# MORE SCHEME

## COMPUTER SCIENCE MENTORS 61A

April 2 to April 4, 2018

## 1 Scheme

1. What will Scheme output? Draw box-and-pointer diagrams to help determine this.

(a) (cons (cons 1 nil) (cons 2 (cons (cons 3 (cons 4 5)) (cons 6 nil))))

```
Solution:
((1) 2 (3 4 . 5) 6)
```

(b) (cons (cons (car '(1 2 3)) (list 2 3 4)) (cons 2 3))

```
Solution: ((1 2 3 4) 2 . 3)
```

(c) (define a 4) ((lambda (x y) (+ a)) 1 2)

```
Solution:
```

(d) ((lambda (x y z) (y x)) 2 / 2)

```
Solution:
0.5
```

(e) ((lambda (x) (x x)) (lambda (y) 4))

```
Solution: 4
```

(f) (**define** boom1 (/ 1 0))

```
Solution: Error: Zero Division
```

(g) boom1

Solution: Error: boom1 not defined

(h) (define boom2 (lambda () (/ 1 0)))

Solution: boom2

(i) (boom2)

**Solution:** Error: Zero Division

(j) How can we rewrite boom 2 without using the lambda operator?

Solution:
(define (boom2) (/ 1 0))

	2.	What	will	Scheme	output?
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```
(a) (if 0 (/ 1 0) 1)
```

#### **Solution:**

Error: Zero Division

(b) (and 1 #f (/ 1 0))

#### **Solution:**

#f

(c) (and 1 2 3)

#### **Solution:**

3

(d) (or #f #f 0 #f (/ 1 0))

#### **Solution:**

0

(e) (or #f #f (/ 1 0) 3 4)

#### **Solution:**

Error: Zero Division

(f) (and (and) (or))

#### **Solution:**

#f

3. **let** is a special form in Scheme which allows you to create local bindings. Consider the example

```
(let ((x 1)) (+ x 1))
```

Here, we assign x to 1, and then evaluate the expression (x + 1) using that binding, returning 2. However, outside of this expression, x would not be bound to anything.

Each let special form has a corresponding lambda equivalent. The equivalent lambda expression for the above example is

```
((lambda (x) (+ x 1)) 1)
```

The following line of code does not work. Why? Write the lambda equivalent of the let expressions.

```
(let ((foo 3)
(bar (+ foo 2)))
(+ foo bar))
```

**Solution:** The above function will error because it is equivalent to:

```
((lambda (foo bar) (+ foo bar)) 3 (+ foo 2))
```

In other words, foo has not been defined in the global frame. When bar is being assigned to (+ foo 2), it will error. The assignment of foo to 3 happens in the lambda?s frame when it's called, not the global frame (this is relevant to the Scheme project – when the interpreter sees lambda, it will call a function to start a new frame).

If we had the line (define foo 3) before the call to let, then it would return 8, because within let, foo would be 3 and bar would be (+ 3 2), since it would use the foo in the Global frame.

1. What is the difference between dynamic and lexical scoping?

#### **Solution:**

- **Lexical:** The parent of a frame is the frame in which a procedure was defined (used in Python).
- **Dynamic:** The parent of a frame is the frame in which a procedure is called (keep an eye out for this in the Scheme project).
- 2. What would this print using lexical scoping? What would it print using dynamic scoping?

```
a = 2
def foo():
    a = 10
    return lambda x: x + a
bar = foo()
bar(10)
```

#### **Solution:**

Lexical: 20Dynamic: 12

3. How would you modify an environment diagram to represent dynamic scoping?

**Solution:** Assign parents when you create a frame (do not set parents when defining functions!). The parent in this case is the frame in which you called this function.

1. Implement waldo. waldo returns #t if the symbol waldo is in a list. You may assume that the list passed in is well-formed.

```
scm> (waldo '(1 4 waldo))
#t
scm> (waldo '())
#f
scm> (waldo '(1 4 9))
#f

(define (waldo lst)
```

)

2. **Extra challenge:** Define waldo so that it returns the index of the list where the symbol waldo was found (if waldo is not in the list, return #f).

```
scm> (waldo '(1 4 waldo))
2
scm> (waldo '())
#f
scm> (waldo '(1 4 9))
#f

(define (waldo lst)
```

)

### 4 Challenge Question

3. **(Optional)** The quicksort sorting algorithm is an efficient and commonly used algorithm to order the elements of a list. We choose one element of the list to be the pivot element and partition the remaining elements into two lists: one of elements less than the pivot and one of elements greater than the pivot. We recursively sort the two lists, which gives us a sorted list of all the elements less than the pivot and all the elements greater than the pivot, which we can then combine with the pivot for a completely sorted list.

Implement quicksort in Scheme. Choose the first element of the list as the pivot. You may assume that all elements are distinct. Hint: you may want to use a helper function.

You may additionally want to use the built-in append function, which takes in two lists and returns a new list containing the elements of the first list followed by the elements of the second list. You can also use filter procedure, which takes in a one-argument function and a list and returns a new list containing only the elements of the original list for which the function returns true, although it is not required.

```
scm> (quicksort (list 5 2 4 3 12 7))
(2 3 4 5 7 12)
```

```
Solution:
(define (quicksort lst)
    (define (helper lst pivot less greater )
        (cond
             ((null? lst)
                 (append (qs less) (list pivot) (qs
                    greater)))
             ((> pivot (car lst))
                 (helper (cdr lst) pivot (append (list (car
                    lst)) less) greater))
             ((< pivot (car lst)</pre>
                 (helper (cdr lst) pivot less (append (list
                    (car lst)) greater)))))
    (if (or (null? lst) (null? (cdr lst)))
                         lst
             (helper (cdr lst) (car lst) nil nil)))
Alternate solution using filter:
(define (quicksort lst)
   (if (null? lst)
       nil
       (let ((less (filter (lambda (x) (< x (car lst)))
          lst))
          (greater (filter (lambda (x) (> x (car lst)))
             lst)))
           (append (append (quicksort less) (list (car
             lst))) (quicksort greater)))))
```