SCFENCE: Automatic Inference of Memory Order Parameters to Obtain SC Behaviors under C/C++11

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Abstract

1. Introduction

2. Motivating Example

3. Technical

3.1 Inference Rules

Previous work takes advantage of model-checking approach to check whether a specific C/C++11 trace is sequentially consistent. By estabishing edges (*sb*, *rf* and *sc* by implication rules) between atomic operations, it judges whether the trace is SC by whether there exists a cycle in that graph.

Under the C/C++11 memory model, inferring the ordering parameters to obtain SC behaviors is essentially a searching problem. In the absence of consume operations, memory order parameters for atomic operations can be only one of the following: *memory_order_relaxed*, *memory_order_release*, *memory_order_acquire*, *memory_order_acq_rel* and *memory_order_seq_cst*. By enumerating all possible memory order parameters, we can guarantee that we can find out all the possible inference of parameters that ensure SC behaviors for a specific test case. However, this naive approach obviously leads to an exponential searching space.

Consider we start from the case where all memory order parameters are *memory_order_relaxed*. Whenever the model-checking approach finds out a cycle in a specific execution, we have to infer some stronger memory orders to eliminate the cycle. In order to guarantee completeness, we propose a search-based approach combined with patterns to reduce searching space. In Figure 1 , we show a number of common patterns that can exist in cycles. We explain what the weakest orders we should impose on operations to eliminate the corresponding cycle as the following.

Synchronization: This pattern considers

Circular reads-from:

Peterson lock:

Release sequence:

Independent reads & independent writes:

Figure 2 shows the core searching algorithm for all possible parameters.

4. Evaluation

5. Related Work

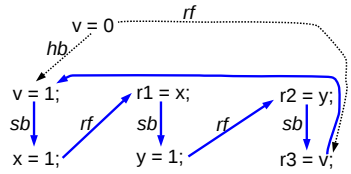
SC [1]

6. Conclusion

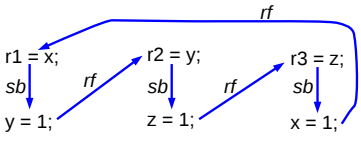
References

[1] L. Lamport. How to make a multiprocessor computer that correctly executes multiprocess programs. *IEEE Transactions on Computers*, 28(9):690–691, Sept. 1979.

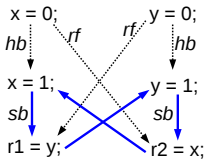
Synchronization



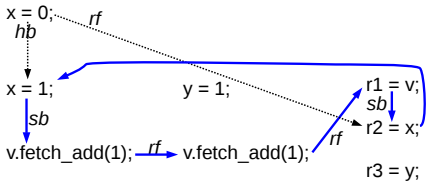
Circular Reads-from



Peterson Lock



Release Sequence



Independent Reads & Independent Writes

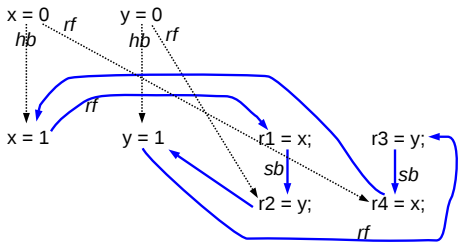


Figure 1. Cycle Patterns for Non-SC Behaviors

```
1:  function INFERPARAMS
2:      candidates := {}
3:      candidate c1 := replace all wildcards with relaxed
4:      candidates += c1
5:      results := {}
6:      while candidates is not empty do
7:          Candidate c := candidates.pop()
8:          Model-check with c and yield a cycle l
9:          if l == NULL then
10:             results += c
11:         else
12:             STRENGTHENPARAM(l, c, candidates)
13:         end if
14:     end while
15:     return results
16: end function
17: procedure STRENGTHENPARAM(cycle, candidate, candidates)
18:     if ∃ a pattern p in c then
19:         Candidate new := strengthen c by pattern p
20:         candidates += new
21:     else
22:         for all wildcard w in c do
23:             if w can be strengthened then
24:                 Candidate new := strengthen w in c
25:                 candidates += new
26:             end if
27:         end for
28:     end if
29: end procedure
```

Figure 2. Algorithm for Searching All Possible Parameters