

Example. Maximize

$$z = 3x_1 + x_2$$

subject to:

$$-x_1 + x_2 \leq 1$$

$$x_1 \leq 3$$

$$2x_1 + x_2 \leq 7$$

$$x_1, x_2 \geq 0$$

Convert to the equality form:

$$\begin{aligned} \text{Constraints: } -x_1 + x_2 + s_1 &= 1 \\ x_1 + s_2 &= 3 \\ 2x_1 + x_2 + s_3 &= 7 \\ x_1, x_2, s_1, s_2, s_3 &\geq 0 \end{aligned}$$

augmented  
matrix

$$\begin{array}{c} \text{free variables} \quad \text{basic variables} \\ \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \\ \begin{array}{ccccc|c} x_1 & x_2 & s_1 & s_2 & s_3 & \\ \hline -1 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 & 0 & 3 \\ 2 & 1 & 0 & 0 & 1 & 7 \end{array} \end{array}$$

We want to maximize

$$z = 3x_1 + x_2 + \underbrace{0s_1 + 0s_2 + 0s_3}_{\substack{\uparrow \text{coefficients of} \\ \text{are all 0.}} \text{ basic variables}}$$

Note: By setting the free variables to 0 we get a basic feasible solution:

$$\begin{cases} x_1 = 0 \\ x_2 = 0 \\ s_1 = 1 \\ s_2 = 3 \\ s_3 = 7 \end{cases} \quad \text{this gives : } z = 0$$

Goal: Look for other basic feasible solutions that make the value of  $z$  larger.

simplex method continued...

Simplex tableau:

	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$		
constraints	-1	1	1	0	0	1	$s_1$
	1	0	0	1	0	3	$s_2$
	2	1	0	0	1	7	$s_3$
the objective function	3	1	0	0	0	$z$	

The current basic feasible solution:

$$\begin{cases} x_1 = 0 \\ x_2 = 0 \\ s_1 = 1 \\ s_2 = 3 \\ s_3 = 7 \end{cases} \rightarrow z = 0$$

The pivot step:

- 1) We can increase the value of  $z$  by increasing either  $x_1$  or  $x_2$ . Let's try to increase  $x_2$ , keep  $x_1 = 0$ .
- 2) Change of  $x_2$  will affect the values of the basic variables  $s_1, s_2, s_3$ . These variables must stay  $\geq 0$ , which restricts possible values of  $x_2$ .

Since by assumption  $x_1 = 0$  we have:

$$x_2 + s_1 = 1 \Rightarrow s_1 = 1 - x_2 \geq 0 \quad \underline{\text{so:}} \quad x_2 \leq 1$$

$$0 \cdot x_2 + s_2 = 3 \Rightarrow s_2 = 3 \geq 0 \quad - \text{no restriction on } x_2$$

$$x_2 + s_3 = 7 \Rightarrow s_3 = 7 - x_2 \geq 0 \quad \underline{\text{so:}} \quad x_2 \leq 7$$

All these conditions are satisfied if  $x_2 \leq 1$ .

For the biggest increase of  $z$  we set  $x_2 = 1$  which makes  $s_1 = 0$ . Since we want free variables to have value 0, and basic variables to have non-zero values, we want to make  $x_2$  into a basic variable and make  $s_1$  free.

simplex method continued...

	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
$(-1)$	-1	1	1	0	0	1
$(-1)$	1	0	0	1	0	3
	2	1	0	0	1	7
	3	1	0	0	0	$z$

	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
	-1	1	1	0	0	1 $x_2$
	1	0	0	1	0	3 $s_2$
	3	0	-1	0	1	6 $s_3$
	4	0	-1	0	0	$z-1$

The new basic feasible solution

$$\begin{array}{l}
 \text{free} \\
 \text{basic}
 \end{array}
 \left[ \begin{array}{l}
 \begin{cases} x_1 = 0 \\ s_1 = 0 \end{cases} \\
 \begin{cases} x_2 = 1 \\ s_2 = 3 \\ s_3 = 6 \end{cases}
 \end{array} \right] \rightarrow \begin{array}{l} z-1 = 0 \\ \text{(or } z = 1) \end{array}$$

The pivot step:

The free variable we can increase are  $x_1$  and  $s_1$ .

Increasing  $s_1$  will decrease the value of  $z$ ,

but increasing  $x_1$  will increase  $z$ .

Thus we will increase the value of  $x_1$ , while keeping  $x_2, s_2, s_3 \geq 0$ .

(Note: we keep  $s_1 = 0$ )

$$-x_1 + x_2 = 1 \Rightarrow x_2 = 1 + x_1 \geq 0 \quad \text{- no restriction on } x_1$$

$$x_1 + s_2 = 3 \Rightarrow s_2 = 3 - x_1 \geq 0 \quad \underline{\text{so:}} \quad x_1 \leq 3$$

$$3x_1 + s_3 = 6 \Rightarrow s_3 = 6 - 3x_1 \geq 0 \quad \underline{\text{so}} \quad 3x_1 \leq 6 \\ x_1 \leq 2$$

The biggest value of  $x_1$  satisfying all conditions is  $x_1 = 2$

Then  $s_3 = 0$ , so  $s_3$  becomes a free variable and  $x_1$  becomes basic.

simplex method continued...

	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
	-1	1	1	0	0	1
	1	0	0	1	0	3
$(\frac{1}{3}) \cdot$	<span style="border: 1px solid red;">3</span>	0	-1	0	<span style="border: 1px solid red;">1</span>	6
	4	0	-1	0	0	$z-1$

→

	$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
	-1	1	1	0	0	1
	1	0	0	1	0	3
	<span style="border: 1px solid red;">1</span>	0	$-\frac{1}{3}$	0	$\frac{1}{3}$	2
	4	0	-1	0	0	$z-1$

Annotations: Red arrows show row operations:  $(-1) \times \text{row 3} \rightarrow \text{row 1}$ ,  $(1) \times \text{row 3} \rightarrow \text{row 2}$ , and  $(-4) \times \text{row 3} \rightarrow \text{row 4}$ .

The new basic feasible solution:

$$\begin{array}{l}
 \text{free} \quad \left\{ \begin{array}{l} s_1 = 0 \\ s_3 = 0 \end{array} \right. \\
 \text{basic} \quad \left\{ \begin{array}{l} x_2 = 3 \\ s_2 = 1 \\ x_1 = 2 \end{array} \right.
 \end{array}
 \rightarrow \begin{array}{l} z - 9 = 0 \\ \text{(or } z = 9) \end{array}$$

$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
0	1	$\frac{2}{3}$	0	$\frac{1}{3}$	3 $x_2$
0	0	$\frac{1}{3}$	1	$-\frac{1}{3}$	1 $s_2$
1	0	$-\frac{1}{3}$	0	$\frac{1}{3}$	2 $x_1$
0	0	$\frac{1}{3}$	0	$-\frac{4}{3}$	$z-9$

Annotations: Red arrows point to  $s_1$  and  $s_3$  with the label "free".

The pivot step:

We can increase  $z$  by increasing  $s_1$  (and keeping  $s_3 = 0$ )

We have:

$$x_2 + \frac{2}{3}s_1 = 3 \Rightarrow x_2 = 3 - \frac{2}{3}s_1 \geq 0 \quad \underline{\text{so:}} \quad \frac{2}{3}s_1 \leq 3$$

$$s_1 \leq \frac{9}{2}$$

$$\frac{1}{3}s_1 + s_2 = 1 \Rightarrow s_2 = 1 - \frac{1}{3}s_1 \geq 0 \quad \underline{\text{so:}} \quad \frac{1}{3}s_1 \leq 1$$

$$s_1 \leq 3$$

$$x_1 - \frac{1}{3}s_1 = 2 \Rightarrow x_1 = 2 + \frac{1}{3}s_1 \geq 0 \quad \text{no restriction on } s_1$$

The largest value of  $s_1$  satisfying all conditions is  $s_1 = 3$ . Then  $s_2 = 0$ , so  $s_2$  becomes a free variable and  $s_1$  basic

simplex method continued...

3.

$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
0	1	$2/3$	0	$1/3$	3
0	0	$1/3$	1	$-1/3$	1
1	0	$-1/3$	0	$1/3$	2
0	0	$1/3$	0	$-4/3$	$z-9$

→

$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
0	1	$2/3$	0	$1/3$	3
0	0	1	3	-1	3
1	0	$-1/3$	0	$1/3$	2
0	0	$1/3$	0	$-4/3$	$z-9$

Annotations:  $(-\frac{2}{3})$  from row 2 to row 1;  $(\frac{1}{3})$  from row 2 to row 3;  $(-\frac{1}{3})$  from row 2 to row 4.

The new basic feasible solution

free

$$\begin{cases} s_2 = 0 \\ s_3 = 0 \end{cases}$$

basic

$$\begin{cases} x_2 = 1 \\ s_1 = 3 \\ x_1 = 3 \end{cases}$$

$$z - 10 = 0 \quad (\text{or: } z = 10)$$

free

$x_1$	$x_2$	$s_1$	$s_2$	$s_3$	
0	1	0	-2	1	1 $x_2$
0	0	1	3	-1	3 $s_1$
1	0	0	1	0	3 $x_1$
0	0	0	-1	-1	$z-10$

Note: The objective function depends on  $s_2$  and  $s_3$  only, but increasing these variables will make  $z$  smaller. This means that we can't increase the value of  $z$  anymore, we reached the maximum.

The final solution of the linear program:

The maximum value is  $z = 10$ .

It is obtained for  $x_1 = 3$ ,  $x_2 = 1$ .

## Geometric interpretation of the simplex method

Recall: Maximize

$$z = 3x_1 + x_2$$

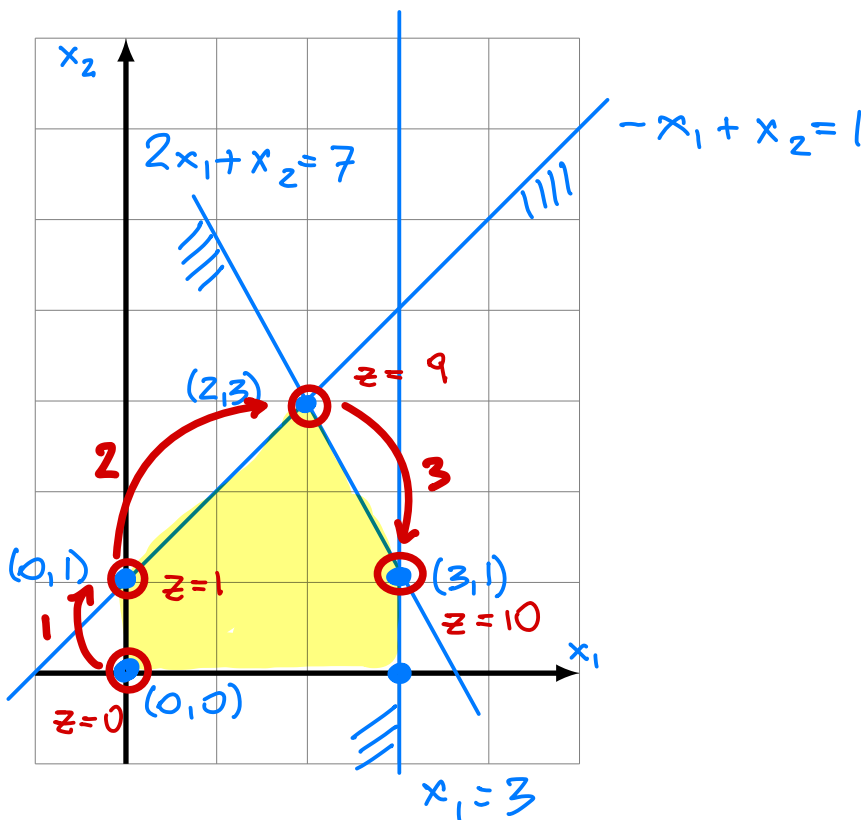
subject to:

$$-x_1 + x_2 \leq 1$$

$$x_1 \leq 3$$

$$2x_1 + x_2 \leq 7$$

$$x_1, x_2 \geq 0$$



Note: It would have been more efficient to increase  $x_1$  instead of  $x_2$  in the first pivot step. In the simplex method there are various pivoting rules which aim to decide which free variable should be used at each step to get to the solution as quickly as possible.