# Output-Oblivious Stochastic Chemical Reaction **Networks**

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### Acknowledgements

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Full paper in Proceedings of OPODIS 2018.

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That's it! (Almost ...)

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- "Programming Language" for molecular programming and DNA computing. Implementation via Strand Displacement Systems.
- Mathematically fun ...



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https://www.memesmonkey.com/topic/chris+pratt

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Computes the function f(n) = 2n - 1.

Function computation is *stable*: will always produce correct answer, no matter the order of computation.

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Essentially, piecewise affine functions. E.g.,

$$f(n) = \begin{cases} 3, & n \le 5 \\ 3n - 2, & n \text{ even}, \ n \ge 5 \\ 2n + 6, & otherwise. \end{cases}$$

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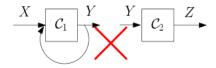
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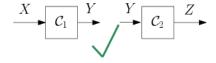
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Leads to incorrect results :(.



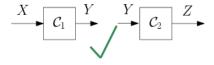
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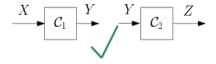
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Call these output-oblivious functions.

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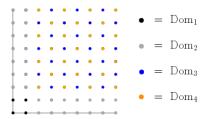
One class of output-oblivious functions: *Grid affine functions*.

These are piecewise affine functions defined on different "grids" (which can be zero-, one-, or two-dimensional).

#### Grid-Affine functions

#### Example:

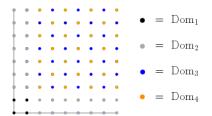
$$f(\mathbf{n}) = \begin{cases} 2, & \mathbf{n} \in \mathsf{Dom}_1 \cup \mathsf{Dom}_2 \\ 2n_1 + 3n_2 + 1, & \mathbf{n} \in \mathsf{Dom}_3 \\ 2n_1 + 3n_2, & \mathbf{n} \in \mathsf{Dom}_4. \end{cases}$$



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Intuition: You can "track" on which grid the input lies, since grids are (roughly) a constant distance from one another.

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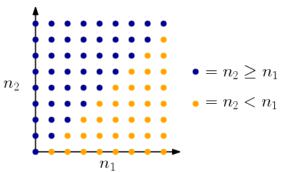
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But not grid-affine!



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Enter *fissure functions*: These are "min-like"; they agree with the min function everywhere except on "fissure lines."

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$$f(\mathbf{n}) = \begin{cases} 3n_1 + 2n_2 + 1, & \mathbf{n} \in \mathsf{Dom}_A, \\ 5n_1, & \mathbf{n} \in \mathsf{Dom}_0, \\ 2n_1 + 3n_2 + 2, & \mathbf{n} \in \mathsf{Dom}_B. \end{cases} \bullet = \mathsf{Dom}_A$$

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Here f agrees with min $\{2n_1 + 3n_2 + 2, 3n_1 + 2n_2 + 1\}$ , except on the "fissure line" Dom<sub>0</sub>.

#### The main result!

If  $f_1$  and  $f_2$  are output-oblivious, then min $\{f_1 + f_2\}$  is output oblivious:

 $C_1$  computes  $f_1$  and produces  $Y_1$ .  $C_2$  computes  $f_2$  and produces  $Y_2$ .

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Therefore, since it's illegal to conclude a theory talk without stating a confusing theorem.

#### Theorem

An increasing, semilinear function  $f: \mathbb{N}^2 \to \mathbb{N}$  is stably computable by an output-oblivious CRN if and only if it is grid-affine or the minimum of finitely many fissure functions.