

Comp416 Report 2 - Wireshark

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68665

Part 1.1 UDP Experiment

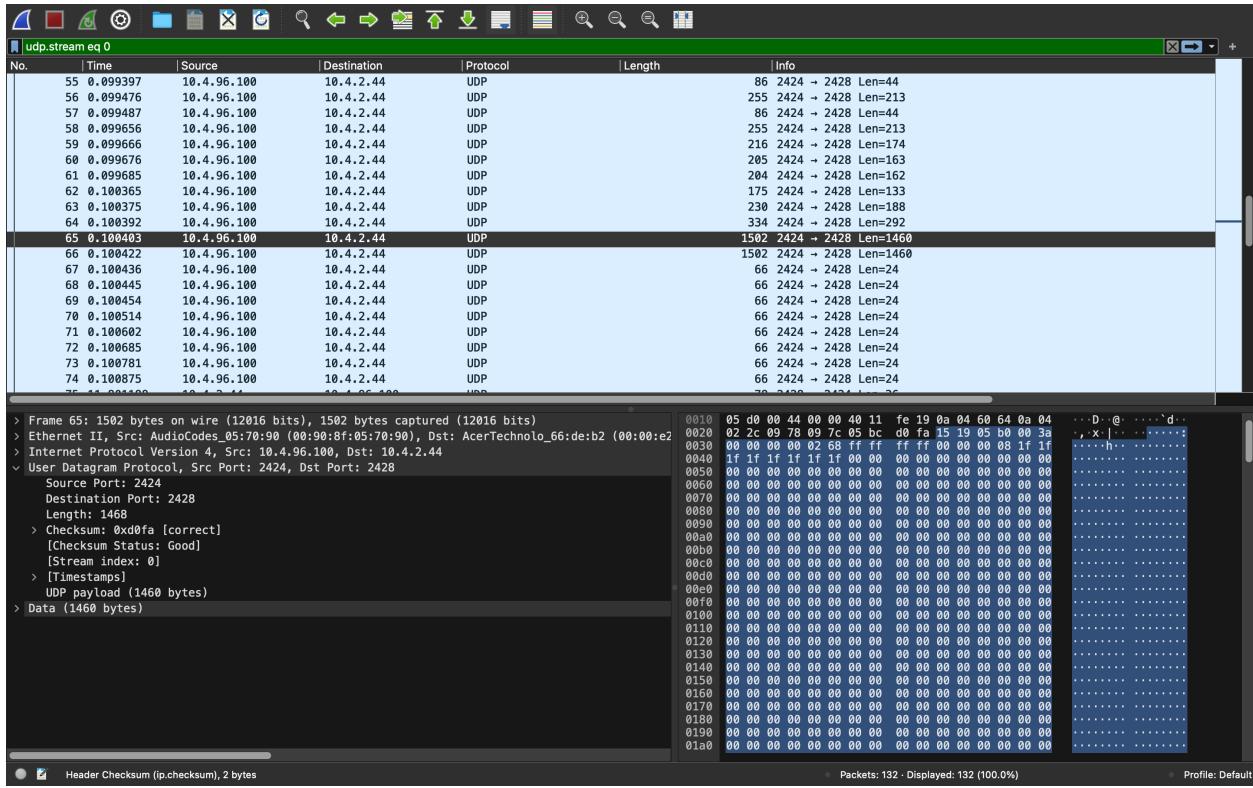


Figure 1

Here is the packet that corresponding the last two digit of my id. 68665 's last two number is 65 so i check the 65.

```
> Frame 65: 1502 bytes on wire (12016 bits), 1502 bytes captured (12016 bits)
> Ethernet II, Src: AudioCodes_05:70:90 (00:90:8f:05:70:90), Dst: AcerTechnolo_66:de:b2 (
  Internet Protocol Version 4, Src: 10.4.96.100, Dst: 10.4.2.44
    0100 .... = Version: 4
    .... 0101 = Header Length: 20 bytes (5)
  > Differentiated Services Field: 0x28 (DSCP: AF11, ECN: Not-ECT)
    Total Length: 1488
    Identification: 0x0044 (68)
    000. .... = Flags: 0x0
      0.... .... = Reserved bit: Not set
      .0.. .... = Don't fragment: Not set
      ..0. .... = More fragments: Not set
      ...0 0000 0000 0000 = Fragment Offset: 0
    Time to Live: 64
    Protocol: UDP (17)
    Header Checksum: 0xfe19 [validation disabled]
      [Header checksum status: Unverified]
    Source Address: 10.4.96.100
    Destination Address: 10.4.2.44
  < User Datagram Protocol, Src Port: 2424, Dst Port: 2428
```

Figure 2

Part 1:

(Figure 2)

Time to live : 64

TTL is a mechanism that limits the lifespan of data in a network. If the TTL value reaches zero before the packet reaches its destination, the packet will be discarded to prevent it from circulating indefinitely.

Part 2:

(Figure 1)

Stream index = 0

The Stream Index column displays a unique number for each stream. Stream indexes are Wireshark-internal. It just uses a number to uniquely identify a TCP stream.

The stream index can help you to identify and follow a specific conversation or stream of packets between two endpoints in a network.

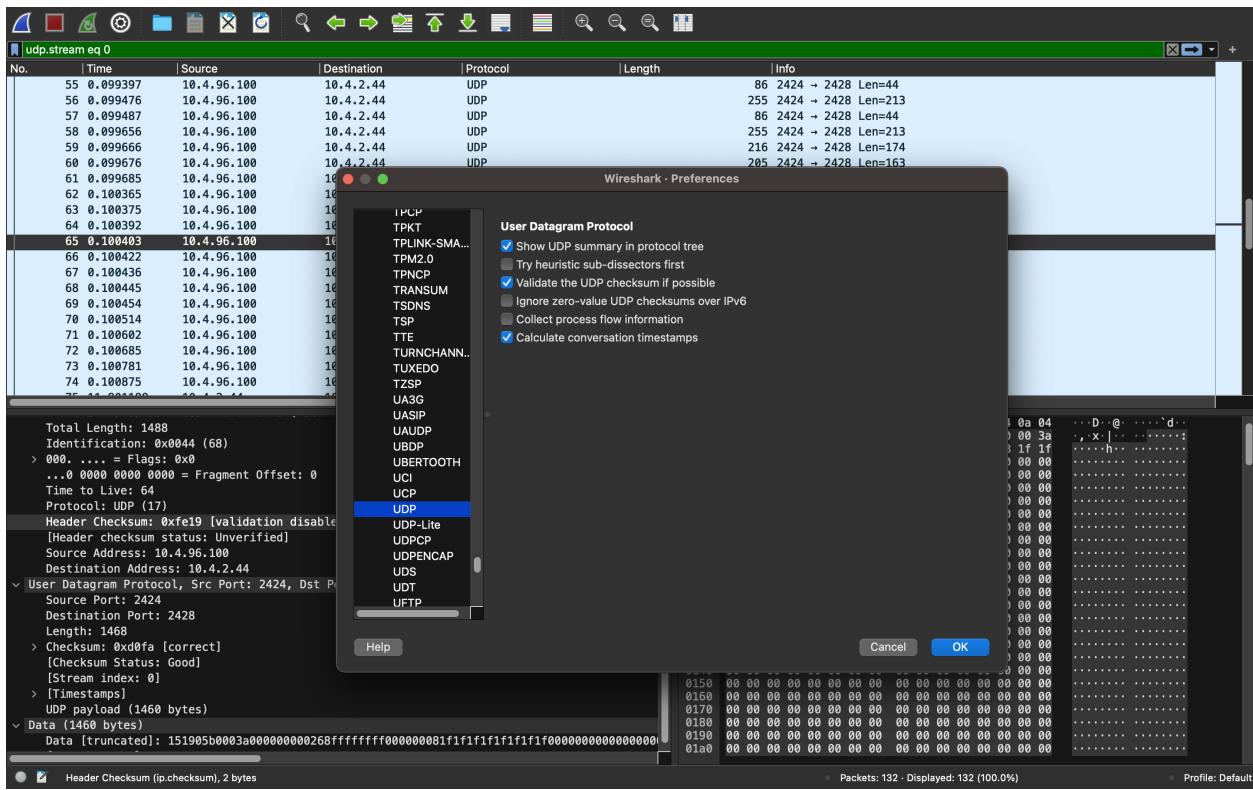
Part 3:

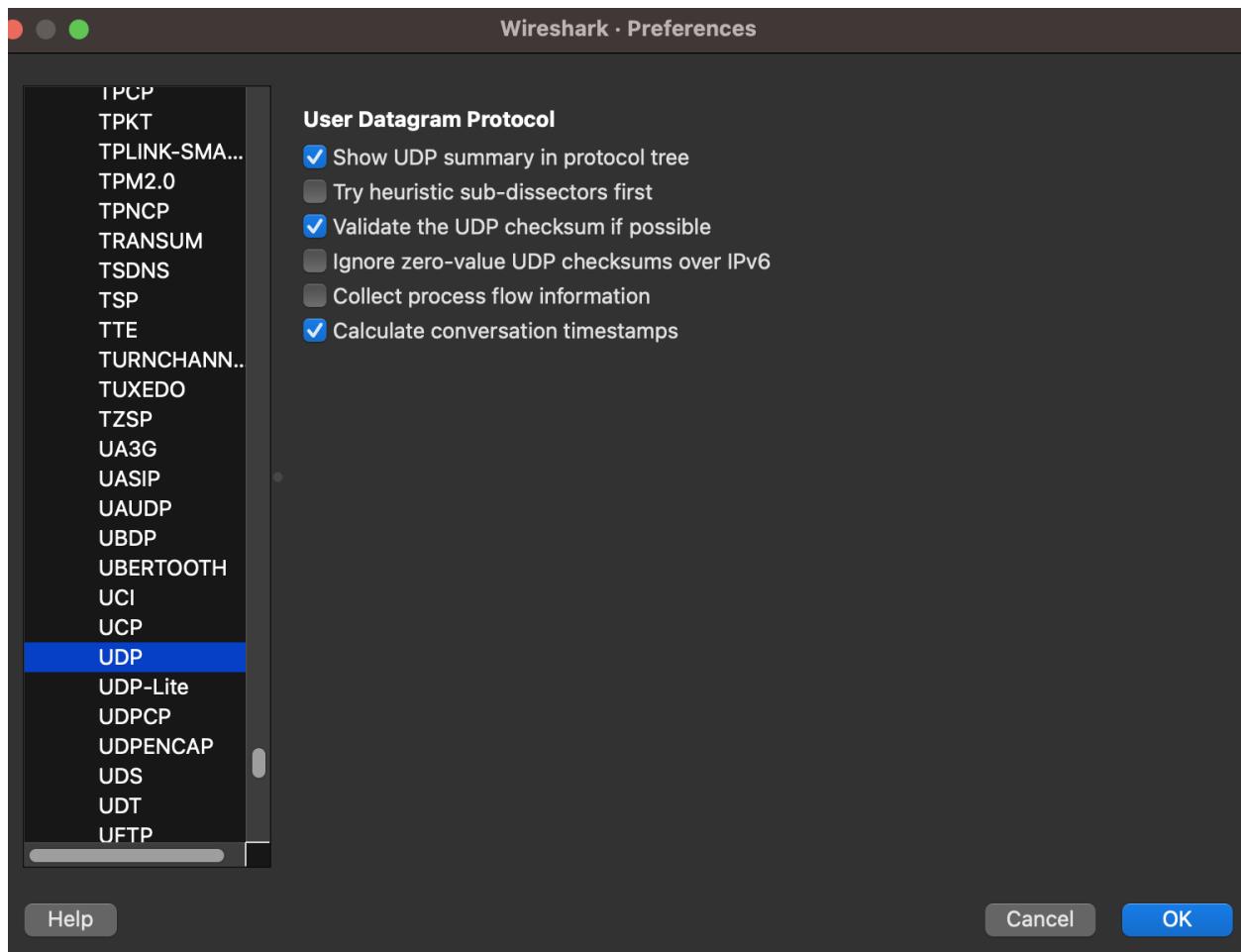
(Figure 2)

The checksum is a mathematical calculation used to ensure the integrity of the data in the packet. It's computed as the sum of all the 16-bit words in the header and data, with any overflow bits added back to the result. If the values was different than data is corrupted.

TCP also has a checksum mechanism. The difference between UDP and TCP in terms of data reliability is, UDP only says this data is corrupted, does not try to retransmit the corrupted packet, while TCP ensures the data transfer reliability therefore it retransmit the corrupted packet.

To activate the checksum validation i did this.





Part 4:

(Figure 2)

ip.flags.rb

0... = Reserved bit: Not set

This is the first bit and it is always set to 0. The purpose of this bit is reserved for future use. However, in practice, it has remained unused and is expected to be 0 in all IPv4 packets. If this bit is set to 1, it's typically considered a malformed packet.

Part 5:

(Figure 2)

The length field specifies the number of bytes in the UDP segment (header plus data).

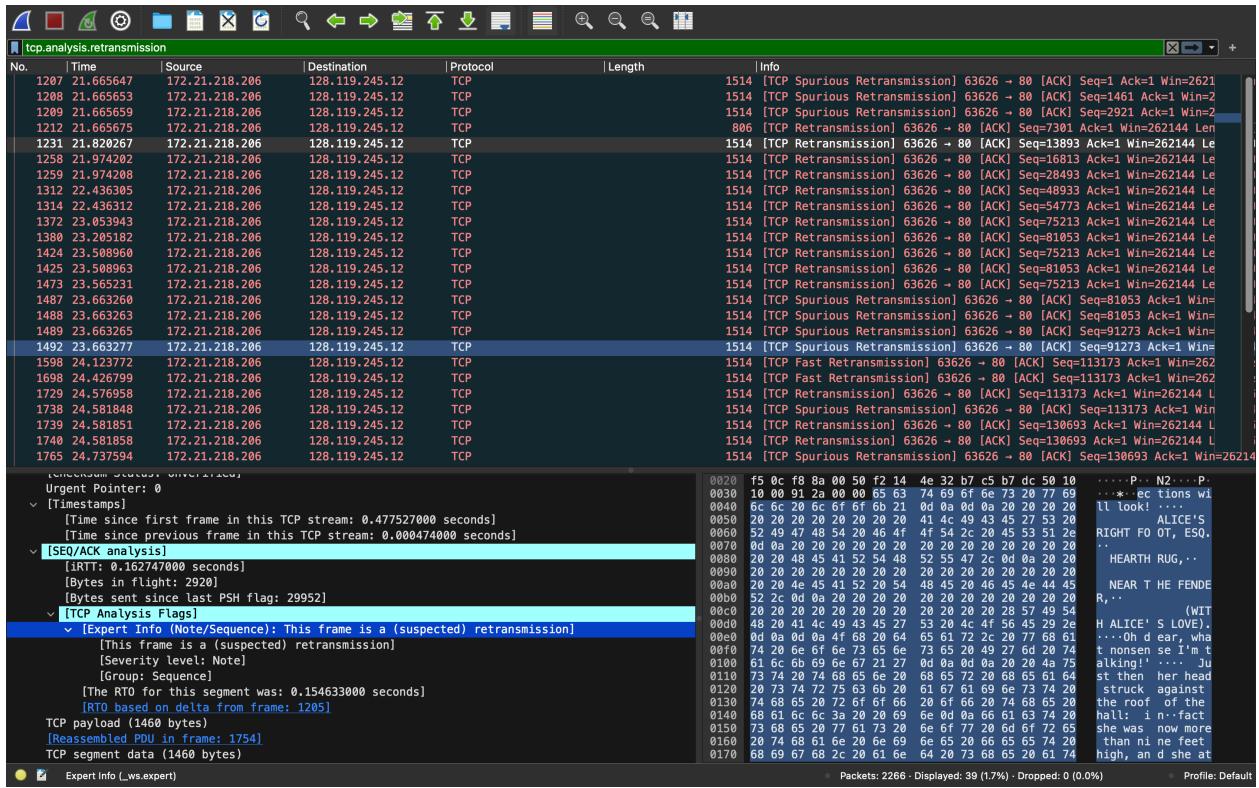
An explicit length value is needed since the size of the data field may differ from one UDP segment to the next. This length includes the UDP header (usually 8 bytes) and the UDP payload. The length of UDP payload for selected packet is 32 bytes. $1468 \text{ bytes} - 8 \text{ bytes} = 1460 \text{ bytes}$.

There is another length information which is written in the packet window. That can be seen in Figure 1. It is set to 1502. This is the total size of the packet as it was captured by Wireshark, including all headers (Ethernet, IP, TCP/UDP) and the data payload. It represents the entire packet size as it exists on the network.

Part 1.2. TCP Experiment

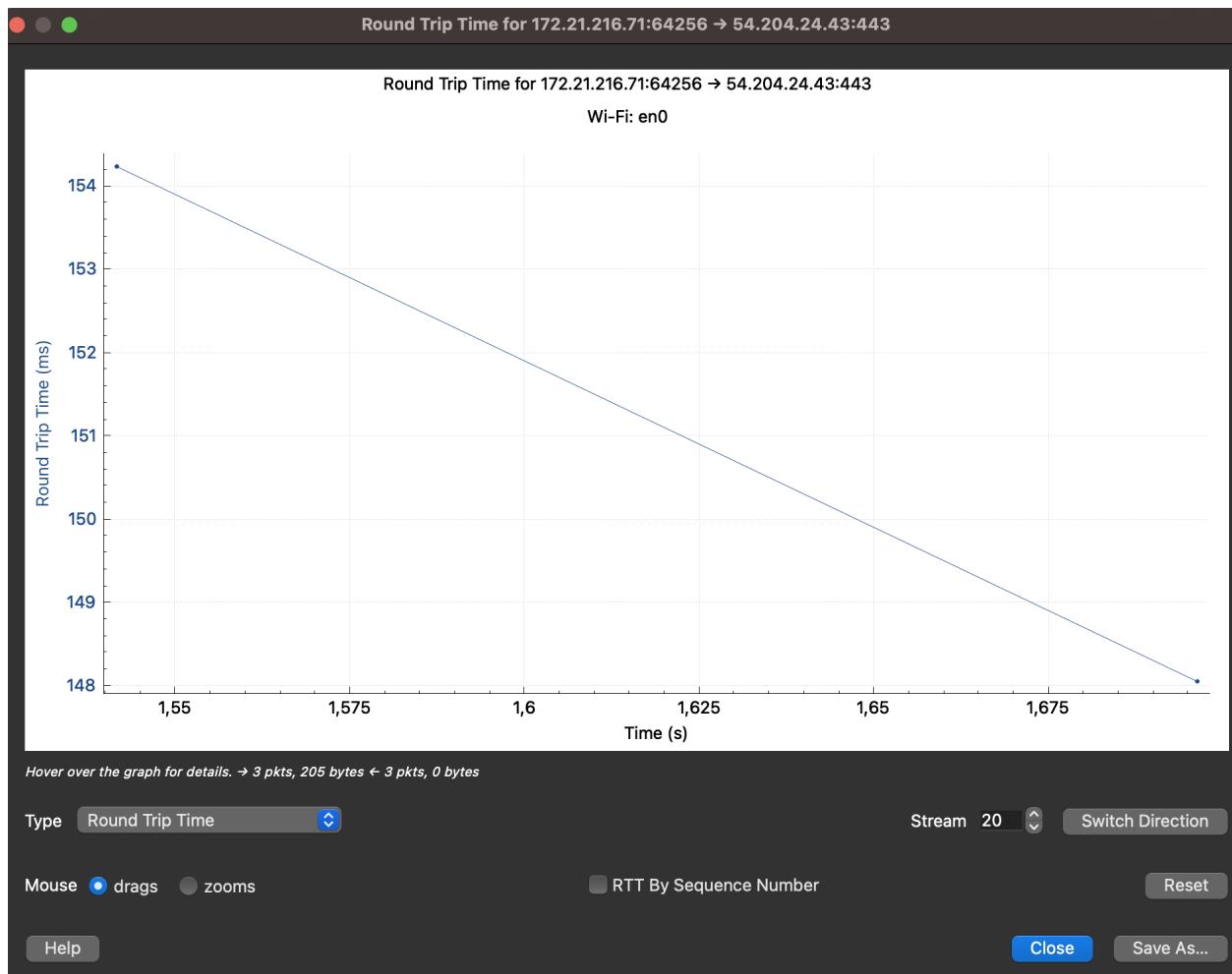
Part 1:

Using `tcp.analysis.retransmission` query on the filter section we can see the TCP Retransmission from the info. This is a post filter, because it is executed after the capture. In the packet details There is a section of [Seq/ack analysis] → [Tcp analysis flags] → [Expert Info (Note/Sequence)]. In this header it says This frame is a [suspected] retransmission. This is a clear indication of segment loss.



Part 2:

Round trip times are differ between 148- 154 ms as we can see the graph from below. There are variations present.



Part 3:

This is udp

```
Destination Address: 10.4.2.44
User Datagram Protocol, Src Port: 2424, Dst Port: 2428
  Source Port: 2424
  Destination Port: 2428
  Length: 209
    > Checksum: 0xfc7f [correct]
      [Checksum Status: Good]
      [Stream index: 0]
    > [Timestamps]
    UDP payload (201 bytes)
Data (201 bytes)
  Data [truncated]: 151900c50016000000000138ffffffff00000008000000000ffffffffff1
  [Length: 201]
```

This is tcp

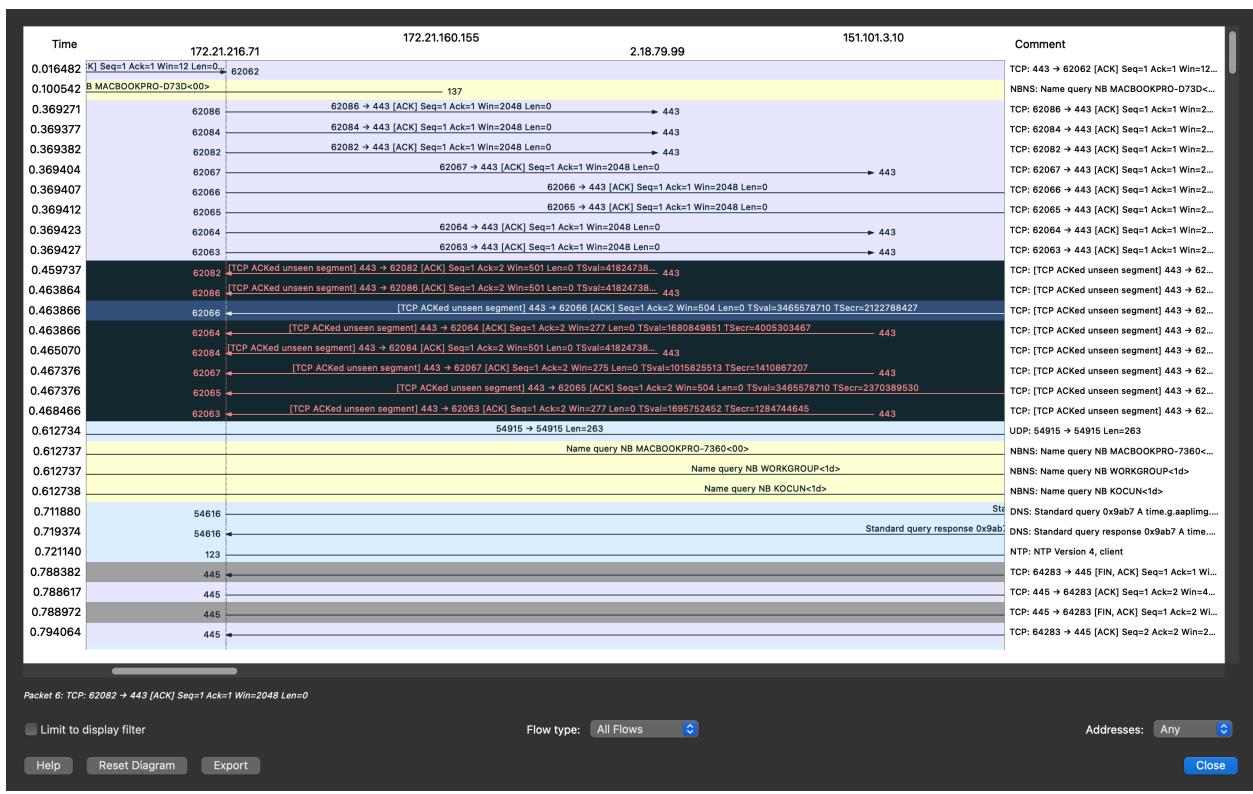
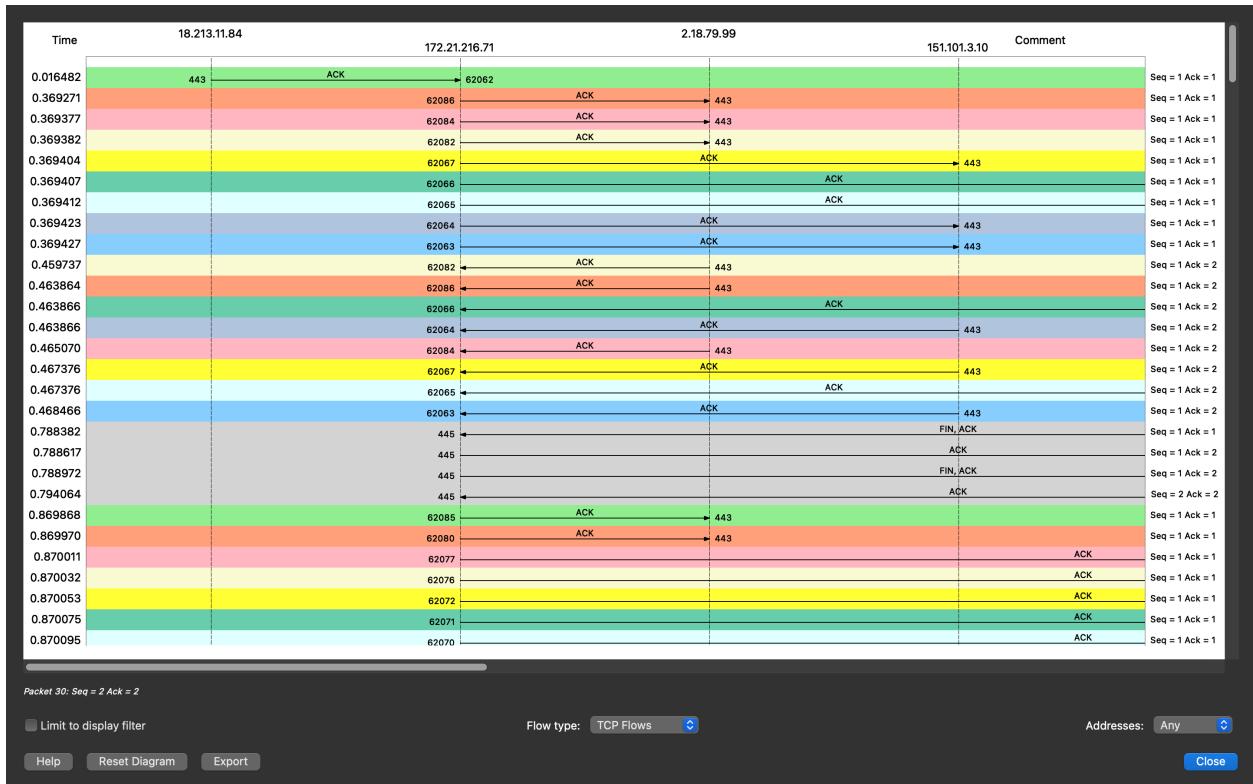
```
▼ Transmission Control Protocol, Src Port: 62067, Dst Port: 443, Seq: 1, Ack: 1, Len: 0
  Source Port: 62067
  Destination Port: 443
  [Stream index: 4]
  > [Conversation completeness: Incomplete (60)]
    [TCP Segment Len: 0]
    Sequence Number: 1      (relative sequence number)
    Sequence Number (raw): 3117365216
    [Next Sequence Number: 1      (relative sequence number)]
    Acknowledgment Number: 1      (relative ack number)
    Acknowledgment number (raw): 1336514025
    0101 .... = Header Length: 20 bytes (5)
  > Flags: 0x010 (ACK)
    Window: 2048
    [Calculated window size: 2048]
    [Window size scaling factor: -1 (unknown)]
    Checksum: 0xbd96 [unverified]
    [Checksum Status: Unverified]
    Urgent Pointer: 0
  > [Timestamps]
```

Window, acknowledgment and sequence numbers are present in the tcp which is not present in the udp. I explain what they signify and their purpose at the next part.

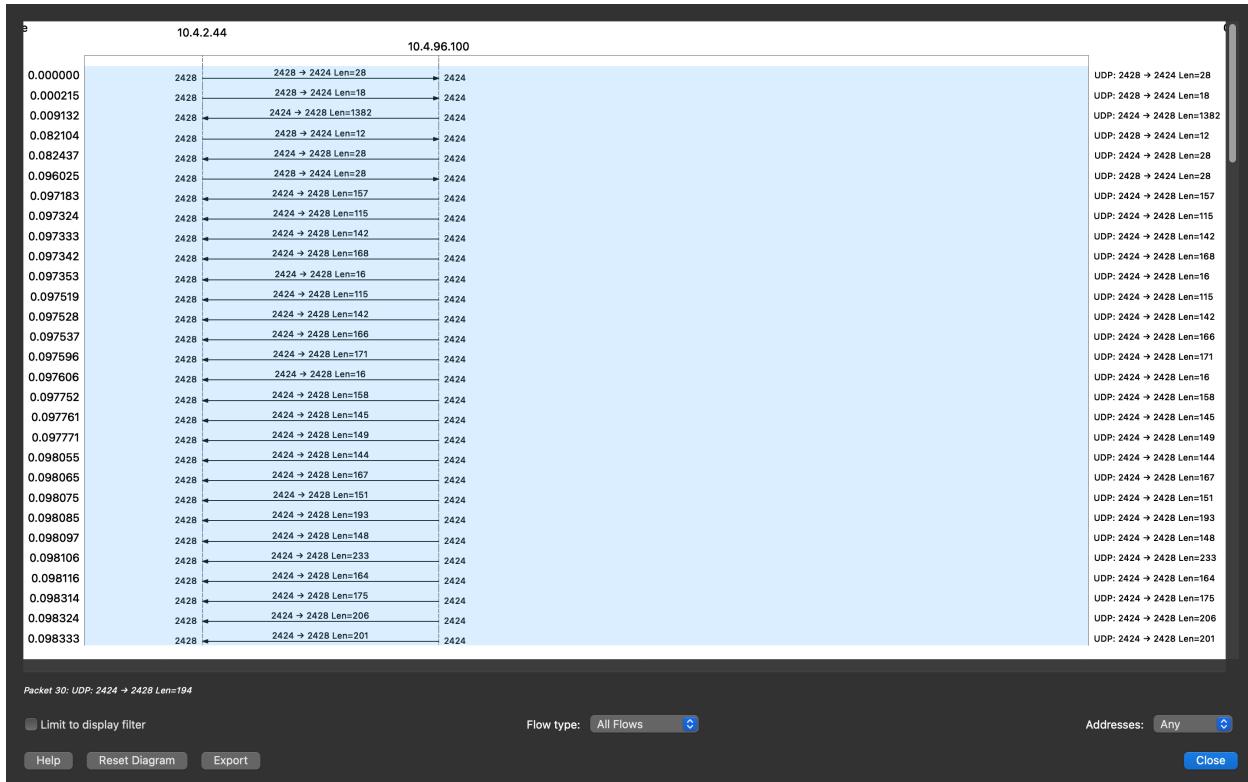
Part 4:

Those are the Flow Graphs from the Statistic Menu. First two of them is taken from the alice example, the last one is from udp_sample example.

Those are TCP flow graphs



This is UDP flow graph



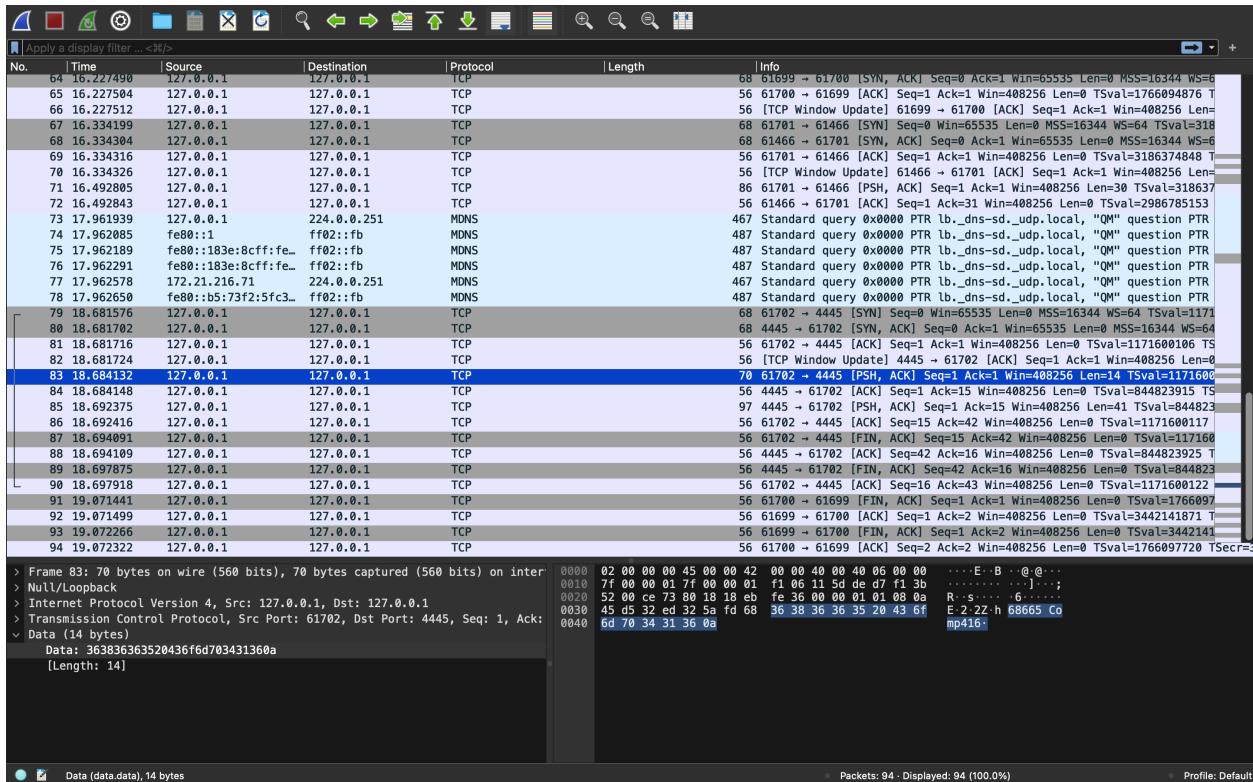
In TCP, we can see that ; seq , ack and win that is not present in the udp. Seq is sequence number which identifies the order of bytes sent over. Ack is the acknowledgment number which confirms the receipt of bytes by the receiver. With using those numbers, TCP knows which data has been received, and know which data needs to be resending. This is the mechanism that makes TCP reliable. Win is a TCP Window size, which determines the amount of data that cant be sent during the connection before receiving an acknowledgment that the data has been successfully received. Those are parts of TCP's flow control mechanism which is not present in the UDP. Udp is connectionless protocol and does not have built-in mechanism for flow control, congestion avoidance and data reliability

Part 2.1 TCP vs SSL Experiments

Part1:

This is the screenshot from captured data of the insecure connection.

As we can see it, protocol is TCP.



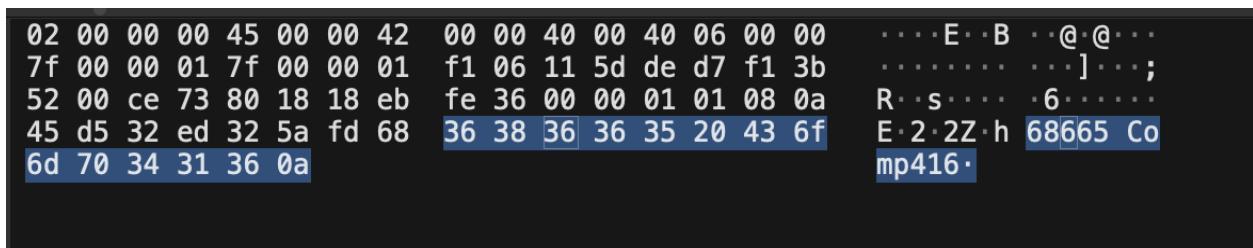
Part2:

When we look at the Packet 83 we can see the sent data within the sent packets. This is the TCP example

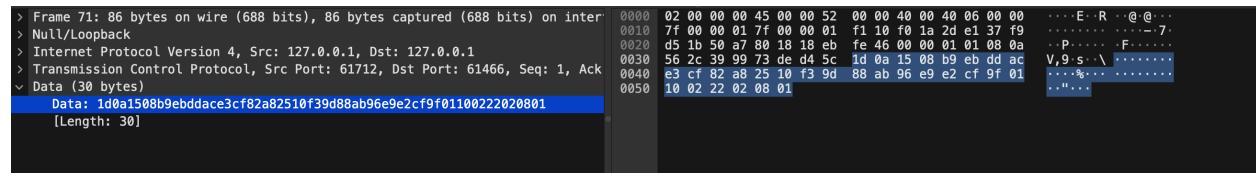
We can also see at the packet 71 that data is sent. This is the SSL example.

Part3:

In the TCP example, we can see the packet contents. Because it use insecure server which is not encrypted.



But in the SSL example the data content is encrypted as a secure connection should be.



A Wireshark screenshot showing a single network frame (Frame 71). The frame details indicate it's 86 bytes on wire (688 bits) and 86 bytes captured (688 bits) on interface Null/Loopback. It's an Internet Protocol Version 4 (IPv4) packet with source 127.0.0.1 and destination 127.0.0.1. The Transmission Control Protocol (TCP) port information shows Source Port 61712 and Destination Port 61466. The sequence number is 1, and the acknowledgement number is 10222020801. The data payload is 30 bytes long and starts with the hex value 1d 0a 15 08 b9 eb dd ace3 cf 82 a8 25 10 f3 9d. The ASCII representation of the data shows encrypted characters like 'V,9 s \%...,"..'. The packet is highlighted in blue.

Part4:

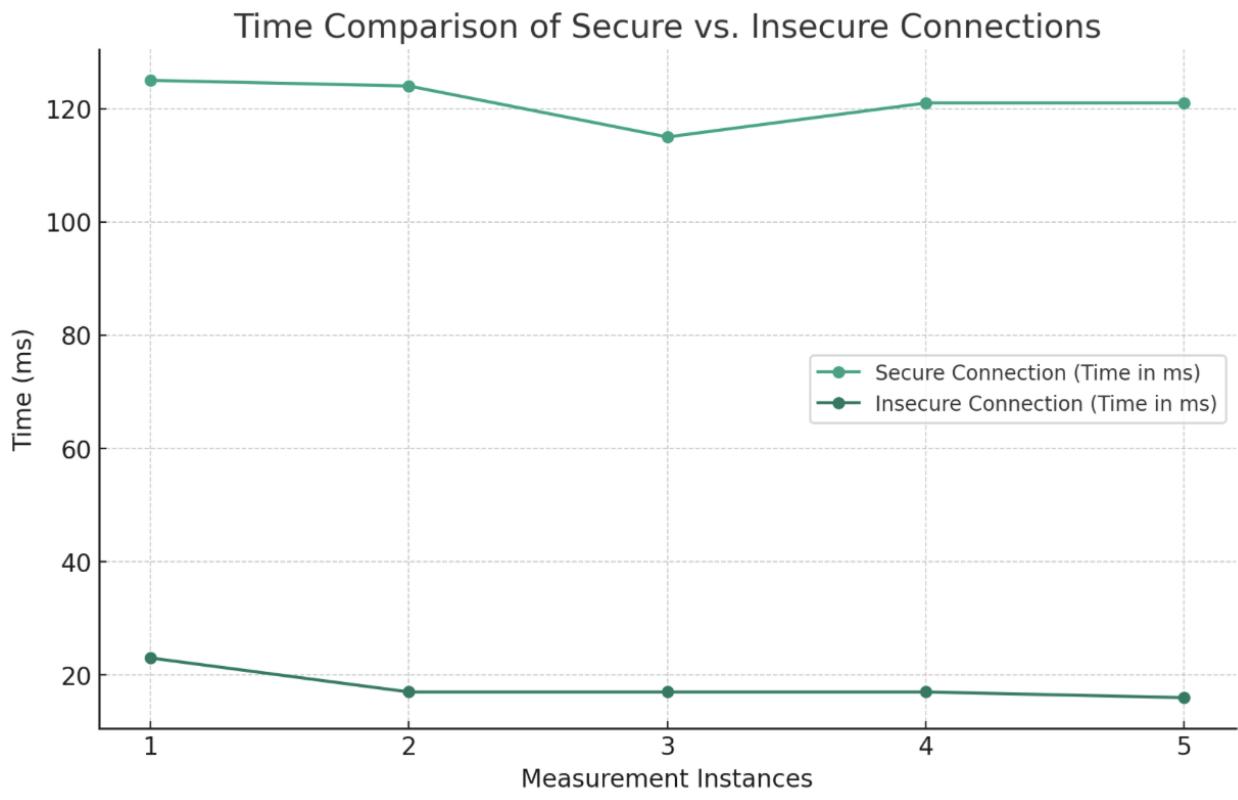
For the Secure Connection 5 different delay data in milliseconds are;

- 125
- 124
- 115
- 121
- 121

For the Insecure Connection 5 different delay data in milliseconds are;

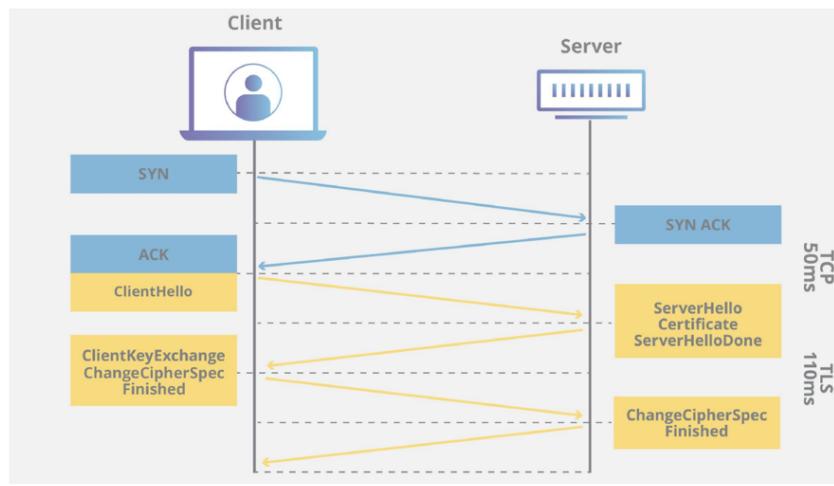
- 23
- 17
- 17
- 17
- 16

Insecure connections is faster because it doesn't do any encryption on the data.



Part5:

TLS Handshake



The packets that are starting with Client Hello and Server Hello packets are from the TLS handshake. As we can see the diagram, initial phase is synchronization packet sent from the

client to server, and acknowledgment of synchronization at the server, and returning to the client. This is the same for the all TCP connection.

When TCP connection is established, TLS handshake begins with ClientHello message to the server. We can see it from the packet's info as well.

74 9.392897	127.0.0.1	127.0.0.1	TCP	68 4446 → 61713 [SYN, ACK] Seq=0 Ack=1 Win=65535 Len=0 MSS=16344 WS=64
75 9.392920	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=1 Ack=1 Win=408256 Len=0 TSval=1905559303 TS
76 9.392935	127.0.0.1	127.0.0.1	TCP	56 [TCP Window Update] 4446 → 61713 [ACK] Seq=1 Ack=1 Win=408256 Len=0
77 9.405766	127.0.0.1	127.0.0.1	TLSv1.3	520 Client Hello
78 9.405795	127.0.0.1	127.0.0.1	TCP	56 4446 → 61713 [ACK] Seq=1 Ack=465 Win=407808 Len=0 TSval=5998312101
79 9.415276	127.0.0.1	127.0.0.1	TLSv1.3	183 Server Hello
80 9.415292	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=465 Ack=128 Win=408128 Len=0 TSval=190555932
81 9.419321	127.0.0.1	127.0.0.1	TLSv1.3	62 Change Cipher Spec
82 9.419353	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=465 Ack=134 Win=408128 Len=0 TSval=190555933
83 9.428538	127.0.0.1	127.0.0.1	TLSv1.3	126 Application Data
84 9.428556	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=465 Ack=204 Win=408064 Len=0 TSval=190555933
85 9.428952	127.0.0.1	127.0.0.1	TLSv1.3	62 Change Cipher Spec
86 9.428965	127.0.0.1	127.0.0.1	TCP	56 4446 → 61713 [ACK] Seq=204 Ack=471 Win=407808 Len=0 TSval=599831225
87 9.421774	127.0.0.1	127.0.0.1	TLSv1.3	980 Application Data
88 9.421788	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=471 Ack=1128 Win=407168 Len=0 TSval=19055593
89 9.435803	127.0.0.1	127.0.0.1	TLSv1.3	358 Application Data
90 9.435835	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=471 Ack=1430 Win=406848 Len=0 TSval=19055593
91 9.436357	127.0.0.1	127.0.0.1	TLSv1.3	146 Application Data
92 9.436386	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=471 Ack=1520 Win=406720 Len=0 TSval=19055593
93 9.439378	127.0.0.1	127.0.0.1	TLSv1.3	146 Application Data
94 9.439403	127.0.0.1	127.0.0.1	TCP	56 4446 → 61713 [ACK] Seq=1520 Ack=561 Win=407680 Len=0 TSval=59983124
95 9.442338	127.0.0.1	127.0.0.1	TLSv1.3	1244 Application Data
96 9.442354	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=561 Ack=2708 Win=405568 Len=0 TSval=19055593
97 9.444452	127.0.0.1	127.0.0.1	TLSv1.3	108 Application Data
98 9.444468	127.0.0.1	127.0.0.1	TCP	56 4446 → 61713 [ACK] Seq=2708 Ack=613 Win=407680 Len=0 TSval=59983124
99 9.444861	127.0.0.1	127.0.0.1	TLSv1.3	104 Application Data
100 9.444874	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=613 Ack=2756 Win=405504 Len=0 TSval=19055593
101 9.448351	127.0.0.1	127.0.0.1	TLSv1.3	96 Application Data
102 9.448383	127.0.0.1	127.0.0.1	TCP	56 61713 → 4446 [ACK] Seq=613 Ack=2796 Win=405504 Len=0 TSval=19055593
103 9.448396	127.0.0.1	127.0.0.1	TLSv1.3	96 Application Data

Part6:

▽ Handshake Protocol: Client Hello
Handshake Type: Client Hello (1)
Length: 455
Version: TLS 1.2 (0x0303)
Random: 3f4bdbbc28c69e005865d835d24a9e3ce525afaf1b891d0e575a7239fa1634d85
Session ID Length: 32
Session ID: 4f2096b3d3271b179a72d2225c1db9d7b0f350d8a2aec34b16145e910888991f
Cipher Suites Length: 98
> Cipher Suites (49 suites)
Compression Methods Length: 1

Version: TLS 1.2 (0x0303)
Length: 122
▽ Handshake Protocol: Server Hello
Handshake Type: Server Hello (2)
Length: 118
Version: TLS 1.2 (0x0303)
Random: b9b0177e2f55187e57f42dd12a1120148585984faef4f78fded0b897757e1eae
Session ID Length: 32
Session ID: 4f2096b3d3271b179a72d2225c1db9d7b0f350d8a2aec34b16145e910888991f
Cipher Suite: TLS_AES_256_GCM_SHA384 (0x1302)
Compression Method: null (0)
Extensions Length: 46
Extensions supported versions (len=2) TLS 1.2

It uses session id's. Server generates as session ID during the handshake and give it to the client. While server recognize the session id, whenever client wants to resume a session it can do it by the session ID.

At the Wireshark, when we click the Client Hello message we can see a session id. The same session id is also present in the Server Hello. That means that client wants to resume this session so it will avoid sending the certificate for each connection to optimization.

Part 2.2. Creating Keystore and Trust Store

Here is the clear screenshot of the information of the generated certificated.

```
(base) batuhanarazat@Batuhan-MacBook-Pro mykeystore % keytool -genkey -alias mykey -keyalg RSA -keypass mykeypass -keystore keystore.jks -storepass mystorepass
Warning: Different store and key passwords not supported for PKCS12 KeyStores. Ignoring user-specified -keypass value.
Enter the distinguished name. Provide a single dot (.) to leave a sub-component empty or press ENTER to use the default value in braces.
What is your first and last name?
[Unknown]: Batuhan Arat
What is the name of your organizational unit?
[Unknown]: Koc University
What is the name of your organization?
[Unknown]: Computer Engineering Department
What is the name of your City or Locality?
[Unknown]: Istanbul
What is the name of your State or Province?
[Unknown]: Turkey
What is the two-letter country code for this unit?
[Unknown]: TR
Is CN=Batuhan Arat, OU=Koc University, O=Computer Engineering Department, L=Istanbul, ST=Turkey, C=TR correct?
[no]: yes
Generating 3,072 bit RSA key pair and self-signed certificate (SHA384withRSA) with a validity of 90 days
    for: CN=Batuhan Arat, OU=Koc University, O=Computer Engineering Department, L=Istanbul, ST=Turkey, C=TR
(base) batuhanarazat@Batuhan-MacBook-Pro mykeystore %
```