



Let  $C$  be a linear code of block length  $n$  over the finite field  $\mathbb{F}_q$ . Then the set

$$C^\perp := \{d \in \mathbb{F}_q^n \mid c \cdot d = 0 \text{ for all } c \in C\}$$

is the *dual code* of  $C$ . Here,  $c \cdot d$  denotes either the standard dot product or the Hermitian dot product.

This definition is reminiscent of orthogonal complements of <http://planetmath.org/node/539> dimensional vector spaces over the real or complex numbers. Indeed,  $C^\perp$  is also a linear code and it is true that if  $k$  is the <http://planetmath.org/node/539> dimension of  $C$ , then the dimension of  $C^\perp$  is  $n - k$ . It is, however, **not** necessarily true that  $C \cap C^\perp = \{0\}$ . For example, if  $C$  is the binary code of block length 2 <http://planetmath.org/node/806> spanned by the codeword  $(1, 1)$  then  $(1, 1) \cdot (1, 1) = 0$ , that is,  $(1, 1) \in C^\perp$ . In fact,  $C$  equals  $C^\perp$  in this case. In general, if  $C = C^\perp$ ,  $C$  is called *self-dual*. Furthermore  $C$  is called *self-orthogonal* if  $C \subseteq C^\perp$ .

Famous examples of self-dual codes are the extended binary Hamming code of block length 8 and the extended binary Golay code of block length 24.