# Teorija Izračunljivosti 2020–21 Homework 2 Computability on $\mathbb{R}$

Homework released: December 23, 2020

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This homework is a paper-and-pencil exercise. Questions 1 and 2 concern computability using ordinary Turing machines. Questions 3 and 4 address real number computation using Type 2 Turing machines exploiting the Cauchy representation of  $\mathbb{R}$  defined in Lecture 12.

If a question asks you to show that a function is computable, you should provide an *informal description* of a TM/T2M algorithm. Do not attempt to explicitly define an actual machine.

The homework must be completed and submitted by 23:59 (CET) on Sunday 17th January 2021. Submission instructions will be announced by email and posted on the course webpage.

### Question 1 [1 mark]

Given a finite alphabet  $\Sigma$ , give a definition in terms of ordinary Turing machines of what it means for a partial function  $f: \Sigma^* \times \mathbb{N} \to \Sigma^*$  to be *computable*.

## Question 2 [4 marks]

Let  $q_d$  be the representation of dyadic rationals by words over the alphabet  $\Sigma_b$ , as defined in Lecture 12. Show that there exists a computable partial function  $f \colon \Sigma_b^* \times \mathbb{N} \rightharpoonup \Sigma_b^*$  that satisfies:

- 1.  $dom(f) = \{u \in dom(q_d) \mid q_d(u) \neq 0\} \times \mathbb{N} \text{ and }$
- 2. for any  $(u, n) \in \mathsf{dom}(f)$ ,  $|q_d(f(u, n)) q_d(u)^{-1}| \le 2^{-n}$ .

(You may assume without proof that addition and multiplication of dyadic rationals are computable as functions on words over  $\Sigma_b$ .)

### Question 3 [4 marks]

Prove that the partial function  $r: \mathbb{R} \to \mathbb{R}$  below is computable.

$$r(x) \simeq \begin{cases} x^{-1} & \text{if } x \neq 0 \\ \uparrow & \text{if } x = 0 \end{cases}$$

## Question 4 [1 mark]

Does there exist any total computable function  $r': \mathbb{R} \to \mathbb{R}$  such that  $r'(x) = x^{-1}$  for all  $x \neq 0$ ?