

**Recall:** The standard way of computing eigenvalues and eigenvectors of an  $n \times n$  matrix  $A$ :

- 1) Compute the characteristic polynomial  $P(\lambda) = \det(A - \lambda I_n)$ .
- 2) Eigenvalues of  $A$  = roots of  $P(\lambda)$ .
- 3) Eigenvectors corresponding to an eigenvalue  $\lambda$  = vectors in  $\text{Nul}(A - \lambda I_n)$ .

**Problems:**

- For large matrices computations of  $P(\lambda)$  are slow.
- Even if we know  $P(\lambda)$ , it is difficult to compute its roots.

**More efficient way:** The power method.

**Assumptions:**

- $A$  is an  $n \times n$  matrix with eigenvalues  $\lambda_1, \dots, \lambda_n$  such that

$$|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$$

- For  $i = 1, \dots, n$ , by  $\mathbf{w}_i$  we will denote an eigenvector corresponding to the eigenvalue  $\lambda_i$ , such that  $\|\mathbf{w}_i\| = 1$ .

## Computing the largest eigenvalue $\lambda_1$ its eigenvector

- Start with a vector  $v \neq 0$ . Since  $A$  is diagonalizable we have  $v = \sum_i c_i w_i$  for some  $c_i \in \mathbb{R}$ . If  $v$  is selected at random, then almost always we will have  $c_1 \neq 0$ .

- Denote:  $v_k = \frac{A^k v}{\|A^k v\|}$

We have:  $A^k v = A^k (\sum c_i w_i) = \sum c_i (A^k w_i) = \sum c_i \lambda_i^k w_i$

This gives:

$$v_k = \frac{\sum c_i \lambda_i^k w_i}{\|\sum c_i \lambda_i^k w_i\|} = \frac{\lambda_1^k (c_1 w_1 + \sum_{i=2}^n c_i (\frac{\lambda_i}{\lambda_1})^k w_i)}{|\lambda_1|^k \cdot \|c_1 w_1 + \sum_{i=2}^n c_i (\frac{\lambda_i}{\lambda_1})^k w_i\|}$$

Since  $|\frac{\lambda_i}{\lambda_1}| < 1$  for  $i \geq 2$  thus  $\lim_{k \rightarrow \infty} (\frac{\lambda_i}{\lambda_1})^k = 0$

Thus for large  $k$  we get

$$v_k \approx \frac{\lambda_1^k c_1 w_1}{|\lambda_1|^k \|c_1 w_1\|} = \frac{\lambda_1^k}{|\lambda_1|^k} \cdot \frac{c_1}{|c_1|} \cdot \frac{w_1}{\|w_1\|} \stackrel{\|w_1\|=1}{=} \pm 1 \cdot w_1$$

Note: Both  $w_1$  and  $-w_1$  are eigenvectors corresponding to the eigenvalue  $\lambda_1$ , so it does not matter if we get  $v_k \approx w_1$  or  $v_k \approx -w_1$

### Computations of $\lambda_1$

Since  $v_k \approx \pm w_1$  thus  $A v_k \approx \lambda_1 v_k$ .

Dividing entries of  $A v_k$  by entries of  $v_k$  we get an approximated value of  $\lambda_1$ .

Note: In practice the vectors  $v_k$  are computed iteratively:

$$v_0 = v, \quad v_1 = \frac{A v_0}{\|A v_0\|}, \quad v_2 = \frac{A v_1}{\|A v_1\|}, \quad \dots, \quad v_k = \frac{A v_{k-1}}{\|A v_{k-1}\|}$$

## Computing the other eigenvalues and eigenvectors

Idea: If we could start the power method with a vector  $v = \sum_i c_i w_i$  such that  $c_1 \neq 0$  then the power method would compute  $w_2$  and  $\lambda_2$ .

Problem: How to find such a vector  $v$ ?

Solution for a symmetric matrix  $A$ :

Recall: For a symmetric matrix eigenvectors for different eigenvalues are orthogonal:

$$w_i \cdot w_j = 0 \text{ for } i \neq j$$

- start with an arbitrary vector  $v = \sum c_i w_i$
- use  $v$  and the power method to compute  $w_1$ .
- we have:

$$\begin{aligned} v \cdot w_1 &= \left( \sum c_i w_i \right) \cdot w_1 = \sum c_i (w_i \cdot w_1) \\ &= c_1 w_1 \cdot w_1 = c_1 \end{aligned}$$

$\uparrow$   
 $\|w_1\| = 1$

- take  $v' = v - (v \cdot w_1) w_1 = \sum c_i w_i - c_1 w_1 = \sum_{i=2}^n c_i w_i$

We can use the power method starting with  $v'$  to compute  $w_2$  and  $\lambda_2$

Note: Due to rounding errors in the power method we need to make the vector  $v'_k$  orthogonal to  $w_1$  at each iteration step.

We can continue this process iteratively to compute all other eigenvectors and their corresponding eigenvalues.