

Appendices Used in the Formal Analysis of Software Library

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Appendix A

Learned Models

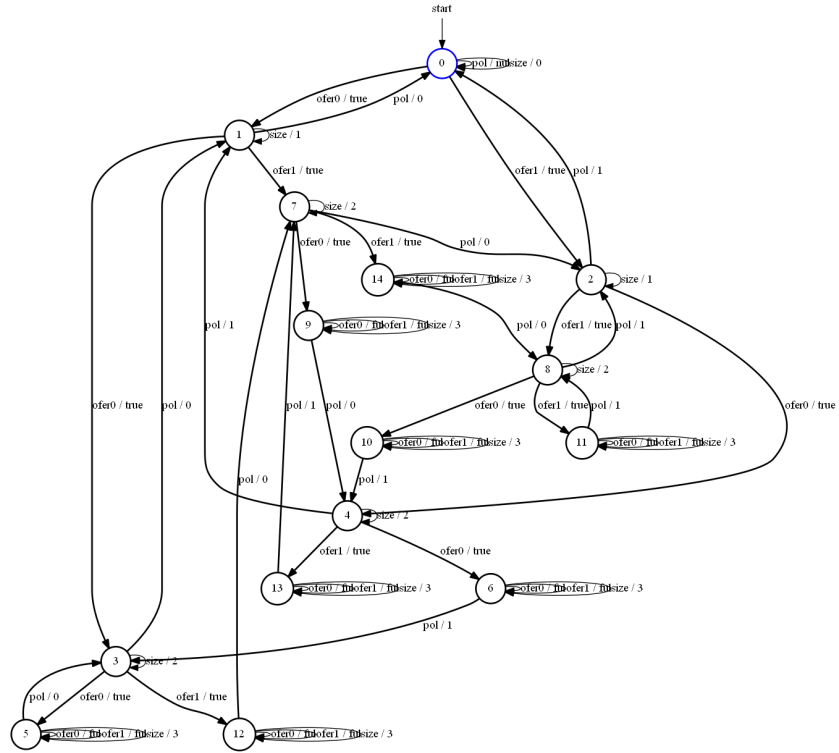


Fig. A.1: Learned Model of Queue Implementation (size=3)

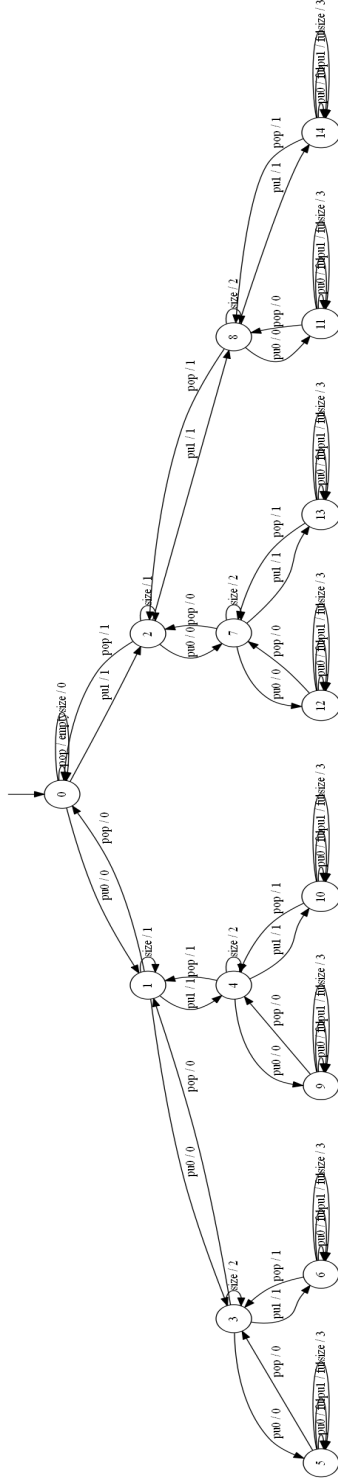


Fig. A.2: Learned Model of Stack Implementation (size=3)

Appendix B

Formal Specifications

Table B.1: Requirement Specifications for SET

Pro- perty	Informal and Formal (LTL/CTL) Specifications
	DEFINE Set-Fixed-Size:= 2 to 10 DEFINE FullSet:= $((inp = add1 inp = add2)\&out = ful)$ DEFINE EmptySet:= $(inp = size\&out = 0) (inp = remov1\&out = empty)$
	Safety Requirement Specifications
P1	It is never the case that the “Set” indicates both “Empty” and “Full” simultaneously. $AG!(((inp = size\&out = 0) ((inp = remov1 inp = remov2)\&out = empty))\&((inp = add1\&out = ful) (inp = add2\&out = ful) (inp = size\&out = Set - Fixed - Size)))$
P2	It is never the case that Set-size becomes “greater than fixed-size” i.e., $(Set - Size > Set - Fixed - Size)$. $AG!(inp = size\&out = (Set - Fixed - Size) + 1)$.
P3	It is never the case that the “Set-size becomes less than zero”. $AG!(inp = size\&out = -1)$ Or.
P4	It never happens that the “Set always remains Empty” (except at initial state). $AG!((state! = s0)\&(inp = size\&out = 0))$
P5	It never happens that the Set always remains Full (except at final state). $G!((inp = size\&out! = Set - Fixed - Size)\&FullSet)$ Or. $AG!((inp = size\&out! = Set - Fixed - Size)\&FullSet)$
P6	It should never be the case that the two operations “Add” and “Remove” occur at the same time. $G!((inp = add1\&out = Tru)\&(inp = remov1\&out = Tru))$ Or. $AG!(inp = add1\&inp = remov1)$
P7	It should never be the case that the two operations “Add0” and “Add1” occur at the same time. $G!((inp = add1\&out = Tru)\&(inp = add2\&out = Tru))$ Or. $AG!(inp = add1\&inp = add2)$
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Table B.1 – continued from previous page

Pro- perty	Informal and Formal (LTL/CTL) Specifications
P8	It never happens that if “Add/Remove” operation occurs then Set move in “Overflows” state. $AG!((inp = remov1 inp = remov2) \& out = (Set - Fixed - Size) + 1)$ Or. $G!(inp = add1 \& out = (Set - Fixed - Size) + 1)$
P9	It never happens that if “Add/Remove” operation occurs then Set move in “Underflow” state. $G!(inp = add1 \& out = -1)$ Or. $AG!((inp = remov1 inp = remov2) \& out = -1)$
Functional Requirement Specification	
P10	Set must not contain duplicated elements. $G(inp = add1 - > X(inp = add1 - > out! = Tru))$ Or $G(inp = add2 - > X(inp = add2 - > out! = Tru))$
P11	It should be ‘Empty’ only at initial state (not in every state). $G(state! = s0 - > !EmptySet)$ Or $G((inp = size \& out! = 0) - > !EmptySet)$
P12	It should be ‘Filled’ only at final state. (not in every state). $G((inp = size \& out! = Set - Fixed - Size) - > !FullSet)$
P13	“Add” operation must increase the Set-size. $G((inp = add1 inp = add2) - > X(inp = size - > out! = 0))$ $G((inp = size \& out = 0) - > X(inp = add1 - > X(inp = add2 - > X(inp = size - > out = 2))))$
P14	“Remove” operation must decrease the Set-size. $G((inp = size \& out = Set - Fixed - Size) - > X(inp = remov1 - > X(inp = remov2 - > X(inp = size - > out = (Set - Fixed - Size) - 2))))$ Or $G((FullSet) - > X((inp = remov1 inp = remov2) - > X((inp = remov1 inp = remov2) - > X(inp = size - > out! = Set - Fixed - Size))))$
P15	Set should not “Overflow” or “Underflow”. $G(!FullSet \& (inp = add1 inp = add2)) - > X(inp = size - > out! = Set - Overflow)$ $G(!EmptySet \& inp = remov1 - > X(inp = size - > out! = Set - Underflow))$
P16	If equal number of “Add” operations followed by same number of “Remove” operations then Set-size should remain same. $G((inp = size \& out = 1) - > X(inp = add1 - > X(!FullSet \& inp = add2 - > X(inp = remov1 - > X(inp = remov1 - > X(inp = size - > out = 1))))))$ $G((FullSet) - > X(inp = add1 - > X(!FullSet \& inp = add2 - > X(inp = remov1 - > X(inp = remov2 - > X(FullSet))))))$
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Table B.1 – continued from previous page

Pro- perty	Informal and Formal (LTL/CTL) Specifications
P17	<p>“Add1” operation must rightly inserts “1” (not any other number like 2) on the Set. And “Add2” must inserts “2” (not 1) on the Set.</p> $G(!FullSet \& inp = add1 - > X(inp = remov1 - > out = 1))$ $G(!FullSet \& inp = add2 - > X(inp = remov2 - > out = 2))$
Liveness Requirement Specifications	
P18	<p>If “Add” operation is done forever on a non-filled Set and “Remove” operation never done then the Set eventually becomes “Full”.</p> $G(!FullSet \& ((inp = add1 inp = add2) U FullSet) - > X(inp = size - > out = Set - Fixed - Size))$ $G(!EmptySet \& (inp = add1 inp = add2) U FullSet \& inp! = remov1) - > X(inp = size - > out = Set - Fixed - Size))$
P19	<p>If “Remove” operation is done forever on non-empty Set and “Push” operation never done then the Set eventually becomes “Empty”.</p> $G(!FullSet - > (((inp = remov1 inp = remov2) U EmptySet) - > F(inp = size - > out = 0)))$ $G(!EmptySet - > ((inp = remov1 inp = remov2) U EmptySet) - > X(inp = size - > out = 0))$
P20	<p>If “Add” operation wants to write on the Set then it eventually did this successfully without “Overflow”.</p> $AG((!FullSet \& (inp = add1 inp = add2)) - > AF(out! = Set - Underflow out! = Set - Overflow))$ $AG((inp = add1 inp = add2) - > AF(out! = Set - Overflow))$
P21	<p>If “Remove” operation wants to remove an element from the Set, then it eventually did this successfully without “Underflow”.</p> $AG(!EmptySet \& inp = remov1 - > AF(out! = Set - Underflow))$

Table B.2: Requirement Specifications for Queue

Pro- perty	Informal and Formal (LTL/CTL) Specifications
<p>DEFINE Queue-Fixed-Size:= 3 to 16</p> <p>DEFINE FullQueue:= $((inp = ofer0 inp = ofer1) \& out = ful)$</p> <p>DEFINE EmptyQueue:= $(inp = size \& out = 0) (inp = pol \& out = null)$</p>	
Safety Requirement Specifications	
P22	<p>It is never the case that the “Queue” indicates both “Empty” and “Full” simultaneously.</p> $AG!(((inp = size \& out = 0) (inp = pol \& out = null)) \& ((inp = ofer0 \& out = ful) (inp = ofer1 \& out = ful) (inp = size \& out = Queue - Fixed - Size)))$
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Pro- perty	Informal and Formal (LTL/CTL) Specifications
P23	It is never the case that Set-size becomes “greater than fixed-size” i.e., ($Set - Size > Set - Fixed - Size$). $AG!(inp = size \& out = (Set - Fixed - Size) + 1)$
P24	It is never the case that the “Queue-size becomes less than zero”. $AG!(inp = size \& out = -1)$
P25	It never happens that the “Queue always remains Empty” (except at initial state). $AG!((state! = s0) \& EmptyQueue)$
P26	It never happens that the “Queue always remains Full” (except at final state). $AG!((inp = size \& out! = Queue - Fixed - Size) \& FullQueue)$
P27	It should never be the case that the two operations “offer” and “pol” occur at the same time. $AG!((inp = offer0 \& out = 0) \& (inp = pol \& out = 0))$
P28	It should never be the case that the two operations “ofer0” and “ofer1” occur at the same time. $AG!((inp = ofer0 \& out = 0) \& (inp = ofer1 \& out = 1))$
P29	It never happens that a “offer/pol” operation occurs but/and the Queue is in “Overflows” state. $AG!(inp = offer0 \& out = Queue - Fixed - Size + 1)$ Or. $AG!(inp = pol \& out = Queue - Fixed - Size + 1)$
P30	It never happens that a “offer/pol” operation occurs but/and the Queue is in “Underflow” state. $AG!(inp = offer0 \& out = -1)$ Or. $AG!(inp = pol \& out = -1)$
P31	It is never the case that the behavior of “offer/pol” operation is different in different state. $G!((inp = offer0 \& out = 0) \& X(!FullQueue \& inp = offer0 \& out = 1))$
Functional Requirement Specification	
P32	Queue should behaves like ‘FIFO’ manner(not like ‘LIFO’). $G(!FullQueue \rightarrow (inp = offer1 \rightarrow X(!FullQueue \& inp = offer0 \rightarrow X(inp = pol \rightarrow out = 1))))$
P33	It should be ‘Empty’ only at initial state (not in every state). $G((inp = size \& out! = 0) \rightarrow !EmptyQueue)$
P34	It should not be ‘Filled’ everywhere except at final state. $G((inp = size \& out! = Queue - Fixed - Size) \rightarrow !FullQueue)$
P35	“Insert (offer)” operation increases the Queue-size. $G((inp = offer0 inp = offer1) \rightarrow X(inp = size \rightarrow out! = 0))$ $G((inp = size \& out = 0) \rightarrow X(inp = offer0 inp = offer1 \rightarrow X(!FullQueue \& inp = offer0 inp = offer1 \rightarrow X(inp = size \rightarrow out = 2))))$
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Pro- perty	Informal and Formal (LTL/CTL) Specifications
P36	<p>“Remove (pol)” operation decreases the Queue-size.</p> $G((inp = size \& out = Queue - Fixed - Size) \rightarrow X(inp = pol \rightarrow X(inp = pol \rightarrow X(inp = size \rightarrow out = (Queue - Fixed - Size) - 2))))$
P37	<p>Queue should not “Overflow” or “Underflow”.</p> $G(!FullQueue \& (inp = offer0 inp = offer1) \rightarrow X(inp = size \rightarrow out! = Queue - Overflow))$ $G(!EmptyQueue \& inp = pol \rightarrow X(inp = size \rightarrow out! = Queue - Underflow))$
P38	<p>If equal number of “offer” operations followed by same number of “poll” operations occur then Queue-size should remain same.</p> $G((inp = size \& out = 1) \rightarrow X(inp = offer0 \rightarrow X(!FullQueue \& inp = offer1 \rightarrow X(inp = pol \rightarrow X(inp = pol \rightarrow X(inp = size \rightarrow out = 1))))))$
P39	<p>“Insert0 (ofer0)” operation must inserts “Underflow” (not 1) on the Queue. And “Insert1 (ofer1)” must inserts “Underflow” (not 0) on the Queue.</p> $inp = offer1 \rightarrow X(((inp = offer0 inp = offer1) U !FullQueue) \rightarrow (inp = pol \rightarrow out = 1))$ $inp = offer0 \rightarrow X(((inp = offer0 inp = offer1) U !FullQueue) \rightarrow (inp = pol \rightarrow out = 0))$
Liveness Requirement Specifications	
P40	<p>If “offer” operation is done forever on a non-filled Queue and “poll” operation never done then the Queue eventually becomes “Full”.</p> $G(!FullQueue \& ((inp = offer0 inp = offer1) U FullQueue) \rightarrow X(inp = size \rightarrow out = Queue - Fixed - Size))$ $G(!EmptyQueue \& (inp = offer0 inp = offer1) U FullQueue) \rightarrow X(inp = size \rightarrow out = Queue - Fixed - Size))$
P41	<p>If “pol” operation is done forever on non-empty Queue and “offer” operation never done then the Queue eventually becomes “Empty”.</p> $G(!FullQueue \rightarrow ((inp = pol U EmptyQueue) \rightarrow F(inp = size \rightarrow out = 0)))$ $G(!EmptyQueue \rightarrow ((inp! = offer0 inp! = offer1) \& inp = pol U EmptyQueue) \rightarrow X(inp = size \rightarrow out = 0))$
P42	<p>If “offer” operation wants to write on the Queue, then it eventually did this successfully without “Overflow”.</p> $AG(!FullQueue \& (inp = offer0 inp = offer1) \rightarrow AF(out! = Queue - Underflow out! = Queue - Overflow))$
P43	<p>If “pol” operation wants to remove an element from the Queue, then it eventually did this successfully without “Underflow”.</p> $AG(!EmptyQueue \& inp = pol \rightarrow AF(out! = Queue - Underflow))$

Table B.3: Requirement Specifications for Stack

Pro- perty	Informal and Formal (LTL/CTL) Specifications
	DEFINE Stack-Fixed-Size:= 3 to 16 DEFINE FullStack:= $((inp = pu0 inp = pu1)\&out = ful)$ DEFINE EmptyStack:= $(inp = size\&out = 0) (inp = pul\&out = null)$
	Safety Requirement Specifications
P44	It never happens that the Stack is “Full” and “Empty” simultaneously. $AG!(((inp = size\&out = 0) (inp = pop\&out = empty))\&((inp = pu0\&out = ful) (inp = pu1\&out = ful) (inp = size\&out = Stack - Fixed - Size)))$
P45	It is never the case that Stack-size becomes “greater than fixed-size” i.e., $(Stack - Size > Stack - Fixed - Size)$. $AG!(inp = size\&out = (Stack - Fixed - Size) + 1)$
P46	It is never the case that the “Stack-size becomes less than zero”. $AG!(inp = size\&out = -1)$
P47	It never happens that the “Stack always remains Empty” (except at initial state). $AG!((state! = s0)\&EmptyStack)$
P48	It never happens that the “Stack always remains Full” (except at final state). $AG!((inp = size\&out! = Stack - Fixed - Size)\&FullStack)$
P49	It should never be the case that the two operations “Push” and “Pop” occur at the same time. $AG!((inp = pu0\&out = 0)\&(inp = pop\&out = 0))$
P50	It should never be the case that the two operations “Push0” and “Push1” occur at the same time. $AG!((inp = pu0\&out = 0)\&(inp = pu1\&out = 1))$
P51	It never happens that a “Push/Pop” operation occurs but/and the Stack is in “Overflows” state. $AG!(inp = pop\&out = Stack - Fixed - Size + 1)$ Or. $AG!(inp = pu0\&out = Stack - Fixed - Size + 1)$
P52	It never happens that a “Push/Pop” operation occurs but/and the Stack is in “Underflow” state. $AG!((inp = pu0 inp = pu1)\&out = -1)$ Or. $AG!(inp = pol\&out = -1)$
P53	It is never the case that the behavior of “Push/Pop” operation is different in different state. $G!((inp = pu0\&out = 0)\&X(!FullStack\&inp = pu0\&out = 1))$
	Functional Requirement Specification
	Continued on next page

Table B.3 – continued from previous page

Pro- perty	Informal and Formal (LTL/CTL) Specifications
P54	Stack should behaves like “LIFO” manner not like “FIFO” manner. $G(!FullStack \rightarrow (inp = pu1 \rightarrow X(!FullStack \& inp = pu0 \rightarrow X(inp = pop \rightarrow out! = 1))))$
P55	It should be ‘Empty’ only at initial state (not in every state). $G((inp = size \& out! = 0) \rightarrow !EmptyStack)$
P56	It should not be ‘Filled’ everywhere except at final state. $G((inp = size \& out! = Stack - Fixed - Size) \rightarrow !FullStack)$
P57	“Push” operation increases the stack-size. $G((inp = size \& out = 1) \rightarrow X(inp = pu0 inp = pu1 \rightarrow X(!FullStack \& inp = pu0 inp = pu1 \rightarrow X(inp = size \rightarrow out = 3))))$
P58	“Pop” operation decreases the stack-size. $G((FullStack) \rightarrow X(inp = pop \rightarrow X(inp = pop \rightarrow X(inp = size \rightarrow out = (Stack - Fixed - Size) - 2))))$
P59	Stack should not “Overflow” or “Underflow”. $G(!FullStack \& (inp = pu0 inp = pu1) \rightarrow X(inp = size \rightarrow out! = Stack - Overflow))$ $G(!EmptyStack \& inp = pop \rightarrow X(inp = size \rightarrow out! = Stack - Underflow))$
P60	If equal number of “Push” operations followed by same number of “Pop” operations occur then stack-size should remain same. $G((inp = size \& out = 2) \rightarrow X(inp = pu0 \rightarrow X(!FullStack \& inp = pu1 \rightarrow X(inp = pop \rightarrow X(inp = pop \rightarrow X(inp = size \rightarrow out = 2))))))$
P61	“Push0” operation must inserts “Underflow” (not 1) on the Stack. And “Push1” must inserts “Underflow” (not 0) on the stack. $G(!FullStack \& inp = pu0 \rightarrow X(inp = pop \rightarrow out = 0))$ $G(!FullStack \& inp = pu1 \rightarrow X(inp = pop \rightarrow out = 1))$
Liveness Requirement Specifications	
P62	If “Push” operation is done forever on a non-filled stack and “Pop” operation never done then the Stack eventually becomes “Full”. $G(!FullStack \& ((inp = pu0 inp = pu1) U FullStack) \rightarrow X(inp = size \rightarrow out = Stack - Fixed - Size))$ $G(!EmptyStack \& (inp = pu0 inp = pu1) U FullStack) \rightarrow X(inp = size \rightarrow out = Stack - Fixed - Size))$
P63	If “Pop” operation is done forever on non-empty stack and “Push” operation never done then the Stack eventually becomes “Empty”. $G(!FullStack \rightarrow ((inp = pop U EmptyStack) \rightarrow (inp = size \rightarrow out = 0)))$ $G(!EmptyStack \rightarrow (inp = pop U EmptyStack) \rightarrow F(inp = size \rightarrow out = 0))$
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Table B.3 – continued from previous page

Pro- perty	Informal and Formal (LTL/CTL) Specifications
P64	<p>If “Push” operation wants to write on the Stack, then it eventually did this successfully without “Overflow”.</p> $AG((!FullStack \& (inp = pu0 inp = pu1)) \rightarrow AF(out! = Stack - Underflow out! = Stack - Overflow))$
P65	<p>If “Pop” operation wants to remove an element from the Stack, then it eventually did this successfully without “Underflow”.</p> $AG(!EmptyStack \& inp = pop \rightarrow AF(out! = Stack - Underflow))$