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ФАКУЛЬТЕТ «Информатика и системы управления»
КАФЕДРА <u>«Программное обеспечение ЭВМ и информационные технологии»</u>
Лабораторная работа № <u>1</u>
Дисциплина Констриурование компиляторов
Тема Распознование цепочек регулярного языка
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# Описание задания

Напишите программу, которая в качестве входа принимает произвольное выражение, и выполняет следующие преобразования:

- 1. По регулярному выраению строит НКА.
- 2. По НКА строит эквивалентный ему ДКА.
- 3. По ДКА строит эквивалентный ему КА, имеющий найменьшее возможное количество состояний. Указание. Возпользоваться алгоритмом, приведенным по адресу <a href="http://neerc.ifmo.ru/wiki/index.php?">http://neerc.ifmo.ru/wiki/index.php?</a>
  <a href="title=Muнимизация ДКА">title=Muнимизация ДКА</a>, алгоритм за O(n^2) с построением пар различимых состояний</a>
- 4. Моделирует минимальный КА для входной цепочки из терминалов исходной грамматики.

# Текст программы

#### Листинг 1 – Класс, представляющий НКА

```
import java.util.ArrayList;
import java.util.Collections;
import java.util.List;
import java.util.stream.Collectors;
import java.util.stream.IntStream;
public class NFA {
    private final ArrayList<Integer> states = new ArrayList<>();
    private final ArrayList<Transition> transitions = new ArrayList<>();
   private int finalState;
   public int getStatesCount() {
        return this.states.size();
    public void setStates(int totalStates) {
        for (int i = 0; i < totalStates; i++) {</pre>
            this.states.add(i);
    }
    public void addTransition(int stateFrom, int stateTo, char symbol) {
        Transition trans = new Transition(stateFrom, stateTo, symbol);
        this.transitions.add(trans);
    }
    public void setFinalState(int finalState) {
        this.finalState = finalState;
    public int getFinalState() {
       return this.finalState;
    public List<Transition> getTransitions() {
        return this.transitions;
    }
    public void display() {
        for (Transition temp : transitions) {
            System.out.println("q" + temp.getFromState() + " " + temp.getSymbol() + "
--> q" + temp.getToState());
        }
```

```
System.out.println("The final state is q" + getFinalState());
    }
   public ArrayList<Character> findPossibleInputSymbols(ArrayList<Integer> states) {
        ArrayList<Character> result = new ArrayList<>();
        for (int stateFrom : states) {
            for (Transition transition: transitions) {
                if (transition.getFromState() == stateFrom &&
transition.getSymbol() != 'e') {
                    result.add(transition.getSymbol());
            }
        }
        return result;
    }
   public ArrayList<Integer> unique(ArrayList<Integer> list) {
        return IntStream
                .range(0, list.size())
                .filter(i -> ((i < list.size() - 1 && !list.get(i).equals(list.get(i))</pre>
+ 1))) || i == list.size() - 1))
.mapToObj(list::get).collect(Collectors.toCollection(ArrayList::new));
   public ArrayList<Integer> eclosure(ArrayList<Integer> states) {
        ArrayList<Integer> result = new ArrayList<>();
        boolean[] visited = new boolean[getStatesCount()];
        for (Integer integer : states) {
            eclosure(integer, result, visited);
        Collections.sort(result);
        return unique(result);
   public void eclosure(int x, ArrayList<Integer> result, boolean[] visited)
//Simple DFS
    {
        result.add(x);
        for (Transition transition: transitions) {
            if (transition.getFromState() == x && transition.getSymbol() == 'e') {
                int y = transition.getToState();
                if (!visited[v]) {
                    visited[y] = true;
                    eclosure(y, result, visited);
                }
            }
        }
   public ArrayList<Integer> move(ArrayList<Integer> T, char symbol) {
        ArrayList<Integer> result = new ArrayList<>();
        for (int t : T) {
            for (Transition transition: transitions) {
                if (transition.getFromState() == t && transition.getSymbol() ==
symbol) {
                    result.add(transition.getToState());
                }
            }
       Collections.sort(result);
        int l1 = result.size();
        unique (result);
        int 12 = result.size();
        if (12 < 11) {
            System.out.println("move(T, a) returns non-unique ArrayList");
```

```
System.exit(1);
}
return result;
}
```

### Листинг 2 – Класс, представляющий ДКА

#### import java.util.ArrayList;

```
import java.util.HashSet;
import java.util.Set;
public class DFA {
    private final ArrayList<Transition> transitions = new ArrayList<>();
    private final ArrayList<ArrayList<Integer>> entries = new ArrayList<>();
    private final ArrayList<Boolean> marked = new ArrayList<>();
    private final ArrayList<Integer> finalStates = new ArrayList<>();
     * Add newly created entry into DFA
    public int addEntry(ArrayList<Integer> entry) {
       entries.add(entry);
        marked.add(false);
        return entries.size() - 1;
    }
     * Return the array position of the next unmarked entry
    public int nextUnmarkedEntryIdx() {
        for (int i = 0; i < marked.size(); i++) {</pre>
            if (!marked.get(i)) {
                return i;
        }
        return -1;
    }
     * mark the entry specified by index as marked (marked = true)
    public void markEntry(int idx) {
       marked.set(idx, true);
    }
    public ArrayList<Integer> entryAt(int i) {
        return entries.get(i);
    }
    public int findEntry(ArrayList<Integer> entry) {
        for (int i = 0; i < entries.size(); i++) {</pre>
            ArrayList<Integer> it = entries.get(i);
            if (it.equals(entry)) {
                return i;
            }
        }
        return -1;
    }
    public void setFinalState(int nfaFinalState) {
        for (int i = 0; i < entries.size(); i++) {</pre>
            ArrayList<Integer> entry = entries.get(i);
            for (int state : entry) {
                if (state == nfaFinalState) {
                    finalStates.add(i);
                }
```

```
}
        }
    }
   public void setMinDfaFinalState(int minDfaFinalState) {
        this.finalStates.add(minDfaFinalState);
   public void setTransition(int fromState, int toState, char symbol) {
        Transition newTransition = new Transition(fromState, toState, symbol);
        transitions.add(newTransition);
   public ArrayList<Integer> getFinalStates() {
       return finalStates;
   public void display() {
        System.out.println();
        for (Transition transition: transitions) {
            System.out.println("q" + transition.getFromState() +
                    " {" + FAUtils.join(entries.get(transition.getFromState()), ",")
                    + "} " + transition.getSymbol() + " --> q" +
transition.getToState() +
                    " {" + FAUtils.join(entries.get(transition.getToState()), ",")
                    + "}");
        System.out.println("The final states are q : " + FAUtils.join(finalStates,
","));
   }
   public void displayMinDFA() {
        System.out.println();
        for (Transition transition: transitions) {
            var printStr = "q" + transition.getFromState() + " " +
transition.getSymbol() +
                    " --> q" + transition.getToState();
            if (finalStates.contains(transition.getToState())) {
                printStr = printStr.concat(" final state ");
            System.out.println(printStr);
        }
   public boolean evaluate(String x) {
        int state = 0;
        for (var i = 0; i < x.length(); i++) {
            char ch = x.charAt(i);
            for (Transition transition: transitions) {
                if (transition.getFromState() == state && transition.getSymbol() ==
ch) {
                    state = transition.getToState();
                    break;
                }
            }
        }
        return finalStates.contains(state);
   public ArrayList<Transition> getTransitions() {
       return transitions;
   public int countStates() {
        var stateCount = new HashSet<Integer>();
        for (Transition transition : this.transitions) {
            stateCount.add(transition.getFromState());
            stateCount.add(transition.getToState());
        }
```

```
return stateCount.size();
}
public Set<Integer> getAllStatesToBySymbol(int to, char symbol) {
    Set<Integer> result = new HashSet<>();
    for (Transition transition : this.transitions) {
        if (transition.getToState() == to && transition.getSymbol() == symbol) {
            result.add(transition.getFromState());
    return result;
private Set<Integer> getAllStatesFrom(int from) {
    Set<Integer> result = new HashSet<>();
    for (Transition transition : this.transitions) {
        if (transition.getFromState() == from) {
            result.add(transition.getToState());
    return result;
private Set<Integer> getStatesFromStartSet(Set<Integer> fromStart) {
    var result = new HashSet<Integer>();
    for (Integer i : fromStart) {
        result.addAll(getAllStatesFrom(i));
    return result;
}
public Set<Integer> getReachableStatesFromStart() {
    var start = 0;
    var fromStart = getAllStatesFrom(start);
    Set<Integer> result = new HashSet<>(fromStart);
    Set<Integer> temp = new HashSet<>(result);
    var containAll = false;
    do {
        temp = getStatesFromStartSet(temp);
        containAll = result.containsAll(temp);
       result.addAll(temp);
    } while (!containAll);
    return result;
}
```

#### Листинг 3 – Класс, создающий НКА по регулярному выражению

# import java.util.\*;

}

```
public class RegexRecognizer {
    //a.b
    private static NFA concat(NFA a, NFA b) {
        NFA result = new NFA();
        //No new state added in concatenation
        result.setStates(a.getStatesCount() + b.getStatesCount());
        //Copy all old transitions of a
        for (Transition transition: a.getTransitions()) {
            result.addTransition(transition.getFromState(), transition.getToState(),
        transition.getSymbol());
        }
        //Creating the link; final state of a will link to initial state of b
        result.addTransition(a.getFinalState(), a.getStatesCount(), 'e');
        //Copy all old transitions of b with offset as a's states have already been
added
```

```
var offset = a.getStatesCount();
        for (Transition transition : b.getTransitions()) {
           result.addTransition(transition.getFromState() + offset,
transition.getToState() + offset, transition.getSymbol());
        //b is the final state of this created NFA
        result.setFinalState(offset + b.getStatesCount() - 1);
        return result;
    }
    //a*
   private static NFA kleene(NFA a) {
        NFA result = addStateBefore(a);
        var oldFinalState = a.getStatesCount();
        var oldInitialState = 1;
        var newInitialState = 0;
       var newFinalState = oldFinalState + 1;
        //Epsilon transition to new final state
       result.addTransition(oldFinalState, newFinalState, 'e');
        //Reverse epsilon transition
        result.addTransition(oldFinalState, oldInitialState, 'e');
        //Forward total epsilon transition
        result.addTransition(newInitialState, newFinalState, 'e');
        //Mark final state
        result.setFinalState(newFinalState);
        return result;
    }
    //a+
   private static NFA plus(NFA a) {
       NFA result = addStateBefore(a);
       var oldFinalState = a.getStatesCount();
        var oldInitialState = 1;
        var newFinalState = oldFinalState + 1;
        //Epsilon transition to new final state
       result.addTransition(oldFinalState, newFinalState, 'e');
        //Reverse epsilon transition
       result.addTransition(oldFinalState, oldInitialState, 'e');
        //Mark final state
       result.setFinalState(newFinalState);
        return result;
    //s0->s1 as result s0->s1->s2 , where s0->s1 epsilon's transition
   private static NFA addStateBefore(NFA a) {
        NFA result = new NFA();
         * +2 because we will have one new initial state with epsilon transition to
a's initial
         * and one new final state with epsilon transition from a's final state and
from the new initial created
        result.setStates(a.getStatesCount() + 2);
        result.addTransition(0, 1, 'e');
        for (Transition transition : a.getTransitions()) {
            result.addTransition(transition.getFromState() + 1,
transition.getToState() + 1, transition.getSymbol());
       return result;
   }
    //a|b
   private static NFA orSelection(ArrayList<NFA> selections, int noOfSelections) {
       NFA result = new NFA();
       var stateCount = 2;
```

```
//Find total states by summing all NFAs
        for (var i = 0; i < noOfSelections; i++) {</pre>
            stateCount += selections.get(i).getStatesCount();
        result.setStates(stateCount);
        var adderTrack = 1;
        for (var i = 0; i < noOfSelections; i++) {</pre>
            //Initial epsilon transition to the first block of 'OR'
            result.addTransition(0, adderTrack, 'e');
            NFA selectedNFA = selections.get(i);
            for (Transition transition : selectedNFA.getTransitions()) {
                result.addTransition(transition.getFromState() + adderTrack,
transition.getToState() + adderTrack, transition.getSymbol());
            adderTrack += selectedNFA.getStatesCount();
            //Add epsilon transition to final state
            result.addTransition(adderTrack - 1, stateCount - 1, 'e');
        result.setFinalState(stateCount - 1);
        return result;
    }
   private static boolean isNotOperator(char currentSymbol) {
        return currentSymbol != '(' && currentSymbol != ')' && currentSymbol != '*'
                && currentSymbol != '|' && currentSymbol != '.' && currentSymbol !=
'+';
   public static NFA regexToNfa(String regex) {
        regex = FAUtils.normalizeInputRegex(regex);
        Stack<Character> operators = new Stack<>();
        Stack<NFA> operands = new Stack<>();
        char operatorSymbol;
        int operatorCount;
        char currentSymbol;
       NFA newSym;
        char[] x = regex.toCharArray();
        for (char value : x) {
            currentSymbol = value;
            if (isNotOperator(currentSymbol)) //Must be a character, so build the
simplest NFA
            {
                newSym = new NFA();
                newSym.setStates(2);
                newSym.addTransition(0, 1, currentSymbol);
                newSym.setFinalState(1);
                operands.push (newSym); //push it back
            } else {
                switch (currentSymbol) {
                    case '*':
                        NFA starSym = operands.pop();
                        operands.push(kleene(starSym));
                        break;
                    case '+':
                        NFA plusSym = operands.pop();
                        operands.push(plus(plusSym));
                        break;
                    case '.':
                    case '|':
                    case '(':
                        operators.push(currentSymbol);
                        break:
                    default:
```

```
operatorCount = 0;
                         operatorSymbol = operators.peek();
                         //Keep searching operands
                         if (operatorSymbol == '(') {
                             continue;
                         //Collect operands
                         do {
                             operators.pop();
                             operatorCount++;
                         } while (operators.peek() != '(');
                         operators.pop();
                         NFA firstOperand;
                         NFA secondOperand;
                         ArrayList<NFA> selections = new ArrayList<>();
                         if (operatorSymbol == '.') {
                             for (int ii = 0; ii < operatorCount; ii++) {</pre>
                                 secondOperand = operands.pop();
                                 firstOperand = operands.pop();
                                 operands.push(concat(firstOperand, secondOperand));
                         } else if (operatorSymbol == '|') {
                             for (int j = 0; j < operatorCount + 1; <math>j++) {
                                 selections.add(new NFA());
                             int tracker = operatorCount;
                             for (int k = 0; k < operatorCount + 1; k++) {
                                 selections.set(tracker, operands.pop());
                                 tracker--;
                             }
                             operands.push(orSelection(selections, operatorCount +
1));
                         break;
                }
            }
        }
        return operands.peek(); //Return the single entity. operands.poll() is also
fine
    }
}
```

#### Листинг 4 – Метод создания ДКА по НКА

```
stateTo = dfa.addEntry(U);
}
dfa.setTransition(stateFrom, stateTo, a);
}
stateFrom = dfa.nextUnmarkedEntryIdx();
}
// The finish states of the DFA are those which contain any
// of the finish states of the NFA.
dfa.setFinalState(nfa.getFinalState());
return dfa;
}
```

# Листинг 5 — Класс, который минимизирует ДКА по алгоритму $O(n^2)$ с построением пар различимых состояний

```
import java.util.*;
/**
 * DFAMinimizer minimizes DFA by O(n^2) algorithm with the
 * construction of pairs of distinguishable states
public class DFAMinimizer {
   /**
     * Step 1
     * build a table of lists of inverse edges of size n \times |\Sigma|
     * n - source DFA's states number
     */
    private static Map<FAEdge, Set<Integer>> buildInverseTransition(DFA dfa) {
        var n = dfa.countStates();
        Map<FAEdge, Set<Integer>> dfaInverseEdges = new HashMap<>();
        for (int i = 0; i < n; i++) {
            for (Character c : getLiterate(dfa)) {
                dfaInverseEdges.put(buildFAEdge(i, c), dfa.getAllStatesToBySymbol(i,
c));
        return dfaInverseEdges;
    }
    /**
     * Step 2
     * build an array of reachability of states from the starting - reachable of
size n
    private static Map<Integer, Boolean> buildReachableStateFromStart(DFA dfa) {
        Map<Integer, Boolean> result = new HashMap<>();
        var reachableStatesFromStart = dfa.getReachableStatesFromStart();
        var dfaStateCount = dfa.countStates();
        for (int i = 0; i < dfaStateCount; i++) {</pre>
            result.put(i, reachableStatesFromStart.contains(i));
        return result;
    }
    /**
     * get all terminal states
    private static boolean[] getTerminalStateArray(DFA dfa) {
       var n = dfa.countStates();
        var finalStates = dfa.getFinalStates();
```

```
boolean[] result = new boolean[n];
        for (var i = 0; i < n; i++) {
            result[i] = finalStates.contains(i);
        return result;
    }
    /**
     * Step 3 and 4
     */
    private static boolean[][] buildTable(DFA dfa) {
        int n = dfa.countStates();
        boolean[] isTerminal = getTerminalStateArray(dfa);
        Map<FAEdge, Set<Integer>> dfaInverseEdges = buildInverseTransition(dfa);
        Stack<StatePair> statePairs = new Stack<>();
        Set<Character> literate = getLiterate(dfa);
        boolean[][] marked = new boolean[n][n];
        //Step 3
        for (var i = 0; i < n; i++) {
            for (var j = 0; j < n; j++) {
                if (!marked[i][j] && isTerminal[i] != isTerminal[j]) {
                    marked[i][j] = marked[j][i] = true;
                    statePairs.push(new StatePair(i, j));
                }
            }
        //Step 4
        while (!statePairs.isEmpty()) {
            var headStatePair = statePairs.pop();
            for (Character c : literate) {
                var rList = dfaInverseEdges.get(buildFAEdge(headStatePair.getI(),
c));
                for (Integer r : rList) {
                    var sList = dfaInverseEdges.get(buildFAEdge(headStatePair.getJ(),
c));
                    for (Integer s : sList) {
                        if (!marked[r][s]) {
                            marked[r][s] = marked[s][r] = true;
                            statePairs.push(new StatePair(r, s));
                    }
                }
            }
        }
        return marked;
    }
    /**
     * Step 6
     * Build the minimized DFA
    private static DFA buildDFA(int[] component, DFA sourceDFA) {
        DFA result = new DFA();
        var oldFinalsState = sourceDFA.getFinalStates();
        var n = sourceDFA.countStates();
        var equivalentStates = getEquivalentState(component, sourceDFA);
        for (var state = 0; state < n; state++) {</pre>
            var currentNewState = component[state];
            for (Transition transition : sourceDFA.getTransitions()) {
                if (transition.getFromState() == state) {
                    var toNewState = component[transition.getToState()];
                    result.setTransition(currentNewState, toNewState,
transition.getSymbol());
```

```
if (oldFinalsState.contains(state)) {
                result.setMinDfaFinalState(currentNewState);
        return result;
    }
    private static Map<Integer, List<Integer>> getEquivalentState(int[] component,
DFA dfa) {
        var n = dfa.countStates();
        Map<Integer, List<Integer>> result = new HashMap<>();
        for (var i = 0; i < n; i++) {
            var index = new ArrayList<Integer>();
            for (var j = i; j < n; j++) {
                if (component[j] == i) {
                    index.add(j);
            if (index.size() >= 2) {
                result.put(i, index);
        }
        return result;
    private static boolean areEquivalentState(int firstState, int secondState,
Map<Integer, List<Integer>> mapEquivalentStates) {
        return (mapEquivalentStates.containsKey(firstState) &&
mapEquivalentStates.get(firstState).contains(secondState))
                | (mapEquivalentStates.containsKey(secondState) &&
mapEquivalentStates.get(secondState).contains(firstState));
    private static FAEdge buildFAEdge(int state, char symbol) {
       return new FAEdge(state, symbol);
    private static Set<Character> getLiterate(DFA dfa) {
        ArrayList<Transition> dfaTransitions = dfa.getTransitions();
        Set<Character> literate = new HashSet<>();
        for (Transition transition : dfaTransitions) {
            if (FAUtils.isInputCharacter(transition.getSymbol())) {
                literate.add(transition.getSymbol());
        return literate;
    }
    /**
     * Main's method
     * Minimizes the source DFA
    public static DFA minimization(DFA dfa) {
        boolean[][] marked = buildTable(dfa);
        var reachable = buildReachableStateFromStart(dfa);
        var n = dfa.countStates();
        //Step 5
        int[] component = new int[n];
        Arrays.fill(component, -1);
        for (var i = 0; i < n; i++) {
            if (!marked[0][i]) {
                component[i] = 0;
            }
        }
```

```
int componentsCount = 0;
        for (var i = 0; i < n; i++) {
            if (!reachable.get(i)) {
                continue;
            if (component[i] == -1) {
                componentsCount++;
                component[i] = componentsCount;
                for (var j = i + 1; j < n; j++) {
                    if (!marked[i][j]) {
                        component[j] = componentsCount;
                }
            }
        }
        //Step 6
       return buildDFA(component, dfa);
   }
}
```

### Листинг 6 – Основной метод выполнения программы

```
public static void recognise() {
    System.out.println("\nThe Thompson's Construction Algorithm takes a regular
expression as " +
            "an input and returns its corresponding Non-Deterministic Finite
Automaton \n");
    System.out.println("The current recognizer supports characters 'a' and 'b',
operations '.', '|', '+' and '*'");
    System.out.println("Metadata '(' and ')' \n\n");
    System.out.println("Enter the regular expression. Ex: (a|b)*aab");
    String regex = new Scanner(System.in).next();
    System.out.println("\nThe required NFA has the transitions: ");
    NFA requiredNfa;
    requiredNfa = regexToNfa(regex);
    requiredNfa.display();
    System.out.println("\nDFA :");
   DFA requiredDfa = nfaToDfa(requiredNfa);
    requiredDfa.display();
    DFA requiredMinDfa = DFAMinimizer.minimization(requiredDfa);
    System.out.println("\nmin DFA :");
    requiredMinDfa.displayMinDFA();
    String eval;
    do {
        System.out.println("Enter string to evaluate " + regex + " or tap 0 to
stop");
        eval = new Scanner(System.in).next();
        if (!eval.equals("0")) {
            if (requiredMinDfa.evaluate(eval)) {
                System.out.println(eval + " is accepted by regex " + regex);
            } else {
                System.out.println(eval + " is rejected by regex " + regex);
    } while (!eval.equals("0"));
}
```

# Набор тестов и ожидаемые результаты для проверки правильности программы

#### Pегулярное выражение (a|b)\*abb

1. Первый тест

входная цепочка abb

выход конечного автомата да

2. Второй тест

входная цепочка abbab

выход конечного автомата нет

3. Третий Тест

входная цепочка bbbbabb

выход конечного автомата да

# Pегулярное выражение b(a|b)+

1. Первый тест

входная цепочка ав

выход конечного автомата нет

2. Второй тест

входная цепочка abab

выход конечного автомата да

3. Третий Тест

входная цепочка bbabb

выход конечного автомата да

# Регулярное выражение b\*

1. Первый тест

входная цепочка пустая цепочка (строка)

выход конечного автомата да

2. Второй тест

входная цепочка abab

выход конечного автомата да

3. Третий Тест

входная цепочка jacques

выход конечного автомата да

#### Результаты выполнения программы

```
Enter the regular expression. Ex: (a|b)*aab
     (alb)*abb
₽
\equiv \downarrow
     The required NFA has the transitions:
q0 e --> q1
而
     q1 e --> q2
     q2 a --> q3
      q3 e --> q6
     q1 e --> q4
      q4 b --> q5
     q5 e --> q6
     q6 e --> q7
      q6 e --> q1
     q0 e --> q7
     q7 e --> q8
     q8 a --> q9
     q9 e --> q10
     q10 b --> q11
     q11 e --> q12
     q12 b --> q13
     The final state is q13
₽
     DFA:
\equiv \downarrow
     q0 {0,1,2,4,7,8} a --> q1 {1,2,3,4,6,7,8,9,10}
q0 {0,1,2,4,7,8} b --> q2 {1,2,4,5,6,7,8}
亩
     q0 {0,1,2,4,7,8} a --> q1 {1,2,3,4,6,7,8,9,10}
     q1 {1,2,3,4,6,7,8,9,10} a --> q1 {1,2,3,4,6,7,8,9,10}
     q1 {1,2,3,4,6,7,8,9,10} b --> q3 {1,2,4,5,6,7,8,11,12}
     q1 {1,2,3,4,6,7,8,9,10} a --> q1 {1,2,3,4,6,7,8,9,10}
     q1 {1,2,3,4,6,7,8,9,10} b --> q3 {1,2,4,5,6,7,8,11,12}
     q2 \{1,2,4,5,6,7,8\} a --> q1 \{1,2,3,4,6,7,8,9,10\}
     q2 {1,2,4,5,6,7,8} b --> q2 {1,2,4,5,6,7,8}
     q2 {1,2,4,5,6,7,8} a --> q1 {1,2,3,4,6,7,8,9,10}
     q3 {1,2,4,5,6,7,8,11,12} a --> q1 {1,2,3,4,6,7,8,9,10}
     q3 {1,2,4,5,6,7,8,11,12} b --> q4 {1,2,4,5,6,7,8,13}
     q3 {1,2,4,5,6,7,8,11,12} a --> q1 {1,2,3,4,6,7,8,9,10}
     q3 {1,2,4,5,6,7,8,11,12} b --> q4 {1,2,4,5,6,7,8,13}
     q4 {1,2,4,5,6,7,8,13} a --> q1 {1,2,3,4,6,7,8,9,10}
     q4 {1,2,4,5,6,7,8,13} b --> q2 {1,2,4,5,6,7,8}
     q4 {1,2,4,5,6,7,8,13} a --> q1 {1,2,3,4,6,7,8,9,10}
     The final states are q : 4
```

```
min DFA:
\downarrow
     q0 a --> q1
=
     q0 b --> q0
     q0 a --> q1
     q1 a --> q1
q1 b --> q2
俞
     q1 a --> q1
     q1 b --> q2
     q2 a --> q1
     q2 b --> q3 final state
     q2 a --> q1
     q2 b --> q3 final state
     q3 a --> q1
     q3 b --> q0
     q3 a --> q1
     Enter string to evaluate (a|b)*abb or tap 0 to stop
     abb
     abb is accepted by regex (a|b)*abb
     Enter string to evaluate (a|b)*abb or tap 0 to stop
     abbab
     abbab is rejected by regex (a|b)*abb
     Enter string to evaluate (a|b)*abb or tap 0 to stop
     bbbbabb
     bbbbabb is accepted by regex (a|b)*abb
     Enter string to evaluate (a|b)*abb or tap 0 to stop
```

Puc. 1: Результат выполения программы - регулярнное выражение (a|b)\*abb

```
Enter the regular expression. Ex: (a|b)*aab
     b(a|b)+
=
     The required NFA has the transitions:
     q0 b --> q1
     q1 e --> q2
q2 e --> q3
     q3 e --> q4
     q4 a --> q5
     q5 e --> q8
     q3 e --> q6
     q6 b --> q7
     q7 e --> q8
     q8 e --> q9
     q8 e --> q3
     The final state is q9
     min DFA:
     q0 b --> q1
     q1 a --> q2 final state
     q1 b --> q2 final state
     q2 a --> q2 final state
     q2 b --> q2 final state
     Enter string to evaluate b(a|b)+ or tap 0 to stop
     ab
     ab is rejected by regex b(a|b)+
     Enter string to evaluate b(a|b)+ or tap 0 to stop
     abab
     abab is accepted by regex b(a|b)+
     Enter string to evaluate b(a|b)+ or tap 0 to stop
     bbabb
     bbabb is accepted by regex b(a|b)+
     Enter string to evaluate b(a|b)+ or tap 0 to stop
```

Рис. 2: Результат выполения программы - регулярнное выражение b(a|b)+

```
Enter the regular expression. Ex: (a|b)*aab
     b*
\downarrow
    The required NFA has the transitions:
=
     q0 e --> q1
\equiv \downarrow
     q1 b --> q2
q2 e --> q3
    q2 e --> q1
而
     q0 e --> q3
     The final state is q3
     DFA:
     q0 {0,1,3} b --> q1 {1,2,3}
     q1 {1,2,3} b --> q1 {1,2,3}
     The final states are q: 0,1
     min DFA:
     q0 b --> q0 final state
     Enter string to evaluate b* or tap 0 to stop
     "" is accepted by regex b*
     Enter string to evaluate b* or tap 0 to stop
     abab
     abab is accepted by regex b*
     Enter string to evaluate b* or tap 0 to stop
     jacques
     jacques is accepted by regex b*
     Enter string to evaluate b* or tap 0 to stop
```

Рис. 3: Результат выполения программы - регулярнное выражение b\*