Emulation Framework for AIoT Federated Learning

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1 Abstract

Federated learning allows a fleet of devices to collaborate towards a globally trained machine learning model. Research has continued to produce novel federated learning algorithms to tackle different issues in FL such as heterogeneity and learning over data from non-identical distributions. Performance of these algorithms depend in part on the system parameters used in FL such as number of clients and number of passes. Furthermore, a realistic benchmark would require one to procure a large fleet of devices. This work seeks to introduce a software emulation framework to streamline the process of building a fleet of clients and allow easy testing of FL system parameters for configuration of optimal values.

2 Acknowledgements

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3 Introduction

Federated learning emerged as a method for solving key issues with the standard centralized learning approach. Some of these issues are (1) Preserving user data privacy: a centralized training approach involves the need for the central machine performing the computation to have full access to all the data. With FL, data never leaves each individual client's device. Instead, only the updated weights are shared to form the global model.(2) Scalability: FL enables leveraging a network to perform computation in parallel.
However, with the distributed nature of FL comes a few key challenges.

- Limited computing resources: individual devices may have constraints on memory, limiting the size of the local model it is able to train. Constraints on computing power can also lead to longer training times. This is particularly so for Internet of Things (IoT) devices such as sensors, where their embedded nature leads to a limited size and power.

- Network limitations: communication speed can become the bottleneck for performance as IoT devices rely on unstable wireless communication networks. Furthermore, data constraints can exist such that the number of bytes sent may be a cost inducing factor.

3.1 Problem

However, prior to effectively deploying FL on resource-constrained mobile devices in large scale, different factors including the convergence rate, energy efficiency and model accuracy should be well studied.

FLSim[**?**] aims to provide a simulation framework for FL by offering users a set of software components which can then be mixed and matched to create a simulator for their use case. Developers need only define the data, model and metrics reported, and FL system parameters can be altered in a JSON configuration file. FLUTE[**?**] adds additional features by allowing users to gain access to cloud based compute and data. However, these simulation solutions are consequently hardware-agnostic. Without taking into account the specific platforms which FL is performed on, these simulators cannot provide insight into whether the system will work in the production environment. Here, this work proposes a new emulation framework zfl, which aims to solve this issue by running FL on emulated hardware.

3.2 QEMU

QEMU[**?**] is an open source machine emulator and virtualizer. It enables system emulation, where it provides a virtual model of an entire machine (CPU, memory and emulated devices) to run a guest OS. In this mode the CPU may be fully emulated, or it may work with a hypervisor to allow the guest to run directly on the host CPU. In zfl, QEMU with full CPU emulation is used without a hypervisor as part of its software architecture to emulate hardware.

3.3 Zephyr OS

To emulate the issue of limited computing resources, it is important to make use of system runtimes used by those devices. In particular, the type of operating system used will help to ensure that the kernel is lightweight and configurable. Zephyr OS[**?**] is one such OS. It is based on a small-footprint kernel designed for use on resource-constrained and embedded systems: from simple embedded environmental sensors and LED wearables to sophisticated embedded controllers, smart watches, and IoT wireless applications. Furthermore, Zephyr is highly configurable, allowing the user to choose only the specific kernel services required, and also delve into lower level memory allocations of the system RAM. In zfl, client code will be written to work in the Zephyr OS environment.

4 Implementation

zfl aims to provide the ability to emulate the traditional FedAvg[**?**] algorithm, with multiple clients communicating with a central server which performs the aggregation. To run on Zephyr OS, the entire software stack is developed in the C programming language. This also makes it suitable to run on embedded devices.

Next the framework needed a method for spawning an arbitrary number of QEMU instances, and allowing these instances to communicate back to the server. When called in client mode, zfl accomplishes this by making use of the ‘fork‘ and ‘exec‘ pattern with the desired number of clients.

However, each instance of QEMU is persists in an isolated network different from the host PC, and is not able to communicate. We can bypass this limitation using a network bridge to act as a virtual network device forwarding packets between connected network devices. The network bridge is set up on the host under the name zfl with a set of network parameters using the `ip` command line utility.

```
ip link add $INTERFACE type bridge
ip addr add $IPV4_ADDR_1 dev $INTERFACE
ip link set enp61s0 master $INTERFACE
ip route add $IPV4_ROUTE_1 dev $INTERFACE > /dev/null 2>&1
ip link set dev $INTERFACE up
```

Listing 1: Network bridge setup

To tell QEMU to use it, we pass the name of the bridge along with a randomly generated MAC address as arguments to the -nic flag

```
-nic bridge,model=e1000,mac=%s,br=zfl
```

Although each QEMU client is now able to communicate with the host via the nic adapter, we still needed a way to monitor the output of each instance. One method is to transmit output and logs over the network. However, this would not allow important crash logs and stacktrace information to be transmitted as the application software would have shutdown. To overcome this, the host creates a named or FIFO pipe [**?**] for each client which is passed in to QEMU through the -serial flag. Now, all standard output goes through the named pipe, ready to be read by the host. Another outcome of this is that the host is now also able to send input to each instance, whose importance will be described in the next section.

```
pid_t child = fork();
if (child < 0) {
    printf("ERROR: could not fork client %d: %s\n", i,
   strerror(errno));
    return 1;
}
...

// generate serial arguments
char serial_arg[80];
snprintf(serial_arg, 80, "pipe:%s", pipe_path);

// generate nic arguments
```

```
13 char nic_arg[100];
14 char *mac = generate_random_mac();
15 snprintf(nic_arg, sizeof(nic_arg), "bridge,model=e1000,mac=%s
      ,br=zfl", mac);
16
17 // start client as new process
18 execlp("qemu-system-i386", "qemu-system-i386",
19
20         "-m", "15", "-cpu", "qemu32,+nx,+pae", "-machine", "
      q35",
21         "-device", "isa-debug-exit,iobase=0xf4,iosize=0x04",
22
23         "-no-reboot", "-nographic", "-no-acpi",
24
25         "-serial", serial_arg,
26
27         "-nic", nic_arg,
28
29         "-kernel", "./zflclient/out/zephyr/zephyr.elf",
30
31         NULL);
```

Listing 2: Client forking process

4.1 Client

In addition to the MAC address configuration, each client needed a unique Internet Protocol Version 4 (IPv4) address in order to establish TCP based connections with the central server. Samples provided by Zephyr OS describe a way to achieve this my setting a compile time configuration flag `CONFIG_NET_CONFIG_MY_IPV4_ADDR`. However, this is impractical to scale to a number of clients, since each client would need a separately compiled binary. Instead, we can assign the IPv4 address dynamically using the built-in Zephyr network function:

`net_if_ipv4_addr_add`

To achieve this, each client starts off as a Zephyr shell instance and a user-defined command is registered as a way to start the main program. The desired IP address is obtained as a command line argument.

```
SHELL_CMD_ARG_REGISTER(run, NULL, "Run with IPv4 address",
    run, 4, 0);
```

Listing 3: Registering user defined command "run" to the function pointer run

```
    char *addr_str = argv[1];
    LOG_INF("instance ipaddr is %s", addr_str);
    struct in_addr addr;
    zsock_inet_pton(AF_INET, addr_str, &addr);
    if (!net_if_ipv4_addr_add(net_if_get_default(), &addr,
    NET_ADDR_MANUAL, UINT32_MAX)) {
        LOG_ERR("failed to add %s to interface", addr_str);
        return -1;
    }
```

Listing 4: the "run" function which performs IP address assignment at runtime

Finally, each client establishes TCP socket connection to the central server and obtains their assigned ID and training data according to that ID. Once done, each client marks itself as ready to begin the training round. The complete initialization process is illustrated below:

4.2 Server

The central server runs on the host machine without needing additional initialization. After assigning each connecting client their ID and training data, it performs the function `start_round` at an interval of 10 seconds. Each time, the server pings all previously connected clients to check if they are ready to start the training round. When enough clients are ready, it sends a HTTP request to the start endpoint of the client, which triggers the training function.

5 Experiments

6 Limitations

6.1 Obtaining client training data