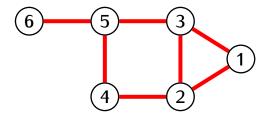
Definition

A *vertex cover* of a graph G is a set S of vertices of G such that every edge of G has at least one end in S.

A minimum vertex cover of G is a vertex cover such that there is no vertex cover with a smaller number of vertices.

Example.

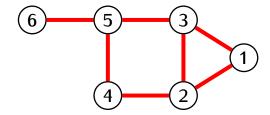


Definition

An *independent set* of a graph G is a set S of vertices of G such that there is no edge between any two elements of S.

A maximum independent set of G is an independent set such that there is no independent set with a larger number of vertices.

Example.



Proposition

Let G = (V, E) be an undirected graph.

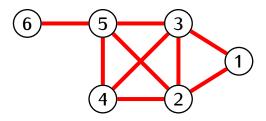
- 1) A set $S \subseteq V$ is independent if and only if the set $V \setminus S$ is vertex cover of G.
- 2) A set $S \subseteq V$ is a maximum independent set if and only if the set $V \setminus S$ is a minimum vertex cover of G.

Definition

Let G be an undirected graph. A *clique* is a set S of vertices of G such that any two vertices are connected by an edge.

A maximum clique of G is a clique such that there is no clique with a bigger number of vertices.

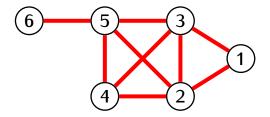
Example.



Definition

Let G be a simple graph. A complement of G is a graph G' such that G' has the same vertices as G, and two vertices are connected by an edge in G' if and only there is no edge between them in G.

Example.



Proposition

Let G be a simple graph and let G' be its complement. A set S of vertices G is a clique in G if and only if S is an independent set in G'.

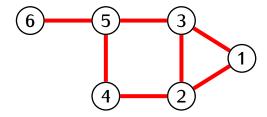
Corollary

Let G = (V, E) be a simple graph, let G' be its complement and let S be a set of vertices of G. The following conditions are equivalent:

- 1) The set S is a maximum clique in G.
- 2) The set S is a maximum independent set in G'.
- 3) The set $V \setminus S$ is a minimum vertex cover in G'.

Problem. Given a graph G = (V, E) find a minimum vertex cover of G.

Integer program reformulation:



An approximated solution of the minimum vertex cover problem:

Proposition

Assume that each minimum vertex cover of a graph G consists of N vertices. Let S_{LP} be a vertex cover selected using the solution of the linear program as described above. Then

$$|S_{LP}| \le 2N$$