

## BUILDING THE TREE

11° We have our sequences, our evolutionary distances, and a metric that corresponds to a phylogenetic tree. How do we actually get our phylogenetic tree? We turn to the *Cherry-Picking Theorem* and the *Neighbor-Joining Algorithm* for the answer. These algorithms are given without too much explanation at this point. Essentially, one begins with a ‘template tree’ or network, then using the above information, one decides whether or not the nodes in the template tree belong or should be ‘pruned’. This process results in a phylogenetic tree  $T$  using our metric  $d_T$ .

**Theorem 2 (Computational Cherry-Picking Theorem).** Let  $d$  be a tree metric on  $[n]$ . For every pair  $i, j \in [n]$  set

$$Q_d(i, j) = (n - 2) \cdot d(i, j) - \sum_{k \neq i} d(i, k) - \sum_{k \neq j} d(j, k).$$

Then the pair  $x, y \in [n]$  that minimizes  $Q_d(x, y)$  is a cherry in the tree.

The above theorem is a computationally superior form of a more theoretically relevant statement of the Cherry-Picking Theorem. Now that we know how to pick a cherry out of a tree, we use the Neighbor-Joining Algorithm to build our tree.

**Algorithm 3 (Neighbor-joining algorithm).**

Input: A dissimilarity map  $d$  on the set  $[n]$ .

Output: A phylogenetic tree  $T$  whose tree metric  $d_T$  is “close” to  $d$ .

*Step 1:* Construct the  $n \times n$  matrix  $Q_d$  whose  $(i, j)$ -entry is given by Theorem 2, and identify the minimum off-diagonal entry  $Q_d(x, y)$ .

*Step 2:* Remove  $x, y$  from the tree, thereby creating a new leaf  $z$ . For each leaf  $k$  among the remaining  $n - 2$  leaves, set

$$d(z, k) = \frac{1}{2}(d(x, k) + d(y, k) - d(x, y)).$$

This replaces the  $n \times n$  matrix  $Q_d$  by an  $(n - 1) \times (n - 1)$  matrix. Return to Step 1 until there are no more leaves to collapse.

*Step 3:* Output the tree  $T$ . The edge lengths of  $T$  are determined recursively: if  $(x, y)$  is a cherry connected to node  $z$  as in Step 2, then the edge from  $x$  to  $z$  has length  $d(x, k) - d_T(z, k)$  and the edge from  $y$  to  $z$  has length  $d(y, k) - d_T(z, k)$ .

Using the Cherry-Picking Theorem and the Neighbor-Joining Algorithm on our metric, we receive the following phylogenetic tree:

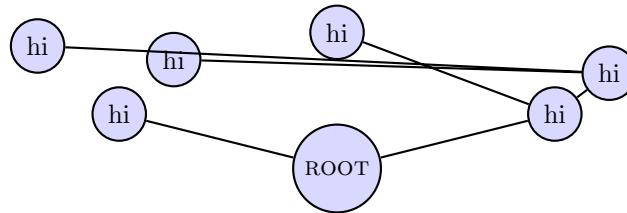
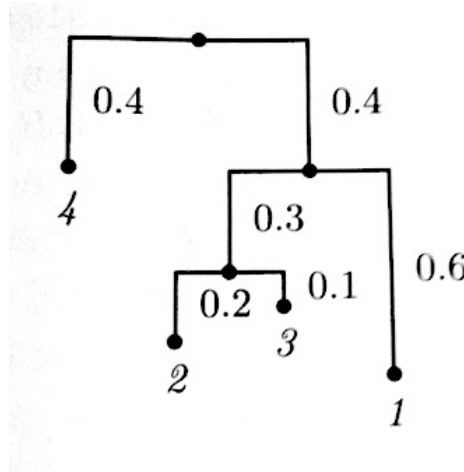


Figure 5: Tree/Form