OS编程作业2

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一、算法并行化问题

1、不考虑线程的话,算法的核心代码如下

```
void merge(int left, int right){
 2
        int mid = (left + right) >> 1;
        int size1 = mid - left + 1;
 3
 4
        int size2 = right - mid;
 5
        int t1[size1];
        int t2[size2];
 6
 7
        memcpy(t1, a+left, sizeof(int) * (mid-left+1));
 8
        memcpy(t2, a+mid+1, sizeof(int) * (right-mid));
 9
10
11
        int i = 0, j = 0;
12
        int k = left;
13
        while (i < size1 && j < size2) {
14
             if (t1[i] <= t2[j]) {
15
                 a[k] = t1[i];
16
                 i++;
17
             }
18
             else {
19
                 a[k] = t2[j];
20
                j++;
21
             }
22
            k++;
23
        }
24
25
        while (i < size1) {
26
             a[k] = t1[i];
27
28
            k++;
             i++;
29
30
        }
31
        while (j < size2) {
32
33
            a[k] = t2[j];
34
            j++;
            k++;
35
36
        }
37
    }
```

```
38
39
40
    void* merge sort(void* arg){
41
        int *argu = (int*)arg;
        int left = argu[0];
42
        int right = argu[1];
43
44
45
        int mid = (left + right) >> 1;
        int arg1[2];
46
47
        int arg2[2];
48
49
        arg1[0] = left;
50
        arg1[1] = mid;
51
52
53
        arg2[0] = mid + 1;
        arg2[1] = right;
54
55
56
        if (left >= right) {
57
            return NULL;
58
        }
59
        merge sort(arg1);
        merge_sort(arg2);
60
61
62
        merge(left, right);
        return NULL;
63
64
    }
65
66
```

上述算法经测试, 结果正确。

2、加入线程机制后, merge_sort的实现如下:

```
1
     pthread_t t2;
 2
        pthread_t t1;
 3
 4
        if (numofThread == maxThreadNumber) {
 5
             flag = 1;
 6
        }
 7
        if (numofThread < maxThreadNumber) {</pre>
 8
9
             numofThread += 1;
             pthread_create(&t1, NULL, merge_sort, (void*)arg1);
10
             pthread_join(t1, NULL);
11
             pthread_exit(NULL);
12
             numofThread -= 1;
13
14
        }
15
        else {
```

```
16
             merge sort(arg1);
17
         }
18
19
         if (numofThread < maxThreadNumber) {</pre>
20
             numofThread += 1;
             pthread create(&t2, NULL, merge sort, arg2);
21
             pthread_join(t2, NULL);
22
             pthread_exit(NULL);
2.3
             numofThread -= 1;
2.4
25
        }
        else {
26
27
             merge_sort(arg2);
28
         }
29
30
        merge(left, right);
```

将数据设置为100个,线程的数量设置0,1,10,20,100,结果如下图所示

```
njucs@njucs-VirtualBox:~/pa2$ ./merge
The running time is 335 microsecond
njucs@njucs-VirtualBox:~/pa2$ gcc merge.c -o merge -pthread
njucs@njucs-VirtualBox:~/pa2$ ./merge
The number of thread that I use : 1
The running time is 225 microsecond
njucs@njucs-VirtualBox:~/pa2$ gcc merge.c -o merge -pthread
njucs@njucs-VirtualBox:~/pa2$ ./merge
The number of thread that I use : 10
The running time is 1480 microsecond
njucs@njucs-VirtualBox:~/pa2$ gcc merge.c -o merge -pthread
njucs@njucs-VirtualBox:~/pa2$ ./merge
The running time is 2292 microsecond
njucs@njucs-VirtualBox:~/pa2$ gcc merge.c -o merge -pthread
njucs@njucs-VirtualBox:~/pa2$ ./merge
The running time is 1809 microsecond
```

1000个数据,执行的次数为二叉树非叶子结点的个数,同时运行的线程个数应小于20,所以线程数量超过20时,将达不到最大线程数,线程的数量会导致开销增大,所以单线程的时间消耗较少,多线程消耗较多。

(Ps:由于每次运行结果都不一样, 所以直接随机取值)

二、信号量与PV操作实现同步问题

实现的方式是写者优先,测试用例如下

```
1 R 3 5
2 W 4 5
3 R 5 2
4 R 6 5
5 W 7 3
```

```
Create the 1
             thread: Reader
Create the 2 thread: Writer
Create the 3 thread: Reader
Create the 4 thread: Reader
Create the 5
             thread: Writer
      1:
          waiting to read
Thread
Thread
      1:
          start reading
Thread 2:
          waiting to write
Thread 3:
          waiting to read
          waiting to read
Thread 4:
          waiting to write
Thread 5:
Thread 1:
          end reading
Thread 2:
          start writing
          end writing
Thread 2:
Thread
      5:
          start writing
Thread 5:
          end writing
          start reading
Thread 3:
Thread 4:
          start reading
Thread 3:
          end reading
Thread 4: end reading
```

由上图分析可得,虽然thread3和4先于5进入等待队列中,但是实现的方式是写者优先,5会在thread2写完后再去写,写操作是可以同步的,所以3、4一起运行读过程。

三、管程机制实现与管程应用问题

使用C++语言和pthread库实现一个有界环形缓冲区类(CircleBuffer), 该类提供一个带参(int K)构造方法,用于指定缓冲区的大小, 提供一个get和一个put方法分别用于从缓冲区取出一个元素和向缓冲区存入一个元素。另外需要编写一个多线程测试用例,测试该有界环形缓冲区类是否能够正确同步。具体实现要求:

- 1. 只能使用pthread库提供的一般信号量(semaphore),参考课堂上介绍的使用信号量与PV操作实现Hoare类型管程,来实现CircleBuffer;
- 2. 使用pthread库提供的互斥信号量(mutex)和条件变量(condition),实现CircleBuffer。

没有去测试,代码如下:

```
#include <iostream>
   #include <ctype.h>
   #include <unistd.h>
   #include <stdio.h>
   #include <queue>
   #include <pthread.h>
 7
   #include <stdlib.h>
   #include <string.h>
   #include <sys/time.h>
9
   #include <time.h>
10
11
    #include <boost.h>
12
    class mesa monitor : boost::noncopyable {
13
14
        typedef boost::unique_lock<mesa_monitor> lock_type;
15
        friend class lock type;
16
        mesa_monitor() : m_notify(0) {}
17
    public:
18
        void lock() const {
19
            m mutex.lock();
            m notify = 0; // 进入管程时要把m_notify归0
2.0
21
22
        void unlock() const {
23
            notify impl(m notify);
            m mutex.unlock();
2.4
2.5
        }
2.6
        bool try lock() const {
27
            bool ret = m_mutex.try_lock();
            if (ret) {
2.8
29
                m_notify = 0;
3.0
            }
31
            return ret;
32
33
        void wait() const {
```

```
34
             notify impl(m notify);
35
             m_cond.wait(m_mutex);
36
             m \text{ notify = 0};
37
        }
38
        void notify_one() {
39
             if (m_notify != -1) {
40
                 ++m_notify;
41
             }
42
        }
43
        void notify_all() {
44
45
             m_notify = -1;
46
        }
47
48
    private:
49
        void notify_impl(int nnotify) const {
50
             if (nnotify != 0) {
51
                 if (nnotify = -1) {
52
                     m_cond.notify_all();
53
                     return;
54
                 } else {
55
                     while (nnotify > 0) {
56
                          m_cond.notify_one();
57
                          --nnotify;
58
                     }
59
                 }
60
             }
61
        }
62
    private:
63
        mutable boost::condition variable any m cond;
        mutable boost::mutex m_mutex;
64
65
        mutable int m_notify;
66
    };
67
68
    template <typename T>
69
    class threadsafe_queue : mesa_monitor {
70
        std::queue<T> m_data;
71
    public:
72
        threadsafe_queue() {}
73
        void pop(T& val) {
74
             mesa_monitor::lock_type lk(*this);
75
             while (m_data.empty()) {
                 wait();
76
77
             }
78
             val = m data.front();
             m data.pop();
79
80
        }
81
        void push(const T& val) {
             mesa_monitor::lock_type lk(*this);
```

四、死锁问题

两个进程分别是打印奇数和偶数,当A线程进入锁定状态是,主线程突然异常将A线程停止,这时将导致B线程也无法继续执行,处于死锁状态,如下图所示:

```
even: 10920
even: 10924
even: 10926
even: 10928
even: 10930
odd: 20795
odd: 20797
odd: 20799
odd: 20801
odd: 20803
odd: 20805
stop todd from the outside
```

贴一段进程创建的代码实现

```
1
     pthread_t todd, teven;
 2
            pthread_mutex_init(&m,0);
 3
            pthread create(&todd,0,runodd,0);
            pthread create(&teven,0,runeven,0);
 4
 5
            sleep(5);
            printf("stop todd from the outside\n");
 6
 7
            pthread_cancel(todd);
 8
            pthread_join(todd,(void**)0);
            pthread join(teven,(void**)0);
 9
10
            pthread_mutex_destroy(&m);
```