



# ***Computer Organization & Assembly Languages***

## ***Data Transfers, Addressing & Arithmetic***

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Adapted from the slides prepared by Kip Irvine for the book,  
Assembly Language for Intel-Based Computers, 5th Ed.



# Chapter Overview

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- **Data Transfer Instructions**
  - Addition and Subtraction
  - Data-Related Operators and Directives
  - Indirect Addressing
  - JMP and LOOP Instructions



# Data Transfer Instructions

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- Operand Types
- Instruction Operand Notation
- Direct Memory Operands
- MOV Instruction
- Zero & Sign Extension
- XCHG Instruction
- Direct-Offset Instructions



# Operand Types

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- Three basic types of operands:
  - Immediate – a constant integer (8, 16, or 32 bits)
    - value is encoded within the instruction
  - Register – the name of a register
    - register name is converted to a number and encoded within the instruction
  - Memory – reference to a location in memory
    - memory address is encoded within the instruction, or a register holds the address of a memory location



# Instruction Operand Notation

Operand	Description
<i>r8</i>	8-bit general-purpose register: AH, AL, BH, BL, CH, CL, DH, DL
<i>r16</i>	16-bit general-purpose register: AX, BX, CX, DX, SI, DI, SP, BP
<i>r32</i>	32-bit general-purpose register: EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP
<i>reg</i>	any general-purpose register
<i>sreg</i>	16-bit segment register: CS, DS, SS, ES, FS, GS
<i>imm</i>	8-, 16-, or 32-bit immediate value
<i>imm8</i>	8-bit immediate byte value
<i>imm16</i>	16-bit immediate word value
<i>imm32</i>	32-bit immediate doubleword value
<i>r/m8</i>	8-bit operand which can be an 8-bit general register or memory byte
<i>r/m16</i>	16-bit operand which can be a 16-bit general register or memory word
<i>r/m32</i>	32-bit operand which can be a 32-bit general register or memory doubleword
<i>mem</i>	an 8-, 16-, or 32-bit memory operand



# Direct Memory Operands

- A direct memory operand is a named reference to storage in memory
- The named reference (label) is automatically dereferenced by the assembler

```
.data
var1 BYTE 10h
.code
mov al,var1           ; AL = 10h
mov al,[var1]         ; AL = 10h
```

alternate format





# MOV Instruction

- Move from source to destination. Syntax:

*MOV destination, source*

- Source and destination have the same size
- No more than one memory operand permitted
- CS, EIP, and IP cannot be the destination
- No immediate to segment moves

```
.data
count BYTE 100
wVal  WORD 2
.code
    mov bl,count
    mov ax,wVal
    mov count,al

    mov al,wVal           ; error
    mov ax,count          ; error
    mov eax,count         ; error
```



## Your Turn . . .

---

Explain why each of the following MOV statements are invalid:

```
.data
bVal  BYTE    100
bVal2 BYTE     ?
wVal  WORD     2
dVal  DWORD    5
.code
    mov ds,45      immediate move to DS not permitted
    mov esi,wVal   size mismatch
    mov eip,dVal   EIP cannot be the destination
    mov 25,bVal    immediate value cannot be destination
    mov bVal2,bVal memory-to-memory move not permitted
```





# Memory to Memory

---

`.data`

`var1 WORD ?`

`var2 WORD ?`

`.code`

`mov ax, var1`

`mov var2, ax`



# Copy Smaller to Larger

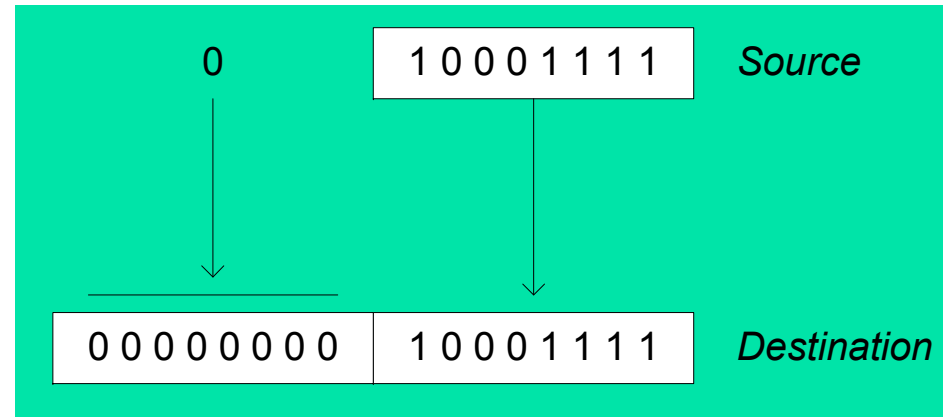
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```
.data
count WORD 1
.code
mov ecx, 0
mov cx, count
```

```
.data
signedVal SWORD -16 ; FFF0h
.code
mov ecx, 0 ; mov ecx, 0FFFFFFFFh
mov cx, signedVal
```

# Zero Extension

When you copy a smaller value into a larger destination, the MOVZX instruction fills (extends) the upper half of the destination with zeros.

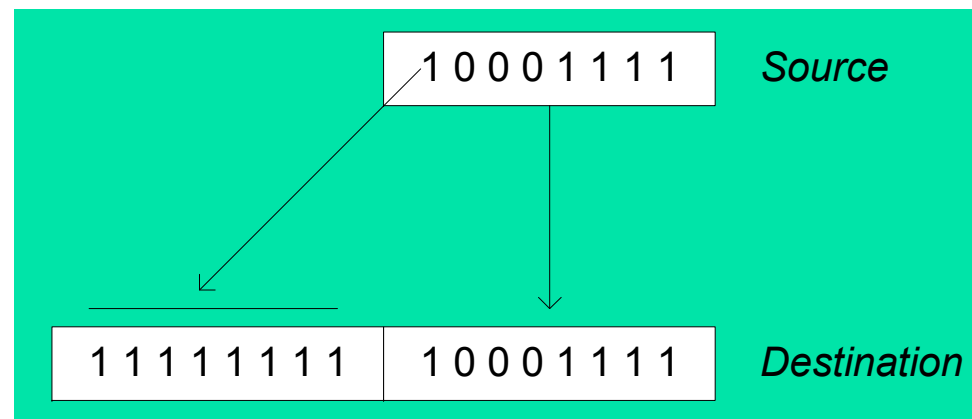


```
mov bl,10001111b  
movzx ax,bl ; zero-extension
```

The destination must be a register.

# Sign Extension

The MOVSX instruction fills the upper half of the destination with a copy of the source operand's sign bit.



```
mov bl,10001111b  
movsx ax,bl           ; sign extension
```

The destination must be a register.



## MOVZX & MOVSX

---

From a smaller location to a larger one

```
mov bx, 0A69Bh
```

```
movzx eax, bx           ; EAX=0000A69Bh
```

```
movzx edx, bl           ; EDX=0000009Bh
```

```
movzx cx, bl            ; EAX=009Bh
```

```
mov bx, 0A69Bh
```

```
movsx eax, bx           ; EAX=FFFA69Bh
```

```
movsx edx, bl           ; EDX=FFFFFF9Bh
```

```
movsx cx, bl            ; EAX=FF9Bh
```



# LAHF & SAHF

---

`.data`

`saveflags BYTE ?`

`.code`

`lahf`

`mov saveflags, ah`

`...`

`mov ah, saveflags`

`sahf`



# XCHG Instruction

---

XCHG exchanges the values of two operands. At least one operand must be a register. No immediate operands are permitted.

```
.data
var1 WORD 1000h
var2 WORD 2000h
.code
xchg ax,bx           ; exchange 16-bit regs
xchg ah,al           ; exchange 8-bit regs
xchg var1,bx         ; exchange mem, reg
xchg eax,ebx         ; exchange 32-bit regs

xchg var1,var2       ; error: two memory operands
```



# Direct-Offset Operands

A constant offset is added to a data label to produce an effective address (EA). The address is dereferenced to get the value inside its memory location.

```
.data
arrayB BYTE 10h,20h,30h,40h
.code
mov al,arrayB+1           ; AL = 20h
mov al,[arrayB+1]         ; alternative notation
```

Q: Why doesn't `arrayB+1` produce 11h?





## Direct-Offset Operands (cont.)

A constant offset is added to a data label to produce an effective address (EA). The address is dereferenced to get the value inside its memory location.

```
.data
arrayW  WORD 1000h,2000h,3000h
arrayD  DWORD 1,2,3,4
.code
mov ax,[arrayW+2]           ; AX = 2000h
mov ax,[arrayW+4]           ; AX = 3000h
mov eax,[arrayD+4]          ; EAX = 00000002h
```

```
; Will the following statements assemble?
mov ax,[arrayW-2]           ; ??
mov eax,[arrayD+16]         ; ??
```

What will happen when they run?



## Your Turn. . .

---

Write a program that rearranges the values of three doubleword values in the following array as: 3, 1, 2.

```
.data  
arrayD DWORD 1,2,3
```

- Step1: copy the first value into EAX and exchange it with the value in the second position.

```
mov  eax,arrayD  
xchg eax,[arrayD+4]
```

- Step 2: Exchange EAX with the third array value and copy the value in EAX to the first array position.

```
xchg eax,[arrayD+8]  
mov  arrayD,eax
```



# Evaluate This . . .

---

- We want to write a program that adds the following three bytes:

```
.data  
myBytes BYTE 80h,66h,0A5h
```

- What is your evaluation of the following code?

```
mov al,myBytes  
add al,[myBytes+1]  
add al,[myBytes+2]
```

- What is your evaluation of the following code?

```
mov ax,myBytes  
add ax,[myBytes+1]  
add ax,[myBytes+2]
```

- Any other possibilities?



## Evaluate This . . . (cont)

---

```
.data  
myBytes BYTE 80h,66h,0A5h
```

- How about the following code. Is anything missing?

```
movzx ax,myBytes  
mov    bl,[myBytes+1]  
add    ax,bx  
mov    bl,[myBytes+2]  
add    ax,bx                ; AX = sum
```

Yes: Move zero to BX before the MOVZX instruction.



# What's Next

---

- Data Transfer Instructions
- **Addition and Subtraction**
- Data-Related Operators and Directives
- Indirect Addressing
- JMP and LOOP Instructions



# Addition and Subtraction

---

- INC and DEC Instructions
- ADD and SUB Instructions
- NEG Instruction
- Implementing Arithmetic Expressions
- Flags Affected by Arithmetic
  - Zero
  - Sign
  - Carry
  - Overflow



# INC and DEC Instructions

---

- Add 1, subtract 1 from destination operand
  - operand may be register or memory
- **INC *destination***
  - Logic:  $destination \leftarrow destination + 1$
- **DEC *destination***
  - Logic:  $destination \leftarrow destination - 1$



# INC and DEC Examples

---

```
.data
myWord  WORD 1000h
myDword DWORD 10000000h
.code

    inc myWord           ; 1001h
    dec myWord           ; 1000h
    inc myDword          ; 10000001h

    mov ax,00FFh
    inc ax                ; AX = 0100h
    mov ax,00FFh
    inc al                ; AX = 0000h
```





## Your Turn...

---

Show the value of the destination operand after each of the following instructions executes:

```
.data
myByte BYTE 0FFh, 0
.code
    mov al,myByte           ; AL = FFh
    mov ah,[myByte+1]       ; AH = 00h
    dec ah                  ; AH = FFh
    inc al                  ; AL = 00h
    dec ax                  ; AX = FEFF
```



# ADD and SUB Instructions

---

- ADD destination, source
  - Logic:  $destination \leftarrow destination + source$
- SUB destination, source
  - Logic:  $destination \leftarrow destination - source$
- Same operand rules as for the MOV instruction



# ADD and SUB Examples

---

```
.data
var1 DWORD 10000h
var2 DWORD 20000h
.code
mov eax,var1           ; ---EAX---
add eax,var2           ; 00010000h
add ax,0FFFFh          ; 00030000h
add eax,1              ; 0003FFFFh
add eax,1              ; 00040000h
sub ax,1               ; 0004FFFFh
```



# NEG (negate) Instruction

---

Reverses the sign of an operand. Operand can be a register or memory operand.

```
.data
valB BYTE -1
valW SWORD +32767
.code
    mov al,valB           ; AL = -1
    neg al                ; AL = +1
    neg valW              ; valW = -32767
```

Suppose AX contains  $-32,768$  and we apply NEG to it. Will the result be valid?



# Implementing Arithmetic Expressions

HLL compilers translate mathematical expressions into assembly language. You can do it also. For example:

$$\text{Rval} = -\text{Xval} + (\text{Yval} - \text{Zval})$$

```
Rval SDWORD ?
Xval SDWORD 26
Yval SDWORD 30
Zval SDWORD 40
.code
    mov eax,Xval
    neg eax                ; EAX = -26
    mov ebx,Yval
    sub ebx,Zval           ; EBX = -10
    add eax,ebx
    mov Rval,eax           ; -36
```



## Your Turn...

---

Translate the following expression into assembly language.  
Do not permit Xval, Yval, or Zval to be modified:

$$Rval = Xval - (-Yval + Zval)$$

Assume that all values are signed doublewords.

```
mov ebx,Yval
neg ebx
add ebx,Zval
mov eax,Xval
sub eax,ebx
mov Rval,eax
```

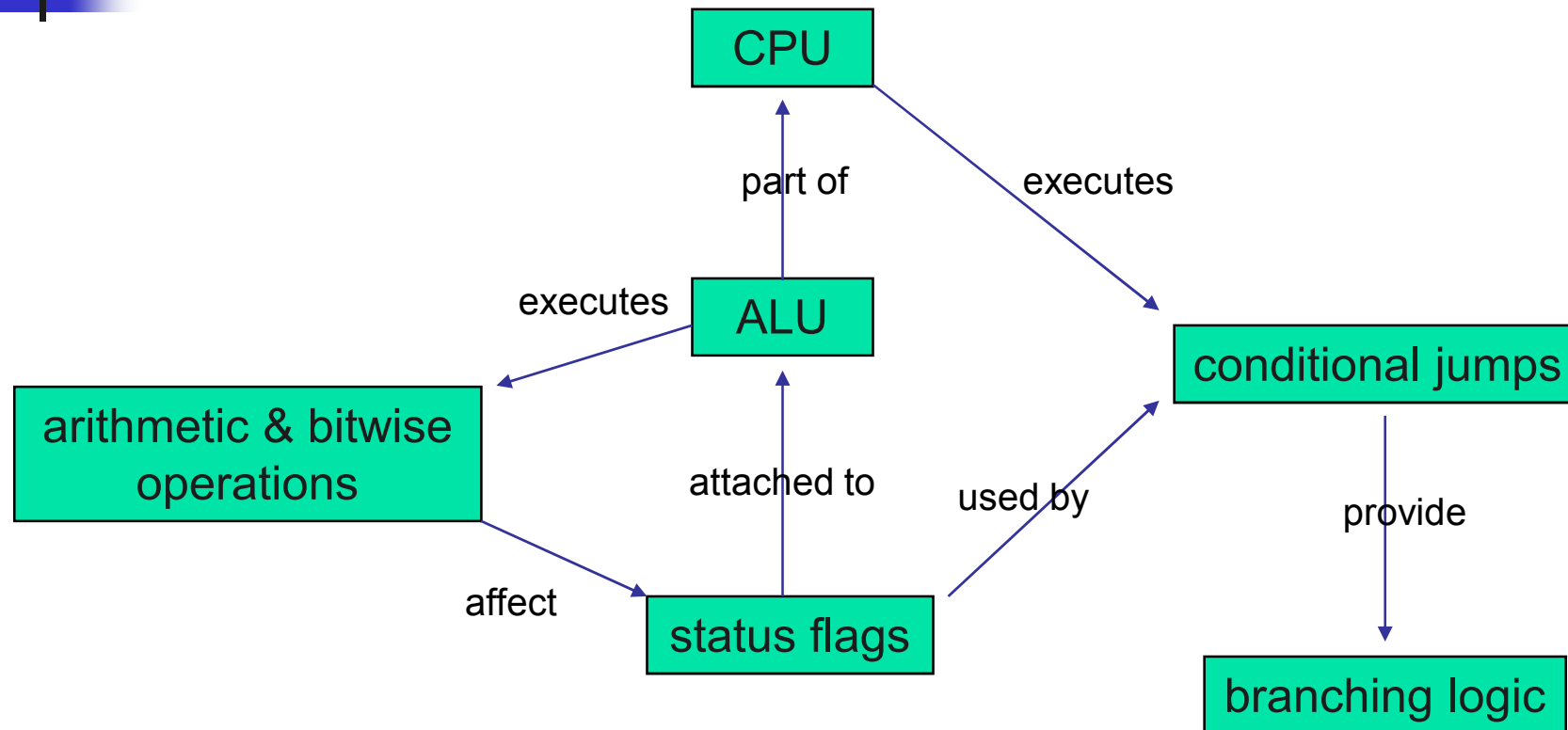


# Flags Affected by Arithmetic

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- The ALU has a number of status flags that reflect the outcome of arithmetic (and bitwise) operations
  - based on the contents of the destination operand
- Essential flags:
  - Zero flag – set when destination equals zero
  - Sign flag – set when destination is negative
  - Carry flag – set when unsigned value is out of range
  - Overflow flag – set when signed value is out of range
- The MOV instruction never affects the flags.

# Concept Map



You can use diagrams such as these to express the relationships between assembly language concepts.





# Zero Flag (ZF)

The Zero flag is set when the result of an operation produces zero in the destination operand.

```
mov cx,1
sub cx,1           ; CX = 0, ZF = 1
mov ax,0FFFFh
inc ax             ; AX = 0, ZF = 1
inc ax             ; AX = 1, ZF = 0
```

Remember...

- A flag is **set** when it equals 1.
- A flag is **clear** when it equals 0.



## Sign Flag (SF)

The Sign flag is set when the destination operand is negative. The flag is clear when the destination is positive.

```
mov cx,0
sub cx,1           ; CX = -1, SF = 1
add cx,2           ; CX = 1, SF = 0
```

The sign flag is a copy of the destination's highest bit:

```
mov al,0
sub al,1           ; AL = 11111111b, SF = 1
add al,2           ; AL = 00000001b, SF = 0
```



# Signed and Unsigned Integers

## A Hardware Viewpoint

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- All CPU instructions operate exactly the same on signed and unsigned integers
- The CPU cannot distinguish between signed and unsigned integers
- YOU, the programmer, are solely responsible for using the correct data type with each instruction



# Overflow and Carry Flags

## A Hardware Viewpoint

- How the **ADD** instruction modifies OF and CF:
  - $OF = (\text{carry out of the MSB}) \text{ XOR } (\text{carry into the MSB})$
  - $CF = (\text{carry out of the MSB})$
- How the **SUB** instruction modifies OF and CF:
  - NEG the source and ADD it to the destination
  - $OF = (\text{carry out of the MSB}) \text{ XOR } (\text{carry into the MSB})$
  - $CF = \text{INVERT} (\text{carry out of the MSB})$

MSB = Most Significant Bit (high-order bit)  
XOR = eXclusive-OR operation  
NEG = Negate (same as SUB 0,operand )



# Carry Flag (CF)

---

The Carry flag is set when the result of an operation generates an **unsigned** value that is out of range (too big or too small for the destination operand).

```
mov al,0FFh
add al,1                ; CF = 1, AL = 00

; Try to go below zero:

mov al,0
sub al,1                ; CF = 1, AL = FF
```

In the second example, we tried to generate a negative value. Unsigned values cannot be negative, so the Carry flag signaled an error condition.



## Your Turn . . .

---

For each of the following marked entries, show the values of the destination operand and the Sign, Zero, and Carry flags:

<code>mov ax, 00FFh</code>	
<code>add ax, 1</code>	; AX= 0100h    SF= 0   ZF= 0   CF= 0
<code>sub ax, 1</code>	; AX= 00FFh    SF= 0   ZF= 0   CF= 0
<code>add al, 1</code>	; AL= 00h       SF= 0   ZF= 1   CF= 1
<code>mov bh, 6Ch</code>	
<code>add bh, 95h</code>	; BH= 01h       SF= 0   ZF= 0   CF= 1
<code>mov al, 2</code>	
<code>sub al, 3</code>	; AL= FFh       SF= 1   ZF= 0   CF= 1



# Overflow Flag (OF)

The Overflow flag is set when the signed result of an operation is invalid or out of range.

```
; Example 1
mov al,+127
add al,1                ; OF = 1,    AL = ??

; Example 2
mov al,7Fh
add al,1                ; OF = 1,    AL = 80h
```

The two examples are identical at the binary level because 7Fh equals +127. To determine the value of the destination operand, it is often easier to calculate in hexadecimal.



# A Rule of Thumb

---

- When adding two integers, remember that the Overflow flag is only set when . . .
  - Two positive operands are added and their sum is negative
  - Two negative operands are added and their sum is positive

What will be the values of the Overflow flag?

```
mov al,80h
add al,92h                ; OF = 1
```

```
mov al,-2
add al,+127              ; OF = 0
```





## Your Turn . . .

---

What will be the values of the given flags after each operation?

```
mov al, -128
neg al                ; CF = 0    OF = 1

mov ax, 8000h
add ax, 2             ; CF = 0    OF = 0

mov ax, 0
sub ax, 2             ; CF = 1    OF = 0

mov al, -5
sub al, +125          ; CF = 0    OF = 1
```



# What's Next

---

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- **Data-Related Operators and Directives**
- Indirect Addressing
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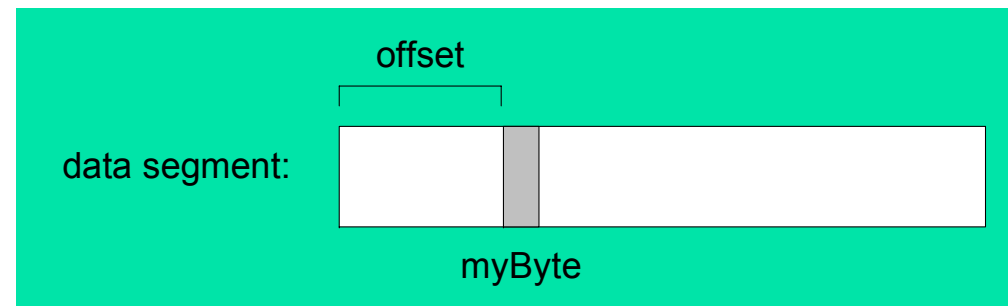
# Data-Related Operators and Directives

---

- OFFSET Operator
- PTR Operator
- TYPE Operator
- LENGTHOF Operator
- SIZEOF Operator
- LABEL Directive

# OFFSET Operator

- OFFSET returns the distance in bytes, of a label from the beginning of its enclosing segment
  - Protected mode: 32 bits
  - Real mode: 16 bits



The Protected-mode programs we write only have a single segment (we use the [flat memory model](#)).



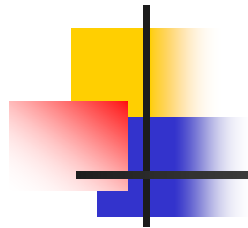
# OFFSET Examples

---

Let's assume that the data segment begins at 00404000h:

```
.data
bVal BYTE ?
wVal WORD ?
dVal DWORD ?
dVal2 DWORD ?

.code
mov esi,OFFSET bVal      ; ESI = 00404000
mov esi,OFFSET wVal      ; ESI = 00404001
mov esi,OFFSET dVal      ; ESI = 00404003
mov esi,OFFSET dVal2     ; ESI = 00404007
```



## Relating to C/C++

The value returned by OFFSET is a pointer. Compare the following code written for both C++ and assembly language:

```
; C++ version:  
char array[1000];  
char * p = array;
```

```
.data  
array BYTE 1000 DUP(?)  
.code  
mov esi,OFFSET array          ; ESI is p
```



# PTR Operator

Overrides the default type of a label (variable).  
Provides the flexibility to access part of a variable.

```
.data
myDouble DWORD 12345678h
.code
mov ax,myDouble           ; error - why?

mov ax,WORD PTR myDouble  ; loads 5678h

mov WORD PTR myDouble,4321h ; saves 4321h
```

Recall that **little endian** order is used when storing data in memory (see Section 3.4.9).



# Little Endian Order

- Little endian order refers to the way Intel stores integers in memory.
- Multi-byte integers are stored in reverse order, with the least significant byte stored at the lowest address
- For example, the doubleword 12345678h would be stored as:

byte	offset
78	0000
56	0001
34	0002
12	0003

When integers are loaded from memory into registers, the bytes are automatically re-reversed into their correct positions.





# PTR Operator Examples

```
.data  
myDouble DWORD 12345678h
```

doubleword	word	byte	offset	
12345678	5678	78	0000	myDouble
		56	0001	myDouble + 1
	1234	34	0002	myDouble + 2
		12	0003	myDouble + 3

```
mov al,BYTE PTR myDouble           ; AL = 78h  
mov al,BYTE PTR [myDouble+1]       ; AL = 56h  
mov al,BYTE PTR [myDouble+2]       ; AL = 34h  
mov ax,WORD PTR myDouble           ; AX = 5678h  
mov ax,WORD PTR [myDouble+2]       ; AX = 1234h
```



## PTR Operator (cont.)

---

PTR can also be used to combine elements of a smaller data type and move them into a larger operand. The CPU will automatically reverse the bytes.

```
.data
myBytes BYTE 12h,34h,56h,78h

.code
mov ax,WORD PTR [myBytes]           ; AX = 3412h
mov ax,WORD PTR [myBytes+2]         ; AX = 7856h
mov eax,DWORD PTR myBytes           ; EAX = 78563412h
```



## Your Turn . . .

---

Write down the value of each destination operand:

```
.data
varB BYTE 65h,31h,02h,05h
varW WORD 6543h,1202h
varD DWORD 12345678h

.code
mov ax,WORD PTR [varB+2]           ; a. 0502h
mov bl,BYTE PTR varD              ; b. 78h
mov bl,BYTE PTR [varW+2]          ; c. 02h
mov ax,WORD PTR [varD+2]          ; d. 1234h
mov eax,DWORD PTR varW            ; e. 12026543h
```



# TYPE Operator

---

The TYPE operator returns the size, in bytes, of a single element of a data declaration.

```
.data
var1 BYTE ?
var2 WORD ?
var3 DWORD ?
var4 QWORD ?

.code
mov eax,TYPE var1      ; 1
mov eax,TYPE var2      ; 2
mov eax,TYPE var3      ; 4
mov eax,TYPE var4      ; 8
```



# LENGTHOF Operator

---

The LENGTHOF operator counts the number of elements in a single data declaration.

	LENGTHOF
<code>.data</code>	
<code>byte1 BYTE 10,20,30</code>	<code>; 3</code>
<code>array1 WORD 30 DUP(?),0,0</code>	<code>; 32</code>
<code>array2 WORD 5 DUP(3 DUP(?))</code>	<code>; 15</code>
<code>array3 DWORD 1,2,3,4</code>	<code>; 4</code>
<code>digitStr BYTE "12345678",0</code>	<code>; 9</code>
<code>.code</code>	
<code>mov ecx,LENGTHOF array1</code>	<code>; 32</code>



# SIZEOF Operator

The SIZEOF operator returns a value that is equivalent to multiplying LENGTHOF by TYPE.

	SIZEOF
<code>.data</code>	
<code>byte1 BYTE 10,20,30</code>	<code>; 3</code>
<code>array1 WORD 30 DUP(?),0,0</code>	<code>; 64</code>
<code>array2 WORD 5 DUP(3 DUP(?))</code>	<code>; 30</code>
<code>array3 DWORD 1,2,3,4</code>	<code>; 16</code>
<code>digitStr BYTE "12345678",0</code>	<code>; 9</code>
<code>.code</code>	
<code>mov ecx,SIZEOF array1</code>	<code>; 64</code>



# Spanning Multiple Lines

---

A data declaration spans multiple lines if each line (except the last) ends with a comma. The `LENGTHOF` and `SIZEOF` operators include all lines belonging to the declaration:

```
.data
array WORD 10,20,
        30,40,
        50,60

.code
mov  eax,LENGTHOF array      ; 6
mov  ebx,SIZEOF array        ; 12
```



## Spanning Multiple Lines (cont.)

---

In the following example, `array` identifies only the first `WORD` declaration. Compare the values returned by `LENGTHOF` and `SIZEOF` here to those in the previous slide:

```
.data
array  WORD 10,20
        WORD 30,40
        WORD 50,60

.code
mov  eax,LENGTHOF array           ; 2
mov  ebx,SIZEOF array             ; 4
```





# LABEL Directive

---

- Assigns an alternate label name and type to an existing storage location
- LABEL does not allocate any storage of its own
- Removes the need for the PTR operator

```
.data
dwList    LABEL DWORD
wordList  LABEL WORD
intList   BYTE 00h,10h,00h,20h
.code
mov  eax,dwList           ; 20001000h
mov  cx,wordList          ; 1000h
mov  dl,intList            ; 00h
```



# What's Next

---

- Data Transfer Instructions
- Addition and Subtraction
- Data-Related Operators and Directives
- **Indirect Addressing**
- JMP and LOOP Instructions



# Indirect Addressing

---

- Indirect Operands
- Array Sum Example
- Indexed Operands
- Pointers



# Indirect Operands

---

An indirect operand holds the address of a variable, usually an array or string. It can be **dereferenced** (just like a pointer).

```
.data
val1 BYTE 10h,20h,30h
.code
mov esi,OFFSET val1
mov al,[esi]                ; dereference ESI (AL = 10h)

inc esi
mov al,[esi]                ; AL = 20h

inc esi
mov al,[esi]                ; AL = 30h
```



## Indirect Operands (cont.)

Use PTR to clarify the size attribute of a memory operand.

```
.data
myCount WORD 0

.code
mov esi,OFFSET myCount
inc [esi]                ; error: ambiguous
inc WORD PTR [esi]      ; ok
```

Should PTR be used here?

```
add [esi],20
```

yes, because [esi] could point to a byte, word, or doubleword



# Array Sum Example

Indirect operands are ideal for traversing an array. Note that the register in brackets must be incremented by a value that matches the array type.

```
.data
arrayW WORD 1000h,2000h,3000h
.code
    mov esi,OFFSET arrayW
    mov ax,[esi]
    add esi,2                ; or: add esi,TYPE arrayW
    add ax,[esi]
    add esi,2
    add ax,[esi]             ; AX = sum of the array
```

ToDo: Modify this example for an array of doublewords.



# Indexed Operands

An indexed operand adds a constant to a register to generate an effective address. There are two notational forms: `[label + reg]`      `label[reg]`

```
.data
arrayW WORD 1000h,2000h,3000h
.code
    mov esi,0
    mov ax,[arrayW + esi]           ; AX = 1000h
    mov ax,arrayW[esi]             ; alternate format
    add esi,2
    add ax,[arrayW + esi]
    etc.
```

ToDo: Modify this example for an array of doublewords.



# Index Scaling

You can scale an indirect or indexed operand to the offset of an array element. This is done by multiplying the index by the array's TYPE:

```
.data
arrayB BYTE 0,1,2,3,4,5
arrayW WORD 0,1,2,3,4,5
arrayD DWORD 0,1,2,3,4,5

.code
mov esi,4
mov al,arrayB[esi*TYPE arrayB]    ; 04
mov bx,arrayW[esi*TYPE arrayW]    ; 0004
mov edx,arrayD[esi*TYPE arrayD]   ; 00000004
```





# Pointers

---

You can declare a **pointer variable** that contains the offset of another variable.

```
.data
arrayW WORD 1000h,2000h,3000h
ptrW DWORD arrayW
.code
    mov esi,ptrW
    mov ax,[esi]           ; AX = 1000h
```

Alternate format:

```
ptrW DWORD OFFSET arrayW
```



# What's Next

---

- Data Transfer Instructions
- Addition and Subtraction
- Data-Related Operators and Directives
- Indirect Addressing
- **JMP and LOOP Instructions**



# JMP and LOOP Instructions

---

- JMP Instruction
- LOOP Instruction
- LOOP Example
- Summing an Integer Array
- Copying a String



# JMP Instruction

---

- JMP is an unconditional jump to a label that is usually within the same procedure.
- Syntax: *JMP target*
- Logic:  $EIP \leftarrow target$
- Example:

```
top:
    .
    .
    jmp top
```

A jump outside the current procedure must be to a special type of label called a [global label](#) (see Section 5.5.2 for details).



# LOOP Instruction

---

- The LOOP instruction creates a counting loop
- Syntax: **LOOP** *target*
- Logic:
  - $ECX \leftarrow ECX - 1$
  - if  $ECX \neq 0$ , jump to *target*
- Implementation:
  - The assembler calculates the distance, in bytes, between the offset of the following instruction and the offset of the target label. It is called the **relative offset**.
  - The relative offset is added to EIP.



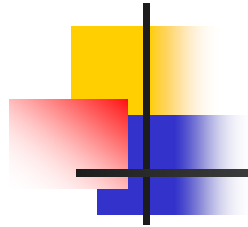
# LOOP Example

The following loop calculates the sum of the integers  
5 + 4 + 3 + 2 + 1:

offset	machine code	source code
00000000	66 B8 0000	mov ax, 0
00000004	B9 00000005	mov ecx, 5
00000009	66 03 C1	L1: add ax, cx
0000000C	E2 FB	loop L1
0000000E		

When LOOP is assembled, the current location = 0000000E (offset of the next instruction). -5 (FBh) is added to the the current location, causing a jump to location 00000009:

$$00000009 \leftarrow 0000000E + FB$$



## Your Turn . . .

---

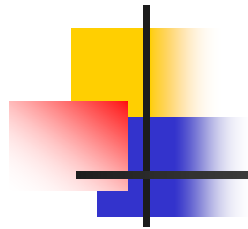
If the relative offset is encoded in a single signed byte,

(a) what is the largest possible backward jump?

(b) what is the largest possible forward jump?

(a) -128

(b) +127



## Your Turn . . .

---

What will be the final value of AX?

10

```
mov ax, 6  
mov ecx, 4  
L1:  
inc ax  
loop L1
```

How many times will the loop execute?

4,294,967,296

```
mov ecx, 0  
x2:  
inc ax  
loop x2
```





# Nested Loop

If you need to code a loop within a loop, you must save the outer loop counter's ECX value.

```
.data
count DWORD ?
.code
    mov ecx,100            ; set outer loop count
L1:
    mov count,ecx          ; save outer loop count
    mov ecx,20             ; set inner loop count
L2: .
    .
    loop L2                ; repeat the inner loop
    mov ecx,count          ; restore outer loop count
    loop L1                ; repeat the outer loop
```



# Summing an Integer Array

The following code calculates the sum of an array of 16-bit integers.

```
.data
intarray WORD 100h,200h,300h,400h
.code
    mov edi,OFFSET intarray      ; address of intarray
    mov ecx,LENGTHOF intarray   ; loop counter
    mov ax,0                     ; zero the accumulator
L1:
    add ax,[edi]                 ; add an integer
    add edi,TYPE intarray        ; point to next integer
    loop L1                     ; repeat until ECX = 0
```



## Your Turn . . .

---

What changes would you make to the program on the previous slide if you were summing a doubleword array?



# Copying a String

The following code copies a string from **source** to **target**:

```
.data
source  BYTE  "This is the source string",0
target  BYTE  sizeof source DUP(0)

.code
    mov     esi,0                ; index register
    mov     ecx, sizeof source ; loop counter
L1:
    mov     al, source[esi]      ; get char from source
    mov     target[esi], al     ; store it in the target
    inc     esi                 ; move to next character
    loop    L1                  ; repeat for entire string
```

good use of  
**sizeof**



## Your Turn . . .

---

Rewrite the program shown in the previous slide, using indirect addressing rather than indexed addressing.





## CMP Instruction (See 6.2.7)

---

- Compares the destination operand to the source operand (both are unsigned)
- Syntax: *CMP destination, source*
- Example: destination == source

```
mov al,5  
cmp al,5                ; Zero flag set
```

- Example: destination < source

```
mov al,4  
cmp al,5                ; Carry flag set
```



# CMP Instruction

---

- Example: destination > source

```
mov al,6  
cmp al,5                ; ZF = 0, CF = 0
```

(both the Zero and Carry flags are clear)





## *Jcond* Instruction

---

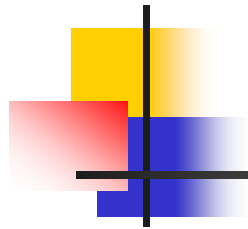
- A conditional jump instruction branches to a label when specific register or flag conditions are met
- Examples:
  - JB, JC jump to a label if the Carry flag is set
  - JE, JZ jump to a label if the Zero flag is set
  - JS jumps to a label if the Sign flag is set
  - JNE, JNZ jump to a label if the Zero flag is clear
  - JECXZ jumps to a label if ECX equals 0



# *Jcond* Ranges

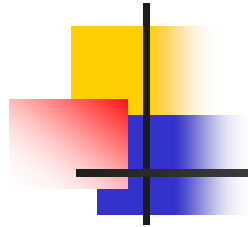
---

- Prior to the 386:
  - jump must be within  $-128$  to  $+127$  bytes from current location counter
- IA-32 processors:
  - 32-bit offset permits jump anywhere in memory



# Jumps Based on Specific Flags

Mnemonic	Description	Flags
JZ	Jump if zero	ZF = 1
JNZ	Jump if not zero	ZF = 0
JC	Jump if carry	CF = 1
JNC	Jump if not carry	CF = 0
JO	Jump if overflow	OF = 1
JNO	Jump if not overflow	OF = 0
JS	Jump if signed	SF = 1
JNS	Jump if not signed	SF = 0
JP	Jump if parity (even)	PF = 1
JNP	Jump if not parity (odd)	PF = 0



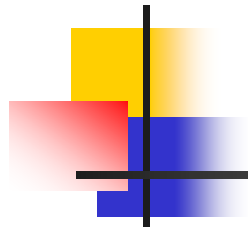
# Jumps Based on Equality

Mnemonic	Description
JE	Jump if equal ( <i>leftOp = rightOp</i> )
JNE	Jump if not equal ( <i>leftOp <math>\neq</math> rightOp</i> )
JCXZ	Jump if CX = 0
JECXZ	Jump if ECX = 0



# Jumps Based on Unsigned Comparisons

Mnemonic	Description
JA	Jump if above (if $leftOp > rightOp$ )
JNBE	Jump if not below or equal (same as JA)
JAE	Jump if above or equal (if $leftOp \geq rightOp$ )
JNB	Jump if not below (same as JAE)
JB	Jump if below (if $leftOp < rightOp$ )
JNAE	Jump if not above or equal (same as JB)
JBE	Jump if below or equal (if $leftOp \leq rightOp$ )
JNA	Jump if not above (same as JBE)



# Jumps Based on Signed Comparisons

Mnemonic	Description
JG	Jump if greater (if <i>leftOp</i> > <i>rightOp</i> )
JNLE	Jump if not less than or equal (same as JG)
JGE	Jump if greater than or equal (if <i>leftOp</i> >= <i>rightOp</i> )
JNL	Jump if not less (same as JGE)
JL	Jump if less (if <i>leftOp</i> < <i>rightOp</i> )
JNGE	Jump if not greater than or equal (same as JL)
JLE	Jump if less than or equal (if <i>leftOp</i> <= <i>rightOp</i> )
JNG	Jump if not greater (same as JLE)



# Block-Structured IF Statements

---

Assembly language programmers can easily translate logical statements written in C++/Java into assembly language. For example:

```
if( op1 == op2 )  
    X = 1;  
else  
    X = 2;
```

```
mov  eax,op1  
cmp  eax,op2  
jne  L1  
mov  X,1  
jmp  L2  
L1:  mov  X,2  
L2:
```



## Your Turn . . .

---

Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx )  
{  
    eax = 5;  
    edx = 6;  
}
```

```
cmp ebx,ecx  
ja  next  
mov eax,5  
mov edx,6  
next:
```

(There are multiple correct solutions to this problem.)





# WHILE Loops

---

A WHILE loop is really an IF statement followed by the body of the loop, followed by an unconditional jump to the top of the loop. Consider the following example:

```
while( eax < ebx)
    eax = eax + 1;
```

This is a possible implementation:

```
top: cmp  eax, ebx           ; check loop condition
     jae  next              ; false? exit loop
     inc  eax               ; body of loop
     jmp  top               ; repeat the loop
next:
```



## Your Turn . . .

---

Implement the following loop, using unsigned 32-bit integers:

```
while( ebx <= val1)
{
    ebx = ebx + 5;
    val1 = val1 - 1
}
```

```
top: cmp ebx, val1          ; check loop condition
     ja next               ; false? exit loop
     add ebx, 5             ; body of loop
     dec val1
     jmp top                ; repeat the loop
next:
```



## DIV Instruction (See 7.4.4)

- The DIV (unsigned divide) instruction performs 8-bit, 16-bit, and 32-bit division on unsigned integers
- A single operand is supplied (register or memory operand), which is assumed to be the divisor
- Instruction formats:

**DIV *r/m8***

**DIV *r/m16***

**DIV *r/m32***

Default Operands:

Dividend	Divisor	Quotient	Remainder
AX	<i>r/m8</i>	AL	AH
DX:AX	<i>r/m16</i>	AX	DX
EDX:EAX	<i>r/m32</i>	EAX	EDX



# DIV Examples

---

Divide 8003h by 100h, using 16-bit operands:

```
mov dx,0                ; clear dividend, high
mov ax,8003h            ; dividend, low
mov cx,100h             ; divisor
div cx                  ; AX = 0080h, DX = 3
```

Same division, using 32-bit operands:

```
mov edx,0               ; clear dividend, high
mov eax,8003h           ; dividend, low
mov ecx,100h            ; divisor
div ecx                 ; EAX = 00000080h, DX = 3
```



## Your Turn . . .

---

What will be the hexadecimal values of DX and AX after the following instructions execute?

```
mov dx,0087h  
mov ax,6000h  
mov bx,100h  
div bx
```

DX = 0000h, AX = 8760h

