



CS320

Syntax and Semantics

Sukyoung Ryu

March 4, 2019

Hands-On Office Hour Tomorrow

March 5th Tuesday
@ E3-1, Room 1501
7:00PM–9:00PM

Bringing your own laptop is recommended.
Slides will be available at the course website.

Don'ts

- Don't import any other libraries.
- Don't use the break statement.
- Don't use the while and for loops.
- Don't use mutation.

Questions and Answers

```
// ... : AE => ...  
  
def ... (ae: AE): ... = ae match {  
  case Num(n) => ...  
  case Add(l, r) => ...  
  case Sub(l, r) => ...  
}
```

- 1 Is Num a constructor or a type?
- 2 Is match a function?
- 3 How is match different from if and switch?
- 4 Why is ae.left.num an error?

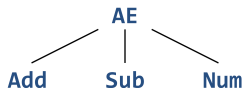
What Is Num?

- It is a constructor.

`Int => AE`

- It is a type.

`Num <: AE`



`Num(3) : Num <: AE`

`Num(3) : AE`

Pattern Matching

```
// interp : AE => Int  
def interp(ae: AE): Int = ae match {  
  case Num(n) => n  
  case Add(l, r) => interp(l) + interp(r)  
  case Sub(l, r) => interp(l) - interp(r)  
}  
interp(ae)
```

Pattern Matching

```
// interp : AE => Int
def interp(ae: AE): Int =
  if (ae.isInstanceOf[Num])
    ae.asInstanceOf[Num].n
  else if (ae.isInstanceOf[Add]) {
    val add = ae.asInstanceOf[Add]
    interp(add.left) + interp(add.right)
  }
  else {
    val sub = ae.asInstanceOf[Sub]
    interp(sub.left) - interp(sub.right)
  }
interp(ae)
```



Type Checking

```
<AE> ::= <num>
      | {+ <AE> <AE>}
      | {- <AE> <AE>}
```

```
trait AE
case class Num(num: Int) extends AE
case class Add(left: AE, right: AE) extends AE
case class Sub(left: AE, right: AE) extends AE
```

```
val ae = Add(Num(3), Sub(Num(8), Num(2)))
ae.left.num // error: value num is not a member of AE
```


Programming Language

A *programming language* is defined by

- syntax
a grammar for programs
- semantics
rules for evaluating any program to produce a result

Arithmetic Expressions: Concrete Syntax

```
<AE> ::= <num>  
      | {+ <AE> <AE>}  
      | {- <AE> <AE>}
```

{+ 3 {- 8 2}}

Arithmetic Expressions: Abstract Syntax

```
<AE> ::= <num>  
      | {+ <AE> <AE>}  
      | {- <AE> <AE>}
```

```
trait AE  
case class Num(n: Int) extends AE  
case class Add(l: AE, r: AE) extends AE  
case class Sub(l: AE, r: AE) extends AE
```

```
Add(Num(3), Sub(Num(8), Num(2)))
```



Parser for Arithmetic Expressions

```
// parser for AE
object AE extends RegexParsers {
  lazy val ae: Parser[AE] =
    int          ^^ { case n      => Num(n)      } |
    wrap("+" ~> ae ~ ae) ^^ { case l ~ r => Add(l, r) } |
    wrap("-" ~> ae ~ ae) ^^ { case l ~ r => Sub(l, r) }
  def apply(str: String): AE =
    parse(ae, str).getOrElse(error(s"bad syntax: $str"))
}
```

- int: receiving an integer value
- str: receiving a string value
- ~: concatenation of rules
- ~>: without receiving the value of the left-hand side
- wrap: enclosed by { and }

Parser for Arithmetic Expressions

```
// parser for AE
object AE extends RegexParsers {
  lazy val ae: Parser[AE] =
    int          ^^ { case n      => Num(n)      } |
    wrap("+" ~> ae ~ ae) ^^ { case l ~ r => Add(l, r) } |
    wrap("-" ~> ae ~ ae) ^^ { case l ~ r => Sub(l, r) }
  def apply(str: String): AE =
    parse(ae, str).getOrElse(error(s"bad syntax: $str"))
}
```

- int: receiving an integer value
- str: receiving a string value
- ~: concatenation of rules
- ~>: without receiving the value of the left-hand side
- wrap: enclosed by { and }

Parser for Arithmetic Expressions

```
// parser for AE
object AE extends RegexParsers {
  lazy val ae: Parser[AE] =
    int          ^^ { case n      => Num(n)      } |
    wrap("+" ~> ae ~ ae) ^^ { case l ~ r => Add(l, r) } |
    wrap("-" ~> ae ~ ae) ^^ { case l ~ r => Sub(l, r) }
  def apply(str: String): AE =
    parse(ae, str).getOrElse(error(s"bad syntax: $str"))
}
```

AE("3")

AE("{+ 3 4}")

AE("{+ {- 3 4} 7}")

AE("{- 5 1 2}")

Parser for Arithmetic Expressions

```
// parser for AE
object AE extends RegexParsers {
  lazy val ae: Parser[AE] =
    int ^^ { case n      => Num(n)      } |
    wrap("+ " ~> ae ~ ae) ^^ { case l ~ r => Add(l, r) } |
    wrap("- " ~> ae ~ ae) ^^ { case l ~ r => Sub(l, r) }
  def apply(str: String): AE =
    parse(ae, str).getOrElse(error(s"bad syntax: $str"))
}
```

AE("3")

AE("{+ 3 4}")

AE("{+ {- 3 4} 7}")

AE("{- 5 1 2}")



Parser for Arithmetic Expressions

```
// parser for AE
object AE extends RegexParsers {
  lazy val ae: Parser[AE] =
    int ^^ { case n      => Num(n)      } |
    wrap("+" ~> ae ~ ae) ^^ { case l ~ r => Add(l, r) } |
    wrap("-" ~> ae ~ ae) ^^ { case l ~ r => Sub(l, r) }
  def apply(str: String): AE =
    parse(ae, str).getOrElse(error(s"bad syntax: $str"))
}
```

Num(3)

Add(Num(3), Num(4))

Add(Sub(Num(3), Num(4)), Num(7))

Sub(Num(5), Num(1), Num(2))

Interpreter for Arithmetic Expressions

```
// interp : AE => Int
def interp(ae: AE): Int = ae match {
  case Num(n) => n
  case Add(l, r) => interp(l) + interp(r)
  case Sub(l, r) => interp(l) - interp(r)
}
```

Interpreter for Arithmetic Expressions

```
// interp : AE => Int
def interp(ae: AE): Int = ae match {
  case Num(n) => n
  case Add(l, r) => interp(l) + interp(r)
  case Sub(l, r) => interp(l) - interp(r)
}
```

```
test(interp(AE("3")), 3)
test(interp(AE("{+ 3 4}")), 7)
test(interp(AE("{+ {- 3 4} 7}")), 6)
```

Operational Semantics

- A method of defining the meaning of programs by describing the actions carried out during a program's execution.
- Many different styles
 - Evaluator semantics
 - Natural semantics, big-step
 - SOS semantics, small-step
 - Reduction semantics
 - Abstract machine semantics

Courtesy by David Van Horn

<https://www.youtube.com/watch?v=TU16mA5-i-g>

Operational Semantics for What?

- Specifying a programming language
- Communicating language design ideas
- Validating claims about languages
- Validating claims about type systems, etc
- Proving correctness of a compiler
- ...



AE: Concrete Syntax

$$\begin{aligned} \langle \text{AE} \rangle &::= \langle \text{num} \rangle \\ &| \{ + \langle \text{AE} \rangle \langle \text{AE} \rangle \} \\ &| \{ - \langle \text{AE} \rangle \langle \text{AE} \rangle \} \end{aligned}$$

AE: Abstract Syntax

```
trait AE
case class Num(n: Int) extends AE
case class Add(l: AE, r: AE) extends AE
case class Sub(l: AE, r: AE) extends AE
```

Syntax of AE (1)

$$n \in \mathbb{Z} \quad \Rightarrow \quad n \in \mathcal{A}$$

$$e_1 \in \mathcal{A} \wedge e_2 \in \mathcal{A} \quad \Rightarrow \quad \{+ \ e_1 \ e_2\} \in \mathcal{A}$$

$$e_1 \in \mathcal{A} \wedge e_2 \in \mathcal{A} \quad \Rightarrow \quad \{- \ e_1 \ e_2\} \in \mathcal{A}$$

Syntax of AE (2)

$$\frac{n \in \mathbb{Z}}{n \in \mathcal{A}}$$

$$\frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{+ \ e_1 \ e_2\} \in \mathcal{A}}$$

$$\frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{- \ e_1 \ e_2\} \in \mathcal{A}}$$

Syntax of AE (2)

Inference rules

$$\frac{H_1 \quad H_2 \quad \dots \quad H_n}{C}$$

Syntax of AE (2)

Inference rules

$$\frac{H_1 \quad H_2 \quad \dots \quad H_n}{C}$$

$$H_1 \wedge H_2 \wedge \dots \wedge H_n \Rightarrow C$$

Syntax of AE (2)

Proof of $\{+ 4 \{- 2 1\}\} \in \mathcal{A}$

$$\frac{n \in \mathbb{Z}}{n \in \mathcal{A}}$$

$$\frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{+ e_1 e_2\} \in \mathcal{A}}$$

$$\frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{- e_1 e_2\} \in \mathcal{A}}$$

Syntax of AE (2)

Proof of $\{+ 4 \{- 2 1\}\} \in \mathcal{A}$

$$\frac{n \in \mathbb{Z}}{n \in \mathcal{A}} \qquad \frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{+ e_1 e_2\} \in \mathcal{A}} \qquad \frac{e_1 \in \mathcal{A} \quad e_2 \in \mathcal{A}}{\{- e_1 e_2\} \in \mathcal{A}}$$

$$\frac{\frac{4 \in \mathbb{Z}}{4 \in \mathcal{A}} \quad \frac{\frac{2 \in \mathbb{Z}}{2 \in \mathcal{A}} \quad \frac{1 \in \mathbb{Z}}{1 \in \mathcal{A}}}{\{- 2 1\} \in \mathcal{A}}}{\{+ 4 \{- 2 1\}\} \in \mathcal{A}}$$

Syntax of AE (3)

$$\begin{aligned} \mathbb{Z} \quad n &::= \dots \mid -1 \mid 0 \mid 1 \mid \dots \\ \mathcal{A} \quad e &::= n \\ &\mid \{+ e e\} \\ &\mid \{- e e\} \end{aligned}$$

Evaluator Semantics of AE

```
// interp : AE => Int
def interp(ae: AE): Int = ae match {
  case Num(n) => n
  case Add(l, r) => interp(l) + interp(r)
  case Sub(l, r) => interp(l) - interp(r)
}
```

Natural Semantics of AE

$$\Rightarrow \subseteq \mathcal{A} \times \mathbb{Z}$$

$$\vdash \{+ 4 \{- 2 1\}\} \Rightarrow 5$$

Natural Semantics of AE

$$\vdash n \Rightarrow n$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2}$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{- e_1 e_2\} \Rightarrow n_1 - n_2}$$

Natural Semantics of AE

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2}$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2 \quad n_3 = n_1 + n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_3}$$

Natural Semantics of AE

Proof of $\vdash \{+ 4 \{- 2 1\}\} \Rightarrow 5$

$$\vdash n \Rightarrow n \qquad \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2}$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{- e_1 e_2\} \Rightarrow n_1 - n_2}$$

Natural Semantics of AE

Proof of $\vdash \{+ 4 \{- 2 1\}\} \Rightarrow 5$

$$\vdash n \Rightarrow n \qquad \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2} \qquad \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{- e_1 e_2\} \Rightarrow n_1 - n_2}$$

$$\frac{\vdash 4 \Rightarrow 4 \quad \frac{\vdash 2 \Rightarrow 2 \quad \vdash 1 \Rightarrow 1}{\vdash \{- 2 1\} \Rightarrow 1}}{\vdash \{+ 4 \{- 2 1\}\} \Rightarrow 5}$$

Natural Semantics of AE

Proof of $\vdash \{+ \{- 3 4\} 7\} \Rightarrow ?$

$$\vdash n \Rightarrow n \qquad \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2}$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{- e_1 e_2\} \Rightarrow n_1 - n_2}$$

Natural Semantics of AE

Proof of $\vdash \{-\ 5\ 1\ 2\} \Rightarrow ?$

$$\vdash n \Rightarrow n \qquad \frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{+ e_1 e_2\} \Rightarrow n_1 + n_2}$$

$$\frac{\vdash e_1 \Rightarrow n_1 \quad \vdash e_2 \Rightarrow n_2}{\vdash \{- e_1 e_2\} \Rightarrow n_1 - n_2}$$



Sukyoung Ryu

sryu.cs@kaist.ac.kr

<http://plrg.kaist.ac.kr>