

Adventures in Group Theory: Rubik's Cube,
Merlin's Machine, and Other Mathematical Toys

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5-15-2008

In mathematics you don't understand things. You just get used to them.

Johann von Neumann

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Preface

This book grew out of a combined fascination with games and mathematics, from a desire to marry ‘play’ with ‘work’ in a sense. It pursues playing with mathematics and working with games. In particular, abstract algebra is developed and used to study certain toys and games from a mathematical modeling perspective. All the abstract algebra needed to understand the mathematics behind the Rubik’s Cube, Lights Out, and many other games is developed here. If you believe in the quote by von Neumann on a previous page, I hope you will enjoy getting used to the mathematics developed here.

Why is it that these games, developed for amusement by non-mathematicians, can be described so well using mathematics? To some extent, I believe it is because many aspects of our experience are universal, crossing cultural boundaries. One can view the games considered in this book, the Rubik’s cube, Lights Out, etc., to be universal in this sense. Mathematics provides a collection of universal analytical methods, which is, I believe, why it works so well to model these games.

This book began as some lecture notes designed to teach discrete mathematics and group theory to students who, though certainly capable of learning the material, had more immediate pressures in their lives than the long-term discipline required to struggle with the abstract concepts involved. My strategy, to tempt them with something irresistible such as the Rubik’s cube, worked and students loved it. Based on my correspondence, my students were not the only ones to enjoy the notes, which have been expanded considerably. I’ve tried to write the book to be interesting to people with a wider variety of backgrounds but at the same time I’ve tried to make it useful as a reference.

This book was truly a labor of love in the sense that I enjoyed every minute of it. I hope the reader derives some fun from it too.

I’ve tried to ensure no mistakes remain. However, if you find any, I’d appreciate hearing about them so they can be fixed in a possible future edition. My e-mail is wdjoyner@gmail.com.

To continue the universal theme mentioned above, half the royalties from this book will go directly to support **SAGE** (<http://www.sagemath.org/>), a free and open-source computer algebra system. **SAGE** installs on all major platforms (Linux, MacOS and Windows) but is also available for use free over the Internet via a web browser, such as Firefox. The other half of the royalties from this book will go directly to the Earth Island Institute, a nonprofit organization dedicated

to environmental projects all over the world; see <http://www.earthisland.org/>. To paraphrase David Brower (1912-2000), the founder of EII, there is but one ocean, one atmosphere, one Earth, and there are no replacements.

Preface to the Second Edition

I've been very gratified by the reception of the first edition and flattered by the many kind e-mails from puzzle enthusiasts and mathematicians who enjoyed the first book. Based on their suggestions, I believe that the time has come for a new edition of the book. Based on popular demand, I have now included material of a more interactive nature, where **SAGE**, a free computer algebra software system, is used to illustrate ideas. I have also taken the opportunity to update and correct all the first edition's typographical, historical, and mathematical errors I'm aware of. I trust that those who read the first edition will welcome the new version – based as it is on the feedback that I've received. I hope that those new to the book will enjoy reading it as much as I've enjoyed writing it.

Acknowledgments

This book owes much to the very interesting books of Christoph Bandelow [B1], Professor of Mathematics at the Ruhr-University Bochum in Germany, and David Singmaster [Si], Professor of Mathematics in the Department of Computing, Information Systems and Mathematics, South Bank University, London. I thank my fine editor at Johns Hopkins University Press, Trevor Lipscombe, for his encouragement and for his many excellent suggestions. Without them as a source of ideas, this book would not exist. This book has benefited greatly from the stimulating discussions, encouraging correspondence, and collaborations with many people. These people include Christoph Bandelow, Dan Hoey, Ann Luers (now Ann Casey – for help with §§13.4, 13.3, and parts of chapter 14), Michael Dunbar (for help with §13.3), Mark Longridge, Jim McShea (for help with §§7.4, 13.5), Justin Montague and G. Gomes (for help with §15.4.2, 13.2), Michael Reid, Andy Southern (for help with §15.3), Dennis Spellman (for help with §10.5.1), Herbert Kociemba (for §15.2.3), several people on the **SAGE** development team (detailed below), and many others.

Some of the graphs given below were produced with the help of **SAGE** [S]. Some of the group-theoretical calculations were determined with the help of GAP [Gap] or **SAGE**. Some of the biographical information included in this book was borrowed from the award-winning Internet site the MacTutor History of Mathematics Archive [MT].

The Rubik's Cube, Pyraminx, Megaminx, Masterball, Lights Out, and other puzzles names mentioned frequently are all trademarked. I shall omit the symbol (tm) after each occurrence for ease of reading. I would also like to thank Hasbro Toys (who market Lights Out) for the opportunity to consult for them.

SAGE credits: **SAGE** was started in 2005 by William Stein, who is its lead developer. It is open source and can be downloaded and installed without cost by anyone with a good internet connection. The command-line interface in the examples in this book uses Ipython [Ip], started by Fernando Perez. Most people actually prefer the GUI notebook interface (I'm in the minority), written primarily by Tom Boothby and William Stein. This is in the downloadable version of **SAGE** but is also available online at <http://www.sagenb.org/>. Most of the commands in this book can be run online there without having to install **SAGE**. For all the group theoretic constructions, **SAGE** uses an interface to GAP, a computer algebra system specializing in group theory which goes back

to the mid-1980's. For the graph theory illustrated in chapter 7, **SAGE** uses the graph package NetworkX [N] and some functionality written primarily by Emily Kirkman and Robert Miller. Bobby Moretti and Robert Miller wrote **SAGE**'s Cayley graph construction, whereas GAP's GRAPH package was written by Leonard Soicher [So]. **SAGE** code to plot and manipulate the Rubik's Cube was written by Robert Bradshaw, Robert Miller, and myself (with help from Tom Boothby, who wrote some of the underlying graphics functions).

Much of **SAGE**'s development is funded by National Science Foundation grants supporting the number theory research of William Stein, the lead developer and architect of **SAGE**. The University of Washington, William Stein's home institution, has also been a generous supporter. Many thanks to William Stein (and the entire **SAGE** development team) for this great piece of software. Credit also must be given to the group-theoretic package GAP, which was originally developed in the 1980's and has been under continuous active development since then. **SAGE** uses GAP extensively for its group-theoretic computations.

This version owes much to the careful reading of the following people: Jamie Adams, Lewis Nowitz, David Youd, Roger Johnson, Jaap Scherphuis, Michael Hoy, Tom Davis, John Rood, Trevor Irwin, Stephen Lepp, Mark Edwards, Carl Patterson, Peter Neumann, Bill Zeno, Herbert Kociemba, Alastair Farrugia, Matthew Lewis, Christopher Paul Tuffley, Robert Bradshaw, and Benjamin Weggenmann. I thank them all.

Last, but certainly not least, I thank my wonderful wife, Elva, for more things than I can possibly list here.

Where to begin...

Speaking personally, I am always fascinated to discover that one topic is connected in a surprising way with another completely different topic. The Rubik's Cube, a mechanical toy that has a reasonably efficient solution using pure mathematics (and no 'strategy'), is a case in point. It's difficult to know where to begin to explain this connection: as Wodehouse's character Bertie Wooster said, 'I don't know if you have the same experience, but the snag I come up against when I'm telling a story is the dashed difficult problem of where to begin it'. Let's begin with Rubik himself.

Erno Rubik was born in the air-raid shelter of a Budapest hospital during World War II. His mother was a published poet, his father an aircraft engineer who started a company to build gliders. Rubik himself, far from being a mathematician, studied architecture and design at the Academy of Applied Arts and Design, remaining there as a professor, teaching interior design. In the mid-1970's, he patented a cube-shaped mechanical puzzle that has since captured the imagination of millions of people worldwide. The Rubik's Cube was born. By 1982, 'Rubik's Cube' was a household term, and became part of the Oxford English Dictionary. More than 100 million cubes have been sold worldwide.

About 150 years earlier, in the late 1820's and early 1830's, a French teenager named Evariste Galois developed a new branch of mathematics: group theory was born from his attempts to understand the solvability of polynomial equations. In high school, students learn that

$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

are the roots of the quadratic equation $ax^2 + bx + c = 0$. There are similar but more complicated formulas for the cubic and quartic polynomials discovered during the Middle Ages (see, for example, [JKT] for more details). His work on group theory was motivated by what was perhaps the main unsolved mathematical problem of the day: does there exist an analogous algebraic formula involving radicals only in the coefficients for an equation of fifth degree or higher? This problem had remained unsolved for centuries despite the efforts of the best and brightest mathematical minds. Galois' ideas succeeded where others had failed (although one must give credit to Ruffini and Abel here as well). Like Rubik, Galois' life story is also interesting (see, for example, Calinger [Ca]

and the article by Rothman [Ro]) and we shall read more of him later in this book.

Groups measure symmetry. (Galois, for example, studied the ‘symmetry group’ of the roots of a polynomial, now called Galois groups.) As the mathematician Hermann Weyl said in his wonderful book [We], symmetry is ‘one idea by which man has tried throughout the ages to comprehend and create order, beauty, and perfection’. Groups play a key role in the study of roots of polynomials, crystallography, elementary particle physics, campanology (or ‘bell-ringing’; see chapter 3 below), cryptography, and the Rubik’s Cube, among others. Surprisingly enough, it turns out that it is possible to use group theory (and only the ‘group-theoretical definition’ of the cube – no knowledge of strategy or special moves) to solve the Rubik’s Cube (see §10.2 below).

This book will develop the basics of group theory and create group-theoretical models of Rubik’s Cube-like puzzles. On the practical side, I also discuss the solution strategy for the Rubik’s Cube in some detail. (For those wanting to see a solution strategy now, see section 15.1.) Some solution strategies are briefly discussed for similar puzzles (the ‘15 Puzzle’, the ‘Rubik tetrahedron’ or Pyraminx, the ‘Rubik dodecahedron’ or Megaminx, the Skewb, the ‘Hockey puck’, and the ‘Masterball’) as well.

The important point to remember, though, is that group theory is a powerful tool with many real-world applications. Solving puzzles happens to be just one of them. Because of our Rubik’s Cube focus, the approach in this book is different from some texts:

- (a) there are a lot of non-standard, though relatively elementary, group theory topics;
- (b) I emphasize permutation groups via examples over general theory (such as Sylow theory);
- (c) I present some of the basic notions algorithmically (as in [Bu]); and
- (d) I include material which is interesting, from both the mathematical and puzzlist’s perspective, while keeping the level as low as possible for as long as possible.

Moreover, most group theory texts prove everything. Here a lot of statements are proven; sometimes only a hint or sketch is provided, and the proof is left to the interested reader; other statements are supported only by an example. When a proof is not provided, a reference for a proof in the literature is given. To keep things as clear as possible, the start of a proof is denoted **Proof:** and the end by \square . I’ve used the number system where a result or example in section a.b has a labeling of the form a.b.c.

Chapters 1, 2 and 3 give some basic mathematical background. Chapter 4 introduces some of the puzzles and some notation we use for them. Except for the chapter on solutions, chapter 15, the remaining chapters discuss these puzzles using group theory, graph theory, linear algebra, or automata (finite

state machines) theory. While the earlier chapters can probably be followed by a good high school student, chapters 12 and 14 are relatively advanced. They are, in my opinion, the most interesting, and illustrate some remarkable ways in which the Rubik's Cube is connected with other branches of mathematics, which is a continuing source of my fascination with the subject.

The last chapter, chapter 16, gives some indications of directions which weren't pursued and the present state of our, or at least my own, ignorance in this area. As it only gives some of the problems and questions that I don't have the answer to, it is not intended to be complete!

A note to teachers: Based on my experience teaching at the U.S. Naval Academy, a reasonable semester course based on this book could aim for the First Fundamental Theorem of Cube Theory, in chapter 9, with lots of time for 'side trips' and other material. It is possible to cover the Second Fundamental Theorem of Cube Theory, in chapter 11, in one semester but it is a race against time. For example, chapter 6 on 'Merlin's Machine', or any uncovered sections, might be used for some very interesting term projects.

A word about **SAGE**: New in this edition are numerous examples which use **SAGE**, a free and open-source computer algebra system. At the **SAGE** website (www.sagemath.org) there is ample documentation (a reference manual, tutorial, e-mail support lists, and so on) but it is worth taking some time here to explain how **SAGE** is used in this book. First, you don't have to know or care about **SAGE** to read the book. You may ignore the **SAGE** boxes if you wish. On the other hand, if you are fairly comfortable with computer software, the examples are added to help you relate to the constructions in a more interactive manner. There are two commonly used interfaces to **SAGE**: one is the command line (you type in a command and hit **Enter** to see the output) and the other is the graphical "notebook" interface (you type a command into a "cell" and hit **Shift-Enter** for the output). Since I prefer the command line, most examples are described that way. To keep publication costs down, **SAGE** plots and graphics are reproduced in greyscale in this book, though of course in **SAGE** they are in color.

Chapter 1

Elementary, my dear Watson

If logic is the hygiene of the mathematician, it is not his source of food;
the great problems furnish the daily bread on which he thrives.

*André Weil, 'The future of mathematics', American Mathematical
Monthly, May 1950*

Think of a scrambled Rubik's Cube as a car you want to fix on your own. You not only need some tools but you need to know how to use them. This chapter, among others, provides you with some of the tools needed to get the job done. As one of our goals is to discuss the mathematics of the Rubik's Cube, and other games, we start with some fundamentals. The basic purpose of this chapter is to introduce standard set theory notation and some basic notions of mathematical logic.

Logic and set theory are as basic to mathematics as light is to the 'real world'. The background presented here hopefully will make some of the terminology and notation introduced later a little easier to follow for those who either haven't seen or may have forgotten the mathematical notation. In any case, this is not intended to be a 'serious' introduction to mathematical logic nor to set theory.

1.1 You have a logical mind if...

The sentence in the section title, intended to be somewhat whimsical, will be finished later in this section!

A **statement** is an assertion which is either true or false. (Of course we assume that this admittedly circular 'definition' is itself a statement.) Sometimes the truth or falsity of a statement is called its **Boolean value**. One can combine several statements into a single statement using the **connectives** 'and' \wedge , 'or' \vee , and 'implies' \Rightarrow . The Boolean value of a statement is changed using the 'negation' \sim . We shall also use 'if and only if' \iff and 'exclusive or' $_ \vee$ (this

1.1. YOU HAVE A LOGICAL MIND IF...

is defined in the table below), but these can be defined in terms of negation \sim and the other three connectives (\vee , \wedge , and \Rightarrow).

Example 1.1.1. *Today is Monday if and only if today is the day before Tuesday. An example of ‘exclusive or’: Either today is Monday or today is not Monday (but not both).*

Ponderable 1.1.1. *Express $\neg \vee$ and \Longleftrightarrow in terms of \sim , \vee , \wedge , and \Rightarrow .*

Solution: The statement $p \Longleftrightarrow q$ is the same as $(p \Rightarrow q) \wedge (q \Rightarrow p)$, and $p \neg \vee q$ is the same as $\sim (p \Longleftrightarrow q)$. \square

LOGIC, n. The art of thinking and reasoning in strict accordance with the limitations and incapacities of the human misunderstanding. The basic of logic is the syllogism, consisting of a major and a minor premise and a conclusion - thus:

Major Premise: Sixty men can do a piece of work sixty times as quickly as one man.

Minor Premise: One man can dig a posthole in sixty seconds; therefore

-

Conclusion: Sixty men can dig a posthole in one second.

Ambrose Bierce, The Devil’s Dictionary

Notation: Let p and q be statements.

Statement	Notation	Terminology
p and q	$p \wedge q$	‘conjunction’
p or q	$p \vee q$	‘disjunction’
p implies q	$p \Rightarrow q$	‘conditional’
$\sim q$ implies $\sim p$	$\sim q \Rightarrow \sim p$	‘contrapositive’ of $p \Rightarrow q$
negate p	$\sim p$	‘negation’
p if and only if q	$p \Longleftrightarrow q$	‘if and only if’
either p or q (not both)	$p \neg \vee q$	‘exclusive or’

The contrapositive is part of any ‘proof by contradiction’, or ‘reductio ad absurdum’, argument.

Truth tables: Given the Boolean values of the statements p, q , we can determine the values of the statements $p \wedge q$, $p \vee q$, $p \Rightarrow q$, $p \Longleftrightarrow q$, $p \neg \vee q$ using the following truth tables:

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

p	q	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

p	q	$p \Longleftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

Example 1.1.2. The connectives “and” and “or” are already part of Python (which SAGE is built on):

SAGE

```
sage: True and True
True
sage: True and False
False
sage: False and True
False
sage: False and False
False
```

You can see how these Boolean values agree with those in the truth table for \wedge above. As an exercise, you can replace **and** by **or** and check the truth table above for \vee .

Here’s a way to implement “implies” as a function in SAGE:

SAGE

```
sage: def implies(p,q): return not(p and not(q))
sage: implies(True,True)
True
sage: implies(True,False)
False
sage: implies(False,True)
True
sage: implies(False,False)
True
```

You can see how these Boolean values agree with those in the truth table for \Rightarrow above. As an exercise, you can try to implement \Longleftrightarrow in Python and use it to check the truth table above for \Longleftrightarrow .

Note that \Longleftrightarrow is analogous to the $=$ sign.

For example, let p be the statement ‘I am a millionaire’ and q the statement ‘I will buy my wife a brand new car.’ I claim that ‘if p then q ’ is true. If I were

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a millionaire I really would buy my wife a new car. But I am not and cannot afford a car, so p and q are *both* false. This supports the idea that ' $F \Rightarrow F$ ' is true. Now let q be the statement 'I will buy my wife some flowers.' I claim that 'if p then q ' is true. The trouble is, I plan to buy my wife flowers in any case, i.e., q is true no matter what p is. This supports the idea that ' $F \Rightarrow T$ ' is true. In particular, be careful about drawing conclusions from false premises because they imply both true and false statements!

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

p	$\sim p$
T	F
F	T

Ponderable 1.1.2. What can you (logically) conclude from the following Chinese proverb?

If there is light in the soul,
 then there will be beauty in the person.
 If there is beauty in the person,
 then there will be harmony in the house.
 If there is harmony in the house,
 then there will be order in the nation.
 If there is order in the nation,
 then there will be peace in the world.
 Chinese Proverb

Definition 1.1.1. 'For all', written \forall , is the **universal quantifier**. 'There exists', written \exists , is the **existential quantifier**.

Ponderable 1.1.3. Determine which of the following statements is true.

- Exactly one of these statements is false.
- Exactly two of these statements are false.
- Exactly three of these statements are false.
- Exactly four of these statements are false.
- Exactly five of these statements are false.

Though a 'set' is defined later (see §1.2), we shall give an example of what can be done with a set containing only two elements.

Definition 1.1.2. Let B be the set containing only the two elements 0 and 1: this is expressed in symbols as $B = \{0, 1\}$. Think of 0 as 'off' and 1 as 'on'. (In applications, B models the two possible states of an electrical circuit.) Give B the two operations **addition** (+) and **multiplication** (*) defined by the following tables.

x	y	$x + y$
1	1	0
1	0	1
0	1	1
0	0	0

x	y	$x * y$
1	1	1
1	0	0
0	1	0
0	0	0

(Note how these mimic the truth tables of ‘exclusive or’ ($_ \vee$) and ‘and’ ($_ \wedge$).) We call B the **Boolean algebra**. (In applications, $+$ and $*$ model two commonly occurring ‘gates’ in an electrical circuit, [Wba].)

With this, we can start to do some mathematics. As Bertrand Russell said, ‘The rules of logic are to mathematics what those of structure are to architecture.’ (**Mysticism and logic and other essays**, 1917, p61)

We can begin with

De Morgan’s laws:

$$(a) \ p \wedge (q \vee r) \iff (p \wedge q) \vee (p \wedge r) ,$$

$$(b) \ p \vee (q \wedge r) \iff (p \vee q) \wedge (p \vee r),$$

and the **laws of negation:**

$$(c) \ \sim (p \wedge q) \iff (\sim p) \vee (\sim q) ,$$

$$(d) \ \sim (p \vee q) \iff (\sim p) \wedge (\sim q).$$

(Hint: It is possible, though tedious, to verify this using truth tables. To check this yourself, you may want to do (a), (c), and (d) first, and then deduce (b) from these. It is easier to prove the set theory analog, given below, using Venn diagrams, and then deduce the above version from that.)

De Morgan’s laws are analogous to the distributive law $a \cdot (b + c) = a \cdot b + a \cdot c$ for the real numbers. A. De Morgan (1806-1871) started out training to be a mathematician but was refused a fellowship at Cambridge University due to his objections over the mandatory theological test. He then began to study for the bar but also (at the age of 21) applied for the chair of mathematics at University College, London. He was appointed and began as a full professor there a year later. He resigned his chair several times on matters of principle.

Of course, to think logically, we need to produce logical arguments. But what are they?

Definition 1.1.3. *Naively, a **logical argument** is a sequence of statements p_1, p_2, \dots, p_n , which imply a statement q (called the **conclusion**). These statements p_i could be hypotheses, axioms (assumed true), or implications of previously made statements. Such a logical argument is also called a **proof** of q .*

In other words, a logical argument is a true statement of the form

$$(p_1 \wedge p_2 \wedge \dots \wedge p_n) \Rightarrow q.$$

Ponderable 1.1.4. *Using truth tables, verify the logical argument*

$$((p \Rightarrow q) \wedge (q \Rightarrow r)) \Rightarrow (p \Rightarrow r).$$

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1.1.1 ‘You talking to me?’

Definition 1.1.4. A **variable** is a letter denoting some (possibly unknown) object. A **constant** is a letter denoting some specific, well-defined object. A **term** is a variable or a constant.

A **predicate** is a rule which associates to each term a logical statement.

Example 1.1.3. ‘Daffy Duck’ is a constant. The statement ‘ x is one of Daffy Duck’s favorite songs’ is a predicate involving a variable and a constant. Here x might be a Miles Davis song, for example.

To finish the whimsical sentence in the section title, you have a logical mind if when you see the famous saying ‘Time waits for no man’ and ‘No man is an island’, you think ‘therefore time waits for an island!’ If we let x denote the constant ‘no man’ and y denote the constant ‘an island’ then these two sayings can be written ‘Time waits for x ’ and ‘ $x = y$ ’. Substituting, we therefore obtain ‘Time waits for y .’ This may seem crazy, but it’s logical.

When creating a model of the Rubik’s Cube, we shall need to convert some ‘everyday statements’ into symbolical form in order to reduce them to a form suitable for mathematical analysis. Let’s illustrate this with some examples.

Example 1.1.4.

- (a) Consider the statement ‘Each friend of mine can solve the Rubik’s Cube’.
Let

$$M(x) = x \text{ can solve the Rubik's Cube}, \quad S(x) = x \text{ is a friend of mine.}$$

The symbolic form is

$$\forall x, S(x) \Rightarrow M(x).$$

- (b) The Rubik’s Cube is subdivided into 27 smaller cubes, which we call **subcubes**. Only 26 of these are visible, since the center of the Rubik’s Cube is completely surrounded. Consider the statement ‘Each of the 26 visible subcubes of the Rubik’s Cube can be moved into any other visible subcube using some move of the Rubik’s Cube.’ This statement is false, since an edge subcube cannot be moved into a corner subcube, but let us convert this statement into more symbolic notation (which we can more easily analyze) anyway. Let M be the set of all possible (sequences of) moves of the Rubik’s Cube. Let S be the set of all 26 visible subcubes of the Rubik’s Cube. The statement above says that $\forall x, y \in S, \exists m \in M$ such that m sends x to y .’ Though, as we already noticed, this is false, if we let S_C denote the set of all corner subcubes of the Rubik’s Cube then the analogous statement $\forall x, y \in S_C, \exists m \in M$ such that m sends x to y ’ is true; given any two corner subcubes there is a Rubik’s Cube move which sends one to the other (and which, of course, may scramble other subcubes as well).

Ponderable 1.1.5. Convert ‘Some friend of mine can solve the Rubik’s Cube’ to symbolic form.

Ponderable 1.1.6. (*M. Gardner [Gar1]*) Professor White, Professor Brown and Professor Black were lunching together. ‘Isn’t it remarkable’, said the lady, ‘that our names are White, Black, and Brown and one of us has black hair, one has brown hair, and one has white hair.’

‘It is indeed’, answered the one with the black hair as Professor Black bit into his sandwich, ‘and have you noticed that not one has hair color to match our name?’

The lady’s hair is not brown. What is the color of Professor Black’s hair?

	White	Brown	Black	Lady
White				
Brown				
Black				
Lady				×

Ponderable 1.1.7. Retrograde chess problem by Christoph Bandelow. Consider the following chess position.

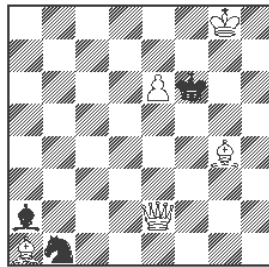


Figure 1.1: What were the last 6 single moves?

1.2 Elements, my dear Watson

Education never ends, Watson.

Sherlock Holmes, ‘The Adventure of The Red Circle’

A **set** is a ‘well-defined’ collection of objects. The objects belonging to a set are the **elements** of the set. If x is an element of a set S then we write $x \in S$. If not, we write $x \notin S$. Let S, T be sets. If all the elements of S also belong to T then we write $S \subset T$ and we say that S is a **subset** of T . If a set S contains only a finite number of distinct elements then we call this number the **cardinality** or **size** of S and denote it by $|S|$. Two sets are said to have the **same cardinality** if there is a one-to-one correspondence between them. A **finite set** is a set which can be enumerated by a finite sequence. An **infinite set** is a set which is not finite. A set containing exactly one element is called a **singleton**.

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There are two common ways to describe a set:

(a) Listing all its elements in between brackets $S = \{a, b, \dots\}$ if the set is finite. For example, assume you have a Rubik's Cube lying on a table in front of you. We denote the set of its **basic moves** by $\{U, D, L, R, F, B\}$, where

- U denotes the move of the Rubik's cube where you turn the upward face clockwise (as you look it from above) one-quarter turn,
- D denotes the move where you turn the downward face clockwise (as you look it from below) one-quarter turn,
- L turns the left face clockwise one-quarter turn,
- R turns the right face clockwise one-quarter turn,
- F turns the front face clockwise one-quarter turn,
- B turns the back face clockwise one-quarter turn.

This shorthand is called the **Singmaster notation**.

(b) Describing the set using properties of its elements using bracket notation: $S = \{x \mid x \text{ has property } P\}$ is pronounced ' S is the set consisting of all x such that x satisfies the property P '. For example, you know a Rubik's Cube has 6 faces (or sides) and each face is subdivided into 9 **facets** (and each facet has a colored sticker glued on it). We could define $S = \{x \mid x \text{ is a facet of the Rubik's cube}\}$. In this case, $|S| = 54$.

Remark 1.2.1. (Russell's paradox) *We must be a little careful when describing sets using properties since some 'self-referential' properties lead to contradictions. Let*

$$R = \{x \mid x \notin x\}.$$

In other words, for all x , $x \in R \iff x \notin x$. In particular, if we take $x = R$ then this becomes

$$R \in R \iff R \notin R,$$

an obvious contradiction. The problem is that R is not 'well-defined' (in the sense that it does not satisfy the set theory axioms, which we will skip here; see [St], for example).

Example 1.2.1. *Here's another way to think of Russell's paradox: A (male) Navy barber is ordered to only shave those (male) Naval personnel who do not shave themselves. Does he shave himself? If yes, then he shaves himself, which is impossible. If no, then he must shave himself, another impossibility. How is this poor barber to avoid being court-martialed for not following orders?!*

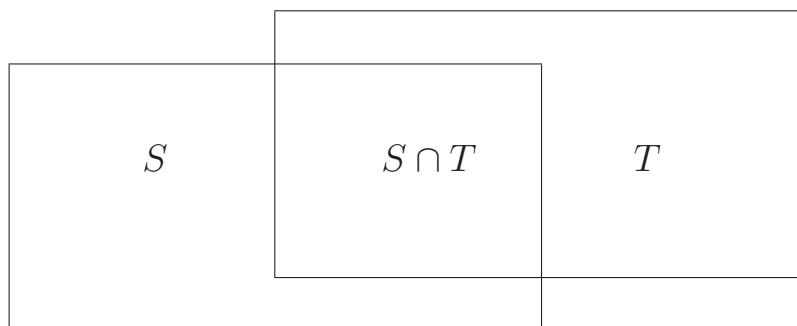
The **empty set** is the set containing no elements, denoted \emptyset .

Notation: Let S and T be sets. Assume that $S \subset X$. The table below defines the terms **intersection**, **union**, **symmetric difference**, and **complement**.

1.2. ELEMENTS, MY DEAR WATSON

Statement	Notation	Terminology
set of elements in S and T	$S \cap T$	intersection
set of elements in S or T	$S \cup T$	union
set of elements in S or in T (not both)	$S \Delta T$	symmetric difference
set of elements not in S	S^c	complement
S is a subset of T	$S \subset T$	subset

We say that two sets S, T are **equal** if $S \subset T$ and $T \subset S$. For example, we can visualize the intersection with the help of a Venn diagram, as in the figure below.



Venn diagrams are named after John Venn (1834-1923), an ordained priest. He was a lecturer at Cambridge University in logic, probability, and set theory.

Example 1.2.2. An example using the computer algebra system **SAGE** [S].

SAGE

```
sage: S1 = Set([1,ZZ,"Zeus"])
sage: S2 = Set([S1,1,"Hera"])
sage: S1; S2
{1, 'Zeus', Integer Ring}
{{1, 'Zeus', Integer Ring}, 1, 'Hera'}
sage: S2.cardinality()
3
sage: S1.difference(S2)
{'Zeus', Integer Ring}
sage: S2.difference(S1)
{{1, 'Zeus', Integer Ring}, 'Hera'}
sage: "Zeus" in S1
True
sage: S2.symmetric_difference(S1)
{{1, 'Zeus', Integer Ring}, 'Hera', 'Zeus', Integer Ring}
```

1.2. ELEMENTS, MY DEAR WATSON

```
sage: S1.symmetric_difference(S2)
{{1, 'Zeus', Integer Ring}, 'Hera', 'Zeus', Integer Ring}
sage: S1.intersection(S2)
{1}
sage: S1.union(S2)
{{1, 'Zeus', Integer Ring}, 1, 'Hera', 'Zeus', Integer Ring}
```

Ponderable 1.2.1. *Using Venn diagrams, verify the De Morgan laws:*

$$(a) S \cap (T \cup U) = (S \cap T) \cup (S \cap U),$$

$$(b) S \cup (T \cap U) = (S \cup T) \cap (S \cup U),$$

and the laws of negation:

$$(c) (S \cap T)^c = S^c \cup T^c,$$

$$(d) (S \cup T)^c = S^c \cap T^c.$$

De Morgan's laws for sets are analogous to De Morgan's laws for logic.

Definition 1.2.1. *We call two sets S, T **disjoint** if they have no elements in common, i.e., if*

$$S \cap T = \emptyset.$$

If

$$S = \cup_{i=1}^n S_i,$$

*where the S_1, S_2, \dots, S_n are pairwise disjoint sets, then we call $\{S_1, S_2, \dots, S_n\}$ a **partition** of S .*

Example 1.2.3. *Let S denote the set of all 54 facets of the Rubik's Cube. Let S_C denote all the facets belonging to a corner subcube, S_E denote all the facets belonging to an edge subcube, and S_Z denote all the facets belonging to a center subcube (we use the subscript Z since the German word for center, zentrum, begins with a Z). Then $S = S_C \cup S_E \cup S_Z$ is a partition.*

The set of all integers will be denoted by \mathbb{Z} . By the way, use of the letter Z (for the German word *Zahlen*, meaning 'numbers') for the integers was apparently first introduced by N. Bourbaki in the 1930's, a pseudonym for a group of French mathematicians aiming to write a thorough, unified account of all mathematics. Earlier, the notation $\overline{3}$ was used by the German number theorist E. Landau (1877-1938) for the integers.

Example 1.2.4. *If*

$$S = \mathbb{Z},$$

$$S_1 = \{\dots, -2, 0, 2, \dots\} = \text{even integers},$$

$$S_2 = \{\dots, -3, -1, 1, 3, \dots\} = \text{odd integers},$$

then $S = S_1 \cup S_2$ is a partition of the integers into the set of even and odd ones.

Logic/set theory analogs: Just as one can use connectives to form new statements from old statements, there are analogous ways to form new sets from old ones using ‘intersection’ (the analog of ‘and’), ‘union’ (the analog of ‘or’), and ‘complement’ (the analog of ‘negation’). The analog of ‘implies’ is ‘subset’. The analog of ‘if and only if’ is ‘equals’.

Set theory	Logic
sets	statements
union	or
intersection	and
subset	implies
symmetric difference	exclusive or
equal	if and only if
Venn diagrams	truth tables
complement	negation

Let P, Q denote statements involving a variable x . We shall often make use of the following facts: if $S = \{x \mid x \text{ has property } P\}$ and if $T = \{x \mid x \text{ has property } Q\}$ then

- $S \cap T = \{x \mid x \text{ has property } P \text{ and property } Q\}$,
- $S \cup T = \{x \mid x \text{ has property } P \text{ or property } Q \text{ (or both)}\}$,
- if $P \Rightarrow Q$ then $S \subset T$,
- if $P \wedge Q$ is false then $S \cap T = \emptyset$,
- if $P \iff Q$ then $S = T$.

We shall prove the first property and leave the other two to the interested reader.

Proof: To show $S \cap T = \{x \mid x \text{ has property } P \text{ and property } Q\}$, we must show $S \cap T \subset \{x \mid x \text{ has property } P \text{ and property } Q\}$, and

$$\{x \mid x \text{ has property } P \text{ and property } Q\} \subset S \cap T.$$

Here are the gory details: First pick an arbitrary element $x \in S \cap T$. x satisfies property P since $x \in S$ and x satisfies property Q since $x \in T$. Thus $x \in \{y \mid y \text{ has property } P \text{ and property } Q\}$. This shows that each element of $S \cap T$ belongs to

$$\{x \mid x \text{ has property } P \text{ and property } Q\},$$

which we can write symbolically (by definition of \subset) as,

$$S \cap T \subset \{x \mid x \text{ has property } P \text{ and property } Q\}.$$

1.2. ELEMENTS, MY DEAR WATSON

Next, pick an arbitrary element

$$x \in \{y \mid y \text{ has property P and property Q}\}.$$

x has property P , so $x \in S$, and x has property Q , so $x \in T$. Therefore (by definition of \cap), $x \in S \cap T$. This shows that each element of

$$\{y \mid y \text{ has property P and property Q}\}$$

belongs to $S \cap T$, so

$$\{y \mid y \text{ has property P and property Q}\} \subset S \cap T.$$

Putting these together, we have shown that

$$S \cap T \subset \{y \mid y \text{ has property P and property Q}\}.$$

This forces the two sets to be equal, and the proof of the first property is complete. \square

Ponderable 1.2.2. (*Joiner's problem [Pe]*) *Cut a 6×6 chessboard and an 8×8 chessboard each into two pieces and join the four pieces into a 10×10 chessboard.*

For more on logic or set theory, see for example [C] or [St].

Chapter 2

‘And you do addition?’

‘And you do Addition?’ the White Queen asked. ‘What’s one and one and one and one and one and one and one and one and one and one?’

Lewis Carroll, Through the looking glass

Just as the White Queen asked Alice about addition, we need to ask ourselves how well we know how to count. As the old pun goes, there are three kinds of people in the world, those who can count and those who can’t.

The mathematics of the Rubik’s Cube involves understanding combinatorial ideas, such as permutations and counting arguments. This chapter will build up some background to introduce some of these notions, and we will end with the binomial theorem. Someone who has read Körner’s *The pleasure of counting* [Ko] may safely skip this chapter (and if you haven’t, it is highly recommended).

2.1 Functions

The type of function we will run across here most frequently is a ‘permutation’, defined precisely later, which is roughly speaking a rule which mixes up and swaps around the elements of a finite set.

Let S and T be finite sets.

Definition 2.1.1. A **function** (sometimes also called a **map** or **transformation**) f from S to T is a rule that associates to each element $s \in S$ exactly one element $t \in T$. We will use the following notation and terminology for this:

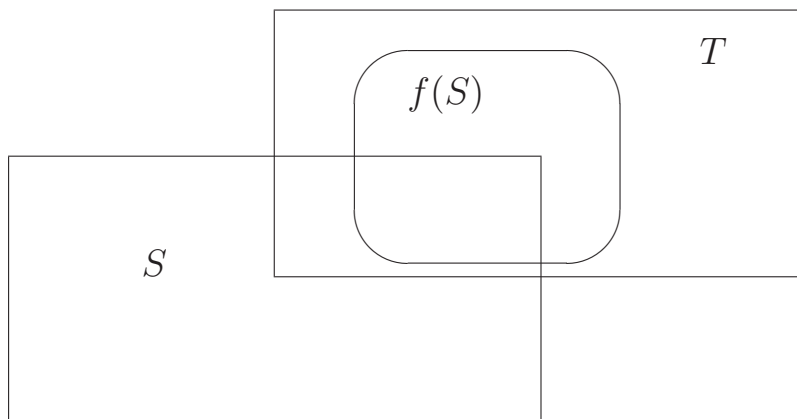
$$\begin{aligned} f : S &\rightarrow T && (f \text{ is a function from } S \text{ to } T), \\ f : s &\longmapsto t && (f \text{ sends } s \text{ in } S \text{ to } t \text{ in } T), \\ t &= f(s) && (t \text{ is the image of } s \text{ under } f). \end{aligned}$$

We call S the **domain** of f , T the **range** of f , and the set

$$f(S) = \{f(s) \in T \mid s \in S\}$$

the **image** of f . The Venn diagram depicting this setup is

2.1. FUNCTIONS



Example 2.1.1.

- One typical example for us will be when S is the set of all 54 facets of the Rubik's Cube (which was introduced in the last chapter) and f is a rule for associating to each facet some other facet determined by a Rubik's Cube move. This example will be given more precisely later.
- Let $\mathbb{N} = \{1, 2, 3, \dots\}$ denote the set of all **natural numbers**. If n is any natural number, let $\phi(n)$ denote the number of natural numbers less than or equal to n which have no prime factor in common with n . (A **prime number** is a natural number greater than 1 which has no factors less than itself except 1. For example, 2, 3, 5, are primes but $10 = 2 \cdot 5$ is not. By the way, a natural number greater than 1 which is not a prime is called a **composite**.) The function $\phi : \mathbb{N} \rightarrow \mathbb{N}$ is sometimes called **Euler's phi function**. We have $\phi(1) = 1$, $\phi(2) = 1$, $\phi(3) = 2$.
- If n is any natural number, let $\pi(n)$ denote the number of prime numbers less than or equal to n . The function $\pi : \mathbb{N} \rightarrow \mathbb{N}$ is simply called the **π function**. We have $\pi(1) = 0$, $\pi(2) = 1$, $\pi(3) = 2$, $\pi(4) = 2$, and so on. Though the rough asymptotic behavior of $\pi(n)$ is known, there are still many unsolved problems regarding $\pi(n)$. For example, it is not yet known if there is an $n > 1$ such that $\pi(n^2 + 2n + 1) = \pi(n^2)$ (i.e., if there is always a prime between any two consecutive squares).

Here is an illustration of how to use **SAGE** to compute with these functions. In fact, **SAGE** (more generally the programming language Python [Py] which **SAGE** uses) allows one to define functions using the 'lambda operator', which does not require a range or domain to define. (The program will return an error

if the function you define is evaluated at points where the definition does not make sense.) I shall give several examples below.

SAGE

```
sage: euler_phi(3)
2
sage: pi = lambda x: pari(x).primepi()
sage: pi(4)
2
sage: f1 = lambda x: sqrt(x)
sage: f1(3)
sqrt(3)
sage: f1(4)
2
sage: f2 = lambda x: RR(sqrt(x))
sage: f2(3)
1.73205080756888
sage: f2(4)
2.000000000000000
```

If you enter `euler_phi(3.1)` or `f1(-3)`, you will get an error since 3.1 is not in the domain of the Euler phi function and -3 is not in the domain of the square-root function.

Ponderable 2.1.1. What is $\phi(16)$? What is $\pi(16)$?

If $f : S \rightarrow T$ is any function and $t \in f(S)$ then we write the **preimage of t** as

$$f^{-1}(t) = \{s \in S \mid f(s) = t\}.$$

(By convention, functional notation is used, even though we don't know if f^{-1} is a function or not. When $f^{-1}(t)$ consists of a single element, say s , we usually call s the **value of $f^{-1}(t)$** and write $s = f^{-1}(t)$.) This is simply the set of all s in S that get sent to t . For example, if $f : S = \mathbb{R} \rightarrow T = \mathbb{R}$ is the function $f(x) = x^2$, and if $t = 4$ then $f^{-1}(t) = f^{-1}(4) = \{2, -2\}$. If $T \subset \mathbb{R}$ then the preimage of 0 is called the set of **zeros** or **roots** of f .

Example 2.1.2. Here is an example of using SAGE to compute preimages.

SAGE

```
sage: f = x^4 - 16
sage: f.roots()
[(2*I, 1), (-2, 1), (-2*I, 1), (2, 1)]
```

In other words, $f^{-1}(0) = \{2i, -2i, 2, -2\}$, each with multiplicity one, where $i = \sqrt{-1}$ is denoted `I` by SAGE.

2.1. FUNCTIONS

Ponderable 2.1.2. If $f : S = \mathbb{R} \rightarrow T = \mathbb{R}$ is the map $f(x) = x^4 - 81$ then find $f^{-1}(0)$.

The **Cartesian product** (or **direct product**) of two sets S, T is the set of pairs of elements taken from these sets:

$$S \times T = \{(s, t) \mid s \in S, t \in T\}.$$

An element of $S \times T$ is simply a list of two things, the first one from S and the second one from T . In the case when $S = T = \mathbb{R}$ is the set of real numbers then sometimes the elements $(x, y) \in \mathbb{R}^2$ are said to have **Cartesian coordinates** or **rectangular coordinates**. This construction of a new set from two given sets is named for the French philosopher René Descartes (1596-1650), whose work **La géométrie** includes the application of algebra to geometry.

Ancient math joke:

Q: What's a rectangular bear?

A: A polar bear after a coordinate transform.

Example 2.1.3. Here is an example of using **SAGE** to compute Cartesian products.

SAGE

```
sage: list(cartesian_product_iterator([[1,2], ['a','b']]))
[(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]
```

More generally, given any collection of m sets, S_1, S_2, \dots, S_m , we can define the m -fold **Cartesian product**, to be the set

$$S_1 \times \dots \times S_m = \{(s_1, \dots, s_m) \mid s_i \in S_i, 1 \leq i \leq m\}.$$

Elements of the Cartesian product $S_1 \times \dots \times S_m$ are called m -**tuples**.

Example 2.1.4. If \mathbb{R} denotes the set of all real numbers then the Cartesian product $\mathbb{R} \times \mathbb{R}$ is simply the set of all pairs of real numbers. In other words, this is the Cartesian plane we are all familiar with from high school.

More generally, we may repeat this process and take the Cartesian product of the set of real numbers with itself n times (where $n > 1$ is any integer) to get $\mathbb{R} \times \dots \times \mathbb{R}$ (n times). This n -fold Cartesian product is denoted more conveniently by \mathbb{R}^n .

The set \mathbb{R}^n has some extra structure, a 'vector space structure', which we shall make use of frequently in this book. The general definition is given next.

Definition 2.1.2. An element of the set \mathbb{R}^n is simply an ordered list of n real numbers. Such a list will be called a **vector** or an n -**vector** to be specific.

More generally, a **real vector space** is a set V (whose elements are called **vectors**) having two operations, **vector addition** $+$ and **scalar multiplication**

\cdot , satisfying the following conditions: (a) if $v, w \in V$ are any two vectors then $v + w = w + v$ is also a vector in V , (b) there is a zero vector $\vec{0}$ such that for any vector $v \in V$, $v + \vec{0} = v$, (c) for any real number $c \in \mathbb{R}$ (called a **scalar**) and any $v \in V$, the product $c \cdot v$ is a vector in V , (d) the distributive laws hold: $(a + b)v = av + bv$ and $c(v + w) = cv + cw$, for all $a, b, c \in \mathbb{R}$ and $v, w \in V$, (e) the associative law holds: $(u + v) + w = u + (v + w)$, for all $u, v, w \in V$, (f) $1 \cdot v = v$ and $a \cdot (b \cdot v) = (ab) \cdot v$, for all $a, b \in \mathbb{R}$ and $v \in V$.

Sometimes it is convenient to ‘picture’ a function as the set of its values inside the Cartesian product $S \times T$. The **graph** of a function $f : S \rightarrow T$ is the subset

$$\{(s, f(s)) \mid s \in S\} \subset S \times T.$$

Example 2.1.5. In high school, if $y = x^2$, one plotted x^2 as a function of x . Here is an example of using **SAGE** to compute preimages.

SAGE

```
sage: f = x^2
sage: P = plot(f, -1, 1)
```

Now type `show(P)` to view the graph below in **SAGE**. Here $f(x) = x^2$ and you plotted $\{(x, f(x)) \mid x \in \mathbb{R}\}$.

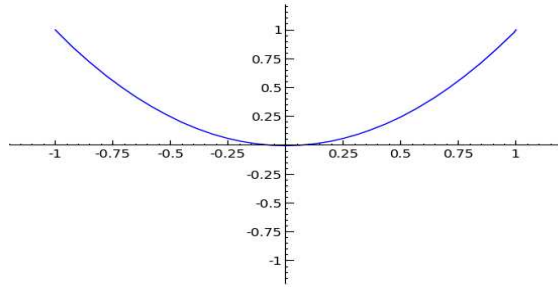


Figure 2.1: **SAGE** plot of $y = x^2$.

It is not possible that every subset of $S \times T$ is the graph of some function from S to T . The following fact classifies exactly which subsets of $S \times T$ can arise as the graph of a function from S to T .

Ponderable 2.1.3. Let $X \subset S \times T$. Show that X is the graph of a function from S to T if and only if (a) for all $(s_1, t_1) \in X$ and $(s_2, t_2) \in X$, whenever $t_1 \neq t_2$ we also have $s_1 \neq s_2$, (b) for all $s \in S$, there is a $t \in T$ such that $(s, t) \in X$.

2.1. FUNCTIONS

(Hint: Let

$$\begin{aligned} pr_1 : S \times T &\rightarrow S \\ (s, t) &\mapsto s \end{aligned}$$

be the projection onto the 1st component. Recall that the graph of a function f has the property that $f^{-1}(s)$ is always a singleton, i.e., a vertical line cannot intersect the graph of a function at two different $f(s)$ values.)

Definition 2.1.3. Let $f : S \rightarrow T$ and $g : T \rightarrow U$ be two functions. We can compose them to get another function, the **composition**, denoted $f \circ g = fg : S \rightarrow U$:

$$\begin{array}{ccccc} x & \mapsto & f(x) & \mapsto & g(f(x)) \\ S & \rightarrow & T & \rightarrow & U. \end{array}$$

Definition 2.1.4. If the image of the function $f : S \rightarrow T$ is all of T , i.e., if $f(S) = T$, then we call f **surjective** (or ‘onto’, or ‘is a surjection’). Equivalently, a function f from S to T is surjective if every $t \in T$ is the image of some $s \in S$ under f .

For example, the map $f : \mathbb{R} \rightarrow \mathbb{R}$ defined by $f(x) = 2x$, for any real number x , is surjective. Another example: let S be the set of all 54 facets of the Rubik’s Cube. Let $f : S \rightarrow S$ be the map which sends a facet to the facet which is diametrically opposite (for instance, the upward center facet would be mapped to the downward center facet). This function is also surjective.

Ponderable 2.1.4. Suppose that $|S| < |T|$. Is there a surjective function $f : S \rightarrow T$? Explain.

Definition 2.1.5. A function $f : S \rightarrow T$ is called **injective** (or ‘one-to-one’ or ‘an injection’) if each element t belonging to the image $f(S)$ is the image of exactly one s in S .

In other words, f is an injection if the condition $f(s_1) = f(s_2)$ (for some $s_1, s_2 \in S$) always forces $s_1 = s_2$.

Ponderable 2.1.5. Suppose that S, T are finite sets and $|S| > |T|$. Is there an injective function $f : S \rightarrow T$? Explain.

Ponderable 2.1.6. Suppose that $|S| = |T| < \infty$. Show that a function $f : S \rightarrow T$ is surjective if and only if it is injective.

Definition 2.1.6. A function $f : S \rightarrow T$ is called a **bijection** if it is both injective and surjective.

Equivalently, a bijection from S to T is a function $f : S \rightarrow T$ for which each $t \in T$ is the image of exactly one $s \in S$.

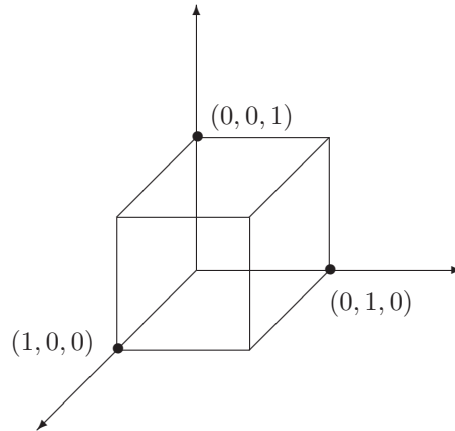
Definition 2.1.7. A set S is called **countable** if there exists a bijection $f : S \rightarrow \mathbb{Z}$ to the set of integers \mathbb{Z} .

Example 2.1.6.

- The set of all odd integers $S = \{\dots, -3, -1, 1, 3, \dots\}$ is countable since the map $f : S \rightarrow \mathbb{Z}$ defined by $f(x) = (x+1)/2$ is a bijection.
- The set of all rational numbers will be denoted by \mathbb{Q} . The set S of all rational numbers within the unit interval $0 < r < 1$ is countable since you can define $f : \mathbb{Z} \rightarrow S$ as follows: $f(1) = 1/2$, $f(2) = 1/3$, $f(3) = 2/3$, $f(4) = 1/4$, $f(5) = 3/4$, $f(6) = 1/5$, $f(7) = 2/5$, $f(8) = 3/5$, and so on. (There are $\phi(n)$ terms of the form m/n , where m is relatively prime to n and $\phi(n)$ denotes the Euler ϕ function in Example 2.1.1 above.).

Example 2.1.7. Let C be the cube in 3-space having vertices at the points $O = (0, 0, 0)$, $P_1 = (1, 0, 0)$, $P_2 = (0, 1, 0)$, $P_3 = (0, 0, 1)$, $P_4 = (1, 1, 0)$, $P_5 = (1, 0, 1)$, $P_6 = (0, 1, 1)$, $P_7 = (1, 1, 1)$. We shall also (to use a notation which will be used more later) denote these by dbl , dfl , dbr , ubl , dfr , ufl , ubr , ufr , respectively. Let $C_0 = \{O, P_1, \dots, P_7\}$ be the set of the eight vertices of C , let $C_1 = \{uf, ur, ub, ul, fr, br, bl, fl, df, dr, db, dl\}$ be the set of the 12 edges of C , and let $C_2 = \{F, B, U, D, L, R\}$ be the set of the 6 faces of C . Let r be the rotation of a point (x, y, z) by 180 degrees about the line passing through the points $(1/2, 1/2, 0)$, $(1/2, 1/2, 1)$. Note that $r : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ is a function which sends the cube C onto itself.

The cube C is pictured as follows:



This function r induces three functions

$$f_0 : C_0 \rightarrow C_0, \quad f_1 : C_1 \rightarrow C_1, \quad f_2 : C_2 \rightarrow C_2,$$

where f_i is the function which sends $x \in C_i$ to its image $r(x)$ under r (which is again in C_i), for $i = 0, 1, 2$. Each f_i is a bijection.

Ponderable 2.1.7. Finish the labeling of the vertices and the faces in the above picture and describe f_1 and f_2 explicitly by filling out the tables

2.2. FUNCTIONS ON VECTORS

v	$f_0(v)$	e	$f_1(e)$	f	$f_2(f)$
$P_1 = dfl$	$P_2 = dbr$	fl	br	F	B

Definition 2.1.8. Suppose $f : S \rightarrow T$ is a bijection. Define f^{-1} to be the rule which associates to $t \in T$ the element $s \in S$,

$$f^{-1}(t) = s \quad \text{if and only if } f(s) = t.$$

This rule defines a function $f^{-1} : T \rightarrow S$. This function satisfies $f(f^{-1}(t)) = t$, for all $t \in T$, and $f^{-1}(f(s)) = s$, for all $s \in S$. In other words, the composition $f \circ f^{-1} : T \rightarrow T$ sends each element of T to itself, and $f^{-1} \circ f : S \rightarrow S$ sends each element of S to itself.

The function f^{-1} constructed above is called the **inverse function** of f .

‘In Riemann, Hilbert or in Banach space
Let superscripts and subscripts go their ways.
Our asymptotes no longer out of phase,
We shall encounter, counting, face to face.’
Stanislaw Lem, The Cyberiad

2.2 Functions on vectors

Later on, we shall need some particular types of functions. This section presents, very briefly, a few basic facts about matrices, which we regard as a certain type of function sending vectors to vectors. For further details, see any text on linear algebra (for example, [Be]).

2.2.1 History

The mathematician who first published a major work that seriously studied matrices and matrix algebra in the western world was Arthur Cayley (1821-1895) of Cambridge, England. He wrote a memoir on the theory of matrices which was published in 1858, and he is often thought of as one of the ‘fathers’ of matrix theory. In fact, it was his friend and colleague James J. Sylvester (1814-1897) who first coined the term ‘matrix’. (Incidentally, Sylvester taught at Johns Hopkins University from 1877 to 1883 and founded the first mathematics journal in the United States, the *American Journal of Mathematics*, put out by the same press that published the book you are reading now.) Though most people learn of matrices and then determinants, historically determinants were studied long before matrices arose.

One of the earliest examples that motivated Cayley's study of matrices was the following. If we have three coordinate systems (x, y) , (x', y') , and (x'', y'') , connected by

$$\begin{cases} x' = x + y \\ y' = x - y \end{cases} \quad \text{and} \quad \begin{cases} x'' = -x' - y' \\ y'' = -x' + y' \end{cases}$$

then the relationship between (x, y) and (x'', y'') is given by

$$\begin{cases} x'' = -x' - y' = -(x + y) - (x - y) = -2x \\ y'' = -x' + y' = -(x + y) + (x - y) = -2y \end{cases}$$

If we do as Cayley did and abbreviate the three changes of coordinates by the square array of the coefficients then we obtain

$$\begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}, \begin{bmatrix} -1 & -1 \\ -1 & 1 \end{bmatrix}, \begin{bmatrix} -2 & 0 \\ 0 & -2 \end{bmatrix}.$$

Cayley noticed that there was an algebraic rule that allows us to determine the third array from the first two, without doing the substitution we did above to relate (x'', y'') to (x, y) . In doing so, Cayley invented 'matrix multiplication'. It corresponds to composing the function sending $(x, y) \mapsto (x', y')$ with the function sending $(x', y') \mapsto (x'', y'')$.

2.2.2 3×3 matrices

A 3×3 matrix is a 3×3 table of real numbers

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}.$$

This may be thought of as a function (which we shall still denote by A) which sends 3-vectors to 3-vectors. This function is defined by the **matrix-vector multiplication formula**,

$$A\vec{v} = \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix} \begin{pmatrix} x \\ y \\ z \end{pmatrix} = \begin{pmatrix} a_{11}x + a_{12}y + a_{13}z \\ a_{21}x + a_{22}y + a_{23}z \\ a_{31}x + a_{32}y + a_{33}z \end{pmatrix}.$$

In other words, A sends the 3-vector (x, y, z) (which we have written above as a column, as is the convention) to the 3-vector $(a_{11}x + a_{12}y + a_{13}z, a_{21}x + a_{22}y + a_{23}z, a_{31}x + a_{32}y + a_{33}z)$. In general, any such 3×3 matrix gives rise to a function $A : \mathbb{R}^3 \rightarrow \mathbb{R}^3$.

Example 2.2.1. *The 3×3 identity matrix,*

$$I_3 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix},$$

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satisfies $I_3 \vec{v} = \vec{v}$ for all $\vec{v} \in \mathbb{R}^3$.

Here is an example of using **SAGE** to compute with matrices.

SAGE

```
sage: A = matrix(3,3,[1,2,3,4,5,6,7,8,9])
sage: (-1)*A
[-1 -2 -3]
[-4 -5 -6]
[-7 -8 -9]
sage: v = vector([1,1,1])
sage: A*v
(6, 15, 24)
sage: x = var("x"); y = var("y"); z = var("z")
sage: w = vector([x,y,z])
sage: A*w
( 3 z + 2 y + x, 6 z + 5 y + 4 x, 9 z + 8 y + 7 x)
```

Ponderable 2.2.1. Show that

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

has the effect of swapping the x - and y -axes.

2.2.3 $m \times n$ matrices

An $m \times n$ **matrix** (of real numbers) is a rectangular array or table of numbers arranged with m rows and n columns. It is usually written:

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & & & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix}.$$

The $(i, j)^{th}$ **entry** of A is a_{ij} . The i **th row** of A is

$$\begin{bmatrix} a_{i1} & a_{i2} & \cdots & a_{in} \end{bmatrix} \quad (1 \leq i \leq m)$$

The j **th column** of A is

$$\begin{bmatrix} a_{1j} \\ a_{2j} \\ \vdots \\ a_{mj} \end{bmatrix} \quad (1 \leq j \leq n)$$

A matrix having as many rows as it has columns ($m = n$) is called a **square matrix**. The entries a_{ii} of an $m \times n$ matrix $A = (a_{ij})$ are called the **diagonal**

entries, the entries a_{ij} with $i > j$ are called the **lower triangular entries**, and the entries a_{ij} with $i < j$ are called the **upper triangular entries**. An $m \times n$ matrix $A = (a_{ij})$ all of whose lower diagonal entries are zero is called an **upper triangular matrix**. This terminology is logical if the matrix is a square matrix, but both the matrices below are called upper triangular

$$\begin{pmatrix} 1 & 2 & 3 & 4 \\ 0 & 5 & 6 & 7 \\ 0 & 0 & 8 & 9 \\ 0 & 0 & 0 & 10 \end{pmatrix}, \quad \begin{pmatrix} 1 & 2 & 3 & 4 \\ 0 & 5 & 6 & 7 \\ 0 & 0 & 8 & 9 \\ 0 & 0 & 0 & 10 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

whether they look triangular or not! A similar definition holds for lower triangular matrices. The square $n \times n$ matrix with 1's on the diagonal and 0's elsewhere,

$$\begin{pmatrix} 1 & & \dots & 0 \\ 0 & \ddots & & \vdots \\ \vdots & & & 0 \\ 0 & \vdots & 0 & 1 \end{pmatrix},$$

is called the $n \times n$ **identity matrix** and denoted I or I_n . This is both upper triangular and lower triangular. (In general, any square matrix which is both upper triangular and lower triangular is called a **diagonal matrix**.)

A square $n \times n$ matrix with exactly one 1 in each row and each column, and 0's elsewhere, is called an $n \times n$ **permutation matrix**. The identity I_n is a permutation matrix. We shall discuss these types of matrices in detail in the next chapter.

A square $n \times n$ matrix with exactly one *non-zero entry* in each row and each column, and 0's elsewhere, is called an $n \times n$ **monomial matrix**. We shall discuss these types of matrices later in the book. They have many properties similar to permutation matrices. Monomial matrices occur in the explicit description of the Rubik's Cube group, which we shall give later.

2.2.4 Multiplication and inverses

Fortunately, we shall not be forced to deal in this book too much with computations of matrix multiplications of large matrices. *Roughly speaking*, we shall eventually see how each move of the Rubik's Cube can be expressed in terms of matrices (more precisely, as a pair of matrices—an 8×8 matrix corresponding to the movement of the 8 corners and a 12×12 matrix corresponding to the movement of the 12 edges). Therefore, a little bit of brief background on matrix multiplication is appropriate.

When you multiply an $m \times n$ matrix A by a $n \times p$ matrix B , you get an $m \times p$ matrix AB . The $(i, j)^{th}$ entry of AB is computed as follows:

1. Let $k = 1$ and $c_0 = 0$.

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2. If $k = m + 1$, you're done and $(AB)_{ij} = c_m$. Otherwise proceed to the next step.
3. Take the k^{th} entry of the i^{th} row of A and multiply it by the k^{th} entry of the j^{th} column of B . Let $c_k = c_{k-1} + a_{ik}b_{kj}$.
4. Increment k by 1 and go to step 2.

In other words, multiply each element of row i in A by the corresponding entry of column j in B , add them up, and put the result in the (i, j) position of the array for AB .

If A is a square $n \times n$ matrix and if there is a matrix B such that $AB = I_n$ then we call B the **inverse** matrix of A , denoted A^{-1} . If you think of A as a function $A: \mathbb{R}^n \rightarrow \mathbb{R}^n$ then A^{-1} is the inverse function. As a practical matter, if n is 'small' (say, $n \leq 3$) then matrix inverses can be computed by pencil and paper. For most larger matrices, computers are needed. We shall discuss easier techniques for inverting permutation matrices in the next chapter and this is the type of matrix we will have to deal with most frequently anyway.

Here is an example of using **SAGE** to compute with matrices.

SAGE

```
sage: A = matrix(3,3,[1,2,3,4,5,6,1,0,0])
sage: det(A)
-3
sage: A^(-1)
[  0   0   1]
[ -2   1  -2]
[ 5/3 -2/3  1]
sage: B = matrix(4,3,[1,2,3,4,5,6,7,8,9,1,1,1])
sage: B*A
[12 12 15]
[30 33 42]
[48 54 69]
[ 6  7  9]
```

Note that AB and $A + B$ are undefined and **SAGE** will return an error if you enter $A*B$ or $A+B$.

Ponderable 2.2.2. *Imagine a chessboard in front of you. You can place at most 8 non-attacking rooks on the chessboard. (Rooks move only horizontally and vertically.) Now imagine you have done this and let $A = (a_{ij})$ be the 8×8 matrix of 0's and 1's (called a $(0,1)$ -**matrix**) where $a_{ij} = 1$ if there is a rook on the square belonging to the i^{th} horizontal down and the j^{th} vertical from the left. Call such a matrix a **rook matrix**. If there are exactly eight 1's in A then*

we shall call A a **full rook matrix**. For example,

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

is a full rook matrix. Show that

- (a) any full rook matrix is invertible,
- (b) the product of any two rook matrices is a rook matrix.

2.2.5 Determinants

The determinant of a matrix A , denoted $\det(A)$, is defined only when A is a square matrix, i.e., the number of rows is equal to the number of columns. If A is an $n \times n$ real matrix then A may be regarded as a function of \mathbb{R}^n , sending points to points. It turns out it will send the unit hypercube in \mathbb{R}^n to a parallelepiped (the n -dimensional analog of a parallelogram). It also turns out that the *absolute value* of the determinant of A measures the volume of that parallelepiped.

If $m = n = 2$ the determinant is easy to define:

$$\det \begin{pmatrix} a & b \\ c & d \end{pmatrix} = ad - bc.$$

Geometrically, $|ad - bc|$ is the area of the parallelogram with vertices $(0, 0)$, (a, c) , (b, d) , $(a + b, c + d)$. This fact, as well as the general definition of the determinant, can be found in most texts on matrix theory or linear algebra. The simplest way to compute the determinant of an arbitrary $n \times n$ matrix $A = (a_{ij})$ is to use the **Laplace cofactor expansion**: for any i with $1 \leq i \leq n$, we have

$$\det(A) = \sum_{j=1}^n (-1)^{i+j} \det(A_{ij}) a_{ij}, \quad (2.1)$$

where A_{ij} is the $(n - 1) \times (n - 1)$ ‘submatrix of A ’ obtained by omitting all the entries in the i^{th} row and the j^{th} column. The matrix A_{ij} is called the (i, j) -**minor** of A and $(-1)^{i+j} \det(A_{ij})$ is called the (i, j) -**cofactor**.

Pierre-Simon Laplace (1749-1827) was one of the most influential scientists in French history. Modesty was not his strong point and he announced in 1780 that he was the greatest mathematician in France. In fact, Laplace made major contributions to many areas of mathematics and physics, especially probability and celestial mechanics (the mathematics of the motion of planets is mostly due to him).

A square matrix A is **singular** if $\det(A) = 0$. Otherwise, A is called **non-singular** or **invertible**.

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Example 2.2.2. Taking $i = 2$,

$$\begin{aligned} & \det \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{pmatrix} \\ &= (-1) \cdot 4 \cdot \det \begin{pmatrix} 2 & 3 \\ 8 & 9 \end{pmatrix} + (+1) \cdot 5 \cdot \det \begin{pmatrix} 1 & 3 \\ 7 & 9 \end{pmatrix} + (-1) \cdot 6 \cdot \det \begin{pmatrix} 1 & 2 \\ 7 & 8 \end{pmatrix} \\ &= (-4)(18 - 24) + (5)(9 - 21) + (-6)(8 - 14) = 0. \end{aligned}$$

This implies A is singular. Indeed, the parallelepiped generated by $(1, 2, 3)$, $(4, 5, 6)$, $(7, 8, 9)$, must be flat (2-dimensional, hence have 0 volume) since $(4, 5, 6) = (1, 2, 3) + (1, 1, 1)$ and $(7, 8, 9) = (1, 2, 3) + 2(1, 1, 1)$.

There is an analogous Laplace expansion like (2.1) for columns as well. See any other book on linear algebra for further details.

Here is an example of using **SAGE** to compute with determinants.

SAGE

```
sage: A = matrix(3,3,[1,2,3,4,5,6,7,8,9])
sage: det(A)
0
sage: a = var("a"); b = var("b"); c = var("c"); d = var("d")
sage: B = matrix(2,2,[a,b,c,d])
sage: det(B)
a*d - b*c
```

An important fact about singular matrices, and one that we will use later, is the following.

Lemma 2.2.1. Suppose A is an $n \times n$ matrix with real entries. Then $\det(A) = 0$ if and only if there is a non-zero vector v such that $Av = 0$, where 0 is the zero vector in \mathbb{R}^n .

Here's a sketch: $\det(A) = 0$ if and only if A is singular if and only if A is not one-to-one. A is not one-to-one if and only if there are vectors $v_1 \neq v_2$ for which $Av_1 = Av_2$, i.e., $Av = 0$, where $v = v_1 - v_2 \neq 0$.

We end this section with one last key property of determinants.

Lemma 2.2.2. If A, B are any two $n \times n$ matrices having real entries then $\det(AB) = \det(A)\det(B)$.

More details on all this material can be found in any textbook on linear algebra.

I think mathematics is a vast territory. The outskirts of mathematics are the outskirts of mathematical civilization. There are certain subjects that people learn about and gather together. Then there is a sort of inevitable development in those fields. You get to a point where a certain

theorem is bound to be proved, independent of any particular individual, because it is just in the path of development.

*William P. Thurston, **More mathematical people**, New York, 1990, p.332*

2.3 Relations

A relation on a set is a generalization of the concept of a function from S to itself. Relations include $=$, $<$, and many other types of ‘comparisons’.

Definition 2.3.1. *Let S be a set. If R is a subset of $S \times S$ then we call R a **relation on S** . If $(x, y) \in R$ then we say that x is **related** to y .*

We may also regard a relation R on S as a function

$$R : S \times S \rightarrow \{0, 1\}.$$

In this form, we say x is **related** to y (for $x, y \in S$) if $R(x, y) = 1$.

This is a very general notion. There are lots and lots of relations in mathematics: inequality symbols, functions, subset symbols are all common examples of relations.

Example 2.3.1. *Let S be any set and let f be a function from S to itself. This function gives rise to the relation R on S defined by the graph of f :*

$$R = \{(x, y) \in S \times S \mid y = f(x), \text{ for } x \in S\}.$$

(It is through this correspondence that we may regard a function as a relation.)

Example 2.3.2. *Let S be the set of all subsets of $\{1, 2, \dots, n\}$. Let R be defined by*

$$R = \{(S_1, S_2) \mid S_1 \subset S_2, S_1 \in S, S_2 \in S\}.$$

Note that R is a relation.

Definition 2.3.2. *Let R be a relation on a set S . We call R an **equivalence relation** if it satisfies the following properties.*

- (a) *Each element $s \in S$ is related to itself (‘reflexive’). In other words, (s, s) belongs to R for all s .*
- (b) *If s is related to t then t is related to s (‘symmetry’). In other words, $(s, t) \in R$ implies $(t, s) \in R$.*
- (c) *If s_1 is related to s_2 and s_2 is related to s_3 then s_1 is related to s_3 (‘transitivity’).*

Example 2.3.3. *The equality symbol $=$ provides an equivalence relation on the real line: let $D = \{(x, x) \mid x \text{ real}\}$. This is an equivalence relation on the real line: note $x = y$ if and only if (x, y) belongs to D .*

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Notation: If R is an equivalence relation on S then we often write $x \sim y$ or $x \equiv y$ in place of $(x, y) \in R$, for simplicity.

Definition 2.3.3. Let R be an equivalence relation on a set S . For $s \in S$, we call the subset

$$[s] = \{t \in S \mid s \sim t\}$$

the **equivalence class** of s in S .

Example 2.3.4. Fix an integer $n > 1$. For integers x, y , define $x \equiv y$ if and only if n divides $x - y$. In this case, we say that x is **congruent** to $y \bmod n$, written $x \equiv y \pmod{n}$. The equivalence class of x is sometimes (for historical reasons) called the **residue class** (or **congruence class**) of $x \pmod{n}$. We shall sometimes abuse notation and denote residue class of $x \bmod n$ simply by x .

This notation was first introduced by Carl F. Gauss (1777-1855), who is regarded by many as one of the top mathematicians of all time. At the age of 23 he wrote **Disquisitiones Arithmeticae**, which started a new era of number theory and introduced this notation.

Ponderable 2.3.1.

- (a) Let $f(x) = x^2$, let S be the real line, and let R be the corresponding relation as in Example 2.3.1. Is R an equivalence relation?
- (b) Let $f(x) = 2x$, let S be the real line, and let R be the corresponding relation as in Example 2.3.1. Is R an equivalence relation?

Ponderable 2.3.2. Let R be the corresponding relation as in the second example. Is R an equivalence relation?

Example 2.3.5. For integers x, y , define $x \equiv y$ if and only if 3 divides $x - y$. the equivalence classes are

$$\begin{aligned}[0] &= \{\dots, -6, -3, 0, 3, 6, \dots\}, \\ [1] &= \{\dots, -5, -2, 1, 4, 7, \dots\}, \\ [2] &= \{\dots, -4, -1, 2, 5, 8, \dots\}.\end{aligned}$$

Using Python's modulo operator `%`, these sets can be computed in **SAGE** as follows:

SAGE

```
sage: C0 = [x for x in range(-10,10) if x%3 == 0]; print C0
[-9, -6, -3, 0, 3, 6, 9]
sage: C1 = [x for x in range(-10,10) if x%3 == 1]; print C1
[-8, -5, -2, 1, 4, 7]
sage: C2 = [x for x in range(-10,10) if x%3 == 2]; print C2
[-10, -7, -4, -1, 2, 5, 8]
```


Ponderable 2.3.3. Show that for any s_1 and s_2 in S , we have either

- (a) $[s_1] = [s_2]$, or
- (b) $[s_1]$ is disjoint from $[s_2]$.

As a consequence of this problem, we see that if R is an equivalence relation on a set S then the equivalence classes of R partition S into disjoint subsets. We state this separately for future reference (we also assume S is finite for simplicity).

Theorem 2.3.1. If S is a finite set and R is an equivalence relation on S then there are subsets

$$S_1 \subset S, S_2 \subset S, \dots, S_k \subset S,$$

satisfying the following properties:

- (1) S is the union

$$S = S_1 \cup S_2 \cup \dots \cup S_k = \cup_i S_i$$

- (2) the S_i 's are disjoint: for $1 \leq i \leq k$, $1 \leq j \leq k$, if $i \neq j$ then $S_i \cap S_j = \emptyset$.
- (3) the S_i 's exhaust the collection of equivalence classes of R : for each i , $1 \leq i \leq k$, there is an $s_i \in S$ such that

$$S_i = [s_i].$$

(This element s_i is called a **representative** of the equivalence class S_i .)

(Properties (1) and (2) say that the S_i 's partition S in the sense of the previous chapter.)

Example 2.3.6. For real numbers x, y , define $x \equiv y$ if and only if $x - y$ is an integer. The equivalence classes are of the form

$$[x] = \{\dots, x - 2, x - 1, x, x + 1, x + 2, \dots\},$$

for x real. Each equivalence class has exactly one representative in the half open interval $[0, 1) = \{x \in \mathbb{R} \mid 0 \leq x < 1\}$.

Remark 2.3.1. Conversely, given a partition as in Theorem 2.3.1(1), there is an equivalence relation R such that $S_i = [s_i]$, for some $s_i \in S$, where

$$[s] = \{x \in S \mid s \sim x\}$$

is the equivalence class of s with respect to R . Indeed, we define

$$R = \cup_{i=1}^k S_i \times S_i.$$

This is an equivalence relation and $s \sim t$ if and only if $s, t \in S_i$, for some $i = 1, 2, \dots, k$.

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Ponderable 2.3.4. Let S be the set \mathbb{Z} of all integers. Let R be the relation defined by $(x, y) \in R$ if and only if $x - y$ is an even number (i.e., an integer multiple of 2).

- (1) Show that this is an equivalence relation,
- (2) Find the sets S_i in the above lemma which partition S ,
- (3) Find a representative of each equivalence class S_i .

2.4 Counting and mathematical induction

Major General Stanley: ‘About binomial theorem I’m teeming with a lot o’ news...’

W. S. Gilbert and A. Sullivan, Pirates of Penzance

This section quickly surveys a few basic counting principles we shall use later.

The Three Stooges, Curly, Moe, and Larry, each know 10 different insults, none in common with each other. For example, Moe knows ‘You knucklehead!’, and 9 others. One of them insults you by calling you a name. How many different insults are possible? Thirty, because each one of the three can insult you in 10 ways. The following principle generalizes this idea.

Theorem 2.4.1. (Addition Principle) Let S_1, \dots, S_n denote disjoint finite sets. Then

$$|S_1 \cup \dots \cup S_n| = |S_1| + \dots + |S_n|.$$

We give a proof of this, mostly as an excuse to tell the corny horse joke below.

The proof proceeds by the method of **mathematical induction** (or simply **induction**, not to be confused with the logical reasoning of induction). This will be presented from a general viewpoint and then applied to proving the Addition Principle.

Definition 2.4.1. Let $P(1), P(2), \dots$ be a sequence of logical statements which you want to prove to be true (i.e., $P(k)$ is a ‘predicate’ with ‘term’ $k = 1, 2, \dots$).

- (1) Show that the case $k = 1$ is true, i.e., prove $P(1)$.
- (2) Assuming the truth of the case $k = n - 1$, i.e., assuming $P(n - 1)$ is true (the **induction hypothesis**), prove $P(n)$.

Once we have these two facts, we can conclude $P(1)$ is true (case (1)), $P(1) \Rightarrow P(2)$ is true (case (2)), so $P(2)$ is true (by the truth tables in the previous chapter), $P(2) \Rightarrow P(3)$ is true (case (2)), so $P(3)$ is true (by the truth tables in the previous chapter), and so on. We deduce that the logical statements $P(k)$ are all true. This is the method of **mathematical induction**.

De Morgan, whom we have met already, was the first to define the term, mathematical induction (though the basic idea was known before him).

As a humorous example, we ‘prove’, using mathematical induction, that horses have an odd number of feet. Let $P(k)$ be the statement that k horses have an odd number of feet. Case $k = 1$ is the statement that one horse has an odd number of feet. Well, any horse with a broken leg has an odd number of feet, so $P(1)$ is true. Now assume $P(k)$ is true and we try to prove $P(k + 1)$. Any k of the $k + 1$ horses have an odd number of feet, so all $k + 1$ do as well, hence $P(k + 1)$ holds as well.

As a more serious example, let $P(k)$ be the logical statement $|S_1 \cup \dots \cup S_k| = |S_1| + \dots + |S_k|$, $1 \leq k \leq n$. We need to

- (1) Prove the claim for $k = 1$, i.e., prove $P(1)$.
- (2) Assuming the truth of the case $k = n - 1$, prove $P(n)$.

Proof: Let $P(k)$ be the logical statement $|S_1 \cup \dots \cup S_k| = |S_1| + \dots + |S_k|$, $1 \leq k \leq n$.

Case $k = 1$. $P(1)$ is the statement $|S_1| = |S_1|$, which is of course true.

Case $k = n - 1$. Assume $|S_1 \cup \dots \cup S_{n-1}| = |S_1| + \dots + |S_{n-1}|$ is true.

Let $S = S_1 \cup \dots \cup S_{n-1}$ and let $T = S_n$. Each element of $S \cup T$ is either an element of S or an element of T but not both since they are disjoint. How many ways can we pick an element from $S \cup T$? $|S| + |T|$ ways, since may pick one from either S or from T . The induction hypothesis implies $|S| = |S_1 \cup \dots \cup S_{n-1}| = |S_1| + \dots + |S_{n-1}|$. Since $S_1 \cup \dots \cup S_n = S \cup T$, this and the previous paragraph together imply $|S_1 \cup \dots \cup S_{n-1} \cup S_n| = |S \cup T| = |S| + |T| = |S_1| + \dots + |S_{n-1}| + |S_n|$. This proves the case $k = n$.

By mathematical induction, the proof of the addition principle is complete.

□

Example 2.4.1. *If there are n bowls, each containing some distinguishable marbles, and if S_i is the set of marbles in the i^{th} bowl then the number of ways to pick a marble from exactly one of the bowls is $|S_1| + \dots + |S_n|$, by the addition principle.*

As a corollary of the above addition principle, we have the following result.

Proposition 2.4.1. (Dirichlet Pigeonhole Principle) *Let d, m, n be positive integers. If there are n objects (pigeons) which must be placed in m boxes (these ‘pigeonholes’ as assumed to be disjoint), where $md < n$, then there is at least one box with at least $d + 1$ objects.*

Proof: We prove this using the **reductio ad absurdum** argument: assume the statement claimed to be true is actually false and derive a contradiction.

Suppose that the i^{th} box S_i has less than or equal to d objects in it, for every i with $1 \leq i \leq m$. Since the boxes are assumed to be disjoint, the hypothesis of the addition principle holds. Therefore, the total number of objects in all the boxes is $|S_1| + \dots + |S_m| \leq d + \dots + d = md$. But there are n such objects, so we must therefore have $n \leq md$. This contradicts the assumption that $md < n$. □

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Example 2.4.2. *If you are in a room with 9 other people then there must be either at least 5 people you know or 5 people you don't know (not counting yourself). This is obvious enough. It is also what the Dirichlet Pigeonhole Principle tells you: In this case, there are $n = 9$ objects (the people in the room) and $m = 2$ boxes (the friend box and the stranger box) so we may take $d = 4$ in the pigeonhole principle.*

Returning to the Three Stooges, who each know 10 different insults, none in common. Each one of them insults you exactly once. How many insults are possible? One thousand, 10^3 , because each one of the three can insult you in 10 ways. The following principle generalizes this idea.

Theorem 2.4.2. (Multiplication Principle) *Let S_1, \dots, S_n denote finite sets. Then*

$$|S_1 \times \dots \times S_n| = |S_1| \cdot \dots \cdot |S_n|.$$

The proof proceeds by the method of induction. Let $P(k)$ be the logical statement $|S_1 \times \dots \times S_k| = |S_1| \cdot \dots \cdot |S_k|$, $1 \leq k \leq n$. Singing our chorus once again, we must

- (1) Prove the claim for $k = 1$, i.e., prove $P(1)$.
- (2) Assuming the truth of the case $k = n - 1$, prove $P(n)$.

Proof: Let $P(k)$ be the logical statement $|S_1 \times \dots \times S_k| = |S_1| \cdot \dots \cdot |S_k|$, $1 \leq k \leq n$.

Case $k = 1$. $P(1)$ is the statement $|S_1| = |S_1|$, which is of course true.

Case $k = n - 1$. Assume $|S_1 \times \dots \times S_{n-1}| = |S_1| \cdot \dots \cdot |S_{n-1}|$ is true.

Let $S = S_1 \times \dots \times S_{n-1}$ and let $T = S_n$. Each element of $S \times T$ is a list of 2 elements, the first from S , the second from T . How many ways can we pick an element from S to put in the first coordinate? $|S|$ ways. Suppose that the first coordinate has been chosen. How many ways can we pick an element from T to put in the second coordinate? $|T|$ ways. Therefore, the number of ways to pick elements for the first coordinate and the second coordinate is $|S| \cdot |T|$. The induction hypothesis implies $|S| = |S_1 \times \dots \times S_{n-1}| = |S_1| \cdot \dots \cdot |S_{n-1}|$. Since $S_1 \times \dots \times S_n = S \times T$, this and the previous paragraph together imply $|S_1 \times \dots \times S_n| = |S \times T| = |S| \cdot |T| = |S_1| \cdot \dots \cdot |S_{n-1}| \cdot |S_n|$. This proves the case $k = n$.

By mathematical induction, the proof of the multiplication principle is complete. \square

Example 2.4.3. *If there are n bowls, each containing some distinguishable marbles and if S_i is the set of marbles in the i^{th} bowl then the number of ways to pick exactly one marble from each of the bowls is $|S_1| \cdot \dots \cdot |S_n|$, by the multiplication principle.*

Corollary 2.4.1. *The number of ordered selections, taken with repetition, of m objects from a set of n objects is n^m .*

An ordered selection is sometimes called an **arrangement**.

Example 2.4.4. *Select from the set of majors,*

Mathematics, Computer Science, Chemistry, Engineering, Physics,

exactly 3 (with repetition allowed). This is the same as selecting 3 objects, with replacement, from 5 jars each containing the 5 phrases above. There are 5^3 possible choices.

Corollary 2.4.2. *The number of ordered selections, taken without repetition, of m objects from a set of n objects ($m < n$) is*

$$\frac{n!}{(n-m)!} = n \cdot (n-1) \cdot \dots \cdot (n-m+1).$$

Proof: Strictly speaking this is not a corollary of the multiplication principle, so a proof is given.

Since the objects selected are ordered, by hypothesis, we may speak of a first selection, a second selection, and so on. There are n ways to make the first selection. Since there is no repetition allowed, this object is removed from the set used to make the second selection. There are $n-1$ ways to make the second selection, leaving $n-2$ objects remaining. This process can be continued m times since $m < n$, by hypothesis. Multiplying the integers $n, n-1, \dots, n-m$ gives the result. \square

Example 2.4.5. *The number of ‘ordered poker hands’, 5-tuples, without repetition, of objects from the set $\{1, 2, \dots, 52\}$ is*

$$\frac{52!}{(47)!} = 52 \cdot 51 \cdot 50 \cdot 49 \cdot 48 = 311875200 \cong 3.1 \times 10^8.$$

Ponderable 2.4.1. *Let C be a set of 6 distinct colors. Fix a cube in space (imagine it sitting in front of you on a table). We call a **coloring** of the cube a choice of exactly one color per side. Let S be the set of all colorings of the cube. We say $x, y \in S$ are **equivalent** if x and y agree after a suitable rotation of the cube.*

(a) *Show that this is an equivalence relation.*

(b) *Count the number of equivalence classes in S .*

The next two results are corollaries of the multiplication principle as well.

Corollary 2.4.3. *The number of unordered selections, taken without repetition, of m objects from a set of n objects ($m < n$) is*

$$C(n, m) = \binom{n}{m} = \frac{n!}{(n-m)!m!} = \frac{n \cdot (n-1) \cdot \dots \cdot (n-m+1)}{m \cdot (m-1) \cdot \dots \cdot 1}.$$

An unordered selection is sometimes called a **combination**.

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Example 2.4.6. The number of (unordered) poker hands, 5-tuples without repetition of objects from the set $\{1, 2, \dots, 52\}$ is

$$C(52, 5) = \frac{52 \cdot 51 \cdot 50 \cdot 49 \cdot 48}{5!} = 2598960 \cong 2.6 \times 10^6.$$

Corollary 2.4.4. The number of unordered selections, taken with repetition allowed, of m objects from a set of n objects ($m < n$) is $C(n + m - 1, m)$

Proof: The proof is rather clever, so we include a sketch.

Imagine m 1's in a row. Order the n objects you will be selecting from as object 1, object 2, Starting from the leftmost 1, count the number of object 1's you will select, then put a | mark to the right of the last 1. Put the | to the left of all the 1's if you don't select any from object 1. If you selected m_1 elements from object 1 then you have m_1 1's to the left of the | and $m - m_1$ 1's to the right. Do the same for object 2, then object 3, You must have inserted n |'s. There are $C(n + m - 1, m)$ ways to do this. \square

This symbol $C(n, m)$, also commonly written $\binom{n}{m}$, is called the **combinatorial** (or **binomial**) **symbol** and pronounced ' n choose m '. It occurs in the binomial theorem and in Pascal's triangle,

				1				
			1		1			
		1		2		1		
	1		3		3		1	
1		4		6		4		1
1	5		10		10		5	1
			...					

where each entry is obtained by adding the two immediately above it. In other words, the defining property of Pascal's triangle is the following general fact, whose proof is left as an exercise for the interested reader.

Lemma 2.4.1. If $0 < m < n$ then

$$C(n, m) + C(n, m + 1) = C(n + 1, m + 1).$$

For example, $C(5, 2) + C(5, 3) = C(6, 3)$. The following **SAGE** computation illustrates this.

SAGE

```

sage: C52 = binomial(5,2)
sage: C53 = binomial(5,3)
sage: C52; C53; C52 + C53
10
10
20
sage: C63 = binomial(6,3)
sage: C63
20
sage: C63 == C52 + C53
True

```

Theorem 2.4.3. (Binomial Theorem)

$$(x + y)^n = \sum_{m=0}^n C(n, m)x^m y^{n-m} \\ = x^n + nx^{n-1}y + \dots + C(n, m)x^m y^{n-m} + \dots + nxy^{n-1} + y^n.$$

We shall not prove this theorem here. For the interested reader: the key ideas in the proof involve using mathematical induction and the defining property of Pascal's triangle—try to find it! By the way, Pascal's triangle was known long before Blaise Pascal (1623-1662) wrote his book **Treatise on the Arithmetical Triangle** on it in the 1600's.

Example 2.4.7. Here is an example of using **SAGE** to compute with binomial coefficients, permutations, and factorials.

For example, to find the number of times x^2y^3 appears in the expansion of $(x + y)^5$, you must count how many “words” can be formed by permuting the letters of $xyyy$. There are 10 such “words”, as **SAGE** shows below, and Pascal's triangle shows above.

SAGE

```

sage: expand((x+y)^(5))
y^5 + 5*x*y^4 + 10*x^2*y^3 + 10*x^3*y^2 + 5*x^4*y + x^5
sage: binomial(5,2)
10
sage: terms = ["x","x","y","y","y"]
sage: number_of_permutations(terms)
10
sage: permutations(terms)
[['x', 'x', 'y', 'y', 'y'],
 ['x', 'y', 'x', 'y', 'y'],
 ['x', 'y', 'y', 'x', 'y'],
 ['x', 'y', 'y', 'y', 'x'],
 ['y', 'x', 'x', 'y', 'y'],

```

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```
[ 'y', 'x', 'y', 'x', 'y'],  
[ 'y', 'x', 'y', 'y', 'x'],  
[ 'y', 'y', 'x', 'x', 'y'],  
[ 'y', 'y', 'x', 'y', 'x'],  
[ 'y', 'y', 'y', 'x', 'x']]  
sage: terms = [1,2,3,4,5]  
sage: number_of_permutations(terms)  
120  
sage: factorial(5)  
120
```

The reader can see that the third term of the binomial expansion is really a sum of the 10 (equal) terms which are given by the output of `permutations(terms)`.

Chapter 3

Bell ringing and other permutations

Mathematics, springing up from the soil of basic human experience with numbers and data and space and motion, builds up a far-flung architectural structure composed of theorems which reveal insights into the reasons behind appearances and of concepts which relate totally disparate concrete ideas.

Sanders MacLane, American Mathematical Monthly, 1954

In the last chapter we discussed counting arguments and some combinatorial notions. This chapter introduces more combinatorial ideas and techniques—more tools for your toolbox, so to speak. The objective is to keep track of rearrangements of objects, such as the facets of a Rubik's Cube, which get rearranged each time you make a move. Some new notation, symbols, and terms must be introduced.

Talk about relating totally disparate ideas—at the end of the chapter, we will see how combinatorial ideas already occurred in (church) bell ringing over 300 years ago!

3.1 Definitions

Suppose you shuffle a pack of cards. You are simply rearranging all 52 of them. The mathematical equivalence of shuffling is a permutation. A more precise general definition is as follows. Let $\mathbb{Z}_n = \{1, 2, \dots, n\}$ be the set of integers from 1 to a fixed positive integer n . A **permutation** of \mathbb{Z}_n is a bijection from \mathbb{Z}_n to itself. (A bijection was defined in Definition 2.1.6.) More generally, if T is any finite set then a permutation of T is a bijection from T to itself. In the case when T has n elements, we shall often label the elements of T by $T = \{t_1, \dots, t_n\}$ and regard a permutation $f : T \rightarrow T$ as a permutation $\phi : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$, where $f(t_i) = t_j$ if and only if $\phi(i) = j$. In other words, once we label the elements of T , there is a 1-1 correspondence between the permutations of T and those of

3.1. DEFINITIONS

\mathbb{Z}_n . Since the elements of \mathbb{Z}_n are easier to write (fewer subscripts!), we often just work with \mathbb{Z}_n .

As an example, on the 3×3 Rubik's Cube there are $9 \cdot 6 = 54$ facets. If you label them $1, 2, \dots, 54$ (in any way you like), then any move of the Rubik's Cube corresponds to a permutation of \mathbb{Z}_{54} . In this chapter we present some basic notation and properties of permutations.

Notation: We may denote a permutation $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ by a $2 \times n$ array:

$$f \leftrightarrow \begin{pmatrix} 1 & 2 & \dots & n \\ f_1 & f_2 & \dots & f_n \end{pmatrix},$$

where f_1, f_2, \dots, f_n is simply a rearrangement (in some order depending on f) of the integers $1, 2, \dots, n$. This notation means that f sends 1 to f_1 , f sends 2 to f_2 , ..., f sends n to f_n . In other words, $f_1 = f(1)$, $f_2 = f(2)$, ..., $f_n = f(n)$, and $f_i \neq f_j$ unless $i = j$.

Example 3.1.1.

- (a) The **identity** permutation, denoted by I , is the permutation which doesn't do anything:

$$I = \begin{pmatrix} 1 & 2 & \dots & n \\ 1 & 2 & \dots & n \end{pmatrix}$$

- (b) The permutation $f : \mathbb{Z}_3 \rightarrow \mathbb{Z}_3$ defined by

$$f = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}$$

simply swaps 1 and 3 and leaves 2 alone.

- (a) The **n -cycle** is a permutation which cyclically permutes the values:

$$\begin{pmatrix} 1 & 2 & \dots & n \\ 2 & 3 & \dots & 1 \end{pmatrix}$$

Imagine an analog clock with the numbers $1, 2, \dots, n$ arranged around the dial. An n -cycle simply moves each number forward (clockwise) by 1 unit. The permutation

$$\begin{pmatrix} 1 & 2 & \dots & n \\ n & 1 & \dots & n-1 \end{pmatrix}$$

is also called an **n -cycle**.

It is easy to construct such permutation in **SAGE**. You simply use the bottom row of the $2 \times n$ array above to make the construction. Here is an example.

— SAGE —

```
sage: p = Permutation([2,3,4,5,1]); p
[2, 3, 4, 5, 1]
sage: p*p*p*p
[5, 1, 2, 3, 4]
```

Definition 3.1.1. Let $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ be a permutation and let

$$e_f(i) = \#\{j > i \mid f(i) > f(j)\}, \quad 1 \leq i \leq n-1.$$

Let

$$\text{swap}(f) = e_f(1) + \dots + e_f(n-1).$$

We call this the **swapping number** (or **length**) of the permutation f since it counts the number of times f swaps the inequality in $i < j$ to $f(i) > f(j)$. If we plot a bar graph of the function f , then $\text{swap}(f)$ counts the number of times the bar at i is higher than the bar at j . We call f **even** if $\text{swap}(f)$ is even and we call f **odd** otherwise.

The number

$$\text{sign}(f) = (-1)^{\text{swap}(f)}$$

is called the sign (or **signum** function) of the permutation f .

Example 3.1.2. Here is an example of SAGE/Python code for implementing the swapping number above. (The third line of the function should be entered all on one line but, for typographical reasons, it has been split in two.)

SAGE

```
def swap(g):
    N = parent(g).degree()
    return sum([len([i2 for i2 in range(i1+1,N+1)\
        if g(i2)<g(i1)]) for i1 in range(1,N)])
```

Here is an example of using this function:

SAGE

```
sage: G = SymmetricGroup(4)
sage: g = G([(1,2,3)]); g
(1,2,3)
sage: swap(g)
2
sage: G = SymmetricGroup(3)
sage: g = G([(1,2)])
sage: swap(g)
1
```

(Note the above manner of inputting a permutation such as $(1,2,3)$ requires the introduction of the ‘symmetric group’. This object will be carefully defined later.)

How is the set of values of the swapping numbers distributed in the collection of all permutations on n letters? SAGE helps us compute this when $n = 6$:

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SAGE

```
sage: G = SymmetricGroup(6)
sage: G.order()
720
sage: swaps = [swap(g) for g in G]
sage: M = max(swaps); M
15
sage: L = [swaps.count(i) for i in range(16)]
sage: list_plot(L).show()
```

Figure 3.1 is a list plot of the swapping numbers vs. their multiplicity for all $6! = 720$ permutations of 6 letters.

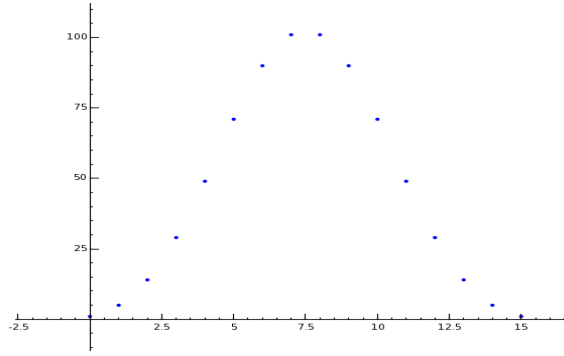


Figure 3.1: Plot of the points $(i, \text{swap}(g_i))$, where i is the i -th element of S_6 .

The reader may verify that the sign function satisfies the following property.

Lemma 3.1.1. *Let $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ be a permutation. Then*

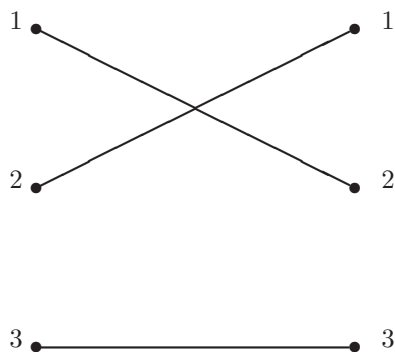
$$\text{sign}(f) = \prod_{i < j \leq n} \frac{f(i) - f(j)}{i - j}.$$

Remark 3.1.1. *Let $S \subset \mathbb{Z}_n$ be a subset. Let $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ be a permutation for which $f(i) = i$, for all $i \in S$. In other words, f permutes the elements of $S^c = \mathbb{Z}_n - S$ but not the elements of S . Define a new permutation g of S^c (not of \mathbb{Z}_n) by $g(i) = f(i)$ for all $i \in S^c$. The above lemma implies that the sign of the permutation $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ is the same as the sign of the permutation $g : S^c \rightarrow S^c$. This will turn out to be a useful fact to keep in mind later.*

Example 3.1.3. *We define the permutation $f : \mathbb{Z}_3 \rightarrow \mathbb{Z}_3$ for which $f(1) = 2, f(2) = 1, f(3) = 3$ by a 2×3 array*

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$$

This swaps 1 and 2 and leaves 3 alone. There are 3 inequalities for \mathbb{Z}_3 : $1 < 2$, $2 < 3$, and $1 < 3$. Applying f to these, we get $f(1) = 2 > f(2) = 1$, $f(2) = 1 < f(3) = 3$, and $f(1) = 2 < f(3) = 3$. Only the first inequality is changed, so the swapping number is $\text{swap}(f) = 1$. Another way to picture f is by a ‘crossing diagram’, where f sends the left-hand column to the right-hand column:



The number of crosses in this diagram is the swapping number of f , from which we can see that f is an odd permutation.

Ponderable 3.1.1. Express $f : \mathbb{Z}_3 \rightarrow \mathbb{Z}_3$ given by $f(1) = 3, f(2) = 1, f(3) = 2$, as (a) a 2×3 array, (b) a crossing diagram. Find its swapping number and sign.

Definition 3.1.2. Let $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ and $g : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ be two permutations. We can compose them to get another permutation, the **composition**, denoted $fg : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$:

$$\begin{array}{ccccc} k & \mapsto & f(k) & \mapsto & g(f(k)) \\ \mathbb{Z}_n & \rightarrow & \mathbb{Z}_n & \rightarrow & \mathbb{Z}_n \end{array}$$

Remark 3.1.2. A Notational Warning! We shall follow standard convention and write our compositions of permutations **left-to-right**. (This contrasts to the right-to-left composition of functions you may have seen in calculus.) When a possible ambiguity may arise, we call this type of composition ‘composition as permutations’ and call ‘right-to-left composition’ the ‘composition as functions’.

When $f = g$ then we write ff as f^2 . In general, we write the n -fold composition $f \dots f$ (n times) as f^n . Every permutation f has the property that there is some integer $N > 0$, which depends on f , such that $f^N = 1$. That is to say, if you repeatedly apply a permutation enough times you will eventually obtain the identity permutation. In particular, if you can perform the same exact shuffle on a deck of cards repeatedly, you will eventually get back to the original ordering of the card deck.

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This is vaguely reminiscent of a theorem due to Persi Diaconis which, roughly speaking, says that you don't need more than 7 shuffles to scramble a card deck. The two situations are almost completely different though, because in one case you are trying to scramble the ordering up and in the other case you are trying to get back to the original ordering.

The mathematician and statistician Persi Diaconis left home at 14 to become a professional magician and dropped out of high school. He received a MacArthur Fellowship ('genius award') in 1979 and is currently Mary V. Sunseri professor of statistics at Stanford University.

Lemma 3.1.2. *Let f and g be permutations $\mathbb{Z}_n \rightarrow \mathbb{Z}_n$. $\text{sign}(fg) = \text{sign}(f)\text{sign}(g)$.*

Ponderable 3.1.2. *For f, g as in Lemma 3.1.1, we have*

$$\prod_{i < j \leq n} \frac{fg(i) - fg(j)}{i - j} = \prod_{i < j \leq n} \frac{g(i) - g(j)}{i - j} \prod_{i < j \leq n} \frac{fg(i) - fg(j)}{g(i) - g(j)}.$$

Use this to prove Lemma 3.1.2.

Definition 3.1.3. *The smallest integer $N > 0$ such that $f^N = 1$ is called the **order** of f .*

For an example from science fiction, suppose you take off in a space ship which goes faster than the speed of light. Suppose it is so fast that in one trip you come back exactly one hour earlier, according to your wrist watch, than everyone else! In other words, your hourly times are the same as everyone else's, except that they have been shifted by 1. (This type of permutation is called a **cyclic permutation**.) How many trips do you need to take before your watch reads the same as everyone else's again? The answer is 24 trips total (assuming everyone is using 'military time').

Example 3.1.4. *Let*

$$f = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} \quad g = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix}$$

be permutations of \mathbb{Z}_3 . We have

$$fg = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 3 & 2 \end{pmatrix}, \quad f^2 = 1, \quad g^3 = 1.$$

Ponderable 3.1.3. *Compute (a) fg and (b) the order of f and the order of g , where*

$$(a) \quad f = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix} \quad g = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix}$$

$$(b) \quad f = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix} \quad g = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$$

Ponderable 3.1.4. *If f, g, h are permutations of \mathbb{Z}_n , is $(fg)h = f(gh)$ (here composition is as functions; see Remark 3.1.2)? In other words, we have three permutations: I have f, g and you have h . First, you perform h and then I perform g then f . Second, I give you g and you perform h then g and then I perform f . Are these the same? Explain.*

3.2 Inverses

‘I don’t understand you’, said Alice. ‘It’s dreadfully confusing!’

‘That’s the effect of living backwards’, the Queen said kindly: ‘it always makes one a little giddy at first—’

‘Living backwards!’ Alice repeated in great astonishment. ‘I never heard of such a thing!’

‘—but there’s one great advantage in it, that one’s memory works both ways.’

Lewis Carroll, Through the looking glass

Try to visualize in your mind the graph of a function $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$. We can think of this as either a collection of points $(i, f(i))$, $i = 1, 2, \dots, n$, or as a bar graph. In any case, we can look at this graph of f and determine

- (a) if f is injective,
- (b) if f is surjective,
- (c) the inverse f^{-1} , if it exists.

How? Well, from the graph of f we can determine the image $f(\mathbb{Z}_n)$ and this determines if f is surjective or not. The inverse exists only if f is injective (hence, in our case, surjective by Ponderable 2.1.6). Its graph is determined by reflecting the graph of f about the diagonal, $x = y$.

Lemma 3.2.1. *The following statements are equivalent:*

- (1) $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ is injective.
- (2) $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ is surjective.
- (3) $|f(\mathbb{Z}_n)| = |\mathbb{Z}_n|$.

Proof: The equivalence of the first two statements is left to the interested reader as an exercise (hints: first show (1) implies (2) using reductio ad absurdum, then show (2) implies (1), again using reductio ad absurdum). Statement (2) is equivalent to (3) by the definition of surjectivity. \square

Example 3.2.1. *The inverse of*

$$\begin{pmatrix} 1 & 2 & 3 \\ 3 & 1 & 2 \end{pmatrix}$$

is obtained by reflecting its graph about the $x = y$ line.

3.2. INVERSES

Ponderable 3.2.1. Graph and determine the inverses of the following permutations:

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}, \begin{pmatrix} 1 & 2 & 3 & 4 \\ 2 & 3 & 4 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 1 & 5 & 3 & 4 \end{pmatrix}$$

There are two more commonly used ways of expressing a permutation. The first is the ‘matrix notation’:

Definition 3.2.1. To a permutation $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$, given by

$$f = \begin{pmatrix} 1 & 2 & \dots & n \\ f(1) & f(2) & \dots & f(n) \end{pmatrix},$$

we associate the matrix $P(f)$ of 0's and 1's defined as follows: the ij^{th} entry of $P(f)$ is 1 if $j = f(i)$ and is 0 otherwise.

Example 3.2.2. The matrix of the permutation f given by

$$f = \begin{pmatrix} 1 & 2 & 3 \\ 1 & 2 & 3 \end{pmatrix}$$

is

$$P(f) = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

Definition 3.2.2. A square matrix which has exactly one 1 per row and per column (as $P(f)$ does) is called a **permutation matrix**.

Lemma 3.2.2. There are $n!$ distinct $n \times n$ permutation matrices and there are $n!$ distinct permutations of the set $\{1, 2, \dots, n\}$.

Ponderable 3.2.2. Prove this using the multiplication principle from Theorem 2.4.2 in the previous chapter.

Example 3.2.3. The matrix of the permutation f given by

$$f = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$$

is

$$P(f) = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

Note that matrix multiplication gives

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 \\ 2 \\ 3 \end{pmatrix} = \begin{pmatrix} 2 \\ 1 \\ 3 \end{pmatrix},$$

from which we can recover the 2×3 array.

It is not hard to do this computation in **SAGE** as well:

SAGE

```

sage: V = VectorSpace(QQ, 3)
sage: v = V([1, 2, 3])
sage: G = SymmetricGroup(3)
sage: g = G([(1, 2)])
sage: v; permutation_action(g, v)
(1, 2, 3)
(2, 1, 3)
sage: g.matrix()

[0 1 0]
[1 0 0]
[0 0 1]

```

How does the permutation function relate to the permutation matrix? The theorem below explains parts of this relationship.

Theorem 3.2.1. *If $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ is a permutation then*

(a)

$$P(f) \begin{pmatrix} 1 \\ 2 \\ \vdots \\ n \end{pmatrix} = \begin{pmatrix} f(1) \\ f(2) \\ \vdots \\ f(n) \end{pmatrix},$$

(b) $P(f)^{-1} = P(f^{-1})$ (the inverse of the permutation matrix is the matrix of the inverse of the permutation),

(c) $P(fg) = P(f)P(g)$ (the permutation matrix of the product is the product of the permutation matrices),

(d) $\text{sign}(f) = \det(P(f))$.

Proof: If \vec{v} is the column vector with entries v_1, v_2, \dots, v_n (the v_i are arbitrary real numbers) then $P(f)\vec{v}$ is the column vector whose i th coordinate is equal to v_j if f sends i to j . Since, in this case, $j = f(i)$ (here we write $f(i)$ to denote the image of i under the permutation f , even though i really gets plugged into f on the left), this implies that $P(f)\vec{v}$ is the column vector with entries $v_{f(1)}, v_{f(2)}, \dots, v_{f(n)}$. This proves (a).

Note that (b) is a consequence of (c) so we need only prove (c). We compute $P(fg)\vec{v}$ and $P(f)P(g)\vec{v}$. By the same reasoning as in (a), we find that the i th coordinate of $P(fg)\vec{v}$ is $v_{(fg)(i)}$. Similarly, the i th coordinate of $P(g)\vec{v}$ is $v'_i = v_{g(i)}$. Therefore, the i th coordinate of $P(f)(P(g)\vec{v})$ is $v'_{f(i)} = v_{g(f(i))} = v_{(fg)(i)}$. This implies $P(fg)\vec{v} = P(f)P(g)\vec{v}$. Since the v_i were arbitrary real numbers, this implies the theorem. \square

3.2. INVERSES

Example 3.2.4. *Let*

$$f = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}, \quad g = \begin{pmatrix} 1 & 2 & 3 \\ 3 & 2 & 1 \end{pmatrix}, \quad h = \begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix},$$

so $f = f^{-1}$, $g = g^{-1}$, $h = fg$. Moreover,

$$P(g) = \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}, \quad P(h) = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix}.$$

A direct matrix calculation verifies that $P(f)P(g) = P(fg) = P(h)$ and $P(h^{-1}) = P(g^{-1}f^{-1}) = P(g^{-1})P(f^{-1}) = P(g)^{-1}P(f)^{-1}$, as predicted by the above theorem.

The matrix can be determined from the graph of the function $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ as follows. First, take the graph of the function f and flip it about the x -axis. In this $n \times n$ grid of integral points (x, y) , the 1st coordinate goes down, the 2nd coordinate goes across, and x and y are integers between 1 and n inclusive. (This is not to be confused with the orientation of the ‘axes’ in the xy -plane.) Now, fill in all the plotted points with 1’s and all the unplots points with 0’s. The resulting $n \times n$ array is the matrix $P(f)$.

Example 3.2.5. *For example, consider the*

$$f = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 1 & 3 & 2 & 1 \end{pmatrix},$$

whose graph and mirror image look like Figure 3.2.

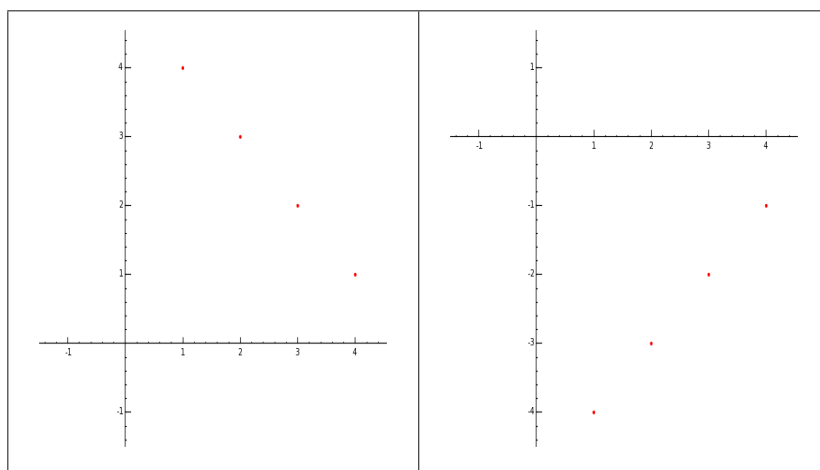


Figure 3.2: graph and mirror image of f .

The matrix associated to this permutation is

$$P(f) = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix},$$

as in the mirror image graph.

The set of all permutations of \mathbb{Z}_n is denoted S_n and is called the **symmetric group** on n letters. When solving the Rubik's Cube it is frequently the case that you are making several moves of the following type: make a move, m_1 say, then another move m_2 , then make the inverse of the first move, m_1^{-1} . For example, $(R^{-1}D^2RB^{-1}U^2B)^2$ is composed of such moves. This move is a two-corner twist: the *urf* corner (in the Singmaster notation, explained in Example 3.2.6 below) is twisted once clockwise, and the corner *bld* is twisted once counterclockwise, as illustrated in Figure 3.3.

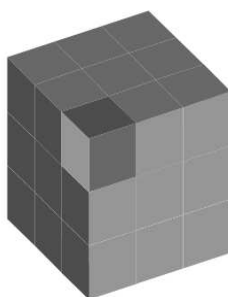


Figure 3.3: $(R^{-1}D^2RB^{-1}U^2B)^2$ permutes edges and faces but not vertices.

In general, to compute the effect of a move like $m_1m_2m_1^{-1}$, you need to know how it moves each vertex and the edge. Rubik cubers will often make implicit use of the the following lemma when solving their cube.

Lemma 3.2.3. *Let $r \in S_n$ denote any permutation and let i, j denote distinct integers belonging to $\{1, 2, \dots, n\}$. Let s denote a permutation sending i to j :*

$$s(i) = j.$$

Then $s^r = r^{-1}sr$ is a permutation sending $r(i)$ to $r(j)$:

$$s^r(r(i)) = r(j).$$

More specifically (and this is the case in which this lemma is most often applied in this book): let i_1, i_2, \dots, i_k denote distinct integers belonging to $\{1, 2, \dots, n\}$. Let s denote the permutation sending i_j to i_{j+1} :

$$s(i_j) = i_{j+1}, \quad 1 \leq j < k, \quad s(i_k) = i_1, \quad s(m) = m, \quad \forall m \notin \{i_1, \dots, i_k\}.$$

3.2. INVERSES

Then $s^r = r^{-1}sr$ is the permutation sending $r(i_j)$ to $r(i_{j+1})$:

$$s^r(r(i_j)) = r(i_{j+1}), \quad 1 \leq j < k, \quad s^r(r(i_k)) = r(i_1),$$

$$s^r(m) = m, \quad \forall m \notin \{r(i_1), \dots, r(i_k)\}.$$

Example 3.2.6. Let us label the 12 edges of the Rubik's Cube using the **Singmaster notation**:

- uf denotes the 'up, front edge',
- ul denotes the 'up, left edge',
- ur denotes the 'up, right edge',
- ub denotes the 'up, back edge',
- df denotes the 'down, front edge',
- dl denotes the 'down, left edge',
- dr denotes the 'down, right edge',
- db denotes the 'down, back edge',
- fl denotes the 'front, left edge',
- fr denotes the 'front, right edge',
- bl denotes the 'back, left edge',
- br denotes the 'back, right edge'.

The corners are denoted similarly (ufr denotes the 'up, front, right corner', and so on).

If you have a Rubik's Cube move s which is a 3-cycle on 3 particular edges, say

$$uf \mapsto ul \mapsto ur \mapsto uf,$$

and another move r which sends these edges somewhere else, say $r = F^2$ so that $r : uf \mapsto df$ but leaves the other edges alone, then $r^{-1}sr$ is the permutation

$$df \mapsto ul \mapsto ur \mapsto df.$$

Try it!

A proof of this lemma will be given in chapter 9 (see Theorem 9.3.1).

3.3 Cycle notation

The most common notation for a permutation is the ‘cycle notation’, due to Cauchy in the mid 1800’s. This notation is more compact than the array notation we’ve been using, and from this point on we will switch over to the cycle notation. If $a_1, a_2, \dots, a_r \in \mathbb{Z}_n$ are distinct then the symbol

$$(a_1, a_2, \dots, a_r)$$

(some r less than or equal to n) denotes the permutation f of \mathbb{Z}_n which sends a_1 to a_2 , sends a_2 to a_3 , ..., sends a_{r-1} to a_r , sends a_r to a_1 , and leaves all the other numbers in \mathbb{Z}_n alone. In other words,

$$f(a_1) = a_2, f(a_2) = a_3, \dots, f(a_r) = a_1,$$

and $f(i) = i$, if i is not equal to one of the a_1, \dots, a_r . Such a permutation is called **cyclic**. The number r is called the **length** of the cycle.

We call two such cycles (a_1, a_2, \dots, a_r) and (b_1, b_2, \dots, b_t) **disjoint** if the sets $\{a_1, a_2, \dots, a_r\}$ and $\{b_1, b_2, \dots, b_t\}$ are disjoint.

Lemma 3.3.1. *If f and g are disjoint cyclic permutations of \mathbb{Z}_n then $fg = gf$.*

Proof: This is true because the permutations f and g of \mathbb{Z}_n affect disjoint collections of integers, so the permutations may be performed in either order. \square

Lemma 3.3.2. *The cyclic permutation (a_1, a_2, \dots, a_r) has order r .*

Proof: Note $f(a_1) = a_2$, $f^2(a_1) = a_3$, ..., $f^{r-1}(a_1) = a_r$, $f^r(a_1) = a_1$, by definition of f . Likewise, for any $i = 1, \dots, r$, we have $f^r(a_i) = a_i$. \square

Definition 3.3.1. A **transposition** is a cycle (i, j) of length 2 which interchanges i and j ($i \neq j$).

Theorem 3.3.1. *Every permutation $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$ is the product of disjoint cyclic permutations. More precisely, if f is a permutation of $\{1, 2, \dots, n\}$ (with $n > 1$) then there are non-empty disjoint subsets of distinct integers*

$$\begin{aligned} A_1 &= \{a_{11}, \dots, a_{1,r_1}\} \subset \{1, 2, \dots, n\}, \\ A_2 &= \{a_{21}, \dots, a_{2,r_2}\} \subset \{1, 2, \dots, n\}, \\ &\vdots \\ A_k &= \{a_{k1}, \dots, a_{k,r_k}\}, \end{aligned}$$

such that

$$\{1, 2, \dots, n\} = A_1 \cup \dots \cup A_k, \quad n = r_1 + r_2 + \dots + r_k,$$

and

$$f = (a_{11}, \dots, a_{1,r_1}) \dots (a_{k1}, \dots, a_{k,r_k}).$$

3.3. CYCLE NOTATION

This product is called a **cycle decomposition** of f . If we rearrange the cardinalities r_i of these sets A_i in decreasing order, say we write this as

$$r'_1 \geq r'_2 \geq \dots \geq r'_k,$$

then the k -tuple (r'_1, \dots, r'_k) is called the **cycle structure** of f and f is called a (r'_1, \dots, r'_k) -**cycle**. For example, $(1, 2)(3, 4, 5)$ is a $(3, 2)$ -cycle in S_5 .

Proof: The proof is constructive.

Let $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ be a permutation. List all the *distinct* elements in \mathbb{Z}_n you get by repeatedly applying f to 1. Call them (for lack of a better notation!), $a_{10}, a_{11}, a_{12}, \dots, a_{1,r_1}$, where,

$$O_1 = \{a_{10} = 1, a_{11} = f(1), a_{12} = f^2(1), \dots, a_{1,r_1} = f^{r_1}(1)\}.$$

(Incidentally, this is called the **orbit of 1 under f** . Like the Earth orbiting the sun, you keep moving 1 around using f until it gets back to itself. We shall discuss orbits later, in §5.10.) Now list the elements in the ‘orbit of 2’:

$$O_2 = \{a_{20} = 2, a_{21} = f(2), a_{22} = f^2(2), \dots, a_{2,r_2} = f^{r_2}(2)\},$$

and so on until we get to the ‘orbit of n ’:

$$O_n = \{a_{n0} = n, a_{n1} = f(n), a_{n2} = f^2(n), \dots, a_{n,r_n} = f^{r_n}(n)\}.$$

An example: let $f : \mathbb{Z}_3 \rightarrow \mathbb{Z}_3$ be the 3-cycle permutation, $f(1) = 2, f(2) = 3, f(3) = 1$. In this case, $O_1 = \{1, 2, 3\}$, $O_2 = \{2, f(2) = 3, f(f(2)) = f(3) = 1\}$, and $O_3 = \{3, f(3) = 1, f(f(3)) = f(1) = 2\}$, so all three orbits are equal to each other in this example.

In general, if you pick any two of these n orbits O_1, \dots, O_n , they will either be the same or disjoint. Denote all the *distinct* orbits by O'_1, \dots, O'_k . (The O'_1, \dots, O'_k are a subsequence of the O_1, \dots, O_n . It doesn't matter what order you write the O'_i 's in or in what order you write the elements in each individual orbit.) Suppose that

$$\begin{aligned} O'_1 &= \{b_{11}, \dots, b_{1,s_1}\} & \text{so } |O'_1| &= s_1, \\ O'_2 &= \{b_{21}, \dots, b_{2,s_2}\} & \text{so } |O'_2| &= s_2, \\ &\vdots \\ &\vdots \\ O'_k &= \{b_{k1}, \dots, b_{k,s_k}\} & \text{so } |O'_k| &= s_k. \end{aligned}$$

(The a_{ij} 's have been relabeled as b_{ij} 's to try to simplify the notation.) In this case,

$$\mathbb{Z}_n = \cup_{i=1}^k O'_i = O'_1 \cup \dots \cup O'_k,$$

and $s_1 + s_2 + \dots + s_k = n$. The restriction of f to O'_1 , denoted $f_{O'_1} : O'_1 \rightarrow O'_1$, is equal to the s_1 -cycle $(b_{11}, b_{12}, \dots, b_{1,s_1})$. In general, the restriction of f to O'_j , denoted $f_{O'_j} : O'_j \rightarrow O'_j$, is equal to the s_j -cycle $(b_{j1}, b_{j2}, \dots, b_{j,s_j})$. Since the O'_j 's partition \mathbb{Z}_n , the definition of f and our construction implies that

$$f = (b_{11}, b_{12}, \dots, b_{1,s_1}) \dots (b_{k1}, b_{k2}, \dots, b_{k,s_k}).$$

□

Example 3.3.1.

- The cycle notation for

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix}$$

is $(1, 2)(3)$ or simply $(1, 2)$. In general, if any of the orbits O_j in the above construction is a singleton, it is often omitted from the notation, with the implicit understanding that f doesn't permute the omitted numbers.

- The cycle notation for

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 3 & 1 \end{pmatrix}$$

is $(1, 2, 3)$.

- The cycle notation for

$$\begin{pmatrix} 1 & 2 & 3 & 4 \\ 3 & 4 & 1 & 2 \end{pmatrix}$$

is $(1, 3)(2, 4) = (2, 4)(1, 3)$.

- The cycle notation for

$$\begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 3 & 4 & 1 & 5 & 2 \end{pmatrix}$$

is $(1, 3)(2, 4, 5) = (4, 5, 2)(1, 3)$.

Here is an example of using **SAGE** to compute with elements in the cycle notation:

SAGE

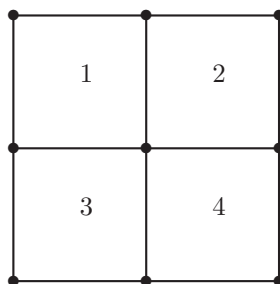
```
sage: G = SymmetricGroup(4)
sage: g1 = G([(2,4),(1,3)])
sage: g2 = G([(1,2,3,4)])
sage: matrix([[1,2,3,4],[g1(i) for i in [1,2,3,4]]])
[1 2 3 4]
[3 4 1 2]
sage: matrix([[1,2,3,4],[g2(i) for i in [1,2,3,4]]])
[1 2 3 4]
[2 3 4 1]
sage: g1*g2
(1,4,3,2)
sage: matrix([[1,2,3,4],[(g1*g2)(i) for i in [1,2,3,4]]])
[1 2 3 4]
[4 1 2 3]
sage: g2*g1
(1,4,3,2)
```

3.3. CYCLE NOTATION

Note the manner of inputting a permutation such as $(2, 4)(1, 3)$: first you must define the group G it belongs to (the `SymmetricGroup(4)` is **SAGE** notation for S_4 , the permutations of $\{1, 2, 3, 4\}$), then enter $(2, 4)(1, 3)$ as a list of disjoint cycles $[(2, 4), (1, 3)]$ which are “coerced” into G using the command `G([(2, 4), (1, 3)])`. (More examples will be given below but you can also find more details in the section “Permutation group elements” in the **SAGE** reference manual [S].)

Ponderable 3.3.1. The disjoint cycle decomposition of $(2, 3, 7)(3, 7, 10)$ is either $(2, 3)(7, 10)$ or $(2, 7)(3, 10)$. Which one is it?

Ponderable 3.3.2. Divide a square into 4 subsquares (‘facets’) and label them 1, 2, 3, 4. For example,



Let r denote counterclockwise rotation by 90 degrees. Then, as a permutation on the facets, $r = (1, 3, 4, 2)$. Let f_x denote reflection about the horizontal line dividing the square in two, let f_y denote reflection about the vertical line dividing the square in two. Use the cycle notation to determine the permutations of the facets (a) r^2 , (b) r^3 , (c) f_x , (d) f_y , (e) $f_x r f_x$, (f) $f_x f_y$.

For those wanting to do this using **SAGE**, here is an example.

SAGE

```
sage: G = SymmetricGroup(4)
sage: r = G([(1,3,4,2)]); r
(1,3,4,2)
sage: r*r
(1,4)(2,3)
```

This tells you that the answer to (a) is $(1, 4)(2, 3)$.

Ponderable 3.3.3. *Label the 24 facets of the $2 \times 2 \times 2$ Rubik's Cube as follows:*

		1	2		
		U			
		3	4		
5	6	9	10	13	14
L		F		R	
7	8	11	12	15	16
		21	22		
		D			
		23	24		
				17	18
				19	20

(You may want to xerox this page and then cut this cube out and tape it together for this.) Let X denote rotation clockwise by 90 degrees of the face labeled x , where $x \in \{r, l, f, b, u, d\}$ (so, for example, if $x = f$ then $X = F$). Use the cycle notation to determine the permutations of the facets given by (a) R , (b) L , (c) F , (d) B , (e) U , (f) D .

Lemma 3.3.3. *A cyclic permutation is even if and only if the length of its cycle is odd. A general permutation $f : \mathbb{Z}_n \rightarrow \mathbb{Z}_n$ is odd if and only if the number of cycles of even length in its disjoint cycle decomposition is odd.*

This follows from the definition of an even/odd permutation (see Definition 3.1.1), the fact that $\text{sign}(p_1 p_2 \dots p_k) = \text{sign}(p_1) \text{sign}(p_2) \dots \text{sign}(p_k)$, for permutations p_i (because of Lemma 3.1.2), and the fact that any k -cycle can be written as a product of $k - 1$ transpositions,

$$(a_1, a_2, \dots, a_k) = (a_1, a_k)(a_1, a_{k-1}) \dots (a_1, a_2). \quad (3.1)$$

(Equation (3.1) can be proven by mathematical induction. The argument is left to the reader. For details, refer to the related results in chapter 1 of Rotman [R] or Theorem 3.3 in Gaglione [G], §3.2, for example.)

3.4. AN ALGORITHM TO LIST ALL THE PERMUTATIONS

Lemma 3.3.4. *The order of a permutation is the least common multiple (lcm) of the lengths r_1, r_2, \dots, r_k of the disjoint cycles in its cycle decomposition.*

Example 3.3.2. *The order of $(1,3)(2,4)$ is 2. It is even. The order of $(1,3)(2,4,5)$ is 6. It is odd.*

Here is an example of using SAGE to compute with orders and signs of permutations.

SAGE

```
sage: G = SymmetricGroup(4)
sage: g1 = G([(2,4),(1,3)])
sage: g2 = G([(1,2,3,4)])
sage: g2*g1
(1,4,3,2)
sage: g1.sign()
1
sage: g2.sign()
-1
sage: (g1*g2).sign()
-1
sage: g1.order()
2
sage: g2.order()
4
sage: (g1*g2).order()
4
```

3.4 An algorithm to list all the permutations

Martin Gardner [Gar1] mentions an algorithm that lists all the permutations of $\{1, 2, \dots, n\}$. This algorithm gives the fastest known method of listing all permutations of $\{1, 2, \dots, n\}$. Rediscovered many times since, the procedure is due originally to the Polish mathematician Hugo Steinhaus (1887-1972). A student of David Hilbert at Göttingen, he did important work on orthogonal series, probability theory, real functions and their applications.

We shall denote each permutation by the second row in its $2 \times n$ array notation. For example, in the case $n = 2$

$$\begin{array}{cc} 1 & 2 \\ 2 & 1 \end{array}$$

are the permutations.

To see the case $n = 3$, the idea is to

3.4. AN ALGORITHM TO LIST ALL THE PERMUTATIONS

(a) write down each row $n = 3$ times each as follows:

```

1  2
1  2
1  2
2  1
2  1
2  1

```

(b) ‘weave’ in a 3 as follows

```

1  2  3
1  3  2
3  1  2
3  2  1
2  3  1
2  1  3

```

In case $n = 4$, the idea is to

(a) write down each row 4 times each as follows (go down the first column then down the second column):

1 2 3	3 2 1
1 2 3	3 2 1
1 2 3	3 2 1
1 2 3	3 2 1
1 3 2	2 3 1
1 3 2	2 3 1
1 3 2	2 3 1
1 3 2	2 3 1
3 1 2	2 1 3
3 1 2	2 1 3
3 1 2	2 1 3
3 1 2	2 1 3

(b) now ‘weave’ a 4 in:

1 2 3 4	4 3 2 1
1 2 4 3	3 4 2 1
1 4 2 3	3 2 4 1
4 1 2 3	3 2 1 4
4 1 3 2	2 3 1 4
1 4 3 2	2 3 4 1
1 3 4 2	2 4 3 1
1 3 2 4	4 2 3 1
3 1 2 4	4 2 1 3
3 1 4 2	2 4 1 3
3 4 1 2	2 1 4 3
4 3 1 2	2 1 3 4

The next subsection (§3.4.1) describes the general procedure. These constructions lead to the following result.

3.4. AN ALGORITHM TO LIST ALL THE PERMUTATIONS

Theorem 3.4.1. (*Steinhaus*) *There is a sequence of (not necessarily distinct) 2-cycles, $(a_1, b_1), \dots, (a_N, b_N)$, where $N = n! - 1$, such that each non-trivial permutation f of $\{1, 2, \dots, n\}$ may be expressed in the form*

$$f = \prod_{i=1}^k (a_i, b_i)$$

for some k , $1 \leq k \leq N$. These transpositions (a_i, b_i) , for $k = 1, 2, \dots, N$, are not necessarily disjoint.

In particular, each permutation may be written as a product of (not necessarily disjoint) 2-cycles. This theorem says much more. It says the set of all permutations can be generated by successive multiplication by a certain transposition. This result will be proven in the next section.

There is an analogous result valid only for *even* permutations: each even permutation may be written as a product of (not necessarily disjoint) 3-cycles. This will be stated more precisely (and proved) in the next section; see Proposition 9.4.1.

3.4.1 Why Steinhaus's algorithm works

The argument is by mathematical induction.

1. First, note that Steinhaus's algorithm works for $n = 2, n = 3$.
2. Write down all the elements of S_{n-1} as a list using Steinhaus's algorithm (by the induction hypothesis, each element differs from the previous one by a suitable transposition in this list). We represent each element in the $2 \times (n-1)$ array notation. In our list, we simply record the 2nd row in the list. i.e., as an $(n-1)$ -tuple. There are $(n-1)!$ elements in this list.
3. For the 1st $(n-1)$ -tuple, write an n at the end, for the 2nd $(n-1)$ -tuple, write an n at the beginning, for the 3rd $(n-1)$ -tuple, write an n at the end, for the 4th $(n-1)$ -tuple, write an n at the beginning, and so on. Note we have a sequence of $(n-1)!$ n -tuples.
4. Suppose the i th n -tuple is $(n, a_1, a_2, \dots, a_{n-1})$. First, act on this by the transposition (n, a_1) , then by the transposition (n, a_2) , then by (n, a_3) , ..., by (n, a_{n-1}) . The result of the last one is $(a_1, a_2, \dots, a_{n-1}, n)$. Suppose the j th n -tuple is $(b_1, b_2, \dots, b_{n-1}, n)$. (Note that by the induction hypothesis, $(a_1, a_2, \dots, a_{n-1}, n)$ and $(b_1, b_2, \dots, b_{n-1}, n)$ 'differ' only by a suitable transposition, in S_{n-1} .) First, act on this by the transposition (b_1, n) , then by the transposition (b_2, n) , then by (b_3, n) , ..., by (b_{n-1}, n) . The result of the last one is $(n, b_1, b_2, \dots, b_{n-1})$.
5. Note that all these (for i, j ranging over even or odd integers from 1 to $n-1$), the resulting $n!$ n -tuples are distinct. This yields a listing of S_n as desired.

3.4.2 A side order of dessert: cake cutting

Though somewhat off the topic, another problem connected with Steinhaus is worth mentioning. In 1944 Steinhaus proposed the problem of dividing a cake into n pieces so that it is ‘proportional’ (i.e., each person is satisfied with his or her share) and ‘envy-free’ (i.e., each person is satisfied that nobody is receiving more than a fair share). For $n = 2$, one person cuts the cake, and the other chooses her piece. Steinhaus found a proportional but *not* envy-free solution for $n = 3$. An envy-free solution to Steinhaus’s problem for $n = 3$ was found in 1962 by John H. Conway (whom we will run into again later in this book) and, independently, by John Selfridge. For general n the problem was solved by Steven Brams and Alan Taylor in 1995. More history and details can be found online in the Wikipedia [Wi] article on ‘fair division’.

3.5 Permutations and bell ringing

No man is an island, entire of itself ... any man’s death diminishes me,
because I am involved in mankind; and therefore never send to know for
whom the bell tolls; it tolls for thee.

John Donne (English clergy-man and poet, 1572-1631)

Since the seventeenth century, and possibly before, cathedral bells in England have been rung by permuting the order of a ‘round’ of bells. The art and study of such bell ringing is referred to as campanology.

Fabian Stedman provided significant contributions to bell ringing. Born in 1640 to Reverend Francis Stedman, Fabian Stedman’s connections with campanology took root at the early age of 15 when he moved to London to work as an apprentice to a master printer. While in London Stedman joined a bell-ringing society known as the ‘Society of Colledg [sic] Youths’, which has since been renamed the Ancient Society of College Youths and is still in existence today. Stedman’s major contributions to campanology are reflected in his efforts in **Tintinnalogia** and **Campanologia**, the first two books published on the subject, in 1668 and 1677, respectively.

Below is a glossary of a few essential terms:

- **Change:** the swapping of one or more disjoint pairs of adjacent. bells
- **Plain change:** the swapping of one pair of adjacent bells only.
- **Cross change:** the swapping of more than one pair of bells.
- **Round:** an ordering of the bells (i.e., a permutation of $(1, 2, 3, \dots, n)$).

In the beginning, change ringing concerned itself with a single row of bells whose order could be denoted by $(1, 2, 3, \dots, n)$. Considering the case where

3.5. PERMUTATIONS AND BELL RINGING

$n = 6$ the concepts of plain and cross changes can be understood more clearly. If we use only plain changes we can generate permutations of the bells as follows:

1	2	3	4	5	6
2	1	3	4	5	6
2	1	4	3	5	6
2	1	4	3	6	5

It should be fairly obvious on inspection that the first plain change swaps 1 and 2, the second swaps 3 and 4, and the third swaps 5 and 6. Considering the same set of six bells acted upon by a cross change, the same result is achieved in one change, as seen below:

1	2	3	4	5	6
2	1	4	3	6	5

More useful and interesting patterns can be generated by combining plain and cross changes. The **plain lead on four bells** is one of the most simplistic patterns and was devised sometime around 1621 by alternating consecutive cross and plain changes as seen below:

1	2	3	4
2	1	4	3
2	4	1	3
4	2	3	1
4	3	2	1
3	4	1	2
3	1	4	2
1	3	2	4
1	2	3	4

The pattern that defines the plain lead on four bells is nothing more than a cross change followed by a plain change on the middle two bells until we reach the round, which is where we started.

Generating the permutations contained in the plain lead on four bells can be easily described using the notation for permutations we have developed. We begin by representing the cross change as $a = (1, 2)(3, 4)$, which swaps the first two and last two bells, and representing the plain change as $b = (2, 3)$, which swaps the middle pair. We begin with the first element, a . To generate the next permutation we multiply this first element by b . To generate the third element we simply multiply this second term, ab , by a to get aba . Continuing on in this manner we multiply alternately by a then b to generate the set $D_4 = \{a, ab, aba, (ab)^2, (ab)^2a, (ab)^3, (ab)^3a, (ab)^4\}$. This is the set of permutations in the plain lead on four bells. (It is denoted D_4 here for reasons which will be explained later.) Since $(ab)^4$ yields the original round, we say $(ab)^4 = 1$ and $D_4 = \{1, a, ab, aba, (ab)^2, (ab)^2a, (ab)^3, (ab)^3a\}$.

Now for a more complex example. We turn our attention now to the composition which is commonly referred to as **Plain Bob Minimus**. Plain Bob

3.5. PERMUTATIONS AND BELL RINGING

Minimus begins at the round and ends at the round $(1, 2, 3, 4)$ and contains all possible permutations of these four bells (in a particular order, described below):

1 2 3 4	1 3 4 2	1 4 2 3
2 1 4 3	3 1 2 4	4 1 3 2
2 4 1 3	3 2 1 4	4 3 1 2
4 2 3 1	2 3 4 1	3 4 2 1
4 3 2 1	2 4 3 1	3 2 4 1
3 4 1 2	4 2 1 3	2 3 1 4
3 1 4 2	4 1 2 3	2 1 3 4
1 3 2 4	1 4 3 2	1 2 4 3
		1 2 3 4

We can now describe this composition using the permutation notation as we did for the plain lead on four bells. Let $a = (1, 2)(3, 4)$ and $b = (2, 3)$ represent possible changes between rows. If we look at the first column of the Plain Bob Minimus composition, we see that it is nothing more than the plain lead on four bells. To generate the second column, we introduce a new permutation $c = (3, 4)$ and we simplify our notation by letting $k = (ab)^3ac$. Multiplying through, we generate the second column:

$$\{k, ka, kab, kaba, k(ab)^2, k(ab)^2a, k(ab)^3, k(ab)^3a\}.$$

Multiplying through by k again, we generate the third column:

$$\{k^2, k^2a, k^2ab, k^2aba, k^2(ab)^2, k^2(ab)^2a, k^2(ab)^3, k^2(ab)^3a\}.$$

This generation of Plain Bob Minimus shows that it can be expressed as the disjoint union of ‘translations’ of the plain lead on four bells! We shall explain this fact using group theory in the next chapter.

Now, since we chose $a = (1, 2)(3, 4)$, $b = (2, 3)$, and $c = (3, 4)$, where b and c are obviously by definition 2-cycle or transpositions and a is the product of two such 2-cycles or transpositions, we have shown a further result, that each permutation of \mathbb{Z}_4 can be written as a product of 2-cycles. More generally, we can state the following theorem originally due to H. Steinhaus.

Theorem 3.5.1. *Let f be a member of S_n , i.e., let f be any permutation of degree n . Then f can be written as a product of transpositions.*

To sketch a proof of this theorem (following [G]) and hence prove Theorem 3.4.1 as promised, we need only to recall that every permutation of S_n can be written uniquely (up to order) as a product of disjoint cycles (Theorem 3.3.1 above).

Proof: Any k -cycle can be written as a product of $k - 1$ transpositions, as in equation (3.1). We see that since any permutation can be written in terms of cycles, and any cycle can be written as product of transposition, it follows that every permutation of \mathbb{Z}_n can be written as a product of transpositions. \square

3.5. PERMUTATIONS AND BELL RINGING

It was a knowledge of the permutations in Kent Treble Bob Major, a ringing of eight bells, which helped Lord Peter Wimsey solve a murder mystery in Dorothy Sayer's **The Nine Tailors** [Sa].

A good reference for this section is A. White [Wh].

Chapter 4

A procession of permutation puzzles

How can it be that mathematics, being after all a product of human thought independent of experience, is so admirably adapted to the objects of reality?

Albert Einstein

I shall describe several permutation puzzles in this chapter. After presenting a lot of mathematics, we shall finally have a chance to have fun attacking some real puzzles! Though Einstein was talking about relativity in his quote, you'll see for yourself how admirably adapted combinatorics (and later group theory) is to describing these puzzles.

First, we nail down for the sake of discussion a definition of what a permutation puzzle is.

A **one-person game** is a sequence of moves that follow certain rules:

- There are finitely many moves at each stage.
- There is a finite sequence of moves that yields a solution.
- There are no chance or random moves.
- There is complete information about each move.
- Each move depends only on the present position, not on the existence or non-existence of a certain previous move (such as chess, where castling is made illegal if the king has been moved previously).

(The reader will find a fascinating discussion of ‘two-person games’ in the book by Berlekamp, Conway, and Guy [BCG].)

A **permutation puzzle** is a one-person game with a finite set T of puzzle pieces satisfying the four properties listed below:

4.1. 15 PUZZLE

- (1) For some $n > 1$ depending only on the puzzle's construction, each move of the puzzle corresponds to a unique permutation of the numbers in \mathbb{Z}_n .
- (2) If the permutation of \mathbb{Z}_n in (1) corresponds to more than one puzzle move, then the two positions reached by those two respective moves must be indistinguishable.
- (3) Each move, say M , must be 'invertible' in the sense that there must exist another move, say M^{-1} , which restores the puzzle to the position it was at before M was performed.
- (4) If M_1 is a move corresponding to a permutation f_1 of T and if M_2 is a move corresponding to a permutation f_2 of T , then $M_1 * M_2$ (the move M_1 followed by the move M_2) is either.
 - not a legal move, or
 - corresponds to the permutation $f_1 * f_2$.

Notation: As in step 4 above, we shall always write successive puzzle moves *left-to-right*.

4.1 15 Puzzle

The 15 Puzzle is often associated with the puzzle inventor and problem composer Sam Loyd (1841-1911). He began publishing chess puzzles at the age of 14, had his own regular magazine column on chess problems by the age of 16, and by the time he died had composed over 10,000 original problems in chess, mathematics, and logic. Though some accounts state that the 15 Puzzle was discovered by Sam Loyd, most authorities dispute this.

The 15 Puzzle is one of the earliest and most popular permutation puzzles. The 'solved position' looks like this:

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	•

These numbered squares represent sliding blocks which can only move into the unoccupied square, denoted by a •. We shall sometimes label the blank square as '16' for convenience. The moves of the puzzle consist of sliding numbered squares (such as 12, for example) into the blank square (e.g., swapping 12 with 16). In this way, each move of this puzzle may be regarded as a permutation of the integers in $\{1, 2, \dots, 16\}$.

Not every permutation of the $\{1, 2, \dots, 16\}$ corresponds to a possible position of the puzzle. For example, the position

1	2	3	4
5	6	7	8
9	10	11	12
13	15	14	•

cannot be attained from the previous position. (The mathematical reason for this is explained in §7.4 below.)

Apparently, Sam Loyd applied for a U.S. patent for the above ‘impossible’ puzzle (the one with the 14, 15 swapped, called the ‘14-15 Puzzle’) but since it could not be ‘solved’—i.e., put in the correct order $1, 2, \dots, 15$ —no working model could be supplied, so his patent was denied. (Patent or no, there were apparently thousands of them on the market already.) Before the Rubik’s Cube was invented, the 15 Puzzle was probably the most popular puzzle of all time.

The moves of the 15 Puzzle may be denoted as follows. Suppose we are in a position such as

*	U	*	*
L	•	R	*
*	D	*	*
*	*	*	*

where the squares labeled with an asterisk can be anything. The possible moves in the above diagram are

$$\begin{aligned}
 R &= R_{u,r,d,l} = (r, 16) = \text{swap } r \text{ and } 16, \\
 L &= L_{u,r,d,l} = (l, 16) = \text{swap } l \text{ and } 16, \\
 U &= U_{u,r,d,l} = (u, 16) = \text{swap } u \text{ and } 16, \\
 D &= D_{u,r,d,l} = (d, 16) = \text{swap } d \text{ and } 16,
 \end{aligned}$$

where ‘16’ denotes the unoccupied space •.

Ponderable 4.1.1. *Verify that the five defining properties of a permutation puzzle are satisfied by the 15 Puzzle.*

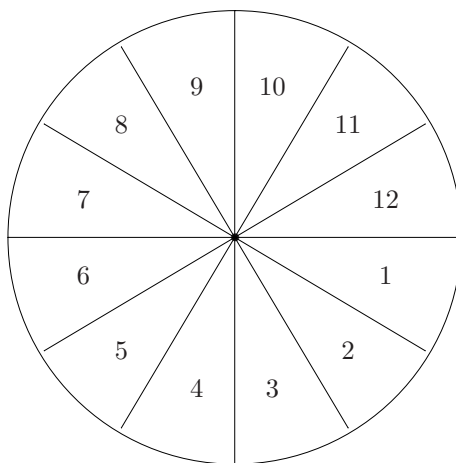
We shall call the 15 Puzzle a **planar puzzle** since all its pieces lie on a flat board. This puzzle is discussed further in §7.4 below.

4.2 The Hockey puck puzzle

The Hockey puck puzzle is a mechanical puzzle invented by András Vègh. Slice a disk with 6 diametrical cuts into 12 congruent ‘pie pieces’. In the solved position, the pieces are denoted $1, 2, \dots, 12$ clockwise. You may make the following moves with this puzzle:

4.3. RAINBOW MASTERBALL

- you may flip over any half of the disk bounded by a cut, leaving the other half unmoved (of course this will change the colors of those pie slices which were flipped);
- you may rotate the disk by any multiple of 30 degrees.



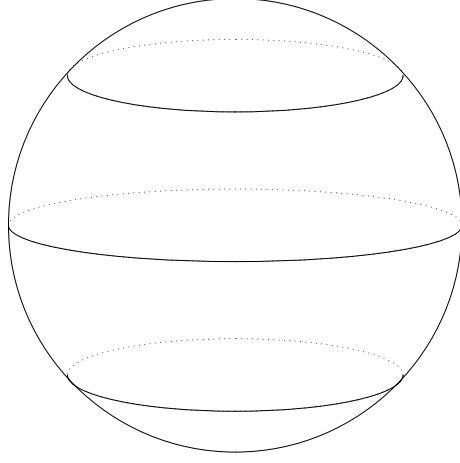
Since the pieces can be flipped over, the top halves of the pieces will be denoted $(1, 0)$, $(2, 0)$, ..., $(12, 0)$ and the corresponding bottom halves by $(1, 1)$, $(2, 1)$, ..., $(12, 1)$.

Basic moves: For the i th radial segment, let f_i be the transformation which flips over that half of the circle which contains the i th pie piece on its edge. Let r denote the rotation of this circle by 30 degrees counterclockwise.

Ponderable 4.2.1. *Verify that the five defining properties of a permutation puzzle are satisfied by this example.*

4.3 Rainbow Masterball

The Rainbow Masterball puzzle will simply be referred to as a ‘Masterball’ in the following. The Masterball is a sphere that has been sliced like an apple along its core into 8 congruent wedges, each having a different color. It has also been sliced three times in the orthogonal direction: viewed as a globe, once along the equator, once along the Tropic of Cancer, and once along the Tropic of Capricorn. (In fact, these lines occur at approximately 23 degree north and south of the equator, not 45 degrees as we shall describe it below.) View this sphere as follows:



A geodesic path from the north pole to the south pole is called a **longitudinal line** and a closed path parallel to the equator is called a **latitudinal line**. There are 8 longitudinal lines and 3 latitudinal lines. In spherical coordinates θ , ϕ , and ρ (θ is the angle from the xz -plane, ϕ is the angle down from the z -axis, ρ is the radius of the Masterball), the longitudinal lines are at the angles which are multiples of $\pi/4$ radians (i.e., at $\theta = n\pi/4$ radians, $n = 1, \dots, 8$) and the latitudinal lines are at $\phi = \pi/4, \pi/2, 3\pi/4$ radians. (Here $\pi = 3.141592\dots$ as usual.)

Without these latitudinal slices, it looks similar to a beachball. Therefore, a Masterball sphere has 32 tiles of 8 distinct colors. We shall assume that the Masterball is in a fixed position in space, centered at the origin.

The sphere shall be oriented by the right-hand rule: the thumb of the right hand wrapping along the polar axis points towards the north pole. We assume that one of the longitudinal lines has been fixed once and for all. This fixed line shall be labeled ‘1’, the next line (with respect to the orientation above) as ‘2’, and so on.

Allowed moves: One may rotate the Masterball east-to-west by multiples of $\pi/4$ along each of the 4 latitudinal bands or by multiples of π along each of the 8 longitudinal lines.

A **facet** will be one of the 32 subdivisions of the Masterball created by these circles. A facet shall be regarded as an immobile position on the sphere and labeled either by an integer $i \in \{1, \dots, 32\}$ or by a pair $(i, j) \in [1, 4] \times [1, 8]$, whichever is more convenient at the time. If a facet has either the north pole or the south pole as a vertex then we call it a **small** (or **polar**) facet. Otherwise, we call a facet **large** (or **middle** or **equatorial**). A **coloring** of the Masterball will be a labeling of each facet by one of the 8 colors in such a way that

- (a) each of the 8 colors occurs exactly twice in the set of the 16 small facets,
- (b) each of the 8 colors occurs exactly twice in the set of the 16 large facets.

4.3. RAINBOW MASTERBALL

A **move** of the Masterball will be a change in the coloring of the Masterball associated with a sequence of maneuvers as described above.

We identify each of the 8 colors with an integer in $\{1, \dots, 8\}$ and identify the collection of facets of the Masterball with a 4×8 array of integers in this range. To **solve** an array one must, by an appropriate sequence of moves corresponding to the above described rotations of the Masterball, put this array into a ‘rainbow’ position so that the matrix entries of each column has the same number. Thus the array

$$\begin{array}{cccccccc} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{array}$$

is ‘solved’. The array

$$\begin{array}{cccccccc} 6 & 7 & 8 & 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{array}$$

corresponds to a rotation of the north pole facets by $3\pi/4$.

Notation. We use matrix notation to denote the 32 facets of the Masterball. The generators for the latitudinal rotations are denoted r_1, r_2, r_3, r_4 . For example, r_1 sends

$$\begin{array}{cccccccc} 11 & 12 & 13 & 14 & 15 & 16 & 17 & 18 \\ 21 & 22 & 23 & 24 & 25 & 26 & 27 & 28 \\ 31 & 32 & 33 & 34 & 35 & 36 & 37 & 38 \\ 41 & 42 & 43 & 44 & 45 & 46 & 47 & 48 \end{array}$$

to

$$\begin{array}{cccccccc} 12 & 13 & 14 & 15 & 16 & 17 & 18 & 11 \\ 21 & 22 & 23 & 24 & 25 & 26 & 27 & 28 \\ 31 & 32 & 33 & 34 & 35 & 36 & 37 & 38 \\ 41 & 42 & 43 & 44 & 45 & 46 & 47 & 48 \end{array}$$

As you look down at the ball from the north pole, this move rotates the top of the ball clockwise. The other moves r_2, r_3, r_4 rotate the associated band of the ball in the same direction—clockwise as viewed from the north pole. The generators for the longitudinal rotations are denoted f_1, f_2, \dots, f_8 . For example, the move $f_1 r_1$ sends

$$\begin{array}{cccccccc} 11 & 12 & 13 & 14 & 15 & 16 & 17 & 18 \\ 21 & 22 & 23 & 24 & 25 & 26 & 27 & 28 \\ 31 & 32 & 33 & 34 & 35 & 36 & 37 & 38 \\ 41 & 42 & 43 & 44 & 45 & 46 & 47 & 48 \end{array}$$

to

44	43	42	41	15	16	17	18
34	33	32	31	25	26	27	28
24	23	22	21	35	36	37	38
14	13	12	11	45	46	47	48

With these rules, one can check the relation

$$f_5 = r_1^4 * r_2^4 * r_3^4 * r_4^4 * f_1 * r_1^4 * r_2^4 * r_3^4 * r_4^4.$$

Ponderable 4.3.1. *Find similar identities for f_6, f_7, f_8 .*

Also, one can check that

$$r_1 = (f_3 * f_7)^{-1} * r_4^{-1} * f_3 * f_7.$$

Ponderable 4.3.2. *There are similar identities for r_2, r_3, r_4 . Find them.*

Identify the facets of the Masterball with the entries of the array

8	7	6	5	4	3	2	1
16	15	14	13	12	11	10	9
24	23	22	21	20	19	18	17
32	31	30	29	28	27	26	25

We may express the generators of the Masterball group in disjoint cycle notation as a subgroup of S_{32} (the symmetric group on 32 letters):

$$\begin{aligned} r_1^{-1} &= (1, 2, 3, 4, 5, 6, 7, 8), \\ r_2^{-1} &= (9, 10, 11, 12, 13, 14, 15, 16), \\ r_3^{-1} &= (17, 18, 19, 20, 21, 22, 23, 24), \\ r_4^{-1} &= (25, 26, 27, 28, 29, 30, 31, 32), \\ f_1 &= (5, 32)(6, 31)(7, 30)(8, 29)(13, 24)(14, 23)(15, 22)(16, 21), \\ f_2 &= (4, 31)(5, 30)(6, 29)(7, 28)(12, 23)(13, 22)(14, 21)(15, 20), \\ f_3 &= (3, 30)(4, 29)(5, 28)(6, 27)(11, 22)(12, 21)(13, 20)(14, 19), \\ f_4 &= (2, 29)(3, 28)(4, 27)(5, 26)(10, 21)(11, 22)(12, 23)(13, 24), \\ f_5 &= (1, 28)(2, 27)(3, 26)(4, 25)(9, 20)(10, 19)(11, 18)(12, 17), \\ f_6 &= (8, 27)(1, 26)(2, 25)(3, 32)(16, 19)(9, 18)(10, 17)(11, 24), \\ f_7 &= (7, 26)(8, 25)(1, 32)(2, 31)(15, 18)(16, 17)(9, 24)(10, 23), \\ f_8 &= (6, 25)(7, 32)(8, 31)(1, 30)(14, 17)(15, 24)(16, 23)(9, 22). \end{aligned}$$

Ponderable 4.3.3. *Verify that the properties of a permutation puzzle are satisfied for the Masterball.*

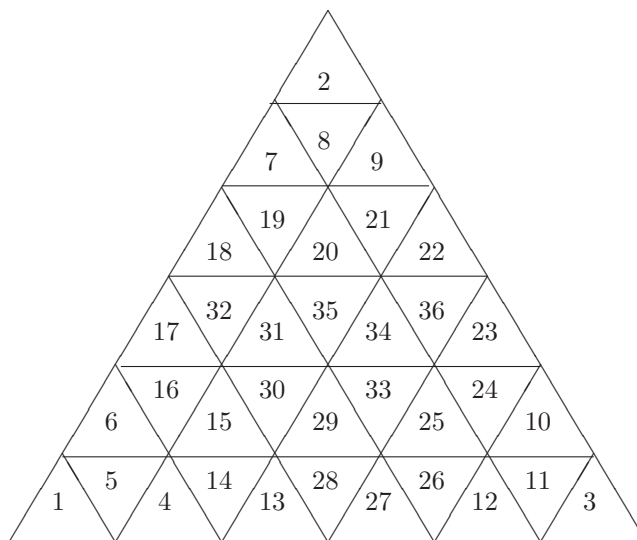
More information on this puzzle is given in §§13.6 and 15.3.

4.4 Pyraminx

The Pyraminx is one of the puzzles invented by Uwe Méffert (1939-), who also markets several of his inventions (see <http://www.mefferts.com> for more information).

A **tetrahedron** is a 4-sided regular platonic solid, all of whose faces are equilateral triangles. In the Pyraminx, each of the 4 faces of the puzzle is divided into 9 triangular facets.

There are a total of $4 \cdot 9 = 36$ facets on the Pyraminx. They will be labeled as in the figure below. (The reader may want to xerox this figure (enlarging it), cut it out, then fold the corners and tape it into a tetrahedron.)



We fix an orientation of the tetrahedron in space so that you are looking at a face which we call the ‘front’. We may also speak of a ‘right’, ‘left’, and ‘down’ face. We label the 4 faces as f (ront), r (ight), l (eft), d (own). We label the vertices U (p), R (ight), L (eft), and B (ack).

The tetrahedron itself has been subdivided into sub-tetrahedrons as follows: to each vertex X (so $X \in \{U, R, L, B\}$) there is an opposing face F of the solid. For each such face, we slice the solid along two planes parallel to the face F and lying in between the face and the vertex. We want these planes, along with the face and the vertex to be spaced apart equally. The sub-tetrahedrons in the slice of the face itself will be called the **face slice** associated with the face F , denoted F_1 ; the sub-tetrahedrons in the middle slice parallel to the face F will be called the **middle slice** associated with that face, denoted F_2 ; and the sub-tetrahedron containing the vertex X to the face **tip** associated with that vertex, denoted F_3 .

To each face labeled F , we have a clockwise rotation by 120 degrees of the first slice F_1 of the face. We shall denote this rotation also by F_1 . This rotation only moves the facets on the slice F_1 . Similarly, we have a clockwise rotation by 120 degrees of the second slice F_2 of the face. We shall denote this rotation also by F_2 . F_3 denotes the clockwise rotation by 120 degrees of the opposing sub-tetrahedron containing the vertex X . These moves permute the labels for the 36 facets, and hence may be regarded as a permutation of the numbers 1, 2, ..., 36.

For example, the clockwise rotation by 120 degrees (looking at the front face) of the sub-tetrahedron opposite to the front face will be denoted F_3 . The disjoint cycle notation for this move, regarded as a permutation, is

$$F_3 = (23, 22, 36).$$

The **basic moves** are given as follows:

$$\begin{aligned} F_1 &= (2, 32, 27)(8, 31, 26)(7, 30, 12)(19, 29, 11) \times \\ &\quad \times (18, 28, 3)(1, 17, 13)(6, 15, 4)(5, 16, 14) \\ F_2 &= (9, 35, 25)(21, 34, 24)(20, 33, 10) \\ F_3 &= (23, 22, 36) \\ R_1 &= (3, 36, 17)(11, 34, 16)(10, 35, 6) \times \\ &\quad \times (24, 31, 5)(23, 32, 1)(2, 22, 18)(9, 20, 7)(8, 21, 19) \\ R_2 &= (12, 33, 15)(26, 29, 14)(25, 30, 4) \\ R_3 &= (27, 28, 13) \\ L_1 &= (1, 28, 22)(5, 29, 21)(4, 33, 9) \times \\ &\quad \times (14, 34, 8)(13, 36, 2)(3, 27, 23)(11, 26, 24)(12, 25, 10) \\ L_2 &= (6, 30, 20)(16, 31, 19)(15, 35, 7) \\ L_3 &= (17, 32, 18) \\ D_1 &= (13, 18, 23)(14, 19, 24)(15, 20, 25) \times \\ &\quad \times (16, 21, 26)(17, 22, 27)(28, 32, 36)(29, 31, 34)(30, 35, 33) \\ D_2 &= (4, 7, 10)(5, 8, 11)(6, 9, 12) \\ D_3 &= (1, 2, 3) \end{aligned}$$

All other moves are obtained by combining these moves sequentially. Indeed, later, we shall want to use moves of the form $F_2 * F_3$, for each face F , but the disjoint cycle notation for these permutations is a little more cumbersome to write down.

Ponderable 4.4.1. *Verify that the properties of a permutation puzzle are satisfied for the Pyraminx, and while doing so, remember Shakespeare's phrase, 'Though this be madness, yet there is method in't.'*

Here's how to construct the group of the Pyraminx using **SAGE**. (Our count is 12 times larger than the usual number of elements, 75582720, since we allow for symmetries; 933120 is the count without 'tip moves'.)

4.5. RUBIK'S CUBES

SAGE

```
sage: S36 = SymmetricGroup(36)
sage: F1 = S36([(2,32,27),(8,31,26),(7,30,12),(19,29,11),\
               (18,28,3),(1,17,13),(6,15,4),(5,16,14)])
sage: F2 = S36([(9, 35, 25),(21, 34, 24),(20, 33, 10)])
sage: F3 = S36([(23, 22, 36)])
sage: R1 = S36([(3,36,17),(11,34,16),(10,35,6),(24,31,5),\
               (23,32,1),(2,22,18),(9,20,7),(8,21,19)])
sage: R2 = S36([(12, 33, 15),(26, 29, 14),(25, 30, 4)])
sage: R3 = S36([(27, 28, 13)])
sage: G = PermutationGroup([F1,F2,F3,R1,R2,R3])
sage: G.order()
906992640
sage: L1 = S36([(1,28,22),(5,29,21),(4,33,9),(14,34,8),\
               (13,36,2),(3,27,23),(11,26,24),(12,25,10)])
sage: L2 = S36([(6,30,20),(16,31,19),(15,35,7)])
sage: L3 = S36([(17,32,18)])
sage: G = PermutationGroup([F1,F2,F3,R1,R2,R3,L1,L2,L3])
sage: G.order()
906992640
sage: 302330880/4
75582720
sage: 933120*81
75582720
sage: 906992640/75582720
12
```

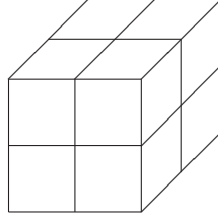
4.5 Rubik's Cubes

We shall consider briefly the $2 \times 2 \times 2$ and $3 \times 3 \times 3$ Rubik's Cubes.

4.5.1 $2 \times 2 \times 2$ Rubik's Cube

The 'pocket' Rubik's Cube has 6 sides, or 'faces', each of which has $2 \cdot 2 = 4$ 'facets', for a total of 24 facets:

Fix an orientation of the Rubik's Cube in space. We may label the 6 sides as f, b, l, r, u, d ('front', 'back', 'left', 'right', 'up', 'down', respectively). The pocket cubet has 8 subcubes. Each face of the cube is associated to a 'slice' of 4 subcubes which share a facet with the face. The face, along with all of the 4 cubes in the 'slice', can be rotated by 90 degrees clockwise. We denote this move by the uppercase letter associated to the lowercase letter denoting the face. For example, F denotes the move which rotates the front face by 90 degrees to clockwise.

Figure 4.1: Facet labeling on the $2 \times 2 \times 2$ Rubik's cube

We label the 24 facets of the 2×2 Rubik's Cube as in Ponderable 3.3.3. The 24 facets will be denoted by xyz , where x is the face on which the facet lives and y, z (or z, y —it doesn't matter) indicate the two edges of the facet. Written in clockwise order, the 'Singmaster notation' is

front face	fru	frd	fld	flu
back face	blu	bld	brd	bru
right face	rbu	rbd	rfd	rfu
left face	lfu	lfd	lbd	lbu
up face	urb	urf	ulf	ulb
down face	drf	drb	dld	dlf

Ponderable 4.5.1. *Verify that the properties of a permutation puzzle are satisfied for this puzzle.*

For future reference, we call this system of notation (which we will also use for the $3 \times 3 \times 3$ and $4 \times 4 \times 4$ Rubik's Cube) the **Singmaster notation**, named after the British mathematician and puzzling enthusiast David Singmaster.

4.5.2 $3 \times 3 \times 3$ Rubik's Cube

In this section I shall, for the most part, simply introduce enough notation (due to Singmaster [Si]) to allow us to check that the puzzle is in fact a permutation puzzle. I shall also introduce a two-person game that is easier to play and learn than solving the cube.

The Rubik's Cube has 6 sides, or 'faces', each of which has $3 \cdot 3 = 9$ 'facets', for a total of 54 facets. Since the center facets are fixed by the basic moves, there are only $54 - 6 = 48$ facets which need labeling. We label these facets $1, 2, \dots, 48$ as in Figure 4.5.2.

The standard generators, corresponding to the six faces of the cube, may then be written in disjoint cycle notation as

4.5. RUBIK'S CUBES

			1	2	3						
			4	U	5						
			6	7	8						
9	10	11	17	18	19	25	26	27	33	34	35
12	L	13	20	F	21	28	R	29	36	B	37
14	15	16	22	23	24	30	31	32	38	39	40
			41	42	43						
			44	D	45						
			46	47	48						

Figure 4.2: Facet labeling on the $3 \times 3 \times 3$ Rubik's cube

$$\begin{aligned}
F &= (17, 19, 24, 22)(18, 21, 23, 20)(6, 25, 43, 16)(7, 28, 42, 13)(8, 30, 41, 11), \\
B &= (33, 35, 40, 38)(34, 37, 39, 36)(3, 9, 46, 32)(2, 12, 47, 29)(1, 14, 48, 27), \\
L &= (9, 11, 16, 14)(10, 13, 15, 12)(1, 17, 41, 40)(4, 20, 44, 37)(6, 22, 46, 35), \\
R &= (25, 27, 32, 30)(26, 29, 31, 28)(3, 38, 43, 19)(5, 36, 45, 21)(8, 33, 48, 24), \\
U &= (1, 3, 8, 6)(2, 5, 7, 4)(9, 33, 25, 17)(10, 34, 26, 18)(11, 35, 27, 19), \\
D &= (41, 43, 48, 46)(42, 45, 47, 44)(14, 22, 30, 38)(15, 23, 31, 39)(16, 24, 32, 40).
\end{aligned}$$

Ponderable 4.5.2. *Check that these six basic moves can indeed be described using the above disjoint cycle notation. (It is helpful to xerox the above diagram, cut it out, and tape together a paper cube for this exercise.)*

The notation for the facets will be similar to the notation used for the $2 \times 2 \times 2$ Rubik's Cube. The corner facets will have the same notation and the edge facets will be denoted by xy , where x is the face on which the facet is located and y is the face the facet borders to. In clockwise order, starting with the upper right-hand corner of each face, the 'Singmaster notation' is

front face	fru	fr	frd	fd	fld	fl	flu	fu
back face	blu	bl	bld	bd	brd	br	bru	bu
right face	rbu	rb	rbd	rd	rfd	rf	rfu	ru
left face	lfu	lf	lfd	ld	lbd	lb	lbu	lu
up face	urb	ur	urf	uf	ulf	ul	ulb	ub
down face	drf	dr	drb	db	dlb	dl	dlf	df

Ponderable 4.5.3. *Verify that the properties of a permutation puzzle are satisfied for the Rubik's Cube.*

Example 4.5.1. *Here is an example of using SAGE to simulate a move on a Rubik's Cube.*

You can enter the Rubik's Cube group as a permutation group generated by the permutations listed above. Permutation groups are objects defined in the next chapter, but for now think of them as sets whose elements are in one-to-one correspondence with the 'states' of the Rubik's Cube, with a fixed orientation in space.

SAGE

```
sage: f= [(17,19,24,22),(18,21,23,20),(6,25,43,16),\
          (7,28,42,13),(8,30,41,11)]
sage: b=[(33,35,40,38),(34,37,39,36),( 3, 9,46,32),\
          ( 2,12,47,29),( 1,14,48,27)]
sage: l=[( 9,11,16,14),(10,13,15,12),( 1,17,41,40),\
          ( 4,20,44,37),( 6,22,46,35)]
sage: r=[(25,27,32,30),(26,29,31,28),( 3,38,43,19),\
          ( 5,36,45,21),( 8,33,48,24)]
sage: u=[( 1, 3, 8, 6),( 2, 5, 7, 4),( 9,33,25,17),\
          (10,34,26,18),(11,35,27,19)]
sage: d=[(41,43,48,46),(42,45,47,44),(14,22,30,38),\
          (15,23,31,39),(16,24,32,40)]
sage: cube = PermutationGroup([f,b,l,r,u,d])
sage: F,B,L,R,U,D = cube.gens()
sage: cube.order()
43252003274489856000
sage: F.order()
4
```

(As an exercise, do the analogous commands for the Masterball, Pyraminx, and $2 \times 2 \times 2$ cube puzzles above.)

Given a 'numerical' labeling as in Figure 4.5.2, you can determine the Singmaster notation for it using SAGE. For example, the facet labeled with the numeral 41 corresponds to the Singmaster notation dlf.

SAGE

```
sage: from sage.groups.perm_gps.cubegroup import *
```

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```
sage: index2singmaster(41)
'dlf'
```

The **SAGE** command `CubeGroup()` encodes a lot of the data associated to the cube: the group, the generators, even the colors.

— SAGE —

```
sage: rubik = CubeGroup()
sage: rubik.move("")
[(), [1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16,
17, 18, 19, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29, 30,
31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 44,
45, 46, 47, 48]]
sage: rubik.faces("R*L")
{'back': [[48, 34, 6], [45, 0, 4], [43, 39, 1]],
'down': [[40, 42, 19], [37, 0, 21], [35, 47, 24]],
'front': [[41, 18, 3], [44, 0, 5], [46, 23, 8]],
'left': [[11, 13, 16], [10, 0, 15], [9, 12, 14]],
'right': [[27, 29, 32], [26, 0, 31], [25, 28, 30]],
'up': [[17, 2, 38], [20, 0, 36], [22, 7, 33]]}
sage: P = rubik.plot_cube("R^2*U^2*R^2*U^2*R^2*U^2")
sage: show(P)
```

This will display the 2-D plot of the Rubik's cube after the move $R^2 * U^2 * R^2 * U^2 * R^2 * U^2$ is made, shown below. (Replace `plot_cube` by `plot3d_cube` for the 3-d plot below, which shows the up, front and right faces.)

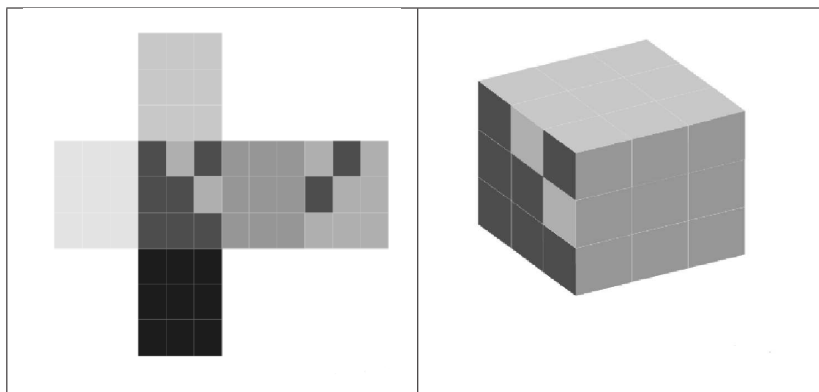


Figure 4.3: $R^2 * U^2 * R^2 * U^2 * R^2 * U^2$.

On the other hand, if you use the **SAGE** notebook interface, then the com-

`mand RubiksCube().move("R^2*U^2*R^2*U^2*R^2*U^2").show3d()` produces the plot shown in Figure 4.4.

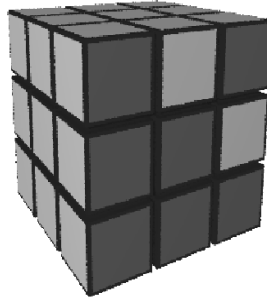


Figure 4.4: $R^2 * U^2 * R^2 * U^2 * R^2 * U^2$.

4.5.3 Some two-player Rubik's Cube games

We collect some two-player games associated with the Rubik's Cube, motivated by [BCG], vol. II.

The superflip game

The position of the Rubik's Cube where every edge is flipped, but all the other subcubes are unaffected, is called the **superflip position**.

To play the game, first choose two particular faces as your up (U) and front (F) faces—say white is up and red is front (assuming you have a cube with adjacent white and red faces). Imagine the cube being placed in space with rectangular coordinate axes in such a way that the bdl corner is at the origin $(0, 0, 0)$, the dl edge is along the x -axis and the bl edge is along the z -axis.

The rules ('slice-superflip game'):

1. Players alternate making moves starting with the cube in the solved position. The first player is determined by (say) a coin toss.
2. A move consists of flipping over exactly two edges. Both edges must lie in a slice. The edge closest to the origin (or, if this is a tie, closest to the x -axis) must be flipped from 'solved' to 'wrong'.
3. The first player to reach the superflip position wins.

Of course, to play this game you must know several edge-flipping moves, such as those in §15.1 below.

4.6. SKEWB

This game is related to a game that might be called a ‘three-dimensional acrostic twins game’. (See [BCG], vol. II, page 441, for a two-dimensional acrostic twins game.)

Several alternate versions of this game may also be played.

The Nonslice-superflip game

The rules are the same except the condition that the two edges belong to the same slice is either dropped altogether or replaced by the condition that the two edges do not belong to the same slice.

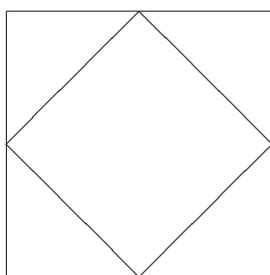
The Möbius-superflip game

The rules for this version are the same except the condition that two edges are flipped is to be replaced by any number of edges fewer than 6 are to be flipped. This game is related to a game that might be called a ‘three-dimensional Möbius’. (See [BCG], vol. II, page 434, for a Möbius game.)

Ponderable 4.5.4. *Play a game!*

4.6 Skewb

The Skewb is a cube-shaped puzzle invented by a London journalist named Tony Durham. The Skewb has been subdivided into regions differently than the Rubik’s Cube: its 8 corners have been sliced in such a way that each edge of the cube is bisected by the slice. The 8 corner pieces are each in the shape of a tetrahedron. First, fix an orientation of the cube in space, so we may talk about a front face, a back face, up, down, left, and right. Each of these six square faces are subdivided into five facets as follows:



The four corner facets are labeled exactly as in the case of the Rubik’s cube (as the lowercase xyz , where x is the label of the face the facet lives on, and y and z are the two neighboring faces).

The moves of the Skewb are different from the Rubik’s Cube as well: Label the corners as XYZ , where xyz is the notation for any of the facets belonging to that corner piece. Pick a corner XYZ of the cube and draw a line L passing

through that corner vertex and the opposite corner vertex ('skewering the cube'). That line defines a 120 degree rotation in the clockwise direction (viewed from the line looking down onto the corner you picked). One move of the Skewb is defined in terms of this rotation as follows: Of course a 120-degree rotation of the entire cube about the line L will preserve the cube but swap some faces and some vertices. The Skewb has a mechanism so that you can actually rotate half (a 'skewed' half) the Skewb by 120 degrees about L and leave the other half completely fixed. This rotation of half the Skewb about L will also be denoted XYZ .

			20	17					
			16	U					
			19	18					
5	2		10	7	25	22	30	27	
	1	L		6	F	21	R	26	B
4	3		9	8	24	23	29	28	
			15	12					
			11	D					
			14	13					

Figure 4.5: Labeling the Skewb facets

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The eight **basic moves** are given by

$$\begin{aligned}
 FUR &= (2, 30, 12)(6, 16, 21)(7, 18, 25)(8, 19, 22)(10, 17, 24) \\
 RUB &= (5, 13, 7)(16, 26, 21)(17, 30, 22)(18, 27, 23)(20, 29, 25) \\
 BUL &= (1, 26, 16)(2, 28, 17)(4, 30, 19)(5, 27, 20)(10, 14, 22) \\
 LUF &= (1, 16, 6)(2, 19, 10)(3, 20, 7)(5, 18, 9)(15, 27, 25) \\
 FDR &= (3, 18, 29)(6, 21, 11)(7, 23, 15)(8, 24, 12)(9, 25, 13) \\
 BDR &= (4, 8, 17)(11, 21, 26)(12, 22, 28)(13, 23, 29)(14, 24, 30) \\
 FDL &= (1, 6, 11)(2, 8, 14)(3, 9, 15)(4, 10, 12)(19, 24, 28) \\
 LDB &= (1, 11, 26)(3, 13, 27)(4, 14, 28)(5, 15, 29)(9, 23, 20).
 \end{aligned}$$

This uses the labeling of the $5 \cdot 6 = 30$ facets as in Figure 4.6. All other moves are obtained by combining these moves sequentially. The set of all moves of the Skewb group can be created in **SAGE** using the following commands.

SAGE

```

sage: FUR = "(2,30,12)(6,16,21)(7,18,25)(8,19,22)(10,17,24)"
sage: RUB = "(5,13,7)(16,26,21)(17,30,22)(18,27,23)(20,29,25)"
sage: BUL = "(1,26,16)(2,28,17)(4,30,19)(5,27,20)(10,14,22)"
sage: LUF = "(1,16,6)(2,19,10)(3,20,7)(5,18,9)(15,27,25)"
sage: FDR = "(3,18,29)(6,21,11)(7,23,15)(8,24,12)(9,25,13)"
sage: BDR = "(4,8,17)(11,21,26)(12,22,28)(13,23,29)(14,24,30)"
sage: FDL = "(1,6,11)(2,8,14)(3,9,15)(4,10,12)(19,24,28)"
sage: LDB = "(1,11,26)(3,13,27)(4,14,28)(5,15,29)(9,23,20)"
sage: G = PermutationGroup([FUR,RUB,BUL,LUF,FDR,BDR,FDL,LDB])
sage: G.order()
37791360

```

Example 4.6.1. Consider the rotation UFR associated to the corner ufr . This move permutes the facets of the Skewb. As a permutation, the disjoint cycle notation for this move is

$$UFR = (6, 16, 21)(7, 18, 25)(10, 17, 24)(8, 19, 22).$$

Note, in particular, that UFR does not move the 9-facet.

Ponderable 4.6.1. Verify that the properties of a permutation puzzle are satisfied for the Skewb.

4.7 Megaminx

This puzzle is in the shape of a **dodecahedron**, a 12-sided regular platonic solid in which each of the 12 faces is a pentagon. We call two faces **neighboring** if they share an edge. There are 20 vertices and 30 edges on a dodecahedron. Each vertex meets exactly three edges and exactly three faces.

Each of the puzzle faces has been subdivided into 11 facets by slicing it with a cut that is parallel to an edge and about one-fifth the way to the opposite vertex. Each face has a total of 5 such cuts, creating 11 facets. There are a total of $11 \cdot 12 = 132$ facets on the puzzle. Each face of the solid is parallel to a face on the opposite side. Fix a face of the dodecahedron and consider a plane parallel to that face slicing through the solid and about one-fifth the way to the opposite face. Now imagine this plane as subdividing, or slicing, the dodecahedron in two. Do this for each of the 12 faces, yielding 12 such slices of the dodecahedron in total. These slices subdivide each of the 12 faces into the same 11 facets as mentioned above.

For each such slice associated to a given face f_i there is a **basic move**, still denoted f_i , of the Megaminx given by rotation of the slice clockwise of the Megaminx by 72 degrees, leaving the rest of the dodecahedron untouched. Such a move effects 26 facets of the Megaminx and leaves the remaining 106 facets completely fixed.

Label the 12 faces of the solid as f_1, f_2, \dots, f_{12} in some fixed way.

Imagine that the dodecahedron is placed in 3-space in such a way that one side lies flat on the xy -plane and is centered along the positive z -axis. Let s denote its height and r the radius of the inscribed circle for any pentagonal face.

Ponderable 4.7.1. *Suppose $r = 1$. Find s . (This is fairly hard. See chapter 8, on Platonic solids, for some ideas.)*

The up face we label as f_1 . The neighboring faces on the upper half of the puzzle may be labeled according to the picture in Figure 4.6, which represents the dodecahedron as viewed from above. The faces on the bottom are labeled similarly: the bottom face is labeled f_{12} and, moving counterclockwise, label the remaining bottom faces f_7, \dots, f_{11} , where f_7 neighbors both f_2 and f_6 .

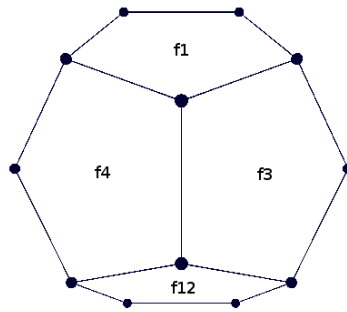


Figure 4.6: The Megaminx labeling.

A more symmetric way to order the faces of the dodecahedron is as follows

4.7. MEGAMINX

(see [B1], exercise 18.35):

f_1	u	f_5	u_3	f_9	d_4
f_2	u_0	f_6	u_4	f_{12}	d
f_3	u_1	f_7	d_2	f_{11}	d_1
f_4	u_2	f_8	d_3	f_{10}	d_0

One property of this labeling is explained in Ponderable 4.7.2.

Ponderable 4.7.2. *Suppose that the permutation $(0, 1, 2, 3, 4)$ of the numbers $\{0, 1, 2, 3, 4\}$ acts on the labels u_0, \dots, u_4 and d_0, \dots, d_4 in the obvious way. Show that this permutation of the faces corresponds to a rotation of the dodecahedron.*

Notice that, like the cube, each vertex is uniquely determined by specifying the three faces it has in common. We use the notation $x.y.z$ for the vertex of the dodecahedron that lies on the three faces x, y, z . Note that the order is irrelevant: $x.y.z$ denotes the same vertex as $y.x.z$ or $z.y.x$.

The facets of the Megaminx may be specified as with the Rubik's Cube: a corner facet may be specified as $[x.y.z]$, where x is the face the facet lives on and y, z are the two neighboring faces of the facet. An edge facet may be specified by $[x.y]$, where x is the face the facet lives on and y is the other neighboring face of the facet. The center facet of f_1 will simply be denoted by $[f_1]$. We will call this label the **intrinsic label**.

We may label the facets of the up face f_1 as follows:

f_1 facet symbol	Numerical label	Intrinsic label
a	1	$[f_1.f_6.f_2]$
b	2	$[f_1.f_2]$
c	3	$[f_1.f_2.f_3]$
d	4	$[f_1.f_3]$
e	5	$[f_1.f_3.f_4]$
f	6	$[f_1.f_4]$
g	7	$[f_1.f_4.f_5]$
h	8	$[f_1.f_5]$
i	9	$[f_1.f_5.f_6]$
j	10	$[f_1.f_6]$
k	11	$[f_1]$

For the next face (the f_2 face), we label the facets in such a way that the abc edge of f_1 joins the ghi edge of f_2 :

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f_2 facet symbol	Numerical label	Intrinsic label
a	12	$[f_2 \cdot f_6 \cdot f_7]$
b	13	$[f_2 \cdot f_7]$
c	14	$[f_2 \cdot f_7 \cdot f_8]$
d	15	$[f_2 \cdot f_8]$
e	16	$[f_2 \cdot f_8 \cdot f_3]$
f	17	$[f_2 \cdot f_3]$
g	18	$[f_2 \cdot f_3 \cdot f_1]$
h	19	$[f_2 \cdot f_1]$
i	20	$[f_1 \cdot f_5 \cdot f_6]$
j	21	$[f_2 \cdot f_6]$
k	22	$[f_2]$

In general, we can label the remaining facets in such a way that the **basic moves** are, as permutations, given by

$$\begin{aligned}
f_1 &= (1, 3, 5, 7, 9)(2, 4, 6, 8, 10)(20, 31, 42, 53, 64) \times \\
&\quad \times (19, 30, 41, 52, 63)(18, 29, 40, 51, 62) \\
f_2 &= (12, 14, 16, 18, 20)(13, 15, 17, 19, 21)(1, 60, 73, 84, 31) \times \\
&\quad \times (3, 62, 75, 86, 23)(2, 61, 74, 85, 32) \\
f_3 &= (23, 25, 27, 29, 31)(24, 26, 28, 30, 32)(82, 95, 42, 3, 16) \times \\
&\quad \times (83, 96, 43, 4, 17)(84, 97, 34, 5, 18) \\
f_4 &= (34, 36, 38, 40, 42)(35, 37, 39, 41, 43)(27, 93, 106, 53, 5) \times \\
&\quad \times (28, 94, 107, 54, 6)(29, 95, 108, 45, 7) \\
f_5 &= (45, 47, 49, 51, 53)(46, 48, 50, 52, 54)(38, 104, 117, 64, 7) \times \\
&\quad \times (39, 105, 118, 65, 8)(40, 106, 119, 56, 9) \\
f_6 &= (56, 58, 60, 62, 64)(57, 59, 61, 63, 65)(49, 115, 75, 20, 9) \times \\
&\quad \times (50, 116, 76, 21, 10)(51, 117, 67, 12, 1) \\
f_7 &= (67, 69, 71, 73, 75)(68, 70, 72, 74, 76)(58, 113, 126, 86, 12) \times \\
&\quad \times (59, 114, 127, 7, 13)(60, 115, 128, 78, 14) \\
f_8 &= (78, 80, 82, 84, 86)(79, 81, 83, 85, 87)(71, 124, 97, 23, 14) \times \\
&\quad \times (72, 125, 98, 24, 15)(73, 126, 89, 25, 16) \\
f_9 &= (89, 91, 93, 95, 97)(90, 92, 94, 96, 98)(80, 122, 108, 34, 25) \times \\
&\quad \times (81, 123, 109, 35, 26)(82, 124, 100, 36, 27) \\
f_{10} &= (100, 102, 104, 106, 108)(101, 103, 105, 107, 109) \times \\
&\quad \times (91, 130, 119, 45, 36)(92, 131, 120, 46, 37)(93, 122, 111, 47, 38) \\
f_{11} &= (111, 113, 115, 117, 119)(112, 114, 116, 118, 120) \times \\
&\quad \times (102, 128, 67, 56, 47)(103, 129, 68, 57, 48)(104, 130, 69, 58, 49) \\
f_{12} &= (122, 124, 126, 128, 130)(123, 125, 127, 129, 131) \times \\
&\quad \times (100, 89, 78, 69, 111)(101, 90, 79, 70, 112)(102, 91, 80, 71, 113)
\end{aligned}$$

4.8 Other permutation puzzles

With regret, I have left out several puzzles. These include such marvelous items as the $4 \times 4 \times 4$ Rubik's Cube, the $5 \times 5 \times 5$ Rubik's Cube; 'Topspin' and

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‘Turnstile’ (planar puzzles); ‘Mozaika’ and ‘Equator’; ‘(Adam) Alexander’s Star’ (a stellated icosahedron), the ‘Impossiball’ (a spherically shaped icosahedron; see [H]), and ‘Mickey’s Challenge’ (a spherically shaped irregular polyhedron with Mickey Mouse figures on it, this puzzle is essentially the same as the Skewb but with some added orientations of faces).

There seems to be no end to the diversity of such puzzles. ‘Christoph’s Jewel’, designed by the German mathematician Christoph Bandelow, based on an invention of Josef Trajber, is essentially a truncated ‘Rubik octahedron’ and may be solved using ‘super-Rubik’s Cube moves’ (see [H]). [For fun, you can take a Rubik’s Cube, strip off all the stickers (using soap and water), and replace them with new stickers modeling a Rubik octahedron. This is because the octahedron is the dual solid of the cube, as described in chapter 8 on symmetry and the Platonic solids, below.]

There are even battery-run puzzles. For example the ‘Orbix’ puzzle, which has 12 buttons, one at each vertex of an icosahedron, that light up. This is a permutation puzzle if you think of a move (which switches certain of the buttons on/off) as permuting the elements of the set of all subsets of the 12 buttons (the subset of buttons which are lit) among themselves. We’ll see this puzzle and other related ones later on in chapter 6 on Merlin’s Machines.

If this isn’t enough, you can find mention of other ‘Rubik’s Cube-like puzzles’ in many other works, such as [Ru], [Si], [H], [B1], [GT], [J2] and Jaap Scherphuis’ excellent site [Sch2].

Chapter 5

What's commutative and purple?

Q: What's commutative and purple?

A: An abelian grape.

Ancient math joke

A story told by Freeman Dyson (1923-), one of the great mathematical physicists of our times, goes as follows: In the early part of the last century mathematician Oswald Veblen and the physicist James Jeans were discussing the reform of the mathematical curriculum at Princeton University. Jeans argued that group theory should be omitted, claiming that group theory was a subject that will never be of any use to physics. Veblen must have won the argument because group theory continued to be taught. It is ironic indeed that group theory not only grew into one of the central themes of physics, but much of the ground-breaking research actually took place at Princeton!

In a rough sense, group theory is the mathematical study of symmetry. The Rubik's Cube displays a remarkable amount of symmetry. This chapter is an introduction to group theory—another useful tool for the Rubik's Cube mechanic wishing to fix his cube, so to speak.

When we studied permutation puzzles in chapter 4, one of the criteria was that each move was 'invertible'. This is, in fact, one of the conditions for the set of all legal moves of a permutation puzzle to form a group.

A group can be defined very generally as a set having a small number of properties. One of these properties is that you need to be able to combine two elements of the set in some particular way to get another element. In the case of the Rubik's Cube, if you combine two moves, you get another move. The precise abstract way to state this type of property is using the notion of a 'binary operation'. A **binary operation** $*$ on a set G is a function that associates with

5.1. THE UNIT QUATERNIONS

each pair of elements (g_1, g_2) of G a single element g_3 , also denoted $g_3 = g_1 * g_2$, in G : $*$: $G \times G \rightarrow G$. A group G is a set with a binary operation $*$ (called the ‘group operation’) satisfying certain properties to be given later. For example, one property is that any element has an ‘inverse element’ associated with it. In the case of the Rubik’s Cube, if you make any move (or sequence of moves), you can always undo the effect of that move by simply reversing every step. This ‘reverse move’ is the ‘inverse’ of the original move, as we will see. Another point to keep in mind is that the notation $*$ for the binary operation is *not* standard. Some people use $+$, some use $*$, some use \cdot , some omit it and simply use juxtaposition of elements. Which notation is used depends on the group in question, but hopefully, this will be unambiguous from the context. In later chapters, we will most often use \cdot or juxtaposition.

Before defining a group, we must decide, just as for sets, on what notation to use to describe or write down a group G . If G is finite then one way is to list all the elements in G and list (or tabulate) all the values of the function $*$. Another method is to describe G in terms of some properties and then define a binary operation $*$ on G . A third method is to give a ‘presentation’ of G (more on this later). Each method has its advantages and disadvantages. We shall eventually use all three approaches.

First, though, we start with some examples.

5.1 The unit quaternions

In the fall of 1843, William R. Hamilton (1805-1865) was walking along the Royal Canal in Ireland with his wife. It was then and there that Hamilton found a generalization of the complex numbers - quaternions. Quaternions are ‘numbers’ of the form $a + bi + cj + dk$, where $a, b, c, d \in \mathbb{R}$ are real numbers, i is the usual $\sqrt{-1}$ and j and k satisfy

$$i^2 = j^2 = k^2 = ijk = -1.$$

In fact, Hamilton could not resist the impulse to carve the formulae for the quaternions in the stone of Brougham Bridge as he and his wife passed it.

Quaternions are not only useful for physics but also for computer graphics animation. They can be added or subtracted easily:

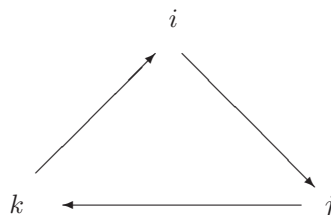
$$(a_1 + b_1i + c_1j + d_1k) + (a_2 + b_2i + c_2j + d_2k) = (a_1 + a_2) + (b_1 + b_2)i + (c_1 + c_2)j + (d_1 + d_2)k.$$

To multiply two quaternions, you use the distributive law and the multiplication rules for i, j, k given above. One important difference between multiplying two quaternions q_1, q_2 together and multiplying two real numbers r_1, r_2 together is that, in general, $q_1q_2 \neq q_2q_1$ (for example, $ij \neq ji$) whereas we always have $r_1r_2 = r_2r_1$ (for example, $\sqrt{2} \cdot 3 = 3 \cdot \sqrt{2}$).

Let Q denote the **quaternion group**:

$$Q = \{1, -1, i, -i, j, -j, k, -k\}.$$

The elements of Q are known as **unit quaternions**. As above, the elements satisfy the rules: $i^2 = j^2 = k^2 = -1$, $ij = k$, $jk = i$, $ki = j$, and in general, $xy = -yx$ for x, y belonging to $\{i, j, k\}$, $x \neq y$. Multiplication for this group can be visualized by the following picture:



This handy mnemonic device indicates that $ij = k$, since the ordering is consistent with the arrows, and $ji = -k$, since the ordering is not consistent with the arrows.

Though Hamilton discovered the quaternions (though Gauss may have known of them earlier), it may have been Cayley who first noticed that the unit quaternions form a group.

5.2 Finite cyclic groups

Consider the set of Rubik's Cube moves $G = \{1, R, R^2, R^3\}$. We make several more or less obvious observations:

- If you make any other rotations of the right face of the cube you won't get any new moves, i.e., ones not already included G . In particular, if you compose any two moves in this set you get another move in this set.
- The move that 'undoes' the effect of R is in this set (in fact, R followed by R^3 is the identity move 1, so R^3 must 'undo' R).
- If $*$: $G \times G \rightarrow G$ is simply the map defined by sending a pair (R^i, R^j) , with $0 \leq i, j \leq 3$, to the composition of these two moves, $R^i * R^j = R^{i+j}$, then $*$ is a binary operation.

It turns out that this set G with this operation $*$ is an example of a group.

Example 5.2.1. Let C_{12} be the set whose elements are $\{0, 1, \dots, 11\}$ and for which the group operation $*$ is simply 'addition mod 12'. This is how one adds time on a clock (except that we call '12 o'clock' '0 o'clock'). Thus $5 + 8 = 1$, $1 + 11 = 0$, and so on.

Questions: What is the ('inverse') element that, when added to 5, gives you 0? What is the inverse element of 1?

This group is called the **cyclic group of order 12**.

5.3. THE DIHEDRAL GROUP

Definition 5.2.1. Let $n > 1$ be an integer and let C_n be the group whose elements are $\{0, 1, \dots, n-1\}$ and for which the group operation is simply ‘addition modulo n ’. This group is called the (additive) **cyclic group of order n** . It is oftentimes also denoted by $\mathbb{Z}/n\mathbb{Z}$.

Geometrically, this may be regarded as the group of all possible rotations of a regular n -gon. This idea is explored in more detail in §5.3 below.

5.3 The dihedral group

Pick an integer $n > 2$ and let R be a regular n -gon centered about the origin in the plane. If $n = 3$ then R is an equilateral triangle, if $n = 4$ then R is a square, if $n = 5$ then R is a pentagon, and so on. Let G denote the set of all linear transformations of the plane to itself that preserve the figure R . The binary operation $\circ : G \times G \rightarrow G$ given by composition of functions gives G the structure of a group. This group is called the **group of symmetries of R** .

Incidentally, if we regard R as a figure in 3-space centered above the origin and let G denote the set of all linear transformations of 3-space, then we obtain a slightly larger group in some cases (see Example 9.3.4 below and [NST] for more details).

Label the vertices of the n -gon as $1, 2, \dots, n$. The group G permutes these vertices amongst themselves; hence, each $g \in G$ may be regarded as a permutation of the set of vertices $V = \{1, 2, \dots, n\}$. In this way, we may regard G as a permutation group, since it is the subgroup of S_n generated by the elements of G .

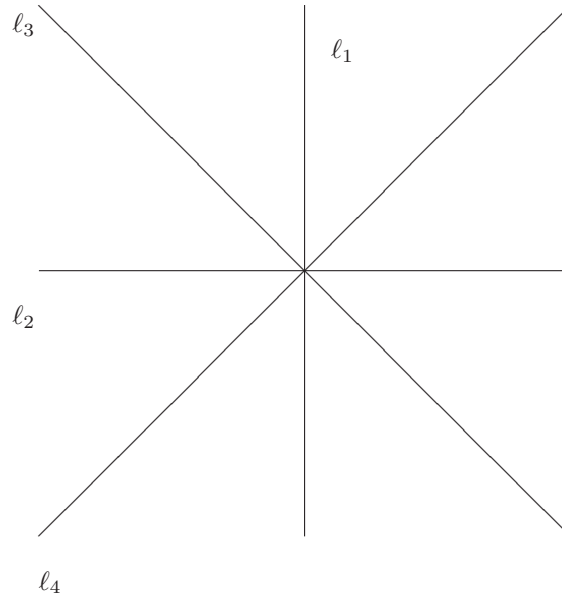
The fact that this group has $2n$ elements follows from a simple counting argument: Let $r \in G$ denote the element that rotates R by $2\pi/n$ radians counter-clockwise about the center. Let L be a line of symmetry of R that bisects the figure into two halves. Let s denote the element of G that is reflection about L . There are n rotations by a multiple of $2\pi/n$ radians about the center in G : $1, r, r^2, \dots, r^{n-1}$. There are n elements of G that are composed of a reflection about L and a rotations by a multiple of $2\pi/n$ radians about the center: $s, s \circ r, s \circ r^2, \dots, s \circ r^{n-1}$. These comprise all the elements of G .

The symmetry group of R is known as the **dihedral group of order $2n$** , denoted D_n . (Note: Some people denote this group instead by D_{2n} .)

Example 5.3.1. Let G be the symmetry group of the square, i.e., the group of symmetries of the square generated by the rigid motions

$$\begin{aligned} g_0 &= 90 \text{ degrees clockwise rotation about } O, \\ g_1 &= \text{reflection about } \ell_1, \\ g_2 &= \text{reflection about } \ell_2, \\ g_3 &= \text{reflection about } \ell_3, \\ g_4 &= \text{reflection about } \ell_4, \end{aligned}$$

where ℓ_1, ℓ_2, ℓ_3 denote the lines of symmetry in the picture below:



The elements of G are

$$1, g_0, g_0^2, g_0^3, g_1, g_2, g_3, g_4.$$

Let X be the set of vertices of the square. Then G acts on X .

5.4 The symmetric group

The art of doing mathematics consists in finding that special case which contains all the germs of generality.

David Hilbert

Before defining anything, we shall provide a little motivation for some general notions that will arise later. Each move of the Rubik's Cube may be regarded as a permutation of the set of 54 facets of the cube. It is natural from our perspective to study in general permutations of a set. As Hilbert advised us, we are looking for a really good special case to focus our attention on, and the Rubik's Cube group is the example we will use.

First, some basic ideas to get us started.

Let X be any finite set and let S_X denote the set of all permutations of X onto itself:

$$S_X = \{f : X \rightarrow X \mid f \text{ is a bijection}\}.$$

This set has the following properties:

1. If f, g belong to S_X then fg (the composition of these permutations) also belongs to S_X ('closed under compositions').

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2. If f, g, h all belong to S_X then $(fg)h = f(gh)$ ('associativity').
3. The identity permutation $I : X \rightarrow X$ belongs to S_X ('existence of the identity').
4. If f belongs to S_X then the inverse permutation f^{-1} also belongs to S_X ('existence of the inverse').

The set S_X is called the **symmetric group of X** . We shall usually take for the set X a set of the form $\{1, 2, \dots, n\}$, in which case we shall denote the symmetric group by S_n . This group is also called the **symmetric group on n letters**.

Example 5.4.1. Suppose $X = \{1, 2, 3\}$. We can describe S_X as

$$S_X = \{I, s_1 = (1, 2), s_2 = (2, 3), s_3 = (1, 3, 2), s_4 = (1, 2, 3), s_5 = (1, 3)\}.$$

We can compute all possible products of two elements of the group and tabulate them in a multiplication table as

	I	s_1	s_2	s_3	s_4	s_5
I	I	s_1	s_2	s_3	s_4	s_5
s_1	s_1	I	s_3	s_2	s_5	s_4
s_2	s_2	s_4	I	s_5	s_1	s_3
s_3	s_3	s_5	s_1	s_4	I	s_2
s_4	s_4	s_2	s_5	I	s_3	s_1
s_5	s_5	s_3	s_4	s_1	s_2	I

Ponderable 5.4.1. Verify the four properties of S_X mentioned above. (Note that the verification of associativity follows from the associative property of the composition of functions—see Ponderable 3.1.4).

5.5 General definitions

We take the above four properties of the symmetric group as the four defining properties of a group. The definition of a group below (or 'abstract group', to be precise) was first formulated by Cayley in 1854. (We met Cayley already in §2.2.1 in connection with his work on matrices.)

Definition 5.5.1. Let G be a set and suppose there is a mapping

$$\begin{aligned} * : G \times G &\rightarrow G \\ (g_1, g_2) &\mapsto g_1 * g_2 \end{aligned}$$

(called the group's **operation**) satisfying

(G1) if g_1, g_2 belong to G then $g_1 * g_2$ belongs to G (' G is closed under $*$ '),

(G2) if g_1, g_2, g_3 belong to G then $(g_1 * g_2) * g_3 = g_1 * (g_2 * g_3)$ ('associativity'),

(G3) that there is an element $1 \in G$ such that $1 * g = g * 1 = g$ for all $g \in G$ ('existence of an identity'),

(G4) if g belongs to G then there is an element $g^{-1} \in G$, called the **inverse of g** such that $g * g^{-1} = g^{-1} * g = 1$ ('existence of inverse').

Then G (along with the operation $*$) is a **group**.

For example, if R is the Rubik's Cube move associated with rotating the right face one quarter turn, then $R^{-1} = R^3$.

Remark 5.5.1. In general, when a set G and a binary operation $*$ are explicitly given and when $(G, *)$ is then expected to be a group, usually the hardest condition to verify is associativity. In the special case when G is a set of permutations of a set X and $*$ is simply the usual multiplication of permutations, then associativity is easy to verify, since all the elements of G are functions (from X to X) and the binary operation is function composition. In this case, associativity is more or less 'obvious'.

Example 5.5.1. Actually, this is a 'non-example'. Let S be the set of all legal moves (one can eventually make from a given legally obtained position) of the 15 Puzzle (as described in chapter 4). In a given position, for example the solved position, there aren't that many possibilities: there are only 2 moves in the solved position and there are never any more than 4 moves possible from any position.

From the solved position one can move $(15, 16)$ and $(12, 16)$ (where 16 denotes the blank square) but not, for example, $(1, 16)$. Since $(15, 16), (12, 16) \in S$ and since $(12, 16)(15, 16)$ is not a legal move, it follows that composition of legal moves is not always legal. This shows that composition is not a binary operation, so property number (G1) fails to hold.

Lemma 5.5.1. (Cancellation Law) Let G be a group and $a, b, c \in G$. If $a * c = b * c$, where $a, b, c \in G$, then $a = b$.

The proof is left to the reader.

In the above definition, we have not assumed that there was exactly one identity element 1 of G because, in fact, one can show that if there is one then it is unique. (Suppose there are two identity elements, denoted e_1, e_2 say. By definition, $e_1 = e_1 * e_2 = e_2$.) Likewise, if G is a group and $g \in G$ then the inverse element of g is unique. The proof of this is left to the reader. There are other properties of a group that can be derived from (G1)-(G4). We shall prove them as needed.

The **multiplication table** (also called 'Cayley tables' after Cayley, who first introduced them) of a finite group G is a tabulation of the values of the binary operation $*$, as in Example 5.4.1. Let $G = \{g_1, \dots, g_n\}$. The multiplication table of G is

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*	g_1	g_2	\dots	g_j	\dots	g_n
g_1						
g_2						
\vdots						
g_i				$g_i * g_j$		
\vdots						
g_n						

Example 5.5.2. Using **SAGE** it is easy to construct a multiplication table:

SAGE

```
sage: G = PermutationGroup(['(1,2,3)', '(2,3)'])
sage: G.cayley_table(names="abcdef")
      [a b c d e f]
      [b a d c f e]
      [c e a f b d]
      [d f b e a c]
      [e c f a d b]
      [f d e b c a]
sage: G.list()
[(), (2,3), (1,2), (1,2,3), (1,3,2), (1,3)]
```

The ordering in the rows and columns of the table is the same as that produced by `G.list()`. In general, the first element is the identity.

Ponderable 5.5.1. Compute the multiplication table for C_3 .

Some properties:

Lemma 5.5.2. (a) Each element $g_k \in G$ occurs exactly once in each row of the table.

(b) Each element $g_k \in G$ occurs exactly once in each column of the table.

(c) If the (i, j) th entry of the table is equal to the (j, i) th entry, then $g_i * g_j = g_j * g_i$.

(d) If the table is symmetric about the diagonal, then $g * h = h * g$ for all $g, h \in G$. (In this case, we call G **abelian**.)

Abelian groups are named after the Norwegian mathematician Niels Abel (1802-1829). Abel's short life was a tragic one, punctuated by poverty and the death of his father in 1820. However, he had great mathematical talent and showed before Galois that the roots of the general 5th degree polynomial,

$$x^5 + a_1x^4 + a_2x^3 + a_3x^2 + a_4x + a_5 = 0,$$

cannot be expressed in terms of radicals alone. (Galois proved a more general result about six years later.) Unfortunately, his health was not good, and during a journey to visit his fiancée over Christmas 1828 he became seriously ill and died. In his short life, Abel made major contributions to other areas of mathematics as well and was awarded the Grand Prix in mathematics by the Paris Academy in 1830.

I can't resist this corny pun: Why didn't Isaac Newton work on group theory? He wasn't Abel.

Definition 5.5.2. Let g and h be two elements of a group G . We say that g **commutes** with h (or that g, h **commute**) if $g * h = h * g$. We call a group **commutative** (or **abelian**) if every pair of elements g, h belonging to G commute. If G is a group that is not necessarily commutative, then we call G **noncommutative** (or **nonabelian**).

Example 5.5.3.

- The elements in $G = \{1, R, R^2, R^3\}$ all commute with each other, so G is an abelian group.
- The integers, with ordinary addition as the group operation, constitutes an abelian group.

Ponderable 5.5.2. Show that any group having exactly 2 elements is abelian.

Convention: When dealing with groups in general we often drop the $*$ and denote multiplication simply by juxtaposition (that is, sometimes we write gh in place of $g * h$) or by a dot \cdot . However, there is one exception. If the group G is abelian, then one often replaces $*$ by $+$ and then $+$ is *not* dropped.

Now that we know the definition of a group, the question arises: how might they be described? The simplest answer is that we describe a group much as we might describe a set: we could list all its elements and give the multiplication table, or we could describe all its elements and their multiplication in terms of some property from which we can verify the four properties of group. Though the first way has the distinct advantage of being explicit, it is this second alternative that is the most common, since it is usually more concise.

Our objective is to introduce terminology and techniques that enable us to analyze permutation puzzles mathematically. The types of groups that arise in this context are defined next.

Definition 5.5.3. Let X be a finite set. Let $S = \{g_1, g_2, \dots, g_n\}$ be a finite set of elements of permutations of X (so that they all belong to S_X). Let G be the set of all possible products of the form

$$g = x_1 * x_2 \cdots * x_m, \quad m > 0,$$

where each of the x_1, \dots, x_m is taken from the set S . The set G , together with the group operation given by composition of permutations, is called a **permutation**

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group with generators g_1, \dots, g_n (or the **permutation group** generated by S). We sometimes write

$$G = \langle g_1, \dots, g_n \rangle \subset S_X.$$

Example 5.5.4.

- If $X = \{1, 2, \dots, 54\}$ and $S = \{(1, 2, 3, \dots, 54)\} \subset S_{54}$ then the permutation group generated by S is a cyclic group G with elements $\{1, g, g^2, \dots, g^{53}\}$.
- Let X be the set of 54 facets of the Rubik's Cube and let $R, L, U, D, F, B \in S_X$ denote the basic moves of the Rubik's Cube, in the notation introduced in the previous chapter. The permutation group

$$G = \langle R, L, U, D, F, B \rangle \subset S_X$$

is called the **Rubik's Cube group**. We shall determine the 'structure' of this group (i.e., its relationship with 'known groups') later in this book.

It is not too hard to justify our terminology.

Lemma 5.5.3. A permutation group is a group.

Proof: Let G be a permutation group as in the above definition.

We shall only prove that each $g \in G$ has an inverse, leaving the remainder of the properties for the reader to verify. The set $\{g^n \mid n \geq 1\} \subset S_X$ is finite. There are $n_1 > 0$, $n_2 > n_1$ such that $g^{n_1} = g^{n_2}$. Then $g^{-1} = g^{n_2 - n_1 - 1}$ since $g \cdot g^{n_2 - n_1 - 1} = 1$. \square

Given generators $g_1, \dots, g_n \in S_X$, how do you construct the permutation group $G = \langle g_1, \dots, g_n \rangle$ generated by them? The algorithm below explains this.

Algorithm:

Input: The generators g_1, \dots, g_n (as permutations in S_X).

Output:

- The elements of G
- $S = \{g_1, \dots, g_n, g_1^{-1}, \dots, g_n^{-1}\}$
- $L = S \cup \{1\}$

Algorithm

```

for g in S do
  for h in L do
    if g*h not in L then L = L union {g*h} endif
  endfor
endfor

```


Note that the size of the list L in the `for` loop changes after each iteration of the loop. The meaning of this is that the `if-then` command is to be executed exactly once for each element of the list.

Definition 5.5.4. If G is a group then the **order** of G , denoted $|G|$, is the cardinality of G . (In other words, if G is a finite set then it is number of elements of G , and if G is not finite then we set $|G| = \infty$.) If g is an element of the group G , then the **order of g** , denoted $\text{ord}(g)$, is the smallest positive integer m such that $g^m = 1$, if it exists. If such an integer m does not exist then we say that g has **infinite order**.

Example 5.5.5.

- A SAGE computation illustrating the above definition.

SAGE

```
sage: rubik = CubeGroup()
sage: G = rubik.group()
sage: N = G.order(); N
43252003274489856000
sage: N.factor()
2^27 * 3^14 * 5^3 * 7^2 * 11
sage: R = rubik.move('R')[0]
sage: R.order()
4
```

- Mathematically, we shall see later why the Rubik's Cube group G does indeed have the property that $|G| = 2^{27}3^{14}5^37^211$ (roughly 4.3×10^{19}).
- There is an odd permutation of order 42 in S_{12} , for example, $(1, 2)(3, 4, 5)(6, 7, 8, 9, 10, 11, 12)$. There is an even permutation of order 15 in S_8 , for example $(1, 2, 3)(4, 5, 6, 7, 8)$.
- There is an element in the Rubik's Cube group whose order is 1260 and no element of higher order. J. Butler found the following move of this order: $m = RU^2D^{-1}BD^{-1}$ (see Bandelow's book [B1], page 51, for another simple move of order 1260).

5.5.1 Cauchy's theorem

We shall be able to make use of the following fact frequently.

Theorem 5.5.1.

- (Cauchy) Let p be a prime dividing $|G|$. There is a $g \in G$ of order p .
- Let n be an integer not dividing $|G|$. There does not exist a $g \in G$ of order n .

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Augustin Louis Cauchy (1789-1857) made significant contributions to many branches of mathematics and was one of the greatest mathematicians of his time.

Part (a) will be proven in §5.11 below (see Corollary 5.11.2) and part (b) is a corollary of Theorem 5.6.1 below.

Application. As the Rubik's Cube group G has the property that $|G| = 2^{27}3^{14}5^37^211$, it follows from this and part (b) of the above theorem that there is no move of the Rubik's Cube of order 13 (since 13 does not divide $|G|$) but there is one of order 11. (That said, just because we know there exists a move of order 11 doesn't mean we know how to find one!) This is equivalent to saying there exists a cyclic subgroup of the Rubik's Cube group of order 11. There is no cyclic subgroup of the Rubik's Cube group of order 13.

Ponderable 5.5.3. Let $X = \{1, 2, 3\}$. We use the notation of Example 5.4.1 above.

- (a) Let G be the permutation group with generator s_1 , $G = \langle s_1 \rangle$. Verify that there are only two elements in G .
- (b) What is the order of s_5 ?
- (c) Let G be the permutation group with generator s_3 : $G = \langle s_3 \rangle$. Verify that there are only three elements in G .
- (d) Find the order of s_3 .
- (e) Show that $S_X = \langle s_1, s_2 \rangle$.

Definition 5.5.5. If G is a group G with only one generator, then we say that G is **cyclic**.

Lemma 5.5.4. If $G = \langle g \rangle$ is a finite cyclic group with generator g then $|G| = \text{ord}(g)$.

Proof: Let $m = \text{ord}(g)$, so $g^m = 1$. We can list all the elements of G as follows:

$$1, g, g^2, \dots, g^{m-1}.$$

There are m elements in this list. \square

5.5.2 The Gordon game

Let $(G, *)$ be a finite group, written

$$G = \{g_0 = \text{the identity}, g_1, \dots, g_n\},$$

so $g_0 = 0$ if $*$ is addition and $g_0 = 1$ if $*$ is multiplication or composition. You and your opponent share a set of **move tokens**, denoted

$$M = \{g_1, \dots, g_n\},$$

and **place tokens**, denoted

$$P = \{g_1, \dots, g_n\}.$$

Rules for play:

- Players alternate turns. Each turn consists of removing one *move* token and one *place* token according to the conditions listed below. The first person who cannot make a legal play loses.
Let $m_0 = p_0 = 1$ (or $m_0 = p_0 = 0$, if G is written additively) and let $i = 1$.
- The first player picks any move token $m_1 \in M$ and the place token $p_1 = m_1 \in P$. These tokens m_1 and p_1 are then removed from M and P , respectively.
- The next player picks a move token m_{i+1} and place token p_{i+1} such that $p_{i+1} = m_{i+1}p_i \in P$. These tokens m_{i+1} and p_{i+1} are then removed from M and P , respectively.
- Increment i and go to the previous step.

Example 5.5.6. *Let*

$$G = \mathbb{Z}/7\mathbb{Z} = \{0, 1, 2, 3, 4, 5, 6\},$$

in the notation of Definition 5.2.1. The moves of a game are determined by recording the move tokens. One possible game is

$$\begin{array}{ccccccc} \bullet & 4 & 1 & & & 3 & 2 \\ 0 & 1 & 2 & 3 & 4 & 5 & 6 \end{array}$$

where the \bullet over the identity element 0 of the group indicates that it isn't moved and the numbers above a group element indicate when it was moved:

player	m	p	P	M
1	$m_1 = 2$	$p_1 = 2$	$\{1, 3, 4, 5, 6\}$	$\{1, 3, 4, 5, 6\}$
2	$m_2 = 4$	$p_2 = 2 + 4 = 6$	$\{1, 3, 4, 5\}$	$\{1, 3, 5, 6\}$
1	$m_3 = 6$	$p_3 = 6 + 6 = 5$	$\{1, 3, 4\}$	$\{1, 3, 5\}$
2	$m_4 = 3$	$p_4 = 3 + 5 = 1$	$\{3, 4\}$	$\{1, 5\}$

The second player wins.

Ponderable 5.5.4. *Play a game!*

Remark 5.5.2. *If $G = \mathbb{Z}/p\mathbb{Z}$ (the cyclic group with p elements) there is a conjecture that the 2nd player has a winning strategy when $p > 5$ (see Isbell's note [I] for more details). For more general groups, strategies are not known. In fact, they haven't even been conjectured.*

Remark 5.5.3. *If you and your opponent both try to drag the game on as long as possible, can you exhaust the set of move tokens and the set of place tokens? The answer is known for abelian groups, dihedral groups, and groups of order < 32 . The general answer seems to be unknown.*

5.6 Subgroups

As already noted, the set of moves of the Rubik's Cube forms a group G under the operation of composition. Suppose now you consider a subset of G that also is a group under composition. Such a subset is called a 'subgroup'. For example, $\{1, R, R^2, R^3\}$ is a subgroup of G .

Definition 5.6.1. *Let G be a group. A **subgroup** of G is a subset H of G such that H , together with the operation $*$ inherited as a subset of G , satisfies the group operations (G1)-(G4) (with G replaced by H everywhere).*

Notation: If G is a group then we will denote the statement ' H is a subgroup of G ' by $H \subset G$.

Example 5.6.1. • $2\mathbb{Z}$ is a subgroup of \mathbb{Z} .

- $2\mathbb{Z}/10\mathbb{Z} = \{0, 2, 4, 6, 8\}$ (with addition mod 10) is a subgroup of $\mathbb{Z}/10\mathbb{Z}$.
- $H = \{1, (1, 2, 3), (1, 3, 2)\}$ is a subgroup of S_3 .
- A permutation group G generated by elements g_1, \dots, g_n belonging to S_X is a subgroup of S_X , i.e., $G \subset S_X$.

One might ask: what are all the subgroups of the Rubik's Cube group? Unfortunately, it turns out that this question is too impractical. There are too many subgroups to list, so a simple answer is not possible. In fact, no one (as far as I am aware) knows exactly how many subgroups the Rubik's Cube group has. However, later, when we have a more useful way of describing a group (using generators and relations (see §10.3), we will explicitly determine some of the subgroups of 'small' order.

We'd like an easy-to-use criterion to determine when a given subgroup H of S_{54} is a subgroup of the Rubik's Cube group G . This type of condition doesn't seem to exist, at least not in a practical form, but the following criterion is very useful. Another version of the following result will be given in Theorem 5.11.1 below.

Theorem 5.6.1. (Lagrange's theorem) *Let H be a subgroup of a finite group G . Then $|H|$ divides $|G|$.*

It is worth remarking that, contrary to what might be suggested by the name of the result, Lagrange's theorem, in the form stated here, is not due to Lagrange. In the form given, Lagrange's theorem (for permutation groups) was probably known to Galois and appeared in works of Serret and Camille Jordan in the 1860's.

Proof: For $x, y \in G$, define $x \sim y$ if $xH = yH$, where

$$xH = \{x * h \mid h \in H\}.$$

This is an equivalence relation. (The curious reader can easily check that the reflexive, symmetry, and transitivity properties hold.) Moreover, the equivalence class of x consists of all elements in G of the form $x * h$, for some $h \in H$,

i.e., $[x] = xH$. Let $g_1, \dots, g_m \in G$ denote a complete set of representatives for the equivalence classes of G . Because of the cancellation law for groups, $|xH| = |H|$ for each $x \in G$. (Exercise: Verify this.) Furthermore, we know that the equivalence classes partition G (see Ponderable 2.3.3), so

$$G = \cup_{i=1}^m [g_i] = \cup_{i=1}^m g_i H.$$

Comparing cardinalities of both sides, we obtain $|G| = |g_1 H| + \dots + |g_m H| = m|H|$. This proves the theorem. \square

Example 5.6.2. Here is an illustration of Lagrange's theorem using SAGE:

SAGE

```
sage: G = SymmetricGroup(6)
sage: g1 = G.random_element()
sage: g2 = G.random_element()
sage: H = G.subgroup([g1, g2])
sage: G.order()/H.order() in ZZ
True
```

Ponderable 5.6.1. Show, as a corollary to Theorem 5.6.1, that Theorem 5.5.1(b) is true.

Definition 5.6.2. If H and G are finite groups and $H \subset G$, then $|G|/|H|$ (which is an integer by Lagrange's theorem 5.6.1, above) is called the **index** of H in G , denoted $[G : H] = |G|/|H|$.

Example 5.6.3. Let

$$A_n = \{g \in S_n \mid g \text{ is even}\}.$$

This is a subgroup of S_n called the **alternating subgroup of degree n** . It is known (and not hard to prove using cosets) that $|A_n| = |S_n|/2$.

Definition 5.6.3. The **center** of a group G is the subgroup $Z(G)$ of all elements that commute with every element of G :

$$Z(G) = \{z \in G \mid z * g = g * z, \text{ for all } g \in G\}.$$

Observe that

- The identity element always belongs to $Z(G)$. (If the identity element is the only element of $Z(G)$, then we say G has a **trivial center**.)
- G is commutative if and only if $G = Z(G)$. (This follows from the definitions.)

Ponderable 5.6.2. Let $G = S_3$. Determine $Z(G)$ using explicit computations. If $n \geq 3$, show that $Z(S_n)$ has a trivial center. (Hint: Consider the elements that commute with all the n -cycles.)

For the center of the Rubik's Cube group, see §5.7.1 and Corollary 11.2.2 below.

5.7 Puzzling examples

There are many puzzles on the market that give rise to interesting groups in mathematics. A few simple examples are given in this section. More complicated examples will be studied later.

Example 5.7.1. *Consider an infinite chessboard, which we imagine being placed on the Cartesian plane.*

	×		×		×		×
×		×		×		×	
	×		×		×		×
×		×		×		×	
	×		×		×		×
×		×		×		×	
	×		×		×		×
×		×		×		×	

Label one square as $(0,0)$ and call it the **origin**. Label the others (m,n) , as one would label the vertices in a lattice in the plane. Place only one chess piece, a king, at $(0,0)$. Label the move one square to the right x , one square to the left x^{-1} , the move one square forward y , one square backwards y^{-1} , and label the other moves xy , $x^{-1}y$, $x^{-1}y^{-1}$, and xy^{-1} , in the obvious way. The set of all possible king's moves may be identified with the set

$$\{x^m y^n \mid m, n \in \mathbb{Z}\}.$$

This is an infinite abelian group under multiplication. The number of ways that the king can reach the square (m,n) in N moves is the coefficient $K_N(m,n)$ of $x^m y^n$ in the expansion of

$$(x + x^{-1} + y + y^{-1} + xy + x^{-1}y + x^{-1}y^{-1} + xy^{-1})^N.$$

Further details on this construction can be found in the chapter ‘Wanderungen von Schachfiguren’ by K. Fabel in [BFR].

For example, using **SAGE** one sees that there are exactly

$$K_{10}(1,1) = 19246920 \approx 1.9 \times 10^7$$

ways to reach $(1,1)$ from the origin in 10 moves.

SAGE

```

sage: R2 = PolynomialRing(QQ,2,"xy")
sage: F2 = FractionField(R2)
sage: x,y = R2.gens()
sage: f = lambda n: F2((x+1/x+y+1/y+x*y+y/x+1/(x*y)+x/y)^n)
sage: a = f(10).numerator()
sage: b = f(10).denominator()
sage: b
x^10*y^10
sage: a.coefficient(x^(11)*y^(11))
19246920

```

Ponderable 5.7.1. Suppose that one day the chess king in the above example comes alive. He starts from $(0,0)$ and begins walking around on his chessboard taking completely random moves at each step. What is the likelihood that he will end up on the square he started on after 3 moves?

Ponderable 5.7.2. Work out a similar result as in the example above using a knight in place of a king.

Example 5.7.2. Let M_R be the middle right slice rotation by 90 degrees (viewed from the right face). Define M_F , for the front face, and M_U similarly. Consider the subgroup H of the Rubik's Cube group generated by the 'square slice moves',

$$H = \langle M_R^2, M_F^2, M_U^2 \rangle.$$

Then $H = \langle M_R^2 \rangle \times \langle M_F^2 \rangle \times \langle M_U^2 \rangle \cong C_2 \times C_2 \times C_2 = C_2^3$, where C_n denotes the cyclic group of order n .

5.7.1 The superflip

The collection of all moves of the Rubik's Cube may be viewed as a subgroup G of S_{48} . We denote multiplication in S_4 (and hence in G) by \cdot . The center of G consists of exactly two elements, the identity and the **superflip** move that has the effect of flipping over every edge, leaving all the corners alone and leaving all the subcubes in their original position ([B1], page 48):

$$Z(G) = \{1, \text{superflip}\}.$$

Later, in §11.2, we shall see that this is a consequence of the 'second fundamental theorem of cube theory'. One move for the superflip is

$$\begin{aligned}
\text{superflip} &= R \cdot L \cdot F \cdot B \cdot U \cdot D \cdot R \cdot L \cdot F \cdot B \cdot U \cdot F^2 \cdot M_R \cdot \\
&\quad \cdot F^2 \cdot U^{-1} \cdot M_R^2 \cdot B^2 \cdot M_R^{-1} \cdot B^2 \cdot U \cdot M_R^2 \cdot D \\
&= R \cdot L \cdot F \cdot B \cdot U \cdot D \cdot R \cdot L \cdot F \cdot B \cdot U \cdot F^2 \cdot R^{-1} \cdot \\
&\quad \cdot L \cdot D^2 \cdot F^{-1} \cdot R^2 \cdot L^2 \cdot D^2 \cdot R \cdot L^{-1} \cdot \\
&\quad \cdot F^2 \cdot D \cdot R^2 \cdot L^2 \cdot D \quad (34 \text{ quarter turns}),
\end{aligned}$$

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where M_R is middle right slice rotation by 90 degrees (viewed from the right face). Other expressions for the superflip are Dik T. Winter's move

$$\begin{aligned} \text{superflip} = & F \cdot B \cdot U^2 \cdot R \cdot F^2 \cdot R^2 \cdot B^2 \cdot U^{-1} \cdot D \cdot F \cdot U^2 \cdot R^{-1} \cdot L^{-1} \cdot U \cdot \\ & \cdot B^2 \cdot D \cdot R^2 \cdot U \cdot B^2 \cdot U \quad (28 \text{ quarter turns}), \end{aligned}$$

and Mike Reid's expression (found with a computer search)

$$\begin{aligned} \text{superflip} = & R^{-1} \cdot U^2 \cdot B \cdot L^{-1} \cdot F \cdot U^{-1} \cdot B \cdot D \cdot F \cdot U \cdot D^{-1} \cdot L \cdot D^2 \cdot F^{-1} \cdot \\ & \cdot R \cdot B^{-1} \cdot D \cdot F^{-1} \cdot U^{-1} \cdot B^{-1} \cdot U \cdot D^{-1} \quad (24 \text{ quarter turns}). \end{aligned} \tag{5.1}$$

Regarding the latter move, Jerry Bryan (in a Feb. 19, 1995, posting to the cube-lover's e-mail list, [CL]) showed that there is no fewer number of quarter-turn moves taken from

$$\{R, R^{-1}, L, L^{-1}, U, U^{-1}, D, D^{-1}, B, B^{-1}, F, F^{-1}\},$$

that will also give the superflip. In the jargon, Mike Reid's move is 'minimal in the quarter-turn metric'.

Remark 5.7.1. *There is a 'longer' element of the Rubik's Cube group, measured in the quarter-turn metric.*

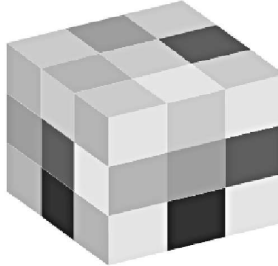


Figure 5.1: Michael Reid's superflip with four-spot.

*The **superflip with four-spot**, depicted in Figure 5.1, is given by 26 quarter turns maneuver*

$$\begin{aligned} \text{superflip4spot} = & U^2 \cdot D^2 \cdot L \cdot F^2 \cdot U^{-1} \cdot D \cdot R^2 \cdot B \cdot U^{-1} \cdot D^{-1} \cdot R \cdot \\ & \cdot L \cdot F^2 \cdot R \cdot U \cdot D^{-1} \cdot R^{-1} \cdot L \cdot U \cdot F^{-1} \cdot B^{-1}. \end{aligned} \tag{5.2}$$

This was proven by Mike Reid (who discovered the move and the above expression with the help of his own computer program) to be minimal in the quarter-turn

metric (see the Aug. 2, 1998, posting in [CL]). It is possible that there is no longer move than Reid's move in the Rubik's Cube group in the quarter turn metric.

Remark 5.7.2. At the time of this writing (spring 2008), there is no 'longer' element than the superflip of the Rubik's Cube group, measured in the face turn metric. It is possible that there is no longer move than the superflip in the Rubik's Cube group in the face turn metric.

5.7.2 Example: The two squares group

Let $H = \langle R^2, U^2 \rangle$ denote the group generated by the two square moves R^2 and U^2 of the Rubik's Cube. (The reader with a cube in hand may want to try the **Singmaster magic grip**: the thumb and forefinger of the right hand are placed on the front and back face of the fr , br edge; and the thumb and forefinger of the left hand are placed on the front and back face of the uf , ub edge. All moves in this group can be made without taking your fingers off the cube.) This group contains the useful 2-pair edge swap move $(R^2 \cdot U^2)^3$.

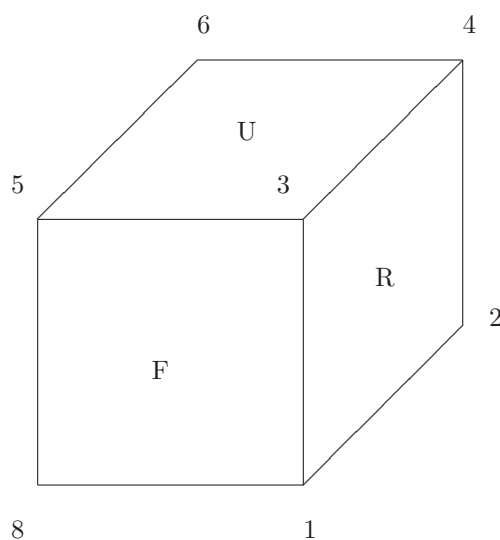


Figure 5.2: Corner labels for the cube

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We can find all the elements in this group fairly easily:

$$H = \{1, R^2, R^2 \cdot U^2, R^2 \cdot U^2 \cdot R^2, (R^2 \cdot U^2)^2, (R^2 \cdot U^2)^2 \cdot R^2, (R^2 \cdot U^2)^3, (R^2 \cdot U^2)^3 \cdot R^2, (R^2 \cdot U^2)^4, (R^2 \cdot U^2)^4 \cdot R^2, (R^2 \cdot U^2)^5, (R^2 \cdot U^2)^5 \cdot R^2\},$$

Therefore, $|H| = 12$. Note that $1 = (R^2 \cdot U^2)^6$, $U^2 = (R^2 \cdot U^2)^5 \cdot R^2$, and $U^2 \cdot R^2 = (R^2 \cdot U^2)^5$.

To discover more about this group, we label the vertices of the cube as in Figure 5.7.2.

The move R^2 acts on the set of vertices by the permutation $(1, 4)(2, 3)$ and the move U^2 acts on the set of vertices by the permutation $(4, 5)(3, 6)$. We label the vertices of a regular hexagon as in Figure 5.7.2.

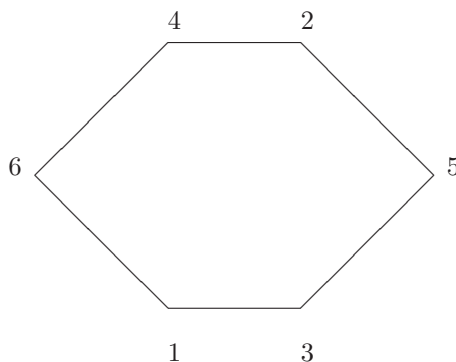


Figure 5.3: Corner labels for the hexagon

The permutation $(1, 4)(2, 3)$ is simply the reflection about the line of symmetry containing both 5 and 6. The permutation $(4, 5)(3, 6)$ is simply the reflection about the line of symmetry containing both 1 and 2. By a fact stated in section 5.3, these two reflections generate the symmetry group of the hexagon.

Example 5.7.3. *This example shows some ways of using SAGE to understand this group.*

SAGE

```
sage: rubik = CubeGroup()
sage: r = rubik.R()
sage: u = rubik.U()
sage: G = PermutationGroup([r^2, u^2])
sage: G.order()
12
```

```

sage: D6 = DihedralGroup(6)
sage: G.is_isomorphic(D6)
True

```

SAGE (which indirectly uses **GAP** for this computation) tells us that the two squares group is isomorphic to the dihedral group with 12 elements.

5.8 Commutators

When ‘playing’ with the Rubik’s Cube, certain operations can occur more often than others. Besides combining two moves together, another operation that takes place frequently is the operation ‘move 1, then move 2, then inverse of move 1 then inverse of move 2’. This type of move is called a ‘commutator’.

Definition 5.8.1. If g, h are two elements of a group G , then we call the element

$$[g, h] = g * h * g^{-1} * h^{-1}$$

the **commutator** of g, h .

Note that $[g, h] = 1$ if and only if g, h commute. Thus the commutator may be regarded as a rough measurement of the lack of commutativity. To put it simply: putting your socks on before the shoes is not the same as putting on shoes before the socks!

Ponderable 5.8.1. Let $G = S_3$, the symmetric group on 3 letters. Compute the commutators

$$[s_1, s_2], \quad [s_2, s_1].$$

Ponderable 5.8.2. Let R, U be as in the notation for the Rubik’s Cube moves introduced in the previous chapter. Determine the order of the move $[R, U]$. (Answer: 6)

Definition 5.8.2. (Singmaster [Si]) Let G be the permutation group generated by the permutations R, L, U, D, F, B regarded as permutations in S_{54} . The **Y commutator** is the element

$$[F, R^{-1}] = F \cdot R^{-1} \cdot F^{-1} \cdot R.$$

The **Z commutator** is the element

$$[F, R] = F \cdot R \cdot F^{-1} \cdot R^{-1}.$$

Ponderable 5.8.3.

(a) Guess why the Y commutator and the Z commutator have their names.

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(b) Find the order of the Y commutator and of the Z commutator.

(c) Find the order of $[R, [F, U]]$.

Example 5.8.1. If x, y are basic moves of the Rubik's Cube associated with faces that share an edge, then

(a) $[x, y]^2$ permutes exactly 3 edges and does not permute any corners;

(b) $[x, y]^3$ permutes exactly 2 pairs of corners and does not permute any edges.

Ponderable 5.8.4. Check to see if the above example is true!

Definition 5.8.3. Let G be any group. The group G' generated by all the commutators

$$\{[g, h] \mid g, h \text{ belong to } G\}.$$

This is called the **commutator subgroup** of G . Sometimes this is also denoted $[G, G]$.

This group may be regarded as a rough measurement of the lack of commutativity of the group G .

Remark 5.8.1. We will see later that the group generated by the basic moves of the Rubik's Cube— R, L, U, D, F, B —has a relatively large commutator subgroup. In other words, roughly speaking ‘most’ moves of the Rubik's Cube can be generated by commutators such as the Y commutator or the Z commutator.

Definition 5.8.4. If we repeatedly take commutator subgroups, we get a series of groups $G, G', G'' = (G')'$, and so on. The **derived series** of a group G is the sequence of subgroups

$$\dots \subset (G')' \subset G' \subset G.$$

A group G is called **solvable** if one of the groups in the derived series is the trivial group consisting only of the identity.

Morally speaking, if abelian groups are the best kind of group, then the next best thing is a solvable group.

The idea of a solvable group first arose and its importance emphasized in the work of Galois. However, it was Camille Jordan (1838–1922) who introduced the concept of a composition series (of which the derived series is a special case).

Jordan is also well-known for the ‘Jordan normal form’ in linear algebra. It is argued in Jordan's entry in [MT] that the theory of finite groups began with Jordan, his 1870 text **Traité des substitutions et des équations algébriques** being the first book on group theory.

Ponderable 5.8.5. Let G be an abelian group. Show that G is solvable.

Example 5.8.2. *This example shows how to use SAGE to test if the two squares subgroup of the cube group is solvable.*

SAGE

```
sage: rubik = CubeGroup()
sage: r = rubik.R()
sage: u = rubik.U()
sage: G = PermutationGroup([r^2, u^2])
sage: G.is_solvable()
True
```

Next, let us consider the commutator $g = [R, U] = R \cdot U \cdot R^{-1} \cdot U^{-1}$. Recall and n -cycle is a cyclic permutation of order n . As Example 5.8.1 predicts, SAGE tells us that, indeed, $[R, U]^2$ is composed of 3-cycles and $[R, U]^3$ is composed of 2-cycles:

SAGE

```
sage: g = r*u*r^(-1)*u^(-1)
sage: g^2
(1, 35, 9)(2, 5, 21)(3, 27, 33)(8, 25, 19)(24, 30, 43)(26, 28, 34)
sage: g^3
(1, 33)(3, 35)(8, 43)(9, 27)(19, 30)(24, 25)
```

5.9 Conjugation

Besides composing two moves together and using the commutator, another operation that takes place frequently on the Rubik's Cube is the operation 'move 1, then move 2, then inverse of move 1'. This type of move is called a 'conjugation'.

Definition 5.9.1. *If g, h are two elements of a group G , then we call the element*

$$g^h = h^{-1} * g * h$$

*the **conjugate** of g by h .*

Note that $g^h = g$ if and only if g, h commute. Thus the conjugates may be regarded as a rough measurement of the lack of commutativity. The exponential notation is justified by the following facts, which are easy to verify.

Ponderable 5.9.1. *Show that*

$$(a) \ (g_1 g_2)^h = g_1^h g_2^h,$$

$$(b) \ g^{h_1 h_2} = (g^{h_1})^{h_2}.$$

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Ponderable 5.9.2. Let $G = S_3$, the symmetric group on 3 letters, in the notation of Example 5.4.1. Compute the conjugations

$$s_1^{s_2}, \quad s_2^{s_1}.$$

Ponderable 5.9.3. Let R, U be as in the notation for the Rubik's Cube moves introduced in the previous chapter. Determine the order of the move R^U . (Answer: 4)

Definition 5.9.2. We say that two elements g_1, g_2 of G are **conjugate** if there is an element $h \in G$ such that $g_2 = g_1^h$.

It turns out that it is easy to see when two permutations $g, h \in S_n$ are conjugate: they are conjugate if and only if the cycles in their respective disjoint cycle decompositions have the same length when arranged from shortest to longest. (This result is due to Cauchy.) For example, the elements

$$g = (6, 9)(1, 3, 4)(2, 5, 7, 8), \quad h = (1, 2)(3, 4, 5)(6, 7, 8, 9)$$

are conjugate in S_9 . We shall leave the details and the proof for later (see §9.3.1).

Ponderable 5.9.4. Show that the notion of conjugate defines an equivalence relation. That is, show that

- (a) any element $g \in G$ is conjugate to itself (**reflexive**);
- (b) if g is conjugate to h (g, h belonging to G), then h is conjugate to g (**symmetry**);
- (c) if g_1 is conjugate to g_2 and g_2 is conjugate to g_3 then g_1 is conjugate to g_3 (**transitivity**).

Notation: The set of equivalence classes of G under the equivalence relation given by conjugation will be denoted G_* .

Ponderable 5.9.5. Let G be a finite group. Show that

- (a) $|G_*| \leq |G|$, and
- (b) $|G| = |G_*|$ if and only if G is abelian.

The polynomial

$$p_G(t) = \sum_{c \in G_*} t^{\text{ord}(c)}, \quad (5.3)$$

is called the **generating polynomial** of the order function on G .

Ponderable 5.9.6. Show two elements that are conjugate must have the same order.

(Hint: $(h^{-1}gh)^n = (h^{-1}gh)(h^{-1}gh) \dots (h^{-1}gh) = h^{-1}g^nh$, for $n = 1, 2, \dots$ and $g, h \in G$.)

Therefore, the expression $\text{ord} : G \rightarrow \mathbb{N}$ can actually be defined on G_* . Let

$$\text{ord}(\{h^{-1}gh \mid h \in G\}) = \text{ord}(g).$$

Notice that if $g_1, g_2 \in G_*$ are both of order d , then $t^{\text{ord}(g_1)} = t^{\text{ord}(g_2)} = t^d$. This tells us that

$$p_G(t) = \sum_{d \geq 1} n_G(d)t^d,$$

where $n_G(d)$ denotes the number of elements in G_* of order d . In other words, this polynomial encodes, via its coefficients, information on the elements of a given order in G . However, since it sums over conjugacy classes (not on elements of G themselves), it does not distinguish between two elements of G if they are conjugate.

In [Si], §5.10D, D. Singmaster asks for the possible orders of the elements of the Rubik's Cube group and how many elements of each order there are. (A method for determining this will be described later in this book; see §9.8.2.) This question of Singmaster motivates the following problem.

Ponderable 5.9.7. (*Hard*) Determine $p_G(t)$ for the Rubik's Cube group.

By hard, I mean you probably won't solve it overnight (at least not unless you have access to a very fast computer).

Example 5.9.1. For D_6 (the dihedral group of order 12), the generating polynomial is

$$2 * t^6 + 2 * t^3 + 7 * t^2 + t;$$

for S_8 , the generating polynomial is

$$t + 4t^2 + 2t^3 + 4t^4 + t^5 + 5t^6 + t^7 + t^8 + t^{10} + t^{12} + t^{15};$$

and for S_{12} , it is

$$\begin{aligned} t + 6t^2 + 4t^3 + 9t^4 + 2t^5 + 16t^6 + t^7 + 4t^8 + 2t^9 + 6t^{10} + \\ t^{11} + 9t^{12} + 2t^{14} + 2t^{15} + t^{18} + 2t^{20} + t^{21} + \\ t^{24} + t^{28} + 3t^{30} + t^{35} + t^{42} + t^{60}. \end{aligned}$$

For example, it follows from the last displayed computation that there exists a permutation of order 42 in S_{12} .

In SAGE, these computations are easy.

SAGE

```
sage: G = DihedralGroup(6)
sage: CG = G.conjugacy_classes_representatives()
sage: t = var("t")
sage: p = sum([t^(g.order()) for g in CG]); p
t^6 + t^3 + 3*t^2 + t
```

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Depending on your computer, the analogous computation for S_8 and S_{12} may take a long time to finish.

Verifying the above examples ‘by hand’ is a tough exercise in itself. **SAGE** comes in handy in these situations. Have fun playing with other groups too!

As was already mentioned, the maximal order (of any element) in the Rubik’s Cube group is 1260. This and the defining equation (5.3) tell us that the degree of the generating polynomial of the Rubik’s Cube group is 1260.

Definition 5.9.3. Fix an element g in a group G . The set

$$Cl(g) = Cl_G(g) = \{h^{-1} * g * h \mid h \in G\}$$

is called the **conjugacy class of g in G** . It is the equivalence class of the element g under the relation given by conjugation.

Example 5.9.2. This example will show how to use **SAGE** to compute all the conjugacy classes of a group and compute the above generating polynomial.

SAGE

```
sage: rubik = CubeGroup()
sage: r = rubik.R()
sage: u = rubik.U()
sage: G = PermutationGroup([r^2, u^2])
sage: CG = G.conjugacy_classes_representatives()
sage: p_G = sum([t^(g.order()) for g in CG]); p_G
t^6 + t^3 + 3*t^2 + t
```

In other words, there are 6 non-conjugate elements in the two squares subgroup of the cube group. Moreover, $p_G(t) = t^6 + t^3 + 3t^2 + t$. Notice that the generating polynomial of this group is the same as that of the group D_6 . In fact, these two groups are actually isomorphic.

Note that if $g_1 \in G$ is conjugate to $g_2 \in G$, then $Cl(g_1) = Cl(g_2)$.

How do you ‘compute’ conjugacy classes?

Algorithm

Input: A set S of generators of a permutation group G and an x belonging to G

Output: The conjugacy class of x , $Cl_G(x)$

Algorithm

```
class = {x}
for y in class do
  for g in S do
    if g*y*g^{-1} not in class then
      class = class union {g*y*g^{-1}}
```



```

        endif
    endfor
endfor

```

Note that the size of the list `class` in the `for` loop changes after each iteration of the loop. As mentioned before, the meaning of this is that the `if-then` command is to be executed exactly once for each element of the list.

Ponderable 5.9.8. Find the elements in $S_4 = \langle (1,2), (2,3), (3,4) \rangle$ that are conjugate to $(1, 2, 3, 4)$.

Theorem 5.9.1. Any finite group may be partitioned into its distinct conjugacy classes,

$$G = \cup_{c \in G_*} Cl(c).$$

Ponderable 5.9.9. Show that Theorem 5.9.1 is a consequence of Theorem 2.3.1.

If H is a subgroup of G and if g is a fixed element of G , then the set

$$H^g = \{g^{-1} * h * g \mid h \in H\}$$

is a subgroup of G . Such a subgroup of G is called a subgroup **conjugate** to H .

Ponderable 5.9.10. Let S be the set of all subgroups of G . We define a relation R on S by

$$R = \{(H_1, H_2) \in S \times S \mid H_1 \text{ is conjugate to } H_2\}.$$

Show that R is an equivalence relation.

Ponderable 5.9.11. Let $G = S_n$ and let $H = \langle g \rangle$ be a cyclic subgroup generated by a permutation g of the set $\{1, 2, \dots, n\}$. With respect to the equivalence relation in the previous problem, show that a subgroup K of G belongs to the equivalence class $[H]$ of H in G if and only if K is cyclic and is generated by an element k of G conjugate to $g \in G$.

5.10 Orbits and actions

Orbits are very important to understand for our purposes. Without them, understanding the mathematics of the Rubik's Cube is a little like trying to fix a car without having a wrench. The concept of an orbit is an essential tool in the toolbox as it helps us both conceptually and computationally with some of the combinatorics that arise later. Richard Feynmann (1918-1988), the famous physicist, was reported to have said something like: 'My job here is not to make

you understand such-and-such but to convince you not to give up. Because no one, not even me, really understands everything about this subject anyway. None-the-less, it is interesting and deserving of our time.’ Likewise, orbits are deserving of our time.

Let’s start our study of orbits with an example. Let G be the Rubik’s Cube group and X the set of facets of the Rubik’s Cube. Imagine having two identical cubes, and disassemble one and spread the facets out on a table. Imagine also that when you move the assembled cube, the facets on the table permute themselves around accordingly. You can see then that each element of G ‘moves’ the elements of X around. The mathematical terminology most commonly used to describe this type of situation is to say ‘ G acts on X ’. The precise general definition is given below.

Definition 5.10.1. *Let X be a set and let G be a group. We call X a **G -set** and we say G acts on X on the right provided the following conditions hold:*

1. *Each g belonging to G gives rise to a function*

$$\phi_g : X \rightarrow X.$$

2. *The identity 1 of the group G defines the identity function on X .*
3. *If g, h belong to G then the composite*

$$\phi_{gh} : X \rightarrow X$$

$$\text{satisfies } \phi_{gh}(x) = \phi_h(\phi_g(x)).$$

We call this a **right action**; see Ponderable 5.10.4 for the reason for this terminology.

One can, in general, define for all $g \in G$, ϕ_g to be the identity map from X to itself. Since this action doesn’t do anything, it is called the **trivial action**. We are interested only in non-trivial actions here.

We define ‘left actions’ similarly.

Definition 5.10.2. *Let X be a set and let G be a group. We say G acts on X provided the following conditions hold:*

1. *Each g belonging to G gives rise to a function*

$$\phi_g : X \rightarrow X.$$

2. *The identity 1 of the group G defines the identity function on X .*
3. *If g, h belong to G then the composite*

$$\phi_{gh} : X \rightarrow X$$

$$\text{satisfies } \phi_{gh}(x) = \phi_g(\phi_h(x)).$$

We call this action a **left action**; see Ponderable 5.10.3 for the reason for this terminology.

Remark 5.10.1.

- (a) We shall see another interpretation of these definitions later in §9.2.
- (b) Given a left action ϕ_g , one can create a right action by defining $\phi'_g = \phi_{g^{-1}}$.

Following the standard convention, the Rubik's Cube will act on the set of facets of the cube *on the right*. For example, start with a basic move, say U , which moves the up face clockwise a quarter turn. Now make another basic move, say F . The two moves together are denoted $U * F$, where F is written *to the right* of U .

Definition 5.10.3. Let G act on a set X . We call the action **transitive** if for each pair x, y belonging to X there is a $g \in G$ such that $y = \phi_g(x)$.

In other words, a group G acts transitively on a set X if *any* element x of X can be sent to *any* other element y of X by some element g of G (depending on x, y).

This notion of transitive is not to be confused with the notion of transitivity of equivalence relations, even though many orbits are equivalence classes.

The transitivity of an action turns out to be a strong condition for a group G acting on a set X to satisfy. This illustrates how unusual the Rubik's Cube is, since it acts transitively on two different sets! For example, it transitively permutes the set of corner subcubes and transitively permutes the set of edge subcubes.

Here are several more examples.

Example 5.10.1. Let X be a finite set and let $G = S_X$ be the symmetric group of X . Then X is a G -set and G acts transitively on X .

Ponderable 5.10.1.

- (a) Show that the action in the previous example is transitive.
- (b) (Hard) If $r = r_{n,d}$ denotes the number of groups of order n acting on a finite set with d elements and $s = s_{n,d}$ denotes the number of groups of order n acting transitively on a finite set with d elements, find a formula for r/s . (It seems likely to me that as n, d become large, r/s tends to 0. In this case, one can quantify the assertion that 'a random group acting on a finite set probably is not transitive'.)

Example 5.10.2. Let G be the group of all 2×2 invertible matrices with real entries, $G = GL_2(\mathbb{R})$. This group acts on the set of column vectors on the left.

Ponderable 5.10.2. Let G be the permutation group generated by the permutations R, L, U, D, F, B , regarded as elements of S_{48} . Let E denote the set of edges of the cube, which we identify with the set of edge subcubes. Let V denote

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the set of vertices of the cube, which we identify with the set of corner subcubes of the cube. Let X be the set of all movable subcubes of the Rubik's Cube (which we may identify as the union of E and V). Then G acts on X , E , and on V .

Questions:

- (a) Is the action of G on X transitive? (Explain.)
- (b) Is the action of G on E transitive? (Explain.)
- (c) Is the action of G on V transitive? (Explain.)

Ponderable 5.10.3. Let G be a group and let $X = G$. Define **left multiplication of G on X** by

$$\begin{aligned}\phi_g : X &\rightarrow X \\ x &\longmapsto \phi_g(x) = g * x.\end{aligned}$$

- (a) Show that left multiplication defines a left action of G on X .
- (b) Show that this action is transitive.
- (c) Show that each $\phi_g : G \rightarrow G$ is a permutation of the set G , so $\phi_g \in S_G$.

Ponderable 5.10.4. Let G be a group and let $X = G$. Define **right multiplication of G on X** by

$$\begin{aligned}\phi_g : X &\rightarrow X \\ x &\longmapsto \phi_g(x) = x * g.\end{aligned}$$

- (a) Show that right multiplication defines a right action of G on X .
- (b) Show that this action is transitive.
- (c) Show that each $\phi_g : G \rightarrow G$ is a permutation of the set G , so $\phi_g \in S_G$.

Ponderable 5.10.5. Let G be a group and let $X = G$. Define **conjugation on X** by

$$\begin{aligned}\phi_g : X &\rightarrow X \\ x &\longmapsto \phi_g(x) = g^{-1} * x * g.\end{aligned}$$

Show that conjugation defines an action of G on X (with X and G as above).

Ponderable 5.10.6. Let G be a group and let X denote the set of all subgroups of G . Define **conjugation on X** by

$$\begin{aligned}\phi_g : X &\rightarrow X \\ x &\longmapsto \phi_g(x) = g^{-1} * x * g.\end{aligned}$$

Show that this defines an action of G on X .

Remark 5.10.2. In general, the conjugation actions in the last two problems are not transitive.

Definition 5.10.4. Let G be a group acting on a set X . For each x belonging to X , the set

$$G * x = \{\phi_g(x) \mid g \in G\}$$

is called the **orbit** of $x \in X$ under G .

Given a group acting on a set X , how do you ‘compute’ the orbits? The following algorithm addresses this.

Algorithm

Input: A set S of generators of a permutation group G and an x belonging to X

Output: The orbit of x , $G * x$

Algorithm

```

orbit = {x}
for y in orbit do
  for g in S do
    if g*y not in orbit then
      orbit = orbit union {g*y}
    endif
  endfor
endfor

```

Note that the size of the list `orbit` in the `for` loop changes after each iteration of the loop. As mentioned before, the meaning of this is that the `if-then` command is to be executed exactly once for each element of the list `orbit`.

Ponderable 5.10.7. Let G be the Rubik’s Cube group and let x be the uf edge subcube. Find the orbit of x under the action of G using the above algorithm. Show each step.

Ponderable 5.10.8. Let G be the group of moves of the Rubik’s Cube and let X be the set of vertices of the cube. Let H be the subgroup of G generated by $U \cdot R$. Find:

- (a) the order of $U \cdot R$, (Answer: 105)
- (b) the orbit (in the Singmaster notation) of the ufr vertex in X under H .

Example 5.10.3. Let X be the set consisting of the 48 facets of the Rubik’s Cube that are not center facets—i.e., the ‘movable’ facets. Let V denote the subset of facets that belong to some corner subcube, and E the subset of facets that belong to some edge subcube. Let G denote the Rubik’s Cube group. As noted above, G acts on X , V , E . The action of G on X induced an equivalence relation as follows: we say that a facet f_1 is ‘equivalent’ to a facet f_2 if there is an element of G (i.e., a move of the Rubik’s Cube) that sends one facet to

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the other. By Ponderable 5.10.2, there are exactly two equivalence classes, or orbits, of G in X : V and E . In particular, the action of G on V is transitive and the action of G on E is transitive.

Let G be the Rubik's Cube group, V the set of vertices, and let $ruf \in V$ be the right upper front vertex. The left face move L clearly does not affect ruf . In this case, we say that L 'stabilizes' or 'fixes' ruf . The general definition is given below.

Definition 5.10.5. Let G be a group acting on a set X with the action denoted by ϕ . For each x belonging to X , the subgroup

$$\text{stab}_G(x) = G_x = \{g \in G \mid \phi_g(x) = x\}$$

is called the **stabilizer** of x in G .

Ponderable 5.10.9. Let G be a group acting on a set X , $\phi_g : X \rightarrow X$, for all $g \in G$. Show that, for all $x \in X$ and all $g \in G$, we have $\text{stab}_G(\phi_g(x)) = g^{-1} * \text{stab}_G(x) * g$.

Example 5.10.4. Let G be the group of symmetries of the square (see the example in §5.3 above), let X be the set of vertices of the square, and let x_0 be the vertex in the lower right hand corner. Then $\text{stab}_G(x_0) = \langle g_3 \rangle$.

Ponderable 5.10.10. Let G be any group and let $X = G$. Let G act on X by left multiplication:

$$\begin{aligned} \phi_g : X &\rightarrow X \\ x &\mapsto \phi_g(x) = g * x. \end{aligned}$$

Show that

$$\text{stab}_G(x) = \{1\},$$

for all x belonging to $X = G$.

Lemma 5.10.1. Let G be any group and let $X = G$. Let G act on X by conjugation:

$$\begin{aligned} \phi_g : X &\rightarrow X \\ x &\mapsto \phi_g(x) = g * x * g^{-1}. \end{aligned}$$

The stabilizer satisfies

$$\text{stab}_G(x) = \{g \in G \mid g * x = x * g\},$$

for all x belonging to $X = G$.

This is relatively easy, so the proof will be omitted.

The 'stabilizer' subgroup for the conjugation action,

$$C_G(x) = \{g \in G \mid g * x = x * g\};$$

is called the **centralizer** of x in G .

5.11 Cosets

Cosets are certain types of subsets of a group G . Like some political parties, they come in two ‘flavors’: left cosets and right cosets. As in politics, you don’t want to get them mixed up. However, they basically look the same (the cosets, that is), so if you just stick with either left cosets all the time or right cosets all the time, you’ll do fine. For typographical reasons, this book will usually use left cosets.

Example 5.11.1. Let G be the Rubik’s Cube group and $H = \{1, D, D^2, D^3\}$ the subgroup of moves of the down face. Let $g \in G$ be a move. The set $gH = \{g, gD, gD^2, gD^3\}$ is called a ‘left coset of H ’. The set $Hg = \{g, Dg, D^2g, D^3g\}$ is called a ‘right coset of H ’. Each move in Hg has the same effect as g , except you might have moved the down face afterwards.

Example 5.11.2. Let $H = S_2 = \{1, (1, 2)\}$ and

$$G = S_3 = \{1, (1, 2), (2, 3), (1, 3), (1, 2, 3), (1, 3, 2)\}.$$

The left cosets of H in G are

$$\begin{aligned} H, \quad (1, 2)H = H, \quad (2, 3)H &= \{(2, 3), (1, 2, 3)\}, \\ (1, 3)H &= \{(1, 3), (1, 3, 2)\}, \quad (1, 2, 3)H = \{(1, 2, 3), (2, 3)\} = (2, 3)H, \\ (1, 3, 2)H &= \{(1, 3, 2), (1, 3)\} = (1, 3)H. \end{aligned}$$

Note that $G = H \cup (2, 3)H \cup (1, 3)H$.

I leave it to the reader to compute the right cosets of H in G .

Ponderable 5.11.1. Let H be a subgroup of G and let $g, g' \in G$. Show that either $gH = g'H$ or $gH \cap g'H = \emptyset$.

Understanding cosets, like orbits, not only helps in counting arguments but also helps us understand the ‘group structure’ of G (i.e., knowing how G is constructed from certain ‘well-understood subgroups’). Notice that if $g \in H$ then $gH = Hg = H$. Conversely, if $gH = Hg = H$ then $g \in H$.

Definition 5.11.1. Let G be a group, written multiplicatively, and H a subgroup of G . For g belonging to G , the subset $g * H$ of G is called a **left coset** of H in G and the subset $H * g$ of G is called a **right coset** of H in G .

If G is an abelian group, written additively, and H is a subgroup, then left cosets and right cosets are the same: $H + g = g + H$.

Whether G is abelian or not, we shall sometimes abuse terminology and drop the word ‘left’ or ‘right’ in front of coset.

If G is non-abelian then the set of left cosets usually differs from the set of right cosets. However, when G is finite there are always the same number of left cosets as right cosets. It was Galois, in the early 1830’s, who first looked closely at subgroups for which the set of left cosets was the same as the set of right cosets. (Such subgroups are now called ‘normal’ and will be studied later.) Galois’ work was applied to the study of roots of polynomials.

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Example 5.11.3. Let $G = \mathbb{Z}$, $H = 12\mathbb{Z}$, under ordinary addition. If $n \in G$ is any integer then $n + H$ is a coset of H in G . In fact, all the cosets of H can be tabulated. Each coset occurs as a column exactly once in the following table:

\vdots	\vdots			\vdots
-12	-11	...	-2	-1
0	1	...	10	11
12	13	...	22	23
\vdots	\vdots			\vdots

Ponderable 5.11.2. If H is finite, show $|H| = |g * H| = |H * g|$.

Ponderable 5.11.3. If X is a left coset of H in G and x is an element of G , show that $x * X$ is also a left coset of H in G .

Notation: The set of all left cosets is written G/H and the set of all right cosets of H in G is denoted $H \backslash G$.

These two sets don't in general inherit a group structure from G but they are useful none-the-less. (G/H is a group with the 'obvious' multiplication $(g_1 * H) * (g_2 * H) = (g_1 g_2) * H$ if and only if H is a 'normal' subgroup of G . I will define 'normal' below.)

As an example of their usefulness, we have the following relationship between the orbits and the cosets of the stabilizers.

Proposition 5.11.1. Let G be a finite group acting on a set X . Then

$$|G * x| = |G/\text{stab}_G(x)|$$

for all x belonging to X .

Proof: The map

$$g * \text{stab}_G(x) \mapsto g * x$$

defines a function $f : G/\text{stab}_G(x) \rightarrow G * x$. The interested reader can easily check that this function is a bijection, since it is both an injection and a surjection. \square

Corollary 5.11.1. Let G be a finite group acting on itself by conjugation. Let $S \subset G$ denote a complete set of representatives of the conjugacy classes G_* in G and let $S' = S - Z(G)$, i.e., the subset of S of those elements that are not central. Then

$$G = \cup_{x \in S} Cl(x) \cong \cup_{x \in S} G/\text{stab}_G(x) = Z(G) \cup \cup_{x \in S'} G/\text{stab}_G(x),$$

for all x belonging to X . In particular,

$$|G| = \sum_{x \in S} |G/\text{stab}_G(x)| = |Z(G)| + \sum_{x \in S'} |G/\text{stab}_G(x)|.$$

Sometimes this is called the **class equation** or **class formula** for groups.

Proof: The first equality, $G = \cup_{x \in S} Cl(x)$, is Theorem 5.9.1. The second equality follows from the isomorphism $G/\text{stab}_G(x) \cong G * x$, established in the proof of above proposition.

By taking cardinalities, the second displayed equation is a consequence of the first. \square

Example 5.11.4. Using SAGE it is easy to illustrate the class equation in the case of S_6 :

SAGE

```

sage: G = SymmetricGroup(6)
sage: CG = G.conjugacy_classes_representatives()
sage: class_eqn = sum([G.order()/(G0.Centralizer(x)).Order()\
                        for x in CG])
sage: class_eqn == G.order()
True

```

(The line defining `class_eqn` had to be split in two for typographical reasons.)

I am finally ready to make good on my promise to prove Cauchy's theorem.

Corollary 5.11.2. *Theorem 5.5.1(a) holds.*

sketch: The argument is by induction on $|G|$.

The result is trivial if $|G| = 1$ (since then no prime divides $|G|$).

Suppose $|G| > 1$ and let p be a prime dividing $|G|$. By the induction hypothesis, we assume that the result is true for all subgroups H of G with $|H| < |G|$.

Suppose $x \in G$ is not central. Then its centralizer $C_G(x)$ is a proper subgroup of G . If $p \mid |C_G(x)|$ then the result follows from the induction hypothesis. If p does not divide $|C_G(x)|$ (for all non-central x), then since $|G| = |C_G(x)||G/\text{stab}_G(x)|$ (note the stabilizer and centralizer are the same in this case), it follows that p divides $|G/\text{stab}_G(x)|$. By Corollary 5.11.1 above, we must have $p \mid |Z(G)|$. If $Z(G)$ is a proper subgroup of G , then we are done by the induction hypothesis.

Thus we may assume $G = Z(G)$ is abelian. If G is cyclic then I leave it to the reader to show that G contains an element of order p . We shall assume that G is not cyclic. Let H be a proper subgroup of G of maximal order (this exists since G is not cyclic) and let $a \in G - H$. Then $G = H \cdot \langle a \rangle$ (else H would not be maximal). This implies p either divides $|H|$ or $|\langle a \rangle|$. The result follows from the induction hypothesis. For further details, see any elementary text on group theory. For example, see Theorem 10.2 (renamed Theorem 10.1.2 in the online version) in [G]. \square

Ponderable 5.11.4. *Let G be the group of symmetries of the square. Using the notation of Example 5.3.1, compute $G/\langle g_3 \rangle$ and $G * x_0$.*

Here is another version of Theorem 5.6.1:

Theorem 5.11.1. (Lagrange's Theorem) *If G is a finite group and H a subgroup then*

$$|G/H| = |G|/|H|.$$

Joseph Louis Lagrange (1736-1813) was an Italian who studied at the University of Turin and later taught mathematics there. He later taught at the University of Berlin and then took a research position at the Académie des Sciences in Paris, though he did teach at the Ecole Polytechnique. He made important contributions to mechanics, applications of calculus, and function theory. Many historians of mathematics believe that finite group theory began with Camille Jordan, who was born 25 years after Lagrange died. So it is fair to say that the original version of 'Lagrange's theorem' was not in the form you see above.

As an immediate consequence of the second version of Lagrange's theorem above, we obtain the first version (Theorem 5.6.1).

proof of theorem: Let X be the set of left cosets of H in G and let G act on X by left multiplication. Apply the previous lemma with $x = H$. \square

Ponderable 5.11.5. *Let $G = S_3$, the symmetric group on 3 letters, and let $H = \langle s_1 \rangle$, in the notation of §5.4 above.*

(a) *Compute $|G/H|$ using Lagrange's theorem.*

(b) *Explicitly write down all the cosets of H in G .*

Definition 5.11.2. *Let H be a subgroup of G and let C be a left coset of H in G . We call an element g of G a **coset representative** of C if $C = g * H$. A **complete set of coset representatives** is a subset of G , x_1, x_2, \dots, x_m , such that*

$$G/H = \{x_1 * H, \dots, x_m * H\}$$

*without repetition (i.e., all the $x_i * H$ are disjoint).*

Ponderable 5.11.6. *For $g_1, g_2 \in G$, define $g_1 \sim g_2$ if and only if g_1 and g_2 belong to the same left coset of H in G .*

(a) *Show that \sim is an equivalence relation.*

(b) *Show that the left cosets of H in G partition G .*

Ponderable 5.11.7. *For $G = S_4$ and $H = S_3$, find a complete set of coset representatives of G/H in G .*

Theorem 5.11.2. *If S is a complete set of coset representatives of G/H , then*

$$G = \cup_{s \in S} s * H$$

*and $|G| = \sum_{s \in S} |s * H|$.*

This follows from the above definition and Theorem 2.3.1, when G is finite.

5.12 Campanology, revisited

Let us put the material on bell-ringing from §3.5 in the context of group theory. The fact that this can be done so easily is quite remarkable, considering that group theory wasn't developed until at least 100 years later!

Generating the plain lead on four bells is analogous algebraically to generating the dihedral group of order 8, D_4 . If $a = (1, 2)(3, 4)$, which swaps the first two and last two bells, and if $b = (2, 3)$, which swaps the middle pair, then the plain lead on four bells corresponds to

$$D_4 = \{1, a, ab, aba, (ab)^2, (ab)^2a, (ab)^3, (ab)^3a\}.$$

Plain Bob Minimus is equivalent algebraically to generating the symmetric group on 4 elements, S_4 . Let a and b be as before. If we look at the first column of the Plain Bob Minimus composition, we see that it is nothing more than the dihedral group, D_4 , which is a subgroup of S_4 . To generate the second column of S_4 we introduce $c = (3, 4)$ and let $k = (ab)^3ac$. The second column corresponds to kD_4 and the third column to k^2D_4 . The generation of the Plain Bob Minimus shows that S_4 can be expressed as the disjoint union of cosets of the subgroup D_4 , that is, the cosets of D_4 in S_4 partition S_4 . This gives an example of the fact that, for any group G and any subgroup H , the cosets of H in G partition G (see also Theorem 5.11.2).

As White [Wh] concludes in his paper, he is not suggesting 'that Fabian Stedman was using group theory explicitly, but rather that group theoretical ideas were implicit in (Stedman's) writings and compositions'.

5.13 Dimino's algorithm

We saw in an earlier chapter a simple algorithm for computing the elements of a permutation group G . Because time is money and computer time is limited, efficiency is an important issue. We shall discuss a much more efficient algorithm in this section.

We recall Dimino's algorithm from Butler [Bu], chapter 3, but first some notation.

Notation: Let $S = \{g_1, g_2, \dots, g_n\}$ be a set of generators for a permutation group G . Let

$$\begin{aligned} S_0 &= \emptyset, \\ S_i &= \{g_1, \dots, g_i\}, \\ G_0 &= \{1\}, \\ G_i &= \langle S_i \rangle = \text{the group generated by the elements in } S_i, \end{aligned}$$

for $1 \leq i \leq n$.

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Example 5.13.1. *Here is an example to illustrate this using SAGE.*

Algorithm

```
sage: G = DihedralGroup(6)
sage: generators = G.gens(); generators
((1,2,3,4,5,6), (1,6)(2,5)(3,4))
sage: g1 = generators[0]; g2 = generators[1]
sage: G0 = G.subgroup([()])
sage: G0.order()
1
sage: G1 = G.subgroup([g1])
sage: G1.order()
6
sage: G2 = G.subgroup([g1, g2])
sage: G2.order()
12
```

Algorithm (inductive step):

Input: The generators S of G and a list L of all the elements of the permutation subgroup G_{i-1} .

Output: A list L of elements of G_i and a list C of coset representatives of G_i/G_{i-1} .

Algorithm

```
C = {1}
for g in C do
  for s in S_i do
    if s*g not in L then
      C = C union {s*g}
      L = L union s*g*G_{i-1}
    endif
  endfor
endfor
```

Algorithm (Dimino):

Input: The generators S of G
Output: A list of elements of G

Algorithm

```

(Initial case  $G_1 = \langle g_1 \rangle$ )
order = 1, element[1] = 1,  $g = g_1$ 
while  $g$  not equal to 1 do
  order = order + 1
  element[order] =  $g$ 
   $g = g * g_1$ 
endwhile

(General case)
for  $i$  from 2 to  $n$  do
  <insert inductive step here>
endfor

```

Example 5.13.2. Let $G = S_3 = \langle s_1, s_2 \rangle$. We use Dimino's algorithm to list all the elements of G . We have

$$G_0 = \{1\} \subset G_1 = \langle s_1 \rangle \subset G_2 = G.$$

First, we list the elements of $G_1 = \langle s_1 \rangle$. Since $s_1 = (1\ 2)$, it is order 2, so

$$G_1 = \{1, s_1\}.$$

This is our list L to which we will apply the 'inductive step' of Dimino's algorithm (with $i = 2$). We start with $C = \{1\}$. Now we look at the left cosets of G_1 in $G_2 = G$. We have (with $g = 1, s = s_1$)

$$s_1 * G_1 = G_1,$$

so we don't increase the size of C or L . Next, we have (with $g = 1, s = s_2$)

$$s_2 * G_1 = \{s_2, s_2 * s_1\} \neq G_1,$$

so $L = \{1, s_1, s_2, s_2 * s_1\}, C = \{1, s_2\}$. Next, we have (with $g = s_2, s = s_1$)

$$s_1 * s_2 * G_1 = \{s_1 * s_2, s_1 * s_2 * s_1\} \neq G_1.$$

(We know $s_1 * s_2 * G_1 \neq G_1$ since neither of the two elements in $s_1 * s_2 * G_1$ is the identity.) Thus, we increase L, C :

$$L = \{1, s_1, s_2, s_2 * s_1, s_1 * s_2, s_1 * s_2 * s_1\},$$

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and $C = \{1, s_2, s_1 * s_2\}$. We know we may stop here since we know $|S_3| = 6$ but the algorithm still has one more statement to execute. Next, we have (with $g = s_2, s = s_2$)

$$s_2 * s_2 * G_1 = G_1,$$

so we don't increase the size of C or L (as expected). This step terminates the algorithm and $S_3 = L$.

Ponderable 5.13.1. Perform Dimino's algorithm on

$$S_4 = \langle s_1 = (1\ 2), s_2 = (2\ 3), s_3 = (3\ 4) \rangle.$$

Chapter 6

Welcome to the machine

...Organization is of the utmost importance for military affairs, as it is ... for other disciplines where the gathering process of practical knowledge exceeds the strength of any individual. In mathematics, however, organizing talent plays a most subordinate role. Here weight is carried only by the individual. The slightest idea of a Riemann or a Weierstrass is worth more than all organizational endeavours. There is no royal road to mathematics.

Georg Frobenius

This chapter expounds upon the mathematical structure in some mathematical models of some games belonging to the ‘Merlin’s Magic’ family.

In particular, we will see that for a randomly chosen initial state, the probability that

- the 3×3 Lights Out can be solved is 1,
- the keychain Lights Out can be solved is 1,
- the 5×5 Lights Out can be solved is 0.25,
- the 6×6 Deluxe Lights Out can be solved is 1,
- the Orbix can be solved is 1.

6.1 Some history

Some years ago, Parker Brothers (also owned by Hasbro) marketed an electronic device called **Merlin’s Magic** (the rules book was published in 1979). This device was marketed as a 3×3 square array of buttons which could turn off or on. Tiger Electronics, a subsidiary of Hasbro, sometime in the late 1990’s marketed several generalizations under the name **Lights Out** (or Lights Out Cube or Lights Out Keychain or ...). These are battery-operated puzzles.

The rule for how each button changes the state is very simple: pressing any button toggles on/off the neighboring buttons (what ‘neighboring’ means will be defined in the next section). Other games in the same family include Rubik’s Clock, Lights Out, and Orbix.

6.2 Merlin’s Machine

Let $\mathbb{Z}/m\mathbb{Z} = \{0, 1, \dots, m-1\}$, where you add, subtract, and multiply mod m . When these numbers correspond to the colors of a game, we shall let zero correspond to ‘off’ and the non-zero numbers to the different lighted colors.

6.2.1 The machine

We shall use some terminology from graph theory, which we have collected in §7.1 of the next chapter. For now, we need only the notion of a ‘connected graph’. Roughly speaking, a connected graph is a collection of edges and vertices (an edge connects two distinct vertices) which is not the disjoint union of two subgraphs (see §7.1).

The graph: Here is the graph-theoretical description of the situation. Let Γ be a finite connected graph. Let V denote the set of vertices and E the set of edges of Γ . The vertex set V represents the set of buttons. Two vertices are called **neighbors** if they are connected by an edge. Suppose $V = \{v_1, v_2, \dots, v_n\}$. The set of all n -tuples of 0’s, 1’s, ..., $m-1$ ’s (we think of these n -tuples as being indexed by the elements of V) will be called the **states** of Γ and denoted $(\mathbb{Z}/m\mathbb{Z})^n$. Let $s_0 \in (\mathbb{Z}/m\mathbb{Z})^n$ denote an ‘initial state’ (and assume s_0 contains at least one non-zero entry). If $v \in V$ is any vertex, let $t(v) \in (\mathbb{Z}/m\mathbb{Z})^n$ denote the state (sometimes called the **toggle vector**) which is 1 at each coordinate that is either indexed by v or by a neighbor of v , and which is 0 otherwise. (In the case of the Orbix puzzle, let $t(v) \in (\mathbb{Z}/m\mathbb{Z})^n$ denote the state (sometimes called the **toggle vector**) which is 1 at each coordinate that is indexed by a neighbor of v (but *not* v itself), and which is 0 otherwise.)

The rules:

Here’s how the game is played.

1. Pick any vertex v .
2. Replace the current state of the game and call it s ($s = s_0$ initially), with $s' = s + t(v)$ (where addition is coordinatewise mod m).
3. If all the coordinates of this new state s' are equal to 0 then stop. Otherwise, go to step 1.

Motivated by §6.4, we shall call this game **Merlin’s Machine**.

It is possible to generalize the above construction slightly, replacing a graph by a digraph with multiple (parallel) labeled edges. This generalization is sufficiently different from the others that it has been omitted. However, Rubik’s Clock falls into this more general framework but with $m = 12$.

In most examples given below, $m = 2$.

6.2.2 The rectangular graph

The general setup was described in the previous subsection. Since the most popular games are rectangular, we shall construct the graph in this special case for easier reference.

Let $M \geq 2, N \geq 2$ be integers. (For example, in the case of Merlin's Magic, we have $M = N = 3$.) Subdivide a square in the plane into $M \cdot N$ equal subsquares whose edges are parallel to the edges of the larger square. Each of these smaller subsquares represents a button in Merlin's Machine. We shall label the $M \cdot N$ buttons on this $M \times N$ array of squares $(S_{i,j})$, as though they were entries of a matrix. Thus, for example, $S_{1,1}$ is in the upper left-hand corner, $S_{1,N}$ is in the upper right-hand corner, $S_{M,1}$ is in the lower left-hand corner, and $S_{M,N}$ is in the lower right-hand corner.

Let Γ denote the graph whose vertices are the $M \cdot N$ squares $S_{i,j}$, where two vertices are connected by an edge if and only if the corresponding squares share a common side. There are $N(M - 1) + M(N - 1)$ edges. This is the graph of the rectangular Merlin's Machine. The rules are as in §6.2.1 above.

6.3 Variants

Several variants of Merlin's Machine have been patented and marketed. They all fit into the framework of §6.2.1.

6.3.1 Merlin's Magic and 3×3 Lights Out

For simplicity, I only describe the 3×3 Lights Out (for Merlin's Magic, which is a slightly different game, see [P], [Sto], [Sch1]). Imagine subdividing a square into 9 equal subsquares in a 3×3 grid. Take $m = 2$. Here the graph Γ of §6.2.1 is the graph of a 3×3 square array with 9 vertices, where two vertices are connected by an edge if and only if the corresponding squares share a boundary. Thus, each interior vertex has degree 4, each edge vertex (not a corner) has degree 3, and each corner vertex has degree 2. The rules are as in §6.2.1.

For the solution strategy, see [P], [Sto], [Sch1].

6.3.2 The Orbix

Here $m = 2$ and the graph Γ of §6.2.1 is the graph of the icosahedron with 12 vertices, each having degree 5. (Actually, Orbix has four different puzzle types, of which the one described is Type 1. The others have different rules for play, as follows. Type 2: Same as type 1, but only lit buttons can be pressed. Type 3: Only unlit buttons can be pressed, and then all buttons *except* the adjacent ones change. Type 4: Only lit buttons can be pressed. If the button opposite the pressed one is on, then the adjacent 5 lights change. If the button opposite the pressed one is off, then all buttons except the pressed one change. These other types might require analysis and ideas beyond what is presented here.).

6.3. VARIANTS

The simplest solution strategy (for the Type 1 puzzle) is given on Jaap's Orbix page at [Sch2]: Simply press a button if an even number of the adjacent buttons are lit. Repeat this until the puzzle is solved.

6.3.3 Keychain Lights Out

Imagine subdividing a square into equal subsquares in a 4×4 grid. Take $m = 2$. Here the graph Γ of §6.2.1 is the graph of a 4×4 array with 16 vertices, where each edge vertex is *also* connected to its mirror-opposite vertex (reflecting about the thick central lines in Figure 6.1 below) on the opposite edge. (The corner vertices actually have two mirror opposites.) Thus, each vertex has degree 4. This may be visualized as the graph of a 4×4 grid on a donut obtained by 'gluing' top and bottom edges and then 'gluing' the left and right edges.

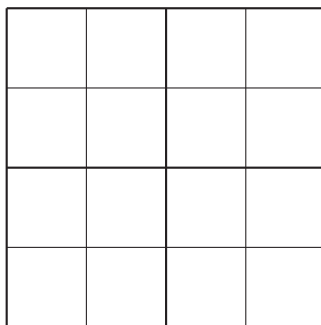


Figure 6.1: Keychain Lights Out Grid.

The rules are as in §6.2.1.

(The 'ordinary' 4×4 game without wraparound has not been analyzed in detail here. However, see Example 6.5.5.)

6.3.4 Lights Out

Imagine subdividing a square into 25 equal subsquares in a 5×5 grid. Take $m = 2$. Here the graph Γ of §6.2.1 is the graph of a 5×5 square array with 25 vertices, where two vertices are connected by an edge if and only if the corresponding squares share a boundary. Thus, each interior vertex has degree 4, each edge vertex (which is not a corner) has degree 3, and each corner vertex has degree 2. The rules are as in §6.2.1.

6.3.5 Deluxe Lights Out

Imagine subdividing a square into equal subsquares in a 6×6 grid. Take $m = 2$. Here the graph Γ of §6.2.1 is the graph of a 6×6 array with 36 vertices, each interior vertex having degree 4, each edge vertex (which is not a corner) having degree 3, and each corner vertex having degree 2. The rules are as in §6.2.1.

Like the Orbix, Deluxe Lights Out has several different types of play, such as ‘lit only’ (where you can punch only lit buttons) and ‘toggle’ (where you must alternately punch lit and then unlit buttons). Here we only describe the easiest ‘classic’ mode of play.

6.3.6 Lights Out Cube

Take $m = 2$. For this puzzle, let Γ be the graph whose vertices are the 54 facets (i.e., the colored tiles) of a Rubik’s Cube. Join two vertices by an edge of the facets share a common boundary. Now play the game following the rules of §6.2.1.

For the solution strategy, see [Ta].

6.3.7 Alien Tiles

Take $m = 4$. (This is the only example where $m \neq 2$.) The Internet site <http://www.alientiles.com>, maintained by Cliff Pickover and Cam Mckechnie (who have also trademarked the term ‘Alien Tiles’), provides one with a game based on an idea similar to Lights Out. Here the game is a 7×7 array of colored buttons. The graph Γ of §6.2.1 is the graph of a 7×7 array with 49 vertices, where two vertices are connected if their corresponding buttons are either in the same row or in the same column. Now play the game following the rules of §6.2.1.

6.3.8 Theoretical generalizations and variants

Besides the Rubik’s Clock already mentioned, other related puzzles include the class of ‘chip-firing’ games investigated, for example, by A. Björner, L. Lovasz, and P. Shor [BLS] and K. Eriksson [E]. Though extremely interesting, their analysis is much more technical and would take us too far afield, so we shall not delve further into them.

6.4 Finite-state machines

This section introduces a precise language which is helpful for comparing electrical devices.

A **finite-state machine** (or **Mealy machine**) is a 5-tuple (S, I, O, f, g) where S is a (finite) set of objects called **states**, I is a finite set of symbols called the **input alphabet**, O is a finite set of symbols called the **output alphabet**, $f : S \times I \rightarrow S$ is the **transition** or **next-state** function, $g : S \times I \rightarrow O$ is the **output** function. (One reference for this is, for example, Grimaldi [Gr].)

Remark 6.4.1. *The states of the machine can be roughly regarded as the ‘memory registers’ of the machine. The input alphabet can be roughly regarded as a*

6.4. FINITE-STATE MACHINES

labeling of the buttons one might press (or moves you make, depending on the constitution of the machine).

Remark 6.4.2. More precisely, f and/or g need not be defined on all of $S \times I$. One or both may only be defined on some subset of $S \times I$ (in other words, f and g need only be what is called a **partial function**). When both f and g are functions (each defined on all of $S \times I$) then the machine is called **complete**.

Define $E_{i,j}$ to be the $M \times N$ matrix all of whose entries are 0 except for the (i, j) th entry, which is equal to 1. An **elementary matrix** is an $M \times N$ matrix which is equal to $E_{i,j}$ for some i, j such that $1 \leq i \leq M, 1 \leq j \leq N$. Let $M_{M \times N}(\mathbb{Z}/2\mathbb{Z})$ denote the set of all $M \times N$ matrices whose entries are either equal to 0 or equal to 1. Clearly, $E_{i,j} \in M_{M \times N}(\mathbb{Z}/2\mathbb{Z})$. We can, if we want, add two elements of $M_{M \times N}(\mathbb{Z}/2\mathbb{Z})$ (adding each corresponding matrix element mod 2) and get another matrix in $M_{M \times N}(\mathbb{Z}/2\mathbb{Z})$.

Example 6.4.1. (*Rubik's Cube*) A **Rubik's Cube** is a cube that has been sliced into 27 subcubes of equal sizes, each of which has 6 smaller faces parallel to the 6 faces of the larger cube. In the 'solved state' each face of the larger cube is colored a different (solid) color. To be concrete, suppose we have fixed such a cube in space (say on a tabletop in front of you, the reader). The 9 facets of the front face are yellow, the 9 facets of the back face are green, the 9 facets of the up (top) face are red, the 9 facets of the down (bottom) face are orange, the 9 facets of the right face are blue, and the 9 facets of the left face are white. There are $9 \cdot 6 = 54$ colored facets total. Each of the 6 faces has 3 parallel slices of 9 subcubes, each of which may be independently rotated at angles of 90 degrees, 180 degrees, and 270 degrees.

With the yellow center facet facing front, and the red center facet facing up, let R denote the 90 degrees clockwise rotation (as you face that side) of the right slice, let L denote the 90 degrees clockwise rotation of the left slice, let F denote the 90 degrees clockwise rotation of the front slice, let B denote the 90 degrees clockwise rotation of the back slice, let D denote the 90 degrees clockwise rotation of the down slice, and let U denote the 90 degrees clockwise rotation of the up slice. Note that each of these rotations leaves the center facets fixed. There are $9 \cdot 6 - 6 = 48$ colored non-central facets total.

Let S denote the set of all possible permutations of the 48 colored facets on the cube (leaving the center facets fixed). Let $I = \{R, L, F, B, U, D\}$. Let $O = S$. Let the transition function $f : S \times I \rightarrow S$ be the function sending (s, i) ($s \in S$ and $i \in I$) which takes the state s (the cube with 48 non-central facets scrambled according to some permutation) and the input i (one of the moves R, L, \dots, U) to the new state s' obtained by performing the move i on the scrambled cube s : $f(s, i) = s'$. Let the output function g be the same as f .

This describes the Rubik's Cube as a finite state machine.

Example 6.4.2. (*rectangular Merlin's Machine*) In this example, let $M, N > 2$ be integers. Let

$$S = M_{M \times N}(\mathbb{Z}/2\mathbb{Z}), \quad I = \{E_{i,j} \mid 1 \leq i \leq M, 1 \leq j \leq N\}, \quad O = M_{M \times N}(\mathbb{Z}/2\mathbb{Z}).$$

We interpret the input $E_{i,j}$ to mean that the (i,j) th button of the array was pressed. The effect of pressing any button is to toggle the button itself and those which are directly north, south, east, or west of the button pressed. There is no wraparound (at least not in the most basic version) so there are

- 3 buttons toggled in case a corner button is pressed,
- 4 buttons toggled in case an edge button is pressed,
- 5 buttons toggled in case a central button is pressed.

To model this as a finite-state machine, let the transition function $f : S \times I \rightarrow S$ be defined by

$$f(s, E) = s + t(E), \quad s \in S, E \in I,$$

where addition is coordinate-wise mod 2 and $t(E)$ is the **toggle matrix**:

$$t(E_{i,j}) = \begin{cases} E_{i,j} + E_{i-1,j} + E_{i,j-1} + E_{i+1,j} + E_{i,j+1}, & 1 < i < M, 1 < j < N, \\ E_{i,j} + E_{i-1,j} + E_{i+1,j} + E_{i,j+1}, & 1 < i < M, j = 1, \\ E_{i,j} + E_{i+1,j} + E_{i,j+1}, & i = 1, j = 1, \\ E_{i,j} + E_{i-1,j} + E_{i,j+1}, & i = M, j = 1, \\ E_{i,j} + E_{i,j-1} + E_{i+1,j} + E_{i,j+1}, & i = 1, 1 < j < N, \\ E_{i,j} + E_{i-1,j} + E_{i,j-1} + E_{i,j+1}, & i = M, 1 < j < N, \\ E_{i,j} + E_{i-1,j} + E_{i,j-1} + E_{i+1,j}, & 1 < i < M, j = N, \\ E_{i,j} + E_{i,j-1} + E_{i+1,j}, & i = 1, j = N, \\ E_{i,j} + E_{i-1,j} + E_{i,j-1}, & i = M, j = N. \end{cases}$$

For example, if $M = N = 3$ and $E = E_{1,2}$ then

$$t(E) = \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \end{pmatrix},$$

and if $E = E_{2,2}$ then

$$t(E) = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 0 \end{pmatrix}.$$

Finally, let the output function $g : S \times I \rightarrow O$ be the same as the transition function $g = f$.

This describes *Lights Out* as a finite-state machine.

6.5 The mathematics of the machine

We shall follow Anderson and Feil's [AF] presentation to begin with. The idea is generalized and extended considerably in the papers of Goldwasser, Klostermeyer, et al. referenced below. The essential idea presented in this section was

posted on the sci.math newsgroup in Feb. 1998, months before [AF] appeared, first by Carsten Haese and then by Dave Rusin [Rus]. It appears that Haese thought of the idea independently of Anderson and Feil.

6.5.1 The square case

Suppose, to begin, that $M = N$. Represent each state $s \in M_N(\mathbb{Z}/2\mathbb{Z})$ as an N^2 -tuple

$$\vec{s} = (s_{1,1}, s_{1,2}, \dots, s_{1,N}, s_{2,1}, \dots, s_{N,N}).$$

Let B denote the $N \times N$ banded matrix with 1's on the diagonal, subdiagonal, and superdiagonal, and 0's elsewhere:

$$B = \begin{pmatrix} 1 & 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 1 & 0 & \dots & 0 \\ 0 & 1 & 1 & 1 & \dots & 0 \\ \vdots & & & & & \\ 0 & 0 & \dots & 0 & 1 & 1 \end{pmatrix}.$$

Let I_N denote the $N \times N$ identity matrix. Let A denote the $N^2 \times N^2$ matrix

$$A = \begin{pmatrix} B & I_N & 0 & 0 & \dots & 0 \\ I_N & B & I_N & 0 & \dots & 0 \\ 0 & I_N & B & I_N & \dots & 0 \\ \vdots & & & & & \\ 0 & 0 & \dots & 0 & I_N & B \end{pmatrix},$$

where (for simplicity of notation) 0 denotes the $N \times N$ matrix of 0's. The i th column of A represents the toggle vector associated with pressing the i th button (in the ordering $(1,1), (1,2), \dots, (1,N), (2,1), \dots, (N,N)$).

Let \vec{s}_0 denote any initial state. The problem of solving the puzzle is that we want to know which buttons to press (i.e., which columns of A to add together) *starting from the solved position* to obtain the given state \vec{s}_0 . (For once we know this, we can work backwards from \vec{s}_0 to solve it.) In other words, we want to add some columns of A together (addition being coordinatewise (mod 2)) and get \vec{s}_0 . This fact may be recorded as follows.

Theorem 6.5.1. *The puzzle with initial state \vec{s}_0 can be solved if and only if the matrix equation $A\vec{x} = \vec{s}_0$ has a solution $\vec{x} \in (\mathbb{Z}/2\mathbb{Z})^{N^2}$.*

With a little knowledge of linear algebra, the following question immediately arises: is $\det(A) \neq 0$ in $\mathbb{Z}/2\mathbb{Z}$? Using the computer algebra program **SAGE**, one finds that the answer is no if $N = 4, 5$ and yes if $N = 3, 6$. This has the following implications.

- No matter what initial states are given for the 3×3 Merlin's Machine and the 6×6 Merlin's Machine, the puzzle can always be solved.

- There are initial states for the 4×4 Merlin's Machine and the 5×5 Merlin's Machine from where the puzzle cannot be solved.

What about other values of N ?

To answer this question, we first introduce a sequence of polynomials. Let

$$f_0(x) = 0, \quad f_1(x) = 1,$$

and for $n > 1$, let

$$f_n(x) = x f_{n-1}(x) + f_{n-2}(x).$$

This defines the polynomials $f_0(x)$, $f_1(x)$, $f_2(x)$, ..., inductively. The polynomial $f_i(x)$ is called the i th **Fibonacci polynomial**.

Example 6.5.1. *Here is an example of some SAGE/Python code which computes the first several Fibonacci polynomials.*

SAGE

```

sage: R = PolynomialRing(QQ, "x")
sage: x = R.gen()
sage: def f(t,n):
        if n<2: return R(n)
        else:
            return f(t,n-2)+t*f(t,n-1)
sage: [f(x,n) for n in [0,1,2,3,4,5]]
[0,1, x, x^2 + 1, x^3 + 2*x, x^4 + 3*x^2 + 1]
```

If you decide to type this in yourself, please note that indentation is very important in Python, so be careful when entering in the function f above.

Let $B_n(x)$ denote the $n \times n$ matrix

$$B_n(x) = \begin{pmatrix} x & 1 & 0 & 0 & \dots & 0 \\ 1 & x & 1 & 0 & \dots & 0 \\ 0 & 1 & x & 1 & \dots & 0 \\ \vdots & & & & & \\ 0 & 0 & \dots & 0 & 1 & x \end{pmatrix}.$$

Lemma 6.5.1. *(Goldwasser, Klostermeyer, Trapp, and Zhang [GKTZ]) For all $n > 1$,*

$$\det(B_n(x)) = f_{n+1}(x).$$

Recall that the determinant was defined in §2.2.5.

The following effectively answers the question of when the $N \times N$ Merlin's Machine can always be solved, no matter what the initial state is.

Theorem 6.5.2. *(Goldwasser, Klostermeyer, and Trapp [GKT]) The $N \times N$ puzzle can be solved with any initial state if and only if the greatest common divisor of $f_{N+1}(x)$ and $f_{N+1}(x+1)$ equals 1.*

6.5.2 Downshifting

We need to slow down and recall a few more mathematical concepts in linear algebra. These are analogs of standard facts for real matrices introduced in Definition 2.1.2 above. However, we shall eventually need them for matrices having coefficients in a *finite* field (finite fields, such as $\mathbb{Z}/p\mathbb{Z}$ (also denoted $GF(p)$), are discussed in §12.2 below).

Let F be either \mathbb{R} or \mathbb{C} or $\mathbb{Z}/p\mathbb{Z} = GF(p)$, where p is a prime.

Definition 6.5.1. *Let V be a set with operations $+$: $V \times V \rightarrow V$ and \cdot : $F \times V \rightarrow V$. V is called a **vector space over F** (or an **F -vector space**) if the following properties hold. For all $\vec{u}, \vec{v}, \vec{w} \in V$ and $a, b \in F$,*

- $\vec{u} + \vec{v} = \vec{v} + \vec{u}$ (commutativity)
- $(\vec{u} + \vec{v}) + \vec{w} = \vec{u} + (\vec{v} + \vec{w})$ (associativity)
- there is a vector $\vec{0} \in V$ which satisfies $\vec{u} + \vec{0} = \vec{u}$ (the zero vector $\vec{0}$ is the ‘additive identity’),
- for each $\vec{v} \in V$ the element $(-1)\vec{v} = -\vec{v} \in V$ satisfies $\vec{v} + (-\vec{v}) = \vec{0}$ (each element \vec{v} has an ‘additive inverse’ $-\vec{v}$)
- $(a + b)\vec{v} = a\vec{v} + b\vec{v}$ and $a(\vec{v} + \vec{w}) = a\vec{v} + a\vec{w}$ (distributive laws)
- $(ab)\vec{v} = a(b\vec{v})$
- $1 \cdot \vec{v} = \vec{v}$

In particular, every vector space must contain at least one element (the zero vector). The elements of V are called **vectors** and the elements of F are called **scalars**.

Any subset $W \subset V$ of a vector space which is closed under vector addition (restricting $+$ from V to W) and scalar multiplication is called a **subspace** of V .

Example 6.5.2. *The ‘mod 2 analog’ of the basic example of the ‘3-space’, $V = \mathbb{R}^3$, is $(\mathbb{Z}/2\mathbb{Z})^3$ consisting of all vectors (x, y, z) in three dimensions whose coordinates (and scalars) belong to $\mathbb{Z}/2\mathbb{Z}$. All arithmetic is mod 2.*

Example 6.5.3. *Another example of a vector space is the space of all $m \times n$ matrices having entries in F , $M_{m \times n}(F)$ (m rows and n columns). Let $m = n > 1$ and let $1 \leq i, j \leq n$. Define the matrix $E_{ij} \in M_{n \times n}(F)$ to have entry 1 at the (i, j) coordinate and 0 at every other coordinate. This is called the (i, j) **th elementary matrix**.*

Suppose V is a vector space over a field F . If there are non-zero vectors $\vec{f}_1, \vec{f}_2, \dots, \vec{f}_k$ such that every $\vec{v} \in V$ can be written as a ‘linear combination’ of these \vec{v}_i ’s, i.e.,

$$\vec{v} = c_1 \vec{f}_1 + c_2 \vec{f}_2 + \dots + c_k \vec{f}_k,$$

for some scalar coefficients $c_i \in F$ ($i = 1, 2, \dots, k$), then we say that V is **spanned** by the elements $\vec{f}_1, \vec{f}_2, \dots, \vec{f}_k$. We also say that elements $\vec{f}_1, \vec{f}_2, \dots, \vec{f}_k$ form a **spanning set** for V .

Definition 6.5.2. *If a vector space V has a finite spanning set, then we say that V is **finite dimensional**. The **dimension** of V over F is the number of vectors in any minimal spanning set of V (it turns out this number does not depend on the minimal spanning set chosen, so this makes sense). The dimension of V is written $\dim(V) = k$. If V is of dimension k , then there are k vectors which span V and any such set of k spanning vectors of V , is called a **basis** of V .*

If W is a subspace of V then the number $\dim(V) - \dim(W)$ is called the **codimension** of W in V .

The dimension of a vector space measures how many degrees of freedom—how ‘big’—the space is. The codimension of a subspace measures how many degrees of freedom V has relative to W .

Let $V = \mathbb{R}^n$ and let A be an $n \times n$ matrix having entries in F . Then A defines a function $A : V \rightarrow V$. If $v \in V$ is a non-zero vector and $\lambda \in \mathbb{C}$ with the property that

$$Av = \lambda v,$$

then we say λ is an **eigenvalue** of A with **eigenvector** v .

The intuitive, geometric way to visualize eigenvalues and eigenvectors is to think of the action of A on a vector as being some sort of combination of ‘rotations’, ‘shears’ (as in the shear force on an airplane), and ‘stretches’. However, if that vector is an eigenvector, then the action of A is simply a stretch (by a factor of λ , at least if λ is real). The following is a basic fact about eigenvalues which is found in any textbook on linear algebra (for example, [Be]).

Lemma 6.5.2. *Assume A is a square matrix having entries in \mathbb{C} .*

- (a) *λ is an eigenvalue of A if and only if $\det(A - \lambda I) = 0$, where I denotes the identity matrix.*
- (b) *If A is an $n \times n$ matrix, then A has exactly n eigenvalues (counted according to multiplicity).*

The result above still holds if \mathbb{C} is replaced by any algebraically closed field. (For a precise definition of ‘algebraic closure’, please see any text on abstract algebra, such as [Ar] or [JKT]. Basically, a field F is algebraically closed if any polynomial $f(x) = a_0 + a_1x + \dots + a_nx^n$ having coefficients a_i in F also has all its roots in F .)

The polynomial

$$p_A(x) = \det(A - xI)$$

is called the **characteristic polynomial** of A . A matrix A for which λ is a root of $p_A(x) = 0$ of multiplicity m is called an eigenvalue of A of **multiplicity** m . This is the multiplicity referred to in part (b) above.

Remark 6.5.1. *If you know the fundamental theorem of algebra and the definition of determinant then you can prove (b) follows from (a). The proof of (a) itself is not too hard. Here's a sketch: $Av = \lambda v$ holds if and only if $(A - \lambda I)v = 0$, where here 0 denotes the 0 vector in the space of v . Since v is non-zero, by assumption, the matrix $A - \lambda I$ cannot be invertible. (And conversely, if $A - \lambda I$ is not one-to-one then $(A - \lambda I)v = 0$ for some non-zero v .) By Lemma 2.2.1, $\det(A - \lambda I) = 0$.*

An odd situation arises if $\lambda = 0$. Matrices that have an eigenvalue of 0 don't behave as 'nicely', in some respects, as matrices for which all the eigenvalues are non-zero. The number of non-zero eigenvalues (counted according to multiplicity) of a square diagonalizable matrix is its 'rank' (defined more generally below). The rank is a measurement of how 'bad' a matrix is, in a sense. (Of course, I do not mean to cast dispersions on the character of low ranking matrices!) The following result enables one to compute the rank of a matrix without having to compute the eigenvalues, which can be difficult if n is large.

Lemma 6.5.3. *Let A be any matrix as in the lemma above. Let r_1, \dots, r_n be the row vectors of A . Let A' be a matrix which is the same as A except that either (a) exactly one row, say the i th one, has been replaced by the sum $r_i + cr_j$, where c is any constant and r_j is any other row (so $i \neq j$) if A , or (b) the rows of A' are the same as A but with exactly two of them swapped. Then*

$$\text{rank}(A) = \text{rank}(A').$$

A matrix A' as in the above lemma is said to be obtained from A by an **elementary row operation**. An elementary column operation may be defined similarly. (The definition is left to the reader.) The analogous result also holds for elementary column operations (see any text on linear algebra, for example Beezer [Be]). The point to this lemma is that one can often choose the constant c so that A' has more zeros than A does. This is an instance of **Gaussian elimination**. Now we shall reformulate this into a little game.

The Gauss Elimination Game

- **Legal moves:** These actually apply to any $m \times n$ matrix A with $m < n$.
 1. $R_i \leftrightarrow R_j$: You can swap row i with row j .
 2. $cR_i \rightarrow R_i$ ($c \neq 0$): You can replace row i with row i multiplied by any non-zero constant c .
 3. $cR_i + R_j \rightarrow R_i$ ($c \neq 0$): You can replace row i with row i multiplied by any non-zero constant c plus row j , $j \neq i$.
- **Goal:** You *win* the game when you can achieve the requirements below. Your goal is to find a sequence of legal moves leading to a matrix B satisfying the following criteria:

1. All rows of B have leading a non-zero term equal to 1 (the position where this leading term in B occurs is called a *pivot position*).
2. The matrix B contains as many zeroes as possible.
3. All entries above and below a pivot position must be 0.
4. No pivot position is below and to the left of any other one. In particular, all entries below the diagonal of B are 0.

This matrix B is unique (a theorem which you can find in any text on elementary matrix theory, such as [Be]) and is called the ‘row-reduced echelon form’ of A . For short, we shall simply say B is **reduced**. The number of all-zero rows of B is called the **nullity** of A . The number of non-zero rows of B is called the **rank** of A . Therefore, ‘rank plus nullity equals m ’. (This last statement is sometimes called ‘the fundamental theory of linear algebra’.) When the rank is as large as possible, the matrix is said to be **full rank**.

The result above is due to Carl F. Gauss (1777-1855), perhaps the greatest mathematician of all time. Probably the best story about Gauss concerns the time when he first entered elementary school at the age of 7. His teacher, no doubt tired and wanting a break, gave his class the following problem: ‘Find the sum of all the integers from 1 to 100: $1 + 2 + \dots + 100 = ?$ ’ He probably expected to get a good bit of peace and quiet, but Gauss instantly saw that you could pair 1 and 100, 2 and 99, ..., 50 and 51, and get the answer $50 \cdot 101 = 5050$. For his 1799 Ph.D. thesis, he gave the first complete proof of the fundamental theorem of algebra. Gauss made major contributions to number theory, astronomy, and geodesy, among other fields, in his lifetime.

Gauss’s recipe to compute the rank

Input: The matrix A .

Output: The rank of A

Using the Gauss elimination game, create from A a matrix B that is reduced. The number of non-zero rows of B is the rank of A .

Gauss was quite a cook!

Example 6.5.4. Here is an illustration of Gauss elimination using SAGE:

SAGE

```
sage: MS = MatrixSpace(QQ, 3, 5)
sage: A1 = MS([[1,2,3,4,5],[1,1,-1,0,1],[6,7,8,9,10]])
sage: A1.echelon_form()
[ 1  0  0  0 -3/2  -3]
[ 0  1  0  0  2    4]
[ 0  0  1  1/2  0]
sage: A2 = MS([[1,2,3,4,5],[1,1,1,1,1],[6,7,8,9,10]])
sage: A2.echelon_form()
[ 1  0 -1 -2 -3]
[ 0  1  2  3  4]
[ 0  0  0  0  0]
```

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The first matrix is full rank (rank 3, which is the largest rank a 3×5 matrix can have). The second matrix is rank 2, which is the number of non-zero rows in its reduced row echelon form.

6.5.3 The rectangular case

We reformulate the problem.

Let $V = V_{M,N} = M_{M \times N}(\mathbb{Z}/2\mathbb{Z})$. This is a vector space over the field $\mathbb{Z}/2\mathbb{Z}$ with two elements. Let $W = W_{M,N}$ denote the subspace spanned by the vectors $t(E)$, in the notation of Example 6.4.2 above, where E runs over all $M \times N$ elementary matrices.

The following question was posed by D. Rusin in [Rus]:

Problem: What is the codimension of W in V ?

This question is important for the solution of the rectangular Merlin's Machine due to the following fact. Let $c_{M,N} = \dim_{\mathbb{Z}/2\mathbb{Z}}(V_{M,N}/W_{M,N})$ denote the codimension of W in V .

Proposition 6.5.1. *The $M \times N$ puzzle can be solved with any initial state if and only if $c_{M,N} = 0$.*

Example 6.5.5. (Rusin [Rus]) Several examples were computed by Rusin. The array of values of these numbers for $0 < m, n < 10$ is

0	1	0	0	1	0	0	1	0
1	0	2	0	1	0	2	0	1
0	2	0	0	3	0	0	2	0
0	0	0	4	0	0	0	0	4
1	1	3	0	2	0	4	1	1
0	0	0	0	0	0	0	6	0
0	2	0	0	4	0	0	2	0
1	0	2	0	1	6	2	0	1
0	1	0	4	1	0	0	1	8

In particular, not every possible initial state for the 3×5 Merlin's Machine puzzle can be solved, since $c_{3,5} \neq 0$.

Proposition 6.5.1 above is useful only if one can compute the $c_{M,N}$. For low values of M, N this can be done with the help of a computer. For larger values, such calculations become more difficult. The ideas of Goldwasser, Klostermeyer, et al., provide a more computationally effective solution. Theorem 6.5.2 above generalizes to the rectangular case as follows.

Theorem 6.5.3. (Goldwasser, Klostermeyer, and Trapp [GKT]) *The $M \times N$ puzzle can be solved with any initial state if and only if the greatest common divisor of $f_{M+1}(x)$ and $f_{N+1}(x+1)$ equals 1.*

6.5.4 Alien Tiles again

This is the same as the rectangular case discussed in the subsection above, except for two things:

- we must replace $\mathbb{Z}/2\mathbb{Z}$, representing on/off, by $\mathbb{Z}/4\mathbb{Z}$ (with addition mod 4), representing red, green, blue, and purple, respectively;
- we must change the ‘toggle matrix’ $t(E)$ in Example 6.4.2 to have a 1 for every matrix entry corresponding to a button in the same row or column as the button pressed (not just the *neighboring* buttons in the same row or column).

The states of this puzzle machine are $M_{M \times N}(\mathbb{Z}/4\mathbb{Z})$. $t'(E)$ is the **toggle matrix** for Alien Tiles,

$$t'(E_{i,j}) = \sum_{i',j'} E_{i',j'},$$

where the sum runs over all $1 \leq i' \leq M$, $1 \leq j' \leq N$, with either $i' = i$ or $j' = j$ or both. If s is the present state, then the next state, s' , is obtained by first choosing an elementary matrix E corresponding to the button you want to press and then computing

$$s' = s + t'(E),$$

where addition is coordinate-wise mod 4.

Let $V = V_{M,N} = M_{M \times N}(\mathbb{Z}/4\mathbb{Z})$. This is a ‘module’ over $\mathbb{Z}/4\mathbb{Z}$ (mathematicians use the word ‘module’ in place of ‘vector space’ when the scalars are taken from a ring such as $\mathbb{Z}/4\mathbb{Z}$). Let $W = W_{M,N}$ denote the submodule (a ‘submodule’ is analogous to a ‘vector subspace’) spanned by the vectors $t'(E)$, in the notation above, where E runs over all $M \times N$ elementary matrices.

The following problem naturally arises.

Ponderable 6.5.1. (*Hard*) *What is the codimension of W in V ?*

A general formula for this codimension, as a function of M and N , was unknown when the first edition of this book appeared. However, Ponderable 6.5.1 has recently been solved by Peter Maier and Werner Nickel in their paper [MN].

6.5.5 Orbix, revisited

Let us number the lights on the Orbix as follows: with the north pole as 1, the south pole as 12, let the numbering of the remaining buttons, as viewed from above, be as in Figure 6.2.

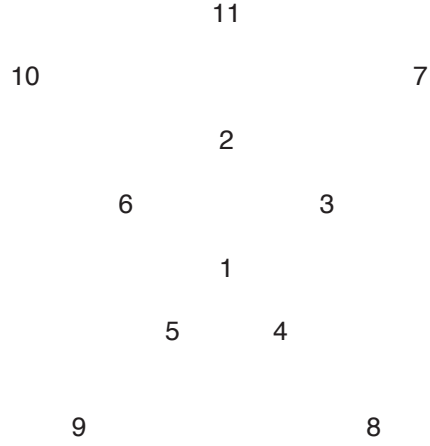


Figure 6.2: Orbix labels.

The labeling is detailed further in the diagram below.

Light number	neighbors	antipodes
1	2,3,4,5,6	1,12
2	1,3,6,7,11	2,9
3	1,2,4,7,8	3,10
4	1,3,5,8,9	4,11
5	1,4,6,9,10	5,7
6	1,2,5,10,11	6,8
7	2,3,8,11,12	
8	3,4,7,9,12	
9	4,5,8,10,12	
10	5,6,9,11,12	
11	2,6,7,10,12	
12	7,8,9,10,11	

The adjacency matrix of the graph in §6.2.1 is the 12×12 matrix $A = (a_{i,j})$, where $a_{i,j} = 1$ if vertex i and vertex j are neighbors (connected by a single

edge) and $a_{i,j} = 0$, otherwise. In this case,

$$A = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 \end{pmatrix}$$

For example, the ‘toggle vector’ for the first button is $(0, 1, 1, 1, 1, 1, 0, 0, 0, 0, 0, 0)$.

Remark 6.5.2.

- The matrix A has determinant $625 = 5^4$ and eigenvalues $5, 5^{1/2}, -5^{1/2}, 5^{1/2}, -5^{1/2}, 5^{1/2}, -5^{1/2}, -1, -1, -1, -1, -1$.
- The state vector $v_0 = (1, 1, \dots, 1)$ is an eigenvector with eigenvalue 5 . This state vector has the property that if you punch all the buttons on an unlit Orbix associated to its non-zero entries of v_0 then you recover v_0 .
- The state vector $v_1 = (0, 0, 1, 0, 1, 0, 1, 0, 0, 1, 0, 0)$ is an eigenvector with eigenvalue -1 . This state vector has the property that if you punch all the buttons on an unlit Orbix associated to its non-zero entries of v_1 then you recover v_1 .

The only other vectors with this property are the eigenvectors

$$\begin{aligned} v_2 &= (0, 1, 0, 0, 1, 0, 1, 0, 1, 0, 0, 0), & v_3 &= (0, 0, 0, 1, 1, 0, 1, 0, 0, 0, 1, 0), \\ v_4 &= (0, 0, 0, 0, 1, 1, 1, 1, 0, 0, 0, 0), & v_5 &= (1, 0, 0, 0, 1, 0, 1, 0, 0, 0, 0, 1), \end{aligned}$$

and their linear combinations.

- A move consists mathematically of choosing a vertex i then adding (modulo 2) row i of the adjacency matrix A to the state vector.
- If you change the rules of the Orbix and decide to toggle the button pressed as well as its neighbors (as in the Lights Out game), then the 12×12 matrix

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of toggle vectors is

$$\begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}$$

This matrix is singular. Its column span (and row span) is only 6-dimensional over $\mathbb{Z}/2\mathbb{Z}$, and there are only 64 (out of $2^{12} = 4096$ possible) solvable states.

Here is an example using **SAGE** to compute the rank of the above matrix.

```

SAGE
sage: MS = MatrixSpace(GF(2), 12, 12)
sage: lights = [[1, 1, 1, 1, 1, 1, 0, 0, 0, 0, 0, 0],\
.....: [1, 1, 1, 0, 0, 1, 1, 0, 0, 0, 1, 0],\
.....: [1, 1, 1, 1, 0, 0, 1, 1, 0, 0, 0, 0],\
.....: [1, 0, 1, 1, 1, 0, 0, 1, 1, 0, 0, 0],\
.....: [1, 0, 0, 1, 1, 1, 0, 0, 1, 1, 0, 0],\
.....: [1, 1, 0, 0, 1, 1, 0, 0, 0, 1, 1, 0],\
.....: [0, 1, 1, 0, 0, 0, 1, 1, 0, 0, 1, 1],\
.....: [0, 0, 1, 1, 0, 0, 1, 1, 1, 0, 0, 1],\
.....: [0, 0, 0, 1, 1, 0, 0, 1, 1, 1, 0, 1],\
.....: [0, 0, 0, 0, 1, 1, 0, 0, 1, 1, 1, 1],\
.....: [0, 1, 0, 0, 0, 1, 1, 0, 0, 1, 1, 1],\
.....: [0, 0, 0, 0, 0, 0, 1, 1, 1, 1, 1, 1]]
sage: A = MS(lights)
sage: A.echelon_form()
[1 0 0 0 0 1 0 1 0 0 0 1]
[0 1 0 0 0 1 0 1 1 0 0 0]
[0 0 1 0 0 1 0 1 0 1 0 0]
[0 0 0 1 0 1 0 1 0 0 1 0]
[0 0 0 0 1 1 0 0 1 1 1 1]
[0 0 0 0 0 0 1 1 1 1 1 1]
[0 0 0 0 0 0 0 0 0 0 0 0]
[0 0 0 0 0 0 0 0 0 0 0 0]
[0 0 0 0 0 0 0 0 0 0 0 0]
[0 0 0 0 0 0 0 0 0 0 0 0]
[0 0 0 0 0 0 0 0 0 0 0 0]
sage: A.rank()
6

```


$$T = \begin{pmatrix} 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \end{pmatrix}$$

```
sage: MS = MatrixSpace(GF(2),16,16)
sage: A = MS([
[1,1,0,0,1,1,0,0,0,0,0,0,0,1,0,0,0],
[1,1,1,0,0,1,0,0,0,0,0,0,0,0,1,0,0],
[0,1,1,1,0,0,1,0,0,0,0,0,0,0,0,1,0],
[1,0,1,1,0,0,0,0,1,0,0,0,0,0,0,0,1],
[1,0,0,0,1,1,0,1,1,0,0,0,0,0,0,0,0],
[0,1,0,0,1,1,1,0,0,1,0,0,0,0,0,0,0],
[0,0,1,0,0,1,1,1,0,0,1,0,0,0,0,0,0],
[0,0,0,1,0,0,1,1,1,0,0,1,0,0,0,0,0],
[0,0,0,1,1,0,1,1,0,0,0,1,0,0,0,0,0],
[0,0,0,0,1,0,0,0,1,1,0,1,1,0,0,0,0],
[0,0,0,0,0,1,0,0,1,1,1,0,0,1,0,0,0],
[0,0,0,0,0,0,1,0,0,1,1,1,0,0,1,0,0],
[0,0,0,0,0,0,0,1,0,0,1,1,1,0,0,1,0],
```

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```
[0,0,0,0,0,0,0,1,1,0,1,1,0,0,0,1],
[1,0,0,0,0,0,0,0,1,0,0,0,1,1,0,1],
[0,1,0,0,0,0,0,0,0,1,0,0,1,1,0,],
[0,0,1,0,0,0,0,0,0,0,1,0,0,1,1,1],
[0,0,0,1,0,0,0,0,0,0,0,1,1,0,1,1]]
sage: p = A.charpoly()
sage: p.factor()
(x + 1)^16
sage: A.eigenspaces()

[
(1, Vector space of degree 16 and dimension 8 over Finite Field of size 2
User basis matrix:
[1 0 0 0 0 0 0 0 1 0 0 0 0 1 0 1]
[0 1 0 0 0 0 0 0 0 1 0 0 1 0 1 0]
[0 0 1 0 0 0 0 0 0 0 1 0 0 1 0 1]
[0 0 0 1 0 0 0 0 0 0 0 1 1 0 1 0]
[0 0 0 0 1 0 0 0 0 1 0 1 1 0 0 0]
[0 0 0 0 0 1 0 0 1 0 1 0 0 1 0 0]
[0 0 0 0 0 0 1 0 0 1 0 1 0 0 1 0]
[0 0 0 0 0 0 0 1 1 0 1 0 0 0 0 1])
]
sage: A.rank()
16
```

Chapter 7

‘God’s algorithm’ and graphs

Unfortunately what is little recognized is that the most worthwhile scientific books are those in which the author clearly indicates what he does not know; for an author most hurts his readers by concealing difficulties.

Evariste Galois

In this chapter we introduce a graphical interpretation of a permutation group, the Cayley graph. This is then interpreted in the special case of a group arising from a permutation puzzle. Along the lines of Galois’ quote above, we shall also discuss the state of our ignorance regarding ‘God’s algorithm’. See also chapter 16.

7.1 In the beginning...

To begin, what’s a graph? A **graph** is a pair of sets (V, E) , where

- V is a set of singleton elements called **vertices**,
- E is a subset of the set of all *unordered* pairs $\{\{v_1, v_2\} \mid v_1, v_2 \in V, v_1 \neq v_2\}$. The elements of E are called **edges**.

Such a graph (one without ‘loops’) is also called a **simple graph** in the literature. A graph is drawn by simply connecting points representing vertices together by a line segment if they belong to the same edge. A **finite graph** is a graph (V, E) with V a finite set.

A **digraph**, or **directed graph**, is a pair of sets (V, E) , where

- V is a countable set of vertices,
- E is a subset of *ordered* pairs $\{(v_1, v_2) \mid v_1, v_2 \in V, v_1 \neq v_2\}$ called **edges**

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A digraph is drawn by simply connecting points representing vertices together by an arrow if they belong to the same edge (v_1, v_2) , the arrow originating at v_1 and arrowhead pointing to v_2 .

If $e = \{v_1, v_2\}$ belongs to E then we say that e is an **edge** from v_1 to v_2 (or from v_2 to v_1) and we say v_1 and v_2 are **neighbors** or **adjacent** in the graph. If v and w are vertices, a **path** from v to w is a finite sequence of edges beginning at v and ending at w :

$$e_0 = \{v, v_1\}, e_1 = \{v_1, v_2\}, \dots, e_n = \{v_n, w\}.$$

If there is a path from v to w then we say v is **connected** to w . We say that a graph (V, E) is **connected** if each pair of vertices is connected. The number of edges emanating from a vertex v is called the **degree** (or **valence**) of v , denoted $\text{degree}(v)$.

Example 7.1.1. *If*

$$V = \{a, b, c\}, \quad E = \{\{a, b\}, \{a, c\}, \{b, c\}\},$$

then we may visualize the graph $\Gamma = (V, E)$ as in Figure 7.1.

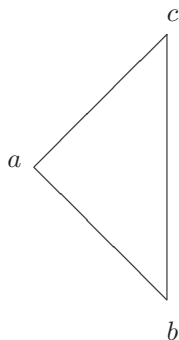


Figure 7.1: Each vertex has valence 2.

Here are illustrations of how to use **SAGE** to construct Γ . The first **SAGE** plot has the (default) labels for the vertices 0, 1, 2. The second has vertices labeled a, b, c .

SAGE

```
sage: G = graphs.CycleGraph(3)
sage: show(G.plot())
sage: G = Graph({"a":["b","c"], "b":["a","c"], "c":["a","b"]})
sage: show(G.plot())
```

Each of these plots illustrates the same graph as in Figure 7.1, but possibly relabeled.

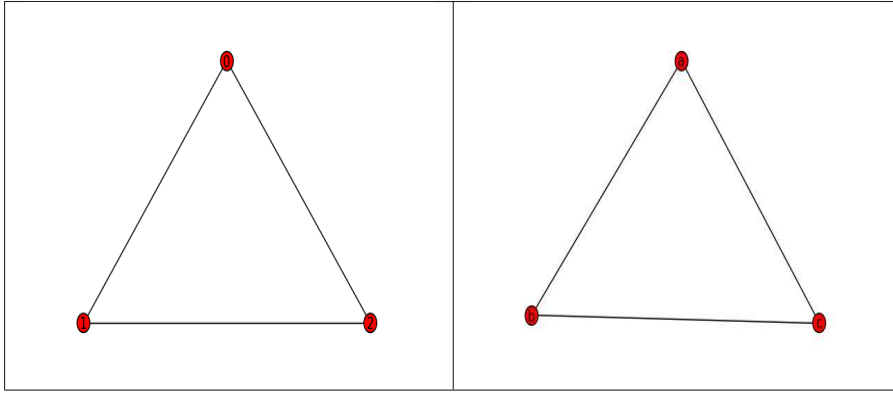


Figure 7.2: Two plots of the graph Γ .

Definition 7.1.1. If v and w are vertices connected to each other in a graph (V, E) , then we define the **distance from v to w** denoted $d(v, w)$, by

$$d(v, w) = \min_{v, w \in V \text{ connected}} \#\{\text{edges in a path from } v \text{ to } w\}$$

By convention, if v and w are not connected, then we set $d(v, w) = \infty$. The **diameter** of a graph is the largest possible distance:

$$\text{diam}((V, E)) = \max_{v, w \in V} d(v, w).$$

In the above example, the diameter is 1.

7.2 Cayley graphs

Let G be a finite group with prescribed generators g_i :

7.2. CAYLEY GRAPHS

$$G = \langle g_1, g_2, \dots, g_n \rangle.$$

Assume that the generating set $X = \{g_1, g_2, \dots, g_n\}$ does not contain the identity and it is closed under taking inverses (i.e., if $x \in X$ then $x^{-1} \in X$). The **Cayley graph** of G with respect to $X = \{g_1, g_2, \dots, g_n\}$ is the graph (V, E) whose vertices V are the elements of G and whose edges are determined by the following condition: if x and y belong to $V = G$ then there is an edge from x to y (or from y to x) if and only if $y = g_i \cdot x$ or $x = g_i \cdot y$, for some $i = 1, 2, \dots, n$.

Cayley graphs are named for Arthur Cayley, whom we met already in §2.2.1. Cayley, though starting out as a lawyer, eventually published over 900 papers and notes covering nearly every aspect of mathematics.

Example 7.2.1. *Here is an example of how to use SAGE and GAP's GRAPE package to construct a graph for the symmetric group S_3 :*

SAGE

```
sage: gens = [(1,2),(2,3)]
sage: gap.eval("gens := "+str(gens))
'[ (1,2), (2,3) ]'
sage: gap.eval("C := CayleyGraph(Group(gens),gens)")
'rec( isGraph := true, order := 6,
      group := Group([ (1,3)(2,4)(5,6), (1,2)(3,5)(4,6) ]),
      schreierVector := [ -1, 2, 1, 1, 2, 1 ],
      adjacencies := [ [ 2, 3 ] ],
      representatives := [ 1 ],
      names := [ (), (2,3), (1,2), (1,2,3), (1,3,2), (1,3) ],
      isSimple := true )'
sage: E = eval(gap.eval("UndirectedEdges(C)"))
sage: V = eval(gap.eval("Elements(Vertices(C))"))
sage: L = [sum([y[i] for y in [x for x in E if v in x]\
                  if y[i]!=v] for i in range(len(gens))],[i])\
           for v in V]
sage: d = dict(zip(V,L))
sage: G = Graph(d)
sage: show(G.plot())
```

(The line defining L was split into three for typographical reasons.) Here is another method which does not use the GRAPE package:

SAGE

```
sage: gens = [(1,2),(2,3)]
sage: G = PermutationGroup(gens)
sage: C = G.cayley_graph()
sage: show(C.plot())
```

This is discussed further in Example 7.2.2 below, which also displays the graph.

Ponderable 7.2.1. Show that the Cayley graph of a finite group is connected.

Lemma 7.2.1. Let $\Gamma_G = (V, E)$ denote the Cayley graph associated with the finite group $G = \langle g_1, g_2, \dots, g_n \rangle$. Let $N = |\{g_1, g_1^{-1}, g_2, g_2^{-1}, \dots, g_n, g_n^{-1}\}|$. Then, for all $v \in V$, $\text{degree}(v) = N$.

Proof: Assume not. Then there is a $v \in V = G$ with either

- (i) $\text{degree}(v) < N$, or
- (ii) $\text{degree}(v) > N$.

First, we note that, for each $h \in \{g_1, g_1^{-1}, g_2, g_2^{-1}, \dots, g_n, g_n^{-1}\}$, the set $\{v, h \cdot v\}$ is an edge of Γ_G . This follows from the definition of the Cayley graph.

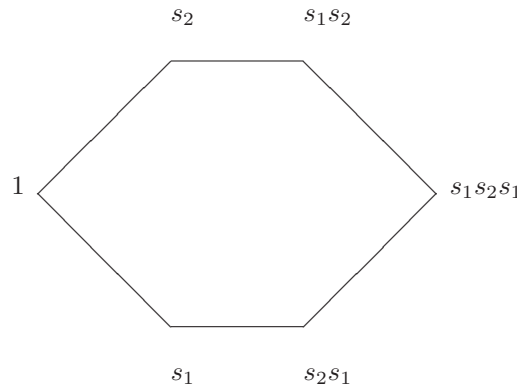
If $r = \text{degree}(v) > N$ then, by definition of the Cayley graph, there are distinct $v_1, \dots, v_r \in V$ with $v = h_i \cdot v_i$, for all $1 \leq i \leq r$, where the h_1, \dots, h_r are distinct elements of $\{g_1, g_1^{-1}, g_2, g_2^{-1}, \dots, g_n, g_n^{-1}\}$. This contradicts the definition of N .

If $r = \text{degree}(v) < N$ then, by definition of the Cayley graph, there are distinct h_i, h_j in $\{g_1, g_1^{-1}, g_2, g_2^{-1}, \dots, g_n, g_n^{-1}\}$ such that $h_i \cdot v = h_j \cdot v$. Since G is a group and $V = G$ (as sets), we may cancel the v from both sides of the equation $h_i \cdot v = h_j \cdot v$, contradicting the assumption that h_i is distinct from h_j . \square

Example 7.2.2. Let

$$G = \langle s_1, s_2 \rangle = S_3,$$

where $s_1 = (1, 2)$, and $s_2 = (2, 3)$. Then the Cayley graph of G with respect to $X = \{s_1, s_2\}$ may be visualized as below (note that $s_1 = s_1^{-1}$ and $s_2 = s_2^{-1}$, so X is closed under taking inverses).



7.3. GOD'S ALGORITHM

Ponderable 7.2.2. *Let*

$$G = \langle s_1, s_2 \rangle = S_3,$$

where $s_1 = (1, 2)$, and $s_2 = (2, 3)$. Find the Cayley digraph of G with respect to $X = \{s_1, s_2\}$.

Ponderable 7.2.3. *Construct the Cayley graph of C_4 , the cyclic group, with respect to the generator $s = (1, 2, 3, 4)$.*

Ponderable 7.2.4. *Construct the Cayley graph of S_4 , the symmetric group on four letters, with respect to the generators $s_1 = (1, 2)$, $s_2 = (2, 3)$ and $s_3 = (3, 4)$.*

Ponderable 7.2.5. *Construct the Cayley digraph of S_3 with respect to the generators $f = (1, 3)$, $r = (1, 2, 3)$. (Show, in particular, that f, r do indeed generate S_3 .)*

Example 7.2.3. *Let*

$$G = \langle R, L, U, D, F, B \rangle \subset S_{54}$$

be the group of the 3×3 Rubik's Cube. Each position of the cube corresponds to an element of the group G (i.e., the move you had to make to get to that position). In other words, each position of the cube corresponds to a vertex of either

- the Cayley graph associated with the **quarter-turn metric** with generating set $\{R, L, U, D, F, B, R^{-1}, L^{-1}, U^{-1}, D^{-1}, F^{-1}, B^{-1}\}$, or
- the Cayley graph associated with the **face-turn metric** with generating set $\{R, R^2, L, L^2, U, U^2, D, D^2, F, F^2, B, B^2, R^{-1}, L^{-1}, U^{-1}, D^{-1}, F^{-1}, B^{-1}\}$.

Each vertex of the first (quarter-turn metric) graph has valence 12, whereas each vertex of the latter (face-turn metric) graph has valence 16.

Ponderable 7.2.6. *Check this.*

Moreover, a solution of the Rubik's Cube is simply a path in the Cayley graph from the vertex associated with the present position of the cube to the vertex associated with the identity element. The number of moves in the shortest possible solution is simply the distance from the vertex associated with the present position of the cube to the vertex associated to the identity element. The diameter of the Cayley graph of G is the number of moves in the best possible solution in the worst possible case. This diameter can be measured either in the quarter-turn metric or the face-turn metric.

7.3 God's algorithm

...O, cursed spite,
that ever I was to set it right!
Shakespeare, Hamlet, Act 1, scene 5

Ponderable 7.3.1. *(Very hard) Let G be the group of a permutation puzzle (with a fixed generating set). Find the diameter of the Cayley graph of G .*

Sometimes this is called ‘God’s algorithm’, though here that term is reserved for a harder version of the problem, stated below. This diameter problem is unsolved for most puzzles (including the 3×3 Rubik’s Cube) and appears to be very difficult computationally in general (see [GJ] for a discussion of similar problems and their complexity). The cases where it is known include (with no attempt at completeness) the following puzzles (the reader is referred to [Lo], [CFS] for more details).

Puzzle	Diameter
Pyraminx	11 (not including tip moves)
Skewb	11
$2 \times 2 \times 2$ Rubik’s Cube	14 (quarter turns)

These diameters were found with the help of a computer.

Many people have worked on God’s algorithm for the Rubik’s Cube. For a while it was guessed that the superflip position is the position which is as far from ‘start’ (the solved position) as possible. The superflip position is 24 quarter turns from start in the quarter turn metric (as was already mentioned, this fact is due to Jerry Bryan and Michael Reid [CL]) and 20 turns from start in the face-turn metric (this is due to Michael Reid and Dik Winter). However, Reid discovered a longer move in the quarter-turn metric—26 quarter-turns away from start (the superflip composed with the four spot, see (5.2)). No longer move has been found in the face-turn metric. At the moment, one might guess that the diameter of the Rubik’s Cube graph is 26 in the quarter-turn metric and 20 moves in the face-turn metric. At the time of this writing (spring 2008), the best-known upper bound is as follows:

- Tom Rokicki proved in 2008 that 25 face turns is enough [Rok]. (Gene Cooperman and Daniel Kunkle proved in 2007 that 26 face turns is enough [CK], but believe that their technique can be pushed to the 25 mark as well.)
- In 2006, Silviu Radu proved that every position can be solved in at most 35 quarter turns [Ra].

For the latest progress, see the Wikipedia article “Optimal solutions for Rubik’s Cube” [Wi]. For a description of the algorithms involved, see §15.2.3 below and Herbert Kociemba’s excellent site [Koc].

Ponderable 7.3.2. *Let G be the group of a permutation puzzle (with a fixed generating set) and let v be a vertex in the Cayley graph of G . Find an effective, practical algorithm for determining a path from v to the vertex v_0 associated to the identity having a length equal to the distance from v to v_0 .*

This problem is even harder! This algorithm is called **God’s algorithm**. It really refers to an optimal method of solution for the permutation puzzle.

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Ponderable 7.3.3. Find the Cayley graph of the sliced squared group

$$G = \langle M_R^2, M_F^2, M_D^2 \rangle,$$

where M_R is the middle-slice move which turns the middle slice parallel to the right face clockwise 90 degrees (with respect to the right face). Find the diameter of this graph.

Let Γ be a graph. A **Hamiltonian circuit** on Γ is a sequence of edges forming a path in Γ which passes through each vertex exactly once. (If you think of the vertices as cities and the edges as roads, then a Hamiltonian circuit is a tour visiting each city exactly once.)

The following unsolved problem was first mentioned in this context by A. Schwenk (in an unpublished paper written while he was at the U.S. Naval Academy).

Ponderable 7.3.4. (Very hard) Let G be the group of the 3×3 Rubik's Cube puzzle. Does the Cayley graph of G have a Hamiltonian circuit? In other words, can we (in principle) 'visit' each possible position of the Rubik's Cube exactly once, by making one move at a time using only the basic generators R, L, U, D, F, B ?

Remark 7.3.1. This is a special case of a more general unsolved problem: For an arbitrary permutation group with a given generating set, is its Cayley graph Hamiltonian? See [CSW] and [CG] for results in this area.

An example of a group (with generators) where its Cayley graph is known to have a Hamiltonian circuit is the symmetric group (with generators the set of all transpositions).

Example 7.3.1. Let G be the group S_n with generators given to be the set of all transpositions:

$$G = S_n, \quad X = \{(i, j) \mid 1 \leq i < j \leq n\}.$$

(There are many more transpositions than necessary to generate S_n since the subset of transpositions of the form $(i, i+1)$, $1 \leq i \leq n-1$, suffice to generate S_n [R].) The algorithm of Steinhaus (see §3.4) shows that there is a Hamiltonian circuit in the Cayley graph of S_n with respect to X .

The reader interested in more examples is referred to the excellent survey article by S. Curran and J. Gallian [CG].

William Hamilton, whom we met already in §5.1, may have originated the problem of finding Hamiltonian paths (hence the name) by patenting a game called the 'Icosian game' or the 'Hamilton game'. The idea is to find a Hamiltonian path around the vertices of the icosahedron. A picture of the original game can be found at the MacTutor History of Mathematics archive [MT]. Hamilton was a prodigy talented in many subjects, who by the age of five, had already learned Latin, Greek, and Hebrew. He eventually turned to mathematics and physics, was knighted in 1835, and was the first foreign member elected to the National Academy of Sciences of the USA.

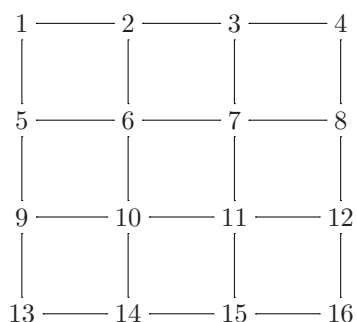
7.4 The graph of the 15 Puzzle

This section discusses the 15 Puzzle from the graph-theoretical point of view following [W].

The 15 Puzzle was introduced in §4.1 above. The object of the puzzle was to order the pieces from 1 to 15 from left to right, top to bottom, as shown in the **solved position** given in §4.1.

To solve a mixed-up puzzle, one would slide the squares around in the puzzle. In order to do this you must slide a numbered square into the place of the space. We could represent this mathematically by saying that this is a transposition of that numbered square and the blank.

If we label each space in the puzzle as a vertex, and label the vertices numerically, then the resulting graph is represented by the figure below.



We will let 16 denote the blank and call this graph Γ . The only legal moves of the puzzle are transpositions of the 16th vertex and a vertex that is adjacent to it. Therefore, any permutation of the vertices produces a labeling on Γ .

7.4.1 General definitions

Now let Γ be a simple graph with the vertex set $V(\Gamma)$ of cardinality N . (In the above example $N = 16$.) By a **labeling** I mean the placement of the numbers 1 through N on distinct vertices of Γ , where N denotes the blank. In other words, a labeling on Γ is a bijective mapping $f : V(\Gamma) \rightarrow \{1, 2, \dots, N\}$. Two labelings f, g on Γ are **adjacent** if and only if g is a result of a single transposition on f of vertex N with a vertex adjacent to N on f . In other words, f and g are adjacent if they differ by one legal move of the puzzle. From Γ , we make a new graph $puz(\Gamma)$ as follows [W]: the vertex set $V(puz(\Gamma))$ contains all labelings on Γ , and two vertices in $puz(\Gamma)$ are joined by an edge if the associated labelings are adjacent.

For example, the labelings in

1	2	3	4
5	6	7	8
9	10	11	
13	14	15	12

7.4. THE GRAPH OF THE 15 PUZZLE

and

1	2	3	4
5	6	7	8
9	10		11
13	14	15	12

are adjacent.

We can consider a **path** p to be a sequence of moves on the graph $puz(\Gamma)$, $p = (x_0, x_1, x_2, \dots, x_n)$, where the x_i 's are vertices of Γ , and (if $n \geq 1$) x_i and x_{i-1} are adjacent in Γ for $1 \leq i \leq n$. Such a path p is said to be from x_0 (its **initial vertex**) to x_n (its **terminal vertex**). The path p is **simple** if $x_0, x_1, x_2, \dots, x_n$ are distinct. If $x_0 = x_n$, then p is called a **closed path based at** x_0 . Let x_0 be a fixed vertex of Γ . If p, p' are closed paths based at x_0 , then the **composition** pp' is the path obtained by juxtaposing the path for p with that for p' . The **inverse** of a path p , denoted p^{-1} , is the path which traverses p in the opposite direction. The **identity path** is the path from x_0 to itself which contains no edges of Γ . The set of closed paths based at x_0 forms a group (under composition of paths) called the **homotopy group** of Γ based at x_0 , denoted $\Gamma(x_0)$.

Now suppose we paint the blocks of the 15 Puzzle in a checkerboard pattern:

	×		×
×		×	
	×		×
×		×	

In this arrangement the blank would start on a white square. If we were to move the blank up, then it is now on a black square. That is one transposition; therefore, the movement is odd. If we then move the blank to the left, the blank would be on a white square. This is a total of two transpositions; therefore, the movement is even. After three transpositions the blank would be on a black square; therefore, it would be an odd permutation. Therefore, if the blank ends on a white square, an even permutation has occurred. If the blank ends on a black square, an odd permutation has occurred.

A **legal position** of the 15 Puzzle is any sequence of legal transpositions starting from the solved position such that the blank ends up in the bottom right-hand corner. Each such position corresponds to a permutation of the 15 numbered vertices and hence to an element of the symmetric group S_{15} . The set of all such permutations (arising as a sequence of transpositions) in S_{15} forms a group called the **(homotopy) group of the 15 Puzzle**.

Note that the group of the 15 Puzzle is isomorphic to the homotopy group of the 15 Puzzle graph based at the 'blank vertex'.

If we assign the number 16 to the blank, then we can see that we can arrange the pieces of the puzzle in $16!$ different ways. However, if we take only legal positions of the 15 Puzzle, then we are fixing one of the pieces. As a result

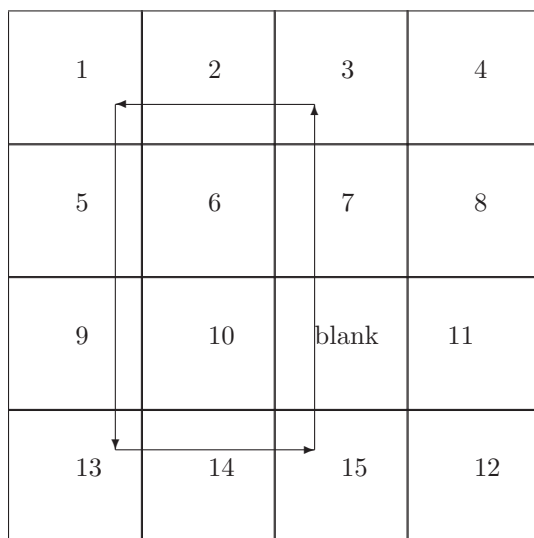
the number of ways to permute the rest of the pieces, with the blank on the white square at the bottom right-hand corner, is at most $15!$. All such permutations have to be even, by the checkerboard analysis above. The number of even permutations of 15 elements (there are an equal number of even and odd permutations) is $15!/2$. From this we see that the 15 Puzzle has at most $15!/2$ possible legal positions. In fact, it has exactly this number:

Theorem 7.4.1. *The positions with the empty space at the bottom right that can be reached from the start position of the 15 Puzzle by shifting tiles are in a bijective correspondence with the $15!/2 = 1,307,674,368,000$ even permutations of the numbers from 1 to 15.*

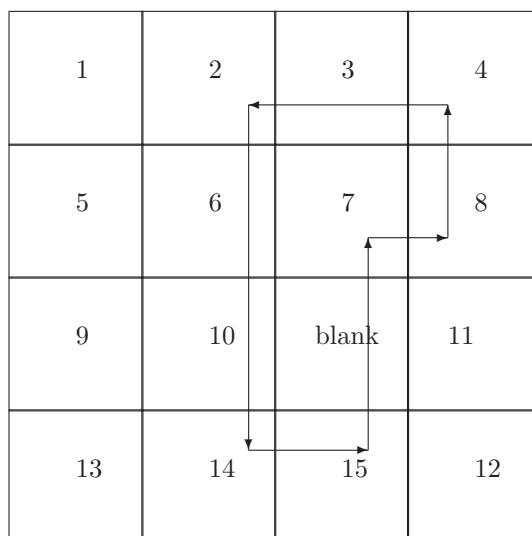
Remark 7.4.1. *The 14-15 Puzzle cannot be solved, because it is an odd permutation. It only has one transposition, 14 interchanged with 15.*

Proof: If we label the empty space 16, then every possible position of the puzzle may be regarded as an element of the symmetric group S_{16} and an element of $\text{puz}(\Gamma)$. There are $16!$ elements in S_{16} , and $16!$ vertices of $\text{puz}(\Gamma)$. With the argument earlier we can show that all legal positions of the puzzle are obtained by an even number of transpositions. Therefore, all legal moves are even permutations of the puzzle.

Now we must show that there is a certain 3-cycle in the group of the 15 Puzzle. For example, if we shift the three pieces surrounding the empty space around in a circle following the order of moves south-east-north-west, then the three-cycle $(11, 12, 15)$ is produced. It can be shown that if you fix the 11 and 12 pieces, then any other piece can take the place of the 15 by following one of the cycles in the figures below.



7.4. THE GRAPH OF THE 15 PUZZLE



By Lemma 9.4.4, such 3-cycles generate A_{15} . This proves the theorem.

□

Remark 7.4.2. *Another proof is possible. Alternatively, with some work we can show that any number can replace the 11 in the three-cycle, and we can show that any other number can replace the 12. From this we can conclude that any three-cycle can be formed. Since every even permutation is a combination of three-cycles (by Proposition 9.4.1), every even permutation of the 15 Puzzle can be reached.*

With this information, we can make a generalization to rectangular puzzles of size $m \times n$ with $m > 1$ and $n > 1$.

Theorem 7.4.2. *(Wilson [W]) The homotopy group of an $m \times n$ rectangular puzzle is the alternating group A_{mn-1} .*

The proof of this is similar to the proof for the 4×4 puzzle, if $m > 3$ and $n > 3$. (The special cases when $1 < m < 4$ or $1 < n < 4$ must be treated separately.) The size of the alternating group is given by $(mn - 1)!/2$.

7.4.2 Remarks on applications

Cayley graphs have been used by computer scientists to model interconnection networks for parallel processors (see [CFS], [CG] for some references).

The problem of finding an efficient algorithm for the shortest solution to the $m \times n$ puzzle is difficult. It amounts to finding the shortest path between two points in a graph which is, in general, a difficult problem computationally [GJ].

Chapter 8

Symmetry and the Platonic solids

Plato said God geometrizes continually.

Plutarch, Convivialium disputationum, liber 8,2

We have encountered regular polygons, such as the hexagon, in §5.3 on the dihedral group. A regular polyhedron is the 3-dimensional analog: a solid figure whose 2-dimensional faces are flat congruent regular polygons. Many puzzles are made by slicing up a regular polyhedron, so it is natural to study them and their symmetries. Such objects also occur in chemistry when the structure of crystals is investigated.

8.1 Descriptions

There are exactly 5 different types of regular polyhedra. We refer to Coxeter [Cox] for a proof.

The **Platonic solids** are the 5 regular polyhedrons:

Polyhedron	# Faces	# Vertices	# Edges	Group	p, q
Tetrahedron	4	4	6	T	3,3
Hexahedron	6	8	12	O	4,3
Octahedron	8	6	12	O	3,4
Dodecahedron	12	20	30	I	5,3
Icosahedron	20	12	30	I	3,5

Here

- p , called the **face degree**, denotes the number of edges bounding each face,
- q , called the **vertex degree**, denotes the number of faces meeting each vertex.

8.1. DESCRIPTIONS

These solids are named after the great Greek philosopher Plato (427 BC-347 BC), who established Plato's Academy, which flourished for 900 years until 529 AD.

A vertex of one of these solids is therefore specified by the q -tuple of faces meeting that vertex. We saw several examples of this already when we specified notation for the movements of the associated Rubik's Cube-like puzzles in chapter 4.

These solids may be drawn in rectangular coordinates using

Polyhedron	Coordinates
Tetrahedron	$(1,1,1), (1,-1,-1), (-1,-1,1), (-1,1,-1)$
Hexahedron	$(1,1,1), (1,1,-1), (1,-1,1), (-1,1,1),$ $(1,-1,-1), (-1,1,-1), (-1,-1,1), (-1,-1,-1)$
Octahedron	$(1,0,0), (0,0,1), (0,1,0),$ $(-1,0,0), (0,-1,0), (0,0,-1)$
Dodecahedron	$(0, \pm\phi^{-1}, \pm\phi), (\pm\phi^{-1}, \pm\phi, 0),$ $(\pm\phi, 0, \pm\phi^{-1}), (\pm 1, \pm 1, \pm 1),$
Icosahedron	$(1,0,\phi), (1,0,-\phi), (-1,0,\phi), (-1,0,-\phi),$ $(0,\phi,1), (0,\phi,-1), (0,-\phi,1), (0,-\phi,-1),$ $(\phi,1,0), (\phi,-1,0), (-\phi,1,0), (-\phi,-1,0)$

where ϕ denotes the golden ratio.

If P_1, P_2, P_3 are three vertices of an icosahedron which form a triangular face, then $(P_1 + P_2 + P_3)/3$ forms a vertex of the dual dodecahedron and every vertex of the dual dodecahedron arises in this way.

The three 'Platonic groups' (the group of 'symmetries' of these figures) will be described below. Their names:

- T = symmetry group of the tetrahedron = **tetrahedral group**,
- O = symmetry group of the octahedron (or cube) = **octahedral group**,
- I = symmetry group of the icosahedron (or dodecahedron) = **icosahedral group**.

8.2 Background on symmetries in 3-space

This subsection presents, with some proofs, background on isometries in 3 dimensions necessary for understanding the symmetry groups of the Platonic solids. Motivated by studies of the structure of crystals, Camille Jordan was the first to consider the classification of symmetries in 3-space. Jordan (1838-1922) was one of the pioneers of group theory and used groups to better understand crystals.

We fix once and for all the ‘right-hand-rule’ orientation in 3-space. We call a distance-preserving transformation in 3-space which fixes the origin a **symmetry of 3-space**. We say that such a symmetry is **orientation preserving** if it preserves the right-hand rule orientation.

Example 8.2.1. Let $s : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ denote the function which takes each vector v belonging to \mathbb{R}^3 and returns its reflection $s(v)$ about the yz -plane. This is not orientation-preserving since it reverses the direction of a counterclockwise moving circular path in the yz -plane. In terms of rectangular coordinates, $s(x, y, z) = (-x, y, z)$.

Let

$$\mathbb{R}^3 = \{(x, y, z) \mid x, y, z \text{ real numbers}\}$$

denote 3-space. We also write this, when convenient, as column vectors

$$\mathbb{R}^3 = \left\{ \begin{pmatrix} x \\ y \\ z \end{pmatrix} \mid x, y, z \text{ real} \right\}.$$

The **distance function** on \mathbb{R}^3 is the function

$$d(\vec{v}_1, \vec{v}_2) = \sqrt{(x_1 - x_2)^2 + (y_1 - y_2)^2 + (z_1 - z_2)^2},$$

where $\vec{v}_1 = (x_1, y_1, z_1)$, $\vec{v}_2 = (x_2, y_2, z_2)$. This may be expressed in terms of the **inner product** $\vec{v}_1 \cdot \vec{v}_2 = x_1x_2 + y_1y_2 + z_1z_2$ as $d(\vec{v}_1, \vec{v}_2) = \sqrt{(\vec{v}_1 - \vec{v}_2) \cdot (\vec{v}_1 - \vec{v}_2)}$. Conversely, the **polarization identity**

$$\vec{v}_1 \cdot \vec{v}_2 = \frac{1}{4}(\|\vec{v}_1 + \vec{v}_2\|^2 - \|\vec{v}_1 - \vec{v}_2\|^2)$$

allows one to recover the value of the inner product from the knowledge of the values of the distance function.

We call a function $f : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ an **isometry** if it satisfies

$$d(f(\vec{v}_1), f(\vec{v}_2)) = d(\vec{v}_1, \vec{v}_2)$$

for all \vec{v}_1 and \vec{v}_2 belonging to \mathbb{R}^3 .

We want to understand isometries a little better since they will preserve distances (and, in particular, preserve the shapes of solids) and therefore provide us with the kinds of symmetries of 3-space we want to consider. We can construct isometries using certain types of 3×3 matrices.

8.2. BACKGROUND ON SYMMETRIES IN 3-SPACE

Lemma 8.2.1. *If A is a 3×3 matrix then the function $A : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ is an isometry if and only if $A^t \cdot A = I_3$, where A^t denotes the transpose matrix (obtained by flipping the entries of A about the diagonal).*

Remark 8.2.1. *In particular, if A is an isometry then $\det(A)^2 = \det(A^t) \det(A) = \det(A^t \cdot A) = \det(I_3) = 1$.*

Proof: The distance function is preserved if and only if the dot product function is preserved. (This is a consequence of the ‘polarization identity’ above.) Let $m(\vec{v}, \vec{w}) = \vec{v} \cdot \vec{w}$, where \cdot denotes the vector dot product. Since $m(A\vec{v}, B\vec{w}) = \vec{v} \cdot (A^t \cdot B)\vec{w}$, we have

$$m(A\vec{v}, A\vec{w}) = m(\vec{v}, \vec{w}), \quad \forall v, w \in \mathbb{R}^3$$

if and only if

$$\vec{v} \cdot (A^t \cdot A)\vec{w} = \vec{v} \cdot \vec{w}, \quad \forall v, w \in \mathbb{R}^3$$

if and only if $A^t \cdot A = I_3$. \square

You may have been wondering how one could construct an isometry. This lemma gives us lots of examples.

Example 8.2.2. *A rotation matrix in 3 dimensions may be written in the form*

$$R(\phi, \theta, \psi) = \begin{pmatrix} r_{11} & r_{12} & \sin(\theta) \sin(\psi) \\ r_{21} & r_{22} & \sin(\theta) \cos(\psi) \\ \sin(\phi) \sin(\theta) & -\sin(\theta) \cos(\phi) & \cos(\theta) \end{pmatrix},$$

where

$$\begin{aligned} r_{11} &= \cos(\phi) \cos(\psi) - \cos(\theta) \sin(\phi) \sin(\psi), \\ r_{12} &= \sin(\phi) \cos(\psi) + \cos(\theta) \cos(\phi) \sin(\psi), \\ r_{21} &= -\cos(\phi) \sin(\psi) - \cos(\theta) \sin(\phi) \cos(\psi), \\ r_{22} &= -\sin(\phi) \sin(\psi) + \cos(\theta) \cos(\phi) \cos(\psi). \end{aligned}$$

and where the angles ϕ, θ, ψ are the ‘Euler angles’. This represents the rotation of 3-space obtained by the following sequence of rotations: rotate by angle ψ about the z -axis, rotate by the angle θ about the x -axis ($0 \leq \theta \leq \pi$), then rotate by angle ϕ about the z -axis again.

Although this is an interesting fact due to its explicitness, we shall not use this expression.

Question: Are there any isometries which do not come from matrices as in the above lemma? Yes: any translation gives rise to an isometry but translation is not linear, so cannot be described as a matrix. Moreover, any reflection about a plane containing the origin is also an isometry. A reflection never belongs to the group generated by the translations and rotations.

Question: Are there any examples of isometries which do not arise from a composition of a translation and an orthogonal matrix? No: the following theorem classifies all the isometries.

Theorem 8.2.1. *A function $f : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ is an isometry fixing the origin if and only if f is left multiplication by an orthogonal matrix.*

This will not be proven here (see Artin [Ar], chapter 4, section 5, Proposition 5.16).

As a consequence of this lemma, we see that if the matrix A gives rise to an isometry then $\det(A)$ is either equal to 1 or -1 (since $\det(A)^2 = \det(A^t \cdot A) = \det(I_3) = 1$). The matrix is invertible since, by definition, $A^{-1} = A^t$.

Lemma 8.2.2. *The set of all 3×3 matrices A such that the function $A : \mathbb{R}^3 \rightarrow \mathbb{R}^3$ is an isometry forms a group under matrix multiplication.*

Ponderable 8.2.1. *Verify the group axioms needed to prove this lemma.*

Notation: This group will be denoted $O_3(\mathbb{R})$ and called the **orthogonal group** of \mathbb{R}^3 . We denote by $SO_3(\mathbb{R})$ the following subset

$$SO_3(\mathbb{R}) = \{A \in O_3(\mathbb{R}) \mid \det(A) = 1\}.$$

which is called the **special orthogonal group** of \mathbb{R}^3 .

Lemma 8.2.3. *$SO_3(\mathbb{R})$ is a subgroup of $O_3(\mathbb{R})$.*

Ponderable 8.2.2. *Verify the group axioms for $SO_3(\mathbb{R})$.*

It is known that the number of cosets in $O_3(\mathbb{R})/SO_3(\mathbb{R})$ is 2. In fact, it is known that

$$O_3(\mathbb{R}) = SO_3(\mathbb{R}) \cup s \cdot SO_3(\mathbb{R}) \quad (\text{disjoint union}), \quad (8.1)$$

where s is the reflection in the above example (this follows from [Ar], chapter 4, section 5).

Lemma 8.2.4. *The isometry A in $O_3(\mathbb{R})$ is orientation-preserving if and only if $\det(A) = 1$.*

We will not prove this lemma here (see for example, [Ar] or [NST]).

8.3 Symmetries of the tetrahedron

Fix a tetrahedron centered at the origin, with one vertex along the z -axis. Each edge has an ‘opposite’ edge on the tetrahedron (which is actually perpendicular to it if you look at it straight on). Each vertex has an ‘opposite’ face.

There are orientation-preserving symmetries (called **rotations**) of the tetrahedron and orientation-reversing symmetries of the tetrahedron. The orientation-preserving symmetries of the tetrahedron will be denoted ST . They are obtained as follows:

- The 4 axes of symmetry through the centers of the faces yield 2 elements each (120 degree clockwise rotation when viewed from outside and a 240 degree rotation), for a total of 8 elements.

(This ‘tetrahedral symmetry’ allows for the mechanical construction of the Pyraminx.)

8.4. SYMMETRIES OF THE CUBE

- The 3 pairs of edges (formed by an edge and its opposite) yield one element each (a 180 degree rotation), for a total of 3 elements.

These, plus the identity, give 12 elements in ST .

Using the coset decomposition (8.1), we have $T = ST \cup s \cdot ST$ (disjoint), so

$$|T| = |ST| + |s \cdot ST| = 12 + 12 = 24,$$

by the addition principle (Theorem 2.4.1).

Remark 8.3.1. *It turns out that ST is essentially the alternating group A_4 of even permutations in S_4 and T is essentially S_4 itself. We shall state the precise result in the next chapter.*

8.4 Symmetries of the cube

We fix a cube centered about the origin in 3-space. The set of centers of the faces of a cube forms a set of vertices of an octahedron drawn inside the cube. This octahedron is called the ‘dual’ polyhedron. These two polyhedra have the same symmetry group, which we denote by O . There are orientation-preserving symmetries, or rotations, of the cube and orientation-reversing symmetries of the cube. The orientation-preserving symmetries of the cube will be denoted SO . They are obtained as follows:

- The 3 axes of symmetry through the centers of the faces yield 3 elements each (90 degree clockwise rotation when viewed from outside, a 180 degree rotation, and a 270 degree rotation), for a total of 9 elements.

(This ‘hexahedral symmetry’ allows for the mechanical construction of the Rubik’s Cube.)

- The 4 axes through the opposing vertices yield 2 elements each (all of order 3), for a total of 8 elements.

(This ‘tetrahedral symmetry’ allows for the construction of the Skewb [H].)

- The 6 axes through the opposing mid-edge points yield 1 element each (of order 2), for a total of 6 elements.

These elements, plus the identity, yield 24 elements. Therefore, I have proven the following result.

Lemma 8.4.1. *There are 24 orientation-preserving elements in O , i.e., $|SO| = 24$.*

Here’s another proof of this result.

Proof: Let V be the set of vertices of the cube. The group SO acts on the set V . Fix a v belonging to V and let $H = \text{stab}_{SO}(v)$. One can check that $|H| = 3$. (Since the only symmetry which fixes v is a rotation g about the line through v and its opposite vertex. Since g is order 3, $H = \langle g \rangle$ is order 3 as

well.) We have $|V| = 8$, so by Proposition 5.11.1 on orbits and stabilizers, we have $|SO/H| = |V|$. By Lagrange's theorem, $|SO| = |SO/H||H| = 8 \cdot 3 = 24$. \square

Now we know SO , what is O ? Note that s , the reflection in the example in the previous section, belongs to O . Using the coset decomposition of the previous section, we have the coset decomposition

$$O = SO \cup s \cdot SO \quad (\text{disjoint union}).$$

We know that $|s \cdot SO| = |SO| = 24$, so the addition principle again comes to the rescue, giving us the following result.

Lemma 8.4.2. *The order of the octahedral group is $|O| = 48$.*

Remark 8.4.1. *It turns out that SO is essentially the symmetric group S_4 and O is 'isomorphic to the direct product' $S_4 \times C_2$. We shall state the precise definitions and result in the next chapter.*

8.5 Symmetries of the dodecahedron

The set of centers of the faces of a dodecahedron forms a set of vertices of an icosahedron drawn inside. This icosahedron is called the 'dual' polyhedron. We fix a dodecahedron in 3-space so that the vertices of the dual icosahedron are as listed at the beginning of this chapter.

Example 8.5.1. *Here is an illustration of how to use **SAGE** to draw an dodecahedron:*

SAGE

```
sage: G = graphs.DodecahedralGraph()
sage: show(G.plot())
```

*This uses **SAGE**'s interface to the graph (as in vertices and edges) package *NetworkX*, so the plot (given in Figure 8.1) includes pre-labeled vertices.*

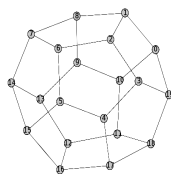


Figure 8.1: A dodecahedron.

8.5. SYMMETRIES OF THE DODECAHEDRON

Let SI denote the group of orientation-preserving symmetries of the dodecahedron. Note that SI is a finite subgroup of $SO_3(\mathbb{R})$. Let I denote the group of all symmetries of the dodecahedron. Note that I is a finite subgroup of $O_3(\mathbb{R})$ and that SI is a subgroup of I . Let F denote the set of faces of the dodecahedron, so $|F| = 12$. SI acts on F .

Lemma 8.5.1. *SI acts on F transitively.*

We won't prove this. If you look at a dodecahedron it follows 'by inspection'. The reason why this is useful is that it tells us that if x is any face, then any other face can be obtained from x by applying some element of SI . In other words, the orbit of x is all of F : $SI \cdot x = F$.

If x is any face then the only orientation-preserving symmetries which don't send x to a different face are rotations by an integer multiple of 72 degrees about the line passing through the center of x and the center of its opposite face. There are, for each face x , exactly 5 distinct rotations of this type. Therefore,

$$|\text{stab}_{SI}(x)| = 5.$$

By Proposition 5.11.1, we have

$$|SI/\text{stab}_{SI}(x)| = |SI \cdot x|,$$

so $|SI| = |\text{stab}_{SI}(x)| |SI \cdot x| = 5 \cdot 12 = 60$.

The elements of SI include

- rotation by $2\pi k/5$ (for $k \in \{0, 1, 2, 3, 4\}$) about the line passing through the center of a face and its opposite (this 'dodecahedral symmetry' allows for the construction of the Megaminx);
- rotation by $2\pi k/3$ (for $k \in \{0, 1, 2\}$) about the line passing through a vertex and its opposite;
- rotation by π about the center of an edge.

Subgroups include the following

- (a) Stabilizer of a vertex. These are all cyclic of order 3, and they are all conjugate. There are 10 distinct such subgroups since a vertex and its opposite share the same stabilizer.
- (b) Stabilizer of a face. These are all cyclic of order 5, and they are all conjugate. There are 6 distinct such subgroups since a face and its opposite share the same stabilizer.
- (c) Stabilizer of an edge. These are all cyclic of order 2, and they are all conjugate. There are 15 distinct such subgroups.

Ponderable 8.5.1. *Verify all these claims in (a)-(c).*

Remark 8.5.1. *It turns out that SI is essentially the alternating group A_5 of even permutations in S_5 and I is 'isomorphic to the direct product' $A_5 \times C_2$. We shall state the precise result in the next chapter.*

For an excellent discussion of the symmetries of the icosahedron, see [Ba].

8.6 Some thoughts on the icosahedron

Six is a number in itself, and not because God created the world in six days; rather the contrary is true. God created the world in six days because this number is perfect, and would remain perfect even if the work of the six days did not exist.

St. Augustine (354-430), The City of God

This section is based on one of John Baez's weekly internet columns [Ba].

A **duad** is a pair of diagonals (a diagonal is a segment from a vertex to its antipodal opposite vertex) of the icosahedron. The top of the icosahedron has 6 vertices and each diagonal must have exactly one of these 6 vertices as an endpoint. There are 12 vertices, hence 6 diagonals, hence

$$\binom{6}{2} = 15$$

different duads. Each duad determines a 'golden rectangle' (i.e., a rectangle whose ratio length/width is either the golden ratio $\phi = (1+\sqrt{5})/2$ or its inverse). We may identify a duad with a pair of distinct integers $\{(i, j) \mid 1 \leq i < j \leq 6\}$, i.e., with a 2-cycle in S_6 .

James J. Sylvester (the same one we met already in §2.2.1) called a partitioning

$$X = X_1 \cup \dots \cup X_n \quad (\text{disjoint})$$

of a set X a **syntheme** if each of the sets $X_i, 1 \leq i \leq n$, has the same number of elements. If we take

$$X = \{\text{set of diagonals of the icosahedron}\}$$

then a syntheme is a set of three duads, no two having a diagonal in common. There are

$$\binom{6}{2} \binom{4}{2} / 3! = 15$$

different syntheses (the $3!$ since there are $3!$ ways to permute the duads among themselves). A syntheme may be represented by a coloring of the vertices on the top of the icosahedron using three colors, each for exactly two vertices. We may identify a syntheme with a product of 3 distinct 2-cycles in S_6 . We partition the set of 15 duads into 5 groups of 3 as follows. (Recall each syntheme is a triple of duads.) First, pick a syntheme, A_1 . Pick another syntheme A_2 , so that A_1, A_2 have no duad in common. Continue on in this way until you pick 5 syntheses A_1, \dots, A_5 , no 2 of which have a duad in common. Such a choice of 5 syntheses is called a **pentad**. There are 6 pentads, which we label P_1, \dots, P_6 in any way you like. (A list of the 6 pentads is given in [R], chapter 7.) Any permutation of the 6 diagonals of the icosahedron gives rise to a permutation of the set of 6 pentads. Hence any permutation of the 6 diagonals of the icosahedron, which may be regarded as an element of S_6 , gives rise to a permutation of the set of 6 pentads, which may also be regarded as an element of S_6 . This gives a map

$$f : S_6 \rightarrow S_6.$$

Example 8.6.1. *Plotting the vertices of an icosahedron and diagonals of a ‘golden rectangle’ using SAGE (which calls the ray tracer Tachyon), a duad may be pictured as in Figure 8.2.*

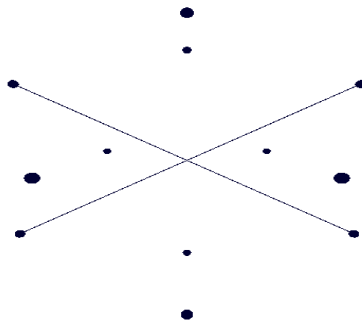


Figure 8.2: A duad

The above graph was obtained by plotting the points given in the table in §8.1 above and drawing lines between the points in the duad.

Each element of the rotation group of the icosahedron (i.e., the group of orientation-preserving symmetries of the icosahedron) must send a duad to a duad. Each duad has 4-fold symmetry, i.e., can be sent to itself in 4 ways. There are 15 duads, so there are $4 \cdot 15 = 60$ ways to send a duad to another. This is precisely the number of orientation-preserving symmetries of the icosahedron.

Example 8.6.2. *Here is another illustration of how to use SAGE to draw an icosahedron:*

SAGE

```
sage: G = graphs.IcosahedralGraph()
sage: len(G.edges())      # the number of edges
30
sage: len(G.vertices())   # the number of vertices
12
sage: show(G.plot())
```

This uses SAGE’s interface to NetworkX, as in the case of the dodecahedron above. The plot is given in Figure 8.3.

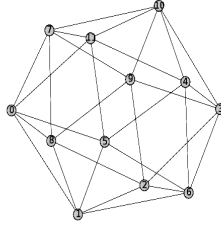


Figure 8.3: An icosahedron.

Note that the pairs $(0, 3)$ and $(4, 8)$ form a duad.

Some of these ideas will be discussed further in §9.3.2. See also [R] and [Ba].

8.7 901083404981813616 cubes

As a simple illustration of the basic idea in this section, consider the following problem. You have 6 colors: red, yellow, blue, orange, white and green. You also have an unpainted cube in front of you. How many ‘different’ ways can you paint the cube with one color per face? By ‘different’, we mean the following: two cubes are ‘colored the same’ if one looks the same as the other after a rotation by some element of the symmetry group of the cube; otherwise, they are ‘different’.

Ponderable 8.7.1. Show that the number of ‘different’ ways can you paint the cube with your 6 fixed colors, one color per face, is 30.

(Hint: Using an element of the symmetry group, we may assume the red face of the cube is on the bottom. There are 5 ways to paint the top face. Show that, up to rotating about the top-bottom axis, there are 6 ways to paint the sides.)

This same notion of ‘different’ is used in the problem discussed in this section.

In a cube-lover’s post from 1994 [Ho], Dan Hoey described the number of ‘different’ cube positions, if you allow for (orientation-preserving) cube symmetries. In other words, the orientation-preserving symmetry group of the cube acts on the space of all possible cube states. Dan, with Jerry Bryan, showed that 901083404981813616 is the number of equivalence classes. The full (possibly non-orientation-preserving) symmetry group of the cube acts also on the space of all possible cube states.

Herbert Kociemba and Mike Godfrey [KG] showed that 450541810590509978 is the number of equivalence classes in that case. Their very intricate computations were done theoretically and using computer software. Theoretical work used the ‘Polya-Burnside Lemma’ [Wbl]. (This is a lemma which goes by many names, including ‘the Lemma not due to Burnside’. It helps count sets on which there is a group action and is a generalization of the class equation we

8.7. 901083404981813616 CUBES

saw earlier.) Computational work used software such as GAP or a modification of Kociemba's Cube Explorer program.

Chapter 9

The illegal cube group

On front: *186,000 mps.*

On back: *It's not just a good idea, it's the law!*

T-shirt humor

Groups in mathematics are analogous to molecules in chemistry. We want to know what groups look like, to know how to describe them, how to compare them, how to make more of them, to know if they fall into families with similar properties, and so on.

Given two molecules, a chemist wants to know how to compare them, to understand their similarities and differences. Given two groups G_1, G_2 , an analogous question is to ask how ‘similar’ are they? (Exactly what is meant by ‘similar’ will be explained later.) We shall, in this chapter, introduce notions and techniques useful for comparing two groups. In a later chapter, we will focus on the 3×3 Rubik’s Cube group by comparing it to ‘better understood’ (or at least more ‘usual’) groups, such as the cyclic groups and the symmetric groups which we have already studied. In this chapter, we shall focus on understanding a somewhat simpler but closely related group, the group of all ‘legal’ and ‘illegal’ moves of the Rubik’s Cube. (In an ‘illegal’ move you can disassemble and reassemble the cube but you cannot peel off and swap individual stickers.) Remarkably, the illegal Rubik’s Cube group is simpler to understand than the Rubik’s Cube group itself, even though it is much larger.

To ‘explicitly construct’ the Rubik’s Cube group, we need to know how to construct

- quotient groups,
- direct products,
- semi-direct products.

These are studied in this chapter. Once the Rubik’s Cube group is constructed, it is a relatively easy matter to find its order, among other things.

9.1 Functions between two groups

To compare two groups, we often first talk about functions between them. If $f : G_1 \rightarrow G_2$ is a function from a group G_1 to a group G_2 then, as you might imagine, if f and G_1 were well understood then some information about G_2 could probably be obtained. It is the purpose of this chapter to find out to what extent this can be made precise.

A homomorphism between two groups is, roughly speaking, a function between them which preserves the (respective) group operations. These special types of functions are the most useful for our purposes.

Definition 9.1.1. *Let G_1, G_2 be groups, with $*_1$ denoting the group operation for G_1 and $*_2$ the group operation for G_2 . A function $f : G_1 \rightarrow G_2$ is a **homomorphism** if and only if, for all $a, b \in G_1$, we have*

$$f(a *_1 b) = f(a) *_2 f(b).$$

Foten, when there is little cause for confusion, we shall denote the binary operation on a group simply by \cdot or juxtaposition, rather than $*$.

Ponderable 9.1.1. *Prove the following: If $f : G_1 \rightarrow G_2$ is a homomorphism of groups then*

$$f(G_1) = \{g \in G_2 \mid g = f(x), \text{ for some } x \in G_1\}$$

is a subgroup of G_2 .

The subgroup $f(G_1) \subset G_2$ is called the **image of f** and is sometimes denoted $\text{im}(f)$.

Example 9.1.1. *Let G be a group and h a fixed element of G . Define $f : G \rightarrow G$ by*

$$f(g) = h^{-1} \cdot g \cdot h, \quad g \in G.$$

Then the following simple trick

$$f(a \cdot b) = h^{-1} \cdot (a \cdot b) \cdot h = h^{-1} \cdot a \cdot h \cdot h^{-1} \cdot b \cdot h = f(a) \cdot f(b)$$

shows that f is a homomorphism. In this case, $\text{im}(f) = G$, i.e., f is surjective.

*Here is an example of how to use **SAGE** to construct a group homomorphism. Basically, you must specify the image of the homomorphism on the generators.*

— SAGE —

```
sage: G = SymmetricGroup(4)
sage: gensG = G.gens()
sage: h = G([(1,3,4,2)])          # a 'random' element of G
sage: gensG_h = [h*g*h^(-1) for g in gensG]
sage: phi = PermutationGroupMorphism_im_gens(G,G,gensG,gensG_h)
sage: phi.image(G)
'Group([ (1,3,2,4), (2,4) ])'
sage: phi.range()
Symmetric group of order 4! as a permutation group
```

In this last line, **SAGE** tells us that the map `phi` describes a one-to-one and onto map from S_4 to itself.

Ponderable 9.1.2. Let

$$A = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}, \quad B = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}.$$

Now, let $G = \langle A, B \rangle$ denote the group of all matrices which can be written as any arbitrary product of these two matrices (in any order and with as many terms as you want). Show that we have

$$G = \{I_3 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{pmatrix}\}.$$

(The interested reader may want to try to check this by regarding each such matrix as a permutation matrix.) Define $f : G \rightarrow S_3$ by

g	$f(g)$
I_3	1
A	s_1
B	s_2
$A \cdot B$	$s_1 \cdot s_2$
$B \cdot A$	$s_2 \cdot s_1$
$A \cdot B \cdot A$	$s_1 \cdot s_2 \cdot s_1$

Show that this is a homomorphism.

Example 9.1.2. The function

$$\text{sign} : S_n \rightarrow \{\pm 1\},$$

which assigns to each permutation its sign, is a homomorphism. This was already proven in §3.1. However, another reason why this is true is because the sign of a permutation g is the determinant of the associated permutation matrix $P(g)$. Since the determinant of the product is the product of the determinants (see §2.2.5), we have

$$\text{sgn}(gh) = \det P(gh) = \det(P(g)P(h)) = \det P(g) \det P(h) = \text{sign}(g)\text{sign}(h),$$

for all $g, h \in S_n$. From this it follows that sign is a homomorphism.

Lemma 9.1.1. If $f : G_1 \rightarrow G_2$ is a homomorphism then

9.1. FUNCTIONS BETWEEN TWO GROUPS

- (a) $f(e_1) = e_2$, where e_1 denotes the identity element of G_1 and e_2 denotes the identity element of G_2 ,
- (b) $f(x^{-1}) = f(x)^{-1}$, for all x belonging to G_1 ,
- (c) $f(y^{-1} *_1 x *_1 y) = f(y)^{-1} *_2 f(x) *_2 f(y)$, for all x, y belonging to G_1 ,

where $*_1$ denotes the group operation for G_1 and $*_2$ the group operation for G_2 .

Proof:

- (a) We have $f(x) = f(x *_1 e_1) = f(x) *_2 f(e_1)$, for any $x \in G_1$. Multiply both sides of this equation on the left by $f(x)^{-1}$.
- (b) We have, by part (a), $e_2 = f(e_1) = f(x *_1 x^{-1}) = f(x) *_2 f(x^{-1})$. Multiply both sides of this equation on the left by $f(x)^{-1}$.
- (c) The proof of this part is left to the reader.

□

Ponderable 9.1.3. Show part (c). (Hint: use the definition and part (b).)

Definition 9.1.2. Let G_1, G_2 be finite groups. We say that G_1 **embeds** (or **injects**) into G_2 if there exists an injective homomorphism $f : G_1 \rightarrow G_2$. A homomorphism $f : G_1 \rightarrow G_2$ is an **isomorphism** if it is a bijection (as a function between sets). In this case, we call G_1 and G_2 **isomorphic** and write $G_1 \cong G_2$. An isomorphism from a group G to itself is called an **automorphism**.

The concept of an isomorphism was first introduced by C. Jordan, whom we have already met. Besides group theory, Jordan is famous in mathematical circles for rigorously proving one of the most intuitively obvious, but difficult to rigorously prove, results: Any simple closed curve divides the plane into two regions. (Here, a **simple curve** in the plane is a continuous path which is not the disjoint union of two or more such paths. A **closed curve** is a continuous path which starts and ends at the same point.)

An isomorphism is the notion we will use when we want to say two groups are ‘essentially the same group’, i.e., one is basically a carbon copy of the other with the elements relabeled and the binary operation modified. For example, if you have any two groups of order 2 then they must be isomorphic. (The interested reader can verify that the map which sends the identity of one group to the identity of the other and the only non-identity of one group to the only non-identity of the other must be a group isomorphism.) In other words, there is only one group of order 2, up to isomorphism.

Let $O(n)$ denote the the number of non-isomorphic groups of order n , so O is a function $O : \mathbb{N} \rightarrow \mathbb{N}$. This is a rather curiously behaving function, jumping up at integers which have lots of prime factors, such as 16, 24, 32, 48,

Ponderable 9.1.4. Show the following:

(a) If $m \geq 1$ then

$$G = \{e^{2\pi i k/m} \mid 0 \leq k \leq m-1\},$$

with group operation given by ordinary multiplication of complex numbers, is isomorphic to the additive group $\mathbb{Z}/m\mathbb{Z}$.

(b) If $m > 1, n > 1$ have no prime divisors in common, i.e., if m, n are relatively prime, then $C_m \times C_n \cong C_{mn}$. (For the definition of the direct product \times , see Definition 9.6.1, below.)

(Hint: (a) Let $a \in C_m$ be a generator. The map $f = f_{m,a} : C_m \rightarrow \mathbb{Z}/m\mathbb{Z}$ which satisfies $f(a) = 1$ extends to a unique homomorphism between these groups. Show that this is an isomorphism. (b) First, show $g(x + mn\mathbb{Z}) = (x + m\mathbb{Z}, x + n\mathbb{Z})$, for all $x \in \mathbb{Z}$ defines a homomorphism $\mathbb{Z}/mn\mathbb{Z} \rightarrow \mathbb{Z}/m\mathbb{Z} \times \mathbb{Z}/n\mathbb{Z}$. Next, using the assumption that m, n are relatively prime, show that this function g is an injection. Conclude, since the domain and range of g both have the same cardinality, that g is also a surjection, by Ponderable 2.1.6. Finally, using part (a) show $C_m \times C_n \cong C_{mn}$.)

9.2 Group actions

We need to study actions, whose discussion was begun in chapter 5 above, in more detail.

Lemma 9.2.1. *Let G be a group and X a finite set. If G acts on X (on the left, respectively, on the right) then there is a homomorphism $G \rightarrow S_X$ given by $g \mapsto \phi_g$. Conversely, if $\phi : G \rightarrow S_X$ is a homomorphism then $\phi(g) : X \rightarrow X$ defines a (left, respectively, right) action of G on X .*

This is an immediate consequence of the definition of an action.

Example 9.2.1. *Let G be the Rubik's Cube group generated by the basic moves R, L, U, D, F, B . For each move $g \in G$, let*

- $\rho(g)$ be the corresponding permutation of the set of vertices V of the cube, and
- $\sigma(g)$ be the corresponding permutation of the set of edges E of the cube.

Then it is possible to show that

(a) $\rho : G \rightarrow S_V$ is a homomorphism,

(b) $\sigma : G \rightarrow S_E$ is a homomorphism.

To check this, the first fundamental theorem of cube theory in §9.6.1 is useful.

9.2. GROUP ACTIONS

Let G act on a set X , denoted by $\phi_g : X \rightarrow X$ for $g \in G$. This action also gives rise to an action on $X \times X$: $\phi_{2,g} : X \times X \rightarrow X \times X$ for $g \in G$, where $\phi_{2,g}(x_1, x_2) = (\phi_g(x_1), \phi_g(x_2))$, for each $(x_1, x_2) \in X \times X$. When G acts transitively on X (see Definition 5.10.3 above), it is remarkable enough. Sometimes, G even acts transitively on $X \times X - \Delta$ ($\Delta = \{(x, x) \mid x \in X\}$, the ‘diagonal’), a far more special circumstance (in this case, the action is called **doubly transitive**). The definition below generalizes this idea.

Definition 9.2.1. *Let a group G act on a set X . We call the action **k -tuply transitive** if for each pair of ordered k -tuples $(x_1, x_2, \dots, x_k), (y_1, y_2, \dots, y_k)$ of distinct elements belonging to X (so $x_i \neq x_j$ and $y_i \neq y_j$), there is a $g \in G$ such that $y_i = \phi_g(x_i)$, for each i with $1 \leq i \leq k$.*

If $k = 2$ then ‘2-tuply transitive’ is the same as ‘doubly transitive’.

Ponderable 9.2.1. *Is the Rubik’s Cube group doubly transitive on the set of edge facets?*

The following result is one illustration of how unique the symmetric group and alternating group are. (Recall that the alternating group A_n was defined in Example 5.6.3 above.)

Theorem 9.2.1. *If $k > 5$ and G is a permutation group acting k -transitively on a finite set X (so G is a subgroup of S_X) then G is isomorphic to S_m or to A_n , for some $m \geq k$ or some $n \geq k + 2$.*

Conversely, S_n acts n -transitively on $\{1, 2, \dots, n\}$ and A_n acts $(n-2)$ -transitively on $\{1, 2, \dots, n\}$.

This is proven in [DiMo], Theorem 7.6A. Some examples are given in §12.4 below.

Remark 9.2.1. *We shall, in conjunction with our group-theoretical determination of the Rubik’s Cube group proven later, be able to deduce from Theorem 9.2.1 the following result.*

Corollary 9.2.1.

- (a) *Suppose c_1, \dots, c_6 are any 6 distinct corner subcubes of the Rubik’s cube. Let c'_1, \dots, c'_6 be any other 6 distinct corner subcubes. There is an element of the Rubik’s Cube group G which fixes all edge subcubes and sends c_i to c'_i , for all $i = 1, \dots, 6$. Loosely speaking, we may express this by saying G acts 6-transitively on the corners, fixing the edges.*
- (b) *Suppose c_1, \dots, c_8 are any 8 distinct corner subcubes of the Rubik’s cube. Let c'_1, \dots, c'_8 be any other 8 distinct corner subcubes. There is an element of the Rubik’s Cube group G which sends c_i to c'_i , for all $i = 1, \dots, 8$ (and might not fix all edge subcubes). Loosely speaking, we may express this by saying G acts 8-transitively on the corners (allowing the edges to be moves).*
- (c) *G acts 10-transitively on the edges, fixing the corners.*
- (d) *G acts 12-transitively on the edges but may permute the corners.*

9.3 When two groups are really the same

This section discusses group isomorphisms.

Example 9.3.1. Let G be the group in Ponderable 9.1.2 and $f : G \rightarrow S_3$ the homomorphism. This is in fact an isomorphism.

SAGE

```
sage: MS = MatrixSpace(GF(2), 3, 3)
sage: A = MS([0, 1, 0, 1, 0, 0, 0, 0, 1])
sage: B = MS([1, 0, 0, 0, 0, 1, 0, 1, 0])
sage: G = MatrixGroup([A,B])
sage: G1 = SymmetricGroup(3)
sage: imG = G._gap_().IsomorphismPermGroup().Image(G._gap_())
sage: gens = imG.GeneratorsOfGroup()
sage: G2 = PermutationGroup(gens)
sage: G2.is_isomorphic(G1)
True
```

Example 9.3.2. Let H be the subgroup of the Rubik's Cube group generated by the basic move R : $H = \langle R \rangle$. Then $H \cong C_4$ (where C_4 denotes the cyclic group of order 4).

Example 9.3.3. Recall that to each permutation g of the set $\{1, 2, \dots, n\}$ we can associate a $n \times n$ permutation matrix $P(g)$ in such a way that

$$P(g) \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} x_{g(1)} \\ x_{g(2)} \\ \vdots \\ x_{g(n)} \end{pmatrix}.$$

(Here the image of i under the permutation g is denoted $g(i)$, though in fact one plugs i into g from the left.) We let P_n denote the set of all $n \times n$ permutation matrices. This is a group under matrix multiplication. The function

$$\begin{aligned} P : S_n &\rightarrow P_n, \\ g &\mapsto P(g) \end{aligned}$$

is an isomorphism. This is a bijection and a homomorphism, by the discussion given earlier in chapter 3 on permutations (see Lemma 3.2.2 and Theorem 3.2.1).

Example 9.3.4. In this example, we consider the full symmetry groups (in $O(3, \mathbb{R})$) of various solids and planar figures, the symmetries being regarded as 3-dimensional transformations (as opposed to 2-dimensional transformations in §5.3).

From [NST], Table 15.7, we have the following table of isomorphisms:

9.3. WHEN TWO GROUPS ARE REALLY THE SAME

Name	Notation	Isomorphic to
symmetry group of tetrahedron	T	S_4
rotation group of tetrahedron	ST	A_4
symmetry group of octahedron	O	$S_4 \times C_2$
rotation group of octahedron	SO	S_4
symmetry group of icosahedron	I	$A_5 \times C_2$
rotation group of icosahedron	SI	A_5
symmetry group of regular n -gon		$D_{2n}, n \text{ odd}$
rotation group of regular n -gon		$D_n \times C_2, n \text{ even}$ D_n

Example 9.3.5. This example may be found in [B1].

Let Q denote the quaternion group,

$$Q = \{1, -1, i, -i, j, -j, k, -k\},$$

where $i^2 = j^2 = k^2 = -1, ij = k, jk = i, ki = j$, as in §5.1. Then Q is isomorphic to the group $Q^* = \langle a, b \rangle \subset G$, where G is the Rubik's Cube group and

$$\begin{aligned} a &= F^2 \cdot M_R \cdot U^{-1} \cdot M_R^{-1} \cdot U^{-1} \cdot M_R \cdot U \cdot M_R^{-1} \cdot U \cdot F^2, \\ b &= F \cdot U^2 \cdot F^{-1} \cdot U^{-1} \cdot L^{-1} \cdot B^{-1} \cdot U^2 \cdot B \cdot U \cdot L, \end{aligned}$$

via the map $f : Q^* \rightarrow Q$ defined by $\phi(a) = i, \phi(b) = k$.

One proof of this claim will be discussed in a later chapter. (The easiest way to prove this theoretically uses ideas I haven't yet introduced.) However, we can use **SAGE** to verify this as follows:

SAGE

```
sage: rubik = CubeGroup()
sage: r = rubik.R()
sage: l = rubik.L()
sage: u = rubik.U()
sage: d = rubik.D()
sage: f = rubik.F()
sage: b = rubik.B()
sage: mr = r^(-1)*l
sage: a = f^2*mr*f^(-1)*mr^(-1)*u^(-1)*mr*f*mr^(-1)*u*f^2
sage: b = f*u^2*f^(-1)*u^(-1)*l^(-1)*b^(-1)*u^2*b*u*l
sage: G = PermutationGroup([a,b])
sage: G.order()
8
sage: G.is_abelian()
False
sage: b*a*b == a
True
sage: b^4
()
sage: a^2 == b^2
True
```

Now we refer to Example 10.4.1 below for further details (actually more along the lines of detailed hints to an exercise).

Lemma 9.3.1. *As above, let (G, \cdot) be a group and h a fixed element of G . Define $c_h : G \rightarrow G$ by*

$$c_h(g) = h \cdot g \cdot h^{-1}, \quad g \in G.$$

This is an automorphism.

Proof: To verify this, we must show that $f = c_h$ is an injective and surjective homomorphism (we drop the subscript for simplicity of notation).

First, we show that f is injective. Suppose $f(g_1) = f(g_2)$. Then $f(g_1 \cdot g_2^{-1}) = 1$, so that $h \cdot g_1 \cdot g_2^{-1} \cdot h^{-1} = 1$. Multiply both sides of this equation on the right by h and on the left by h^{-1} . We obtain $g_1 \cdot g_2^{-1} = 1$. This implies $g_1 = g_2$, so f is injective.

Now we show f is surjective. Let x be an arbitrary but fixed element of G . Let $y = h^{-1} \cdot x \cdot h$. Then

$$f(y) = f(h^{-1} \cdot x \cdot h) = h \cdot h^{-1} \cdot x \cdot h \cdot h^{-1} = x.$$

Therefore, f is surjective.

It was previously verified that f is a homomorphism. \square

Definition 9.3.1. *An automorphism as in Lemma 9.3.1 is called **inner**. An automorphism of G which is not of this form, for some $h \in G$, is called **outer**.*

The concepts of inner and outer automorphisms were introduced by Otto Hölder (1859-1937). Hölder also classified all the finite simple groups, up to isomorphism, of order ≤ 200 . (The definition of ‘simple’ is given in Definition 9.5.1 below.) The classification of *all* finite simple groups has only recently been finished and is said to require over 3000 pages of details.

Notation: The set of all automorphisms of a group G is denoted $\text{Aut}(G)$. The subset of inner automorphisms is denoted

$$\text{Inn}(G) := \{f \in \text{Aut}(G) \mid f = c_h, \text{ some } h \in G\},$$

in the notation of the Lemma 9.3.1.

Ponderable 9.3.1.

- (a) Show $\text{Aut}(G)$ is a group with composition as the group operation.
- (b) Show that $\text{Inn}(G)$ is a subgroup of $\text{Aut}(G)$.

9.3.1 Conjugation in S_n

In the early part of the 1900’s a brilliant Indian mathematician, Srinivasa Ramanujan (1887-1920), lay in hospital in England very ill. The great British mathematician, G. H. Hardy (1877-1947), visited him and remarked that he

took taxi number 1729, a ‘boring number’. ‘On the contrary,’ Ramanujan argued, ‘it’s a very interesting number, being the smallest integer which is the sum of two cubes in two different ways.’ One of the most interesting and original mathematicians of all time, Ramanujan worked (often in collaboration with Hardy) on elliptic functions, continued fractions, and infinite series [MT]. He knew an astonishing number of unusual facts about integers, such as this one.

What does Ramanujan have to do with conjugation? Read on, and you’ll find out!

As promised in §5.9, we will find a simple criteria for determining when two permutations (in S_n) are conjugate (in S_n).

The following result will be of importance to us in a later chapter.

Lemma 9.3.2. *Suppose $f : S_n \rightarrow S_n$ is an inner automorphism. If $g \in S_n$ is a disjoint product of cycles of length k_1, \dots, k_r then $f(g)$ is a disjoint product of cycles of length k_1, \dots, k_r .*

In other words, an inner automorphism (i.e., conjugation) must ‘preserve the cycle structure’.

Proof: Since f is inner, let it be conjugation by some element $h \in S_n$ say, so $f(g) = h^{-1}gh$, for all $g \in S_n$. Let $(i)g \in \{1, \dots, n\}$ denote the image of $i \in \{1, \dots, n\}$ under $g \in S_n$. The lemma is a consequence of the following simple calculation: if $(i)g = j$ then, for all i with $1 \leq i \leq n$, we have

$$((i)h)(h^{-1}gh) = (j)h. \quad (9.1)$$

In other words, if g sends $i \mapsto j$ then $h^{-1}gh$ sends $(i)h \mapsto (j)h$. It follows that g and $h^{-1}gh$ have the cycle structure. \square

Theorem 9.3.1. *Two elements $g, g' \in S_n$ are conjugate if and only if they have the same cycle structure.*

Proof: The Lemma proves the ‘only if’ direction of this equivalence. Suppose that $g, g' \in S_n$ have the same cycle structure. Write their disjoint cycle decompositions using the lexicographical ordering imposed on the lengths of the cycles occurring in the decomposition: say

$$\begin{aligned} g &= (i_1, \dots, i_{n_1})(i_1, \dots, i_{n_2}) \dots (i_1, \dots, i_{n_k}), \\ g' &= (i'_1, \dots, i'_{n_1})(i'_1, \dots, i'_{n_2}) \dots (i'_1, \dots, i'_{n_k}), \end{aligned}$$

where $1 \leq n_1 \leq \dots \leq n_k \leq n$ and $n = n_1 + \dots + n_r$. Pick an $h \in S_n$ such that $h : i_j \mapsto i'_j$, for all j with $1 \leq j \leq n$. Then $g' = h^{-1}gh$, by (9.1). \square

A **partition** of n is an r -tuple, (n_1, \dots, n_r) of positive integers such that $1 \leq n_1 \leq \dots \leq n_r \leq n$ and $n = n_1 + \dots + n_r$. From the results above, it follows that each conjugacy class of S_n corresponds to one and only one partition of n . In particular, the number of conjugacy classes of S_n is equal to the number $p(n)$ of distinct partitions of n . An asymptotic formula for this number was found by Srinivasa Ramanujan. His formula is too complicated for this book, but a special case of it is the following beautiful formula:

$$\lim_{n \rightarrow \infty} \frac{\log p(n)}{\sqrt{n}} = \pi \sqrt{\frac{2}{3}}.$$

Example 9.3.6. Here is an example of how to use **SAGE** to show that the number of conjugacy classes of S_5 is equal to the number $p(5)$ of distinct partitions of 5.

SAGE

```
sage: G = SymmetricGroup(5)
sage: CG = G.conjugacy_classes_representatives()
sage: len(CG)
7
sage: number_of_partitions(5, algorithm='bober')
7
sage: number_of_partitions(2000, algorithm='bober')
4720819175619413888601432406799959512200344166
```

This uses Jonathon Bober's very fast implementation to compute $p(5)$ and $p(2000)$.

9.3.2 ... and a side order of automorphisms, please

Though we shall not need it here, the following fact is interesting since it illustrates what a unique role the symmetric group S_6 plays in the family of all symmetric groups.

Theorem 9.3.2. If $n \neq 2, 6$ then the homomorphism $\phi : S_n \rightarrow \text{Aut}(S_n)$ defined by $\phi(g) = c_g$ (where c is as in Lemma 9.3.1 above) is an isomorphism:

$$S_n \cong \text{Aut}(S_n).$$

If $n = 6$ then $|\text{Aut}(S_6)| = 2 \cdot |S_6|$.

The following example continues the discussion from §8.6.

Example 9.3.7. Any permutation of the 6 diagonals of the icosahedron, which may be regarded as an element of S_6 , gives rise to a permutation of the set of 6 pentads, which may also be regarded as an element of S_6 . This gives a map

$$f : S_6 \rightarrow S_6,$$

which is in fact a homomorphism. This homomorphism is injective so it is actually an automorphism.

However, a 2-cycle on the set of 6 diagonals (i.e., swapping exactly 2 diagonals) does not induce a 2-cycle on the set of these 6 pentads. In fact, a 2-cycle on the set of diagonals gives rise to a product of three disjoint 2-cycles on the set of these 6 pentads. Therefore, by the above theorem (which says that an inner automorphism must preserve the cycle structure) this automorphism f cannot be an inner automorphism.

9.4 Kernels are normal, some subgroups are not

Let $f : G_1 \rightarrow G_2$ be a homomorphism between two groups. Let

$$\ker(f) = \{g \in G_1 \mid f(g) = e_2\},$$

where e_2 is the identity element of G_2 . This set is called the **kernel** of f .

Lemma 9.4.1. *$\ker(f)$ is a subgroup of G_1 .*

The interested reader can check this for him/herself by using the definition of a homomorphism.

Example 9.4.1. *Let*

$$\operatorname{sgn} : S_n \rightarrow \{\pm 1\}$$

denote the homomorphism which associates with a permutation either 1, if it is even, or -1 , if it is odd. Then $A_n = \ker(\operatorname{sgn}) \subset S_n$.

The following properties of the kernel are useful.

Lemma 9.4.2. *Let $f : G_1 \rightarrow G_2$ be a homomorphism between two groups.*

- (a) *The homomorphism f is injective if and only if $\ker(f) = \{e_1\}$.*
- (b) *If g belongs to the kernel and x is any element of G_1 , then $x^{-1} \cdot g \cdot x$ must also belong to the kernel.*

Proof:

- (a) That f is injective if and only if $f(g_1) = f(g_2)$ implies $g_1 = g_2$ ($g_1, g_2 \in G_1$). Note that $f(g_1) = f(g_2)$ is true if and only if $f(g_1 \cdot g_2^{-1}) = e_2$. If $\ker(f) = \{e_1\}$ then $f(g_1 \cdot g_2^{-1}) = e_2$ implies $g_1 \cdot g_2^{-1} = e_2$, which implies $g_1 = g_2$, which implies f is injective.

Therefore, if $\ker(f) = \{e_2\}$ then f is injective. Conversely, if f is injective then $f(x) = f(e_1) (= e_2)$ implies $x = e_1$ ($x \in G_1$). This implies $\ker(f) = \{e_1\}$.

- (b) Multiply both sides of $e_2 = f(g)$ on the left by $f(x)^{-1}$ and on the right by $f(x)$. We get

$$e_2 = f(x)^{-1} \cdot e_2 \cdot f(x) = f(x^{-1}) \cdot f(g) \cdot f(x) = f(x^{-1} \cdot g \cdot x),$$

as desired.

□

Definition 9.4.1. *Let H be a subgroup of G . We say that H is a **normal subgroup** if, for each $g \in G$, $g^{-1} \cdot H \cdot g = H$ (i.e., for each $g \in G$ and each $h \in H$, $g^{-1} \cdot h \cdot g$ belongs to H).*

Notation: Sometimes we denote ‘ H is a normal subgroup of G ’ by

$$H \triangleleft G$$

9.4.1 Examples of non-normal subgroups

We saw above that examples of normal subgroups are easy to construct. For instance, the kernel of any homomorphism $f : G \rightarrow H$ is a normal subgroup of G . For example, the kernel of the sgn homomorphism is normal: $A_n \triangleleft S_n$ (see Example 9.4.1). More generally, if G is any group then the commutator subgroup, G' , is a normal subgroup of G . (The commutator subgroup of S_n is A_n .)

What about non-normal subgroups? Examples of subgroups which are *not* normal are easy to come by.

Example 9.4.2.

- (a) If $n > 5$ and H is any non-trivial proper subgroup of A_n (for example, any non-trivial cyclic subgroup) then H is not normal in A_n (see Theorem 9.4.1 below).
- (b) Any subgroup of S_6 of order 2 is non-normal.
- (c) Here is a way to get **SAGE** (rather **GAP**, via **SAGE**) to return the normal subgroups of a group:

```

SAGE
sage: G = SymmetricGroup(5)
sage: GG = gap(G)
sage: GG.NormalSubgroups()
[ Group( () ), AlternatingGroup( [ 1 .. 5 ] ), SymmetricGroup( [ 1 .. 5 ] ) ]
sage: G = DihedralGroup(5)
sage: GG = gap(G)
sage: GG.NormalSubgroups()
[ Group( () ), Group( [ (1,2,3,4,5) ] ), Group( [ (1,2,3,4,5), (1,5)(2,4) ] ) ]
```

This tells us the following facts.

1. Any subgroup of S_5 other than itself, A_5 , and the trivial group, is not normal.
2. Any subgroup of D_5 other than itself, $C_5 = \langle (1,2,3,4,5) \rangle$, and the trivial group, is not normal.

We have already shown the following result.

Lemma 9.4.3. *If $f : G_1 \rightarrow G_2$ is a homomorphism between two groups then $\ker(f)$ is a normal subgroup of G_1 .*

Example 9.4.3. *Here is an example of how to use **SAGE** to compute the kernel of a group homomorphism.*

9.4. KERNELS ARE NORMAL, SOME SUBGROUPS ARE NOT

SAGE

```
sage: G = CyclicPermutationGroup(10)
sage: g0 = G.gens()[0]
sage: phi = PermutationGroupMorphism_im_gens(G,G,[g0],[g0^2])
sage: phi.kernel()
Group([ (1,6)(2,7)(3,8)(4,9)(5,10) ])
sage: ker = PermutationGroup(["(1,6)(2,7)(3,8)(4,9)(5,10)"])
sage: ker.order()
2
```

This tells us the kernel of the homomorphism of abelian groups, $\phi : C_{10} \rightarrow C_{10}$ given by $\phi(x) = x^2$, is $\{1, (1,6)(2,7)(3,8)(4,9)(5,10)\}$ (a subgroup of order 2).

9.4.2 The alternating group

The following remarkable result about the alternating group will not be needed to understand the structure of the Rubik's Cube. However, the theorem below is interesting because of its connection with the fact (due to Ruffini and Abel) that you cannot solve the general polynomial of degree 5 or higher using radicals, i.e., that there is no analog of the quadratic formula for polynomials of degree 5 or higher. Explaining this connection would, unfortunately, take us too far from our main topic.

Theorem 9.4.1. *If X has 5 elements or greater, then A_X has no non-trivial proper normal subgroups. In other words, if $H \triangleleft A_X$ is a normal subgroup, then either $H = \{1\}$ or $H = A_X$.*

This will not be proven here. (For a proof, see, for example, [R] or [Ar], chapter 14.)

However, the next fact about the alternating group *will* be needed later in our determination of the structure of the Rubik's Cube group. This fact also arose in connection with our discussion of the 'legal positions' of the 15 Puzzle in chapter 4 (see §4.1).

Proposition 9.4.1. *Let H be the subgroup of S_n generated by all the 3-cycles in S_n ; then $H = A_n$.*

Proof: Since $\text{sgn} : S_n \rightarrow \{\pm 1\}$ is a homomorphism, and since any 3-cycle is even, any product of 3-cycles must also be even. Therefore, $H \subset A_n$. If $g \in A_n$ then g must swap an even number of the inequalities $1 < 2 < \dots < n-1 < n$, by Definition 3.1.1. Therefore (since any permutation may be written as a product of 2-cycles, Theorem 3.4.1), g must be composed of permutations of the form $(i, j)(k, l)$ or $(i, j)(j, k)$. But $(i, j)(k, l) = (i, j, k)(j, k, l)$ and $(i, j)(j, k) = (i, j, k)$. Therefore, $g \in H$. This implies $A_n \subset H$, so $A_n = H$. \square

The following more precise result is very useful for the purposes of the analysis of permutation puzzles.

Lemma 9.4.4. ([W]) *Let X be a finite set, $|X| \geq 3$ and fix u, v as elements in X . Then the 3-cycles (u, v, x) , as x runs over all elements of $X - \{u, v\}$, generate A_X .*

This lemma gives an even stronger statement than the previous claim. Now, instead of only one element being fixed, there are two elements fixed and the alternating group is still generated.

9.5 Quotient groups

QUOTIENT, n. A number showing how many times a sum of money belonging to one person is contained in the pocket of another—usually about as many times as it can be got there.

Ambrose Bierce, The Devil's Dictionary

One of the most useful facts about normal subgroups is the following result/definition.

Lemma 9.5.1. *If H is a normal subgroup of G then the coset space G/H with the binary operation,*

$$aH \cdot bH = (ab)H, \quad (aH)^{-1} = a^{-1}H,$$

for all a, b belonging to G , is a group. The identity element of this group is the trivial coset H .

This group G/H is called the **quotient group** of G by H and is sometimes pronounced ‘ $G \bmod H$ ’. Older texts may use the terminology ‘factor group’ instead. This notion, for an abstract group, was first introduced by Otto Hölder, whom we have already met. He originally started out training to become an engineer. Hölder made important contributions not only to group theory but also to Fourier series.

Example 9.5.1. *Here is an example of how to use SAGE to show that A_5 is a simple group, normal as a subgroup of S_5 , and that the quotient S_5/A_5 is a group of order 2.*

SAGE

```
sage: G = SymmetricGroup(5)
sage: H = AlternatingGroup(5)
sage: H.is_subgroup(G)
True
sage: H.is_normal(G)
True
sage: H.is_simple()
True
sage: G.quotient_group(H)
Permutation Group with generators [(1,2)]
```

9.5. QUOTIENT GROUPS

Recall from §9.4.1 that the commutator subgroup is normal. Another example of a quotient group which arises naturally is called the **abelianization** of G : the quotient group $G_{ab} = G/G'$, where $G' = [G, G]$ denotes the commutator subgroup.

Example 9.5.2. *If $f : G_1 \rightarrow G_2$ is a homomorphism between two groups, then $G_1/\ker(f)$ is a quotient group.*

Next we introduce an important idea first emphasized by E. Galois, the French mathematician mentioned in the introduction. Though Galois attended secondary school, he had trouble entering the French university system. Sadly, by the time he entered the École Normale Supérieure in November 1829, he also began to get caught up in the activities of the French Revolution of 1830 which toppled Charles X. His political activities resulted in his being expelled from school in December 1830. Though apparently he had been a rather rebellious teenager, this only frustrated him further. He died in a fight with another Frenchman.

Galois' work was the first to begin to illuminate the 'basic building blocks' of the collection of finite groups (i.e., the analog of the idea that atoms are the 'basic building blocks' of molecules). Galois introduced the ideas of solvable groups and normal subgroups, in the context of permutation groups. Roughly speaking, one may say that the basic building blocks of finite groups are those groups which have no proper non-trivial normal subgroups. Intuitively, this is because a non-trivial quotient group (by a normal subgroup) is closely related to the original group but smaller in size (and hence perhaps subject to analysis by an inductive argument of some type). These basic building blocks are called 'simple' groups.

Definition 9.5.1. *A simple group is a group with no proper normal subgroups other than the trivial subgroup $\{1\}$.*

Example 9.5.3. *If p is a prime then C_p (the cyclic group having p elements) is simple. In fact, if G is any group which is both abelian and simple then there is a prime p such that $G \cong C_p$. If $n > 4$ then A_n is simple (as was stated above in Theorem 9.4.1). These facts are proven in [R].*

Simple groups are not very abundant. In fact, the first non-abelian simple group is of order 60 (it's A_5).

The following basic result describes the quotient group $G_1/\ker(f)$.

Theorem 9.5.1. (First isomorphism theorem) *If $f : G_1 \rightarrow G_2$ is a homomorphism between two groups, then $G_1/\ker(f)$ is isomorphic to $f(G_1)$.*

Proof: $\ker(f)$ is a normal subgroup of G_1 , so $G_1/\ker(f)$ is a group. We must show that this group is isomorphic to the group $f(G_1)$. Define $\phi : G_1/\ker(f) \rightarrow f(G_2)$ by $\phi(g \cdot \ker(f)) = f(g)$, for $g \in G_1$. We must show

(a) ϕ is well-defined,

(b) ϕ is a homomorphism,

(c) ϕ is a bijection.

If $g \cdot \ker(f) = g' \cdot \ker(f)$ then $g^{-1}g' \in \ker(f)$, since $\ker(f)$ is a group. This implies $f(g^{-1}g') \in f(\ker(f)) = \{1\}$, so $f(g) = f(g')$. This implies ϕ is well-defined.

Since $\ker(f)$ is normal, $(g \cdot \ker(f))(g' \cdot \ker(f)) = gg'(g'^{-1}\ker(f)g')\ker(f) = gg' \cdot \ker(f)$. Therefore $\phi((g \cdot \ker(f))(g' \cdot \ker(f))) = \phi(gg' \cdot \ker(f)) = f(gg') = f(g)f(g') = \phi(g \cdot \ker(f))\phi(g' \cdot \ker(f))$, for all $g, g' \in G$. This implies ϕ is a homomorphism.

It is clear that ϕ is surjective. To show that ϕ is a bijection, it suffices to prove ϕ is an injection. Suppose that $\phi(g \cdot \ker(f)) = \phi(g' \cdot \ker(f))$, for some $g, g' \in G$. Then $f(g) = f(g')$, so $f(g^{-1}g') = 1$. By definition of the kernel, this implies $g^{-1}g' \in \ker(f)$, so $g \cdot \ker(f) = g' \cdot \ker(f)$. This implies ϕ is injective.

□

The other isomorphism theorems will not be needed but will be stated to help to illustrate the usefulness of the notion of normality.

Theorem 9.5.2. (Second isomorphism theorem) *If H, N are subgroups of a group G and if N is normal then*

(a) $H \cap N$ is normal in H ,

(b) there is an isomorphism

$$H/(H \cap N) \cong NH/N.$$

Ponderable 9.5.1. *Let $n \geq 4$ be an even integer, $H = D_n \subset S_n$ and $N = A_n$. Use the second isomorphism theorem to show that $HN = S_n$.*

Here is an illustration using SAGE:

SAGE

```
sage: n = 6
sage: H = DihedralGroup(n)
sage: N = AlternatingGroup(n)
sage: gens = H.gens()+N.gens()
sage: G = PermutationGroup(gens)
sage: G == SymmetricGroup(n)
True
```

Theorem 9.5.3. (Third isomorphism theorem) *If N_1, N_2 are subgroups of a group G , if $N_1 \subset N_2$, and if N_1 and N_2 are both normal then*

(a) N_2/N_1 is normal in G/N_1 ,

(b) there is an isomorphism

$$(G/N_1)/(N_2/N_1) \cong G/N_2.$$

We shall not prove these results here; see [G] or [R]. These are due to Emmy Noether (1882-1935), one of the first women to obtain a Ph.D. in mathematics in Germany (or anywhere for that matter). Though extremely talented, the fact that she was Jewish and a woman made her progress in mathematics difficult. D. Hilbert and F. Klein, two of the world's leading mathematicians at the time, supported her and she conducted research and lectured at the University of Göttingen until the Nazis came to power in the 1930's. She left for Bryn Mawr College in the USA but died a few years later. She made many very important contributions to modern algebra.

9.6 Dabbling in direct products

Given two integers n_1 and n_2 , you can always form another one from them, using multiplication, denoted $n_1 n_2$. Analogously, given two groups H_1 and H_2 , you can always form another group from them using the Cartesian product construction (which we saw already in the context of sets in §2.1), denoted $H_1 \times H_2$.

Definition 9.6.1. *Let H_1, H_2 be two subgroups. We say that a group G is the direct product of H_1 with H_2 , written*

$$G = H_1 \times H_2,$$

if

- (a) $G = H_1 \times H_2$ (Cartesian product, as sets),
- (b) the group operation on G is given ‘coordinate-wise’ (still denoted ‘ \cdot ’ for simplicity):

$$(x_1, y_1) \cdot (x_2, y_2) = (x_1 \cdot x_2, y_1 \cdot y_2),$$

for $x_1, x_2 \in H_1$, $y_1, y_2 \in H_2$ (where \cdot denotes multiplication in H_1, H_2 , and G).

Example 9.6.1. *Let G be (as a set) the Cartesian product $G = C_2 \times C_3$, where C_n denotes the cyclic group of order n (with addition mod n as the operation, $n = 2, 3$). Define addition on G coordinate-wise (mod 2 in the first coordinate, mod 3 in the second coordinate):*

$$(m_1, n_1) + (m_2, n_2) = (m_1 + m_2, n_1 + n_2),$$

where $0 \leq m_i \leq 1$, $0 \leq n_j \leq 2$, for $i = 1, 2$, $j = 1, 2$.

You can use **SAGE** to construct this direct product as illustrated below. The direct product in **SAGE** constructs a 5-tuple: (1) the direct product $C_2 \times C_3$, (2) the injection $C_2 \rightarrow C_2 \times C_3$, (3) the injection $C_3 \rightarrow C_2 \times C_3$, (4) the projection $C_2 \times C_3 \rightarrow C_2$, (5) the projection $C_2 \times C_3 \rightarrow C_3$.

SAGE

```

sage: C2 = CyclicPermutationGroup(2)
sage: C3 = CyclicPermutationGroup(3)
sage: G = C2.direct_product(C3)
sage: G
(Permutation Group with generators [(1,2), (3,4,5)],
 Homomorphism : Cyclic group of order 2 as a permutation group
--> Permutation Group with generators [(1,2), (3,4,5)],
 Homomorphism : Cyclic group of order 3 as a permutation group
--> Permutation Group with generators [(1,2), (3,4,5)],
 Homomorphism : Permutation Group with generators [(1,2), (3,4,5)]
--> Cyclic group of order 2 as a permutation group,
 Homomorphism : Permutation Group with generators [(1,2), (3,4,5)]
--> Cyclic group of order 3 as a permutation group)
sage: G[0].order()
6

```

Example 9.6.2. *The symmetry group O of the octahedron is isomorphic to $S_4 \times C_2$. The symmetry group I of the icosahedron is isomorphic to $A_5 \times C_2$. (This is not isomorphic to S_5 , despite the fact that they both have the same number of elements and they both contain A_5 as a normal subgroup.)*

9.6.1 First fundamental theorem of cube theory

The following basic fact is implicitly used in some of the examples which follow. We shall frequently need the following fact.

Theorem 9.6.1. *Beginning with a solved cube, label the following facets with an invisible ‘+’ (i.e., mark the spatial position of the facet on the cube with a ‘+’):*

- *U facet of the uf edge subcube*
- *U facet of the ur edge subcube*
- *F facet of the fr edge subcube*
- *all facets which can be obtained from these by a move of the slice group.*

*Label the U and D facets of each corner subcube with an invisible ‘+’. These ‘+’ signs are called the **standard reference markings**. Each move g of the Rubik’s Cube yields a new collection of ‘+’ labels, called the **markings relative to g** . A position of the Rubik’s Cube is determined by the following decision process:*

- (a) *How are the edge subcubes permuted?*
- (b) *How are the center subcubes permuted?*
- (c) *How are the corner subcubes permuted?*

(d) Which of the relative edge markings are flipped (relative to the standard reference markings)?

(e) Which of the relative corner markings are rotated from the standard reference markings and, if so, by how much ($2\pi/3$ or $4\pi/3$ radians clockwise, relative to the standard reference markings)?

This is labeled as a theorem because of its relative importance for us, not because of its difficulty! This is the **First Fundamental Theorem of Rubik's Cube theory**. As an exercise, the reader should convince him or herself that this theorem is correct.

Whether or not the above theorem is 'obvious' or not reminds me of the following anecdote regarding the famous British mathematician G.H. Hardy, whom we met earlier in connection with Ramanujan. Hardy was known to be a bit of an eccentric character in his day. The story goes that he was giving a mathematics lecture and wrote down a theorem, then said: 'This is obvious.' After a pause, he went on: 'Hmm ... is it really obvious?' After thinking silently for a few moments he left the room to consider the point. He returned 20 minutes later, saying: 'Yes, it is obvious.' He then proceeded to finish his lecture. (The story is attributed to Prof. Jonathan R. Partington.) So, when I say the above theorem is 'obvious', don't worry if you have to think about it for 20 minutes. It might help to have a cube in front of you.

9.6.2 Example: cube twists and flips

We recall some notation:

- X is the set of 48 facets of the Rubik's Cube which are not center facets.
- V denotes the subset of facets which belong to some corner subcube.
- E is the subset of facets which belong to some edge subcube.
- Let G denote the Rubik's Cube group.
- Let F be the group generated by all the moves of the Rubik's Cube group which do not permute any corner or edge subcubes but may twist or flip them.
- Let S_X, S_V, S_E , denote the symmetric group on X, V, E , respectively. We may regard F, G , as subgroups of S_X . We may also regard S_V, S_E as subgroups of S_X (for example, S_V is the subgroup of S_X which leaves all the elements of E fixed).
- Let

$$G_V = S_V \cap G, \quad G_E = S_E \cap G, \quad F_V = S_V \cap F, \quad F_E = S_E \cap F.$$

Note that the action of G on X induces an equivalence relation as follows: we say that a facet f_1 is **equivalent** to a facet f_2 if there is a move of the Rubik's

Cube which sends one facet to the other. There are exactly two equivalence classes, or orbits, of G in X : namely, V and E . In particular, the action of G on V is transitive and the action of G on E is transitive. On the other hand, F leaves each vertex (respectively, edge) fixed, though it may permute the corner facets (respectively, edge facets) associated to a vertex (respectively, edge).

Ponderable 9.6.1. *Show that*

- (a) *the set \overline{V} of equivalence classes of F acting on V is in bijective correspondence with the set of all vertices of the cube;*
- (b) *the set \overline{E} of equivalence classes of F acting on E is in bijective correspondence with the set of all edges of the cube.*

The interesting thing is that we have

$$F = F_V \times F_E \quad (\text{direct product}).$$

Ponderable 9.6.2. *With notation as above,*

- (a) *show $F = F_V \times F_E$ (direct product);*
- (b) *is $S_X = S_V \times S_E$ (direct product)?*

A harder question, which we will answer later in the negative: Is $G = G_V \times G_E$ (direct product)?

9.6.3 Example: the slice group of the cube

It is worthwhile studying a subgroup of the Rubik's Cube group which is easier to analyze than the Rubik's Cube group itself. This way, some of the basic ideas and setup are introduced in a simpler setting.

Some of the material below can also be found in [BH], [Si].

Let H be the group $\langle M_R, M_F, M_U \rangle$ generated by the middle slice moves. This group is called the 'slice group'. Let E be the set of edges of the cube (which we identify with the set of edge subcubes), let C be the set of center facets of the cube, and let $X = E \cup C$. Note that H acts on X and that H does not move the elements of C (i.e., the corner vertices of the cube).

Ponderable 9.6.3.

- (a) *Is the action of H on X transitive?*
- (b) *Is the action of H on C transitive?*
- (c) *Is the action of H on E transitive?*

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Solution: The answer to (a) is no; for example, an edge subcube cannot be sent to a center facet. Consequently, there are elements of X which cannot be sent to each by an element of H . The answer to (b) is yes; any center facet can be sent to any other center facet by an element of H . The answer to (c) is no; for example, the uf edge subcube cannot be sent to the ur edge subcube by a slice move, so there is an element of E which cannot be sent to any other element of E by an element of H . This completes the solution.

The answer of ‘no’ to (c) brings about the following question.

Ponderable 9.6.4. *What are the orbits of H on E ?*

Solution: The answer may be phrased in various ways, but let us look at it in the following way: suppose we call two edge subcubes **equivalent** if one can be sent to the other by a slice move (i.e., an element of H). There are 3 disjoint equivalence classes: all the subcubes in the middle RL -slice are equivalent, all the subcubes in the middle FB -slice are equivalent, and all the subcubes in the middle UD -slice are equivalent. The distinct orbits of H acting on E are the following:

- the middle RL -slice, denoted by E_{RL} ,
- the middle FB -slice, denoted by E_{FB} ,
- the middle UD -slice, denoted by E_{UD} .

Note that

$$E = E_{RL} \cup E_{FB} \cup E_{UD},$$

is a partitioning of E into the distinct equivalence classes defined by the action of H on E . This completes the solution.

Each element of H determines an element in S_X . We have a homomorphism

$$f : H \rightarrow S_X.$$

This is another way of saying that H acts on the set X , which we already know. Note that each basic slice move M (so M is either M_R , M_F , or M_U) is, as an element of S_X , of the following form:

$$M = (4\text{-cycle in } S_E)(4\text{-cycle in } S_C).$$

Conversely, does an element of S_X uniquely determine an element of H ? In other words, is f injective (i.e., an embedding)?

To answer this, fix an $h \in H$ and think about what $f(h)$ tells us: $f(h)$ tells us which subcube moves to which other subcube but it doesn’t tell us, for example, how a subcube is flipped or rotated.

The first fundamental theorem 9.6.1 of the cube inspires the following question.

Ponderable 9.6.5. *Can an element of H flip, but not permute, an edge subcube (and possibly permuting or flipping other subcubes of the cube)?*

Solution: The answer is no. The reason is that the slice moves can only rotate a given edge subcube within the slice it belongs to. This completes the solution.

It follows, therefore, that the permutations of the edge subcube and centers determine a unique element of the slice group. In other words, we have proven the following

Proposition 9.6.1. *The homomorphism*

$$f : H \rightarrow S_X$$

is an embedding.

Ponderable 9.6.6. *The analog of this for the Rubik's Cube group is false! Why?*

H acts on the set E_{RL} , so we have a homomorphism

$$r_{RL} : H \rightarrow S_{E_{RL}}$$

and similarly, $r_{UD} : H \rightarrow S_{E_{UD}}$, $r_{FB} : H \rightarrow S_{E_{FB}}$.

H acts on each of the sets E and C , so we have homomorphisms

$$r = r_{RL} \times r_{UD} \times r_{FB} : H \rightarrow S_{E_{RL}} \times S_{E_{UD}} \times S_{E_{FB}} \subset S_E, \quad s : H \rightarrow S_C,$$

which we can put together to obtain an injective homomorphism

$$r \times s : H \rightarrow S_{E_{RL}} \times S_{E_{FB}} \times S_{E_{UD}} \times S_C.$$

To determine H , we determine the image of H in $S_{E_{RL}} \times S_{E_{FB}} \times S_{E_{UD}} \times S_C$.

To do this, we first look at the image of H in each of $S_{E_{RL}}$, $S_{E_{FB}}$, and $S_{E_{UD}}$. This is easy enough:

- the image of H in $S_{E_{RL}}$ is $\langle M_R \rangle \cong C_4$,
- the image of H in $S_{E_{FB}}$ is $\langle M_F \rangle \cong C_4$,
- the image of H in $S_{E_{UD}}$ is $\langle M_U \rangle \cong C_4$.

Later, we shall want to think of C_4 as $\{0, 1, 2, 3\}$, with addition mod 4, and the image of an element $h \in H$ under one of the homomorphisms above, $r_{RL} : H \rightarrow S_{E_{RL}}$, say, as an integer $0 \leq r_{RL}(h) \leq 3$.

Next, we must determine the image of H in S_C . This is easy if it's looked at in the right way. As far as the movements of the center facets are concerned, the slice moves may be replaced by their corresponding rotations of the entire cube about an axis of symmetry. In this case, we see that the image of H in S_C is the same as the image of the orientation-preserving symmetry group of the cube! This we know, by the discussion in Example 9.3.4 above, is isomorphic to S_4 .

Putting all this together, we see that the image of H in $S_{E_{RL}} \times S_{E_{FB}} \times S_{E_{UD}} \times S_C$ is isomorphic to a subgroup of

$$C_4^3 \times S_4.$$

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We may represent the elements of H , therefore, as 4-tuples (h_1, h_2, h_3, h_4) , with $h_1, h_2, h_3 \in C_4$ and $h_4 \in S_4$. Since each of the generating moves of H (namely, M_R , M_U , and M_F) satisfies

$$\text{sgn}(r(h)) = \text{sgn}(s(h)),$$

for all $h \in H$, the image of H cannot be all of $C_4^3 \times S_4$.

Proposition 9.6.2. *The image of H in $C_4^3 \times S_4$ is isomorphic to the kernel of the map*

$$\begin{aligned} t : C_4^3 \times S_4 &\rightarrow \{\pm 1\} \\ (h_1, h_2, h_3, h_4) &\mapsto \text{sgn}(h_1) \cdot \text{sgn}(h_2) \cdot \text{sgn}(h_3) \cdot \text{sgn}(h_4), \end{aligned}$$

where each sgn is the sign of the permutation, regarded as an element of S_X .

Ponderable 9.6.7. *Show that $|\ker(t)| = (4^3 \cdot 4!)/2 = 768$.*

Solution: We have shown that H is isomorphic to a subgroup of $C_4^3 \times S_4$. In fact, we know that the basic slice moves M_R, M_U, M_F (which generate H) all belong to the kernel of t , so H is isomorphic to a subgroup of $\ker(t) \subset C_4^3 \times S_4$.

It remains to show that every element in $\ker(t)$ belongs to H . To do this, we consider the projection homomorphism

$$p : H \rightarrow S_4$$

obtained by composing the homomorphism $r \times s : H \rightarrow C_4^3 \times S_4$ constructed above with the projection homomorphism $C_4^3 \times S_4 \rightarrow S_4$. We have shown that p is surjective. Our next objective is to compute the kernel of p and use the first isomorphism theorem to determine H .

Claim: The kernel of p is

$$\ker(p) = \{h \in H \mid s(h) = 1, r_{RL}(h) + r_{UD}(h) + r_{FB}(h) \equiv 0 \pmod{2}\}.$$

Note that $\ker(p)$ is a subgroup of H so the sign of the permutation $s(h)$ is equal to the sign of the permutation $r(h)$:

$$\text{sgn}(s(h)) = \text{sgn}(r(h)) = \text{sgn}(r_{RL}(h)) \cdot \text{sgn}(r_{UD}(h)) \cdot \text{sgn}(r_{FB}(h)).$$

This implies that

$$\ker(p) \subset \{h \in H \mid s(h) = 1, r_{RL}(h) + r_{UD}(h) + r_{FB}(h) \equiv 0 \pmod{2}\}.$$

Conversely, pick an $h \in H$ such that $s(h) = 1$ and $r_{RL}(h) + r_{UD}(h) + r_{FB}(h) \equiv 0 \pmod{2}$. We may represent this element h as a 4-tuple (n_1, n_2, n_3, s) , with $0 \leq n_1, n_2, n_3 \leq 3$ and $s = 1 \in S_4$.

For example,

- the element $M_1 = M_R \cdot M_F^{-1} \cdot M_D \cdot M_F$ is represented by the 4-tuple $(1, 1, 0, 1)$,

- the element $M_2 = M_R \cdot M_D \cdot M_F \cdot M_D^{-1}$ is represented by the 4-tuple $(0, 1, 1, 1)$,
- the element $M_3 = M_F \cdot M_D \cdot M_R^{-1} \cdot M_D^{-1} \cdot M_F \cdot M_D^{-1} \cdot M_R \cdot M_D$ is represented by the 4-tuple $(0, 0, 2, 1)$.

These elements generate all the elements of the group

$$\{(a, b, c, 1) \mid a, b, c \in C_4, a + b + c \equiv 0 \pmod{2}\}.$$

Note that the group

$$\{(a, b, c) \mid a, b, c \in C_4, a + b + c \equiv 0 \pmod{2}\}$$

is, in turn, the kernel of the map $C_4^3 \rightarrow C_2$ given by $(a, b, c) \mapsto a + b + c \equiv 0 \pmod{2}$.

Ponderable 9.6.8. *Show that $|\{(a, b, c) \mid a, b, c \in C_4, a + b + c \equiv 0 \pmod{2}\}| = 4^3/2 = 32$.*

Therefore, the element h chosen above must be expressible as a ‘word’ in these three elements M_1, M_2, M_3 . This shows that

$$\{h \in H \mid s(h) = 1, r_{RL}(h) + r_{UD}(h) + r_{FB}(h) \equiv 0 \pmod{2}\} \subset \ker(p),$$

which implies the claim.

To summarize what we have so far: we have a surjective homomorphism $p : H \rightarrow S_4$ with kernel $\{h \in H \mid s(h) = 1, r_{RL}(h) + r_{UD}(h) + r_{FB}(h) \equiv 0 \pmod{2}\}$. The kernel has $|\ker(p)| = 32$ elements and the image has $|\text{im}(p)| = |S_4| = 4! = 24$ elements. Since, by the first isomorphism theorem,

$$H/\ker(p) \cong S_4,$$

we have $|H| = 32 \cdot 24 = 768$. But the kernel of the homomorphism $t : C_4^3 \times S_4 \rightarrow \{\pm 1\}$, which we know contains H as a subgroup, also has 768 elements. This forces $H = \ker(t)$. This completes the solution.

9.6.4 Example: the slice group of the Megaminx

The subgroup S of the Megaminx group generated by all elements of the form $x \cdot y^{-1}$, where x, y correspond to opposite faces is called the **slice group** of the Megaminx.

Ponderable 9.6.9. *(Mark Longridge): Determine S .*

This group has recently (via an email sent Aug. 31, 2001) been completely determined by Mike Reid. It has

$$(15!/2) \cdot 4^{14} \cdot (10!/2) \cdot 6^9 \cdot 3 = 9627755206121277812101663948800000$$

elements.

9.7 A smorgasbord of semi-direct products

To ‘explicitly construct’ the Rubik’s Cube group, we need to know how to construct semi-direct products. We study semi-direct products in this section. A semi-direct product is a specific type of construction, more general than the direct product, of a new group, G , from two groups, H_1 and H_2 , having certain properties. It will turn out that H_1 and H_2 will be subgroups of G .

If a group G contains two subgroups H_1 and H_2 , with $H_1 \triangleleft G$ normal, such that each element of G can be written uniquely as a product $h_1 h_2$, with $h_1 \in H_1$ and $h_2 \in H_2$ then we say that G is the **semi-direct product** of H_1 and H_2 . In this situation, H_2 is called a **complement** of H_1 . In this section, we shall see another way how a semi-direct product can be expressed. Later, we shall see that the Rubik’s Cube group can easily be described using semi-direct products.

Definition 9.7.1. *Now suppose that H_1, H_2 are both subgroups of a group G . We say that G is the **semi-direct product** of H_1 by H_2 , written*

$$G = H_1 \rtimes H_2,$$

if

- (a) $G = H_1 \cdot H_2$,
- (b) H_1 and H_2 only have 1, the identity of G , in common,
- (c) H_1 is normal in G .

This is the ‘internal version’ of the semi-direct product.

Of course, if we define anything using two apparently different definitions, we’d better be sure that they are equivalent! The fact that they are is a theorem (Theorem 7.23 in [R]) which we won’t prove here.

Note that the multiplication rule in G doesn’t have to be mentioned since we are assuming here that G is given.

The ‘external version’ is defined by a construction as follows:

Definition 9.7.2. *Assume we have a homomorphism*

$$\phi : H_2 \rightarrow \text{Aut}(H_1).$$

Define multiplication on the set $H_1 \times H_2$ by

$$(x_1, y_1) \cdot (x_2, y_2) = (x_1 \cdot \phi(y_1)(x_2), y_1 \cdot y_2).$$

This defines a group operation. This group, denoted $H_1 \rtimes_{\phi} H_2$, is the (external) **semi-direct product**.

This definition will be used with the example of $H_1 = C_m^n$ and $H_2 = S_n$ in the next chapter.

These last two definitions are equivalent by Theorems 7.22-7.23 in [R].

As a set, $H_1 \rtimes H_2$ is simply the Cartesian product $H_1 \times H_2$.

Example 9.7.1. Let \mathbb{R}^2 denote the direct product of the additive group of real numbers with itself:

$$\mathbb{R}^2 = \{(x, y) \mid x, y \text{ real}\},$$

with the group operation being addition performed componentwise. Let C_2 denote the multiplicative cyclic group with 2 elements, whose elements we write (somewhat abstractly) as $C_2 = \{1, s\}$. (We may think of s as being equal to -1 but there is a reason for this notation which will be made clear soon.) Define an action of C_2 on \mathbb{R}^2 by

$$1(x, y) = (x, y), \quad s(x, y) = (y, x), \quad (x, y) \in \mathbb{R}^2.$$

Let G be the set

$$G = \mathbb{R}^2 \times C_2.$$

Define the binary operation $\cdot : G \times G \rightarrow G$ by

$$(g_1, z_1) \cdot (g_2, z_2) = (g_1 + z_1(g_2), z_1 \cdot z_2),$$

for all $g_1, g_2 \in G$ and all $z_1, z_2 \in C_2$. This is a group—the semi-direct product of \mathbb{R}^2 with C_2 .

To see this, we must answer some questions:

- closed under the operation? Yes
- existence of identity? Yes, $e = ((0, 0), 1)$
- existence of inverse? Yes, $((x, y), 1)^{-1} = ((-x, -y), 1)$, and $((x, y), s)^{-1} = ((-y, -x), s)$
- associative? This is the hard one:

$$\begin{aligned} ((g_1, z_1) \cdot (g_2, z_2)) \cdot (g_3, z_3) &= (g_1 + z_1(g_2), z_1 \cdot z_2) \cdot (g_3, z_3) \\ &= (g_1 + z_1(g_2) + (z_1 \cdot z_2)(g_3), (z_1 \cdot z_2) \cdot z_3) \\ (g_1, z_1) \cdot ((g_2, z_2) \cdot (g_3, z_3)) &= (g_1, z_1) \cdot (g_2 + z_2(g_3), z_2 \cdot z_3) \\ &= (g_1 + z_1(g_2 + z_2(g_3)), z_1 \cdot (z_2 \cdot z_3)) \end{aligned}$$

This implies associativity.

Ponderable 9.7.1. Let G be the semi-direct product constructed in the above example. Show that $H = \{((x, y), 1) \mid (x, y) \in \mathbb{R}^2\}$ is a normal subgroup of G . Show $H \cong \mathbb{R}^2$.

Example 9.7.2. Let

$$S_3 = \{1, s_1, s_2, s_1 \cdot s_2, s_2 \cdot s_1, s_1 \cdot s_2 \cdot s_1\}, \quad H_1 = \{1, s_2, s_1 \cdot s_2 \cdot s_1\}, \quad H_2 = \{1, s_1\}.$$

Let

$$\phi : H_2 \rightarrow \text{Aut}(H_1)$$

be defined by

$$\begin{aligned} \phi(1) &= 1 \quad (\text{the identity automorphism}) \\ \phi(s_1)(h) &= s_1^{-1} \cdot h \cdot s_1 = s_1 \cdot h \cdot s_1 \end{aligned}$$

(since $s_1^{-1} = s_1$), $h \in H_1$.

Denote the (external) semi-direct product of H_1, H_2 by $G = H_1 \rtimes_{\phi} H_2$.

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Ponderable 9.7.2. Let G be the semi-direct product constructed in the above example. Show that H_1 (which we identify with the subgroup $\{(h, 1) \mid h \in H_1\}$ of G) is a normal subgroup.

There is of course a close relationship between internally defined semi-direct products and externally defined ones. The following lemma, which is proven in [R], explains this connection:

Lemma 9.7.1. If G is the (internal) semi-direct product of H_1 by H_2 (so H_1 is a normal subgroup of G) then there is a homomorphism

$$\phi : H_2 \rightarrow \text{Aut}(H_1)$$

such that $G \cong H_1 \rtimes_{\phi} (H_2)$.

Example 9.7.3. Let C_d be the cyclic group of order d , which we may regard as a set as $C_d = \{0, 1, \dots, d-1\}$, with addition performed mod d . Let $N = C_d^n$, which we regard as the group of n -vectors with ‘coefficients in C_d ’. Let $H = S_n$ be the symmetric group on n letters, i.e., the group of all permutations

$$p : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}.$$

The group H acts on N by permuting the indices, i.e., the coordinates of the vectors. For $f \in S_n$, define $f^* : C_d^n \rightarrow C_d^n$ by

$$f^*(v) = (x_{f^{-1}(1)}, \dots, x_{f^{-1}(n)}) = P(f^{-1})v^t, \quad v = (x_1, \dots, x_n) \in C_d^n, \quad (9.2)$$

where $P(g)$ is the permutation matrix associated to a permutation g and v^t denotes the transpose (a column vector) of v . Now, for $p, q \in S_n$ and $v, w \in C_d^n$, define

$$(p, v) \cdot (q, w) = (pq, w + q^*(v)).$$

This defines a semi-direct product $C_d^n \rtimes S_n$, the **generalized symmetric group**. A **monomial matrix** is a matrix which contains exactly one non-zero entry for each row and column. (It is known that $C_d^n \rtimes S_n$ is isomorphic to the group of all ‘ C_n -valued $n \times n$ monomial matrices’, [R], Exercise 7.33.)

Here is an example of how to construct $C_3^8 \rtimes S_8$ using SAGE:

SAGE

```
sage: C3 = CyclicPermutationGroup(3)
sage: S8 = SymmetricGroup(8)
sage: W = C3._gap_().WreathProduct(S8._gap_())
sage: W.Order()
264539520
sage: W.Order() == 3^8*factorial(8)
True
```

9.8 A reification of wreath products

We briefly survey the basic facts about wreath products. Wreath products are generalizations of semi-direct products. They occur naturally in the theory of the Rubik's Cube group (see [Si], for example). I shall be brief because the wreath products which do occur in the theory of the Rubik's Cube may be, if desired, reformulated in terms of semi-direct products. In fact, it suffices to consider generalized symmetric groups (see Theorem 9.8.1 below), which are even easier to deal with and themselves special types of semi-direct products.

Let G_1 be a group, and let G_2 be a group acting on a finite set X_2 . Fix some labeling of X_2 as say $X_2 = \{h_1, h_2, \dots, h_m\}$, where $m = |X_2|$ and let $G_1^{X_2}$ denote the direct product of G_1 with itself m times, with the coordinates labeled by the elements of X_2 .

Definition 9.8.1. *The wreath product of G_1, G_2 is the group*

$$G_1 \text{ wr } G_2 = G_1^{X_2} \rtimes G_2,$$

where the action of G_2 on $G_1^{X_2}$ is via its action on X_2 .

In particular, to each $t \in G_1 \text{ wr } G_2$ there is a $g_2 \in G_2$. We denote this **projection** by $g_2 = pr(t)$. Define the **base** of the wreath product by

$$B = \{t \in G_1 \text{ wr } G_2 \mid pr(t) = 1\},$$

so $B = G_1^{X_2}$.

Example 9.8.1. *Let \mathbb{R}^n denote the direct product of the additive group of reals with itself n times. The group operation on \mathbb{R}^n is componentwise addition. Let S_n denote the symmetric group. This acts on \mathbb{R}^n by permuting coordinates as in (9.2). This action respects addition:*

$$r^*(x_1 + y_1, \dots, x_n + y_n) = r^*(x_1, \dots, x_n) + r^*(y_1, \dots, y_n),$$

$(x_1, \dots, x_n), (y_1, \dots, y_n) \in \mathbb{R}^n$. (Incidentally, it also preserves scalar multiplication:

$$r^*(a \cdot \vec{v}) = a \cdot r^*\vec{v},$$

for $\vec{v} \in \mathbb{R}^n$ and $a \in \mathbb{R}$. Therefore, r defines an invertible linear transformation on \mathbb{R}^n . In fact, there is a homomorphism $S_n \rightarrow \text{Aut}(\mathbb{R}^n)$, where $\text{Aut}(\mathbb{R}^n)$ denotes the group of all invertible linear transformations on \mathbb{R}^n . Do you recognize this homomorphism? It has occurred previously.)

Let G be the set

$$G = \mathbb{R}^n \times S_n$$

and define a binary operation $\cdot : G \times G \rightarrow G$ by

$$(\vec{v}_1, p_1) \cdot (\vec{v}_2, p_2) = (\vec{v}_1 + p_1^*(\vec{v}_2), p_1 \cdot p_2),$$

for all $\vec{v}_1, \vec{v}_2 \in \mathbb{R}^n$ and all $p_1, p_2 \in S_n$. This is a group. To see this, we must answer some questions:

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- Is G closed under the operation \cdot ? Yes.
- Does there exist an identity? Yes, $e = (\vec{0}, 1)$.
- Does there exist an inverse? Yes, $(\vec{v}, p)^{-1} = (-(p^{-1})^* \vec{v}, p^{-1})$.
- Is G associative? This is the hard one:

$$\begin{aligned} ((\vec{v}_1, p_1) \cdot (\vec{v}_2, p_2)) \cdot (\vec{v}_3, p_3) &= (\vec{v}_1 + p_1^*(\vec{v}_2), p_1 \cdot p_2) \cdot (\vec{v}_3, p_3) \\ &= (\vec{v}_1 + p_1^*(\vec{v}_2) + (p_1 \cdot p_2)^*(\vec{v}_3), (p_1 \cdot p_2) \cdot p_3) \\ (\vec{v}_1, p_1) \cdot ((\vec{v}_2, p_2) \cdot (\vec{v}_3, p_3)) &= (\vec{v}_1, p_1) \cdot (\vec{v}_2 + p_2^*(\vec{v}_3), p_2 \cdot p_3) \\ &= (\vec{v}_1 + p_1^*(\vec{v}_2 + p_2^*(\vec{v}_3)), p_1 \cdot (p_2 \cdot p_3)) \end{aligned}$$

This implies associativity.

This group is the wreath product of \mathbb{R} with S_n ,

$$G = \mathbb{R} \text{ wr } S_n,$$

where \mathbb{R}^n is the base.

Lemma 9.8.1.

- (a) The base B , which is isomorphic to the direct product $G_1^{|X_2|}$, is a normal subgroup of $G_1 \text{ wr } G_2$.
- (b) $(G_1 \text{ wr } G_2)/B$ is isomorphic to G_2 .

This is a consequence of the first isomorphism theorem of group theory. The proof is omitted. (Refer to chapter 8 of [NST] or to [R] for further details.)

9.8.1 The illegal Rubik's Cube group

Let H be the **illegal Rubik's Cube group** of all legal and illegal moves of the 3×3 Rubik's Cube. In other words, in addition to the usual basic moves (namely, R, L, U, D, F, B), we allow you to take apart the cube and reassemble the corner and edge subcubes (but you cannot remove stickers from the facets). Let C_3 denote the group of all rotations of a particular corner subcube by a 120 degree angle. (Actually, this group depends on the corner being rotated, but since these groups are all isomorphic, we drop the dependence from the notation.) Let C_2 denote the group of all flips of a particular edge subcube. (Again, this group depends on the edge being flipped but since these groups are all isomorphic, we drop the dependence from the notation.) We shall prove later that

$$H = (C_3 \text{ wr } S_V) \times (C_2 \text{ wr } S_E),$$

where V is the set of corner subcubes and E is the set of edge subcubes.

9.8.2 Elements of order d in $C_m \wr S_n$

In this section, we will apply our understanding of wreath products to increasing our ability to determine all the elements of order d in the Rubik's Cube group. (The case $d = 2$ will be examined in some detail later.) In some cases (e.g., in the cases we deal with here), wreath products turn out to be relatively concrete and familiar groups.

Let $S(n, m)$ denote the group of all $n \times n$ monomial matrices with entries in C_m . We begin with the following result, which allows us to identify $S(n, m)$ with the generalized symmetric group in Example 9.7.3.

Theorem 9.8.1. *There is an isomorphism between $C_m \wr S_n$ and the group $S(n, m)$ which sends an element $(\vec{v}, f) \in C_m \wr S_n$, $\vec{v} = (v_1, v_2, \dots, v_n) \in C_m^n$ and $f \in S_n$, to the vector $P(f)\text{diag}(\vec{v})$.*

This follows from a result (see Theorem 10.5.1) which will be proven in the next chapter. It is also a special case of an exercise in [R].

Example 9.8.2. *If $n = 3$ and $m = 2$ then*

- $((-1, 1, -1), (1, 2))$ corresponds to

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} -1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & -1 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 0 \\ -1 & 0 & 0 \\ 0 & 0 & -1 \end{pmatrix};$$

- $((1, -1, 1), (1, 2))$ corresponds to

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{pmatrix} = \begin{pmatrix} 0 & -1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

Here is an example of how to use **SAGE** to construct a group multiplication in this wreath product.

```

SAGE

sage: P = Matrix([[0,1,0], [1,0,0], [0,0,1]])
sage: v = Matrix([[-1,0,0], [0,1,0], [0,0,-1]])
sage: A = P*v
sage: A

[ 0  1  0]
[-1  0  0]
[ 0  0 -1]
sage: w = Matrix([[1,0,0], [0,-1,0], [0,0,1]])
sage: B = P*w
sage: B

[ 0 -1  0]
[ 1  0  0]
[ 0  0  1]
```

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```
sage: A*B
[ 1  0  0]
[ 0  1  0]
[ 0  0 -1]
sage: v*P*w*P^(-1)
[ 1  0  0]
[ 0  1  0]
[ 0  0 -1]
```

Let $\vec{v} = (v_1, v_2, \dots, v_n) \in C_m^n$ and $f \in S_n$. It is clear from this theorem that an element $(\vec{v}, f) \in C_m \wr S_n$ is order d only if the permutation matrix $P(f)$ is order d . Indeed,

$$\begin{aligned} (\vec{v}, f)^2 &= (\vec{v} + f^*(\vec{v}), f^2), \\ (\vec{v}, f)^3 &= (\vec{v} + f^*(\vec{v}) + (f^*)^2(\vec{v}), f^3), \\ &\vdots \\ (\vec{v}, f)^k &= (\vec{v} + \dots + (f^*)^{k-1}(\vec{v}), f^k). \end{aligned}$$

We conclude with the following classification of the elements of order d in the wreath product.

Proposition 9.8.1. *An element $(\vec{v}, f) \in C_m \wr S_n$, $\vec{v} = (v_1, v_2, \dots, v_n) \in C_m^n$ and $f \in S_n$, is order d if and only if $f^d = 1$ and $\vec{v} + \dots + (f^*)^{d-1}(\vec{v}) = 0$, where f^* is defined in (9.2) on p. 194.*

This result can, in principle, be used in conjunction with the explicit determination of the Rubik's Cube group given later to determine all the elements of a given order. See §11.3.

Chapter 10

Words which move

Everyone knows what a curve is, until he has studied enough mathematics to become confused through the countless number of possible exceptions.

Felix Klein

Felix Klein (1849–1925), one of the most important figures in group theory, was right on the money with his comment about curves. However, as far as this chapter is concerned, we might replace ‘curve’ with ‘word’. We originally used ‘words’ and ‘relations’ to describe groups arising in the topological study of surfaces. This chapter touches on a broad area of group theory called ‘combinatorial group theory’, a subject having its origins in the late 1800’s and early 1900’s. Henri Poincaré (1854–1912), one of the greatest mathematical geniuses of all time, introduced the ‘fundamental group’ of a surface, which turns out to be naturally defined in terms of ‘generators’ (the ‘letters’ of our ‘words’) and ‘relations’ (‘grammatical rules’ our ‘words’ satisfy). Although topology is well outside the scope of this book, it is interesting to note how the study of one area can ultimately lead to applications of something completely different.

Let us begin with an example. *RUBDUBDUB* is both a word (well, sort of, anyway) and a move of the Rubik’s Cube. Is that move different than *DUBRUBRUB*? Just because they look different doesn’t mean that they are. (R^4 and D^4 look different but they both have the same effect as doing nothing at all!) Basically, (one formulation of) the ‘word problem’ is the problem of determining when two words (i.e., strings of generators of a group) are the same or not in a given group.

The word problem is related to the solution of the Rubik’s Cube since if we knew which ‘words’ in the generators R, L, U, D, F, B , described a given position of the Rubik’s Cube, then we could use one of them to solve the position. The problem is that this ‘word problem’ is usually quite hard to solve in practice. It is remarkable that there are GAP commands (and also in **SAGE**) which can always find a word in the generators R, L, U, D, F, B which describes a given position of the Rubik’s Cube. These commands use no knowledge of the Rubik’s

Cube beyond the definition of the Rubik's Cube group, yet allow one to solve the cube using group theory. It is the purpose of this chapter to introduce some of the mathematical ideas which arise when trying to study the Rubik's Cube group from this perspective.

Roughly speaking, this cube-solving algorithm does the following:

- (a) constructs the free group on the 6 face move generators and then computes a reasonable set of relations which they satisfy;
- (b) computes a homomorphism from the cube group to this free group quotient (we discuss such quotients in this chapter);
- (c) takes the cube position, regarded as a group element, and maps it over to the free group quotient;
- (d) using those relations, and tricks from combinatorial group theory (stabilizer chains, which go beyond the level of this book), solves the 'word problem' for that element;
- (e) uses Python string parsing functions to rewrite that in cube notation.

The Rubik's cube group has about 4.3×10^{19} elements, so this process is time-consuming, even for simple moves.

We give an example using **SAGE** below, but see [Sch] at the GAP website for an interesting discussion of some GAP code analyzing the Rubik's Cube.

10.1 Words in free groups

To be precise and make sure everyone's on the same page, let's record a few definitions and introduce some notation.

Definition 10.1.1. *Given a list L of questions, a **decision algorithm** for L is a uniform set of unambiguous instructions which, when applied to any question in L , gives the correct answer 'yes' or 'no' after a finite number of steps.*

Let $X = \{x_1, \dots, x_n\}$ denote a set and X^{-1} a set disjoint from X whose elements we denote by $\{x_1^{-1}, \dots, x_n^{-1}\}$. Assume that the map $x \mapsto x^{-1}$ defines a bijection $X \rightarrow X^{-1}$. It will be convenient to let $x^1 = x, x_i^0 = 1$, where 1 is an element not belonging to $X \cup X^{-1}$ which we will call the **identity element**. A **word on X** is a sequence

$$w = (a_1, a_2, \dots, a_N),$$

where $N > 0$ is some integer and each a_i belongs to

$$X \cup X^{-1} \cup \{1\}.$$

The sequence of all 1's is called the **empty word**. The **inverse** of the word w is the word

$$w^{-1} = (a_N^{-1}, \dots, a_1^{-1}).$$

If $a_i = y_i^{e_i}$, where e_i is in $\{0, 1, -1\}$ and $y_i \in X$, then we shall write the word w as $w = y_1^{e_1} \dots y_N^{e_N}$.

Example 10.1.1. Let $X = \{R, L, U, D, F, B\}$. The set of words on X are in a bijective correspondence with the set of sequences of basic moves you can make on the Rubik's Cube.

We call a word $w = y_1^{e_1} \dots y_N^{e_N}$ on X **reduced** if either w is empty or if the exponents e_i are non-zero and if there are no $x \in X$ with x, x^{-1} adjacent in w .

Definition 10.1.2. The **free group** $F_n = F_X$ on the generators x_1, \dots, x_n is the group of all reduced words on X .

The notion of a free group was first formulated by Walther von Dyck (1856-1934). Von Dyck made important contributions to topology, function theory, and group theory.

The proof that F_n is a group is not entirely easy (verifying the associativity property is perhaps the hardest part); see Theorem 11.1 in [R].

10.1.1 Length

If, in the notation above, $w = y_1^{e_1} \dots y_N^{e_N}$ is a reduced word (so $e_i \in \{1, -1\}$) then we call N the **length** (or **reduced length**, to be more precise) of w .

If $G = \langle g_1, \dots, g_k \rangle \subset S_n$ is a finite permutation group generated by permutations g_1, \dots, g_k , then we may still define the notion of length:

Definition 10.1.3. Suppose $g \in G$ is not the identity, where G is a permutation group as above. Then g may be written

$$g = y_1^{e_1} \dots y_N^{e_N},$$

where each $y_i \in \{g_1, \dots, g_k\}$ and where $e_i \in \{0, 1, -1\}$. The number N and the sets $\{y_1, \dots, y_N\}$, $\{e_1, \dots, e_N\}$ may not be unique for a given g but among all such possibilities there is at least one such that the value of N is minimum. We call this the **length** of g , denoted $\ell(g)$. In other words, the length of g is the size of a “minimal” decomposition of g as a product of transpositions.

Let

$$P_G(t) = \sum_{g \in G} t^{\ell(g)}.$$

This is called the **Poincaré polynomial** of G .

The length of g is the distance in the Cayley graph between the vertex g and the vertex 1. As was mentioned in chapter 7, the problem of determining the largest possible distance in the Cayley graph of the Rubik's Cube group is known as ‘God's algorithm’ and is currently unsolved.

Example 10.1.2. Let $G = S_n$ with generators $g_i = (i, i+1)$, $i = 1, \dots, n-1$. The Poincaré polynomial is known:

$$\prod_{k=1}^{n-1} \frac{t^{k+1} - 1}{t - 1},$$

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by [Hum], §§1.11, 3.15.

In case $n = 3$, this is

$$t^3 + 2t^2 + 2t + 1 = (t + 1)(t^2 + t + 1)$$

Of course, this is very easy to compute using SAGE:

SAGE

```
sage: p = lambda n: \
        prod([(t^(j+1)-1)/(t-1) for j in range(1,n+1)])
sage: p(2)
t^3 + 2*t^2 + 2*t + 1
```

If you define the “swapping number” as in Definition 3.1.1 then you can show that the swapping number of an element $g \in S_n$ is the same as its length with respect to the generators $s_i = (i, i + 1)$, $i = 1, \dots, n - 1$.

Using the code in Example 3.1.2, we can compute the Poincaré polynomial in another way:

SAGE

```
sage: G = SymmetricGroup(3)
sage: p = sum([t^swap(g) for g in G])
sage: p
t^3 + 2*t^2 + 2*t + 1
```

Ponderable 10.1.1. What is the longest possible length of an element of the Rubik’s Cube group (with respect to the generators R, L, U, D, F, B)?

Note this is the same as asking for longest possible length of an element of the Rubik’s Cube group with respect to the generators

$$R, L, U, D, F, B, R^{-1}, L^{-1}, U^{-1}, D^{-1}, F^{-1}, B^{-1}.$$

In other words, it is simply a reformulation of ‘God’s algorithm’ in the quarter-turn metric.

Ponderable 10.1.2. What is the Poincaré polynomial of the Rubik’s Cube group (with respect to R, L, U, D, F, B)?

This is unknown, though I refer the reader interested in pursuing this to Thompson’s paper [Th].

10.1.2 Trees

We may represent the free group graphically as follows. We define the Cayley graph of F_X inductively:

- Draw a vertex for each element of $X \cup X^{-1}$ (these are the vertices, V_1 say, for the words of length 1).
- Suppose that you have already drawn all the vertices for the words of length $k - 1$, V_{k-1} let's call them. For each $x \in X \cup X^{-1}$ and each $v \in V_{k-1}$, draw a vertex for each word of length k obtained by multiplying v by x on the right, $v \cdot x$, and connect v and $v \cdot x$ by an edge.

There are infinitely many vertices, each of which has degree $|X \cup X^{-1}|$. Moreover, this graph has no circuits or loops (i.e., no path of edges crosses back over onto itself). Such a graph is called a **tree**.

Example 10.1.3. Let $X = \{R, L, U, D, F, B\}$. The elements of the free group F_X correspond to the mechanically different sequences of basic moves you can make on the Rubik's Cube that do not contain a basic move followed by its inverse. Of course, different sequences of moves may yield the same position of the Rubik's Cube (e.g., R^4 and 1 are the same position, but the sequence of moves used to attain them are distinct).

There are infinitely many vertices of the Cayley graph of F_X , each of which corresponds to a mechanically distinct move of the Rubik's Cube.

Further details on the background material given here may be found in [R], [MKS].

10.2 The word problem

There is a way to list all the elements of F_n , called the **lexicographic ordering**. I shall describe a way to list all the words in F_n as though they were in a dictionary. I shall give an algorithm for determining if a word $w \in F_n$ occurs before a word w' , in which case we write $w < w'$. In the dictionary below, we shall, for example, distinguish between the identity 1 and the 'non-reduced' word $x_1 \cdot x_1^{-1}$.

The first element in this lexicographically ordered list is the word 1, and the next $2n$ words are the words

$$x_1 < x_1^{-1} < \dots < x_n < x_n^{-1}.$$

In general, we define $y_1 \dots y_M < z_1 \dots z_N$ if either (a) $M < N$, or (b) $M = N$ and $y_1 < z_1$, or (c) $y_1 = z_1$ and $y_2 < z_2$, or (d) $y_1 = z_1$ and $y_2 = z_2$ and $y_3 < z_3$, or ...

List all the elements of F_n as

$$F_n = \{w_1, w_2, \dots\},$$

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so $w_1 = 1, w_2 = x_1, \dots$. Let G be a subgroup of F_n or a permutation group $G = \langle g_1, \dots, g_n \rangle$. If G is a permutation group then we regard a word w_k as an element of G by substituting g_i for each $x_i, 1 \leq i \leq n$.

Definition 10.2.1. Fix a $g \in G$. We say that G has a **solvable word problem** if there is a decision algorithm for the list L of questions of the form: Is $w_k = g$ in G ?

The word problem was first formulated by Max Dehn (1878–1952), a topologist who worked extensively on group theory. He was a student of David Hilbert and wrote several important papers. However, after he emigrated to the USA in 1940 Dehn had a hard time finding a good teaching position. He taught at several colleges in the States, including St. Johns University in Annapolis, Maryland, and then settled down at Black Mountain College, a small school in North Carolina. There was no mathematics department, yet he remained there until his death. The college itself closed several years later.

Lemma 10.2.1. ([R], chapter 12) F_n has a solvable word problem.

Proof: (Sketch) The following decision algorithm yields a solvable word problem.

1. If w is a word not equal to 1, then underline the first occurrence, if any, of the expression $x_i \cdot x_i^{-1}$. If no such expression occurs in w , then go to step 3; otherwise, go to step 2.
2. Delete the expression $x_i \cdot x_i^{-1}$ from w and go to step 1.
3. If $w = 1$ then stop and return ‘yes’; otherwise, stop and return ‘no’.

□

Remark 10.2.1. There are examples of groups G for which there is no decision algorithm to solve the word problem. We shall not discuss them here (see [R], chapter 12, for more details). Sometimes this is referred to in the literature by saying ‘the word problem for G is undecidable’.

There is another use of the word ‘undecidable’ in mathematics. For example, Georg Cantor (1845–1918), a brilliant German mathematician who suffered from mental illness during his last years, is considered by some to be the father of modern set theory. Using Cantor’s ‘ordinal arithmetic’, we know that the ‘cardinality’ of the real numbers \mathbb{R} is not the same as the ‘cardinality’ of the rational numbers \mathbb{Q} . Cantor conjectured, but could not prove, the Continuum Hypothesis: ‘there are no cardinal numbers between the cardinality of \mathbb{Q} and of \mathbb{R} ’. On the other hand, it has been shown that the Continuum Hypothesis cannot be proven from the axioms of set theory and that its truth or falsity is consistent with the axioms of set theory. This was proven by Paul Cohen (1934–2007), an American mathematician who won the Fields medal in 1966 (the ‘Nobel Prize’ of mathematics) for his solution. Paul Cohen’s result is sometimes referred to in the literature as saying ‘the Continuum Hypothesis is undecidable.’ These two uses of the word ‘undecidable’ are very different, so be careful!

Lemma 10.2.2. *Each finite permutation group has a solvable word problem.*

Ponderable 10.2.1. *Why?*

Theorem 10.2.1. *A decision algorithm for the word problem for the Rubik's Cube group with generators R, L, U, D, F, B is the same as an algorithm for solving the Rubik's Cube.*

In fact, there is an algorithm for writing each element in the Rubik's Cube group (or more generally any permutation group G) as a word in the generators R, L, U, D, F, B (or more generally the given generators of G). This has been implemented in GAP [Gap]. Though implementing this algorithm is a remarkable piece of work, it should be pointed out that the words one obtains in this way from GAP are often of length 60 or more.

Example 10.2.1. *In this example, we illustrate the syntax of SAGE's cube solver command. To use it, we must enter the state of the Rubik's cube as a permutation group element. This is easiest if we simply scramble the Cube and see if SAGE (indirectly using GAP, in this case) can unscramble it. To keep things simple, let's scramble the cube with a simple move like $R \cdot U \cdot R^{-1}$:*

SAGE

```
sage: rubik = CubeGroup()
sage: b = rubik.B()
sage: d = rubik.D()
sage: f = rubik.F()
sage: l = rubik.L()
sage: r = rubik.R()
sage: u = rubik.U()
sage: r*u*r^(-1)
(1,19,24,6)(2,21,7,4)(8,30,17,9)(10,34,28,18)(11,35,25,43)
sage: G = PermutationGroup([b,d,f,l,r,u])
sage: g = G(str(r*u*r^(-1)))
sage: g.word_problem([b,d,f,l,r,u])
x5*x6*x5^-1
[[ '(25,27,32,30)(26,29,31,28)(3,38,43,19)(5,36,45,21)
  *(8,33,48,24)', 1],
  ['(1,3,8,6)(2,5,7,4)(9,33,25,17)(10,34,26,18)
  *(11,35,27,19)', 1],
  ['(25,27,32,30)(26,29,31,28)(3,38,43,19)(5,36,45,21)
  *(8,33,48,24)', -1]]
'(25,27,32,30)(26,29,31,28)(3,38,43,19)(5,36,45,21)
*(8,33,48,24)(1,3,8,6)(2,5,7,4)(9,33,25,17)(10,34,26,18)
*(11,35,27,19)(25,27,32,30)(26,29,31,28)(3,38,43,19)
*(5,36,45,21)(8,33,48,24)^-1'
```

(For typographical reasons, the above SAGE session had to be slightly reformat-
ted.) The 'solution' is printed as the first line of output after the word_problem

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command: $x_5 * x_6 * x_5^{-1}$. Note that **SAGE** rewrites the ordered list of generators b, d, f, l, r, u as $x_1, x_2, x_3, x_4, x_5, x_6$. This implies that **SAGE** says to unscramble the cube, make the inverse of the move $ru r^{-1}$, which we know to be correct.

The `word_problem` command works with any permutation group. The Rubik's Cube group has a specialized `solve` command that has a more familiar-looking output.

— SAGE —

```
sage: rubik = CubeGroup()
sage: state = rubik.faces("R*U")
sage: rubik.solve(state)
'R U'
sage: state = rubik.faces("R*U*R^(-1)")
sage: state
{'back': [[33, 28, 25], [36, 0, 37], [38, 39, 40]],
 'down': [[41, 42, 11], [44, 0, 45], [46, 47, 48]],
 'front': [[9, 10, 24], [20, 0, 7], [22, 23, 6]],
 'left': [[8, 34, 35], [12, 0, 13], [14, 15, 16]],
 'right': [[43, 26, 27], [18, 0, 29], [17, 31, 32]],
 'up': [[19, 21, 3], [2, 0, 5], [1, 4, 30]]}
sage: rubik.solve(state)
'R U R'
```

You can apply `solve` to any legal state of the Rubik's Cube. The easiest way to construct an example of a cube state is to actually make a legal move and read off its state using the `faces` method, as shown above. This particular solver is rather slow (it was explained in the introduction to this chapter) and has an impractical interface, being designed more for pedagogical reasons than to actually solve cubes.

Here is another way to use **SAGE** to solve the cube.

— SAGE —

```
sage: C = RubiksCube("R U F L D")
sage: C.solve(algorithm = 'kociemba')
'R U F L D'
```

This uses an implementation by Dik T. Winter of an algorithm of Herbert Kociemba. There are several other Cube solvers included with **SAGE**. Please see the **SAGE** reference manual for more details.

10.3 Presentations and Plutonian robots

Here's a hypothetical situation: You and a friend each have a robot on the planet Pluto with a scrambled Rubik's Cube. (We will call the robots R^2D^2 and R^2B^2 if you don't mind.) These robots have manual dexterity but no preprogramming on how to solve the cube. Furthermore, assume it is very expensive to program different moves, so you want to teach R^2D^2 and R^2B^2 the smallest number of separate moves that you can. On the other hand, the moves need not be basic moves (U , R , ...) since we will assume it costs roughly the same to teach the robot the move R as the longer move $R \cdot U^2 \cdot R^{-1}$, for example. Your solution will be a 'word' in these taught moves. Again, to minimize the cost of transmission, you want the 'word' to be absolutely as short as possible. Suppose that a prize of one million dollars has been set up for the first of you who can get his/her robot to solve its cube.

In other words, we want to solve the word problem for the cube *and* we want to do it as efficiently as possible. Suppose we know we need n generators and we know that this is the smallest number. How do we make a 'word' as short as possible? To make a word in these generators as small as possible, we must know all the 'relations' between these generators so we can, if necessary, substitute them into the word and perform some cancelation. This is what this section is about.

Let X be a finite set, say $n = |X|$. Let Y be a set of reduced words on X . Let R be the smallest normal subgroup of F_n containing Y . Since R is normal, the quotient F_n/R is a group.

Definition 10.3.1. *Let G be a group. We say that G has **generators** X and **relations** Y if G is isomorphic to F_n/R . A collection of generators and relations defining a group is called a **presentation** of the group. If X and Y are both finite, then G is said to be **finitely presented**.*

This concept of a group presentation is due to von Dyck, whom we met in connection with free groups. As in the situation with free groups, von Dyck was motivated by considerations in topology.

Example 10.3.1. *If G is the Rubik's Cube group, then X is a set of moves which can solve any cube, no matter how mixed up it is. The symbol R is the standard notation used for the smallest normal subgroup of F_n generated by the relations satisfied by the generators g_1, \dots, g_n of a finitely presented group G . Don't confuse the set of relations R with the Rubik's Cube right-face move R !*

If $X = \{R, L, D, U, F, B\}$ then R might contain $\{R^4, L^4, D^4, U^4, F^4, B^4\}$, among other things, standing for the identities $R^4 = 1$, $L^4 = 1$,

Example 10.3.2. *The Rubik's Cube group G is generated by 2 elements. In fact, Martin Schönert used GAP to compute the subgroups generated by 300 random pairs of elements of G . He found that 151 of those pairs generated the entire group, so the probability that two randomly chosen elements generate G is about $\frac{1}{2}$ (see the Dec. 4, 1995, posting to [CL]).*

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In [B1], the following two generators discovered by F. Barnes, are given: $g_1 = UBLUL^{-1}U^{-1}B^{-1}$ and $g_2 = R^2FLD^{-1}R^{-1}$. In other words, as a permutation group, $G = \langle g_1, g_2 \rangle$. We know that there is a finite set of relations R_1, \dots, R_k such that there is an isomorphism $\phi : G \rightarrow \langle x, y \mid R_1(x, y) = 1, \dots, R_k(x, y) = 1 \rangle$, and moreover the isomorphism satisfies $\phi(g_1) = x$ and $\phi(g_2) = y$. What is a collection of relations $R_1(x, y) = 1, \dots, R_k(x, y) = 1$ where k is as small as possible and the total word lengths of the $R_1(x, y), \dots, R_k(x, y)$ is as small as possible?

As a matter of notation, an element $r \in R$ is written as an equation $r = 1$ in G .

Remark 10.3.1. For those with a background in topology: Serre [Se], §3.3, gives a topological interpretation of R as ‘the fundamental group of the Cayley graph of G with respect to X ’.

Example 10.3.3. The cyclic group of order 3, C_3 , has one generator x and one relation $x^3 = 1$, so

$$X = \{x\}, \quad R = \{(x^3)^k \mid k \in \mathbb{Z}\} \subset F_1 = \{x^k \mid k \in \mathbb{Z}\}.$$

Here the cosets of F_1/R are R, xR, x^2R . The set of these three cosets is closed under multiplication. For example, $(xR)(x^2R) = x(Rx^2)R = xx^2RR = x^3R = R$, so the inverse of the element xR is x^2R .

More generally, C_n has the presentation

$$C_n = \{x \mid x^n = 1\}.$$

Ponderable 10.3.1. By constructing moves of the Rubik’s Cube of order 2, 3, 4, show that the Rubik’s Cube ‘contains’ the subgroups C_2, C_3, C_4 .

Lemma 10.3.1. The group $C_m \times C_n$ is presented by

$$C_m \times C_n = \{a, b \mid a^m = 1, b^n = 1, ab = ba\} = \{a, b \mid a^m = 1, b^n = 1, aba^{-1}b^{-1} = 1\}.$$

Proposition 10.3.1. In general, if G is generated by g_1, \dots, g_m with relations $R_i(g_1, \dots, g_m) = 1$ and if H is generated by h_1, \dots, h_n with relations $S_i(h_1, \dots, h_n) = 1$ then $G \times H$ is the group generated by the g_i, h_j , with relations $R_i(g_1, \dots, g_m) = 1, S_i(h_1, \dots, h_n) = 1$, and $g_i h_j = h_j g_i$.

We shall not prove this but refer to [MKS], §4.1, Exercise 13.

Ponderable 10.3.2. By constructing moves of the Rubik’s Cube of order 2, 3 which commute show that the Rubik’s Cube ‘contains’ the subgroups $C_2 \times C_3 \cong C_6$.

Example 10.3.4. If $G = \langle R \rangle$ and $H = \langle L^2 \rangle$ then

$$G \cong \langle x \mid x^4 = 1 \rangle, \quad H \cong \langle y \mid y^2 = 1 \rangle,$$

and

$$G \times H \cong \langle x \mid x^4 = 1, y^2 = 1, xy = yx \rangle.$$

10.4 Generators, relations for groups of order < 26

This section gives a table of all the (non-isomorphic) finite groups of order ≤ 25 and their generators and relations. Figure 10.1 below is included to give a feeling for how complicated the table below can get for larger n . It is a histogram plot of the function $O(n)$, where $O(n)$ is the number of non-isomorphic groups of order n . It was obtained using **SAGE** and **GAP**.

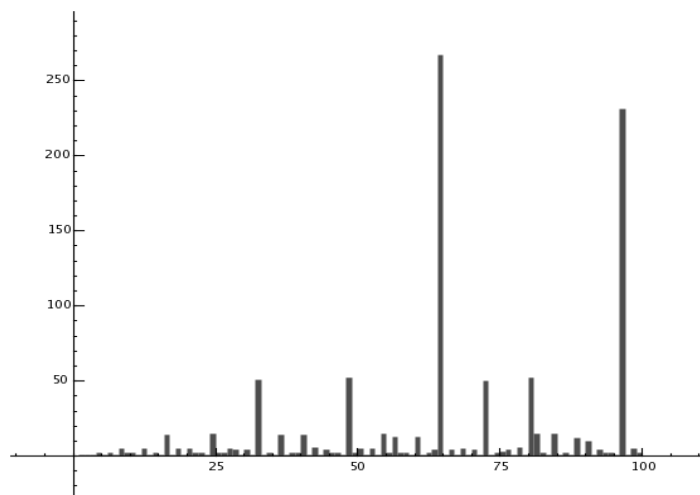


Figure 10.1: Group orders, $O(n)$, $1 \leq n \leq 100$

Notation:

- C_n = cyclic group of order n ,
- D_n = dihedral group of order $2n$ = symmetry group of the regular n -gon,
- Q = quaternion group = $\{-1, 1, -i, i, -j, j, -k, k\}$,
- Q_{2m} = generalized quaternion group = $\langle a, b \mid aba = b, b^2 = a^m, a^{2m} = 1 \rangle$,
- S_n = symmetric group of permutations of $\{1, 2, \dots, n\}$,
- A_n = alternating group of even permutations of $\{1, 2, \dots, n\}$,
- \mathbb{F}_q = finite field with q elements (q =power of a prime),
- $\mathbb{Z}/n\mathbb{Z}$ = integers modulo n .

The following table was obtained from the tables in [TW].

10.4. GENERATORS, RELATIONS FOR GROUPS OF ORDER < 26

Order	Group G	Generators	Relations	Notes
2	C_2	a	$a^2 = 1$	
3	C_3	a	$a^3 = 1$	$G \cong A_3$
4	C_4	a	$a^4 = 1$	
4	$C_2 \times C_2$	a, b	$a^2 = 1, b^2 = 1,$ $ab = ba$	Klein 4-group $Aut(G) \cong GL(2, \mathbb{F}_2)$
5	C_5	a	$a^5 = 1$	
6	$C_6 = C_2 \times C_3$	a	$a^6 = 1$	
6	S_3	a, b	$a^3 = 1, b^2 = 1,$ $aba = b$	$Aut(G) = G$ $G \cong GL(2, \mathbb{F}_2)$
7	C_7	a	$a^7 = 1$	
8	C_8	a	$a^8 = 1$	
8	$C_2 \times C_4$	a, b	$a^2 = 1, b^4 = 1,$ $ab = ba$	
8	$C_2 \times C_2 \times C_2$	a, b, c	$a^2 = 1, b^2 = 1,$ $c^2 = 1, ab = ba,$ $bc = cb, ac = ca$	$Aut(G) \cong GL(3, \mathbb{F}_2)$
8	D_4	a, b	$a^4 = 1, b^2 = 1,$ $aba = b$	
8	Q	a, b	$a^4 = 1, b^2 = a^2,$ $aba = b$	
9	C_9	a	$a^9 = 1$	
9	$C_3 \times C_3$	a, b	$a^3 = 1, b^3 = 1,$ $ab = ba$	$Aut(G) \cong GL(2, \mathbb{F}_3)$
10	$C_{10} = C_2 \times C_5$	a	$a^{10} = 1$	
10	D_5	a, b	$a^5 = 1, b^2 = 1,$ $aba = b$	
11	C_{11}	a	$a^{11} = 1$	
12	$C_{12} = C_3 \times C_4$	a	$a^{12} = 1$	
12	$C_2 \times C_6$ $= C_2 \times C_2 \times C_3$	a, b	$a^2 = 1, b^6 = 1,$ $ab = ba$	
12	D_6	a, b	$a^6 = 1, b^2 = 1,$ $aba = b$	
12	A_4	a, b	$a^2 = 1, b^3 = 1,$ $(ba)^3 = 1$	$Aut(G) = G$
12	Q_6	a, b	$a^6 = 1, b^2 = a^3,$ $aba = b$	‘dicyclic’

10.4. GENERATORS, RELATIONS FOR GROUPS OF ORDER < 26

Order	Group G	Generators	Relations	Notes
13	C_{13}	a	$a^{13} = 1$	
14	$C_{14} = C_2 \times C_7$	a	$a^{14} = 1$	
14	D_7	a, b	$a^7 = 1, b^2 = 1,$ $aba = b$	$Aut(G) = G$
15	$C_{15} = C_3 \times C_5$	a	$a^{15} = 1$	
16	C_{16}	a	$a^{16} = 1$	
16	$C_2 \times C_8$	a, b	$a^2 = 1, b^8 = 1,$ $ab = ba$	
16	$C_4 \times C_4$	a, b	$a^4 = 1, b^4 = 1,$ $ab = ba$	$Aut(G)$ $\cong GL(2, \mathbb{Z}/4\mathbb{Z})$
16	$C_2^2 \times C_4$	a, b, c	$a^2 = 1, b^2 = 1,$ $c^2 = 1, ab = ba,$ $ac = ca, bc = cb$	
16	C_2^4	a, b, c, d	$a^2 = 1, b^2 = 1,$ $c^2 = 1, d^2 = 1,$ $ab = ba, ac = ca,$ $bc = cb, bd = db,$ $ad = da, cd = dc$	$Aut(G)$ $\cong GL(4, \mathbb{F}_2)$
16	$D_4 \times C_2$	a, b, c	$a^4 = 1, b^2 = 1,$ $c^2 = 1, aba = b,$ $ac = ca, bc = cb$	
16	$Q \times C_2$	a, b, c	Ponderable	
16		a, b, c	$a^2 = 1, b^2 = 1, c^2 = 1,$ $abc = bca = cab$	
16		a, b	$a^2 = 1, b^2 = 1,$ $(ab)^2 = 1, (a^{-1}b)^2 = 1$	
16		a, b	$a^4 = 1, b^4 = 1,$ $aba = b$	semidirect product of C_4 with C_4
16		a, b	$a^8 = 1, b^2 = 1,$ $ab = ba^5$	semidirect product of C_8 with C_2 (C_8 normal)
16		a, b	$a^8 = 1, b^2 = 1,$ $ab = ba^3$	semidirect product of C_8 with C_2 (C_8 normal)
16	D_8	a, b	$a^8 = 1, b^2 = 1,$ $aba = b$	semidirect product of C_8 with C_2 (C_8 normal)
16	Q_8	a, b	$a^8 = 1, b^2 = a^4,$ $aba = b$	
17	C_{17}	a	$a^{17} = 1$	
18	$C_{18} = C_2 \times C_9$	a	$a^{18} = 1$	

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Order	Group G	Generators	Relations	Notes
18	$C_3 \times C_6$ $= C_3 \times C_3 \times C_2$	a, b	$a^3 = 1, b^6 = 1,$ $ab = ba$	
18	$S_3 \times C_3$	a, b, c	$a^3 = 1, b^2 = 1, c^3 = 1$ $aba = b, ac = ca, bc = cb$	
18	D_9	a, b	$a^9 = 1, b^2 = 1,$ $aba = b$	semidirect product of C_9 with C_2 (C_9 normal), $Aut(G) = G$
18		a, b, c	$a^3 = 1, b^3 = 1, c^2 = 1$ $ab = ba, aca = c, bcb = c$	
19	C_{19}	a	$a^{19} = 1$	
20	$C_{20} = C_4 \times C_5$	a	$a^{20} = 1$	
20	$C_2 \times C_{10}$	a, b	$a^2 = 1, b^{10} = 1,$ $ab = ba$	
20	D_{10}	a, b	$a^{10} = 1, b^2 = 1,$ $aba = b$	
20	Q_{10}	a, b	$a^{10} = 1, b^2 = a^5,$ $aba = b$	
20		a, b	$a^5 = 1, b^4 = 1,$ $ab = ba^3$	semidirect product of C_5 with C_4 (C_5 normal), $Aut(G) = G$
21	$C_{21} = C_3 \times C_7$	a	$a^{21} = 1$	
21		a, b	$a^7 = 1, b^3 = 1,$ $ab = ba^4$	semidirect product of C_7 with C_3 (C_7 normal), $Aut(G) = G$
22	$C_{22} = C_2 \times C_{11}$	a	$a^{22} = 1$	
22	D_{11}	a, b	$a^{11} = 1, b^2 = 1,$ $aba = b$	$Aut(G) = G$
23	C_{23}	a	$a^{23} = 1$	
24	$C_{24} = C_3 \times C_8$	a	$a^{24} = 1$	
24	$C_2 \times C_{12}$ $= C_2 \times C_3 \times C_4$	a, b	$a^2 = 1, b^{12} = 1,$ $ab = ba$	
24	$C_2^2 \times C_6$	a, b, c	$a^2 = 1, b^2 = 1, c^6 = 1,$ $ab = ba, ac = ca, bc = cb$	

10.4. GENERATORS, RELATIONS FOR GROUPS OF ORDER < 26

Order	Group G	Generators	Relations	Notes
24	$D_6 \times C_2$	a, b, c	$a^6 = 1, b^2 = 1, c^2 = 1,$ $aba = b, ac = ca, bc = cb$	
24	$A_4 \times C_2$	a, b, c	Ponderable	
24	$Q_6 \times C_2$	a, b, c	Ponderable	
24	$D_4 \times C_3$	a, b, c	Ponderable	
24	$Q \times C_3$	a, b, c	Ponderable	
24	$S_3 \times C_4$	a, b, c	Ponderable	
24	D_{12}	a, b	$a^{12} = 1, b^2 = 1,$ $aba = b$	
24	Q_{12}	a, b	$a^{12} = 1, b^2 = a^6,$ $aba = b$	
24	S_4	a, b	$a^4 = 1, b^2 = 1,$ $(ab)^3 = 1$	$Aut(G) = G$
24	$SL(2, \mathbb{F}_3)$	a, b, c	$a^4 = 1, b^2 = a^2,$ $c^3 = 1, aba = b,$ $ac = cb, bc = cab$	$Aut(G) \cong Aut(Q)$ $= S_4,$ semidirect product of Q with C_3 (Q normal)
24		a, b	$a^3 = 1, b^8 = 1,$ $aba = b$	semidirect product of C_3 with C_8 (C_3 normal)
24		a, b, c	$a^3 = 1, b^4 = 1,$ $c^2 = 1, bcb = c, aba = b,$ $ac = ca$	semidirect product of C_3 with D_4 (C_3 normal)
25	C_{25}	a	$a^{25} = 1$	
25	C_5^2	a, b	Ponderable	

Ponderable 10.4.1. *Extend this Table to order 26.*

Ponderable 10.4.2. *Find a subgroup of the Rubik's cube group isomorphic to the Klein 4-group.*

10.5. THE PRESENTATION PROBLEM

Example 10.4.1. Let G denote the Rubik's cube group and let Q denote the quaternion group:

$$Q = \{1, -1, i, -i, j, -j, k, -k\},$$

where i, j, k are as in §5.1. Let

$$Q^* = \langle g_1, g_2 \rangle \subset G$$

where

$$\begin{aligned} g_1 &= F^2 \cdot M_R \cdot U^{-1} \cdot M_R^{-1} \cdot U^{-1} \cdot M_R \cdot U \cdot M_R^{-1} \cdot U \cdot F^2, \\ g_2 &= F \cdot U^2 \cdot F^{-1} \cdot U^{-1} \cdot L^{-1} \cdot B^{-1} \cdot U^2 \cdot B \cdot U \cdot L. \end{aligned}$$

Then $\phi : Q \rightarrow Q^*$ is an isomorphism, where $\phi(i) = g_1$ and $\phi(k) = g_2$.

This was stated without proof in Example 9.3.5. To prove it, we use the fact that we now know the generators and relations for Q . It suffices to check that the above generators g_i satisfy the relations defining Q given in the table above. This is left as an exercise for the interested reader with a Rubik's Cube.

Ponderable 10.4.3. Which of these groups tabulated above is (isomorphic to) a subgroup of the Rubik's Cube group?

Ponderable 10.4.4. Of those which are subgroups, find moves associated to the generators given which satisfy the relations given. In other words, find an explicit isomorphism in terms of the generators.

Ponderable 10.4.4 is probably much harder than Ponderable 10.4.3 but possibly still 'doable', with the help of a program such as **SAGE**.

10.5 The presentation problem

The expectations of life depend upon diligence; the mechanic that would perfect his work must first sharpen his tools.

– Confucius, Chinese philosopher

The following problem is apparently unsolved at this time.

Ponderable 10.5.1. (Singmaster [Si]) Let G be the Rubik's Cube group. Find a set of generators and relations for G of minimal cardinality (i.e., $|X| + |Y|$ is of minimal cardinality).

Dan Hoey's post of Dec. 17, 1995 to the cube-lover's list [CL] may describe the best-known results. He suggests that G has a set X of 5 generators and a set Y of 44 relations such that the total length of all the reduced words in Y is 605. However, to find a set of generators for G of minimal cardinality, but with any number of relations, is easy. Indeed, it is known that the Rubik's Cube group can be generated by two elements [Si].

10.5.1 A presentation for $C_m^n \rtimes S_{n+1}$

Motivated by the problem of D. Singmaster mentioned above, we shall find a presentation for the generalized symmetric group. From this, one can derive a presentation for the illegal Rubik's Cube group.

We can identify the $C_m^n \rtimes S_{n+1}$ with the group of $(n+1) \times (n+1)$ invertible monomial matrices g with coefficients in C_m having the following condition on the determinant: if we write $g = p \cdot d$, where p is a permutation matrix and d is a diagonal matrix, then $\det(d) = 1$ (this determinant 1 condition is a condition corresponding to the 'conservation of twists' for the moves of the Rubik's Cube).

We may identify C_m^n with the subgroup

$$\{(x_1, \dots, x_{n+1}) \mid x_1 x_2 \dots x_{n+1} = 1, x_i \in C_m\}$$

and $C_m^n \rtimes S_{n+1}$ as a subgroup of the wreath product $S_{n+1} \wr C_m$.

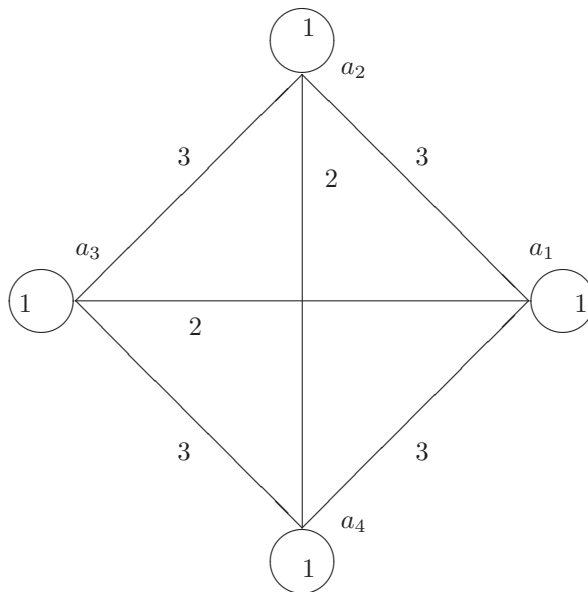
Consider $G = C_m^{n+1} \rtimes S_{n+1}$. The group S_{n+1} has the presentation

$$S_{n+1} = \langle a_1, \dots, a_n \mid (a_i a_j)^{m_{ij}} = 1, \text{ provided } 1 \leq i, j \leq n \rangle,$$

where

$$m_{ij} = \begin{cases} 3, & j = i \pm 1, \\ 2, & |i - j| > 1, \\ 1, & i = j. \end{cases}$$

The following diagram may help to visualize the exponents m_{ij} in the case $n = 4$:



As a group of $(n+1) \times (n+1)$ monomial matrices, we identify a_i with the permutation matrix,

$$s_i = \begin{pmatrix} 1 & 0 & & \dots & & 0 \\ & \ddots & & & & \\ & & 1 & & & \\ & & 0 & 1 & & \\ & & 1 & 0 & & \\ & & & & 1 & \\ & & & & & \ddots & \\ 0 & \dots & & & & 0 & 1 \end{pmatrix}.$$

If I is the $(n+1) \times (n+1)$ identity matrix and if E_{ij} denotes the matrix which is 0 in every entry except the ij entry, which is 1, then

$$s_i = I - E_{ii} - E_{i+1,i+1} + E_{i,i+1} + E_{i+1,i}.$$

The group C_m has the presentation

$$C_m = \langle h \mid h^m = 1 \rangle.$$

The group C_m^n has the presentation

$$C_m^n = \langle h_1, \dots, h_n \mid h_i^m = 1, h_i h_j = h_j h_i, \text{ provided } 1 \leq i, j \leq n+1 \rangle.$$

We identify C_m^n with the Cartesian product

$$\{(h_1(x_1), h_2(x_2), \dots, h_n(x_n)) \mid x_i \in C_m\},$$

where $h_i(t)$ is the diagonal matrix

$$h_i(t) = I - E_{ii} - E_{i+1,i+1} + tE_{i,i} + t^{-1}E_{i+1,i+1}.$$

There are the following identities between the s_i and the $h_j(t)$:

$$\begin{aligned} s_i h_j(t) s_i^{-1} &= h_j(t), \quad |i-j| > 1, \\ s_i h_i(t) s_i^{-1} &= h_i(t)^{-1}, \\ s_{i\pm 1} h_j(t) s_{i\pm 1}^{-1} &= h_i(t) h_{i\pm 1}(t). \end{aligned}$$

This motivates the formulation of the following statement.

Theorem 10.5.1. *The group $C_m^n \rtimes S_{n+1}$ is given by generators $a_1, \dots, a_n, h_1, \dots, h_n$ and relations*

$$\begin{aligned} (a_i a_j)^{m_{ij}} &= 1, \text{ provided } 1 \leq i, j \leq n, \\ h_i^m &= 1, \quad h_i h_j = h_j h_i, \text{ provided } 1 \leq i, j \leq n \\ a_i h_j a_i^{-1} &= h_j, \text{ provided } |i-j| > 1, \\ a_i h_i a_i^{-1} &= h_i^{-1}, \\ a_{i\pm 1} h_j a_{i\pm 1}^{-1} &= h_i h_{i\pm 1}. \end{aligned}$$

Remark 10.5.1. *Essentially the same presentation may be found in the paper [DM] by Davies and Morris. I only sketch the proof since it is in [DM].*

10.5.2 Idea of the proof

Here is a sketch of the proof (for which I thank Dennis Spellman) of the above theorem. Let P denote the group presented in the above theorem. The claim is, of course, that $C_m^n \rtimes S_{n+1} \cong P$. There is a surjective homomorphism $f : P \rightarrow C_m^n \rtimes S_{n+1}$ given by sending the generators to the generators. The problem is to show that this is injective. Let $K = \ker(f)$, so

$$|P| = |P/K| |K| \geq |P/K| = |C_m^n \rtimes S_{n+1}| = m^n (n+1)!.$$

Note that $H = \langle h_1, \dots, h_n \mid h_i^m = 1, h_i h_j = h_j h_i \rangle \subset P$ is a normal subgroup of P since each a_i sends a generator of H to a product of them or their inverses. Also, note that $H \cong C_m^n$.

We claim that $P/H \cong S_{n+1}$. From this it will follow that $|P| = m^n (n+1)!$, proving that $|K| = 1$, as desired.

To establish $P/H \cong S_{n+1}$, we verify that the presentation on P/H one gets from Theorem 2.1 in [MKS] is the same as that of S_{n+1} . Indeed, if

$W(a_1, \dots, a_n, h_1, \dots, h_n) = 1$ is a relation in the presentation, we must determine the word $W(a_1, \dots, a_n, h_1, \dots, h_n)H$ in P/H . Note that every relation ‘collapses’ and becomes trivial except for the relations $(a_i a_j)^{m_{ij}} = 1$. These relations define the presentation for S_{n+1} , as desired. \square

Ponderable 10.5.2. *How many generators and relations are specified by this theorem when*

- (a) $n = 7$ and $m = 3$,
- (b) $n = 11$ and $m = 2$?

Chapter 11

The (legal) Rubik's Cube group

The advantage is that mathematics is a field in which one's blunders tend to show very clearly and can be corrected or erased with a stroke of the pencil. It is a field which has often been compared with chess, but differs from the latter in that it is only one's best moments that count and not one's worst.

Norbert Wiener, Ex-prodigy: my childhood and youth

In this chapter, we build on the material in the previous chapters to describe mathematically the group of (legal) moves of the 3×3 Rubik's Cube. For first-timers, it might seem like we've just finished a long steep climb to the top of a hill. However, we are not done! Though the terrain may flatten out, there is still a bit of a hike until we get to our destination.

11.1 Mathematical description of the $3 \times 3 \times 3$ cube moves

In this section, we describe mathematically the moves of the $3 \times 3 \times 3$ Rubik's Cube. As we will see, this will lead eventually to the description of the Rubik's Cube group as a subgroup of index 12 of a direct product of two wreath products.

11.1.1 Notation

First, orient all the corners and edges as in Theorem 9.6.1.

11.1. MATHEMATICAL DESCRIPTION OF MOVES

The edge orientations are depicted as follows on the F , R , and U faces:

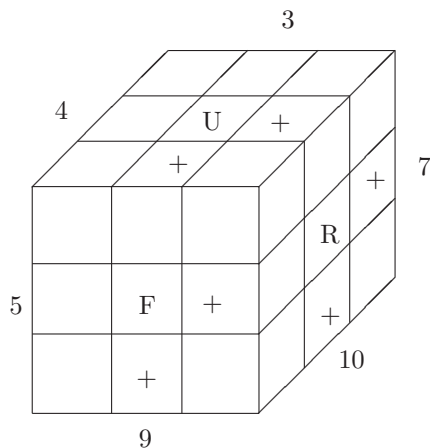


Figure 11.1: Edge orientations and labels for the cube

For the other faces: put a '+' on the dl , db , lf , lu , bl , and bu edges. This fixes the edge orientations. For the edge labels, order the edges uf , ur , ub , ul , lf , fr , rb , bl , df , dr , db , dl , and label them, in order, 1 through 12 (so, for example, the db edge is 11).

For the corner orientations, put a '+' sign on all the top corner facets and on all the bottom corner facets. The corner labels and orientations are depicted on the F , R , and D faces as in Figure 11.1.1 below. (The same labeling was used in §5.7.2 above.)

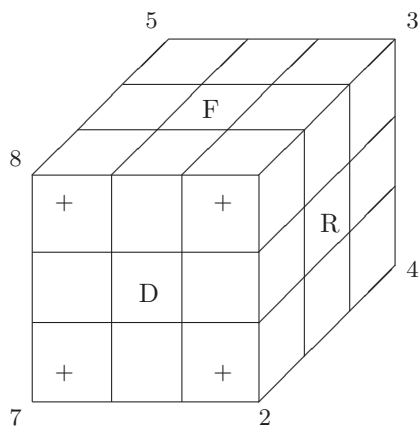


Figure 11.2: Corner orientations and labels for the cube

Let $G = \langle R, L, U, D, F, B \rangle$ be the group of the $3 \times 3 \times 3$ Rubik's Cube and let

H be the ‘enlarged’ group generated by R, L, U, D, F, B and all the ‘illegal’ moves (where one is allowed to disassemble and reassemble the cube but not remove any facets). Let V denote the set of vertices of the cube (which we identify with the set of corner subcubes of the Rubik’s Cube) and let

$$\rho : H \rightarrow S_V$$

denote the homomorphism which associates with each move of the Rubik’s Cube the corresponding permutation of the vertices (see Example 9.2.1). Let E denote the set of edges of the cube (which we identify with the set of edge subcubes of the Rubik’s Cube) and let

$$\sigma : H \rightarrow S_E$$

denote the homomorphism which associates with each move of the Rubik’s Cube the corresponding permutation of the edges (see Example 9.2.1).

These homomorphisms ρ and σ will be reformulated in a slightly different way in §11.2.2 below.

11.1.2 Corner orientations

Let $v : H \rightarrow C_3^8$ be the function which associates with each move $g \in H$ the corresponding corner orientations. More precisely, let $g \in H$ and say g moves corner i to corner j . Then $v_i(g) \in C_3$ is the orientation to which the i th vertex is sent by g , where the vertices are labeled as in the diagram shown and where the orientation is the number of 120 degrees clockwise twists required to turn the relative reference ‘+’ obtained by moving corner i to j using the move g into the standard reference ‘+’ on corner j .

Example 11.1.1. *We have*

X	$\vec{v}(X)$
F	$(2, 0, 1, 0, 1, 0, 0, 2)$
U	$(0, 0, 0, 0, 0, 0, 0, 0)$
D	$(0, 0, 0, 0, 0, 0, 0, 0)$
B	$(0, 1, 0, 2, 0, 2, 1, 0)$
R	$(1, 2, 2, 1, 0, 0, 0, 0)$
L	$(0, 0, 0, 0, 1, 2, 1, 2)$

The labeling of the corners used in this table are as in Theorem 9.6.1.

Remark 11.1.1. *The effect of a move $g \in H$ on the corner orientations may also be regarded as a relabeling of the ‘+’ markings.*

Note that a move $g \in H$ has two effects on the corners:

- (a) a permutation $\rho(g) \in S_V$ of the vertices,
- (b) a reorientation of the vertices moves in (a).

11.1. MATHEMATICAL DESCRIPTION OF MOVES

In particular, for $g, h \in H$, the orientation $\vec{v}(gh)$ can only differ from $\vec{v}(g)$ in the coordinates corresponding to the vertices permuted by h .

We shall now verify that the ‘relative’ orientation $\vec{v}(gh) - \vec{v}(g)$ is the same as the orientation $\vec{v}(h)$, provided one takes into account the effect of g on the vertices: $\vec{v}(h) = \rho(g)(\vec{v}(gh) - \vec{v}(g))$.

Lemma 11.1.1. $\vec{v}(gh) = \vec{v}(g) + \rho(g)^{-1}(\vec{v}(h))$.

Proof: The move gh orients the i th corner subcube by $v_i(gh)$ and permutes the vertices by $\rho(gh)$, by definition.

On the other hand, gh will first act by g then h . The move g will reorient the i th corner subcube by $v_i(g)$ and send the i th vertex to the $\rho(g)(i)$ th vertex.

To study the subsequent effect of h on this, let us subtract $\vec{v}(g)$ from $\vec{v}(gh)$, so that we are back to our original orientation (we will add $\vec{v}(g)$ back in later). Call this position the **modified cube** for now.

The move h first orients the j th corner subcube of the modified cube by $v_j(h)$ and permutes it to vertex $\rho(h)(j)$. The i th subcube of the modified cube comes from (via g) the $\rho(g)^{-1}(i)$ th subcube of the original cube. Thus the i th corner subcube of the modified cube is, by means of h , reoriented by $v_{\rho(g)^{-1}(i)}(h)$. To this we must add in $v_i(g)$ to get the total effect of gh on the i th vertex of the original:

$$v_i(gh) = v_i(g) + v_{\rho(g)^{-1}(i)}(h),$$

for each $1 \leq i \leq 8$, which implies Lemma 11.1.1. \square

11.1.3 Edge orientations

Let $w : H \rightarrow C_2^{12}$ be the function which associates to each move $g \in H$ the corresponding edge orientations. More precisely, let $g \in H$ and say g moves edge i to edge j . Then $w_i(g) \in C_2$ is the orientation which the i th edge gets sent to by g , where the edges are labeled as in the diagram shown and where the orientation is the number of 180 degrees flips required to turn the relative reference ‘+’ obtained by moving edge i to j using the move g into the standard reference ‘+’ on edge j .

Example 11.1.2. *We have*

X	$\vec{w}(X)$
F	$(1, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 0)$
U	$(1, 0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0)$
$F \cdot U$	$(1, 0, 1, 0, 1, 0, 0, 0, 1, 0, 0, 0)$
$U \cdot F$	$(1, 1, 1, 0, 0, 0, 0, 0, 1, 0, 0, 0)$
B	$(0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0)$
D	$(0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 1)$
R	$(0, 1, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0)$
L	$(0, 0, 0, 0, 1, 0, 0, 0, 1, 0, 0, 0)$

The labeling of the edges used in this table are as in Theorem 9.6.1.

Remark 11.1.2. *The effect of a move $g \in H$ on the edge orientations may also be regarded as a relabeling of the ‘+’ markings.*

Note that a move $g \in H$ has two effects on the edges:

- (a) a permutation $\sigma(g) \in S_E$ of the edges,
- (b) a reorientation of the edges which were moved in (a).

In particular, for $g, h \in H$, the orientation $\vec{w}(gh)$ can differ from $\vec{w}(g)$ only in the coordinates corresponding to the edges permuted by h .

We shall now claim that

$$\vec{w}(gh) = \vec{w}(g) + \sigma(g)^{-1}(\vec{w}(h)), \quad (11.1)$$

i.e., that

$$w_i(gh) = w_i(g) + w_{\sigma(g)^{-1}(i)}(h),$$

for each $1 \leq i \leq 12$. The proof of this is so similar to the proof of Lemma 11.1.1 that I leave it to the student to modify its proof to verify (11.1).

11.1.4 The semi-direct product

Consider the following direct product of two semi-direct products:

$$H' = (C_3^8 \rtimes S_V) \times (C_2^{12} \rtimes S_E).$$

Remark 11.1.3. *This may also be written in the notation of wreath products as the following direct product of two wreath products:*

$$H' = (C_3 \text{ wr } S_V) \times (C_2 \text{ wr } S_E).$$

As a set, we think of H as belonging to $C_3^8 \times S_V \times C_2^{12} \times S_E$. If we represent elements h, h' of H as $h = (v, r, w, s), h' = (v', r', w', s') \in C_3^8 \times S_V \times C_3^8 \times S_V$ then the group operation will be given by

$$h \cdot h' = (v, r, w, s) \cdot (v', r', w', s') = (v + P(r)(v'), rr', w + P(s)(w'), ss').$$

Consider the function

$$\begin{aligned} \iota : H &\rightarrow (C_3^8 \rtimes S_V) \times (C_2^{12} \rtimes S_E) \\ g &\longmapsto (v(g), \rho(g), w(g), \sigma(g)). \end{aligned}$$

Proposition 11.1.1. *The homomorphism ι is a group isomorphism, $H \cong H'$.*

Proof: Since

$$\begin{aligned} &(\vec{v}(g), \rho(g), \vec{w}(g), \sigma(g)) \cdot (\vec{v}(h), \rho(h), \vec{w}(h), \sigma(h)) \\ &= (\vec{v}(g) + P(\rho(g))(\vec{v}(h)), \rho(g)\rho(h), \vec{w}(g) + P(\sigma(g))(\vec{w}(h)), \sigma(g)\sigma(h)), \end{aligned}$$

the map ι is a homomorphism. Since any reorientation and permutation can be achieved by some illegal move, ι must be surjective. By theorem 9.6.1, the kernel of ι is trivial (this is just a fancy way of saying that if no subcube is permuted or reoriented then the cube doesn't change!). \square

11.2 Structure of the cube group

The scientist does not study nature because it is useful; he studies it because he delights in it, and he delights in it because it is beautiful. If nature were not beautiful, it would not be worth knowing, and if nature were not worth knowing, life would not be worth living.

Henri Poincaré

The main result of this section is the ‘second fundamental theorem of cube theory.’ The result of this section is both beautiful and, we hope the reader agrees, is worth knowing.

Some preliminaries. We identify, as in §11.1, each $g \in G$ with a 4-tuple

$$(\vec{v}(g), \rho(g), \vec{w}(g), \sigma(g)),$$

where

- $\rho(g)$ is the corresponding permutation of the set of vertices V of the cube,
- $\sigma(g)$ is the corresponding permutation of the set of edges E of the cube,
- $v(g), w(g)$ are ‘orientations’ defined in §11.1.

Remark 11.2.1. Let S_n denote the symmetric group on n letters and identify S_V with S_8 , S_E with S_{12} . By example 9.2.1, we know that

- (a) $\rho : G \rightarrow S_8$ is a homomorphism,
- (b) $\sigma : G \rightarrow S_{12}$ is a homomorphism.

11.2.1 The second fundamental theorem of cube theory

Question: Given a 4-tuple (v, r, w, s) , where r, s are permutations of the corners, resp. edges, as above and

$$v \in C_3^8, \quad w \in C_2^{12},$$

what conditions on r, s, v, w ensure that it corresponds to a possible position of the Rubik’s Cube?

The following result answers this question. It is, according to [BCG], due to Ann Scott.

Theorem 11.2.1. (Second fundamental theorem of cube theory) *A 4-tuple (\vec{v}, r, \vec{w}, s) as above ($r \in S_8, s \in S_{12}, \vec{v} \in C_3^8, \vec{w} \in C_2^{12}$) corresponds to a possible position of the Rubik’s Cube if and only if*

- (a) $\text{sgn}(r) = \text{sgn}(s)$, (‘equal parity as permutations’)
- (b) $v_1 + \dots + v_8 \equiv 0 \pmod{3}$, (‘conservation of total twists’)
- (c) $w_1 + \dots + w_{12} \equiv 0 \pmod{2}$, (‘conservation of total flips’).

Proof: For simplicity, we write v for \vec{v} .

First we prove the ‘only if’ part. That is, we assume that $(v, r, w, s) \in S_V \times S_E \times C_3^8 \times C_2^{12}$ represents a (legally obtained!) position of the Rubik’s Cube. From this we want to prove (a)–(c).

Let $g \in G$ be the element which moves the Rubik’s Cube from the solved position to the position associated with this 4-tuple. Then $r = \rho(g)$ and $s = \sigma(g)$. We know that g may be written as a word in the basic moves R, L, U, D, F, B , say $g = X_1 \dots X_k$, where each X_i is equal to one of the R, L, U, D, F, B . Observe that if X is any one of these basic moves, then $\text{sgn}(\rho(X)) = \text{sgn}(\sigma(X))$. Since sgn , ρ , and σ are homomorphisms, it follows that

$$\text{sgn}(r) = \text{sgn}(\rho(g)) = \prod_{i=1}^k \text{sgn}(\rho(X_i)) = \prod_{i=1}^k \text{sgn}(\sigma(X_i)) = \text{sgn}(\sigma(X)) = \text{sgn}(s).$$

This proves (a).

We have verified (b) for the basic moves in example 11.1.1 above. Note that

- (i) the conservation of twists condition in (b) is true for (v_1, \dots, v_8) if and only if it is true for any permutation $P(p)(v) = (v_{(1)p}, \dots, v_{(8)p})$;
- (ii) if (v_1, \dots, v_8) and (v'_1, \dots, v'_8) each satisfy the conservation of twists condition in (b) then their sum also satisfies it.

As above, write g as a word in the basic moves R, L, U, D, F, B , say $g = X_1 \dots X_k$, where each X_i is equal to one of the R, L, U, D, F, B . We assume that this expression is minimal in the sense that we choose X_i so that k is as small as possible. This k is called the **length** of g . (This length is the same as the distance from g to the identity in the Cayley graph of G .)

We now prove (b) by induction on the length. We have already checked it for all words of length $k = 1$.

Assume $k > 1$. By the formula giving the orientation of the product of two moves in terms of the two orientations of the moves, we have

$$\vec{v}(X_1 \dots X_{k-1} X_k) = \rho(X_1 \dots X_{k-1})^{-1}(\vec{v}(X_k)) + \vec{v}(X_1 \dots X_{k-1}).$$

The term $\rho(X_1 \dots X_{k-1})^{-1}(\vec{v}(X_k))$ satisfies the conservation of twists condition in (b) by (i) above. The term $\vec{v}(X_1 \dots X_{k-1})$ satisfies the conservation of twists condition in (b) by the induction hypothesis. Their sum satisfies the conservation of twists condition in (b) by (ii) above. This proves (b).

The proof of (c) is very similar to the proof of (b), except that we use Example 11.1.2 in place of Example 11.1.1.

Ponderable 11.2.1. *Provide the details.*

Now, we must prove the theorem in the ‘if’ direction. In other words, assuming (a), (b), and (c) we must show that there is a corresponding legal position of the Rubik’s Cube. This part of the proof is constructive.

11.2. STRUCTURE OF THE CUBE GROUP

First, we prove a special case. Assume that r and s are both the identity and that $(w_1, \dots, w_{12}) = (0, \dots, 0)$.

There is a move which twists exactly two corners and preserves the orientations and positions of all other subcubes. For example, the move $g = (R^{-1}D^2RB^{-1}U^2B)^2$ twists the ufr corner by 120 degrees clockwise, the bdl corner by 240 degrees clockwise, and preserves the orientations and positions of all other subcubes. This move can be easily modified, by a suitable conjugation, to obtain a move which twists any pair of corners, and preserves the orientations and positions of all other subcubes. These moves generate all possible 8-tuples satisfying the conservation of twists condition in (b). This proves the ‘if’ part of the theorem in the case that r and s are both the identity and that $(w_1, \dots, w_{12}) = (0, \dots, 0)$.

Next, we prove another special case. Assume that r and s are both the identity and that $(v_1, \dots, v_8) = (0, \dots, 0)$.

There is a move which flips exactly two edges and preserves the orientations and positions of all other subcubes. For example, the move

$$g = LFR^{-1}F^{-1}L^{-1}U^2RURU^{-1}R^2U^2R$$

(found in [B1], page 112) flips the uf edge, the ur edge, and preserves the orientations and positions of all other subcubes. This move can be easily modified, by a suitable conjugation, to obtain a move which flips any pair of edges, and preserves the orientations and positions of all other subcubes. These moves generate all possible 12-tuples satisfying the conservation of flips condition in (c). This proves the ‘if’ part of the theorem in the case that r and s are both the identity and that $(v_1, \dots, v_8) = (0, \dots, 0)$.

As a consequence of these last two special cases, it follows that the ‘if’ part of the theorem is true in the case that r and s are both the identity.

Finally, we prove our last special case. Assume that $(v_1, \dots, v_8) = (0, \dots, 0)$ and that $(w_1, \dots, w_{12}) = (0, \dots, 0)$. Consider the following three claims.

- Given any three edge subcubes, there is a move which is a 3-cycle on these edges and preserves the orientations and positions of all other subcubes.
- Given any three corners, there is a move which is a 3-cycle on these corners and preserves the orientations and positions of all other subcubes.
- Given any pair of edges and any pair of corners, there is a move which is a 2-cycle on these edges, a 2-cycle on these corners, and preserves the orientations and positions of all other subcubes.

Ponderable 11.2.2. *Verify these three claims.*

By proposition 9.4.1, we know that A_E is generated by the edge 3-cycles above and that A_V is generated by the corner 3-cycles above. In other words, we can construct a position of the Rubik’s Cube associated with any 4-tuple $(r, s, 0, 0)$, provided $r \in A_V$ and $s \in A_E$. The subgroup $A_E \times A_V$ is index 4 in $S_E \times S_V$ since $|S_n/A_n| = 2$. The third type of move, the edge-corner 2-cycles above,

does not correspond to an element of the subset $A_E \times A_V$ of the Rubik's Cube group because an edge 2-cycle is an odd permutation of the edges. Therefore, if we consider the subgroup of $S_E \times S_V$ generated by all three types of moves we will obtain either all of $S_E \times S_V$ or some subgroup of index 2 which properly contains $A_E \times A_V$. The first possibility can be ruled out since it contradicts the parity condition in (a). The only subgroup of $S_E \times S_V$ of index 2 which properly contains $A_E \times A_V$ is the subgroup of elements satisfying the parity condition in (a).

It follows that the 'if' part of the theorem is true in the case that v and w are both zero.

The theorem is a consequence of these special cases because of the following.

Claim: No matter what position the Rubik's Cube is in, there is always a move which does not permute any subcubes but 'solves' the orientation of the cube so that v and w are both zero.

Ponderable 11.2.3. *Prove this claim.*

Finally, the proof of the theorem is finished. \square

Corollary 11.2.1. *The Rubik's Cube group is given by*

$$G = \{g = (\vec{v}, r, \vec{w}, s) \in H \mid (a), (b), (c) \text{ hold}\},$$

where (a), (b), (c) are as in the above theorem.

Now we show that the center of the Rubik's Cube group consists of the identity and the superflip.

Corollary 11.2.2. *The center of G consists of two elements: the identity and the 'superflip' element $z = (\vec{v}, r, \vec{w}, s)$ where $\vec{w} = (1, 1, \dots, 1) \in C_2^{12}$, $\vec{v} = \vec{0} \in C_3^8$, $s = 1 \in S_{12}$ and $r = 1 \in S_8$.*

Proof: The proof given here is essentially the same as that of Bandelow [B1].

Again, write v in place of \vec{v} , for simplicity. Recall that the center of S_n , $n > 2$, is trivial (see Ponderable 5.6.2 above). Fix $(v, r, w, s) \in G$. We have

$$\begin{aligned} (v, r, w, s) \cdot (v', r', w', s') &= (v + P(r)(v'), rr', w + P(s)(w'), ss') \\ &= (v', r', w', s') \cdot (v, r, w, s) = (v' + P(r')(v), r'r, w' + P(s')(w), s's), \end{aligned}$$

for all $(v', r', w', s') \in G$, only if $r = 1$ and $s = 1$. This implies that $v + v' = v' + P(r')v$, for all $r' \in S_8$. This forces v to be either $(0, 0, \dots, 0)$ or $(1, 1, \dots, 1)$ or $(2, 2, \dots, 2)$. Since it must satisfy conservation of twists, it can't be $(1, 1, \dots, 1)$ or $(2, 2, \dots, 2)$. Similarly, $w + w' = w' + P(s')w$, for all $s' \in S_{12}$ forces w to be either $(0, 0, \dots, 0)$ or $(1, 1, \dots, 1)$. Either of these choices is okay since they both satisfy conservation of flips. \square

11.2.2 Some consequences

We shall now reformulate the above fact about the Rubik's cube group from a different point of view. This new perspective helps us determine the size of this group. Let

$$\begin{aligned} G_0 = \{(\vec{v}, r, \vec{w}, s) \mid r \in S_8, s \in S_{12}, \\ \vec{v} = (v_1, v_2, \dots, v_8), v_i \in \{0, 1, 2\}, v_1 + \dots + v_8 \equiv 0 \pmod{3}, \\ \vec{w} = (w_1, w_2, \dots, w_{12}), w_i \in \{0, 1\}, w_1 + \dots + w_{12} \equiv 0 \pmod{2}\}. \end{aligned}$$

In other words, G_0 contains the permutations of the corners, permutations of the edges, the orientations of the corners, and the orientations of the edges. There is a conservation of flips condition and a conservation of twists condition but no parity condition on the permutations. Define a binary operation $\cdot : G_0 \times G_0 \rightarrow G_0$ by

$$(\vec{v}, r, \vec{w}, s) \cdot (\vec{v}', r', \vec{w}', s') = (\vec{v} + P(r)(\vec{v}'), r \cdot r', \vec{w} + P(s)(\vec{w}'), s \cdot s'),$$

where P is the permutation matrix (Definition 3.2.2). This defines a group structure on G_0 . This is a subgroup of the illegal Rubik's Cube group of index 6.

Theorem 11.2.2. *There is an isomorphism*

$$G_0 \cong (C_3^7 \rtimes S_8) \times (C_2^{11} \rtimes S_{12}),$$

where C_n is the cyclic group with n elements and \rtimes denotes the semi-direct product and where C_n^k ($n = 2, 3, k = 7, 11$) is identified with the subgroup of C_n^{k+1} defined by

$$\{\vec{v} = (v_1, v_2, \dots, v_k) \mid v_i \in \{0, 1, n-1\}, v_1 + \dots + v_k \equiv 0 \pmod{n}\}.$$

In particular,

$$|G_0| = |S_8| |S_{12}| |C_2^{11}| |C_3^7| = 8! \cdot 12! \cdot 2^{11} \cdot 3^7.$$

Proof: This follows from the definition of semi-direct product and the Multiplication Principle (Theorem 2.4.2). \square

Corollary 11.2.3. *The Rubik's Cube group G is the kernel of the homomorphism*

$$\begin{aligned} \phi : G_0 &\rightarrow \{1, -1\} \\ (\vec{v}, r, \vec{w}, s) &\mapsto \text{sgn}(r) \text{sgn}(s). \end{aligned}$$

In particular, $G < G_0$ is normal of index 2 and

$$|G| = 8! \cdot 12! \cdot 2^{10} \cdot 3^7.$$

Proof: This follows from the Theorem and the first isomorphism theorem of group theory (Theorem 9.5.1). \square

Recall that the commutator subgroup G_1 of G is the subgroup consisting of all finite products of commutators, $[g, h] = g \cdot h \cdot g^{-1} \cdot h^{-1}$, where g, h are arbitrary elements of G .

Theorem 11.2.3. $G_1 = \{g \in G \mid \text{sgn}(\rho(g)) = \text{sgn}(\sigma(g)) = 1\}$.

Proof: It is known that the commutator subgroup of S_n is A_n , for $n > 4$ (see chapter 3 of [R]; in fact, for $n > 4$, A_n is the only proper non-trivial normal subgroup of S_n).

Consider the projection map $\phi : C_k^n \rtimes S_n \rightarrow S_n$. The kernel of ϕ is C_k^n . Moreover,

$$\phi([g, h]) = \phi(g \cdot h \cdot g^{-1} \cdot h^{-1}) = \phi(g) \cdot \phi(h) \cdot \phi(g^{-1}) \cdot \phi(h^{-1}) = [\phi(g), \phi(h)],$$

for any $g, h \in C_k^n \rtimes S_n$. This implies that ϕ maps the commutator subgroup $(C_k^n \rtimes S_n)_1$ to the commutator subgroup $(S_n)_1 = A_n$. From this it follows that $(C_k^n \rtimes S_n)_1 = C_k^n \rtimes A_n$. This implies the theorem. \square

Corollary 11.2.4. $|G_1| = |G|/2$.

Proof: The above theorem implies that $|G/G_1| = 2$, by using the first isomorphism theorem. \square

11.3 The moves of order 2

In this section I present, as promised, a method for determining the moves of a given order in the Rubik's Cube group. We shall only compute the number of moves of order 2, though the method should, in principle, work more generally. However, perhaps it is possible for the curious and enterprising reader to use a computer find the moves of order d , where $d > 2$, in a manner similar to what we do here.

The set of elements of order 2 may be divided into 3 disjoint subsets:

- (a) elements moving corner pieces only,
- (b) elements moving edge pieces only,
- (c) elements moving corner and edge pieces.

Let us look at a subset of (a), consisting only of the even permutations of order 2 of the 8 corner pieces. Obviously those elements can only consist of orientation-preserving swaps of corner pieces where either 2 pairs (4 corner pieces) or 4 pairs (all 8 corner pieces) are affected. Given a pair of corner pieces, there are always 3 orientation-preserving swaps.

The remainder of (a) and (b) (yes, do not forget (b)) has to be calculated analogously.

Recall from §9.8.2 that an element $(\vec{v}, f) \in C_m \text{ wr } S_n$, $\vec{v} = (v_1, v_2, \dots, v_n) \in C_m^n$ and $f \in S_n$, is order d if and only if $f^d = 1$ and $\vec{v} + \dots + f^{d-1}(\vec{v}) = 0$.

In particular, elements of order 2 in $C_3^7 \rtimes S_8$ are those (\vec{v}, r) with r of order 2 (hence a product of distinct 2-cycles) and \vec{v} satisfying $v_1 + \dots + v_8 \cong 0 \pmod{3}$ and $v_i + v_{(i)r} \cong 0 \pmod{3}$. *Not all such elements correspond to a legal move*

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of the *Rubik's Cube*. Using the ‘multiplication principle of counting’ (Theorem 2.4.2), we compute the number of elements of order 2 in $C_3^7 \rtimes S_8$ to be

$$\begin{aligned} & \frac{1}{4!} \binom{8}{2} \binom{6}{2} \binom{4}{2} 3^4 + \\ & \frac{1}{3!} \binom{8}{2} \binom{6}{2} \binom{4}{2} 3^3 + \\ & \frac{1}{2!} \binom{8}{2} \binom{6}{2} 3^2 + \binom{8}{2} 3 = 21819. \end{aligned} \quad (11.2)$$

The ‘3³’ term and the ‘3’ term do not count legal Rubik’s Cube moves, if taken on their own, since they are not even (hence do not satisfy the ‘conservation of parity’ condition). The number of elements (\vec{v}, r) of order 2 in $C_3^7 \rtimes S_8$ with r even is

$$\frac{1}{4!} \binom{8}{2} \binom{6}{2} \binom{4}{2} 3^4 + \frac{1}{2!} \binom{8}{2} \binom{6}{2} 3^2 = 10395. \quad (11.3)$$

The number of elements (\vec{v}, r) of order 2 in $C_3^7 \rtimes S_8$ with r odd is

$$\frac{1}{3!} \binom{8}{2} \binom{6}{2} \binom{4}{2} 3^3 + \binom{8}{2} 3 = 11424. \quad (11.4)$$

Similarly, using the ‘multiplication principle of counting’ (Theorem 2.4.2), we compute the number of elements of order 2 in $C_2^{11} \rtimes S_{12}$ to be

$$\begin{aligned} & \frac{1}{6!} \binom{12}{2} \binom{10}{2} \cdots \binom{4}{2} 2^6 + \\ & \frac{1}{5!} \binom{12}{2} \binom{10}{2} \cdots \binom{4}{2} 2^6 + \frac{1}{4!} \binom{12}{2} \cdots \binom{6}{2} 2^7 \\ & + \frac{1}{3!} \binom{12}{2} \cdots \binom{8}{2} 2^8 + \frac{1}{2!} \binom{12}{2} \binom{10}{2} 2^9 + \binom{12}{2} 2^{10} \\ & + \binom{12}{0} 2^{11} - 1 = 15687871. \end{aligned} \quad (11.5)$$

Again, not all these count legal Rubik’s Cube moves, if taken on their own. The number of elements (\vec{w}, s) of order 2 in $C_2^{11} \rtimes S_{12}$ with s even is

$$\begin{aligned} & \frac{1}{6!} \binom{12}{2} \binom{10}{2} \cdots \binom{4}{2} 2^6 + \frac{1}{4!} \binom{12}{2} \cdots \binom{6}{2} 2^7 + \frac{1}{2!} \binom{12}{2} \binom{10}{2} 2^9 \\ & = 8080447. \end{aligned} \quad (11.6)$$

The number of elements (\vec{w}, s) of order 2 in $C_2^{11} \rtimes S_{12}$ with s odd is

$$\begin{aligned} & \frac{1}{5!} \binom{12}{2} \binom{10}{2} \cdots \binom{4}{2} 2^6 + \frac{1}{3!} \binom{12}{2} \cdots \binom{8}{2} 2^8 + \binom{12}{2} 2^{10} \\ & = 7607424. \end{aligned} \quad (11.7)$$

Example 11.3.1. Here is an example of how to use **SAGE** to compute one of these terms, in case you wish to verify it yourself.

SAGE

```
sage: C2 = [binomial(i,2) for i in range(13)]
sage: C2
[0, 0, 1, 3, 6, 10, 15, 21, 28, 36, 45, 55, 66]
sage: C2[12] == binomial(12,2)
True
sage: a = prod(C2[i] for i in [4,6,8,10,12])*2^6/factorial(5)
sage: b = prod(C2[i] for i in [8,10,12])*2^8/factorial(3)
sage: c = prod(C2[i] for i in [12])*2^10/factorial(1)
sage: a+b+c
7607424
```

The others can be checked similarly.

The second fundamental theorem of cube theory (§11.2.1) implies that an element (\vec{v}, r, \vec{w}, s) of the Rubik's Cube group is order 2 if and only if it is an element of order 2 in the group $H = (C_3^7 \rtimes S_8) \times (C_2^{11} \rtimes S_{12})$ and $\text{sgn}(r) = \text{sgn}(s)$. A non-trivial element $(h_1, h_2) \in H$, with $h_1 = (\vec{v}, r)$ and $h_2 = (\vec{w}, s)$, is order 2 if and only if exactly one of the following mutually exclusive cases occurs:

- (a) $h_1 \neq 1, h_2 = 1, h_1^2 = 1, \text{sgn}(r) = 1$,
- (b) $h_1 = 1, h_2 \neq 1, h_2^2 = 1, \text{sgn}(s) = 1$,
- (c) $h_1 \neq 1, h_2 \neq 1, h_1^2 = h_2^2 = 1, \text{sgn}(r) = \text{sgn}(s) = 1$,
- (d) $h_1 \neq 1, h_2 \neq 1, h_1^2 = h_2^2 = 1, \text{sgn}(r) = \text{sgn}(s) = -1$.

We count each case. Case (a) is counted in (11.3), case (b) is counted in (11.6), case (c) is counted in (11.3) and (11.6) (which you multiply together to get the total), and case (d) is counted in (11.4) and (11.7) (again, you multiply together to get the total). Totalling these up, we find that the number of elements of order 2 in the Rubik's Cube group is

$$(11.3) + (11.6) + (11.3) \cdot (11.6) + (11.4) \cdot (11.7) = 170911549183 = (1.7\ldots) \times 10^{11}.$$

Remark 11.3.1. Why all this theory to solve this problem? Why not make a computer do it? How long would it take a computer to just count up the number of moves of order 2 in the Rubik's Cube group? Suppose that your computer can check if an element g in the Rubik's Cube group G is order 2 or not in 0.001 seconds (this is actually faster than the computer I'm using to type this). By 'looping' over all (roughly) 4.3×10^{19} elements, this means that we can count the moves in order 2 in about 4.3×10^{15} seconds. There are about 3.1×10^6

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seconds in a year, so this means that it would take about 136 million years to do this by this method!

A lot of mathematical problems arising in applications require a lot of theory and significant computational horsepower. This is an example.

Chapter 12

Squares, two-faces, and other subgroups

In arctic and tropical times,
the integers, addition, and times,
taken mod p will yield
a full finite field,
as p ranges over the primes.

Ancient math poem

Rather than solve the Rubik's Cube for an arbitrary scrambled position, what if you noticed that the scrambler only used certain types of moves closed under composition (such as the slice group)? Could you solve the cube using only those types of moves?

It is remarkable that several 'familiar' groups may be embedded into the Rubik's Cube group, and hence be regarded as a subgroup of the cube group. For example, we have seen in an earlier chapter how to embed the group of quaternions $Q = \{1, -1, i, -i, j, -j, k, -k\}$ inside the Rubik's Cube group. The slice (sub)group is another such example. We shall meet another subgroup—the 'two faces' group—here, which allows us to introduce finite fields.

The subgroup method, discussed in chapter 15, is a method for solving the Rubik's Cube which can be implemented on a computer. One of the groups arising in this method is the group studied in the next section.

12.1 The squares subgroup

Let G denote the subgroup of the Rubik's Cube group generated by the *squares* of the basic moves,

$$G = \langle U^2, D^2, R^2, L^2, F^2, B^2 \rangle,$$

12.1. THE SQUARES SUBGROUP

called the **squares subgroup**. We will verify below that the order of this group is

$$|G| = 2^{13}3^4. \quad (12.1)$$

(By the way, as a consequence of this and Burnside's theorem [R], ch. 5, it follows that G is a solvable group. We will not need this fact.) In this section, we will investigate the group structure of G using the same method which was used to determine the structure of the Rubik's Cube group.

The group G acts on the set of edges and the set of vertices of the cube. There is a choice of orientation of the edges (resp., corners) similar to that in §§10.1.2-3 such that each element of G *preserves the edge (resp., corner) orientations*.

The action ϕ of G on the edges E of the cube has exactly 3 orbits: the middle slice parallel to the right face E_R , the middle slice parallel to the front face E_F , the middle slice parallel to the up face E_U . In particular, the group G acts (by restriction) on E_R , E_F , and E_U . The action ψ of G on the set of vertices V has exactly 2 orbits: $V_1 = \{ufr, ubl, dfl, drb\}$, $V_2 = \{ufl, ubr, dfr, dlb\}$. Therefore, the group G acts (by restriction) on V_1 and V_2 . These actions yield associated homomorphisms:

$$\begin{aligned} \phi &: G \rightarrow S_E, \\ \phi_{E_R} &: G \rightarrow S_{E_R}, \\ \phi_{E_F} &: G \rightarrow S_{E_F}, \\ \phi_{E_U} &: G \rightarrow S_{E_U}, \\ \psi &: G \rightarrow S_V, \\ \psi_1 &: G \rightarrow S_{V_1}, \\ \psi_2 &: G \rightarrow S_{V_2}. \end{aligned}$$

Proposition 12.1.1. $G = \phi(G) \times \psi(G)$.

The proof of this is left as an exercise. (Hint: use the second fundamental theorem of Rubik's Cube theory.)

Lemma 12.1.1. *If $g \in G$ then*

$$\text{sgn}(\phi_{E_R}(g))\text{sgn}(\phi_{E_F}(g))\text{sgn}(\phi_{E_U}(g)) = 1.$$

Conversely, if $(p_1, p_2, p_3) \in S_{E_R} \times S_{E_F} \times S_{E_U}$ then there is a $g \in G$ such that $p_1 = \phi_{E_R}(g)$, $p_2 = \phi_{E_F}(g)$, $p_3 = \phi_{E_U}(g)$ if and only if $\text{sgn}(p_1)\text{sgn}(p_2)\text{sgn}(p_3) = 1$.

As a consequence, we find that

$$\begin{aligned} \phi(G) &= \ker(\text{sgn} \times \text{sgn} \times \text{sgn} : \phi_{E_R}(G) \times \phi_{E_F}(G) \times \phi_{E_U}(G) \rightarrow \{\pm 1\}) \\ &\cong (S_4 \times S_4 \times S_4)/C_2. \end{aligned}$$

In particular, $|\phi(G)| = (4!)^3/2 = 2^8 3^3$.

It remains to determine $\psi(G)$. We denote this group by H for notational simplicity. We may label the vertices of the cube $1, 2, \dots, 8$ in such a way that $H = \langle u, d, l, r, f, b \rangle$, where $u = \psi(U^2)$, ..., and

$$\begin{aligned} u &= (1, 3)(2, 4), & f &= (1, 8)(4, 5), & d &= (5, 7)(6, 8), \\ b &= (3, 6)(2, 7), & r &= (2, 5)(1, 6), & l &= (4, 7)(3, 8). \end{aligned} \quad (12.2)$$

The action of H on the set of vertices of the cube has two orbits ($\{1, 3, 6, 8\}$ and $\{2, 4, 5, 7\}$ in our labeling above), which we denote for simplicity by V_1 and V_2 . There are homomorphisms

$$\psi_1 : H \rightarrow S_{V_1}, \quad \psi_2 : H \rightarrow S_{V_2},$$

but I shall not say much about these. Instead, we shall determine more about H . According to **SAGE**, this group has $|H| = 96$ elements and 10 conjugacy classes (by the way, ‘under the hood’ **SAGE** is using [Gap]):

Size	Representative
1	1
12	$d = (5, 7)(6, 8)$
32	$l \cdot d = (3, 6, 8)(4, 5, 7)$
3	$d \cdot l \cdot b \cdot l = (2, 4)(5, 7)$
12	$d \cdot b \cdot l = (2, 4, 7, 5)(3, 6)$
3	$l \cdot b \cdot l \cdot u = (1, 3)(6, 8)$
12	$b \cdot l \cdot u = (1, 3, 6, 8)(4, 7)$
3	$u \cdot d = (1, 3)(2, 4)(5, 7)(6, 8)$
12	$l \cdot d \cdot u = (1, 3, 6, 8)(2, 4, 5, 7)$
6	$u \cdot r \cdot d \cdot f = (1, 3)(2, 5)(4, 7)(6, 8)$

SAGE

```
sage: u2 = '(1,3)(2,4)'; f2 = '(1,8)(4,5)'; d2 = '(5,7)(6,8)'
sage: b2 = '(3,6)(2,7)'; r2 = '(2,5)(1,6)'; l2 = '(4,7)(3,8)'
sage: H = PermutationGroup([u2,f2,d2,b2,r2,l2])
sage: H.order()
96
sage: H.conjugacy_classes_representatives()

[(),
 (5,7)(6,8),
 (3,6,8)(4,5,7),
 (2,4)(5,7),
 (2,4,7,5)(3,6),
 (1,3)(6,8),
 (1,3,6,8)(4,7),
 (1,3)(2,4)(5,7)(6,8),
 (1,3,6,8)(2,4,5,7),
 (1,3)(2,5)(4,7)(6,8)]
```

The stabilizer in H of any vertex $v \in V$, written H_v , is a subgroup of order 24 isomorphic to the symmetric group S_4 .

Furthermore, H has a normal subgroup N of order 48 (and index 2), where N is a semidirect product of C_3 by C_2^4 , with C_2^4 normal in N .

This is all I shall say about H .

The order of G is therefore $G = |\phi(G)| \cdot |H| = 96 \cdot (4!)^3 / 2 = 2^{13} 3^4$, as claimed in (12.1) above.

12.2 Fast-forwarding though finite fields

In this subsection, I introduce very briefly fields and especially finite fields. There are many excellent books on these subjects, so the treatment here will give the reader only a quick introduction to the basic ideas.

Finite fields are especially useful for communication theory since they are finite, and hence can represent the letters of an alphabet inside a computer, and since they carry a lot of structure, which can aid in making certain types of calculations faster. We shall focus mostly on finite fields in this section.

12.2.1 The general definition of a field

A field is a set F with an addition law $+$ and a multiplication law \cdot which obeys a list of properties similar to those for the field of real numbers \mathbb{R} . More precisely, we call $(F, +, \cdot)$ a **field** if

- (F1) $(F, +)$ is an abelian group, with an identity element denoted 0 ('the additive group of the field'),
- (F2) for all $x, y, z \in F$, $(x + y)z = xz + yz$ ('distributive law'),
- (F3) $(F - \{0\}, \cdot)$ is an abelian group, with an identity element denoted 1 ('the multiplicative group of the field').

Lemma 12.2.1. *If F is a finite field, then not only is $(F - \{0\}, \cdot)$ an abelian group, but it is actually a cyclic group.*

We shall not prove this here (see chapter 13 of Artin [Ar] for a proof). (The basic idea is that all non-zero elements x of a finite field must satisfy $x^q = 1$, for some fixed prime power q , and the set of all such elements is a cyclic group.)

Definition 12.2.1. *Let F_1, F_2 be two fields. A function $f : F_1 \rightarrow F_2$ is called a **field isomorphism** if*

- *when f is restricted to the additive group $(F_1, +)$, call this restriction f again, it yields an isomorphism of groups $f : (F_1, +) \rightarrow (F_2, +)$,*
- *when f is restricted to the multiplicative group $(F_1 - \{0\}, \cdot)$, call this restriction f again, it yields an isomorphism of groups $f : (F_1 - \{0\}, \cdot) \rightarrow (F_2 - \{0\}, \cdot)$.*

12.2.2 A construction of \mathbb{F}_p

For each $k \in \mathbb{Z}$, let \bar{k} denote the residue class of $k \bmod p$. Let \mathbb{F}_p denote the **finite field with p elements**, so \mathbb{F}_p is, as a set,

$$\mathbb{F}_p = \{\bar{0}, \bar{1}, \dots, \overline{p-1}\},$$

with addition and multiplication being performed mod p .

Example 12.2.1. When $p = 5$, \mathbb{F}_5 will denote the finite field with 5 elements, so \mathbb{F}_5 is, as a set,

$$\mathbb{F}_5 = \{\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}\},$$

with addition and multiplication being performed mod 5.

It is a general fact that if F is any finite field, then there is a prime number p such that $px = 0$ for all $x \in F$. This prime number is called the **characteristic** of F . The easiest example of a finite field with characteristic p is the finite field constructed above having p elements, \mathbb{F}_p . It is not hard to see that any other finite field F of characteristic p must be isomorphic (as a vector space) to a finite dimensional vector space over \mathbb{F}_p . (Even if you've never seen a 'vector space over \mathbb{F}_p ' defined before, if you know what a 'real vector space' is from Definition 2.1.2, then you've got the right idea.) The dimension of the vector space F is called the **degree** of the field F over \mathbb{F}_p , denoted $d = [F : \mathbb{F}_p]$. In this case, we say F is a **field extension** of \mathbb{F}_p (of degree d).

Lemma 12.2.2. For fixed p, d there is, up to isomorphism, only one such field extension of degree d and characteristic p .

We shall not prove this. See chapter 13 of Artin [Ar] for a proof.

Next, we shall show how to *construct* such extensions.

12.2.3 A construction of finite fields

First, some general remarks on how one might begin to implement finite fields on a computer.

Let F be a finite field. Since F is a finite dimensional vector space containing \mathbb{F}_p , it has a vector space basis which we label as

$$e_1 = 1, e_2, \dots, e_d.$$

Thus F is, as a set, the collection of elements of the form

$$x_1 e_1 + \dots + x_n e_n, \quad x_i \in \mathbb{F}_p.$$

Since F is a field, there are $c_{ij}^k \in \mathbb{F}_p$, which we call **structure constants**, such that

$$e_i e_j = \sum_{k=1}^n c_{ij}^k e_k.$$

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There are $d_i^k \in \mathbb{F}_p$, which we call **inversion constants**, such that

$$e_i^{-1} = \sum_{k=1}^n d_i^k e_k.$$

These $p^3 + p^2$ constants determine how to multiply and divide elements of F . All a computer needs to do the arithmetic operations of F can be stored in memory.

Example 12.2.2. Let $p = 5$, so $\mathbb{F}_5 = \{\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}\}$. The set of squares is given by

$$\{x^2 \mid x \in \mathbb{F}_5\} = \{\bar{0}, \bar{1}, \bar{4}\}.$$

In particular, $\bar{2}, \bar{3}$ are not squares in this field. Let $e_2 = \sqrt{\bar{2}}$ be a formal symbol for some element which satisfies $e_2^2 = \bar{2}$. This is a root of the polynomial $x^2 - \bar{2} = 0$.

The vector space F over \mathbb{F}_5 with basis $\{e_1 = 1, e_2\}$ is 2-dimensional over \mathbb{F}_5 . Two elements $x_1 e_1 + x_2 e_2 = x_1 + x_2 \sqrt{\bar{2}}$ and $y_1 e_1 + y_2 e_2 = y_1 + y_2 \sqrt{\bar{2}}$ are multiplied by the rule

$$(x_1 + x_2 \sqrt{\bar{2}}) \cdot (y_1 + y_2 \sqrt{\bar{2}}) = x_1 y_1 + \bar{2} x_2 y_2 + (x_1 y_2 + y_1 x_2) \sqrt{\bar{2}}.$$

It is a degree 2 field extension since

$$c_{11}^1 = \bar{1}, \quad c_{11}^2 = \bar{0}, \quad c_{12}^1 = \bar{0}, \quad c_{12}^2 = \bar{1}, \quad c_{21}^1 = \bar{2}, \quad c_{21}^2 = \bar{0},$$

and

$$d_1^1 = \bar{1}, \quad d_1^2 = \bar{0}, \quad d_2^1 = \bar{0}, \quad d_2^2 = \bar{3}.$$

Using **SAGE**, this extension field is constructed as follows (in fact, I give two ways).

SAGE

```
sage: P = PolynomialRing(GF(5), "x")
sage: x = P.gen()
sage: K.<a> = GF(5^2, name='a', modulus=x^2 - 2 )
sage: K
Finite Field in a of size 5^2
sage: L = GF(25, "b")
sage: L
Finite Field in b of size 5^2
```

The construction used in the above example may be summarized more generally as follows:

1. Pick an element $\bar{m} \in \mathbb{F}_p$ which is not the square of another element, if such an element exists.

2. Let $e_1 = 1$ and $e_2 = \sqrt{\bar{m}}$ be a formal symbol for some element which satisfies $e_2^2 = \bar{m}$.
3. As a set, let $F = \{xe_1 + ye_2 \mid x, y \in \mathbb{F}_p\}$. To define F as a field, let $+$ be ‘componentwise addition’ mod p and let \cdot be defined by

$$(x_1 + x_2\sqrt{\bar{m}}) \cdot (y_1 + y_2\sqrt{\bar{m}}) = x_1y_1 + \bar{m}x_2y_2 + (x_1y_2 + y_1x_2)\sqrt{\bar{m}}.$$

A finite field F constructed in this way is called a **quadratic extension** of \mathbb{F}_p .

Definition 12.2.2. Let $f(x)$ be a polynomial having coefficients in a field F . If $f(x)$ does not factor into polynomials of smaller degree having coefficients in F , then we call $f(x)$ an **irreducible polynomial over F** . A polynomial that is not irreducible is called **reducible over F** , or we say it **factors over F** .

Remark 12.2.1. For this definition to make sense, it is important that F be a field. For example, consider the following bizarre computation with integer polynomials mod 6: $(3x + 2)(2x + 3) = 6x^2 + 13x + 6 = x$. The degree of the product actually goes down!

Let $d > 1$ be an integer.

1. Assume $g(x) = x^{d-1} + x^{d-2} + \dots + x + 1$ is irreducible over \mathbb{F}_p . Pick an element $\bar{m} \in \mathbb{F}_p$ which is not the k th power of another element in \mathbb{F}_p (for all $1 < k$ dividing d). (Since we assumed $g(x)$ was irreducible, such an element must exist.)
2. Let $e_1 = 1$, let $e_2 = \bar{m}^{1/d}$ be a formal symbol that satisfies $e_2^d = \bar{m}$, and (if $d > 2$) let $e_i = e_2^{i-1}$ for $i = 3, \dots, d$.
3. As a set, let $F = \{x_1e_1 + \dots + x_de_d \mid x_i \in \mathbb{F}_p\}$. To define F as a field, let $+$ be ‘componentwise addition’ mod p and let \cdot be defined by expanding and collecting $(x_1e_1 + \dots + x_de_d) \cdot (y_1e_1 + \dots + y_de_d)$.

A finite field F constructed in this way has p^d elements.

What if $x^{d-1} + x^{d-2} + \dots + x + 1$ is reducible? (This can happen, for example, when $p = 2$.) In this case, we must replace the formal symbol $e_2 = \bar{m}^{1/d}$ above by a root of an arbitrary but fixed irreducible polynomial $g(x)$ of degree d having coefficients in \mathbb{F}_p . (In fact, **SAGE** and **GAP** use special irreducible polynomials, so-called Conway polynomials [L], for the purpose of constructing finite field extensions. Such polynomials always exist but, as of the date of this writing, there is no known simple general formula for writing them down.) Again, we let $F = \{x_1e_1 + \dots + x_de_d \mid x_i \in \mathbb{F}_p\}$, where $e_i = e_2^{i-1}$ for $i = 3, \dots, d$. This is a field for the same reason as before.

Let $\mathbb{F}_p[x]$ denote the collection of all polynomials having coefficients in \mathbb{F}_p . This set is closed under ordinary polynomial addition and multiplication. Given any fixed polynomial $g(x)$ which is not a constant, let $\mathbb{F}_p[x]/g(x)\mathbb{F}_p[x]$ denote the

12.3. $\text{PGL}(2, \mathbb{F}_5)$ AND TWO FACES OF THE CUBE

set of all cosets $f(x) + g(x)\mathbb{F}_p[x]$ (for $f \in \mathbb{F}_p[x]$). Define $+$, \cdot on $\mathbb{F}_p[x]/g(x)\mathbb{F}_p[x]$ by $(f_1(x) + g(x)\mathbb{F}_p[x]) + (f_2(x) + g(x)\mathbb{F}_p[x]) = (f_1(x) + f_2(x)) + g(x)\mathbb{F}_p[x]$ and $(f_1(x) + g(x)\mathbb{F}_p[x]) \cdot (f_2(x) + g(x)\mathbb{F}_p[x]) = (f_1(x) \cdot f_2(x)) + g(x)\mathbb{F}_p[x]$.

Lemma 12.2.3. *If $g(x)$ is an irreducible polynomial in $\mathbb{F}_p[x]$ then $\mathbb{F}_p[x]/g(x)\mathbb{F}_p[x]$, with addition and multiplication defined above, is a field.*

See Artin [Ar], chapter 13, for a proof and more details on finite fields.

In fact, the field constructed above is isomorphic (as a field) to $\mathbb{F}_p[x]/g(x)\mathbb{F}_p[x]$.

12.3 $\text{PGL}(2, \mathbb{F}_5)$ and two faces of the cube

This section is devoted to ‘determining’ the **two-face group** generated by only two basic moves, $\langle F, U \rangle$. Singmaster [Si] has shown that

$$\langle \sigma(F), \sigma(U) \rangle \cong S_7, \quad \langle \rho(F), \rho(U) \rangle \cong \text{PGL}_2(\mathbb{F}_5),$$

where $\text{PGL}_2(\mathbb{F}_5)$ (sometimes also denoted $\text{PGL}(2, 5)$) is a group of order 120 which is defined below. Here S_7 arises from the action of the two-face group on the edges and $\text{PGL}_2(\mathbb{F}_5)$ arises from the action on the corners. In this section, we focus on the action on the corners.

Example 12.3.1. *Here is some SAGE code which illustrates Singmaster’s result. First of all, here is the two-face group:*

SAGE

```
sage: rubik = CubeGroup()
sage: f = rubik.F()
sage: u = rubik.U()
sage: T = PermutationGroup([f,u])
sage: T.order()
73483200
```

Using the corner labeling in (12.2) above, we see that $\sigma(F)$ acts by $(1, 4, 8, 5)$ and $\sigma(U)$ acts by $(1, 2, 3, 4)$:

SAGE

```
sage: P = PGL(2,5)
sage: sT = PermutationGroup(['(1,8,5,3)', '(4,3,5,6)'])
sage: sT.order()
120
sage: sT.is_isomorphic(P)
True
```

Using the labeling in §11.1.1, we see that $\rho(F)$ acts by $(1, 6, 9, 5)$ and $\rho(U)$ acts by $(1, 4, 3, 2)$:

SAGE

```
sage: rT = PermutationGroup(['(1,6,9,5)', '(1,4,3,2)'])
sage: rT.is_isomorphic(S7)
True
```

12.3.1 Möbius transformations

Definition 12.3.1. *The projective plane*

$$\mathbb{P}^1(\mathbb{F}_p) = \{\bar{0}, \bar{1}, \dots, \overline{p-1}, \infty\}$$

is defined to be the set of lines through the origin in the Cartesian plane \mathbb{F}_p^2 , associating each number (including ∞) with the slope of the corresponding line.

If $a, b, c, d \in \mathbb{F}_p$ are given numbers (not all equal to zero) then we define the **Möbius transformation** f by:

$$f : \mathbb{P}^1(\mathbb{F}_p) \rightarrow \mathbb{P}^1(\mathbb{F}_p) \\ x \mapsto \frac{ax+b}{cx+d}.$$

This is also called a **linear fractional transformation** (or a fractional linear transformation). This function is named after August Möbius (1790-1868), best known for his work in topology, especially for his conception of the Möbius strip, a two-dimensional surface with only one side. (Can't resist the riddle: What's one-sided and swims in the sea? Möbius Dick.)

Theorem 12.3.1. *The mapping f is a bijection if and only if $\det \begin{pmatrix} a & b \\ c & d \end{pmatrix} \neq 0$.*

Before proving this, we need the following definition.

Definition 12.3.2. *For any prime number p , define*

$$GL(2, \mathbb{F}_p) = \left\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} \mid a, b, c, d \in \mathbb{F}_p, \ ad - bc \neq 0 \right\}.$$

This set $GL(2, \mathbb{F}_p)$ is a group under ordinary matrix multiplication (check this). Furthermore, the group $G = GL(2, \mathbb{F}_p)$ acts on the set $\mathbb{P}^1(\mathbb{F}_p)$ by means of Möbius transformations, thus defining a function

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} : \mathbb{P}^1(\mathbb{F}_p) \rightarrow \mathbb{P}^1(\mathbb{F}_p),$$

for each $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in G$.

Lemma 12.3.1.

- (a) The center of $GL(2, \mathbb{F}_p)$ (i.e., the subgroup of all elements which commute with every element in $GL(2, \mathbb{F}_p)$) is given by

$$Z(GL(2, \mathbb{F}_p)) = \left\{ \begin{pmatrix} a & 0 \\ 0 & a \end{pmatrix} \mid a \in \mathbb{F}_p, a \neq 0 \right\}.$$

- (b) This subgroup is normal in $GL(2, \mathbb{F}_p)$.

- (c) There is an isomorphism

$$Z(GL(2, \mathbb{F}_p)) \cong \mathbb{F}_p^\times.$$

Proof: (a) Since

$$\begin{pmatrix} r & 0 \\ 0 & r \end{pmatrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix} = \begin{pmatrix} r & 0 \\ 0 & r \end{pmatrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

we can conclude that

$$\left\{ \begin{pmatrix} a & 0 \\ 0 & a \end{pmatrix} \mid a \in \mathbb{F}_p, a \neq 0 \right\} \subset Z(GL(2, \mathbb{F}_p)).$$

To show that

$$Z(GL(2, \mathbb{F}_p)) \subset \left\{ \begin{pmatrix} a & 0 \\ 0 & a \end{pmatrix} \mid a \in \mathbb{F}_p, a \neq 0 \right\},$$

assume that

$$\begin{pmatrix} r & s \\ u & v \end{pmatrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix} = \begin{pmatrix} r & s \\ u & v \end{pmatrix} \begin{pmatrix} a & b \\ c & d \end{pmatrix}$$

for all a, b, c, d . This implies $bu = cs$ for all b, c . This is impossible unless $u = s = 0$. This in turn forces $cr = cv$, for all c . This implies that $r = v$. This proves the desired inclusion.

The proof of parts (b) and (c) are left as an exercise for the reader. \square

Definition 12.3.3. The quotient group, denoted

$$PGL(2, \mathbb{F}_p) = GL(2, \mathbb{F}_p)/Z(GL(2, \mathbb{F}_p)),$$

is called the **projective linear group**. (This is a group since the center is a normal subgroup by the lemma above.)

The quotient group, denoted

$$PSL(2, \mathbb{F}_p) = SL(2, \mathbb{F}_p)/Z(SL(2, \mathbb{F}_p)),$$

is called the **projective special linear group**.

Lemma 12.3.2. This group $PGL(2, \mathbb{F}_p)$ acts on the set $\mathbb{P}^1(\mathbb{F}_p)$ by means of the linear fractional transformations.

Remark 12.3.1. *In fact, the action of $\text{PGL}(2, \mathbb{F}_p)$ on the set $\mathbb{P}^1(\mathbb{F}_p)$ is 3-transitive. This 3-transitivity claim is not hard to prove, but I leave it to the interested reader to look up in [R] (see Theorem 9.48).*

Proof: First, we show that the $GL(2, \mathbb{F}_p)$ acts on the set $\mathbb{P}^1(\mathbb{F}_p)$ by means of the linear fractional transformations. In other words, if

$$\phi \begin{pmatrix} a & b \\ c & d \end{pmatrix} (x) = \frac{ax + b}{cx + d}$$

then

$$(a) \quad \phi \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} (x) = x, \text{ for all } x \text{ (i.e., the linear fractional transformation } \phi \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \text{ is the identity map),}$$

$$(b) \quad \phi(A) \circ \phi(B) = \phi(AB), \text{ for all } A, B \in GL(2, \mathbb{F}_p).$$

I leave the verification of (a) to the reader, if he or she wishes, and only check

(b). Let $A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$ and $B = \begin{pmatrix} r & s \\ u & v \end{pmatrix}$. Then

$$AB = \begin{pmatrix} ar + bu & as + bv \\ cr + du & cs + dv \end{pmatrix},$$

so

$$\phi(AB)(x) = \frac{(ar + bu)x + as + bv}{(cr + du)x + cs + dv}.$$

On the other hand, $\phi(A)(\phi(B)(x))$ is equal to

$$\phi(A)\left(\frac{rx + s}{ux + v}\right) = \frac{a\left(\frac{rx + s}{ux + v}\right) + b}{c\left(\frac{rx + s}{ux + v}\right) + d}.$$

Simplifying this, we see that the last two displayed equations are equal. This verifies (b).

Therefore, $GL(2, \mathbb{F}_p)$ acts on the set $\mathbb{P}^1(\mathbb{F}_p)$.

Let $Z = Z(GL(2, \mathbb{F}_p))$. Since $\phi \begin{pmatrix} a & b \\ c & d \end{pmatrix} = \phi \begin{pmatrix} ra & rb \\ rc & rd \end{pmatrix}$, for all non-zero r , it follows that we may define an action of $\text{PGL}(2, \mathbb{F}_p)$ (still denoted by ϕ) on the set $\mathbb{P}^1(\mathbb{F}_p)$ by $\phi(A \cdot Z) = \phi(A)$, for all $A \in GL(2, \mathbb{F}_p)$.

□

Proof (of Theorem 12.3.1): Let ϕ be as above and let f be as in the statement of the theorem.

(\Leftarrow): Since $GL(2, \mathbb{F}_p)$ acts on the set $\mathbb{P}^1(\mathbb{F}_p)$, we have $1 = \phi(AA^{-1}) = \phi(A)\phi(A^{-1})$, so $\phi(A)$ is invertible. This implies that f is a bijection.

(\Rightarrow): We prove the contrapositive. Suppose that $\det \begin{pmatrix} a & b \\ c & d \end{pmatrix} = 0$. The row vectors of this matrix are therefore linearly dependent. In other words, there is

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an $r \in \mathbb{F}_p$ such that either $(a, b) = r \cdot (c, d)$ or $(c, d) = r \cdot (a, b)$. In either case, the quotient $f(x) = \frac{ax+b}{cx+d}$ is a constant independent of x , so cannot be surjective. This proves that f is not a bijection, which verifies the contrapositive.

□

12.3.2 The main isomorphism

Let G denote the Rubik's Cube group. Let H be the subgroup generated by F and U :

$$H = \langle F, U \rangle.$$

Ponderable 12.3.1. *Show that the group H acts on the set of vertices above (via the Rubik's Cube group).*

To better understand the action of H on the vertices, we describe how to label the six vertices on the 'up' and 'front' faces of the cube,

$$fru, flu, dfl, dfr, bru, blu,$$

with the elements in the projective plane

$$\mathbb{P}^1(\mathbb{F}_5) = \{\bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}, \infty\}$$

in a certain way. More precisely, we will label the six vertices above with elements of $\mathbb{P}^1(\mathbb{F}_5)$ in such a way that (a), (b) of the following theorem hold true.

Theorem 12.3.2. *There are $a_0, a_1, b_0, b_1, c_0, c_1, d_0, d_1 \in \mathbb{F}_5$ (given explicitly below) such that*

- (a) *the action of F (the usual rotation of the front face) on these vertices is the same as the action of some linear fractional transformation*

$$f_F(x) = \frac{a_0x + b_0}{c_0x + d_0};$$

- (b) *the action of U (the usual rotation of the up face) on these vertices is the same as the action of some linear fractional transformation*

$$f_U(x) = \frac{a_1x + b_1}{c_1x + d_1}.$$

In other words, the basic moves F, U may be regarded as linear fractional transformations over a finite field!

Theorem 12.3.3. $\text{PGL}(2, \mathbb{F}_5) \cong \langle f_F, f_U \rangle$.

Remark 12.3.2. $\text{PGL}(2, \mathbb{F}_5)$ is isomorphic to S_5 (this is part of Exercise 9.25 in [R]).

We shall prove theorem 12.3.3 below.

12.3.3 The labeling

Label the up, front, and right vertices as in §5.7.2 above.

Let

$$f_F(x) = \frac{x-1}{x+1}, \quad f_U(x) = 3x+3.$$

The map

$$\phi : F \mapsto f_F, \quad U \mapsto f_U,$$

extends to a surjective homomorphism of groups

$$\phi : \langle F, U \rangle \rightarrow \langle f_F, f_U \rangle \subset PGL(2, \mathbb{F}_5).$$

Ponderable 12.3.2. *Verify Theorem 12.3.2 above.*

12.3.4 Proof of $\mathbf{PGL}(2, \mathbb{F}_5) \cong \langle \mathbf{f}_F, \mathbf{f}_U \rangle$

Let

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix}_* \in PGL(2, \mathbb{F}_F)$$

denote the image of

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \in GL(2, \mathbb{F}_F)$$

under the natural map $GL(2, \mathbb{F}_5) \rightarrow PGL(2, \mathbb{F}_5)$, $g \mapsto \mathbb{F}_5^\times \cdot g$.

Since

$$f_U^2 = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}_*$$

we have

$$\begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}_* \in \langle f_U, f_F \rangle.$$

Since

$$f_F \cdot f_U^5 = \begin{pmatrix} -1 & 0 \\ -1 & -1 \end{pmatrix}_*$$

it follows that

$$\begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}_* \text{ belongs to } \langle f_F, f_U \rangle.$$

Conjugating this matrix by f_U^2 , we find that

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$$\begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix}_* \text{ belongs to } \langle f_F, f_U \rangle.$$

It is known that $SL(2, \mathbb{F}_5)$ is generated by elementary transvections (see [R]). Therefore,

$$PSL(2, \mathbb{F}_5) \subset \langle f_F, f_U \rangle \subset PGL(2, \mathbb{F}_5).$$

It is also known (see [R]) that

$$|PSL(2, \mathbb{F}_5)| = 60 \quad \text{and} \quad |PGL(2, \mathbb{F}_5)| = 120.$$

It remains to show that there is an element of $\langle f_F, f_U \rangle$ which does not belong to $PSL(2, \mathbb{F}_5)$. We claim that such an element is f_U . Note that $\det(f_U)$ belongs to the set

$$3(\mathbb{F}_5^\times)^2 = \{3x^2 \mid x \in \mathbb{F}_5^\times\}.$$

But an element of $PSL(2, \mathbb{F}_5)$ must have determinant 1. Since $3^{-1} = 2 \pmod{5}$ is not a square mod 5, there is no element of \mathbb{F}_5 which satisfies $1 = 3x^2$. Thus f_U does not belong to $PSL(2, \mathbb{F}_5)$. \square

Ponderable 12.3.3. *As an application of Theorem 12.3.3, use the example in section 5.7.2 to show that there exists an embedding $D_{12} \hookrightarrow PGL(2, \mathbb{F}_5)$ of the symmetry group of the hexagon into $PGL(2, \mathbb{F}_5)$.*

12.4 The cross groups

Define a **cross move** of the cube to be a move of the form $X \cdot Y^{-1}$, where $X, Y \in \{R, L, U, D, F, B\}$. The subgroup of the Rubik's Cube group generated by the cross moves will be called the **cross group**.

The cross moves permute the set V of vertices of the cube and therefore generate a subgroup of S_V . This is called the **vertex cross group**. The cross moves permute the set E of edges of the cube and therefore generate a subgroup of S_E . This is called the **edge cross group**.

All the entries in the following table are, as far as I am aware, new except for the M_{12} entry.

Rubik polyhedra	Edge cross group	Vertex cross group
Tetrahedron	$A_5 \cong PSL_2(\mathbb{F}_5)$	$C_2 \times C_2$
Cube	A_{12}	$PSL_2(\mathbb{F}_7)$
Octahedron	A_{12}	$PSL_2(\mathbb{F}_5)$
Dodecahedron	A_{30}	A_{20}
Icosahedron	A_{30}	A_{12}
Rubicon	A_{30}	M_{12}

In fact, the subgroup of the dodecahedral edge cross group generated by a subset of the cross moves can yield (smaller but still simple) alternating groups.

Ponderable 12.4.1. (*F. Dyson*) *Work out the analogous cross groups of the ‘Rubicized’ 4-dimensional regular polyhedra.*

Freeman Dyson (1923-), who has contributed important work to both mathematics and physics, is a professor at the Institute for Advanced Studies in Princeton, New Jersey.

Regarding Dyson’s problem, we conjecture that

- if E denotes the set of edges of a 4-dimensional regular polyhedron, then the edge cross group is isomorphic to A_E ;
- if V denotes the set of vertices of a 4-dimensional regular polyhedron, then the vertex cross group is isomorphic to A_V ;
- if F denotes the set of faces of a 4-dimensional regular polyhedron, then the face cross group is isomorphic to A_F .

Some cases of this conjecture have been verified using the computer algebra software package GAP [Gap].

Note that these cross groups are, with only one exception so far, all simple.

Ponderable 12.4.2. *Are any of the analogous cross groups of the ‘Rubicized’ 3-dimensional Archimedean polyhedra simple?*

According to some calculations performed using GAP, the vertex cross group of the regular truncated cube is *not* simple.

12.4.1 $PSL(2, \mathbb{F}_7)$ and crossing the cube

Let C denote the vertex cross group of the Rubik’s Cube.

Theorem 12.4.1. $C \cong PSL_2(\mathbb{F}_7)$.

The first proof of this is by computer!

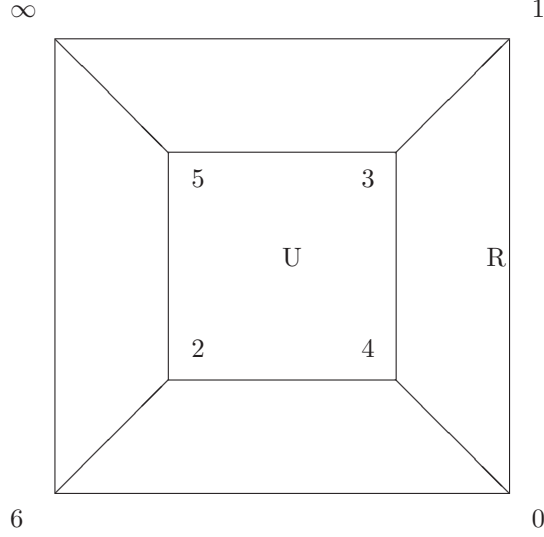
First Proof: GAP [Gap] gives that C is a simple group of order 168. By the classification of simple groups (or, more simply, Exercise 9.26 in [R]), C must be isomorphic to $PSL_2(\mathbb{F}_7)$. \square

The second proof is from [CD].

Second Proof: $PSL_2(\mathbb{F}_7)$ can be generated by the three matrices:

$$f_1 = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix}, f_2 = \begin{pmatrix} 2 & 1 \\ 0 & 1 \end{pmatrix}, f_3 = \begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix}.$$

Label the vertices of the cube as below.



With this labeling, we can show the following:

- The image of the move $m_1 = (UD^{-1})^2$ will permute the vertices in this way: $(\infty, 0)(1, 6)(2, 3)(4, 5)$. The same permutation is given by the Möbius transform $(0 \cdot x - 1)/(x + 0)$ acting on $P^1(\mathbb{F}_7)$.
- The move $m_2 = UR^{-1}$ gives us the permutation $(0, 1, 3)(2, 5, 4)$, which is given by $(2x + 1)/(0 \cdot x + 1)$ acting on $P^1(\mathbb{F}_7)$.
- The move $m_3 = BU^{-1}LB^{-1}$ gives us the permutation $(1, 2, 4)(3, 6, 5)$, which is given by $(2x + 0)/(0 \cdot x + 1)$ acting on $P^1(\mathbb{F}_7)$.

You should notice that if the constants in these Möbius transformations (a, b, c, d) are written in matrix form, they correspond to the generators of $PSL_2(\mathbb{F}_7)$. Now we will define a homomorphism $q : C \rightarrow PSL_2(\mathbb{F}_7)$, such that $q(m_1) = f_1, q(m_2) = f_2, q(m_3) = f_3$. We want to show that our q is an isomorphism.

To do this we will first show that it is surjective. Let f be a matrix in $PSL_2(\mathbb{F}_7)$, which can be written as a product of generators f_1, f_2, f_3 (where $q(m_1) = f_1, q(m_2) = f_2, q(m_3) = f_3$). Now take f as some element of $PSL_2(\mathbb{F}_7)$. The element f can be broken down as a product of its generators, f_1, f_2, f_3 , we'll say

$$f = \prod_{k=1}^n f_{i_k}^{e_k}.$$

Since we have a homomorphism, we can write it as a product of the images of the generators of C . Again we can rewrite it as $f = q(\prod_{k=1}^n m_{i_k}^{e_k})$. Therefore q is surjective.

To show that q is one-to-one we need to know that $PSL_2(\mathbb{F}_7)$ has order 168 [R], and that the order of the cross group is also 168. (This fact was proven by computer, using GAP.) We will prove by contradiction that q is one-to-one.

Now we assume that c_1 and c_2 are elements of C , such that $q(c_1) = q(c_2)$, and $c_1 \neq c_2$. $|PSL_2(\mathbb{F}_7)| = |q(C)|$. We now subtract c_2 from C , and $|q(C)| = |q(C - c_2)|$ because $q(c_1) = q(c_2)$. Now we can say that $|q(C - c_2)| < |C - c_2|$ because we know that q is surjective.

Since we have taken c_2 out of C , we know $|C - c_2| < |C|$, which by transitivity implies $|PSL_2(\mathbb{F}_7)| < |C|$. This is a contradiction because we know $|PSL_2(\mathbb{F}_7)| = |C|$. Therefore q is injective.

Now that we have shown that q is both surjective and injective, it is bijective and an isomorphism. \square

The above proof of the theorem tells us explicitly that there exists a labeling of the vertices V of the cube by the elements of the projective line

$$\mathbb{P}^1(\mathbb{F}) = \{\infty, \bar{0}, \bar{1}, \bar{2}, \bar{3}, \bar{4}, \bar{5}, \bar{6}\},$$

with the property that there is a move $c : V \rightarrow V$ in C if and only if there is a Möbius transformation $f : \mathbb{P}^1(\mathbb{F}) \rightarrow \mathbb{P}^1(\mathbb{F})$ in $PSL_2(\mathbb{F}_7)$.

Because the group of Möbius transformations in $PSL_2(\mathbb{F}_7)$ acts 2-transitively on the projective line $\mathbb{P}^1(\mathbb{F})$ (see [R], Theorem 9.45), it follows that we have the following result.

Corollary 12.4.1. *C acts 2-transitively on V . In other words, for any ordered pairs (v_1, v_2) , (v'_1, v'_2) of distinct vertices there is an element $c \in C$ sending v_i to v'_i , for $i = 1, 2$.*

12.4.2 Klein's 4-group and crossing the Pyraminx

We leave the main result of this section as an exercise—actually more of a project—for the interested reader.

Ponderable 12.4.3. *Show that the subgroup of S_V generated by the twist-untwist moves of the Pyraminx is isomorphic to the Klein 4-group $C_2 \times C_2$.*

The determination of the cross group for the Megaminx is due to J. Conway. It will be presented in chapter 14.

As a kind of a hint, the following **SAGE** code is given. Notice that it produces the 'right' group.

SAGE

```
sage: S4 = SymmetricGroup(4)
sage: F = S4([(1,3,2)]); R = S4([(1,4,3)])
sage: D = S4([(2,3,4)]); L = S4([(1,2,4)])
sage: G = PermutationGroup([F,R,D,L])
sage: G.order()
12
sage: c1 = F*R^-1; c2 = D*R^-1; c3 = L*F^-1
```

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```
sage: c4 = L*R^-1; c5 = D*L^-1; c6 = D*R^-1
sage: C = PermutationGroup([c1,c2,c3,c4,c5,c6])
sage: C.order()
4
sage: KleinFourGroup().is_isomorphic(C)
True
```

Can you solve ponderable 12.4.3 now?

Chapter 13

Other Rubik-like puzzle groups

The youthful impulse is rather to prowl about the problem for awhile, looking perhaps not for an open window or a door with a weak lock, for if these were available some earlier malefactor would have discovered them, but for some wall that can be scaled or some unsuspected underground access.

Robert Langlands

We’ve spent a lot of time thinking about the Rubik’s Cube and how to solve it—we scaled that wall using group theory. What about the Megaminx, or other puzzles? This chapter says something about the group-theoretical structure of some of the other permutation puzzle groups.

13.1 A uniform approach

This section provides a *uniform* discussion of the Pyraminx, the 3×3 Rubik’s Cube, and the Megaminx.

Notation: Let

- G_p (resp., G_R , G_m) denote the permutation puzzle group generated by the basic moves of the Pyraminx (resp., the Rubik’s Cube, Megaminx),
- V_p (resp., V_R , V_m) denote the set of vertex pieces of the Pyraminx (resp., the Rubik’s Cube, Megaminx),
- E_p (resp., E_R , E_m) denote the set of edge pieces of the Pyraminx (resp., the Rubik’s Cube, Megaminx),
- F_p (resp., F_R , F_m) denote the set of facets of the movable pieces of the Pyraminx (resp., the Rubik’s Cube, Megaminx).

13.1.1 General remarks

Let G, V, E, F (resp.) denote either G_p, E_p, V_p, F_p (resp.), or G_R, E_R, V_R, F_R (resp.), or G_m, E_m, V_m, F_m (resp.).

Lemma 13.1.1. *G acts on the set V (resp., E, F).*

If g is any move in G , then, since g acts on the sets V , E , and F , we may regard g

- as an element of the symmetric group S_V of V ,
- as an element of the symmetric group S_E of E , or
- as an element of the symmetric group S_F of F .

These groups S_V , S_E , and S_F are different, so to distinguish these three ways of regarding g , let us write

- g_V for the element of S_V corresponding to g ,
- g_E for the element of S_E corresponding to g ,
- g_F for the element of S_F corresponding to g .

What is the kernel of f_V ? What is its image? To answer this question (actually, we shall not answer this precise question but one similar to it) we introduce a certain subgroup of the symmetric group.

Recall the alternating group A_X is the subgroup of all even permutations of X (in the sense of Example 5.6.3 above).

13.1.2 Parity conditions

Consider the function

$$\begin{aligned} f_{VE} : G &\rightarrow S_V \times S_E \\ g &\longmapsto (g_V, g_E) \end{aligned}$$

It is easy to check that this is a homomorphism.

Theorem 13.1.1. *Let P denote a puzzle which is either the Pyraminx, Megaminx, or (Rubik's) Cube. The image $f_{VE}(G)$ of f_{VE} is isomorphic to*

$$\begin{cases} A_V \times A_E, & \text{if } P \in \{\text{Pyraminx, Megaminx}\}, \\ \{(x, y) \in S_V \times S_E \mid \text{sgn}(x) = \text{sgn}(y)\}, & \text{if } P \in \{\text{Cube}\}, \end{cases}$$

where sgn denotes the sign of the permutation.

This is a consequence of a result proven below in the case of the Rubik's Cube. To see what this theorem means, we look at an example.

Example 13.1.1. Let $G = G_R$.

Question: Can you find a move of the Rubik's Cube which flips a single-edge subcube over, leaving the rest of the puzzle pieces unmoved?

If so, then the image of f_{VE} would have to contain an element (x, y) with $x = 1$ (since moving an edge only does not effect the vertices) and where y is a 2-cycle. But $x = 1$ is even and a 2-cycle is odd. This contradicts the theorem, which says that x, y are either both even or both odd. Therefore, the answer is no: a single-edge flip is impossible.

Next, some more **notation**: let

$$K = \ker(f_{VE}) \triangleleft G$$

denote the kernel of the map f_{VE} introduced above. This is a normal subgroup of G .

Example 13.1.2. In the case of the Rubik's Cube, this subgroup K is the set of moves which may reorient (i.e., flip or rotate) a subcube but does not swap it with some other subcube. For example, the move

$$(R^{-1} \cdot D^2 \cdot R \cdot B^{-1} \cdot U^2 \cdot B)^2,$$

which twists the ufr corner clockwise and the bld corner counterclockwise, belongs to K .

Theorem 13.1.2. (Gold and Turner [GT]) G is a semi-direct product of K with $f_{VE}(G)$.

This is a consequence of a result proven above in the case of the Rubik's Cube. In the case of the $3 \times 3 \times 3$ Rubik's Cube, some more details are given in chapter 11. See also [GT], chapter 2 of [B1], or [NST], chapter 19.

13.2 On the group structure of the Skewb

Notation: We fix an orientation of the cube and label the sides by R, L, U, D, F, B as in the case of the Rubik's Cube. The 120 degree clockwise rotation of a corner is denoted by a 3-letter juxtaposition of the letters abbreviating the 3 faces which the corner meets. (When you twist a corner of the Skewb you must permute three other corners but the opposite side of the Skewb is unaffected.) Such a move will be called a **basic move**; there are 8 of them, though twisting about a corner and twisting about the antipodal opposite corner is basically the same move (up to a rotation of the entire cube.) For example, FRU denotes the 120 degree clockwise rotation of the front-right-up corner, leaving the rest of the cube alone.

Let C denote the set of square center facets and V the set of vertices of the cube.

Let

$$G = \langle FRU, FLU, BRU, BLU, DFR, DFL, BDR, BDL \rangle$$

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denote the group of all (legal) Skewb moves. Let G^* denote the group of all legal and ‘illegal’ moves (where disassembly and then reassembly is allowed). The commands for entering this into **SAGE** were given in §4.6.

On each square center facet of the Skewb we may choose a vertex with the following property: if we draw an arrow pointing from the chosen vertex of the square to the diametrically opposite vertex on the square, then the moves of the Skewb permute these arrows among themselves, except that some arrows may possibly be reversed. This determines an orientation of each center facet. Call this puzzle the **super Skewb**. For this new puzzle, let

$$G_{super} = \langle FRU, FLU, BRU, BLU, DFR, DFL, BDR, BDL \rangle$$

denote the group of all (legal) super Skewb moves. Let G_{super}^* denote the group of all legal and ‘illegal’ moves (where disassembly is allowed).

We orient the corners as in the case for the Rubik’s Cube. Let $y(g) \in C_3^8 = \{0, 1, 2\}^8$ denote the orientation for the corners.

For the super Skewb, we orient the center facets similarly. Let $z(g) \in C_4^6 = \{0, 1, 2, 3\}^6$ denote the orientation for the centers.

Let S_C denote the symmetric group on the set C and S_V the symmetric group on the set V .

Claim: There are homomorphisms

$$\rho : G \rightarrow S_C, \quad \sigma : G \rightarrow S_V,$$

given, for each move $g \in G$, by

$$\rho(g) = \text{permutation of the center facets associated to } g,$$

and

$$\sigma(g) = \text{permutation of the vertices associated to } g.$$

Let

$$H = C_3^8 \times S_C \times S_V$$

and define $\cdot : H \times H \rightarrow H$ by

$$(\vec{y}, r, s) \cdot (\vec{y}', r', s') = (rr', ss', r^{-1}(\vec{y}') + \vec{y}).$$

Let

$$H_{super} = C_3^8 \times S_V \times C_4^6 \times S_C$$

and define $\cdot : H_{super} \times H_{super} \rightarrow H_{super}$ by

$$(\vec{y}, r, \vec{z}, s) \cdot (\vec{y}', r', \vec{z}', s') = (r^{-1}(\vec{y}') + \vec{y}, rr', s^{-1}(\vec{z}') + \vec{z}, ss').$$

Observation: There is an embedding of G into H and an embedding of G_{super} into H_{super} .

Let G_C be the group that acts only on the center facets of the Skewb, and G_V the group that acts only on the vertices. Now,

$$G = G_C \times G_V.$$

Every generator of G is a 3-cycle on the center facets. This means that r is an element of A_C . It is a fact that the elements (i, j, k) of S_n generate A_n (for i, j, k elements of $\{1, 2, \dots, n\}$). Therefore, $G_C = A_6$.

The group that acts on the vertices of the Skewb is slightly more complicated. Unlike the Rubik's Cube, there is no condition like conservation of twists which applies to the entire vertex set. Instead, we must split the vertices of the Skewb into two 4-corner orbits. This idea is borrowed from Christoph Bandelow's booklet [B2] on Mickey's Challenge (a puzzle similar to the Skewb). An orbit is constructed by starting with one corner and including the opposite corner of each face that meets at the first corner. Referring back to our original labeling of the Skewb, the orbits are the odd corners $\{1, 3, 5, 7\}$, and the even corners $\{2, 4, 6, 8\}$. Let the orbit of odd corners be denoted by $V(\text{odd})$, and let $V(\text{even})$ denote the orbit of even corners. We now partition G_V so that

$$G_V = G_{V(\text{odd})} \times G_{V(\text{even})}.$$

We know that each orbit maps to a permutation on 4 vertices and an orientation on 4 vertices. So

$$G_V \text{ is a subgroup of } (C_3^4 \rtimes S_4) \times (C_3^4 \rtimes S_4).$$

Let $h = (s, u(h), t, v(h))$ be an element of G_V .

First we will examine the permutations of the vertices of both orbits. Each generator produces a 3-cycle on the vertices, whether they are in the odd or even orbit. Therefore, we can say that the permutations of each orbit generate A_4 by the same argument used for the center permutations. Now,

$$G_V \text{ is a subgroup of } (C_3^4 \rtimes A_4) \times (C_3^4 \rtimes A_4).$$

Claim 1: There exist h , such that s is an element of A_4 , $u(h) = (0, 0, 0, 0)$, t is an element of A_4 , and $v(h) = (0, 0, 0, 0)$. We know this is true because there are Skewb moves which permute only 3 vertices (and no edges).

Claim 2: Given any permutation u' of $(1, 2, 0, 0)$ by an element of A_4 and any permutation v' of $(1, 2, 0, 0)$ by an element of A_4 , there exist h , such that $s = 1$, $u(h) = u'$, $t = 1$, and $v(h) = v'$. This is true because there are Skewb moves which twist only vertices (and no edges).

If we combine the moves of Claims 1 and 2, we should generate all of the possible moves of G_V . The condition on each of the vertex 4-tuples will drop them in dimension to elements of C_3^3 . So we can conclude that G_V is a subgroup of index 9 of

$$(C_3^4 \rtimes A_4) \times (C_3^4 \rtimes A_4).$$

Note: This claim is verified by GAP.

Since $G_C = A_6$, and $G_V = (C_3^3 \rtimes A_4) \times (C_3^3 \rtimes A_4)$, we can conclude that

$$G = A_6 \times (C_3^3 \rtimes A_4) \times (C_3^3 \rtimes A_4)$$

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and

$$|G| = (6!/2)(4!/2)^2 3^6 = 37,791,360.$$

In conclusion, it is interesting to note that if we let G' denote the illegal Skewb group—where reassembly is permitted—then

$$G' = S_6 \times S_8 \times C_3^8.$$

and

$$|G|/|G'| = 0.0001984127...$$

This means that if you could take apart the Skewb and reassemble it however you wanted (leaving the stickers intact), then only about 0.02 percent of all possible reassemblies would be solvable. The analogous percentage for the Rubik's Cube is about 8.3 percent. As Martin Schönert points out in a post to [CL], this makes the Skewb harder to solve than the Rubik's Cube, in some sense.

Permutation and Orientation Tables

Move	Center permutation	Vertex permutation
UFR	(1, 5, 2)	(2, 6, 4)
UFL	(1, 4, 5)	(1, 7, 3)
DFR	(1, 2, 6)	(1, 5, 7)
DFL	(1, 6, 4)	(4, 6, 8)
BRU	(2, 5, 3)	(1, 3, 5)
BLU	(3, 5, 4)	(2, 4, 8)
DBR	(2, 3, 6)	(2, 8, 6)
DBL	(3, 4, 6)	(3, 7, 5)

Move	Vertex orientation
UFR	(1, 2, 0, 2, 0, 2, 0, 0)
UFL	(2, 0, 2, 1, 0, 0, 2, 0)
DFR	(2, 0, 0, 0, 2, 1, 2, 0)
DFL	(0, 0, 0, 2, 0, 2, 1, 2)
BRU	(2, 1, 2, 0, 2, 0, 0, 0)
BLU	(0, 2, 1, 2, 0, 0, 0, 2)
DBR	(0, 2, 0, 0, 1, 2, 0, 2)
DBL	(0, 0, 2, 0, 2, 0, 2, 1)
UFR·UFL	(0, 2, 2, 2, 0, 0, 2, 0)
DFR·DFL	(2, 0, 0, 2, 0, 0, 2, 2)

Note: The orientations for the generator moves contain two repeated orbits—permutations of (1, 0, 0, 0) and permutations of (2, 2, 2, 0).

This section owes much to Martin Schönert's contribution on this topic in [CL].

13.3 Mathematical description of the $2 \times 2 \times 2$ cube

This section derives the group structure of the $2 \times 2 \times 2$ Rubik's Cube.

A position on the $2 \times 2 \times 2$ cube is determined by

- (a) a permutation of the vertices, and
- (b) the orientation of the corner sub-cubes.

An **illegal move** on the $2 \times 2 \times 2$ cube is a reassembly of the corners.

Let

$$H = \langle R, L, U, D, F, B, \text{ and all the illegal moves} \rangle.$$

This will be called the **illegal $2 \times 2 \times 2$ cube group**. Let $G = \langle R, L, U, D, F, B \rangle$. G is contained in H .

Let $C_3^8 = \{0, 1, 2\}^8$ be the group of 8-tuples with coordinate-wise addition mod 3. Let $\vec{v} : H \rightarrow C_3^8$ be defined as follows: Assume $h \in H$ sends the i th corner to the j th corner. Define $v_i(h)$ to be that number in $C_3 = \{0, 1, 2\}$ which describes the orientation that the standard reference marking of the i th corner is sent to relative to the standard reference marking of the j th corner. The values of v are tabulated in (11.1.1).

Let S_V be the group of permutation of corner sub-cubes. We may identify S_V with S_8 since we have labeled the corners $1, \dots, 8$. The group H is a subset of the Cartesian product $S_V \times C_3^8$.

Let $p(h)$ denote the permutation of the vertices of the cube associated with $h \in H$. We have

$$\begin{aligned} (v, r) \cdot (v', r') &= (v + r(v'), r \cdot r') \\ (\vec{v}(g), p(g)) \cdot (\vec{v}(h), p(h)) &= (\vec{v}(g) + p(g)\vec{v}(h), p(g) \cdot p(h)) \\ &= (\vec{v}(g \cdot h), p(g \cdot h)). \end{aligned}$$

It is not hard to show, based on the results of the previous section, that $H = C_3^8 \rtimes S_8 = \{(v, r) \mid r \in S_V, v \in C_3^8\}$. In other words, H is the wreath product of S_8 and C_3 (with base C_3).

Theorem 13.3.1. *A pair $(\vec{v}, r) \in C_3^8 \times S_V$ corresponds to a legal position if and only if $v_1 + \dots + v_8 \equiv 0 \pmod{3}$ (conservation of twists).*

Proof: PART 1: In this part, we show that any pair (v, r) as in the theorem (where v satisfies conservation of twists) corresponds to a legal move g in such a way that $r = \rho(g)$ and $v = \vec{v}(g)$.

Case 1: Assume $r = 1$ and \vec{v} is arbitrary. From the solved position, any two corners, corner i and corner j say, can be twisted so that corner i has orientation 1, corner j has orientation 2, and all other corners have orientation 0. Call such a move $e_{i,j}$. Example: $(R^{-1} \cdot D^2 \cdot R \cdot B^{-1} \cdot U^2 \cdot B)^2$ is $e_{2,7}$.

Let $y = a_1 \cdot e_{1,8} + \dots + a_7 \cdot e_{7,8}$, where $a_i \in \{0, 1, 2\}$. This is a move of the 2×2 Rubik's Cube of the form $(v, 1)$ —in other words, it permutes nothing but

may twist some corners. By construction, all moves of this form are legal. For each $a_i \cdot e_{i,8}$ there are three different possible positions (independent of all other $a_j \cdot e_{j,8}$). Since there are three choices for each a_i , there are a total of 3^7 distinct moves of the form y as above.

On the other hand, there are exactly 3^7 possible moves of the form $(v, 1)$ which satisfy the conservation of twists. (Proof: If $\vec{v} = (v_1, \dots, v_8)$, $v_i \in \{0, 1, 2\}$, and $v_1 + \dots + v_8 \equiv 0 \pmod{3}$, then there are 3 ways to choose each of v_1, \dots, v_7 , but then once these are fixed, the conservation of twists condition leaves no choice for v_8 . This leaves a total of 3^7 choices.) These 3^7 possible moves include, of course, the legal moves of the form y above. Thus every move of the form $(1, v)$, with v satisfying conservation of twists, is legal.

Case 2: Assume $\vec{v} = \vec{0}$ and r is arbitrary. Recall S_8 is generated by the two-cycles (see Steinhaus's Theorem 3.4.1 above).

Claim 1: Given any pair of corners, there is a 2-cycle move which swaps them. (Example: $F^{-1} \cdot U \cdot B \cdot U^{-1} \cdot F \cdot U^2 \cdot B^{-1} \cdot U \cdot B \cdot U^2 \cdot B^{-1}$.) Once two corners have been swapped, you may correct the orientation of any sub-cube by Case 1. Thus, any permutation which preserves orientations is a legal move.

Case 3: Assume \vec{v} and r are both arbitrary but satisfy conservation of twists. By case 2, we may make a legal move that changes (\vec{v}, r) to $(\vec{v}, 1)$. By case 1, $(\vec{v}, 1)$ is a legal move.

PART 2: In this part, we show that any legal move satisfies conservation of twists.

Assume $(\vec{v}, r) \in C_3^8 \times S_V$ is a legal move.

Define the **length** of a move $g \in G$ to be the smallest number (counted according to multiplicity) n of generators needed to create the move, written $\text{length}(g) = n$.

Induction hypothesis: If a move is length n , it satisfies conservation of twists.

Step $n = 1$: Every $\vec{v}(x)$ where $x \in \{R, L, U, D, F, B\}$ satisfies the conservation of twists.

Step $n > 1$: Assume the induction hypothesis is true for all lengths $\leq n-1$. Let x be length n and write $x = x_1 \cdot x_2$, where $\text{length}(x_1) \leq n-1$ and $\text{length}(x_2) = 1$. Then $\vec{v}(x) = \vec{v}(x_1) + p(x_1)\vec{v}(x_2)$, by the group operation. Furthermore, $\vec{v}(x_1)$ satisfies conservation of twists by the induction hypothesis. Since $p(x_1)$ simply permutes the coordinates of $\vec{v}(x_2)$, $\vec{v}(x_2)$ still satisfies the conservation of twists. The sum of moves satisfying conservation of twists still satisfies conservation of twists.

Conclusion: By induction, any move $(v, r) \in C_3^8 \times S_V$ that is a legal move satisfies the conservation of twists. \square

13.4 On the group structure of the Pyraminx

Notation: Let

- V denote the vertices of the tetrahedron (which we identify with the set of corner pieces of the Pyraminx),

- E denote the edges of the tetrahedron (which we identify with the set of edge pieces of the Pyraminx),
- C denote the set of interior pieces of the tetrahedron (i.e., movable pieces of the Pyraminx not in E or V),
- S_V denote the permutation group of V ,
- A_V denote the alternating group of V ,
- S_E denote the permutation group of E ,
- A_E denote the alternating group of E .

Assume that the tetrahedron is lying on a flat surface in front of you, with the triangle base pointing away from you. The corners are denoted L (left), R (right), U (up), and B (back).

Basic Moves: Opposite each corner or vertex there are three **layers**: the **tip**, the **middle layer**, and the **opposite face**. Let

- l denote the 120 degree clockwise rotation of the tip containing the left corner,
- L denote the 120 degree clockwise rotation of the tip/middle layer containing the left corner,
- r denote the 120 degree clockwise rotation of the tip containing the right corner,
- R denote the 120 degree clockwise rotation of the tip/middle layer containing the right corner,
- u denote the 120 degree clockwise rotation of the tip containing the up corner,
- U denote the 120 degree clockwise rotation of the tip/middle layer containing the up corner,
- b denote the 120 degree clockwise rotation of the tip containing the back corner,
- B denote the 120 degree clockwise rotation of the tip/middle layer containing the back corner.

Let $G = \langle R, L, U, B, r, l, u, b \rangle$ denote the **Pyraminx group**.

Each move $g \in G$ induces a permutation of E denoted $\sigma(g)$. Note that G does not permute the vertices. Furthermore, the tip moves r, l, u, b do not affect the edges.

Lemma 13.4.1. $\sigma : G \rightarrow S_E$ is a group homomorphism.

Example 13.4.1. $\rho(L)$ is a 3-cycle in S_V , $\sigma(L)$ is a 3-cycle in S_E .

13.4.1 Orientations

Assume for the moment that the Pyraminx is fixed in space as above and is in the ‘solved’ position. For each corner or edge piece, choose once and for all one facet on that piece. There are three possible choices for each corner piece and two for the edges. Mark each of these chosen facets with an imaginary ‘+’, leaving the other facets unmarked. For the rest of this section, we shall make the following choices for the marked facets (with reference to the numbering in §4.4):

- marked edge facets: 4, 6, 10, 15, 20, 25
- marked corner facets: 1, 13, 17, 23

For each edge piece, assign to a move $g \in G$ either

- a ‘0’ if the ‘+ facet’ for that piece when it was in the solved position is sent to the ‘+ facet’ for that piece in the present position,
- a ‘1’ otherwise.

Since there are 6 edge pieces, this yields a 6-tuple of 0’s and 1’s: $\vec{w}(g) = (w_1, w_2, \dots, w_6)$.

Example 13.4.2. *We compute the effect of the basic twist moves on the edge orientations in Figure 13.1.*

X	$\vec{w}(X)$
B	(0,0,0,1,0,1)
R	(0,0,1,0,1,0)
L	(1,0,1,0,0,0)
U	(0,1,0,1,0,0)
$R \cdot U^{-1} \cdot R^{-1} \cdot U$	(1,1,0,0,0,0)

Figure 13.1: Corner orientations for the Pyraminx

For each corner piece, assign to a move $g \in G$ either

- a ‘0’ if the ‘+ facet’ for that piece when it was in the solved position is sent to the ‘+ facet’ for that piece in the present position,
- a ‘1’ if the ‘+ facet’ for that piece when it was in the solved position is sent to the facet which is a 120 degree rotation about its vertex from the ‘+ facet’ for that piece in the present position,
- a ‘2’ otherwise.

Since there are 4 corner pieces, this yields a 4-tuple of 0’s, 1’s, and 2’s: $\vec{v}(g) = (v_1, v_2, v_3, v_4)$.

X	$\vec{v}(X)$
B	$(0,0,0,2)$
R	$(0,0,2,0)$
L	$(0,2,0,0)$
U	$(2,0,0,0)$

Figure 13.2: Corner orientations for the Pyraminx

Example 13.4.3. *We compute the effect of the basic twist moves on the corner orientations in Figure 13.2.*

Proposition 13.4.1. *If $\vec{w}(g) = (w_1, w_2, \dots, w_6)$ corresponds to a move $g \in G$, then*

$$w_1 + w_2 + \dots + w_6 \equiv 0 \pmod{2}.$$

Observation: There is no corresponding condition for the v_1, \dots, v_4 , since corner moves move them around freely.

Proof: The proof uses the following lemma, but is otherwise essentially the same as the corresponding fact (Theorem 11.2.1 (c)) which we proved for the Rubik's Cube. The modifications required for the proof are left to students as an exercise to test their understanding of the argument. \square

Lemma 13.4.2. *For $g, h \in G$, we have*

$$\vec{w}(g \cdot h) = \sigma(g)^{-1}(\vec{w}(h)) + \vec{w}(g).$$

Proof: The argument for the proof is essentially the same as the corresponding fact (Lemma 11.1.1) which we proved for the Rubik's Cube. The modifications required for the proof are left to the interested reader. \square

Lemma 13.4.3. *For $g, h \in G$, we have*

$$\vec{v}(g \cdot h) = \vec{v}(h) + \vec{v}(g).$$

Ponderable 13.4.1. *Try to verify this lemma.*

Let H denote the **illegal Pyraminx group** generated by G and the ‘illegal edge moves’ (that is, one may physically remove the edge pieces and reassemble the Pyraminx). Illegal center or corner moves are not allowed in H . Let

$$H^* = \{(s, v, w) \mid r \in S_V, s \in S_E, v \in C_3^4, w \in C_2^6\}$$

and define $\cdot : H^* \times H^* \rightarrow H^*$ by

$$(s, v, w) \cdot (s', v', w') = (s \cdot s', v + v', s'(w) + w').$$

Theorem 13.4.1.

- H^* is, with this operation, a group.

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- *There is an isomorphism*

$$H \cong H^* \cong C_3^4 \times (C_2^6 \rtimes S_E),$$

and hence between H and the direct product of the tip moves C_3^4 and the wreath product

$$C_3^4 \times (C_2 \text{ wr } S_E).$$

- *The map $G \rightarrow H^*$ defined by*

$$g \longmapsto (\sigma(g), v(g), w(g)),$$

is a homomorphism.

13.4.2 Center pieces

Each corner piece has 3 center pieces neighboring it. Facts:

- The center pieces associated with a corner can never be moved into a middle layer associated with another corner.
- The center pieces associated with a corner can always be color-aligned with the colors of the corner piece by a corner twist move.

The last fact says, in other words, that the center pieces can always be ‘solved’ by a corner piece.

13.4.3 The group structure

Theorem 13.4.2. *G is isomorphic to*

$$\{(s, v, w) \in H^* \mid s \text{ even}, w_1 + w_2 + \dots + w_6 \equiv 0 \pmod{2}\}.$$

Proof: The idea to prove this is to show that

- $A_E = \langle \sigma(R), \sigma(L), \sigma(U), \sigma(B) \rangle$,
- the map

$$\begin{aligned} G &\rightarrow A_E, \\ g &\longmapsto \sigma(g), \end{aligned}$$

is surjective, and the map

- $G \rightarrow \{(w_1, \dots, w_6) \in C_2^6 \mid w_1 + \dots + w_6 \equiv 0 \pmod{2}\}, g \longmapsto \vec{w}(g)$, is surjective (as a map of sets).

Here's the proof of the first point: We can label (as in §5.10) the edges $1, 2, \dots, 6$ so that the edges on the front face are $1, 2, 3$, resp., where the fl edge is 1 , the fr edge is 2 , and the fd edge is 3 (here f, r, l , and d denote the front face, right face, left face, and down face, resp.). The move $[R, U^{-1}] = R \cdot U^{-1} \cdot R^{-1} \cdot U$ is the counterclockwise cyclic permutation $(1, 3, 2)$ of the edges on the f face. (This move does not affect any corners but does flip some edges, a fact which we may ignore for now since we are concerned only with the permutations now.) In particular, $(1, 3, 2)$ may be written as a product of the generators in $\{\sigma(R), \sigma(L), \sigma(U), \sigma(B)\}$. Now pick any $i \in \{4, 5, 6\}$ and let $s \in G$ denote a move which sends edge i to edge 2 and does not move edge 1 or 3 . The move $s \cdot [R, U^{-1}] \cdot s^{-1}$ is the 3-cycle $(1, 3, i)$. It does not affect any corners or other edges. By Lemma 9.4.4, these permutations generate $A_6 \cong A_E$.

The second point follows immediately from the first point proven above.

Here's the proof of the last point: The move $g = [R, U^{-1}]$ has the following effect on the orientation: $w(g) = (1, 1, 0, 0, 0, 0)$. The group

$$\{(w_1, \dots, w_6) \in C_2^6 \mid w_1 + \dots + w_6 \equiv 0 \pmod{2}\} \cong C_2^5$$

is a vector space over \mathbb{F}_2 . The 5 vectors listed in the table of Figure 13.1 for the values for w are all independent. It follows from this and the group law for G that the map $g \mapsto \vec{w}(g)$ is surjective.

Theorem 13.4.2 above is thus proven. \square

13.5 The homotopy group of the Square 1

The Square 1 puzzle is described, for example, on Jaap Scherphuis's web page [Sch2]. It is cube-shaped (at least in the solved position), with three parallel slices from top to bottom, just like the Rubik's Cube has. However, the vertical slices are not parallel to the faces, nor to each other. This means that after some moves the puzzle is no longer cube-shaped.

Here we study the group theoretic properties of the collection G of all 'words' in the basic moves of the Square 1 puzzle which preserve the cube shape. This collection forms a group which, motivated by Wilson [W], we call the **homotopy group of the Square 1 puzzle**. The list of shapes which the Square 1 puzzle can make is given in [Sch2] (also [Sn2], though it is out of print). There are 10 such shapes, including the cube shape. It is not hard to see that the homotopy group of any one of these other shapes is conjugate to G_{\square} .

A by-product of the proof is an collection of moves which can be used to solve the Square 1 puzzle once it is put in the cube shape.

13.5.1 The main result

Let S_n denote the symmetric group of degree n , i.e., the group of permutations of $\{1, 2, \dots, n\}$. Let $\text{sgn} : S_n \rightarrow \{\pm 1\}$ denote the homomorphism which assigns to each permutation its sign, as in Definition 3.1.1 (the sign of a cyclic permutation of length r is $(-1)^{r+1}$, for example).

13.5. THE HOMOTOPY GROUP OF THE SQUARE 1

Let G_{\square} denote the **illegal Square 1 group** generated by all legal moves preserving the cube shape and all illegal moves (i.e., disassembly and reassembly is allowed) preserving the cube shape. It is clear that

$$G_{\square} \cong S_8 \times S_8.$$

We shall see that the size of the homotopy group G of the Square 1 is about 0.8 billion.

Theorem 13.5.1. *Let G_{\square} be as above and let G be the subgroup generated by the ‘square moves’ d^3 , u^3 , $B(u^3, 1)$, $B(1, d^3)$, $T(u^3, 1)$, and $T(1, d^3)$, as in §13.5.2 below. The group G_{\square} is isomorphic to $S_8 \times S_8$, and G is isomorphic to the kernel of index 2 in $G_{\square} \cong S_8 \times S_8$ of the homomorphism $f : S_8 \times S_8 \rightarrow \{\pm 1\}$ defined by $f(g_1, g_2) = \text{sgn}(g_1)\text{sgn}(g_2)$. Consequently, $|G_{\square}| = 2^{14}3^45^27^2 = 1625702400$ and $|G| = 2^{13}3^45^27^2 = 812851200$.*

As a corollary of the proof of this theorem, given below, we shall see that any even permutation of the corners is possible and any even permutation of the wedges is possible.

13.5.2 Some notation

We shall assume that the puzzle is in the solved position with the ‘Square 1’ side in front, right side up.

Definition 13.5.1. *Let*

- u denote rotation of the up face by 30 degrees clockwise,
- d denote rotation of the down face by 30 degrees clockwise,
- R denote rotation of the cube by 180 degrees though one of the skew-diagonal cuts (in a given position, at most one such move is possible, so this is unambiguous).

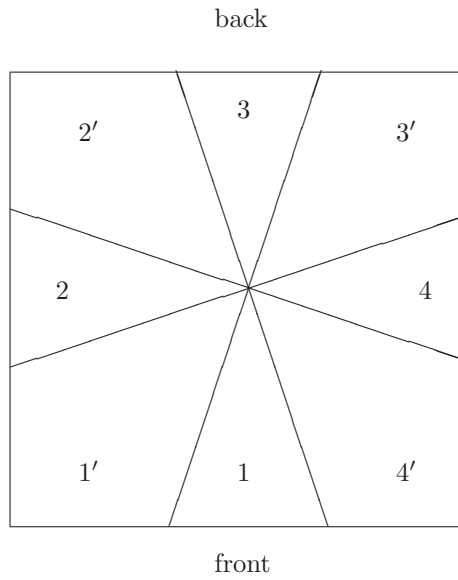
Like the 15 Puzzle, and unlike the Rubik’s Cube, not any sequence of u, d , and R ’s is possible.

Let

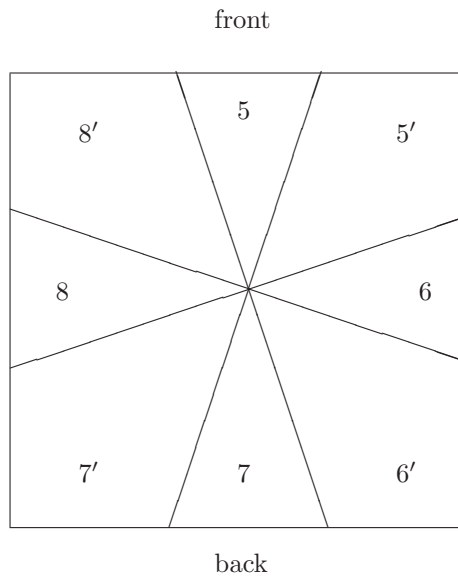
$$T(x, y) = u \cdot R \cdot x \cdot y \cdot R \cdot u^{-1}, \quad B(x, y) = d^{-1} \cdot R \cdot x \cdot y \cdot R \cdot d,$$

where x, y are moves of the Square 1 puzzle.

13.5. THE HOMOTOPY GROUP OF THE SQUARE 1



The up facets of the 'Square 1' puzzle.



The down facets of the 'Square 1' puzzle.

In the notation of these diagrams, we have

$$\begin{aligned} uRu^{-1}d^{-1}Rd &= (2, 8)(4, 6) \\ T(u^3, 1) &= (1', 6', 7', 4')(1, 6, 7, 4) \\ T(1, d^3) &= (2', 3', 8', 5')(2, 3, 8, 5) \\ B(u^3, 1) &= (1, 2, 7, 8)(1', 6', 7', 4') \\ B(1, d^3) &= (3, 4, 5, 6)(2', 3', 8', 5'). \end{aligned}$$

13.5.3 Two subgroups

Let

$$G_u = \langle T(u^3, 1), T(1, d^3) \rangle$$

and

$$G_d = \langle B(u^3, 1), B(1, d^3) \rangle.$$

Lemma 13.5.1. *G_u and G_d are each isomorphic to $C_4 \times C_4$.*

Proof: We have $T(u^3, 1)T(1, d^3) = T(1, d^3)T(u^3, 1)$. Moreover, $T(u^3, 1)$ and $T(1, d^3)$ are each of order 4. Since

$$C_4 \times C_4 = \langle a, b \mid a^4 = 1, b^4 = 1, ab = ba \rangle,$$

the lemma follows. \square

The homotopy group of the Square 1 puzzle is

$$G = \langle d^3, u^3, B(u^3, 1), B(1, d^3), T(u^3, 1), T(1, d^3) \rangle.$$

We use the labelings in Definition 13.5.1 to describe the moves of the Square 1 puzzle.

13.5.4 Proof of the theorem

We shall prove the theorem in the following steps:

- Show that the wedge 3-cycle $(1, 2, 3)$ and the corner 3-cycle $(1', 2', 3')$ each belong to G .
- Show that any wedge 3-cycle $(1, 2, i)$ and each corner 3-cycle $(1', 2', i')$ belong to G .
- Show that there is an injective homomorphism $\phi : G \rightarrow S_8 \times S_8$ where the image $\phi(G)$ contains $A_8 \times A_8$.
- Conclude that $G \cong S_8 \times S_8 / \{\pm 1\}$.

Step 1: First, we claim that $(1, 2, 3)$ belongs to G . In fact, the 3-cycle $(1, 2, 3)$ is obtained from the move M_1 , defined to be

$$(B(u^3, 1) \cdot d^3) \cdot ((B(u^3, 1) \cdot d^{-3}) \cdot (B(u^{-3}, 1) \cdot T(1, d^{-3}) \cdot d^6))^4 \cdot (B(u^3, 1) \cdot d^3)^{-1}.$$

(Incidentally, this 80-move maneuver may be verified using **SAGE** or **GAP** [Gap]. See also [Sn2].)

Next, we claim that $(1', 2', 3')$ belongs to G . In fact,

$$M_2 = Ru^3 Rd^{-3} Ru^3 (Ru^{-3})^2 d^3 Ru^{-3}$$

is the product of 2-cycles $(2', 3')(3, 4)$. (This move was found in [Sn2].) Therefore, $u^3 M_2 u^{-3}$ is the product of 2-cycles $(1', 2')(2, 3)$. The product of these is $(1', 2', 3')(2, 3, 4)$. Since $(2, 3, 4)$ is obtained from $u^{-3} M_1 u^3$, we see that $(1', 2', 3')$ is in G .

Step 2: Let g be any move in G which sends wedges 3 to wedge i , resp., and does not move wedges 1, 2 (it may permute other wedges and corners). Then $(1, 2, i) = g \cdot (1, 2, 3) \cdot g^{-1}$. Thus $(1, 2, i) \in G$.

The proof that each $(1', 2', i') \in G$ is similar.

Step 3: It is clear from our definition that there is an injection $G \rightarrow S_8 \times S_8$ as sets. The verification that this is a homomorphism is straightforward.

Step 4: The group A_8 is generated by the 3-cycles $(1, 2, i)$ (see Lemma 9.4.4). Since these all belong to G_\square , all even wedge permutations are possible. Similarly, all even corner permutations are possible. Thus $A_8 \times A_8 \subset G$.

Let $p_1 : S_8 \times S_8 \rightarrow S_8$ denote the projection onto the first factor. Let p_2 denote the projection onto the second factor. For each generator $g \in \{d^3, u^3, B(u^3, 1), B(1, d^3), T(u^3, 1), T(1, d^3)\}$ of G we have $\text{sgn}(p_1(g)) = \text{sgn}(p_2(g))$. Thus the image $\phi(G)$ is strictly contained in $S_8 \times S_8$. In fact, this shows that $\phi(G)$ is contained in the kernel $\ker(f)$ of the homomorphism $f : S_8 \times S_8 \rightarrow \{\pm 1\}$ defined in the statement of the theorem. Since

$$A_8 \times A_8 \subset G \subset \ker(f),$$

$[\ker(f) : A_8 \times A_8] = 2$, and $T(u^3, 1) \notin A_8 \times A_8$, the theorem follows. \square

13.6 The Masterball group

Recall the notation from §4.3. The Masterball group,

$$G = \langle f_1, f_2, f_3, f_4, f_5, f_6, f_7, f_8, r_1, r_2, r_3, r_4 \rangle,$$

is a subset of S_{32} . We can provide a simple proof to show that G is a subset of S_{32} , by showing the following result.

Theorem 13.6.1. $G = S_{16} \times S_{16}$.

Proof: Since the Masterball is made up of 32 facets, the permutation group G , on these facets, must be a subset of S_{32} . If we agree that we can take the ball apart and put the pieces back in any position, we have generated S_{32} , saying that any swap is legal. However, as stated before, equator facets can be swapped only with equator facets and polar facets can be swapped only with polar facets. Since there are 16 polar facets, the group of polar swaps, G_p is a subset of S_{16} .

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Similarly, since there are 16 equatorial facets, the group of equatorial swaps, G_e , is a subset of S_{16} . To prove that G_p and G_e do each generate S_{16} , we need only to recall Steinhaus's Theorem 3.4.1.

If then we can provide a 2-cycle which swaps any two polar facets and a 2-cycle which swaps any two equatorial facets (neglecting setup moves), we have shown that G_p and G_e , both products of 2-cycles, have generated S_{16} . The following polar swap can be represented by the 2-cycle, (18, 48):

$$\begin{aligned} & r_4 r_1 f_1 r_1 f_1 r_4 f_1 r_1^{-1} f_1 r_4^{-2} f_1 r_1 f_1 r_4 f_1 r_1^{-1} f_1 r_4^4 f_1 r_1 f_1 r_4 \times \\ & \times f_1 r_1^{-1} f_1 r_4^{-2} f_1 r_1 f_1 r_4^2 f_1 r_1^{-1} f_1 r_4^2 f_1 r_1 f_1 r_4^{-1} f_1 r_1^{-1} f_1 r_4^{-1} f_1 \times \\ & \times r_1 f_1 r_4 f_1 r_1^{-1} f_1 r_1 r_4^{-2} f_1 r_1^{-1} f_1 r_4^{-1} f_1 r_1 f_1 r_1^{-1} r_4 f_1 r_1 f_1 r_1^{-1} \times \\ & \times f_1 r_1 f_1 r_4^2 f_1 r_1^{-1} f_1 r_1 f_1 r_1^{-1} f_1 r_1 r_4^{-1} f_1 r_1^{-1} f_1 r_4 f_1 r_1 f_1 r_1^{-2} r_4^{-2}. \end{aligned}$$

The following equator swap can be represented by the 2-cycle, (28, 38):

$$\begin{aligned} & r_3 r_2 f_1 r_2 f_1 r_3 f_1 r_2^{-1} f_1 r_3^{-2} f_1 r_2 f_1 r_3 f_1 r_2^{-1} f_1 r_3^4 f_1 r_2 f_1 r_3 f_1 r_2^{-1} \times \\ & \times f_1 r_3^{-2} f_1 r_2 f_1 r_3^2 f_1 r_2^{-1} f_1 r_3^2 f_1 r_2 f_1 r_3^{-1} f_1 r_2^{-1} f_1 r_3^{-1} f_1 r_2 f_1 r_3 \times \\ & \times f_1 r_2^{-1} f_1 r_2 r_3^{-2} f_1 r_2^{-1} f_1 r_3^{-1} f_1 r_2 f_1 r_2^{-1} r_3 f_1 r_2 f_1 r_2^{-1} f_1 r_2 f_1 r_3^2 \times \\ & \times f_1 r_2^{-1} f_1 r_2 f_1 r_2^{-1} f_1 r_2 r_3^{-1} f_1 r_2^{-1} f_1 r_3 f_1 r_2 f_1 r_2^{-2} r_3^{-2}. \end{aligned}$$

Since any polar facet can be placed in the 18 and 48 positions, and any equator facet can be placed in the 28 and 38 positions (with the aid of setup moves), we have shown that G_p and G_e can be written as products of 2-cycles and therefore generate S_{16} . Then it follows, since $G = G_e \times G_p$, then $G = S_{16} \times S_{16}$ and is therefore a subgroup of S_{32} . \square

Chapter 14

Crossing the Rubicon

Mathematical structures are among the most beautiful discoveries by the human mind. The best of these discoveries have tremendous metaphorical and explanatory power.

Douglas Hofstadter, Metamathematical themas, 1985

The title of this chapter is ‘borrowed’ from a similarly worded title of an article by D. Hofstadter [H].

This chapter is a little more advanced than some of the others. The reader is assumed to be familiar with some topics covered in a course in linear algebra and elementary number theory or coding theory. I will also assume some results from Rotman [R], though for the understanding of the material in this chapter the reader may simply take those results on faith.

Let g_1, \dots, g_{12} denote the basic moves of the Rubik icosahedron. A surprising result of Conway states that the group generated by $g_i \cdot g_j^{-1}$ is the simple ‘sporadic’ group M_{12} . (This is stated more precisely below.) I shall describe, in this chapter, what M_{12} is and some of its remarkable properties. They form a basis for my opinion, which I hope you will agree with, that M_{12} is one of the most interesting objects in mathematics.

Let p be a prime unless otherwise stated and let q be a power of p . **We shall assume that $p > 3$.**

14.1 Doing the Mongean shuffle

Consider a deck of 12 cards labeled 0, 1, ..., 11. Let r, s be the permutations

$$r(t) = 11 - t, \quad s(t) = \min(2t, 23 - 2t).$$

The permutation r simply reverses the cards around; the permutation s is called the ‘Mongean shuffle’, named after Gaspard Monge (1746-1818), a French mathematician known mostly for his work in geometry. Monge also worked on the

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calculus of variations, partial differential equations, and combinatorics. To perform the reverse shuffle, simply take a stack of cards (face down, say) in your left hand and put them in your right hand one at a time (face down). To perform the Mongean shuffle, take the same stack of cards and, one at a time, put them alternately into one of two piles: the first card face up into the first pile, the second card face down into the second pile, the third card face up into the first pile, the fourth card face down into the second pile, and so on until the pile is exhausted. Now pick up the first pile of face-up cards, flip the entire pile over so that they are all face down, and put this first pile on top of the second pile.

Cards	Reverse shuffle	Mongean shuffle
0	11	0
1	10	2
2	9	4
3	8	6
4	7	8
5	6	10
6	5	11
7	4	9
8	3	7
9	2	5
10	1	3
11	0	1

Definition 14.1.1. *The Mathieu group M_{12} is defined to be the permutation group $M_{12} = \langle r, s \rangle \subset S_{12}$.*

Emile Mathieu (1835–1890) actually discovered five ‘sporadic’ simple groups, of which M_{12} is only one, a year after his Docteur es Sciences thesis in 1859. (Most simple groups belong to certain infinite families of simple groups; ‘sporadic’ simply means that it is not from one of these infinite families.) Mathieu’s main work was in mathematical physics and special (hypergeometric) functions.

Example 14.1.1. *Here we use SAGE to explore the Mathieu group M_{12} .*

SAGE

```
sage: r = lambda t: 11-t
sage: s = lambda t: min(2*t,23-2*t)
sage: R = [str((i+1,r(i)+1)) for i in range(12)]
sage: S = [str((i+1,s(i)+1)) for i in range(12)]
sage: print R
['(1, 12)', '(2, 11)', '(3, 10)', '(4, 9)', '(5, 8)',
 '(6, 7)', '(7, 6)', '(8, 5)', '(9, 4)', '(10, 3)',
 '(11, 2)', '(12, 1)']
sage: print S
['(1, 1)', '(2, 3)', '(3, 5)', '(4, 7)', '(5, 9)',
```

```

'(6, 11)', '(7, 12)', '(8, 10)', '(9, 8)', '(10, 6)',
'(11, 4)', '(12, 2)']
sage: g1 = G('(1, 12)(2, 11)(3, 10)(4, 9)(5, 8)(6, 7)')
sage: g2 = G('(2, 3, 5, 9, 8, 10, 6, 11, 4, 7, 12)')
sage: M = PermutationGroup([g1, g2])
sage: M.order()
95040
sage: MathieuGroup(12).order()
95040
sage: MathieuGroup(12).is_isomorphic(M)
True

```

14.2 Background on \mathbf{PSL}_2

We need a few basic facts about the projective special linear group of degree 2. We have already discussed the related group $GL(2, \mathbb{F}_p)$ in §12.3.1, so I refer you there for more details.

Möbius transformations, introduced in §12.3.1 are bijections from the projective line to itself, so we may interpret each Möbius transformation as an element of the symmetric group S_X , where $X = \mathbb{P}^1(\mathbb{F}_q)$ (and therefore also of S_n , where $n = |\mathbb{P}^1(\mathbb{F}_q)| = q + 1$).

Example 14.2.1. *Let $p = 11$ and let $f(x) = -1/x$. Then*

x	∞	0	1	2	3	4	5	6	7	8	9	10
$f(x)$	0	∞	10	5	7	8	2	9	3	4	6	1

Therefore, as a permutation, $f = (\infty, 0)(1, 10)(2, 5)(3, 7)(4, 8)(6, 9)$.

The following facts are known about the projective special linear group.

Theorem 14.2.1. *If $q > 3$ then $PSL_2(\mathbb{F}_q)$ is a simple group. Moreover, for all prime powers q ,*

$$|PSL_2(\mathbb{F}_q)| = (q^2 - 1)q / \gcd(2, q - 1).$$

(A simple group was defined in Definition 9.5.1.) This theorem is over 100 years old. It is proven, for example, in [R].

Theorem 14.2.2. *Choose a $k \in \mathbb{F}_q$ such that $\langle k \rangle = \mathbb{F}_q^\times$. Let*

$$f_1(x) = x + 1, \quad f_2(x) = k \cdot x, \quad f_3(x) = -1/x.$$

Then $PSL_2(\mathbb{F}_q)$ is generated by f_1, f_2 , and f_3 . In particular, the action of $PSL_2(\mathbb{F}_q)$ on the projective line $X = \mathbb{P}^1(\mathbb{F}_q)$ yields an injective homomorphism $PSL_2(\mathbb{F}_q) \rightarrow S_X$.

Basically, this is proven in [R] as well.

14.3 Galois' last dream

Supposedly, the night before he died in a duel, Galois wrote a letter to a friend containing the following remarkable theorem:

Theorem 14.3.1. (Galois) *Assume $p > 11$. Then $PSL_2(\mathbb{F}_p)$ has no embedding into a symmetric group S_n with $n \leq p$.*

Example 14.3.1. *Although this example doesn't prove anything, the following SAGE calculation is consistent with Galois' theorem.*

SAGE

```
sage: G = PGL(2,13)
sage: G._gap_().LargestMovedPoint()
14
sage: G._gap_().SmallerDegreePermutationRepresentation()
IdentityMapping( Group([ (3,14,13,12,11,10,9,8,7,6,5,4),
(1,2,9)(3,8,10)(4,5,12)(6,13,14) ] ) )
```

The following isomorphisms (for $q \leq 11$) are known:

$$PSL_2(\mathbb{F}_q) \cong \begin{cases} A_4, & q = 3, \\ A_5, & q = 5, \\ A_6, & q = 9. \end{cases}$$

If $p = 7$ or $p = 11$ then explicit embeddings of $PSL_2(\mathbb{F}_p)$ into A_7 ($p = 7$), A_{11} ($p = 11$) are known (see [CS], ch 10, or [K] for an excellent discussion of this).

14.4 The M_{12} generation

One of the most amazing aspects about M_{12} is its close relationship with other 'interesting' groups.

Definition 14.4.1. *Define the permutation f_4 of the set $\mathbb{P}^1(\mathbb{F}_p) = \{\infty, 0, 1, \dots, p-1\}$, for $3 \leq p \leq 11$, as follows*

$$f_4 = \begin{cases} 1, & p = 3, \\ (1, 2)(3, 4), & p = 5, \\ (1, 2)(3, 6), & p = 7, \\ (2, 10)(3, 4)(5, 9)(6, 7), & p = 11. \end{cases}$$

We have run across the group S_6 before, when studying the symmetries of the icosahedron. We have also seen that S_6 is rather an interesting group because it is the only non-abelian symmetric group S_n which has an outer automorphism. One rather interesting connection between M_{12} and S_6 is given by the following result.

Theorem 14.4.1.

(a) If $p = 5$ then $S_6 = \langle f_1, f_2, f_3, f_4 \rangle$.

(b) If $p = 11$ then $M_{12} = \langle f_1, f_2, f_3, f_4 \rangle$.

We shall see another interpretation of M_{12} below using coding theory!

Definition 14.4.2. Let

$$\delta(x) = \begin{cases} x^3/9, & x \in (\mathbb{F}_{23})^2 - 0, \\ 9x^3, & x \in \mathbb{P}^1(\mathbb{F}_{23}) - (\mathbb{F}_{23})^2. \end{cases}$$

This is an element of S_X , where $X = \mathbb{P}^1(\mathbb{F}_{23})$. Define the **Mathieu group** M_{24} by

$$\langle f, \delta \mid f \in PSL_2(\mathbb{F}_{23}) \rangle.$$

This is a permutation group in S_X , where $X = \mathbb{P}^1(\mathbb{F}_{23})$.

By the way,

(a) $\{x^2 \mid x \in \mathbb{F}_{23}\} = \{0, 1, 2, 3, 4, 6, 8, 9, 12, 13, 16, 18\}$,

(b) $M_{12} = 8 \cdot 9 \cdot 10 \cdot 11 \cdot 12 = 95040$,

(c) $|M_{24}| = 244823040$.

Here is how to verify this using **SAGE**:

```

SAGE
sage: Set([x^2 for x in GF(23)])
{0, 1, 2, 3, 4, 6, 8, 9, 12, 13, 16, 18}
sage: MathieuGroup(12).order()
95040
sage: MathieuGroup(24).order()
244823040

```

14.5 Coding the Golay way

The Golay code is probably the most important of all codes for both practical and theoretical reasons.

MacWilliams and Sloane [MS], p. 64

Codes are used in everyday life, from ISBN numbers on books to bar codes on food products to music CDs to satellite transmissions. There are many types of codes, some more efficient than others, some with better error-correcting ability than others, some more practical than others, and so on. We shall concern ourselves only with aspects which are related to (in one way or another) permutation puzzles.

Let $q = p^k$ be a prime power.

Definition 14.5.1. A **q-ary code** is a subset C of a finite dimensional vector space V , provided a fixed basis over the finite field \mathbb{F}_q . The number of coordinates (i.e., the dimension of V) is called the **length** of the code word. If $q = 2$ then the code is called **binary** (instead of 2-ary) and if $q = 3$ then the code is called **ternary** (instead of 3-ary).

A code C which is a vector space over \mathbb{F}_q is called a **linear code**. The elements (or vectors) in this vector space C are called **codewords**.

Two linear codes $C, C' \subset \mathbb{F}_q^n$ are **equivalent** if there is a bijection $F; C \rightarrow C'$ which is obtained by composing a permutation of the positions of the codewords $c = (c_1, c_2, \dots, c_n)$ (via an element p of S_n sending c to $(c_{p(1)}, c_{p(2)}, \dots, c_{p(n)})$) with a map of the form $c = (c_1, c_2, \dots, c_n) \mapsto c' = (a_1 c_1, a_2 c_2, \dots, a_n c_n)$, for some $a_1, \dots, a_n \in \mathbb{F}_q^\times$.

Imagine we have a long message to send: ‘Hi Mom! I just learned how to solve the Rubik’s Cube ...’ First, chop this message into blocks of a specific size, say each block contains k consecutive letters. (This depends on the code you want to use.) Next, each letter, space and punctuation mark can be translated into a number. (It doesn’t matter how.) There are only finitely many letters and punctuation marks, so we can choose q larger than every number arising this way. Now associate with each number arising this way a corresponding element of \mathbb{F}_q , so that different numbers are associated with different elements. In this way, our message can be regarded as a sequence of elements of \mathbb{F}_q^k . For this reason, the vector space \mathbb{F}_q^k over \mathbb{F}_q is often called the **information space**.

Suppose that we use a linear code to transmit our message. A linear code of length n may be regarded as the image of an injective linear transformation $G : \mathbb{F}_q^k \rightarrow \mathbb{F}_q^n$. Give \mathbb{F}_q^k and \mathbb{F}_q^n their standard bases. The encoding map G , with respect to these bases, is called the **generating matrix** of C . The rows of G form a basis for the vector space C over \mathbb{F}_q .

Example 14.5.1. Here is an example of such a code using SAGE.

SAGE

```
sage: C = HammingCode(3,GF(2))
sage: C
Linear code of length 7, dimension 4 over Finite Field of size 2
sage: C.gen_mat()
[1 1 1 0 0 0 0]
[1 0 0 1 1 0 0]
[0 1 0 1 0 1 0]
[1 1 0 1 0 0 1]
sage: C.check_mat()
[0 1 1 1 1 0 0]
[1 0 1 1 0 1 0]
[1 1 0 1 0 0 1]
```

One other interesting thing about linear codes is that there is a surjective linear map $H : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^{n-k}$ with the property that $Hv = \vec{0}$ if and only if $v \in C$. In other words, H checks whether or not your received vector v contains any

errors (i.e., if v is actually a codeword). The map H with respect to the standard basis is called the **parity check matrix**, or simply the **check matrix**.

Ponderable 14.5.1. Show that for all $w \in \mathbb{F}_q^k$, $HGw = \vec{0}$.

Example 14.5.2. $V = \mathbb{F}_q^n = \mathbb{F}_q \times \dots \times \mathbb{F}_q$ (n times) is a code.

If $v = (v_1, v_2, \dots, v_n)$, $w = (w_1, w_2, \dots, w_n)$ are vectors in V , then we define

$$d(v, w) = |\{i \mid 1 \leq i \leq n, v_i \neq w_i\}|$$

to be the **Hamming distance** between v and w .

Definition 14.5.2. We say that C is **e -error correcting** if the following property always holds: given any $w \in F^n$ whose Hamming distance from some $c \in C$ is $\leq e$ then any $c' \in C$, $c \neq c'$, must satisfy $d(w, c') > e$ (in other words, a codeword at most distance e from w is unique if w has at most e errors).

Definition 14.5.3. Let $\text{Mon}_n(\mathbb{F}_q)$ denote the group of all $n \times n$ matrices which have exactly one non-zero entry from \mathbb{F}_q per row and per column. An element of $\text{Mon}_n(\mathbb{F}_q)$ is called a **monomial matrix**.

Ponderable 14.5.2.

(a) Show $\text{Mon}_n(\mathbb{F}_q)$ is a group.

(b) Show $|\text{Mon}_n(\mathbb{F}_q)| = (q - 1)^n \cdot n!$.

Definition 14.5.4. The set of all $A \in \text{Mon}_n(\mathbb{F}_q)$ such that $A \cdot C = C$ (i.e., are the same code) is called the **automorphism group of C** , denoted $\text{Aut}(C)$.

We shall see some examples below.

Definition 14.5.5. If w is a codeword in \mathbb{F}_q^n then the number of non-zero coordinates of w is called the **weight** of w , denoted $\text{wt}(w)$. The set of locations where the coordinates of w are non-zero, $\{i \mid w_i \neq 0\}$, is called the **support** of w , denoted $\text{supp}(w)$.

If C is a linear code then the smallest non-zero weight of any codeword is called the **minimum distance** of C .

A **cyclic code** is a code which has the property that whenever $(c_0, c_1, \dots, c_{n-1})$ is a codeword then so is $(c_{n-1}, c_0, \dots, c_{n-2})$. If $c = (c_0, c_1, \dots, c_{n-1})$ is a non-zero codeword in a cyclic code C with the property that the degree of

$$g_c(x) = c_0 + c_1x + \dots + c_{n-1}x^{n-1}$$

is minimum (among all such polynomials, as c ranges over all the non-zero codewords in C), then we call g_c a **generator polynomial for C**

Example 14.5.3. *The code with generator matrix*

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}$$

is a binary code of length 7 and minimum distance 3. It is equivalent to a cyclic code having a generator polynomial

$$g(x) = 1 + x + x^3.$$

(For the proof of this statement, see for example Huffman and Pless [HP].) This code is called a ‘Hamming code’ and has many interesting properties which, to describe, would take us too far afield. The interested reader is referred to [CS], ch. 3, [MS], and [HP].)

Ponderable 14.5.3. *Check that the code in Example 14.5.3 is 1-error correcting.*

Definition 14.5.6. *Let n be a positive integer relatively prime to q and let α be a primitive n -th root of unity. Each generator polynomial g of a cyclic code C of length n has a factorization of the form*

$$g(x) = (x - \alpha^{k_1}) \dots (x - \alpha^{k_r}),$$

where $\{k_1, \dots, k_r\} \subset \{0, \dots, n-1\}$. The numbers α^{k_i} , $1 \leq i \leq r$, are called the **zeros** of the code C . They do not depend on the choice of g .

Definition 14.5.7. *Let p and n be distinct primes and assume that p is a square mod n . The **quadratic residue code** of length n over \mathbb{F}_p is the cyclic code whose generator polynomial has zeros*

$$\{\alpha^k \mid k \text{ is a square (mod } n)\}.$$

The **binary Golay code** GC_{23} is the quadratic residue code of length 23 over \mathbb{F}_2 . The **binary Golay code** GC_{24} is the code of length 24 over \mathbb{F}_2 obtained by appending onto GC_{23} a zero-sum check digit (i.e., appending a field element which makes the codeword entries sum to 0 in the field).

The **ternary Golay code** GC_{11} is the quadratic residue code of length 11 over \mathbb{F}_3 . The **ternary Golay code** GC_{12} is the code of length 12 over \mathbb{F}_3 obtained by appending onto GC_{11} a zero-sum check digit.

A generator matrix of GC_{11} is

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 2 & 0 & 1 & 2 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 2 & 2 & 2 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 1 & 1 & 0 & 2 & 2 \\ 0 & 0 & 0 & 0 & 1 & 0 & 2 & 1 & 2 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 2 & 1 & 2 & 2 \end{pmatrix}$$

and a generator matrix of GC_{23} is

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$

These were obtained using the following commands.

SAGE

```
sage: C = TernaryGolayCode(); C
Linear code of length 11, dimension 6 over Finite Field of size 3
sage: C = BinaryGolayCode(); C
Linear code of length 23, dimension 12 over Finite Field of size 2
```

Now simply enter the command `C.gen_mat()` to obtain the matrices above.

The Golay codes are among the most important of all error-correcting codes. They are named after Marcel Golay (1902-1989), a physicist born in Switzerland, and educated in Zürich and then at the University of Chicago. Right after receiving his diploma, he worked for the U. S. Army Signal Corps Labs for almost 25 years. It was during this time that he discovered, in 1949, the Golay codes and the non-binary generalization of the Hamming codes. (Golay knew only of the $[7, 4]$ -binary Hamming code, discovered the binary Hamming codes independently, and then generalized them to the non-binary case.)

The following result illustrates how the Mathieu groups arise in coding theory.

Theorem 14.5.1.

- (a) *There is a normal subgroup N of $\text{Aut}(GC_{12})$ of order 2 such that $\text{Aut}(GC_{12})/N$ is isomorphic to M_{12} .*
- (b) $\text{Aut}(GC_{24}) = M_{24}$.

These can be found in the excellent book Huffman and Pless [HP] and the wonderful **Atlas of finite groups** [A] (look under M_{12} and M_{24} , respectively).

Since the Mathieu groups are so large, this theorem above indicates that the Golay codes GC_{12} and GC_{24} have a lot of symmetry.

It is a basic rule of thumb in mathematics that whenever you find something displaying a lot of symmetry, it will quite often have other interesting properties.

With this philosophy spurring us on, let us turn to some of the other properties of these codes.

Lemma 14.5.1. *Any two codewords in GC_{24} differ by 8 bits. The code GC_{24} detects 4 errors (per 24 bits) and corrects 3 errors.*

When you compare that with the correcting ability of bar codes or ISBN codes (which merely have a check-digit), GC_{24} is much better.

Lemma 14.5.2. *If w is a codeword in GC_{24} then $wt(w)$ is either 0, 8, 12, 16, 24.*

In other words, the weight distribution vector $(A_0, A_1, \dots, A_{24})$, where A_i is the number of codewords of weight i , only has $A_i > 0$ for i in $\{0, 8, 12, 16, 24\}$.

Definition 14.5.8. *The codewords of weight 12 in GC_{24} are called **dodecads**.*

Example 14.5.4. *We can use **SAGE** to get a handle on this concept of a dodecad.*

```

SAGE
sage: C = ExtendedBinaryGolayCode(); C
Linear code of length 24, dimension 12 over
      Finite Field of size 2
sage: print C.weight_distribution()
[1, 0, 0, 0, 0, 0, 0, 0, 0, 759, 0, 0, 0, 2576, 0, 0, 0, 759, 0,
 0, 0, 0, 0, 0, 0, 1]
sage: c = C.random(); c
(0, 0, 1, 0, 1, 0, 0, 0, 1, 1, 0, 1, 0, 0, 1, 1, 0, 1,
 0, 1, 0, 1, 1, 1)
sage: hamming_weight(C.random())
12
sage: c in C
True
sage: d = c.nonzero_positions(); d
[2, 4, 8, 9, 11, 14, 15, 17, 19, 21, 22, 23]
```

The above codeword c is an example of a dodecad. You can see from the weight distribution that most codewords in GC_{24} are dodecads.

We may identify a codeword $w = (c_0, c_1, \dots, c_{23})$ with the set of indices i of the non-zero coordinates $c_i \neq 0$.

Theorem 14.5.2. *M_{12} is the stabilizer in M_{24} of a dodecad, regarded as a set of indices.*

There are also fascinating connections between the code GC_{24} and the constructions of certain designs.

A **block design** is a pair (X, B) , where X is a non-empty finite set of $v > 0$ elements called **points**, and B is a non-empty finite multiset of size b whose

elements are called **blocks**, such that each block is a non-empty finite multiset of k points. A design without repeated blocks is called a **simple** block design. If every subset of points of size t is contained in exactly λ blocks, then the block design is called a $t - (v, k, \lambda)$ **design** (or simply a t -design when the parameters are not specified). When $\lambda = 1$ then the block design is called a $S(t, k, v)$ **Steiner system**. (See chapter 8 of [HP], or almost any textbook on combinatorics, for further details on block designs.)

Theorem 14.5.3. (*Assmus and Mattson theorem, §8.4, p. 303 of [HP]*) Let A_0, A_1, \dots, A_n be the weights of the codewords in a binary linear $[n, k, d]$ code C , and let $A_0^\perp, A_1^\perp, \dots, A_n^\perp$ be the weights of the codewords in its dual $[n, n - k, d^\perp]$ code C^\perp . Fix a t , $0 < t < d$, and let

$$s = |\{i \mid A_i^\perp \neq 0, 0 < i \leq n - t\}|.$$

Assume $s \leq d - t$.

- If $A_i \neq 0$ and $d \leq i \leq n$ then $C_i = \{c \in C \mid wt(c) = i\}$ holds a simple t -design.
- If $A_i^\perp \neq 0$ and $d^\perp \leq i \leq n - t$ then $C_i^\perp = \{c \in C^\perp \mid wt(c) = i\}$ holds a simple t -design.

In the Assmus and Mattson theorem, X is the set $\{1, 2, \dots, n\}$ of coordinate locations and $B = \{\text{supp}(c) \mid c \in C_i\}$ is the set of supports of the codewords of C of weight i . Therefore, the parameters of the t -design for C_i are:

- t (given),
- $v = n$,
- $k = i$ (k is not to be confused with $\dim(C)$),
- $b = A_i$
- $\lambda = b \cdot C(k, t) / C(v, t)$, where $C(a, b)$ denotes the binomial coefficient ‘ a choose b ’,

by Theorem 8.1.6, p 294, in [HP].

Example 14.5.5. The following SAGE example below means that the binary $[24, 12, 8]$ -code C has the property that the (support of the) codewords of weight 8 (resp, 12, 16) form a 5-design. Similarly for its dual code C^\perp (of course $C = C^\perp$ in this case, so this information is extraneous).

First, we enter the binary $[24, 12, 8]$ Golay code and then use the **verbose** option of SAGE’s `assmus_mattson_designs` function. For typographical reasons, SAGE uses C^* instead of C^\perp in its output.

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SAGE

```
sage: C = ExtendedBinaryGolayCode()
sage: C.assmus_mattson_designs(5, mode="verbose")

The weight 8 codewords of C form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 8 , lambda = 1
    There are 759 blocks of this design.
The weight 12 codewords of C form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 12 , lambda = 48
    There are 2576 blocks of this design.
The weight 16 codewords of C form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 16 , lambda = 78
    There are 759 blocks of this design.
The weight 24 codewords of C form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 24 , lambda = 1
    There are 1 blocks of this design.
The weight 8 codewords of C* form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 8 , lambda = 1
    There are 759 blocks of this design.
The weight 12 codewords of C* form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 12 , lambda = 48
    There are 2576 blocks of this design.
The weight 16 codewords of C* form a t-(v,k,lambda) design,
where
    t = 5 , v = 24 , k = 16 , lambda = 78
    There are 759 blocks of this design.

['weights from C: ',
 [8, 12, 16, 24],
 'designs from C: ',
 [[5, (24, 8, 1)], [5, (24, 12, 48)], [5, (24, 16, 78)],
 [5, (24, 24, 1)]],
 'weights from C*: ',
 [8, 12, 16],
 'designs from C*: ',
 [[5, (24, 8, 1)], [5, (24, 12, 48)], [5, (24, 16, 78)]]]
```

14.6 M_{12} is crossing the Rubicon

The result of this section was mentioned briefly in §12.3 above.

Let f_1, f_2, \dots, f_{12} denote the basic moves ($2\pi/5$ degree turns of a pentagon about a vertex) of the Rubik icosahedron, regarded as elements of S_V , where V denotes the set of 12 vertices of the Rubik icosahedron ('Rubicon').

The following remarkable result is due to John Conway [CS]. John Conway (1937-) is a very interesting character in the evolution of group theory, or even in mathematics in general. In the late 1960's, at a time when some experts thought all finite simple groups had been classified, Conway found another one, of order 8,315,553,613,086,720,000. This group he found had the nice property that two of its subgroups were also undiscovered finite simple groups! Thompson [Tho] has a fascinating account of this discovery.

Theorem 14.6.1. $M_{12} = \langle x \cdot y^{-1} \mid x, y \in \{f_1, \dots, f_{12}\} \rangle$.

In other words, the Mathieu group M_{12} (which, recall, is a finite simple group discovered by Mathieu) is generated by the twist-untwist moves of the Rubik icosahedron, or Rubicon. If we call a ‘twist-untwist’ move of the form $x \cdot y^{-1}$ (with x, y as in the theorem above) a **cross** move, then, with apologies to Caesar, the theorem above says that M_{12} is generated by the crosses of the Rubicon.

In fact, C acts 5-transitively on the set of vertices of the Rubicon (this is implicit in [CS]).

14.7 An aside: A pair of cute facts

It’s hard to resist stating some more interesting facts about the Mathieu groups.

14.7.1 Hadamard matrices

Let $A = (a_{ij})_{1 \leq i, j \leq n}$ denote a real $n \times n$ matrix. The following question seems quite natural in a course in advanced vector calculus or real analysis:

Question: What is the maximum value of $|\det(A)|$, where the entries of A range over all real numbers $|a_{ij}| \leq 1$?

From vector calculus we know that the absolute value of the determinant of a real square matrix equals the volume of the parallelepiped spanned by the row (or column) vectors of the matrix. The volume of a parallelepiped with sides of a *fixed* length depends on the angles the row vectors make with each other. This volume is maximized when the row vectors are mutually orthogonal, i.e., when the parallelepiped is a cube in \mathbb{R}^n . Suppose now that the row vectors of A are all orthogonal. The row vectors of A , $|a_{ij}| \leq 1$, are longest when each $a_{ij} = \pm 1$, which implies that the length of each row vector is \sqrt{n} . Suppose, in addition, that the row vectors of A are all of length \sqrt{n} . Such a matrix is called a **Hadamard matrix of order n** . Then $|\det(A)| = (\sqrt{n})^n = n^{n/2}$, since the cube has n sides of length \sqrt{n} . Now, if A is any matrix as in the above question then we must have $|\det(A)| \leq n^{n/2}$. This inequality is called **Hadamard’s determinant inequality**.

Jacques Hadamard (1865-1963) was a prolific mathematician who worked in many areas, but he is most famous for giving one of the first proofs of the prime number theorem (in 1896). (The prime number theorem, ‘known’ to Gauss though not proven, roughly states that the number of prime numbers less than N is about $N/\log(N)$, as N grows to infinity.)

The above question is unsolved for arbitrary n in the sense that it is not yet (as of this writing) known for which n ’s Hadamard matrices exist. The **Hadamard conjecture** states that for each n which is a multiple of 4 a Hadamard matrix exists. (This conjecture might, in fact, be due to Raymond Paley, a brilliant mathematician who contributed to several areas of mathematics, including two constructions of Hadamard matrices using the theory of finite fields. Sadly, he

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died in 1933 at the age of 26 in a skiing accident.) At the time of this writing, the smallest order for which no Hadamard matrix is presently known is 668.

What might be surprising at first sight is that there does not always exist a Hadamard matrix—for some n 's they exist and for other n 's they don't. For example, there is a 2×2 Hadamard matrix but not a 3×3 one. What is perhaps even more surprising is that, in spite of the fact that the above question arose from an analytic perspective, Hadamard matrices are related more to coding theory, number theory, and combinatorics [vLW]! In fact, a linear code constructed from a Hadamard matrix was used in the 1971 Mariner 9 mission to Mars [vL].

Example 14.7.1. *Consider the 12×12 matrix*

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ -1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & -1 & -1 & -1 & -1 \\ -1 & -1 & 1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 \\ -1 & 1 & -1 & 1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 \\ -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 \\ -1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 & 1 & 1 & -1 \\ -1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 & 1 & 1 \\ -1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 & 1 \\ -1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 & 1 \\ -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 & -1 \\ -1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 & 1 \\ -1 & 1 & -1 & 1 & 1 & 1 & -1 & -1 & -1 & 1 & -1 & 1 \end{pmatrix}.$$

This is a Hadamard matrix of order 12. A different Hadamard matrix of order 12 can be computed using SAGE:

```

SAGE
sage: MS = MatrixSpace(QQ,12,12)
sage: A = MS(sage_eval(gap.eval("HadamardMat( 12 )")))
sage: A
[ 1  1  1  1  1  1  1  1  1  1  1  1]
[ 1 -1  1 -1  1  1  1 -1 -1 -1  1 -1]
[ 1 -1 -1  1 -1  1  1  1 -1 -1 -1  1]
[ 1  1 -1 -1  1 -1  1  1  1 -1 -1 -1]
[ 1 -1  1 -1 -1  1 -1  1  1  1 -1 -1]
[ 1 -1 -1  1 -1 -1  1 -1  1  1  1 -1]
[ 1 -1 -1 -1  1 -1 -1  1 -1  1  1  1]
[ 1  1 -1 -1 -1  1 -1 -1  1 -1  1  1]
[ 1  1  1 -1 -1 -1  1 -1 -1  1 -1  1]
[ 1  1  1  1 -1 -1 -1  1 -1 -1  1 -1]
[ 1 -1  1  1  1 -1 -1 -1  1 -1 -1  1]
[ 1  1 -1  1  1  1 -1 -1 -1  1 -1 -1]
sage: det(A)
-2985984
sage: sqrt(12^12)
2985984

```

Here **SAGE** is calling a GAP package GUAVA [Gu], included as part of **SAGE**, which implements some methods for constructing Hadamard matrices. The last lines check that $\det(A) = n^{n/2}$.

Ponderable 14.7.1. Show that

- (a) if you swap two rows or columns of a Hadamard matrix, you will get another Hadamard matrix;
- (b) if you multiply any row or column of a Hadamard matrix by -1 , you will get another Hadamard matrix;
- (c) if you multiply any Hadamard matrix on the left by a signed permutation matrix (that is, a matrix with exactly one ± 1 per row and column), you will get another Hadamard matrix;
- (d) if you multiply any Hadamard matrix on the left by a signed permutation matrix (that is, a matrix with exactly one ± 1 per row and column), you will get another Hadamard matrix.

Ponderable 14.7.2. Let A, B be two Hadamard matrices of order n . Call A, B **left equivalent** if there is an $n \times n$ signed permutation matrix P such that $A = PB$. Show that this defines an equivalence relation on the set of all Hadamard matrices of order n .

Ponderable 14.7.3. Let A be a Hadamard matrix of order n . Let $\text{Aut}(A)$ denote the set of all $n \times n$ signed permutation matrices Q such that A is left equivalent to AQ . Show that $\text{Aut}(A)$, called the **automorphism group of A** , is a group under matrix multiplication.

The following result is yet another indication of the unique role of these Mathieu groups in mathematics.

Theorem 14.7.1. (Assmus and Mattson [AM]) Let A be the Hadamard matrix of order 12 in the above example. Then $\text{Aut}(A) \cong M_{12}$.

14.7.2 5-transitivity

The following result exemplifies once more the unique role of these Mathieu groups in group theory.

Theorem 14.7.2. If G is a subgroup of S_X for some finite set X and if G acts 5-transitively on X , then exactly one of the following must be true:

- (a) $G \cong S_n$, for some $n > 4$,
- (b) $G \cong A_m$, for some $m > 6$,

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$$(c) \ G \cong M_{12},$$

$$(d) \ G \cong M_{24}.$$

Furthermore, each of the groups in (a)-(d) acts 5-transitively on some finite set.

For a proof of this, see [CS] and [DiMo], Theorem 7.6A.

Chapter 15

Some solution strategies

I had a feeling once about Mathematics—that I saw it all. Depth beyond depth was revealed to me—the Byss and Abyss. I saw—as one might see the transit of Venus or even the Lord Mayor’s Show—a quantity passing through infinity and changing its sign from plus to minus. I saw exactly why it happened and why the tergiversation was inevitable but it was after dinner and I let it go.’

Sir Winston Churchill

It is a person far better than I who can look at the Rubik’s Cube and solve it behind his back (there are such people, I believe). They remind me of Churchill’s amusing quote—they can ‘see it all’. Speaking for myself, I need strategies, and the simpler the better. This chapter includes some strategies for solving the Rubik’s Cube, the Masterball, the Skewb, and the Pyraminx. This is good hands-on practice: in the words of Aristotle, ‘What we have to learn to do, we learn by doing.’ In addition, I discuss some of the mathematical ideas behind the algorithms used to solve the Rubik’s Cube. My favorite method involves a little bit of group theory (see §15.2 below).

For unexplained notation used in some of the sections below, see chapter 4. A great online source of solutions of lots of permutation puzzles is Jaap Scherphuis’ web site, [Sch2].

15.1 A strategy for solving the Rubik’s Cube

I recommend that the person who has never ‘played’ with the Rubik’s Cube before and wants to follow the strategy indicated below first learn how to ‘play a move’ without looking at the cube. Why? Think of the cube as a musical instrument like a guitar. To play music on the guitar, you should know how to look at the song sheet and play without having to constantly keep looking at your fingers. The cube is similar: in the beginning, you want to be able to

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make a move such as $F^2 \cdot L^2 \cdot U^2 \cdot (F^2 \cdot L^2)^3 \cdot U^2 \cdot L^2 \cdot F^2$ without having to constantly keep looking at your fingers. This will help make you faster and cut down on frustrating mistakes.

One way to solve it is the **corner-edge method**. The basic idea is as follows.

- Ignoring orientations, i.e., twists and flips, first solve the corner subcubes. This means, make sure that the corners will, up to possible twisting, match the center subcubes.
- Ignoring orientations, i.e., twists and flips, secondly solve the edge subcubes. This means, make sure that the edges will, up to possible flipping, match the center subcubes.
- Fix the corner orientations, i.e., twists of the corner subcubes. Now all the corners are completely solved.
- Fix the edge orientations, i.e., flips of the edge subcubes. Now the cube is solved.

For completeness, I mention another method. The **layer method** solution strategy is composed of 3 stages.

- Solve the top face and top edges.
- Solve the middle edges (and bottom edges as best as possible).
- Solve the bottom corners (and bottom edges if necessary).

There is nothing wrong with this method, but we shall not discuss this method further. (The reader who wishes to become a fast Rubik's Cube solver should learn both methods.)

15.1.1 Strategy for solving the cube

Remember, the basic moves are denoted R, L, U, D, F, B . Recall M_R denotes the clockwise (with respect to the **R**ight side) quarter turn of the **M**iddle slice parallel to the right side.

To 'solve' the corners (ignoring orientations) I first describe a simple method, which I will call 'fishing', for achieving this. (Mathematically, this amounts to performing some carefully chosen commutators.) Hold the cube in front of you. Suppose you want the lfd corner to be in the urf position without messing up the rest of the cube too much. First perform R^{-1} (this is 'baiting the hook'), then D ('setting the hook'), R ('reeling in the hook'), and finally D^{-1} . This last move can be omitted since after $R^{-1}DR$ you have moved lfd to urf .

By fishing you can either (a) solve all the corners (if you are lucky), (b) solve all except two which need to be swapped, or (c) solve all but 3 which need to be cyclically permuted. To finish, we need a corner swap move in case (b) and a corner 3-cycle in case (c). In case (b), the move

$$U \cdot F \cdot [R, U]^3 \cdot F^{-1} = \\ U \cdot F \cdot R \cdot U \cdot R^{-1} \cdot U^{-1} \cdot R \cdot U \cdot R^{-1} \cdot U^{-1} \cdot R \cdot U \cdot R^{-1} \cdot U^{-1} \cdot F^{-1}$$

swaps the corners (ubr, ufl) and permutes the edges (uf, ul, ub, ur), leaving the rest of the cube alone. (If the corners you need to ‘fix’ are not in the ubr, ufl positions, then you first need the following trick. Rotate the cube so that one of them is in the ubr position, say, then move (‘bait’) the other into the ufl position, make the move above, then unmoved (‘cut bait’).) In case (c), the move

$$[R \cdot D \cdot R^{-1}, U] = R \cdot D \cdot R^{-1} \cdot U \cdot R \cdot D^{-1} \cdot R^{-1} \cdot U^{-1}$$

is a corner 3-cycle (brd, urb, ulb), leaving the rest of the cube alone. (Again, if the corners you need to ‘fix’ are not in the brd, urb, ulb positions, you will need the ‘bait/cut bait’ trick.)

Once the corners have been solved, to ‘solve’ the edges (ignoring orientations), one simply has to use the edge 3-cycle

$$M_R^2 \cdot U^{-1} \cdot M_R^{-1} \cdot U^2 \cdot M_R \cdot U^{-1} \cdot M_R^2$$

repeatedly. This move is a 3-cycle on (uf, ul, ur) and leaves the rest of the cube alone. (This is not quick, and it is not at all obvious that only this move (along with the ‘bait/cut bait’ trick) works, but it can be proven using that fact that 3-cycles generate the alternating group.)

To twist the corners correctly, you need a corner twisting move. The move

$$(R^{-1} \cdot D^2 \cdot R \cdot B^{-1} \cdot U^2 \cdot B)^2$$

twists the ufr clockwise and the bld corner counterclockwise, leaving the rest of the cube alone. This move (along with the ‘bait/cut bait’ trick) will solve all the corners.

Finally, to flip the edges correctly, use

$$(M_R \cdot U)^3 \cdot U \cdot (M_R^{-1} \cdot U)^3 \cdot U,$$

which flips the top edges uf, ub (leaving the rest of the cube alone), and the move

$$(M_R \cdot U)^4,$$

which flips ub, ul and flips df, db (leaving the rest of the cube alone).

Now practice!

To shorten the notation and emphasize the group-theoretical symmetry of some of the moves given below, let $x^y = y^{-1} \cdot x \cdot y$ denote conjugation and $[x, y] = x \cdot y \cdot x^{-1} \cdot y^{-1}$ denote the commutator, for x, y group elements.

A table of ‘clean’ edge and corner moves (i.e., moves which perform as advertised and don’t mess up the rest of your cube!) is given below.

15.2. THE SUBGROUP METHOD

Move	Effect
$M_R^2 \cdot U^{-1} \cdot M_R^{-1} \cdot U^2 \cdot M_R \cdot U^{-1} \cdot M_R^2$	edge 3-cycle (uf,ul,ur)
$(M_R \cdot U)^3 \cdot U \cdot (M_R^{-1} \cdot U)^3 \cdot U$	flips the top edges uf, ub
$(R^2 \cdot U^2)^3$	permutes (uf,ub)(fr,br)
$(M_R \cdot U)^4$	flips ub,ul and flips df,db
$(R^{-1} \cdot D^2 \cdot R \cdot B^{-1} \cdot U^2 \cdot B)^2$	$ufr+$, $bld+$
$[R, U]^3 = (R \cdot U \cdot R^{-1} \cdot U^{-1})^3$	permutes $(ufr, dfr)(ubr, ubl)$
$F^2 \cdot L^2 \cdot U^2 \cdot (F^2 \cdot L^2)^3 \cdot U^2 \cdot L^2 \cdot F^2$	permutes $(uf, ub)(ur, ul)$
$(D^2 \cdot R^2 \cdot D^2 \cdot (F^2 \cdot R^2)^2 \cdot U)^2$	permutes $(ufl, ubr)(dfr, dbl)$
$(M_R^2 \cdot U \cdot M_R^2 \cdot U^2)^2$	permutes $(ufl, ubr)(ufr, ubl)$
$R \cdot D \cdot R^{-1} \cdot U \cdot R \cdot D^{-1} \cdot R^{-1} \cdot U^{-1}$	corner 3-cycle (brd, urb, ulb)

These moves were compiled with help from the books [Si], [B1], and [Sn1].

15.1.2 Catalog of 3×3 Rubik's 'supercube' moves

The supercube is the Rubik's Cube with each center facet marked with a short line through it and an adjoining edge (i.e., the center facets are oriented).

- $(M_R^2 \cdot U^{-1} \cdot M_R^{-1} \cdot U^2 \cdot M_R \cdot U^{-1} \cdot M_R^2)^2$ is the top edge 3-cycle (uf, ur, ul) .
- $(R^{-1} \cdot D^2 \cdot R \cdot B^{-1} \cdot U^2 \cdot B)^2$ twists the ufr corner clockwise and the bld corner counterclockwise (and does not twist any centers).
- $M_R^{-1} \cdot M_D^{-1} \cdot M_R \cdot U^{-1} \cdot M_R^{-1} \cdot M_D \cdot M_R \cdot U$ is the center twist $u+$, $r-$.

These moves were compiled from [Si] and [Sn1].

15.2 The subgroup method

One approach to solving the Rubik's Cube using a computer has been to construct a certain sequence of subgroups

$$G_n = \{1\} \subset G_{n-1} \subset \dots \subset G_1 \subset G_0 = G,$$

where $G = \langle R, L, F, B, U, D \rangle$ is the Rubik's Cube group, which allows the following strategy to be implemented:

- Represent a given position of the Rubik's Cube by an element $g_0 \in G$.
- Determine a complete set of coset representatives of G_{k+1}/G_k :

$$G_{k+1}/G_k = \cup_{i=1}^{r_k} g_{k+1,i} G_{k+1}, \quad \text{some } r_k > 1, \forall 0 \leq k < n$$

(note $m_{n-1} = 1, g_{n,1} = 1$).

- (Step 1) If $g_0 \in g_{1,i} G_1$ (where $i \in \{1, \dots, n_1\}$) then let $g_1 = g_{1,i}$ and $g'_1 = g_1^{-1} g_0$ (note $g'_1 \in G_1$).

- (Inductive step) If $g'_k \in G_k$ has been defined and if $g'_k \in g_{k+1,j}G_k$ (where $j \in \{1, \dots, n_1\}$), then let $g_{k+1} = g_{k+1,j}$ and $g'_{k+1} = g_{k+1}^{-1}g'_k$ (note $g'_{k+1} \in G_{k+1}$).
- Putting all these together, we obtain $1 = g_n^{-1}g_{n-1}^{-1}g_{n-2}^{-1}\dots g_1^{-1}g_0$, so

$$g_0 = g_1g_2\dots g_{n-1}g_n.$$

The hope is to be able to choose the sequence of subgroups G_i in such a way that the coset representatives are short, relatively simple moves on the Rubik's Cube so that the 'solution' $g_0 = g_1g_2\dots g_{n-1}g_n$ is not too long.

15.2.1 Example: The corner-edge method

I now present an example—a fairly unsophisticated one—but you will get the idea.

Let G_1 denote the subgroup which does not permute any corners, let G_2 denote the subgroup which does not permute any corners or edges, let G_3 denote the subgroup which does not permute any corners or edges and does not reorient any corners, and let $G_4 = \{1\}$:

$$G_4 = \{1\} \subset G_3 \subset G_2 \subset G_1 \subset G_0 = G.$$

This choice of subgroups roughly corresponds to the corner-edge method described in Singmaster [Si].

The idea is simple.

1. Represent a given position of the Rubik's Cube by an element $g_0 \in G$.
2. Let g_1 denote the move which moves all the corners into the correct positions (i.e., permutes them into the solved position and possibly twists them), so $g_1^{-1}g_0 \in G_1$. Let $g'_1 = g_1^{-1}g_0$.
3. Let g_2 denote the move which moves all the edges into the correct positions (i.e., permutes them into the solved position and possibly reorients corners and edges) and leaves all other pieces unpermuted, so $g_2^{-1}g'_1 \in G_2$. Let $g'_2 = g_2^{-1}g'_1$.
4. Let g_3 denote the move which 'solves' all the corners (i.e., twists them all into the correct orientation and may flip some edges) but does not permute any pieces, so $g_3^{-1}g'_2 \in G_3$. Let $g'_3 = g_3^{-1}g'_2$.
5. Let g_4 denote the move which 'solves' all the edges (i.e., flips them all into the correct orientation) and leaves all other facets alone.
6. The 'solution' is $g_0 = g_1g_2g_3g_4$.

15.2.2 Example: Thistlethwaite's method

Morwen Thistlethwaite (a mathematician and formerly a colleague of David Singmaster, now at the University of Tennessee) developed one of the best subgroup methods for solving the cube [FS]. He takes

$$\begin{aligned} G_1 &= \langle R, L, F, B, U^2, D^2 \rangle, & G_2 &= \langle R, L, F^2, B^2, U^2, D^2 \rangle, \\ G_3 &= \langle R^2, L^2, F^2, B^2, U^2, D^2 \rangle, & G_4 &= \{1\}. \end{aligned}$$

G_2 is isomorphic to the 'Rubik's $3 \times 3 \times 2$ -domino' group. Its order is $(8!)^2 \cdot 12$, according to [FS], §7.6. G_3 is the 'squares' group. Its order is $2^{13} \cdot 3^4$, according to [FS], §7.6.

Thistlethwaite has shown (using a computer to help with some of the work) the following facts.

- There is a complete set of coset representatives $\{g_{1,i} \mid 1 \leq i \leq n_1\}$ of G/G_1 such that each $g_{1,i}$ is at most 7 moves long (and $n_1 = 2048$). This set of moves flips edges only.
- There is a complete set of coset representatives $\{g_{2,i} \mid 1 \leq i \leq n_2\}$ of G_1/G_2 such that each $g_{2,i}$ is at most 13 moves long (and $n_2 = 1082565$). This set of moves twists corners only.
- There is a complete set of coset representatives $\{g_{3,i} \mid 1 \leq i \leq n_3\}$ of G_2/G_3 such that each $g_{3,i}$ is at most 15 moves long (and $n_3 = 29400$). This set of moves puts all the edge subcubes and corner subcubes in the correct position.
- There is a complete set of coset representatives $\{g_{4,i} \mid 1 \leq i \leq n_4\}$ of G_3/G_4 such that each $g_{4,i}$ is at most 17 moves long (and $n_4 = 663552$).

Therefore, the Rubik's Cube can be solved in at most $7 + 13 + 15 + 17 = 52$ moves.

More recent improvements on this method have reduced this number (see the Wikipedia page [Wi] on "Optimal solutions for Rubik's Cube" and [Lo] for details and recent updates).

15.2.3 Example: Kociemba's method

In the early 1990's, Herbert Kociemba improved Thistlethwaite's algorithm, calling his algorithm the 'two-phase-algorithm'. He reduced the number of intermediate groups to only two:

- $G_0 = \langle L, R, F, B, U, D \rangle$,
- $G_1 = \langle L, R, F^2, B^2, U^2, D^2 \rangle$,
- $G_2 = 1$.

As with Thistlethwaite's algorithm, he would search through the coset space G_0/G_1 to take the cube to group G_1 ('phase 1'). Next he searched the optimal solution for group G_1 ('phase 2').

In phase 1, the algorithm looks for maneuvers which will transform a scrambled cube to G_1 . That is, the orientations of corners and edges have to be constrained and the edges of the UD-slice have to be transferred into that slice. To find this goal state the program uses a search algorithm which is sometimes called *iterative deepening A^** (or IDA*). In the case of the cube, this means that it iterates through all maneuvers of increasing length using lookup tables.

In phase 2 the algorithm restores the cube in the subgroup G_1 , using only moves of this subgroup. It restores the permutation of the 8 corners, the permutation of the 8 edges of the U -face and D -face, and the permutation of the 4 UD -slice edges.

In 1995, Michael Reid showed that

- the search in G_0/G_1 needs at most 12 moves;
- the search in G_1 at most 18 moves;
- using these two groups, every position can be solved in at most 29 face turns, or in 42 quarter turns. (This result was improved by Silviu Radu [Ra] in 2005 to 40.)

For further details on this algorithm and software, please see Herbert Kociemba's page [Koc] or Michael Reid's page [Re].

15.3 Rainbow Masterball

The solution strategy for the Masterball puzzle is as follows.

Step 1 : The idea is to first get all the middle bands aligned first, so you get the Masterball corresponding to a matrix of the form

$$\begin{array}{cccccccc} * & * & * & * & * & * & * & * \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \\ * & * & * & * & * & * & * & * \end{array}$$

Here, $*$ denotes any color. I have labeled the colors on the Masterball as 1, 2, ..., 8 in order of occurrence.

I describe a method which I call 'fishing' for achieving this. (Mathematically, this amounts to performing some carefully chosen commutators.) Without too much trouble you can always assume that we have one column aligned. You may need to flip or rotate the ball a little to do this. Call this aligned column 'column 1' and call the color in column 1, 'color

1'. We want to get the middle two entries in column 2 aligned. Call the color in the $(2, 3)$ -entry 'color 2'.

We want to get color 2 in the $(2, 2)$ -entry. The remaining large color 2 tile is what we will 'fish' for. Hold the ball in front of you in such a way that column 2 is slightly to the left of center and column 3 is slightly to the right of center. There are 4 facets in the right upper middle band, 4 facets in the left upper middle band, 4 facets in the right lower middle band, and 4 facets in the left lower middle band. A flip about the center on the right half (i.e., perform f_2) exchanges these. We may assume that color 2 is on one of the 4 facets in the right lower middle band. (If it isn't, you need to apply f_2 first). Now perform $r_2^{-1} \cdot f_2^{-1} \cdot r_2 \cdot f_2$: first perform r_2^{-1} (this is 'baiting the hook'), then f_2^{-1} ('putting the hook in the water'), then r_2 ('setting the hook'), and finally f_2 ('reeling in the hook'). You may or may not have color 2 in the $(2, 2)$ place, but the color 1 stripe is intact. If necessary, try again. After at most 4 tries you'll be successful.

Step 2 : Repeat this 'fishing' strategy to get color 2 in the $(1, 2)$ position (using $r_1^{-1} \cdot f_2^{-1} \cdot r_1 \cdot f_2$ in place of $r_2^{-1} \cdot f_2^{-1} \cdot r_2 \cdot f_2$). Now, by turning the ball over if necessary, repeat this idea to get color 2 in the $(4, 2)$ position. Now you have two 'aligned' stripes on your ball—color 1 in column 1 and color 2 in column 2. We say, in this case, that columns 1 and 2 have been 'solved'.

Step 3 : Repeat the analogs of Steps 1 and 2 for columns 3 and 4.

Step 4 : Use the moves in the catalog in Figure 15.3.1 below to finish the puzzle. (I believe the only moves needed are the 'equator2swap36' and the 'polar2swap36' below, along with suitable, cleverly chosen 'set-up moves'.)

15.3.1 A catalog of Masterball moves

Column moves

We number the columns as $1, \dots, 8$. We will use a signed cycle notation to denote an action of a move on the columns of the Masterball.

Example 15.3.1. *A move which switches the 1st and 3rd column but flips both of them over will be denoted by $(1, 3)_-$.*

A move which sends the 4th column to the 6th column, the 6th column to the 5th column, and switches the 2nd and 3rd column but flips both of them over will be denoted by $(2, 3)_- (6, 5, 4)$. We have used a $-$ as a subscript to indicate the flip.

Finally, $(f_1 \cdot f_2 \cdot f_3 \cdot f_4)^2 \cdot r_1 \cdot r_2 \cdot r_3 \cdot r_4$ swaps the 7,8 columns and leaves all the others fixed but flipped over.

Move	Cycle
f_1	$(1, 4)_{-}(2, 3)_{-}$
f_2	$(2, 5)_{-}(3, 4)_{-}$
f_3	$(3, 6)_{-}(4, 5)_{-}$
f_4	$(4, 7)_{-}(5, 6)_{-}$
f_5	$(5, 8)_{-}(6, 7)_{-}$
f_6	$(1, 6)_{-}(7, 8)_{-}$
f_7	$(2, 7)_{-}(1, 8)_{-}$
f_8	$(3, 8)_{-}(1, 2)_{-}$
$f_1 \cdot f_2 \cdot f_1$	$(1, 2)_{-}(3, 5)$
$f_1 \cdot f_2 \cdot f_1 \cdot f_2$	$(5, 4, 3, 2, 1)_{-}$
$f_1 \cdot f_3 \cdot f_1$	$(1, 5)(2, 6)$
$f_2 \cdot f_3 \cdot f_2$	$(2, 3)_{-}(6, 5, 4)$
$f_1 \cdot f_4 \cdot f_1$	$(1, 7)(5, 6)$
$f_1 \cdot f_5 \cdot f_1$	$(5, 8)_{-}(6, 7)_{-}$
$f_1 \cdot f_8 \cdot f_1$	$(2, 8)(3, 4)_{-}$
$f_8 \cdot f_1 \cdot f_8$	$(1, 8)_{-}(4, 3, 2)$
$f_2 \cdot f_1 \cdot f_2$	$(1, 3)(4, 5)_{-}$
$f_3 \cdot f_1 \cdot f_3$	$(1, 3)(4, 8)_{-}$
$f_8 \cdot f_1 \cdot f_2$	$(1, 4)_{-}(2, 3, 8, 5)_{-}$

Move	Cycle
$x = r_1 \cdot f_4 \cdot r_1^{-1} \cdot r_4 \cdot f_4 \cdot r_4^{-1}$	$(41, 84)(44, 81)$
$x \cdot r_1^4 \cdot x \cdot r_4^4$	$(41, 81)(44, 84)$
$f_1 \cdot r_1 \cdot f_4 \cdot r_1^{-1} \cdot r_4 \cdot f_4 \cdot r_4^{-1} \cdot f_1$	$(14, 84)(11, 81)$
polar2swap36	$(11, 14)(31, 61)$
polar2swap18	$(61, 64)(11, 81)$
equator2swap36	$(12, 13)(32, 62)$
equator2swap18	$(62, 63)(12, 82)$

Some products of 2-cycles on the facets

The polar2swap and equator2swap were obtained by trying variations of some moves of Andrew Southern's on a software implementation of the Masterball [J1] (see also [JS]).

We number the facets in the i th column, north-to-south, as $i1, i2, i3, i4$ (where $i = 1, 2, \dots, 8$).

where

- polar2swap36 is

$$f_1 \cdot r_3^{-1} \cdot r_4^{-1} \cdot f_1 \cdot f_2 \cdot r_1 \cdot r_4^{-1} \cdot f_2 \cdot r_4^4 \cdot f_2 \cdot r_1^{-1} \cdot r_4 \cdot f_2 \cdot r_4^4 \cdot f_1 \cdot r_3 \cdot r_4 \cdot f_1;$$

- polar2swap18 is

$$f_1 \cdot r_3^{-1} \cdot r_4^{-1} \cdot f_3 \cdot f_4 \cdot r_1 \cdot r_4^{-1} \cdot f_4 \cdot r_4^4 \cdot f_4 \cdot r_1^{-1} \cdot r_4 \cdot f_4 \cdot r_4^4 \cdot f_3 \cdot r_3 \cdot r_4 \cdot f_1;$$

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- equator2swap36 is

$$f_1 \cdot r_4^{-1} \cdot r_3^{-1} \cdot f_1 \cdot f_2 \cdot r_2 \cdot r_3^{-1} \cdot f_2 \cdot r_3^4 \cdot f_2 \cdot r_2^{-1} \cdot r_3 \cdot f_2 \cdot r_3^4 \cdot f_1 \cdot r_4 \cdot r_3 \cdot f_1;$$

- and equator2swap18 is

$$f_1 \cdot r_4^{-1} \cdot r_3^{-1} \cdot f_3 \cdot f_4 \cdot r_2 \cdot r_3^{-1} \cdot f_4 \cdot r_3^4 \cdot f_4 \cdot r_2^{-1} \cdot r_3 \cdot f_4 \cdot r_3^4 \cdot f_3 \cdot r_4 \cdot r_3 \cdot f_1.$$

For further details on the rainbow puzzle, see [JS], [J1].

15.4 The Skewb

15.4.1 Strategy

The goal here is to collect enough moves to support the following solution strategy: fix the centers and solve the corners using ‘clean corner moves’ (i.e., moves which do not affect the centers).

The basic moves are twists by 120 degrees clockwise about each of the 6 corners $FRU, FLU, BRU, BLU, BDR, BDL, DFR, DFL$.

15.4.2 A catalog of Skewb moves

The descriptions below were discovered with the help of a simulation of the Skewb written for MAPLE [J2].

1. $FRU \cdot BLU \cdot FRU^{-1}$ is order 3.
2. $[FRU \cdot FLU]^3$ twists 6 corners clockwise by 120 degrees. The 2 corners not twisted are those opposite the FRU, FLU corners: the BDR, BDL corners. The centers are all fixed. $(FRU \cdot FLU)^3 = (FLU \cdot FRU)^3$ rotates all the corners except for the bd corners. It does not permute any facets.
3. The move $[FRU \cdot BLU]^5$ fixes all the centers and the 2 ‘opposite’ corners: DFL, BDR . It twists the 3 corners FLU, BLU , and FRU . On the remaining 3 corners, it acts as the permutation (DFR, BRU, BDL) .
4. $FRU = BDL$. (Actually, they are equal only up to a rotation of the entire cube.) In general, a corner move is equal to the opposite corner move up to a rotation of the entire cube.
5. $(FRU \cdot FLU)^3 \cdot (BDL \cdot BDR)^6$ rotates all the corners except for the bu and the df corners. The uf corners are rotated clockwise and the bd corners counterclockwise. It does not permute any facets. $(FRU \cdot FLU)^6 \cdot (BDL \cdot BDR)^3$ is the same move, but rotates in the opposite direction.
6. $(FRU \cdot FLU)^6 \cdot (BDL \cdot BDR)^6$ rotates the corners as follows:
 - the uf corners counterclockwise,

- the *db* corners counterclockwise,
- the *df* corners clockwise,
- the *ub* corners clockwise.

It does not permute any facets.

7. Let $bottomspin = (FRU \cdot FLU)^6 \cdot (BDL \cdot BDR)^3 \cdot (DFR \cdot DFL)^3$. This move rotates the 4 bottom corners (the *df* corners clockwise and the *db* counterclockwise). It does not permute any facets.
8. $(BRU \cdot FLU)^9$ is a 5-cycle on the center facets (F, R, B, U, L) . It fixes the bottom and does not affect any corners. $(BLU \cdot FRU)^9$ is a 5-cycle on the center facets. It fixes the bottom and does not affect any corners.
9. $(BLU \cdot FRU)^9 \cdot (BRU \cdot FLU)^9$ is a product of 2 transpositions on the center facets, swapping front/back and up/right. It fixes the bottom and does not affect any corners.
10. Let U denote the clockwise (with respect to the up face) rotation of entire cube by 90 degrees. Then $bottomspin \cdot U \cdot bottomspin$ rotates but does not swap 2 corners (the *DFR* and *BDL*) and does not affect any other corners or faces.

15.5 The Pyraminx

Assume that the tetrahedron is lying on a flat surface in front of you, with the triangle base pointing away from you. The corners are denoted L (left), R (right), U (up), and B (back).

Basic moves: Let

- L denote the 120 degree clockwise rotation of the 2-level subtetrahedron containing the left corner,
- R denote the 120 degree clockwise rotation of the 2-level subtetrahedron containing the right corner,
- U denote the 120 degree clockwise rotation of the 2-level subtetrahedron containing the up corner,
- B denote the 120 degree clockwise rotation of the 2-level subtetrahedron containing the back corner.

First, get the ‘center’ facets solved, then twist the corner tips to solve them and the center facets. Finally, to solve the edge facets, use the following moves (given in [EK]):

- $[R, U^{-1}]$ is a 3-cycle of edge pieces on the URL face,
- $[R, U^{-1}] * [R^{-1}, L]$ is a flip of two edges (UR edge and UL edge) on the URL face.

15.6 The Megaminx

The strategy here is the same as for the $3 \times 3 \times 3$ Rubik's Cube.

- Place the corners correctly first (ignoring correct corner orientation).
- Place the edges correctly first (ignoring correct edge orientation).
- Twist the corners if necessary.
- Flip the edges if necessary.

Moves useful for carrying out these steps are included in the following catalog.

First, some notation. We label the faces f_1, f_2, \dots, f_6 on top and label the bottom faces f_7, f_8, \dots, f_{12} as in chapter 4. The same notation is used to indicate the move of the Megaminx given by rotating that face of the Megaminx by 72 degrees clockwise.

Catalog of moves:

- $f_1^{-1} * f_2^{-1} * f_1 * f_2 * f_1^{-1} * f_2^{-1} * f_1 * f_2 * f_1^{-1} * f_2^{-1} * f_1 * f_2 = [f_1, f_2]^3$ - swaps the $f_1.f_3$ and the $f_2.f_6$ corners: $(f_1.f_2.f_3, f_1.f_3.f_4)(f_1.f_2.f_6, f_2.f_6.f_7)$
- $m = f_3 * f_6^{-1} * f_4 * f_2^{-1} * f_5 * f_3^{-1} * f_6 * f_4^{-1} * f_2 * f_5^{-1}; f_6 * f_1 * m^6 * f_1^{-1} * f_6^{-1}$ - swaps 2 pairs of corners on the f_1 face: $(f_1.f_2.f_6, f_1.f_3.f_4)(f_1.f_2.f_3, f_1.f_5.f_6)$
- $f_1 * f_6 * f_1^{-1} * f_2 * f_1 * f_6^{-1} * f_1^{-1} * f_2^{-1}$ 3-cycle on corners and 3-cycle on edges $(f_1.f_2.f_6, f_1.f_7.f_6, f_2.f_6.f_7)(f_1.f_2, f_6.f_7, f_2.f_6)$
- $M2 = f_6 * f_2 * f_1 * f_2^{-1} * f_1^{-1} * f_6^{-1} * f_3^{-1} * f_1^{-1} * f_2^{-1} * f_1 * f_2 * f_3$ (Mark Longridge) - edge 3-cycle $(f_1.f_2, f_2.f_3, f_2.f_6)$
- $(f_6^{-1} * f_2^{-1} * f_3^{-1} * f_6 * f_2 * f_3)^6$ - triple corner twister, ccw twists of $f_1.f_2.f_3, f_1.f_2.f_6, f_1.f_5.f_6$
- $M3 = f_3^{-2} * f_6^2 * f_2 * f_1^{-1} * f_6 * f_1 * f_3^2 * f_6^2 * f_1 * f_6^{-2} * f_3^{-2} * f_1^{-1} * f_6^{-1} * f_1 * f_2^{-1} * f_6^{-2} * f_3^2 * f_1^{-1}$ (Mark Longridge) - edge 2-flip of $f_1.f_2, f_1.f_6$
- $M3a = f_6^{-1} * f_2^{-1} * f_1 * f_3^{-1} * f_1^{-1} * f_3 * f_2 * f_6 * f_3 * f_2 * f_1^{-1} * f_6 * f_1 * f_6^{-1} * f_2^{-1} * f_3^{-1}$ (Mark Longridge) - edge 2-flip of $f_1.f_2, f_1.f_3$

For further details, see the Internet sites [Sch2], [J2], or [Lo].

Chapter 16

Coda: Questions and other directions

[Lefschetz and Einstein] had a running debate for many years. Lefschetz insisted that there was difficult mathematics. Einstein said that there was no difficult mathematics, only stupid mathematicians. I think that the history of mathematics is on the side of Einstein.

*Richard Bellman, **Eye of the hurricane**, 1984*

In the spirit of Galois' quote in chapter 7, this book will conclude with some of the (many) things which I don't know related to the mathematics of the Rubik's Cube-like puzzles. Some of these have already appeared previously in the book.

Let G be the group of all legal moves of a permutation puzzle such as the Rubik's Cube, Skewb, Pyraminx, Megaminx, or Masterball.

1. (God's algorithm, Singmaster) Find the diameter of the Cayley graph of G .

Is it 26 in the quarter-turn metric (i.e., relative to the usual generators R, L, U, D, F, B with their inverses)? Is Mike Reid's move 'superflip composed with the four-spot' a longest move of the Rubik's Cube in the quarter-turn metric?

Is it 20 in the face-turn metric (i.e., relative to the generators $R, R^2, L, L^2, U, U^2, D, D^2, F, F^2, B, B^2$ with their inverses)? Is the 'superflip' a longest move of the Rubik's Cube in the face-turn metric?

In either metric, how many moves are there of maximal distance from start?

2. (Schwenk) Is the Cayley graph of G a Hamiltonian graph?
3. With the exception of the vertex cross group of the Pyraminx, are all the cross groups simple?

-
4. (Dyson) What are the cross groups of the 4-dimensional regular polyhedra?
 5. Are all the cross groups of the 4-dimensional regular polyhedra alternating groups?
 6. Let H be a group in the table in §9.4. If H is a subgroup of G , find explicit puzzle moves which satisfy the relations given for H in the table.
 7. (Singmaster) Find the number of elements of each order in G .
 8. Compute the Poincaré polynomial of G with respect to the generators $\{R, R^{-1}, L, L^{-1}, U, U^{-1}, D, D^{-1}, F, F^{-1}, B, B^{-1}\}$.
 9. Compute the generating polynomial of G .
 10. Find the ‘simplest’ presentation of G .
 11. If G is the Rubik’s Cube group generated by $\{U, D, R, L, F, B\}$, what is the average length of a word in G ?

There are other group-theoretical questions related to puzzles and games which we have not touched on. The reader interested in pursuing these notions further is referred (for example) to **Winning ways** [BCG], Conway and Sloane’s book [CS], their paper on lexicographic codes [CS2], their paper on Gray codes [CSW], M. Gardner [Gar2], Luers’ thesis [Lu], Stiller’s thesis [Sti], Elkies’ paper on the combinatorial game theory aspects of chess [El], and Eriksson’s paper [E].

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