{est,i+1} & \text{otherwise} \\ D_{est,$$

- δ_1 and δ_2 : decrease and increase steps (1 ms and 2 ms)
- ΔD : delay difference between $D_{max,i}$ and average $D_{max,i-1}$
- $D_{max,i}$: maximum delay during last epoch

Goal 1: Tracking fast channel changes

Slow start:

- Every ACK: add a point (W, delay)

Build delay curve:

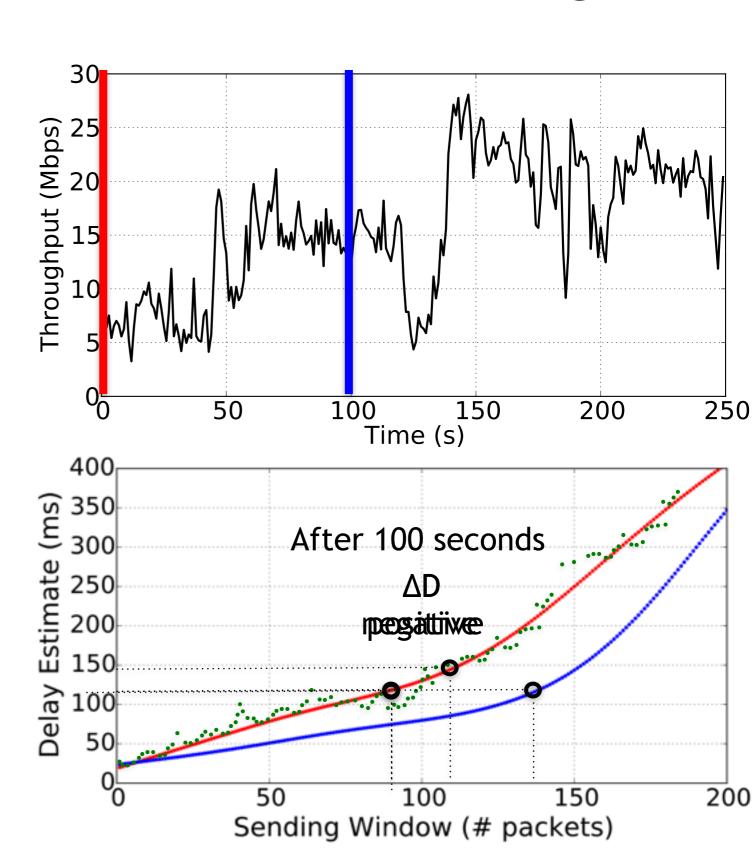
- Cubic spline interpolation

Verus control loop:

- every epoch 5 ms

Rebuild delay curve:

- every 1 second



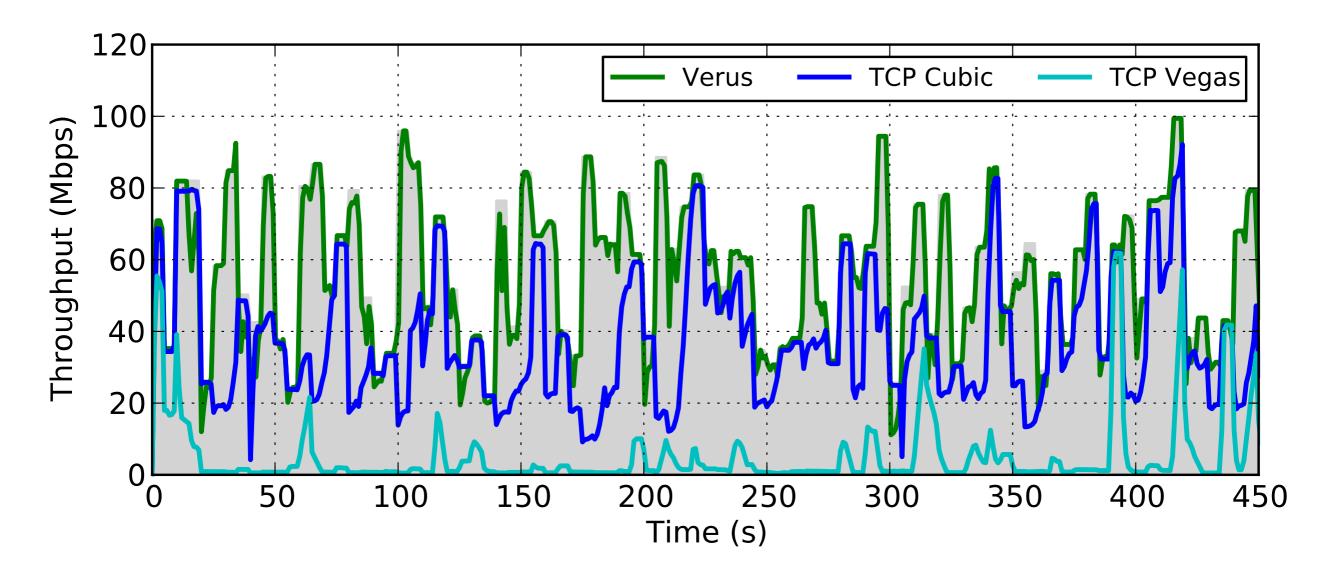
Goal 1: Tracking fast channel changes

Every 5 sec:

Link: 10-100 Mbps

Round trip time: 10-100 ms

Losses: 0-1 %



Goal 2: Balance throughput and delay

• Verus D_{est} is upper bounded by R

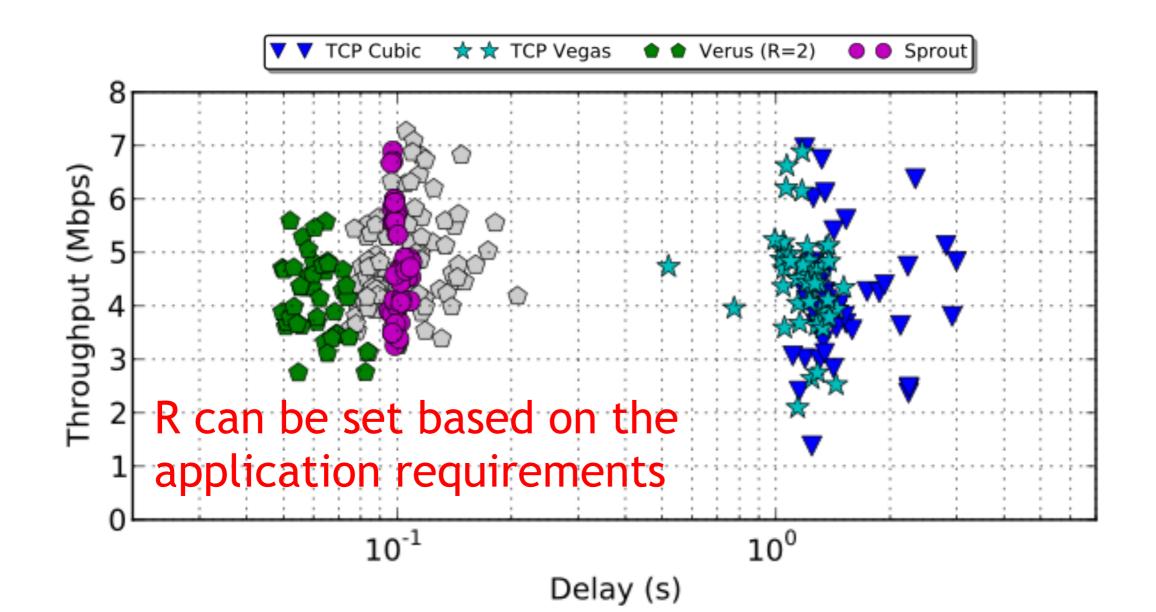
$$D_{est,i+1} = egin{cases} D_{est,i} - \delta_2 & ext{if } rac{D_{max,i}}{D_{min}} > R \ \\ D_{est,i} - \delta_1 & ext{elif } \Delta D_i > 0 \ \\ D_{est,i} + \delta_2 & ext{otherwise} \end{cases}$$

- R defines the ratio between max and min delay of the network
- Setting R specifies the trade-off between throughput and delay

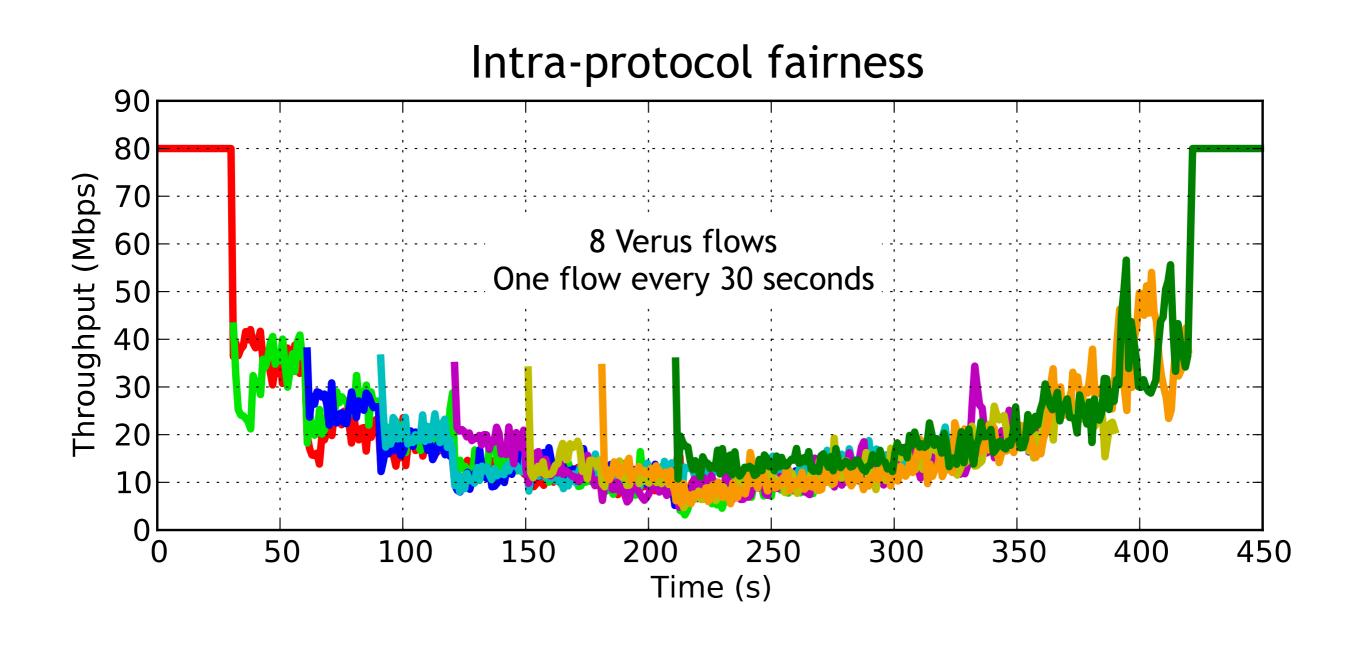
Goal 2: Balance throughput and delay

Experiments over real LTE network:

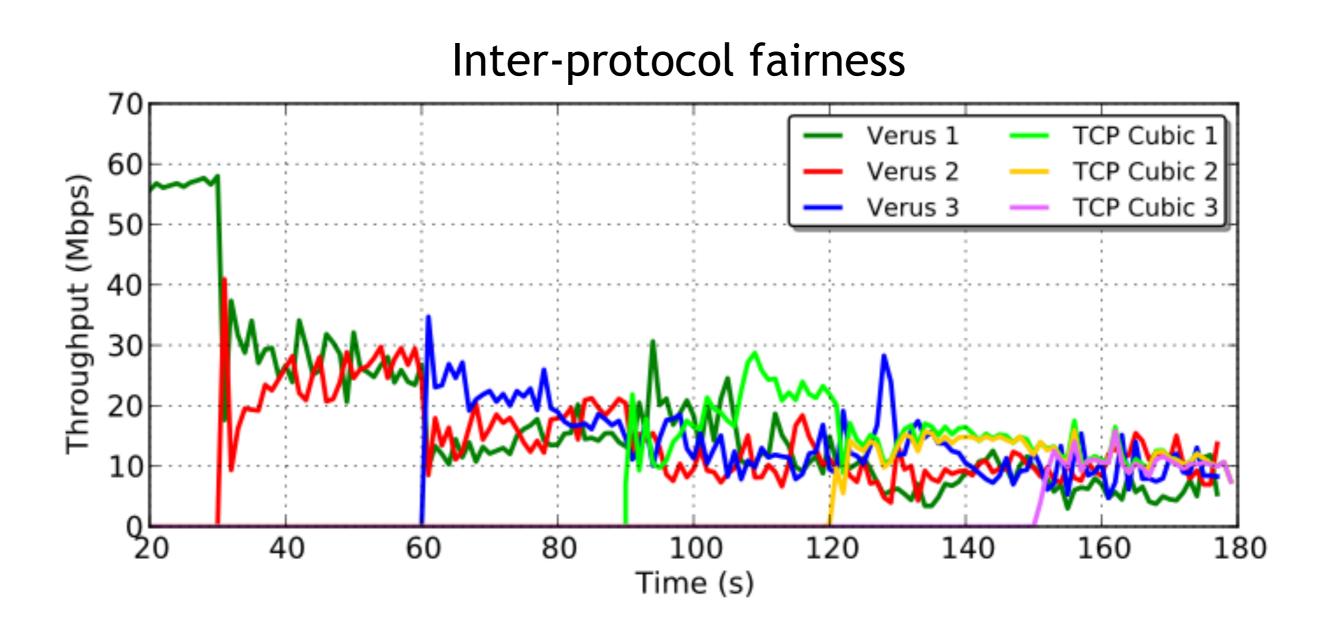
- Stationary scenario
- 3 phones each running 3 flows
- Repeated 5 times each



Goal 3: Provide fairness between competing flows



Goal 3: Provide fairness between competing flows



Verus source code is open source and available on github: http://yzaki.github.io/verus/