

Spec*Bench: Benchmarking Specification Generation with Automated Verification

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The Problem

Spec checking lacks guarantees (soundness, completeness, independence).

 Property-based testing (PBT) compares specs by examples, providing no correctness guarantees [4]:

```
Equivalence requires: \forall x \in D, \ S_{\text{gen}}(x) = S_{\text{ref}}(x)
           Samples are evidence, not a guarantee.
```

 LLM-assisted ITP proofs (e.g., Lean, Rocq) offload substantial effort to LLMs, resulting in 0% accuracy on state-of-the-art models [2]:

```
-- Goal: forall x, Sgen x <-> Sref x
theorem equiv : forall x, Sgen x <-> Sref x := by
 -- Generated equivalence proof - LLMs achieve 0% accuracy!
 sorry -- proof not found => no guarantee
```

• Conclusion: Sample agreement \Rightarrow equivalence; we need $\forall x. S_{\text{gen}}(x) \equiv S_{\text{ref}}(x)$ or a counterexample. LLM + ITP checking is not feasible.

Contributions

 Spec*Bench introduces the first benchmark and protocol for spec generation with verifiable and automated checking, considering only specifications whose proofs can be offloaded to an automated verifier.

Method	Sound?	LLM-independent Verification?	Successful in Practice?
LLM as a Judge	×	×	?
PBT	×	\checkmark	\checkmark
LLM + ITP	\checkmark	×	×
Spec*Bench		\checkmark	\checkmark

Table 1. Trade-offs among existing specification generation benchmarks.

Background

- Specification. A predicate that formalizes a natural-language description into precise mathematical conditions.
- Correctness. A generated specification $S_{\rm gen}$ is correct if the base specification $S_{\rm ref}$ can be formally proven equivalent to it using an automated (ATP) or interactive (ITP) theorem prover (e.g., SMT solvers such as Z3, or Coq).
- Soundness. ITPs and ATPs are sound, meaning if they accept a proof of equivalence between $S_{\rm gen}$ and $S_{\rm ref}$, then that equivalence is indeed valid.

Github: github.com/bbayazit16/specbench

Method

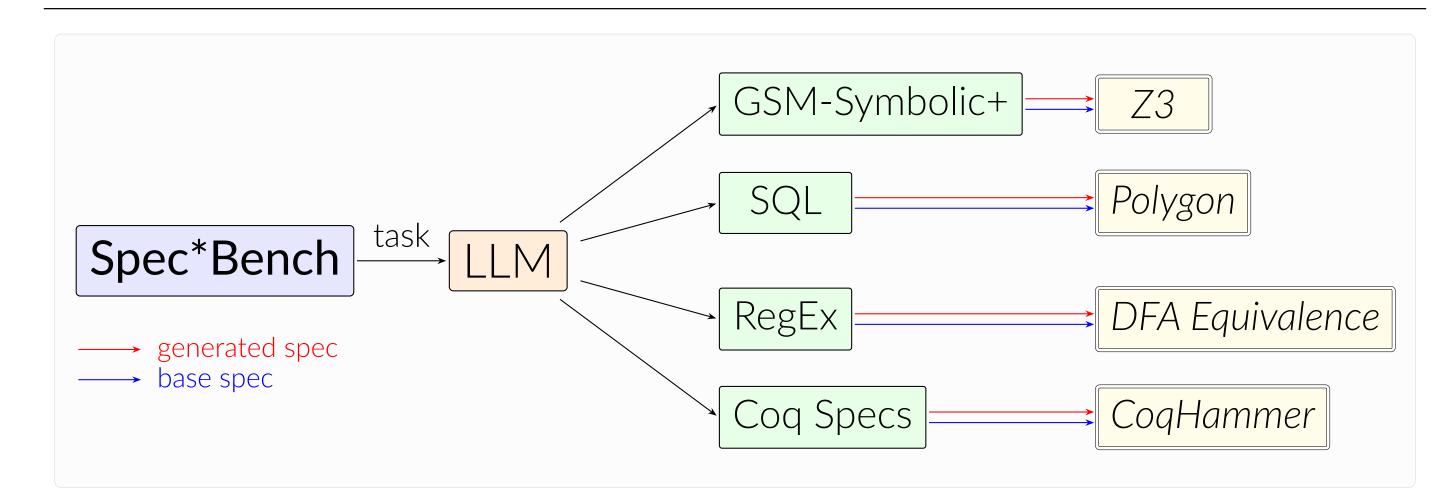
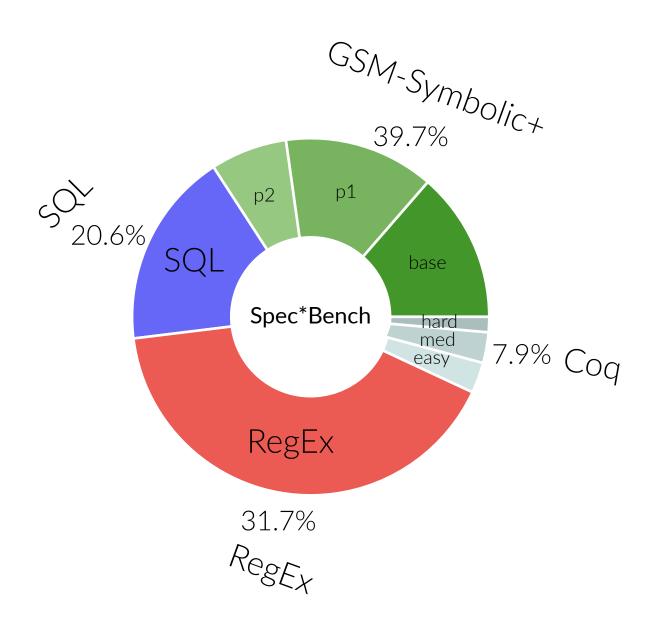
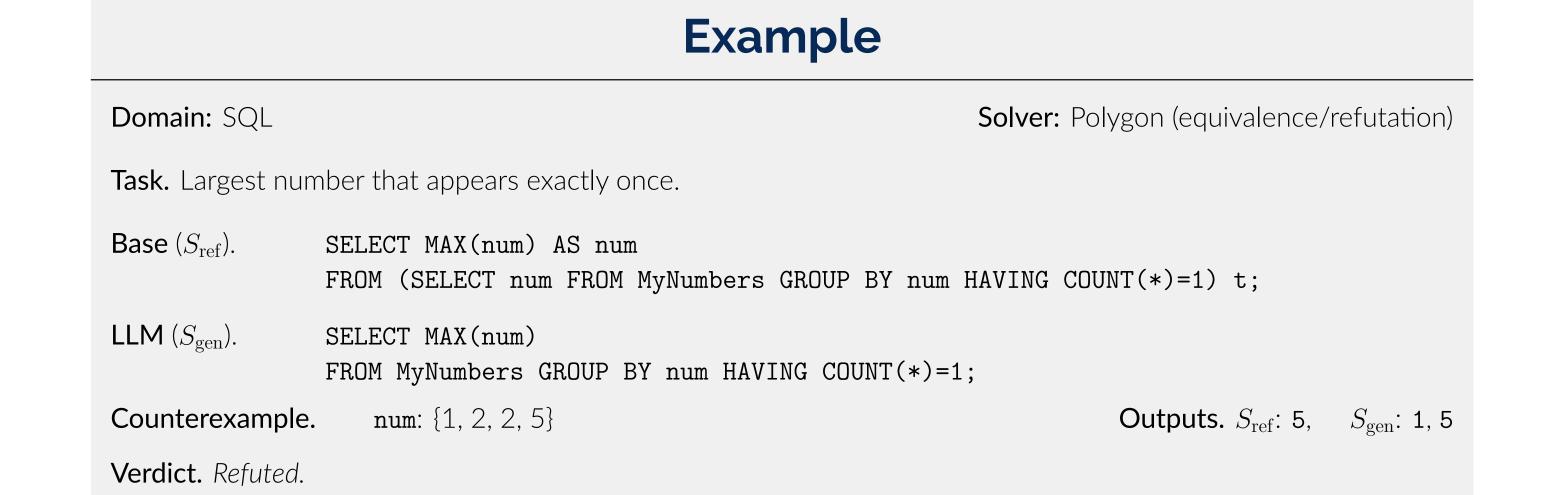


Figure 1. Overview of Spec*Bench methodology. Each natural language task is given to an LLM, which produces a candidate specification. The generated specification is compared against the base specification using domain-specific verifiers.



- Input. A natural-language task picked from one of the domains and a hidden base specification $S_{\rm ref}$.
- Generation. The LLM proposes a candidate specification $S_{\rm gen}$.
- Verification. Check $S_{\rm gen} \equiv S_{\rm ref}$ with the respective solver.
- Verdict. Verified, Refuted (counterexample), or *Unknown* (timeout/failure).

Figure 2. Spec*Bench task composition across four domains. The chart shows both main domain percentages and their sub-domain breakdown: GSM-Symbolic+ contains three variants (base, p1, p2), while Coq Specifications have three difficulty levels (easy, medium, hard).



Evaluation

- **Protocol.** For each task \times model, generate top-k candidates, and verify each with the corresponding solver under fixed time limits. With r optional refinement rounds, if Refuted, feed the counterexample to the LLM, and repeat.
- Unknown policy. If the limits are exceeded or a proof is not found/reconstructed, the verdict is Unknown. Unknowns are reported separately and not as correct.
- Reporting. The Verdict, Time2Verdict, Resource Limits, and Counterexamples.
- Overall Accuracy. Accuracy is calculated by total Success verdicts divided by total items across domains.

Results

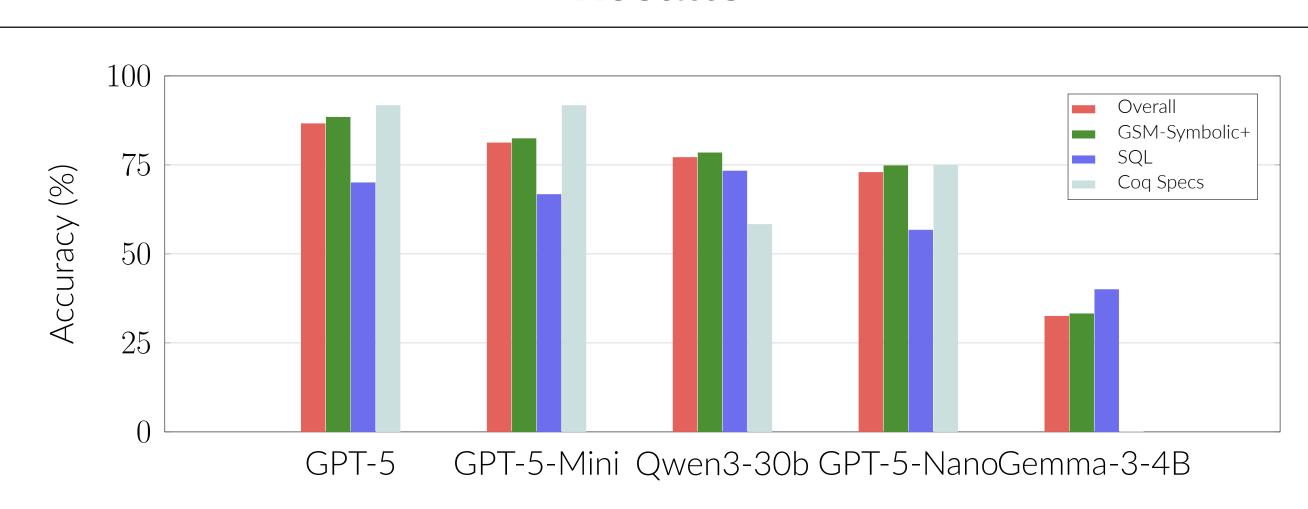


Figure 3. Spec*Bench performance across models and domains for a subset of the benchmark, with no refinement rounds and a 4 second timeout. RegEx was omitted for this experiment.

Analysis

- **SQL evaluation gap.** Most SQL failures are *still accepted* by LeetCode. Counterexamples are valid and fail LeetCode's own "Run Test." \Rightarrow **Tests are** insufficient to evaluate specifications.
- Numbers vs. variables. Accuracy drops from numeric variations are known [1], but GSM-Symbolic+ reveals a further decline when templating with variables.
- Scaling for Coq. Smaller models underperform at writing Coq specifications.

References

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