## **Project 6 Report**

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## Part I

Q: Explain how you implement your code clearly with screenshots of your functions.

Step 1. After reading the list of data, I split the list into two halves and create two threads running function 'multi\_thread\_morge\_sort' to sort them respectively.

```
// Split the list into two halves and sort them in parallel
pthread_t threads[2];
ThreadData thread_data[2];

thread_data[0].thread_id = 0;
thread_data[0].l = 0;
thread_data[0].r = list_size / 2 - 1;
pthread_create(&threads[0], NULL, multi_thread_merge_sort, &thread_data[0]);

thread_data[1].thread_id = 1;
thread_data[1].l = list_size / 2;
thread_data[1].r = list_size - 1;
pthread_create(&threads[1], NULL, multi_thread_merge_sort, &thread_data[1]);
```

Step 2. When running 'multi\_thread\_morge\_sort', I use merge sort algorithm to sort the list.

```
void *multi_thread_merge_sort(void *arg)
{
    ThreadData *data = (ThreadData *)arg;
    merge_sort(data→l, data→r);
    pthread_exit(NULL);
}
```

Step 3. Use 'pthread\_join' to wait for the two threads to finish sorting.

```
// Wait for the two halves to be sorted
pthread_join(threads[0], NULL);
pthread_join(threads[1], NULL);
```

Step 4. Merge the two sorted lists into one sorted list.

```
// Merge the two sorted halves
merge(0, list_size / 2 - 1, list_size - 1);
```

Step 5. Write the sorted list into the output file.

Following is the screenshot of merge & merge\_sort function: merge:

```
void merge(int l, int m, int r)
    int i, j, k;
int n1 = m - l + 1; // size of left subarray
        R[j] = list[m + 1 + j];
        if (L[i] \leq R[j]) // Change this to L[i] > R[j] for descending order
            list[k] = L[i];
            list[k] = R[j];
   while (i < n1)
        list[k] = R[j];
```

merge\_sort:

```
void merge_sort(int l, int r)
{
    if (l < r)
    {
        int m = l + (r - l) / 2; // Same as (l+r)/2, but avoids overflow for large l and r
        merge_sort(l, m);
        merge_sort(m + 1, r);
        merge(l, m, r);
    }
}</pre>
```

## Part II

Q1: Describe which lines of code you modify with detailed explanations.

A1: I mainly modify the function "calculate\_innner\_product" as following:

```
// Function to calculate the inner product for a portion of the vectors
void* calculate_inner_product(void* arg) {
    ThreadData* data = (ThreadData*)arg;

    int temp_inner_product = 0;
    for (int i = data → start; i < data → end; i++) {
        temp_inner_product += vector1[i] * vector2[i];
    }

    // Lock the mutex before updating the global inner product
    pthread_mutex_lock(&mutex);
    inner_product += temp_inner_product;
    pthread_mutex_unlock(&mutex);

    pthread_exit(NULL);
}</pre>
```

The original code add the result of multiplication of each element directly to global variable "inner\_product". I modify the code to add the result of multiplication of each element to local variable "temp\_inner\_product". After finishing the loop, I lock the mutex and add "temp\_inner\_product" to "inner\_product" then unlock the mutex. I also add some code to initialize and destroy the mutex as following: In global domain:

```
13 pthread_mutex_t mutex;
14
```

In main function:

```
// Initialize the mutex
45 pthread_mutex_init(&mutex, NULL);

75
76 pthread_mutex_destroy(&mutex);
```

Q2: After implementing modifications to the sample code in Part II, consider whether the multi-threaded version will outperform the serial version provided below. What factors contribute to the difference in performance?

A2: The multi-threaded version will outperform the serial version. In this case, NUM\_THREADS = 8, so the multi-threaded version will split the two vectors into 8 parts and calculate the inner product of each part respectively. The multi-threaded version will add the results of each part to final result only after finishing the loop and get the mutex lock. So the multi-threaded version will be 8 times faster than the serial version ideally (ignore overhead).