

SCFENCE: Automatic Inference of Memory Order Parameters to Obtain SC Behaviors under C/C++11

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Abstract

1. Introduction
2. Motivating Example
3. Technical
 - 3.1 Inference Rules

Previous work takes advantage of model-checking approach to check whether a specific C/C++11 trace is sequentially consistent. By establishing edges (*sb*, *rf* and *sc* by implication rules) between atomic operations, it judges whether the trace is SC by whether there exists a cycle in that graph.

Under the C/C++11 memory model, inferring the ordering parameters to obtain SC behaviors is essentially a searching problem. In the absence of consume operations, memory order parameters for atomic operations can be only one of the following: *memory_order_relaxed*, *memory_order_release*, *memory_order_acquire*, *memory_order_acq_rel* and *memory_order_seq_cst*. By enumerating all possible memory order parameters, we can guarantee that we can find out all the possible inference of parameters that ensure SC behaviors for a specific test case. However, this naive approach obviously leads to an exponential searching space.

Consider we start from the case where all memory order parameters are *memory_order_relaxed*. Whenever the model-checking approach finds out a cycle in a specific execution, we have to infer some stronger memory orders to eliminate the cycle. In order to guarantee completeness, we propose a search-based approach combined with patterns to reduce searching space. In Figure 1, we show a number of common patterns that can exist in cycles. We explain what the weakest orders we should impose on operations to eliminate the corresponding cycle as the following.

Text running away

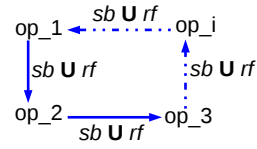
Old value read I:

Old value read II:

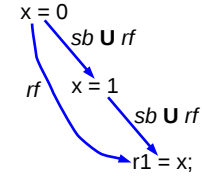
Reading future value:

Figure 2 shows the core searching algorithm for all possible parameters.

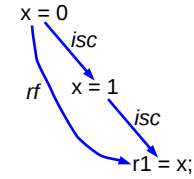
Circular *sb* & *rf*



Old Value Read I



Old Value Read II



Future Value Read

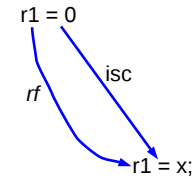


Figure 1. Cycle Patterns for Non-SC Behaviors

4. Evaluation

5. Related Work

SC [1]

```

1: function INFERPARAMS
2:   candidates := {}
3:   candidate c1 := replace all wildcards with relaxed
4:   candidates += c1
5:   results := {}
6:   while candidates is not empty do
7:     Candidate  $c := \text{candidates.pop}()$ 
8:     Model-check with  $c$  and yield a cycle  $l$ 
9:     if  $l == \text{NULL}$  then
10:       results +=  $c$ 
11:     else
12:       STRENGTHENPARAM( $l, c, \text{candidates}$ )
13:     end if
14:   end while
15:   return results
16: end function
17: procedure STRENGTHENPARAM( $\text{cycle}, \text{candidate},$ 
   candidates)
18:   if  $\exists$  a pattern  $p$  in  $c$  then
19:     Candidate new := strengthen  $c$  by pattern  $p$ 
20:     candidates += new
21:   else
22:     for all wildcard  $w$  in  $c$  do
23:       if  $w$  can be strengthened then
24:         Candidate new := strengthen  $w$  in  $c$ 
25:         candidates += new
26:       end if
27:     end for
28:   end if
29: end procedure

```

Figure 2. Algorithm for Searching All Possible Parameters

6. Conclusion

References

- [1] L. Lamport. How to make a multiprocessor computer that correctly executes multiprocess programs. *IEEE Transactions on Computers*, 28(9):690–691, Sept. 1979.