# **ACM ICPC Reference**

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#### 1 Data Structures

## Treap (balanced binary search tree)

```
struct node {
  int v, key, size;
  node *c[2];
  void resize() { size = c[0]->size + c[1]->size + 1; }
node *newNode(int _v, node *n) {
  pool[ref].v = _v, pool[ref].c[0] = pool[ref].c[1] = n, pool[ref].size = 1,
      pool[ref].key = rand();
  return &pool[ref];
struct Treap {
  node *root, *nil;
  void rotate(node *&t, int d) {
    node *c = t - > c[d];
   t - c[d] = c - c[!d];
   c \rightarrow c[!d] = t;
   t->resize(); c->resize();
    t = c;
  void insert(node *&t, int x) {
    if (t == nil) t = newNode(x, nil);
    else {
      if (x == t->v) return;
      int d = x > t -> v;
      insert(t->c[d], x);
      if (t->c[d]->key < t->key) rotate(t, d);
      else t->resize();
   }
  void remove(node *&t, int x) {
```

```
if (t == nil) return:
  if (t->v == x) {
    int d = t - c[1] - key < t - c[0] - key;
    if (t->c[d] == nil) {
      t = nil:
      return;
    rotate(t, d);
    remove(t \rightarrow c[!d], x);
    int d = x > t -> v;
    remove(t->c[d], x);
  t->resize();
int rank(node *t, int x) {
  if (t == nil) return 0;
  int r = t -> c[0] -> size;
  if (x == t -> v) return r + 1;
  if (x < t->v) return rank(t->c[0], x);
  return r + 1 + rank(t->c[1], x);
int select(node *t, int k) {
  int r = t -> c[0] -> size;
  if (k == r + 1) return t \rightarrow v;
  if (k <= r) return select(t->c[0], k);
  return select(t->c[1], k - r - 1);
int size() {
  return root->size;
void init(int *a, int n) {
  nil = newNode(0, 0);
  nil->size = 0, nil->key = ~0U >> 1;
  root = nil;
```

## 2 Geometry

};

### Welzl's algorithm (minimum enclosing circle

```
// Minimum enclosing circle, Welzl's algorithm
// Expected linear time
// If there are any duplicate points in the input, be sure to remove them
   first.
struct point {
  double x;
 double y;
struct circle {
  double x:
  double y;
  double r:
  circle() {}
  circle(double x, double y, double r): x(x), y(y), r(r) {}
circle b md(vector < point > R) {
 if (R.size() == 0) {
   return circle(0, 0, -1);
 } else if (R.size() == 1) {
    return circle(R[0].x, R[0].y, 0);
 } else if (R.size() == 2) {
    return circle((R[0].x+R[1].x)/2.0,
                  (R[0].y+R[1].y)/2.0,
            hypot (R[0].x-R[1].x, R[0].y-R[1].y)/2.0);
 } else {
    double D = (R[0].x - R[2].x)*(R[1].y - R[2].y) - (R[1].x - R[2].x)*(R[0].y
         - R[2].v);
   double p0 = (((R[0].x - R[2].x)*(R[0].x + R[2].x) + (R[0].y - R[2].y)*(R[0].x + R[2].x)
        [0].y + R[2].y) / 2 * (R[1].y - R[2].y) - ((R[1].x - R[2].x)*(R[1].x
       + R[2].x) + (R[1].y - R[2].y)*(R[1].y + R[2].y)) / 2 * (R[0].y - R[2].
       y))/D;
   double p1 = (((R[1].x - R[2].x)*(R[1].x + R[2].x) + (R[1].y - R[2].y)*(R[1].x)
        [1].y + R[2].y) / 2 * (R[0].x - R[2].x) - ((R[0].x - R[2].x)*(R[0].x
       + R[2].x) + (R[0].y - R[2].y)*(R[0].y + R[2].y)) / 2 * (R[1].x - R[2].y)
       x))/D;
   return circle(p0, p1, hypot(R[0].x - p0, R[0].y - p1));
 }
```

```
circle b_minidisk(vector<point>& P, int i, vector<point> R) {
   if (i == P.size() || R.size() == 3) {
      return b_md(R);
   } else {
      circle D = b_minidisk(P, i+1, R);
      if (hypot(P[i].x-D.x, P[i].y-D.y) > D.r) {
        R.push_back(P[i]);
      D = b_minidisk(P, i+1, R);
   }
   return D;
}

// Call this function.
circle minidisk(vector<point> P) {
   random_shuffle(P.begin(), P.end());
   return b_minidisk(P, 0, vector<point>());
}
```

## Monotone chain (convex hull) and rotating calipers (farthest pair)

```
typedef long double gtype;
const gtype pi = M_PI;
typedef complex < gtype > point;
#define x real()
#define y imag()
#define polar(r, t) polar((gtype) (r), (t))
// vector
#define rot(v, t) ((v) * polar(1, t))
#define crs(a, b) ((conj(a) * (b)).v)
#define dot(a, b) ( (conj(a) * (b)).x )
#define pntLinDist(a, b, p) ( abs(crs((b)-(a), (p)-(a)) / abs((b)-(a))) )
bool cmp_point(point const& p1, point const& p2) {
    return p1.x == p2.x ? (p1.y < p2.y) : (p1.x < p2.x);
}
// O(n.log(n)) - monotone chain
vector < point > mcH;
void monotoneChain(vector<point> &ps) {
    vector < point > p(ps.begin(), ps.end() - 1);
    int n = p.size(), k = 0;
    mcH = vector < point > (2 * n);
```

```
sort(p.begin(), p.end(), cmp_point);
   for (int i = 0; i < n; i++) {
        while (k \ge 2 \&\& crs(mcH[k - 1] - mcH[k - 2], p[i] - mcH[k - 2]) \le 0)
            k--:
        mcH[k++] = p[i];
   for (int i = n - 2, t = k + 1; i \ge 0; i--) {
        while (k \ge t \&\& crs(mcH[k - 1] - mcH[k - 2], p[i] - mcH[k - 2]) \le 0)
            k - - ;
        mcH[k++] = p[i];
   mcH.resize(k);
// O(n) - rotating calipers (works on a ccw closed convex hull)
gtype rotatingCalipers(vector<point> &ps) {
   int aI = 0, bI = 0;
   for (size_t i = 1; i < ps.size(); ++i)</pre>
        aI = (ps[i].y < ps[aI].y ? i : aI), bI = (ps[i].y > ps[bI].y ? i : bI)
    gtype minWidth = ps[bI].y - ps[aI].y, aAng, bAng;
    point aV = point(1, 0), bV = point(-1, 0);
   for (gtype ang = 0; ang < pi; ang += min(aAng, bAng)) {</pre>
        aAng = acos(dot(ps[aI + 1] - ps[aI], aV)
            / abs(aV) / abs(ps[aI + 1] - ps[aI]));
       bAng = acos(dot(ps[bI + 1] - ps[bI], bV)
            / abs(bV) / abs(ps[bI + 1] - ps[bI]));
        aV = rot(aV, min(aAng, bAng)), bV = rot(bV, min(aAng, bAng));
       if (aAng < bAng)
            minWidth = min(minWidth, pntLinDist(ps[aI], ps[aI] + aV, ps[bI]))
            , aI = (aI + 1) \% (ps.size() - 1);
        else
            minWidth = min(minWidth, pntLinDist(ps[bI], ps[bI] + bV, ps[aI]))
            , bI = (bI + 1) \% (ps.size() - 1);
   }
    return minWidth;
```

#### 3d Convex Hull

```
int n, bf[maxn][maxn], fcnt;
point3_t pt[maxn];
struct face_t {
  int a, b, c;
```

```
bool vis;
} fc[maxn << 5]; /* Number of Faces(Unknown) */</pre>
bool remove(int p, int b, int a) {
 int f = bf[b][a];
  face_t ff;
  if (fc[f].vis) {
    if (dblcmp(volume(pt[p], pt[fc[f].a], pt[fc[f].b], pt[fc[f].c])) >= 0) {
      return true;
   } else {
      ff.a = a, ff.b = b, ff.c = p;
      bf[ff.a][ff.b] = bf[ff.b][ff.c] = bf[ff.c][ff.a] = ++fcnt;
     ff.vis = true:
     fc[fcnt] = ff:
 }
  return false:
void dfs(int p, int f) {
 fc[f].vis = false;
  if (remove(p, fc[f].b, fc[f].a)) dfs(p, bf[fc[f].b][fc[f].a]);
  if (remove(p, fc[f].c, fc[f].b)) dfs(p, bf[fc[f].c][fc[f].b]);
  if (remove(p, fc[f].a, fc[f].c)) dfs(p, bf[fc[f].a][fc[f].c]);
}
void hull3d() {
  for (int i = 2; i <= n; ++i) {
    if (dblcmp((pt[i] - pt[1]).length()) > 0) swap(pt[i], pt[2]);
 }
  for (int i = 3; i <= n; ++i) {
    if (dblcmp(fabs(area(pt[1], pt[2], pt[i]))) > 0) swap(pt[i], pt[3]);
  }
  for (int i = 4; i <= n; ++i) {
    if (dblcmp(fabs(volume(pt[1], pt[2], pt[3], pt[i]))) > 0) swap(pt[i], pt
        [4]);
  zm(fc), fcnt = 0, zm(bf);
  for (int i = 1; i <= 4; ++i) {
    face t f:
    f.a = i + 1, f.b = i + 2, f.c = i + 3;
    if (f.a > 4) f.a -= 4:
    if (f.b > 4) f.b -= 4;
    if (f.c > 4) f.c -= 4:
    if (dblcmp(volume(pt[i], pt[f.a], pt[f.b], pt[f.c])) > 0) swap(f.a, f.b);
    f.vis = true:
```

```
bf[f.a][f.b] = bf[f.b][f.c] = bf[f.c][f.a] = ++fcnt;
fc[fcnt] = f;
}
random_shuffle(pt + 5, pt + 1 + n);
for (int i = 5; i <= n; ++i) {
    for (int j = 1; j <= fcnt; ++j) {
        if (!fc[j].vis) continue;
        if (dblcmp(volume(pt[i], pt[fc[j].a], pt[fc[j].b], pt[fc[j].c])) >= 0) {
        dfs(i, j);
        break;
    }
}
for (int i = 1; i <= fcnt; ++i) if (!fc[i].vis) swap(fc[i--], fc[fcnt--]);
}</pre>
```

### 3 Graph

#### Min cost flow

```
/* Min cost max flow (Edmonds-Karp relabelling + fast heap Dijkstra)
* Based on code by Frank Chu and Igor Naverniouk
* (http://shygypsy.com/tools/mcmf4.cpp)
* COMPLEXITY:
        - Worst case: O(min(m*log(m)*flow, n*m*log(m)*fcost))
* FIELD TESTING:
        - Valladolid 10594: Data Flow
* REFERENCE:
        Edmonds, J., Karp, R. "Theoretical Improvements in Algorithmic
           Efficieincy for Network Flow Problems".
        This is a slight improvement of Frank Chu's implementation.
#define Inf (LLONG_MAX/2)
#define BUBL { \
   t = q[i]; q[i] = q[j]; q[j] = t; \
   t = inq[q[i]]; inq[q[i]] = inq[q[j]]; inq[q[j]] = t; }
#define Pot(u,v) (d[u] + pi[u] - pi[v])
struct MinCostMaxFlow {
 typedef long long LL;
 int n, qs;
 vector < vector < LL > > cap, cost, fnet;
```

```
vector < vector < int > > adj;
vector < LL > d, pi;
vector<int> deg, par, q, inq;
// n = number of vertices
MinCostMaxFlow(int n): n(n), qs(0), deg(n+1), par(n+1), d(n+1), q(n+1), inq(
    n+1), pi(n+1), cap(n+1), vector < LL > (n+1)), cost(cap), fnet(cap), adj(n+1),
     vector < int > (n+1)) {}
// call to add a directed edge. vertices are 0-based
// ALL COSTS MUST BE NON-NEGATIVE
void AddEdge(int from, int to, LL cap_, LL cost_) {
  cap[from][to] = cap_; cost[from][to] = cost_;
}
bool dijkstra( int s, int t ) {
  fill(d.begin(), d.end(), 0x3f3f3f3f3f3f3f3f1LL);
  fill(par.begin(), par.end(), -1);
  fill(inq.begin(), inq.end(), -1);
  d[s] = qs = 0;
  inq[q[qs++] = s] = 0;
  par[s] = n;
  while( qs ) {
   int u = q[0]; inq[u] = -1;
   q[0] = q[--qs];
    if(qs) inq[q[0]] = 0;
    for( int i = 0, j = 2*i + 1, t; j < qs; i = j, j = 2*i + 1 ) {
      if (j + 1 < qs && d[q[j + 1]] < d[q[j]]) j++;
      if( d[q[i]] >= d[q[i]] ) break;
      BUBL;
   }
    for ( int k = 0, v = adj[u][k]; k < deg[u]; <math>v = adj[u][++k] ) {
      if( fnet[v][u] && d[v] > Pot(u,v) - cost[v][u] )
        d[v] = Pot(u,v) - cost[v][par[v] = u];
      if ( fnet [u][v] < cap[u][v] && d[v] > Pot(u,v) + cost[u][v] )
        d[v] = Pot(u,v) + cost[par[v] = u][v];
      if( par[v] == u ) {
        if ( inq[v] < 0 ) { inq[q[qs] = v] = qs; qs++; }
        for ( int i = inq[v], j = ( i - 1 )/2, t;
           d[q[i]] < d[q[j]]; i = j, j = (i - 1)/2)
           BUBL;
      }
   }
  for( int i = 0; i < n; i++ ) if( pi[i] < Inf ) pi[i] += d[i];</pre>
  return par[t] >= 0:
```

```
}
  // Returns: (flow, total cost) between source s and sink t
  // Call this once only. fnet[i][j] contains the flow from i to j. Careful,
      fnet[i][j] and fnet[j][i] could both be positive.
  pair<LL, LL> mcmf4(int s, int t) {
   for( int i = 0; i < n; i++ )
     for ( int j = 0; j < n; j++ )
        if( cap[i][j] || cap[j][i] ) adj[i][deg[i]++] = j;
   LL flow = 0; LL fcost = 0;
    while( dijkstra( s, t ) ) {
      LL bot = LLONG_MAX;
     for( int v = t, u = par[v]; v != s; u = par[v = u] )
       bot = min(bot, fnet[v][u] ? fnet[v][u] : ( cap[u][v] - fnet[u][v] ));
      for( int v = t, u = par[v]; v != s; u = par[v = u] )
        if( fnet[v][u] ) { fnet[v][u] -= bot; fcost -= bot * cost[v][u]; }
        else { fnet[u][v] += bot: fcost += bot * cost[u][v]: }
      flow += bot;
    return make_pair(flow, fcost);
 }
};
```

# Dinic's ( $VE^2$ max flow)

```
// edge in adjacency list so that
              // we can quickly find it.
};
// Residual Graph
class Graph
{
    int V; // number of vertex
    int *level ; // stores level of a node
    vector < Edge > *adj;
public :
    Graph(int V)
        adj = new vector < Edge > [V];
        this -> V = V;
        level = new int[V];
    // add edge to the graph
    void addEdge(int u, int v, int C)
        // Forward edge : 0 flow and C capacity
        Edge a{v, 0, C, adj[v].size()};
        // Back edge : 0 flow and 0 capacity
        Edge b{u, 0, 0, adj[u].size()};
        adj[u].push_back(a);
        adj[v].push_back(b); // reverse edge
    }
    bool BFS(int s, int t);
    int sendFlow(int s, int flow, int t, int ptr[]);
    int DinicMaxflow(int s, int t);
};
// Finds if more flow can be sent from s to t.
// Also assigns levels to nodes.
bool Graph::BFS(int s, int t)
    for (int i = 0; i < V; i++)
        level[i] = -1;
    level[s] = 0; // Level of source vertex
    // Create a queue, enqueue source vertex
```

```
// and mark source vertex as visited here
    // level[] array works as visited array also.
    list < int > q;
    g.push_back(s);
    vector < Edge > :: iterator i ;
    while (!q.empty())
       int u = q.front();
        q.pop_front();
        for (i = adj[u].begin(); i != adj[u].end(); i++)
            Edge &e = *i;
            if (level[e.v] < 0 && e.flow < e.C)
                // Level of current vertex is,
                // level of parent + 1
                level[e.v] = level[u] + 1;
                q.push_back(e.v);
            }
        }
    // IF we can not reach to the sink we
    // return false else true
    return level[t] < 0 ? false : true ;</pre>
// A DFS based function to send flow after BFS has
// figured out that there is a possible flow and
// constructed levels. This function called multiple
// times for a single call of BFS.
// flow : Current flow send by parent function call
// start[] : To keep track of next edge to be explored.
             start[i] stores count of edges explored
             from i.
// u : Current vertex
// t : Sink
int Graph::sendFlow(int u, int flow, int t, int start[])
   // Sink reached
    if (u == t)
        return flow;
   // Traverse all adjacent edges one -by - one.
```

```
for ( ; start[u] < adj[u].size(); start[u]++)</pre>
        // Pick next edge from adjacency list of u
        Edge &e = adj[u][start[u]];
        if (level[e.v] == level[u]+1 && e.flow < e.C)
            // find minimum flow from u to t
            int curr_flow = min(flow, e.C - e.flow);
            int temp_flow = sendFlow(e.v, curr_flow, t, start);
            // flow is greater than zero
            if (temp_flow > 0)
                // add flow to current edge
                e.flow += temp_flow;
                // subtract flow from reverse edge
                // of current edge
                adj[e.v][e.rev].flow -= temp_flow;
                return temp_flow;
        }
    return 0;
}
// Returns maximum flow in graph
int Graph::DinicMaxflow(int s, int t)
{
    // Corner case
    if (s == t)
        return -1;
    int total = 0: // Initialize result
    // Augment the flow while there is path
    // from source to sink
    while (BFS(s, t) == true)
        // store how many edges are visited
        // from V { 0 to V }
        int *start = new int[V+1]:
```

```
// while flow is not zero in graph from S to D
       while (int flow = sendFlow(s, INT_MAX, t, start))
           // Add path flow to overall flow
           total += flow;
   }
    // return maximum flow
    return total;
// Driver program to test above functions
int main()
   Graph g(6);
   g.addEdge(0, 1, 16);
   g.addEdge(0, 2, 13);
   g.addEdge(1, 2, 10);
   g.addEdge(1, 3, 12);
   g.addEdge(2, 1, 4);
   g.addEdge(2, 4, 14);
   g.addEdge(3, 2, 9);
   g.addEdge(3, 5, 20);
   g.addEdge(4, 3, 7);
   g.addEdge(4, 5, 4);
   // next exmp
    /*g.addEdge(0, 1, 3);
     g.addEdge(0, 2, 7);
     g.addEdge(1, 3, 9);
     g.addEdge(1, 4, 9);
     g.addEdge(2, 1, 9);
     g.addEdge(2, 4, 9);
     g.addEdge(2, 5, 4);
     g.addEdge(3, 5, 3);
     g.addEdge(4, 5, 7);
     g.addEdge(0, 4, 10);
    // next exp
    g.addEdge(0, 1, 10);
    g.addEdge(0, 2, 10);
    g.addEdge(1, 3, 4);
    g.addEdge(1, 4, 8);
    g.addEdge(1, 2, 2);
    g.addEdge(2, 4, 9);
    g.addEdge(3, 5, 10);
```

```
g.addEdge(4, 3, 6);
g.addEdge(4, 5, 10); */

cout << "Maximum_flow_" << g.DinicMaxflow(0, 5);
return 0;
}</pre>
```

#### Edmond's algorithm (unweighted general matching)

```
// Unweighted general matching.
// Vertices are numbered from 1 to V.
// G is an adjlist.
// G[x][0] contains the number of neighbours of x.
// The neighbours are then stored in G[x][1] .. G[x][G[x][0]].
// Mate[x] will contain the matching node for x.
// V and E are the number of edges and vertices.
// Slow Version (2x on random graphs) of Gabow's implementation
// of Edmonds' algorithm (O(V^3)).
const int MAXV = 250:
int G[MAXV][MAXV];
int VLabel[MAXV];
int Queue[MAXV]:
int Mate[MAXV];
int Save[MAXV];
int Used[MAXV]:
     Up, Down;
int
int
              ۷;
void ReMatch(int x, int y)
 int m = Mate[x]; Mate[x] = y;
 if (Mate[m] == x)
    {
      if (VLabel[x] <= V)</pre>
          Mate[m] = VLabel[x];
          ReMatch(VLabel[x], m);
        }
      else
          int a = 1 + (VLabel[x] - V - 1) / V;
          int b = 1 + (VLabel[x] - V - 1) % V:
          ReMatch(a, b); ReMatch(b, a);
```

```
}
void Traverse(int x)
  for (int i = 1; i <= V; i++) Save[i] = Mate[i];</pre>
  ReMatch(x, x);
  for (int i = 1; i <= V; i++)</pre>
      if (Mate[i] != Save[i]) Used[i]++;
      Mate[i] = Save[i];
   }
}
void ReLabel(int x, int y)
  for (int i = 1; i <= V; i++) Used[i] = 0;</pre>
  Traverse(x); Traverse(y);
  for (int i = 1: i <= V: i++)
      if (Used[i] == 1 && VLabel[i] < 0)</pre>
          VLabel[i] = V + x + (y - 1) * V;
          Queue[Up++] = i;
        }
   }
// Call this after constructing G
void Solve()
  for (int i = 1; i <= V; i++)</pre>
    if (Mate[i] == 0)
     {
        for (int j = 1; j <= V; j++) VLabel[j] = -1;
        VLabel[i] = 0; Down = 1; Up = 1; Queue[Up++] = i;
        while (Down != Up)
            int x = Queue[Down++];
            for (int p = 1; p <= G[x][0]; p++)
              {
                int y = G[x][p];
                if (Mate[y] == 0 && i != y)
                     Mate[y] = x; ReMatch(x, y);
```

```
Down = Up; break;
                  }
                if (VLabel[y] >= 0)
                  {
                    ReLabel(x, y);
                    continue;
                  }
                if (VLabel[Mate[y]] < 0)</pre>
                    VLabel[Mate[y]] = x;
                    Queue[Up++] = Mate[y];
                  }
              }
          }
     }
// Call this after Solve(). Returns number of edges in matching (half the
    number of matched vertices)
int Size()
 int Count = 0;
 for (int i = 1; i <= V; i++)
   if (Mate[i] > i) Count++;
 return Count;
```

8

#### Link cut tree

```
struct node_t {
   node_t();
   node_t *ch[2], *p;
   int size, root;
   int dir() { return this == p->ch[1]; }
   void setc(node_t *c, int d) { ch[d] = c, c->p = this; }
   void update() { size = ch[0]->size + ch[1]->size + 1; }
} s[maxn], *nil = s;

node_t::node_t() {
   size = 1, root = true;
   ch[0] = ch[1] = p = nil;
}
```

```
void rotate(node_t *t) {
  node_t *p = t->p;
  int d = t - > dir();
  if (!p->root) {
    p->p->setc(t, p->dir());
    p->root = false, t->root = true;
    t \rightarrow p = p \rightarrow p; // Path Parent
  p->setc(t->ch[!d], d);
  t->setc(p, !d);
  p->update(), t->update();
void splay(node_t *t) {
  // t->update(); // tag!
  while (!t->root) {
    // if (!t->p->root) t->p->update(); t->p->update(); // !
    if (!t->p->root) rotate(t->dir() == t->p->dir() ? t->p : t);
    rotate(t):
  }
}
void access(node_t *x) { // Ask u, v: access(u), access(v, true), x = LCA
  node_t *y = nil;
  while (x != nil) {
    splay(x);
    // if (x-p == nil) at second call, x-ch[1](rev) + (x)_single + y
    x \rightarrow ch[1] \rightarrow root = true;
    x \rightarrow ch[1] = y, y \rightarrow root = false;
    x->update();
    y = x, x = x->p;
  }
void cut(node_t *x) {
  access(x);
  splay(x);
  x - ch[0] - root = true;
  x - ch[0] - p = nil;
  x - ch[0] = nil;
void link(node_t *x, node_t *y) {
  access(y);
  splay(y);
```

```
y->p = x;
access(y);
}
void init() { nil->size = 0; }
```

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### 4 Strings

## KMP (linear string search)

```
typedef vector<int> VI;
void buildTable(string& w, VI& t)
  t = VI(w.length());
  int i = 2, j = 0;
  t[0] = -1; t[1] = 0;
  while(i < w.length())</pre>
    if(w[i-1] == w[j]) { t[i] = j+1; i++; j++; }
    else if(j > 0) j = t[j];
    else { t[i] = 0; i++; }
 }
}
int KMP(string& s, string& w)
  int m = 0, i = 0;
  VI t;
  buildTable(w, t);
  while(m+i < s.length())</pre>
    if(w[i] == s[m+i])
      i++;
     if(i == w.length()) return m;
    else
      m += i-t[i]:
      if(i > 0) i = t[i];
```

```
}
return s.length();
}
```

## Manacher (max palindrome substring)

```
// Manacher's algorithm: finds maximal palindrome lengths centered around each
// position in a string (including positions between characters) and returns
// them in left-to-right order of centres. Linear time.
// Ex: "opposes" -> [0, 1, 0, 1, 4, 1, 0, 1, 0, 1, 0, 3, 0, 1, 0]
vector < int > fastLongestPalindromes(string str) {
   int i=0,j,d,s,e,lLen,palLen=0;
   vector < int > res;
    while (i < str.length()) {</pre>
        if (i > palLen && str[i-palLen-1] == str[i]) {
            palLen += 2; i++; continue;
       res.push_back(palLen);
        s = res.size()-2;
        e = s-palLen;
        bool b = true;
        for (j=s; j>e; j--) {
            d = i - e - 1;
            if (res[j] == d) { palLen = d; b = false; break; }
            res.push_back(min(d, res[j]));
       }
        if (b) { palLen = 1; i++; }
   res.push_back(palLen);
   lLen = res.size();
   s = 1Len-2:
   e = s-(2*str.length()+1-lLen);
   for (i=s; i>e; i--) { d = i-e-1; res.push_back(min(d, res[i])); }
    return res:
```

# Suffix array

```
// Suffix array construction in O(L log^2 L) time. Routine for
```

```
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
11
// INPUT: string s
11
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
            of substring s[i...L-1] in the list of sorted suffixes.
11
            That is, if we take the inverse of the permutation suffix[],
            we get the actual suffix array.
struct SuffixArray {
  const int L:
  string s;
  vector < vector < int > > P;
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s): L(s.length()), s(s), P(1, vector<int>(L, 0)),
       M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>
    for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push back(vector<int>(L. 0));
      for (int i = 0; i < L; i++)
             M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level
                 -1][i + skip] : -1000), i):
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)
             P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P
                 [level][M[i-1].second] : i;
  }
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {
      if (P[k][i] == P[k][j]) {
        i += 1 << k:
        j += 1 << k;
        len += 1 << k:
    return len;
  }
```

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};

#### 5 Math

#### Chinese Remainder Theorem

```
/* Extended Euclidean Algorithm
* find x, y s.t. ax+by=gcd(a,b)
*/
void eea(int a, int b, int &x, int &y) {
   int r[3] = \{a, b\}, s[3] = \{1, 0\}, t[3] = \{0, 1\};
   while (r[1]) {
       int q = r[0] / r[1];
       r[2] = r[0] - q * r[1];
       s[2] = s[0] - q * s[1];
       t[2] = t[0] - q * t[1];;
       r[0] = r[1]; r[1] = r[2];
       s[0] = s[1]; s[1] = s[2];
       t[0] = t[1]; t[1] = t[2];
   x = s[0]; y = t[0];
/* Chinese Remainder Theorem
* find x s.t. x = a[i] mod b[i]
*/
int crt(int *a, int *b, int n) {
   int B = 1;
   for (int i = 0; i < n; ++i)
       B *= b[i];
   int x = 0;
   for (int i = 0; i < n; ++i) {
       int c, d;
       eea(b[i], B / b[i], c, d);
       x = (x + B / b[i] * d * a[i]) % B:
   x = (x + B) \% B;
    return x;
```

#### **Fast Fourier Transform**

```
typedef complex <double > cd;
int const NMAX = 1 << 9;</pre>
double const PI2 = atan(1.0) * 8;
cd a[NMAX], b[NMAX];
// fft(src, num, stride, dst, nth root of unity)
// e.g. fft(a, n, 1, b, polar(1.0, -PI2 / n))
void fft(cd *a, int n, int s, cd *b, cd unit) {
    if (n == 1) {
        *b = *a:
        return;
    }
    int nh = n / 2;
    fft(a , nh, s * 2, b , unit * unit);
    fft(a + s, nh, s * 2, b + nh, unit * unit);
    cd coef = 1;
    for (int i = 0; i < nh; ++i) {</pre>
        cd ofs = coef * b[i + nh];
        b[i + nh] = b[i] - ofs;
        b[i]
                 = b[i] + ofs;
        coef *= unit;
    }
}
```

## Simplex Algorithm (Linear programming)

```
// Two-phase simplex algorithm for solving linear programs of the form
11
11
       maximize
                    c^T x
       subject to Ax <= b
11
                    x >= 0
11
// INPUT: A -- an m x n matrix
          b -- an m-dimensional vector
          c -- an n-dimensional vector
11
          x -- a vector where the optimal solution will be stored
// OUTPUT: value of the optimal solution (infinity if unbounded
11
           above, nan if infeasible)
11
```

```
// To use this code, create an LPSolver object with A, b, and c as
// arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector < DOUBLE > VD;
typedef vector < VD > VVD;
typedef vector <int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
 int m, n;
 VI B, N;
  VVD D:
 LPSolver (const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2))  {
   for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D[i][j] = A[i][j];
   for (int i = 0; i < m; i++) { B[i] = n+i; D[i][n] = -1; D[i][n+1] = b[i];
   for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j]; }
   N[n] = -1; D[m+1][n] = 1;
 }
  void Pivot(int r, int s) {
   DOUBLE inv = 1.0 / D[r][s];
   for (int i = 0; i < m+2; i++) if (i != r)
     for (int j = 0; j < n+2; j++) if (j != s)
        D[i][j] -= D[r][j] * D[i][s] * inv;
   for (int j = 0; j < n+2; j++) if (j != s) D[r][j] *= inv;
   for (int i = 0; i < m+2; i++) if (i != r) D[i][s] *= -inv;
   D[r][s] = inv;
   swap(B[r], N[s]);
 }
 bool Simplex(int phase) {
   int x = phase == 1 ? m+1 : m;
    while (true) {
     int s = -1;
     for (int j = 0; j <= n; j++) {
       if (phase == 2 && N[j] == -1) continue;
       if (s == -1 \mid | D[x][j] < D[x][s] \mid | D[x][j] == D[x][s] && N[j] < N[s])
             s = j;
      if (s < 0 \mid \mid D[x][s] > -EPS) return true;
      int r = -1:
```

```
for (int i = 0; i < m; i++) {
        if (D[i][s] < EPS) continue;</pre>
        if (r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r][s] ||
            D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i] < B[r]) r = i;
      if (r == -1) return false;
      Pivot(r, s);
  }
  DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r = i;
    if (D[r][n+1] \leftarrow -EPS) {
      Pivot(r, n);
      if (!Simplex(1) || D[m+1][n+1] < -EPS) return -numeric_limits < DOUBLE >::
          infinity():
      for (int i = 0; i < m; i++) if (B[i] == -1) {
        int s = -1:
        for (int j = 0; j <= n; j++)
          if (s == -1 \mid \mid D[i][i] < D[i][s] \mid \mid D[i][i] == D[i][s] && N[i] < N[s]
              ]) s = j;
        Pivot(i, s):
      }
    if (!Simplex(2)) return numeric_limits < DOUBLE >:: infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][n+1];
    return D[m][n+1];
};
```

#### Gaussian elimination

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
// (1) solving systems of linear equations (AX=B)
// (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
//
// Running time: O(n^3)
//
// INPUT: a[][] = an nxn matrix
```

```
11
             b[][] = an nxm matrix
             A MUST BE INVERTIBLE!
// OUTPUT: X
                    = an nxm matrix (stored in b[][])
             A^{-1} = an nxn matrix (stored in a[][])
             returns determinant of a[][]
const double EPS = 1e-10;
typedef vector <int> VI;
typedef double T;
typedef vector <T> VT;
typedef vector < VT > VVT;
T GaussJordan(VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
 VI irow(n), icol(n), ipiv(n);
 T det = 1:
  for (int i = 0; i < n; i++) {</pre>
   int pj = -1, pk = -1;
   for (int j = 0; j < n; j++) if (!ipiv[j])
     for (int k = 0; k < n; k++) if (!ipiv[k])
       if (pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
   if (fabs(a[pj][pk]) < EPS) { return 0; }</pre>
   ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
   icol[i] = pk;
   T c = 1.0 / a[pk][pk];
   det *= a[pk][pk];
   a[pk][pk] = 1.0;
   for (int p = 0; p < n; p++) a[pk][p] *= c;
   for (int p = 0; p < m; p++) b[pk][p] *= c;
   for (int p = 0; p < n; p++) if (p != pk) {
     c = a[p][pk];
      a[p][pk] = 0;
     for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
     for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
 }
```

```
for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
}
return det;</pre>
```

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### Karatsuba multiplication

```
class Karatsuba {
 typedef typename vector<T>::iterator vTi;
 void convolve_naive(vTi a, vTi b, vTi c, int n) {
   int n2 = n * 2;
   for (int i = 0; i < n2; ++i)
     c[i] = 0:
   for (int i = 0; i < n; ++i)</pre>
     for (int j = 0; j < n; ++ j)
        c[i + j] += a[i] * b[j];
 }
     | al * bl | ah * bh | as * bs |
   * ^x0 ^xh ^x1
                         ^x2
 void karatsuba(vTi a, vTi b, vTi c, int n) {
   if(n <= cut) {
      convolve_naive(a, b, c, n);
     return;
   int nh = n / 2;
   vTi al = a, ah = a + nh, as = c + nh * 10;
   vTi bl = b, bh = b + nh, bs = c + nh * 11;
   vTi x0 = c, x1 = c + n, x2 = c + n * 2, xh = c + nh;
   for (int i = 0; i < nh; ++i) {</pre>
     as[i] = al[i] + ah[i];
      bs[i] = bl[i] + bh[i];
   karatsuba(al, bl, x0, nh);
   karatsuba(ah, bh, x1, nh):
   karatsuba(as, bs, x2, nh);
```

```
for (int i = 0; i < n; ++i) x2[i] -= x0[i] + x1[i];
  for (int i = 0; i < n; ++i) xh[i] += x2[i];
}
public:
  Karatsuba(int _cut = 1 << 5) : cut(_cut) {}
  vector<T> convolve(vector<T> &_a, vector<T> &_b) {
    vector<T> a = _a, b = _b, c;
    int sz = max(a.size(), b.size()), sz2;
    for (sz2 = 1; sz2 < sz; sz2 *= 2);
    a.resize(sz2); b.resize(sz2); c.resize(sz2 * 6);
    karatsuba(a.begin(), b.begin(), c.begin(), sz2);
    c.resize(_a.size() + _b.size() - 1);
    return c;
}
};</pre>
```

## Miller-Rabin (probabilistic primality testing)

```
// modulo(a,b,c) = (a^b) % c
// \text{ mulmod(a,b,c)} = (a*b) \% c
bool Miller(long long p,int iteration){
    if(p<2){
        return false;
    if(p!=2 && p%2==0){
        return false;
    long long s=p-1;
    while (s\%2==0) {
        s/=2;
    }
    for(int i=0;i<iteration;i++){</pre>
        long long a=rand()%(p-1)+1,temp=s;
        long long mod=modulo(a,temp,p);
        while(temp!=p-1 && mod!=1 && mod!=p-1){
            mod=mulmod(mod,mod,p);
            temp *= 2;
        }
        if (mod!=p-1 && temp%2==0) {
            return false;
        }
    return true;
```

}

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