

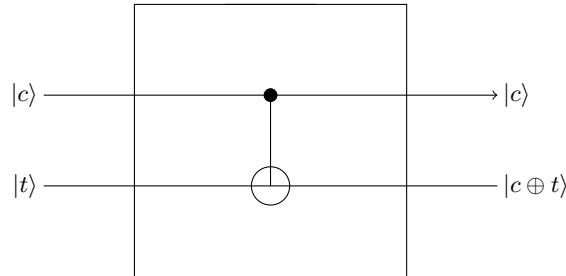
# CS 506 Lecture Notes

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## 1 Phase Kickback

Consider a C-NOT gate:



When  $|c\rangle = |0\rangle$  and  $|t\rangle = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$ , we have

$$\begin{aligned}
 & \text{C-NOT} \left( |0\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \\
 &= \text{C-NOT} \left( \frac{|0\rangle|0\rangle}{\sqrt{2}} \right) - \text{C-NOT} \left( \frac{|0\rangle|1\rangle}{\sqrt{2}} \right) \\
 &= \frac{|0\rangle|0\rangle}{\sqrt{2}} - \frac{|0\rangle|1\rangle}{\sqrt{2}} \\
 &= |0\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}
 \end{aligned}$$

where the C-NOT gate acts like  $I$ .

When  $|c\rangle = |1\rangle$  and  $|t\rangle = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$ , we have

$$\begin{aligned}
 & \text{C-NOT} \left( |1\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \\
 &= \text{C-NOT} \left( \frac{|1\rangle|0\rangle}{\sqrt{2}} \right) - \text{C-NOT} \left( \frac{|1\rangle|1\rangle}{\sqrt{2}} \right) \\
 &= \frac{|1\rangle|1\rangle}{\sqrt{2}} - \frac{|1\rangle|0\rangle}{\sqrt{2}} \\
 &= -|1\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}
 \end{aligned}$$

In general, for  $b \in \{0, 1\}$ ,

$$\begin{aligned}
 & \text{C-NOT} \left( |b\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \\
 &= (-1)^b |b\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} \quad \text{where } (-1)^b \text{ is the phase.} \\
 & \text{C-NOT} \left( (\alpha_0|0\rangle + \alpha_1|1\rangle) \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \\
 &= (\alpha_0|0\rangle - \alpha_1|1\rangle) \frac{|0\rangle - |1\rangle}{\sqrt{2}}
 \end{aligned}$$

So that the eigenvalues of C-NOT are  $(-1)^b = \pm 1$  with respect to the same eigenvector

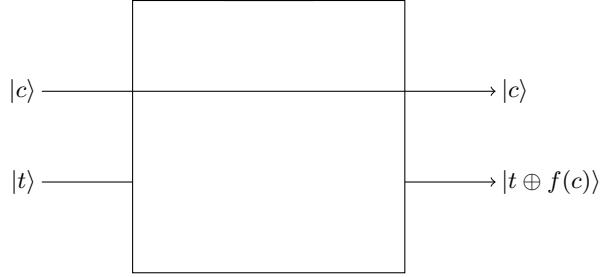
$$\frac{|0\rangle - |1\rangle}{\sqrt{2}},$$

and the eigenvalues are *kicked back* from the target register to the control register.



**Scribe's Note** I felt very confused about how the control qubit  $|c\rangle$  can be regarded as  $-|c\rangle$  as well. After reviewing the Measurement Postulate, I learned that we cannot physically distinguish between them, as they only differ by a phase.

Next, consider a general Control-f gate ( $u_f$ -gate):



where  $f : \{0, 1\} \rightarrow \{0, 1\}$ , and may be irreversible. Then

$$u_f|x\rangle|y\rangle = |x\rangle|y \oplus f(x)\rangle$$

What are the eigenvalues and eigenvectors of  $u_f$ ?

$$\begin{aligned} & u_f|x\rangle \left( \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \\ &= \frac{1}{\sqrt{2}}u_f(|0\rangle|x\rangle) - \frac{1}{\sqrt{2}}u_f(|1\rangle|x\rangle) \\ &= \frac{1}{\sqrt{2}}|x\rangle|0 \oplus f(x)\rangle - \frac{1}{\sqrt{2}}|x\rangle|1 \oplus f(x)\rangle \\ &= |x\rangle \left( \frac{|0 \oplus f(x)\rangle - |1 \oplus f(x)\rangle}{\sqrt{2}} \right) \end{aligned}$$

Case 1:  $f(x) = 0$

$$u_f|x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} = |x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

Case 1:  $f(x) = 1$

$$u_f|x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} = -|x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

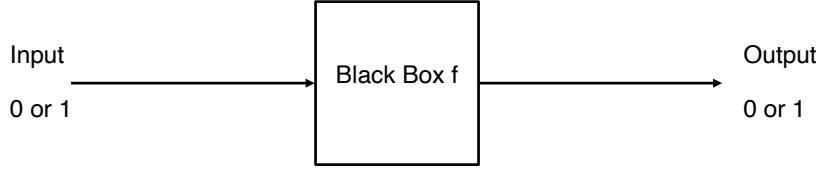
Hence there is a eigenvalues of  $u_f$ ,  $(-1)^{f(x)}$ , with eigenvector  $|x\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}$ .

In general

$$u_f(\alpha_0|0\rangle + \alpha_1|1\rangle) \frac{|0\rangle - |1\rangle}{\sqrt{2}} = \left( (-1)^{f(0)}\alpha_0|0\rangle + (-1)^{f(1)}\alpha_1|1\rangle \right) \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$

## 2 Deutsch Problem

Given a black box



where  $f : \{0, 1\} \rightarrow \{0, 1\}$ , how can we decide if  $f$  is constant or balanced?

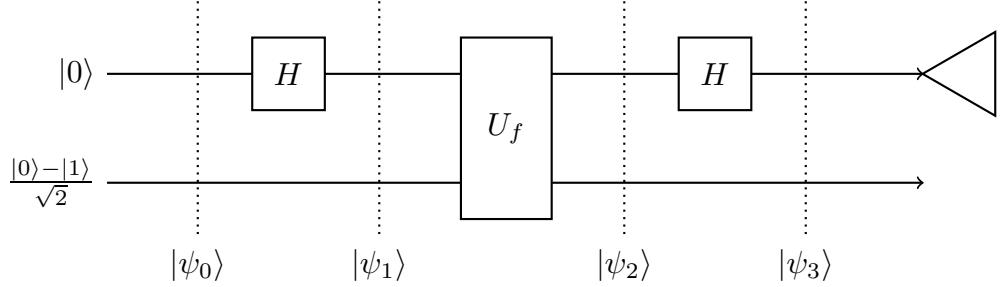
The classical way needs to use 2 queries,  $f(0)$  and  $f(1)$ :

If  $f(0) = f(1)$  then  $f$  is constant;

If  $f(0) \neq f(1)$  then  $f$  is balanced.

However, it only requires only 1 query in quantum computing.

Consider this circuit:



where we  $U_f$  is a Control-f ( $u_f$ -gate) defined previously.

$$u_f|x>|y> := |x>|y \oplus f(x)>.$$

Then we have:

$$\begin{aligned}
 |\psi_0> &= |0>\frac{|0>-|1>}{\sqrt{2}} \\
 |\psi_1> &= (H|0>)\frac{|0>-|1>}{\sqrt{2}} \\
 &= \frac{|0>+|1>}{\sqrt{2}}\frac{|0>-|1>}{\sqrt{2}} \\
 &= \frac{1}{\sqrt{2}}|0>\frac{|0>-|1>}{\sqrt{2}} + \frac{1}{\sqrt{2}}|1>\frac{|0>-|1>}{\sqrt{2}} \\
 |\psi_2> &= u_f|\psi_1> \\
 &= u_f\left(\frac{1}{\sqrt{2}}|0>\frac{|0>-|1>}{\sqrt{2}} + \frac{1}{\sqrt{2}}|1>\frac{|0>-|1>}{\sqrt{2}}\right) \\
 &= u_f\left(\frac{1}{\sqrt{2}}|0>\frac{|0>-|1>}{\sqrt{2}}\right) + u_f\left(\frac{1}{\sqrt{2}}|1>\frac{|0>-|1>}{\sqrt{2}}\right) \\
 &= \frac{1}{\sqrt{2}}(-1)^{f(0)}|0>\frac{|0>-|1>}{\sqrt{2}} + \frac{1}{\sqrt{2}}(-1)^{f(1)}|1>\frac{|0>-|1>}{\sqrt{2}} \\
 &= \frac{(-1)^{f(0)}|0> + (-1)^{f(1)}|1>}{\sqrt{2}}\frac{|0>-|1>}{\sqrt{2}} \\
 &= (-1)^{f(0)}\frac{|0>}{\sqrt{2}} + (-1)^{f(1)-f(0)}\frac{|1>}{\sqrt{2}}\frac{|0>-|1>}{\sqrt{2}} \\
 &= \frac{(-1)^{f(0)}|0> + (-1)^{f(0)\oplus f(1)}|1>}{\sqrt{2}}\frac{|0>-|1>}{\sqrt{2}}
 \end{aligned}$$

- Case  $f(0) \oplus f(1) = 0$ :

$$|\psi_2> = (-1)^{f(0)}\frac{|0> + |1>}{\sqrt{2}}\frac{|0>-|1>}{\sqrt{2}}.$$

Since  $H \frac{|0\rangle + |1\rangle}{\sqrt{2}} = |0\rangle$

$$|\psi_3\rangle = (-1)^{f(0)} |0\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}.$$

Therefore observing the first qubit of  $|\psi_3\rangle$  will always result 0.

- Case  $f(0) \oplus f(1) = 1$ :

$$|\psi_2\rangle = (-1)^{f(0)} \frac{|0\rangle - |1\rangle}{\sqrt{2}} \frac{|0\rangle - |1\rangle}{\sqrt{2}}.$$

Since  $H \frac{|0\rangle + |1\rangle}{\sqrt{2}} = |1\rangle$

$$|\psi_3\rangle = (-1)^{f(0)} |1\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}}.$$

Therefore, observing the first qubit of  $|\psi_3\rangle$  will always result in 1.

Hence, the circuit we built above can tell if  $f$  is constant or balanced by querying  $f$  only once.



**Scribe's Note** This quantum solution is not appealing to me at all, since we are querying one qubit to  $f$  instead of two classic bits. I cannot tell if it has any better performance. However, it may lead to a significant improvement in a more general problem in the next section.

### 3 Deutsch-Jozsa Problem

Consider a general version of the Deutsch Problem:

Given a function  $f : \{0, 1\}^n \rightarrow \{0, 1\}$ , we are granted that  $f$  is either constant or balanced, tell which case it actually is.

For classic computation, we need to query  $f$  at least  $2^{n-1} + 1$  times. But there is a way to query it only once using quantum computation, which will be introduced next class.