

Recall: The general form of a linear program

For the objective variables x_1, \dots, x_n find the minimum (or the maximum) of the objective function

$$Z = c_1x_1 + \dots + c_nx_n$$

subject to the constraints:

$$a_{i1}x_1 + \dots + a_{in}x_n \begin{matrix} \leq \\ = \\ \geq \end{matrix} b_i$$

for $i = 1, \dots, m$, and possibly $x_j \geq 0$ for $j = 1, \dots, n$.

The *equality* (or *standard*) form of a linear program:

- we are looking for the maximum;
- all constraints are of the form

$$a_{i1}x_1 + \dots + a_{in}x_n = b_i$$

- we require that $x_j \geq 0$ for $j = 1, \dots, n$.

Fact

Every linear program can be converted to the equality form.

- finding minimum of $z = c_1x_1 + c_2x_2 + \dots + c_nx_n$
||
finding maximum of $z = -c_1x_1 - c_2x_2 - \dots - c_nx_n$
- if we have a constraint
$$a_{i1}x_1 + \dots + a_{in}x_n \leq b_i$$
we can replace it by
$$a_{i1}x_1 + \dots + a_{in}x_n + s_i = b_i$$
where $s_i \geq 0$ is a new slack variable
- a constraint of the form
$$a_{i1}x_1 + \dots + a_{in}x_n \geq b_i$$
is the same as
$$-a_{i1}x_1 - \dots - a_{in}x_n \leq -b_i$$
so we can use a slack variable again:
$$-a_{i1}x_1 - \dots - a_{in}x_n + s_i = -b_i$$
$$s_i \geq 0.$$
- if x_j is an unrestricted variable (i.e. can be any real number), then we can replace it by $x_j = x_j^+ - x_j^-$ where $x_j^+, x_j^- \geq 0$ - new variables.

Example. Convert the following linear program to the equality form.

Minimize the function

$$z = 6x_1 - 10x_2$$

subject to the constraints:

$$5x_1 + 7x_2 \leq 8$$

$$4x_1 + 2x_2 \geq 10$$

$$x_1 \geq 0$$

$$x_2 \in \mathbb{R}$$

Solution:

Maximize

$$z = -6x_1 + 10x_2$$

Constraints:

$$5x_1 + 7x_2 + s_1 = 8$$

$$-4x_1 - 2x_2 + s_2 = -10$$

$$x_1 \geq 0$$

$$s_1 \geq 0$$

$$s_2 \geq 0$$

$$x_2 \in \mathbb{R}$$

Maximize:

$$z = -6x_1 + 10x_2^+ - 10x_2^- + 0s_1 + 0s_2$$

$$5x_1 + 7x_2^+ - 7x_2^- + s_1 = 8$$

$$-4x_1 - 2x_2^+ + 2x_2^- + s_2 = -10$$

$$x_1 \geq 0$$

$$x_2^+ \geq 0$$

$$x_2^- \geq 0$$

$$s_1 \geq 0$$

$$s_2 \geq 0$$

$$x_2 = x_2^+ - x_2^-$$

the equality form

The *inequality* form of a linear program:

- we are looking for the maximum;
- all constraints are of the form

$$a_{i1}x_1 + \dots + a_{in}x_n \leq b_i$$

- we require that $x_j \geq 0$ for $j = 1, \dots, n$.

Fact

Every linear program can be converted to the inequality form.

- if we have a constraint of the form

$$a_{i1}x_1 + \dots + a_{in}x_n = b_i$$

then we can replace it by

$$a_{i1}x_1 + \dots + a_{in}x_n \leq b_i$$

$$a_{i1}x_1 + \dots + a_{in}x_n \geq b_i$$