CS1632: Static Analysis, Part 3

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Symbolic Model Checking

- Model checking can be categorized into:
 - 1. Enumerative model checking
 - What we learned in the last chapter
 - Hard to escape state explosion
 - 2. Symbolic model checking
 - What we will learn in this chapter
 - Model checking using symbolic execution
 - Can fundamentally solve the state explosion problem

Symbolic Execution

- Symbolic value: Math symbol that stands for an input value
 - A, B, ..., X, Y, Z
- Symbolic expression: An expression using symbolic values
 - A + 1, A * B, ...
- Symbolic execution: Assigning symbolic expressions instead of actual values to variables during execution
 - Instead of assigning values like: x = 10, y = 20
 - Assign symbolic expressions like: x = A + 1, y = A * B
- Example: let's say x == A+1 and y == A+2 at a given point
 - We can prove that assert (x < y) always passes. Why? assert (x < y) \equiv assert (A+1 < A+2) \equiv assert (0 < 1) for all possible input values A

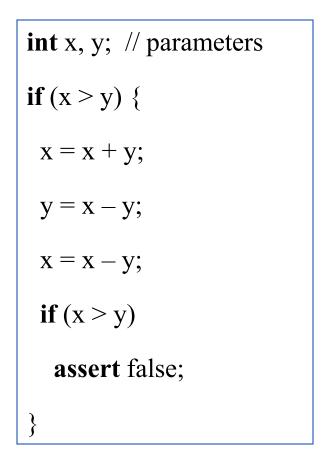
Notation We Will Use

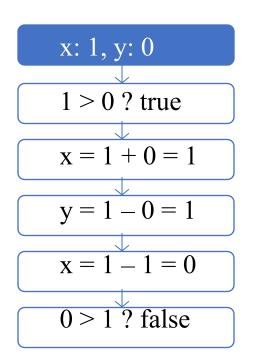
- Program variables: lower case
 - •int x, y, z;
- Symbolic values: UPPER CASE
 - A, B, ..., X, Y, Z
 - Represent input values that are given to the code
 - Can be values from user input
 - Can be values from command line arguments
 - Can be values passed into method parameters

Example: Enumerative Model Checking

Code that swaps 2 integers

Execution Path for x=1, y=0



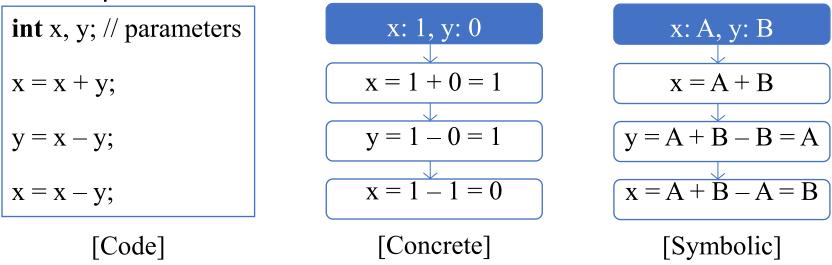


- Must do this for all values of x and y.
- But is that how a human would do it?

Symbolic Model Checking

- Trace through a program like a human being would
- In a symbolic execution:
 - Inputs are symbolic values instead of concrete data values

• Example:



Symbolic execution proves that the swap works for all A and B!

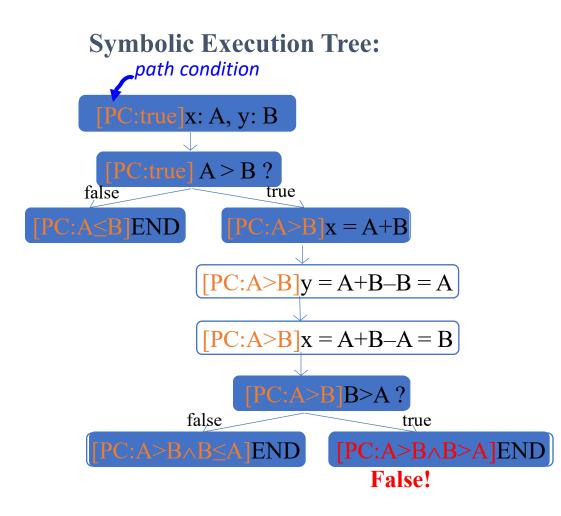
Symbolic Model Checking

- What if there is path divergence?
 - if statement
 - for loop
 - while loop
- For each path, build a Path Condition (PC)
 - Condition on symbolic values (the As and the Bs)

Example: Symbolic Execution

Code that swaps 2 integers:

```
int x, y;
if (x > y) {
 x = x + y;
 y = x - y;
 x = x - y;
 if (x > y)
  assert false;
```



Is the Path Condition Feasible?

Each path condition is checked using a constraint solver

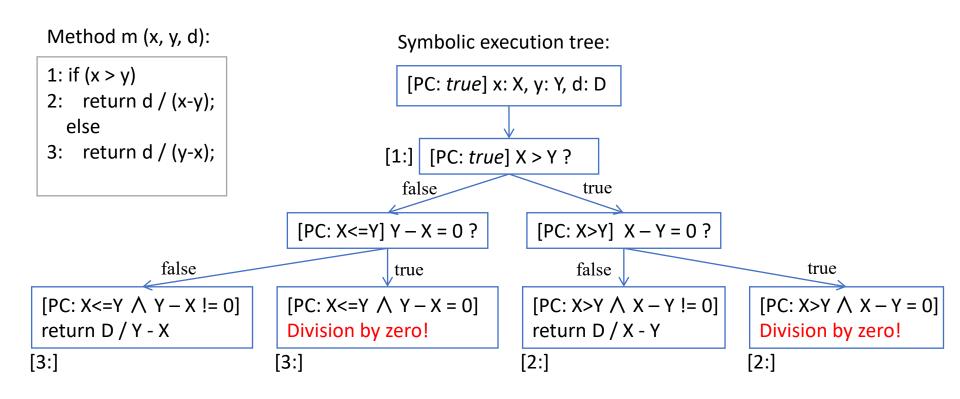


- If path is infeasible, does not continue down that path
 - Hence, assert false is never reached

Symbolic Model Checking Uses

- Prove a program correct
 - Much less state explosion than enumerative checking
- Generate test vectors that exercise each path
 - For the purposes of increasing test coverage of code
 - The output sets of the constraint solver are the test vectors
- Generate program invariants
 - Invariants enhance programmer's understanding of code
 - The path conditions themselves *are* the invariants

Generating Test Cases out of Path Conditions



Solve path conditions \rightarrow test inputs

Auto-generated JUnit Tests

- Constraint solver returns values satisfying each Path Condition
- Can achieve full code coverage by exercising each path!

```
@Test public void t1() { Pass \checkmark PC: X<=Y \land Y - X!= 0 \Leftrightarrow X < Y
  m(10, 20, 0);
                                            (X=10, Y=20, D=0) \in PC
@Test public void t2() { Fail X PC: X <= Y \land Y - X = 0 \Leftrightarrow X = Y
                                            (X=10, Y=10, D=0) \in PC
  m(10, 10, 0);
                                  Pass \checkmark PC: X>Y \land X - Y !=0 \Leftrightarrow X > Y
@Test public void t3() {
                                            (X=20, Y=10, D=0) \in PC
  m(20, 10, 0);
                                  Skip \bigcirc PC: X>Y \bigwedge X – Y =0 \Leftrightarrow { }
```

Symbolic Model Checking Challenges

- Symbolic model checking does have challenges
 - ... Or everyone would be using symbolic model checking
- Some examples are:
 - Complex math constraints
 - Complex data structures

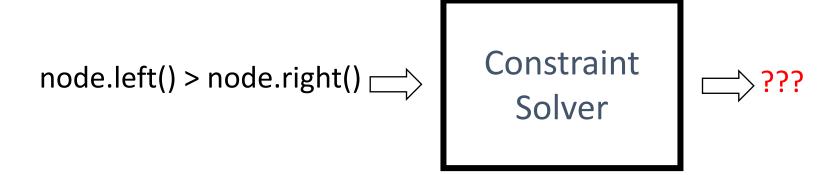
Challenges: Complex Math Constraints

Constraint solvers are not particularly good at math

- Constraint solving is NP-Complete in general
 - Research resulted in some subproblems being proved polynomial
 - State-of-the-art solvers have trouble when constraints are non-linear

Challenges: Complex Data Structures

• Say "node" is a tree node with two children, left() and right()



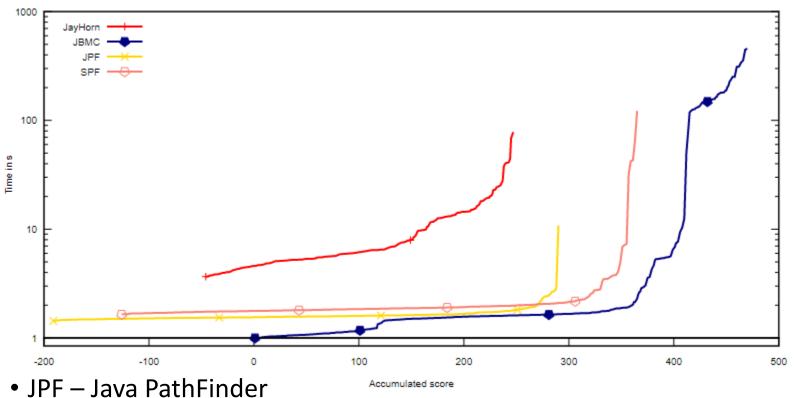
- In order to solve above constraint, solver must know:
 - What the tree data structure looks like
 - Where in the tree "node" is
- Solvers understand some data structures, but not many

The Best of Both Worlds

- Symbolic Model Checking (Symbolic Execution)
 - + Much less state explosion
 - Hard time dealing with loops, math, data structures
- Enumerative Model Checking (Concrete Execution)
 - Serious state explosion
 - + No problems with complex math or data structures
- The best of both worlds: Concolic Execution
 - Concolic = Concrete + Symbolic
 - Symbolic Java Path Finder does exactly this!
 - Chooses between the two depending on the method

Model Checking is Getting Better Every Year

https://sv-comp.sosy-lab.org/2019/results/results-verified/



- SPF Symbolic Java PathFinder (JPF with symbolic execution)
- JBMC Java Bounded Model Checker (2018 newcomer)

References

- Ranjit Jhala and Rupak Majumdar. 2009. "Software model checking". ACM Computing Surveys: https://people.mpi-sws.org/~rupak/Papers/SoftwareModelChecking.pdf
- Cristian Cadar and Koushik Sen. 2013. "Symbolic execution for software testing: three decades later". Communications of the ACM: https://people.eecs.berkeley.edu/~ksen/papers/cacm13.pdf
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