# PageRank

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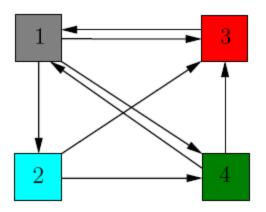
#### Motivation

- Problem: the internet has a lot of information.
- ► A lot of the information out there either isn't relevant to us, or is incorrect.
- ► Motivation: we want a program that, when provided a phrase, returns webpages with information relevant to the input
- Intuitive solution: return back websites that contain that phrase in content of the pages
- ➤ This helps us find more relevant pages, but we can't know of we're getting the best information, much less accurate information, without manually going through each result
- ► Want a stronger solution

### Formalizing the PageRank problem

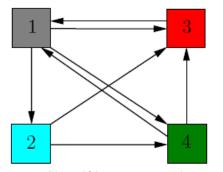
- lacktriangle We're going to construct a directed graph G=(V,E)
- ightharpoonup For each website we consider, we construct a node  $v_i \in V$
- ▶ For two distinct nodes  $v_i$ ,  $v_j \in V$ , the *directed* edge  $v_i v_j \in E$  iff there is a link on website i that goes to website j.
- ▶ If  $v_i$  and  $v_j$  are not distinct (a website is linking to itself), we ignore the link and do not construct a loop edge.
  - ► *G* is not a psuedograph
- ► Multiple hyperlinks on page *i* to page *j* are all represented by the single, directed edge
  - ightharpoonup G is not a multigraph.

# Visual Representation



- ▶ We have a set of 4 websites
- ► Each edge represents a hyperlink from the origin node to the destinationx node

# Adjacency Matrix



$$A = \begin{pmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

No self loops means the main diagonal is all zeros

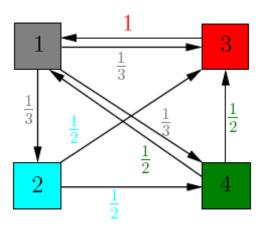
### Applying PageRank Values and Edge Weights

- At first, we assign
- For each *vertex*  $v_i$ , we assign some PageRank value  $PR(v_i)$ 
  - ► TODO: How to calculate?
- ightharpoonup Consider the set of vertices that  $v_i$  has an edge to
  - $V_i = \{v_i | v_i v_i \in E\}$
- ► For each vertex in the set, the weight of the edge to that vertex will be the PageRank value of the source node divided by the number of outbound edges it has

$$\forall v_j \in V_i : w(v_i v_j) = \frac{PR(v_i)}{|V_i|}$$

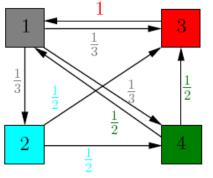
In other words, all the outbound edges from a particular vertex will have the same weight

# Visual Representation



▶ In this case, all nodes have a PageRank value of 1.

#### Transition Matrix



$$T = \begin{pmatrix} 0 & 1/3 & 1/3 & 1/3 \\ 0 & 0 & 1/2 & 1/2 \\ 1 & 0 & 0 & 0 \\ 1/2 & 0 & 1/2 & 0 \end{pmatrix}$$

- ▶ All entries of the transition matrix are non-negative
- If  $v_i$  has at least one outgoing edge, the sum of the entries in row i is 1
  - Else, the sum is 0.
- ► This is a column stochastic matrix