

Game Theory: Week 3

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Mixed strategies

Matching pennies game

	H	T
H	$1, -1$	$-1, 1$
T	$-1, 1$	$1, -1$

- last digit even: row player
- last digit odd: column player

What is going on here?

- No equilibrium
- First reaction: equilibrium sucks as a prediction tool
- Instead: “lack of equilibrium” is a feature not a bug
- When no (s_1, s_2, \dots, s_n) is an equilibrium...
 - players must be **unpredictable**
 - there is an equilibrium
 - but it involves acting in unpredictable ways
 - cannot be described by (s_1, s_2, \dots, s_n)

Mixed strategies

- Given a strategy space S_i
 - let \tilde{S}_i be the set of all probability distributions on S_i
- Payoffs $u_i(\mathbf{s})$ imply payoffs $u_i(\tilde{\mathbf{s}})$

Definition

A strategy profile $\tilde{\mathbf{s}}^*$ is a *mixed strategy Nash equilibrium* if for each player i , $u_i(\tilde{s}_i^*, \tilde{\mathbf{s}}_{-i}^*) \geq u_i(\tilde{s}'_i, \tilde{\mathbf{s}}_{-i}^*)$ for all $\tilde{s}'_i \in \tilde{S}_i$.

- What we studied last week, we will call *pure strategy* Nash equilibrium

Matching pennies

	H	T
H	$1, -1$	$-1, 1$
T	$-1, 1$	$1, -1$

- Unique equilibrium: each player plays H with probability $\frac{1}{2}$
- Column player wants to be as unpredictable as possible
 - hence plays two strategies with equal probability
- A best response for the row player:
 - play two strategies with equal probability

Matching pennies in the world

- Shooting penalty kicks
- Hide-and-seek
- Zara vs. high-end design
- Costly monitoring
 - checking up on an employee
 - auditing a tax payer

Costly monitoring

- Tax payer pays or cheats
- IRS audits or not
- Tax payer wants to cheat iff IRS does not audit
- IRS wants to audit iff taxpayer cheats

	<i>audit</i>	<i>not</i>
<i>pay</i>	1, -1	-1, 1
<i>cheat</i>	-1, 1	1, -1

- Symmetry not really plausible

Costly monitoring

- Tax payer pays or cheats
- IRS audits or not
- Tax payer wants to cheat iff IRS does not audit
- IRS wants to audit iff taxpayer cheats

	<i>audit</i>	<i>not</i>
<i>pay</i>	$0, -1$	$0, 0$
<i>cheat</i>	$-x, 1$	$1, -1$

Costly monitoring = generalized matching pennies

	<i>audit</i>	<i>not</i>
<i>pay</i>	0, -1	0, 0
<i>cheat</i>	-x, 1	1, -1

- x is the prison term for the cheating tax payer
 - lowers tax payer's utility
 - does not give direct benefits to the IRS
- Like matching pennies
 - IRS wants to match actions (audit/cheat, not/pay)
 - tax payer wants to mismatch (pay/audit, cheat/not)
- No pure strategy equilibrium
 - if IRS audits, tax payer pays \rightarrow IRS shouldn't audit
 - if IRS doesn't audit, tax payer cheats \rightarrow IRS should audit

Costly monitoring, version 1

	<i>audit</i>	<i>not</i>
<i>pay</i>	0, -1	0, 0
<i>cheat</i>	-2, 1	1, -1

- last digit even: row player
- last digit odd: column player

Costly monitoring, version 2

	<i>audit</i>	<i>not</i>
<i>pay</i>	0, -1	0, 0
<i>cheat</i>	-3, 1	1, -1

- last digit even: row player
- last digit odd: column player

Solving for a mixed strategy equilibrium

- Players must be playing both strategies
- Hence they must be **indifferent** between the two strategies

Solving for the mixed strategy equilibrium

- Let m be the probability of auditing
- Let p be the probability of paying
- Tax payer indifferent:

$$0 = -xm + (1 - m)$$

- IRS indifferent:

$$-p + (1 - p) = -(1 - p)$$

- Equilibrium:

$$\begin{aligned} m &= \frac{1}{1+x} \\ p &= \frac{2}{3} \end{aligned}$$

Penalties and behavior

- Equilibrium:

$$m = \frac{1}{1+x}$$
$$p = \frac{2}{3}$$

- Remarkable feature:

- x affects only the tax payer's payoff
- cheating prevalence independent of x
- affects only IRS' behavior
- **increasing the penalty for a crime does not reduce its prevalence**

Entry games

Entry game, version 1

- 10 players choose whether to *enter* or *not*
- Payoff to each player
 - if enter: $5 - \#$ firms who entered
 - if not: 0
- In other words,
 - market can support up to 5 firms
 - if 5 firms enter, profits are 0
 - if fewer enter, profits are positive
 - if more enter, profits are negative

Entry game, version 2

- 10 players choose whether to *enter* or *not*
- Payoff to each player
 - if enter: $3 - \#$ firms who entered
 - if not: 0
- In other words,
 - market can support up to 3 firms
 - if 3 firms enter, profits are 0
 - if fewer enter, profits are positive
 - if more enter, profits are negative

Revisiting Hawk-Dove

Recall Hawk-Dove

	<i>hawk</i>	<i>dove</i>
<i>hawk</i>	$-1, -1$	$2, 0$
<i>dove</i>	$0, 2$	$1, 1$

- There are two pure strategy equilibria
- Is there a mixed strategy equilibrium?
- Yes: everyone plays hawk with probability $\frac{1}{2}$
- Let h be the probability of hawk
- Indifference:

$$\begin{aligned} -h + 2(1 - h) &= 1 - h \\ h &= \frac{1}{2} \end{aligned}$$

Zero-sum games

Two player zero-sum games

- Many games are zero sum: $u_i = -u_j$
 - attack game, location game, matching pennies, etc.
- There is an easy way to solve zero-sum games

A pessimistic approach

- Suppose each player believes in Murphy's law:
 - whatever he does, the worst possible thing will happen
- Say s_i^m is a MaxMin strategy if it maximizes $\min_{s_{-i}} u_i(s_i, s_{-i})$

MaxMin vs. Nash

- MaxMin well defined in any game, zero sum or not
- Generally distinct from Nash
 - players do not pay attention to other players' payoffs
- Main result on two player zero-sum games:

Theorem

In any two player zero-sum game, a strategy profile is a Nash equilibrium if and only if both players play MaxMin strategies.

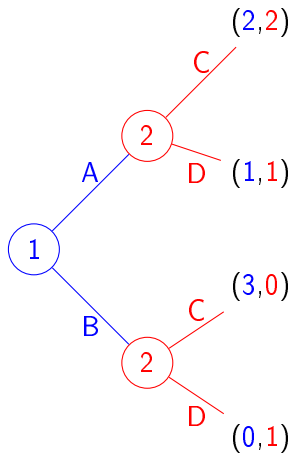
MaxMin Theorem

- MaxMin Theorem makes it easy to find equilibria of zero-sum games
- We can determine what will happen **player-by-player**
 - usually only possible with dominant strategies
- Think about the location game

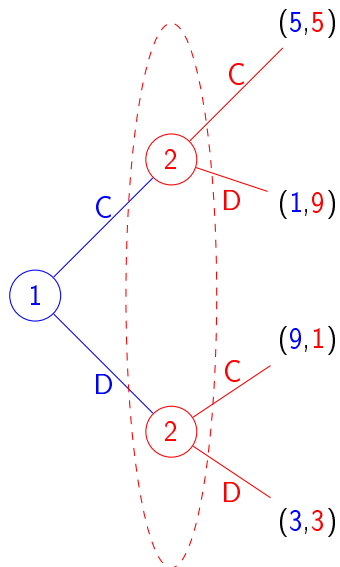
- Mixed strategies
 - unpredictability can be the only internally consistent behavior
 - does not mean we cannot make any predictions
 - mixed strategy equilibria often match the data well
 - counterintuitive comparative statics (costly monitoring)
- Zero-sum games and the MaxMin strategy
 - can treat a game as independent decision problems of pessimists

Dynamic games of complete information

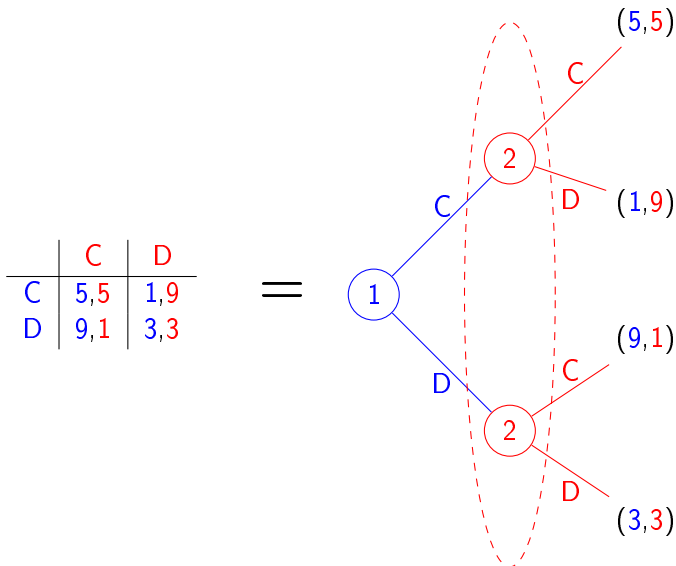
A game tree



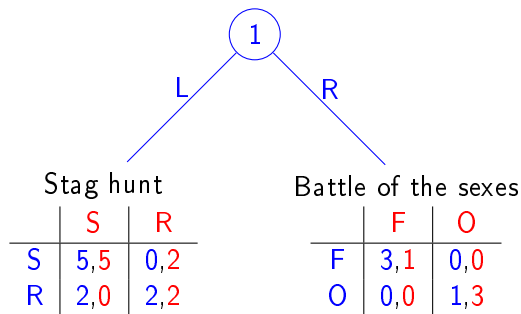
Information sets



Static games as game trees

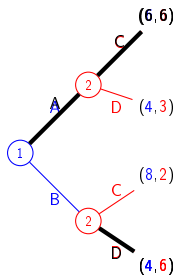


Choosing a game



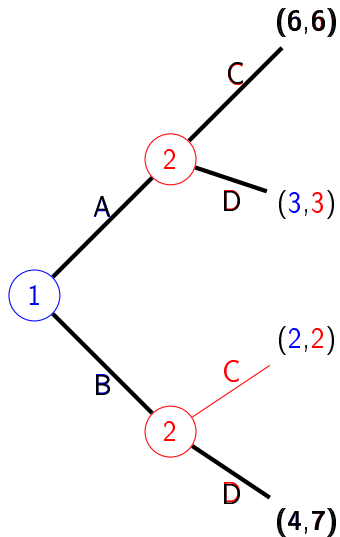
Subgame perfect equilibrium

Subgame perfect equilibrium

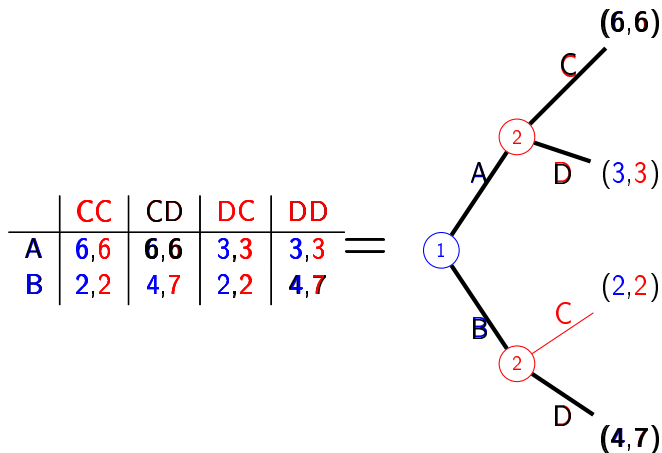


- Not enough to specify what happens in equilibrium
- Must specify what each player does at each potential node
- Only this way sure i is best-responding to j
- Related to, but distinct from, backward induction
 - A node can lead to a matrix; we assume NE in that matrix

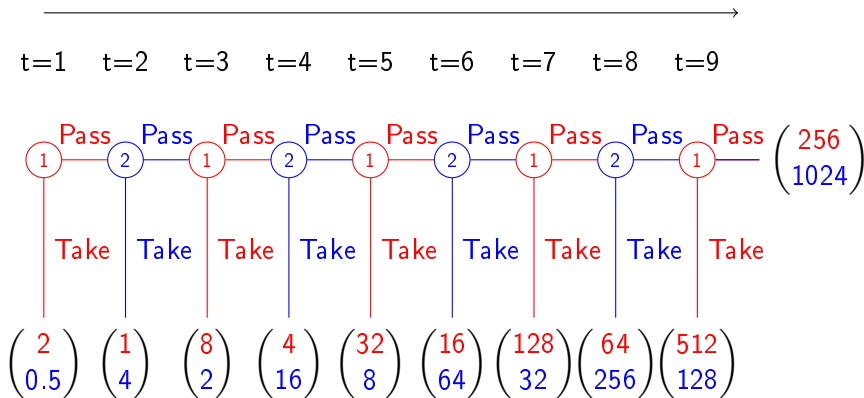
Nash vs. SPE



Nash vs. SPE

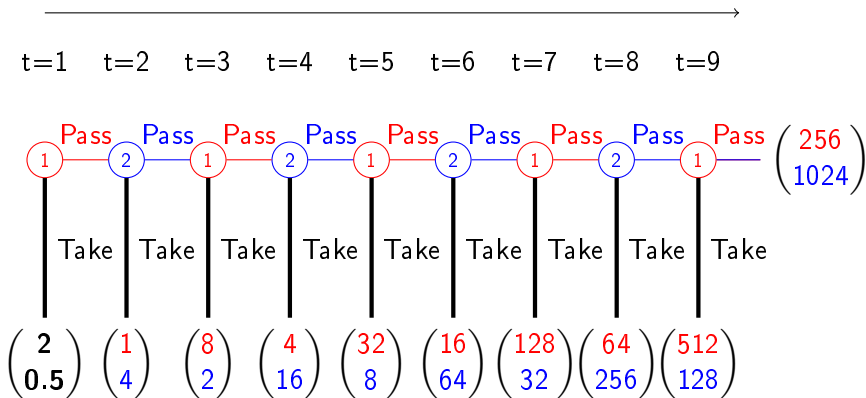


Centipede game



- If always Pass, select 0
- If Take at some node(s), select *first* node where you Take

Centipede game



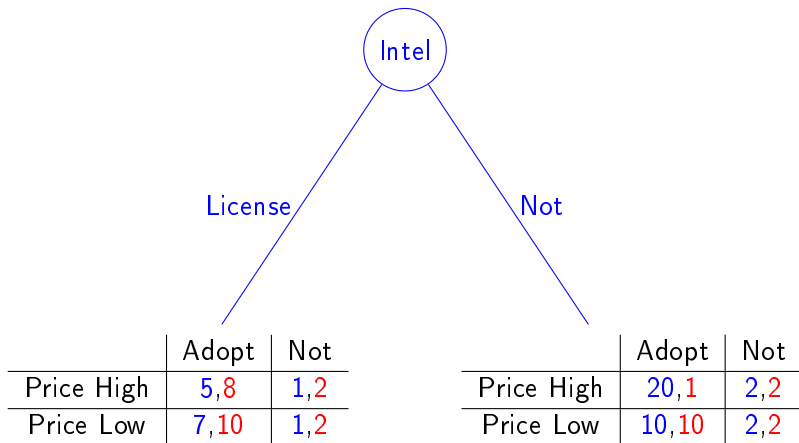
- As with IEDS, Subgame Perfect equilibrium requires common knowledge of rationality

Credible weakness

Second source licensing

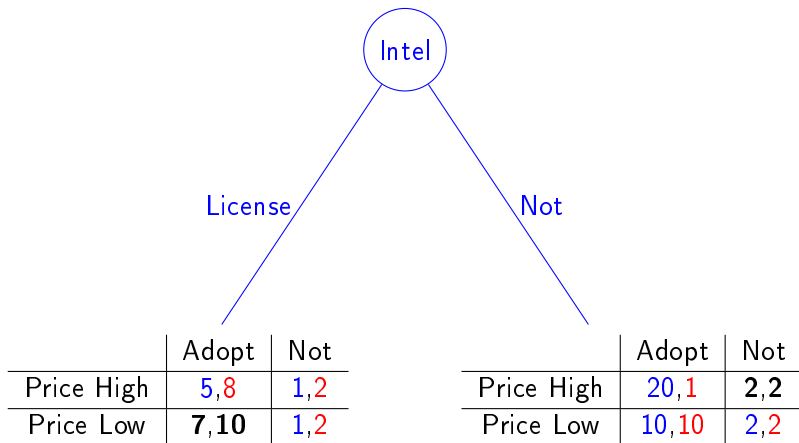
- 1978: Intel develops the 8086 microprocessor
 - grants second-source licensing to IBM, AMD, and 10 foreign firms
- 1987: Intel has only a 30% market share of 8086
- Intel could have been a monopolist
 - why did it voluntarily give away its market power?

Credible weakness



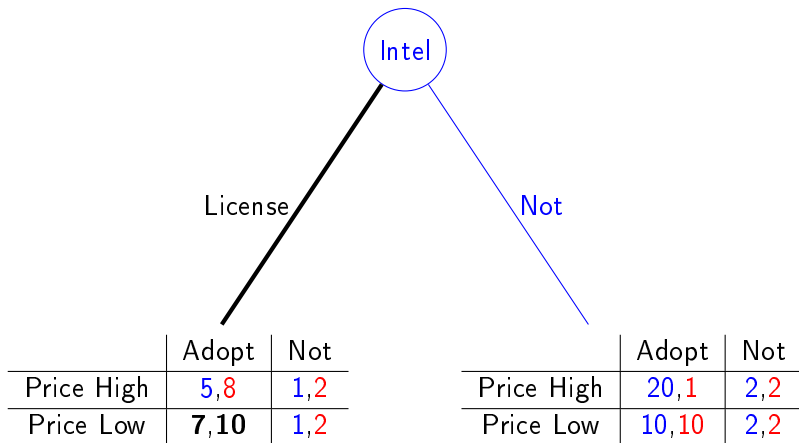
- If it licenses, Intel's payoff strictly lower for every possible outcome
- Hence, it was a mistake to license?
 - No

Credible weakness



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Credible weakness



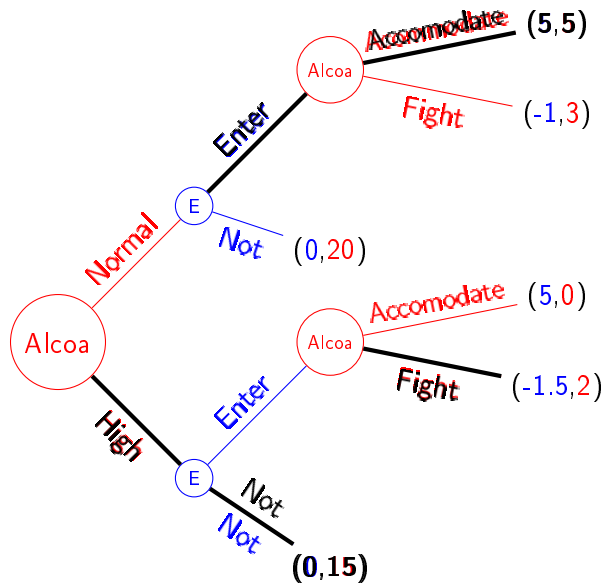
- If it licenses, Intel's payoff strictly lower for every possible outcome
- Hence, it was a mistake to license?
 - No

Credible belligerence

Excess capacity

- 1930s-40s Alcoa controlled the virgin aluminum market
- Built capacity well above the demand
- Building unused capacity is expensive
- Was it a mistake to have so much capacity?

Capacity game



Dynamic pricing

Dynamic pricing

- A monopolist sells a durable widget with zero marginal cost
- Consumers are patient and have unit demand
 - $\frac{1}{2}$ high valuation consumers willing to pay \$100
 - $\frac{1}{2}$ low valuation consumers willing to pay \$30
- What price should the monopolist charge?
 - Suppose he charges \$100
 - What price should he charge tomorrow?
 - But high valuation consumers will expect that
- No price greater than \$30 possible in a subgame perfect equilibrium!
- Monopolist's future self is a nasty competitor

Imperfectly patient consumers

- Suppose consumers have a discount rate of 10%
- Monopolist still cannot charge more than \$33.33

Solving the problem

- Blow up the company after selling to the first batch of consumers
- Price protection plans
 - GE in sales of electric-turbine generators
- Lease-only policy
 - United Shoe Machinery
- Planned obsolescence

Take aways

- Mixed strategies
 - unpredictability can be the only internally consistent behavior
 - does not mean we cannot make any predictions
 - mixed strategy equilibria often match the data well
 - counterintuitive comparative statics (costly monitoring)
- Zero-sum games and the MaxMin strategy
 - can treat a game as independent decision problems of pessimists

Take aways, continued

- We use game trees to analyze dynamic games
- Subgame perfection captures the notion of credibility
- Lowering your payoffs can increase your equilibrium payoff
- Becoming weaker can make it easier to be trusted
 - Intel and the 8086 microprocessor
- Becoming more irritable can make it easier to be feared
 - Alcoa and excess capacity
- Dynamic pricing
 - monopolist's tomorrow self can be his worst competitor
 - variety of strategies for alleviating the problem

Thank you