The LATEX3 Interfaces

The LATEX Project*
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Abstract

This is the reference documentation for the <code>expl3</code> programming environment. The <code>expl3</code> modules set up an experimental naming scheme for LaTeX commands, which allow the LaTeX programmer to systematically name functions and variables, and specify the argument types of functions.

The TEX and ε -TEX primitives are all given a new name according to these conventions. However, in the main direct use of the primitives is not required or encouraged: the expl3 modules define an independent low-level LATEX3 programming language.

The expl3 modules are designed to be loaded on top of LATEX 2ε . With an up-to-date LATEX 2ε kernel, this material is loaded as part of the format. The fundamental programming code can also be loaded with other TeX formats, subject to restrictions on the full range of functionality.

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Part I Introduction

Chapter 1

Introduction to expl3 and this document

This document is intended to act as a comprehensive reference manual for the expl3 language. A general guide to the LATEX3 programming language is found in expl3.pdf.

1.1 Naming functions and variables

IATEX3 does not use **@** as a "letter" for defining internal macros. Instead, the symbols _ and : are used in internal macro names to provide structure. The name of each function is divided into logical units using _, while : separates the name of the function from the argument specifier ("arg-spec"). This describes the arguments expected by the function. In most cases, each argument is represented by a single letter. The complete list of arg-spec letters for a function is referred to as the signature of the function.

Each function name starts with the *module* to which it belongs. Thus apart from a small number of very basic functions, all expl3 function names contain at least one underscore to divide the module name from the descriptive name of the function. For example, all functions concerned with comma lists are in module clist and begin \clist_.

Every function must include an argument specifier. For functions which take no arguments, this will be blank and the function name will end:. Most functions take one or more arguments, and use the following argument specifiers:

- N and n These mean no manipulation, of a single token for N and of a set of tokens given in braces for n. Both pass the argument through exactly as given. Usually, if you use a single token for an n argument, all will be well.
- c This means <code>csname</code>, and indicates that the argument will be turned into a csname before being used. So \foo:c {ArgumentOne} will act in the same way as \foo:N \ArgumentOne. All macros that appear in the argument are expanded. An internal error will occur if the result of expansion inside a c-type argument is not a series of character tokens.
- V and v These mean value of variable. The V and v specifiers are used to get the content of a variable without needing to worry about the underlying TEX structure containing the data. A V argument will be a single token (similar to N), for example

\foo:V \MyVariable; on the other hand, using v a csname is constructed first, and then the value is recovered, for example \foo:v {MyVariable}.

- o This means *expansion once*. In general, the V and v specifiers are favoured over o for recovering stored information. However, o is useful for correctly processing information with delimited arguments.
- **x** The **x** specifier stands for *exhaustive expansion*: every token in the argument is fully expanded until only unexpandable ones remain. The TEX \edef primitive carries out this type of expansion. Functions which feature an **x**-type argument are *not* expandable.
- e The e specifier is in many respects identical to x, but with a very different implementation. Functions which feature an e-type argument may be expandable. The drawback is that e is extremely slow (often more than 200 times slower) in older engines, more precisely in non-LuaTeX engines older than 2019.
- f The f specifier stands for full expansion, and in contrast to x stops at the first non-expandable token (reading the argument from left to right) without trying to expand it. If this token is a \(space token \), it is gobbled, and thus won't be part of the resulting argument. For example, when setting a token list variable (a macro used for storage), the sequence

```
\tl_set:Nn \l_mya_tl { A }
\tl_set:Nn \l_myb_tl { B }
\tl_set:Nf \l_mya_tl { \l_mya_tl \l_myb_tl }
```

will leave \l_mya_tl with the content A\l_myb_tl, as A cannot be expanded and so terminates expansion before \l_myb_tl is considered.

- T and F For logic tests, there are the branch specifiers T (true) and F (false). Both specifiers treat the input in the same way as n (no change), but make the logic much easier to see.
- ${\tt p}$ The letter ${\tt p}$ indicates TEX parameters. Normally this will be used for delimited functions as ${\sf expl3}$ provides better methods for creating simple sequential arguments.
- w Finally, there is the w specifier for weird arguments. This covers everything else, but mainly applies to delimited values (where the argument must be terminated by some specified string).
- D The D stands for **Do not use**. All of the T_EX primitives are initially \let to a D name, and some are then given a second name. These functions have no standardized syntax, they are engine dependent and their name can change without warning, thus their use is *strongly discouraged* in package code: programmers should instead use the interfaces documented in interface3.pdf¹.

Notice that the argument specifier describes how the argument is processed prior to being passed to the underlying function. For example, \foo:c will take its argument, convert it to a control sequence and pass it to \foo:N.

¹If a primitive offers a functionality not yet in the kernel, programmers and users are encouraged to write to the LaTeX-L mailing list (mailto:LATEX-L@listserv.uni-heidelberg.de) describing their use-case and intended behaviour, so that a possible interface can be discussed. Temporarily, while an interface is not provided, programmers may use the procedure described in the l3styleguide.pdf.

Variables are named in a similar manner to functions, but begin with a single letter to define the type of variable:

- c Constant: global parameters whose value should not be changed.
- g Parameters whose value should only be set globally.
- 1 Parameters whose value should only be set locally.

Each variable name is then build up in a similar way to that of a function, typically starting with the module² name and then a descriptive part. Variables end with a short identifier to show the variable type:

clist Comma separated list.

dim "Rigid" lengths.

fp Floating-point values;

int Integer-valued count register.

muskip "Rubber" lengths for use in mathematics.

seq "Sequence": a data-type used to implement lists (with access at both ends) and stacks.

skip "Rubber" lengths.

str String variables: contain character data.

tl Token list variables: placeholder for a token list.

Applying V-type or v-type expansion to variables of one of the above types is supported, while it is not supported for the following variable types:

bool Either true or false.

box Box register.

coffin A "box with handles" — a higher-level data type for carrying out box alignment operations.

flag Integer that can be incremented expandably.

fparray Fixed-size array of floating point values.

intarray Fixed-size array of integers.

ior/iow An input or output stream, for reading from or writing to, respectively.

prop Property list: analogue of dictionary or associative arrays in other languages.

regex Regular expression.

²The module names are not used in case of generic scratch registers defined in the data type modules, e.g., the int module contains some scratch variables called $\l_{tmpa_int}, \l_{tmpb_int}$, and so on. In such a case adding the module name up front to denote the module and in the back to indicate the type, as in $\l_{int_tmpa_int}$ would be very unreadable.

1.1.1 Scratch variables

Modules focussed on variable usage typically provide four scratch variables, two local and two global, with names of the form $\c cope \t mpa_{type}/\c scope \t mpb_{type}$. These are never used by the core code. The nature of TeX grouping means that as with any other scratch variable, these should only be set and used with no intervening third-party code.

1.1.2 Terminological inexactitude

A word of warning. In this document, and others referring to the expl3 programming modules, we often refer to "variables" and "functions" as if they were actual constructs from a real programming language. In truth, TEX is a macro processor, and functions are simply macros that may or may not take arguments and expand to their replacement text. Many of the common variables are also macros, and if placed into the input stream will simply expand to their definition as well — a "function" with no arguments and a "token list variable" are almost the same. On the other hand, some "variables" are actually registers that must be initialised and their values set and retrieved with specific functions.

The conventions of the expl3 code are designed to clearly separate the ideas of "macros that contain data" and "macros that contain code", and a consistent wrapper is applied to all forms of "data" whether they be macros or actually registers. This means that sometimes we will use phrases like "the function returns a value", when actually we just mean "the macro expands to something". Similarly, the term "execute" might be used in place of "expand" or it might refer to the more specific case of "processing in TeX's stomach" (if you are familiar with the TeXbook parlance).

If in doubt, please ask; chances are we've been hasty in writing certain definitions and need to be told to tighten up our terminology.

1.2 Documentation conventions

This document is typeset with the experimental l3doc class; several conventions are used to help describe the features of the code. A number of conventions are used here to make the documentation clearer.

Each group of related functions is given in a box. For a function with a "user" name, this might read:

\ExplSyntaxOn \ExplSyntaxOff

The textual description of how the function works would appear here. The syntax of the function is shown in mono-spaced text to the right of the box. In this example, the function takes no arguments and so the name of the function is simply reprinted.

For programming functions, which use _ and : in their name there are a few additional conventions: If two related functions are given with identical names but different argument specifiers, these are termed *variants* of each other, and the latter functions are printed in grey to show this more clearly. They will carry out the same function but will take different types of argument:

 $^{^3\}mbox{TeXnically},$ functions with no arguments are $\mbox{\sc long}$ while token list variables are not.

\seq_new:N \seq_new:N \sequence

\seq_new:c

When a number of variants are described, the arguments are usually illustrated only for the base function. Here, (sequence) indicates that \seq new: N expects the name of a sequence. From the argument specifier, \seq_new:c also expects a sequence name, but as a name rather than as a control sequence. Each argument given in the illustration should be described in the following text.

Fully expandable functions Some functions are fully expandable, which allows them to be used within an x-type or e-type argument (in plain T_FX terms, inside an \edef or \expanded), as well as within an f-type argument. These fully expandable functions are indicated in the documentation by a star:

 $\cs_{to_str:N} \star \cs_{to_str:N} \langle cs \rangle$

As with other functions, some text should follow which explains how the function works. Usually, only the star will indicate that the function is expandable. In this case, the function expects a $\langle cs \rangle$, shorthand for a $\langle control \ sequence \rangle$.

Restricted expandable functions A few functions are fully expandable but cannot be fully expanded within an f-type argument. In this case a hollow star is used to indicate this:

\seq map function:NN 🌣

\seq_map_function:NN \langle seq \langle function \rangle

Conditional functions Conditional (if) functions are normally defined in three variants, with T, F and TF argument specifiers. This allows them to be used for different "true"/"false" branches, depending on which outcome the conditional is being used to test. To indicate this without repetition, this information is given in a shortened form:

\sys_if_engine_xetex: TF * \sys_if_engine_xetex: TF {\langle true code \rangle} {\langle false code \rangle}

The underlining and italic of TF indicates that three functions are available:

- \sys_if_engine_xetex:T
- \sys_if_engine_xetex:F
- \sys_if_engine_xetex:TF

Usually, the illustration will use the TF variant, and so both $\langle true\ code \rangle$ and $\langle false\ code \rangle$ will be shown. The two variant forms T and F take only \(\lambda true \code \rangle \) and \(\lambda \lambda lse \code \rangle \rangle \). respectively. Here, the star also shows that this function is expandable. With some minor exceptions, all conditional functions in the expl3 modules should be defined in this way.

Variables, constants and so on are described in a similar manner:

\l_tmpa_tl A short piece of text will describe the variable: there is no syntax illustration in this case.

In some cases, the function is similar to one in $IAT_{FX} 2_{\varepsilon}$ or plain T_{FX} . In these cases, the text will include an extra "TeXhackers note" section:

 $\verb|\token_to_str:N| \star \verb|\token_to_str:N| \langle token| \rangle$

The normal description text.

TEXhackers note: Detail for the experienced TEX or \LaTeX 2 ε programmer. In this case, it would point out that this function is the TEX primitive \string.

Changes to behaviour When new functions are added to expl3, the date of first inclusion is given in the documentation. Where the documented behaviour of a function changes after it is first introduced, the date of the update will also be given. This means that the programmer can be sure that any release of expl3 after the date given will contain the function of interest with expected behaviour as described. Note that changes to code internals, including bug fixes, are not recorded in this way *unless* they impact on the expected behaviour.

1.3 Formal language conventions which apply generally

As this is a formal reference guide for IATEX3 programming, the descriptions of functions are intended to be reasonably "complete". However, there is also a need to avoid repetition. Formal ideas which apply to general classes of function are therefore summarised here.

For tests which have a TF argument specification, the test if evaluated to give a logically TRUE or FALSE result. Depending on this result, either the $\langle true\ code \rangle$ or the $\langle false\ code \rangle$ will be left in the input stream. In the case where the test is expandable, and a predicate (_p) variant is available, the logical value determined by the test is left in the input stream: this will typically be part of a larger logical construct.

1.4 TeX concepts not supported by LaTeX3

The TEX concept of an "\outer" macro is not supported at all by LATEX3. As such, the functions provided here may break when used on top of LATEX 2_{ε} if \outer tokens are used in the arguments.

 $\begin{array}{c} {\rm Part~II} \\ {\bf Bootstrapping} \end{array}$

Chapter 2

The **I3bootstrap** package Bootstrap code

2.1Using the LATEX3 modules

The modules documented in source3 are designed to be used on top of \LaTeX and are loaded all as one with the usual \usepackage{expl3} or \RequirePackage{expl3} instructions.

As the modules use a coding syntax different from standard LATEX 2_{ε} it provides a few functions for setting it up.

\ExplSyntaxOn \ExplSyntaxOff \ExplSyntaxOn \(code \) \ExplSyntaxOff

The \ExplSyntaxOn function switches to a category code regime in which spaces and Updated: 2011-08-13 new lines are ignored, and in which the colon (:) and underscore (_) are treated as "letters", thus allowing access to the names of code functions and variables. Within this environment, ~ is used to input a space. The \ExplSyntaxOff reverts to the document category code regime.

> TeXhackers note: Spaces introduced by ~ behave much in the same way as normal space characters in the standard category code regime: they are ignored after a control word or at the start of a line, and multiple consecutive ~ are equivalent to a single one. However, ~ is not ignored at the end of a line.

\ProvidesExplClass \ProvidesExplFile

\ProvidesExplPackage \RequirePackage{expl3}

These functions act broadly in the same way as the corresponding \LaTeX Z_{ε} kernel func-Updated: 2017-03-19 tions \ProvidesPackage, \ProvidesClass and \ProvidesFile. However, they also implicitly switch \ExplSyntaxOn for the remainder of the code with the file. At the end of the file, \ExplSyntaxOff will be called to reverse this. (This is the same concept as $\text{ETEX } 2_{\varepsilon}$ provides in turning on \makeatletter within package and class code.) The $\langle date \rangle$ should be given in the format $\langle year \rangle / \langle month \rangle / \langle day \rangle$ or in the ISO date format $\langle year \rangle - \langle month \rangle - \langle day \rangle$. If the $\langle version \rangle$ is given then it will be prefixed with v in the package identifier line.

\GetIdInfo

\RequirePackage{13bootstrap}

Updated: 2012-06-04

\GetIdInfo \\$Id: $\langle SVN \ info \ field \rangle$ \\$ \{\description\}\}

Extracts all information from a SVN field. Spaces are not ignored in these fields. The information pieces are stored in separate control sequences with \ExplFileName for the part of the file name leading up to the period, \ExplFileDate for date, \ExplFileVersion for version and \ExplFileDescription for the description.

To summarize: Every single package using this syntax should identify itself using one of the above methods. Special care is taken so that every package or class file loaded with \RequirePackage or similar are loaded with usual LATEX 2_{ε} category codes and the LATEX3 category code scheme is reloaded when needed afterwards. See implementation for details. If you use the \GetIdInfo command you can use the information when loading a package with

\ProvidesExplPackage{\ExplFileName}
{\ExplFileDate}{\ExplFileVersion}{\ExplFileDescription}

Chapter 3

The l3names package Namespace for primitives

3.1 Setting up the LATEX3 programming language

This module is at the core of the LATEX3 programming language. It performs the following tasks:

- defines new names for all TEX primitives;
- emulate required primitives not provided by default in LuaT_EX;
- switches to the category code régime for programming;

This module is entirely dedicated to primitives (and emulations of these), which should not be used directly within LATEX3 code (outside of "kernel-level" code). As such, the primitives are not documented here: The TeXbook, TeX by Topic and the manuals for pdfTeX, XaTeX, LuaTeX, pTeX and upTeX should be consulted for details of the primitives. These are named tex_{name} :D, typically based on the primitive's name in pdfTeX and omitting a leading pdf when the primitive is not related to pdf output.

Part III
Programming Flow

Chapter 4

The I3basics package Basic definitions

As the name suggest this package holds some basic definitions which are needed by most or all other packages in this set.

Here we describe those functions that are used all over the place. With that we mean functions dealing with the construction and testing of control sequences. Furthermore the basic parts of conditional processing are covered; conditional processing dealing with specific data types is described in the modules specific for the respective data types.

4.1 No operation functions

\prg_do_nothing: * \prg_do_nothing:

An expandable function which does nothing at all: leaves nothing in the input stream after a single expansion.

\scan_stop: \scan_stop:

A non-expandable function which does nothing. Does not vanish on expansion but produces no typeset output.

4.2 Grouping material

\group_begin: \group_begin:

\group_end:

\group_end:

These functions begin and end a group for definition purposes. Assignments are local to groups unless carried out in a global manner. (A small number of exceptions to this rule will be noted as necessary elsewhere in this document.) Each \group_begin: must be matched by a \group_end:, although this does not have to occur within the same function. Indeed, it is often necessary to start a group within one function and finish it within another, for example when seeking to use non-standard category codes.

 $\verb|\group_insert_after:N \qroup_insert_after:N \qtoken||$

Adds $\langle token \rangle$ to the list of $\langle tokens \rangle$ to be inserted when the current group level ends. The list of $\langle tokens \rangle$ to be inserted is empty at the beginning of a group: multiple applications of \group_insert_after:N may be used to build the inserted list one $\langle token \rangle$ at a time. The current group level may be closed by a \group_end: function or by a token with category code 2 (close-group), namely a } if standard category codes apply.

\group_show_list: \group_show_list: \group_log_list:

New: 2021-05-11 Display (to the terminal or log file) a list of the groups that are currently opened. This is intended for tracking down problems.

TrXhackers note: This is a wrapper around the \showgroups primitive.

4.3 Control sequences and functions

As T_EX is a macro language, creating new functions means creating macros. At point of use, a function is replaced by the replacement text ("code") in which each parameter in the code (#1, #2, etc.) is replaced the appropriate arguments absorbed by the function. In the following, $\langle code \rangle$ is therefore used as a shorthand for "replacement text".

Functions which are not "protected" are fully expanded inside an x expansion. In contrast, "protected" functions are not expanded within x expansions.

4.3.1 Defining functions

Functions can be created with no requirement that they are declared first (in contrast to variables, which must always be declared). Declaring a function before setting up the code means that the name chosen is checked and an error raised if it is already in use. The name of a function can be checked at the point of definition using the \cs_new... functions: this is recommended for all functions which are defined for the first time.

There are three ways to define new functions. All classes define a function to expand to the substitution text. Within the substitution text the actual parameters are substituted for the formal parameters ($\#1, \#2, \ldots$).

- new Create a new function with the new scope, such as \cs_new:Npn. The definition is global and results in an error if it is already defined.
- set Create a new function with the set scope, such as \cs_set:Npn. The definition is restricted to the current TeX group and does not result in an error if the function is already defined.
- gset Create a new function with the gset scope, such as \cs_gset:Npn. The definition is global and does not result in an error if the function is already defined.

Within each set of scope there are different ways to define a function. The differences depend on restrictions on the actual parameters and the expandability of the resulting function.

nopar Create a new function with the nopar restriction, such as \cs_set_nopar:Npn.

The parameter may not contain \par tokens.

protected Create a new function with the protected restriction, such as \cs_set_protected: Npn. The parameter may contain \par tokens but the function will not expand within an x-type or e-type expansion.

Finally, the functions in Subsections 4.3.2 and 4.3.3 are primarily meant to define base functions only. Base functions can only have the following argument specifiers:

N and n No manipulation.

T and F Functionally equivalent to n (you are actually encouraged to use the family of \prg_new_conditional: functions described in Section 9.1).

p and w These are special cases.

The \cs new: functions below (and friends) do not stop you from using other argument specifiers in your function names, but they do not handle expansion for you. You should define the base function and then use \cs_generate_variant: Nn to generate custom variants as described in Section 5.2.

4.3.2Defining new functions using parameter text

 $\verb|\cs_new:Npn \cs_new:Npn \defunction| \de$

\cs_new:cpn

\cs_new:Npx

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_nopar:cpn

\cs_new_nopar:Npx \cs_new_nopar:cpx

\cs_new_nopar:Npn \cs_new_nopar:Npn \(\lambda function\) \(\lambda parameters\) \{\(\lambda code\)\}

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_protected:cpn \cs_new_protected:Npx

\cs_new_protected:cpx

\cs_new_protected:Npn \cs_new_protected:Npn \function \ \(\range parameters \) \ \{\(\code \) \}

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type or or e-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_protected_nopar:Npn \cs_new_protected_nopar:Npn \(\lambda function\) \(\lambda parameters\) \{\(\lambda code\)\} \cs_new_protected_nopar:cpn

\cs_new_protected_nopar:Npx

\cs_new_protected_nopar:cpx

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.

 $\cs_{set:Npn \cs_{set:Npn \defunction} \defunction} \defunction \defunction$

\cs_set:cpn Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the \cs_set:cpx \langle parameters \rangle (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current T_FX group level.

\cs_set_nopar:Npn \cs_set_nopar:Npn \(function \) \(\rangle parameters \) \{ \((code \) \)}

\cs_set_nopar:Npx

\cs_set_nopar:cpx

\cs_set_nopar:cpn Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current T_EX group level.

\cs_set_protected:cpn \cs_set_protected:Npx

\cs_set_protected:cpx

\cs_set_protected:Npn \cs_set_protected:Npn \(\frac{function}{\range}\) \(\frac{code}{\range}\)

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current TeX group level. The $\langle function \rangle$ will not expand within an x-type or e-type argument.

\cs_set_protected_nopar:Npn \cs_set_protected_nopar:Npn \(\lambda function\) \(\lambda parameters\) \{\(\lambda code\)\}

\cs_set_protected_nopar:cpn

\cs_set_protected_nopar:Npx

\cs_set_protected_nopar:cpx

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current T_FX group level. The (function) will not expand within an x-type or e-type argument.

 $\verb|\cs_gset:Npn \cs_gset:Npn \| \langle function \rangle \| \langle parameters \rangle \| \{\langle code \rangle\}|$

\cs_gset:cpn Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, \cs_gset:cpx the \(\rho parameters\) (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is not restricted to the current TEX group level: the assignment is global.

 $\cs_gset_nopar:Npn \cs_gset_nopar:Npn \defunction \end{argmath} \ \defunction \end{argmath} \ \defunction \end{argmath} \ \defunction \end{argmath} \defunction \end{argmath$

\cs_gset_nopar:cpn \cs_gset_nopar:Npx \cs_gset_nopar:cpx

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The assignment of a meaning to the $\langle function \rangle$ is not restricted to the current TEX group level: the assignment is global.

\cs_gset_protected:cpn

\cs_gset_protected:Npx

\cs_gset_protected:cpx

\cs_gset_protected:Npn \cs_gset_protected:Npn \function\ \((code \) \}

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the (parameters) (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is not restricted to the current T_FX group level: the assignment is global. The $\langle function \rangle$ will not expand within an x-type or e-type argument.

```
\cs_gset_protected_nopar:Npn \cs_gset_protected_nopar:Npn \function \ (parameters) \ \{\code\}\\cs_gset_protected_nopar:Cpn \\cs_gset_protected_nopar:Npx \\cs_gset_protected_nopar:Cpx
```

Globally sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The assignment of a meaning to the $\langle function \rangle$ is not restricted to the current TEX group level: the assignment is global. The $\langle function \rangle$ will not expand within an x-type or e-type argument.

4.3.3 Defining new functions using the signature

 $\cs_new:Nn \\ \cs_new:(cn|Nx|cx)$

 $\cs_new:Nn \ \langle function \rangle \ \{\langle code \rangle\}$

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_nopar:Nn \cs_new_nopar:(cn|Nx|cx) $\cs_new_nopar:Nn \langle function \rangle \{\langle code \rangle\}$

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain \par tokens. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_protected:Nn \cs_new_protected:(cn|Nx|cx) $\cs_new_protected:Nn \langle function \rangle \{\langle code \rangle\}$

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_new_protected_nopar:Nn \cs_new_protected_nopar:Nn $\langle function \rangle \{\langle code \rangle\}$ \cs_new_protected_nopar:(cn|Nx|cx)

Creates $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain $\langle par$ tokens. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The definition is global and an error results if the $\langle function \rangle$ is already defined.

\cs_set:Nn
\cs_set:(cn|Nx|cx)

 $\cs_set:Nn \langle function \rangle \{\langle code \rangle\}$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current TEX group level.

\cs_set_nopar:Nn
\cs_set_nopar:(cn|Nx|cx)

 $\cs_set_nopar:Nn \langle function \rangle \{\langle code \rangle\}$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain $\langle parameters \rangle$ are assignment of a meaning to the $\langle function \rangle$ is restricted to the current $\langle function \rangle$ is restri

\cs_set_protected:Nn \cs_set_protected:(cn|Nx|cx)

 $\verb|\cs_set_protected:Nn| \langle function \rangle | \{\langle code \rangle\}|$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current TFX group level.

\cs_set_protected_nopar:Nn
\cs_set_protected_nopar:(cn|Nx|cx)

 $\verb|\cs_set_protected_nopar:Nn| \langle \textit{function} \rangle | \{ \langle \textit{code} \rangle \}|$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain $\langle parameters \rangle$ absorbed cannot contain $\langle parameters \rangle$ will not expand within an x-type or e-type argument. The assignment of a meaning to the $\langle function \rangle$ is restricted to the current TeX group level.

\cs_gset:Nn
\cs_gset:(cn|Nx|cx)

 $\cs_gset:Nn \langle function \rangle \{\langle code \rangle\}$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The assignment of a meaning to the $\langle function \rangle$ is global.

\cs_gset_nopar:Nn
\cs_gset_nopar:(cn|Nx|cx)

\cs_gset_nopar:Nn \(\langle function \rangle \{\langle code \rangle \}\)

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle function \rangle$ absorbed cannot contain $\langle function \rangle$ are assignment of a meaning to the $\langle function \rangle$ is global.

\cs_gset_protected:Nn
\cs_gset_protected:(cn|Nx|cx)

 $\cs_gset_protected:Nn \langle function \rangle \{\langle code \rangle\}$

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. The $\langle function \rangle$ will not expand within an x-type or e-type argument. The assignment of a meaning to the

 $\langle function \rangle$ is global.

Sets $\langle function \rangle$ to expand to $\langle code \rangle$ as replacement text. Within the $\langle code \rangle$, the number of $\langle parameters \rangle$ is detected automatically from the function signature. These $\langle parameters \rangle$ (#1, #2, etc.) will be replaced by those absorbed by the function. When the $\langle function \rangle$ is used the $\langle parameters \rangle$ absorbed cannot contain $\langle parameters \rangle$ absorbed cannot contain $\langle parameters \rangle$ will not expand within an x-type or e-type argument. The assignment of a meaning to the $\langle function \rangle$ is global.

Uses the $\langle creator \rangle$ function (which should have signature Npn, for example \cs_new:Npn) to define a $\langle function \rangle$ which takes $\langle number \rangle$ arguments and has $\langle code \rangle$ as replacement text. The $\langle number \rangle$ of arguments is an integer expression, evaluated as detailed for \int_eval:n.

4.3.4 Copying control sequences

Control sequences (not just functions as defined above) can be set to have the same meaning using the functions described here. Making two control sequences equivalent means that the second control sequence is a *copy* of the first (rather than a pointer to it). Thus the old and new control sequence are not tied together: changes to one are not reflected in the other.

In the following text "cs" is used as an abbreviation for "control sequence".

```
\label{eq:new_eq:NN} $$ \cs_new_eq:NN $$ \cs_new_eq:NN
```

Globally creates $\langle control\ sequence_1 \rangle$ and sets it to have the same meaning as $\langle control\ sequence_2 \rangle$ or $\langle token \rangle$. The second control sequence may subsequently be altered without affecting the copy.

Sets $\langle control\ sequence_1 \rangle$ to have the same meaning as $\langle control\ sequence_2 \rangle$ (or $\langle token \rangle$). The second control sequence may subsequently be altered without affecting the copy. The assignment of a meaning to the $\langle control\ sequence_1 \rangle$ is restricted to the current TeX group level.

```
\label{eq:nn} $$ \cs_gset_eq:NN $$ \cs_gset_eq
```

Globally sets $\langle control\ sequence_1 \rangle$ to have the same meaning as $\langle control\ sequence_2 \rangle$ (or $\langle token \rangle$). The second control sequence may subsequently be altered without affecting the copy. The assignment of a meaning to the $\langle control\ sequence_1 \rangle$ is not restricted to the current T_FX group level: the assignment is global.

Deleting control sequences

There are occasions where control sequences need to be deleted. This is handled in a very simple manner.

 $\verb|\cs_undefine:N | cs_undefine:N | \langle control | sequence \rangle|$ \cs_undefine:c Sets $\langle control \ sequence \rangle$ to be globally undefined. Updated: 2011-09-15

Showing control sequences 4.3.6

\cs_meaning:N * \cs_meaning:N \(\)control sequence \(\)

 $\underline{\hspace{0.2cm}}$ This function expands to the *meaning* of the $\langle control\ sequence \rangle$ control sequence. For a Updated: 2011-12-22 macro, this includes the $\langle replacement \ text \rangle$.

> TEXhackers note: This is TEX's \meaning primitive. For tokens that are not control sequences, it is more logical to use \token_to_meaning:N. The c variant correctly reports undefined arguments.

\cs_show:N \(control sequence\) \cs_show:N \cs_show:c

Undated: 2017-02-14

Displays the definition of the $\langle control \ sequence \rangle$ on the terminal.

TEX hackers note: This is similar to the TEX primitive \show, wrapped to a fixed number of characters per line.

\cs_log:N \(control \) sequence \(\) \cs_log:N

\cs_log:c

Writes the definition of the (control sequence) in the log file. See also \cs_show:N which New: 2014-08-22 displays the result in the terminal. Updated: 2017-02-14

4.3.7Converting to and from control sequences

\use:c * \use:c {\langle control sequence name \rangle}

Expands the (control sequence name) until only characters remain, and then converts this into a control sequence. This process requires two expansions. As in other c-type arguments the (control sequence name) must, when fully expanded, consist of character tokens, typically a mixture of category code 10 (space), 11 (letter) and 12 (other).

As an example of the \use:c function, both

\use:c { a b c } and \tl_new:N \l_my_tl \tl_set:Nn \l_my_tl { a b c } \use:c { \tl_use:N \l_my_tl } would be equivalent to

\abc

after two expansions of \use:c.

```
\cs_if_exist_use:N
\cs_if_exist_use:c
\cs_if_exist_use:NTF \star
\cs_if_exist_use:c<u>TF</u>
```

* \cs_if_exist_use:N \(\)control sequence \(\) $\star \cs_{if_exist_use:NTF} \control sequence \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}$

Tests whether the (control sequence) is currently defined according to the conditional \cs_if_exist:NTF (whether as a function or another control sequence type), and if it is New: 2012-11-10 inserts the $\langle control\ sequence \rangle$ into the input stream followed by the $\langle true\ code \rangle$. Otherwise the $\langle false\ code \rangle$ is used.

```
\cs:w
\cs_end: *
```

* \cs:w \(control \) sequence name \(\) \cs_end:

Converts the given (control sequence name) into a single control sequence token. This process requires one expansion. The content for $\langle control \ sequence \ name \rangle$ may be literal material or from other expandable functions. The $\langle control\ sequence\ name \rangle$ must, when fully expanded, consist of character tokens which are not active: typically of category code 10 (space), 11 (letter) or 12 (other), or a mixture of these.

TEXhackers note: These are the TEX primitives \csname and \endcsname.

As an example of the \cs:w and \cs_end: functions, both

```
\cs:w a b c \cs end:
```

and

```
\tl_new:N \l_my_tl
\tl_set:Nn \l_my_tl { a b c }
\cs:w \tl_use:N \l_my_tl \cs_end:
```

would be equivalent to

\abc

after one expansion of \cs:w.

\cs_to_str:N * \cs_to_str:N \(\)control sequence \(\)

Converts the given $\langle control\ sequence \rangle$ into a series of characters with category code 12 (other), except spaces, of category code 10. The result does not include the current escape token, contrarily to \token_to_str:N. Full expansion of this function requires exactly 2 expansion steps, and so an x-type or e-type expansion, or two o-type expansions are required to convert the $\langle control\ sequence \rangle$ to a sequence of characters in the input stream. In most cases, an f-expansion is correct as well, but this loses a space at the start of the result.

4.4 Analysing control sequences

\cs_split_function:N * \cs_split_function:N \(\) function \(\)

New: 2018-04-06 Splits the $\langle function \rangle$ into the $\langle name \rangle$ (i.e. the part before the colon) and the $\langle signature \rangle$ (i.e. after the colon). This information is then placed in the input stream in three parts: the $\langle name \rangle$, the $\langle signature \rangle$ and a logic token indicating if a colon was found (to differentiate variables from function names). The $\langle name \rangle$ does not include the escape character, and both the $\langle name \rangle$ and $\langle signature \rangle$ are made up of tokens with category code 12 (other).

> The next three functions decompose T_FX macros into their constituent parts: if the $\langle token \rangle$ passed is not a macro then no decomposition can occur. In the latter case, all three functions leave \scan_stop: in the input stream.

\cs_prefix_spec:N * \cs_prefix_spec:N \langle token \rangle

New: 2019-02-27 If the $\langle token \rangle$ is a macro, this function leaves the applicable T_FX prefixes in input stream as a string of tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1~y #2 }
\cs_prefix_spec:N \next:nn
```

leaves \long in the input stream. If the $\langle token \rangle$ is not a macro then \scan_stop: is left in the input stream.

TEXhackers note: The prefix can be empty, \long, \protected or \protected\long with backslash replaced by the current escape character.

\cs_parameter_spec:N * \cs_parameter_spec:N \(\(\text{token}\)\)

New: 2022-06-24 If the $\langle token \rangle$ is a macro, this function leaves the primitive TEX parameter specification in input stream as a string of character tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1 y #2 }
\cs_parameter_spec:N \next:nn
```

leaves #1#2 in the input stream. If the $\langle token \rangle$ is not a macro then \scan_stop: is left in the input stream.

TeXhackers note: If the parameter specification contains the string ->, then the function produces incorrect results.

```
\verb|\cs_replacement_spec:N| \star \verb|\cs_replacement_spec:N| \langle token \rangle|
```

New: 2019-02-27 If the $\langle token \rangle$ is a macro, this function leaves the replacement text in input stream as a string of character tokens of category code 12 (with spaces having category code 10). Thus for example

```
\cs_set:Npn \next:nn #1#2 { x #1~y #2 }
\cs_replacement_spec:N \next:nn
```

leaves $x#1_{\cup}y#2$ in the input stream. If the $\langle token \rangle$ is not a macro then \scan_stop: is left in the input stream.

TeXhackers note: If the parameter specification contains the string ->, then the function produces incorrect results.

4.5 Using or removing tokens and arguments

Tokens in the input can be read and used or read and discarded. If one or more tokens are wrapped in braces then when absorbing them the outer set is removed. At the same time, the category code of each token is set when the token is read by a function (if it is read more than once, the category code is determined by the situation in force when first function absorbs the token).

```
\{\langle group_1 \rangle\}
\use:n
        * \use:n
                  \{\langle \textit{group}_1 \rangle\} \ \{\langle \textit{group}_2 \rangle\}
        * \use:nn
\use:nn
        \use:nnn
```

As illustrated, these functions absorb between one and four arguments, as indicated by the argument specifier. The braces surrounding each argument are removed and the remaining tokens are left in the input stream. The category code of these tokens is also fixed by this process (if it has not already been by some other absorption). All of these functions require only a single expansion to operate, so that one expansion of

```
\use:nn { abc } { { def } }
results in the input stream containing
    abc { def }
```

i.e. only the outer braces are removed.

TEXhackers note: The \use:n function is equivalent to LATEX 2ε 's \Offirstofone.

```
\use_i:nn
\use_ii:nn *
```

These functions absorb two arguments from the input stream. The function \use_i:nn discards the second argument, and leaves the content of the first argument in the input stream. \use_ii:nn discards the first argument and leaves the content of the second argument in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

TEXhackers note: These are equivalent to LATEX 2ε 's \Offirstoftwo and \Osecondoftwo.

```
\use_i:nnn
      \use_ii:nnn
\use_iii:nnn *
```

These functions absorb three arguments from the input stream. The function $\use_i:nnn$ discards the second and third arguments, and leaves the content of the first argument in the input stream. \use ii:nnn and \use iii:nnn work similarly, leaving the content of second or third arguments in the input stream, respectively. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

```
\use_i:nnnn
\use_ii:nnnn
\use_iii:nnnn *
\use_iv:nnnn
```

 $\star \text{ \use_i:nnnn } \{\langle arg_1 \rangle\} \ \{\langle arg_2 \rangle\} \ \{\langle arg_3 \rangle\} \ \{\langle arg_4 \rangle\}$

These functions absorb four arguments from the input stream. The function \use_i:nnnn discards the second, third and fourth arguments, and leaves the content of the first argument in the input stream. \use_ii:nnnn, \use_iii:nnnn and \use_iv:nnnn work similarly, leaving the content of second, third or fourth arguments in the input stream, respectively. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the functions to take effect.

```
\use_i_i:nnn * \use_i_i:nnn {\langle arg_1 \rangle} {\langle arg_2 \rangle} {\langle arg_3 \rangle}
```

This function absorbs three arguments and leaves the content of the first and second in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the function to take effect. An example:

```
\use_i_ii:nnn { abc } { { def } } { ghi }
```

results in the input stream containing

```
abc { def }
```

i.e. the outer braces are removed and the third group is removed.

```
\use_{ii_i:nn} \star \use_{ii_i:nn} \{\langle arg_1 \rangle\} \{\langle arg_2 \rangle\}
```

New: 2019-06-02 This function absorbs two arguments and leaves the content of the second and first in the input stream. The category code of these tokens is also fixed (if it has not already been by some other absorption). A single expansion is needed for the function to take effect.

```
\star \text{ \use_none:n } \{\langle group_1 \rangle\}
\use_none:n
\use_none:nn
                         These functions absorb between one and nine groups from the input stream, leaving
\use_none:nnn
                         nothing on the resulting input stream. These functions work after a single expansion.
\use_none:nnnn
                         One or more of the n arguments may be an unbraced single token (i.e. an N argument).
\use_none:nnnn
\use_none:nnnnn
                              TEXhackers note: These are equivalent to LATEX 2\varepsilon's \@gobble, \@gobbletwo, etc.
\use_none:nnnnnn
\use none:nnnnnnn
\use_none:nnnnnnn
                       * \use:e {\(\langle\) expandable tokens\(\rangle\)}
           \use:e
```

New: 2018-06-18 Fully expands the $\langle token \ list \rangle$ in an x-type manner, but the function remains fully expandable, and parameter character (usually #) need not be doubled.

TEXhackers note: \use:e is a wrapper around the primitive \expanded where it is available: it requires two expansions to complete its action. When \expanded is not available this function is very slow.

```
\use:x \use:x \(\left(\expandable tokens\right)\)\\
\text{Updated: 2011-12-31} \text{Fully expands the \(\left(\expandable tokens\right)\) and inserts the result into the input stream at the current location. Any hash characters (#) in the argument must be doubled.
```

4.5.1 Selecting tokens from delimited arguments

A different kind of function for selecting tokens from the token stream are those that use delimited arguments.

Absorb the $\langle balanced\ text \rangle$ from the input stream delimited by the marker given in the function name, leaving nothing in the input stream.

Absorb the $\langle balanced\ text \rangle$ from the input stream delimited by the marker given in the function name, leaving $\langle inserted\ tokens \rangle$ in the input stream for further processing.

4.6 Predicates and conditionals

LATEX3 has three concepts for conditional flow processing:

Branching conditionals Functions that carry out a test and then execute, depending on its result, either the code supplied as the $\langle true\ code \rangle$ or the $\langle false\ code \rangle$. These arguments are denoted with T and F, respectively. An example would be

```
\cs_if_free:cTF \{abc\} \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}
```

a function that turns the first argument into a control sequence (since it's marked as c) then checks whether this control sequence is still free and then depending on the result carries out the code in the second argument (true case) or in the third argument (false case).

These type of functions are known as "conditionals"; whenever a TF function is defined it is usually accompanied by T and F functions as well. These are provided for convenience when the branch only needs to go a single way. Package writers are free to choose which types to define but the kernel definitions always provide all three versions.

Important to note is that these branching conditionals with $\langle true\ code \rangle$ and/or $\langle false\ code \rangle$ are always defined in a way that the code of the chosen alternative can operate on following tokens in the input stream.

These conditional functions may or may not be fully expandable, but if they are expandable they are accompanied by a "predicate" for the same test as described below.

Predicates "Predicates" are functions that return a special type of boolean value which can be tested by the boolean expression parser. All functions of this type are expandable and have names that end with _p in the description part. For example,

```
\cs_if_free_p:N
```

would be a predicate function for the same type of test as the conditional described above. It would return "true" if its argument (a single token denoted by \mathbb{N}) is still free for definition. It would be used in constructions like

```
\bool_if:nTF {
  \cs_if_free_p:N \l_tmpz_tl || \cs_if_free_p:N \g_tmpz_tl
} {\langle true code \rangle} {\langle false code \rangle}
```

For each predicate defined, a "branching conditional" also exists that behaves like a conditional described above.

Primitive conditionals There is a third variety of conditional, which is the original concept used in plain T_EX and $I^AT_EX 2_E$. Their use is discouraged in expl3 (although still used in low-level definitions) because they are more fragile and in many cases require more expansion control (hence more code) than the two types of conditionals described above.

\c_true_bool
\c_false_bool

Constants that represent ${\tt true}$ and ${\tt false}$, respectively. Used to implement predicates.

4.6.1 Tests on control sequences

```
\label{eq:cs_if_eq_p:NN $$ $$ \langle cs_1 \rangle \langle cs_2 \rangle $$ $$ \langle s_1 \rangle \langle cs_2 \rangle $$ $$ \langle s_1 \rangle \langle cs_2 \rangle {\true code} {\true code} $$ $$ {\true code} $$ $$
```

Compares the definition of two $\langle control\ sequences \rangle$ and is logically true if they are the same, *i.e.* if they have exactly the same definition when examined with \cs_show:N.

```
\label{eq:cs_if_exist_p:N * cs_if_exist_p:N * control sequence} $$ \cs_if_exist_p:x * \cs_if_exist:NTF & control sequence} $$ \{\true code\} $$ {\cs_if_exist:NTF * Tests whether the $$ \control sequence}$ is currently defined (whether as a function or another $$ \cs_if_exist:cTF * control sequence type). Any definition of $$ \control sequence$$ other than $$ \cs_if_exist:cTF * control sequence$$ as true.
```

```
\cs_if_free_p:N \  \  \cs_if_free_p:N \  \control\  sequence \rangle \  \cs_if_free_p:c \  \  \cs_if_free:NTF \  \control\  sequence \rangle \  \{\cs_if_free:NTF \  \control\  sequence \rangle \  \} \  \{\cs_if_free:NTF \  \control\  sequence \rangle \  \cs_if_free:CTF \  \  \control\  sequence \rangle \  \cs_if_exist:NTF).}
```

4.6.2 Primitive conditionals

The ε -TEX engine itself provides many different conditionals. Some expand whatever comes after them and others don't. Hence the names for these underlying functions often contains a :w part but higher level functions are often available. See for instance \int_compare_p:nNn which is a wrapper for \if_int_compare:w.

Certain conditionals deal with specific data types like boxes and fonts and are described there. The ones described below are either the universal conditionals or deal with control sequences. We prefix primitive conditionals with \if_.

TEXhackers note: These are equivalent to their corresponding TEX primitive conditionals; $\text{reverse_if:} N \text{ is } \varepsilon\text{-TEX's } \text{ unless.}$

```
\if_meaning:w * \if_meaning:w \arg_1 \arg_2 \arg_2 \true code \else: \false code \fi:
```

\if_meaning:w executes $\langle true\ code \rangle$ when $\langle arg_1 \rangle$ and $\langle arg_2 \rangle$ are the same, otherwise it executes $\langle false\ code \rangle$. $\langle arg_1 \rangle$ and $\langle arg_2 \rangle$ could be functions, variables, tokens; in all cases the unexpanded definitions are compared.

TeXhackers note: This is TeX's \ifx.

* These conditionals expand any following tokens until two unexpandable tokens are left. If you wish to prevent this expansion, prefix the token in question with \exp_not:N. \if_catcode:w tests if the category codes of the two tokens are the same whereas \if:w tests if the character codes are identical. \if_charcode:w is an alternative name for \if:w.

Check if $\langle cs \rangle$ appears in the hash table or if the control sequence that can be formed from $\langle tokens \rangle$ appears in the hash table. The latter function does not turn the control sequence in question into \scan_stop:! This can be useful when dealing with control sequences which cannot be entered as a single token.

```
\if_mode_horizontal: \( \tau \) code\\ \else: \( \false \) code\\ \fi:
\if_mode_vertical: \( \tau \) code\\ if currently in horizontal mode, otherwise execute \( \false \) code\\. Similar for the other functions.
```

4.7 Starting a paragraph

\mode_leave_vertical: \mode_leave_vertical:

Ensures that TEX is not in vertical (inter-paragraph) mode. In horizontal or math mode this command has no effect, in vertical mode it switches to horizontal mode, and inserts a box of width \parindent, followed by the \everypar token list.

TEXhackers note: This results in the contents of the \everypar token register being inserted, after \mode_leave_vertical: is complete. Notice that in contrast to the LATEX 2ε \leavevmode approach, no box is used by the method implemented here.

4.8 Debugging support

\debug_on:n \debug_off:n

```
\debug_on:n { \( \comma-separated list \) }
\debug_off:n { \( comma-separated list \) }
```

New: 2017-07-16 Turn on and off within a group various debugging code, some of which is also available Updated: 2017-08-02 as expl3 load-time options. The items that can be used in the $\langle list \rangle$ are

- check-declarations that checks all expl3 variables used were previously declared and that local/global variables (based on their name or on their first assignment) are only locally/globally assigned;
- check-expressions that checks integer, dimension, skip, and muskip expressions are not terminated prematurely;
- deprecation that makes soon-to-be-deprecated commands produce errors;
- log-functions that logs function definitions;
- all that does all of the above.

Providing these as switches rather than options allows testing code even if it relies on other packages: load all other packages, call \debug_on:n, and load the code that one is interested in testing. These functions can only be used in LATEX 2ε package mode loaded with enable-debug or another option implying it.

\debug_resume:

\debug_suspend: \debug_suspend: ... \debug_resume:

Suppress (locally) errors and logging from debug commands, except for the deprecation errors or warnings. These pairs of commands can be nested. This can be used around pieces of code that are known to fail checks, if such failures should be ignored. See for instance I3coffins.

Chapter 5

The **I3expan** package Argument expansion

This module provides generic methods for expanding TeX arguments in a systematic manner. The functions in this module all have prefix exp.

Not all possible variations are implemented for every base function. Instead only those that are used within the LATEX3 kernel or otherwise seem to be of general interest are implemented. Consult the module description to find out which functions are actually defined. The next section explains how to define missing variants.

5.1 Defining new variants

The definition of variant forms for base functions may be necessary when writing new functions or when applying a kernel function in a situation that we haven't thought of before.

Internally preprocessing of arguments is done with functions of the form \exp_-.... They all look alike, an example would be \exp_args:NNo. This function has three arguments, the first and the second are a single tokens, while the third argument should be given in braces. Applying \exp_args:NNo expands the content of third argument once before any expansion of the first and second arguments. If \seq_gpush:No was not defined it could be coded in the following way:

```
\exp_args:NNo \seq_gpush:Nn
\g_file_name_stack
{ \l_tmpa_tl }
```

In other words, the first argument to $\exp_{args:NNo}$ is the base function and the other arguments are preprocessed and then passed to this base function. In the example the first argument to the base function should be a single token which is left unchanged while the second argument is expanded once. From this example we can also see how the variants are defined. They just expand into the appropriate \exp_{args} function followed by the desired base function, e.g.

```
\cs_generate_variant:Nn \seq_gpush:Nn { No }
results in the definition of \seq_gpush:No
```

```
\cs_new:Npn \seq_gpush:No { \exp_args:NNo \seq_gpush:Nn }
```

Providing variants in this way in style files is safe as the \cs_generate_variant:Nn function will only create new definitions if there is not already one available. Therefore adding such definition to later releases of the kernel will not make such style files obsolete.

The steps above may be automated by using the function \cs_generate_-variant:Nn, described next.

5.2 Methods for defining variants

We recall the set of available argument specifiers.

- N is used for single-token arguments while c constructs a control sequence from its name and passes it to a parent function as an N-type argument.
- Many argument types extract or expand some tokens and provide it as an n-type argument, namely a braced multiple-token argument: V extracts the value of a variable, v extracts the value from the name of a variable, n uses the argument as it is, o expands once, f expands fully the front of the token list, e and x expand fully all tokens (differences are explained later).
- A few odd argument types remain: T and F for conditional processing, otherwise
 identical to n-type arguments, p for the parameter text in definitions, w for arguments with a specific syntax, and D to denote primitives that should not be used
 directly.

\cs_generate_variant:Nn \cs_generate_variant:cn

 $\cs_generate_variant:Nn \langle parent control sequence \rangle \{\langle variant argument specifiers \rangle\}$

This function is used to define argument-specifier variants of the $\langle parent\ control\ sequence \rangle$ Updated: 2017-11-28 for LATEX3 code-level macros. The \(\text{parent control sequence} \) is first separated into the $\langle base\ name \rangle$ and $\langle original\ argument\ specifier \rangle$. The comma-separated list of $\langle variant\$ argument specifiers is then used to define variants of the (original argument specifier) if these are not already defined. For each $\langle variant \rangle$ given, a function is created that expands its arguments as detailed and passes them to the (parent control sequence). So for example

```
\cs_set:Npn \foo:Nn #1#2 { code here }
\cs_generate_variant:Nn \foo:Nn { c }
```

creates a new function \foo:cn which expands its first argument into a control sequence name and passes the result to \foo:Nn. Similarly

```
\cs_generate_variant:Nn \foo:Nn { NV , cV }
```

generates the functions \foo:NV and \foo:cV in the same way. The \cs_generate_variant: Nn function can only be applied if the (parent control sequence) is already defined. If the $\langle parent\ control\ sequence \rangle$ is protected or if the $\langle variant \rangle$ involves any x argument, then the $\langle variant\ control\ sequence \rangle$ is also protected. The $\langle variant \rangle$ is created globally, as is any \exp_args:N(variant) function needed to carry out the expansion.

Only n and N arguments can be changed to other types. The only allowed changes are

- c variant of an N parent;
- o, V, v, f, e, or x variant of an n parent;
- N, n, T, F, or p argument unchanged.

This means the $\langle parent \rangle$ of a $\langle variant \rangle$ form is always unambiguous, even in cases where both an n-type parent and an N-type parent exist, such as for \tl_count:n and \tl_count: N.

For backward compatibility it is currently possible to make n, o, V, v, f, e, or x-type variants of an N-type argument or N or c-type variants of an n-type argument. Both are deprecated. The first because passing more than one token to an N-type argument will typically break the parent function's code. The second because programmers who use that most often want to access the value of a variable given its name, hence should use a V-type or v-type variant instead of c-type. In those cases, using the lower-level \exp_args:No or \exp_args:Nc functions explicitly is preferred to defining confusing variants.

5.3 Introducing the variants

The V type returns the value of a register, which can be one of tl, clist, int, skip, dim, muskip, or built-in T_FX registers. The v type is the same except it first creates a control sequence out of its argument before returning the value.

In general, the programmer should not need to be concerned with expansion control. When simply using the content of a variable, functions with a V specifier should be used. For those referred to by (cs)name, the v specifier is available for the same purpose. Only when specific expansion steps are needed, such as when using delimited arguments, should the lower-level functions with o specifiers be employed.

The e type expands all tokens fully, starting from the first. More precisely the expansion is identical to that of TeX's \message (in particular # needs not be doubled). It was added in May 2018. In recent enough engines (starting around 2019) it relies on the primitive \expanded hence is fast. In older engines it is very much slower. As a result it should only be used in performance critical code if typical users will have a recent installation of the TeX ecosystem.

The x type expands all tokens fully, starting from the first. In contrast to e, all macro parameter characters # must be doubled, and omitting this leads to low-level errors. In addition this type of expansion is not expandable, namely functions that have x in their signature do not themselves expand when appearing inside x or e expansion.

The f type is so special that it deserves an example. It is typically used in contexts where only expandable commands are allowed. Then x-expansion cannot be used, and f-expansion provides an alternative that expands the front of the token list as much as can be done in such contexts. For instance, say that we want to evaluate the integer expression 3+4 and pass the result 7 as an argument to an expandable function $\exp(n)$. For this, one should define a variant using $\csc(n)$ to $\exp(n)$ then do

```
\ensuremath{\mbox{example:f { \int eval:n { 3 + 4 } }}
```

Note that x-expansion would also expand \int_eval:n fully to its result 7, but the variant \example:x cannot be expandable. Note also that o-expansion would not expand \int_eval:n fully to its result since that function requires several expansions. Besides the fact that x-expansion is protected rather than expandable, another difference between f-expansion and x-expansion is that f-expansion expands tokens from the beginning and stops as soon as a non-expandable token is encountered, while x-expansion continues expanding further tokens. Thus, for instance

```
\example:f { \int_eval:n { 1 + 2 } , \int_eval:n { 3 + 4 } }
results in the call
\example:n { 3 , \int_eval:n { 3 + 4 } }
```

while using \example:x or \example:e instead results in

```
\epsilon \ ( 3 , 7 )
```

at the cost of being protected (for x type) or very much slower in old engines (for e type). If you use f type expansion in conditional processing then you should stick to using TF type functions only as the expansion does not finish any if... fi: itself!

It is important to note that both f- and o-type expansion are concerned with the expansion of tokens from left to right in their arguments. In particular, o-type expansion applies to the first *token* in the argument it receives: it is conceptually similar to

```
\exp_after:wN <base function> \exp_after:wN { <argument> }
```

At the same time, f-type expansion stops at the first non-expandable token. This means for example that both

```
\tl_set:No \l_tmpa_tl { { \g_tmpb_tl } }
```

```
\tl_set:Nf \l_tmpa_tl { { \g_tmpb_tl } }
```

leave \g_tmpb_tl unchanged: { is the first token in the argument and is non-expandable. It is usually best to keep the following in mind when using variant forms.

- Variants with x-type arguments (that are fully expanded before being passed to the n-type base function) are never expandable even when the base function is. Such variants cannot work correctly in arguments that are themselves subject to expansion. Consider using f or e expansion.
- In contrast, e expansion (full expansion, almost like x except for the treatment of #) does not prevent variants from being expandable (if the base function is). The drawback is that e expansion is very much slower in old engines (before 2019). Consider using f expansion if that type of expansion is sufficient to perform the required expansion, or x expansion if the variant will not itself need to be expandable.
- Finally f expansion only expands the front of the token list, stopping at the first non-expandable token. This may fail to fully expand the argument.

When speed is essential (for functions that do very little work and whose variants are used numerous times in a document) the following considerations apply because internal functions for argument expansion come in two flavours, some faster than others.

- Arguments that might need expansion should come first in the list of arguments.
- Arguments that should consist of single tokens N, c, V, or v should come first among these.
- Arguments that appear after the first multi-token argument n, f, e, or o require slightly slower special processing to be expanded. Therefore it is best to use the optimized functions, namely those that contain only N, c, V, and v, and, in the last position, o, f, e, with possible trailing N or n or T or F, which are not expanded. Any x-type argument causes slightly slower processing.

5.4 Manipulating the first argument

These functions are described in detail: expansion of multiple tokens follows the same rules but is described in a shorter fashion.

 $[\]ensuremath{\texttt{exp_args:Nc}} \ \ \, \ensuremath{\texttt{\texp_args:Nc}} \ \ \, \ensuremath{\texttt{\langle tokens}\rangle} \}$

[\]texp_args:cc \times This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are expanded until only characters remain, and are then turned into a control sequence. The result is inserted into the input stream after reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

The :cc variant constructs the $\langle function \rangle$ name in the same manner as described for the $\langle tokens \rangle$.

\exp_args:No \star \exp_args:No $\langle function \rangle \{\langle tokens \rangle\} \dots$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are expanded once, and the result is inserted in braces into the input stream after reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

\exp_args:NV * \exp_args:NV \(function \) \((variable \)

This function absorbs two arguments (the names of the $\langle function \rangle$ and the $\langle variable \rangle$). The content of the $\langle variable \rangle$ are recovered and placed inside braces into the input stream after reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

 $\verb|\exp_args:Nv * \exp_args:Nv & function| & \{ tokens | \}$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are expanded until only characters remain, and are then turned into a control sequence. This control sequence should be the name of a $\langle variable \rangle$. The content of the $\langle variable \rangle$ are recovered and placed inside braces into the input stream after reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

 $\ensuremath{\texttt{exp_args:Ne}} \ \ \, \ensuremath{\texttt{\tokens}} \ensuremath{\texttt{\tokens}}) \}$

New: 2018-05-15 This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$) and exhaustively expands the $\langle tokens \rangle$. The result is inserted in braces into the input stream after reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

TEXhackers note: This relies on the \expanded primitive when available (in LuaTEX and starting around 2019 in other engines). Otherwise it uses some fall-back code that is very much slower. As a result it should only be used in performance-critical code if typical users have a recent installation of the TEX ecosystem.

 $\ensuremath{\texttt{\exp_args:Nf}} \ \ensuremath{\texttt{\exp_args:Nf}} \ \ensuremath{\texttt{\exp_args:Nf}}$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$). The $\langle tokens \rangle$ are fully expanded until the first non-expandable token is found (if that is a space it is removed), and the result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

 $\ensuremath{\texttt{\exp_args:Nx}} \ensuremath{\texttt{\exp_args:Nx}} \ensuremath{\texttt{$

This function absorbs two arguments (the $\langle function \rangle$ name and the $\langle tokens \rangle$) and exhaustively expands the $\langle tokens \rangle$. The result is inserted in braces into the input stream *after* reinsertion of the $\langle function \rangle$. Thus the $\langle function \rangle$ may take more than one argument: all others are left unchanged.

5.5 Manipulating two arguments

```
\verb|\exp_args:NNc * \exp_args:NNc $\langle token_1 \rangle $ $\langle token_2 \rangle $ {\langle tokens \rangle} }
\exp_args:NNV *
               detailed by their argument specifier. The first argument of the function is then the next
\exp_args:NNv *
\langle \exp\_args:NNe  \star item on the input stream, followed by the expansion of the second and third arguments.
\exp_args:NNf *
\exp_args:Ncc *
\exp_args:Nco *
\exp_args:NcV *
\exp_args:Ncv *
\exp_args:Ncf
\exp_args:NVV *
Updated: 2018-05-15
\texttt{\exp\_args:Nnc} \; \star \; \texttt{\exp\_args:Noo} \; \langle token \rangle \; \{\langle tokens_1 \rangle\} \; \{\langle tokens_2 \rangle\}
\exp args:NnV *
               their argument specifier. The first argument of the function is then the next item on
\exp_args:Nnv *
\exp_args:Nne * the input stream, followed by the expansion of the second and third arguments. These
\exp_args:Nnf * functions need slower processing.
\exp args:Noc *
\exp_args:Noo *
\exp_args:Nof *
\exp_args:NVo *
\exp_args:Nfo *
\exp_args:Nff *
\exp_args:Nee *
Updated: 2018-05-15
  \verb|\exp_args:NNx | exp_args:NNx | \langle token_1 \rangle | \langle token_2 \rangle | \{\langle tokens \rangle\}|
  \exp_args:Nnx
                their argument specifier. The first argument of the function is then the next item on
  \exp_args:Nox
               the input stream, followed by the expansion of the second and third arguments. These
  \exp_args:Nxo
  \exp_args:Nxx functions are not expandable due to their x-type argument.
```

5.6 Manipulating three arguments

```
\texttt{\exp\_args:NNcf} \;\; \star \; \texttt{\exp\_args:NNoo} \; \langle token_1 \rangle \; \langle token_2 \rangle \; \{\langle token_3 \rangle\} \; \{\langle tokens \rangle\}
\exp_args:NNno *
                                                    These functions absorb four arguments and expand the second, third and fourth as de-
\exp_args:NNnV *
                                                    tailed by their argument specifier. The first argument of the function is then the next
\exp_args:NNVV * item on the input stream, followed by the expansion of the second argument, etc. These
\exp_args:NNoo *
\exp_args:Ncno * functions need slower processing.
\exp_args:NcnV
\exp_args:Ncoo
\exp_args:NcVV
\exp_args:Nnnc
\exp_args:Nnno
\exp_args:Nnnf
\exp_args:Nnff
\exp_args:Nooo
\exp_args:Noof *
\exp_args:Nffo *
\exp_args:Neee *
      \verb|\exp_args:NNNx \exp_args:NNnx \doken_1| \doken_2| \floor| 
      \exp_args:NNnx These functions absorb four arguments and expand the second, third and fourth as de-
      \exp_args:NNox
                                                     tailed by their argument specifier. The first argument of the function is then the next
      \exp_args:Nccx
                                                    item on the input stream, followed by the expansion of the second argument, etc.
      \exp_args:Ncnx
      \exp_args:Nnnx
      \exp_args:Nnox
      \exp_args:Noox
               New: 2015-08-12
```

5.7 Unbraced expansion

```
\ensuremath{\texttt{\colored}} \ensuremath{\texttt{\colo
\exp_last_unbraced:No
\exp_last_unbraced:NV
                                                                                              These functions absorb the number of arguments given by their specification, carry out
\exp_last_unbraced:Nv
                                                                                               the expansion indicated and leave the results in the input stream, with the last argument
\exp_last_unbraced:Ne
                                                                                               not surrounded by the usual braces. Of these, the :Nno, :Noo, :Nfo and :NnNo variants
\exp_last_unbraced:Nf
                                                                                              need slower processing.
\exp_last_unbraced:NNo
\exp_last_unbraced:NNV
                                                                                                               TeXhackers note: As an optimization, the last argument is unbraced by some of those
\exp last unbraced:NNf
                                                                                                functions before expansion. This can cause problems if the argument is empty: for instance,
\exp_last_unbraced:Nco
                                                                                               \exp_last_unbraced: Nf \foo_bar: w { } \q_stop leads to an infinite loop, as the quark is f-
\exp_last_unbraced:NcV
                                                                                               expanded.
\exp_last_unbraced:Nno
\exp_last_unbraced:Noo
\exp_last_unbraced:Nfo
\exp_last_unbraced:NNNo
\exp_last_unbraced:NNNV
\exp_last_unbraced:NNNf
\exp_last_unbraced:NnNo
\exp_last_unbraced:NNNNo
\exp_last_unbraced:NNNNf
                                         Updated: 2018-05-15
```

 $\verb|\exp_last_unbraced:Nx \exp_last_unbraced:Nx \exp_last_unbraced$

This function fully expands the $\langle tokens \rangle$ and leaves the result in the input stream after reinsertion of the $\langle function \rangle$. This function is not expandable.

This function absorbs three arguments and expands the second and third once. The first argument of the function is then the next item on the input stream, followed by the expansion of the second and third arguments, which are not wrapped in braces. This function needs special (slower) processing.

\exp_after:wN * \exp_after:wN \langle token_1 \rangle \taken_2 \rangle

Carries out a single expansion of $\langle token_2 \rangle$ (which may consume arguments) prior to the expansion of $\langle token_1 \rangle$. If $\langle token_2 \rangle$ has no expansion (for example, if it is a character) then it is left unchanged. It is important to notice that $\langle token_1 \rangle$ may be *any* single token, including group-opening and -closing tokens ($\{ \text{ or } \} \}$ assuming normal TEX category codes). Unless specifically required this should be avoided: expansion should be carried out using an appropriate argument specifier variant or the appropriate $\{ \}$

TEXhackers note: This is the TEX primitive \expandafter renamed.

5.8 Preventing expansion

Despite the fact that the following functions are all about preventing expansion, they're designed to be used in an expandable context and hence are all marked as being 'expandable' since they themselves disappear after the expansion has completed.

\exp_not:N * \exp_not:N \langle token \rangle

Prevents expansion of the $\langle token \rangle$ in a context where it would otherwise be expanded, for example an x-type argument or the first token in an o or e or f argument.

TEX hackers note: This is the TEX \noexpand primitive. It only prevents expansion. At the beginning of an f-type argument, a space $\langle token \rangle$ is removed even if it appears as \exp_not:N \c_space_token. In an x-expanding definition (\cs_new:Npx), a macro parameter introduces an argument even if it appears as \exp_not:N # 1. This differs from \exp_not:n.

 $\verb|\exp_not:c * \exp_not:c {$\langle tokens \rangle$}|$

Expands the $\langle tokens \rangle$ until only characters remain, and then converts this into a control sequence. Further expansion of this control sequence is then inhibited using $\ensuremath{\texttt{\converts}}$ into a control sequence.

$\ensuremath{\mbox{exp_not:n}} \ensuremath{\mbox{$\langle tokens \rangle$}}$

Prevents expansion of the $\langle tokens \rangle$ in an e or x-type argument. In all other cases the $\langle tokens \rangle$ continue to be expanded, for example in the input stream or in other types of arguments such as c, f, v. The argument of \exp_not:n must be surrounded by braces.

TEXhackers note: This is the ε -TEX \unexpanded primitive. In an x-expanding definition (\cs_new:Npx), \exp_not:n {#1} is equivalent to ##1 rather than to #1, namely it inserts the two characters # and 1. In an e-type argument \exp_not:n {#} is equivalent to #, namely it inserts the character #.

\exp_not:o \star \exp_not:o $\{\langle tokens \rangle\}$

Expands the $\langle tokens \rangle$ once, then prevents any further expansion in x-type or e-type arguments using \exp_not:n.

\exp_not:V * \exp_not:V \(\text{variable} \)

Recovers the content of the $\langle variable \rangle$, then prevents expansion of this material in x-type or e-type arguments using \exp not:n.

$\ensuremath{\texttt{exp_not}:v} \ \star \ensuremath{\texttt{exp_not}:v} \ \{\langle tokens \rangle\}$

Expands the $\langle tokens \rangle$ until only characters remains, and then converts this into a control sequence which should be a $\langle variable \rangle$ name. The content of the $\langle variable \rangle$ is recovered, and further expansion in x-type or e-type arguments is prevented using \exp_not:n.

\exp_not:e \star \exp_not:e $\{\langle tokens \rangle\}$

Expands $\langle tokens \rangle$ exhaustively, then protects the result of the expansion (including any tokens which were not expanded) from further expansion in e or x-type arguments using \exp_not:n. This is very rarely useful but is provided for consistency.

Expands $\langle tokens \rangle$ fully until the first unexpandable token is found (if it is a space it is removed). Expansion then stops, and the result of the expansion (including any tokens which were not expanded) is protected from further expansion in x-type or e-type arguments using \exp_not:n.

```
\exp_stop_f: * \foo_bar:f { \langle tokens \rangle \exp_stop_f: \langle more tokens \rangle }
```

Updated: 2011-06-03 This function terminates an f-type expansion. Thus if a function \foo_bar:f starts an f-type expansion and all of \(\lambda tokens\rangle\) are expandable \exp_stop_f: terminates the expansion of tokens even if $\langle more\ tokens \rangle$ are also expandable. The function itself is an implicit space token. Inside an x-type or e-type expansion, it retains its form, but when typeset it produces the underlying space (\Box) .

5.9 Controlled expansion

The expl3 language makes all efforts to hide the complexity of T_FX expansion from the programmer by providing concepts that evaluate/expand arguments of functions prior to calling the "base" functions. Thus, instead of using many \expandafter calls and other trickery it is usually a matter of choosing the right variant of a function to achieve a desired result.

Of course, deep down TFX is using expansion as always and there are cases where a programmer needs to control that expansion directly; typical situations are basic data manipulation tools. This section documents the functions for that level. These commands are used throughout the kernel code, but we hope that outside the kernel there will be little need to resort to them. Instead the argument manipulation methods document above should usually be sufficient.

While \exp_after: wN expands one token (out of order) it is sometimes necessary to expand several tokens in one go. The next set of commands provide this functionality. Be aware that it is absolutely required that the programmer has full control over the tokens to be expanded, i.e., it is not possible to use these functions to expand unknown input as part of $\langle expandable\text{-}tokens \rangle$ as that will break badly if unexpandable tokens are encountered in that place!

\exp:w \exp_end:

* \exp:w \(\left(\text{expandable tokens}\right)\) \(\text{exp_end:}\)

Expands $\langle expandable\text{-}tokens \rangle$ until reaching \exp_end: at which point expansion stops. New: 2015-08-23 The full expansion of $\langle expandable\ tokens\rangle$ has to be empty. If any token in $\langle expandable$ tokens) or any token generated by expanding the tokens therein is not expandable the expansion will end prematurely and as a result \exp_end: will be misinterpreted later $on.^4$

> In typical use cases the \exp_end: is hidden somewhere in the replacement text of (expandable-tokens) rather than being on the same expansion level than \exp:w, e.g., you may see code such as

```
\exp:w \00 case:NnTF #1 {#2} { } { }
```

where somewhere during the expansion of \@@_case:NnTF the \exp_end: gets generated.

TeXhackers note: The current implementation uses \romannumeral hence ignores space tokens and explicit signs + and - in the expansion of the (expandable tokens), but this should not be relied upon.

⁴Due to the implementation you might get the character in position 0 in the current font (typically "'") in the output without any error message!

```
\exp:w *
\exp_end_continue_f:w *
```

New: 2015-08-23

```
\verb|\exp:w| \langle expandable-tokens \rangle \\ | exp_end_continue_f: w| \langle further-tokens \rangle \\ |
```

Expands \(\left(\text{expandable-tokens} \right) \) until reaching \(\text{exp_end_continue_f:w} \) at which point expansion continues as an f-type expansion expanding \(\left(\frac{further-tokens}{r} \right) \) until an unexpandable token is encountered (or the f-type expansion is explicitly terminated by \\exp_-\stop_f:). As with all f-type expansions a space ending the expansion gets removed.

The full expansion of $\langle expandable\text{-}tokens\rangle$ has to be empty. If any token in $\langle expandable\text{-}tokens\rangle$ or any token generated by expanding the tokens therein is not expandable the expansion will end prematurely and as a result $\langle exp_end_continue_f:w$ will be misinterpreted later on.⁵

In typical use cases $\langle expandable\text{-}tokens \rangle$ contains no tokens at all, e.g., you will see code such as

```
\exp_after:wN { \exp:w \exp_end_continue_f:w #2 }
```

where the \exp_after:wN triggers an f-expansion of the tokens in #2. For technical reasons this has to happen using two tokens (if they would be hidden inside another command \exp_after:wN would only expand the command but not trigger any additional f-expansion).

You might wonder why there are two different approaches available, after all the effect of

```
\exp:w \(\langle expandable-tokens \rangle \)\exp_end:
```

can be alternatively achieved through an f-type expansion by using \exp_stop_f:, i.e.

```
\exp:w \exp_end_continue_f:w \( expandable-tokens \) \exp_stop_f:
```

The reason is simply that the first approach is slightly faster (one less token to parse and less expansion internally) so in places where such performance really matters and where we want to explicitly stop the expansion at a defined point the first form is preferable.

```
\exp:w *
\exp_end_continue_f:nw *
```

```
\star \  \, \langle \texttt{exp:w} \  \, \langle \texttt{expandable-tokens} \, \, \rangle \  \, \\ \texttt{exp\_end\_continue\_f:nw} \  \, \langle \texttt{further-tokens} \, \rangle \\
```

The difference to \exp_end_continue_f:w is that we first we pick up an argument which is then returned to the input stream. If $\langle further\text{-}tokens \rangle$ starts with space tokens then these space tokens are removed while searching for the argument. If it starts with a brace group then the braces are removed. Thus such spaces or braces will not terminate the f-type expansion.

⁵In this particular case you may get a character into the output as well as an error message.

5.10 Internal functions

```
\::n \cs_new:Npn \exp_args:Ncof { \::c \::o \::f \::: }
       \::p LATEX3 approach as this makes them more readily visible in the log and so forth. They
       \:: o should not be used outside this module.
       \::e
       \::f
       \::x
       \::v
       \::V
       \:::
\::o_unbraced \cs_new:Npn \exp_last_unbraced:Nno { \::n \::o_unbraced \::: }
\::e_unbraced
           Internal forms for the expansion types which leave the terminal argument unbraced.
\:::f_umbraced
\:::x_umbraced
\:::x_umbraced
\::v_unbraced readily visible in the log and so forth. They should not be used outside this module.
\::V_unbraced
```

Chapter 6

The **I3sort** package Sorting functions

6.1 Controlling sorting

LATEX3 comes with a facility to sort list variables (sequences, token lists, or comma-lists) according to some user-defined comparison. For instance,

results in \l_{foo_clist} holding the values { -2 , 01 , +1 , 3 , 5 } sorted in non-decreasing order.

The code defining the comparison should call \sort_return_swapped: if the two items given as #1 and #2 are not in the correct order, and otherwise it should call \sort_return_same: to indicate that the order of this pair of items should not be changed.

For instance, a $\langle comparison\ code \rangle$ consisting only of \sort_return_same: with no test yields a trivial sort: the final order is identical to the original order. Conversely, using a $\langle comparison\ code \rangle$ consisting only of \sort_return_swapped: reverses the list (in a fairly inefficient way).

TeXhackers note: The current implementation is limited to sorting approximately 20000 items (40000 in LuaTeX), depending on what other packages are loaded.

Internally, the code from l3sort stores items in $\$ registers allocated locally. Thus, the $\langle comparison\ code \rangle$ should not call $\$ registers. On the other hand, altering the value of a previously allocated $\$ register is not a problem.

```
\sort_return_same:
\sort_return_swapped:
```

```
\scalebox{seq\_sort:Nn } \langle seq\ var \rangle
  { ... \sort_return_same: or \sort_return_swapped: ... }
```

 $_{\text{New: 2017-02-06}}$ Indicates whether to keep the order or swap the order of two items that are compared in the sorting code. Only one of the $\scalebox{sort_return_}\dots$ functions should be used by the code, according to the results of some tests on the items #1 and #2 to be compared.

Chapter 7

The I3tl-analysis package: Analysing token lists

This module provides functions that are particularly useful in the 13regex module for mapping through a token list one $\langle token \rangle$ at a time (including begin-group/end-group tokens). For \tl_analysis_map_inline: Nn or \tl_analysis_map_inline: nn, the token list is given as an argument; the analogous function \peek_analysis_map_inline:n documented in l3token finds tokens in the input stream instead. In both cases the user provides $\langle inline\ code \rangle$ that receives three arguments for each $\langle token \rangle$:

- $\langle tokens \rangle$, which both o-expand and x-expand to the $\langle token \rangle$. The detailed form of $\langle tokens \rangle$ may change in later releases.
- $\langle char\ code \rangle$, a decimal representation of the character code of the $\langle token \rangle$, -1 if it is a control sequence.
- $\langle catcode \rangle$, a capital hexadecimal digit which denotes the category code of the $\langle token \rangle$ (0: control sequence, 1: begin-group, 2: end-group, 3: math shift, 4: alignment tab, 6: parameter, 7: superscript, 8: subscript, A: space, B: letter, C: other, D: active). This can be converted to an integer by writing " $\langle catcode \rangle$.

In addition, there is a debugging function \tl_analysis_show:n, very similar to the \ShowTokens macro from the ted package.

```
\tl_analysis_log:N
\tl_analysis_log:n
```

```
\t_{analysis\_show:N \t_{analysis\_show:n \{\langle token\ list \rangle\}}
\t_{analysis\_show:n \tl_analysis\_log:n \{\langle token \ list \rangle\}}
```

Displays to the terminal (or log) the detailed decomposition of the $\langle token \ list \rangle$ into tokens, showing the category code of each character token, the meaning of control sequences and $New:\,2021-05-11\,$ active characters, and the value of registers.

\tl_analysis_map_inline:Nn

 $\tilde{l} = 1$ analysis_map_inline:nn $\tilde{l} = 1$ analysis_map_inline:nn $\tilde{l} = 1$ analysis_map_inline:nn $\tilde{l} = 1$

Applies the $\langle inline\ function \rangle$ to each individual $\langle token \rangle$ in the $\langle token\ list \rangle$. The $\langle inline\ function \rangle$ New: 2018-04-09 function receives three arguments as explained above. As all other mappings the mapping is done at the current group level, i.e. any local assignments made by the \(\lambda inline\) function remain in effect after the loop.

Chapter 8

The l3regex package: Regular expressions in T_EX

The l3regex package provides regular expression testing, extraction of submatches, splitting, and replacement, all acting on token lists. The syntax of regular expressions is mostly a subset of the PCRE syntax (and very close to POSIX), with some additions due to the fact that TEX manipulates tokens rather than characters. For performance reasons, only a limited set of features are implemented. Notably, back-references are not supported.

Let us give a few examples. After

```
\tl_set:Nn \l_my_tl { That~cat. }
\regex_replace_once:nnN { at } { is } \l_my_tl
```

the token list variable \l_my_tl holds the text "This cat.", where the first occurrence of "at" was replaced by "is". A more complicated example is a pattern to emphasize each word and add a comma after it:

```
\ensuremath{\mbox{regex_replace\_all:nnN { \w+ } { \c{emph}\cB{ \0 \cE{} } , } \label{eq:nny_tl}} \ensuremath{\mbox{vegex_replace\_all:nnN { \w+ } { \c{emph}\cB{} } \cB{} } \ensuremath{\mbox{vegex_replace\_all:nnN { \w+ } { \c{emph}\cB{} } \cB{} } \ensuremath{\mbox{vegex_replace\_all:nnN { \w+ } { \center{all:nnN { \w+ } { \center{all:nnN { \w+ } { \center{all:nnN { \cent
```

The \w sequence represents any "word" character, and + indicates that the \w sequence should be repeated as many times as possible (at least once), hence matching a word in the input token list. In the replacement text, $\0$ denotes the full match (here, a word). The command \e is inserted using \c{emph} , and its argument $\0$ is put between braces \c and \c .

If a regular expression is to be used several times, it can be compiled once, and stored in a regex variable using \regex_set:Nn. For example,

```
\regex_new:N \l_foo_regex
\regex_set:Nn \l_foo_regex { \c{begin} \cB. (\c[^BE].*) \cE. }
```

stores in \l_foo_regex a regular expression which matches the starting marker for an environment: \begin, followed by a begin-group token (\cB.), then any number of tokens which are neither begin-group nor end-group character tokens (\c[^BE].*), ending with an end-group token (\cE.). As explained in the next section, the parentheses "capture" the result of \c[^BE].*, giving us access to the name of the environment when doing replacements.

8.1 Syntax of regular expressions

8.1.1 Regular expression examples

We start with a few examples, and encourage the reader to apply \encourage to these regular expressions.

- Cat matches the word "Cat" capitalized in this way, but also matches the beginning of the word "Cattle": use \bCat\b to match a complete word only.
- [abc] matches one letter among "a", "b", "c"; the pattern (a|b|c) matches the same three possible letters (but see the discussion of submatches below).
- [A-Za-z]* matches any number (due to the quantifier *) of Latin letters (not accented).
- \c{[A-Za-z]*} matches a control sequence made of Latin letters.
- _[^_]*_ matches an underscore, any number of characters other than underscore, and another underscore; it is equivalent to _.*?_ where . matches arbitrary characters and the lazy quantifier *? means to match as few characters as possible, thus avoiding matching underscores.
- [\+\-]?\d+ matches an explicit integer with at most one sign.
- [\+\-_]*(\d+|\d*\.\d+)_* matches an explicit integer or decimal number; using [.,] instead of \. would allow the comma as a decimal marker.
- [\+\-_]*(\d+\\d*\.\d+)_*((?i)pt|in|[cem]m|ex|[bs]p|[dn]d|[pcn]c)_* matches an explicit dimension with any unit that TEX knows, where (?i) means to treat lowercase and uppercase letters identically.
- [\+\-_]*((?i)nan|inf|(\d+|\d*\.\d+)(_*e[\+\-_]*\d+)?)_* matches an explicit floating point number or the special values nan and inf (with signs and spaces allowed).
- [\+\-_]*(\d+|\cC.)_* matches an explicit integer or control sequence (without checking whether it is an integer variable).
- \G.*?\K at the beginning of a regular expression matches and discards (due to \K) everything between the end of the previous match (\G) and what is matched by the rest of the regular expression; this is useful in \regex_replace_all:nnN when the goal is to extract matches or submatches in a finer way than with \regex_-extract_all:nnN.

8.1.2 Characters in regular expressions

Most characters match exactly themselves, with an arbitrary category code. Some characters are special and must be escaped with a backslash (e.g., * matches a star character). Some escape sequences of the form backslash—letter also have a special meaning (for instance \d matches any digit). As a rule,

- every alphanumeric character (A-Z, a-z, 0-9) matches exactly itself, and should not be escaped, because \A, \B, ... have special meanings;
- non-alphanumeric printable ascii characters can (and should) always be escaped: many of them have special meanings $(e.g., use \setminus (, \setminus), \cdot?, \setminus., \cdot^)$;
- spaces should always be escaped (even in character classes);
- any other character may be escaped or not, without any effect: both versions match exactly that character.

Note that these rules play nicely with the fact that many non-alphanumeric characters are difficult to input into T_EX under normal category codes. For instance, \\abc\\% matches the characters \abc\% (with arbitrary category codes), but does not match the control sequence \abc followed by a percent character. Matching control sequences can be done using the \c{\(regex\)\} syntax (see below).

Any special character which appears at a place where its special behaviour cannot apply matches itself instead (for instance, a quantifier appearing at the beginning of a string), after raising a warning.

Characters.

```
\x{hh...} Character with hex code hh...
```

\xhh Character with hex code hh.

 \a Alarm (hex 07).

\e Escape (hex 1B).

\f Form-feed (hex 0C).

\n New line (hex 0A).

\r Carriage return (hex 0D).

\t Horizontal tab (hex 09).

8.1.3 Characters classes

Character properties.

. A single period matches any token.

\d Any decimal digit.

\h Any horizontal space character, equivalent to [\ \^^I]: space and tab.

\s Any space character, equivalent to $[\ \^{I}^-J^-L^-M]$.

- \v Any vertical space character, equivalent to [\^^J\^^K\^^L\^^M]. Note that \^^K is a vertical space, but not a space, for compatibility with Perl.
- \w Any word character, *i.e.*, alphanumerics and underscore, equivalent to the explicit class [A-Za-z0-9].
- \D Any token not matched by \d.
- \H Any token not matched by h.
- \N Any token other than the \n character (hex 0A).
- \S Any token not matched by \s .
- \V Any token not matched by \v .
- \W Any token not matched by \w.
- Of those, ., \D , \H , \N , \S , \V , and \W match arbitrary control sequences. Character classes match exactly one token in the subject.
- [...] Positive character class. Matches any of the specified tokens.
- [^...] Negative character class. Matches any token other than the specified characters.
 - x-y Within a character class, this denotes a range (can be used with escaped characters).
- [:\langle name \rangle:] Within a character class (one more set of brackets), this denotes the POSIX character class \langle name \rangle, which can be alnum, alpha, ascii, blank, cntrl, digit, graph, lower, print, punct, space, upper, word, or xdigit.
- [:^\(name\):] Negative POSIX character class.

For instance, [a-oq-z\cC.] matches any lowercase latin letter except p, as well as control sequences (see below for a description of \c).

In character classes, only [, ^, -,], \ and spaces are special, and should be escaped. Other non-alphanumeric characters can still be escaped without harm. Any escape sequence which matches a single character (\d, \D, etc.) is supported in character classes. If the first character is ^, then the meaning of the character class is inverted; ^ appearing anywhere else in the range is not special. If the first character (possibly following a leading ^) is] then it does not need to be escaped since ending the range there would make it empty. Ranges of characters can be expressed using -, for instance, [\D 0-5] and [^6-9] are equivalent.

8.1.4 Structure: alternatives, groups, repetitions

Quantifiers (repetition).

- ? 0 or 1, greedy.
- ?? 0 or 1, lazy.
 - * 0 or more, greedy.
- *? 0 or more, lazy.
- + 1 or more, greedy.

- +? 1 or more, lazy.
- $\{n\}$ Exactly n.
- $\{n,\}$ n or more, greedy.
- $\{n,\}$? n or more, lazy.
- $\{n, m\}$ At least n, no more than m, greedy.
- $\{n, m\}$? At least n, no more than m, lazy.

For greedy quantifiers the regex code will first investigate matches that involve as many repetitions as possible, while for lazy quantifiers it investigates matches with as few repetitions as possible first.

Alternation and capturing groups.

- A|B|C Either one of A, B, or C, investigating A first.
- (...) Capturing group.
- (?:...) Non-capturing group.
- (?|...) Non-capturing group which resets the group number for capturing groups in each alternative. The following group is numbered with the first unused group number.

Capturing groups are a means of extracting information about the match. Parenthesized groups are labelled in the order of their opening parenthesis, starting at 1. The contents of those groups corresponding to the "best" match (leftmost longest) can be extracted and stored in a sequence of token lists using for instance \regex_extract_-once:nnNTF.

The \K escape sequence resets the beginning of the match to the current position in the token list. This only affects what is reported as the full match. For instance,

```
\regex_extract_all:nnN { a \K . } { a123aaxyz } \l_foo_seq
```

results in \l_foo_seq containing the items {1} and {a}: the true matches are {a1} and {aa}, but they are trimmed by the use of \K . The \K command does not affect capturing groups: for instance,

```
\regex_extract_once:nnN { (. \K c)+ \d } { acbc3 } \l_foo_seq
```

results in \l_foo_seq containing the items {c3} and {bc}: the true match is {acbc3}, with first submatch {bc}, but \K resets the beginning of the match to the last position where it appears.

8.1.5 Matching exact tokens

The \c escape sequence allows to test the category code of tokens, and match control sequences. Each character category is represented by a single uppercase letter:

- C for control sequences;
- B for begin-group tokens;
- E for end-group tokens;

- M for math shift;
- T for alignment tab tokens;
- P for macro parameter tokens;
- U for superscript tokens (up);
- D for subscript tokens (down);
- S for spaces;
- L for letters;
- 0 for others; and
- A for active characters.

The \c escape sequence is used as follows.

- $\c{\langle regex \rangle}$ A control sequence whose csname matches the $\langle regex \rangle$, anchored at the beginning and end, so that \c{begin} matches exactly \begin , and nothing else.
 - \cX Applies to the next object, which can be a character, escape character sequence such as \x{0A}, character class, or group, and forces this object to only match tokens with category X (any of CBEMTPUDSLOA. For instance, \cL[A-Z\d] matches uppercase letters and digits of category code letter, \cC. matches any control sequence, and \cO(abc) matches abc where each character has category other.⁶
 - \c[XYZ] Applies to the next object, and forces it to only match tokens with category X, Y, or Z (each being any of CBEMTPUDSLOA). For instance, \c[LSO](...) matches two tokens of category letter, space, or other.
 - \c[^XYZ] Applies to the next object and prevents it from matching any token with category X, Y, or Z (each being any of CBEMTPUDSLOA). For instance, \c[^0]\d matches digits which have any category different from other.

The category code tests can be used inside classes; for instance, \c 0\d \c[L0][A-F]] matches what TEX considers as hexadecimal digits, namely digits with category other, or uppercase letters from A to F with category either letter or other. Within a group affected by a category code test, the outer test can be overridden by a nested test: for instance, \cL(ab\c0*cd) matches ab*cd where all characters are of category letter, except * which has category other.

The \u escape sequence allows to insert the contents of a token list directly into a regular expression or a replacement, avoiding the need to escape special characters. Namely, $\u\{\langle var\ name\rangle\}\$ matches the exact contents (both character codes and category codes) of the variable $\\langle var\ name\rangle$, which are obtained by applying $\ensuremath{\mbox{exp_not:v}}\$ { $\ensuremath{\langle var\ name\rangle}$ } at the time the regular expression is compiled. Within a $\ensuremath{\mbox{c}\{...\}}\$ control sequence matching, the \u escape sequence only expands its argument once, in effect performing \t 1_to_str:v. Quantifiers are supported.

The \ur escape sequence allows to insert the contents of a regex variable into a larger regular expression. For instance, A\ur{1_tmpa_regex}D matches the tokens A and

⁶This last example also captures "abc" as a regex group; to avoid this use a non-capturing group \c0(?:abc).

D separated by something that matches the regular expression \l_tmpa_regex. This behaves as if a non-capturing group were surrounding \l_tmpa_regex, and any group contained in \l_tmpa_regex is converted to a non-capturing group. Quantifiers are supported.

For instance, if \l_tmpa_regex has value B|C, then A\ur{1_tmpa_regex}D is equivalent to A(?:B|C)D (matching ABD or ACD) and not to AB|CD (matching AB or CD). To get the latter effect, it is simplest to use TEX's expansion machinery directly: if $\l_tmymodule_BC_tl$ contains B|C then the following two lines show the same result:

```
\regex_show:n { A \u{l_mymodule_BC_tl} D }
\regex_show:n { A B | C D }
```

8.1.6 Miscellaneous

Anchors and simple assertions.

- \b Word boundary: either the previous token is matched by \w and the next by \W, or the opposite. For this purpose, the ends of the token list are considered as \W.
- \B Not a word boundary: between two \w tokens or two \W tokens (including the boundary).

or \A Start of the subject token list.

- $\$, $\$ or $\$ End of the subject token list.
 - \G Start of the current match. This is only different from ^ in the case of multiple matches: for instance \regex_count:nnN { \G a } { aaba } \l_tmpa_int yields 2, but replacing \G by ^ would result in \l_tmpa_int holding the value 1.

The option (?i) makes the match case insensitive (treating A-Z and a-z as equivalent, with no support yet for Unicode case changing). This applies until the end of the group in which it appears, and can be reverted using (?-i). For instance, in (?i)(a(?-i)b|c)d, the letters a and d are affected by the i option. Characters within ranges and classes are affected individually: (?i)[\?-B] is equivalent to [\?@ABab] (and differs from the much larger class [\?-b]), and (?i)[^aeiou] matches any character which is not a vowel. The i option has no effect on \c{...}, on \u{...}, on character properties, or on character classes, for instance it has no effect at all in (?i)\u{l_foo_tl}\d\d[:lower:]].

8.2 Syntax of the replacement text

Most of the features described in regular expressions do not make sense within the replacement text. Backslash introduces various special constructions, described further below:

- \0 is the whole match;
- \1 is the submatch that was matched by the first (capturing) group (...); similarly for $\2, \ldots, \9$ and $\g\{\langle number \rangle\}$;

- \a, \e, \f, \n, \r, \t, \xhh, \x{hhh} correspond to single characters as in regular expressions;
- $\c{\langle cs \ name \rangle}$ inserts a control sequence;
- $\c\c$ (see below);
- $\{\langle tl \ var \ name \rangle\}$ inserts the contents of the $\langle tl \ var \rangle$ (see below).

Characters other than backslash and space are simply inserted in the result (but since the replacement text is first converted to a string, one should also escape characters that are special for TeX, for instance use \#). Non-alphanumeric characters can always be safely escaped with a backslash.

For instance,

```
\tl_set:Nn \l_my_tl { Hello,~world! }
\regex_replace_all:nnN { ([er]?l|o) . } { (\0--\1) } \l_my_tl
results in \l_my_tl holding H(ell--el)(o,--o) w(or--o)(ld--l)!
```

The submatches are numbered according to the order in which the opening parenthesis of capturing groups appear in the regular expression to match. The n-th submatch is empty if there are fewer than n capturing groups or for capturing groups that appear in alternatives that were not used for the match. In case a capturing group matches several times during a match (due to quantifiers) only the last match is used in the replacement text. Submatches always keep the same category codes as in the original token list.

By default, the category code of characters inserted by the replacement are determined by the prevailing category code regime at the time where the replacement is made, with two exceptions:

- space characters (with character code 32) inserted with \u or \x20 or \x{20} have category code 10 regardless of the prevailing category code regime;
- if the category code would be 0 (escape), 5 (newline), 9 (ignore), 14 (comment) or 15 (invalid), it is replaced by 12 (other) instead.

The escape sequence \c allows to insert characters with arbitrary category codes, as well as control sequences.

- \cX(...) Produces the characters "..." with category X, which must be one of CBEMTPUDSLOA as in regular expressions. Parentheses are optional for a single character (which can be an escape sequence). When nested, the innermost category code applies, for instance \cL(Hello\cS\ world)! gives this text with standard category codes.
- $\c{\langle text \rangle}$ Produces the control sequence with csname $\langle text \rangle$. The $\langle text \rangle$ may contain references to the submatches 0, 1, and so on, as in the example for u below.

The escape sequence \u{\(var name \)} allows to insert the contents of the variable with name \(var name \) directly into the replacement, giving an easier control of category codes. When nested in \c{...} and \u{...} constructions, the \u and \c escape sequences perform \tl_to_str:v, namely extract the value of the control sequence and turn it into a string. Matches can also be used within the arguments of \c and \u. For instance,

```
\tl_set:Nn \l_my_one_tl { first }
\tl_set:Nn \l_my_two_tl { \emph{second} }
\tl_set:Nn \l_my_tl { one , two , one , one }
\regex_replace_all:nnN { [^,]+ } { \u{l_my_\0_tl} } \l_my_tl
```

results in \l_my_tl holding first,\emph{second},first,first.

Regex replacement is also a convenient way to produce token lists with arbitrary category codes. For instance

```
\tl_clear:N \l_tmpa_tl
\regex_replace_all:nnN { } { \cU\% \cA\~ } \l_tmpa_tl
```

results in \l_tmpa_tl containing the percent character with category code 7 (superscript) and an active tilde character.

8.3 Pre-compiling regular expressions

If a regular expression is to be used several times, it is better to compile it once rather than doing it each time the regular expression is used. The compiled regular expression is stored in a variable. All of the l3regex module's functions can be given their regular expression argument either as an explicit string or as a compiled regular expression.

```
\regex_new:N \regex_new:N \regex var \rangle
           New: 2017-05-26 Creates a new \langle regex\ var \rangle or raises an error if the name is already taken. The declaration
                                                  is global. The \langle regex\ var \rangle is initially such that it never matches.
                                                  \ensuremath{\mbox{regex\_set:Nn}} \langle \ensuremath{\mbox{regex}} \ensuremath{\mbox{var}} \rangle \ \{\langle \ensuremath{\mbox{regex}} \rangle\}
   \regex_set:Nn
   \regex_gset:Nn
                                                  Stores a compiled version of the \langle regular\ expression \rangle in the \langle regex\ var \rangle. The assignment
           New: 2017-05-26 is local for \regex_set:Nn and global for \regex_gset:Nn. For instance, this function
                                                  can be used as
                                                                \regex_new:N \l_my_regex
                                                                \regex_set:Nn \l_my_regex { my\ (simple\ )? reg(ex|ular\ expression) }
New: 2017-05-26 Creates a new constant \langle regex\ var \rangle or raises an error if the name is already taken. The
                                                  value of the \langle regex\ var \rangle is set globally to the compiled version of the \langle regular\ expression \rangle.
  \regex_show:N
                                                  \rcspace \
                                                  \ensuremath{\mbox{regex\_log:n }} \langle \ensuremath{\mbox{regex}} \rangle
  \regex_show:n
  \regex_log:N
                                                  Displays in the terminal or writes in the log file (respectively) how |3regex interprets the
  \regex_log:n
                                                  \langle regex \rangle. For instance, \regex_show:n {\A X|Y} shows
           New: 2021-04-26
                                                                +-branch
 Updated: 2021-04-29
                                                                       anchor at start (\A)
                                                                       char code 88 (X)
```

+-branch

char code 89 (Y)

indicating that the anchor \A only applies to the first branch: the second branch is not anchored to the beginning of the match.

8.4 Matching

All regular expression functions are available in both :n and :N variants. The former require a "standard" regular expression, while the later require a compiled expression as generated by \regex_set:Nn.

New: 2017-05-26

 $\label{lem:normatch$

```
\regex_match:nnTF { b [cde]* } { abecdcx } { TRUE } { FALSE }
\regex_match:nnTF { [b-dq-w] } { example } { TRUE } { FALSE }
```

leaves TRUE then FALSE in the input stream.

\regex_count:NnN

```
\label{limit_equal_count:nnN} $$\operatorname{count:nnN} {\langle regex \rangle} {\langle token \; list \rangle} \; \langle int \; var \rangle$$
```

New: 2017-05-26

Sets (int var) within the current TFX group level equal to the number of times (regular expression appears in $\langle token \ list \rangle$. The search starts by finding the left-most longest match, respecting greedy and lazy (non-greedy) operators. Then the search starts again from the character following the last character of the previous match, until reaching the end of the token list. Infinite loops are prevented in the case where the regular expression can match an empty token list: then we count one match between each pair of characters. For instance,

```
\int_new:N \l_foo_int
    \regex_count:nnN { (b+|c) } { abbababcbb } \l_foo_int
results in \l_foo_int taking the value 5.
```

\regex_match_case:nn \regex_match_case:nn*TF*

New: 2022-01-10

```
\regex_match_case:nnTF
   {
        \{\langle regex_1 \rangle\}\ \{\langle code\ case_1 \rangle\}
        \{\langle regex_2 \rangle\}\ \{\langle code\ case_2 \rangle\}
        \{\langle regex_n \rangle\}\ \{\langle code\ case_n \rangle\}
   } {\langle token list\rangle}
    \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
```

Determines which of the $\langle regular\ expressions \rangle$ matches at the earliest point in the $\langle token$ list, and leaves the corresponding $\langle code_i \rangle$ followed by the $\langle true\ code \rangle$ in the input stream. If several $\langle regex \rangle$ match starting at the same point, then the first one in the list is selected and the others are discarded. If none of the $\langle regex \rangle$ match, the $\langle false\ code \rangle$ is left in the input stream. Each $\langle regex \rangle$ can either be given as a regex variable or as an explicit regular expression.

In detail, for each starting position in the $\langle token \ list \rangle$, each of the $\langle regex \rangle$ is searched in turn. If one of them matches then the corresponding $\langle code \rangle$ is used and everything else is discarded, while if none of the $\langle reqex \rangle$ match at a given position then the next starting position is attempted. If none of the $\langle regex \rangle$ match anywhere in the $\langle token \ list \rangle$ then nothing is left in the input stream. Note that this differs from nested \regex_match:nnTF statements since all $\langle regex \rangle$ are attempted at each position rather than attempting to match $\langle regex_1 \rangle$ at every position before moving on to $\langle regex_2 \rangle$.

8.5 Submatch extraction

```
\regex_extract_once:nnN
\verb|\regex_extract_once:nnN| $\underline{TF}$|
\regex_extract_once:NnN
\regex_extract_once:NnNTF
```

```
\c \ensuremath{\verb|regex_extract_once:nnN||} \{\langle extract_once:nnN|| \{\langle extract_once:nnN|| \{\langle extract_once:nnN|| | \langle extract_once:nnN|| | \langle
\label{lem:nntf} $$\operatorname{code}_{nn} TF {\langle regex \rangle} {\langle token \; list \rangle} \; \langle seq \; var \rangle \; {\langle true \; code \rangle} \; {\langle false \; true \; false \; true \; code \rangle} \; {\langle false \; true \; false \; false \; true \; false \; true \; false \; false \; true \; false \; false \; false \; true \; false \; fa
     code \}
```

Finds the first match of the $\langle regular\ expression \rangle$ in the $\langle token\ list \rangle$. If it exists, the match New: 2017-05-26 is stored as the first item of the $\langle seq \, var \rangle$, and further items are the contents of capturing groups, in the order of their opening parenthesis. The $\langle seq \ var \rangle$ is assigned locally. If there is no match, the $\langle seq \ var \rangle$ is cleared. The testing versions insert the $\langle true \ code \rangle$ into the input stream if a match was found, and the $\langle false\ code \rangle$ otherwise.

For instance, assume that you type

```
\regex_extract_once:nnNTF { \A(La)?TeX(!*)\Z } { LaTeX!!! } \l_foo_seq
  { true } { false }
```

Then the regular expression (anchored at the start with A and at the end with Z) must match the whole token list. The first capturing group, (La)?, matches La, and the second capturing group, (!*), matches !!!. Thus, \l_foo_seq contains as a result the items {LaTeX!!!}, {La}, and {!!!}, and the true branch is left in the input stream. Note that the n-th item of \l_foo_seq, as obtained using \seq_item:Nn, correspond to the submatch numbered (n-1) in functions such as $\regex_replace_once:nnN$.

```
\regex_extract_all:nnN
\regex_extract_all:nnNTF
\regex_extract_all:NnN
\regex_extract_all:NnNTF
```

```
\c \c extract_all:nnN {\langle regex\rangle} {\langle token \ list\rangle} {\langle seq \ var\rangle}
\label{lem:list} $\operatorname{\code} \ {\code} \ {\cod
```

Finds all matches of the $\langle regular \ expression \rangle$ in the $\langle token \ list \rangle$, and stores all the sub-New: 2017-05-26 match information in a single sequence (concatenating the results of multiple \regex_extract_once:nnN calls). The $\langle seq \ var \rangle$ is assigned locally. If there is no match, the $\langle seq \ var \rangle$ is cleared. The testing versions insert the $\langle true \ code \rangle$ into the input stream if a match was found, and the $\langle false\ code \rangle$ otherwise. For instance, assume that you type

```
\regex_extract_all:nnNTF { \w+ } { Hello,~world! } \l_foo_seq
  { true } { false }
```

Then the regular expression matches twice, the resulting sequence contains the two items {Hello} and {world}, and the true branch is left in the input stream.

```
\regex_split:nnN
\regex_split:NnN
```

```
\rcspace{1.5cm} \rcspace{1.5
\{\langle false\ code \rangle\}
```

New: 2017-05-26 expression. If the $\langle regular \ expression \rangle$ has capturing groups, then the token lists that they match are stored as items of the sequence as well. The assignment to $\langle seq \, var \rangle$ is local. If no match is found the resulting $\langle seq \ var \rangle$ has the $\langle token \ list \rangle$ as its sole item. If the $\langle regular \ expression \rangle$ matches the empty token list, then the $\langle token \ list \rangle$ is split into single tokens. The testing versions insert the $\langle true\ code \rangle$ into the input stream if a match was found, and the $\langle false\ code \rangle$ otherwise. For example, after

```
\seq_new:N \l_path_seq
\regex_split:nnNTF { / } { the/path/for/this/file.tex } \l_path_seq
  { true } { false }
```

the sequence \l_path_seq contains the items {the}, {path}, {for}, {this}, and {file.tex}, and the true branch is left in the input stream.

Replacement 8.6

\regex_replace_once:nnN \regex_replace_once:nnNTF \regex_replace_once:NnN \regex_replace_once:NnNTF

 $\ensuremath{\mbox{regex_replace_once:nnN}} \{\langle \ensuremath{\mbox{regular expression}} \rangle\} \{\langle \ensuremath{\mbox{replace_ment}} \rangle\} \langle \ensuremath{\mbox{tl var}} \rangle$ $\rule = \rule = \rul$ code} { $\langle false \ code \rangle$ }

Searches for the $\langle regular \ expression \rangle$ in the contents of the $\langle tl \ var \rangle$ and replaces the first New: 2017-05-26 match with the $\langle replacement \rangle$. In the $\langle replacement \rangle$, \0 represents the full match, \1 represent the contents of the first capturing group, $\$ 2 of the second, etc. The result is assigned locally to $\langle tl \ var \rangle$.

\regex_replace_all:nnN \regex_replace_all:nnNTF \regex_replace_all:NnN \regex_replace_all:NnNTF

code)} {\(false code \) }

New: 2017-05-26

Replaces all occurrences of the $\langle regular\ expression \rangle$ in the contents of the $\langle tl\ var \rangle$ by the $\langle replacement \rangle$, where \0 represents the full match, \1 represent the contents of the first capturing group, \2 of the second, etc. Every match is treated independently, and matches cannot overlap. The result is assigned locally to $\langle tl \ var \rangle$.

Replaces the earliest match of the regular expression $(?|\langle regex_1\rangle|...|\langle regex_n\rangle)$ in the $\langle token\ list\ variable \rangle$ by the $\langle replacement \rangle$ corresponding to which $\langle regex_i \rangle$ matched, then leaves the $\langle true\ code \rangle$ in the input stream. If none of the $\langle regex \rangle$ match, then the $\langle tl\ var \rangle$ is not modified, and the $\langle false\ code \rangle$ is left in the input stream. Each $\langle regex \rangle$ can either be given as a regex variable or as an explicit regular expression.

In detail, for each starting position in the $\langle token\ list \rangle$, each of the $\langle regex \rangle$ is searched in turn. If one of them matches then it is replaced by the corresponding $\langle replacement \rangle$ as described for $\regex_replace_once:nnN$. This is equivalent to checking with $\regex_-match_case:nn$ which $\langle regex \rangle$ matches, then performing the replacement with $\regex_-replace_once:nnN$.

```
\label{eq:local_case_all:nNTF} $$\operatorname{\ensuremath{$\sim$}} $$ \operatorname{\ensuremath{$\sim$}} $$ $$\operatorname{\ensuremath{$\sim$}} $$$ $$\operatorname{\ensuremath{$\sim$}} $$$ $$\operatorname{\ensuremath{$\sim$}} $$$$ $$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$$\operatorname{\ensuremath{$\sim$}} $$$$
```

Replaces all occurrences of all $\langle regex \rangle$ in the $\langle token\ list \rangle$ by the corresponding $\langle replacement \rangle$. Every match is treated independently, and matches cannot overlap. The result is assigned locally to $\langle tl\ var \rangle$, and the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ is left in the input stream depending on whether any replacement was made or not.

In detail, for each starting position in the $\langle token\ list \rangle$, each of the $\langle regex \rangle$ is searched in turn. If one of them matches then it is replaced by the corresponding $\langle replacement \rangle$, and the search resumes at the position that follows this match (and replacement). For instance

results in \l_tmpa_tl having the contents ''Hello'',——[,][_]''world'',——[!]. Note in particular that the word-boundary assertion \b did not match at the start of words because the case [A-Za-z]+ matched at these positions. To change this, one could simply swap the order of the two cases in the argument of \regex_replace_case_all:nN.

8.7 Scratch regular expressions

\l_tmpa_regex Scratch regex for local assignment. These are never used by the kernel code, and so are \l_tmpb_regex safe for use with any LATEX3-defined function. However, they may be overwritten by New: 2017-12-11 other non-kernel code and so should only be used for short-term storage.

\g_tmpa_regex Scratch regex for global assignment. These are never used by the kernel code, and so \g_tmpb_regex are safe for use with any LATEX3-defined function. However, they may be overwritten by New: 2017-12-11 other non-kernel code and so should only be used for short-term storage.

Bugs, misfeatures, future work, and other possi-8.8 bilities

The following need to be done now.

- Rewrite the documentation in a more ordered way, perhaps add a BNF? Additional error-checking to come.
- Clean up the use of messages.
- Cleaner error reporting in the replacement phase.
- Add tracing information.
- Detect attempts to use back-references and other non-implemented syntax.
- Test for the maximum register \c_max_register_int.
- Find out whether the fact that \W and friends match the end-marker leads to bugs. Possibly update __regex_item_reverse:n.
- The empty cs should be matched by \c{\}, not by \c{csname.?endcsname\s?}. Code improvements to come.
- Shift arrays so that the useful information starts at position 1.
- Only build \c{...} once.
- Use arrays for the left and right state stacks when compiling a regex.
- Should __regex_action_free_group:n only be used for greedy {n,} quantifier? (I think not.)
- Quantifiers for \u and assertions.
- When matching, keep track of an explicit stack of curr_state and curr_submatches.
- If possible, when a state is reused by the same thread, kill other subthreads.

- Use an array rather than \g_regex_balance_tl to build the function _regex_replacement_balance_one_match:n.
- Reduce the number of epsilon-transitions in alternatives.
- Optimize simple strings: use less states (abcade should give two states, for abc and ade). [Does that really make sense?]
- Optimize groups with no alternative.
- Optimize states with a single __regex_action_free:n.
- Optimize the use of __regex_action_success: by inserting it in state 2 directly instead of having an extra transition.
- Optimize the use of \int_step_... functions.
- Groups don't capture within regexes for csnames; optimize and document.
- Better "show" for anchors, properties, and catcode tests.
- Does \K really need a new state for itself?
- When compiling, use a boolean in cs and less magic numbers.
- Instead of checking whether the character is special or alphanumeric using its character code, check if it is special in regexes with \cs_if_exist tests.

The following features are likely to be implemented at some point in the future.

- General look-ahead/behind assertions.
- Regex matching on external files.
- Conditional subpatterns with look ahead/behind: "if what follows is $[\dots]$, then $[\dots]$ ".
- (*..) and (?..) sequences to set some options.
- UTF-8 mode for pdfT_FX.
- Newline conventions are not done. In particular, we should have an option for . not to match newlines. Also, \A should differ from $\$, and \Z , \z and $\$ should differ.
- Unicode properties: \p{..} and \P{..}; \X which should match any "extended" Unicode sequence. This requires to manipulate a lot of data, probably using tree-boxes.

The following features of PCRE or Perl may or may not be implemented.

• Callout with (?C...) or other syntax: some internal code changes make that possible, and it can be useful for instance in the replacement code to stop a regex replacement when some marker has been found; this raises the question of a potential \regex_break: and then of playing well with \tl_map_break: called from within the code in a regex. It also raises the question of nested calls to the regex machinery, which is a problem since \fontdimen are global.

- Conditional subpatterns (other than with a look-ahead or look-behind condition): this is non-regular, isn't it?
- Named subpatterns: TeX programmers have lived so far without any need for named macro parameters.

The following features of PCRE or Perl will definitely not be implemented.

- Back-references: non-regular feature, this requires backtracking, which is prohibitively slow.
- Recursion: this is a non-regular feature.
- Atomic grouping, possessive quantifiers: those tools, mostly meant to fix catastrophic backtracking, are unnecessary in a non-backtracking algorithm, and difficult to implement.
- Subroutine calls: this syntactic sugar is difficult to include in a non-backtracking algorithm, in particular because the corresponding group should be treated as atomic.
- Backtracking control verbs: intrinsically tied to backtracking.
- \ddd, matching the character with octal code ddd: we already have \x{...} and the syntax is confusingly close to what we could have used for backreferences (\1, \2, ...), making it harder to produce useful error message.
- \cx, similar to TEX's own \^^x.
- Comments: TEX already has its own system for comments.
- $\Q...\E$ escaping: this would require to read the argument verbatim, which is not in the scope of this module.
- \C single byte in UTF-8 mode: X_HT_EX and LuaT_EX serve us characters directly, and splitting those into bytes is tricky, encoding dependent, and most likely not useful anyways.

Chapter 9

The l3prg package Control structures

Conditional processing in LaTeX3 is defined as something that performs a series of tests, possibly involving assignments and calling other functions that do not read further ahead in the input stream. After processing the input, a *state* is returned. The states returned are $\langle true \rangle$ and $\langle false \rangle$.

If TeX3 has two forms of conditional flow processing based on these states. The first form is predicate functions that turn the returned state into a boolean $\langle true \rangle$ or $\langle false \rangle$. For example, the function \cs_if_free_p:N checks whether the control sequence given as its argument is free and then returns the boolean $\langle true \rangle$ or $\langle false \rangle$ values to be used in testing with \if_predicate:w or in functions to be described below. The second form is the kind of functions choosing a particular argument from the input stream based on the result of the testing as in \cs_if_free:NTF which also takes one argument (the N) and then executes either true or false depending on the result.

 T_EX hackers note: The arguments are executed after exiting the underlying $\inf... fi:$ structure.

9.1 Defining a set of conditional functions

\prg_new_conditional:Npnn \prg_set_conditional:Npnn \prg_gset_conditional:Npnn \prg_new_conditional:Nnn \prg_set_conditional:Nnn \prg_gset_conditional:Nnn

Updated: 2022-11-01

 $\prg_new_conditional:Npnn \end{arg spec} \end{arg$

These functions create a family of conditionals using the same $\{\langle code \rangle\}$ to perform the test created. Those conditionals are expandable if $\langle code \rangle$ is. The new versions check for existing definitions and perform assignments globally $(cf. \cs_new:Npn)$ whereas the set versions do no check and perform assignments locally $(cf. \cs_set:Npn)$. The conditionals created are dependent on the comma-separated list of $\langle conditions \rangle$, which should be one or more of p, T, F and TF.

Updated: 2012-02-06

These functions create a family of protected conditionals using the same $\{\langle code \rangle\}$ to perform the test created. The $\langle code \rangle$ does not need to be expandable. The new version check for existing definitions and perform assignments globally $(cf. \cs_new:Npn)$ whereas the set version do not $(cf. \cs_set:Npn)$. The conditionals created are depended on the comma-separated list of $\langle conditions \rangle$, which should be one or more of T, F and TF (not p).

The conditionals are defined by \prg_new_conditional:Npnn and friends as:

- \\name_p:\langle arg spec \rangle a predicate function which will supply either a logical true or logical false. This function is intended for use in cases where one or more logical tests are combined to lead to a final outcome. This function cannot be defined for protected conditionals.
- $\langle name \rangle : \langle arg \ spec \rangle T$ a function with one more argument than the original $\langle arg \ spec \rangle$ demands. The $\langle true \ branch \rangle$ code in this additional argument will be left on the input stream only if the test is true.
- $\langle name \rangle : \langle arg \ spec \rangle F$ a function with one more argument than the original $\langle arg \ spec \rangle$ demands. The $\langle false \ branch \rangle$ code in this additional argument will be left on the input stream only if the test is false.
- $\mbox{\ensuremath{\langle name \rangle : \langle arg~spec \rangle TF}}$ a function with two more argument than the original $\mbox{\ensuremath{\langle arg~spec \rangle}}$ demands. The $\mbox{\ensuremath{\langle true~branch \rangle}}$ code in the first additional argument will be left on the input stream if the test is true, while the $\mbox{\ensuremath{\langle false~branch \rangle}}$ code in the second argument will be left on the input stream if the test is false.

The $\langle code \rangle$ of the test may use $\langle parameters \rangle$ as specified by the second argument to $prg_{set_conditional:Npnn:}$ this should match the $\langle argument\ specification \rangle$ but this is not enforced. The Nnn versions infer the number of arguments from the argument specification given $(cf. \cs_new:Nn,\ etc.)$. Within the $\langle code \rangle$, the functions $prg_return_true:$ and $prg_return_false:$ are used to indicate the logical outcomes of the test.

An example can easily clarify matters here:

This defines the function \foo_if_bar_p:NN, \foo_if_bar:NNTF and \foo_if_bar:NNT but not \foo_if_bar:NNF (because F is missing from the \(\chiconditions \rangle \) list). The return statements take care of resolving the remaining \else: and \fi: before returning the state. There must be a return statement for each branch; failing to do so will result in erroneous output if that branch is executed.

```
\label{lem:lem:new_eq_conditional:NNn } $$ \operatorname{prg_new_eq_conditional:NNn } \langle \operatorname{name_1} \rangle : \langle \operatorname{arg spec_1} \rangle : \langle \operatorname{arg spec_2} \rangle \operatorname{prg_set_eq_conditional:NNn} \{\langle \operatorname{conditions} \rangle \}$$
```

These functions copy a family of conditionals. The new version checks for existing definitions (cf. \cs_new_eq:NN) whereas the set version does not (cf. \cs_set_eq:NN). The conditionals copied are depended on the comma-separated list of $\langle conditions \rangle$, which should be one or more of p, T, F and TF.

```
\prg_return_true: * \prg_return_true:
\prg_return_false: * \prg_return_false:
```

These "return" functions define the logical state of a conditional statement. They appear within the code for a conditional function generated by \prg_set_conditional:Npnn, etc, to indicate when a true or false branch should be taken. While they may appear multiple times each within the code of such conditionals, the execution of the conditional must result in the expansion of one of these two functions exactly once.

The return functions trigger what is internally an f-expansion process to complete the evaluation of the conditional. Therefore, after \prg_return_true: or \prg_return_false: there must be no non-expandable material in the input stream for the remainder of the expansion of the conditional code. This includes other instances of either of these functions.

```
\frac{\texttt{\prg\_generate\_conditional\_variant:Nnn \nm\nmember\normal}}{\texttt{\prg\_generate\_conditional\_variant:Nnn \nm\nmember\normal}} \ \{\langle variant \ argument \ specifiers \rangle\} \ \{\langle condition \ spe
```

Defines argument-specifier variants of conditionals. This is equivalent to running \cs_generate_variant: \Nn \langle conditional \rangle \{\langle variant argument specifiers \rangle \} \) on each \langle conditional \rangle \decorpoonders \rangle conditionals \rangle \text{ are obtained from the } \langle name \rangle \text{ and } \langle arg spec \rangle \text{ as described for \prg_new_conditional: Npnn, and they should be defined.}

9.2 The boolean data type

This section describes a boolean data type which is closely connected to conditional processing as sometimes you want to execute some code depending on the value of a switch (e.g., draft/final) and other times you perhaps want to use it as a predicate function in an \if_predicate:w test. The problem of the primitive \if_false: and \if_true: tokens is that it is not always safe to pass them around as they may interfere with scanning for termination of primitive conditional processing. Therefore, we employ two canonical booleans: \c_true_bool or \c_false_bool. Besides preventing problems as described above, it also allows us to implement a simple boolean parser supporting the logical operations And, Or, Not, etc. which can then be used on both the boolean type and predicate functions.

All conditional \bool_ functions except assignments are expandable and expect the input to also be fully expandable (which generally means being constructed from predicate functions and booleans, possibly nested).

TEXhackers note: The bool data type is not implemented using the \iffalse/\iffrue primitives, in contrast to \newif, etc., in plain TEX, LATEX 2ε and so on. Programmers should not base use of bool switches on any particular expectation of the implementation.

```
\bool_new:N \bool_new:N \boolean \
                                              \bool_new:c
                                                                                              Creates a new \langle boolean \rangle or raises an error if the name is already taken. The declaration
                                                                                               is global. The \langle boolean \rangle is initially false.
                                  \verb|\bool_const:Nn \bool_const:Nn \aligned boolean| \{\aligned boolean \aligned \} | \{\aligned boolean \aligned boolean \aligne
                                  \bool_const:cn
                                                                                               Creates a new constant \langle boolean \rangle or raises an error if the name is already taken. The
                                            New: 2017-11-28
                                                                                             value of the \langle boolean \rangle is set globally to the result of evaluating the \langle boolean \rangle.
                  \bool_set_false:N
                                                                                             \bool_set_false:N \langle boolean \rangle
                  \bool_set_false:c
                                                                                              Sets \(\langle boolean \rangle \) logically false.
                  \bool_gset_false:N
                  \bool_gset_false:c
                      \bool_set_true:N
                                                                                              \bool_set_true:N \langle boolean \rangle
                      \bool_set_true:c
                                                                                              Sets \langle boolean \rangle \logically true.
                     \bool_gset_true:N
                      \bool_gset_true:c
\bool_set_eq:NN
                                                                                               \bool_set_eq:NN \ \langle boolean_1 \rangle \ \langle boolean_2 \rangle
\bool_set_eq:(cN|Nc|cc)
                                                                                              Sets \langle boolean_1 \rangle to the current value of \langle boolean_2 \rangle.
\bool_gset_eq:NN
\bool_gset_eq:(cN|Nc|cc)
                                                                                               \verb|\bool_set:Nn| \langle boolean \rangle | \{\langle boolexpr \rangle\}|
                                \bool_set:Nn
                                \bool_set:cn
                                                                                              Evaluates the \(\langle boolean \) expression\\ as described for \\\bool_if:nTF, and sets the \(\langle boolean \rangle \)
                                \bool_gset:Nn
                                                                                               variable to the logical truth of this evaluation.
                                \bool_gset:cn
                                Updated: 2017-07-15
                                \bool_if_p:N * \bool_if_p:N \langle boolean \rangle
                                \bool_{if_p:c} \star \bool_{if:NTF} \bool_{if_n:code} {\true\ code} \ {\true\ code} \
                               \begin{tabular}{ll} \beg
                                \bool_if:cTF *
                                Updated: 2017-07-15
                         \verb|\bool_to_str:N| \star \verb|\bool_to_str:N| \langle boolean \rangle|
                         \bool_to_str:c * \bool_to_str:n \langle boolean expression \rangle
                         \underline{\underline{\text{bool\_to\_str:n}}} Expands to the letters true or false depending on the logical truth of the \underline{\langle boolean \rangle} or
                                            New: 2021-11-01 \langle boolean \ expression \rangle.
```

```
\bool_show:N \langle boolean \rangle
                          \bool_show: N
                          \bool_show:c
                                                                                               Displays the logical truth of the \langle boolean \rangle on the terminal.
                                         New: 2012-02-09
                          Updated: 2021-04-29
                          \bool_show:n
                                                                                               \bool_show:n {\boolean expression}}
                                        New: 2012-02-09 Displays the logical truth of the (boolean expression) on the terminal.
                          Updated: 2017-07-15
                          \bool_log:N
                                                                                               \bool_log:N \langle boolean \rangle
                          \bool_log:c
                                                                                               Writes the logical truth of the \langle boolean \rangle in the log file.
                                        New: 2014-08-22
                          Updated: 2021-04-29
                                                                                               \bool_log:n {\boolean expression}}
                          \bool_log:n
                                       New: 2014-08-22 Writes the logical truth of the (boolean expression) in the log file.
                          Updated: 2017-07-15
\bool_if_exist_p:N * \bool_if_exist_p:N \langle boolean \rangle
\label{local_if_exist_p:c} $$ \bool_if_exist:NTF $$ \boolean $$ {\langle true\ code \rangle} $$ $$ {\langle false\ code \rangle}$$
\verb|\bool_if_exist:N$ $\underline{\textit{TF}}$ $^{\star}$ Tests whether the $\langle boolean \rangle$ is currently defined. This does not check that the $\langle boolean \rangle$ is currently defined. The $\langle boolean \rangle$ is currently defined at $\langle boolean \rangle$ is currently 
 \bool_if_exist:cTF \star
                                                                                             really is a boolean variable.
                                       New: 2012-03-03
```

9.2.1Scratch booleans

\l_tmpa_bool A scratch boolean for local assignment. It is never used by the kernel code, and so is \l_tmpb_bool safe for use with any LATEX3-defined function. However, it may be overwritten by other non-kernel code and so should only be used for short-term storage.

\g_tmpb_bool

\g_tmpa_bool A scratch boolean for global assignment. It is never used by the kernel code, and so is safe for use with any LATEX3-defined function. However, it may be overwritten by other non-kernel code and so should only be used for short-term storage.

9.3 Boolean expressions

As we have a boolean datatype and predicate functions returning boolean $\langle true \rangle$ or $\langle false \rangle$ values, it seems only fitting that we also provide a parser for $\langle boolean \ expressions \rangle$.

A boolean expression is an expression which given input in the form of predicate functions and boolean variables, return boolean $\langle true \rangle$ or $\langle false \rangle$. It supports the logical operations And, Or and Not as the well-known infix operators && and || and prefix! with their usual precedences (namely, && binds more tightly than ||). In addition to this, parentheses can be used to isolate sub-expressions. For example,

```
\int_compare_p:n { 1 = 1 } &&
  (
     \int_compare_p:n { 2 = 3 } ||
     \int_compare_p:n { 4 <= 4 } ||
     \str_if_eq_p:nn { abc } { def }
  ) &&
! \int_compare_p:n { 2 = 4 }</pre>
```

is a valid boolean expression.

Contrarily to some other programming languages, the operators && and || evaluate both operands in all cases, even when the first operand is enough to determine the result. This "eager" evaluation should be contrasted with the "lazy" evaluation of \bool_lazy_-... functions.

TeXhackers note: The eager evaluation of boolean expressions is unfortunately necessary in TeX. Indeed, a lazy parser can get confused if && or || or parentheses appear as (unbraced) arguments of some predicates. For instance, the innocuous-looking expression below would break (in a lazy parser) if #1 were a closing parenthesis and \l_tmpa_bool were true.

```
( \l_tmpa_bool || \token_if_eq_meaning_p:NN X #1 )
```

Minimal (lazy) evaluation can be obtained using the conditionals \bool_lazy_-all:nTF, \bool_lazy_and:nnTF, \bool_lazy_any:nTF, or \bool_lazy_or:nnTF, which only evaluate their boolean expression arguments when they are needed to determine the resulting truth value. For example, when evaluating the boolean expression

the line marked with skipped is not expanded because the result of \bool_lazy_any_-p:n is known once the second boolean expression is found to be logically true. On the other hand, the last line is expanded because its logical value is needed to determine the result of \bool_lazy_and_p:nn.

```
\bool_if_p:n \  \  \bool_if_p:n \  \{\langle boolean \  expression\rangle\} \  \bool_if:nTF \  \bool_if:nTF \  \{\langle boolean \  expression\rangle\} \  \{\langle true \  code\rangle\} \  \{\langle false \  code\rangle\} \  \bool_if:nTF \  \b
```

Tests the current truth of \(\begin{aligned} \begin{aligned} \leftarrow \text{topdated: 2017-07-15} \]

Tests the current truth of \(\begin{aligned} \leftarrow \text{boolean expression} \rightarrow \) should consist of a series of predicates or boolean variables with the logical relationship between these defined using && ("And"), || ("Or"), ! ("Not") and parentheses. The logical Not applies to the next predicate or group.

```
\begin{subarray}{ll} $$ \beg
  \begin{tabular}{ll} $$ \bool_lazy_all:nTF { $$ (boolexpr_1)$ } {(boolexpr_2)$} & \cdots {(boolexpr_N)}$ } {(true \ code)}$ \\ \end{tabular} $$ \bool_lazy_all:nTF ( (boolexpr_1)) } $$ (code)$ } $$ \bool_lazy_all:nTF ( (boolexpr_2)) } & \cdots {(boolexpr_N)}$ } $$ (code)$ } $$ (code)$ } $$
                           Updated: 2017-07-15 Implements the "And" operation on the (boolean expressions), hence is true if all of
                                                              them are true and false if any of them is false. Contrarily to the infix operator &&,
                                                              only the (boolean expressions) which are needed to determine the result of \bool_lazy_-
                                                               all:nTF are evaluated. See also \bool lazy and:nnTF when there are only two \boolean
                                                               expressions \rangle.
\bool_lazy_and_p:nn \star \bool_lazy_and_p:nn {\langle boolexpr_1 \rangle} {\langle boolexpr_2 \rangle}
\begin{tabular}{ll} $$ \bool_lazy_and:nnTF $$ \bool_expr_1$ } $$ {\bool_expr_2$ } $$ {\true code} $$ {\false code}$ } $$
                           New: 2015-11-15 Implements the "And" operation between two boolean expressions, hence is true if both
                  Updated: 2017-07-15 are true. Contrarily to the infix operator &&, the \langle boolexpr_2 \rangle is only evaluated if it is
                                                              needed to determine the result of \bool_lazy_and:nnTF. See also \bool_lazy_all:nTF
                                                               when there are more than two \langle boolean \ expressions \rangle.
  \begin{subarray}{ll} \begin{
  \begin{tabular}{ll} $$ \bool_lazy_any:nTF $ { (boolexpr_1) } { (boolexpr_2) } \cdots { (boolexpr_N) } } { (true \ code) } \end{tabular} $$ $$ (true \ code) }
                                                             \{\langle false\ code \rangle\}
                           New: 2015-11-15
                  Updated: 2017-07-15 Implements the "Or" operation on the (boolean expressions), hence is true if any of
                                                               them is true and false if all of them are false. Contrarily to the infix operator ||,
                                                               only the (boolean expressions) which are needed to determine the result of \bool lazy -
                                                               any:nTF are evaluated. See also \bool_lazy_or:nnTF when there are only two \boolean
                                                               expressions \rangle.
  \bool_lazy_or_p:nn * \bool_lazy_or_p:nn {\langle boolexpr_1 \rangle} {\langle boolexpr_2 \rangle}
  \begin{tabular}{ll} $$ \bool_lazy_or:nnTF $$ \bool_lazy_or:nnTF $$ (\boolexpr_1)$ } $$ {\codeyr_2}$ } $$ {\codeyr_2}$ } $$ {\codeyr_2}$ } $$
                            New: 2015-11-15 Implements the "Or" operation between two boolean expressions, hence is true if either
                  Updated: 2017-07-15 one is true. Contrarily to the infix operator | \cdot |, the \langle boolexpr_2 \rangle is only evaluated if it is
                                                              needed to determine the result of \bool_lazy_or:nnTF. See also \bool_lazy_any:nTF
                                                               when there are more than two \langle boolean \ expressions \rangle.
                \verb|\bool_not_p:n * \bool_not_p:n {| \langle boolean \ expression \rangle }|
                  Updated: 2017-07-15 Function version of ! (\langle boolean\ expression \rangle) within a boolean expression.
              \bool\_xor\_p:nn * \bool\_xor\_p:nn {\langle boolexpr_1 \rangle} {\langle boolexpr_2 \rangle}
              \verb|\bool_xor:nnTF| $$ \bool_xor:nnTF {$\langle boolexpr_1 \rangle$} {\langle boolexpr_2 \rangle$} {\langle true\ code \rangle} {\langle false\ code \rangle} $
                           New: 2018-05-09 Implements an "exclusive or" operation between two boolean expressions. There is no
                                                              infix operation for this logical operation.
```

9.4 Logical loops

Loops using either boolean expressions or stored boolean values.

```
\bool_do_until:\n \( \triangle \) \bool_do_until:\n \( \lambda \) \lambda \( \code \rangle \) \\
\triangle \bool_do_until:\cn \( \frac{\pi}{\pi} \) \text{Places the } \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \) in the input stream for TeX to procupate the \( \code \rangle \rangle \) in the input stream for TeX to procupate the \( \code \rangle \rangle \) in the input stream for TeX to procupate the \( \code \rangle \rangle \) in the input stream for TeX to procupate the \( \code \rangle \rangle \) in the input stream for \( \code \rangle \rangle \) in the input stream for \( \code \rangle \rangle \rangle \) in the input stream for \( \code \rangle \
```

Places the $\langle code \rangle$ in the input stream for T_EX to process, and then checks the logical value of the $\langle boolean \rangle$. If it is false then the $\langle code \rangle$ is inserted into the input stream again and the process loops until the $\langle boolean \rangle$ is true.

 $\bool_do_while:Nn \Leftrightarrow \bool_do_while:Nn \boolean \ \{\code\}\}$ \bool_do_while:cn ☆ Places the $\langle code \rangle$ in the input stream for T_FX to process, and then checks the logical value of the $\langle boolean \rangle$. If it is true then the $\langle code \rangle$ is inserted into the input stream again and the process loops until the $\langle boolean \rangle$ is false. $\bool_until_do: Nn \boolean \ \{\code\}\}$ \bool until do:Nn ☆ \bool_until_do:cn ☆ This function firsts checks the logical value of the $\langle boolean \rangle$. If it is false the $\langle code \rangle$ is Updated: 2017-07-15 placed in the input stream and expanded. After the completion of the $\langle code \rangle$ the truth of the $\langle boolean \rangle$ is re-evaluated. The process then loops until the $\langle boolean \rangle$ is true. $\bool_while_do:Nn \Leftrightarrow \bool_while_do:Nn \boolean \ \{\code\}\}$ \bool_while_do:cn ☆ This function firsts checks the logical value of the $\langle boolean \rangle$. If it is true the $\langle code \rangle$ is Updated: 2017-07-15 placed in the input stream and expanded. After the completion of the $\langle code \rangle$ the truth of the $\langle boolean \rangle$ is re-evaluated. The process then loops until the $\langle boolean \rangle$ is false. $\begin{tabular}{ll} \verb|\doc| do_until:nn | $\langle boolean \ expression \rangle \} & \{\langle code \rangle \} \\ \end{tabular}$ Updated: 2017-07-15 Places the $\langle code \rangle$ in the input stream for T_FX to process, and then checks the logical value of the (boolean expression) as described for \bool_if:nTF. If it is false then the $\langle code \rangle$ is inserted into the input stream again and the process loops until the $\langle boolean \rangle$ expression \rangle evaluates to true. \bool_do_while:nn ☆ \bool_do_while:nn {\boolean expression}} {\cdot code} Updated: 2017-07-15 Places the $\langle code \rangle$ in the input stream for TFX to process, and then checks the logical value of the (boolean expression) as described for \bool if:nTF. If it is true then the $\langle code \rangle$ is inserted into the input stream again and the process loops until the $\langle boolean \rangle$ expression evaluates to false. Updated: 2017-07-15 This function firsts checks the logical value of the (boolean expression) (as described for \bool_if:nTF). If it is false the $\langle code \rangle$ is placed in the input stream and expanded. After the completion of the $\langle code \rangle$ the truth of the $\langle boolean \ expression \rangle$ is re-evaluated. The process then loops until the $\langle boolean \ expression \rangle$ is true. \bool_while_do:nn ☆ \bool_while_do:nn {\boolean expression}} {\code} Updated: 2017-07-15 This function firsts checks the logical value of the (boolean expression) (as described for

9.5 Producing multiple copies

The process then loops until the $\langle boolean \ expression \rangle$ is false.

\bool_if:nTF). If it is **true** the $\langle code \rangle$ is placed in the input stream and expanded. After the completion of the $\langle code \rangle$ the truth of the $\langle boolean\ expression \rangle$ is re-evaluated.

Detecting T_EX's mode 9.6

```
\mode_if_horizontal_p: * \mode_if_horizontal_p:
\verb|\mode_if_horizontal|: $\underline{\mathit{TF}} \; \star \; \mode_if_horizontal: $\mathsf{TF} \; \{\langle \mathit{true} \; \mathit{code} \rangle\} \; \{\langle \mathit{false} \; \mathit{code} \rangle\} \; \\
                                                                                                                     Detects if T<sub>F</sub>X is currently in horizontal mode.
                       \mode_if_inner_p: * \mode_if_inner_p:
                       \verb|\mode_if_inner: TF | & \mode_if_inner: TF
                                                                                                                     Detects if T<sub>F</sub>X is currently in inner mode.
                            \mbox{\code_if_math_p: } \star \mbox{\code_if_math:TF } {\code} \} \ {\code} \}
                           Updated: 2011-09-05
          \mode_if_vertical_p: * \mode_if_vertical_p:
         \mbox{\code_if_vertical:} \underline{TF} \ \star \mbox{\code_if_vertical:} TF \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}
                                                                                                                     Detects if T<sub>F</sub>X is currently in vertical mode.
```

Primitive conditionals 9.7

```
\verb|\if_predicate:w| \star \verb|\if_predicate:w| \langle predicate \rangle | \langle true| code \rangle \\ \verb|\ext{lese}: \langle false| code \rangle \\ \verb|\fi:|
```

This function takes a predicate function and branches according to the result. (In practice this function would also accept a single boolean variable in place of the $\langle predicate \rangle$ but to make the coding clearer this should be done through \if bool:N.)

```
\if_bool:N * \if_bool:N \langle boolean \rangle \true code \rangle \langle lse: \langle false code \rangle \fi:
```

This function takes a boolean variable and branches according to the result.

9.8 Nestable recursions and mappings

There are a number of places where recursion or mapping constructs are used in expl3. At a low-level, these typically require insertion of tokens at the end of the content to allow "clean up". To support such mappings in a nestable form, the following functions are provided.

```
\proonup \
```

New: 2018-03-26 Used to mark the end of a recursion or mapping: the functions $\langle type \rangle$ _map_break: and \\type_map_break:n use this to break out of the loop (see \prg_map_break:Nn for how to set these up). After the loop ends, the $\langle code \rangle$ is inserted into the input stream. This occurs even if the break functions are not applied: \prg_break_point:Nn is functionally-equivalent in these cases to \use_ii:nn.

```
\project{projection} \projec
```

New: 2018-03-26

```
\prg_break_point: \propty $$ \qrowvert in $$
```

Breaks a recursion in mapping contexts, inserting in the input stream the $\langle user\ code \rangle$ after the $\langle ending\ code \rangle$ for the loop. The function breaks loops, inserting their $\langle ending\$ code, until reaching a loop with the same $\langle type \rangle$ as its first argument. This $\langle type \rangle_{-}$ map break: argument must be defined; it is simply used as a recognizable marker for the $\langle type \rangle$.

For types with mappings defined in the kernel, $\langle type \rangle_{map_break}$: and $\langle type \rangle_{-}$ map_break:n are defined as $\prg_map_break:Nn \times \$ _map_break: {} and the same with {} omitted.

9.8.1Simple mappings

In addition to the more complex mappings above, non-nestable mappings are used in a number of locations and support is provided for these.

\prg_break_point: * This copy of \prg_do_nothing: is used to mark the end of a fast short-term recursion: New: 2018-03-27 the function \prg_break:n uses this to break out of the loop.

\prg_break: $\prg_break:n *$

 $\star \project{\projection} \project{\projecti$

Breaks a recursion which has no $\langle ending\ code \rangle$ and which is not a user-breakable mapping New: 2018-03-27 (see for instance \prop get:Nn), and inserts the $\langle code \rangle$ in the input stream.

9.9 Internal programming functions

\group_align_safe_begin: * \group_align_safe_begin:

\group_align_safe_end:

\group_align_safe_end:

Updated: 2011-08-11

These functions are used to enclose material in a TFX alignment environment within a specially-constructed group. This group is designed in such a way that it does not add brace groups to the output but does act as a group for the & token inside \halign. This is necessary to allow grabbing of tokens for testing purposes, as TFX uses group level to determine the effect of alignment tokens. Without the special grouping, the use of a function such as \peek_after: Nw would result in a forbidden comparison of the internal \endtemplate token, yielding a fatal error. Each \group_align_safe_begin: must be matched by a \group_align_safe_end:, although this does not have to occur within the same function.

Chapter 10

The l3sys package: System/runtime functions

The name of the job 10.1

\c_sys_jobname_str Constant that gets the "job name" assigned when TEX starts.

New: 2015-09-19 Updated: 2019-10-27

TrXhackers note: This copies the contents of the primitive \jobname. For technical reasons, the string here is not of the same internal form as other, but may be manipulated using normal string functions.

10.2 Date and time

\c_sys_hour_int \c_sys_day_int \c_sys_year_int

\c_sys_minute_int The date and time at which the current job was started: these are all reported as integers.

TeXhackers note: Whilst the underlying primitives can be altered by the user, this \c_sys_month_int interface to the time and date is intended to be the "real" values.

New: 2015-09-22

10.3 Engine

\sys_if_engine_luatex_p: $\sys_if_engine_pdftex:TF \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}$ \sys_if_engine_luatex: TF Conditionals which allow engine-specific code to be used. The names follow naturally \sys_if_engine_pdftex_p: from those of the engine binaries: note that the (u)ptex tests are for ε -pTFX and ε -upTFX \sys_if_engine_pdftex: TF as expl3 requires the ε -T_EX extensions. Each conditional is true for exactly one supported \sys_if_engine_ptex_p: engine. In particular, \sys_if_engine_ptex_p: is true for ε -pTFX but false for ε -upTFX. \sys_if_engine_ptex: <u>TF</u> \sys_if_engine_uptex_p: \sys_if_engine_uptex: TF \sys_if_engine_xetex_p: \sys_if_engine_xetex: TF New: 2015-09-07 The current engine given as a lower case string: one of luatex, pdftex, ptex, uptex or New: 2015-09-19 xetex. \c_sys_engine_exec_str The name of the standard executable for the current TFX engine given as a lower case string: one of luatex, luahbtex, pdftex, eptex, euptex or xetex. New: 2020-08-20 \c_sys_engine_format_str The name of the preloaded format for the current TFX run given as a lower case string: New: 2020-08-20 one of lualatex (or dvilualatex), pdflatex (or latex), platex, uplatex or xelatex for LATEX, similar names for plain TEX (except pdfTEX in DVI mode yields etex), and cont-en for ConTFXt (i.e. the \fmtname). \sys_timer: \star \sys_timer: New: 2020-09-24 Expands to the current value of the engine's timer clock, a non-negative integer. This function is only defined for engines with timer support. This command measures not just CPU time but real time (including time waiting for user input). The unit are scaled seconds $(2^{-16} \text{ seconds}).$

10.4 Output format

New: 2015-09-19

\sys_if_output_dvi_p: * \sys_if_output_dvi:TF {\langle code \rangle} {\langle false code \rangle} \sys_if_output_dvi:TF * Conditionals which give the current output mode the TeX run is operating in. This is always one of two outcomes, DVI mode or PDF mode. The two sets of conditionals are thus complementary and are both provided to allow the programmer to emphasise the most appropriate case.

\[\sum_{New: 2015-09-19} \]
\[\scale_{c_sys_output_str} \]
The current output mode given as a lower case string: one of dvi or pdf.

73

10.5 Platform

Conditionals which allow platform-specific code to be used. The names follow the Lua os.type() function, *i.e.* all Unix-like systems are unix (including Linux and MacOS).

\(\c_sys_platform_str\)
New: 2018-07-27

The current platform given as a lower case string: one of unix, windows or unknown.

10.6 Random numbers

 $\sys_rand_seed: * \sys_rand_seed:$

New: 2017-05-27 Expands to the current value of the engine's random seed, a non-negative integer. In engines without random number support this expands to 0.

 $\verb|\sys_gset_rand_seed:n \sys_gset_rand_seed:n \{\langle intexpr \rangle\}|$

New: 2017-05-27 Globally sets the seed for the engine's pseudo-random number generator to the \(\lambda integer expression \rangle\). This random seed affects all \(\ldot\)._rand functions (such as \int_rand:nn or \(\clist_rand_item:n\)) as well as other packages relying on the engine's random number generator. In engines without random number support this produces an error.

TeXhackers note: While a 32-bit (signed) integer can be given as a seed, only the absolute value is used and any number beyond 2^{28} is divided by an appropriate power of 2. We recommend using an integer in $[0, 2^{28} - 1]$.

10.7 Access to the shell

Defines $\langle tl \ var \rangle$ to the text returned by the $\langle shell \ command \rangle$. The $\langle shell \ command \rangle$ is converted to a string using $\tl_to_str:n$. Category codes may need to be set appropriately via the $\langle shell \ command \rangle$ argument, which is run just before running the $\langle shell \ command \rangle$ (in a group). If shell escape is disabled, the $\langle tl \ var \rangle$ will be set to \tl_no_value in the non-branching version. Note that quote characters (") cannot be used inside the $\langle shell \ command \rangle$. The $\sl_sys_get_shell:nnNTF$ conditional inserts the true code if the shell is available and no quote is detected, and the false code otherwise.

\c_sys_shell_escape_int This variable exposes the internal triple of the shell escape status. The possible values New: 2017-05-27

0 Shell escape is disabled

- 1 Unrestricted shell escape is enabled
- 2 Restricted shell escape is enabled

```
\sys_if_shell_p: * \sys_if_shell_p:
```

 $\sys_if_shell: TF \star \sys_if_shell: TF \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}$

New: 2017-05-27 Performs a check for whether shell escape is enabled. This returns true if either of restricted or unrestricted shell escape is enabled.

```
\sys_if_shell_unrestricted_p: * \sys_if_shell_unrestricted_p:
\verb|\sys_if_shell_unrestricted: $\underline{\mathit{TF}} \ \star \sys_if_shell_unrestricted: $\mathsf{TF} \ \{\langle \mathit{true} \ \mathit{code} \rangle\} \ \{\langle \mathit{false} \ \mathit{code} \rangle\} \ \} \ \{\langle \mathit{false} \ \mathit{code} \rangle\} \ \} \ \{\langle \mathit{false} \ \mathit{code} \rangle\} \ \} \ \}
                                                                                                                                                                                                                                                                                                                   New: 2017-05-27
```

Performs a check for whether *unrestricted* shell escape is enabled.

```
\sys_if_shell_restricted_p: * \sys_if_shell_restricted_p:
\sps_{if\_shell\_restricted:} TF \ \ \sps_{if\_shell\_restricted:} TF \ \{\t code\}\} \ \{\t false \ code\}\}
                    New: 2017-05-27
```

Performs a check for whether restricted shell escape is enabled. This returns false if unrestricted shell escape is enabled. Unrestricted shell escape is not considered a superset of restricted shell escape in this case. To find whether any shell escape is enabled use \sys_if_shell:.

```
\sys\_shell\_now:n \sys\_shell\_now:n {\langle tokens \rangle}
\frac{\text{\colored}_{\text{\colored}} x}{\text{\colored}} Execute \langle tokens \rangle through shell escape immediately.
       New: 2017-05-27
```

 $\sys_{shell_shipout:n \sys_shell_shipout:n }{ \langle tokens \rangle }$

New: 2017-05-27

Loading configuration data 10.8

 $\sys_load_backend:n \sys_load_backend:n {\langle backend \rangle}$

New: 2019-09-12 Loads the additional configuration file needed for backend support. If the $\langle backend \rangle$ is empty, the standard backend for the engine in use will be loaded. This command may only be used once.

\sys_ensure_backend: \sys_ensure_backend:

New: 2022-07-29 Ensures that a backend has been loaded by calling \sys_load_backend:n if required.

\c_sys_backend_str Set to the name of the backend in use by \sys_load_backend:n when issued. Possible values are

- pdftex
- luatex
- xetex
- dvips
- dvipdfmx
- dvisvgm

\sys_load_debug: \sys_load_debug:

 $New:\ 2019-09-12$ Load the additional configuration file for debugging support.

10.8.1 Final settings

\sys_finalise: \sys_finalise:

New: 2019-10-06 Finalises all system-dependent functionality: required before loading a backend.

Chapter 11

The **I3msg** package Messages

Messages need to be passed to the user by modules, either when errors occur or to indicate how the code is proceeding. The l3msg module provides a consistent method for doing this (as opposed to writing directly to the terminal or log).

The system used by l3msg to create messages divides the process into two distinct parts. Named messages are created in the first part of the process; at this stage, no decision is made about the type of output that the message will produce. The second part of the process is actually producing a message. At this stage a choice of message class has to be made, for example error, warning or info.

By separating out the creation and use of messages, several benefits are available. First, the messages can be altered later without needing details of where they are used in the code. This makes it possible to alter the language used, the detail level and so on. Secondly, the output which results from a given message can be altered. This can be done on a message class, module or message name basis. In this way, message behaviour can be altered and messages can be entirely suppressed.

11.1 Creating new messages

Messages may be subdivided by one level using the / character. This is used within the message filtering system to allow for example the LATEX kernel messages to belong to the module LaTeX while still being filterable at a more granular level. Thus for example

```
\msg_new:nnnn { mymodule } { submodule / message } ...
```

will allow to filter out specifically messages from the submodule.

\msg_new:nnnn \msg_new:nnn

 $\mbox{\constraints} $$\max_{new:nnnn} {\langle module \rangle} {\langle message \rangle} {\langle text \rangle} {\langle more\ text \rangle}$$

Updated: 2011-08-16

Creates a $\langle message \rangle$ for a given $\langle module \rangle$. The message is defined to first give $\langle text \rangle$ and then $\langle more\ text \rangle$ if the user requests it. If no $\langle more\ text \rangle$ is available then a standard text is given instead. Within $\langle text \rangle$ and $\langle more\ text \rangle$ four parameters (#1 to #4) can be used: these will be supplied at the time the message is used. An error is raised if the $\langle message \rangle$ already exists.

\msg_set:nnnn \msg_set:nnn

 $\mbox{\constraints} $$ \mbox{\constraints} {\constraints} {\cons$

\msg_gset:nnnn \msg_gset:nnn

Sets up the text for a $\langle message \rangle$ for a given $\langle module \rangle$. The message is defined to first give $\langle text \rangle$ and then $\langle more\ text \rangle$ if the user requests it. If no $\langle more\ text \rangle$ is available then a standard text is given instead. Within $\langle text \rangle$ and $\langle more\ text \rangle$ four parameters (#1 to #4) can be used: these will be supplied at the time the message is used.

 $\mbox{\sc msg_if_exist_p:nn $$\times \mbox{\sc msg_if_exist_p:nn } {\mbox{\sc module}$} {\mbox{\sc message}$}}$

 $\label{localization} $$\max_{if_exist:nnTF \ {\module}} {\module} \ {\module} \$

New: 2012-03-03 Tests whether the $\langle message \rangle$ for the $\langle module \rangle$ is currently defined.

Customizable information for message modules 11.2

New: 2018-10-10 Expands to the public name of the \(\lambda\) as defined by \g_msg_module_name_prop (or otherwise leaves the $\langle module \rangle$ unchanged).

 $\label{local_module_type:n * local_module_type:n { (module)}} $$ \mbox{msg_module_type:n } {\mbox{module_type:n } {\mbox{module_type:n$

New: 2018-10-10 Expands to the description which applies to the $\langle module \rangle$, for example a Package or Class. The information here is defined in \g_msg_module_type_prop, and will default to Package if an entry is not present.

New: 2018-10-10

\g_msg_module_name_prop Provides a mapping between the module name used for messages, and that for documentation.

\g_msg_module_type_prop Provides a mapping between the module name used for messages, and that type of New: 2018-10-10 module. For example, for LATEX3 core messages, an empty entry is set here meaning that they are not described using the standard Package text.

Contextual information for messages 11.3

\msg_line_context: ☆ \msg_line_context:

Prints the current line number when a message is given, and thus suitable for giving context to messages. The number itself is proceeded by the text on line.

\msg_line_number: * \msg_line_number:

Prints the current line number when a message is given.

 $\mbox{\mbox{$\mbox{msg_fatal_text:n}$}}$

Produces the standard text

Fatal Package (module) Error

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included.

 $\verb|\msg_critical_text:n * \msg_critical_text:n {| module |}$

Produces the standard text

Critical Package (module) Error

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included.

Produces the standard text

Package \(\text{module} \) Error

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included.

\msg_warning_text:n * \msg_warning_text:n {\(\lambda odule \) \}

Produces the standard text

Package (module) Warning

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included. The $\langle type \rangle$ of $\langle module \rangle$ may be adjusted: Package is the standard outcome: see \msg_module_type:n.

 $\label{local_msg_info_text:n + local} $$\max_{i=1}^{\infty} info_{text:n \{(module)\}}$$$

Produces the standard text:

Package \(\text{module} \) Info

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included. The $\langle type \rangle$ of $\langle module \rangle$ may be adjusted: Package is the standard outcome: see \msg_module_type:n.

Produces the standard text

See the \(\lambda module \rangle \) documentation for further information.

This function can be redefined to alter the language in which the message is given, using #1 as the name of the $\langle module \rangle$ to be included. The name of the $\langle module \rangle$ is produced using \msg_module_name:n.

11.4 Issuing messages

Messages behave differently depending on the message class. In all cases, the message may be issued supplying 0 to 4 arguments. If the number of arguments supplied here does not match the number in the definition of the message, extra arguments are ignored, or empty arguments added (of course the sense of the message may be impaired). The four arguments are converted to strings before being added to the message text: the x-type variants should be used to expand material. Note that this expansion takes place with the standard definitions in effect, which means that shorthands such as $\$ or $\$ are not available; instead one should use $\$ iow_char:N $\$ and $\$ iow_newline:, respectively. The following message classes exist:

- fatal, ending the TeX run;
- critical, ending the file being input;
- error, interrupting the TEX run without ending it;
- warning, written to terminal and log file, for important messages that may require corrections by the user;
- note (less common than info) for important information messages written to the terminal and log file;
- info for normal information messages written to the log file only;
- term and log for un-decorated messages written to the terminal and log file, or to the log file only;
- none for suppressed messages.

```
\msg_fatal:nnnnnn \msg_fatal:n \\msg_fatal:n \\msg_fatal:nnxxxx \\msg_fatal:nnxxxx \\msg_fatal:nnxxxx \\msg_fatal:nnxxxx \\msg_fatal:nnxxx \\msg_fatal:nnxxx \\msg_fatal:nnxx \\msg_fatal:nnx \\msg_fatal:nnx \\msg_fatal:nnx
```

\msg_fatal:nn

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg\ one \rangle$ to $\langle arg\ four \rangle$ to the text-creating functions. After issuing a fatal error the T_EX run halts. No PDF file will be produced in this case (DVI mode runs may produce a truncated DVI file).

Updated: 2012-08-11

```
\msg_critical:nnnnnn \msg_critical:nsxxx \( \arg four \) \\ \msg_critical:nnxxx \\ \msg_critical:nnxxx \\ \msg_critical:nnnn \\ \msg_critical:nnxx \\ \msg_critical:nnx \\ \msg_c
```

\msg_critical:nn

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg\ one \rangle$ to $\langle arg\ four \rangle$ to the text-creating functions. After issuing a critical error, TEX stops reading the current input file. This may halt the TEX run (if the current file is the main file) or may abort reading a sub-file.

TEXhackers note: The TEX \endinput primitive is used to exit the file. In particular, the rest of the current line remains in the input stream.

```
\msg_error:nnnnnn \msg_error:n \msg_error:n \msg_error:n \msg_error:n \msg_error:n \msg_error:n \msg_error:nnnn \msg_error:nnnn \msg_error:nnnx \msg_error:nnn \msg_error:nnn \msg_error:nnn \msg_error:nnn \msg_error:nnn
```

Updated: 2012-08-11

Updated: 2012-08-11

 $\label{lem:msg_error:nnnnn} $$\max_{error:nnnnn} {\langle module \rangle} {\langle message \rangle} {\langle arg one \rangle} {\langle arg two \rangle} {\langle arg three \rangle} $$\mbox{$| msg_error:nnxxx} {\langle arg four \rangle}$$

Issues $\langle module \rangle$ error $\langle message \rangle$, passing $\langle arg\ one \rangle$ to $\langle arg\ four \rangle$ to the text-creating functions. The error interrupts processing and issues the text at the terminal. After user input, the run continues.

 $\label{lem:msg_warning:nnnnnn} $$\max_{\operatorname{module}} {\langle \operatorname{module} \rangle} {\langle \operatorname{message} \rangle} {\langle \operatorname{arg one} \rangle} {\langle \operatorname{arg two} \rangle} {\langle \operatorname{arg three} \rangle} $$\mbox{ } {\langle \operatorname{arg four} \rangle} $$$

Issues $\langle module \rangle$ warning $\langle message \rangle$, passing $\langle arg\ one \rangle$ to $\langle arg\ four \rangle$ to the text-creating functions. The warning text is added to the log file and the terminal, but the TEX run is not interrupted.

Updated: 2012-08-11

```
\label{localization} $$\max_{n\to\infty} \operatorname{localize} {\langle message \rangle} {\langle arg one \rangle} {\langle arg two \rangle} {\langle arg three \rangle} {\langle arg three \rangle} 
\msg_note:nnxxxx four \}
\msg_note:nnnn
                      \label{eq:continuous_section} $$\max_{i=1,\dots,i=1} {(arg one)} {(arg two)} {(arg three)} {(arg three)} $$
\msg_note:nnxxx
                      four \}
\msg_note:nnnn
                      Issues \langle module \rangle information \langle message \rangle, passing \langle arg\ one \rangle to \langle arg\ four \rangle to the text-creating
\msg_note:nnxx
                      functions. For the more common \msg_info:nnnnn, the information text is added to
\msg_note:nnn
                      the log file only, while \msg_note:nnnnn adds the info text to both the log file and the
\msg note:nnx
                      terminal. The T<sub>E</sub>X run is not interrupted.
\msg note:nn
\msg_info:nnnnn
\msg_info:nnxxxx
\msg_info:nnnnn
\msg_info:nnxxx
\msg_info:nnnn
\msg_info:nnxx
\msg_info:nnn
\msg_info:nnx
\msg_info:nn
      New: 2021-05-18
\label{lem:nnnnn} $$\max_{term:nnnnn} {\mbox{wsg_term:nnnnn} {\mbox{wodule}}} {\mbox{wssage}} {\mbox{arg one}} {\mbox{darg two}} {\mbox{darg three}} {\mbox{darg three}} 
\msg_term:nnxxxx four \}
                      \label{log:nnnnnn} $$ \{\mbox{module}\} $$ {\mbox{message}} $$ {\arg one}$ $$ {\arg two}$ $$ {\arg three}$ $$ $$ $$
\msg_term:nnnnn
                      four \rangle \}
\msg_term:nnxxx
\msg_term:nnnn
                      Issues \langle module \rangle information \langle message \rangle, passing \langle arg\ one \rangle to \langle arg\ four \rangle to the text-creating
\msg_term:nnxx
                      functions. The output is briefer than \msg_info:nnnnn, omitting for instance the mod-
\msg_term:nnn
                      ule name. It is added to the log file by \msg_log:nnnnnn while \msg_term:nnnnn also
\msg_term:nnx
                      prints it on the terminal.
\msg_term:nn
\msg_log:nnnnn
\msg_log:nnxxxx
\msg_log:nnnnn
\msg_log:nnxxx
\msg_log:nnnn
\msg_log:nnxx
\msg_log:nnn
\msg_log:nnx
\msg_log:nn
  Updated: 2012-08-11
\label{localization} $$\max_{none:nnnnn} {\sigma(u)} {\langle uessage \rangle} {\langle arg one \rangle} {\langle arg two \rangle} {\langle arg three \rangle} {\langle arg three \rangle} $$
\msg_none:nnxxxx four \}
\msg_none:nnnnn
                      Does nothing: used as a message class to prevent any output at all (see the discussion of
\msg none:nnxxx
                      message redirection).
\msg_none:nnnn
\msg_none:nnxx
\msg_none:nnn
\msg_none:nnx
\msg_none:nn
```

Updated: 2012-08-11

11.4.1 Messages for showing material

```
\msg_show:nnnnnn \msg_s \forall function \msg_show:nnnx \forall function \msg_show:nnn \msg_show:nnn \msg_show:nnn \msg_show:nnn \msg_show:nnn \msg_show:nnn \msg_show:nn end. In
```

Issues $\langle module \rangle$ information $\langle message \rangle$, passing $\langle arg\ one \rangle$ to $\langle arg\ four \rangle$ to the text-creating functions. The information text is shown on the terminal and the TeX run is interrupted in a manner similar to \tl_show:n. This is used in conjunction with \msg_show_item:n and similar functions to print complex variable contents completely. If the formatted text does not contain >~ at the start of a line, an additional line >~. will be put at the end. In addition, a final period is added if not present.

New: 2017-12-04

11.4.2 Expandable error messages

In very rare cases it may be necessary to produce errors in an expansion-only context. The functions in this section should only be used if there is no alternative approach using \msg_error:nnnnnn or other non-expandable commands from the previous section. Despite having a similar interface as non-expandable messages, expandable errors must be handled internally very differently from normal error messages, as none of the tools to print to the terminal or the log file are expandable. As a result, short-hands such as \{ or \\ do not work, and messages must be very short (with default settings, they are truncated after approximately 50 characters). It is advisable to ensure that the message is understandable even when truncated, by putting the most important information up front. Another particularity of expandable messages is that they cannot be redirected or turned off by the user.

```
\msg_expandable_error:nnnnnn * \msg_expandable_error:nnnnnn {\module\} {\module\} {\arg one\} {\arg on
```

Issues an "Undefined error" message from T_EX itself using the undefined control sequence $\::=$ error then prints "! $\langle module \rangle$: " $\langle error \ message \rangle$, which should be short. With default settings, anything beyond approximately 60 characters long (or bytes in some engines) is cropped. A leading space might be removed as well.

11.5 Redirecting messages

Each message has a "name", which can be used to alter the behaviour of the message when it is given. Thus we might have

```
\msg_new:nnnn { module } { my-message } { Some~text } { Some~more~text }
```

to define a message, with

```
\msg_error:nn { module } { my-message }
```

when it is used. With no filtering, this raises an error. However, we could alter the behaviour with

```
\msg redirect class:nn { error } { warning }
```

to turn all errors into warnings, or with

```
\msg_redirect_module:nnn { module } { error } { warning }
```

to alter only messages from that module, or even

```
\msg redirect name:nnn { module } { my-message } { warning }
```

to target just one message. Redirection applies first to individual messages, then to messages from one module and finally to messages of one class. Thus it is possible to select out an individual message for special treatment even if the entire class is already redirected.

Multiple redirections are possible. Redirections can be cancelled by providing an empty argument for the target class. Redirection to a missing class raises an error immediately. Infinite loops are prevented by eliminating the redirection starting from the target of the redirection that caused the loop to appear. Namely, if redirections are requested as $A \to B$, $B \to C$ and $C \to A$ in this order, then the $A \to B$ redirection is cancelled.

 $\mbox{\sc msg_redirect_class:nn } {\c class:nn }$

Updated: 2012-04-27 Changes the behaviour of messages of (class one) so that they are processed using the code for those of $\langle class \ two \rangle$. Each $\langle class \rangle$ can be one of fatal, critical, error, warning, note, info, term, log, none.

 $\mbox{\mbox{$\mbox{$\mbox{$\mbox{$\mbox{$\mbox{$}\mb$

Updated: 2012-04-27 Redirects message of $\langle class \ one \rangle$ for $\langle module \rangle$ to act as though they were from $\langle class \ one \rangle$ two). Messages of $\langle class\ one \rangle$ from sources other than $\langle module \rangle$ are not affected by this redirection. This function can be used to make some messages "silent" by default. For example, all of the warning messages of $\langle module \rangle$ could be turned off with:

```
\msg_redirect_module:nnn { module } { warning } { none }
```

```
\msg_redirect_name:nnn \msg_redirect_name:nnn {\module\} {\message\} {\langle class\}
```

Updated: 2012-04-27 Redirects a specific $\langle message \rangle$ from a specific $\langle module \rangle$ to act as a member of $\langle class \rangle$ of messages. No further redirection is performed. This function can be used to make a selected message "silent" without changing global parameters:

```
\msg_redirect_name:nnn { module } { annoying-message } { none }
```

Chapter 12

The I3file package File and I/O operations

This module provides functions for working with external files. Some of these functions apply to an entire file, and have prefix \file_..., while others are used to work with files on a line by line basis and have prefix \ior_... (reading) or \iow_... (writing).

It is important to remember that when reading external files TEX attempts to locate them using both the operating system path and entries in the TEX file database (most TEX systems use such a database). Thus the "current path" for TEX is somewhat broader than that for other programs.

For functions which expect a $\langle file\ name \rangle$ argument, this argument may contain both literal items and expandable content, which should on full expansion be the desired file name. Active characters (as declared in $\lower_{charactive_seq}$) are not expanded, allowing the direct use of these in file names. Quote tokens (") are not permitted in file names as they are reserved for internal use by some T_{EX} primitives.

Spaces are trimmed at the beginning and end of the file name: this reflects the fact that some file systems do not allow or interact unpredictably with spaces in these positions. When no extension is given, this will trim spaces from the start of the name only.

12.1 Input-output stream management

As TEX engines have a limited number of input and output streams, direct use of the streams by the programmer is not supported in LATEX3. Instead, an internal pool of streams is maintained, and these are allocated and deallocated as needed by other modules. As a result, the programmer should close streams when they are no longer needed, to release them for other processes.

Note that I/O operations are global: streams should all be declared with global names and treated accordingly.

\ior_new:N \ior_new:N \(stream \) \ior_new:c $\operatorname{ imes iow_new:N}\ \langle \mathit{stream}\
angle$

\iow_new:N \iow_new:c

Globally reserves the name of the $\langle stream \rangle$, either for reading or for writing as appropriate. The \(\stream\) is not opened until the appropriate \\\\..__open:\(\mathbb{Nn}\) function is used. New: 2011-09-26 Attempting to use a $\langle stream \rangle$ which has not been opened is an error, and the $\langle stream \rangle$

Updated: 2011-12-27

will behave as the corresponding \c_term_....

\ior_open:Nn \ior_open:cn $ior_open:Nn \langle stream \rangle \{\langle file name \rangle\}$

Updated: 2012-02-10

Opens $\langle file\ name \rangle$ for reading using $\langle stream \rangle$ as the control sequence for file access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The $\langle stream \rangle$ is available for access immediately and will remain allocated to \(file name \) until a \ior_close: N instruction is given or the TFX run ends. If the file is not found, an error is raised.

\ior_open:NnTF \ior_open:cn<u>TF</u>

 $\colon{constraint} $$ \operatorname{NnTF} \langle \operatorname{stream} \rangle \ \{\langle \operatorname{file name} \rangle\} \ \{\langle \operatorname{true code} \rangle\} \ \{\langle \operatorname{false code} \rangle\} $$$

New: 2013-01-12

Opens $\langle file\ name \rangle$ for reading using $\langle stream \rangle$ as the control sequence for file access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The $\langle stream \rangle$ is available for access immediately and will remain allocated to \(file name \) until a \ior_close: N instruction is given or the T_{FX} run ends. The $\langle true\ code \rangle$ is then inserted into the input stream. If the file is not found, no error is raised and the $\langle false\ code \rangle$ is inserted

into the input stream.

\iow_open:Nn \iow_open:cn $\iow_{pen:Nn \ (stream) \ \{(file name)\}\}$

Updated: 2012-02-09

Opens $\langle file\ name \rangle$ for writing using $\langle stream \rangle$ as the control sequence for file access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The $\langle stream \rangle$ is available for access immediately and will remain allocated to (file name) until a \iow close: N instruction is given or the TFX run ends. Opening a file for writing clears any existing content in the file (i.e. writing is not additive).

\ior_close:N \ior_close:N \(stream \) \ior_close:c \iow_close:N \(stream \)

\iow_close:N \iow_close:c

Closes the (stream). Streams should always be closed when they are finished with as this ensures that they remain available to other programmers.

Updated: 2012-07-31

\ior_show:N \ior_show:N \stream> \ior_show:c \ior_log:N \stream\

\ior_log:N \iow_show:N \(\langle stream \rangle \) \ior_log:c \iow_log:N \(stream \)

\iow_show:N \iow_show:c

Display (to the terminal or log file) the file name associated to the (read or write) $\langle stream \rangle$.

\iow_log:N \iow_log:c

New: 2021-05-11

```
\ior_show_list: \ior_show_list:
\ior_log_list:
                \ior_log_list:
\iow_show_list: \iow_show_list:
\iow_log_list: \iow_log_list:
```

New: 2017-06-27 Display (to the terminal or log file) a list of the file names associated with each open (read or write) stream. This is intended for tracking down problems.

Reading from files

Reading from files and reading from the terminal are separate processes in expl3. The functions \ior_get:NN and \ior_str_get:NN, and their branching equivalents, are designed to work with files.

```
\ior_get:NN
\ior_get:NNTF
```

```
\ior_get:NN \( \stream \) \( \taken list variable \)
\ior_get:NNTF \(\stream\) \(\tau\) \(\tau\) \(\tau\) \(\tau\) \(\tau\) \(\tau\) \(\tau\) \(\tau\) \(\tau\)
```

New: 2012-06-24 Function that reads one or more lines (until an equal number of left and right braces are Updated: 2019-03-23 found) from the file input $\langle stream \rangle$ and stores the result locally in the $\langle token\ list \rangle$ variable. The material read from the $\langle stream \rangle$ is tokenized by T_FX according to the category codes and \endlinechar in force when the function is used. Assuming normal settings, any lines which do not end in a comment character % have the line ending converted to a space, so for example input

```
ab c
```

results in a token list $a_{\sqcup}b_{\sqcup}c_{\sqcup}$. Any blank line is converted to the token \par. Therefore, blank lines can be skipped by using a test such as

```
\ior_get:NN \l_my_stream \l_tmpa_tl
\tl_set:Nn \l_tmpb_tl { \par }
\tl_if_eq:NNF \l_tmpa_tl \l_tmpb_tl
```

Also notice that if multiple lines are read to match braces then the resulting token list can contain \par tokens. In the non-branching version, where the $\langle stream \rangle$ is not open the $\langle tl \ var \rangle$ is set to \q_no_value.

TEXhackers note: This protected macro is a wrapper around the TEX primitive \read. Regardless of settings, T_FX replaces trailing space and tab characters (character codes 32 and 9) in each line by an end-of-line character (character code \endlinechar, omitted if \endlinechar is negative or too large) before turning characters into tokens according to current category codes. With default settings, spaces appearing at the beginning of lines are also ignored.

```
\ior_str_get:NN
```

```
\ior_str_get:NN \( \stream \) \( \token list variable \)
\operatorname{list}_{\operatorname{get}:\operatorname{NN}\overline{TF}} \operatorname{lior\_str\_get}:\operatorname{NNTF} \langle \operatorname{stream} \rangle \langle \operatorname{token} \operatorname{list} \operatorname{variable} \rangle \langle \operatorname{true} \operatorname{code} \rangle \langle \operatorname{false} \operatorname{code} \rangle
```

New: 2016-12-04 Function that reads one line from the file input $\langle stream \rangle$ and stores the result locally in Updated: 2019-03-23 the $\langle token \ list \rangle$ variable. The material is read from the $\langle stream \rangle$ as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). Multiple whitespace characters are retained by this process. It always only reads one line and any blank lines in the input result in the $\langle token\ list\ variable \rangle$ being empty. Unlike \ior_get:NN, line ends do not receive any special treatment. Thus input

ab c

results in a token list a b c with the letters a, b, and c having category code 12. In the non-branching version, where the $\langle stream \rangle$ is not open the $\langle tl \ var \rangle$ is set to \q_no_value.

TeXhackers note: This protected macro is a wrapper around the ε -TeX primitive \readline. Regardless of settings, TFX removes trailing space and tab characters (character codes 32 and 9). However, the end-line character normally added by this primitive is not included in the result of \ior_str_get:NN.

All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

\ior_map_inline:Nn \ior_map_inline:Nn \stream \ {\langle inline function \}}

New: 2012-02-11 Applies the (inline function) to each set of (lines) obtained by calling \ior get:NN until reaching the end of the file. TEX ignores any trailing new-line marker from the file it reads. The $\langle inline\ function \rangle$ should consist of code which receives the $\langle line \rangle$ as #1.

\ior_str_map_inline:Nn \ior_str_map_inline:Nn \stream \ {\langle inline function \}

New: 2012-02-11 Applies the $\langle inline\ function \rangle$ to every $\langle line \rangle$ in the $\langle stream \rangle$. The material is read from the (stream) as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). The (inline function) should consist of code which receives the (line) as #1. Note that TFX removes trailing space and tab characters (character codes 32 and 9) from every line upon input. TEX also ignores any trailing new-line marker from the file it reads.

 $\operatorname{ior_map_variable:NNn \setminus ior_map_variable:NNn } \langle stream \rangle \langle tl \ var
angle \ \{\langle code
angle\}\}$

New: 2019-01-13 For each set of $\langle lines \rangle$ obtained by calling \ior_get:NN until reaching the end of the file, stores the $\langle lines \rangle$ in the $\langle tl \ var \rangle$ then applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last set of $\langle lines \rangle$, or its original value if the $\langle stream \rangle$ is empty. T_FX ignores any trailing new-line marker from the file it reads. This function is typically faster than \ior_map_inline:Nn.

```
\ior_str_map_variable:NNn \ior_str_map_variable:NNn \stream \rangle \text{variable} \{\langle code \rangle \}
```

New: 2019-01-13 For each $\langle line \rangle$ in the $\langle stream \rangle$, stores the $\langle line \rangle$ in the $\langle variable \rangle$ then applies the $\langle code \rangle$. The material is read from the $\langle stream \rangle$ as a series of tokens with category code 12 (other), with the exception of space characters which are given category code 10 (space). The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle line \rangle$, or its original value if the (stream) is empty. Note that TFX removes trailing space and tab characters (character codes 32 and 9) from every line upon input. TEX also ignores any trailing new-line marker from the file it reads. This function is typically faster than \ior_str_map_inline:Nn.

\ior_map_break: \ior_map_break:

New: 2012-06-29 Used to terminate a \ior_map_... function before all lines from the \(stream \) have been processed. This normally takes place within a conditional statement, for example

```
\ior_map_inline:Nn \l_my_ior
    \str_if_eq:nnTF { #1 } { bingo }
      { \ior_map_break: }
        % Do something useful
      }
```

Use outside of a \ior_map_... scenario leads to low level T_FX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

 $ior_map_break:n ior_map_break:n {\langle code \rangle}$

New: 2012-06-29 Used to terminate a \ior_map_... function before all lines in the \(stream \) have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\ior_map_inline:Nn \l_my_ior
    \str if eq:nnTF { #1 } { bingo }
      { \ior_map_break:n { <code> } }
        % Do something useful
  }
```

Use outside of a \ior_map_... scenario leads to low level TEX errors.

TEXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

```
\ior_if_eof_p:N * \ior_if_eof_p:N \( stream \)
```

Updated: 2012-02-10 Tests if the end of a file $\langle stream \rangle$ has been reached during a reading operation. The test also returns a true value if the $\langle stream \rangle$ is not open.

Writing to files 12.1.2

\iow_now:Nn \iow_now:(Nx|cn|cx) $\iow_now:Nn \langle stream \rangle \{\langle tokens \rangle\}$

This functions writes $\langle tokens \rangle$ to the specified $\langle stream \rangle$ immediately (i.e. the write oper-Updated: 2012-06-05 ation is called on expansion of \iow_now:Nn).

 $\verb|\iow_log:n \iow_log:n {$\langle tokens \rangle$}|$

 $\frac{\log x}{\log x}$ This function writes the given $\langle tokens \rangle$ to the log (transcript) file immediately: it is a dedicated version of \iow_now:Nn.

 $\verb|\iow_term:n \iow_term:n {$\langle tokens \rangle$}|$

\iow_term:x

This function writes the given $\langle tokens \rangle$ to the terminal file immediately: it is a dedicated version of \iow_now:Nn.

\iow_shipout:Nn

\iow_shipout:(Nx|cn|cx)

 $\iow_shipout:Nn \langle stream \rangle \{\langle tokens \rangle\}$

This functions writes $\langle tokens \rangle$ to the specified $\langle stream \rangle$ when the current page is finalised (i.e. at shipout). The x-type variants expand the $\langle tokens \rangle$ at the point where the function is used but not when the resulting tokens are written to the $\langle stream \rangle$ (cf. \iow_shipout_x:Nn).

TEXhackers note: When using expl3 with a format other than IATEX, new line characters inserted using \iow_newline: or using the line-wrapping code \iow_wrap:nnnN are not recognized in the argument of \iow_shipout:Nn. This may lead to the insertion of additional unwanted line-breaks.

\iow_shipout_x:Nn \iow_shipout_x:(Nx|cn|cx) $\in \sl \$ \iow_shipout_x:Nn \(stream \) \(\{ tokens \) \}

Updated: 2012-09-08

This functions writes $\langle tokens \rangle$ to the specified $\langle stream \rangle$ when the current page is finalised (i.e. at shipout). The $\langle tokens \rangle$ are expanded at the time of writing in addition to any expansion when the function is used. This makes these functions suitable for including material finalised during the page building process (such as the page number integer).

TEX hackers note: This is a wrapper around the TEX primitive \write. When using expl3 with a format other than LATEX, new line characters inserted using \iow_newline: or using the line-wrapping code \iow_wrap:nnnN are not recognized in the argument of \iow_shipout:Nn. This may lead to the insertion of additional unwanted line-breaks.

$\in \color=0.000$ \iow_char:N $\in \color=0.000$

Inserts $\langle char \rangle$ into the output stream. Useful when trying to write difficult characters such as %, $\{$, $\}$, etc. in messages, for example:

```
\iow_now:Nx \g_my_iow { \iow_char:N \{ text \iow_char:N \} }
```

The function has no effect if writing is taking place without expansion (e.g. in the second argument of $\iow_now:Nn$).

\iow_newline: * \iow_newline:

Function to add a new line within the $\langle tokens \rangle$ written to a file. The function has no effect if writing is taking place without expansion (e.g. in the second argument of \iow_-now:Nn).

TEXhackers note: When using expl3 with a format other than LATEX, the character inserted by \iow_newline: is not recognized by TEX, which may lead to the insertion of additional unwanted line-breaks. This issue only affects \iow_shipout:Nn, \iow_shipout_x:Nn and direct uses of primitive operations.

12.1.3 Wrapping lines in output

\iow_wrap:nnnN \iow_wrap:nxnN $\label{low_wrap:nnnN} $$ \{\langle \text{run-on text} \rangle\} $$ {\langle \text{set up} \rangle} $$ \langle \text{function} \rangle$$$

New: 2012-06-28 Updated: 2017-12-04

This function wraps the $\langle text \rangle$ to a fixed number of characters per line. At the start of each line which is wrapped, the $\langle run\text{-}on \ text \rangle$ is inserted. The line character count targeted is the value of \l iow line count int minus the number of characters in the $\langle run\text{-}on \ text \rangle$ for all lines except the first, for which the target number of characters is simply $l_{iow_{int}}$ and run-on text. The $\langle text \rangle$ and run-ontext are exhaustively expanded by the function, with the following substitutions:

- \\ or \iow_newline: may be used to force a new line,
- _ may be used to represent a forced space (for example after a control sequence),
- $\$, $\$, $\$, $\$, $\$ may be used to represent the corresponding character,
- \iow_allow_break: may be used to allow a line-break without inserting a space (this is experimental),
- \iow_indent:n may be used to indent a part of the $\langle text \rangle$ (not the $\langle run\text{-}on\ text \rangle$).

Additional functions may be added to the wrapping by using the $\langle set\ up \rangle$, which is executed before the wrapping takes place: this may include overriding the substitutions listed.

Any expandable material in the $\langle text \rangle$ which is not to be expanded on wrapping should be converted to a string using \token_to_str:N, \tl_to_str:N, \tl_to_str:N, etc.

The result of the wrapping operation is passed as a braced argument to the (function), which is typically a wrapper around a write operation. The output of \iow wrap:nnnN (i.e. the argument passed to the $\langle function \rangle$) consists of characters of category "other" (category code 12), with the exception of spaces which have category "space" (category code 10). This means that the output does not expand further when written to a file.

TEXhackers note: Internally, \iow_wrap:nnnN carries out an x-type expansion on the $\langle \textit{text} \rangle$ to expand it. This is done in such a way that $\ensuremath{\texttt{\ensuremath{\texttt{exp_not:}}}} \ensuremath{\texttt{\ensuremath{\texttt{exp_not:}}}} \ensuremath{\texttt{\ensuremath{\texttt{could}}}} \ensuremath{\texttt{\ensuremath{\texttt{exp_not:}}}} \ensuremath{\texttt{\ensu$ to prevent expansion of material. However, this is less conceptually clear than conversion to a string, which is therefore the supported method for handling expandable material in the $\langle text \rangle$.

 $\iow_indent:n \iow_indent:n {\langle text \rangle}$

New: 2011-09-21 In the first argument of \iow_wrap:nnnN (for instance in messages), indents \langle text \rangle by four spaces. This function does not cause a line break, and only affects lines which start within the scope of the $\langle text \rangle$. In case the indented $\langle text \rangle$ should appear on separate lines from the surrounding text, use $\$ to force line breaks.

\l_iow_line_count_int The maximum number of characters in a line to be written by the \iow_wrap:nnnN New: 2012-06-24 function. This value depends on the TEX system in use: the standard value is 78, which is typically correct for unmodified TFX Live and MiKTFX systems.

12.1.4Constant input-output streams, and variables

\g_tmpb_ior

\g_tmpa_ior Scratch input stream for global use. These are never used by the kernel code, and so are safe for use with any LATEX3-defined function. However, they may be overwritten by $_{\text{New: }2017-12-11}$ other non-kernel code and so should only be used for short-term storage.

\c_log_iow \c_term_iow

Constant output streams for writing to the log and to the terminal (plus the log), respec-

\g_tmpb_iow

\g_tmpa_iow Scratch output stream for global use. These are never used by the kernel code, and so are safe for use with any LATEX3-defined function. However, they may be overwritten by New: 2017-12-11 other non-kernel code and so should only be used for short-term storage.

12.1.5 Primitive conditionals

\if eof:w *

\if_eof:w \(stream \) ⟨true code⟩ \else: ⟨false code⟩

Tests if the $\langle stream \rangle$ returns "end of file", which is true for non-existent files. The **\else**: branch is optional.

TEXhackers note: This is the TEX primitive \ifeof.

12.2 File operation functions

\g_file_curr_dir_str \g_file_curr_ext_str

Contain the directory, name and extension of the current file. The directory is empty if \g_file_curr_name_str the file was loaded without an explicit path (i.e. if it is in the TEX search path), and does not end in / other than the case that it is exactly equal to the root directory. The $\langle name \rangle$ New: 2017-06-21 and $\langle ext \rangle$ parts together make up the file name, thus the $\langle name \rangle$ part may be thought of as the "job name" for the current file. Note that TFX does not provide information on the $\langle ext \rangle$ part for the main (top level) file and that this file always has an empty $\langle dir \rangle$ component. Also, the (name) here will be equal to \c sys jobname str, which may be different from the real file name (if set using --jobname, for example).

\l_file_search_path_seq Each entry is the path to a directory which should be searched when seeking a file. Each New: 2017-06-18 path can be relative or absolute, and should not include the trailing slash. The entries are not expanded when used so may contain active characters but should not feature any variable content. Spaces need not be quoted.

> **TeXhackers note:** When working as a package in E^{T} EX 2ε , expl3 will automatically append the current \input@path to the set of values from \l_file_search_path_seq.

 $file_if_exist:nTF \file_if_exist:nTF \{\langle file name \rangle\} \{\langle true code \rangle\} \{\langle false code \rangle\}$

Updated: 2012-02-10 Searches for (file name) using the current TFX search path and the additional paths controlled by \l_file_search_path_seq.

\file_get:nnN \file_get:nnNTF

 $file_get:nnN {\langle filename \rangle} {\langle setup \rangle} {\langle t1 \rangle}$ $file_get:nnNTF \{\langle filename \rangle\} \{\langle setup \rangle\} \langle t1 \rangle \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}$

New: 2019-01-16 Defines $\langle tl \rangle$ to the contents of $\langle filename \rangle$. Category codes may need to be set appropri-Updated: 2019-02-16 ately via the $\langle setup \rangle$ argument. The non-branching version sets the $\langle tl \rangle$ to q_no_value if the file is not found. The branching version runs the $\langle true\ code \rangle$ after the assignment to $\langle tl \rangle$ if the file is found, and $\langle false\ code \rangle$ otherwise.

\file_get_full_name:nN \file_get_full_name:VN \file_get_full_name:nNTF $file_get_full_name:nN {\langle file name \rangle} \langle t1 \rangle$ $\label{lem:norm} $$ \left(\text{file name} \right) \ \langle t1 \rangle \ \left(\text{true code} \right) \ \left(\text{false code$

\file_get_full_name:VNTF

Updated: 2019-02-16

Searches for \(file name \) in the path as detailed for \file_if_exist:nTF, and if found sets the $\langle tl \ var \rangle$ the fully-qualified name of the file, i.e. the path and file name. This includes an extension .tex when the given $\langle file\ name \rangle$ has no extension but the file found has that extension. In the non-branching version, the $\langle tl \ var \rangle$ will be set to \q_no_value in the case that the file does not exist.

\file_full_name:n \file_full_name:V ☆

\file_full_name:n {\langle file name \rangle}

Searches for (file name) in the path as detailed for \file if exist:nTF, and if found New: 2019-09-03 leaves the fully-qualified name of the file, i.e. the path and file name, in the input stream. This includes an extension .tex when the given \(\frac{file name}{} \) has no extension but the file found has that extension. If the file is not found on the path, the expansion is empty.

```
\file_parse_full_name:VNNN
```

 $file_parse_full_name:nNNN \file_parse_full_name:nNNN \{\langle full\ name
angle\} \ \langle dir
angle \ \langle name
angle \ \langle ext
angle$

Updated: 2020-06-24

Parses the \(\frac{full name}\) and splits it into three parts, each of which is returned by setting New: 2017-06-23 the appropriate local string variable:

- The $\langle dir \rangle$: everything up to the last / (path separator) in the $\langle file\ path \rangle$. As with system PATH variables and related functions, the $\langle dir \rangle$ does not include the trailing / unless it points to the root directory. If there is no path (only a file name), $\langle dir \rangle$
- The $\langle name \rangle$: everything after the last / up to the last ., where both of those characters are optional. The $\langle name \rangle$ may contain multiple . characters. It is empty if $\langle full\ name \rangle$ consists only of a directory name.
- The $\langle ext \rangle$: everything after the last . (including the dot). The $\langle ext \rangle$ is empty if there is no . after the last /.

Before parsing, the $\langle full\ name \rangle$ is expanded until only non-expandable tokens remain, except that active characters are also not expanded. Quotes (") are invalid in file names and are discarded from the input.

New: 2020-06-24 Parses the $\langle full\ name \rangle$ as described for \file_parse_full_name: nNNN, and leaves $\langle dir \rangle$, $\langle name \rangle$, and $\langle ext \rangle$ in the input stream, each inside a pair of braces.

```
\file_parse_full_name_apply:nN * \file_parse_full_name_apply:nN {\( full name \) \\ \( function \)
                       New: 2020-06-24
```

Parses the $\langle full\ name \rangle$ as described for \file_parse_full_name:nNNN, and passes $\langle dir \rangle$, $\langle name \rangle$, and $\langle ext \rangle$ as arguments to $\langle function \rangle$, as an n-type argument each, in this order.

```
\file_hex_dump:n
\file_hex_dump:nnn ☆
```

☆ \file_hex_dump:n {⟨file name⟩}

 $file_hex_dump:nnn {\langle file name \rangle} {\langle start index \rangle} {\langle end index \rangle}$

New: 2019-11-19 Searches for (file name) using the current TEX search path and the additional paths dump of the file content in the input stream. The file is read as bytes, which means that in contrast to most T_FX behaviour there will be a difference in result depending on the line endings used in text files. The same file will produce the same result between different engines: the algorithm used is the same in all cases. When the file is not found, the result of expansion is empty. The $\{\langle start\ index \rangle\}$ and $\{\langle end\ index \rangle\}$ values work as described for \str_range:nnn.

```
\file_get_hex_dump:nN
\file_get_hex_dump:nNTF
\file_get_hex_dump:nnnN
\file_get_hex_dump:nnnNTF
```

```
\file_get_hex_dump:nN {\langle file name \rangle} \langle t1 var \rangle
```

 $file_{det_{name}} = file_{det_{name}} \{ (start index) \} \{ (end index) \} \langle t1 var \rangle$

Sets the $\langle tl \ var \rangle$ to the result of applying \file_hex_dump:n/\file_hex_dump:nnn to the $\langle file \rangle$. If the file is not found, the $\langle tl \ var \rangle$ will be set to \q_no_value.

New: 2019-11-19

 $file_mdfive_hash:n \Leftrightarrow file_mdfive_hash:n {\langle file_name \rangle}$

New: 2019-09-03 Searches for \(file \ name \) using the current TFX search path and the additional paths controlled by \l_file_search_path_seq. It then expands to leave the MD5 sum generated from the contents of the file in the input stream. The file is read as bytes, which means that in contrast to most T_FX behaviour there will be a difference in result depending on the line endings used in text files. The same file will produce the same result between different engines: the algorithm used is the same in all cases. When the file is not found, the result of expansion is empty.

\file_get_mdfive_hash:nN

\file_get_mdfive_hash:nN {\file name\} \langle tl var \rangle

\file_get_mdfive_hash:nNTF

Sets the $\langle tl \ var \rangle$ to the result of applying \file_mdfive_hash:n to the $\langle file \rangle$. If the file New: 2017-07-11 is not found, the $\langle tl \ var \rangle$ will be set to \q_no_value.

Updated: 2019-02-16

\file_size:n ☆ \file_size:n {\langle file name \rangle}

New: 2019-09-03 Searches for (file name) using the current TEX search path and the additional paths controlled by \l_file_search_path_seq. It then expands to leave the size of the file in bytes in the input stream. When the file is not found, the result of expansion is empty.

\file_get_size:nN \file_get_size:nNTF

 $file_get_size:nN {\langle file name \rangle} \langle tl var \rangle$

Sets the $\langle tl \ var \rangle$ to the result of applying \file_size:n to the $\langle file \rangle$. If the file is not New: 2017-07-09 found, the $\langle tl \ var \rangle$ will be set to \q_no_value . This is not available in older versions of Updated: 2019-02-16 XATEX.

 $file_timestamp:n \Leftrightarrow file_timestamp:n {\langle file name \rangle}$

New: 2019-09-03 Searches for $\langle file\ name \rangle$ using the current TFX search path and the additional paths controlled by \l_file_search_path_seq. It then expands to leave the modification timestamp of the file in the input stream. The timestamp is of the form $D: \langle year \rangle \langle month \rangle \langle day \rangle \langle hour \rangle \langle minute \rangle \langle second \rangle \langle offset \rangle$, where the latter may be Z (UTC) or $\langle plus-minus\rangle \langle hours\rangle \langle minutes\rangle$. When the file is not found, the result of expansion is empty. This is not available in older versions of XATEX.

\file_get_timestamp:nN \file_get_timestamp:nN<u>TF</u>

 $file_get_timestamp:nN {\langle file name \rangle} \langle tl var \rangle$

New: 2017-07-09 Updated: 2019-02-16

Sets the $\langle tl \ var \rangle$ to the result of applying \file_timestamp:n to the $\langle file \rangle$. If the file is not found, the $\langle tl \ var \rangle$ will be set to \q no value. This is not available in older versions of X₇T_FX.

```
file\_compare\_timestamp\_p:nNn * file\_compare\_timestamp:nNn {\file-1\} {\comparator} {\cap file-2\} {\cap true}
file\_compare\_timestamp:nNn\underline{TF} \star code {false\ code}
```

New: 2019-05-13 Updated: 2019-09-20

> Compares the file stamps on the two $\langle files \rangle$ as indicated by the $\langle comparator \rangle$, and inserts either the $\langle true\ code \rangle$ or $\langle false\ case \rangle$ as required. A file which is not found is treated as older than any file which is found. This allows for example the construct

```
\file_compare_timestamp:nNnT { source-file } > { derived-file }
     \mbox{\ensuremath{\mbox{\%}}} Code to regenerate derived file
```

to work when the derived file is entirely absent. The timestamp of two absent files is regarded as different. This is not available in older versions of X₇T_FX.

\file_input:n

\file_input:n $\{\langle file name \rangle\}$

Updated: 2017-06-26 Searches for \(file name \) in the path as detailed for \file_if_exist:nTF, and if found reads in the file as additional LATEX source. All files read are recorded for information and the file name stack is updated by this function. An error is raised if the file is not found.

```
\file_if_exist_input:n
```

```
\verb|\file_if_exist_input:n \{ \langle \textit{file name} \rangle \}|
\label{linear_stable} $$ \left( i_{exist_input:nF} \left( i_{exist_input
```

New: 2014-07-02 Searches for (file name) using the current TFX search path and the additional paths included in \l_file_search_path_seq. If found then reads in the file as additional ETeX source as described for $file_input:n$, otherwise inserts the $(false\ code)$. Note that these functions do not raise an error if the file is not found, in contrast to \file_input:n.

\file_input_stop: \file_input_stop:

New: 2017-07-07 Ends the reading of a file started by \file_input:n or similar before the end of the file is reached. Where the file reading is being terminated due to an error, \msg_critical:nn(nn) should be preferred.

> TEXhackers note: This function must be used on a line on its own: TEX reads files line-by-line and so any additional tokens in the "current" line will still be read.

> This is also true if the function is hidden inside another function (which will be the normal case), i.e., all tokens on the same line in the source file are still processed. Putting it on a line by itself in the definition doesn't help as it is the line where it is used that counts!

\file_show_list: \file_show_list: \file_log_list:

\file_log_list:

These functions list all files loaded by $\LaTeX 2_{\mathcal{E}}$ commands that populate \complement by \file_input:n. While \file_show_list: displays the list in the terminal, \file_log_list: outputs it to the log file only.

Chapter 13

The I3luatex package: LuaT_FX-specific functions

The LuaTEX engine provides access to the Lua programming language, and with it access to the "internals" of T_EX. In order to use this within the framework provided here, a family of functions is available. When used with pdfTFX, pTFX, upTFX or XFTFX these raise an error: use \sys_if_engine_luatex:T to avoid this. Details on using Lua with the LuaTEX engine are given in the LuaTEX manual.

13.1Breaking out to Lua

 $\displaystyle \sum_{now:n} \star \displaystyle (\langle token \ list \rangle)$

 $\frac{\text{lua_now:e} \star}{\text{The } \langle token \ list \rangle}$ is first tokenized by TeX, which includes converting line ends to spaces in New: 2018-06-18 the usual TeX manner and which respects currently-applicable TeX category codes. The resulting $\langle Lua\ input \rangle$ is passed to the Lua interpreter for processing. Each \lua_now:n block is treated by Lua as a separate chunk. The Lua interpreter executes the $\langle Lua \rangle$ *input*\(\right\) immediately, and in an expandable manner.

> TrXhackers note: \lua_now:e is a macro wrapper around \directlua: when LuaTrX is in use two expansions are required to yield the result of the Lua code.

\lua_shipout:n

The (token list) is first tokenized by T_FX, which includes converting line ends to spaces in $_{\mbox{\scriptsize New: 2018-06-18}}$ the usual TeX manner and which respects currently-applicable TeX category codes. The resulting $\langle Lua\ input \rangle$ is passed to the Lua interpreter when the current page is finalised (i.e. at shipout). Each \lua_shipout:n block is treated by Lua as a separate chunk. The Lua interpreter will execute the $\langle Lua\ input \rangle$ during the page-building routine: no TeX expansion of the $\langle Lua\ input \rangle$ will occur at this stage.

> In the case of the \lua_shipout_e:n version the input is fully expanded by TFX in an e-type manner during the shipout operation.

TEXhackers note: At a TeX level, the \(Lua input \) is stored as a "whatsit".

```
\lua_escape:e
```

 $\lambda = \frac{\star \langle token \ list \rangle}{}$

Converts the $\langle token \ list \rangle$ such that it can safely be passed to Lua: embedded backslashes, double and single quotes, and newlines and carriage returns are escaped. This is done by prepending an extra token consisting of a backslash with category code 12, and for the line endings, converting them to \n and \r , respectively.

TrXhackers note: \lua_escape:e is a macro wrapper around \luaescapestring: when LuaTFX is in use two expansions are required to yield the result of the Lua code.

\lua_load_module:n \lua_load_module:n {\(Lua module name \) \}

New: 2022-05-14 Loads a Lua module into the Lua interpreter.

 $\label{lua_now:n} \ passes its {$\langle token \ list\rangle$} \ argument to the Lua interpreter as a single line,$ with characters interpreted under the current catcode regime. These two facts mean that \lua_now:n rarely behaves as expected for larger pieces of code. Therefore, package authors should **not** write significant amounts of Lua code in the arguments to \lua_now:n. Instead, it is strongly recommended that they write the majorty of their Lua code in a separate file, and then load it using \lua_load_module:n.

TEXhackers note: This is a wrapper around the Lua call require $\langle module \rangle$.

13.2 Lua interfaces

As well as interfaces for T_FX, there are a small number of Lua functions provided here.

ltx.utils Most public interfaces provided by the module are stored within the ltx.utils table.

```
ltx.utils.filedump \langle dump \rangle = ltx.utils.filedump(\langle file \rangle, \langle offset \rangle, \langle length \rangle)
```

Returns the uppercase hexadecimal representation of the content of the $\langle file \rangle$ read as bytes. If the $\langle length \rangle$ is given, only this part of the file is returned; similarly, one may specify the $\langle offset \rangle$ from the start of the file. If the $\langle length \rangle$ is not given, the entire file is read starting at the $\langle offset \rangle$.

```
ltx.utils.filemd5sum \langle hash \rangle = ltx.utils.filemd5sum(\langle file \rangle)
```

Returns the MD5 sum of the file contents read as bytes; note that the result will depend on the nature of the line endings used in the file, in contrast to normal TFX behaviour. If the $\langle file \rangle$ is not found, nothing is returned with no error raised.

```
ltx.utils.filemoddate \langle date \rangle = ltx.utils.filemoddate(\langle file \rangle)
```

Returns the date/time of last modification of the $\langle file \rangle$ in the format

```
D: \langle year \rangle \langle month \rangle \langle day \rangle \langle hour \rangle \langle minute \rangle \langle second \rangle \langle offset \rangle
```

where the latter may be Z (UTC) or $\langle plus-minus\rangle \langle hours\rangle$, $\langle minutes\rangle$. If the $\langle file\rangle$ is not found, nothing is returned with no error raised.

 ${\tt ltx.utils.filesize \ size = ltx.utils.filesize(} \langle file \rangle {\tt)}$

Returns the size of the $\langle file \rangle$ in bytes. If the $\langle file \rangle$ is not found, nothing is returned with no error raised.

Chapter 14

The **I3legacy** package Interfaces to legacy concepts

There are a small number of TFX or LATFX 2ε concepts which are not used in expl3 code but which need to be manipulated when working as a LATEX 2ε package. To allow these to be integrated cleanly into expl3 code, a set of legacy interfaces are provided here.

 $\lceil nTF \rceil \star$

 $\label{legacy_if_p:n * legacy_if:nTF {(name)} {(true \ code)} {(false \ code)}}$

Tests if the \LaTeX 2ε /plain TeX conditional (generated by \newif) if true or false and branches accordingly. The $\langle name \rangle$ of the conditional should *omit* the leading if.

\legacy_if_set_true:n \legacy_if_set_false:n \legacy_if_gset_true:n \legacy_if_gset_false:n

 $\lceil \cdot \rceil = \{ \langle name \rangle \}$ $\lceil \lceil \rceil \rceil = \lceil \lceil \rceil \rceil$

Sets the LATEX 2ε /plain TEX conditional \if $\langle name \rangle$ (generated by \newif) to be true

New: 2021-05-10

\legacy_if_gset:nn

 $\lceil \lceil \rceil \rceil = \lceil \rceil \rceil$

Sets the LATEX 2ε /plain TeX conditional \if $\langle name \rangle$ (generated by \newif) to the result New: 2021-05-10 of evaluating the $\langle boolean \ expression \rangle$.

 $\begin{array}{c} {\rm Part~IV} \\ {\bf Data~types} \end{array}$

Chapter 15

The **|3t|** package Token lists

TEX works with tokens, and LATEX3 therefore provides a number of functions to deal with lists of tokens. Token lists may be present directly in the argument to a function:

```
\foo:n { a collection of \tokens }
```

or may be stored in a so-called "token list variable", which have the suffix tl: a token list variable can also be used as the argument to a function, for example

```
\foo:N \l_some_tl
```

In both cases, functions are available to test and manipulate the lists of tokens, and these have the module prefix t1. In many cases, functions which can be applied to token list variables are paired with similar functions for application to explicit lists of tokens: the two "views" of a token list are therefore collected together here.

A token list (explicit, or stored in a variable) can be seen either as a list of "items", or a list of "tokens". An item is whatever $\use:n$ would grab as its argument: a single non-space token or a brace group, with optional leading explicit space characters (each item is thus itself a token list). A token is either a normal N argument, or \sqcup , $\{$, or $\}$ (assuming normal TeX category codes). Thus for example

```
{ Hello } ~ world
```

contains six items (Hello, w, o, r, 1 and d), but thirteen tokens ($\{$, H, e, 1, 1, o, $\}$, \sqcup , w, o, r, 1 and d). Functions which act on items are often faster than their analogue acting directly on tokens.

15.1 Creating and initialising token list variables

 $\verb|\tl_new:N \tl_new:N \ \langle tl \ var \rangle|$

Creates a new $\langle tl \ var \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle tl \ var \rangle$ is initially empty.

```
\t! const:Nn \langle tl var \rangle \{\langle token list \rangle\}
   \tl_const:Nn
   \t: (Nx|cn|cx)
                              Creates a new constant \langle tl \ var \rangle or raises an error if the name is already taken. The value
                              of the \langle tl \ var \rangle is set globally to the \langle token \ list \rangle.
            \tl_clear:N \tl_clear:N \tl var>
            \tl_clear:c
                              Clears all entries from the \langle tl \ var \rangle.
            \tl_gclear:N
            \tl_gclear:c
      \tl_clear_new:N \tl_clear_new:N \tl var>
      \tl_clear_new:c
                              Ensures that the \(\langle tl var \rangle \) exists globally by applying \t1_new:N if necessary, then applies
      \tl_gclear_new:N
                               \t_{g} clear: N to leave the \langle tl \ var \rangle empty.
      \tl_gclear_new:c
                              \t_{set_eq:NN} \langle t1 \ var_1 \rangle \langle t1 \ var_2 \rangle
\tl_set_eq:NN
\tl_set_eq:(cN|Nc|cc)
                              Sets the content of \langle tl \ var_1 \rangle equal to that of \langle tl \ var_2 \rangle.
\tl_gset_eq:NN
\t_gset_eq:(cN|Nc|cc)
       \tl_concat:NNN
                              \t_{concat:NNN} \langle t1 \ var_1 \rangle \langle t1 \ var_2 \rangle \langle t1 \ var_3 \rangle
        \tl_concat:ccc
                               Concatenates the content of \langle tl \ var_2 \rangle and \langle tl \ var_3 \rangle together and saves the result in
        \tl_gconcat:NNN
                              \langle tl \ var_1 \rangle. The \langle tl \ var_2 \rangle is placed at the left side of the new token list.
        \tl_gconcat:ccc
             New: 2012-05-18
   \tilde{tl}_{if}_{exist}_{p:c} \star \tilde{tl}_{if}_{exist}:NTF \langle tl \ var \rangle \{\langle true \ code \rangle\} \{\langle false \ code \rangle\}
   \t^{\text{li_-exist:N}}_{\overline{zz}} * Tests whether the \langle tl \ var \rangle is currently defined. This does not check that the \langle tl \ var \rangle
   \tl_if_exist:c<u>TF</u> *
                              really is a token list variable.
            New: 2012-03-03
```

15.2 Adding data to token list variables

```
\tl_set:Nn \tl_set:Nn \tl_set:Nn \tl_var\ {\langle tokens\}}
\tl_set:(NV|Nv|No|Nf|Nx|cn|cV|cv|co|cf|cx)
\tl_gset:Nn
\tl_gset:(NV|Nv|No|Nf|Nx|cn|cV|cv|co|cf|cx)

Sets \langle tl var\rangle to contain \langle tokens\rangle, removing any previous content from the variable.
\tl_put_left:Nn \tl_put_left:Nn \langle tl var\rangle {\langle tokens\rangle}}
\tl_put_left:(NV|No|Nx|cn|cV|co|cx)
\tl_gput_left:(NV|No|Nx|cn|cV|co|cx)
```

Appends $\langle tokens \rangle$ to the left side of the current content of $\langle tl \ var \rangle$.

```
\tl_put_right:Nn \tl_put_right:Nn \tl_put_right:Nn \tl var \rangle \{\tokens\}\\tl_put_right:Nn \tl_put_right:Nn \\tl_put_right:Nn \\tl_put_right:Nn \\tl_put_right:Nn \\tl_put_right:Nn \\tl_put_right:(NV|No|Nx|cn|cV|co|cx)
```

Appends $\langle tokens \rangle$ to the right side of the current content of $\langle tl \ var \rangle$.

15.3 Token list conditionals

```
\tl_if_blank_p:n
                          \star \til_{if\_blank\_p:n \{\langle token \ list \rangle\}}
                          \tl_if_blank_p:(e|V|o)
\t:nTF
                          ^{\star} Tests if the \langle token\ list \rangle consists only of blank spaces (i.e. contains no item). The test is
\t: (e|V|o) \underline{TF} \star
                            true if \langle token\ list \rangle is zero or more explicit space characters (explicit tokens with character
          Updated: 2019-09-04 code 32 and category code 10), and is false otherwise.
     \verb|\tl_if_empty_p:N * \tl_if_empty_p:N & tl_if_empty_p:N & tl_var||
     \label{limit} $$ \tilde{code} : TF \ \langle tl \ var \rangle \ {\langle true \ code \rangle} \ {\langle false \ code \rangle} $$
     \tl_if_empty:cTF *
                          * \tl_if_empty_p:n {\langle token list \rangle}
 \tl_if_empty_p:n
 \t_if_empty_p:(V|o) * \t_if_empty:nTF {\langle token \ list \rangle} {\langle true \ code \rangle} {\langle false \ code \rangle}
 \tl_if_empty:nTF
                            Tests if the \langle token \ list \rangle is entirely empty (i.e. contains no tokens at all).
 \t_i = \text{V}_i \cdot (V | o) \underline{TF}
             New: 2012-05-24
         Updated: 2012-06-05
\tl_if_eq_p:NN
                          \star \tl_if_eq_p:NN \langle tl \ var_1 \rangle \langle tl \ var_2 \rangle
\t_i = eq_p: (Nc|cN|cc) * \t_i = eq:NNTF \langle tl var_1 \rangle \langle tl var_2 \rangle \{\langle true code \rangle\} \{\langle false code \rangle\}
\tl_if_eq:NNTF
                            Compares the content of two (token list variables) and is logically true if the two contain
\t_i = (Nc|cN|cc)TF
                            the same list of tokens (i.e. identical in both the list of characters they contain and the
                             category codes of those characters). Thus for example
                                  \tl_set:Nn \l_tmpa_tl { abc }
                                  \tl_set:Nx \l_tmpb_tl { \tl_to_str:n { abc } }
                                  \tl_if_eq:NNTF \l_tmpa_tl \l_tmpb_tl { true } { false }
                             yields false. See also \str_if_eq:nnTF for a comparison that ignores category codes.
          \label{limit} $$ \tilde{f}_eq:NnTF \ \langle tl \ var_1\rangle \ \{\langle token \ list_2\rangle\} \ \{\langle true \ code\rangle\} \ \{\langle false \ code\rangle\} $$
```

```
\til_{if_eq:nn}TF \til_{if_eq:nn}TF \{\langle token\ list_1 \rangle\} \{\langle token\ list_2 \rangle\} \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}
```

Tests if $\langle token \ list_1 \rangle$ and $\langle token \ list_2 \rangle$ contain the same list of tokens, both in respect of character codes and category codes. This conditional is not expandable: see \tl_if_eq:NNTF for an expandable version when token lists are stored in variables, or \str_if_eq:nnTF if category codes are not important.

```
\tilde{t}_{in:Nn} T_{in:Nn} T
```

 $\frac{\text{tl_if_in:cn}\underline{TF}}{\text{Tests}}$ Tests if the $\langle token\ list \rangle$ is found in the content of the $\langle tl\ var \rangle$. The $\langle token\ list \rangle$ cannot contain the tokens {, } or # (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

\tl_if_in:nnTF $\til_{if_in}(Vn|on|no)$

```
\tilde{1}_{in:nnTF} {\langle token \ list_1 \rangle} {\langle token \ list_2 \rangle} {\langle true \ code \rangle} {\langle false \ code \rangle}
```

Tests if $\langle token \ list_2 \rangle$ is found inside $\langle token \ list_1 \rangle$. The $\langle token \ list_2 \rangle$ cannot contain the tokens {, } or # (more precisely, explicit character tokens with category code 1 (begingroup) or 2 (end-group), and tokens with category code 6). The search does not enter brace (category code 1/2) groups.

```
\t_i^n = \frac{1}{2} \cdot \frac{1}{
    \til_{if_novalue:nTF} \star \til_{if_novalue:nTF} \{\langle token\ list \rangle\} \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}
```

New: 2017-11-14 Tests if the $\langle token\ list \rangle$ is exactly equal to the special \c_novalue_tl marker. This function is intended to allow construction of flexible document interface structures in which missing optional arguments are detected.

```
\t_i=single_p:N \star tl_if_single_p:N \langle tl var \rangle
\tilde{tl_if_single_p:c} \star \tilde{tl_if_single:NTF} \langle tl var \rangle \{\langle true code \rangle\} \{\langle false code \rangle\}
```

 $\t^{if}_{single:NTF} \star \text{ Tests if the content of the } \langle tl \ var \rangle \text{ consists of a single } \langle item \rangle, \ i.e. \text{ is a single normal token}$

 $\frac{\text{tl_if_single:cTF} \star \text{ (neither an explicit space character nor a begin-group character) or a single brace group,}{\text{transformation of the state of the sta$ Updated: 2011-08-13 surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to \tl count: N.

```
\tilde{\ } \tilde{\ 
\t_if_single:nTF * \t_if_single:nTF {\langle token \ list \rangle} {\langle true \ code \rangle} {\langle false \ code \rangle}
```

 $v_{pdated: 2011-08-13}$ Tests if the $\langle token\ list \rangle$ has exactly one $\langle item \rangle$, i.e. is a single normal token (neither an explicit space character nor a begin-group character) or a single brace group, surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to \tl_count:n.

```
\t_if_single_token_p:n * \t_if_single_token_p:n {$\langle token \ list \rangle$}
```

```
\t = \frac{TF}{single\_token:nTF} + \frac{f(token\ list)}{f(true\ code)}
```

Tests if the token list consists of exactly one token, i.e. is either a single space character or a single normal token. Token groups $(\{...\})$ are not single tokens.

```
\tl_case:NnTF \(\text\) token list variable \(\)
\tl case:Nn
\tl_case:cn
\tl_case:NnTF
                                 \langle token\ list\ variable\ case_1 \rangle\ \{\langle code\ case_1 \rangle\}
\tl_case:cn<u>TF</u>
                                 \langle token\ list\ variable\ case_2 \rangle\ \{\langle code\ case_2 \rangle\}
      New: 2013-07-24
                                 \langle token\ list\ variable\ case_n \rangle\ \{\langle code\ case_n \rangle\}
                             \{\langle true\ code \rangle\}
                             {\false code\}
```

This function compares the $\langle test\ token\ list\ variable \rangle$ in turn with each of the $\langle token\ list\ variable \rangle$ list variable cases. If the two are equal (as described for \tl_if_eq:NNTF) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The function \tl_case: Nn, which does nothing if there is no match, is also available.

Testing the first token 15.3.1

```
\t l_if_head_eq_catcode_p:nN \star \t l_if_head_eq_catcode_p:nN \ \{\langle token\ list \rangle\} \ \langle test\ token \rangle
\verb|\tl_if_head_eq_catcode_p:oN * \verb|\tl_if_head_eq_catcode:nNTF = {(token list)}| (test token)| |
\{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
\tl_if_head_eq_catcode:oN<u>TF</u> *
                   Updated: 2012-07-09
```

Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ has the same category code as the $\langle test | token \rangle$. In the case where the $\langle token \ list \rangle$ is empty, the test is always false.

```
\t = \frac{1}{2} \cdot \frac{1}{2} \cdot
\tl_if_head_eq_charcode_p:fN * \tl_if_head_eq_charcode:nNTF {\langle token list \rangle} \langle test token \rangle
\tl_if_head_eq_charcode:nNTF *
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                 \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
  \tl_if_head_eq_charcode:fNTF
                                                                                                                                                                                                                              Updated: 2012-07-09
```

Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ has the same character code as the $\langle test | token \rangle$. In the case where the $\langle token \ list \rangle$ is empty, the test is always false.

```
\tl_if_head_eq_meaning_p:nN \star \tl if_head_eq_meaning_p:nN \{\langle token\ list \rangle\} \langle test\ token \rangle
\tl_if_head_eq_meaning:nNTF * \tl_if_head_eq_meaning:nNTF {\textit{token list}} \textit{test token}
                                              \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
                     Updated: 2012-07-09
```

Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ has the same meaning as the $\langle test | token \rangle$. In the case where $\langle token \ list \rangle$ is empty, the test is always false.

```
\t = \int_{a}^{b} \int_{a}^{b}
\tilde{f} = \frac{TF}{2} \cdot \tilde{f} = \frac{T
```

New: 2012-07-08 Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ is an explicit begin-group character (with category code 1 and any character code), in other words, if the $\langle token \ list \rangle$ starts with a brace group. In particular, the test is false if the $\langle token \ list \rangle$ starts with an implicit token such as \c_group_begin_token, or if it is empty. This function is useful to implement actions on token lists on a token by token basis.

```
\tilde{tl_if_head_is_N_type_p:n} \star \tilde{tl_if_head_is_N_type_p:n} \{\langle token\ list \rangle\}
\t = \frac{f(t)}{f(t)} \{ (t) \} 
             New: 2012-07-08
```

Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ is a normal N-type argument. In other words, it is neither an explicit space character (explicit token with character code 32 and category code 10) nor an explicit begin-group character (with category code 1 and any character code). An empty argument yields false, as it does not have a normal first token. This function is useful to implement actions on token lists on a token by token basis.

```
\t_if_head_is_space_p:n * \tl_if_head_is_space_p:n {\langle token \ list \rangle}
\t = \frac{TF}{\sqrt{token \ list}} {\langle true \ code \rangle} {\langle false \ code \rangle}
```

Updated: 2012-07-08 Tests if the first $\langle token \rangle$ in the $\langle token | list \rangle$ is an explicit space character (explicit token with character code 32 and category code 10). In particular, the test is false if the (token list) starts with an implicit token such as \c_space_token, or if it is empty. This function is useful to implement actions on token lists on a token by token basis.

15.4 Working with token lists as a whole

15.4.1 Using token lists

```
\t_{to_str:n} \star \t_{to_str:n} \{\langle token \ list \rangle\}
\tl_to_str:V *
```

Converts the $\langle token \ list \rangle$ to a $\langle string \rangle$, leaving the resulting character tokens in the input stream. A $\langle strinq \rangle$ is a series of tokens with category code 12 (other) with the exception of spaces, which retain category code 10 (space). This function requires only a single expansion. Its argument must be braced.

TEXhackers note: This is the ε -TEX primitive \detokenize. Converting a $\langle token \ list \rangle$ to a $\langle string \rangle$ yields a concatenation of the string representations of every token in the $\langle token \ list \rangle$. The string representation of a control sequence is

- an escape character, whose character code is given by the internal parameter \escapechar, absent if the \escapechar is negative or greater than the largest character code;
- the control sequence name, as defined by \cs_to_str:N;
- a space, unless the control sequence name is a single character whose category at the time of expansion of \tl_to_str:n is not "letter".

The string representation of an explicit character token is that character, doubled in the case of (explicit) macro parameter characters (normally #). In particular, the string representation of a token list may depend on the category codes in effect when it is evaluated, and the value of the \escapechar: for instance \tl_to_str:n {\a} normally produces the three character "backslash", "lower-case a", "space", but it may also produce a single "lower-case a" if the escape character is negative and a is currently not a letter.

```
\t_t_s \times \
                    \tl_to_str:c *
```

Converts the content of the $\langle tl \ var \rangle$ into a series of characters with category code 12 (other) with the exception of spaces, which retain category code 10 (space). This $\langle string \rangle$ is then left in the input stream. For low-level details, see the notes given for \t1_to_str:n.

* \tl_use:N \langle tl var \rangle \tl_use:N

Recovers the content of a $\langle tl \ var \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Note that it is possible to use a $\langle tl \ var \rangle$ directly without an accessor function.

15.4.2Counting and reversing token lists

 $\star \t1_count:n \{\langle tokens \rangle\}$ \tl_count:n $\t: (V|o) \star$

Counts the number of $\langle items \rangle$ in $\langle tokens \rangle$ and leaves this information in the input stream. New: 2012-05-13 Unbraced tokens count as one element as do each token group $(\{...\})$. This process ignores any unprotected spaces within $\langle tokens \rangle$. See also \tl_count:N. This function requires three expansions, giving an $\langle integer\ denotation \rangle$.

\tl count:N \tl_count:N \langle tl var \rangle

\tl_count:c *

Counts the number of $\langle items \rangle$ in the $\langle tl \ var \rangle$ and leaves this information in the input New: 2012-05-13 stream. Unbraced tokens count as one element as do each token group $\{\{...\}\}$. This process ignores any unprotected spaces within the $\langle tl \ var \rangle$. See also \t1_count:n. This function requires three expansions, giving an $\langle integer\ denotation \rangle$.

 $\t_{count_tokens:n} \star \t_{count_tokens:n} \{\langle tokens \rangle\}$

New: 2019-02-25 Counts the number of TeX tokens in the $\langle tokens \rangle$ and leaves this information in the input stream. Every token, including spaces and braces, contributes one to the total; thus for instance, the token count of a~{bc} is 6.

* \tl_reverse:n {\langle token list \rangle} \tl_reverse:n

\tl_reverse:(V|o) Updated: 2012-01-08

Reverses the order of the $\langle items \rangle$ in the $\langle token \ list \rangle$, so that $\langle item_1 \rangle \langle item_2 \rangle \langle item_3 \rangle$ $\ldots \langle item_n \rangle$ becomes $\langle item_n \rangle \ldots \langle item_3 \rangle \langle item_2 \rangle \langle item_1 \rangle$. This process preserves unprotected space within the $\langle token\ list \rangle$. Tokens are not reversed within braced token groups, which keep their outer set of braces. In situations where performance is important, consider \tl reverse items:n. See also \tl reverse:N.

TEXhackers note: The result is returned within \unexpanded, which means that the token list does not expand further when appearing in an x-type or e-type argument expansion.

\tl_reverse:N

\tl_reverse:N \(t1 \ var \)

\tl_reverse:c \tl_greverse:N

\tl_greverse:c

Updated: 2012-01-08

Sets the $\langle tl \ var \rangle$ to contain the result of reversing the order of its $\langle items \rangle$, so that $\langle item_1 \rangle \langle item_2 \rangle \langle item_3 \rangle \dots \langle item_n \rangle$ becomes $\langle item_n \rangle \dots \langle item_3 \rangle \langle item_2 \rangle \langle item_1 \rangle$. This process preserves unprotected spaces within the \(\langle token list variable \rangle\). Braced token groups are copied without reversing the order of tokens, but keep the outer set of braces. This is equivalent to a combination of an assignment and \tl_reverse: V. See also \tl_reverse_items:n for improved performance.

 $\t_reverse_items:n \star \t_reverse_items:n \{\langle token \ list \rangle\}$

New: 2012-01-08 Reverses the order of the $\langle items \rangle$ stored in $\langle tl \ var \rangle$, so that $\{\langle item_1 \rangle\} \{\langle item_2 \rangle\} \{\langle item_3 \rangle\}$ $\ldots \{\langle item_n \rangle\}$ becomes $\{\langle item_n \rangle\} \ldots \{\langle item_3 \rangle\} \{\langle item_2 \rangle\} \{\langle item_1 \rangle\}.$ This process removes any unprotected space within the $\langle token \ list \rangle$. Braced token groups are copied without reversing the order of tokens, and keep the outer set of braces. Items which are initially not braced are copied with braces in the result. In cases where preserving spaces is important, consider the slower function \tl_reverse:n.

> TEXhackers note: The result is returned within \unexpanded, which means that the token list does not expand further when appearing in an x-type or e-type argument expansion.

 $\t_t_{\text{trim_spaces:n}} \star \t_{\text{trim_spaces:n}} \{\langle token \ list \rangle\}$

Updated: 2012-06-25 stream.

 $\til_{trim_spaces:o} \star \text{Removes any leading and trailing explicit space characters (explicit tokens with characters)}$ New: 2011-07-09 code 32 and category code 10) from the $\langle token \ list \rangle$ and leaves the result in the input

> TEXhackers note: The result is returned within \unexpanded, which means that the token list does not expand further when appearing in an x-type or e-type argument expansion.

\tl_trim_spaces_apply:oN *

 $\verb|\tl_trim_spaces_apply:nN| * \verb|\tl_trim_spaces_apply:nN| \{ \langle token \ list \rangle \} \ \langle function \rangle|$

Removes any leading and trailing explicit space characters (explicit tokens with character New: 2018-04-12 code 32 and category code 10) from the $\langle token \ list \rangle$ and passes the result to the $\langle function \rangle$ as an n-type argument.

\tl_trim_spaces:c \tl_gtrim_spaces:N \tl_gtrim_spaces:c

 $\t_t_{trim_spaces:N} \t_t_{trim_spaces:N} \t var$

Sets the $\langle tl \ var \rangle$ to contain the result of removing any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from its contents.

New: 2011-07-09

Viewing token lists 15.4.3

\tl_show:N

\tl_show:N \langle tl var \rangle

\tl_show:c

Displays the content of the $\langle tl \ var \rangle$ on the terminal.

Updated: 2021-04-29

TeXhackers note: This is similar to the TeX primitive \show, wrapped to a fixed number of characters per line.

\tl_show:n

 $\t: \{\langle token \ list \rangle\}$

Updated: 2015-08-07 Displays the $\langle token \ list \rangle$ on the terminal.

TEX hackers note: This is similar to the ε -TEX primitive \showtokens, wrapped to a fixed number of characters per line.

\tl_log:N

\tl_log:N \(tl var \)

\tl_log:c

Writes the content of the $\langle tl \ var \rangle$ in the log file. See also \tl_show: N which displays the New: 2014-08-22 result in the terminal.

Updated: 2021-04-29

\tl_log:n

 $\tilde{\beta} = \frac{\langle token \ list \rangle}{\langle token \ list \rangle}$

Updated: 2015-08-07 the terminal.

New: 2014-08-22 Writes the $\langle token \ list \rangle$ in the log file. See also \t1_show:n which displays the result in

15.5 Manipulating items in token lists

15.5.1 Mapping over token lists

All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

\tl_map_function:NN ☆ \tl_map_function:cN ☆

\tl_map_function:NN \langle tl var \rangle \function \rangle

Applies $\langle function \rangle$ to every $\langle item \rangle$ in the $\langle tl \ var \rangle$. The $\langle function \rangle$ receives one argument Updated: 2012-06-29 for each iteration. This may be a number of tokens if the (item) was stored within braces. Hence the $\langle function \rangle$ should anticipate receiving n-type arguments. See also \tl_map_function:nN.

 $\t = \mathbf{1}$ $\t = \mathbf{1}$ $\t = \mathbf{1}$

Updated: 2012-06-29 Applies $\langle function \rangle$ to every $\langle item \rangle$ in the $\langle token\ list \rangle$, The $\langle function \rangle$ receives one argument for each iteration. This may be a number of tokens if the $\langle item \rangle$ was stored within braces. Hence the $\langle function \rangle$ should anticipate receiving n-type arguments. See also \tl_map_function:NN.

 $\tilde{\ }$ \tl_map_inline:Nn \tl_map_inline:Nn $\langle tl \ var \rangle \ \{\langle inline \ function \rangle \}$

\tl_map_inline:cn

Updated: 2012-06-29

Applies the $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle tl\ var \rangle$. The $\langle inline\ function \rangle$ function should consist of code which receives the (item) as #1. See also \tl_map_function: NN.

 $\tilde{\theta} = \tilde{\theta}$ (inline function)

Updated: 2012-06-29 Applies the $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle token\ list \rangle$. The $\langle inline\ function \rangle$ function should consist of code which receives the (item) as #1. See also \tl_map_function:nN.

```
\t_map_tokens:Nn \Leftrightarrow \t_map_tokens:Nn \langle tl var \rangle \{\langle code \rangle\}
\t_{map\_tokens:cn} \not\simeq \t_{map\_tokens:nn} {\langle tokens \rangle} {\langle code \rangle}
```

\t1_map_tokens:nn 🌣 Analogue of \t1_map_function:NN which maps several tokens instead of a single func-New: 2019-09-02 tion. The $\langle code \rangle$ receives each $\langle item \rangle$ in the $\langle tl \ var \rangle$ or in $\langle tokens \rangle$ as a trailing brace group. For instance,

```
\tl_map_tokens:Nn \l_my_tl { \prg_replicate:nn { 2 } }
```

expands to twice each $\langle item \rangle$ in the $\langle tl \ var \rangle$: for each $\langle item \rangle$ in \l_my_tl the function \prg_replicate:nn receives 2 and \(\lambda item\rangle\) as its two arguments. The function \t1_map inline: Nn is typically faster but is not expandable.

 $\t = \mbox{ } \t = \mbox{ }$

Updated: 2012-06-29

 $\frac{\text{tl_map_variable:cNn}}{\text{stores each }\langle item\rangle}$ Stores each $\langle item\rangle$ of the $\langle tl\ var\rangle$ in turn in the (token list) $\langle variable\rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle tl \ var \rangle$, or its original value if the $\langle tl \ var \rangle$ is blank. See also \tl map inline: Nn.

 $\tilde{\zeta} = \tilde{\zeta}$ variable:nNn $\tilde{\zeta} = \tilde{\zeta}$ (code)

Updated: 2012-06-29 Stores each $\langle item \rangle$ of the $\langle token \ list \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle tl \ var \rangle$, or its original value if the $\langle tl \ var \rangle$ is blank. See also \tl_map_inline:nn.

\tl_map_break: ☆ \tl_map_break:

Updated: 2012-06-29 Used to terminate a \t1_map_... function before all entries in the \(\lambda to terminate a \t1_map_...\) have been processed. This normally takes place within a conditional statement, for example

```
\tl_map_inline:Nn \l_my_tl
    \str_if_eq:nnT { #1 } { bingo } { \tl_map_break: }
   % Do something useful
 }
```

See also \tl_map_break:n. Use outside of a \tl_map_... scenario leads to low level

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle tokens \rangle$ are inserted into the input stream. This depends on the design of the mapping function.

```
\t_map_break:n \Leftrightarrow \t_map_break:n \{\langle code \rangle\}
```

Updated: 2012-06-29 Used to terminate a \tl_map_... function before all entries in the \(\lambda token list variable\rangle\) have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\tl_map_inline:Nn \l_my_tl
 {
    \str_if_eq:nnT { #1 } { bingo }
     { \tl_map_break:n { <code> } }
    % Do something useful
 }
```

Use outside of a \tl_map_... scenario leads to low level TeX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

15.5.2 Head and tail of token lists

Functions which deal with either only the very first item (balanced text or single normal token) in a token list, or the remaining tokens.

```
\tl head:N
\tl_head:n
```

```
* \tl_head:n {\langle token list \rangle}
```

Leaves in the input stream the first $\langle item \rangle$ in the $\langle token \ list \rangle$, discarding the rest of the (token list). All leading explicit space characters (explicit tokens with character code 32 ${\tt Updated:2012-09-09}$ and category code 10) are discarded; for example

```
\tl_head:n { abc }
```

and

```
\tl_head:n { ~ abc }
```

both leave a in the input stream. If the "head" is a brace group, rather than a single token, the braces are removed, and so

```
\tl head:n { ~ { ~ ab } c }
```

yields ⊔ab. A blank ⟨token list⟩ (see \tl_if_blank:nTF) results in \tl_head:n leaving nothing in the input stream.

TEXhackers note: The result is returned within \exp_not:n, which means that the token list does not expand further when appearing in an x-type or e-type argument expansion.

```
\t_head: w \star tl_head: w \langle token \ list \rangle \  \   \   \   \
```

Leaves in the input stream the first $\langle item \rangle$ in the $\langle token \ list \rangle$, discarding the rest of the (token list). All leading explicit space characters (explicit tokens with character code 32 and category code 10) are discarded. A blank $\langle token \ list \rangle$ (which consists only of space characters) results in a low-level T_EX error, which may be avoided by the inclusion of an empty group in the input (as shown), without the need for an explicit test. Alternatively, \tl_if_blank:nF may be used to avoid using the function with a "blank" argument. This function requires only a single expansion, and thus is suitable for use within an o-type expansion. In general, \t1_head:n should be preferred if the number of expansions is not critical.

```
\tl_tail:N
\tl_tail:n
\tl_tail:(V|v|f)
```

```
* \tl_tail:n {\langle token list \rangle}
```

Discards all leading explicit space characters (explicit tokens with character code 32 and category code 10) and the first $\langle item \rangle$ in the $\langle token \ list \rangle$, and leaves the remaining tokens Updated: 2012-09-01 in the input stream. Thus for example

```
\tl_tail:n { a ~ {bc} d }
and
    \tl_tail:n { ~ a ~ {bc} d }
```

both leave | {bc}d in the input stream. A blank \(\langle token \ list \rangle \) (see \t1 if blank:nTF) results in \tl_tail:n leaving nothing in the input stream.

 $\textbf{T}_{\!\!\textbf{E}}\!\textbf{X}\textbf{hackers note:} \text{ The result is returned within } \texttt{\tt exp_not:n}, \text{ which means that the token}$ list does not expand further when appearing in an x-type or e-type argument expansion.

If you wish to handle token lists where the first token may be a space, and this needs to be treated as the head/tail, this can be accomplished using \tl if head is space:nTF, for example

```
\exp_last_unbraced:NNo
  \cs_new:Npn \__mypkg_gobble_space:w \c_space_tl { }
\cs_new:Npn \mypkg_tl_head_keep_space:n #1
  {
    \tl_if_head_is_space:nTF {#1}
      { ~ }
      { \tl head:n {#1} }
\cs_new:Npn \mypkg_tl_tail_keep_space:n #1
    \tl_if_head_is_space:nTF {#1}
      { \exp_not:o { \__mypkg_gobble_space:w #1 } }
      { \tl_tail:n {#1} }
  }
```

Items and ranges in token lists

```
\tl_item:Nn *
\tl_item:cn *
```

 $\tilde{\theta} \to \tilde{\theta}$ {\tilem:nn \tau\tilem:nn \{\tileq token list\}} {\tileq integer expression\}}

Indexing items in the $\langle token\ list \rangle$ from 1 on the left, this function evaluates the $\langle integer$ expression and leaves the appropriate item from the $\langle token \ list \rangle$ in the input stream. New: 2014-07-17 If the $\langle integer\ expression \rangle$ is negative, indexing occurs from the right of the token list, starting at -1 for the right-most item. If the index is out of bounds, then the function expands to nothing.

> TEX hackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the \(\langle item \rangle \) does not expand further when appearing in an x-type or e-type argument expansion.

```
\tl_rand_item:N * \tl_rand_item:N \( tl \ var \)
\t_{rand_item:c} \star \t_{rand_item:n} \{\langle token \ list \rangle\}
```

New: 2016-12-06 is empty. This is not available in older versions of $X_{\overline{1}}T_{\overline{1}}X$.

> TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

```
\t_range:Nnn * \t_range:Nnn \t var \{\langle start index \rangle\} \{\langle end index \rangle\}}
\tilde{\langle tl_range:nnn * \langle tl_range:nnn {\langle token list \rangle} {\langle start index \rangle} {\langle end index \rangle}
```

New: 2017-02-17 Leaves in the input stream the items from the $\langle start\ index \rangle$ to the $\langle end\ index \rangle$ inclusive. Updated: 2017-07-15 Spaces and braces are preserved between the items returned (but never at either end of the list). Here $\langle start\ index \rangle$ and $\langle end\ index \rangle$ should be $\langle integer\ expressions \rangle$. For describing in detail the functions' behavior, let m and n be the start and end index respectively. If either is 0, the result is empty. A positive index means 'start counting from the left end', and a negative index means 'from the right end'. Let l be the count of the token list.

> The actual start point is determined as M=m if m>0 and as M=l+m+1if m < 0. Similarly the actual end point is N = n if n > 0 and N = l + n + 1 if n < 0. If M > N, the result is empty. Otherwise it consists of all items from position M to position N inclusive; for the purpose of this rule, we can imagine that the token list extends at infinity on either side, with void items at positions s for s < 0 or s > l.

> Spaces in between items in the actual range are preserved. Spaces at either end of the token list will be removed anyway (think to the token list being passed to \tl_trim_spaces:n to begin with.

> Thus, with l=7 as in the examples below, all of the following are equivalent and result in the whole token list

```
\tl_range:nnn { abcd~{e{}}fg } { 1 } { 7 }
\tl_range:nnn { abcd~{e{}}fg } { 1 } { 12 }
\tl_range:nnn { abcd~{e{}}fg } { -7 } { 7 }
\tl_range:nnn { abcd~{e{}}fg } { -12 } { 7 }
```

Here are some more interesting examples. The calls

```
\iow_term:x { \tl_range:nnn { abcd{e{}}fg } { 2 } { 5 } }
\iow_term:x { \tl_range:nnn { abcd{e{}}fg } { 2 } { -3 } }
\iow_term:x { \tl_range:nnn { abcd{e{}}fg } { -6 } { 5 } }
\iow_term:x { \tl_range:nnn { abcd{e{}}fg } { -6 } { -3 } }
```

are all equivalent and will print bcd{e{}} on the terminal; similarly

```
\iow_term:x { \tl_range:nnn { abcd~{e{}}fg } { 2 } { 5 } }
\iow_term:x { \tl_range:nnn { abcd~{e{}}fg } { 2 } { -3 } }
\iow term:x { \tl range:nnn { abcd~{e{}}fg } { -6 } { 5 } }
\iow_term:x { \tl_range:nnn { abcd~{e{}}fg } { -6 } { -3 } }
```

are all equivalent and will print bcd {e{}} on the terminal (note the space in the middle). To the contrary,

```
\tl_range:nnn { abcd~{e{}}f } { 2 } { 4 }
```

will discard the space after 'd'.

If we want to get the items from, say, the third to the last in a token list <tl>, the call is \t1_range:nnn { <t1> } { 3 } { -1 }. Similarly, for discarding the last item, we can do \tl_range:nnn { <tl> } { 1 } { -2 }.

For better performance, see \tl_range_braced:nnn and \tl_range_unbraced:nnn.

TrXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the \(\lambda item\rangle\) does not expand further when appearing in an x-type or e-type argument expansion.

15.5.4 Sorting token lists

\tl_sort:Nn \tl_sort:cn \tl_gsort:Nn \tl_gsort:cn

\tl_sort:Nn \langletl var \rangle \langle \comparison code \rangle \rangle

Sorts the items in the $\langle tl \ var \rangle$ according to the $\langle comparison \ code \rangle$, and assigns the result to $\langle tl \ var \rangle$. The details of sorting comparison are described in Section 6.1.

New: 2017-02-06

 $\t:nN \star tl_sort:nN \{\langle token \ list \rangle\} \langle conditional \rangle$

New: 2017-02-06 Sorts the items in the $\langle token\ list \rangle$, using the $\langle conditional \rangle$ to compare items, and leaves the result in the input stream. The (conditional) should have signature:nnTF, and return true if the two items being compared should be left in the same order, and false if the items should be swapped. The details of sorting comparison are described in Section 6.1.

> TEXhackers note: The result is returned within \exp_not:n, which means that the token list does not expand further when appearing in an x-type or e-type argument expansion.

15.6 Manipulating tokens in token lists

15.6.1Replacing tokens

Within token lists, replacement takes place at the top level: there is no recursion into brace groups (more precisely, within a group defined by a category code 1/2 pair).

\tl_replace_once:Nnn \tl_replace_once:cnn \tl_greplace_once:Nnn \tl_greplace_once:cnn

 $\verb|\tl_replace_once:Nnn| \langle tl var \rangle | \{\langle old tokens \rangle\} | \{\langle new tokens \rangle\}|$

Replaces the first (leftmost) occurrence of $\langle old\ tokens \rangle$ in the $\langle tl\ var \rangle$ with $\langle new\ tokens \rangle$. (Old tokens) cannot contain {, } or # (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

\tl_replace_all:cnn \tl_greplace_all:Nnn

Updated: 2011-08-11

Updated: 2011-08-11

Replaces all occurrences of $\langle old\ tokens \rangle$ in the $\langle tl\ var \rangle$ with $\langle new\ tokens \rangle$. $\langle Old\ tokens \rangle$ \tl_greplace_all:cnn cannot contain {, } or # (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern (old tokens) may remain after the replacement (see \tl_remove_all:Nn for an example).

\tl_remove_once:Nn \tl_remove_once:cn \tl_gremove_once:Nn \tl_gremove_once:cn

Updated: 2011-08-11

Removes the first (leftmost) occurrence of $\langle tokens \rangle$ from the $\langle tl \ var \rangle$. $\langle Tokens \rangle$ cannot contain {, } or # (more precisely, explicit character tokens with category code 1 (begingroup) or 2 (end-group), and tokens with category code 6).

```
\tl_remove_all:Nn
\tl_remove_all:cn
\tl_gremove_all:Nn
\tl_gremove_all:cn
Updated:2011-08-11
```

```
\t! remove_all:Nn \langle tl \ var \rangle \ \{\langle tokens \rangle\}
```

Removes all occurrences of $\langle tokens \rangle$ from the $\langle tl \ var \rangle$. $\langle Tokens \rangle$ cannot contain $\{$, $\}$ or # (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern $\langle tokens \rangle$ may remain after the removal, for instance,

```
\tl_set:Nn \l_tmpa_tl {abbccd} \tl_remove_all:Nn \l_tmpa_tl {bc}
results in \l_tmpa_tl containing abcd.
```

15.6.2 Reassigning category codes

These functions allow the rescanning of tokens: re-apply TEX's tokenization process to apply category codes different from those in force when the tokens were absorbed. Whilst this functionality is supported, it is often preferable to find alternative approaches to achieving outcomes rather than rescanning tokens (for example construction of token lists token-by-token with intervening category code changes or using \char_generate:nn).

Sets $\langle tl\,var\rangle$ to contain $\langle tokens\rangle$, applying the category code régime specified in the $\langle setup\rangle$ before carrying out the assignment. (Category codes applied to tokens not explicitly covered by the $\langle setup\rangle$ are those in force at the point of use of $\verb+\tlue{tl_set_rescan:Nnn.}$) This allows the $\langle tl\,var\rangle$ to contain material with category codes other than those that apply when $\langle tokens\rangle$ are absorbed. The $\langle setup\rangle$ is run within a group and may contain any valid input, although only changes in category codes, such as uses of $\verb+\cute_ctab_select:N|$, are relevant. See also $\verb+\tl_rescan:nn|$.

TEXhackers note: The $\langle tokens \rangle$ are first turned into a string (using \tl_to_str:n). If the string contains one or more characters with character code \newlinechar (set equal to \endlinechar unless that is equal to 32, before the user $\langle setup \rangle$), then it is split into lines at these characters, then read as if reading multiple lines from a file, ignoring spaces (catcode 10) at the beginning and spaces and tabs (character code 32 or 9) at the end of every line. Otherwise, spaces (and tabs) are retained at both ends of the single-line string, as if it appeared in the middle of a line read from a file.

\tl_rescan:nn

 $\t!$ rescan:nn { $\langle setup \rangle$ } { $\langle tokens \rangle$ }

Updated: 2015-08-11 Rescans $\langle tokens \rangle$ applying the category code régime specified in the $\langle setup \rangle$, and leaves the resulting tokens in the input stream. (Category codes applied to tokens not explicitly covered by the \(\setup\) are those in force at the point of use of \tl_rescan:nn.\) The (setup) is run within a group and may contain any valid input, although only changes in category codes, such as uses of \cctab_select:N, are relevant. See also \tl_set_rescan: Nnn, which is more robust than using \tl_set: Nn in the \(\lambda tokens \rangle\) argument of \tl_rescan:nn.

> **TEXhackers note:** The $\langle tokens \rangle$ are first turned into a string (using $tl_to_str:n$). If the string contains one or more characters with character code \newlinechar (set equal to \endlinechar unless that is equal to 32, before the user $\langle setup \rangle$), then it is split into lines at these characters, then read as if reading multiple lines from a file, ignoring spaces (catcode 10) at the beginning and spaces and tabs (character code 32 or 9) at the end of every line. Otherwise, spaces (and tabs) are retained at both ends of the single-line string, as if it appeared in the middle of a line read from a file.

> Contrarily to the \scantokens primitive, \tl_rescan:nn tokenizes the whole string in the same category code regime rather than one token at a time, so that directives such as \verb that rely on changing category codes will not function properly.

Constant token lists 15.7

\c_empty_tl Constant that is always empty.

 $\c_novalue_tl$ A marker for the absence of an argument. This constant tl can safely be typeset (cf. \q_- New: 2017-11-14 nil), with the result being -NoValue-. It is important to note that \c_novalue_tl is constructed such that it will not match the simple text input -NoValue-, i.e. that

```
\tl_if_eq:NnTF \c_novalue_tl { -NoValue- }
```

is logically false. The \c_novalue_tl marker is intended for use in creating documentlevel interfaces, where it serves as an indicator that an (optional) argument was omitted. In particular, it is distinct from a simple empty t1.

\c_space_tl An explicit space character contained in a token list (compare this with \c_space_token). For use where an explicit space is required.

15.8 Scratch token lists

\l_tmpa_tl Scratch token lists for local assignment. These are never used by the kernel code, and so \l_tmpb_tl are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Chapter 16

The l3str package: Strings

TEX associates each character with a category code: as such, there is no concept of a "string" as commonly understood in many other programming languages. However, there are places where we wish to manipulate token lists while in some sense "ignoring" category codes: this is done by treating token lists as strings in a TEX sense.

A TEX string (and thus an expl3 string) is a series of characters which have category code 12 ("other") with the exception of space characters which have category code 10 ("space"). Thus at a technical level, a TEX string is a token list with the appropriate category codes. In this documentation, these are simply referred to as strings.

String variables are simply specialised token lists, but by convention should be named with the suffix ...str. Such variables should contain characters with category code 12 (other), except spaces, which have category code 10 (blank space). All the functions in this module which accept a token list argument first convert it to a string using \t1_to_str:n for internal processing, and do not treat a token list or the corresponding string representation differently.

As a string is a subset of the more general token list, it is sometimes unclear when one should be used over the other. Use a string variable for data that isn't primarily intended for typesetting and for which a level of protection from unwanted expansion is suitable. This data type simplifies comparison of variables since there are no concerns about expansion of their contents.

The functions \cs_to_str:N, \tl_to_str:n, \tl_to_str:N and \token_to_str:N (and variants) generate strings from the appropriate input: these are documented in l3basics, l3tl and l3token, respectively.

Most expandable functions in this module come in three flavours:

- \str_...:N, which expect a token list or string variable as their argument;
- \str_...:n, taking any token list (or string) as an argument;
- \str_..._ignore_spaces:n, which ignores any space encountered during the operation: these functions are typically faster than those which take care of escaping spaces appropriately.

16.1 Creating and initialising string variables

```
\str_new:N \langle str var \rangle
                                   \str_new:N
                                   \str_new:c
                                                                   Creates a new \langle str \, var \rangle or raises an error if the name is already taken. The declaration
                                  New: 2015-09-18 is global. The \langle str \; var \rangle is initially empty.
                                                                   \str_const:Nn \langle str var \rangle \{\langle token list \rangle\}
\str_const:Nn
\str_const:(NV|Nx|cn|cV|cx)
                                                                   Creates a new constant \langle str \ var \rangle or raises an error if the name is already taken. The
                                  New: 2015-09-18 value of the \langle str \ var \rangle is set globally to the \langle token \ list \rangle, converted to a string.
                          Updated: 2018-07-28
                              \str_clear:N \str_clear:N \str var >
                               \str_clear:c
                                                                   Clears the content of the \langle str \ var \rangle.
                              \str_gclear:N
                               \str_gclear:c
                                  New: 2015-09-18
                      \str_clear_new:N \str_clear_new:N \str var \
                       \str_clear_new:c
                                                                  Ensures that the \( \str var \rangle \) exists globally by applying \str_new:N if necessary, then
                                  New: 2015-09-18 applies \str_(g)clear: N to leave the \langle str \ var \rangle empty.
       \str_set_eq:NN
                                                                   \str_set_eq:NN \langle str var_1 \rangle \langle str var_2 \rangle
        \str_set_eq:(cN|Nc|cc)
                                                                   Sets the content of \langle str \ var_1 \rangle equal to that of \langle str \ var_2 \rangle.
       \str_gset_eq:NN
        \str_gset_eq:(cN|Nc|cc)
                                  New: 2015-09-18
                      \str\_concat:NNN \str\_concat:NNN \str\_var_1 \slash \slash
                      \sl_{\text{str\_concat:NNN}} Concatenates the content of \langle str\ var_2 \rangle and \langle str\ var_3 \rangle together and saves the result in
                      \str_gconcat:ccc \langle str \ var_1 \rangle. The \langle str \ var_2 \rangle is placed at the left side of the new string variable. The
                                                                   \langle str \ var_2 \rangle and \langle str \ var_3 \rangle must indeed be strings, as this function does not convert their
                                  New: 2017-10-08 contents to a string.
             \str_if_exist_p:N * \str_if_exist_p:N \langle str var \rangle
              \str_if_exist_p:c \star \str_if_exist:NTF \str var \ {\langle true code \rangle} \ {\langle false code \rangle}
             \str_{if_exist:N} = \star Tests whether the \langle str \ var \rangle is currently defined. This does not check that the \langle str \ var \rangle
              \str_{if} = \frac{cTF}{cTF} \star really is a string.
                                  New: 2015-09-18
```

16.2 Adding data to string variables

```
\str_set:Nn \str_s
```

Converts the $\langle token \ list \rangle$ to a $\langle string \rangle$, and prepends the result to $\langle str \ var \rangle$. The current contents of the $\langle str \ var \rangle$ are not automatically converted to a string.

Converts the $\langle token \ list \rangle$ to a $\langle string \rangle$, and appends the result to $\langle str \ var \rangle$. The current contents of the $\langle str \ var \rangle$ are not automatically converted to a string.

16.3 String conditionals

```
\str_if_empty_p:N \ \str_if_empty_p:N \ \str_if_empty_p:N \ \str_if_empty_p:N \ \str_if_empty:NTF \ \str_if_empty:NTF \ \str_if_empty:NTF \ \str_if_empty:NTF \ \str_if_empty:NTF \ \str_if_empty:CTF \ \str_if_empty_p:n \ \str_if_empty:nTF \ \ \str_if_eq_p:NN \ \str_if_eq_p:NN \ \str_if_eq_p:NN \ \str_if_eq_p:NN \ \str_if_eq_p:NN \ \str_if_eq:NNTF \ \s
```

Compares the two $\langle token \ lists \rangle$ on a character by character basis (namely after converting them to strings), and is **true** if the two $\langle strings \rangle$ contain the same characters in the same order. Thus for example

```
\str_if_eq_p:no { abc } { \tl_to_str:n { abc } }
```

is logically true. See $\t_if_eq:nnTF$ to compare tokens (including their category codes) rather than characters.

```
\label{eq:linear_case:nn} $$ \operatorname{test\ string} $$ \operatorname{case:nnTF} \{\langle \operatorname{test\ string} \rangle \} $$ \operatorname{test\ case:(Vn|on|nV|nv)} $$ \times $$ \{ \operatorname{test\ string} \rangle \} $$ \times \operatorname{test\ case:(Vn|on|nV|nv)} $$ \times $$ \{ \langle \operatorname{string\ case_1} \rangle \} $$ \{ \langle \operatorname{code\ case_2} \rangle \} $$ \times \operatorname{test\ case:(Vn|on|nV|nv)} $$ \times $$ \{ \langle \operatorname{string\ case_2} \rangle \} $$ \{ \langle \operatorname{code\ case_2} \rangle \} $$ \times \operatorname{test\ case:(Nn\ TF)} $$ \times $$ \{ \langle \operatorname{string\ case_n} \rangle \} $$ \{ \langle \operatorname{code\ case_n} \rangle \} $$ \} $$ \times \operatorname{test\ code} $$
```

Compares the $\langle test\ string \rangle$ in turn with each of the $\langle string\ cases \rangle$ (all token lists are converted to strings). If the two are equal (as described for $\str_if_eq:nnTF$) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The function $\str_case:nn$, which does nothing if there is no match, is also available.

This set of functions performs no expansion on each $\langle string \; case \rangle$ argument, so any variable in there will be compared as a string. If expansion is needed in the $\langle string \; cases \rangle$, then $\str_{case_e:nn}(TF)$ should be used instead.

```
\str_case_e:nnTF {\langle test string \rangle}
\str_case_e:nn
\str_case_e:nn<u>TF</u>
                                           \{\langle string \ case_1 \rangle\} \ \{\langle code \ case_1 \rangle\}
            New: 2018-06-19
                                           \{\langle string \ case_2 \rangle\} \ \{\langle code \ case_2 \rangle\}
                                           \{\langle string \ case_n \rangle\}\ \{\langle code \ case_n \rangle\}
                                       {\langle true code \rangle}
                                       {\false code\}
```

Compares the full expansion of the $\langle test \ strinq \rangle$ in turn with the full expansion of the (string cases) (all token lists are converted to strings). If the two full expansions are equal (as described for \str_if_eq:nnTF) then the associated \(code \) is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the \(false \) code\() is inserted. The function \\str_case_e:nn, \) which does nothing if there is no match, is also available. The $\langle test \ string \rangle$ is expanded in each comparison, and must always yield the same result: for example, random numbers must not be used within this string.

```
\str_compare:nNnTF
\str_compare:eNe<u>TF</u> *
```

```
\str_compare_p:nNn * \str_compare_p:nNn {\langle tl_1 \rangle} {\langle relation \rangle} {\langle tl_2 \rangle}
\str\_compare\_p:eNe * \str\_compare:nNnTF {(tl_1)} (relation) {(tl_2)} {(true code)} {(false code)}
```

Compares the two $\langle token\ lists \rangle$ on a character by character basis (namely after converting them to strings) in a lexicographic order according to the character codes of the New: 2021-05-17 characters. The $\langle relation \rangle$ can be <, =, or > and the test is true under the following conditions:

- for <, if the first string is earlier than the second in lexicographic order;
- for =, if the two strings have exactly the same characters;
- for >, if the first string is later than the second in lexicographic order.

Thus for example the following is logically true:

```
\str_compare_p:nNn { ab } < { abc }
```

TEXhackers note: This is a wrapper around the TEX primitive \((pdf)\)strcmp. It is meant for programming and not for sorting textual contents, as it simply considers character codes and not more elaborate considerations of grapheme clusters, locale, etc.

16.4 Mapping over strings

All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

```
\str_map_function:nN & \str_map_function:nN {\langle token list \rangle} \langle function \rangle
\str_map_function:NN ☆ \str_map_function:NN ⟨str var⟩ ⟨function⟩
```

New: 2017-11-14 $\langle string \rangle$ including spaces.

```
\operatorname{str}_{map}_{inline:nn} \operatorname{str}_{map}_{inline:nn} \{\langle token\ list \rangle\} \{\langle inline\ function \rangle\}
\str_map_inline:Nn \str_map_inline:Nn \str \var \ {\langle inline \ function \rangle}
\overline{\text{Var_map_inline:cn}} Converts the \langle token\ list \rangle to a \langle string \rangle then applies the \langle inline\ function \rangle to every
          New: 2017-11-14 \langle character \rangle in the \langle str \ var \rangle including spaces. The \langle inline \ function \rangle should consist of
                             code which receives the \langle character \rangle as #1.
```

```
\str_map\_tokens:Nn \Leftrightarrow \str_map\_tokens:Nn \langle str var \rangle \{\langle code \rangle\}
```

 $\$ Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then applies $\langle code \rangle$ to every $\langle character \rangle$ in the New: 2021-05-05 $\langle strinq \rangle$ including spaces. The $\langle code \rangle$ receives each character as a trailing brace group. This is equivalent to $\operatorname{str}_{map}_{function}$:nN if the $\langle code \rangle$ consists of a single function.

 $\str_map_variable:nNn \str_map_variable:nNn {$\langle token\ list\rangle$} \ \langle variable\rangle$ {$\langle code\rangle$}$

 $\overline{\text{Var_map_variable:cNn}}$ Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then stores each $\langle character \rangle$ in the $\langle string \rangle$ [in-New: 2017-11-14 cluding spaces) in turn in the (string or token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle character \rangle$ in the $\langle string \rangle$, or its original value if the \(\string \) is empty. See also \str_map_inline: Nn.

\str_map_break: ☆ \str_map_break:

New: 2017-10-08 Used to terminate a \str_map_... function before all characters in the \string\ have been processed. This normally takes place within a conditional statement, for example

```
\str_map_inline:Nn \l_my_str
    \str_if_eq:nnT { #1 } { bingo } { \str_map_break: }
   % Do something useful
```

See also \str_map_break:n. Use outside of a \str_map_... scenario leads to low level T_EX errors.

TEXhackers note: When the mapping is broken, additional tokens may be inserted before continuing with the code that follows the loop. This depends on the design of the mapping function.

```
\str_map_break:n \Leftrightarrow \str_map_break:n {\langle code \rangle}
```

New: 2017-10-08 Used to terminate a \str_map_... function before all characters in the \(\string \) have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\str_map_inline:Nn \l_my_str
  {
    \str_if_eq:nnT { #1 } { bingo }
      { \str_map_break:n { <code> } }
    % Do something useful
 }
```

Use outside of a \str_map_... scenario leads to low level TeX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

16.5 Working with the content of strings

```
\str_use:N \star \str_use:N \langle str \ var \rangle
```

\str_use:c \star Recovers the content of a $\langle str \ var \rangle$ and places it directly in the input stream. An error New: 2015-09-18 is raised if the variable does not exist or if it is invalid. Note that it is possible to use a $\langle str \rangle$ directly without an accessor function.

```
\str_count:N
                                  \star \str_count:n {\langle token list\rangle}
\str_count:c
\str_count:n
\str_count_ignore_spaces:n
                    New: 2015-09-18
```

Leaves in the input stream the number of characters in the string representation of $\langle token \rangle$ list, as an integer denotation. The functions differ in their treatment of spaces. In the case of \str_count:N and \str_count:n, all characters including spaces are counted. The \str_count_ignore_spaces:n function leaves the number of non-space characters in the input stream.

```
\str\_count\_spaces:N * \str\_count\_spaces:n {\langle token list \rangle}
\str_count_spaces:c *
```

\str_count_spaces:n *

New: 2015-09-18 variant.

Leaves in the input stream the number of space characters in the string representation of (token list), as an integer denotation. Of course, this function has no _ignore_spaces

```
\label{list} $$ \left( \begin{array}{ccc} & \star & \star \\ str_head:n & \star \\ str_head:c & \star \\ str_head:n & \star \\ & \star \\ \hline & str_head\_ignore\_spaces:n & \star \\ \hline & New: 2015-09-18 \\ \end{array} \right) $$
```

Converts the $\langle token\ list\rangle$ into a $\langle string\rangle$. The first character in the $\langle string\rangle$ is then left in the input stream, with category code "other". The functions differ if the first character is a space: $\str_head:N$ and $\str_head:n$ return a space token with category code 10 (blank space), while the $\str_head_ignore_spaces:n$ function ignores this space character and leaves the first non-space character in the input stream. If the $\langle string\rangle$ is empty (or only contains spaces in the case of the <code>_ignore_spaces</code> function), then nothing is left on the input stream.

```
\label{limit} $$ \left( \begin{array}{cccc} token \ list \\ \\ tr_tail:n & \star \\ \\ tr_tail:n & \star \\ \\ tr_tail\_ignore\_spaces:n & \star \\ \\ \hline \\ New: 2015-09-18 \\ \end{array} \right)$
```

Converts the \(\lambda to ken \ list \rangle \) to a \(\lambda string \rangle \), removes the first character, and leaves the remaining characters (if any) in the input stream, with category codes 12 and 10 (for spaces). The functions differ in the case where the first character is a space: \str_tail:N and \str_tail:n only trim that space, while \str_tail_ignore_spaces:n removes the first non-space character and any space before it. If the \(\lambda to ken \ list \rangle \) is empty (or blank in the case of the _ignore_spaces variant), then nothing is left on the input stream.

```
\label{limin} $$ \left( integer \ expression \right) \ \left( integer \ expresion \right) \ \left( integer \ expression \right) \ \left( integer \ expression \right)
```

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and leaves in the input stream the character in position $\langle integer\ expression \rangle$ of the $\langle string \rangle$, starting at 1 for the first (left-most) character. In the case of $\str_item:Nn$ and $\str_item:nn$, all characters including spaces are taken into account. The $\str_item_ignore_spaces:nn$ function skips spaces when counting characters. If the $\langle integer\ expression \rangle$ is negative, characters are counted from the end of the $\langle string \rangle$. Hence, -1 is the right-most character, etc.

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and leaves in the input stream the characters from the $\langle start\ index \rangle$ to the $\langle end\ index \rangle$ inclusive. Spaces are preserved and counted as items (contrast this with $\t_{range:nnn}$ where spaces are not counted as items and are possibly discarded from the output).

Here $\langle start\ index \rangle$ and $\langle end\ index \rangle$ should be integer denotations. For describing in detail the functions' behavior, let m and n be the start and end index respectively. If either is 0, the result is empty. A positive index means 'start counting from the left end', a negative index means 'start counting from the right end'. Let l be the count of the token list.

The actual start point is determined as M=m if m>0 and as M=l+m+1 if m<0. Similarly the actual end point is N=n if n>0 and N=l+n+1 if n<0. If M>N, the result is empty. Otherwise it consists of all items from position M to position N inclusive; for the purpose of this rule, we can imagine that the token list extends at infinity on either side, with void items at positions s for $s\leq 0$ or s>l. For instance,

```
\iow_term:x { \str_range:nnn { abcdef } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abcdef } { -4 } { -1 } }
\iow_term:x { \str_range:nnn { abcdef } { -2 } { -1 } }
\iow_term:x { \str_range:nnn { abcdef } { 0 } { -1 } }
```

prints bcde, cdef, ef, and an empty line to the terminal. The $\langle start\ index \rangle$ must always be smaller than or equal to the $\langle end\ index \rangle$: if this is not the case then no output is generated. Thus

```
\iow_term:x { \str_range:nnn { abcdef } { 5 } { 2 } }
\iow_term:x { \str_range:nnn { abcdef } { -1 } { -4 } }
```

both yield empty strings.

The behavior of \str_range_ignore_spaces:nnn is similar, but spaces are removed before starting the job. The input

```
\iow_term:x { \str_range:nnn { abcdefg } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abcdefg } { 2 } { -3 } }
\iow_term:x { \str_range:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range:nnn { abcdefg } { -6 } { -3 } }
\iow_term:x { \str_range:nnn { abc~efg } { 2 } { 5 } }
\iow_term:x { \str_range:nnn { abc~efg } { 2 } { -3 } }
\iow_term:x { \str_range:nnn { abc~efg } { 2 } { -3 } }
\iow_term:x { \str_range:nnn { abc~efg } { -6 } { 5 } }
\iow_term:x { \str_range:nnn { abc~efg } { -6 } { -3 } }
\iow_term:x { \str_range:nnn { abc~efg } { 2 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { 2 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { 2 } { -3 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { -6 } { 6 } { 6 } }
\invertext{ } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcdefg } { 6 } { 6 } { 6 } }
\invertext{ } }
\invertext{ } }
\invertext{ } \left\{ \str_range_ignore_spaces:nnn } { \str_range_ignore_spaces:nnn } }
```

```
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { 2 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { 2 } { -3 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { -6 } { 5 } }
\iow_term:x { \str_range_ignore_spaces:nnn { abcd~efg } { -6 } { -3 } }
```

will print four instances of bcde, four instances of bc e and eight instances of bcde.

16.6 Modifying string variables

\str_replace_once:Nnn \str_replace_once:cnn \str_greplace_once:Nnn \str_greplace_once:cnn $\verb|\str_replace_once:Nnn| \langle str| var \rangle \ \{\langle old \rangle\} \ \{\langle new \rangle\}$

Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces the first (leftmost) occurrence of $\langle old \ string \rangle$ in the $\langle str \ var \rangle$ with $\langle new \ string \rangle$.

New: 2017-10-08

\str_replace_all:Nnn \str_replace_all:cnn \str_greplace_all:Nnn \str_greplace_all:cnn

 $\str_replace_all:Nnn \langle str var \rangle \{ \langle old \rangle \} \{ \langle new \rangle \}$

Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces all occurrences of $\langle old \ string \rangle$ in the $\langle str \ var \rangle$ with $\langle new \ string \rangle$. As this function operates from left to right, the pattern $\langle old \ string \rangle$ may remain after the replacement (see \str_remove_all:Nn for an example).

New: 2017-10-08

 $\str_remove_once:Nn \langle str var \rangle \{\langle token list \rangle\}$

\str_remove_once:Nn \str_remove_once:cn \str_gremove_once:Nn \str_gremove_once:cn

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then removes the first (leftmost) occurrence of $\langle string \rangle$ from the $\langle str\ var \rangle$.

New: 2017-10-08

\str_remove_all:Nn \str_remove_all:cn

\str_gremove_all:Nn
\str_gremove_all:cn

New: 2017-10-08

 $\str_remove_all:Nn \langle str var \rangle \{\langle token list \rangle\}$

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then removes all occurrences of $\langle string \rangle$ from the $\langle str\ var \rangle$. As this function operates from left to right, the pattern $\langle string \rangle$ may remain after the removal, for instance,

results in \1_tmpa_str containing abcd.

16.7 String manipulation

```
\str_lowercase:n * \str_lowercase:n {\langle tokens \rangle}
\str_uppercase:f *
```

```
\str_lowercase: f * \str_uppercase: n {\langle tokens \rangle}
```

 $\str_uppercase:n \star Converts the input <math>\langle tokens \rangle$ to their string representation, as described for $\tl_to_$ str:n, and then to the lower or upper case representation using a one-to-one mapping New: 2019-11-26 as described by the Unicode Consortium file UnicodeData.txt.

These functions are intended for case changing programmatic data in places where upper/lower case distinctions are meaningful. One example would be automatically generating a function name from user input where some case changing is needed. In this situation the input is programmatic, not textual, case does have meaning and a languageindependent one-to-one mapping is appropriate. For example

```
\cs_new_protected:Npn \myfunc:nn #1#2
    \cs_set_protected:cpn
        user
        \str_uppercase:f { \tl_head:n {#1} }
        \str_lowercase:f { \tl_tail:n {#1} }
      }
      { #2 }
 }
```

would be used to generate a function with an auto-generated name consisting of the upper case equivalent of the supplied name followed by the lower case equivalent of the rest of the input.

These functions should *not* be used for

- Caseless comparisons: use \str_casefold:n for this situation (case folding is distinct from lower casing).
- Case changing text for typesetting: see the \text_lowercase:n(n), \text_uppercase:n(n) and \text_titlecase:n(n) functions which correctly deal with context-dependence and other factors appropriate to text case changing.

```
\str_casefold:n *
\str_casefold:V *
```

 $\str_casefold:n {\langle tokens \rangle}$

Converts the input \(\lambda to kens\rangle\) to their string representation, as described for \tl_to_str:n, New: 2022-10-16 and then folds the case of the resulting $\langle strinq \rangle$ to remove case information. The result of this process is left in the input stream.

> String folding is a process used for material such as identifiers rather than for "text". The folding provided by \str_casefold:n follows the mappings provided by the Unicode Consortium, who state:

Case folding is primarily used for caseless comparison of text, such as identifiers in a computer program, rather than actual text transformation. Case folding in Unicode is based on the lowercase mapping, but includes additional changes to the source text to help make it language-insensitive and consistent. As a result, case-folded text should be used solely for internal processing and generally should not be stored or displayed to the end user.

The folding approach implemented by \str_casefold:n follows the "full" scheme defined by the Unicode Consortium (e.g. SSfolds to SS). As case-folding is a language-insensitive process, there is no special treatment of Turkic input (i.e. I always folds to i and not to 1).

16.8 Viewing strings

\str_show:N \str_show:c \str_show:n $\str_show:N \langle str var \rangle$

Displays the content of the $\langle str \ var \rangle$ on the terminal.

New: 2015-09-18 Updated: 2021-04-29

\str_log:N \str_log:c \str_log:N \str var \rangle

Writes the content of the $\langle str \ var \rangle$ in the log file.

\str_log:n New: 2019-02-15 Updated: 2021-04-29

Constant strings 16.9

\c_ampersand_str \c_atsign_str \c_backslash_str \c_left_brace_str \c_right_brace_str \c_circumflex_str \c_colon_str \c_dollar_str \c_hash_str \c_percent_str \c_{tilde_str} \c_underscore_str \c_zero_str

> New: 2015-09-19 Updated: 2020-12-22

Scratch strings 16.10

\l_tmpa_str Scratch strings for local assignment. These are never used by the kernel code, and so \l_tmpb_str are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Constant strings, containing a single character token, with category code 12.

\g_tmpa_str Scratch strings for global assignment. These are never used by the kernel code, and so \g_tmpb_str are safe for use with any IATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Deprecated functions 16.11

 $\str_foldcase:n * \str_foldcase:n {\langle tokens \rangle}$

 $\frac{\texttt{\sc str_foldcase:V} \ \, \star}{\texttt{A previous name for the functionally-identical \sc str_casefold:n.}}$

New: 2019-11-26

Chapter 17

The l3str-convert package: String encoding conversions

17.1 Encoding and escaping schemes

Traditionally, string encodings only specify how strings of characters should be stored as bytes. However, the resulting lists of bytes are often to be used in contexts where only a restricted subset of bytes are permitted (e.g., PDF string objects, URLs). Hence, storing a string of characters is done in two steps.

- The code points ("character codes") are expressed as bytes following a given "encoding". This can be UTF-16, ISO 8859-1, etc. See Table 1 for a list of supported encodings.⁷
- Bytes are translated to T_EX tokens through a given "escaping". Those are defined for the most part by the pdf file format. See Table 2 for a list of escaping methods supported.⁷

 $^{^7\}mathrm{Encodings}$ and escapings will be added as they are requested.

Table 1: Supported encodings. Non-alphanumeric characters are ignored, and capital letters are lower-cased before searching for the encoding in this list.

$\langle Encoding \rangle$	description	
utf8	UTF-8	
utf16	UTF-16, with byte-order mark	
utf16be	UTF-16, big-endian	
utf16le	UTF-16, little-endian	
utf32	UTF-32, with byte-order mark	
utf32be	UTF-32, big-endian	
utf32le	UTF-32, little-endian	
iso88591, latin1	ISO 8859-1	
iso88592, latin2	ISO 8859-2	
iso88593, latin3	ISO 8859-3	
iso88594, latin4	ISO 8859-4	
iso88595	ISO 8859-5	
iso88596	ISO 8859-6	
iso88597	ISO 8859-7	
iso88598	ISO 8859-8	
iso88599, latin5	ISO 8859-9	
iso885910, latin6	ISO 8859-10	
iso885911	ISO 8859-11	
iso885913, latin7	ISO 8859-13	
iso885914, latin8	ISO 8859-14	
iso885915, latin9	ISO 8859-15	
iso885916, latin10	ISO 8859-16	
clist	comma-list of integers	
$\langle empty angle$	native (Unicode) string	
default	like utf8 with 8-bit engines, and like native with unicode-engines	

Table 2: Supported escapings. Non-alphanumeric characters are ignored, and capital letters are lower-cased before searching for the escaping in this list.

$\langle Escaping \rangle$	description
bytes, or empty	arbitrary bytes
hex, hexadecimal	byte = two hexadecimal digits
name	$\operatorname{see} \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $
string	$\operatorname{see} \setminus \mathtt{pdfescapestring}$
url	encoding used in URLs

17.2 Conversion functions

\str_set_convert:Nnnn \str_gset_convert:Nnnn

```
\str_set_convert:Nnnn \langle str var \rangle \{\langle string \rangle\} \{\langle name 1 \rangle\} \{\langle name 2 \rangle\}
```

This function converts the $\langle strinq \rangle$ from the encoding given by $\langle name\ 1 \rangle$ to the encoding given by $\langle name 2 \rangle$, and stores the result in the $\langle str \ var \rangle$. Each $\langle name \rangle$ can have the form $\langle encoding \rangle$ or $\langle encoding \rangle / \langle escaping \rangle$, where the possible values of $\langle encoding \rangle$ and $\langle escaping \rangle$ are given in Tables 1 and 2, respectively. The default escaping is to input and output bytes directly. The special case of an empty $\langle name \rangle$ indicates the use of "native" strings, 8-bit for pdfTFX, and Unicode strings for the other two engines.

For example,

```
\str_set_convert:Nnnn \l_foo_str { Hello! } { } { utf16/hex }
```

results in the variable \l_foo_str holding the string FEFF00480065006C006C006F0021. This is obtained by converting each character in the (native) string Hello! to the UTF-16 encoding, and expressing each byte as a pair of hexadecimal digits. Note the presence of a (big-endian) byte order mark "FEFF, which can be avoided by specifying the encoding utf16be/hex.

An error is raised if the $\langle strinq \rangle$ is not valid according to the $\langle escapinq 1 \rangle$ and $\langle encoding 1 \rangle$, or if it cannot be reencoded in the $\langle encoding 2 \rangle$ and $\langle escaping 2 \rangle$ (for instance, if a character does not exist in the $\langle encodinq 2 \rangle$). Erroneous input is replaced by the Unicode replacement character "FFFD, and characters which cannot be reencoded are replaced by either the replacement character "FFFD if it exists in the $\langle encoding 2 \rangle$, or an encoding-specific replacement character, or the question mark character.

\str_set_convert:Nnnn*TF* \str_gset_convert:NnnnTF

```
\verb|\str_set_convert:NnnnTF| \langle str| var \rangle \ \{\langle string \rangle\} \ \{\langle name | 1 \rangle\} \ \{\langle name | 2 \rangle\} \ \{\langle true| code \rangle\} 
\{\langle false\ code \rangle\}
```

As $\str_set_convert: Nnnn, converts the <math>\langle string \rangle$ from the encoding given by $\langle name\ 1 \rangle$ to the encoding given by $\langle name 2 \rangle$, and assigns the result to $\langle str \ var \rangle$. Contrarily to $\texttt{\sc t_set_convert:Nnn}$, the conditional variant does not raise errors in case the $\langle string \rangle$ is not valid according to the $\langle name\ 1 \rangle$ encoding, or cannot be expressed in the $\langle name\ 2 \rangle$ encoding. Instead, the $\langle false\ code \rangle$ is performed.

17.3 Conversion by expansion (for PDF contexts)

A small number of expandable functions are provided for use in PDF string/name contexts. These assume UTF-8 and no escaping in the input.

\str_convert_pdfname:n * \str_convert_pdfname:n \string\

As \str_set_convert: Nnnn, converts the \(\string \) on a byte-by-byte basis with non-ASCII codepoints escaped using hashes.

17.4Possibilities, and things to do

Encoding/escaping-related tasks.

- In X_TT_EX/LuaT_EX, would it be better to use the ^^^.... approach to build a string from a given list of character codes? Namely, within a group, assign 0-9a-f and all characters we want to category "other", then assign ^ the category superscript, and use \scantokens.
- Change \str_set_convert:Nnnn to expand its last two arguments.
- Describe the internal format in the code comments. Refuse code points in ["D800, "DFFF] in the internal representation?
- Add documentation about each encoding and escaping method, and add examples.
- The hex unescaping should raise an error for odd-token count strings.
- Decide what bytes should be escaped in the url escaping. Perhaps the characters !'()*-./0123456789_ are safe, and all other characters should be escaped?
- Automate generation of 8-bit mapping files.
- Change the framework for 8-bit encodings: for decoding from 8-bit to Unicode, use 256 integer registers; for encoding, use a tree-box.
- More encodings (see Heiko's stringenc). CESU?
- More escapings: ASCII85, shell escapes, lua escapes, etc.?

Chapter 18

The **I3quark** package Quarks

Two special types of constants in LATeX3 are "quarks" and "scan marks". By convention all constants of type quark start out with $\q_$, and scan marks start with $\s_$.

18.1 Quarks

Quarks are control sequences (and in fact, token lists) that expand to themselves and should therefore *never* be executed directly in the code. This would result in an endless loop!

They are meant to be used as delimiter in weird functions, the most common use case being the 'stop token' (i.e. \q_stop). For example, when writing a macro to parse a user-defined date

```
\date_parse:n {19/June/1981}
one might write a command such as
\cs_new:Npn \date_parse:n #1 { \date_parse_aux:w #1 \q_stop }
\cs_new:Npn \date_parse_aux:w #1 / #2 / #3 \q_stop
{ <do something with the date> }
```

Quarks are sometimes also used as error return values for functions that receive erroneous input. For example, in the function \prop_get:NnN to retrieve a value stored in some key of a property list, if the key does not exist then the return value is the quark \q_no_value. As mentioned above, such quarks are extremely fragile and it is imperative when using such functions that code is carefully written to check for pathological cases to avoid leakage of a quark into an uncontrolled environment.

Quarks also permit the following ingenious trick when parsing tokens: when you pick up a token in a temporary variable and you want to know whether you have picked up a particular quark, all you have to do is compare the temporary variable to the quark using \tl_if_eq:NNTF. A set of special quark testing functions is set up below. All the quark testing functions are expandable although the ones testing only single tokens are much faster.

18.2 Defining quarks

 $\verb|\quark_new:N \quark_new:N \quark|$

Creates a new $\langle quark \rangle$ which expands only to $\langle quark \rangle$. The $\langle quark \rangle$ is defined globally, and an error message is raised if the name was already taken.

\q_stop Used as a marker for delimited arguments, such as

```
\cs_set:Npn \tmp:w #1#2 \q_stop {#1}
```

\q_mark Used as a marker for delimited arguments when \q_stop is already in use.

\q_nil Quark to mark a null value in structured variables or functions. Used as an end delimiter when this may itself need to be tested (in contrast to \q_stop, which is only ever used as a delimiter).

A canonical value for a missing value, when one is requested from a data structure. This is therefore used as a "return" value by functions such as \prop_get:NnN if there is no data to return.

18.3 Quark tests

The method used to define quarks means that the single token (N) tests are faster than the multi-token (n) tests. The latter should therefore only be used when the argument can definitely take more than a single token.

```
Tests if the \langle token \rangle is equal to q_nil.
  \quark_if_nil_p:n
                            * \quark_if_nil_p:n {\langle token list \rangle}
                           \star \qquad tin_nin \{ (token list) \} \{ (true code) \} \{ (false code) \}
  \quark_if_nil_p:(o|V)
  \quark_if_nil:n<u>TF</u>
                              Tests if the \langle token\ list \rangle contains only \neqnil (distinct from \langle token\ list \rangle being empty or
  \quark_if_nil:(o|V) <u>TF</u>
                              containing \q_{nil} plus one or more other tokens).
\quark_if_no_value_p:N * \quark_if_no_value_p:N \(\lambda token\rangle)
                            * \quark_if_no_value:NTF \langle token \rangle \langle true code \rangle \rangle \langle false code \rangle \rangle
\quark_if_no_value_p:c
\quark_if_no_value:NTF
                              Tests if the \langle token \rangle is equal to q_no_value.
\quark_if_no_value:c<u>TF</u>
\label{eq:continuous} $$\operatorname{quark_if_no\_value_p:n } {\langle token \; list \rangle}$
\quark_if_no_value:n<u>TF</u>
                               \verb|\quark_if_no_value:nTF {|\langle token \ list \rangle}| \ \{\langle true \ code \rangle\}| \ \{\langle false \ code \rangle\}|
```

Tests if the $\langle token \ list \rangle$ contains only \q_no_value (distinct from $\langle token \ list \rangle$ being empty or containing \q_no_value plus one or more other tokens).

18.4 Recursion

This module provides a uniform interface to intercepting and terminating loops as when one is doing tail recursion. The building blocks follow below and an example is shown in Section 18.4.1.

\q_recursion_tail This quark is appended to the data structure in question and appears as a real element there. This means it gets any list separators around it.

\q_recursion_stop This quark is added after the data structure. Its purpose is to make it possible to terminate the recursion at any point easily.

Tests if \(\lambda token \rangle\) contains only the marker \(\mathbb{q}\)_recursion_tail, and if so uses \(\mathbb{u}\)se_none_delimit_by_q_recursion_stop:w to terminate the recursion that this belongs to. The recursion input must include the marker tokens \q_recursion_tail and \q_recursion_stop as the last two items.

```
\quark_if_recursion_tail_stop:n * \quark_if_recursion_tail_stop:n {\langle token list\rangle}
\quark_if_recursion_tail_stop:o *
                     Updated: 2011-09-06
```

Tests if the \(\langle token \ list \rangle \) contains only \q_recursion_tail, and if so uses \use_none_delimit_by_q_recursion_stop:w to terminate the recursion that this belongs to. The recursion input must include the marker tokens \q_recursion_tail and \q_recursion_stop as the last two items.

```
\quark_if_recursion_tail_stop_do: Nn * \quark_if_recursion_tail_stop_do: Nn $\langle token \rangle$ {\langle insertion \rangle}
```

Tests if \(\lambda token \rangle \) contains only the marker \(\mathbb{q}\)_recursion_tail, and if so uses \(\mathbb{u}\)se_i_delimit_by_q_recursion_stop:w to terminate the recursion that this belongs to. The recursion input must include the marker tokens \q_recursion_tail and \q_recursion stop as the last two items. The $\langle insertion \rangle$ code is then added to the input stream after the recursion has ended.

```
\quark_if_recursion_tail\_stop_do:nn * \quark_if_recursion_tail\_stop_do:nn \{\langle token \ list \rangle\} \{\langle insertion \rangle\}
\quark_if_recursion_tail_stop_do:on *
                            Updated: 2011-09-06
```

Tests if the \(\langle token \ list \rangle \) contains only \q_recursion_tail, and if so uses \use_i_delimit_by_q_recursion_stop:w to terminate the recursion that this belongs to. The recursion input must include the marker tokens \q_recursion_tail and \q_recursion_stop as the last two items. The $\langle insertion \rangle$ code is then added to the input stream after the recursion has ended.

Tests if $\langle token \ list \rangle$ contains only $\q_recursion_tail$, and if so terminates the recursion using $\t vpe \mbox{\ensuremath{\mbox{contains}}} \mbox{\ensuremath{\mbox{c$

18.4.1 An example of recursion with quarks

Quarks are mainly used internally in the expl3 code to define recursion functions such as $\tl_map_inline:nn$ and so on. Here is a small example to demonstrate how to use quarks in this fashion. We shall define a command called $\mbox{map_dbl:nn}$ which takes a token list and applies an operation to every pair of tokens. For example, $\mbox{my_map_dbl:nn {abcd} {[--#1--#2--]^}}$ would produce "[-a-b-] [-c-d-] ". Using quarks to define such functions simplifies their logic and ensures robustness in many cases.

Here's the definition of \my_map_dbl:nn. First of all, define the function that does the processing based on the inline function argument #2. Then initiate the recursion using an internal function. The token list #1 is terminated using \q_recursion_tail, with delimiters according to the type of recursion (here a pair of \q_recursion_tail), concluding with \q_recursion_stop. These quarks are used to mark the end of the token list being operated upon.

```
\cs_new:Npn \my_map_dbl:nn #1#2
{
   \cs_set:Npn \__my_map_dbl_fn:nn ##1 ##2 {#2}
   \_my_map_dbl:nn #1 \q_recursion_tail \q_recursion_tail
   \q_recursion_stop
}
```

The definition of the internal recursion function follows. First check if either of the input tokens are the termination quarks. Then, if not, apply the inline function to the two arguments.

```
\cs_new:Nn \__my_map_dbl:nn
{
    \quark_if_recursion_tail_stop:n {#1}
    \quark_if_recursion_tail_stop:n {#2}
    \__my_map_dbl_fn:nn {#1} {#2}

Finally, recurse:
    \__my_map_dbl:nn
}
```

Note that contrarily to LATEX3 built-in mapping functions, this mapping function cannot be nested, since the second map would overwrite the definition of __my_map_dbl_fn:nn.

18.5 Scan marks

Scan marks are control sequences set equal to \scan_stop:, hence never expand in an expansion context and are (largely) invisible if they are encountered in a typesetting context.

Like quarks, they can be used as delimiters in weird functions and are often safer to use for this purpose. Since they are harmless when executed by TEX in non-expandable contexts, they can be used to mark the end of a set of instructions. This allows to skip to that point if the end of the instructions should not be performed (see I3regex).

\scan_new:N \scan_new:N \scan mark \

New: 2018-04-01 Creates a new \(\scan \text{mark}\) which is set equal to \scan_stop:. The \(\scan \text{mark}\) is defined globally, and an error message is raised if the name was already taken by another scan mark.

\s_stop

Used at the end of a set of instructions, as a marker that can be jumped to using \use_- ${\tt New:\,2018-04-01}\ \, {\tt none_delimit_by_s_stop:w}.$

\use_none_delimit_by_s_stop:w * \use_none_delimit_by_s_stop:w \land tokens \rangle \s_stop New: 2018-04-01

> Removes the \(\lambda to kens\rangle\) and \s_stop from the input stream. This leads to a low-level TeX error if \s_stop is absent.

Chapter 19

The **I3seq** package Sequences and stacks

LATEX3 implements a "sequence" data type, which contain an ordered list of entries which may contain any $\langle balanced\ text \rangle$. It is possible to map functions to sequences such that the function is applied to every item in the sequence.

Sequences are also used to implement stack functions in IATEX3. This is achieved using a number of dedicated stack functions.

19.1 Creating and initialising sequences

```
\seq_new:N \seq_new:N \sequence
                             Creates a new \langle sequence \rangle or raises an error if the name is already taken. The declaration
                             is global. The \langle sequence \rangle initially contains no items.
           \seq_clear:N \seq_clear:N \sequence \
           \seq_clear:c
                             Clears all items from the \langle sequence \rangle.
           \seq_gclear:N
           \seq_gclear:c
      \seq_clear_new:N \seq_clear_new:N \sequence
     \seq_clear_new:c \seq_gclear_new:N Ensures that the \( sequence \) exists globally by applying \seq_new:N if necessary, then
      \seq_gclear_new: applies \seq_(g) clear: N to leave the \( \sequence \) empty.
\seq_set_eq:NN
                             \ensuremath{\mbox{\sc sequence}_1}\ensuremath{\mbox{\sc sequence}_2}\ensuremath{\mbox{\sc sequence}_2}
\seq_set_eq:(cN|Nc|cc)
                             Sets the content of \langle sequence_1 \rangle equal to that of \langle sequence_2 \rangle.
\seq_gset_eq:NN
\seq_gset_eq:(cN|Nc|cc)
```

```
\seq_set_from_clist:NN
                                   \seq_set_from_clist:NN \langle sequence \rangle \langle comma-list \rangle
\seq_set_from_clist:(cN|Nc|cc)
\seq_set_from_clist:Nn
\seq_set_from_clist:cn
\seq_gset_from_clist:NN
\seq_gset_from_clist:(cN|Nc|cc)
\seq_gset_from_clist:Nn
\seq_gset_from_clist:cn
```

New: 2014-07-17

Converts the data in the $\langle comma | list \rangle$ into a $\langle sequence \rangle$: the original $\langle comma | list \rangle$ is

\seq_const_from_clist:cn New: 2017-11-28

 $\scalebox{$\scalebox{\sim}} \scalebox{\sim} \scaleb$

Creates a new constant $\langle seq \ var \rangle$ or raises an error if the name is already taken. The $\langle seq \ var \rangle$ is set globally to contain the items in the $\langle comma \ list \rangle$.

\seq_set_split:NnV \seq_gset_split:Nnn \seq_gset_split:NnV

Updated: 2012-07-02

 $\seq_set_split:Nnn \seq_set_split:Nnn \seq_set_sp$

Splits the $\langle token\ list\rangle$ into $\langle items\rangle$ separated by $\langle delimiter\rangle$, and assigns the result to the $\langle sequence \rangle$. Spaces on both sides of each $\langle item \rangle$ are ignored, then one set of outer braces is removed (if any); this space trimming behaviour is identical to that of I3clist functions. $^{\text{New: 2011-08-15}}$ Empty $\langle items \rangle$ are preserved by seq_set_split:Nnn , and can be removed afterwards using \seq_remove_all: $Nn \langle sequence \rangle$ {}. The $\langle delimiter \rangle$ may not contain {, } or # (assuming TEX's normal category code régime). If the $\langle delimiter \rangle$ is empty, the $\langle token \rangle$ list is split into $\langle items \rangle$ as a $\langle token \ list \rangle$. See also \seq_set_split_keep_spaces:Nnn, which omits space stripping.

\seq_set_split_keep_spaces:NnV \seq_gset_split_keep_spaces:Nnn \seq_gset_split_keep_spaces:NnV New: 2021-03-24

 $\seq_set_split_keep_spaces:Nnn \seq_set_split_keep_spaces:Nnn \seq_set_split_keep_spaces:Nn$

Splits the $\langle token \ list \rangle$ into $\langle items \rangle$ separated by $\langle delimiter \rangle$, and assigns the result to the (sequence). One set of outer braces is removed (if any) but any surrounding spaces are retained: any braces inside one or more spaces are therefore kept. Empty (items) are preserved by \seq_set_split_keep_spaces:Nnn, and can be removed afterwards using \seq_remove_all:Nn $\langle sequence \rangle$ {}. The $\langle delimiter \rangle$ may not contain {, } or # (assuming T_FX's normal category code régime). If the $\langle delimiter \rangle$ is empty, the $\langle token \rangle$ list into (items) as a (token list). See also \seq_set_split:Nnn, which removes spaces around the delimiters.

\seq_concat:NNN \seq_concat:ccc \seq_gconcat:NNN \seq_gconcat:ccc

 $\seq_{concat:NNN} \ \langle sequence_1 \rangle \ \langle sequence_2 \rangle \ \langle sequence_3 \rangle$

Concatenates the content of $\langle sequence_2 \rangle$ and $\langle sequence_3 \rangle$ together and saves the result in $\langle sequence_1 \rangle$. The items in $\langle sequence_2 \rangle$ are placed at the left side of the new sequence.

```
\label{eq:code} $$ \left(\frac{false\ code}{seq\_if\_exist\_p:N\ } \times \left(\frac{false\ code}{seq\_if\_exist:NTF\ } \times (\frac{false\ code}{seq\_if\_ex
```

19.2 Appending data to sequences

Appends the $\langle item \rangle$ to the right of the $\langle sequence \rangle$.

19.3 Recovering items from sequences

Items can be recovered from either the left or the right of sequences. For implementation reasons, the actions at the left of the sequence are faster than those acting on the right. These functions all assign the recovered material locally, *i.e.* setting the $\langle token \ list \ variable \rangle$ used with $tl_set:Nn$ and $never \ tl_gset:Nn$.

```
\seq_get_left:NN \seq_get_left:NN \sequence \sequence \int token list variable \seq_get_left:cn
\text{Updated: 2012-05-14} Stores the left-most item from a \sequence \int the \sequence \int the \sequence \int the \sequence \int the \sequence \int token list variable \int is assigned locally. If \sequence \int the \seq_uence \int is empty the \seq_get_right:NN \seq_get_r
```

```
\seq_pop_left:NN \seq_pop_left:NN \seq_pop_left:NN \seq_pop_left:cN \seq_p
```

Pops the left-most item from a $\langle sequence \rangle$ into the $\langle token \ list \ variable \rangle$, i.e. removes the item from the sequence and stores it in the $\langle token \ list \ variable \rangle$. Both of the variables are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token \ list \ variable \rangle$ is set to the special marker q_no_value .

\seq_gpop_left:cN

\seq_gpop_left:NN \seq_gpop_left:NN \sequence \daggeright \taken list variable \daggeright

Updated: 2012-05-14

Pops the left-most item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$, i.e. removes the item from the sequence and stores it in the $\langle token\ list\ variable \rangle$. The $\langle sequence \rangle$ is modified globally, while the assignment of the $\langle token\ list\ variable \rangle$ is local. If $\langle sequence \rangle$ is empty the \(\langle token list variable \rangle\) is set to the special marker \q_no_value.

\seq_pop_right:cN

Pops the right-most item from a \(\sequence \) into the \(\text{token list variable} \), i.e. removes the Updated: 2012-05-19 item from the sequence and stores it in the $\langle token\ list\ variable \rangle$. Both of the variables are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token\ list\ variable \rangle$ is set to the special marker \q_no_value.

\seq_gpop_right:NN \seq_gpop_right:cN

\seq_gpop_right:NN \(\sequence \) \(\taken list variable \)

Pops the right-most item from a (sequence) into the (token list variable), i.e. removes Updated: 2012-05-19 the item from the sequence and stores it in the $\langle token\ list\ variable \rangle$. The $\langle sequence \rangle$ is modified globally, while the assignment of the $\langle token\ list\ variable \rangle$ is local. If $\langle sequence \rangle$ is empty the \(\langle token list variable \rangle\) is set to the special marker \q_no_value.

 $\seq_item:Nn \star \seq_item:Nn \seq_uence$ {\langle integer expression \rangle}

\seq_item:cn *

Indexing items in the (sequence) from 1 at the top (left), this function evaluates the New: 2014-07-17 (integer expression) and leaves the appropriate item from the sequence in the input stream. If the (integer expression) is negative, indexing occurs from the bottom (right) of the sequence. If the (integer expression) is larger than the number of items in the \(\langle sequence \rangle \) (as calculated by \seq_count:N) then the function expands to nothing.

> TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

\seq_rand_item:N

\seq_rand_item:N \langle seq_var \rangle

\seq_rand_item:c *

Selects a pseudo-random item of the $\langle sequence \rangle$. If the $\langle sequence \rangle$ is empty the result is New: 2016-12-06 empty. This is not available in older versions of XTFX.

> TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

19.4 Recovering values from sequences with branching

The functions in this section combine tests for non-empty sequences with recovery of an item from the sequence. They offer increased readability and performance over separate testing and recovery phases.

\seq_get_left:NNTF

 $\label{lem:list_variable} $$ \left(\text{true code} \right) $$ \left(\text{false code} \right) $$$

\seq_get_left:cN<u>TF</u>

New: 2012-05-14

Updated: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the left-most item from the $\langle sequence \rangle$ in the $\langle token\ list$ variable without removing it from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

\seq_get_right:NNTF \seq_get_right:cNTF

\seq_get_right:NNTF \(\sequence \) \(\tau \) \(\tau \) \(\lambda \) \(\tau \) \(\t

New: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the right-most item from the $\langle sequence \rangle$ in the $\langle token \rangle$ list variable without removing it from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

\seq_pop_left:cNTF

 $\ \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \end{array} \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right) \left(\begin{array}{c} seq_pop_left:NNTF \ \end{array} \right$

New: 2012-05-14 Updated: 2012-05-19 If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the left-most item from the $\langle sequence \rangle$ in the $\langle token\ list$ variable, i.e. removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. Both the $\langle sequence \rangle$ and the $\langle token\ list\ variable \rangle$ are assigned locally.

\seq_gpop_left:cN<u>TF</u>

 $\label{lem:code} $$ \left(\frac{pop_left:NNTF}{seq_pop_left:NNTF} \left(\frac{code}{code}\right) \left(\frac{code}{code}\right) \right] $$$

Updated: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of New: 2012-05-14 the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the left-most item from the $\langle sequence \rangle$ in the $\langle token$ list variable, i.e. removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle sequence \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

\seq_pop_right:cN<u>TF</u>

New: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the right-most item from the $\langle sequence \rangle$ in the $\langle token\ list$ variable, i.e. removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. Both the $\langle sequence \rangle$ and the $\langle token\ list\ variable \rangle$ are assigned locally.

\seq_gpop_right:NNTF \seq_gpop_right:cNTF

 $\ensuremath{\verb|seq_gpop_right:NNTF||} \langle sequence \rangle \ensuremath{|} \langle token \ensuremath{|} list \ensuremath{|} variable \rangle \ensuremath{|} \langle true \ensuremath{|} code \rangle \} \ensuremath{|} \langle false \ensuremath{|} code \rangle \}$

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of New: 2012-05-19 the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the right-most item from the $\langle sequence \rangle$ in the $\langle token$ list variable, i.e. removes the item from the $\langle sequence \rangle$, then leaves the $\langle true\ code \rangle$ in the input stream. The $\langle sequence \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

19.5 Modifying sequences

While sequences are normally used as ordered lists, it may be necessary to modify the content. The functions here may be used to update sequences, while retaining the order of the unaffected entries.

\seq_remove_duplicates:N
\seq_remove_duplicates:C
\seq_gremove_duplicates:N
\seq_gremove_duplicates:c

 $\verb|\seq_remove_duplicates:N| \langle sequence \rangle|$

Removes duplicate items from the $\langle sequence \rangle$, leaving the left most copy of each item in the $\langle sequence \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for $\t_i=eq:nnTF$.

T_EXhackers note: This function iterates through every item in the $\langle sequence \rangle$ and does a comparison with the $\langle items \rangle$ already checked. It is therefore relatively slow with large sequences.

\seq_remove_all:Nn
\seq_remove_all:cn
\seq_gremove_all:Nn
\seq_gremove_all:cn

\seq_remove_all:Nn \langle sequence \rangle \langle \tau item \rangle \}

Removes every occurrence of $\langle item \rangle$ from the $\langle sequence \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for $\t = if_eq:nnTF$.

\seq_reverse:N
\seq_reverse:C
\seq_greverse:N

\seq_reverse:N \langle sequence \rangle

Reverses the order of the items stored in the (sequence).

\seq_greverse:c

New: 2014-07-18

\seq_sort:cn
\seq_gsort:Nn
\seq_gsort:cn

New: 2017-02-06

Sorts the items in the $\langle sequence \rangle$ according to the $\langle comparison\ code \rangle$, and assigns the result to $\langle sequence \rangle$. The details of sorting comparison are described in Section 6.1.

\seq_shuffle:N \seq_shuffle:c

\seq_shuffle:N \(seq var \)

\seq_gshuffle:N \seq_gshuffle:c Sets the $\langle seq \ var \rangle$ to the result of placing the items of the $\langle seq \ var \rangle$ in a random order. Each item is (roughly) as likely to end up in any given position.

New: 2018-04-29

TEXhackers note: For sequences with more than 13 items or so, only a small proportion of all possible permutations can be reached, because the random seed \sys_rand_seed: only has 28-bits. The use of \toks internally means that sequences with more than 32767 or 65535 items (depending on the engine) cannot be shuffled.

19.6 Sequence conditionals

```
\label{eq:code} $$ \left(\frac{f_{empty_p:N} \times \left(\frac{f_{empty_p:N} \left(\frac{f_{empty_p:N}}{f_{empty_p:C} \times \left(\frac{f_{empty_p:N}}{f_{empty_p:C} \times \left(\frac{f_{empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_empty_em
```

```
\seq_if_in:NnTF
                                                                               \ensuremath{\mbox{seq\_if\_in:NnTF}} \ensuremath{\mbox{\mbox{sequence}}} \ensuremath{\mbox{\mbox{\mbox{$\langle$ true\ code}$\rangle$}} \ensuremath{\mbox{$\langle$ false\ code}$\rangle$}
\sqrt{\sqrt{Nv|Nv|No|Nx|cn|cV|cv|co|cx}}
```

Tests if the $\langle item \rangle$ is present in the $\langle sequence \rangle$.

19.7 Mapping over sequences

All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

\seq_map_function:NN 🌣 \seq_map_function:cN ☆ Updated: 2012-06-29

\seq_map_function:NN \langle sequence \rangle \langle function \rangle

Applies $\langle function \rangle$ to every $\langle item \rangle$ stored in the $\langle sequence \rangle$. The $\langle function \rangle$ will receive one argument for each iteration. The $\langle items \rangle$ are returned from left to right. To pass further arguments to the $\langle function \rangle$, see \seq_map_tokens: Nn. The function \seq_map_inline: Nn is faster than \seq_map_function: NN for sequences with more than about 10 items.

\seq_map_inline: Nn \seq_map_inline: Nn \sequence \ {\langle inline function \}

\seq_map_inline:cn

Updated: 2012-06-29

Applies $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence \rangle$. The $\langle inline\ function \rangle$ function should consist of code which will receive the $\langle item \rangle$ as #1. The $\langle items \rangle$ are returned from left to right.

\seq_map_tokens:cn ☆

Analogue of \seq_map_function: NN which maps several tokens instead of a single func-New: 2019-08-30 tion. The $\langle code \rangle$ receives each item in the $\langle sequence \rangle$ as a trailing brace group. For instance,

```
\seq_map_tokens:Nn \l_my_seq { \prg_replicate:nn { 2 } }
```

expands to twice each item in the \(sequence \): for each item in \\1_my_seq the function \prg_replicate:nn receives 2 and \(\lambda item\rangle\) as its two arguments. The function \seq_map inline: Nn is typically faster but it is not expandable.

\seq_map_variable:NNn \seq_map_variable:(Ncn|cNn|ccn)

Updated: 2012-06-29

Stores each $\langle item \rangle$ of the $\langle sequence \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle sequence \rangle$, or its original value if the $\langle sequence \rangle$ is empty. The $\langle items \rangle$ are returned from left to right.

\seq_map_indexed_function:NN & \seq_map_indexed_function:NN \langle seq_var \rangle \langle function \rangle

New: 2018-05-03

Applies $\langle function \rangle$ to every entry in the $\langle sequence\ variable \rangle$. The $\langle function \rangle$ should have signature: nn. It receives two arguments for each iteration: the $\langle index \rangle$ (namely 1 for the first entry, then 2 and so on) and the $\langle item \rangle$.

 $\ensuremath{\texttt{Seq_map_indexed_inline}}$ \seq_map_indexed_inline: Nn $\ensuremath{\texttt{Seq\ var}}\ \{\ensuremath{\texttt{(inline\ function)}}\}$

New: 2018-05-03 Applies $\langle inline\ function \rangle$ to every entry in the $\langle sequence\ variable \rangle$. The $\langle inline\ function \rangle$ should consist of code which receives the $\langle index \rangle$ (namely 1 for the first entry, then 2 and so on) as #1 and the $\langle item \rangle$ as #2.

\seq_map_break: ☆ \seq_map_break:

Updated: 2012-06-29 Used to terminate a \searrow eq_map_... function before all entries in the \langle sequence \rangle have been processed. This normally takes place within a conditional statement, for example

```
\seq_map_inline:Nn \l_my_seq
    \str_if_eq:nnTF { #1 } { bingo }
      { \seq_map_break: }
      {
        % Do something useful
 }
```

Use outside of a \seq_map_... scenario leads to low level TEX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

 $\scalebox{ } seq_map_break:n $$ \scalebox{ } seq_map_break:n $$ (code)$$

Updated: 2012-06-29 Used to terminate a \seq_map_... function before all entries in the \(sequence \) have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\seq_map_inline: Nn \l_my_seq
    \str_if_eq:nnTF { #1 } { bingo }
      { \seq_map_break:n { <code> } }
        % Do something useful
  }
```

Use outside of a \seq_map_... scenario leads to low level TFX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

```
\seq_set_map:NNn
\seq_gset_map:NNn
```

```
\seq_set_map:NNn \ \langle sequence_1 \rangle \ \langle sequence_2 \rangle \ \{\langle inline\ function \rangle\}
```

Applies $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence_2 \rangle$. The $\langle inline\ function \rangle$ New: 2011-12-22 function should consist of code which will receive the $\langle item \rangle$ as #1. The sequence result-Updated: 2020-07-16 ing applying $\langle inline\ function \rangle$ to each $\langle item \rangle$ is assigned to $\langle sequence_1 \rangle$.

> TrXhackers note: Contrarily to other mapping functions, \seq map break: cannot be used in this function, and would lead to low-level TFX errors.

\seq_set_map_x:NNn $\seq_gset_map_x:NNn$

```
\seq_set_map_x:NNn \seq_uence_1 \seq_uence_2 \ {\langle inline \ function \rangle}
```

New: 2020-07-16

Applies $\langle inline\ function \rangle$ to every $\langle item \rangle$ stored within the $\langle sequence_2 \rangle$. The $\langle inline\ function \rangle$ function should consist of code which will receive the $\langle item \rangle$ as #1. The sequence resulting from x-expanding $\langle inline\ function \rangle$ applied to each $\langle item \rangle$ is assigned to $\langle sequence_1 \rangle$. As such, the code in $\langle inline\ function \rangle$ should be expandable.

TEXhackers note: Contrarily to other mapping functions, \seq_map_break: cannot be used in this function, and would lead to low-level TEX errors.

\seq_count:c *

\seq_count:N * \seq_count:N \langle sequence \rangle

Leaves the number of items in the $\langle sequence \rangle$ in the input stream as an $\langle integer \rangle$ New: 2012-07-13 denotation). The total number of items in a $\langle sequence \rangle$ includes those which are empty and duplicates, *i.e.* every item in a $\langle sequence \rangle$ is unique.

19.8 Using the content of sequences directly

```
\verb|\seq_use:Nnnn| \star \verb|\seq_use:Nnnn| \langle seq \ var \rangle \ \{ \langle separator \ between \ two \rangle \}
```

 $\ensuremath{\mbox{seq_use:cnnn}} \star \{\ensuremath{\mbox{\langle separator between more than two}}\} \$

New: 2013-05-26 Places the contents of the $\langle seq \, var \rangle$ in the input stream, with the appropriate $\langle separator \rangle$ between the items. Namely, if the sequence has more than two items, the \(\separator \) between more than two is placed between each pair of items except the last, for which the $\langle separator\ between\ final\ two\rangle$ is used. If the sequence has exactly two items, then they are placed in the input stream separated by the $\langle separator\ between\ two \rangle$. If the sequence has a single item, it is placed in the input stream, and an empty sequence produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\seq_set_split: Nnn \l_tmpa_seq { | } { a | b | c | {de} | f }
\seq_use:Nnnn \l_tmpa_seq { ~and~ } { ,~ } { ,~and~ }
```

inserts "a, b, c, de, and f" in the input stream. The first separator argument is not used in this case because the sequence has more than 2 items.

TeXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the (items) do not expand further when appearing in an x-type or e-type argument expansion.

```
\seq_use:Nn
\seq_use:cn *
```

```
\star \sl (seq\_use:Nn \sl (seq\_var) \sl (separator))
```

Places the contents of the $\langle seq \ var \rangle$ in the input stream, with the $\langle separator \rangle$ between New: 2013-05-26 the items. If the sequence has a single item, it is placed in the input stream with no (separator), and an empty sequence produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\seq_set_split:Nnn \l_tmpa_seq { | } { a | b | c | {de} | f }
\seq_use: Nn \l_tmpa_seq { ~and~ }
```

inserts "a and b and c and de and f" in the input stream.

TrXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the (items) do not expand further when appearing in an x-type or e-type argument expansion.

19.9Sequences as stacks

Sequences can be used as stacks, where data is pushed to and popped from the top of the sequence. (The left of a sequence is the top, for performance reasons.) The stack functions for sequences are not intended to be mixed with the general ordered data functions detailed in the previous section: a sequence should either be used as an ordered data type or as a stack, but not in both ways.

\seq_get:NN \seq_get:cN \seq_get:NN \langle sequence \rangle \token list variable \rangle

Reads the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$ without removing it Updated: 2012-05-14 from the $\langle sequence \rangle$. The $\langle token\ list\ variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the \(\lambda token list variable\)\) is set to the special marker \q_no_value.

\seq_pop:NN \seq_pop:cN \seq_pop:NN \langle sequence \rangle \token list variable \rangle

Pops the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$. Both of the variables Updated: 2012-05-14 are assigned locally. If $\langle sequence \rangle$ is empty the $\langle token\ list\ variable \rangle$ is set to the special marker \q_no_value.

\seq_gpop:NN \seq_gpop:cN \seq_gpop:NN \(\langle sequence \rangle \) \(\tag{token list variable} \)

Pops the top item from a $\langle sequence \rangle$ into the $\langle token\ list\ variable \rangle$. The $\langle sequence \rangle$ is Updated: 2012-05-14 modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally. If $\langle sequence \rangle$ is empty the \(\langle token list variable \rangle\) is set to the special marker \q no value.

\seq_get:NNTF \seq_get:cNTF \seq_get:NNTF \langle sequence \rangle \tau token list variable \rangle \langle \tau true code \rangle \rangle \langle \family \langle \tau true code \rangle \rangle \langle \family \langle \tau true code \rangle \rangle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \langle \family \langle \tau true code \rangle \rangle \langle \langle \langle \family \langle \langle \family \langle \langle \langle \langle \langle \langle \langle \family \langle \l

New: 2012-05-14 Updated: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, stores the top item from a $\langle sequence \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle sequence \rangle$. The $\langle token\ list\ variable \rangle$ is assigned locally.

\seq_pop:NNTF \seq_pop:cNTF

\seq_pop:NNTF \langle sequence \rangle \tau token list variable \rangle \langle \tau token \rangle \rangle \tau token \rangle \rangle \tau token \rangle \rangle \rangle \tau token \rangle \rangle \rangle \rangle \tau token \rangle \rangle

Updated: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of New: 2012-05-14 the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the top item from the $\langle sequence \rangle$ in the $\langle token\ list$ variable, i.e. removes the item from the $\langle sequence \rangle$. Both the $\langle sequence \rangle$ and the $\langle token$ *list variable*\) are assigned locally.

\seq_gpop:NNTF \seq_gpop:cNTF

 $\ensuremath{\verb|seq_gpop:NNTF||} \langle sequence \rangle \ensuremath{|} \langle token \ensuremath{|} list \ensuremath{|} variable \rangle \ensuremath{|} \langle true \ensuremath{|} code \rangle \} \ensuremath{|} \langle false \ensuremath{|} code \rangle \}$

Updated: 2012-05-19

If the $\langle sequence \rangle$ is empty, leaves the $\langle false\ code \rangle$ in the input stream. The value of New: 2012-05-14 the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle sequence \rangle$ is non-empty, pops the top item from the $\langle sequence \rangle$ in the $\langle token \ list$ variable, i.e. removes the item from the $\langle sequence \rangle$. The $\langle sequence \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

\seq_push:Nn

\seq_push:Nn \langle sequence \rangle \langle (item \rangle)

 $\scalebox{Neq_push:} (NV|Nv|No|Nx|cn|cV|cv|co|cx)$

\seq_gpush:Nn

 $\scalebox{ } seq_gpush: (NV|Nv|No|Nx|cn|cV|cv|co|cx)$

Adds the $\{\langle item \rangle\}$ to the top of the $\langle sequence \rangle$.

19.10 Sequences as sets

Sequences can also be used as sets, such that all of their items are distinct. Usage of sequences as sets is not currently widespread, hence no specific set function is provided. Instead, it is explained here how common set operations can be performed by combining several functions described in earlier sections. When using sequences to implement sets, one should be careful not to rely on the order of items in the sequence representing the set.

Sets should not contain several occurrences of a given item. To make sure that a ⟨sequence variable⟩ only has distinct items, use \seq_remove_duplicates:N ⟨sequence variable. This function is relatively slow, and to avoid performance issues one should only use it when necessary.

Some operations on a set $\langle seq \, var \rangle$ are straightforward. For instance, \seq count: N $\langle seg \ var \rangle$ expands to the number of items, while \seq if in:NnTF $\langle seg \ var \rangle$ { $\langle item \rangle$ } tests if the $\langle item \rangle$ is in the set.

Adding an $\langle item \rangle$ to a set $\langle seq \ var \rangle$ can be done by appending it to the $\langle seq \ var \rangle$ if it is not already in the $\langle seq var \rangle$:

```
\ensuremath{\mbox{seq\_if\_in:NnF}} \langle \ensuremath{\mbox{seq}} \ensuremath{\mbox{var}} \rangle \ \{\langle \ensuremath{\mbox{item}} \rangle\}
{ \seq_put_right: Nn \langle seq var \rangle \langle \item\rangle} }
```

Removing an $\langle item \rangle$ from a set $\langle seq var \rangle$ can be done using \seq remove all:Nn,

```
\sq_remove_all:Nn \langle seq var \rangle \{\langle item \rangle\}
```

The intersection of two sets $\langle seq \ var_1 \rangle$ and $\langle seq \ var_2 \rangle$ can be stored into $\langle seq \ var_3 \rangle$ by collecting items of $\langle seq \ var_1 \rangle$ which are in $\langle seq \ var_2 \rangle$.

The union of two sets $\langle seq \ var_1 \rangle$ and $\langle seq \ var_2 \rangle$ can be stored into $\langle seq \ var_3 \rangle$ through

```
\label{eq:seq_concat:NNN} $$ \langle seq\ var_3 \rangle \ \langle seq\ var_1 \rangle \ \langle seq\ var_2 \rangle $$ \\ $ \langle seq\_remove\_duplicates: N \ \langle seq\ var_3 \rangle $$ $$
```

or by adding items to (a copy of) $\langle seq \ var_1 \rangle$ one by one

```
\label{eq:nn} $\langle seq\ var_3\rangle\ \langle seq\ var_1\rangle$ $$ \eq_map_inline: Nn $\langle seq\ var_2\rangle$ $$ $\{ \seq_if_in: NnF $\langle seq\ var_3\rangle$ $$ $\{\#1\}$ $$ $\{ \seq_put_right: Nn $\langle seq\ var_3\rangle$ $$ $\{\#1\}$ $$ $$ $$
```

The second approach is faster than the first when the $\langle seq \ var_2 \rangle$ is short compared to $\langle seq \ var_1 \rangle$.

The difference of two sets $\langle seq \ var_1 \rangle$ and $\langle seq \ var_2 \rangle$ can be stored into $\langle seq \ var_3 \rangle$ by removing items of the $\langle seq \ var_2 \rangle$ from (a copy of) the $\langle seq \ var_1 \rangle$ one by one.

```
\ensuremath{\verb|seq_set_eq:NN||} \langle seq \ensuremath{ var_3} \rangle \langle seq \ensuremath{ var_1} \rangle \\ \ensuremath{|seq_remove_all:Nn||} \langle seq \ensuremath{ var_2} \rangle \\ \ensuremath{|seq_remove_all:Nn||} \langle seq \ensuremath{ var_3} \rangle \\ \ensuremath{|seq_remove_all:Nn||} \langle
```

The symmetric difference of two sets $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ can be stored into $\langle seq\ var_3 \rangle$ by computing the difference between $\langle seq\ var_1 \rangle$ and $\langle seq\ var_2 \rangle$ and storing the result as $\l_-\langle pkg \rangle$ _internal_seq, then the difference between $\langle seq\ var_2 \rangle$ and $\langle seq\ var_1 \rangle$, and finally concatenating the two differences to get the symmetric differences.

```
\eq_set_eq:NN \l__\langle pkg\rangle_internal\_seq \enskip seq\_var_1\rangle $$ \eq_map_inline:Nn \enskip seq\_var_2\rangle $$ {\enskip seq\_remove\_all:Nn \l__\langle pkg\rangle_internal\_seq $\{\#1\}$ }$$ \eq_set_eq:NN \enskip seq\_var_2\rangle $$ \eq_map_inline:Nn \enskip seq\_var_1\rangle $$ {\enskip seq\_remove\_all:Nn \enskip seq\_var_3\rangle $$ $\{\#1\}$ }$$ \esq\_concat:NNN \enskip seq\_var_3\rangle \enskip seq\_var_3\rangle \l__\langle pkg\rangle_internal\_seq\_var_3\rangle $$
```

19.11 Constant and scratch sequences

\c_empty_seq Constant that is always empty.

New: 2012-07-02

\1_	_tmpa_	seq
	tmpb	

Scratch sequences for local assignment. These are never used by the kernel code, and so are safe for use with any LATEX3-defined function. However, they may be overwritten by $_{\text{New: }2012-04-26}$ other non-kernel code and so should only be used for short-term storage.

\g_tmpa_seq Scratch sequences for global assignment. These are never used by the kernel code, and \g_tmpb_seq so are safe for use with any LATEX3-defined function. However, they may be overwritten $_{\mbox{\scriptsize New: }2012-04-26}$ by other non-kernel code and so should only be used for short-term storage.

Viewing sequences 19.12

\seq_show: N

 $\scalebox{seq_show:N} \langle sequence \rangle$

\seq_show:c

Displays the entries in the $\langle sequence \rangle$ in the terminal.

Updated: 2021-04-29

\seq_log:N

\seq_log:N \(\sequence \)

\seq_log:c

Writes the entries in the $\langle sequence \rangle$ in the log file.

New: 2014-08-12 Updated: 2021-04-29

Chapter 20

The l3int package Integers

Calculation and comparison of integer values can be carried out using literal numbers, int registers, constants and integers stored in token list variables. The standard operators +, -, / and * and parentheses can be used within such expressions to carry arithmetic operations. This module carries out these functions on *integer expressions* ("intexpr").

20.1 Integer expressions

Throughout this module, (almost) all n-type argument allow for an \(\intexpr \) argument with the following syntax. The \(\int integer expression \) should consist, after expansion, of +, -, *, /, (,) and of course integer operands. The result is calculated by applying standard mathematical rules with the following peculiarities:

- / denotes division rounded to the closest integer with ties rounded away from zero;
- there is an error and the overall expression evaluates to zero whenever the absolute value of any intermediate result exceeds $2^{31} 1$, except in the case of scaling operations a*b/c, for which a*b may be arbitrarily large (but the operands a, b, c are still constrained to an absolute value at most $2^{31} 1$);
- parentheses may not appear after unary + or -, namely placing +(or -(at the start of an expression or after +, -, *, / or (leads to an error.

Each integer operand can be either an integer variable (with no need for \int_use:N) or an integer denotation. For example both

```
\int_show:n { 5 + 4 * 3 - ( 3 + 4 * 5 ) }
and
\tl_new:N \l_my_tl
\tl_set:Nn \l_my_tl { 5 }
\int_new:N \l_my_int
\int_set:Nn \l_my_int { 4 }
\int_show:n { \l_my_tl + \l_my_int * 3 - ( 3 + 4 * 5 ) }
```

show the same result -6 because 1_my_t1 expands to the integer denotation 5 while the integer variable \l_my_int takes the value 4. As the (integer expression) is fully expanded from left to right during evaluation, fully expandable and restricted-expandable functions can both be used, and \exp_not:n and its variants have no effect while \exp_not:N may incorrectly interrupt the expression.

TeXhackers note: Exactly two expansions are needed to evaluate \int_eval:n. The result is not an $\langle internal\ integer \rangle$, and therefore should be terminated by a space if used in \int_value:w or in a TEX-style integer assignment.

As all TeX integers, integer operands can also be: $\$ value{ $\langle \cancel{ET}_{EX} 2_{\varepsilon} \$ counter \rangle }; dimension or skip variables, converted to integers in sp; the character code of some character given as '\(char\) or '\\(char\); octal numbers given as ' followed by digits from 0 to 7; or hexadecimal numbers given as " followed by digits and upper case letters from A to F.

\int_eval:n * \int_eval:n {\langle integer expression \rangle}

Evaluates the $\langle integer\ expression \rangle$ and leaves the result in the input stream as an integer denotation: for positive results an explicit sequence of decimal digits not starting with 0, for negative results - followed by such a sequence, and 0 for zero.

\int_eval:w * \int_eval:w \(\text{integer expression} \)

New: 2018-03-30 Evaluates the (integer expression) as described for \int_eval:n. The end of the expression is the first token encountered that cannot form part of such an expression. If that token is \scan_stop: it is removed, otherwise not. Spaces do not terminate the expression. However, spaces terminate explict integers, and this may terminate the expression: for instance, $int eval: w 1_{i-1} + 1_{i-1} = 0$ (with explicit space tokens inserted using \sim in a code setting) expands to 29 since the digit 9 is not part of the expression.

New: 2018-11-03 Evaluates the (integer expression) then leaves 1 or 0 or -1 in the input stream according to the sign of the result.

\int_abs:n * \int_abs:n {\(\(\)integer expression \(\)\}

Updated: 2012-09-26 Evaluates the (integer expression) as described for \int_eval:n and leaves the absolute value of the result in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.

Updated: 2012-09-26 Evaluates the two (integer expressions) as described earlier, then divides the first value by the second, and rounds the result to the closest integer. Ties are rounded away from zero. Note that this is identical to using / directly in an $\langle integer\ expression \rangle$. The result is left in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.

$\int_div_truncate:nn * \int_div_truncate:nn {\langle intexpr_1 \rangle} {\langle intexpr_2 \rangle}$

Updated: 2012-02-09 Evaluates the two (integer expressions) as described earlier, then divides the first value by the second, and rounds the result towards zero. Note that division using / rounds to the closest integer instead. The result is left in the input stream as an (integer denotation) after two expansions.

```
\star \left( int_{max:nn} \left( intexpr_1 \right) \right) \left( intexpr_2 \right) \right)
\int_min:nn
                              \star \left( int_min: nn \left( \langle intexpr_1 \rangle \right) \right) \left( \langle intexpr_2 \rangle \right)
```

Updated: 2012-09-26 Evaluates the (integer expressions) as described for \int_eval:n and leaves either the larger or smaller value in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.

 $\star \left(int_{mod:nn} \left(intexpr_1 \right) \right) \left(intexpr_2 \right) \right)$ \int_mod:nn

Updated: 2012-09-26 Evaluates the two (integer expressions) as described earlier, then calculates the integer remainder of dividing the first expression by the second. This is obtained by subtracting $\int \int div_{truncate:nn} \{ (intexpr_1) \} \{ (intexpr_2) \}$ times $\langle intexpr_2 \rangle \}$ from $\langle intexpr_1 \rangle$. Thus, the result has the same sign as $\langle intexpr_1 \rangle$ and its absolute value is strictly less than that of $\langle intexpr_2 \rangle$. The result is left in the input stream as an $\langle integer\ denotation \rangle$ after two expansions.

20.2Creating and initialising integers

\int_new:N \int_new:N \integer \

\int_new:c

Creates a new $\langle integer \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle integer \rangle$ is initially equal to 0.

\int_const:Nn \langle integer \rangle \langle \integer expression \rangle \rangle \int_const:Nn

\int_const:cn

Updated: 2011-10-22

Creates a new constant $\langle integer \rangle$ or raises an error if the name is already taken. The value of the $\langle integer \rangle$ is set globally to the $\langle integer \ expression \rangle$.

\int_zero:N \int_zero:N \integer \int zero:c Sets $\langle integer \rangle$ to 0. \int_gzero:N \int_gzero:c

\int_zero_new:N \langle integer \rangle \int_zero_new:N

\int_zero_new:c \int_gzero_new:N

\int_gzero_new:c

New: 2011-12-13

Ensures that the \(\langle integer\rangle\) exists globally by applying \\int_new:N if necessary, then applies \inf_{g} int_g zero: N to leave the $\langle integer \rangle$ set to zero.

\int_set_eq:NN \int_set_eq:(cN|Nc|cc) \int_gset_eq:NN

Sets the content of $\langle integer_1 \rangle$ equal to that of $\langle integer_2 \rangle$.

\int_gset_eq:(cN|Nc|cc)

```
\int_if_exist_p:N * \int_if_exist_p:N \langle int \rangle
\int_{int_i} exist_p: c * \int_{int_i} exist:NTF \langle int \rangle \{\langle true \ code \rangle\} \{\langle false \ code \rangle\}
```

 $\int int_if_exist:cTF \star$

New: 2012-03-03

an integer variable.

20.3 Setting and incrementing integers

 $\int \int add: Nn \langle integer \rangle \{\langle integer expression \rangle\}$ \int_add:Nn \int_add:cn Adds the result of the $\langle integer\ expression \rangle$ to the current content of the $\langle integer \rangle$. \int_gadd:Nn \int_gadd:cn Updated: 2011-10-22 \int_decr:N \int_decr:N \integer\ \int_decr:c Decreases the value stored in $\langle integer \rangle$ by 1. \int_gdecr:N \int_gdecr:c \int_incr:N \int_incr:N \integer\ \int_incr:c Increases the value stored in $\langle integer \rangle$ by 1. \int_gincr:N \int_gincr:c \int_set:Nn $\displaystyle \operatorname{int_set:Nn} \ \langle \operatorname{integer} \rangle \ \{\langle \operatorname{integer} \ \operatorname{expression} \rangle \}$ \int_set:cn Sets $\langle integer \rangle$ to the value of $\langle integer\ expression \rangle$, which must evaluate to an integer (as \int_gset:Nn described for \int_eval:n). \int_gset:cn Updated: 2011-10-22 \int_sub:Nn $\verb|\int_sub:Nn| \langle integer \rangle \ \{ \langle integer \ expression \rangle \}$ \int_sub:cn Subtracts the result of the $\langle integer\ expression \rangle$ from the current content of the $\langle integer \rangle$. \int_gsub:Nn \int_gsub:cn Updated: 2011-10-22

20.4 Using integers

 T_EX hackers note: $\\int_use:N$ is the T_EX primitive the: this is one of several E^TEX3 names for this primitive.

20.5 Integer expression conditionals

This function first evaluates each of the $\langle integer\ expressions \rangle$ as described for \int_- eval:n. The two results are then compared using the $\langle relation \rangle$:

Equal = Greater than > Less than <

This function is less flexible than \int_compare:nTF but around 5 times faster.

This function evaluates the $\langle integer\ expressions \rangle$ as described for $\int_{eval:n}\ and\ compares\ consecutive\ result\ using the\ corresponding\ \langle relation \rangle$, namely it compares $\langle intexpr_1 \rangle$ and $\langle intexpr_2 \rangle$ using the $\langle relation_1 \rangle$, then $\langle intexpr_2 \rangle$ and $\langle intexpr_3 \rangle$ using the $\langle relation_2 \rangle$, until finally comparing $\langle intexpr_N \rangle$ and $\langle intexpr_{N+1} \rangle$ using the $\langle relation_N \rangle$. The test yields true if all comparisons are true. Each $\langle integer\ expression \rangle$ is evaluated only once, and the evaluation is lazy, in the sense that if one comparison is false, then no other $\langle integer\ expression \rangle$ is evaluated and no other comparison is performed. The $\langle relations \rangle$ can be any of the following:

```
Equal = or ==
Greater than or equal to >=
Greater than >=
Creater than Creater than >=
Creater than Creater than >=
Creater than C
```

This function is more flexible than \int_compare:nNnTF but around 5 times slower.

This function evaluates the $\langle test\ integer\ expression \rangle$ and compares this in turn to each of the $\langle integer\ expression\ cases \rangle$. If the two are equal then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The function $\int_case:nn$, which does nothing if there is no match, is also available. For example

leaves "Medium" in the input stream.

```
\int_if_even_p:n * \int_if_odd_p:n {\( \lambda \) int_if_even:n \( TF \) * \int_if_odd:nTF {\( \lambda \) int_if_odd_p:n * {\( \lambda \) true \( \cdot \) } \\ \( \lambda \) int_if_odd:nTF * \( \lambda \) int_if_odd_p:n {\( \lambda \)
```

This function first evaluates the *(integer expression)* as described for **\int_eval:n**. It then evaluates if this is odd or even, as appropriate.

20.6 Integer expression loops

Places the $\langle code \rangle$ in the input stream for T_EX to process, and then evaluates the relationship between the two $\langle integer\ expressions \rangle$ as described for \int_compare:nNnTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.

Places the $\langle code \rangle$ in the input stream for TEX to process, and then evaluates the relationship between the two $\langle integer\ expressions \rangle$ as described for \int_compare:nNnTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.

int_until_do:nNnn ☆	$\label{lem:local_local_local} $$ \int_{0}^{\infty} {\langle intexpr_1 \rangle} \ \langle relation \rangle \ \{\langle intexpr_2 \rangle\} \ \{\langle code \rangle\} $$$
	Evaluates the relationship between the two $\langle integer\ expressions \rangle$ as described for \intcompare:nNnTF, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by TEX the test is repeated, and a loop occurs until the test is true.
int_while_do:nNnn ☆	$\int_{\infty} \left(\frac{1}{\sqrt{1 + (1 - 1)^2}} \right) \left(\frac{1}{$
	Evaluates the relationship between the two $\langle integer\ expressions \rangle$ as described for \intcompare:nNnTF, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed by $T_E X$ the test is repeated, and a loop occurs until the test is false.
\int_do_until:nn ☆	$\label{limit_do_until:nn} $$ \left(integer\ relation \right) $$ \left(\left(code \right) \right) $$$
Updated: 2013-01-13	Places the $\langle code \rangle$ in the input stream for TeX to process, and then evaluates the $\langle integer\ relation \rangle$ as described for \int_compare:nTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.
\int_do_while:nn ☆	$\label{limit_do_while:nn} $$ \left(integer\ relation \right) $$ \left(\left(code \right) \right) $$$
Updated: 2013-01-13	Places the $\langle code \rangle$ in the input stream for TeX to process, and then evaluates the $\langle integer\ relation \rangle$ as described for \int_compare:nTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.
\int_until_do:nn ☆	$\label{limit_until_do:nn} $$ \left(integer\ relation \right) $$ \left(\left(code \right) \right) $$$
Updated: 2013-01-13	Evaluates the $\langle integer\ relation \rangle$ as described for $\int_compare:nTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by TEX the test is repeated, and a loop occurs until the test is true.
\int_while_do:nn ☆	$\int_{\infty} {\langle integer\ relation \rangle} \ {\langle code \rangle}$
	Evaluates the \(\lambda integer relation\rangle\) as described for \int_compare:nTF, and then places the \(\lambda code\rangle\) in the input stream if the \(\lambda relation\rangle\) is true. After the \(\lambda code\rangle\) has been processed

 $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed

by TEX the test is repeated, and a loop occurs until the test is false.

20.7Integer step functions

```
\int_step_function:nN
\int_step_function:nnN
```

```
☆ \int_step_function:nN {\( final value \) } \( \frac{function}{} \)
                          ☆ \int_step_function:nnN {⟨initial value⟩} {⟨final value⟩} ⟨function⟩
\int_step_function:nnnN ☆ \int_step_function:nnnN {⟨initial value⟩} {⟨step⟩} {⟨final value⟩} ⟨function⟩
```

New: 2012-06-04 This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should Updated: 2018-04-22 be integer expressions. The $\langle function \rangle$ is then placed in front of each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$). The $\langle step \rangle$ must be non-zero. If the $\langle step \rangle$ is positive, the loop stops when the $\langle value \rangle$ becomes larger than the $\langle final\ value \rangle$. If the $\langle step \rangle$ is negative, the loop stops when the $\langle value \rangle$ becomes smaller than the $\langle final\ value \rangle$. The $\langle function \rangle$ should absorb one numerical argument. For example

```
\cs_set:Npn \my_func:n #1 { [I~saw~#1] \quad }
   \int_step_function:nnnN { 1 } { 1 } { 5 } \my_func:n
would print
```

```
[I \text{ saw } 1]
                  [I \text{ saw } 2]
                                    [I saw 3]
                                                       [I \text{ saw } 4]
                                                                         [I saw 5]
```

The functions \int_step_function:nN and \int_step_function:nN both use a fixed $\langle step \rangle$ of 1, and in the case of \int_step_function:nN the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

\int_step_inline:nn \int_step_inline:nnn

```
\int_step_inline:nnn {\langle initial \ value \rangle} {\langle final \ value \rangle} {\langle code \rangle}
```

New: 2012-06-04 This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which Updated: 2018-04-22 should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ $value\rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with #1 replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (#1).

The functions \int_step_inline:nn and \int_step_inline:nnn both use a fixed $\langle step \rangle$ of 1, and in the case of \int_step_inline:nn the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

\int_step_variable:nNn \int_step_variable:nnNn

```
\{\langle code \rangle\}
```

New: 2012-06-04 Updated: 2018-04-22

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\$ $value\rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl \ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl \ var \rangle$.

The functions \int_step_variable:nNn and \int_step_variable:nnNn both use a fixed $\langle step \rangle$ of 1, and in the case of \int_step_variable:nNn the $\langle initial\ value \rangle$ is also fixed as 1. These functions are provided as simple short-cuts for code clarity.

20.8 Formatting integers

Integers can be placed into the output stream with formatting. These conversions apply to any integer expressions.

\int_to_arabic:n * \int_to_arabic:n {\langle integer expression \rangle}

Updated: 2011-10-22 Places the value of the (integer expression) in the input stream as digits, with category code 12 (other).

\int_to_Alph:n *

\int_to_alph:n * \int_to_alph:n {\(\lambda \) integer expression \(\rangle \)}

Evaluates the (integer expression) and converts the result into a series of letters, which Updated: 2011-09-17 are then left in the input stream. The conversion rule uses the 26 letters of the English alphabet, in order, adding letters when necessary to increase the total possible range of representable numbers. Thus

```
\int_to_alph:n { 1 }
```

places a in the input stream,

```
\int to alph:n { 26 }
```

is represented as z and

```
\int_to_alph:n { 27 }
```

is converted to aa. For conversions using other alphabets, use \int_to_symbols:nnn to define an alphabet-specific function. The basic \int_to_alph:n and \int_to_Alph:n functions should not be modified. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).

\int_to_symbols:nnn *

Updated: 2011-09-17

```
\int_to_symbols:nnn
   {\langle integer expression \rangle} \{\langle total symbols \rangle}
    \{\langle value\ to\ symbol\ mapping \rangle\}
```

This is the low-level function for conversion of an $\langle integer\ expression \rangle$ into a symbolic form (often letters). The $\langle total\ symbols \rangle$ available should be given as an integer expression. Values are actually converted to symbols according to the $\langle value \ to \ symbol \ mapping \rangle$. This should be given as \(\lambda total \) symbols\(\rangle\) pairs of entries, a number and the appropriate symbol. Thus the \int_to_alph:n function is defined as

```
\cs_new:Npn \int_to_alph:n #1
    \int_to_symbols:nnn {#1} { 26 }
        { 1 } { a }
        { 2 } { b }
        { 26 } { z }
 }
```

\int_to_bin:n * \int_to_bin:n {\langle integer expression \rangle}

New: 2014-02-11 Calculates the value of the (integer expression) and places the binary representation of the result in the input stream.

 $\verb|\int_to_hex:n * \int_to_hex:n { (integer expression) } |$

 $\underline{\text{Vint_to_Hex:n}}$ Calculates the value of the $\langle integer\ expression \rangle$ and places the hexadecimal (base 16) New: 2014-02-11 representation of the result in the input stream. Letters are used for digits beyond 9: lower case letters for \int_to_hex:n and upper case ones for \int_to_Hex:n. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).

\int_to_oct:n * \int_to_oct:n {\(\lambda\) integer expression\(\rangle\)}

New: 2014-02-11 Calculates the value of the (integer expression) and places the octal (base 8) representation of the result in the input stream. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).

 $\int \int_{\infty}^{\infty} \int_{\infty}^{\infty} \left(\int_{\infty}^{\infty} \int_{\infty}^{\infty}$

\int_to_Base:nn *

Calculates the value of the *(integer expression)* and converts it into the appropriate Updated: 2014-02-11 representation in the $\langle base \rangle$; the later may be given as an integer expression. For bases greater than 10 the higher "digits" are represented by letters from the English alphabet: lower case letters for \int_to_base:n and upper case ones for \int_to_Base:n. The maximum $\langle base \rangle$ value is 36. The resulting tokens are digits with category code 12 (other) and letters with category code 11 (letter).

TeXhackers note: This is a generic version of \int_to_bin:n, etc.

\int_to_roman:n ☆ \int_to_roman:n {⟨integer expression⟩} \int_to_Roman:n ☆

Places the value of the $\langle integer\ expression \rangle$ in the input stream as Roman numerals, Updated: 2011-10-22 either lower case (\int_to_roman:n) or upper case (\int_to_Roman:n). If the value is negative or zero, the output is empty. The Roman numerals are letters with category code 11 (letter). The letters used are mdclxvi, repeated as needed: the notation with bars (such as \bar{v} for 5000) is not used. For instance \int_to_roman:n { 8249 } expands to mmmmmmmccxlix.

20.9 Converting from other formats to integers

Updated: 2014-08-25 Converts the (letters) into the integer (base 10) representation and leaves this in the input stream. The $\langle letters \rangle$ are first converted to a string, with no expansion. Lower and upper case letters from the English alphabet may be used, with "a" equal to 1 through to "z" equal to 26. The function also accepts a leading sign, made of + and -. This is

the inverse function of \int_to_alph:n and \int_to_Alph:n.

\int_from_bin:n * \int_from_bin:n {\langle binary number \rangle}

New: 2014-02-11 Converts the $\langle binary\ number \rangle$ into the integer (base 10) representation and leaves this in Updated: 2014-08-25 the input stream. The $\langle binary\ number \rangle$ is first converted to a string, with no expansion.

The function accepts a leading sign, made of + and -, followed by binary digits. This is

the inverse function of \int_to_bin:n.

 $\displaystyle \inf_{\text{from_hex:n}} \star \displaystyle \inf_{\text{from_hex:n}} \{\langle \text{hexadecimal number} \rangle \}$

New: 2014-02-11 Converts the $\langle hexadecimal\ number \rangle$ into the integer (base 10) representation and leaves Updated: 2014-08-25 this in the input stream. Digits greater than 9 may be represented in the $\langle hexadecimal \rangle$

 $number\rangle$ by upper or lower case letters. The $\langle hexadecimal\ number\rangle$ is first converted to a string, with no expansion. The function also accepts a leading sign, made of + and -.

This is the inverse function of \int_to_hex:n and \int_to_Hex:n.

 $\verb|\int_from_oct:n * \int_from_oct:n {\langle octal number \rangle}|$

New: 2014-02-11 Converts the $\langle octal\ number \rangle$ into the integer (base 10) representation and leaves this in Updated: 2014-08-25 the input stream. The $\langle octal\ number \rangle$ is first converted to a string, with no expansion.

The function accepts a leading sign, made of + and -, followed by octal digits. This is

the inverse function of \int_to_oct:n.

\int_from_roman:n * \int_from_roman:n {\langle roman numeral \rangle}

Updated: 2014-08-25 Converts the $\langle roman \ numeral \rangle$ into the integer (base 10) representation and leaves this in the input stream. The $\langle roman \ numeral \rangle$ is first converted to a string, with no expansion.

The $\langle roman \ numeral \rangle$ may be in upper or lower case; if the numeral contains characters besides mdclxvi or MDCLXVI then the resulting value is -1. This is the inverse function

of \int_to_roman:n and \int_to_Roman:n.

Updated: 2014-08-25 Converts the $\langle number \rangle$ expressed in $\langle base \rangle$ into the appropriate value in base 10. The $\langle number \rangle$ is first converted to a string, with no expansion. The $\langle number \rangle$ should consist

 $\langle number \rangle$ is first converted to a string, with no expansion. The $\langle number \rangle$ should consist of digits and letters (either lower or upper case), plus optionally a leading sign. The maximum $\langle base \rangle$ value is 36. This is the inverse function of $\int \int abs ds$ and $\int as$

to_Base:nn.

20.10 Random integers

 $\int \int rand:nn + \int rand:nn {\langle intexpr_1 \rangle} {\langle intexpr_2 \rangle}$

New: 2016-12-06 Evaluates the two $\langle integer\ expressions \rangle$ and produces a pseudo-random number between

Updated: 2018-04-27 the two (with bounds included). This is not available in older versions of XATEX.

\int_rand:n * \int_rand:n {\langle intexpr \rangle}

New: 2018-05-05 Evaluates the $\langle integer\ expression \rangle$ then produces a pseudo-random number between 1

and the $\langle intexpr \rangle$ (included). This is not available in older versions of X_HT_EX.

20.11 Viewing integers

 $\verb|\int_show:N \ | int_show:N \ | \langle integer \rangle |$

 $\underline{\underline{\quad}}$ Displays the value of the $\langle integer \rangle$ on the terminal.

 $\verb|\int_show:n | \{ (integer expression) \} |$

New: 2011-11-22 Displays the result of evaluating the (integer expression) on the terminal.

Updated: 2015-08-07

 $\label{log:N} $$ \inf_{\log : \mathbb{N}} \ \langle int_{ger} \rangle $$$

New: 2014-08-22 Updated: 2015-08-03

\int_log:n \int_log:n {\(\lambda\) integer expression\(\rangle\)}

New: 2014-08-22 Writes the result of evaluating the (integer expression) in the log file.

Updated: 2015-08-07

20.12 Constant integers

\c_zero_int Integer values used with primitive tests and assignments: their self-terminating nature \c_one_int makes these more convenient and faster than literal numbers.

New: 2018-05-07

\c_max_int The maximum value that can be stored as an integer.

\c_max_register_int Maximum number of registers.

\c_max_char_int Maximum character code completely supported by the engine.

20.13 Scratch integers

\g_tmpa_int Scratch integer for global assignment. These are never used by the kernel code, and so are safe for use with any IATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

20.14 Direct number expansion

```
\int_value:w * \int_value:w \langle integer \\
New: 2018-03-27 \int_value:w \langle integer denotation \rangle \langle optional space \rangle
```

Expands the following tokens until an $\langle integer \rangle$ is formed, and leaves a normalized form (no leading sign except for negative numbers, no leading digit 0 except for zero) in the input stream as category code 12 (other) characters. The $\langle integer \rangle$ can consist of any number of signs (with intervening spaces) followed by

- an integer variable (in fact, any TFX register except \toks) or
- explicit digits (or by '\(\dot\) octal digits\(\right\) or "\(\lambda\) hexadecimal digits\(\right\) or '\(\lambda\) character\(\right\).

In this last case expansion stops once a non-digit is found; if that is a space it is removed as in f-expansion, and so \exp_stop_f: may be employed as an end marker. Note that protected functions *are* expanded by this process.

This function requires exactly one expansion to produce a value, and so is suitable for use in cases where a number is required "directly". In general, \int_eval:n is the preferred approach to generating numbers.

TeXhackers note: This is the TeX primitive \number.

20.15 Primitive conditionals

Compare two integers using $\langle relation \rangle$, which must be one of =, < or > with category code 12. The **\else**: branch is optional.

TEXhackers note: These are both names for the TEX primitive \ifnum.

Selects a case to execute based on the value of the $\langle integer \rangle$. The first case $(\langle case_0 \rangle)$ is executed if $\langle integer \rangle$ is 0, the second $(\langle case_1 \rangle)$ if the $\langle integer \rangle$ is 1, etc. The $\langle integer \rangle$ may be a literal, a constant or an integer expression (e.g. using \int_eval:n).

TEXhackers note: These are the TEX primitives \ifcase and \or.

Expands $\langle tokens \rangle$ until a non-numeric token or a space is found, and tests whether the resulting $\langle integer \rangle$ is odd. If so, $\langle true\ code \rangle$ is executed. The **\else**: branch is optional.

TEXhackers note: This is the TEX primitive \ifodd.

Chapter 21

The **I3flag** package: Expandable flags

Flags are the only data-type that can be modified in expansion-only contexts. This module is meant mostly for kernel use: in almost all cases, booleans or integers should be preferred to flags because they are very significantly faster.

A flag can hold any non-negative value, which we call its $\langle height \rangle$. In expansion-only contexts, a flag can only be "raised": this increases the $\langle height \rangle$ by 1. The $\langle height \rangle$ can also be queried expandably. However, decreasing it, or setting it to zero requires non-expandable assignments.

Flag variables are always local. They are referenced by a $\langle flag\ name \rangle$ such as str_missing. The $\langle flag\ name \rangle$ is used as part of \use:c constructions hence is expanded at point of use. It must expand to character tokens only, with no spaces.

A typical use case of flags would be to keep track of whether an exceptional condition has occurred during expandable processing, and produce a meaningful (non-expandable) message after the end of the expandable processing. This is exemplified by l3str-convert, which for performance reasons performs conversions of individual characters expandably and for readability reasons produces a single error message describing incorrect inputs that were encountered.

Flags should not be used without carefully considering the fact that raising a flag takes a time and memory proportional to its height. Flags should not be used unless unavoidable.

21.1 Setting up flags

 $flag_new:n \{flag_new:n \{\langle flag_name \rangle\}\}$

Creates a new flag with a name given by $\langle flag\ name \rangle$, or raises an error if the name is already taken. The $\langle flag\ name \rangle$ may not contain spaces. The declaration is global, but flags are always local variables. The $\langle flag \rangle$ initially has zero height.

 $\frac{\flag_clear:n \flag_clear:n \{\langle flag name \rangle\}}{}$

The $\langle flag \rangle$'s height is set to zero. The assignment is local.

\flag_clear_new:n \flag_clear_new:n \{\(flag name \)\}

Ensures that the \(flag \) exists globally by applying \flag_new:n if necessary, then applies \flag_clear:n, setting the height to zero locally.

\flag_show:n \{\(flag name \)\}

Displays the \(flag \) in the terminal.

\flag_log:n \flag_log:n \{\(flag name \)\}

Writes the \(flag \) is height to the log file.

21.2 Expandable flag commands

\flag_if_exist_p:n * \flag_if_exist:n \{\langle flag name \rangle}\}
This function returns true if the \langle flag name \rangle references a flag that has been defined previously, and false otherwise.

\flag_if_raised_p:n * \flag_if_raised:n \{\langle flag name \rangle}\}
This function returns true if the \langle flag \rangle has non-zero height, and false if the \langle flag \rangle has zero height.

\flag_height:n * \flag_height:n \{\langle flag name \rangle}\}
Expands to the height of the \langle flag \rangle as an integer denotation.

\flag_raise:n * \flag_raise:n \{\langle flag name \rangle}\}
The \langle flag raise:n \{\langle flag name \rangle}\}
The \langle flag raise:n \{\langle flag name \rangle}\}
The \langle flag raise:n \tag{flag name}\}

Chapter 22

The **I3clist** package Comma separated lists

Comma lists (in short, clist) contain ordered data where items can be added to the left or right end of the list. This data type allows basic list manipulations such as adding/removing items, applying a function to every item, removing duplicate items, extracting a given item, using the comma list with specified separators, and so on. Sequences (defined in l3seq) are safer, faster, and provide more features, so they should often be preferred to comma lists. Comma lists are mostly useful when interfacing with LATEX $2_{\mathcal{E}}$ or other code that expects or provides items separated by commas.

Several items can be added at once. To ease input of comma lists from data provided by a user outside an \ExplSyntaxOn ... \ExplSyntaxOff block, spaces are removed from both sides of each comma-delimited argument upon input. Blank arguments are ignored, to allow for trailing commas or repeated commas (which may otherwise arise when concatenating comma lists "by hand"). In addition, a set of braces is removed if the result of space-trimming is braced: this allows the storage of any item in a comma list. For instance,

```
\clist_new:N \l_my_clist
\clist_put_left:Nn \l_my_clist { ~a~ , ~{b}~ , c~\d }
\clist_put_right:Nn \l_my_clist { ~{e~} , , {{f}} , }
```

results in \l_my_clist containing a,b,c~\d,{e~},{{f}} namely the five items a, b, c~\d, e~ and {f}. Comma lists normally do not contain empty or blank items so the following gives an empty comma list:

```
\clist_clear_new:N \l_my_clist
\clist_set:Nn \l_my_clist { , ~ , , }
\clist_if_empty:NTF \l_my_clist { true } { false }
```

and it leaves **true** in the input stream. To include an "unsafe" item (empty, or one that contains a comma, or starts or ends with a space, or is a single brace group), surround it with braces.

Any n-type token list is a valid comma list input for l3clist functions, which will split the token list at every comma and process the items as described above. On the other hand, N-type functions expect comma list variables, which are particular token list variables in which this processing of items (and removal of blank items) has already

occurred. Because comma list variables are token list variables, expanding them once yields their items separated by commas, and l3tl functions such as \t1_show:N can be applied to them. (These functions often have l3clist analogues, which should be preferred.)

Almost all operations on comma lists are noticeably slower than those on sequences so converting the data to sequences using \seq_set_from_clist:Nn (see |3seq) may be advisable if speed is important. The exception is that \clist_if_in:NnTF and \clist_-remove_duplicates:N may be faster than their sequence analogues for large lists. However, these functions work slowly for "unsafe" items that must be braced, and may produce errors when their argument contains {, } or # (assuming the usual TEX category codes apply). The sequence data type should thus certainly be preferred to comma lists to store such items.

22.1 Creating and initialising comma lists

\clist_new:N
\clist_new:c

 $\clist_new:N \clist_new:N \cl$

Creates a new $\langle comma \ list \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle comma \ list \rangle$ initially contains no items.

\clist_const:Nn \clist_const:(Nx|cn|cx) New: 2014-07-05 $\verb|\clist_const:Nn| \langle clist| var \rangle | \{\langle comma| list \rangle\}|$

Creates a new constant $\langle clist \ var \rangle$ or raises an error if the name is already taken. The value of the $\langle clist \ var \rangle$ is set globally to the $\langle comma \ list \rangle$.

\clist_clear:N
\clist_clear:c
\clist_gclear:N
\clist_gclear:c

\clist_clear:N \(comma list \)

Clears all items from the $\langle comma \ list \rangle$.

\clist_clear_new:N
\clist_clear_new:c
\clist_gclear_new:N
\clist_gclear_new:c

\clist_clear_new:N \(comma list \)

Ensures that the \(\langle comma \ list \rangle \) exists globally by applying \clist_new:N if necessary, then applies \clist_(g) clear:N to leave the list empty.

\clist_set_eq:NN
\clist_set_eq:(cN|Nc|cc)
\clist_gset_eq:NN
\clist_gset_eq:(cN|Nc|cc)

 $\clist_set_eq:NN \ \langle comma \ list_1 \rangle \ \langle comma \ list_2 \rangle$

Sets the content of $\langle comma \; list_1 \rangle$ equal to that of $\langle comma \; list_2 \rangle$. To set a token list variable equal to a comma list variable, use $\t l_set_eq:NN$. Conversely, setting a comma list variable to a token list is unadvisable unless one checks space-trimming and related issues.

```
\clist_set_from_seq:NN \clist_set_from_seq:NN \comma list\) \sequence \clist_set_from_seq:(cN|Nc|cc) \clist_gset_from_seq:(cN|Nc|cc) \clist_gset_from_seq:(cN|Nc|cc) \quad New: 2014-07-17
```

Converts the data in the $\langle sequence \rangle$ into a $\langle comma\ list \rangle$: the original $\langle sequence \rangle$ is unchanged. Items which contain either spaces or commas are surrounded by braces.

```
\clist_concat:NNN
\clist_concat:ccc
```

 $\clist_{concat:NNN} \langle comma \ list_1 \rangle \langle comma \ list_2 \rangle \langle comma \ list_3 \rangle$

Concatenates the content of $\langle comma \; list_2 \rangle$ and $\langle comma \; list_3 \rangle$ together and saves the \clist_gconcat: ccc result in $\langle comma \ list_1 \rangle$. The items in $\langle comma \ list_2 \rangle$ are placed at the left side of the new comma list.

```
\clist_if_exist:NTF \star
\clist_if_exist:cTF
           New: 2012-03-03
```

\clist_if_exist_p:N * \clist_if_exist_p:N \(comma list \) $\clist_if_exist_p:c \star \clist_if_exist:NTF \comma list) {\langle true code \rangle} {\langle false code \rangle}$

> Tests whether the $\langle comma \ list \rangle$ is currently defined. This does not check that the $\langle comma \ list \rangle$ list really is a comma list.

22.2Adding data to comma lists

```
\clist_set:Nn
                                               \clist_set:Nn \langle comma \ list \rangle \ \{\langle item_1 \rangle, \ldots, \langle item_n \rangle\}
\clist_set:(NV|No|Nx|cn|cV|co|cx)
\clist_gset:Nn
\clist_gset:(NV|No|Nx|cn|cV|co|cx)
                            New: 2011-09-06
```

Sets $\langle comma \ list \rangle$ to contain the $\langle items \rangle$, removing any previous content from the variable. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To store some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: $\clist_set:Nn \langle comma \ list \rangle \ \{ \ \{ \langle tokens \rangle \} \ \}.$

```
\clist_put_left:Nn \langle comma list \rangle \{\langle item_1 \rangle, \ldots, \langle item_n \rangle\}
\clist_put_left:Nn
\clist_put_left:(NV|No|Nx|cn|cV|co|cx)
\clist_gput_left:Nn
\clist_gput_left:(NV|No|Nx|cn|cV|co|cx)
                             Updated: 2011-09-05
```

Appends the $\langle items \rangle$ to the left of the $\langle comma\ list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the \(\langle tokens\rangle\) contain commas or spaces, add a set of braces: \clist put left:\(\mathbb{N}\mathbb{n}\) \(\langle\) comma list { { $\langle tokens \rangle$ } }.

```
\clist_put_right: Nn \langle comma \ list \rangle \ \{\langle item_1 \rangle, \ldots, \langle item_n \rangle\}
\clist_put_right:Nn
\clist_put_right:(NV|No|Nx|cn|cV|co|cx)
\clist_gput_right:Nn
\clist_gput_right: (NV|No|Nx|cn|cV|co|cx)
                               Updated: 2011-09-05
```

Appends the $\langle items \rangle$ to the right of the $\langle comma\ list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the \(\langle tokens \rangle \) contain commas or spaces, add a set of braces: \clist_put_right: \(\mathbb{N}\mathbb{n} \) \(\chi \) comma list { { $\langle tokens \rangle$ } }.

22.3 Modifying comma lists

While comma lists are normally used as ordered lists, it may be necessary to modify the content. The functions here may be used to update comma lists, while retaining the order of the unaffected entries.

```
\clist_remove_duplicates:N
                                  \verb|\clist_remove_duplicates:N| & \langle \textit{comma list} \rangle \\
\clist_remove_duplicates:c
\clist_gremove_duplicates:N
\clist_gremove_duplicates:c
```

Removes duplicate items from the $\langle comma | list \rangle$, leaving the left most copy of each item in the $\langle comma | list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for $tl_$ if_eq:nnTF.

T_EXhackers note: This function iterates through every item in the $\langle comma\ list \rangle$ and does a comparison with the (items) already checked. It is therefore relatively slow with large comma lists. Furthermore, it may fail if any of the items in the (comma list) contains {, }, or # (assuming the usual TEX category codes apply).

```
\clist_remove_all:Nn \langle comma list \rangle \{\langle item \rangle\}
\clist_remove_all:Nn
\clist_remove_all:(cn|NV|cV)
\clist_gremove_all:Nn
\clist_gremove_all:(cn|NV|cV)
                Updated: 2011-09-06
```

Removes every occurrence of $\langle item \rangle$ from the $\langle comma \ list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for \tl_if_eq:nnTF.

TEXhackers note: The function may fail if the $\langle item \rangle$ contains $\{,\}$, or # (assuming the usual TeX category codes apply).

\clist_reverse:N \clist_reverse:c \clist_greverse:N \clist_greverse:c

\clist_reverse:N \(comma list \)

Reverses the order of items stored in the $\langle comma \ list \rangle$.

New: 2014-07-18

 $\verb|\clist_reverse:n \clist_reverse:n \{\langle comma \ list \rangle\}|$

New: 2014-07-18 Leaves the items in the $\langle comma | list \rangle$ in the input stream in reverse order. Contrarily to other what is done for other n-type (comma list) arguments, braces and spaces are preserved by this process.

> TEXhackers note: The result is returned within \unexpanded, which means that the comma list does not expand further when appearing in an x-type or e-type argument expansion.

```
\clist_sort:Nn
\clist sort:cn
\clist_gsort:Nn
\clist_gsort:cn
```

New: 2017-02-06

\clist_sort:Nn \(clist var \) \{\(comparison code \) \}

Sorts the items in the $\langle clist \ var \rangle$ according to the $\langle comparison \ code \rangle$, and assigns the result to $\langle clist \ var \rangle$. The details of sorting comparison are described in Section 6.1.

22.4 Comma list conditionals

```
\clist_if_empty_p:N \star \clist_if_empty_p:N \c
 \verb|\clist_if_empty_p:c * \clist_if_empty:NTF | \langle comma | list \rangle | \{\langle true | code \rangle\} | \{\langle false | code \rangle\}| 
\verb|\clist_if_empty:NTF|| \star \text{ Tests if the } \langle comma \ list \rangle \text{ is empty (containing no items)}.
 \clist_if_empty:cTF *
```

```
\clist_if_empty_p:n \star \clist_if_empty_p:n \{\langle comma \ list \rangle\}
\clist_if_empty:nTF \ \clist_if_empty:nTF
```

New: 2014-07-05 Tests if the $\langle comma \ list \rangle$ is empty (containing no items). The rules for space trimming are as for other n-type comma-list functions, hence the comma list {~,~,~,,~} (without outer braces) is empty, while {~,{}}, (without outer braces) contains one element, which happens to be empty: the comma-list is not empty.

```
\clist_if_in:Nn<u>TF</u>
                                               \clist_if_in:NnTF \ \langle comma \ list \rangle \ \{\langle item \rangle\} \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}
\clist_{if_in:(NV|No|cn|cV|co)}
\clist_if_in:nnTF
\clist_if_in:(nV|no)<u>TF</u>
                      Updated: 2011-09-06
```

Tests if the $\langle item \rangle$ is present in the $\langle comma\ list \rangle$. In the case of an n-type $\langle comma\ list \rangle$, the usual rules of space trimming and brace stripping apply. Hence,

```
\clist_if_in:nnTF { a , {b}~ , {b} , c } { b } {true} {false}
yields true.
```

T_EXhackers note: The function may fail if the $\langle item \rangle$ contains $\{,\}$, or # (assuming the usual TeX category codes apply).

22.5 Mapping over comma lists

The functions described in this section apply a specified function to each item of a comma list. All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

When the comma list is given explicitly, as an n-type argument, spaces are trimmed around each item. If the result of trimming spaces is empty, the item is ignored. Otherwise, if the item is surrounded by braces, one set is removed, and the result is passed to the mapped function. Thus, if the comma list that is being mapped is $\{a_{\sqcup,\sqcup}\{\{b\}_{\sqcup}\},_{\sqcup},\{\},_{\sqcup}\{c\},\}$ then the arguments passed to the mapped function are 'a', '{b} $_{\perp}$ ', an empty argument, and 'c'.

When the comma list is given as an N-type argument, spaces have already been trimmed on input, and items are simply stripped of one set of braces if any. This case is more efficient than using n-type comma lists.

```
\clist_map_function:NN ☆ \clist_map_function:NN ⟨comma list⟩ ⟨function⟩
\clist_map_function:cN ☆
                               Applies \langle function \rangle to every \langle item \rangle stored in the \langle comma \ list \rangle. The \langle function \rangle receives one
\clist_map_function:nN ☆
                              argument for each iteration. The (items) are returned from left to right. The function
            Updated: 2012-06-29 \clist_map_inline: Nn is in general more efficient than \clist_map_function: NN.
      \clist_map_inline: Nn \clist_map_inline: Nn \comma list \ {\langle inline function \}
      \clist map inline:cn
                               Applies \langle inline\ function \rangle to every \langle item \rangle stored within the \langle comma\ list \rangle. The \langle inline\ function \rangle
      \clist_map_inline:nn
                               function should consist of code which receives the \langle item \rangle as #1. The \langle items \rangle are returned
            Updated: 2012-06-29 from left to right.
  \clist_map\_variable:NNn \clist_map\_variable:NNn \comma list \ \variable \ \{\code\}\}
  \clist_map_variable:cNn
                               Stores each \langle item \rangle of the \langle comma\ list \rangle in turn in the (token list) \langle variable \rangle and applies
  \clist_map_variable:nNn
                               the \langle code \rangle. The \langle code \rangle will usually make use of the \langle variable \rangle, but this is not enforced.
            Updated: 2012-06-29 The assignments to the \langle variable \rangle are local. Its value after the loop is the last \langle item \rangle in
                               the \langle comma \ list \rangle, or its original value if there were no \langle item \rangle. The \langle items \rangle are returned
                               from left to right.
  \clist_map_tokens:Nn ☆ \clist_map_tokens:Nn ⟨clist var⟩ {⟨code⟩}
  \clist_map\_tokens:cn \cdot \clist_map\_tokens:nn {(comma list)} {(code)}
  New: 2021-05-05 each \langle item \rangle as a trailing brace group. If the \langle code \rangle consists of a single function this is
                               equivalent to \clist map function:nN.
      \clist_map_break: ☆ \clist_map_break:
```

Updated: 2012-06-29 Used to terminate a \clist_map_... function before all entries in the \(comma \ list \) have been processed. This normally takes place within a conditional statement, for example

```
\clist_map_inline:Nn \l_my_clist
 {
    \str_if_eq:nnTF { #1 } { bingo }
      { \clist_map_break: }
        % Do something useful
      }
 }
```

Use outside of a \clist_map_... scenario leads to low level TeX errors.

TEXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

```
\clist_map_break:n \( \text{\clist_map_break:n } \( \clist_map_break:n \)
```

Updated: 2012-06-29 Used to terminate a \clist_map_... function before all entries in the \(comma \ list \) have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\clist_map_inline:Nn \l_my_clist
    \str if eq:nnTF { #1 } { bingo }
      { \clist_map_break:n { <code> } }
        % Do something useful
 }
```

Use outside of a \clist_map_... scenario leads to low level TeX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

```
\clist_count:c *
\clist_count:n *
```

\clist_count:N * \clist_count:N \(\comma \) list \

Leaves the number of items in the $\langle comma \ list \rangle$ in the input stream as an $\langle integer \rangle$ denotation). The total number of items in a $\langle comma | list \rangle$ includes those which are New: 2012-07-13 duplicates, i.e. every item in a $\langle comma~list \rangle$ is counted.

22.6 Using the content of comma lists directly

```
\clist_use:Nnnn * \clist_use:Nnnn \langle clist var \rangle \langle \separator between two \rangle \.
\clist use:cnnn \star {\langle separator between more than two\} {\langle separator between final two\}
```

New: 2013-05-26 Places the contents of the $\langle clist \ var \rangle$ in the input stream, with the appropriate $\langle separator \rangle$ between the items. Namely, if the comma list has more than two items, the \(\separator \) between more than two is placed between each pair of items except the last, for which the (separator between final two) is used. If the comma list has exactly two items, then they are placed in the input stream separated by the $\langle separator\ between\ two\rangle$. If the comma list has a single item, it is placed in the input stream, and a comma list with no items produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\clist_set:Nn \l_tmpa_clist { a , b , , c , {de} , f }
\clist_use:Nnnn \l_tmpa_clist { ~and~ } { ,~ } { ,~and~ }
```

inserts "a, b, c, de, and f" in the input stream. The first separator argument is not used in this case because the comma list has more than 2 items.

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the \(\langle items \rangle \) do not expand further when appearing in an x-type or e-type argument expansion.

```
\clist_use:Nn \star \clist_use:Nn \clist \var \clist \var \clist \clist \var \clist \cli
```

New: 2013-05-26 the items. If the comma list has a single item, it is placed in the input stream, and a comma list with no items produces no output. An error is raised if the variable does not exist or if it is invalid.

For example,

```
\clist_set:Nn \l_tmpa_clist { a , b , , c , {de} , f }
\clist use:Nn \l tmpa clist { ~and~ }
```

inserts "a and b and c and de and f" in the input stream.

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the (items) do not expand further when appearing in an x-type or e-type argument expansion.

```
\clist_use:nn
```

```
\clist_use:nnnn * \clist_use:nnnn \( comma list \) \{\( (separator between two \) \}
                    * {\separator between more than two\} {\separator between final two\}
                      \clist_use:nn \langle comma \ list \rangle \{\langle separator \rangle\}
```

New: 2021-05-10

Places the contents of the \(\comma \) list\\ in the input stream, with the appropriate (separator) between the items. As for \clist_set:Nn, blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. The (separators) are then inserted in the same way as for \clist_use:Nnnn and \clist_use:Nn, respectively.

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the (items) do not expand further when appearing in an x-type or e-type argument expansion.

22.7 Comma lists as stacks

Comma lists can be used as stacks, where data is pushed to and popped from the top of the comma list. (The left of a comma list is the top, for performance reasons.) The stack functions for comma lists are not intended to be mixed with the general ordered data functions detailed in the previous section: a comma list should either be used as an ordered data type or as a stack, but not in both ways.

```
\clist_get:NN
\clist_get:cN
\clist_get:NN<u>TF</u>
\clist_get:cNTF
```

```
\clist_get:NN \( comma list \) \( \taken list variable \)
```

Stores the left-most item from a $\langle comma \ list \rangle$ in the $\langle token \ list \ variable \rangle$ without removing it from the $\langle comma \; list \rangle$. The $\langle token \; list \; variable \rangle$ is assigned locally. In the nonbranching version, if the $\langle comma \; list \rangle$ is empty the $\langle token \; list \; variable \rangle$ is set to the marker value \q_no_value.

New: 2012-05-14 Updated: 2019-02-16

```
\clist_pop:NN
\clist_pop:cN
```

```
\clist_pop:NN \( comma list \) \( \text{token list variable} \)
```

Pops the left-most item from a $\langle comma\ list \rangle$ into the $\langle token\ list\ variable \rangle$, i.e. removes the Updated: 2011-09-06 item from the comma list and stores it in the $\langle token\ list\ variable \rangle$. Both of the variables are assigned locally.

\clist_gpop:NN \clist_gpop:NN \comma list \ \doken list variable \

\clist_gpop:cN

Pops the left-most item from a $\langle comma \ list \rangle$ into the $\langle token \ list \ variable \rangle$, i.e. removes the item from the comma list and stores it in the $\langle token\ list\ variable \rangle$. The $\langle comma\ list \rangle$ is modified globally, while the assignment of the $\langle token\ list\ variable \rangle$ is local.

\clist_pop:NNTF \clist_pop:cNTF

 $\left(\text{clist_pop:NNTF} \left(\text{comma list} \right) \left(\text{token list variable} \right) \left\{ \left(\text{true code} \right) \right\} \left\{ \left(\text{false code} \right) \right\}$

If the $\langle comma \ list \rangle$ is empty, leaves the $\langle false \ code \rangle$ in the input stream. The value of New: 2012-05-14 the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle comma \; list \rangle$ is non-empty, pops the top item from the $\langle comma \; list \rangle$ in the $\langle token \; list \rangle$ variable, i.e. removes the item from the $\langle comma \ list \rangle$. Both the $\langle comma \ list \rangle$ and the $\langle token\ list\ variable \rangle$ are assigned locally.

\clist_gpop:cNTF

 $\clist_gpop:NNTF \clist_gpop:NNTF \comma list \comma list \variable \code \clist \code \clist \code \clist \code \clist \clist$

If the $\langle comma \; list \rangle$ is empty, leaves the $\langle false \; code \rangle$ in the input stream. The value of New: 2012-05-14 the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle comma \; list \rangle$ is non-empty, pops the top item from the $\langle comma \; list \rangle$ in the $\langle token \; list \rangle$ variable, i.e. removes the item from the $\langle comma \ list \rangle$. The $\langle comma \ list \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

\clist_push:Nn

 $\clist_push:Nn \langle comma \ list \rangle \{\langle items \rangle\}$

\clist_push:(NV|No|Nx|cn|cV|co|cx)

\clist_gpush:Nn

\clist_gpush:(NV|No|Nx|cn|cV|co|cx)

Adds the $\{\langle items \rangle\}$ to the top of the $\langle comma \; list \rangle$. Spaces are removed from both sides of each item as for any n-type comma list.

22.8 Using a single item

 $\clist_item:Nn \star \clist_item:Nn \clist\} {\clist_item:Nn \clist}$

\clist_item:cn * \clist_item:nn *

Indexing items in the $\langle comma \ list \rangle$ from 1 at the top (left), this function evaluates the (integer expression) and leaves the appropriate item from the comma list in the input New: 2014-07-17 stream. If the $\langle integer\ expression \rangle$ is negative, indexing occurs from the bottom (right) of the comma list. When the $\langle integer\ expression \rangle$ is larger than the number of items in the (comma list) (as calculated by \clist_count:N) then the function expands to nothing.

> TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

```
\label{list_rand_item:N} $$ \clist_rand_item:N & \clist_rand_item:N &
```

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle item \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

22.9 Viewing comma lists

\clist_show:N \clist_show:c	$\label{list_show:N} $$ \langle comma \; list \rangle$$ Displays the entries in the $$ \langle comma \; list \rangle$ in the terminal.$
\clist_show:n Updated: 2013-08-03	$\clist_{show:n} {\langle tokens \rangle} \label{list_show:n}$ Displays the entries in the comma list in the terminal.
\clist_log:N \clist_log:c New: 2014-08-22 Updated: 2021-04-29	$\label{log:N comma list} $$ \writes the entries in the $$ \langle comma list \rangle$ in the log file. See also $$ \clist_show:N which displays the result in the terminal.$
	$\label{log:n} $$ \clist_log:n {$\langle tokens \rangle$} $$ Writes the entries in the comma list in the log file. See also $$ \clist_show:n which displays the result in the terminal.$

22.10 Constant and scratch comma lists

\c_empty_clist	Constant that is always empty.
New: 2012-07-02	
\l_tmpa_clist	Scratch comma lists for local assignment. These are never used by the kernel code, and
\l_tmpb_clist	so are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.
\a tmpo clist	Contain annual lists for alchel and mount. There are accounted by the learned and and
\g_tmpb_clist	Scratch comma lists for global assignment. These are never used by the kernel code, and so are safe for use with any LATEX3-defined function. However, they may be overwritten
New: 2011-09-06	by other non-kernel code and so should only be used for short-term storage.

Chapter 23

The **I3token** package Token manipulation

This module deals with tokens. Now this is perhaps not the most precise description so let's try with a better description: When programming in TEX, it is often desirable to know just what a certain token is: is it a control sequence or something else. Similarly one often needs to know if a control sequence is expandable or not, a macro or a primitive, how many arguments it takes etc. Another thing of great importance (especially when it comes to document commands) is looking ahead in the token stream to see if a certain character is present and maybe even remove it or disregard other tokens while scanning. This module provides functions for both and as such has two primary function categories: \token_ for anything that deals with tokens and \peek_ for looking ahead in the token stream.

Most functions we describe here can be used on control sequences, as those are tokens as well.

It is important to distinguish two aspects of a token: its "shape" (for lack of a better word), which affects the matching of delimited arguments and the comparison of token lists containing this token, and its "meaning", which affects whether the token expands or what operation it performs. One can have tokens of different shapes with the same meaning, but not the converse.

For instance, $\if:w$, $\if:charcode:w$, and $\tex_if:D$ are three names for the same internal operation of T_EX , namely the primitive testing the next two characters for equality of their character code. They have the same meaning hence behave identically in many situations. However, T_EX distinguishes them when searching for a delimited argument. Namely, the example function $\slashed{show_until_if:w}$ defined below takes everything until $\if:w$ as an argument, despite the presence of other copies of $\slashed{show_until_if:w}$ under different names.

```
\cs_new:Npn \show_until_if:w #1 \if:w { \tl_show:n {#1} }
\show_until_if:w \tex_if:D \if_charcode:w \if:w
```

A list of all possible shapes and a list of all possible meanings are given in section 23.7.

23.1Creating character tokens

\char_set_active_eq:NN \char_set_active_eq:Nc \char_gset_active_eq:NN \char_gset_active_eq:Nc

Sets the behaviour of the $\langle char \rangle$ in situations where it is active (category code 13) to be equivalent to that of the $\langle function \rangle$. The category code of the $\langle char \rangle$ is unchanged by this process. The $\langle function \rangle$ may itself be an active character.

Updated: 2015-11-12

\char_set_active_eq:nN \char_set_active_eq:nc \char_gset_active_eq:nN \char_gset_active_eq:nc

\char_set_active_eq:nN {\langle integer expression \rangle} \langle function \rangle

Sets the behaviour of the $\langle char \rangle$ which has character code as given by the $\langle integer \rangle$ expression) in situations where it is active (category code 13) to be equivalent to that of the $\langle function \rangle$. The category code of the $\langle char \rangle$ is unchanged by this process. The $\langle function \rangle$ may itself be an active character.

New: 2015-11-12

New: 2015-09-09 Generates a character token of the given $\langle charcode \rangle$ and $\langle catcode \rangle$ (both of which may be Updated: 2019-01-16 integer expressions). The $\langle catcode \rangle$ may be one of

- 1 (begin group)
- 2 (end group)
- 3 (math toggle)
- 4 (alignment)
- 6 (parameter)
- 7 (math superscript)
- 8 (math subscript)
- 10 (space)
- 11 (letter)
- 12 (other)
- 13 (active)

and other values raise an error. The $\langle charcode \rangle$ may be any one valid for the engine in use, except that for $\langle catcode \rangle$ 10, $\langle charcode \rangle$ 0 is not allowed. Active characters cannot be generated in older versions of XATFX. Another way to build token lists with unusual category codes is $\regex_replace:nnN {.*} {\langle replacement \rangle} \langle tl \ var \rangle$.

TeXhackers note: Exactly two expansions are needed to produce the character.

\c_catcode_other_space_tl Token list containing one character with category code 12, ("other"), and character code New: 2011-09-05 32 (space).

23.2 Manipulating and interrogating character tokens

```
\char_set_catcode_letter:N \( \character \)
\char_set_catcode_escape:N
\char_set_catcode_group_begin:N
\char_set_catcode_group_end:N
\char_set_catcode_math_toggle:N
\char_set_catcode_alignment:N
\char_set_catcode_end_line:N
\char_set_catcode_parameter:N
\char_set_catcode_math_superscript:N
\char_set_catcode_math_subscript:N
\char_set_catcode_ignore:N
\char_set_catcode_space:N
\char_set_catcode_letter:N
\char_set_catcode_other:N
\char_set_catcode_active:N
\char_set_catcode_comment:N
\char_set_catcode_invalid:N
```

Updated: 2015-11-11

Sets the category code of the $\langle character \rangle$ to that indicated in the function name. Depending on the current category code of the $\langle token \rangle$ the escape token may also be needed:

\char_set_catcode_other:N \%

The assignment is local.

```
\verb|\char_set_catcode_letter:n {| (integer expression)|}|
\char_set_catcode_escape:n
\char_set_catcode_group_begin:n
\char_set_catcode_group_end:n
\char_set_catcode_math_toggle:n
\char_set_catcode_alignment:n
\char_set_catcode_end_line:n
\char_set_catcode_parameter:n
\char_set_catcode_math_superscript:n
\char_set_catcode_math_subscript:n
\char_set_catcode_ignore:n
\char_set_catcode_space:n
\char_set_catcode_letter:n
\char_set_catcode_other:n
\char_set_catcode_active:n
\char_set_catcode_comment:n
\char_set_catcode_invalid:n
```

Updated: 2015-11-11

Sets the category code of the $\langle character \rangle$ which has character code as given by the $\langle integer\ expression \rangle$. This version can be used to set up characters which cannot otherwise be given (cf. the N-type variants). The assignment is local.

 $\c set_catcode:nn \c set_catcode:nn \{\langle intexpr_1 \rangle\} \} \{\langle intexpr_2 \rangle\}$

Updated: 2015-11-11 These functions set the category code of the (character) which has character code as given by the $\langle integer\ expression \rangle$. The first $\langle integer\ expression \rangle$ is the character code and the second is the category code to apply. The setting applies within the current T_FX group. In general, the symbolic functions \c set_catcode \c should be preferred, but there are cases where these lower-level functions may be useful.

Expands to the current category code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$.

\char_show_value_catcode:n \char_show_value_catcode:n {\langle integer expression \rangle}

Displays the current category code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$ on the terminal.

Updated: 2015-08-06 Sets up the behaviour of the (character) when found inside \text_lowercase:n, such that $\langle character_1 \rangle$ will be converted into $\langle character_2 \rangle$. The two $\langle characters \rangle$ may be specified using an (integer expression) for the character code concerned. This may include the TFX '(character) method for converting a single character into its character code:

```
\char_set_lccode:nn { '\A } { '\a } % Standard behaviour
\frac{ (A + 32)}{ }
\char_set_lccode:nn { 50 } { 60 }
```

The setting applies within the current T_FX group.

\char_value_lccode:n * \char_value_lccode:n {\(integer expression \) \}

Expands to the current lower case code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$.

\char_show_value_lccode:n \char_show_value_lccode:n {\langle integer expression \rangle}

Displays the current lower case code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$ on the terminal.

Updated: 2015-08-06 Sets up the behaviour of the (character) when found inside \text_uppercase:n, such that $\langle character_1 \rangle$ will be converted into $\langle character_2 \rangle$. The two $\langle characters \rangle$ may be specified using an (integer expression) for the character code concerned. This may include the TeX ' $\langle character \rangle$ method for converting a single character into its character code:

```
\char_set_uccode:nn { '\a } { '\A } % Standard behaviour
\char_set_uccode:nn { '\A } { '\A - 32 }
\char_set_uccode:nn { 60 } { 50 }
```

The setting applies within the current TFX group.

\char_value_uccode:n * \char_value_uccode:n {\(\langle integer expression \rangle \rangle \)}

Expands to the current upper case code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$.

Displays the current upper case code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$ on the terminal.

Updated: 2015-08-06 This function sets up the math code of $\langle character \rangle$. The $\langle character \rangle$ is specified as an (integer expression) which will be used as the character code of the relevant character. The setting applies within the current T_FX group.

Expands to the current math code of the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$.

 $\verb|\char_show_value_mathcode:n \char_show_value_mathcode:n \cdot \{ \langle integer \ expression \rangle \}|$

Displays the current math code of the (character) with character code given by the $\langle integer\ expression \rangle$ on the terminal.

Updated: 2015-08-06 This function sets up the space factor for the $\langle character \rangle$. The $\langle character \rangle$ is specified as an (integer expression) which will be used as the character code of the relevant character. The setting applies within the current T_FX group.

\char_value_sfcode:n * \char_value_sfcode:n {\langle integer expression \rangle}

Expands to the current space factor for the $\langle character \rangle$ with character code given by the $\langle integer\ expression \rangle$.

Displays the current space factor for the (character) with character code given by the $\langle integer\ expression \rangle$ on the terminal.

\l_char_active_seq Used to track which tokens may require special handling at the document level as they New: 2012-01-23 are (or have been at some point) of category $\langle active \rangle$ (catcode 13). Each entry in the Updated: 2015-11-11 sequence consists of a single escaped token, for example \~. Active tokens should be added to the sequence when they are defined for general document use.

\l_char_special_seq Used to track which tokens will require special handling when working with verbatim-

New: 2012-01-23 like material at the document level as they are not of categories (letter) (catcode 11)

 $U_{polated: 2015-11-11}$ or $\langle other \rangle$ (catcode 12). Each entry in the sequence consists of a single escaped token, for example \\ for the backslash or \{ for an opening brace. Escaped tokens should be added to the sequence when they are defined for general document use.

23.3 Generic tokens

\c_group_begin_token \c_group_end_token \c_math_toggle_token \c_alignment_token \c_parameter_token \c_math_superscript_token \c_math_subscript_token \c_space_token

These are implicit tokens which have the category code described by their name. They are used internally for test purposes but are also available to the programmer for other

\c_catcode_other_token

\c_catcode_letter_token These are implicit tokens which have the category code described by their name. They are used internally for test purposes and should not be used other than for category code

\c_catcode_active_tl A token list containing an active token. This is used internally for test purposes and should not be used other than in appropriately-constructed category code tests.

23.4 Converting tokens

\token_to_meaning:c *

 $\token_{to_meaning:N} \star \token_{to_meaning:N} \langle token \rangle$

Inserts the current meaning of the $\langle token \rangle$ into the input stream as a series of characters of category code 12 (other). This is the primitive T_{FX} description of the $\langle token \rangle$, thus for example both functions defined by \cs_set_nopar:Npn and token list variables defined using \tl_new:N are described as macros.

TEXhackers note: This is the TEX primitive \meaning. The $\langle token \rangle$ can thus be an explicit space tokens or an explicit begin-group or end-group character token ({ or } when normal TFX category codes apply) even though these are not valid N-type arguments.

\token_to_str:N * \token_to_str:N \(\token \)

\token_to_str:c *

Converts the given $\langle token \rangle$ into a series of characters with category code 12 (other). If the $\langle token \rangle$ is a control sequence, this will start with the current escape character with category code 12 (the escape character is part of the $\langle token \rangle$). This function requires only a single expansion.

TEXhackers note: $\texttt{token_to_str:N}$ is the TEX primitive string renamed. The (token)can thus be an explicit space tokens or an explicit begin-group or end-group character token ({ or } when normal TEX category codes apply) even though these are not valid N-type arguments.

23.5 Token conditionals

```
\token_if_group_begin_p:N * \token_if_group_begin_p:N \token\
                         Tests if \langle token \rangle has the category code of a begin group token ({ when normal TEX
                                                                                                                               category codes are in force). Note that an explicit begin group token cannot be tested in
                                                                                                                                this way, as it is not a valid N-type argument.
\token_if_group_end_p:N * \token_if_group_end_p:N \langle token \rangle
\token_if_group_end:NTF \star \token_if_group_end:NTF \token) {\token} {\token} {\token} {\token}
                                                                                                                                Tests if \langle token \rangle has the category code of an end group token () when normal T<sub>E</sub>X category
                                                                                                                               codes are in force). Note that an explicit end group token cannot be tested in this way,
                                                                                                                               as it is not a valid N-type argument.
                         \token_if_math_toggle_p:N * \token_if_math_toggle_p:N \( \token \)
                         \label{locality} $$ \ \left( \norm{if_math_toggle:NTF} \ (\norm{token} \ (\norm{token}) \ (\n
                                                                                                                                Tests if \langle token \rangle has the category code of a math shift token ($ when normal TeX category
                                                                                                                               codes are in force).
\token_if_alignment_p:N * \token_if_alignment_p:N \token\
\token_if_alignment:NTF \star \token_if_alignment:NTF \token) {\token} {\token} {\token}
                                                                                                                               Tests if \langle token \rangle has the category code of an alignment token (& when normal T<sub>E</sub>X category
                                                                                                                               codes are in force).
\verb|\token_if_parameter_p:N * \token_if_parameter_p:N & | token_if_parameter_p:N & | token_if_paramete
\token_if_parameter:NTF * \token_if_parameter:NTF \token \ {\token \ {\times code}} {\times code}}
                                                                                                                               Tests if \langle token \rangle has the category code of a macro parameter token (# when normal T<sub>E</sub>X
                                                                                                                               category codes are in force).
                         \token_if_math_superscript_p:N * \token_if_math_superscript_p:N \(\token\)
                         \t \token_if_math_superscript:NTF \star \token_if_math_superscript:NTF \langle token \rangle {\langle true\ code \rangle} {\langle false\ code \rangle}
                                                                                                                               Tests if \langle token \rangle has the category code of a superscript token (^ when normal TFX category
                                                                                                                               codes are in force).
                         \token_if_math_subscript_p:N * \token_if_math_subscript_p:N \(\lambda\)
                         Tests if \langle token \rangle has the category code of a subscript token (_ when normal T<sub>F</sub>X category
                                                                                                                               codes are in force).
                   \verb|\token_if_space_p:N * \token_if_space_p:N & \\token_if_space_p:N & \\token_if_space_p:
                   \label{token_if_space:NTF} $$ \token_if_space:NTF $$ \token \ {\token} \ {\token} \ {\token} \ $$ \token \ \t
```

Tests if $\langle token \rangle$ has the category code of a space token. Note that an explicit space token with character code 32 cannot be tested in this way, as it is not a valid N-type argument.

```
\token_if_letter_p:N \star \token_if_letter_p:N \token
                   \token_if_letter:NTF 	imes \token_if_letter:NTF \token \{ (true code) \} \{ (false code) \}}
                                                                                                                                      Tests if \langle token \rangle has the category code of a letter token.
                        \token_if_other_p:N * \token_if_other_p:N \(\lambda token\rangle)
                        \time TF \star \time TF \to \time TF 
                                                                                                                                      Tests if \langle token \rangle has the category code of an "other" token.
                   \token_if_active_p:N * \token_if_active_p:N \langle token \rangle
                    \time \tim
                                                                                                                                      Tests if \langle token \rangle has the category code of an active character.
                              \token_if_eq_catcode_p:NN \star \token_if_eq_catcode_p:NN \token_i \token_i \token_i
                              Tests if the two \langle tokens \rangle have the same category code.
                              \verb|\token_if_eq_charcode_p:NN| \star \\ \verb|\token_if_eq_charcode_p:NN| \\ \langle token_1 \rangle \\ \langle token_2 \rangle
                              \verb|\token_if_eq_charcode:NNTF| $$ \token_if_eq_charcode:NNTF| $$ \token_i$ \ $$ \token_i$ \token_i$ \ $$ \toke
                                                                                                                                     Tests if the two \langle tokens \rangle have the same character code.
                              \token_if_eq_meaning_p:NN \star \token_if_eq_meaning_p:NN \token_i \
                              Tests if the two \langle tokens \rangle have the same meaning when expanded.
                        \verb|\token_if_macro_p:N * \token_if_macro_p:N & token|
                        \verb|\token_if_macro:NTF| * \token_if_macro:NTF| \langle token \rangle \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\} 
                                                         Updated: 2011-05-23 Tests if the \langle token \rangle is a TeX macro.
                                      \verb|\token_if_cs_p:N * \token_if_cs_p:N & token_if_cs_p:N & token|
                                      \verb|\token_if_cs:NTF| & \texttt|\token_if_cs:NTF| & \texttt|\token| & \{\true| code\} \} & \{\true| code\} \} \\
                                                                                                                                      Tests if the \langle token \rangle is a control sequence.
\token_if_expandable_p:N \star \token_if_expandable_p:N \token \
\time TF \star \time TF \star \time TF \star \time TF \star \time TF \t
                                                                                                                                      Tests if the \langle token \rangle is expandable. This test returns \langle false \rangle for an undefined token.
\verb|\token_if_long_macro_p:N * \token_if_long_macro_p:N & token|
\token_if_long_macro:NTF \star \token_if_long_macro:NTF \token) {\token} {\token} {\token} {\token} {\token}
                                                         Updated: 2012-01-20 Tests if the \langle token \rangle is a long macro.
                              \verb|\token_if_protected_macro_p:N * \token_if_protected_macro_p:N & token|
                              \time T_{protected_macro:NTF} \star \time T_{protected_macro:NTF} \time T_{token} {\time code} {\time Code} \
                                                                                                                 Updated: 2012-01-20
```

Tests if the $\langle token \rangle$ is a protected macro: for a macro which is both protected and long this returns false.

```
\token_if_protected_long_macro_p:N \token_if_protected_long_macro_p:N \token\\
\token_if_protected_long_macro:NTF \token\\
\token_if_protected_long_macro:
```

TEXhackers note: Booleans, boxes and small integer constants are implemented as \chardefs.

Tests if the $\langle token \rangle$ is defined to be a mathchardef.

Tests if the $\langle token \rangle$ is defined to be a font selection command.

```
\label{token_if_dim_register_p:N } $$ \token_if_dim_register_p:N \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register:NTF \token_if_dim_register_p:N \token_i
```

Tests if the $\langle token \rangle$ is defined to be a dimension register.

Tests if the $\langle token \rangle$ is defined to be a integer register.

TEXhackers note: Constant integers may be implemented as integer registers, **\chardefs**, or **\mathchardefs** depending on their value.

```
\label{token_if_muskip_register_p:N } $$ \token_if_muskip_register_p:N \ \langle token_\rangle $$ \token_if_muskip_register:NTF \ \langle token_\rangle \ \{\langle true\ code_\rangle\} \ \{\langle false\ code_\rangle\} $$ \end{tabular} $$ New: 2012-02-15 $$
```

Tests if the $\langle token \rangle$ is defined to be a muskip register.

```
\token_if_skip_register_p:N * \token_if_skip_register_p:N \langle token \rangle
     \token_if_skip\_register:NTF * \token\_if\_skip\_register:NTF \{token} {\times code}\} {\times code}
                     Updated: 2012-01-20
                            Tests if the \langle token \rangle is defined to be a skip register.
     \token_if_toks_register_p:N * \token_if_toks_register_p:N \(\token\)
     Updated: 2012-01-20
                            Tests if the \langle token \rangle is defined to be a toks register (not used by LATEX3).
\token_if_primitive_p:N * \token_if_primitive_p:N \langle token \rangle
Updated: 2020-09-11 Tests if the \langle token \rangle is an engine primitive. In LuaTFX this includes primitive-like com-
                            mands defined using {token.set_lua}.
                                    \token_case_meaning:NnTF \(\langle test token \rangle \)
     \token case catcode:Nn
     \token_case_catcode:NnTF
                                      {
                                         \langle token \ case_1 \rangle \ \{\langle code \ case_1 \rangle\}
     \token_case_charcode:Nn
                                        \langle token \ case_2 \rangle \ \{\langle code \ case_2 \rangle\}
     \token_case_charcode:NnTF
     \token_case_meaning:Nn
     \token_case_meaning:NnTF
                                         \langle token \ case_n \rangle \ \{\langle code \ case_n \rangle\}
                                      }
```

This function compares the $\langle test\ token \rangle$ in turn with each of the $\langle token\ cases \rangle$. If the two are equal (as described for \token_if_eq_catcode:NNTF, \token_if_eq_- charcode:NNTF and \token_if_eq_meaning:NNTF, respectively) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The functions \token_case_catcode:Nn, \token_case_charcode:Nn, and \token_case_meaning:Nn, which do nothing if there is no match, are also available.

23.6 Peeking ahead at the next token

There is often a need to look ahead at the next token in the input stream while leaving it in place. This is handled using the "peek" functions. The generic \peek_after:Nw is provided along with a family of predefined tests for common cases. As peeking ahead does not skip spaces the predefined tests include both a space-respecting and space-skipping version. In addition, using \peek_analysis_map_inline:n, one can map through the following tokens in the input stream and repeatedly perform some tests.

```
\peek_after:Nw \peek_after:Nw \function \ \token \
```

New: 2020-12-03

 $\{\langle true \ code \rangle\}$ $\{\langle false \ code \rangle\}$

Locally sets the test variable \locall _peek_token equal to $\langle token \rangle$ (as an implicit token, not as a token list), and then expands the $\langle function \rangle$. The $\langle token \rangle$ remains in the input stream as the next item after the $\langle function \rangle$. The $\langle token \rangle$ here may be \Box , { or } (assuming normal TeX category codes), i.e. it is not necessarily the next argument which would be grabbed by a normal function.

 $\verb|\peek_gafter:Nw| \land peek_gafter:Nw| \land function \rangle \land token \rangle$

Globally sets the test variable \g_peek_token equal to $\langle token \rangle$ (as an implicit token, not as a token list), and then expands the $\langle function \rangle$. The $\langle token \rangle$ remains in the input stream as the next item after the $\langle function \rangle$. The $\langle token \rangle$ here may be \Box , { or } (assuming normal TeX category codes), i.e. it is not necessarily the next argument which would be grabbed by a normal function.

\l_peek_token Token set by \peek_after:Nw and available for testing as described above.

\g_peek_token Token set by \peek_gafter:Nw and available for testing as described above.

 $\verb|\peek_catcode:NTF| \langle test token \rangle \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}|$

Updated: 2012-12-20 Tests if the next $\langle token \rangle$ in the input stream has the same category code as the $\langle test token \rangle$ (as defined by the test \token_if_eq_catcode:NNTF). Spaces are respected by the test and the $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).

 $\label{eq:catcode_remove:NTF} $$\operatorname{catcode_remove:NTF} \ {\ test \ token} \ {\ true \ code} \ {\ false \ code} \ $$$

Updated: 2012-12-20 Tests if the next $\langle token \rangle$ in the input stream has the same category code as the $\langle test token \rangle$ (as defined by the test \token_if_eq_catcode:NNTF). Spaces are respected by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).

 $\label{eq:node:ntf} $$ \operatorname{code} NTF \ \operatorname{test \ token} \ {\langle true \ code \rangle} \ {\langle false \ code \rangle} $$$

Updated: 2012-12-20 Tests if the next $\langle token \rangle$ in the input stream has the same character code as the $\langle token \rangle$ (as defined by the test \token_if_eq_charcode: NNTF). Spaces are respected by the test and the $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).

 $\ensuremath{\mbox{\sc harcode_remove:NTF}}\$ $\ensuremath{\mbox{\sc harcode_remove:NTF}}\$ $\ensuremath{\mbox{\sc token}}\$ $\ensuremath{\mbox{\sc harcode_remove:NTF}}\$

Tests if the next $\langle token \rangle$ in the input stream has the same character code as the $\langle test token \rangle$ (as defined by the test \token_if_eq_charcode:NNTF). Spaces are respected by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).

 $\verb|\peek_meaning:NTF| \langle test token \rangle \ \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}|$

Updated: 2011-07-02 Tests if the next $\langle token \rangle$ in the input stream has the same meaning as the $\langle test\ token \rangle$ (as defined by the test $\token_if_eq_meaning:NNTF$). Spaces are respected by the test and the $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).

 $\ensuremath{\mbox{peek_meaning_remove:NTF}}\$ $\ensuremath{\mbox{test token}}\$ $\ensuremath{\mbox{\{\langle false code}\rangle\}}\$

Updated: 2011-07-02 Tests if the next $\langle token \rangle$ in the input stream has the same meaning as the $\langle test\ token \rangle$ (as defined by the test \token_if_eq_meaning:NNTF). Spaces are respected by the test and the $\langle token \rangle$ is removed from the input stream if the test is true. The function then places either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ in the input stream (as appropriate to the result of the test).

 $\percent{peek_remove_spaces:n } \peek_remove_spaces:n } {\code}$

New: 2018-10-01 Peeks ahead and detect if the following token is a space (category code 10 and character code 32). If so, removes the token and checks the next token. Once a non-space token is found, the $\langle code \rangle$ will be inserted into the input stream. Typically this will contain a peek operation, but this is not required.

 $\verb|\peek_remove_filler:n \peek_remove_filler:n \end{} \{\langle code \rangle\}$

New: 2022-01-10 Peeks ahead and detect if the following token is a space (category code 10) or has meaning equal to \scan_stop:. If so, removes the token and checks the next token. If neither of these cases apply, expands the next token using f-type expansion, then checks the resulting leading token in the same way. If after expansion the next token is neither of the two test cases, the $\langle code \rangle$ will be inserted into the input stream. Typically this will contain a peek operation, but this is not required.

> TEX hackers note: This is essentially a macro-based implementation of how TEX handles the search for a left brace after for example \everypar, except that any non-expandable token cleanly ends the $\langle filler \rangle$ (i.e. it does not lead to a T_FX error).

> In contrast to T_FX's filler removal, a construct \exp_not:N \foo will be treated in the same way as \foo.

\peek_N_type: <u>TF</u>

 $\perbox{$\operatorname{N}_{type}:TF {\langle true\ code \rangle} {\langle false\ code \rangle}$}$

Updated: 2012-12-20 Tests if the next $\langle token \rangle$ in the input stream can be safely grabbed as an N-type argument. The test is $\langle false \rangle$ if the next $\langle token \rangle$ is either an explicit or implicit begin-group or endgroup token (with any character code), or an explicit or implicit space character (with character code 32 and category code 10), or an outer token (never used in LATEX3) and $\langle true \rangle$ in all other cases. Note that a $\langle true \rangle$ result ensures that the next $\langle token \rangle$ is a valid N-type argument. However, if the next $\langle token \rangle$ is for instance \c _space_token, the test takes the $\langle false \rangle$ branch, even though the next $\langle token \rangle$ is in fact a valid N-type argument. The $\langle token \rangle$ is left in the input stream after the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ (as appropriate to the result of the test).

```
{\bf ek\_analysis\_map\_inline:n \neq analysis\_map\_inline:n \{\langle inline\ function \rangle\}}
```

New: 2020-12-03 Updated: 2022-10-03

> Repeatedly removes one $\langle token \rangle$ from the input stream and applies the $\langle inline\ function \rangle$ to it, until \peek_analysis_map_break: is called. The \(\int interpretation \) receives three arguments for each $\langle token \rangle$ in the input stream:

- $\langle tokens \rangle$, which both o-expand and x-expand to the $\langle token \rangle$. The detailed form of ⟨tokens⟩ may change in later releases.
- $\langle char\ code \rangle$, a decimal representation of the character code of the $\langle token \rangle$, -1 if it is a control sequence.
- $\langle catcode \rangle$, a capital hexadecimal digit which denotes the category code of the $\langle token \rangle$ (0: control sequence, 1: begin-group, 2: end-group, 3: math shift, 4: alignment tab, 6: parameter, 7: superscript, 8: subscript, A: space, B: letter, C: other, D: active). This can be converted to an integer by writing " $\langle catcode \rangle$.

These arguments are the same as for \tl_analysis_map_inline:nn defined in I3tlanalysis. The $\langle char \ code \rangle$ and $\langle catcode \rangle$ do not take the meaning of a control sequence or active character into account: for instance, upon encountering the token \c_group_begin_token in the input stream, \peek_analysis_map_inline:n calls the (inline function) with #1 being \exp_not:n { \c_group_begin_token } (with the current implementation), #2 being -1, and #3 being 0, as for any other control sequence. In contrast, upon encountering an explicit begin-group token $\{$, the $\langle inline\ function \rangle$ is called with arguments \exp_after:wN { \if_false: } \fi:, 123 and 1.

The mapping is done at the current group level, i.e. any local assignments made by the (inline function) remain in effect after the loop. Within the code, \l_peek_token is set equal (as a token, not a token list) to the token under consideration.

\peek_analysis_map_break:

```
\peek_analysis_map_inline:n
```

New: 2020-12-03 Stops the \peek_analysis_map_inline:n loop from seeking more tokens, and inserts $\langle code \rangle$ in the input stream (empty for \peek_analysis_map_break:).

```
\peek_regex:NTF
```

 $\perb{peek_regex:nTF } {\code} {\code} {\cde} {\cde} {\cde}$

Tests if the $\langle tokens \rangle$ that follow in the input stream match the $\langle regular\ expression \rangle$. Any New: 2020-12-03 $\langle tokens \rangle$ that have been read are left in the input stream after the $\langle true\ code \rangle$ or $\langle false$ of regular expressions. The $\langle regular \ expression \rangle$ is implicitly anchored at the start, so for instance \peek_regex:nTF { a } is essentially equivalent to \peek_charcode:NTF a.

> TEXhackers note: Implicit character tokens are correctly considered by \peek regex:nTF as control sequences, while functions that inspect individual tokens (for instance \peek_charcode: NTF) only take into account their meaning.

> The \peek_regex:nTF function only inspects as few tokens as necessary to determine whether the regular expression matches. For instance \peek_regex:nTF { abc | [a-z] } { } { } abc will only inspect the first token a even though the first branch abc of the alternative is preferred in functions such as \peek_regex_remove_once:n. This may have an effect on tokenization if the input stream has not yet been tokenized and category codes are changed.

New: 2020-12-03

Tests if the $\langle tokens \rangle$ that follow in the input stream match the $\langle regex \rangle$. If the test is true, the $\langle tokens \rangle$ are removed from the input stream and the $\langle true\ code \rangle$ is inserted, while if the test is false, the $\langle false\ code \rangle$ is inserted followed by the $\langle tokens \rangle$ that were originally in the input stream. See |3regex for documentation of the syntax of regular expressions. The $\langle regular\ expression \rangle$ is implicitly anchored at the start, so for instance \peek_regex_remove_once:nTF { a } is essentially equivalent to \peek_charcode_remove:NTF a.

TEXhackers note: Implicit character tokens are correctly considered by \peek_regex_-remove_once:nTF as control sequences, while functions that inspect individual tokens (for instance \peek_charcode:NTF) only take into account their meaning.

```
\label{lem:code} $$ \operatorname{\ensurement} \ \operatorname{\ensurement} \ {\ensurement} \ {\ensu
```

If the $\langle tokens \rangle$ that follow in the input stream match the $\langle regex \rangle$, replaces them according to the $\langle replacement \rangle$ as for $\regex_replace_once:nnN$, and leaves the result in the input stream, after the $\langle true\ code \rangle$. Otherwise, leaves $\langle false\ code \rangle$ followed by the $\langle tokens \rangle$ that were originally in the input stream, with no modifications. See I3regex for documentation of the syntax of regular expressions and of the $\langle replacement \rangle$: for instance $\0$ in the $\langle replacement \rangle$ is replaced by the tokens that were matched in the input stream. The $\langle regular\ expression \rangle$ is implicitly anchored at the start. In contrast to $\regex_replace_once:nnN$, no error arises if the $\langle replacement \rangle$ leads to an unbalanced token list: the tokens are inserted into the input stream without issue.

TeXhackers note: Implicit character tokens are correctly considered by \peek_regex_replace_once:nnTF as control sequences, while functions that inspect individual tokens (for instance \peek_charcode:NTF) only take into account their meaning.

23.7 Description of all possible tokens

Let us end by reviewing every case that a given token can fall into. This section is quite technical and some details are only meant for completeness. We distinguish the meaning of the token, which controls the expansion of the token and its effect on TEX's state, and its shape, which is used when comparing token lists such as for delimited arguments. Two tokens of the same shape must have the same meaning, but the converse does not hold.

A token has one of the following shapes.

• A control sequence, characterized by the sequence of characters that constitute its name: for instance, \use:n is a five-letter control sequence.

- An active character token, characterized by its character code (between 0 and 1114111 for LuaT_FX and X_TT_FX and less for other engines) and category code 13.
- A character token, characterized by its character code and category code (one of 1, 2, 3, 4, 6, 7, 8, 10, 11 or 12 whose meaning is described below).

There are also a few internal tokens. The following list may be incomplete in some engines.

- Expanding \the\font results in a token that looks identical to the command that was used to select the current font (such as \tenrm) but it differs from it in shape.
- A "frozen" \relax, which differs from the primitive in shape (but has the same meaning), is inserted when the closing \fi of a conditional is encountered before the conditional is evaluated.
- Expanding \noexpand \langle token \rangle \text{ (when the \langle token \rangle is expandable) results in an internal token, displayed (temporarily) as \notexpanded: \langle token \rangle, whose shape coincides with the \langle token \rangle and whose meaning differs from \rangle relax.
- An \outer endtemplate: can be encountered when peeking ahead at the next token; this expands to another internal token, end of alignment template.
- Tricky programming might access a frozen \endwrite.
- Some frozen tokens can only be accessed in interactive sessions: \cr, \right, \endgroup, \fi, \inaccessible.
- In LuaTeX, there is also the strange case of "bytes" ^^^1100xy where x, y are any two lowercase hexadecimal digits, so that the hexadecimal number ranges from \text{110000}=1114112\$ to~\$1100ff = 1114367. These are used to output individual bytes to files, rather than UTF-8. For the purposes of token comparisons they behave like non-expandable primitive control sequences (not characters) whose \meaning is the_character_\mu followed by the given byte. If this byte is in the range 80-ff this gives an "invalid utf-8 sequence" error: applying \token_to_str:N or \token_to_meaning:N to these tokens is unsafe. Unfortunately, they don't seem to be detectable safely by any means except perhaps Lua code.

The meaning of a (non-active) character token is fixed by its category code (and character code) and cannot be changed. We call these tokens *explicit* character tokens. Category codes that a character token can have are listed below by giving a sample output of the TEX primitive \meaning, together with their IATEX3 names and most common example:

```
1 begin-group character (group_begin, often {),
```

- 2 end-group character (group_end, often }),
- 3 math shift character (math_toggle, often \$),
- 4 alignment tab character (alignment, often &),
- 6 macro parameter character (parameter, often #),
- 7 superscript character (math_superscript, often ^),

```
8 subscript character (math_subscript, often _),
```

- 10 blank space (space, often character code 32),
- 11 the letter (letter, such as A),
- 12 the character (other, such as 0).

Category code 13 (active) is discussed below. Input characters can also have several other category codes which do not lead to character tokens for later processing: 0 (escape), 5 (end line), 9 (ignore), 14 (comment), and 15 (invalid).

The meaning of a control sequence or active character can be identical to that of any character token listed above (with any character code), and we call such tokens *implicit* character tokens. The meaning is otherwise in the following list:

- a macro, used in LATEX3 for most functions and some variables (tl, fp, seq, ...),
- a primitive such as \def or \topmark, used in LATEX3 for some functions,
- a register such as \count123, used in IATEX3 for the implementation of some variables (int, dim, ...),
- a constant integer such as \char"56 or \mathchar"121,
- a font selection command,
- undefined.

Macros can be \protected or not, \long or not (the opposite of what LATEX3 calls nopar), and \outer or not (unused in LATEX3). Their \meaning takes the form

```
\langle prefix \rangle macro: \langle argument \rangle -> \langle replacement \rangle
```

where $\langle prefix \rangle$ is among \protected\long\outer, $\langle argument \rangle$ describes parameters that the macro expects, such as #1#2#3, and $\langle replacement \rangle$ describes how the parameters are manipulated, such as \int eval:n{#2+#1*#3}.

Now is perhaps a good time to mention some subtleties relating to tokens with category code 10 (space). Any input character with this category code (normally, space and tab characters) becomes a normal space, with character code 32 and category code 10.

When a macro takes an undelimited argument, explicit space characters (with character code 32 and category code 10) are ignored. If the following token is an explicit character token with category code 1 (begin-group) and an arbitrary character code, then TeX scans ahead to obtain an equal number of explicit character tokens with category code 1 (begin-group) and 2 (end-group), and the resulting list of tokens (with outer braces removed) becomes the argument. Otherwise, a single token is taken as the argument for the macro: we call such single tokens "N-type", as they are suitable to be used as an argument for a function with the signature : N.

When a macro takes a delimited argument T_EX scans ahead until finding the delimiter (outside any pairs of begin-group/end-group explicit characters), and the resulting list of tokens (with outer braces removed) becomes the argument. Note that explicit space characters at the start of the argument are *not* ignored in this case (and they prevent brace-stripping).

Deprecated functions 23.8

```
\char_lowercase:N
\char_uppercase:N
\char_titlecase:N
\char_foldcase:N
\char_str_foldcase:N * character code changes).
```

 \star \char_lowercase:N $\langle char \rangle$

Converts the $\langle char \rangle$ to the equivalent case-changed character as detailed by the function name (see \str_foldcase:n and \text_titlecase:n for details of these terms). \char_str_lowercase:N * The case mapping is carried out with no context-dependence (cf. \text_uppercase:n, \char_str_uppercase: N * etc.) The str versions always generate "other" (category code 12) characters, whilst the $\c str_{titlecase:N} \star standard versions generate characters with the category code of the <math>\langle char \rangle$ (i.e. only the

New: 2020-01-09

Chapter 24

The **I3prop** package Property lists

expl3 implements a (property list) data type, which contain an unordered list of entries each of which consists of a $\langle key \rangle$ and an associated $\langle value \rangle$. The $\langle key \rangle$ and $\langle value \rangle$ may both be any $\langle balanced text \rangle$, the $\langle key \rangle$ is processed using \tl_to_str:n, meaning that category codes are ignored. It is possible to map functions to property lists such that the function is applied to every key-value pair within the list.

Each entry in a property list must have a unique $\langle key \rangle$: if an entry is added to a property list which already contains the $\langle key \rangle$ then the new entry overwrites the existing one. The \(\langle keys \rangle \) are compared on a string basis, using the same method as \str_if_eq:nn.

Property lists are intended for storing key-based information for use within code. This is in contrast to key-value lists, which are a form of *input* parsed by the l3keys module.

Creating and initialising property lists 24.1

\prop_new:N \prop_new:N \property list \

\prop_new:c

Creates a new $\langle property | list \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle property \ list \rangle$ initially contains no entries.

\prop_clear:c

\prop_clear:N \prop_clear:N \property list \

\prop_gclear:N \prop_gclear:c

Clears all entries from the $\langle property \ list \rangle$.

\prop_clear_new:c

\prop_clear_new:N \prop_clear_new:N \quad property list \rangle

\prop_gclear_new:N

Ensures that the (property list) exists globally by applying \prop_new:N if necessary, \prop_gclear_new:c then applies \prop_(g)clear:N to leave the list empty.

 $\frac{\text{Updated: 2021-11-07}}{\text{keys appear only the last of the values is kept.}} \text{ Sets } \langle property \ list \rangle \text{ to contain key-value pairs given in the second argument. If duplicate}$

Spaces are trimmed around every $\langle key \rangle$ and every $\langle value \rangle$, and if the result of trimming spaces consists of a single brace group then a set of outer braces is removed. This enables both the $\langle key \rangle$ and the $\langle value \rangle$ to contain spaces, commas or equal signs. The $\langle key \rangle$ is then processed by **\tl_to_str:n**. This function correctly detects the = and , signs provided they have the standard category code 12 or they are active.

Notice that in contrast to most keyval lists (e.g. those in l3keys), each key here must be followed with an = sign.

Creates a new constant $\langle property\ list \rangle$ or raises an error if the name is already taken. The $\langle property\ list \rangle$ is set globally to contain key-value pairs given in the second argument, processed in the way described for $prop_set_from_keyval:Nn$. If duplicate keys appear only the last of the values is kept. This function correctly detects the = and , signs provided they have the standard category code 12 or they are active.

Notice that in contrast to most key val lists (e.g. those in |3keys), each key here must be followed with an = sign.

24.2 Adding and updating property list entries

Adds an entry to the $\langle property \ list \rangle$ which may be accessed using the $\langle key \rangle$ and which has $\langle value \rangle$. If the $\langle key \rangle$ is already present in the $\langle property \ list \rangle$, the existing entry is overwritten by the new $\langle value \rangle$. Both the $\langle key \rangle$ and $\langle value \rangle$ may contain any $\langle balanced \ text \rangle$. The $\langle key \rangle$ is stored after processing with $\t1_to_str:n$, meaning that category codes are ignored.

\prop_put_if_new:Nnn \prop_put_if_new:cnn \prop_gput_if_new:Nnn \prop_gput_if_new:cnn

```
\prop_put_if_new: Nnn \property list \property \prop \prop
```

If the $\langle key \rangle$ is present in the $\langle property | list \rangle$ then no action is taken. Otherwise, a new entry is added as described for \prop_put:Nnn.

\prop_concat:NNN \prop_concat:ccc \prop_gconcat:NNN

New: 2021-05-16

 $\prop_concat:NNN \property \ list_1 \property \ list_2 \property \ list_3$

Combines the key-value pairs of $\langle property \ list_2 \rangle$ and $\langle property \ list_3 \rangle$, and saves the result in $\langle pproperty \ list_1 \rangle$. If a key appears in both $\langle property \ list_2 \rangle$ and $\langle property \ list_3 \rangle$ then the last value, namely the value in $\langle property \ list_3 \rangle$ is kept.

\prop_put_from_keyval:Nn \prop_put_from_keyval:cn $\prop_gput_from_keyval:Nn \langle key1 \rangle = \langle value1 \rangle$,

```
\prop_put_from_keyval:Nn \(\rangle property list \rangle \)
                                        {
\prop_gput_from_keyval:cn \langle key2 \rangle = \langle value2 \rangle , ...
```

New: 2021-05-16

Updated: 2021-11-07 Updates the (property list) by adding entries for each key-value pair given in the second argument. The addition is done through \prop put:Nnn, hence if the \(\lambda property \ list \rangle \) already contains some of the keys, the corresponding values are discarded and replaced by those given in the key-value list. If duplicate keys appear in the key-value list then only the last of the values is kept.

> The function is equivalent to storing the key-value pairs in a temporary property list using \prop_set_from_keyval:Nn, then combining \(\rho property \ list \rangle \) with the temporary variable using \prop_concat:NNN. In particular, the $\langle keys \rangle$ and $\langle values \rangle$ are spacetrimmed and unbraced as described in \prop_set_from_keyval:Nn. This function correctly detects the = and, signs provided they have the standard category code 12 or they are active.

Recovering values from property lists 24.3

 $\prop_get: NnN \property list \property \prop \prop$ \prop_get:NnN \prop_get:(NVN|NvN|NoN|cnN|cVN|cvN|coN)

Updated: 2011-08-28

Recovers the $\langle value \rangle$ stored with $\langle key \rangle$ from the $\langle property | list \rangle$, and places this in the $\langle token\ list\ variable \rangle$. If the $\langle key \rangle$ is not found in the $\langle property\ list \rangle$ then the $\langle token\ list \rangle$ variable is set to the special marker \q_no_value. The $\langle token\ list\ variable \rangle$ is set within the current TFX group. See also \prop_get:NnNTF.

\prop_pop:NnN \prop_pop:(NoN|cnN|coN)

 $\prop_pop:NnN \property list \prop_table \prop_table$

Updated: 2011-08-18

Recovers the $\langle value \rangle$ stored with $\langle key \rangle$ from the $\langle property \ list \rangle$, and places this in the $\langle token\ list\ variable \rangle$. If the $\langle key \rangle$ is not found in the $\langle property\ list \rangle$ then the $\langle token\ list \rangle$ list variable is set to the special marker q_no_value . The $\langle key \rangle$ and $\langle value \rangle$ are then deleted from the property list. Both assignments are local. See also \prop_pop:NnNTF.

```
\prop_gpop:NnN
```

\prop_gpop:(NoN|cnN|coN)

Updated: 2011-08-18

 $\prop_gpop: \prop_gpop: \prop_grop: \prop_grop: \prop_grop: \prop_grop : \prop_grop: \pr$

Recovers the $\langle value \rangle$ stored with $\langle key \rangle$ from the $\langle property \ list \rangle$, and places this in the $\langle token\ list\ variable \rangle$. If the $\langle key \rangle$ is not found in the $\langle property\ list \rangle$ then the $\langle token\ list$ variable is set to the special marker $\neq no_value$. The $\langle key \rangle$ and $\langle value \rangle$ are then deleted from the property list. The $\langle property | list \rangle$ is modified globally, while the assignment of the (token list variable) is local. See also \prop_gpop:NnNTF.

 $\label{limits} $$ \operatorname{Prop_item:Nn} \ \langle property \ list \rangle \ \{\langle key \rangle\}$ $$$

New: 2014-07-17 missing, this has an empty expansion.

 $\underline{\hspace{0.2cm}}$ Expands to the $\langle value \rangle$ corresponding to the $\langle key \rangle$ in the $\langle property \ list \rangle$. If the $\langle key \rangle$ is

TeXhackers note: This function is slower than the non-expandable analogue \prop_get:NnN. The result is returned within the \unexpanded primitive (\exp_not:n), which means that the $\langle value \rangle$ does not expand further when appearing in an x-type or e-type argument expansion.

\prop_count:N * \prop_count:N \langle property list \rangle

\prop_count:c *

Leaves the number of key-value pairs in the $\langle property | list \rangle$ in the input stream as an $\langle integer\ denotation \rangle$.

\prop_to_keyval:N * \prop_to_keyval:N \langle property list \rangle

Expands to the $\langle property | list \rangle$ in a key-value notation. Keep in mind that a $\langle property |$ list) is unordered, while key-value interfaces don't necessarily are, so this can't be used for arbitrary interfaces.

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the key-value list does not expand further when appearing in an x-type or e-type argument expansion. It also needs exactly two steps of expansion.

24.4 Modifying property lists

\prop_remove:Nn

 $\prop_remove:(NV|cn|cV)$

New: 2012-05-12

\prop_gremove:Nn

 $\prop_gremove:(NV|cn|cV)$

 $\verb|\prop_remove:Nn| \langle property list \rangle | \{\langle key \rangle\}|$

Removes the entry listed under $\langle key \rangle$ from the $\langle property \ list \rangle$. If the $\langle key \rangle$ is not found in the $\langle property \ list \rangle$ no change occurs, i.e there is no need to test for the existence of a key before deleting it.

24.5Property list conditionals

\prop_if_exist_p:N * \prop_if_exist_p:N \langle property list \rangle

 $\prop_if_exist_p:c * \prop_if_exist:NTF \property list \property \property$

\prop_if_exist:c<u>TF</u> *

 $\prop_{if}_{exist}: NTF \times Tests$ whether the $\langle property | list \rangle$ is currently defined. This does not check that the ⟨property list⟩ really is a property list variable.

New: 2012-03-03

```
\prop_if_empty_p:N * \prop_if_empty_p:N \langle property list \rangle
\displaystyle \begin{array}{l} \prop_if_empty_p:c \; \star \; prop_if_empty:NTF \; \langle property \; list 
angle \; \{\langle true \; code 
angle \} \; \{\langle false \; code 
angle \} \; \} \end{array}
\label{eq:prop_if_empty:NTF} $$ \text{Tests if the } \langle \textit{property list} \rangle$ is empty (containing no entries).
\prop_if_empty:cTF
                                               \star \prop_{if_in:NnTF} \property \ list \end{code} \{\langle true \ code \rangle\} \ \{\langle false \ code \rangle\}
\prop_if_in_p:Nn
\prop_if_in_p:(NV|No|cn|cV|co)
\prop_if_in:NnTF
\prop_if_in:(NV|No|cn|cV|co)TF
                         Updated: 2011-09-15
```

Tests if the $\langle key \rangle$ is present in the $\langle property | list \rangle$, making the comparison using the method described by \str_if_eq:nnTF.

TEXhackers note: This function iterates through every key-value pair in the (property $list\rangle$ and is therefore slower than using the non-expandable \prop_get:NnNTF.

24.6 Recovering values from property lists with branching

The functions in this section combine tests for the presence of a key in a property list with recovery of the associated valued. This makes them useful for cases where different cases follow dependent on the presence or absence of a key in a property list. They offer increased readability and performance over separate testing and recovery phases.

```
\prop_get:NnNTF
                                                              \prop_get:NnNTF \(\rangle property list \rangle \{\langle key \rangle \} \\ \taken list variable \rangle \)
\prop_get:(NVN|NvN|NoN|cnN|cVN|cvN|coN)TF
                                                                 \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
                                    Updated: 2012-05-19
```

If the $\langle key \rangle$ is not present in the $\langle property | list \rangle$, leaves the $\langle false | code \rangle$ in the input stream. The value of the $\langle token\ list\ variable \rangle$ is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property | list \rangle$, stores the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$ without removing it from the $\langle property\ list \rangle$, then leaves the $\langle true$ code in the input stream. The $\langle token\ list\ variable \rangle$ is assigned locally.

```
prop_pop:NnNTF \prop_pop:NnNTF \property list \ {\langle key \rangle} \ \ \ token list variable \ {\langle true \ code \rangle}
\prop_pop:cnNTF {\langle false code \rangle}
```

New: 2011-08-18 If the $\langle key \rangle$ is not present in the $\langle property \ list \rangle$, leaves the $\langle false \ code \rangle$ in the input stream. Updated: 2012-05-19 The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property \ list \rangle$, pops the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$, i.e. removes the item from the $\langle property\ list \rangle$. Both the $\langle property\ list \rangle$ list and the $\langle token\ list\ variable \rangle$ are assigned locally.

 $\prop_gpop:cnNTF {\langle false code \rangle}$

New: 2011-08-18 If the $\langle key \rangle$ is not present in the $\langle property \ list \rangle$, leaves the $\langle false \ code \rangle$ in the input stream. Updated: 2012-05-19 The value of the (token list variable) is not defined in this case and should not be relied upon. If the $\langle key \rangle$ is present in the $\langle property \ list \rangle$, pops the corresponding $\langle value \rangle$ in the $\langle token\ list\ variable \rangle$, i.e. removes the item from the $\langle property\ list \rangle$. The $\langle property\ list \rangle$ is modified globally, while the $\langle token\ list\ variable \rangle$ is assigned locally.

24.7 Mapping over property lists

All mappings are done at the current group level, i.e. any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

\prop_map_function:NN \prop_map_function:cN

\prop_map_function:NN \(\text{property list} \) \(\text{function} \)

Updated: 2013-01-08

Applies $\langle function \rangle$ to every $\langle entry \rangle$ stored in the $\langle property\ list \rangle$. The $\langle function \rangle$ receives two arguments for each iteration: the $\langle key \rangle$ and associated $\langle value \rangle$. The order in which (entries) are returned is not defined and should not be relied upon. To pass further arguments to the $\langle function \rangle$, see \prop_map_tokens:Nn.

\prop_map_inline:Nn \prop_map_inline:cn

\prop_map_inline: Nn \(\rhoperty list \) \{ \(\lambda \text{inline function} \) \}

Updated: 2013-01-08

Applies $\langle inline\ function \rangle$ to every $\langle entry \rangle$ stored within the $\langle property\ list \rangle$. The $\langle inline\ function \rangle$ function should consist of code which receives the $\langle key \rangle$ as #1 and the $\langle value \rangle$ as #2. The order in which $\langle entries \rangle$ are returned is not defined and should not be relied upon.

\prop_map_tokens:Nn ☆ \prop_map_tokens:cn ☆

 $prop_map_tokens:Nn \langle property \ list \rangle \{\langle code \rangle\}$

Analogue of \prop_map_function: NN which maps several tokens instead of a single function. The $\langle code \rangle$ receives each key-value pair in the $\langle property | list \rangle$ as two trailing brace groups. For instance,

```
\prop_map_tokens:Nn \l_my_prop { \str_if_eq:nnT { mykey } }
```

expands to the value corresponding to mykey: for each pair in \l_my_prop the function $\mathsf{str}_{\mathsf{if}} = \mathsf{q:nnT}$ receives mykey, the $\langle key \rangle$ and the $\langle value \rangle$ as its three arguments. For that specific task, \prop_item:Nn is faster.

\prop_map_break: ☆ \prop_map_break:

Updated: 2012-06-29 Used to terminate a \prop_map_... function before all entries in the \(\lambda property \ list \rangle \) have been processed. This normally takes place within a conditional statement, for example

```
\prop_map_inline:Nn \l_my_prop
    \str_if_eq:nnTF { #1 } { bingo }
      { \prop_map_break: }
        % Do something useful
      }
```

Use outside of a \prop_map_... scenario leads to low level TEX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before further items are taken from the input stream. This depends on the design of the mapping function.

 $\verb|\prop_map_break:n| \Leftrightarrow \verb|\prop_map_break:n| \{\langle code \rangle\}$

Updated: 2012-06-29 Used to terminate a $prop_map_...$ function before all entries in the $property\ list$ have been processed, inserting the $\langle code \rangle$ after the mapping has ended. This normally takes place within a conditional statement, for example

```
\prop_map_inline:Nn \l_my_prop
    \str_if_eq:nnTF { #1 } { bingo }
      { \prop_map_break:n { <code> } }
        % Do something useful
     }
 }
```

Use outside of a \prop_map_... scenario leads to low level TEX errors.

TeXhackers note: When the mapping is broken, additional tokens may be inserted before the $\langle code \rangle$ is inserted into the input stream. This depends on the design of the mapping function.

24.8 Viewing property lists

\prop_show: N

\prop_show:N \(\rhoperty list\)

\prop_show:c

Displays the entries in the $\langle property \ list \rangle$ in the terminal.

Updated: 2021-04-29

\prop_log:N \prop_log:c $\verb|\prop_log:N| \langle property \ list \rangle|$

New: 2014-08-12

Writes the entries in the $\langle property | list \rangle$ in the log file.

Updated: 2021-04-29

Scratch property lists 24.9

\l_tmpa_prop Scratch property lists for local assignment. These are never used by the kernel code, and \1_tmpb_prop so are safe for use with any LATEX3-defined function. However, they may be overwritten New: 2012-06-23 by other non-kernel code and so should only be used for short-term storage.

\g_tmpa_prop Scratch property lists for global assignment. These are never used by the kernel code, and \g_tmpb_prop so are safe for use with any LATEX3-defined function. However, they may be overwritten $_{\text{New: 2012-06-23}}$ by other non-kernel code and so should only be used for short-term storage.

Constants 24.10

\c_empty_prop A permanently-empty property list used for internal comparisons.

Chapter 25

New: 2012-01-07

The **I3skip** package Dimensions and skips

IATEX3 provides two general length variables: dim and skip. Lengths stored as dim variables have a fixed length, whereas skip lengths have a rubber (stretch/shrink) component. In addition, the muskip type is available for use in math mode: this is a special form of skip where the lengths involved are determined by the current math font (in mu). There are common features in the creation and setting of length variables, but for clarity the functions are grouped by variable type.

25.1 Creating and initialising dim variables

| dim_new:N | dim_new:N | dimension |
| Creates a new | dimension | or raises an error if the name is already taken. The declaration is global. The | dimension | sinitially equal to 0 pt.
| dim_const:Nn | dim_const:Nn | dimension | sinitially equal to 0 pt.
dim_const:Cn	Creates a new constant	dimension	expression	or raises an error if the name is already taken. The value of the	dimension	or raises an error if the name is already taken. The
dim_zero:N	dim_zero:N	dimension	expression	dim_zero:N	dimension	or raises an error if the name is already taken. The
dim_zero:N	dim_zero:N	dimension	expression	dim_zero:N	dimension	or raises an error if the name is already taken. The
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dim_zero:N	dimension	or raises an error if the name is already taken. The			
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				
dim_zero:N	dimension	or raises an error if the name is already taken. The				

```
\label{lem:dim_if_exist_p:N dim_if_exist_p:N dimension} $$ \dim_{if_exist_p:c} \star \dim_{if_exist:NTF} \langle dimension \rangle {\langle true\ code \rangle} {\langle false\ code \rangle} $$$ \dim_{if_exist:NTE} \star $$$ Tests whether the $\langle dimension \rangle$ is currently defined. This does not check that the $$$$ Mew: 2012-03-03 $$$$ \langle dimension \rangle$ really is a dimension variable.
```

25.2 Setting dim variables

```
\dim_add:Nn
                             \verb|\dim_add:Nn| \langle dimension \rangle \{\langle dimension| expression \rangle\}|
         \dim_add:cn
                             Adds the result of the \langle dimension \ expression \rangle to the current content of the \langle dimension \rangle.
         \dim_gadd:Nn
         \dim_gadd:cn
         Updated: 2011-10-22
                             \dim_set:Nn \dimension \ {\dimension expression \}
         \dim_set:Nn
         \dim set:cn
                             Sets \langle dimension \rangle to the value of \langle dimension \ expression \rangle, which must evaluate to a length
         \dim_gset:Nn
                             with units.
         \dim_gset:cn
         Updated: 2011-10-22
\dim_set_eq:NN
                             \dim_{\text{set\_eq:NN}} \langle dimension_1 \rangle \langle dimension_2 \rangle
\dim_set_eq:(cN|Nc|cc)
                             Sets the content of \langle dimension_1 \rangle equal to that of \langle dimension_2 \rangle.
\dim_gset_eq:NN
\dim_gset_eq:(cN|Nc|cc)
                             \dim_sub:Nn \dimension \ {\dimension expression \}
         \dim_sub:Nn
         \dim_sub:cn
                             Subtracts the result of the (dimension expression) from the current content of the
         \dim_gsub:Nn
                             \langle dimension \rangle.
         \dim_gsub:cn
         Updated: 2011-10-22
```

25.3 Utilities for dimension calculations

```
\frac{\mbox{\sc dim\_abs:n} \quad \star \quad \mbox{\sc dim\_abs:n} \quad \{\mbox{\sc dimexpr}\} \ \mbox{\sc dimexpr} \mbox{\sc dimexp} \ \mbox{\sc dimexp} \ \mbox{\sc dimexp} \
```

```
\verb|\dim_ratio:nn| \Leftrightarrow \verb|\dim_ratio:nn| \{\langle dimexpr_1 \rangle\} | \{\langle dimexpr_2 \rangle\}|
```

Updated: 2011-10-22 Parses the two (dimension expressions) and converts the ratio of the two to a form suitable for use inside a $\langle dimension \ expression \rangle$. This ratio is then left in the input stream, allowing syntax such as

```
\dim_set:Nn \l_my_dim
  { 10 pt * \dim_ratio:nn { 5 pt } { 10 pt } }
```

The output of \dim_ratio:nn on full expansion is a ratio expression between two integers, with all distances converted to scaled points. Thus

```
\tl_set:Nx \l_my_tl { \dim_ratio:nn { 5 pt } { 10 pt } }
\tl_show:N \l_my_tl
```

displays 327680/655360 on the terminal.

25.4 Dimension expression conditionals

```
\dim_{compare_p:nNn} \star \dim_{compare_p:nNn} \{\langle dimexpr_1 \rangle\} \langle relation \rangle \{\langle dimexpr_2 \rangle\}
\verb|\dim_compare:nNn] $\underline{TF} \star \dim_compare:nNnTF|
                                                 \{\langle \texttt{dimexpr}_1 \rangle\} \ \langle \texttt{relation} \rangle \ \{\langle \texttt{dimexpr}_2 \rangle\}
                                                 \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
```

This function first evaluates each of the \(dimension expressions \) as described for \(\dim_-\) eval:n. The two results are then compared using the $\langle relation \rangle$:

> Equal Greater than Less than

This function is less flexible than \dim_compare:nTF but around 5 times faster.

This function evaluates the $\langle dimension \; expressions \rangle$ as described for $\langle dim_eval:n$ and compares consecutive result using the corresponding $\langle relation \rangle$, namely it compares $\langle dimexpr_1 \rangle$ and $\langle dimexpr_2 \rangle$ using the $\langle relation_1 \rangle$, then $\langle dimexpr_2 \rangle$ and $\langle dimexpr_3 \rangle$ using the $\langle relation_2 \rangle$, until finally comparing $\langle dimexpr_N \rangle$ and $\langle dimexpr_{N+1} \rangle$ using the $\langle relation_N \rangle$. The test yields true if all comparisons are true. Each $\langle dimension \; expression \rangle$ is evaluated only once, and the evaluation is lazy, in the sense that if one comparison is false, then no other $\langle dimension \; expression \rangle$ is evaluated and no other comparison is performed. The $\langle relations \rangle$ can be any of the following:

```
Equal = or ==
Greater than or equal to >=
Greater than >=
Less than or equal to <=
Less than <
Not equal !=
```

This function is more flexible than \dim_compare:nNnTF but around 5 times slower.

This function evaluates the $\langle test\ dimension\ expression \rangle$ and compares this in turn to each of the $\langle dimension\ expression\ cases \rangle$. If the two are equal then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted. The function $\dim_case:nn$, which does nothing if there is no match, is also available. For example

leaves "Medium" in the input stream.

25.5 Dimension expression loops

```
\label{linear_continuous} $$\dim_{o} \min_{\alpha} \dim_{o} \pi_1:Nnn {\langle dimexpr_1 \rangle} {\langle relation \rangle} {\langle dimexpr_2 \rangle} {\langle code \rangle}$$
```

Places the $\langle code \rangle$ in the input stream for TEX to process, and then evaluates the relationship between the two $\langle dimension \ expressions \rangle$ as described for \dim_compare:nNnTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.

```
\label{limited} $$\dim_{\phi} = \min_{\phi} \dim_{\phi} \min_{\phi} \|\phi_{\phi}\|_{L^{2}(GMe)} + \|\phi_{\phi}\|_{L^{2}(GMe)} +
```

Places the $\langle code \rangle$ in the input stream for TEX to process, and then evaluates the relationship between the two $\langle dimension\ expressions \rangle$ as described for \dim_compare:nNnTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.

```
\verb|\dim_until_do:nNnn| \Leftrightarrow \verb|\dim_until_do:nNnn| \{ \langle dimexpr_1 \rangle \} | \langle relation \rangle | \{ \langle dimexpr_2 \rangle \} | \{ \langle code \rangle \} | \} |
```

Evaluates the relationship between the two $\langle dimension \ expressions \rangle$ as described for $\dim_compare:nNnTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by T_EX the test is repeated, and a loop occurs until the test is true.

 $\dim_{\text{while_do:nNnn}} \Leftrightarrow \dim_{\text{while_do:nNnn}} \{\langle dimexpr_1 \rangle\} \langle relation \rangle \{\langle dimexpr_2 \rangle\} \{\langle code \rangle\}$

Evaluates the relationship between the two $\langle dimension \ expressions \rangle$ as described for $\dim_{compare:nNnTF}$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed by T_FX the test is repeated, and a loop occurs until the test is false.

 $\dim_{\operatorname{do_until:nn}} \Leftrightarrow \dim_{\operatorname{do_until:nn}} {\langle \operatorname{dimension} \ \operatorname{relation} \rangle} {\langle \operatorname{code} \rangle}$

Updated: 2013-01-13 Places the $\langle code \rangle$ in the input stream for T_FX to process, and then evaluates the (dimension relation) as described for \dim compare:nTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.

\dim_do_while:nn ☆ \dim_do_while:nn {\dimension relation}} {\code}}

Updated: 2013-01-13 Places the $\langle code \rangle$ in the input stream for T_FX to process, and then evaluates the (dimension relation) as described for \dim_compare:nTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.

\dim_until_do:nn ☆ \dim_until_do:nn {\dimension relation}} {\code}}

Updated: 2013-01-13 Evaluates the $\langle dimension \ relation \rangle$ as described for \dim_compare:nTF, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by TEX the test is repeated, and a loop occurs until the test is true.

\dim_while_do:nn ☆ \dim_while_do:nn {\dimension relation}} {\code}}

Updated: 2013-01-13 Evaluates the $\langle dimension \ relation \rangle$ as described for \dim_compare:nTF, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed by T_FX the test is repeated, and a loop occurs until the test is false.

25.6 Dimension step functions

\dim_step_function:nnnN ☆ \dim_step_function:nnnN {⟨initial value⟩} {⟨step⟩} {⟨final value⟩} ⟨function⟩

New: 2018-02-18 This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be dimension expressions. The $\langle function \rangle$ is then placed in front of each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$). The $\langle step \rangle$ must be non-zero. If the $\langle step \rangle$ is positive, the loop stops when the $\langle value \rangle$ becomes larger than the $\langle final\ value \rangle$. If the $\langle step \rangle$ is negative, the loop stops when the $\langle value \rangle$ becomes smaller than the $\langle final\ value \rangle$. The $\langle function \rangle$ should absorb one argument.

 $\label{limin} $$\dim_{step_{inline:nnnn} {\langle initial\ value \rangle} {\langle step \rangle} {\langle final\ value \rangle} {\langle code \rangle} $$$

New: 2018-02-18 This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be dimension expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with #1 replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (#1).

\dim_step_variable:nnnNn \dim_step_variable:nnnNn

New: 2018-02-18

```
{\langle initial \ value \rangle} \ {\langle step \rangle} \ {\langle final \ value \rangle} \ {\langle tl \ var \rangle} \ {\langle code \rangle}
```

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be dimension expressions. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl \ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl \ var \rangle$.

Using dim expressions and variables 25.7

 $\verb|\dim_eval:n * \\dim_eval:n { | (dimension expression) | }$

Updated: 2011-10-22 Evaluates the $\langle dimension \ expression \rangle$, expanding any dimensions and token list variables within the \(\lambda expression \rangle \) to their content (without requiring \dim_use: N/\tl_use: N) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a $\langle dimension \ denotation \rangle$ after two expansions. This is expressed in points (pt), and requires suitable termination if used in a TEX-style assignment as it is not an $\langle internal \ dimension \rangle$.

 $\dim_{sign:n} \star \dim_{sign:n} \{\langle dimexpr \rangle\}$

New: 2018-11-03 Evaluates the $\langle dimexpr \rangle$ then leaves 1 or 0 or -1 in the input stream according to the sign of the result.

\dim_use:c *

\dim_use:N * \dim_use:N \dimension \

Recovers the content of a $\langle dimension \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a $\langle dimension \rangle$ is required (such as in the argument of \dim_{e} val:n).

TEXhackers note: \dim_use: N is the TEX primitive \the: this is one of several LATEX3 names for this primitive.

 $\dim_{to_decimal:n} \star \dim_{to_decimal:n} \{\langle dimexpr \rangle\}$

New: 2014-07-15 Evaluates the $\langle dimension \ expression \rangle$, and leaves the result, expressed in points (pt) in the input stream, with no units. The result is rounded by TEX to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

\dim_to_decimal:n { 1bp }

leaves 1.00374 in the input stream, i.e. the magnitude of one "big point" when converted to (TeX) points.

 $\dim_{to_decimal_in_bp:n } \star \dim_{to_decimal_in_bp:n } {\langle dimexpr \rangle}$

New: 2014-07-15 Evaluates the (dimension expression), and leaves the result, expressed in big points (bp) in the input stream, with no units. The result is rounded by TEX to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

```
\dim_to_decimal_in_bp:n { 1pt }
```

leaves 0.99628 in the input stream, i.e. the magnitude of one (T_FX) point when converted to big points.

```
\dim_{to\_decimal_in\_sp:n} \star \dim_{to\_decimal_in\_sp:n} \{\langle dimexpr \rangle\}
```

New: 2015-05-18 Evaluates the $\langle dimension \ expression \rangle$, and leaves the result, expressed in scaled points (sp) in the input stream, with no units. The result is necessarily an integer.

```
\label{limit} $$\dim_{t^{-1}} = \dim_{t^{-1}} \left( \dim_{t^{-1}} \operatorname{decimal_{in}}_{t^{-1}} \left( \dim_{t^{-1}} \right) \right) $$
                             New: 2014-07-15
```

Evaluates the $\langle dimension \ expressions \rangle$, and leaves the value of $\langle dimexpr_1 \rangle$, expressed in a unit given by $\langle dimexpr_2 \rangle$, in the input stream. The result is a decimal number, rounded by T_FX to four or five decimal places. If the decimal part of the result is zero, it is omitted, together with the decimal marker.

For example

```
\dim to decimal in unit:nn { 1bp } { 1mm }
```

leaves 0.35277 in the input stream, i.e. the magnitude of one big point when converted to millimetres.

Note that this function is not optimised for any particular output and as such may give different results to \dim_to_decimal_in_bp:n or \dim_to_decimal_in_sp:n. In particular, the latter is able to take a wider range of input values as it is not limited by the ability to calculate a ratio using ε -TeX primitives, which is required internally by \dim_to_decimal_in_unit:nn.

```
\dim_{to_{fp:n}} \star \dim_{to_{fp:n}} {\langle dimexpr \rangle}
```

New: 2012-05-08 Expands to an internal floating point number equal to the value of the $\langle dimexpr \rangle$ in pt. Since dimension expressions are evaluated much faster than their floating point equivalent, \dim_to_fp:n can be used to speed up parts of a computation where a low precision and a smaller range are acceptable.

25.8 Viewing dim variables

\dim_show:N \dim_show:N \dimension \

 $\underline{\dim_{\underline{\text{show}:c}}}$ Displays the value of the $\langle dimension \rangle$ on the terminal.

\dim_show:n

\dim_show:n {\dimension expression\}

New: 2011-11-22 Displays the result of evaluating the $\langle dimension \; expression \rangle$ on the terminal.

Updated: 2015-08-07

\dim_log:N

\dim_log:N \dimension \

\dim_log:c

Writes the value of the $\langle dimension \rangle$ in the log file. New: 2014-08-22

Updated: 2015-08-03

\dim_log:n

\dim_log:n {\dimension expression}}

New: 2014-08-22 Writes the result of evaluating the $\langle dimension \ expression \rangle$ in the log file.

Updated: 2015-08-07

Constant dimensions 25.9

\c_max_dim The maximum value that can be stored as a dimension. This can also be used as a component of a skip.

\c_zero_dim A zero length as a dimension. This can also be used as a component of a skip.

Scratch dimensions 25.10

\l_tmpa_dim Scratch dimension for local assignment. These are never used by the kernel code, and so \1_tmpb_dim are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

\g_tmpa_dim Scratch dimension for global assignment. These are never used by the kernel code, and so are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Creating and initialising skip variables 25.11

\skip_new:N \skip_new:N \skip)

 $\$ Creates a new $\langle skip \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle skip \rangle$ is initially equal to 0 pt.

 $\slip_const:Nn \slip_const:Nn \slip \{\langle skip \ expression \rangle\}$ \skip_const:cn Creates a new constant $\langle skip \rangle$ or raises an error if the name is already taken. The value New: 2012-03-05 of the $\langle skip \rangle$ is set globally to the $\langle skip \ expression \rangle$. \skip_zero:N \skip_zero:N \langle skip \rangle \skip_zero:c Sets $\langle skip \rangle$ to 0 pt. \skip_gzero:N \skip_gzero:c \skip_zero_new:N \skip_zero_new:N \(skip \) \skip_zero_new:c Ensures that the $\langle skip \rangle$ exists globally by applying \skip_new:N if necessary, then applies \skip_gzero_new:N $\$ iskip_(g)zero:N to leave the $\langle skip \rangle$ set to zero. \skip_gzero_new:c New: 2012-01-07 $\sin_{if}_{exist_p:c} \star \sin_{if}_{exist:NTF} \langle skip \rangle \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}$ $\sl p_{if}=xist:NTF \times Tests$ whether the $\langle skip \rangle$ is currently defined. This does not check that the $\langle skip \rangle$ really \slash is a skip variable. New: 2012-03-03

25.12 Setting skip variables

\skip_add:Nn \skip_add:cn Adds the result of the $\langle skip \; expression \rangle$ to the current content of the $\langle skip \rangle$. \skip_gadd:Nn \skip_gadd:cn Updated: 2011-10-22 \skip_set:Nn \skip_set:cn Sets $\langle skip \rangle$ to the value of $\langle skip \ expression \rangle$, which must evaluate to a length with units \skip_gset:Nn and may include a rubber component (for example 1 cm plus 0.5 cm. \skip_gset:cn Updated: 2011-10-22 $\sline \sline \sline$ \skip_set_eq:NN \skip_set_eq:(cN|Nc|cc) Sets the content of $\langle skip_1 \rangle$ equal to that of $\langle skip_2 \rangle$. \skip_gset_eq:NN \skip_gset_eq:(cN|Nc|cc) \skip_sub:Nn \skip_sub:cn Subtracts the result of the $\langle skip \; expression \rangle$ from the current content of the $\langle skip \rangle$. \skip_gsub:Nn \skip_gsub:cn Updated: 2011-10-22

25.13 Skip expression conditionals

```
\slip_if_eq_p:nn * \slip_if_eq_p:nn {\langle skipexpr_1 \rangle} {\langle skipexpr_2 \rangle}
\sin TF \star \sin TF = \sin TF
                                   \{\langle skipexpr_1 \rangle\}\ \{\langle skipexpr_2 \rangle\}
                                   \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
```

This function first evaluates each of the (skip expressions) as described for \skip_eval:n. The two results are then compared for exact equality, i.e. both the fixed and rubber components must be the same for the test to be true.

```
\sline \frac{1}{skip_if_finite_p:n} \star \frac{1}{skip_if_finite_p:n} \{\langle skip_expr \rangle\}
\left[ \frac{TF}{t} + \left( \frac{TF}{t} \right) \right]
```

New: 2012-03-05 Evaluates the (skip expression) as described for \skip_eval:n, and then tests if all of its components are finite.

Using skip expressions and variables 25.14

Updated: 2011-10-22 Evaluates the $\langle skip \; expression \rangle$, expanding any skips and token list variables within the ⟨expression⟩ to their content (without requiring \skip_use:N/\tl_use:N) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a $\langle glue\ denotation \rangle$ after two expansions. This is expressed in points (pt), and requires suitable termination if used in a TeX-style assignment as it is not an \(\internal glue \).

 $\sline \sline \sline$

Recovers the content of a $\langle skip \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a $\langle dimension \rangle$ or $\langle skip \rangle$ is required (such as in the argument of \skip_eval:n).

TEXhackers note: \skip_use:N is the TEX primitive \the: this is one of several LATEX3 names for this primitive.

Viewing skip variables 25.15

\skip_show: N \skip_show:N \(skip \)

\skip_show:c Updated: 2015-08-03

Displays the value of the $\langle skip \rangle$ on the terminal.

 $\sin {\langle skip expression \rangle}$ \skip_show:n

New: 2011-11-22 Displays the result of evaluating the $\langle skip \; expression \rangle$ on the terminal.

Updated: 2015-08-07

\skip_log:N

\skip_log:N \(\skip \)

\skip_log:c

Writes the value of the $\langle skip \rangle$ in the log file.

New: 2014-08-22

Updated: 2015-08-03

\skip_log:n

\skip_log:n {\langle skip expression \rangle}

New: 2014-08-22 Writes the result of evaluating the $\langle skip \; expression \rangle$ in the log file.

Updated: 2015-08-07

25.16 Constant skips

\c_max_skip

The maximum value that can be stored as a skip (equal to \c_max_dim in length), with Updated: 2012-11-02 no stretch nor shrink component.

\c_zero_skip

A zero length as a skip, with no stretch nor shrink component.

Updated: 2012-11-01

Scratch skips 25.17

\l_tmpa_skip Scratch skip for local assignment. These are never used by the kernel code, and so are \l_tmpb_skip safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

\g_tmpa_skip Scratch skip for global assignment. These are never used by the kernel code, and so are \g_tmpb_skip safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

25.18Inserting skips into the output

\skip_horizontal:N \skip_horizontal:N \skip\

\skip_horizontal:c \skip_horizontal:n {\langle skipexpr\rangle}

\skip_horizontal:n Inserts a horizontal $\langle skip \rangle$ into the current list. The argument can also be a $\langle dim \rangle$.

Updated: 2011-10-22

TEXhackers note: \skip_horizontal: N is the TEX primitive \hskip renamed.

```
\skip_vertical:N \skip_vertical:N \skip\
\skip_vertical:c \skip_vertical:n \{\skipexpr\}\
\skip_vertical:n \skip_vertical:n \text{skipexpr}\}
\text{Inserts a vertical \(\skip\) into the current list. The argument can also be a \(\skip\).

\text{TFXhackers note: \skip_vertical:N is the TFX primitive \vskip renamed.}
```

25.19 Creating and initialising muskip variables

```
\muskip_new:N \muskip_new:N \muskip)
           \muskip_new:c
                            Creates a new \langle muskip \rangle or raises an error if the name is already taken. The declaration
                            is global. The \langle muskip \rangle is initially equal to 0 mu.
       \muskip_const:Nn \muskip_const:Nn \muskip\ {\muskip expression\}
        \muskip_const:cn
                            Creates a new constant \langle muskip \rangle or raises an error if the name is already taken. The
             New: 2012-03-05 value of the \langle muskip \rangle is set globally to the \langle muskip \; expression \rangle.
        \muskip_zero:N \skip_zero:N \muskip\
        \muskip_zero:c
                            Sets \langle muskip \rangle to 0 mu.
        \muskip_gzero:N
         \muskip_gzero:c
    \muskip_zero_new:N \muskip_zero_new:N \muskip\
    \muskip_zero_new:c
                           Ensures that the \langle muskip \rangle exists globally by applying \muskip_new: N if necessary, then
    \muskip_gzero_new:N
    \muskip_gzero_new:c applies \muskip_(g)zero:N to leave the \langle muskip \set to zero.
             New: 2012-01-07
\mbox{\mbox{\tt muskip\_if\_exist\_p:N}} \ \star \mbox{\mbox{\tt muskip\_if\_exist\_p:N}} \ \langle \mbox{\tt muskip} \rangle
\mbox{\tt muskip\_if\_exist:N} \mbox{\tt TF} \star \mbox{\tt Tests} whether the \langle muskip \rangle is currently defined. This does not check that the \langle muskip \rangle
\mbox{muskip\_if\_exist:c} \ \ \star
                            really is a muskip variable.
             New: 2012-03-03
```

25.20 Setting muskip variables

\muskip_set:Nn

\muskip_set:Nn \langle muskip \langle \langle muskip expression \rangle \rangle

\muskip_set:cn

\muskip_gset:Nn

\muskip_gset:cn

Updated: 2011-10-22

Sets $\langle muskip \rangle$ to the value of $\langle muskip \ expression \rangle$, which must evaluate to a math length with units and may include a rubber component (for example 1 mu plus 0.5 mu.

\muskip_set_eq:NN

\muskip_set_eq:(cN|Nc|cc)

\muskip_gset_eq:NN

\muskip_gset_eq:(cN|Nc|cc)

 $\mbox{\mbox{\tt muskip_set_eq:NN}} \ \langle \mbox{\tt muskip}_1 \rangle \ \langle \mbox{\tt muskip}_2 \rangle$

Sets the content of $\langle muskip_1 \rangle$ equal to that of $\langle muskip_2 \rangle$.

\muskip_sub:Nn

\muskip_sub:cn \muskip_gsub:Nn

\muskip_gsub:cn

Updated: 2011-10-22

Subtracts the result of the $\langle muskip | expression \rangle$ from the current content of the $\langle muskip \rangle$.

25.21Using muskip expressions and variables

\muskip_eval:n * \muskip_eval:n {\muskip expression}}

Updated: 2011-10-22 Evaluates the $\langle muskip \ expression \rangle$, expanding any skips and token list variables within the \(\langle expression \rangle \) to their content (without requiring \(\muskip_use: \mathbb{N} \\tal_use: \mathbb{N}\) and applying the standard mathematical rules. The result of the calculation is left in the input stream as a (muglue denotation) after two expansions. This is expressed in mu, and requires suitable termination if used in a TeX-style assignment as it is not an \(\lambda internal \) $muglue \rangle$.

\muskip_use:N * \muskip_use:N \langle muskip \rangle

\muskip_use:c *

Recovers the content of a $\langle skip \rangle$ and places it directly in the input stream. An error is raised if the variable does not exist or if it is invalid. Can be omitted in places where a ⟨dimension⟩ is required (such as in the argument of \muskip_eval:n).

TEXhackers note: \muskip_use: N is the TEX primitive \the: this is one of several LATEX3 names for this primitive.

25.22Viewing muskip variables

\muskip_show:N \muskip_show:N \muskip \

\muskip_show:c

Displays the value of the $\langle muskip \rangle$ on the terminal.

Updated: 2015-08-03

\muskip_show:n \muskip_show:n {\muskip expression}} New: 2011-11-22 Displays the result of evaluating the $\langle muskip \ expression \rangle$ on the terminal. Updated: 2015-08-07 \muskip_log:N \dagma muskip \dagma \muskip_log:N \muskip_log:c Writes the value of the $\langle muskip \rangle$ in the log file. New: 2014-08-22 Updated: 2015-08-03 \muskip_log:n {\muskip expression}} New: 2014-08-22 Writes the result of evaluating the $\langle muskip \ expression \rangle$ in the log file.

25.23Constant muskips

\c_max_muskip The maximum value that can be stored as a muskip, with no stretch nor shrink component.

\c_zero_muskip A zero length as a muskip, with no stretch nor shrink component.

25.24 Scratch muskips

\l_tmpa_muskip Scratch muskip for local assignment. These are never used by the kernel code, and so \l_tmpb_muskip are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

Updated: 2015-08-07

\g_tmpa_muskip Scratch muskip for global assignment. These are never used by the kernel code, and so \g_tmpb_muskip are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

25.25Primitive conditional

```
⟨true code⟩
     \else:
     (false)
     \fi:
```

Compare two dimensions. The $\langle relation \rangle$ is one of $\langle \cdot, = \text{ or } \rangle$ with category code 12.

TeXhackers note: This is the TeX primitive \ifdim.

Chapter 26

The l3keys package Key-value interfaces

The key–value method is a popular system for creating large numbers of settings for controlling function or package behaviour. The system normally results in input of the form

```
\MyModuleSetup{
   key-one = value one,
   key-two = value two
}

or

\MyModuleMacro[
   key-one = value one,
   key-two = value two
]{argument}
```

for the user.

The high level functions here are intended as a method to create key–value controls. Keys are themselves created using a key–value interface, minimising the number of functions and arguments required. Each key is created by setting one or more *properties* of the key:

```
\keys_define:nn { mymodule }
    {
       key-one .code:n = code including parameter #1,
       key-two .tl_set:N = \l_mymodule_store_tl
    }
```

These values can then be set as with other key–value approaches:

```
\keys_set:nn { mymodule }
    {
       key-one = value one,
       key-two = value two
    }
```

As illustrated, keys are created inside a $\langle module \rangle$: a set of related keys, typically those for a single module/IAT_FX 2ε package. See Section for suggestions on how to divide large numbers of keys for a single module.

At a document level, \keys_set:nn is used within a document function, for example

```
\DeclareDocumentCommand \MyModuleSetup { m }
  { \keys_set:nn { mymodule } { #1 } }
\DeclareDocumentCommand \MyModuleMacro { o m }
  {
    \group_begin:
      \keys_set:nn { mymodule } { #1 }
      % Main code for \MyModuleMacro
    \group_end:
```

Key names may contain any tokens, as they are handled internally using \t1 to str:n. As discussed in section 26.2, it is suggested that the character / is reserved for sub-division of keys into logical groups. Functions and variables are not expanded when creating key names, and so

```
\tl_set:Nn \l_mymodule_tmp_tl { key }
\keys_define:nn { mymodule }
    \label{local_mymodule_tmp_tl} \code:n = code
```

creates a key called \l_mymodule_tmp_tl, and not one called key.

Creating keys 26.1

 $\enskip \label{list} $$ \enskip \ens$

Updated: 2017-11-14 Parses the $\langle keyval \ list \rangle$ and defines the keys listed there for $\langle module \rangle$. The $\langle module \rangle$ name is treated as a string. In practice the $\langle module \rangle$ should be chosen to be unique to the module in question (unless deliberately adding keys to an existing module).

> The $\langle keyval | list \rangle$ should consist of one or more key names along with an associated key property. The properties of a key determine how it acts. The individual properties are described in the following text; a typical use of \keys_define:nn might read

```
\keys_define:nn { mymodule }
   keyname .code:n = Some~code~using~#1,
   keyname .value_required:n = true
```

where the properties of the key begin from the . after the key name.

The various properties available take either no arguments at all, or require one or more arguments. This is indicated in the name of the property using an argument specification. In the following discussion, each property is illustrated attached to an arbitrary $\langle key \rangle$, which when used may be supplied with a $\langle value \rangle$. All key definitions are local.

Key properties are applied in the reading order and so the ordering is significant. Key properties which define "actions", such as .code:n, .tl_set:N, etc., override one another. Some other properties are mutually exclusive, notably .value_required:n and .value_forbidden:n, and so they replace one another. However, properties covering non-exclusive behaviours may be given in any order. Thus for example the following definitions are equivalent.

Note that with the exception of the special .undefine: property, all key properties define the key within the current TeX scope.

```
.bool_set:N
.bool_set:C
.bool_gset:N
.bool_gset:c
```

\langle key \rangle .bool_set:N = \langle boolean variable \rangle

Defines $\langle key \rangle$ to set $\langle boolean\ variable \rangle$ to $\langle value \rangle$ (which must be either "true" or "false"). If the variable does not exist, it will be created globally at the point that the key is set up.

.bool_set_inverse:N
.bool_set_inverse:C
.bool_gset_inverse:C
.bool_gset_inverse:C

 $\langle key \rangle$.bool_set_inverse:N = $\langle boolean \ variable \rangle$

Defines $\langle key \rangle$ to set $\langle boolean \ variable \rangle$ to the logical inverse of $\langle value \rangle$ (which must be either "true" or "false"). If the $\langle boolean \ variable \rangle$ does not exist, it will be created globally at the point that the key is set up.

New: 2011-08-28 Updated: 2013-07-08

.choice: $\langle key \rangle$.choice:

Sets $\langle key \rangle$ to act as a choice key. Each valid choice for $\langle key \rangle$ must then be created, as discussed in section 26.3.

```
.choices:nn .choices:(Vn|on|xn)
```

 $\langle key \rangle$.choices:nn = $\{\langle choices \rangle\}\ \{\langle code \rangle\}$

New: 2011-08-21 Updated: 2013-07-10

```
.clist_set:N
.clist_set:c
.clist_gset:N
.clist_gset:c
```

New: 2011-09-11

 $\langle \texttt{key} \rangle \ . \texttt{clist_set:N} = \langle \textit{comma list variable} \rangle$

Defines $\langle key \rangle$ to set $\langle comma \ list \ variable \rangle$ to $\langle value \rangle$. Spaces around commas and empty items will be stripped. If the variable does not exist, it is created globally at the point that the key is set up.

```
\langle key \rangle .code:n = \{\langle code \rangle\}
.code:n
```

Updated: 2013-07-10 Stores the $\langle code \rangle$ for execution when $\langle key \rangle$ is used. The $\langle code \rangle$ can include one parameter (#1), which will be the $\langle value \rangle$ given for the $\langle key \rangle$.

```
\langle key \rangle .cs_set:Np = \langle control sequence \rangle \langle arg. spec. \rangle
.cs_set:Np
.cs_set:cp
                                Defines \langle key \rangle to set \langle control\ sequence \rangle to have \langle arg.\ spec. \rangle and replacement text \langle value \rangle.
.cs_set_protected:Np
.cs_set_protected:cp
.cs_gset:Np
.cs_gset:cp
.cs_gset_protected:Np
.cs\_gset\_protected:cp
             New: 2020-01-11
```

```
.default:n
.default:(V|o|x)
```

 $\langle key \rangle$.default:n = $\{\langle default \rangle\}$

Creates a $\langle default \rangle$ value for $\langle key \rangle$, which is used if no value is given. This will be used Updated: 2013-07-09 if only the key name is given, but not if a blank (value) is given:

```
\keys_define:nn { mymodule }
  {
   key .code:n
                   = Hello~#1.
   key .default:n = World
 }
\keys set:nn { mymodule }
   key = Fred, % Prints 'Hello Fred'
   key,
                % Prints 'Hello World'
    key =
                % Prints 'Hello '
 }
```

The default does not affect keys where values are required or forbidden. Thus a required value cannot be supplied by a default value, and giving a default value for a key which cannot take a value does not trigger an error.

```
.dim_set:N
                         \langle key \rangle .dim_set:N = \langle dimension \rangle
.dim_set:c
```

Defines $\langle key \rangle$ to set $\langle dimension \rangle$ to $\langle value \rangle$ (which must a dimension expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

Updated: 2020-01-17

.dim_gset:N

.dim_gset:c

```
.fp_set:N
                       \langle key \rangle .fp_set:N = \langle floating point \rangle
.fp_set:c
.fp_gset:N
```

Defines $\langle key \rangle$ to set $\langle floating\ point \rangle$ to $\langle value \rangle$ (which must a floating point expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

```
Updated: 2020-01-17
```

.fp_gset:c

```
\langle key \rangle .groups:n = \{\langle groups \rangle\}
.groups:n
```

New: 2013-07-14 Defines $\langle key \rangle$ as belonging to the $\langle qroups \rangle$ declared. Groups provide a "secondary axis" for selectively setting keys, and are described in Section 26.7.

```
\langle key \rangle .inherit:n = \{\langle parents \rangle\}
                  .inherit:n
                  New: 2016-11-22 Specifies that the \langle key \rangle path should inherit the keys listed as any of the \langle parents \rangle (a
                                   comma list), which can be a module or a subgroup. For example, after setting
                                          \keys_define:nn { foo } { test .code:n = \tl_show:n {#1} }
                                         \keys_define:nn { } { bar .inherit:n = foo }
                                   setting
                                          \keys_set:nn { bar } { test = a }
                                   will be equivalent to
                                          \keys_set:nn { foo } { test = a }
                                   Inheritance applies at point of use, not at definition, thus keys may be added to the
                                   \langle parent \rangle after the use of .inherit:n and will be active. If more than one \langle parent \rangle is
                                   specified, the presence of the \langle key \rangle will be tested for each in turn, with the first successful
                                   hit taking priority.
              .initial:n
                                   \langle key \rangle .initial:n = \{\langle value \rangle\}
              .initial:(V|o|x)
                                   Initialises the \langle key \rangle with the \langle value \rangle, equivalent to
              Updated: 2013-07-09
                                           \keys set:nn \{\langle module \rangle\} \{\langle key \rangle = \langle value \rangle\}
              .int_set:N
                                    \langle key \rangle .int_set:N = \langle integer \rangle
              .int_set:c
                                   Defines \langle key \rangle to set \langle integer \rangle to \langle value \rangle (which must be an integer expression). If the
              .int_gset:N
                                   variable does not exist, it is created globally at the point that the key is set up. The key
              .int_gset:c
                                   will require a value at point-of-use unless a default is set.
              Updated: 2020-01-17
.legacy_if_set:n
                                   \left(key\rightarrow .legacy_if_set:n = \left(switch\right)
.legacy_if_gset:n
                                   Defines \langle key \rangle to set legacy \if \langle switch \rangle to \langle value \rangle (which must be either "true" or
.legacy_if_set_inverse:n
                                   "false"). The \langle switch \rangle is the name of the switch without the leading \if.
.legacy_if_gset_inverse:n
                                         The inverse versions will set the \langle switch \rangle to the logical opposite of the \langle value \rangle.
              Updated: 2022-01-15
                                   \langle key \rangle .meta:n = \{\langle keyval \ list \rangle\}
              .meta:n
              Updated: 2013-07-10 Makes \langle key \rangle a meta-key, which will set \langle keyval \ list \rangle in one go. The \langle keyval \ list \rangle can refer
                                   as #1 to the value given at the time the \langle key \rangle is used (or, if no value is given, the \langle key \rangle's
                                   default value).
                                   \langle key \rangle .meta:nn = \{\langle path \rangle\} \{\langle keyval \ list \rangle\}
```

the current one. The $\langle keyval \ list \rangle$ can refer as #1 to the value given at the time the $\langle key \rangle$

New: 2013-07-10 Makes $\langle key \rangle$ a meta-key, which will set $\langle keyval \ list \rangle$ in one go using the $\langle path \rangle$ in place of

is used (or, if no value is given, the $\langle key \rangle$'s default value).

.meta:nn

.multichoice: $\langle key \rangle$.multichoice:

New: 2011-08-21 Sets $\langle key \rangle$ to act as a multiple choice key. Each valid choice for $\langle key \rangle$ must then be created, as discussed in section 26.3.

.multichoices:nn

 $\langle key \rangle$.multichoices:nn $\{\langle choices \rangle\}$ $\{\langle code \rangle\}$.multichoices:(Vn|on|xn)

Updated: 2013-07-10

Sets $\langle key \rangle$ to act as a multiple choice key, and defines a series $\langle choices \rangle$ which are im-New: 2011-08-21 plemented using the $\langle code \rangle$. Inside $\langle code \rangle$, \lambda_keys_choice_tl will be the name of the choice made, and \l_keys_choice_int will be the position of the choice in the list of $\langle choices \rangle$ (indexed from 1). Choices are discussed in detail in section 26.3.

.muskip_set:N $\langle key \rangle$.muskip_set:N = $\langle muskip \rangle$

.muskip_set:c

.muskip_gset:N

Defines $\langle key \rangle$ to set $\langle muskip \rangle$ to $\langle value \rangle$ (which must be a muskip expression). If the variable does not exist, it is created globally at the point that the key is set up. The key .muskip_gset:c will require a value at point-of-use unless a default is set.

New: 2019-05-05

Updated: 2020-01-17

.prop_put:N

.prop put:c .prop_gput:N .prop_gput:c

Defines $\langle key \rangle$ to put the $\langle value \rangle$ onto the $\langle property | list \rangle$ stored under the $\langle key \rangle$. If the variable does not exist, it is created globally at the point that the key is set up.

New: 2019-01-31

 $\langle key \rangle$.skip_set:N = $\langle skip \rangle$.skip_set:N

.skip_set:c

.skip_gset:N

.skip_gset:c

Defines $\langle key \rangle$ to set $\langle skip \rangle$ to $\langle value \rangle$ (which must be a skip expression). If the variable does not exist, it is created globally at the point that the key is set up. The key will require a value at point-of-use unless a default is set.

Updated: 2020-01-17

 $\langle key \rangle$.str_set:N = $\langle string \ variable \rangle$.str_set:N

.str set:c

.str_gset:N

Defines $\langle key \rangle$ to set $\langle string\ variable \rangle$ to $\langle value \rangle$. If the variable does not exist, it is created globally at the point that the key is set up. .str_gset:c

New: 2021-10-30

 $\langle key \rangle$.str_set_x:N = $\langle string \ variable \rangle$.str_set_x:N

.str set x:c

.str_gset_x:N

Defines $\langle key \rangle$ to set $\langle string\ variable \rangle$ to $\langle value \rangle$, which will be subjected to an x-type .str_gset_x:c expansion (i.e. using \str_set:Nx). If the variable does not exist, it is created globally at the point that the key is set up.

New: 2021-10-30

 $\langle key \rangle$.tl_set:N = $\langle token\ list\ variable \rangle$.tl set:N

Defines $\langle key \rangle$ to set $\langle token\ list\ variable \rangle$ to $\langle value \rangle$. If the variable does not exist, it is .tl_gset:c created globally at the point that the key is set up.

26.2 Sub-dividing keys

When creating large numbers of keys, it may be desirable to divide them into several sub-groups for a given module. This can be achieved either by adding a sub-division to the module name:

As illustrated, the best choice of token for sub-dividing keys in this way is /. This is because of the method that is used to represent keys internally. Both of the above code fragments set the same key, which has full name mymodule/subgroup/key.

As illustrated in the next section, this subdivision is particularly relevant to making multiple choices.

26.3 Choice and multiple choice keys

The l3keys system supports two types of choice key, in which a series of pre-defined input values are linked to varying implementations. Choice keys are usually created so that the various values are mutually-exclusive: only one can apply at any one time. "Multiple" choice keys are also supported: these allow a selection of values to be chosen at the same time.

Mutually-exclusive choices are created by setting the .choice: property:

```
\keys_define:nn { mymodule }
    { key .choice: }
```

For keys which are set up as choices, the valid choices are generated by creating sub-keys of the choice key. This can be carried out in two ways.

In many cases, choices execute similar code which is dependent only on the name of the choice or the position of the choice in the list of all possibilities. Here, the keys can share the same code, and can be rapidly created using the .choices:nn property.

```
\keys_define:nn { mymodule }
 {
   key .choices:nn =
      { choice-a, choice-b, choice-c }
      {
        You~gave~choice~'\tl_use:N \l_keys_choice_tl',~
        which~is~in~position~\int_use:N \l_keys_choice_int \c_space_tl
        in~the~list.
      }
 }
```

The index \l_keys_choice_int in the list of choices starts at 1.

\l_keys_choice_tl

\l_keys_choice_int Inside the code block for a choice generated using .choices:nn, the variables \l_keys_choice_tl and \l_keys_choice_int are available to indicate the name of the current choice, and its position in the comma list. The position is indexed from 1. Note that, as with standard key code generated using .code:n, the value passed to the key (i.e. the choice name) is also available as #1.

> On the other hand, it is sometimes useful to create choices which use entirely different code from one another. This can be achieved by setting the .choice: property of a key, then manually defining sub-keys.

```
\keys_define:nn { mymodule }
 {
   key .choice:,
   key / choice-a .code:n = code-a,
   key / choice-b .code:n = code-b,
   key / choice-c .code:n = code-c,
```

It is possible to mix the two methods, but manually-created choices should not use \l_keys_choice_tl or \l_keys_choice_int. These variables do not have defined behaviour when used outside of code created using .choices:nn (i.e. anything might happen).

It is possible to allow choice keys to take values which have not previously been defined by adding code for the special unknown choice. The general behavior of the unknown key is described in Section 26.6. A typical example in the case of a choice would be to issue a custom error message:

```
\keys_define:nn { mymodule }
   key .choice:,
   key / choice-a .code:n = code-a,
   key / choice-b .code:n = code-b,
   key / choice-c .code:n = code-c,
```

Multiple choices are created in a very similar manner to mutually-exclusive choices, using the properties .multichoice: and .multichoices:nn. As with mutually exclusive choices, multiple choices are define as sub-keys. Thus both

```
\keys_define:nn { mymodule }
      key .multichoices:nn =
        { choice-a, choice-b, choice-c }
        {
          You~gave~choice~'\tl_use:N \l_keys_choice_tl',~
          which ~ is ~ in ~ position ~
          \int_use:N \l_keys_choice_int \c_space_tl
          in~the~list.
    }
and
  \keys_define:nn { mymodule }
      key .multichoice:,
      key / choice-a .code:n = code-a,
      key / choice-b .code:n = code-b,
      key / choice-c .code:n = code-c,
    }
are valid.
    When a multiple choice key is set
  \keys_set:nn { mymodule }
      key = { a , b , c } % 'key' defined as a multiple choice
```

each choice is applied in turn, equivalent to a clist mapping or to applying each value individually:

Thus each separate choice will have passed to it the \l_keys_choice_tl and \l_keys_-choice_int in exactly the same way as described for .choices:nn.

26.4 Key usage scope

Some keys will be used as settings which have a strictly limited scope of usage. Some will be only available once, others will only be valid until typesetting begins. To allow formats to support this in a structured way, I3keys allows this information to be specified using the .usage:n property.

 $\langle key \rangle$.usage:n = $\langle scope \rangle$.usage:n

New: 2022-01-10 Defines the $\langle key \rangle$ to have usage within the $\langle scope \rangle$, which should be one of general, preamble or load.

\l_keys_usage_load_prop \l_keys_usage_preamble_prop

New: 2022-01-10

13keys itself does not attempt to redefine keys based on the usage scope. Rather, this information is made available with these two property lists. These hold an entry for each module (prefix); the value of each entry is a comma-separated list of the usage-restricted key(s).

Setting keys 26.5

 $\ensuremath{\verb|keys_set:nn||} \{\ensuremath{\verb|keys_set:nn||} \{\ensuremath{\verb|keys_set:nn||} \} \}$

\keys_set:nn $\keys_set:(nV|nv|no)$

Parses the $\langle keyval \ list \rangle$, and sets those keys which are defined for $\langle module \rangle$. The behaviour Updated: 2017-11-14 on finding an unknown key can be set by defining a special unknown key: this is illustrated later.

\l_keys_key_str \l_keys_value_tl

Updated: 2020-02-08

For each key processed, information of the full path of the key, the name of the key and \l_keys_path_str the value of the key is available within three token list variables. These may be used within the code of the key.

> The value is everything after the =, which may be empty if no value was given. This is stored in \l_keys_value_tl, and is not processed in any way by \keys_set:nn.

> The path of the key is a "full" description of the key, and is unique for each key. It consists of the module and full key name, thus for example

\keys_set:nn { mymodule } { key-a = some-value }

has path mymodule/key-a while

```
\keys set:nn { mymodule } { subset / key-a = some-value }
```

has path mymodule/subset/key-a. This information is stored in \l_keys_path_str.

The *name* of the key is the part of the path after the last /, and thus is not unique. In the preceding examples, both keys have name key-a despite having different paths. This information is stored in \l_keys_key_str.

26.6 Handling of unknown keys

If a key has not previously been defined (is unknown), \keys_set:nn looks for a special unknown key for the same module, and if this is not defined raises an error indicating that the key name was unknown. This mechanism can be used for example to issue custom error texts.

These functions set keys which are known for the $\langle module \rangle$, and simply ignore other keys. The $\ensuremath{\mbox{keys_set_known:nn}}$ function parses the $\langle keyval \ list \rangle$, and sets those keys which are defined for $\langle module \rangle$. Any keys which are unknown are not processed further by the parser. In addition, $\ensuremath{\mbox{keys_set_known:nnN}}$ stores the key-value pairs in the $\langle tl \rangle$ in comma-separated form (i.e. an edited version of the $\langle keyval \ list \rangle$). When a $\langle root \rangle$ is given ($\ensuremath{\mbox{keys_set_known:nnnN}}$), the key-value entries are returned relative to this point in the key tree. When it is absent, only the key name and value are provided. The correct list is returned by nested calls.

26.7 Selective key setting

In some cases it may be useful to be able to select only some keys for setting, even though these keys have the same path. For example, with a set of keys defined using

the use of \keys_set:nn attempts to set all four keys. However, in some contexts it may only be sensible to set some keys, or to control the order of setting. To do this, keys may be assigned to *groups*: arbitrary sets which are independent of the key tree. Thus modifying the example to read

```
\keys_define:nn { mymodule }
  {
    key-one
              .code:n
                         = \{ \my_func:n \{\#1\} \},
    key-one
              .groups:n = { first }
    key-two
              .tl_set:N = \l_my_a_tl
              .groups:n = { first }
    key-two
    key-three .tl set: N = \l my b tl
    key-three .groups:n = { second }
    kev-four
             .fp_set:N = \l_my_a_fp
  }
```

assigns key-one and key-two to group first, key-three to group second, while key-four is not assigned to a group.

Selective key setting may be achieved either by selecting one or more groups to be made "active", or by marking one or more groups to be ignored in key setting.

```
\label{lem:nnn} $$ \eys_{set_filter:nnn} $$ \eys_{set_filter:nnn} {\module} {\groups} {\eys_{list}} $$ \eys_{set_filter:nnnN} {\module} {\groups} {\eys_{list}} $$ \eys_{set_filter:nnnN} {\module} {\groups} {\eys_{list}} $$ \eys_{set_filter:nnnN} {\module} {\groups} {\eys_{list}} $$ \eys_{list} $$ \eys_{
```

Activates key filtering in an "opt-out" sense: keys assigned to any of the $\langle groups \rangle$ specified are ignored. The $\langle groups \rangle$ are given as a comma-separated list. Unknown keys are not assigned to any group and are thus always set. The key-value pairs for each key which is filtered out are stored in the $\langle tl \rangle$ in a comma-separated form (*i.e.* an edited version of the $\langle keyval\ list \rangle$). The \keys_set_filter:nnn version skips this stage.

Use of \keys_set_filter:nnnN can be nested, with the correct residual $\langle keyval \ list \rangle$ returned at each stage. In the version which takes a $\langle root \rangle$ argument, the key list is returned relative to that point in the key tree. In the cases without a $\langle root \rangle$ argument, only the key names and values are returned.

```
\label{list} $$ \ensuremath{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}}{\mbox{\colored}{\mbox{\colored}{\mbox{\colored}}{\mbox{\colored}{\mbox{\colored}}{\mbox{\colored}}}}}}}}} $$ \end{tabular} $\mbox{\colored}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}{\mbox{\colored}}}}} }} $$ $$ \end{tabular} $$ $$ \end{tabular
```

Activates key filtering in an "opt-in" sense: only keys assigned to one or more of the $\langle groups \rangle$ specified are set. The $\langle groups \rangle$ are given as a comma-separated list. Unknown keys are not assigned to any group and are thus never set.

26.8 Digesting keys

```
\label{list} $$ \ensuremath{$\operatorname{keys\_precompile:nnN} \ \{\langle module \rangle\} \ \{\langle keyval \ list \rangle\} \ \langle t1 \rangle $} $$
```

New: 2022-03-09 Parses the $\langle keyval \ list \rangle$ as for \keys_set:nn, placing the resulting code for those which set variables or functions into the $\langle tl \rangle$. Thus this function "precompiles" the keyval list into a set of results which can be applied rapidly.

26.9 Utility functions for keys

```
\ensuremath{$\langle \rangle \in \mathcal{K}(\mathbb{R}^n) \times \mathcal{K}(\mathbb{R}^n) \in \mathcal{K}(\mathbb{R}^n) \times \mathcal{K}(\mathbb{R}^n) \in 
\label{eq:lemma:temma:tem} $$ \ker = \operatorname{Inn}_{\overline{Z}} * $$ Tests if the $$ \langle key \rangle$ exists for $$ \langle module \rangle$, i.e. if any code has been defined for $$ \langle key \rangle$. }
\keys_if_exist:neTF *
                                      Updated: 2022-01-10
       \label{eq:condition} $$ \ensuremath{$\mathbb{F}$} $$ \ensuremath{$\mathbb{F}$} $$ \ensuremath{$\mathbb{F}$} $$ \ensuremath{$\mathbb{F}$} $$
                                                                                                                                                                                       {\langle false code \rangle}
                                                                                                              New: 2011-08-21
                                                                                            Updated: 2017-11-14
                                                                                                                                Tests if the \langle choice \rangle is defined for the \langle key \rangle within the \langle module \rangle, i.e. if any code has been
                                                                                                                                defined for \langle key \rangle / \langle choice \rangle. The test is false if the \langle key \rangle itself is not defined.
                                       \keys_show:nn
                                                                                                                                 \ensuremath{\verb|keys_show:nn|} \{\langle module \rangle\} \{\langle key \rangle\}
                                      Updated: 2015-08-09 Displays in the terminal the information associated to the \langle key \rangle for a \langle module \rangle, including
                                                                                                                                 the function which is used to actually implement it.
                                                                                                                                 \ensuremath{\mbox{keys\_log:nn } \{\langle module \rangle\} \ \{\langle key \rangle\}}
                                       \keys_log:nn
```

New: 2014-08-22 Writes in the log file the information associated to the $\langle key \rangle$ for a $\langle module \rangle$. See also Updated: 2015-08-09 \keys_show:nn which displays the result in the terminal.

26.10 Low-level interface for parsing key-val lists

To re-cap from earlier, a key-value list is input of the form

```
KeyOne = ValueOne ,
KeyTwo = ValueTwo ,
KeyThree
```

where each key-value pair is separated by a comma from the rest of the list, and each key-value pair does not necessarily contain an equals sign or a value! Processing this type of input correctly requires a number of careful steps, to correctly account for braces, spaces and the category codes of separators.

While the functions described earlier are used as a high-level interface for processing such input, in special circumstances you may wish to use a lower-level approach. The low-level parsing system converts a $\langle key_value\ list\rangle$ into $\langle keys\rangle$ and associated $\langle values\rangle$. After the parsing phase is completed, the resulting keys and values (or keys alone) are available for further processing. This processing is not carried out by the low-level parser itself, and so the parser requires the names of two functions along with the key—value list. One function is needed to process key—value pairs (it receives two arguments), and a second function is required for keys given without any value (it is called with a single argument).

The parser does not double # tokens or expand any input. Active tokens = and , appearing at the outer level of braces are converted to category "other" (12) so that the

parser does not "miss" any due to category code changes. Spaces are removed from the ends of the keys and values. Keys and values which are given in braces have exactly one set removed (after space trimming), thus

```
key = {value here},
and
  key = value here,
are treated identically.
```

New: 2020-12-19 Parses the $\langle key-value\ list \rangle$ into a series of $\langle keys \rangle$ and associated $\langle values \rangle$, or keys alone Updated: 2021-05-10 (if no $\langle value \rangle$ was given). $\langle code_1 \rangle$ receives each $\langle key \rangle$ (with no $\langle value \rangle$) as a trailing brace group, whereas $\langle code_2 \rangle$ is appended by two brace groups, the $\langle key \rangle$ and $\langle value \rangle$. The order of the $\langle keys \rangle$ in the $\langle key-value\ list \rangle$ is preserved. Thus

```
\keyval_parse:nnn
      { \use_none:nn { code 1 } }
      { \use none:nnn { code 2 } }
      { key1 = value1 , key2 = value2, key3 = , key4 }
is converted into an input stream
    \use_none:nnn { code 2 } { key1 } { value1 }
```

```
\use_none:nnn { code 2 } { key2 } { value2 }
\use_none:nnn { code 2 } { key3 } { }
\use none:nn { code 1 } { key4 }
```

Note that there is a difference between an empty value (an equals sign followed by nothing) and a missing value (no equals sign at all). Spaces are trimmed from the ends of the $\langle key \rangle$ and $\langle value \rangle$, then one *outer* set of braces is removed from the $\langle key \rangle$ and $\langle value \rangle$ as part of the processing. If you need exactly the output shown above, you'll need to either x-type or e-type expand the function.

TeXhackers note: The result of each list element is returned within \exp_not:n, which means that the converted input stream does not expand further when appearing in an x-type or e-type argument expansion.

Updated: 2021-05-10 Parses the $\langle key-value\ list \rangle$ into a series of $\langle keys \rangle$ and associated $\langle values \rangle$, or keys alone (if no $\langle value \rangle$ was given). $\langle function_1 \rangle$ should take one argument, while $\langle function_2 \rangle$ should absorb two arguments. After $\ensuremath{\mbox{keyval_parse:NNn}}$ has parsed the $\langle key_value\ list \rangle$, $\langle function_1 \rangle$ is used to process keys given with no value and $\langle function_2 \rangle$ is used to process keys given with a value. The order of the $\langle keys \rangle$ in the $\langle key-value\ list \rangle$ is preserved. Thus

```
\keyval_parse:NNn \function:n \function:nn
  \{ \text{ key1 = value1 , key2 = value2, key3 = , key4 } \}
```

is converted into an input stream

```
\function:nn { key1 } { value1 }
\function:nn { key2 } { value2 }
\function:nn { key3 } { }
\function:n { key4 }
```

Note that there is a difference between an empty value (an equals sign followed by nothing) and a missing value (no equals sign at all). Spaces are trimmed from the ends of the $\langle key \rangle$ and $\langle value \rangle$, then one outer set of braces is removed from the $\langle key \rangle$ and $\langle value \rangle$ as part of the processing.

This shares the implementation of \keyval_parse:nnn, the difference is only semantically.

TEXhackers note: The result is returned within \exp_not:n, which means that the converted input stream does not expand further when appearing in an x-type or e-type argument expansion.

Chapter 27

The **I3**intarray package: Fast global integer arrays

27.1**13intarray** documentation

For applications requiring heavy use of integers, this module provides arrays which can be accessed in constant time (contrast 13seq, where access time is linear). These arrays have several important features

- The size of the array is fixed and must be given at point of initialisation
- The absolute value of each entry has maximum $2^{30} 1$ (i.e. one power lower than the usual \c_max_int ceiling of $2^{31} - 1$)

The use of intarray data is therefore recommended for cases where the need for fast access is of paramount importance.

 $\operatorname{Intarray_new:Nn \setminus intarray_new:Nn \langle intarray \, var \rangle \, \{\langle size \rangle\}}$

\intarray_new:cn

Evaluates the integer expression $\langle size \rangle$ and allocates an $\langle integer\ array\ variable \rangle$ with that New: 2018-03-29 number of (zero) entries. The variable name should start with \g _ because assignments are always global.

\intarray_count:N * \intarray_count:N \(\intarray \, var \)

Expands to the number of entries in the (integer array variable). Contrarily to \seq_-New: 2018-03-29 count: N this is performed in constant time.

\intarray_gset:cnn

\intarray_gset:Nnn \intarray_gset:Nnn \(intarray var \) \(\(\lambda o ition \) \} \(\lambda value \) \}

Stores the result of evaluating the integer expression (value) into the (integer array New: 2018-03-29 variable at the (integer expression) $\langle position \rangle$. If the $\langle position \rangle$ is not between 1 and the \intarray_count:N, or the $\langle value \rangle$'s absolute value is bigger than $2^{30} - 1$, an error occurs. Assignments are always global.

```
\intarray_const_from_clist: Nn \intarray_const_from_clist: Nn \intarray var \rangle \intexpr clist \rangle
    \intarray_const_from_clist:cn
                         New: 2018-05-04
                             Creates a new constant (integer array variable) or raises an error if the name is already
                             taken. The \langle integer\ array\ variable \rangle is set (globally) to contain as its items the results of
                             evaluating each \langle integer\ expression \rangle in the \langle comma\ list \rangle.
       \intarray_gzero:N \intarray_gzero:N \intarray var \rangle
       \intarray_gzero:c
                            Sets all entries of the (integer array variable) to zero. Assignments are always global.
              New: 2018-05-04
    \intarray_item:Nn ★ \intarray_item:Nn ⟨intarray var⟩ {⟨position⟩}
    \intarray_item:cn *
                            Expands to the integer entry stored at the (integer expression) (position) in the (integer
              New: 2018-03-29 array variable). If the \(\lambda position\rangle\) is not between 1 and the \intarray_count:N, an error
\intarray_rand_item:N * \intarray_rand_item:N \langle intarray var \rangle
\intarray_rand_item:c *
                             Selects a pseudo-random item of the \langle integer\ array \rangle. If the \langle integer\ array \rangle is empty,
              New: 2018-05-05 produce an error.
        \intarray_show:N \intarray_show:N \intarray var \)
        \intarray_show:c \intarray_log:N \( intarray var \)
        \intarray_log:N
                             Displays the items in the \langle integer\ array\ variable \rangle in the terminal or writes them in the
         \intarray_log:c
                             log file.
              New: 2018-05-04
```

27.1.1 Implementation notes

It is a wrapper around the \fontdimen primitive, used to store arrays of integers (with a restricted range: absolute value at most $2^{30} - 1$). In contrast to |3seq sequences the access to individual entries is done in constant time rather than linear time, but only integers can be stored. More precisely, the primitive \fontdimen stores dimensions but the |3intarray package transparently converts these from/to integers. Assignments are always global.

While LuaTEX's memory is extensible, other engines can "only" deal with a bit less than 4×10^6 entries in all \fontdimen arrays combined (with default TEX Live settings).

Chapter 28

The l3fp package: Floating points

A decimal floating point number is one which is stored as a significand and a separate exponent. The module implements expandably a wide set of arithmetic, trigonometric, and other operations on decimal floating point numbers, to be used within floating point expressions. Floating point expressions support the following operations with their usual precedence.

- Basic arithmetic: addition x + y, subtraction x y, multiplication x * y, division x/y, square root \sqrt{x} , and parentheses.
- Comparison operators: x < y, x <= y, x > ? y, x ! = y etc.
- Boolean logic: sign sign x, negation !x, conjunction x && y, disjunction x || y, ternary operator x ? y : z.
- Exponentials: $\exp x$, $\ln x$, x^y , $\log x$.
- Integer factorial: fact x.
- Trigonometry: $\sin x$, $\cos x$, $\tan x$, $\cot x$, $\sec x$, $\csc x$ expecting their arguments in radians, and $\sin dx$, $\cos dx$, $\tan dx$, $\cot dx$, $\sec dx$, $\csc dx$ expecting their arguments in degrees.
- Inverse trigonometric functions: $a\sin x$, $a\cos x$, $a\tan x$, $a\cot x$, $a\sec x$, $a\csc x$ giving a result in radians, and $a\sin dx$, $a\cos dx$, $a\tan dx$, $a\cot dx$, $a\sec dx$, $a\sec dx$ giving a result in degrees.

(not yet) Hyperbolic functions and their inverse functions: $\sinh x$, $\cosh x$, $\tanh x$, $\coth x$, $\operatorname{sech} x$, $\operatorname{csch} x$, and $\operatorname{asinh} x$, $\operatorname{acosh} x$, atanh x, acoth x, asech x, acch x.

- Extrema: $\max(x_1, x_2, ...), \min(x_1, x_2, ...), abs(x)$.
- Rounding functions, controlled by two optional values, n (number of places, 0 by default) and t (behavior on a tie, nan by default):
 - $-\operatorname{trunc}(x,n)$ rounds towards zero,
 - floor(x, n) rounds towards $-\infty$,

- $\operatorname{ceil}(x, n)$ rounds towards $+\infty$,
- round(x, n, t) rounds to the closest value, with ties rounded to an even value by default, towards zero if t = 0, towards $+\infty$ if t > 0 and towards $-\infty$ if t < 0.

And (not yet) modulo, and "quantize".

- Random numbers: rand(), randint(m, n).
- Constants: pi, deg (one degree in radians).
- Dimensions, automatically expressed in points, e.g., pc is 12.
- Automatic conversion (no need for \\tauture _use:N) of integer, dimension, and skip variables to floating point numbers, expressing dimensions in points and ignoring the stretch and shrink components of skips.
- Tuples: (x_1, \ldots, x_n) that can be stored in variables, added together, multiplied or divided by a floating point number, and nested.

Floating point numbers can be given either explicitly (in a form such as 1.234e-34, or -.0001), or as a stored floating point variable, which is automatically replaced by its current value. A "floating point" is a floating point number or a tuple thereof. See section 28.9.1 for a description of what a floating point is, section 28.9.2 for details about how an expression is parsed, and section 28.9.3 to know what the various operations do. Some operations may raise exceptions (error messages), described in section 28.7.

An example of use could be the following.

The operation round can be used to limit the result's precision. Adding +0 avoids the possibly undesirable output -0, replacing it by +0. However, the l3fp module is mostly meant as an underlying tool for higher-level commands. For example, one could provide a function to typeset nicely the result of floating point computations.

See the documentation of siunitx for various options of \num.

Creating and initialising floating point variables 28.1

\fp_new:N

 $fp_new:N \langle fp var \rangle$

\fp_new:c

Creates a new $\langle fp \ var \rangle$ or raises an error if the name is already taken. The declaration is Updated: 2012-05-08 global. The $\langle fp \ var \rangle$ is initially +0.

\fp_const:Nn

 $fp_const:Nn \langle fp \ var \rangle \{\langle floating \ point \ expression \rangle\}$

\fp_const:cn

Creates a new constant $\langle fp \ var \rangle$ or raises an error if the name is already taken. The Updated: 2012-05-08 $\langle fp \ var \rangle$ is set globally equal to the result of evaluating the $\langle floating \ point \ expression \rangle$.

\fp_zero:N

\fp_zero:N \langle fp var \rangle

\fp_zero:c

Sets the $\langle fp \ var \rangle$ to +0. \fp_gzero:N

\fp_gzero:c

Updated: 2012-05-08

\fp_zero_new:N

\fp_zero_new:N \langle fp var \rangle

\fp_gzero_new:N

\fp_zero_new:c

\fp_gzero_new:c

Ensures that the $\langle fp \ var \rangle$ exists globally by applying \fp_new:N if necessary, then applies \fp_(g)zero:N to leave the $\langle fp \ var \rangle$ set to +0.

Updated: 2012-05-08

Setting floating point variables 28.2

\fp_set:Nn

\fp_set:cn \fp_gset:Nn

\fp_gset:cn

 $fp_set:Nn \langle fp \ var \rangle \{\langle floating \ point \ expression \rangle\}$

Sets $\langle fp \ var \rangle$ equal to the result of computing the $\langle floating \ point \ expression \rangle$.

Updated: 2012-05-08

\fp_set_eq:NN

 $fp_set_eq:NN \langle fp \ var_1 \rangle \langle fp \ var_2 \rangle$

\fp_set_eq:(cN|Nc|cc) \fp_gset_eq:NN

 $\fp_gset_eq:(cN|Nc|cc)$

Sets the floating point variable $\langle fp \ var_1 \rangle$ equal to the current value of $\langle fp \ var_2 \rangle$.

Updated: 2012-05-08

\fp_add:Nn

\fp_add:cn

\fp_gadd:Nn

\fp_gadd:cn

\fp_add:Nn \langle fp var \rangle \langle floating point expression \rangle \rangle

Adds the result of computing the $\langle floating\ point\ expression \rangle$ to the $\langle fp\ var \rangle$. This also applies if $\langle fp \ var \rangle$ and $\langle floating \ point \ expression \rangle$ evaluate to tuples of the same size.

Updated: 2012-05-08

\fp_sub:Nn \fp_sub:cn \fp_gsub:Nn \fp_gsub:cn \fp_sub:Nn \langle fp var \rangle \langle floating point expression \rangle \rangle \rangle floating point expression \rangle \rangle \rangle floating point expression \rangle \rangle floating point expression \rangle \rangle floating floating point expression \rangle \rangle floating floating point expression \rangle \rangle floating floating

Subtracts the result of computing the $\langle floating\ point\ expression \rangle$ from the $\langle fp\ var \rangle$. This also applies if $\langle fp \ var \rangle$ and $\langle floating \ point \ expression \rangle$ evaluate to tuples of the same size.

Updated: 2012-05-08

28.3 Using floating points

\fp_eval:n

* \fp_eval:n {\langle floating point expression \rangle}

New: 2012-05-08 Evaluates the (floating point expression) and expresses the result as a decimal number Updated: 2012-07-08 with no exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Non-significant trailing zeros are trimmed, and integers are expressed without a decimal separator. The values $\pm \infty$ and nan trigger an "invalid operation" exception. For a tuple, each item is converted using \fp_eval:n and they are combined as $(\langle fp_1\rangle, \sqcup \langle fp_2\rangle, \sqcup \ldots \langle fp_n\rangle)$ if n>1 and $(\langle fp_1\rangle,)$ or () for fewer items. This function is identical to $\int p_to_decimal:n$.

 $fp_sign:n * fp_sign:n {\langle fpexpr \rangle}$

New: 2018-11-03 Evaluates the \(\langle fpexpr \rangle \) and leaves its sign in the input stream using \fp_eval:n $\{sign(\langle result \rangle)\}: +1 \text{ for positive numbers and for } +\infty, -1 \text{ for negative numbers and }$ for $-\infty$, ± 0 for ± 0 . If the operand is a tuple or is nan, then "invalid operation" occurs and the result is 0.

 $fp_{to_decimal:N * fp_{to_decimal:N } \langle fp \ var \rangle$

\fp_to_decimal:c * \fp_to_decimal:n {\langle floating point expression \rangle}

 $\underline{\text{Vfp_to_decimal:n}}$ Evaluates the $\langle floating\ point\ expression \rangle$ and expresses the result as a decimal number New: 2012-05-08 with no exponent. Leading or trailing zeros may be inserted to compensate for the Updated: 2012-07-08 exponent. Non-significant trailing zeros are trimmed, and integers are expressed without a decimal separator. The values $\pm \infty$ and nan trigger an "invalid operation" exception. For a tuple, each item is converted using \fp_to_decimal:n and they are combined as $(\langle fp_1\rangle, \cup \langle fp_2\rangle, \cup \ldots \langle fp_n\rangle)$ if n>1 and $(\langle fp_1\rangle,)$ or () for fewer items.

 $fp_{to_dim:N} * fp_{to_dim:N} \langle fp var \rangle$

\fp_to_dim:c

* \fp_to_dim:n {\langle floating point expression \rangle}

Updated: 2016-03-22 suitable for use in dimension expressions. The output is identical to \fp_to_decimal:n, with an additional trailing pt (both letter tokens). In particular, the result may be outside the range $[-2^{14} + 2^{-17}, 2^{14} - 2^{-17}]$ of valid T_EX dimensions, leading to overflow errors if used as a dimension. Tuples, as well as the values $\pm \infty$ and nan, trigger an "invalid operation" exception.

 $fp_{to_int:N} \star fp_{to_int:N} \langle fp \ var \rangle$

* \fp_to_int:n {\langle floating point expression \rangle}

 $\underline{\text{ \ \ }}$ Evaluates the $\langle floating\ point\ expression \rangle$, and rounds the result to the closest integer, Updated: 2012-07-08 rounding exact ties to an even integer. The result may be outside the range $[-2^{31} +$ $1, 2^{31} - 1$] of valid T_FX integers, leading to overflow errors if used in an integer expression. Tuples, as well as the values $\pm \infty$ and nan, trigger an "invalid operation" exception.

```
fp_to_scientific:N * fp_to_scientific:N < fp var
\fp_to_scientific:c * \fp_to_scientific:n {\langle floating point expression \rangle}
\overline{\text{hfp\_to\_scientific:n}} Evaluates the \langle floating\ point\ expression \rangle and expresses the result in scientific notation:
              New: 2012-05-08
                                         \langle optional - \rangle \langle digit \rangle. \langle 15 \ digits \ranglee\langle optional \ sign \rangle \langle exponent \rangle
         Updated: 2016-03-22
```

The leading $\langle digit \rangle$ is non-zero except in the case of ± 0 . The values $\pm \infty$ and nan trigger an "invalid operation" exception. Normal category codes apply: thus the e is category code 11 (a letter). For a tuple, each item is converted using \fp_to_scientific:n and they are combined as $(\langle fp_1 \rangle, \cup \langle fp_2 \rangle, \cup ... \langle fp_n \rangle)$ if n > 1 and $(\langle fp_1 \rangle,)$ or () for fewer items.

```
\fp_to_tl:N
                 * \fp_to_tl:N \langle fp var \rangle
                 * \fp_to_tl:n {\langle floating point expression \rangle}
\fp_to_tl:c
\fp_to_tl:n
```

 \star Evaluates the (floating point expression) and expresses the result in (almost) the shortest Updated: 2016-03-22 possible form. Numbers in the ranges $(0, 10^{-3})$ and $[10^{16}, \infty)$ are expressed in scientific notation with trailing zeros trimmed and no decimal separator when there is a single significant digit (this differs from \fp_to_scientific:n). Numbers in the range $[10^{-3}, 10^{16})$ are expressed in a decimal notation without exponent, with trailing zeros trimmed, and no decimal separator for integer values (see \fp to decimal:n. Negative numbers start with -. The special values ± 0 , $\pm \infty$ and nan are rendered as 0, -0, inf, -inf, and nan respectively. Normal category codes apply and thus inf or nan, if produced, are made up of letters. For a tuple, each item is converted using \fp_to_tl:n and they are combined as $(\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \ldots \langle fp_n \rangle)$ if n > 1 and $(\langle fp_1 \rangle,)$ or () for fewer items.

```
* \fp_use:N \langle fp var \rangle
\fp_use:N
\fp_use:c
```

 $\stackrel{\star}{=}$ Inserts the value of the $\langle fp \ var \rangle$ into the input stream as a decimal number with no Updated: 2012-07-08 exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Nonsignificant trailing zeros are trimmed. Integers are expressed without a decimal separator. The values $\pm \infty$ and nan trigger an "invalid operation" exception. For a tuple, each item is converted using $fp_to_decimal:n$ and they are combined as $(\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \ldots \langle fp_n \rangle)$ if n>1 and $(\langle fp_1\rangle_{\bullet})$ or () for fewer items. This function is identical to \fp_to_decimal:N.

28.4 Floating point conditionals

```
fp_if_exist_p:N \star fp_if_exist_p:N \langle fp var \rangle
\frac{\texttt{fp\_if\_exist:N}\underline{\mathit{TF}} \ \ \star}{\texttt{fp\_if\_exist:c}\underline{\mathit{TF}} \ \ \star} \text{ Tests whether the } \langle \mathit{fp\ var} \rangle \text{ is currently defined. This does not check that the } \langle \mathit{fp\ var} \rangle
     Updated: 2012-05-08
```

```
\folightharpunce \fol
fp_compare:nNnTF * fp_compare:nNnTF {\langle fpexpr_1 \rangle} \langle relation \rangle {\langle fpexpr_2 \rangle} {\langle true \ code \rangle} {\langle false \ code \rangle}
```

Updated: 2012-05-08 Compares the $\langle fpexpr_1 \rangle$ and the $\langle fpexpr_2 \rangle$, and returns true if the $\langle relation \rangle$ is obeyed. Two floating points x and y may obey four mutually exclusive relations: x < y, x = y, x > y, or x > y ("not ordered"). The last case occurs exactly if one or both operands is nan or is a tuple, unless they are equal tuples. Note that a nan is distinct from any value, even another nan, hence x = x is not true for a nan. To test if a value is nan, compare it to an arbitrary number with the "not ordered" relation.

```
\fp_compare:nNnTF { <value> } ? { 0 }
  \{\ \} % <value> is nan
  \{\ \} % <value> is not nan
```

Tuples are equal if they have the same number of items and items compare equal (in particular there must be no nan). At present any other comparison with tuples yields? (not ordered). This is experimental.

This function is less flexible than \fp_compare:nTF but slightly faster. It is provided for consistency with \int_compare:nNnTF and \dim_compare:nNnTF.

```
\fp_compare_p:n
                                        \fp_compare_p:n
\fp_compare:n<u>TF</u>
                                                  \langle \texttt{fpexpr}_1 \rangle \ \langle \texttt{relation}_1 \rangle
     Updated: 2013-12-14
                                                  \langle 	extsf{fpexpr}_N 
angle \, \langle 	extsf{relation}_N 
angle
                                                  \langle 	extit{fpexpr}_{N+1} 
angle
                                         \fp_compare:nTF
                                             {
                                                  \langle fpexpr_1 \rangle \langle relation_1 \rangle
                                                  \langle fpexpr_N \rangle \langle relation_N \rangle
                                                  \langle 	extit{fpexpr}_{N+1} 
angle
                                             }
                                             \{\langle true\ code \rangle\}\ \{\langle false\ code \rangle\}
```

Evaluates the (floating point expressions) as described for \fp_eval:n and compares consecutive result using the corresponding $\langle relation \rangle$, namely it compares $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$ using the $\langle relation_1 \rangle$, then $\langle fpexpr_2 \rangle$ and $\langle fpexpr_3 \rangle$ using the $\langle relation_2 \rangle$, until finally comparing $\langle fpexpr_N \rangle$ and $\langle fpexpr_{N+1} \rangle$ using the $\langle relation_N \rangle$. The test yields true if all comparisons are true. Each (floating point expression) is evaluated only once. Contrarily to \int_compare:nTF, all \(\)floating point expressions \(\) are computed, even if one comparison is false. Two floating points x and y may obey four mutually exclusive relations: x < y, x = y, x > y, or x?y ("not ordered"). The last case occurs exactly if one or both operands is nan or is a tuple, unless they are equal tuples. Each $\langle relation \rangle$ can be any (non-empty) combination of <, =, >, and ?, plus an optional leading! (which negates the $\langle relation \rangle$, with the restriction that the $\langle relation \rangle$ may not start with ?, as this symbol has a different meaning (in combination with:) within floating point expressions. The comparison $x \langle relation \rangle y$ is then true if the $\langle relation \rangle$ does not start with! and the actual relation ($\langle =, >, \text{ or } \rangle$) between x and y appears within the $\langle relation \rangle$, or on the contrary if the $\langle relation \rangle$ starts with ! and the relation between x and y does not appear within the $\langle relation \rangle$. Common choices of $\langle relation \rangle$ include >= (greater or equal), != (not equal), !? or \leftarrow (comparable).

This function is more flexible than \fp_compare:nNnTF and only slightly slower.

28.5Floating point expression loops

New: 2012-08-16 Places the $\langle code \rangle$ in the input stream for TEX to process, and then evaluates the relationship between the two (floating point expressions) as described for \fp_compare:nNnTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.

 $f_0 = 0$ while: $f_0 = 0$ hile: f_0

New: 2012-08-16 Places the $\langle code \rangle$ in the input stream for T_EX to process, and then evaluates the relationship between the two (floating point expressions) as described for \fp_compare:nNnTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.

\fp_until_do:nNnn 🌣	$\footnote{the posterior} $$ \int_{\mathbb{R}^n} {\langle fpexpr_1 \rangle} \ \langle relation \rangle \ \{\langle fpexpr_2 \rangle\} \ \{\langle code \rangle\} $$$
New: 2012-08-16	Evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for $fp_compare:nNnTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by T_EX the test is repeated, and a loop occurs until the test is true.
\fp_while_do:nNnn \six	$\fo_{\rm while_do:nNnn} \ \{\langle {\it fpexpr}_1 \rangle\} \ \langle {\it relation} \rangle \ \{\langle {\it fpexpr}_2 \rangle\} \ \{\langle {\it code} \rangle\}$
New: 2012-08-16	Evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for $fp_compare:nNnTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed by TEX the test is repeated, and a loop occurs until the test is false.
\fp_do_until:nn ☆	lem:lem:lem:lem:lem:lem:lem:lem:lem:lem:
	Places the $\langle code \rangle$ in the input stream for T _E X to process, and then evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for \fp_compare:nTF. If the test is false then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is true.
\fp_do_while:nn ☆	$\label{lem:norm} $$ \P_0_{\min} $
	Places the $\langle code \rangle$ in the input stream for TeX to process, and then evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for \fp_compare:nTF. If the test is true then the $\langle code \rangle$ is inserted into the input stream again and a loop occurs until the $\langle relation \rangle$ is false.
\fp_until_do:nn ☆	lem:lem:lem:lem:lem:lem:lem:lem:lem:lem:
	Evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for $fp_compare:nTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is false. After the $\langle code \rangle$ has been processed by T_EX the test is repeated, and a loop occurs until the test is true.
\fp_while_do:nn ☆	lem:lem:lem:lem:lem:lem:lem:lem:lem:lem:
	Evaluates the relationship between the two $\langle floating\ point\ expressions \rangle$ as described for $fp_compare:nTF$, and then places the $\langle code \rangle$ in the input stream if the $\langle relation \rangle$ is true. After the $\langle code \rangle$ has been processed by TEX the test is repeated, and a loop occurs until the test is false.

```
\fp_step_function:nnnN
\fp_step_function:nnnc
```

 $\label{lem:nnnn} $$ \left(initial\ value \right) $$ \left(\left(step \right) \right) $$$

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, each of which should be a floating point expression evaluating to a floating point number, not a tuple. Updated: 2016-12-06 The $\langle function \rangle$ is then placed in front of each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ $value\rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$). The $\langle step \rangle$ must be non-zero. If the $\langle step \rangle$ is positive, the loop stops when the $\langle value \rangle$ becomes larger than the $\langle final\ value \rangle$. If the $\langle step \rangle$ is negative, the loop stops when the $\langle value \rangle$ becomes smaller than the $\langle final \rangle$ value). The $\langle function \rangle$ should absorb one numerical argument. For example

```
\cs_set:Npn \my_func:n #1 { [I~saw~#1] \quad }
\fp_step_function:nnnN { 1.0 } { 0.1 } { 1.5 } \my_func:n
```

would print

```
[I saw 1.1]
              [I saw 1.2]
                             [I saw 1.3]
                                           [I saw 1.4]
                                                          [I saw 1.5]
```

TeXhackers note: Due to rounding, it may happen that adding the $\langle step \rangle$ to the $\langle value \rangle$ does not change the (value); such cases give an error, as they would otherwise lead to an infinite loop.

 $f_{\text{ode}} = \inf \{\langle code \}\}$

New: 2016-11-21 This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which Updated: 2016-12-06 should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with #1 replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (#1).

\fp_step_variable:nnnNn \fp_step_variable:nnnNn

New: 2017-04-12

 $\{\langle \mathtt{initial\ value} \rangle\}\ \{\langle \mathtt{step} \rangle\}\ \{\langle \mathtt{final\ value} \rangle\}\ \langle \mathtt{t1\ var} \rangle\ \{\langle \mathtt{code} \rangle\}$

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$, the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl \ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl \ var \rangle$.

28.6 Some useful constants, and scratch variables

\c_zero_fp \c_minus_zero_fp Zero, with either sign.

New: 2012-05-08

\c_one_fp

One as an fp: useful for comparisons in some places.

New: 2012-05-08

\c.	_inf_fp)	
\c.	_minus_	_inf	_fp

Infinity, with either sign. These can be input directly in a floating point expression as inf and -inf.

New: 2012-05-08

 c_e_fp

The value of the base of the natural logarithm, $e = \exp(1)$.

Updated: 2012-05-08

\c_pi_fp

The value of π . This can be input directly in a floating point expression as pi.

Updated: 2013-11-17

\c_one_degree_fp The value of 1° in radians. Multiply an angle given in degrees by this value to obtain a $_{\text{New: }2012-05-08}$ result in radians. Note that trigonometric functions expecting an argument in radians or Updated: 2013-11-17 in degrees are both available. Within floating point expressions, this can be accessed as deg.

\l_tmpa_fp Scratch floating points for local assignment. These are never used by the kernel code, and \l_tmpb_fp so are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

\g_tmpa_fp Scratch floating points for global assignment. These are never used by the kernel code, \g_tmpb_fp and so are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

28.7Floating point exceptions

The functions defined in this section are experimental, and their functionality may be altered or removed altogether.

"Exceptions" may occur when performing some floating point operations, such as 0 / 0, or 10 ** 1e9999. The relevant IEEE standard defines 5 types of exceptions, of which we implement 4.

- Overflow occurs whenever the result of an operation is too large to be represented as a normal floating point number. This results in $\pm \infty$.
- Underflow occurs whenever the result of an operation is too close to 0 to be represented as a normal floating point number. This results in ± 0 .
- Invalid operation occurs for operations with no defined outcome, for instance 0/0 or $\sin(\infty)$, and results in a nan. It also occurs for conversion functions whose target type does not have the appropriate infinite or nan value $(e.g., f_{\text{to_dim:n}})$.
- Division by zero occurs when dividing a non-zero number by 0, or when evaluating functions at poles, e.g., $\ln(0)$ or $\cot(0)$. This results in $\pm \infty$.

(not yet) Inexact occurs whenever the result of a computation is not exact, in other words, almost always. At the moment, this exception is entirely ignored in LATEX3.

To each exception we associate a "flag": fp_overflow, fp_underflow, fp_invalid_operation and fp_division_by_zero. The state of these flags can be tested and modified with commands from I3flag

By default, the "invalid operation" exception triggers an (expandable) error, and raises the corresponding flag. Other exceptions raise the corresponding flag but do not trigger an error. The behaviour when an exception occurs can be modified (using \fp_trap:nn) to either produce an error and raise the flag, or only raise the flag, or do nothing at all.

\fp_trap:nn

 $\footnote{The points of the property of the$

New: 2012-07-19 All occurrences of the $\langle exception \rangle$ (overflow, underflow, invalid_operation or Updated: 2017-02-13 division_by_zero) within the current group are treated as $\langle trap\ type \rangle$, which can be

- none: the \(\langle exception\rangle\) will be entirely ignored, and leave no trace;
- flag: the \(\langle exception \rangle \) will turn the corresponding flag on when it occurs;
- error: additionally, the *(exception)* will halt the TEX run and display some information about the current operation in the terminal.

This function is experimental, and may be altered or removed.

flag_fp_overflow flag_fp_underflow flag_fp_invalid_operation flag_fp_division_by_zero Flags denoting the occurrence of various floating-point exceptions.

28.8 Viewing floating points

 $fp_show:N \qquad fp_show:N \ \langle fp \ var \rangle$

\fp_show:c \fp_show:n {\(floating point expression \) }

\fp_show:n Evaluates the \(\forall floating point expression\) and displays the result in the terminal.

New: 2012-05-08 Updated: 2021-04-29

 $\fp_log:N \qquad \fp_log:N \qquad \fp_log:N \qquad \fp_log:N \qquad \$

\fp_log:c \fp_log:n {\(floating point expression \) }

\fp_log:n Evaluates the \(\(\floating point expression \)\) and writes the result in the log file.

New: 2014-08-22 Updated: 2021-04-29

28.9 Floating point expressions

28.9.1 Input of floating point numbers

We support four types of floating point numbers:

- $\pm m \cdot 10^n$, a floating point number, with integer $1 \le m \le 10^{16}$, and $-10000 \le n \le 10000$;
- ± 0 , zero, with a given sign;
- $\pm \infty$, infinity, with a given sign;
- nan, is "not a number", and can be either quiet or signalling (not yet: this distinction is currently unsupported);

Normal floating point numbers are stored in base 10, with up to 16 significant figures. On input, a normal floating point number consists of:

- $\langle sign \rangle$: a possibly empty string of + and characters;
- (significand): a non-empty string of digits together with zero or one dot;
- $\langle exponent \rangle$ optionally: the character **e** or **E**, followed by a possibly empty string of + and tokens, and a non-empty string of digits.

The sign of the resulting number is + if $\langle sign \rangle$ contains an even number of -, and - otherwise, hence, an empty $\langle sign \rangle$ denotes a non-negative input. The stored significand is obtained from $\langle significand \rangle$ by omitting the decimal separator and leading zeros, and rounding to 16 significant digits, filling with trailing zeros if necessary. In particular, the value stored is exact if the input $\langle significand \rangle$ has at most 16 digits. The stored $\langle exponent \rangle$ is obtained by combining the input $\langle exponent \rangle$ (0 if absent) with a shift depending on the position of the significand and the number of leading zeros.

A special case arises if the resulting $\langle exponent \rangle$ is either too large or too small for the floating point number to be represented. This results either in an overflow (the number is then replaced by $\pm \infty$), or an underflow (resulting in ± 0).

The result is thus ± 0 if and only if $\langle significand \rangle$ contains no non-zero digit (*i.e.*, consists only in characters 0, and an optional period), or if there is an underflow. Note that a single dot is currently a valid floating point number, equal to +0, but that is not guaranteed to remain true.

The $\langle significand \rangle$ must be non-empty, so e1 and e-1 are not valid floating point numbers. Note that the latter could be mistaken with the difference of "e" and 1. To avoid confusions, the base of natural logarithms cannot be input as e and should be input as exp(1) or \c_e_fp (which is faster).

Special numbers are input as follows:

- inf represents $+\infty$, and can be preceded by any $\langle sign \rangle$, yielding $\pm \infty$ as appropriate.
- nan represents a (quiet) non-number. It can be preceded by any sign, but that sign is ignored.
- Any unrecognizable string triggers an error, and produces a nan.
- Note that commands such as \infty, \pi, or \sin do not work in floating point expressions. They may silently be interpreted as completely unexpected numbers, because integer constants (allowed in expressions) are commonly stored as mathematical characters.

28.9.2 Precedence of operators

We list here all the operations supported in floating point expressions, in order of decreasing precedence: operations listed earlier bind more tightly than operations listed below them.

- Function calls (sin, ln, etc).
- Binary ** and ^ (right associative).
- Unary +, -, !.
- Implicit multiplication by juxtaposition (2pi) when neither factor is in parentheses.
- Binary * and /, implicit multiplication by juxtaposition with parentheses (for instance 3(4+5)).
- Binary + and -.
- Comparisons >=, !=, <?, etc.
- Logical and, denoted by &&.
- Logical or, denoted by ||.
- Ternary operator ?: (right associative).
- Comma (to build tuples).

The precedence of operations can be overridden using parentheses. In particular, the precedence of juxtaposition implies that

$$\begin{aligned} 1/2 \text{pi} &= 1/(2\pi), \\ 1/2 \text{pi}(\text{pi} + \text{pi}) &= (2\pi)^{-1}(\pi + \pi) \simeq 1, \\ \sin &2 \text{pi} &= \sin(2)\pi \neq 0, \\ 2^2 \text{max}(3,5) &= 2^2 \max(3,5) = 20, \\ 1 \text{in}/1 \text{cm} &= (1 \text{in})/(1 \text{cm}) = 2.54. \end{aligned}$$

Functions are called on the value of their argument, contrarily to TFX macros.

28.9.3 Operations

We now present the various operations allowed in floating point expressions, from the lowest precedence to the highest. When used as a truth value, a floating point expression is false if it is ± 0 , and true otherwise, including when it is nan or a tuple such as (0,0). Tuples are only supported to some extent by operations that work with truth values (?:, | |, &&, !), by comparisons (!<=>?), and by +, -, *, /. Unless otherwise specified, providing a tuple as an argument of any other operation yields the "invalid operation" exception and a nan result.

```
?: fp_eval:n \{ \langle operand_1 \rangle ? \langle operand_2 \rangle : \langle operand_3 \rangle \}
```

The ternary operator ?: results in $\langle operand_2 \rangle$ if $\langle operand_1 \rangle$ is true (not ± 0), and $\langle operand_3 \rangle$ if $\langle operand_1 \rangle$ is false (± 0) . All three $\langle operand_3 \rangle$ are evaluated in all cases; they may be tuples. The operator is right associative, hence

```
\fp_eval:n
{
    1 + 3 > 4 ? 1 :
    2 + 4 > 5 ? 2 :
    3 + 5 > 6 ? 3 : 4
}
```

first tests whether 1+3>4; since this isn't true, the branch following: is taken, and 2+4>5 is compared; since this is true, the branch before: is taken, and everything else is (evaluated then) ignored. That allows testing for various cases in a concise manner, with the drawback that all computations are made in all cases.

```
| | fp_eval:n \{ \langle operand_1 \rangle | \langle operand_2 \rangle \}
```

If $\langle operand_1 \rangle$ is true (not ± 0), use that value, otherwise the value of $\langle operand_2 \rangle$. Both $\langle operands \rangle$ are evaluated in all cases; they may be tuples. In $\langle operand_1 \rangle \mid \mid \langle operand_2 \rangle \mid \mid \dots \mid \mid \langle operands_n \rangle$, the first true (nonzero) $\langle operand \rangle$ is used and if all are zero the last one (± 0) is used.

```
&& fp_eval:n \{ \langle operand_1 \rangle \&\& \langle operand_2 \rangle \}
```

If $\langle operand_1 \rangle$ is false (equal to ± 0), use that value, otherwise the value of $\langle operand_2 \rangle$. Both $\langle operands \rangle$ are evaluated in all cases; they may be tuples. In $\langle operand_1 \rangle$ && $\langle operand_2 \rangle$ && ... && $\langle operands_n \rangle$, the first false (± 0) $\langle operand \rangle$ is used and if none is zero the last one is used.

Each $\langle relation \rangle$ consists of a non-empty string of $\langle , =, \rangle$, and ?, optionally preceded by !, and may not start with ?. This evaluates to +1 if all comparisons $\langle operand_i \rangle$ $\langle relation_i \rangle$ $\langle operand_{i+1} \rangle$ are true, and +0 otherwise. All $\langle operands \rangle$ are evaluated (once) in all cases. See $fp_compare:nTF$ for details.

```
+ \fp_eval:n { \langle operand_1 \rangle + \langle operand_2 \rangle }
- \fp_eval:n { \langle operand_1 \rangle - \langle operand_2 \rangle }
```

Computes the sum or the difference of its two $\langle operands \rangle$. The "invalid operation" exception occurs for $\infty - \infty$. "Underflow" and "overflow" occur when appropriate. These operations supports the itemwise addition or subtraction of two tuples, but if they have a different number of items the "invalid operation" exception occurs and the result is nan.

```
* \fp_eval:n { \langle operand_1 \rangle * \langle operand_2 \rangle } / \fp_eval:n { \langle operand_1 \rangle / \langle operand_2 \rangle }
```

Computes the product or the ratio of its two $\langle operands \rangle$. The "invalid operation" exception occurs for ∞/∞ , 0/0, or $0 * \infty$. "Division by zero" occurs when dividing a finite non-zero number by ± 0 . "Underflow" and "overflow" occur when appropriate. When $\langle operand_1 \rangle$ is a tuple and $\langle operand_2 \rangle$ is a floating point number, each item of $\langle operand_1 \rangle$ is multiplied or divided by $\langle operand_2 \rangle$. Multiplication also supports the case where $\langle operand_1 \rangle$ is a floating point number and $\langle operand_2 \rangle$ a tuple. Other combinations yield an "invalid operation" exception and a nan result.

```
+ \fp_eval:n { + \langle operand \rangle }
- \fp_eval:n { - \langle operand \rangle }
! \fp_eval:n { ! \langle operand \rangle }
```

The unary + does nothing, the unary - changes the sign of the $\langle operand \rangle$ (for a tuple, of all its components), and ! $\langle operand \rangle$ evaluates to 1 if $\langle operand \rangle$ is false (is ± 0) and 0 otherwise (this is the not boolean function). Those operations never raise exceptions.

```
** \fp_eval:n { \langle operand_1 \rangle ** \langle operand_2 \rangle } 
^ \fp_eval:n { \langle operand_1 \rangle ^ \langle operand_2 \rangle }
```

Raises $\langle operand_1 \rangle$ to the power $\langle operand_2 \rangle$. This operation is right associative, hence 2 ** 2 ** 3 equals $2^{2^3} = 256$. If $\langle operand_1 \rangle$ is negative or -0 then: the result's sign is + if the $\langle operand_2 \rangle$ is infinite and $(-1)^p$ if the $\langle operand_2 \rangle$ is $p/5^q$ with p, q integers; the result is +0 if abs $(\langle operand_1 \rangle) \hat{\langle} operand_2 \rangle$ evaluates to zero; in other cases the "invalid operation" exception occurs because the sign cannot be determined. "Division by zero" occurs when raising ± 0 to a finite strictly negative power. "Underflow" and "overflow" occur when appropriate. If either operand is a tuple, "invalid operation" occurs.

```
abs \fp_eval:n { abs( \langle fpexpr \rangle ) }
```

Computes the absolute value of the $\langle fpexpr \rangle$. If the operand is a tuple, "invalid operation" occurs. This operation does not raise exceptions in other cases. See also $fp_abs:n$.

```
exp \fp_eval:n { exp( \langle fpexpr \rangle ) }
```

Computes the exponential of the $\langle fpexpr \rangle$. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
fact \fp_eval:n { fact( \langle fpexpr \rangle ) }
```

Computes the factorial of the $\langle fpexpr \rangle$. If the $\langle fpexpr \rangle$ is an integer between -0 and 3248 included, the result is finite and correctly rounded. Larger positive integers give $+\infty$ with "overflow", while fact($+\infty$) = $+\infty$ and fact(nan) = nan with no exception. All other inputs give nan with the "invalid operation" exception.

```
ln \leq ln ( \langle fpexpr \rangle ) }
```

Computes the natural logarithm of the $\langle fpexpr \rangle$. Negative numbers have no (real) logarithm, hence the "invalid operation" is raised in that case, including for $\ln(-0)$. "Division by zero" occurs when evaluating $\ln(+0) = -\infty$. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
* \fp_eval:n { logb( \langle fpexpr \rangle ) }
logb
```

New: 2018-11-03 Determines the exponent of the $\langle fpexpr \rangle$, namely the floor of the base-10 logarithm of its absolute value. "Division by zero" occurs when evaluating $logb(\pm 0) = -\infty$. Other special values are $logb(\pm \infty) = +\infty$ and logb(nan) = nan. If the operand is a tuple or is nan, then "invalid operation" occurs and the result is nan.

```
max \fp_eval:n { max( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle , ... ) }
min \fp_eval:n { min( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle , ... ) }
```

Evaluates each $\langle fpexpr \rangle$ and computes the largest (smallest) of those. If any of the $\langle fpexpr \rangle$ is a nan or tuple, the result is nan. If any operand is a tuple, "invalid operation" occurs; these operations do not raise exceptions in other cases.

```
fp_eval:n { round ( \langle fpexpr \rangle ) }
round
trunc
ceil
floor
```

```
\fp_eval:n { round ( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle ) }
\fp_eval:n { round ( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle , \langle fpexpr_3 \rangle ) }
```

Only round accepts a third argument. Evaluates $\langle fpexpr_1 \rangle = x$ and $\langle fpexpr_2 \rangle = n$ and New: 2013-12-14 $\langle fpexpr_3 \rangle = t$ then rounds x to n places. If n is an integer, this rounds x to a multiple Updated: 2015-08-08 of 10^{-n} ; if $n=+\infty$, this always yields x; if $n=-\infty$, this yields one of $\pm 0, \pm \infty$, or nan; if n = nan, this yields nan; if n is neither $\pm \infty$ nor an integer, then an "invalid operation" exception is raised. When $\langle fpexpr_2 \rangle$ is omitted, n = 0, i.e., $\langle fpexpr_1 \rangle$ is rounded to an integer. The rounding direction depends on the function.

- round yields the multiple of 10^{-n} closest to x, with ties (x half-way between two such multiples) rounded as follows. If t is nan (or not given) the even multiple is chosen ("ties to even"), if $t = \pm 0$ the multiple closest to 0 is chosen ("ties to zero"), if t is positive/negative the multiple closest to $\infty/-\infty$ is chosen ("ties towards positive/negative infinity").
- floor yields the largest multiple of 10^{-n} smaller or equal to x ("round towards negative infinity");
- ceil yields the smallest multiple of 10^{-n} greater or equal to x ("round towards positive infinity");
- trunc yields a multiple of 10^{-n} with the same sign as x and with the largest absolute value less than that of x ("round towards zero").

"Overflow" occurs if x is finite and the result is infinite (this can only happen if $\langle fpexpr_2 \rangle <$ -9984). If any operand is a tuple, "invalid operation" occurs.

```
sign \fp_eval:n { sign( \langle fpexpr \rangle ) }
```

Evaluates the $\langle fpexpr \rangle$ and determines its sign: +1 for positive numbers and for $+\infty$, -1 for negative numbers and for $-\infty$, ± 0 for ± 0 , and nan for nan. If the operand is a tuple, "invalid operation" occurs. This operation does not raise exceptions in other cases.

```
fp_eval:n { sin( \langle fpexpr \rangle ) }
sin
cos
                       \fp_eval:n { cos( \langle fpexpr \rangle }
tan
                       fp_eval:n { tan( \langle fpexpr 
angle ) }
cot
                      fp_eval:n { cot( \langle fpexpr \rangle ) }
                      fp_eval:n { csc( \langle fpexpr \rangle ) }
csc
                      fp_eval:n { sec( \langle fpexpr \rangle ) }
sec
```

Updated: 2013-11-17 Computes the sine, cosine, tangent, cotangent, cosecant, or secant of the $\langle fpexpr \rangle$ given in radians. For arguments given in degrees, see sind, cosd, etc. Note that since π is irrational, $\sin(8pi)$ is not quite zero, while its analogue $\sin d(8 \times 180)$ is exactly zero. The trigonometric functions are undefined for an argument of $\pm \infty$, leading to the "invalid operation" exception. Additionally, evaluating tangent, cotangent, cosecant, or secant at one of their poles leads to a "division by zero" exception. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
sind
                  \fp_eval:n { sind( \langle fpexpr \rangle ) }
cosd
                  \fp_eval:n { cosd( \langle fpexpr \rangle ) }
                  \fp_eval:n { tand( \langle fpexpr \rangle ) }
tand
cotd
                  fp_eval:n { cotd( \langle fpexpr \rangle ) }
                  fp_eval:n { cscd( \langle fpexpr \rangle ) }
cscd
                  fp_eval:n { secd( \langle fpexpr \rangle ) }
secd
```

New: 2013-11-02 Computes the sine, cosine, tangent, cotangent, cosecant, or secant of the $\langle fpexpr \rangle$ given in degrees. For arguments given in radians, see sin, cos, etc. Note that since π is irrational, $\sin(8pi)$ is not quite zero, while its analogue $\sin d(8 \times 180)$ is exactly zero. The trigonometric functions are undefined for an argument of $\pm \infty$, leading to the "invalid operation" exception. Additionally, evaluating tangent, cotangent, cosecant, or secant at one of their poles leads to a "division by zero" exception. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
asin
                  fp_eval:n { asin( \langle fpexpr \rangle ) }
acos
                  fp_eval:n { acos( \langle fpexpr \rangle ) }
                  fp_eval:n { acsc( \langle fpexpr \rangle ) }
acsc
asec
                  \fp_eval:n { asec( \langle fpexpr \rangle ) }
```

New: 2013-11-02 Computes the arcsine, arccosine, arccosecant, or arcsecant of the $\langle fpexpr \rangle$ and returns the result in radians, in the range $[-\pi/2, \pi/2]$ for asin and acsc and $[0, \pi]$ for acos and asec. For a result in degrees, use asind, etc. If the argument of asin or acos lies outside the range [-1,1], or the argument of acsc or asec inside the range (-1,1), an "invalid operation" exception is raised. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
asind
                 fp_eval:n { asind( \langle fpexpr \rangle ) }
acosd
                 \fp_eval:n { acosd( \langle fpexpr \rangle ) }
                 fp_eval:n { acscd( \langle fpexpr \rangle ) }
acscd
                 fp_eval:n { asecd( \langle fpexpr \rangle ) }
asecd
```

New: 2013-11-02 Computes the arcsine, arccosine, arccosecant, or arcsecant of the $\langle fpexpr \rangle$ and returns the result in degrees, in the range [-90, 90] for asin and acsc and [0, 180] for acos and asec. For a result in radians, use asin, etc. If the argument of asin or acos lies outside the range [-1,1], or the argument of accc or asec inside the range (-1,1), an "invalid operation" exception is raised. "Underflow" and "overflow" occur when appropriate. If the operand is a tuple, "invalid operation" occurs.

```
atan \fp_eval:n { atan( \langle fpexpr \rangle ) } acot \fp_eval:n { atan( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle ) } \fp_eval:n { acot( \langle fpexpr \rangle ) } \fp_eval:n { acot( \langle fpexpr \rangle ) }
```

Those functions yield an angle in radians: at and acotd are their analogs in degrees. The one-argument versions compute the arctangent or arccotangent of the $\langle fpexpr \rangle$: arctangent takes values in the range $[-\pi/2,\pi/2]$, and arccotangent in the range $[0,\pi]$. The two-argument arctangent computes the angle in polar coordinates of the point with Cartesian coordinates $(\langle fpexpr_2 \rangle, \langle fpexpr_1 \rangle)$: this is the arctangent of $\langle fpexpr_1 \rangle/\langle fpexpr_2 \rangle$, possibly shifted by π depending on the signs of $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$. The two-argument arccotangent computes the angle in polar coordinates of the point $(\langle fpexpr_1 \rangle, \langle fpexpr_2 \rangle)$, equal to the arccotangent of $\langle fpexpr_1 \rangle/\langle fpexpr_2 \rangle$, possibly shifted by π . Both two-argument functions take values in the wider range $[-\pi,\pi]$. The ratio $\langle fpexpr_1 \rangle/\langle fpexpr_2 \rangle$ need not be defined for the two-argument arctangent: when both expressions yield ± 0 , or when both yield $\pm \infty$, the resulting angle is one of $\{\pm \pi/4, \pm 3\pi/4\}$ depending on signs. The "underflow" exception can occur. If any operand is a tuple, "invalid operation" occurs.

```
atand \fp_eval:n { atand( \langle fpexpr \rangle ) } acotd \fp_eval:n { atand( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle ) } \fp_eval:n { acotd( \langle fpexpr \rangle ) } \fp_eval:n { acotd( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle ) }
```

Those functions yield an angle in degrees: at and acotd are their analogs in radians. The one-argument versions compute the arctangent or arccotangent of the $\langle fpexpr \rangle$: arctangent takes values in the range [-90,90], and arccotangent in the range [0,180]. The two-argument arctangent computes the angle in polar coordinates of the point with Cartesian coordinates $(\langle fpexpr_2 \rangle, \langle fpexpr_1 \rangle)$: this is the arctangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by 180 depending on the signs of $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$. The two-argument arccotangent computes the angle in polar coordinates of the point $(\langle fpexpr_1 \rangle, \langle fpexpr_2 \rangle)$, equal to the arccotangent of $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$, possibly shifted by 180. Both two-argument functions take values in the wider range [-180,180]. The ratio $\langle fpexpr_1 \rangle / \langle fpexpr_2 \rangle$ need not be defined for the two-argument arctangent: when both expressions yield ± 0 , or when both yield $\pm \infty$, the resulting angle is one of $\{\pm 45, \pm 135\}$ depending on signs. The "underflow" exception can occur. If any operand is a tuple, "invalid operation" occurs.

```
sqrt \fp_eval:n { sqrt( \langle fpexpr \rangle ) }
```

New: 2013-12-14 Computes the square root of the $\langle fpexpr \rangle$. The "invalid operation" is raised when the $\langle fpexpr \rangle$ is negative or is a tuple; no other exception can occur. Special values yield $\sqrt{-0} = -0, \ \sqrt{+0} = +0, \ \sqrt{+\infty} = +\infty$ and $\sqrt{\text{nan}} = \text{nan}$.

rand

\fp_eval:n { rand() }

New: 2016-12-05 Produces a pseudo-random floating-point number (multiple of 10^{-16}) between 0 included and 1 excluded. This is not available in older versions of X-TFX. The random seed can be queried using \sys_rand_seed: and set using \sys_gset_rand_seed:n.

> TeXhackers note: This is based on pseudo-random numbers provided by the engine's primitive \pdfuniformdeviate in pdfTFX, pTFX, upTFX and \uniformdeviate in LuaTFX and X_HT_EX. The underlying code is based on Metapost, which follows an additive scheme recommended in Section 3.6 of "The Art of Computer Programming, Volume 2".

> While we are more careful than \uniformdeviate to preserve uniformity of the underlying stream of 28-bit pseudo-random integers, these pseudo-random numbers should of course not be relied upon for serious numerical computations nor cryptography.

randint New: 2016-12-05

```
\fp_eval:n { randint( \langle fpexpr \rangle ) }
\fp_eval:n { randint( \langle fpexpr_1 \rangle , \langle fpexpr_2 \rangle ) }
```

Produces a pseudo-random integer between 1 and $\langle fpexpr \rangle$ or between $\langle fpexpr_1 \rangle$ and $\langle fpexpr_2 \rangle$ inclusive. The bounds must be integers in the range $(-10^{16}, 10^{16})$ and the first must be smaller or equal to the second. See rand for important comments on how these pseudo-random numbers are generated.

inf The special values $+\infty$, $-\infty$, and nan are represented as inf, -inf and nan (see $\c_ ^{\tt nan} \ \, {\tt inf_fp}, \verb|\c_minus_inf_fp| \ \, {\tt and} \ \, \verb|\c_nan_fp|.$

Pi The value of π (see \c_pi_fp).

deg The value of 1° in radians (see \c_one_degree_fp).

Those units of measurement are equal to their values in pt, namely

ex in
$$1 \text{ in} = 72.27 \, \text{pt}$$
 $1 \text{pt} = 1 \, \text{pt}$ $1 \text{pt} = 1 \, \text{pt}$ $1 \text{pc} = 12 \, \text{pt}$ 1mm $1 \text{cm} = \frac{1}{2.54} \, \text{in} = 28.45275590551181 \, \text{pt}$ 1cc $1 \text{mm} = \frac{1}{25.4} \, \text{in} = 2.845275590551181 \, \text{pt}$ 1mc $1 \text{dd} = 0.376065 \, \text{mm} = 1.07000856496063 \, \text{pt}$ $1 \text{cc} = 12 \, \text{dd} = 12.84010277952756 \, \text{pt}$ $1 \text{nd} = 0.375 \, \text{mm} = 1.066978346456693 \, \text{pt}$ $1 \text{nc} = 12 \, \text{nd} = 12.80374015748031 \, \text{pt}$ $1 \text{pp} = \frac{1}{72} \, \text{in} = 1.00375 \, \text{pt}$ $1 \text{pp} = 2^{-16} \, \text{pt} = 1.52587890625 \times 10^{-5} \, \text{pt}$.

The values of the (font-dependent) units em and ex are gathered from TFX when the surrounding floating point expression is evaluated.

true Other names for 1 and +0. false

```
* \fp_abs:n {\langle floating point expression \rangle}
```

em

New: 2012-05-14 Evaluates the \(\frac{floating point expression} \) as described for \frac{fp_eval:n}{} and leaves the Updated: 2012-07-08 absolute value of the result in the input stream. If the argument is $\pm \infty$, nan or a tuple, "invalid operation" occurs. Within floating point expressions, abs() can be used; it accepts $\pm \infty$ and nan as arguments.

```
fp_{max:nn * fp_{max:nn } {\langle fp \ expression \ 1 \rangle} {\langle fp \ expression \ 2 \rangle}
```

 $\frac{\text{fp_min:nn} \star}{\text{Evaluates the } \langle \text{floating point expressions} \rangle}$ as described for $\frac{\text{fp_eval:n}}{\text{fp_min:nn}}$ New: 2012-09-26 resulting larger (max) or smaller (min) value in the input stream. If the argument is a tuple, "invalid operation" occurs, but no other case raises exceptions. Within floating point expressions, max() and min() can be used.

28.10 Disclaimer and roadmap

The package may break down if the escape character is among 0123456789_+, or if it receives a TEX primitive conditional affected by \exp_not:N.

The following need to be done. I'll try to time-order the items.

- Function to count items in a tuple (and to determine if something is a tuple).
- Decide what exponent range to consider.

- Support signalling nan.
- Modulo and remainder, and rounding function quantize (and its friends analogous to trunc, ceil, floor).
- \fp_format:nn {\langle fpexpr\rangle} \{\langle format\rangle}, but what should \langle format\rangle be? More general pretty printing?
- Add and, or, xor? Perhaps under the names all, any, and xor?
- Add log(x, b) for logarithm of x in base b.
- hypot (Euclidean length). Cartesian-to-polar transform.
- Hyperbolic functions cosh, sinh, tanh.
- Inverse hyperbolics.
- Base conversion, input such as OxAB.CDEF.
- Factorial (not with !), gamma function.
- Improve coefficients of the sin and tan series.
- Treat upper and lower case letters identically in identifiers, and ignore underscores.
- Add an array(1,2,3) and i=complex(0,1).
- Provide an experimental map function? Perhaps easier to implement if it is a single character, @sin(1,2)?
- Provide an isnan function analogue of \fp_if_nan:nTF?
- Support keyword arguments?

Pgfmath also provides box-measurements (depth, height, width), but boxes are not possible expandably.

Bugs, and tests to add.

- Check that functions are monotonic when they should.
- Add exceptions to ?:, !<=>?, &&, ||, and !.
- Logarithms of numbers very close to 1 are inaccurate.
- When rounding towards $-\infty$, \dim_to_fp:n {0pt} should return -0, not +0.
- The result of $(\pm 0) + (\pm 0)$, of x + (-x), and of (-x) + x should depend on the rounding mode.
- Subnormals are not implemented.

Possible optimizations/improvements.

- Document that ||3trial/||3fp-types introduces tools for adding new types.
- In subsection 28.9.1, write a grammar.

- It would be nice if the parse auxiliaries for each operation were set up in the corresponding module, rather than centralizing in l3fp-parse.
- Some functions should get an _o ending to indicate that they expand after their result
- More care should be given to distinguish expandable/restricted expandable (auxiliary and internal) functions.
- The code for the ternary set of functions is ugly.
- There are many ~ missing in the doc to avoid bad line-breaks.
- The algorithm for computing the logarithm of the significand could be made to use a 5 terms Taylor series instead of 10 terms by taking $c = 2000/(\lfloor 200x \rfloor + 1) \in [10, 95]$ instead of $c \in [1, 10]$. Also, it would then be possible to simplify the computation of t. However, we would then have to hard-code the logarithms of 44 small integers instead of 9.
- Improve notations in the explanations of the division algorithm (I3fp-basics).
- Understand and document __fp_basics_pack_weird_low:NNNNw and __fp_basics_pack_weird_high:NNNNNNNw better. Move the other basics_pack auxiliaries to l3fp-aux under a better name.
- Find out if underflow can really occur for trigonometric functions, and redoc as appropriate.
- Add bibliography. Some of Kahan's articles, some previous TEX fp packages, the international standards,...
- Also take into account the "inexact" exception?
- Support multi-character prefix operators (e.g., @/ or whatever)?

Chapter 29

The **I3fparray** package: Fast global floating point arrays

29.1**13fparray** documentation

For applications requiring heavy use of floating points, this module provides arrays which can be accessed in constant time (contrast 13seq, where access time is linear). The interface is very close to that of l3intarray. The size of the array is fixed and must be given at point of initialisation

 $\sigma. \$ New: 2018-05-05 Evaluates the integer expression $\langle size \rangle$ and allocates an $\langle floating\ point\ array\ variable \rangle$ with that number of (zero) entries. The variable name should start with \g_ because assignments are always global. \fparray_count:N * \fparray_count:N \langle fparray var \rangle New: 2018-05-05 Expands to the number of entries in the (floating point array variable). This is performed in constant time. $\footnote{Mon parray_gset:Nnn $$ {parray var} {\langle position \rangle} {\langle value \rangle} }$ New: 2018-05-05 Stores the result of evaluating the floating point expression $\langle value \rangle$ into the $\langle floating\ point$ array variable at the (integer expression) $\langle position \rangle$. If the $\langle position \rangle$ is not between 1 and the \fparray_count:N, an error occurs. Assignments are always global. \fparray_gzero:N \fparray_gzero:N \fparray var \ New: 2018-05-05 Sets all entries of the (floating point array variable) to +0. Assignments are always global.

\fparray_item:Nn

\fparray_item_to_tl:Nn * Applies \fp_use:N or \fp_to_tl:N (respectively) to the floating point entry stored at New: 2018-05-05 the (integer expression) $\langle position \rangle$ in the $\langle floating\ point\ array\ variable \rangle$. If the $\langle position \rangle$ is not between 1 and the \fparray_count:N, an error occurs.

Chapter 30

The **I3cctab** package Category code tables

A category code table enables rapid switching of all category codes in one operation. For LuaT_FX, this is possible over the entire Unicode range. For other engines, only the 8-bit range (0-255) is covered by such tables.

30.1Creating and initialising category code tables

\cctab_new:N

\cctab_new:N \(category \) code table \(\)

\cctab_new:c

Creates a new (category code table) variable or raises an error if the name is already Updated: 2020-07-02 taken. The declaration is global. The $\langle category \ code \ table \rangle$ is initialised with the codes as used by iniT_FX.

\cctab_const:Nn \cctab_const:Nn \category code table \ {\category code set up\}

\cctab_const:cn

Creates a new $\langle category\ code\ table \rangle$, applies (in a group) the $\langle category\ code\ set\ up \rangle$ on Updated: 2020-07-07 top of iniT_FX settings, then saves them globally as a constant table. The $\langle category \ code \$ $set up \rangle$ can include a call to \cctab_select:N.

\cctab_gset:cn

\cctab_gset:Nn \cctab_gset:Nn \category code table \ {\category code set up\}

Starting from the iniT_FX category codes, applies (in a group) the $\langle category \ code \ set \ up \rangle$, Updated: 2020-07-07 then saves them globally in the $\langle category \ code \ table \rangle$. The $\langle category \ code \ set \ up \rangle$ can include a call to \cctab_select:N.

Using category code tables 30.2

\cctab_begin:N \cctab_begin:N \category code table \

 $\cdot = \cdot = \cdo$ Updated: 2020-07-02 The prevailing codes before the function is called are added to a stack, for use with \cctab_end:. This function does not start a TeX group.

\cctab_end:	\cctab_end:
Updated: 2020-07-02	Ends the scope of a $\langle category\ code\ table \rangle$ started using $\cctab_begin:N$, returning the codes to those in force before the matching $\cctab_begin:N$ was used. This must be used within the same T_EX group (and at the same T_EX group level) as the matching $\cctab_begin:N$.
\cctab_select:c New: 2020-05-19	$\label{lem:cctab_select:N} $$ \category \ code \ table $$ $$ Selects the $$ \category \ code \ table $$ $$ for the scope of the current group. This is in particular useful in the $$ \setup $$ $$ arguments of $$tl_set_rescan:Nnn, $$tl_rescan:nn, $$ $$ const:Nn, and $$ \cctab_gset:Nn. $$$
\cctab_item:cn *	$\label{lem:nn category code table} $$ { \langle integer expression \rangle }$$ Determines the $$ \langle character \rangle$ with character code given by the $$ \langle integer expression \rangle$ and expands to its category code specified by the $$ \langle category code table \rangle$.$

30.3 Category code table conditionals

	30.4 Constant category code tables
\c_code_cctab Updated: 2020-07-10	Category code table for the $\exp 3$ code environment; this does not include $@$, which is retained as an "other" character.
	Category code table for a standard LATEX document, as set by the LATEX kernel. In particular, the upper-half of the 8-bit range will be set to "active" with pdfTEX only. No babel shorthands will be activated.
\c_initex_cctab	Category code table as set up by iniTeX.
\c_other_cctab Updated: 2020-07-02	Category code table where all characters have category code 12 (other).
\c_str_cctab Updated: 2020-07-02	Category code table where all characters have category code 12 (other) with the exception of spaces, which have category code 10 (space).

 ${f Part\ V}$ Text manipulation

Chapter 31

The l3unicode package: Unicode support functions

This module provides Unicode-specific functions along with loading data from a range of Unicode Consortium files. Most of the code here is internal, but there are a small set of public functions. These work with Unicode $\langle codepoints \rangle$ and are designed to give useable results with both Unicode-aware and 8-bit engines.

\codepoint_generate:nn * \codepoint_generate:nn {\langle codepoint \rangle {\langle catcode \rangle}

New: 2022-10-09 Generates one or more character tokens representing the $\langle codepoint \rangle$. With Unicode Updated: 2022-11-09 engines, exactly one character token will be generated, and this will have the (catcode) specified as the second argument:

- 1 (begin group)
- 2 (end group)
- 3 (math toggle)
- 4 (alignment)
- 6 (parameter)
- 7 (math superscript)
- 8 (math subscript)
- 10 (space)
- 11 (letter)
- 12 (other)
- 13 (active)

For 8-bit engines, between one and four character tokens will be produced: these will be the bytes of the UTF-8 representation of the $\langle codepoint \rangle$. For all codepoints outside of the classical ASCII range, the generated character tokens will be active (category code 13); for codepoints in the ASCII range, the given $\langle catcode \rangle$ will be used. To allow the result of this function to be used inside a expansion context, the result is protected by \exp_not:n.

 $\verb|\codepoint_str_generate:n| * \verb|\codepoint_str_generate:n| \{ \langle codepoint \rangle \}|$

New: 2022-10-09

Generates one or more character tokens representing the $\langle codepoint \rangle$. With Unicode engines, exactly one character token will be generated. For 8-bit engines, between one and four character tokens will be produced: these will be the bytes of the UTF-8 representation of the (codepoint). All of the generated character tokens will be of category code 12, except any spaces (codepoint 32), which will be category code 10.

\codepoint_to_nfd:n * \codepoint_to_nfd:n {\langle codepoint \rangle}

New: 2022-10-09 Converts the $\langle codepoint \rangle$ to the Unicode Normalization Form Canonical Decomposition. The generated character(s) will have the current category code as they would if typed in directly for Unicode engines; for 8-bit engines, active characters are used for all codepoints outside of the ASCII range.

Chapter 32

The l3text package: Text processing

This module deals with manipulation of (formatted) text; such material is comprised of a restricted set of token list content. The functions provided here concern conversion of textual content for example in case changing, generation of bookmarks and extraction to tags. All of the major functions operate by expansion. Begin-group and end-group tokens in the $\langle text \rangle$ are normalized and become { and }, respectively.

Expanding text 32.1

 $\texttt{\text_expand:n } \star \texttt{\text_expand:n } \{\langle text \rangle\}$

New: 2020-01-02 Takes user input $\langle text \rangle$ and expands the content. Protected commands (typically format-Updated: 2023-02-07 ting) are left in place, and no processing takes place of math mode material (as delimited by pairs given in \l_text_math_delims_tl or as the argument to commands listed in \1 text math arg t1). Commands which are neither engine- nor IATEX protected are expanded exhaustively. Any commands listed in \l_text_expand_exclude_tl are excluded from expansion.

 $\text{text_declare_expand_equivalent:Nn } \text{text_declare_expand_equivalent:Nn } \{\text{cmd}\} \{\text{creplacement}\}$ \text_declare_expand_equivalent:cn New: 2020-01-22

> Declares that the $\langle replacement \rangle$ tokens should be used whenever the $\langle cmd \rangle$ (a single token) is encountered. The $\langle replacement \rangle$ tokens should be expandable. A token can be "replaced" by itself if the defined replacement wraps it in \exp_not:n, for example

\text_declare_expand_equivalent:Nn \' { \exp_not:n { \' } }

32.2Case changing

 $\text{text_uppercase:n} \{\langle tokens \rangle\}$

 $\text{text_uppercase:nn } \{\langle language \rangle\} \{\langle tokens \rangle\}$

```
\text_lowercase:n
\text_uppercase:n
\text_titlecase:n
\text_titlecase_first:n
\text_lowercase:nn
\text_uppercase:nn
\text_titlecase:nn
```

New: 2019-11-20

Updated: 2022-10-13

Takes user input $\langle text \rangle$ first applies \text_expand, then transforms the case of character tokens as specified by the function name. The category code of letters are not changed by this process when Unicode engines are used; in 8-bit engines, case changed charters in the ASCII range will have the current prevailing category code, while those outside of \text_titlecase_first:nn * it will be represented by active characters.

> Upper- and lowercase have the obvious meanings. Titlecasing may be regarded informally as converting the first character of the $\langle tokens \rangle$ to uppercase and the rest to lowercase. However, the process is more complex than this as there are some situations where a single lowercase character maps to a special form, for example ij in Dutch which becomes IJ. The titlecase_first variant does not attempt any case changing at all after the first letter has been processed.

> Importantly, notice that these functions are intended for working with user text for typesetting. For case changing programmatic data see the l3str module and discussion there of \str_lowercase:n, \str_uppercase:n and \str_casefold:n.

Case changing does not take place within math mode material so for example

```
\text_uppercase:n { Some~text~$y = mx + c$~with~{Braces} }
```

becomes

```
SOME TEXT y = mx + c WITH {BRACES}
```

The first mandatory argument of commands listed in \l_text_case_exclude_arg_tl is excluded from case changing; the latter are entirely non-textual content (such as labels).

The standard mappings here follow those defined by the Unicode Consortium in UnicodeData.txt and SpecialCasing.txt. For pTFX, only the ASCII range is covered as the engine treats input outside of this range as east Asian.

Language-sensitive conversions are enabled using the $\langle language \rangle$ argument, and follow Unicode Consortium guidelines. Currently, the languages recognised for special handling are as follows.

- Armenian (hy and hy-x-yiwn) The setting hy maps the codepoint U+0587, the ligature of letters ech and yiwn, to the codepoints for capital ech and vew when uppercasing: this follows the spelling reform which is used in Armenia. The alternative hy-x-yiwn maps U+0587 to capital ech and yiwn on uppercasing (also the output if Armenian is not selected at all).
- Azeri and Turkish (az and tr). The case pairs I/i-dotless and I-dot/i are activated for these languages. The combining dot mark is removed when lowercasing I-dot and introduced when upper casing i-dotless.
- German (de-x-eszett). An alternative mapping for German in which the lowercase Eszett maps to a großes Eszett. Since there is a T1 slot for the großes Eszett in T1, this tailoring is available with pdfTFX as well as in the Unicode TFX engines.

- Greek (el). Removes accents from Greek letters when uppercasing; titlecasing leaves accents in place. A variant el-x-iota is available which converts the *ypoge-grammeni* (subscript muted iota) to capital iota when uppercasing: the standard version retains the subscript versions. capital iota.
- Lithuanian (1t). The lowercase letters i and j should retain a dot above when the accents grave, acute or tilde are present. This is implemented for lowercasing of the relevant uppercase letters both when input as single Unicode codepoints and when using combining accents. The combining dot is removed when uppercasing in these cases. Note that *only* the accents used in Lithuanian are covered: the behaviour of other accents are not modified.
- Medieval Latin (la-xmedieval). The characters u and V are interchanged on case changing.
- Dutch (n1). Capitalisation of ij at the beginning of titlecased input produces IJ rather than Ij. The output retains two separate letters, thus this transformation is available using pdfT_EX.

For titlecasing, note that there are two functions available. The function \text_titlecase:n applies (broadly) uppercasing to the first letter of the input, then lowercasing to the remainder. In contrast, \text_titlecase_first:n only carries out the uppercasing operation, and leaves the balance of the input unchanged. Determining whether non-letter characters at the start of text should switch from upper- to lowercasing is controllable. When \l_text_titlecase_check_letter_bool is true, characters which are not letters (category code 11) are left unchanged and "skipped": the first letter is uppercased. (With 8-bit engines, this is extended to active characters which form part of a multi-byte letter codepoint.) When \l_text_titlecase_check_letter_bool is false, the first character is uppercased, and the rest lowercased, irrespective of the nature of the character.

```
\text_declare_case_equivalent:Nn \text_declare_case_equivalent:Nn \( cmd \) \{\rangle replacement \)} \text_declare_case_equivalent:cn

New: 2022-07-04
```

Declares that the $\langle replacement \rangle$ tokens should be used whenever the $\langle cmd \rangle$ (a single token) is encountered during case changing.

```
\text{text\_case\_switch:nnnn } \star \text{text\_case\_switch:nnnn } {\langle normal \rangle} {\langle upper \rangle} {\langle lower \rangle} {\langle title \rangle}
```

Nov. 2022-07-0/

New: 2022-07-04 Context-sensitive function which will expand to one of the $\langle normal \rangle$, $\langle upper \rangle$, $\langle lower \rangle$ or $\langle title \rangle$ tokens depending on the current case changing operation. Outside of case changing, the $\langle normal \rangle$ tokens are produced. Within case changing, the appropriate mapping tokens are inserted.

Removing formatting from text 32.3

 $\text{text_purify:n} \star \text{text_purify:n} \{\langle text \rangle\}$

New: 2020-03-05 Takes user input $\langle text \rangle$ and expands as described for \text_expand:n, then removes all Updated: 2020-05-14 functions from the resulting text. Math mode material (as delimited by pairs given in \l_text_math_delims_tl or as the argument to commands listed in \l_text_math_arg_t1) is left contained in a pair of \$ delimiters. Non-expandable functions present in the $\langle text \rangle$ must either have a defined equivalent (see \text_declare_purify_equivalent: Nn) or will be removed from the result. Implicit tokens are converted to their explicit equivalent.

\text_declare_purify_equivalent:Nn \text_declare_purify_equivalent:Nn \cmd\ {\text_declare_purify_equivalent:Nn \cmd\} \text_declare_purify_equivalent:Nx

New: 2020-03-05

Declares that the $\langle replacement \rangle$ tokens should be used whenever the $\langle cmd \rangle$ (a single token) is encountered. The $\langle replacement \rangle$ tokens should be expandable.

Control variables 32.4

 $\label{lem:lemmath_arg_tl}$ Lists commands present in the $\langle text \rangle$ where the argument of the command should be treated as math mode material. The treatment here is similar to \l_text_math_delims_tl but for a command rather than paired delimiters.

\l_text_math_delims_tl Lists pairs of tokens which delimit (in-line) math mode content; such content may be excluded from processing.

\l_text_case_exclude_arg_tl

Lists commands where the first mandatory argument is excluded from case changing.

\l_text_expand_exclude_tl Lists commands which are excluded from expansion. This protection includes everything up to and including their first braced argument.

\l_text_titlecase_check_letter_bool

Controls how the start of titlecasing is handled: when true, the first letter in text is considered. The standard setting is true.

32.5 Mapping to graphemes

Grapheme splitting is implemented using the algorithm described in Unicode Standard Annex #29. This includes support for extended grapheme clusters. Text starting with a line feed or carriage return character will drop this due to standard TFX processing. At present extended pictograms are not supported: these may be added in a future release.

\text_map_function:nN ☆

\text map function:nN \langle text \rangle \langle \text \langle \langle \langle \text \langle \langle \langle \text \langle \langle \langle \text \langle \langle \langle \text \langle \la

New: 2022-08-04 Takes user input \(\lambda text\rangle\) and expands as described for \text_expand:n, then maps over the graphemes within the result, passing each grapheme to the $\langle function \rangle$. Broadly a grapheme is a "user perceived character": the Unicode Consortium describe the decomposition of input to graphemes in depth, and the approach used here implements that algorithm. The $\langle function \rangle$ should accept one argument as $\langle balanced\ text \rangle$: this may be comprise codepoints or may be a control sequence. With 8-bit engines, the codepoint(s) themselves may of course be made up of multiple bytes: the mapping will pass the correct codepoints independent of the engine in use. See also \text_map_inline:nn.

\text_map_inline:nn \text_map_inline:nn \text\ {\langle inline function \rangle}

New: 2022-08-04 Takes user input \(\lambda text\rangle\) and expands as described for \text expand:n, then maps over the graphemes within the result, passing each grapheme to the $\langle inline\ function \rangle$. Broadly a grapheme is a "user perceived character": the Unicode Consortium describe the decomposition of input to graphemes in depth, and the approach used here implements that algorithm. The $\langle inline\ function \rangle$ should consist of code which receives the grapheme as (balanced text): this may be comprise codepoints or may be a control sequence. With 8-bit engines, the codepoint(s) themselves may of course be made up of multiple bytes: the mapping will pass the correct codepoints independent of the engine in use. See also \text_map_function:nN.

```
\text_map_break:
```

☆ \text_map_break:

\text_map_break:n ☆

 $\text{text_map_break:n } \{\langle code \rangle\}$

New: 2022-08-04 Used to terminate a \text_map_... function before all entries in the $\langle text \rangle$ have been processed. This normally takes place within a conditional statement.

 $egin{array}{c} ext{Part VI} \ ext{Typesetting} \end{array}$

Chapter 33

The **I3box** package Boxes

Box variables contain typeset material that can be inserted on the page or in other boxes. Their contents cannot be converted back to lists of tokens. There are three kinds of box operations: horizontal mode denoted with prefix \hbox_, vertical mode with prefix \vbox_, and the generic operations working in both modes with prefix \box_. For instance, a new box variable containing the words "Hello, world!" (in a horizontal box) can be obtained by the following code.

```
\box_new:N \l_hello_box
\hbox_set:Nn \l_hello_box { Hello, ~ world! }
```

The argument is typeset inside a TeX group so that any variables assigned during the construction of this box restores its value afterwards.

Box variables from l3box are compatible with those of LATEX 2_{ε} and plain TEX and can be used interchangeably. The l3box commands to construct boxes, such as \hbox:n or \hbox_set:Nn, are "color-safe", meaning that

```
\hbox:n { \color_select:n { blue } Hello, } ~ world!
```

will result in "Hello," taking the color blue, but "world!" remaining with the prevailing color outside the box.

33.1 Creating and initialising boxes

 $\verb|\box_new:N| \box_new:N| \aligned box|$

\box_new:c

Creates a new $\langle box \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle box \rangle$ is initially void.

\box_clear:N \box_clear:N \box_clear:N \
box_clear:c Cl

\box_gclear: Clears the content of the $\langle box \rangle$ by setting the box equal to \c_empty_box.

\box_gclear:c

```
\box_clear_new:N
\box_clear_new:c
\box_gclear_new:N
\box_gclear_new:c

\box_gclear_new:c

\box_set_eq:NN
\box_set_eq:(cN|Nc|cc)
\box_gset_eq:(cN|Nc|cc)
```

```
\box_clear_new:N \box_clear_new:N \box>
```

Ensures that the $\langle box \rangle$ exists globally by applying \box_new:N if necessary, then applies \box_(g) clear:N to leave the $\langle box \rangle$ empty.

```
\box_set_eq:NN \langle box_1 \rangle \langle box_2 \rangle
Sets the content of \langle box_1 \rangle equal to that of \langle box_2 \rangle.
```

33.2 Using boxes

\box_use:N \box_use:N \box_use:N \box_use:C

\box_use:c Inserts the current content of the $\langle box \rangle$ onto the current list for typesetting. An error is raised if the variable does not exist or if it is invalid.

TEXhackers note: This is the TEX primitive \copy.

\box_move_right:nn \box_move_left:nn

```
\box_move_right:nn \box_move_right:nn {\dimexpr\} {\dimexpr\}
```

This function operates in vertical mode, and inserts the material specified by the $\langle box function \rangle$ such that its reference point is displaced horizontally by the given $\langle dimexpr \rangle$ from the reference point for typesetting, to the right or left as appropriate. The $\langle box function \rangle$ should be a box operation such as $\box_use:N \c)$ or a "raw" box specification such as $\box_use:N \c)$.

\box_move_up:nn \box_move_down:nn

```
\verb|\box_move_up:nn| \{\langle dimexpr \rangle\} \ \{\langle box\ function \rangle\}|
```

This function operates in horizontal mode, and inserts the material specified by the $\langle box function \rangle$ such that its reference point is displaced vertically by the given $\langle dimexpr \rangle$ from the reference point for typesetting, up or down as appropriate. The $\langle box function \rangle$ should be a box operation such as $\box_use:N \c)$ or a "raw" box specification such as $\box_nin \{ xyz \}$.

Measuring and setting box dimensions 33.3

 $\verb|\box_dp:N \box_dp:N \label{local_power}| \label{local_power}$

\box_dp:c Calculates the depth (below the baseline) of the $\langle box \rangle$ in a form suitable for use in a $\langle dimension \ expression \rangle$.

TeXhackers note: This is the TeX primitive \dp.

\box_ht:N \box_ht:N \box_h

Calculates the height (above the baseline) of the $\langle box \rangle$ in a form suitable for use in a $\langle dimension \ expression \rangle$.

TEXhackers note: This is the TEX primitive \ht.

\box_wd:N \box_wd:N \box}

 $\begin{tabular}{ll} \begin{tabular}{ll} \beg$

TEXhackers note: This is the TEX primitive \wd.

\box_ht_plus_dp:N \box_ht_plus_dp:N \box>

 $\underline{\text{box_ht_plus_dp:c}}$ Calculates the total vertical size (height plus depth) of the $\langle box \rangle$ in a form suitable for

New: 2021-05-05 use in a $\langle dimension \ expression \rangle$.

 $\verb|\box_set_dp:Nn \ \box_set_dp:Nn \ \ \box| \ \{ \langle dimension \ expression \rangle \}$

\box_gset_dp:\n Set the depth (below the baseline) of the $\langle box \rangle$ to the value of the $\{\langle dimension\}\}$

\box_gset_dp:cn expression\}.

Updated: 2019-01-22

 $\verb|\box_set_ht:Nn \ \box_set_ht:Nn \ \box| \{ \langle dimension \ expression \rangle \}$

\box_set_ht:cn

Set the height (above the baseline) of the $\langle box \rangle$ to the value of the $\{\langle dimension\}\}$ \box_gset_ht:Nn \box_gset_ht:cn expression\}.

Updated: 2019-01-22

 $\box_set_wd:Nn \box_set_wd:Nn \box\} \{\langle dimension \ expression \rangle\}$

\box_set_wd:cn

Set the width of the $\langle box \rangle$ to the value of the $\{\langle dimension \ expression \rangle\}$. \box_gset_wd:Nn

\box_gset_wd:cn Updated: 2019-01-22

Box conditionals 33.4

```
\box_if_empty_p:N \star \box_if_empty_p:N \langle box \rangle
                                         \verb|\box_if_empty_p:c * \box_if_empty:NTF $$\langle box \rangle $ \{\langle true \ code \rangle \} $$\{\langle false \ code \rangle \} $$
                                        \verb|\box_if_empty:N| $\underline{\mathit{TF}}$ * Tests if $\langle box \rangle$ is a empty (equal to \verb|\c_empty_box|).
                                         \box_if_empty:cTF *
  \box_if_horizontal_p:N * \box_if_horizontal_p:N \langle box \rangle
  \begin{cal}{l} \beg
\label{eq:loss_if_horizontal:N} $$ \text{Tests if } \langle box \rangle$ is a horizontal box.
 \box_if_horizontal:cTF
                \verb|\box_if_vertical_p:N * \verb|\box_if_vertical_p:N| & box_if_vertical_p:N| & box_if_vertical
                 \sum_{i=1}^{\infty} \frac{1}{i} \operatorname{vertical_p:c} \star \operatorname{vertical:NTF} \langle box \rangle \{\langle true\ code \rangle\} \{\langle false\ code \rangle\}
                 \box_if_vertical:N\overline{TF} * Tests if \langle box \rangle is a vertical box.
                 \box_if_vertical:cTF *
```

33.5The last box inserted

\box_set_to_last:N \box_set_to_last:N \box>

 $\begin{array}{l} \texttt{NDOX_SEL_LO_LAST:C} \\ \texttt{Nbox_gset_to_last:N} \end{array} \text{ Sets the } \langle box \rangle \text{ equal to the last item (box) added to the current partial list, removing the} \end{array}$ \box_gset_to_last:c item from the list at the same time. When applied to the main vertical list, the $\langle box \rangle$ is always void as it is not possible to recover the last added item.

33.6 Constant boxes

\c_empty_box

This is a permanently empty box, which is neither set as horizontal nor vertical.

Updated: 2012-11-04

TeXhackers note: At the TeX level this is a void box.

Scratch boxes 33.7

\l_tmpa_box \l_tmpb_box

Scratch boxes for local assignment. These are never used by the kernel code, and so are safe for use with any IATEX3-defined function. However, they may be overwritten by Updated: 2012-11-04 other non-kernel code and so should only be used for short-term storage.

\g_tmpa_box Scratch boxes for global assignment. These are never used by the kernel code, and so \g_tmpb_box are safe for use with any LATEX3-defined function. However, they may be overwritten by other non-kernel code and so should only be used for short-term storage.

33.8 Viewing box contents

33.9 Boxes and color

All LATEX3 boxes are "color safe": a color set inside the box stops applying after the end of the box has occurred.

33.10 Horizontal mode boxes

\hbox:n	$\hbox:n \ \{\langle contents \rangle\}$
Updated: 2017-04-05	Type sets the $\langle \mathit{contents} \rangle$ into a horizontal box of natural width and then includes this box in the current list for type setting.
\hbox_to_wd:nn	$\label{local_state} $$ \ \ $
Updated: 2017-04-05	Type sets the $\langle contents \rangle$ into a horizontal box of width $\langle dimexpr \rangle$ and then includes this box in the current list for type setting.
\hbox_to_zero:n	$\hbox_to_zero:n \ \{\langle contents \rangle\}$
Updated: 2017-04-05	Type sets the $\langle contents \rangle$ into a horizontal box of zero width and then includes this box in the current list for type setting.
\hbox_set:Nn	$\label{local_norm_local} $$ \box_set:Nn $$ \langle box \rangle $$ {\langle contents \rangle}$$$
<pre>\hbox_set:cn \hbox_gset:Nn \hbox_gset:cn</pre>	Type sets the $\langle contents \rangle$ at natural width and then stores the result in side the $\langle box \rangle$.
Updated: 2017-04-05	

 $\begin{tabular}{ll} $$ \box_set_to_wd:Nnn $$\langle box\rangle $ {\langle dimexpr\rangle} $ {\langle contents\rangle}$ \end{tabular}$ \hbox_set_to_wd:Nnn \hbox_set_to_wd:cnn Typesets the $\langle contents \rangle$ to the width given by the $\langle dimexpr \rangle$ and then stores the result \hbox_gset_to_wd:Nnn inside the $\langle box \rangle$. \hbox_gset_to_wd:cnn Updated: 2017-04-05 $\verb|\hbox_overlap_center:n \hbox_overlap_center:n \{\langle contents \rangle\}|$ New: 2020-08-25 Typesets the (contents) into a horizontal box of zero width such that material protrudes equally to both sides of the insertion point. $\verb|\hbox_overlap_right:n \hbox_overlap_right:n {| \langle contents \rangle|}$ Updated: 2017-04-05 Typesets the (contents) into a horizontal box of zero width such that material protrudes to the right of the insertion point. Updated: 2017-04-05 Typesets the (contents) into a horizontal box of zero width such that material protrudes to the left of the insertion point. \hbox_set:Nw \langle box \langle contents \langle \hbox_set_end: \hbox_set:Nw \hbox_set:cw Typesets the $\langle contents \rangle$ at natural width and then stores the result inside the $\langle box \rangle$. In \hbox_set_end: contrast to \hbox_set:Nn this function does not absorb the argument when finding the \hbox_gset:Nw $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple \hbox_gset:cw argument. \hbox_gset_end: Updated: 2017-04-05

\hbox_set_to_wd:Nnw
\hbox_set_to_wd:cnw
\hbox_gset_to_wd:Nnw
\hbox_gset_to_wd:cnw

 $\verb|\hbox_set_to_wd:Nnw| \langle box \rangle | \{\langle dimexpr \rangle\} | \langle contents \rangle | \\ | hbox_set_end: \\ | hbox_set_end:$

Typesets the $\langle contents \rangle$ to the width given by the $\langle dimexpr \rangle$ and then stores the result inside the $\langle box \rangle$. In contrast to $\hbox_set_to_wd:\hnn$ this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument

...

\hbox_unpack:c

New: 2017-06-08

\hbox_unpack:N \hbox_unpack:N \dox\

Unpacks the content of the horizontal $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set.

TEXhackers note: This is the TEX primitive \unhcopy.

33.11 Vertical mode boxes

Vertical boxes inherit their baseline from their contents. The standard case is that the baseline of the box is at the same position as that of the last item added to the box. This means that the box has no depth unless the last item added to it had depth. As a result most vertical boxes have a large height value and small or zero depth. The exception are

_top boxes, where the reference point is that of the first item added. These tend to have a large depth and small height, although the latter is typically non-zero.

\vbox:n	$\wdots:n {(contents)}$		
Updated: 2017-04-05	Type sets the $\langle contents \rangle$ into a vertical box of natural height and includes this box in the current list for type setting.		
\vbox_top:n	$\wbox_top:n \ \{\langle contents \rangle\}$		
Updated: 2017-04-05	Typesets the $\langle contents \rangle$ into a vertical box of natural height and includes this box in the current list for typesetting. The baseline of the box is equal to that of the <i>first</i> item added to the box.		
\vbox_to_ht:nn	$\begin{tabular}{ll} $$ \begin{tabular}{ll} $$ \begin{tabular}{ll} $ & (contents) $ \end{tabular} \end{tabular}$		
Updated: 2017-04-05	Typesets the $\langle contents \rangle$ into a vertical box of height $\langle dimexpr \rangle$ and then includes this box in the current list for typesetting.		
\vbox_to_zero:n	$\wbox_to_zero:n \ \{\langle contents \rangle\}$		
Updated: 2017-04-05	Type sets the $\langle contents \rangle$ into a vertical box of zero height and then includes this box in the current list for type setting.		
\vbox_set:Nn \vbox_set:cn \vbox_gset:Nn \vbox_gset:cn Updated: 2017-04-05	lem:lem:lem:lem:lem:lem:lem:lem:lem:lem:		
	$\begin{tabular}{ll} \verb&vbox_set_top:Nn & $\langle box \rangle$ & {\langle contents \rangle$} \end{tabular}$		
\vbox_set_top:cn \vbox_gset_top:Nn \vbox_gset_top:cn	Typesets the $\langle contents \rangle$ at natural height and then stores the result inside the $\langle box \rangle$. The baseline of the box is equal to that of the first item added to the box.		
Updated: 2017-04-05			
<pre>\vbox_set_to_ht:Nnn \vbox_set_to_ht:cnn \vbox_gset_to_ht:nn \vbox_gset_to_ht:cnn</pre>	lem:lem:lem:lem:lem:lem:lem:lem:lem:lem:		
Updated: 2017-04-05			

\vbox_set:Nw
\vbox_set:cw
\vbox_set:Nw
\vbox_gset:Nw
\vbox_gset:cw
\vbox_gset_end:

 $\verb|\vbox_set:Nw| \langle box \rangle \ \langle contents \rangle \ \verb|\vbox_set_end:|$

Typesets the $\langle contents \rangle$ at natural height and then stores the result inside the $\langle box \rangle$. In contrast to $\vbox_set:Nn$ this function does not absorb the argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument.

Updated: 2017-04-05

```
\vbox_set_to_ht:Nnw
\vbox_set_to_ht:cnw
\vbox_gset_to_ht:Nnw
```

New: 2017-06-08

```
\vbox_set_to_ht:Nnw \langle box \ \{\dimexpr\}\ \contents \ \vbox_set_end:
```

Typesets the $\langle contents \rangle$ to the height given by the $\langle dimexpr \rangle$ and then stores the result $\color=0.05$ inside the $\color=0.05$ inside the $\color=0.05$. In contrast to $\color=0.05$ inside the $\color=0.05$ i argument when finding the $\langle content \rangle$, and so can be used in circumstances where the $\langle content \rangle$ may not be a simple argument

```
\vbox_set_split_to_ht:NNn
                                                 \verb|\vbox_set_split_to_ht:NNn| \langle box_1 \rangle | \langle box_2 \rangle | \{\langle dimexpr \rangle\}|
\vbox_set_split_to_ht:(cNn|Ncn|ccn)
\vbox_gset_split_to_ht:NNn
\vbox_gset_split_to_ht:(cNn|Ncn|ccn)
                           Updated: 2018-12-29
```

Sets $\langle box_1 \rangle$ to contain material to the height given by the $\langle dimexpr \rangle$ by removing content from the top of $\langle box_2 \rangle$ (which must be a vertical box).

```
\vbox_unpack:c
```

```
\vbox_unpack:N \vbox_unpack:N \box\
```

Unpacks the content of the vertical $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set.

TEXhackers note: This is the TEX primitive \unvcopy.

33.12 Using boxes efficiently

The functions above for using box contents work in exactly the same way as for any other expl3 variable. However, for efficiency reasons, it is also useful to have functions which drop box contents on use. When a box is dropped, the box becomes empty at the group level where the box was originally set rather than necessarily at the current group level. For example, with

```
\hbox_set:Nn \l_tmpa_box { A }
\group_begin:
  \hbox_set:Nn \l_tmpa_box { B }
   \group_begin:
   \box_use_drop:N \l_tmpa_box
  \group_end:
  \box_show:N \l_tmpa_box
\group end:
\box show: N \l tmpa box
```

the first use of \box_show:N will show an entirely cleared (void) box, and the second will show the letter A in the box.

These functions should be preferred when the content of the box is no longer required after use. Note that due to the unusual scoping behaviour of drop functions they may be applied to both local and global boxes: the latter will naturally be set and thus cleared at a global level.

```
\box_use_drop:N \box_use_drop:N \box_box_use_drop: N \box_use_drop: N \box
```

Inserts the current content of the $\langle box \rangle$ onto the current list for typesetting then drops the box content. An error is raised if the variable does not exist or if it is invalid. This function may be applied to local or global boxes.

TeXhackers note: This is the \box primitive.

\box_set_eq_drop:NN \box_set_eq_drop:(cN Nc cc) New: 2019-01-17	\box_set_eq_drop:NN $\langle box_1 \rangle$ $\langle box_2 \rangle$ Sets the content of $\langle box_1 \rangle$ equal to that of $\langle box_2 \rangle$, then drops $\langle box_2 \rangle$.
\box_gset_eq_drop:NN \box_gset_eq_drop:(cN Nc cc) New: 2019-01-17	\box_gset_eq_drop:NN $\langle box_1 \rangle$ $\langle box_2 \rangle$ Sets the content of $\langle box_1 \rangle$ globally equal to that of $\langle box_2 \rangle$, then drops $\langle box_2 \rangle$.
\hbox_unpack_drop:c	\hbox_unpack_drop:N $\langle box \rangle$ Unpacks the content of the horizontal $\langle box \rangle$, retaining any stretching or shrinking applied when the $\langle box \rangle$ was set. The original $\langle box \rangle$ is then dropped.

TeXhackers note: This is the TeX primitive \unhbox.

TeXhackers note: This is the TeX primitive \unvbox.

33.13 Affine transformations

Affine transformations are changes which (informally) preserve straight lines. Simple translations are affine transformations, but are better handled in TeX by doing the translation first, then inserting an unmodified box. On the other hand, rotation and resizing of boxed material can best be handled by modifying boxes. These transformations are described here.

```
\box_autosize_to_wd_and_ht:Nnn \box_autosize_to_wd_and_ht:Nnn \box_autosize_to_wd_and_ht:nnn \box_autosize_to_wd_and_ht:nnn \box_gautosize_to_wd_and_ht:nnn \b
```

Resizes the $\langle box \rangle$ to fit within the given $\langle x\text{-}size \rangle$ (horizontally) and $\langle y\text{-}size \rangle$ (vertically); both of the sizes are dimension expressions. The $\langle y\text{-}size \rangle$ is the height only: it does not include any depth. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. The final size of the $\langle box \rangle$ is the smaller of $\{\langle x\text{-}size \rangle\}$ and $\{\langle y\text{-}size \rangle\}$, i.e. the result fits within the dimensions specified. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vice versa.

```
\box_autosize_to_wd_and_ht_plus_dp:Nnn \box_autosize_to_wd_and_ht_plus_dp:cnn \box_gautosize_to_wd_and_ht_plus_dp:nnn \box_gautosize_to_wd_and_ht_plus_dp:cnn \box_gautosize_to_wd_and_ht_plus_dp:cnn \box_gautosize_to_wd_and_ht_plus_dp:cnn \box_2017-04-04 \box_2019-01-22
```

Resizes the $\langle box \rangle$ to fit within the given $\langle x\text{-}size \rangle$ (horizontally) and $\langle y\text{-}size \rangle$ (vertically); both of the sizes are dimension expressions. The $\langle y\text{-}size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. The final size of the $\langle box \rangle$ is the smaller of $\{\langle x\text{-}size \rangle\}$ and $\{\langle y\text{-}size \rangle\}$, i.e. the result fits within the dimensions specified. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vice versa.

```
\box_resize_to_ht:Nn
\box_resize_to_ht:cn
\box_gresize_to_ht:Nn
\box_gresize_to_ht:cn
```

 $\verb|\box_resize_to_ht:Nn| \langle box \rangle | \{\langle y\text{-}size \rangle\}|$

Resize the $\langle box \rangle$ to $\langle y\text{-}size \rangle$ (vertically), scaling the horizontal size by the same amount; $\langle y\text{-}size \rangle$ is a dimension expression. The $\langle y\text{-}size \rangle$ is the height only: it does not include any depth. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle y\text{-}size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vice versa.

```
\box_resize_to_ht_plus_dp:Nn \box_resize_to_ht_plus_dp:Nn \box_resize_to_ht_plus_dp:Cn \box_gresize_to_ht_plus_dp:Nn \box_gresize_to_ht_plus_dp:Cn \box_gres
```

Resizes the $\langle box \rangle$ to $\langle y\text{-}size \rangle$ (vertically), scaling the horizontal size by the same amount; $\langle y\text{-}size \rangle$ is a dimension expression. The $\langle y\text{-}size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle y\text{-}size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vicev ersa.

\box_resize_to_wd:Nn \box_resize_to_wd:cn \box_gresize_to_wd:Nn \box_gresize_to_wd:cn

 $\verb|\box_resize_to_wd:Nn| \langle box \rangle | \{\langle x\text{-}size \rangle\}|$

Resizes the $\langle box \rangle$ to $\langle x\text{-}size \rangle$ (horizontally), scaling the vertical size by the same amount; $\langle x\text{-}size \rangle$ is a dimension expression. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. A negative $\langle x\text{-}size \rangle$ causes the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle x\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and $vice\ versa$.

```
\box_resize_to_wd_and_ht:Nnn \box_resize_to_wd_and_ht:Nnn \box_size\} {\langle y-size \rangle} {
```

Resizes the $\langle box \rangle$ to $\langle x\text{-}size \rangle$ (horizontally) and $\langle y\text{-}size \rangle$ (vertically): both of the sizes are dimension expressions. The $\langle y\text{-}size \rangle$ is the height only and does not include any depth. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vice versa.

```
\label{localize_to_wd_and_ht_plus_dp:Nnn} $$ \ \c = to_wd_and_ht_plus_dp:Nnn \ \c = to_wd_and_ht_plus_dp:Cnn \ \c = to_wd_an
```

Resizes the $\langle box \rangle$ to $\langle x\text{-}size \rangle$ (horizontally) and $\langle y\text{-}size \rangle$ (vertically): both of the sizes are dimension expressions. The $\langle y\text{-}size \rangle$ is the total vertical size (height plus depth). The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the resizing is applied. Negative sizes cause the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}size \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and $vice\ versa$.

```
\box_rotate:Nn
\box_rotate:cn
\box_grotate:Nn
```

```
\box_rotate:Nn \langle box \rangle \{\langle angle \rangle\}
```

Rotates the $\langle box \rangle$ by $\langle angle \rangle$ (in degrees) anti-clockwise about its reference point. The \box_grotate:cn reference point of the updated box is moved horizontally such that it is at the left side of the smallest rectangle enclosing the rotated material. The updated $\langle box \rangle$ is an hbox, $\overset{\tt Updated:\,2019-01-22}{---}$ irrespective of the nature of the $\langle box \rangle$ before the rotation is applied.

```
\box_scale:cnn
\box_gscale:Nnn
```

```
\box_scale:Nnn \box_scale:Nnn \box \ \{\langle x-scale \rangle\} \ \{\langle y-scale \rangle\}
```

Scales the $\langle box \rangle$ by factors $\langle x\text{-}scale \rangle$ and $\langle y\text{-}scale \rangle$ in the horizontal and vertical directions, \box_gscale: cnn respectively (both scales are integer expressions). The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the scaling is applied. Negative scalings cause Updated: 2019-01-22 the material in the $\langle box \rangle$ to be reversed in direction, but the reference point of the $\langle box \rangle$ is unchanged. Thus a negative $\langle y\text{-}scale \rangle$ results in the $\langle box \rangle$ having a depth dependent on the height of the original and vice versa.

33.14 Primitive box conditionals

```
\if_hbox:N * \if_hbox:N \langle box\
                     \langle true code \rangle
                   \else:
                     ⟨false code⟩
```

Tests is $\langle box \rangle$ is a horizontal box.

TeXhackers note: This is the TeX primitive \ifhbox.

```
⟨true code⟩
            \else:
              ⟨false code⟩
            \fi:
            Tests is \langle box \rangle is a vertical box.
```

TeXhackers note: This is the TeX primitive \ifvbox.

```
\if_box_empty:N *
                          \ightharpoonup (box_empty:N \langle box \rangle)
                              ⟨true code⟩
                              ⟨false code⟩
                           \fi:
                          Tests is \langle box \rangle is an empty (void) box.
```

TEXhackers note: This is the TEX primitive \ifvoid.

Chapter 34

\coffin_new:N \coffin_new:N \coffin

The **I3coffins** package Coffin code layer

The material in this module provides the low-level support system for coffins. For details about the design concept of a coffin, see the xcoffins module (in the l3experimental bundle).

34.1 Creating and initialising coffins

```
\coffin_new:c
                                                                                                                                                                       Creates a new \langle coffin \rangle or raises an error if the name is already taken. The declaration is
                                                                                       New: 2011-08-17 global. The \langle coffin \rangle is initially empty.
                                                        \coffin_clear:N \coffin_clear:N \coffin
                                                         \coffin_clear:c
                                                                                                                                                                        Clears the content of the \langle coffin \rangle.
                                                         \coffin_gclear:N
                                                          \coffin_gclear:c
                                                                                       New: 2011-08-17
                                                                 Updated: 2019-01-21
                                                                                                                                                                          \coffin\_set\_eq:NN \ \langle coffin_1 \rangle \ \langle coffin_2 \rangle
\coffin_set_eq:NN
\coffin_set_eq:(Nc|cN|cc)
                                                                                                                                                                         Sets both the content and poles of \langle coffin_1 \rangle equal to those of \langle coffin_2 \rangle.
\coffin_gset_eq:NN
\coffin_gset_eq:(Nc|cN|cc)
                                                                                       New: 2011-08-17
                                                                 Updated: 2019-01-21
              \verb|\coffin_if_exist_p:N * \verb|\coffin_if_exist_p:N | & coffin_if_exist_p:N | & 
               \verb|\coffin_if_exist_p:c * \verb|\coffin_if_exist:NTF| | & coffin| & {\coffin_if_exist:NTF| | coffin| & {\code} & {\code
               \coffin_if_exist:NTF \times Tests whether the \langle coffin \rangle is currently defined.
               \coffin_if_exist:cTF *
                                                                                     New: 2012-06-20
```

34.2Setting coffin content and poles

\hcoffin_set:Nn \hcoffin_set:cn $\coffin_set: \verb"Nn $$\langle coffin \rangle $$ \{\langle material \rangle \}$$

\hcoffin_gset:Nn \hcoffin_gset:cn

Typesets the $\langle material \rangle$ in horizontal mode, storing the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material.

New: 2011-08-17 Updated: 2019-01-21

\hcoffin_set:Nw \(coffin \) \(\material \) \hcoffin_set_end:

\hcoffin_set:cw \hcoffin_set_end: \hcoffin_gset:Nw \hcoffin_gset:cw \hcoffin_gset_end:

\hcoffin_set:Nw

Typesets the $\langle material \rangle$ in horizontal mode, storing the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material. These functions are useful for setting the entire contents of an environment in a coffin.

New: 2011-09-10

Updated: 2019-01-21

\vcoffin set:Nnn \vcoffin_set:cnn $\coeffin_set:Nnn \ \langle coffin \rangle \ \{\langle width \rangle\} \ \{\langle material \rangle\}$

\vcoffin_gset:Nnn \vcoffin_gset:cnn

Typesets the $\langle material \rangle$ in vertical mode constrained to the given $\langle width \rangle$ and stores the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material.

New: 2011-08-17 Updated: 2019-01-21

\vcoffin_set:Nnw $\coffin_set:Nnw \coffin\coffin\colored{width} \dots\coffin_set_end:$

\vcoffin_set:cnw \vcoffin_set_end: \vcoffin_gset:Nnw \vcoffin_gset:cnw \vcoffin_gset_end:

Typesets the $\langle material \rangle$ in vertical mode constrained to the given $\langle width \rangle$ and stores the result in the $\langle coffin \rangle$. The standard poles for the $\langle coffin \rangle$ are then set up based on the size of the typeset material. These functions are useful for setting the entire contents of an environment in a coffin.

New: 2011-09-10 Updated: 2019-01-21

\coffin_set_horizontal_pole:Nnn \coffin_set_horizontal_pole:cnn \coffin_gset_horizontal_pole:Nnn

\coffin_gset_horizontal_pole:cnn

 $\coffin_set_horizontal_pole:Nnn \langle coffin \rangle$ $\{\langle pole \rangle\} \ \{\langle offset \rangle\}$

New: 2012-07-20 Updated: 2019-01-21

> Sets the $\langle pole \rangle$ to run horizontally through the $\langle coffin \rangle$. The $\langle pole \rangle$ is placed at the $\langle offset \rangle$ from the baseline of the $\langle coffin \rangle$. The $\langle offset \rangle$ should be given as a dimension expression.

```
\coffin_set_vertical_pole:Nnn \coffin_set_vertical_pole:Nnn \coffin \ {\langle pole \} {\langle offset \} \coffin_set_vertical_pole:Nnn \coffin_gset_vertical_pole:Nnn \coffin_gset_vertical_pole:cnn \ New: 2012-07-20 \ Updated: 2019-01-21
```

Sets the $\langle pole \rangle$ to run vertically through the $\langle coffin \rangle$. The $\langle pole \rangle$ is placed at the $\langle offset \rangle$ from the left-hand edge of the bounding box of the $\langle coffin \rangle$. The $\langle offset \rangle$ should be given as a dimension expression.

34.3 Coffin affine transformations

\coffin_resize:Nnn
\coffin_resize:cnn
\coffin_gresize:Nnn
\coffin_gresize:cnn

 $\verb|\coffin_resize:Nnn| & \langle coffin \rangle | \{ \langle width \rangle \} | \{ \langle total-height \rangle \}|$

Resized the $\langle coffin \rangle$ to $\langle width \rangle$ and $\langle total\text{-}height \rangle$, both of which should be given as dimension expressions.

Updated: 2019-01-23

\coffin_rotate:Nn
\coffin_rotate:cn
\coffin_grotate:Nn
\coffin_grotate:cn

 $\coffin_rotate:Nn \langle coffin \rangle \{\langle angle \rangle\}$

Rotates the $\langle coffin \rangle$ by the given $\langle angle \rangle$ (given in degrees counter-clockwise). This process rotates both the coffin content and poles. Multiple rotations do not result in the bounding box of the coffin growing unnecessarily.

\coffin_scale:Nnn
\coffin_scale:cnn
\coffin_gscale:Nnn
\coffin_gscale:cnn

 $\verb|\coffin_scale:Nnn| \langle coffin \rangle | \{\langle x-scale \rangle\} | \{\langle y-scale \rangle\}|$

Scales the $\langle coffin \rangle$ by a factors $\langle x\text{-}scale \rangle$ and $\langle y\text{-}scale \rangle$ in the horizontal and vertical directions, respectively. The two scale factors should be given as real numbers.

Updated: 2019-01-23

34.4 Joining and using coffins

Updated: 2019-01-22

This function attaches $\langle coffin_2 \rangle$ to $\langle coffin_1 \rangle$ such that the bounding box of $\langle coffin_1 \rangle$ is not altered, *i.e.* $\langle coffin_2 \rangle$ can protrude outside of the bounding box of the coffin. The alignment is carried out by first calculating $\langle handle_1 \rangle$, the point of intersection of $\langle coffin_1 \text{-}pole_1 \rangle$ and $\langle coffin_1 \text{-}pole_2 \rangle$, and $\langle handle_2 \rangle$, the point of intersection of $\langle coffin_2 \text{-}pole_1 \rangle$ and $\langle coffin_2 \text{-}pole_2 \rangle$. $\langle coffin_2 \rangle$ is then attached to $\langle coffin_1 \rangle$ such that the relationship between $\langle handle_1 \rangle$ and $\langle handle_2 \rangle$ is described by the $\langle x\text{-}offset \rangle$ and $\langle y\text{-}offset \rangle$. The two offsets should be given as dimension expressions.

```
\coffin_join:NnnNnnnn
                                                                          \coffin_join:NnnNnnnn
\coffin_join:(cnnNnnnn|Nnncnnnn|cnncnnnn)
                                                                              \langle coffin_1 \rangle \ \{\langle coffin_1 - pole_1 \rangle\} \ \{\langle coffin_1 - pole_2 \rangle\}
\coffin_gjoin:NnnNnnnn
                                                                              \langle 	ext{coffin}_2 
angle \; \{ \langle 	ext{coffin}_2	ext{-pole}_1 
angle \} \; \{ \langle 	ext{coffin}_2	ext{-pole}_2 
angle \}
\coffin_gjoin:(cnnNnnnn|Nnncnnnn|cnncnnnn)
                                                                              \{\langle x\text{-offset}\rangle\}\ \{\langle y\text{-offset}\rangle\}
                                              Updated: 2019-01-22
```

This function joins $\langle coffin_2 \rangle$ to $\langle coffin_1 \rangle$ such that the bounding box of $\langle coffin_1 \rangle$ may expand. The new bounding box covers the area containing the bounding boxes of the two original coffins. The alignment is carried out by first calculating $\langle handle_1 \rangle$, the point of intersection of $\langle coffin_1 - pole_1 \rangle$ and $\langle coffin_1 - pole_2 \rangle$, and $\langle handle_2 \rangle$, the point of intersection of $\langle coffin_2 - pole_1 \rangle$ and $\langle coffin_2 - pole_2 \rangle$. $\langle coffin_2 \rangle$ is then attached to $\langle coffin_1 \rangle$ such that the relationship between $\langle handle_1 \rangle$ and $\langle handle_2 \rangle$ is described by the $\langle x\text{-offset} \rangle$ and $\langle y\text{-offset}\rangle$. The two offsets should be given as dimension expressions.

```
\coffin_typeset:cnnnn
                \{\langle x-offset \rangle\}\ \{\langle y-offset \rangle\}
```

Updated: 2012-07-20 Typesetting is carried out by first calculating $\langle handle \rangle$, the point of intersection of $\langle pole_1 \rangle$ and $\langle pole_2 \rangle$. The coffin is then typeset in horizontal mode such that the relationship between the current reference point in the document and the $\langle handle \rangle$ is described by the $\langle x\text{-offset}\rangle$ and $\langle y\text{-offset}\rangle$. The two offsets should be given as dimension expressions. Typesetting a coffin is therefore analogous to carrying out an alignment where the "parent" coffin is the current insertion point.

Measuring coffins 34.5

```
\coffin_dp:N \coffin_dp:N \coffin\
\coffin_dp:c
                Calculates the depth (below the baseline) of the \langle coffin \rangle in a form suitable for use in a
                 \langle dimension \ expression \rangle.
\coffin_ht:N \coffin_ht:N \coffin
\coffin_ht:c
                Calculates the height (above the baseline) of the \langle coffin \rangle in a form suitable for use in a
                 \langle dimension \ expression \rangle.
\coffin_wd:N \coffin_wd:N \coffin
\coffin_wd:c
                Calculates the width of the \langle coffin \rangle in a form suitable for use in a \langle dimension \ expression \rangle.
```

34.6 Coffin diagnostics

```
\coffin display handles: Nn \coffin display handles: Nn \langle coffin \rangle \{\langle color \rangle\}
\coffin_display_handles:cn
```

This function first calculates the intersections between all of the $\langle poles \rangle$ of the $\langle coffin \rangle$ to give a set of $\langle handles \rangle$. It then prints the $\langle coffin \rangle$ at the current location in the source, with the position of the $\langle handles \rangle$ marked on the coffin. The $\langle handles \rangle$ are labelled as part of this process: the locations of the $\langle handles \rangle$ and the labels are both printed in the $\langle color \rangle$ specified.

\coffin_mark_handle:cnnn

 $\coffin_mark_handle:Nnnn \coffin_mark_handle:Nnnn \coffin \ {\langle pole_1 \rangle} \ {\langle color \rangle}$

This function first calculates the $\langle handle \rangle$ for the $\langle coffin \rangle$ as defined by the intersection Updated: 2011-09-02 of $\langle pole_1 \rangle$ and $\langle pole_2 \rangle$. It then marks the position of the $\langle handle \rangle$ on the $\langle coffin \rangle$. The $\langle handle \rangle$ are labelled as part of this process: the location of the $\langle handle \rangle$ and the label are both printed in the $\langle color \rangle$ specified.

\coffin_show_structure:c

 $\verb|\coffin_show_structure:N \coffin_show_structure:N \coffn_show_structure:N \coffn_show_s$

This function shows the structural information about the $\langle coffin \rangle$ in the terminal. The Updated: 2015-08-01 width, height and depth of the typeset material are given, along with the location of all of the poles of the coffin.

> Notice that the poles of a coffin are defined by four values: the x and y co-ordinates of a point that the pole passes through and the x- and y-components of a vector denoting the direction of the pole. It is the ratio between the later, rather than the absolute values, which determines the direction of the pole.

\coffin_log_structure:N \coffin_log_structure:N \coffin\

This function writes the structural information about the $\langle coffin \rangle$ in the log file. See also New: 2014-08-22 \coffin_show_structure: N which displays the result in the terminal.

Updated: 2015-08-01

\coffin_show:N \coffin_show:N \coffin\ \coffin_show:c \coffin_log:N \coffin \

\coffin_log:N \coffin_log:c

Shows full details of poles and contents of the $\langle coffin \rangle$ in the terminal or log file. See \coffin_show_structure: N and \box_show: N to show separately the pole structure and New: 2021-05-11 the contents.

\coffin_log:Nnn

\coffin_log:cnn

 $\verb|\coffin_show:Nnn | coffin_show:Nnn | \langle coffin \rangle | \{\langle intexpr_1 \rangle\} | \{\langle intexpr_2 \rangle\}|$ $\coffin_show:cnn \coffin_log:Nnn \coffin\ \{\langle intexpr_1 \rangle\} \ \{\langle intexpr_2 \rangle\}$

Shows poles and contents of the $\langle coffin \rangle$ in the terminal or log file, showing the first $\langle intexpr_1 \rangle$ items in the coffin, and descending into $\langle intexpr_2 \rangle$ group levels. See \coffin_-New: 2021-05-11 show_structure: N and \box_show: Nnn to show separately the pole structure and the contents.

34.7 Constants and variables

\c_empty_coffin A permanently empty coffin.

\l_tmpa_coffin Scratch coffins for local assignment. These are never used by the kernel code, and so \l_tmpb_coffin are safe for use with any LATFX3-defined function. However, they may be overwritten by $_{\text{New: 2012-06-19}}$ other non-kernel code and so should only be used for short-term storage.

\g_tmpa_coffin Scratch coffins for global assignment. These are never used by the kernel code, and so \g_tmpb_coffin are safe for use with any LATEX3-defined function. However, they may be overwritten by $_{\text{New: 2019-01-24}}$ other non-kernel code and so should only be used for short-term storage.

Chapter 35

The **I3color** package Color support

Color in boxes 35.1

Controlling the color of text in boxes requires a small number of control functions, so that the boxed material uses the color at the point where it is set, rather than where it is used.

\color_group_begin: \color_group_begin: \color_group_end:

New: 2011-09-03

\color_group_end:

Creates a color group: one used to "trap" color settings. This grouping is built in to for example \hbox_set:Nn.

\color_ensure_current: \color_ensure_current:

New: 2011-09-03 Ensures that material inside a box uses the foreground color at the point where the box is set, rather than that in force when the box is used. This function should usually be used within a \color_group_begin: ...\color_group_end: group.

Color models 35.2

A color *model* is a way to represent sets of colors. Different models are particularly suitable for different output methods, e.g. screen or print. Parameter-based models can describe a very large number of unique colors, and have a varying number of axes which define a color space. In contrast, various proprietary models are available which define spot colors (more formally separations).

Core models are used to pass color information to output; these are "native" to 13color. Core models use real numbers in the range [0,1] to represent values. The core models supported here are

- gray Grayscale color, with a single axis running from 0 (fully black) to 1 (fully white)
- rgb Red-green-blue color, with three axes, one for each of the components

cmyk Cyan-magenta-yellow-black color, with four axes, one for each of the components

There are also interface models: these are convenient for users but have to be manipulated before storing/passing to the backend. Interface models are primarily integer-based: see below for more detail. The supported interface models are

- Gray Grayscale color, with a single axis running from 0 (fully black) to 15 (fully white)
- hsb Hue-saturation-brightness color, with three axes, all real values in the range [0, 1] for hue saturation and brightness
- Hsb Hue-saturation-brightness color, with three axes, integer in the range [0, 360] for hue, real values in the range [0, 1] for saturation and brightness
- HSB Hue-saturation-brightness color, with three axes, integers in the range [0, 240] for hue, saturation and brightness
- HTML HTML format representation of RGB color given as a single six-digit hexadecimal number
- RGB Red-green-blue color, with three axes, one for each of the components, values as integers from 0 to 255
- wave Light wavelength, a real number in the range 380 to 780 (nanometres)

All interface models are internally stored as rgb.

Finally, there are a small number of models which are parsed to allow data transfer from xcolor but which should not be used by end-users. These are

- cmy Cyan-magenta-yellow color with three axes, one for each of the components; converted to cmyk
- tHsb "Tuned" hue-saturation-brightness color with three axes, integer in the range [0, 360] for hue, real values in the range [0, 1] for saturation and brightness; converted to rgb using the standard tuning map defined by xcolor
- &spot Spot color tint with one value; treated as a gray tint as spot color data is not available for extraction

To allow parsing of data from xcolor, any leading model up the first: will be discarded; the approach of selecting an internal form for data is not used in I3color.

Additional models may be created to allow mixing of separation colors with each other or with those from other models. See Section 35.9 for more detail of color support for additional models.

When color is selected by model, the $\langle values \rangle$ given are specified as a comma-separated list. The length of the list will therefore be determined by the detail of the model involved.

Color models (and interconversion) are complex, and more details are given in the manual to the \LaTeX 2 ε xcolor package and in the $PostScript\ Language\ Reference\ Manual$, published by Addison–Wesley.

35.3 Color expressions

In addition to allowing specification of color by model and values, |3color also supports color expressions. These are created by combining one or more color names, with the amount of each specified as a percentage. The latter is given between ! symbols in the expression. Thus for example

red!50!green

is a mixture of $50\,\%$ red and $50\,\%$ green. A trailing percentage is interpreted as implicitly followed by white, and so

red!25

specifies $25\,\%$ red mixed with $75\,\%$ white.

Where the models for the mixed colors are different, the model of the first color is used. Thus

red!50!cyan

will result in a color specification using the rgb model, made up of 50% red and 50% of cyan expressed in rgb. This may be important as color model interconversion is not exact.

The one exception to the above is where the first model in an expression is gray. In this case, the order of mixing is "swapped" internally, so that for example

black!50!red

has the same result as

red!50!black

(the predefined colors black and white use the gray model).

Where more than two colors are mixed in an expression, evaluation takes place in a stepwise fashion. Thus in

cyan!50!magenta!10!yellow

the sub-expression

cyan!50!magenta

is first evaluated to give an intermediate color specification, before the second step

<intermediate>!10!yellow

where <intermediate> represents this transitory calculated value.

Within a color expression, . may be used to represent the color active for type setting (the current color). This allows for example

.!50

to mean a mixture of $50\,\%$ of current color with white.

(Color expressions supported here are a subset of those provided by the LATEX 2ε xcolor package. At present, only such features as are clearly useful have been added here.)

35.4 Named colors

Color names are stored in a single namespace, which makes them accessible as part of color expressions. Whilst they are not reserved in a technical sense, the names black, white, red, green, blue, cyan, magenta and yellow have special meaning and should not be redefined. Color names should be made up of letters, numbers and spaces only: other characters are reserved for use in color expressions. In particular, . represents the current color at the start of a color expression.

\[\color_set:nn \color_set:nn \color_set:nn \color_set:nn \color_expression \rangle \ and stores the resulting color specification as the \langle name \rangle. \]
\[\color_set:nnn \color_set:nnnnn \color_set:nnn \color_set:nnn \color_set:nnnn \color_set:nnn \

```
\label{localization} $$  \color_show:n {\langle name \rangle} $$  \color_log:n {\langle name \rangle} $$
```

New: 2021-05-11 Displays the color specification stored in the $\langle name \rangle$ on the terminal or log file.

35.5 Selecting colors

General selection of color is safe when split across pages: a stack is used to ensure that the correct color is re-selected on the new page.

These commands set the current color (.): other more specialised functions such as fill and stroke selectors do not adjust this value.

 $\verb|\color_select:n \color_select:n \{ (color expression) \} }|$

Parses the $\langle color\ expression \rangle$ and then activates the resulting color specification for type-set material.

 $\verb|\color_select:nn \color_select:nn {| (model(s))|} {| (value(s))|} |$

Activates the color specification equivalent to the $\langle model(s) \rangle$ and $\langle value(s) \rangle$ for typeset material.

When this is set to a non-empty value, colors will be converted to the specified model when they are selected. Note that included images and similar are not influenced by this setting.

35.6 Colors for fills and strokes

Colors for drawing operations and so forth are split into strokes and fills (the latter may also be referred to as non-stroke color). The fill color is used for text under normal circumstances. Depending on the backend, stroke color may use a stack, in which case it exhibits the same page breaking behavior as general color. However, dvips/dvisvgm do not support this, and so color will need to be contained within a scope, such as \draw begin:/\draw end:.

\color_fill:n \color_stroke:n \color_fill:n {\langle color expression \rangle}

Parses the (color expression) and then activates the resulting color specification for filling or stroking.

\color_fill:nn \color_stroke:nn

 $\color_fill:nn {\langle model(s) \rangle} {\langle value(s) \rangle}$

Activates the color specification equivalent to the $\langle model(s) \rangle$ and $\langle value(s) \rangle$ for filling or stroking.

color.sc

When using dvips, this PostScript variables hold the stroke color.

35.6.1Coloring math mode material

Coloring math mode material using \color_select:nn(n) has some restrictions and often leads to spacing issues and/or poor input syntax. Avoiding generating \mathord atoms whilst coloring only those parts of the input which are required needs careful handling. The functionality here covers this important use case.

\color_math:nn \color_math:nnn

```
\color_math:nn {\langle color expression \rangle} {\langle content \rangle}
\color_math:nnn {\langle model(s) \rangle} {\langle value(s) \rangle} {\langle content \rangle}
```

New: 2022-01-26 Works as for \color_select:n(n) but applies color only to the math mode \(\content \). The function does not generate a group and the $\langle content \rangle$ therefore retains its math atom states. Sub/superscripts are also properly handled.

\l_color_math_active_tl This list controls which tokens are considered as math active and should therefore be $_{\text{New: 2022-01-26}}$ replaced by their definition during searching for sub/superscripts.

35.7 Multiple color models

When selecting or setting a color with an explicit model, it is possible to give values for more than one model at one time. This is particularly useful where automated conversion between models does not give the desired outcome. To do this, the list of models and list of values are both subdivided using / characters (as for the similar function in xcolor). For example, to save a color with explicit cmyk and rgb values, one could use

```
\color set:nnn { foo } { cmyk / rgb }
  { 0.1 , 0.2 , 0.3 , 0.4 / 0.1, 0.2 , 0.3 }
```

The manually-specified conversion will be used in preference to automated calculation whenever the model(s) listed are used: both in expressions and when a fixed model is active.

Similarly, the same syntax can be applied to directly selecting a color.

```
\color_select:nn { cmyk / rgb }
    { 0.1 , 0.2 , 0.3 , 0.4 / 0.1, 0.2 , 0.3 }
```

Again, this list is used when a fixed model is active: the first entry is used unless there is a fixed model matching one of the other entries.

35.8 Exporting color specifications

The major use of color expressions is in setting typesetting output, but there are other places in which some form of color information is required. These may need data in a different format or using a different model to the internal representation. Thus a set of functions are available to export colors in different formats.

Valid export targets are

- backend Two brace groups: the first containing the model, the second containing space-separated values appropriate for the model; this is the format required by backend functions of expl3
- comma-sep-cmyk Comma-separated cyan-magenta-yellow-black values
- comma-sep-rgb Comma-separated red-green-blue values suitable for use as a PDF annotation color
- \bullet HTML Uppercase two-digit hexa decimal values, expressing a red-green-blue color; the digits are not separated
- space-sep-cmyk Space-separated cyan-magenta-yellow-black values
- space-sep-rgb Space-separated red-green-blue values suitable for use as a PDF annotation color

```
\verb|\color_export:nnN| $$ \{\langle color_export:nnN| \{\langle color_export:nnN| \{\langle tl\rangle\}\} \}
```

Parses the $\langle color\ expression \rangle$ as described earlier, then converts to the $\langle format \rangle$ specified and assigns the data to the $\langle tl \rangle$.

```
\verb|\color_export:nnnN| $$ \{\langle model \rangle\} $$ {\langle value(s) \rangle} $$ {\langle format \rangle} $$ {\langle tl \rangle} $$
```

Expresses the combination of $\langle model \rangle$ and $\langle value(s) \rangle$ in an internal representation, then converts to the $\langle format \rangle$ specified and assigns the data to the $\langle tl \rangle$.

35.9 Creating new color models

Additional color models are required to support specialist workflows, for example those involving separations (see https://helpx.adobe.com/indesign/using/spot-process-colors. html for details of the use of separations in print). Color models may be split into families; for the standard device-based color models (DeviceCMYK, DeviceRGB, DeviceGray), these are synonymous. This is not generally the case: see the PDF reference for more details. (Note that | 3color uses the shorter names cmyk, etc.)

 $\verb|\color_model_new:nnn \color_model_new:nnn {| (model)|} {| (family)|} {| (params)|}$

Creates a new $\langle model \rangle$ which is derived from the color model $\langle family \rangle$. The latter should be one of

- DeviceN
- ICCBased
- Separation

(The $\langle family \rangle$ may be given in mixed case as-in the PDF reference: internally, case of these strings is folded.) Depending on the $\langle family \rangle$, one or more $\langle params \rangle$ are mandatory

For a Separation space, there are three compulsory keys.

- name The name of the Separation, for example the formal name of a spot color ink. Such a $\langle name \rangle$ may contain spaces, etc., which are not permitted in the $\langle model \rangle$.
- alternative-model An alternative device colorspace, one of cmyk, rgb, gray or CIELAB. The three parameter-based models work as described above; see below for details of CIELAB colors.
- alternative-values A comma-separated list of values appropriate to the alternative-model. This information is used by the PDF application if the Separation is not available.

CIELAB color separations are created using the alternative-model = CIELAB setting. These colors must also have an illuminant key, one of a, c, e, d50, d55, d65 or d75. The alternative-values in this case are the three parameters L*, a* and b* of the CIELAB model. Full details of this device-independent color approach are given in the documentation to the colorspace package.

CIELAB colors cannot be converted into other device-dependent color spaces, and as such, mixing can only occur if colors set up using the CIELAB model are also given with an alternative parameter-based model. If that is not the case, I3color will fallback to using black as the colorant in any mixing.

For a DeviceN space, there is one *compulsory* key.

• names The names of the components of the DeviceN space. Each should be either the $\langle name \rangle$ of a Separation model, a process color name (cyan, etc.) or the special name none.

For a ICCBased space, there is one *compulsory* key.

• file The name of the file containing the profile.

35.9.1Color profiles

Color profiles are used to ensure color accuracy by linking to collaboration. Applying a profile can be used to standardise color which is otherwise device-dependence.

 $\verb|\color_profile_apply:nn \color_profile_apply:nn \ \{\langle profile \rangle\} \ \{\langle model \rangle\}|$

New: 2021-02-23 This function applies a $\langle profile \rangle$ to one of the device $\langle models \rangle$. The profile will then apply to all color of the selected $\langle model \rangle$. The $\langle profile \rangle$ should specify an ICC profile file. The $\langle model \rangle$ has to be one the standard device models: cmyk, gray or rgb.

Chapter 36

The **I3pdf** package Core PDF support

36.1 Objects

| \pdf_object_new:n \pdf_object_new:n \{\langle object\}\} | Declares \langle object\ \as a PDF object. The object may be referenced from this point on, and written later using \pdf_object_write:nnn.

| \pdf_object_write:nnn \pdf_object_write:nn \{\langle object\}\ \{\langle type\}\ \{\langle content\}\} \text{Writes the \langle content\}\ \as a content of the \langle object\}. Depending on the \langle type\}\ \text{declared for the object}\ \text{new: 2022-08-23} \text{dict Key-value pairs in the form \langle key}\ \langle value\}\ \text{fstream Two brace groups: \langle file name} \text{ and \langle file content}\}\ \text{stream Two brace groups: \langle file name}\ \text{and \langle file contents}\}\ \text{New: 2021-02-10} \text{Inserts the appropriate information to reference the \langle object\}\ \text{in for example page resource}\ \text{allocation}

```
\pdf_object_unnamed_write:nx
                           New: 2021-02-10
                                Writes the \langle content \rangle as content of an anonymous object. Depending on the \langle type \rangle, the
                                format required for the \langle data \rangle will vary
                               array A space-separated list of values
                                dict Key-value pairs in the form / (key) (value)
                            fstream Two brace groups: \( \alpha ttributes \) (dictionary) \( \rangle \) and \( \lambda file \) name \( \rangle \)
                             stream Two brace groups: \( \alpha ttributes \) (dictionary) \( \rangle \) and \( \stream \) contents \( \rangle \)
   \pdf_object_ref_last: * \pdf_object_ref_last:
                 New: 2021-02-10 Inserts the appropriate information to reference the last \langle object \rangle created. This is partic-
                                ularly useful for anonymous objects.
   \pdf_pageobject_ref:n * \pdf_pagobject_ref:n {\langle pageobject \rangle}
                 New: 2021-02-10 Inserts the appropriate information to reference the \langle pageobject \rangle.
\pdf_object_if_exist_p:n * \pdf_object_if_exist_p:n {\langle object \rangle}
\pdf_object_if_exist:nTF * \pdf_object_if_exist:nTF {\langle object \rangle}
                 New: 2020-05-15 Tests whether an object with name \{\langle object \rangle\} has been defined.
                                            Version
                                36.2
```

```
\pdf_version_compare:NnTF ★ code \}
     New: 2021-02-10
```

Compares the version of the PDF being created with the $\langle version \rangle$ string specified, using the $\langle comparator \rangle$. Either the $\langle true\ code \rangle$ or $\langle false\ code \rangle$ will be left in the output stream.

```
\pdf_version_gset:n
\pdf_version_min_gset:n
```

 $\verb| \df_version_gset:n {| \langle version \rangle |}$

Sets the (version) of the PDF being created. The min version will not alter the output New: 2021-02-10 version unless it is currently lower than the $\langle version \rangle$ requested.

This function may only be used up to the point where the PDF file is initialised. With dvips it sets \pdf_version_major: and \pdf_version_minor: and allows to compare the values with \pdf_version_compare: Nn, but the PDF version itself still has to be set with the command line option -dCompatibilityLevel of ps2pdf.

```
⋆ \pdf_version:
\pdf_version:
\pdf_version_major:
                       Expands to the currently-active PDF version.
\pdf_version_minor: *
          New: 2021-02-10
```

Page (media) size 36.3

New: 2023-01-14 Sets the page size (mediabox) of the PDF being created to the $\langle width \rangle$ and $\langle height \rangle$, both of which are $\langle dimexpr \rangle$.

36.4 Compression

\pdf_uncompress: \pdf_uncompress:

New: 2021-02-10 Disables any compression of the PDF, where possible.

This function may only be used up to the point where the PDF file is initialised.

36.5 **Destinations**

Destinations are the places a link jumped too. Unlike the name may suggest they don't described an exact location in the PDF. Instead a destination contains a reference to a page along with an instruction how to display this page. The normally used "XYZ top left zoom" for example instructs the viewer to show the page with the given zoom and the top left corner at the top left coordinates—which then gives the impression that there is an anchor at this position.

If an instruction takes a coordinate, it is calculated by the following commands relative to the location the command is issued. So to get a specific coordinate one has to move the command to the right place.

 $\verb| \df_destination:nn \df_destination:nn {$\langle name \rangle$} {\langle type \ or \ integer \rangle$}$

New: 2021-01-03 This creates a destination. $\{\langle type\ or\ integer\rangle\}\$ can be one of fit, fith, fith, fith, fitbh, fitbv, fitr, xyz or an integer representing a scale factor in percent. fitr here gives only a lightweight version of /FitR: The backend code defines fitr so that it will with pdfLATEX and LuaLATEX use the coordinates of the surrounding box, with dvips and dvipdfmx it falls back to fit. For full control use \pdf_destination:nnnn.

The keywords match to the PDF names as described in the following tabular.

Keyword	PDF	Remarks
fit	/Fit	Fits the page to the window
fith	/FitH top	Fits the width of the page to the window
fitv	/FitV $left$	Fits the height of the page to the window
fitb	/FitB	Fits the page bounding box to the window
fitbh	/FitBH top	Fits the width of the page bounding box to the window.
fitbv	/FitBV $left$	Fits the height of the page bounding box to the window.
fitr	/FitR left bottom right top	Fits the rectangle specified by the four coordinates to the window (see above for the restrictions)
xyz	/XYZ left top null	Sets a coordinate but doesn't change the zoom.
$\{\langle integer \rangle \}$	/XYZ left top zoom	Sets a coordinate and a zoom meaning $\{\langle integer \rangle\}\%$.

 $\positive \positive \pos$

New: 2021-01-17 This creates a destination with /FitR type with the given dimensions relative to the current location. The destination is in a box of size zero, but it doesn't switch to horizontal mode.

36.6 Deprecated functions

 $\verb| \df_object_new:nn \df_object_new:nn \ \{\langle object\rangle\} \ \{\langle type\rangle\}|$

New: 2021-02-10 Declares $\langle object \rangle$ as a PDF object of $\langle type \rangle$, which should be one of

- array
- dict
- fstream
- stream

The object may be referenced from this point on, and written later using \pdf_object_-write:nn.

Deprecated in favor of \pdf_object_new:n.

 $\verb|\pdf_object_write:nn \pdf_object_write:nn {|\langle object \rangle| } {|\langle content \rangle|}$

\pdf_object_write:nx

Writes the $\langle content \rangle$ as content of the $\langle object \rangle$. Depending on the $\langle type \rangle$ declared for the New: 2021-02-10 object, the format required for the $\langle data \rangle$ will vary

array A space-separated list of values

dict Key-value pairs in the form / (key) (value)

fstream Two brace groups: \(\file \ name \rangle \) and \(\file \ content \rangle \)

stream Two brace groups: $\langle attributes\ (dictionary)\rangle$ and $\langle stream\ contents\rangle$

Deprecated in favor of \pdf_object_write:nnn.

Part VII Additions and removals

Chapter 37

The **I3candidates** package Experimental additions to **I3kernel**

37.1 Important notice

This module provides a space in which functions can be added to l3kernel (expl3) while still being experimental.

As such, the functions here may not remain in their current form, or indeed at all, in I3kernel in the future.

In contrast to the material in [3experimental, the functions here are all *small* additions to the kernel. We encourage programmers to test them out and report back on the LaTeX-L mailing list.

Thus, if you intend to use any of these functions from the candidate module in a public package offered to others for productive use (e.g., being placed on CTAN) please consider the following points carefully:

- Be prepared that your public packages might require updating when such functions are being finalized.
- Consider informing us that you use a particular function in your public package, e.g., by discussing this on the LaTeX-L mailing list. This way it becomes easier to coordinate any updates necessary without issues for the users of your package.
- Discussing and understanding use cases for a particular addition or concept also helps to ensure that we provide the right interfaces in the final version so please give us feedback if you consider a certain candidate function useful (or not).

We only add functions in this space if we consider them being serious candidates for a final inclusion into the kernel. However, real use sometimes leads to better ideas, so functions from this module are **not necessarily stable** and we may have to adjust them!

37.2 Additions to 13box

\box_clip:N \box_clip:c \box_gclip:N \box_gclip:c \box_clip:N \langle box \rangle

Updated: 2019-01-23

Clips the $\langle box \rangle$ in the output so that only material inside the bounding box is displayed in the output. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the clipping is applied.

These functions require the LATEX3 native drivers: they do not work with the LATEX 2ε graphics drivers!

TeXhackers note: Clipping is implemented by the driver, and as such the full content of the box is placed in the output file. Thus clipping does not remove any information from the raw output, and hidden material can therefore be viewed by direct examination of the file.

\box_set_trim:Nnnnn \box_set_trim:cnnnn \box_gset_trim:Nnnnn \box_gset_trim:cnnnn $\box_set_trim:Nnnnn \ \langle box \rangle \ \{\langle left \rangle\} \ \{\langle bottom \rangle\} \ \{\langle right \rangle\} \ \{\langle top \rangle\}$

New: 2019-01-23

Adjusts the bounding box of the $\langle box \rangle$ $\langle left \rangle$ is removed from the left-hand edge of the bounding box, \(\langle right\rangle\) from the right-hand edge and so fourth. All adjustments are $\langle dimension \ expressions \rangle$. Material outside of the bounding box is still displayed in the output unless $\box_{clip:N}$ is subsequently applied. The updated $\langle box \rangle$ is an abox, irrespective of the nature of the $\langle box \rangle$ before the trim operation is applied. The behavior of the operation where the trims requested is greater than the size of the box is undefined.

\box_set_viewport:Nnnnn \box_set_viewport:cnnn \box_gset_viewport:Nnnnn \box_gset_viewport:cnnnn

 $\box_set_viewport:Nnnn \ \langle box \rangle \ \{\langle 11x \rangle\} \ \{\langle 11y \rangle\} \ \{\langle urx \rangle\} \ \{\langle ury \rangle\}$

Adjusts the bounding box of the $\langle box \rangle$ such that it has lower-left co-ordinates ($\langle llx \rangle$, $\langle lly \rangle$) and upper-right co-ordinates ($\langle urx \rangle$, $\langle ury \rangle$). All four co-ordinate positions are $\langle dimension$ expressions). Material outside of the bounding box is still displayed in the output unless \box clip: N is subsequently applied. The updated $\langle box \rangle$ is an hbox, irrespective of the nature of the $\langle box \rangle$ before the viewport operation is applied.

Additions to **I3expan** 37.3

\exp_args_generate:n \exp_args_generate:n {\variant argument specifiers\}}

New: 2019-01-23

New: 2018-04-04 Defines \exp_args: N(variant) functions for each (variant) given in the comma list Updated: 2019-02-08 {\langle variant argument specifiers \rangle}. Each \langle variant \rangle should consist of the letters N, c, n, V, v, o, f, e, x, p and the resulting function is protected if the letter x appears in the $\langle variant \rangle$. This is only useful for cases where $\cs_generate_variant:Nn$ is not applicable.

37.4 Additions to 13fp

\fp_if_nan_p:n

 $\star \fp_{if_nan:n} \{\langle fpexpr \rangle\}$

 $\frac{\mathbf{fp_if_nan:n}\underline{TF} \star}{\mathbf{Evaluates}}$ Evaluates the $\langle fpexpr \rangle$ and tests whether the result is exactly nan. The test returns false New: 2019-08-25 for any other result, even a tuple containing nan.

Additions to 13file 37.5

\iow_allow_break: \iow_allow_break:

New: 2018-12-29 In the first argument of \iow wrap:nnnN (for instance in messages), inserts a break-point that allows a line break. In other words this is a zero-width breaking space.

\ior_get_term:nN \ior_str_get_term:nN

\ior_get_term:nN \(\rho prompt \rangle \) \(\taken list variable \)

Function that reads one or more lines (until an equal number of left and right braces New: 2019-03-23 are found) from the terminal and stores the result locally in the $\langle token\ list \rangle$ variable. Tokenization occurs as described for \ior_get:NN or \ior_str_get:NN, respectively. When the $\langle prompt \rangle$ is empty, T_FX will wait for input without any other indication: typically the programmer will have provided a suitable text using e.g. \iow term:n. Where the $\langle prompt \rangle$ is given, it will appear in the terminal followed by an =, e.g.

prompt=

 $\verb|\ior_shell_open:Nn \ior_shell_open:Nn \| \{stream\} \| \{shell \| command\} \}|$

New: 2019-05-08 Opens the pseudo-file created by the output of the $\langle shell\ command \rangle$ for reading using $\langle stream \rangle$ as the control sequence for access. If the $\langle stream \rangle$ was already open it is closed before the new operation begins. The \(\stream \) is available for access immediately and will remain allocated to (shell command) until a \ior_close:N instruction is given or the TeX run ends. If piped system calls are disabled an error is raised.

For details of handling of the \(shell \) command \(\), see \sys_get_shell:nnNTF.

37.6 Additions to 13flag

New: 2018-04-02 Ensures the $\langle flag \rangle$ is raised by making its height at least 1, locally.

Additions to **I3intarray**

\intarray_gset_rand:Nn

\intarray_gset_rand:cn

New: 2018-05-05

 $\displaystyle \operatorname{Intarray_gset_rand:Nnn \ (intarray_gset_rand:Nnn \ (intarray\ var) \ \{\langle minimum \rangle\} \ \{\langle maximum \rangle\}}$ \intarray_gset_rand:cnn \intarray_gset_rand:Nn \(intarray \ var \) {\((maximum \)}

> Evaluates the integer expressions $\langle minimum \rangle$ and $\langle maximum \rangle$ then sets each entry (independently) of the $\langle integer\ array\ variable \rangle$ to a pseudo-random number between the two (with bounds included). If the absolute value of either bound is bigger than $2^{30} - 1$, an error occurs. Entries are generated in the same way as repeated calls to \int_rand:nn or \int rand:n respectively, in particular for the second function the $\langle minimum \rangle$ is 1. Assignments are always global. This is not available in older versions of X¬T¬X.

\intarray_to_clist:N ☆ \intarray_to_clist:N ⟨intarray var⟩

New: 2018-05-04 Converts the $\langle intarray \rangle$ to integer denotations separated by commas. All tokens have category code other. If the $\langle intarray \rangle$ has no entry the result is empty; otherwise the result has one fewer comma than the number of items.

Additions to 13msg 37.8

```
\msg_log_eval:Nn
```

 $\label{lem:msg_show_eval:Nn and function} $$\max_{\text{expression}} $$$

Shows or logs the \(\left(expression \right) \) (turned into a string), an equal sign, and the result of New: 2017-12-04 applying the $\langle function \rangle$ to the $\{\langle expression \rangle\}$ (with f-expansion). For instance, if the $\langle function \rangle$ is $\forall int_eval:n$ and the $\langle expression \rangle$ is 1+2 then this logs > 1+2=3.

```
\msg_show_item:n
                               * \seq_map_function:NN \langle seq \msg_show_item:n
                               \star \prop_map_function:NN \langle prop \rangle \msg_show_item:nn
\msg_show_item_unbraced:n
\msg_show_item:nn
\msg_show_item_unbraced:nn *
```

Used in the text of messages for \msg_show:nnxxxx to show or log a list of items or key-value pairs. The one-argument functions are used for sequences, clist or token lists and the others for property lists. These functions turn their arguments to strings.

37.9 Additions to 13prg

\bool_set_inverse:c \bool_gset_inverse:N \bool_gset_inverse:c current value.

\bool_set_inverse:N \bool_set_inverse:N \boolean \}

Toggles the $\langle boolean \rangle$ from true to false and conversely: sets it to the inverse of its

New: 2018-05-10

```
\bool_case_true:n
                                  * \bool_case_true:nTF
\bool_case_true:nTF
\bool_case_false:n
                                            \{\langle boolexpr case_1 \rangle\} \{\langle code case_1 \rangle\}
                                            \{\langle boolexpr case_2 \rangle\} \{\langle code case_2 \rangle\}
\bool_case_false:nTF
                 New: 2019-02-10
                                            \{\langle boolexpr\ case_n \rangle\}\ \{\langle code\ case_n \rangle\}
                                         {\langle true code \rangle}
                                         \{\langle false\ code \rangle\}
```

Evaluates in turn each of the (boolean expression cases) until the first one that evaluates to true or to false, for \bool_case_true:n and \bool_case_false:n, respectively. The $\langle code \rangle$ associated to this first case is left in the input stream, followed by the $\langle true \rangle$ code, and other cases are discarded. If none of the cases match then only the $\langle false \rangle$ code) is inserted. The functions \bool_case_true:n and \bool_case_false:n, which do nothing if there is no match, are also available. For example

```
\bool_case_true:nF
  {
    { \dim_compare_p:n { \l__mypkg_wd_dim <= 10pt } }
        { Fits }
    { \int_compare_p:n { \l_mypkg_total_int >= 10 } }
        { Many }
    { \l_mypkg_special_bool }
        { Special }
  { No idea! }
```

leaves "Fits" or "Many" or "Special" or "No idea!" in the input stream, in a way similar to some other language's "if ... elseif ... elseif ... else ... ".

37.10 Additions to 13prop

```
\prop_rand_key_value:N * \prop_rand_key_value:N \langle property list \rangle
\prop_rand_key_value:c *
```

Selects a pseudo-random key-value pair from the $\langle property | list \rangle$ and returns $\{\langle key \rangle\}$ and New: 2016-12-06 $\{\langle value \rangle\}$. If the $\langle property \ list \rangle$ is empty the result is empty. This is not available in older versions of X₇T_FX.

> TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the \(\value \rangle \) does not expand further when appearing in an x-type or e-type argument expansion.

37.11 Additions to **|3seq**|

```
\label{lem:nnn} $$ \eq_mapthread_function:NNN $$ $$ \eq_mapthread_function:NNN $$ $$ \eq_napthread_function:(NcN|cNN|ccN) $$ $$
```

Applies $\langle function \rangle$ to every pair of items $\langle seq_1\text{-}item \rangle - \langle seq_2\text{-}item \rangle$ from the two sequences, returning items from both sequences from left to right. The $\langle function \rangle$ receives two n-type arguments for each iteration. The mapping terminates when the end of either sequence is reached (i.e. whichever sequence has fewer items determines how many iterations occur).

\seq_set_filter:NNn
\seq_gset_filter:NNn

```
\verb|\seq_set_filter:NNn| \langle sequence_1 \rangle | \langle sequence_2 \rangle | \{\langle inline| boolexpr \rangle \}|
```

Evaluates the $\langle inline\ boolexpr \rangle$ for every $\langle item \rangle$ stored within the $\langle sequence_2 \rangle$. The $\langle inline\ boolexpr \rangle$ receives the $\langle item \rangle$ as #1. The sequence of all $\langle items \rangle$ for which the $\langle inline\ boolexpr \rangle$ evaluated to true is assigned to $\langle sequence_1 \rangle$.

TeXhackers note: Contrarily to other mapping functions, \seq_map_break: cannot be used in this function, and would lead to low-level TeX errors.

New: 2018-04-06

Sets the $\langle seq\ var \rangle$ equal to a sequence whose items are obtained by x-expanding $\langle loop\ code \rangle$ $\langle function \rangle$. This expansion must result in successive calls to the $\langle function \rangle$ with no nonexpandable tokens in between. More precisely the $\langle function \rangle$ is replaced by a wrapper function that inserts the appropriate separators between items in the sequence. The $\langle loop\ code \rangle$ must be expandable; it can be for example $\texttt{tl_map_function:NN}\ \langle tl\ var \rangle$ or $\texttt{clist_map_function:NN}\ \{\langle clist \rangle\}$ or $\texttt{int_step_function:nnnN}\ \{\langle initial\ value \rangle\}$ $\{\langle step \rangle\}\ \{\langle final\ value \rangle\}$.

New: 2018-04-06

Sets the $\langle seq\ var \rangle$ equal to a sequence whose items are obtained by x-expanding $\langle loop\ code \rangle$ applied to a $\langle function \rangle$ derived from the $\langle inline\ code \rangle$. A $\langle function \rangle$ is defined, that takes one argument, x-expands the $\langle inline\ code \rangle$ with that argument as #1, then adds appropriate separators to turn the result into an item of the sequence. The x-expansion of $\langle loop\ code \rangle\ \langle function \rangle\$ must result in successive calls to the $\langle function \rangle\$ with no nonexpandable tokens in between. The $\langle loop\ code \rangle\$ must be expandable; it can be for example \t1_map_function:NN $\langle tl\ var \rangle$ or \clist_map_function:nN $\{\langle clist \rangle\}$ or \int_step_function:nnnN $\{\langle initial\ value \rangle\}$ $\{\langle step \rangle\}$ $\{\langle final\ value \rangle\}$, but not the analogous "inline" mappings.

```
\seq_set_item:Nnn
\seq_set_item:cnn
\seq_set_item:Nnn<u>TF</u>
\seq_set_item:cnnTF
\seq_gset_item:Nnn
\seq_gset_item:cnn
\seq_gset_item:Nnn<u>TF</u>
\seq_gset_item:cnnTF
         New: 2021-04-29
```

```
\seq_set_item:Nnn \langle seq var \rangle \{\langle intexpr \rangle\} \{\langle item \rangle\}
\ \left( item : NnnTF \ (seq var) \ {(intexpr)} \ {(item)} \ {(true code)} \ {(false code)}
```

Removes the item of $\langle sequence \rangle$ at the position given by evaluating the $\langle integer \rangle$ expression and replaces it by $\langle item \rangle$. Items are indexed from 1 on the left/top of the $\langle sequence \rangle$, or from -1 on the right/bottom. If the $\langle integer\ expression \rangle$ is zero or is larger (in absolute value) than the number of items in the sequence, the $\langle sequence \rangle$ is not modified. In these cases, \seq_set_item:Nnn raises an error while \seq_set_item:NnnTF runs the $\langle false\ code \rangle$. In cases where the assignment was successful, $\langle true\ code \rangle$ is run afterwards.

```
\seq_pop_item:NnN
\seq_pop_item:cnN
\seq_pop_item:NnN<u>TF</u>
\seq_pop_item:cnNTF
\seq_gpop_item:NnN
```

```
\ensuremath{\verb|Seq_pop_item:NnN|} \ensuremath{ \langle seq \ var \rangle } \ensuremath{ \langle \langle intexpr \rangle \} } \ensuremath{ \langle \langle t1| \ var \rangle \} }
\label{lem:nntf} $$ \left( \operatorname{seq} \operatorname{var} \right) \left( \left( \operatorname{intexpr} \right) \right) \left( \left( \operatorname{true} \operatorname{code} \right) \right) \left( \left( \operatorname{false} \operatorname{code} \right) \right) \right) $$
```

\seq_gpop_item:cnN \seq_gpop_item:NnNTF \seq_gpop_item:cnNTF

Removes the $\langle item \rangle$ at position $\langle integer\ expression \rangle$ in the $\langle sequence \rangle$, and places it in the $\langle token\ list\ variable \rangle$. Items are indexed from 1 on the left/top of the $\langle sequence \rangle$, or from -1 on the right/bottom. If the position is zero or is larger (in absolute value) than the number of items in the sequence, the $\langle seq \ var \rangle$ is not modified, the $\langle token \ list \rangle$ is set to the special marker \q_no_value, and the \(false \code \) is left in the input stream; otherwise the $\langle true\ code \rangle$ is. The $\langle token\ list \rangle$ assignment is local while the $\langle sequence \rangle$ is assigned locally for pop or globally for gpop functions.

New: 2021-04-28

New: 2018-05-02

37.12 Additions to **13sys**

\c_sys_engine_version_str The version string of the current engine, in the same form as given in the banner issued when running a job. For pdfTFX and LuaTFX this is of the form

```
\langle major \rangle. \langle minor \rangle. \langle revision \rangle
```

For $X_{\overline{A}}T_{\overline{E}}X$, the form is

 $\langle major \rangle. \langle minor \rangle$

For pT_FX and upT_FX, only releases since T_FX Live 2018 make the data available, and the form is more complex, as it comprises the pTFX version, the upTFX version and the e-pT_FX version.

```
p\langle major \rangle. \langle minor \rangle. \langle revision \rangle - u\langle major \rangle. \langle minor \rangle - \langle epTeX \rangle
```

where the u part is only present for upT_FX.

```
\sys_if_rand_exist_p: * \sys_if_rand_exist_p:
\verb|\sys_if_rand_exist: $\underline{\mathit{TF}} \; \star \; \sys_if_rand_exist: $\mathsf{TF} \; \{\langle \mathit{true} \; \mathit{code} \rangle \} \; \{\langle \mathit{false} \; \mathit{code} \rangle \} \; \\
```

New: 2017-05-27 Tests if the engine has a pseudo-random number generator. Currently this is the case in pdfTfX, LuaTfX, pTfX, upTfX and recent releases of XfTfX.

Additions to 13tl 37.13

```
\tl_range_braced:Nnn
\tl_range_braced:cnn
\tl_range_braced:nnn
\tl_range_unbraced:cnn
\tl_range_unbraced:nnn *
```

```
\star \t1_range_braced:Nnn \langle tl var \rangle \{\langle start index \rangle\} \{\langle end index \rangle\}
                   \star \t1_range\_braced:nnn {\langle token \ list \rangle} {\langle start \ index \rangle} {\langle end \ index \rangle}
```

Leaves in the input stream the items from the $\langle start\ index \rangle$ to the $\langle end\ index \rangle$ inclusive, using the same indexing as \t1_range:nnn. Spaces are ignored. Regardless of whether New: 2017-07-15 items appear with or without braces in the \(token list\), the \tl_range_braced:nnn function wraps each item in braces, while \tl_range_unbraced:nnn does not (overall it removes an outer set of braces). For instance,

```
\iow_term:x { \tl_range_braced:nnn { abcd~{e{}}f } { 2 } { 5 } }
\iow_term:x { \tl_range_braced:nnn { abcd~{e{}}f } { -4 } { -1 } }
\iow_term:x { \tl_range_braced:nnn { abcd~{e{}}f } { -2 } { -1 } }
\iow_term:x { \tl_range_braced:nnn { abcd~{e{}}f } { 0 } { -1 } }
```

prints $\{b\}\{c\}\{d\}\{e\{\}\}, \{c\}\{d\}\{e\{\}\}\{f\}, \{e\{\}\}\{f\}, and an empty line to the terminal,$ while

```
\iow_term:x { \tl_range_unbraced:nnn { abcd~{e{}}f } { 2 } { 5 } }
\label{limits} $$ \operatorname{tl\_range\_unbraced:nnn { abcd~{e{}}f } { -4 } { -1 } } $$
\label{limits} $$ \operatorname{tl\_range\_unbraced:nnn { abcd~{e{}}f } { -2 } { -1 } }
\iow_term:x { \tl_range_unbraced:nnn { abcd~{e{}}f } { 0 } { -1 } }
```

prints bcde{}, cde{}f, e{}f, and an empty line to the terminal. Because braces are removed, the result of \tl_range_unbraced:nnn may have a different number of items as for \tl_range:nnn or \tl_range_braced:nnn. In cases where preserving spaces is important, consider the slower function \tl_range:nnn.

TEXhackers note: The result is returned within the \unexpanded primitive (\exp_not:n), which means that the \(\lambda item\rangle\) does not expand further when appearing in an x-type or e-type argument expansion.

\tl_build_begin:N \tl_build_gbegin:N \tl_build_begin:N \langle tl var \rangle

Clears the $\langle tl \ var \rangle$ and sets it up to support other \t_{build} ... functions, which allow New: 2018-04-01 accumulating large numbers of tokens piece by piece much more efficiently than standard | 13tl functions. Until \t1_build_end:N \(tl var \) is called, applying any function from | 13tl other than \tl_build_... will lead to incorrect results. The begin and gbegin functions must be used for local and global $\langle tl \ var \rangle$ respectively.

\tl_build_clear:N \tl_build_gclear:N

\tl_build_clear:N \(t1 \) var \)

Clears the $\langle tl \ var \rangle$ and sets it up to support other $\t \$ build... functions. The clear New: 2018-04-01 and gclear functions must be used for local and global $\langle tl \ var \rangle$ respectively.

\tl_build_put_left:Nn \tl_build_put_left:Nx \tl_build_gput_left:Nn \tl_build_gput_left:Nx \tl_build_put_right:Nn \tl_build_put_right:Nx \tl_build_gput_right:Nn \tl_build_gput_right:Nx

 $\tilde{tl_build_put_left:Nn} \langle tl\ var \rangle \{\langle tokens \rangle\}$ $\t!$ build_put_right:Nn $\langle tl \ var \rangle \ \{\langle tokens \rangle\}$

Adds $\langle tokens \rangle$ to the left or right side of the current contents of $\langle tl \, var \rangle$. The $\langle tl \, var \rangle$ must have been set up with \tl_build_begin:N or \tl_build_gbegin:N. The put and gput functions must be used for local and global $\langle tl \ var \rangle$ respectively. The right functions are about twice faster than the left functions.

New: 2018-04-01

 $\tilde{t}l_build_get:NN \tl_build_get:NN \tl_var_1 \ \langle tl\ var_2 \rangle$

New: 2018-04-01 Stores the contents of the $\langle tl \ var_1 \rangle$ in the $\langle tl \ var_2 \rangle$. The $\langle tl \ var_1 \rangle$ must have been set up with $\t1_build_begin:N$ or $\t1_build_gbegin:N$. The $\langle tl\ var_2\rangle$ is a "normal" token list variable, assigned locally using \tl_set:Nn.

\tl_build_end:N \tl_build_gend:N

 $\t!$ \tl_build_end:N $\langle tl \ var \rangle$

Gets the contents of $\langle tl \ var \rangle$ and stores that into the $\langle tl \ var \rangle$ using \tl_set:Nn or New: 2018-04-01 \tl_gset:Nn. The $\langle tl \ var \rangle$ must have been set up with \tl_build_begin:N or \tl_build_gbegin:N. The end and gend functions must be used for local and global $\langle tl \ var \rangle$ respectively. These functions completely remove the setup code that enabled $\langle tl \ var \rangle$ to be used for other \tl_build_... functions.

37.14Additions to **I3token**

\c_catcode_active_space_tl Token list containing one character with category code 13, ("active"), and character code New: 2017-08-07 32 (space).

New: 2018-09-23

Collects and removes tokens from the input stream until finding a token that does not match the $\langle test\ token \rangle$ (as defined by the test \token_if_eq_catcode:NNTF or \token_if_eq_meaning:NNTF). The collected tokens are passed to the $\langle inline\ code \rangle$ as #1. When begin-group or end-group tokens (usually { or }) are collected they are replaced by implicit \c_group_begin_token and \c_group_end_token, and when spaces (including \c_space_token) are collected they are replaced by explicit spaces.

For example the following code prints "Hello" to the terminal and leave ", world!" in the input stream.

```
\peek_catcode_collect_inline: Nn A { \iow_term:n {#1} } Hello, ~world!
```

Another example is that the following code tests if the next token is *, ignoring intervening spaces, but putting them back using #1 if there is no *.

```
\peek_meaning_collect_inline:Nn \c_space_token
{ \peek_charcode:NTF * { star } { no~star #1 } }
```

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