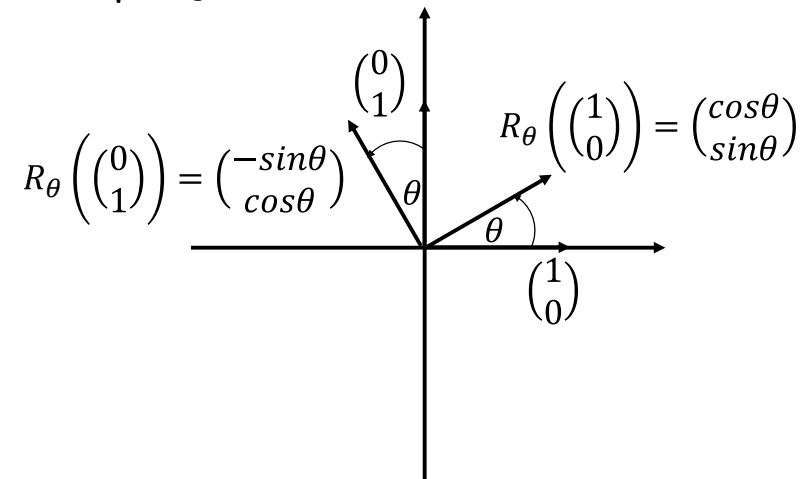


Matrix-Matrix Multiplication

Jirasak Sittigorn
Department of Computer Engineering
Faculty of Engineering
King Mongkut's Institute of Technology Ladkrabang

Opening Remarks

- Composing Rotations

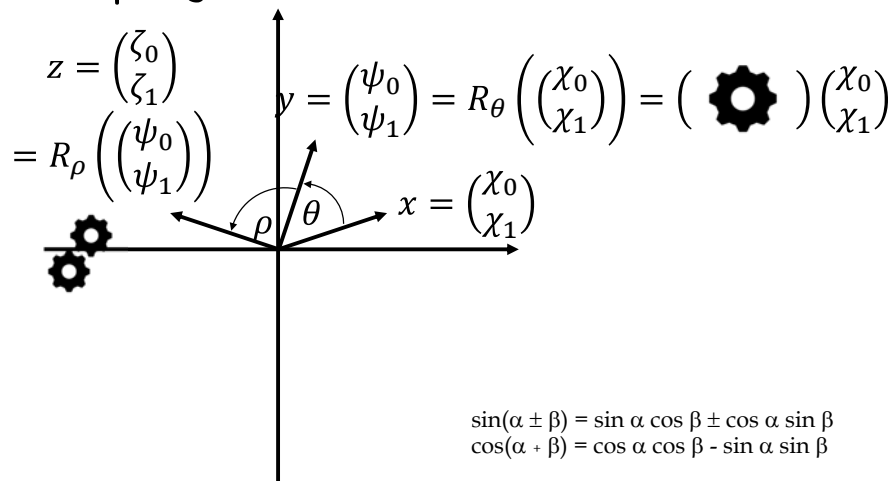


Robert van de Geijn and Maggie Myers. *Linear Algebra - Foundations to Frontiers*. <https://www.edx.org/>

2

Opening Remarks

- Composing Rotations



$$\sin(\alpha \pm \beta) = \sin \alpha \cos \beta \pm \cos \alpha \sin \beta$$

$$\cos(\alpha + \beta) = \cos \alpha \cos \beta - \sin \alpha \sin \beta$$

Robert van de Geijn and Maggie Myers. *Linear Algebra - Foundations to Frontiers*. <https://www.edx.org/>

3

Observations

- Partitioned Matrix-Matrix Multiplication
- Properties Transposing a Product of Matrices
- Matrix-Matrix Multiplication with Special Matrices

Robert van de Geijn and Maggie Myers. *Linear Algebra - Foundations to Frontiers*. <https://www.edx.org/>

4

Observations

- Partitioned Matrix-Matrix Multiplication

$$\begin{pmatrix} C_{00} & C_{01} \\ C_{10} & C_{11} \end{pmatrix} = \begin{pmatrix} A_{00} & A_{01} \\ A_{10} & A_{11} \end{pmatrix} \begin{pmatrix} B_{00} & B_{01} \\ B_{10} & B_{11} \end{pmatrix} = \begin{pmatrix} A_{00}B_{00} + A_{01}B_{10} & A_{00}B_{01} + A_{01}B_{11} \\ A_{10}B_{00} + A_{11}B_{10} & A_{10}B_{01} + A_{11}B_{11} \end{pmatrix}$$

$$\begin{pmatrix} \gamma_{0,0} & \gamma_{0,1} & \gamma_{0,2} & \gamma_{0,3} \\ \gamma_{1,0} & \gamma_{1,1} & \gamma_{1,2} & \gamma_{1,3} \\ \gamma_{2,0} & \gamma_{2,1} & \gamma_{2,2} & \gamma_{2,3} \\ \gamma_{3,0} & \gamma_{3,1} & \gamma_{3,2} & \gamma_{3,3} \end{pmatrix} = \begin{pmatrix} \alpha_{0,0} & \alpha_{0,1} & \alpha_{0,2} & \alpha_{0,3} \\ \alpha_{1,0} & \alpha_{1,1} & \alpha_{1,2} & \alpha_{1,3} \\ \alpha_{2,0} & \alpha_{2,1} & \alpha_{2,2} & \alpha_{2,3} \\ \alpha_{3,0} & \alpha_{3,1} & \alpha_{3,2} & \alpha_{3,3} \end{pmatrix} \begin{pmatrix} \beta_{0,0} & \beta_{0,1} & \beta_{0,2} & \beta_{0,3} \\ \beta_{1,0} & \beta_{1,1} & \beta_{1,2} & \beta_{1,3} \\ \beta_{2,0} & \beta_{2,1} & \beta_{2,2} & \beta_{2,3} \\ \beta_{3,0} & \beta_{3,1} & \beta_{3,2} & \beta_{3,3} \end{pmatrix} =$$

$$\begin{pmatrix} \alpha_{0,0} & \alpha_{0,1} & \alpha_{0,2} & \alpha_{0,3} \\ \alpha_{1,0} & \alpha_{1,1} & \alpha_{1,2} & \alpha_{1,3} \\ \alpha_{2,0} & \alpha_{2,1} & \alpha_{2,2} & \alpha_{2,3} \\ \alpha_{3,0} & \alpha_{3,1} & \alpha_{3,2} & \alpha_{3,3} \end{pmatrix} \begin{pmatrix} \beta_{0,0} & \beta_{0,1} \\ \beta_{1,0} & \beta_{1,1} \\ \beta_{2,0} & \beta_{2,1} \\ \beta_{3,0} & \beta_{3,1} \end{pmatrix} + \begin{pmatrix} \alpha_{0,2} & \alpha_{0,3} \\ \alpha_{1,2} & \alpha_{1,3} \\ \alpha_{2,2} & \alpha_{2,3} \\ \alpha_{3,2} & \alpha_{3,3} \end{pmatrix} \begin{pmatrix} \beta_{2,2} & \beta_{2,3} \\ \beta_{3,2} & \beta_{3,3} \end{pmatrix}$$

Observations

- Partitioned Matrix-Matrix Multiplication

Theorem 5.1 Let $C \in \mathbb{R}^{m \times n}$, $A \in \mathbb{R}^{m \times k}$, and $B \in \mathbb{R}^{k \times n}$. Let

- $m = m_0 + m_1 + \dots + m_{M-1}$, $m_i \geq 0$ for $i = 0, \dots, M-1$;

- $n = n_0 + n_1 + \dots + n_{N-1}$, $n_j \geq 0$ for $j = 0, \dots, N-1$; and

- $k = k_0 + k_1 + \dots + k_{K-1}$, $k_p \geq 0$ for $p = 0, \dots, K-1$.

Partition

$$C = \begin{pmatrix} C_{0,0} & C_{0,1} & \dots & C_{0,N-1} \\ C_{1,0} & C_{1,1} & \dots & C_{1,N-1} \\ \vdots & \vdots & \ddots & \vdots \\ C_{M-1,0} & C_{M-1,1} & \dots & C_{M-1,N-1} \end{pmatrix}, A = \begin{pmatrix} A_{0,0} & A_{0,1} & \dots & A_{0,K-1} \\ A_{1,0} & A_{1,1} & \dots & A_{1,K-1} \\ \vdots & \vdots & \ddots & \vdots \\ A_{M-1,0} & A_{M-1,1} & \dots & A_{M-1,K-1} \end{pmatrix},$$

$$\text{and } B = \begin{pmatrix} B_{0,0} & B_{0,1} & \dots & B_{0,N-1} \\ B_{1,0} & B_{1,1} & \dots & B_{1,N-1} \\ \vdots & \vdots & \ddots & \vdots \\ B_{K-1,0} & B_{K-1,1} & \dots & B_{K-1,N-1} \end{pmatrix},$$

with $C_{i,j} \in \mathbb{R}^{m_i \times n_j}$, $A_{i,p} \in \mathbb{R}^{m_i \times k_p}$, and $B_{p,j} \in \mathbb{R}^{k_p \times n_j}$. Then $C_{i,j} = \sum_{p=0}^{K-1} A_{i,p} B_{p,j}$.

Observations

- Partitioned Matrix-Matrix Multiplication

Example 5.2 Consider

$$A = \begin{pmatrix} -1 & 2 & 4 & 1 \\ 1 & 0 & -1 & -2 \\ 2 & -1 & 3 & 1 \\ 1 & 2 & 3 & 4 \end{pmatrix}, B = \begin{pmatrix} -2 & 2 & -3 \\ 0 & 1 & -1 \\ -2 & -1 & 0 \\ 4 & 0 & 1 \end{pmatrix}, \text{ and } AB = \begin{pmatrix} -2 & -4 & 2 \\ -8 & 3 & -5 \\ -6 & 0 & -4 \\ 8 & 1 & -1 \end{pmatrix}.$$

If

$$A_0 = \begin{pmatrix} -1 & 2 \\ 1 & 0 \\ 2 & -1 \\ 1 & 2 \end{pmatrix}, A_1 = \begin{pmatrix} 4 & 1 \\ -1 & -2 \\ 3 & 1 \\ 3 & 4 \end{pmatrix}, B_0 = \begin{pmatrix} -2 & 2 & -3 \\ 0 & 1 & -1 \end{pmatrix}, \text{ and } B_1 = \begin{pmatrix} -2 & -1 & 0 \\ 4 & 0 & 1 \end{pmatrix}.$$

Observations

- Partitioned Matrix-Matrix Multiplication

Then

$$AB = \begin{pmatrix} A_0 & A_1 \end{pmatrix} \begin{pmatrix} B_0 \\ B_1 \end{pmatrix} = A_0 B_0 + A_1 B_1.$$

$$\begin{pmatrix} -1 & 2 & 4 & 1 \\ 1 & 0 & -1 & -2 \\ 2 & -1 & 3 & 1 \\ 1 & 2 & 3 & 4 \end{pmatrix} \begin{pmatrix} -2 & 2 & -3 \\ 0 & 1 & -1 \\ -2 & -1 & 0 \\ 4 & 0 & 1 \end{pmatrix} = \begin{pmatrix} -1 & 2 \\ 1 & 0 \\ 2 & -1 \\ 1 & 2 \end{pmatrix} \begin{pmatrix} -2 & 2 & -3 \\ 0 & 1 & -1 \end{pmatrix} + \begin{pmatrix} 4 & 1 \\ -1 & -2 \\ 3 & 1 \\ 3 & 4 \end{pmatrix} \begin{pmatrix} -2 & -1 & 0 \\ 4 & 0 & 1 \end{pmatrix}$$

$$= \begin{pmatrix} -2 & 2 & -3 \\ 0 & 1 & -1 \\ -2 & -1 & 0 \\ 4 & 0 & 1 \end{pmatrix} + \begin{pmatrix} -8 & -4 & 2 \\ -2 & -1 & 0 \\ 8 & 1 & -1 \end{pmatrix}$$

$$= \begin{pmatrix} -2 & 0 & 1 \\ -2 & 2 & -3 \\ -4 & 3 & -5 \\ -2 & 4 & -5 \end{pmatrix} + \begin{pmatrix} -4 & -4 & 1 \\ -6 & 1 & -2 \\ -2 & -3 & 1 \\ 10 & -3 & 4 \end{pmatrix} = \begin{pmatrix} -2 & -4 & 2 \\ -8 & 3 & -5 \\ -6 & 0 & -4 \\ 8 & 1 & -1 \end{pmatrix}.$$

Observations

- Properties Transposing a Product of Matrices

Homework 5.2.3.1 Let $A = \begin{pmatrix} 2 & 0 & 1 \\ -1 & 1 & 0 \\ 1 & 3 & 1 \\ -1 & 1 & 1 \end{pmatrix}$ and $B = \begin{pmatrix} 2 & 1 & 2 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$. Compute

- $A^T A =$
- $AA^T =$
- $(AB)^T =$
- $A^T B^T =$
- $B^T A^T =$

Homework 5.2.3.2 Let $A \in \mathbb{R}^{m \times k}$ and $B \in \mathbb{R}^{k \times n}$. $(AB)^T = B^T A^T$.

Always/Sometimes/Never

Homework 5.2.3.3 Let A , B , and C be conformal matrices so that ABC is well-defined. Then $(ABC)^T = C^T B^T A^T$.

Always/Sometimes/Never

Observations

- Matrix-Matrix Multiplication with Special Matrices

Homework 5.2.4.9 Compute the following, using what you know about partitioned matrix-matrix multiplication:

$$\begin{pmatrix} 1 & -1 & -2 \\ 0 & 2 & 3 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} -2 & 1 & -1 \\ 0 & 1 & 2 \\ 0 & 0 & 1 \end{pmatrix} =$$

Homework 5.2.4.14 Evaluate

$$\begin{pmatrix} -1 \\ 1 \\ 2 \end{pmatrix} \begin{pmatrix} -1 & 1 & 2 \end{pmatrix} =$$

$$\begin{pmatrix} 2 \\ 0 \\ -1 \end{pmatrix} \begin{pmatrix} 2 & 0 & -1 \end{pmatrix} =$$

$$\begin{pmatrix} -1 & 2 \\ 1 & 0 \\ 2 & -1 \end{pmatrix} \begin{pmatrix} -1 & 1 & 2 \\ 2 & 0 & -1 \end{pmatrix} =$$

$$\begin{pmatrix} 1 \\ -2 \\ 2 \end{pmatrix} \begin{pmatrix} 1 & -2 & 2 \end{pmatrix} =$$

$$\begin{pmatrix} -1 & 2 & 1 \\ 1 & 0 & -2 \\ 2 & -1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 1 & 2 \\ 2 & 0 & -1 \\ 1 & -2 & 2 \end{pmatrix} =$$

Algorithms for Computing Matrix-Matrix Multiplication

- Lots of Loops
- Matrix-Matrix Multiplication by Columns
- Matrix-Matrix Multiplication by Rows
- Matrix-Matrix Multiplication with Rank-1 Updates

Algorithms for Computing Matrix-Matrix Multiplication

- Lots of Loops

$$C = AB \quad \gamma_{i,j} = \sum_{p=0}^{k-1} \alpha_{i,p} \beta_{p,j}$$

- Consider the MATLAB function

```
for i = 1 : r_A
    for j = 1 : c_B
        for p = 1 : c_A
            C(i, j) = A(i, p) * B(p, j) + C(i, j);
        end
    end
end
```

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Columns

In Theorem 5.1 let us partition C and B by columns and not partition A . In other words, let $M = 1, m_0 = m; N = n, n_j = 1, j = 0, \dots, n-1$; and $K = 1, k_0 = k$. Then

$$C = (c_0 \mid c_1 \mid \dots \mid c_{n-1}) \quad \text{and} \quad B = (b_0 \mid b_1 \mid \dots \mid b_{n-1})$$

so that

$$(c_0 \mid c_1 \mid \dots \mid c_{n-1}) = C = AB = A (b_0 \mid b_1 \mid \dots \mid b_{n-1}) = (Ab_0 \mid Ab_1 \mid \dots \mid Ab_{n-1}).$$

Homework 5.3.2.3

$$\cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 \\ 2 \\ 1 \end{pmatrix} =$$

$$\cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 \\ 2 & 1 \\ 1 & -1 \end{pmatrix} =$$

$$\cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \\ 2 & 1 & -1 \\ 1 & -1 & 2 \end{pmatrix} =$$



Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Columns

By moving the loop indexed by j to the outside in the algorithm for computing $C = AB + C$ we observe that

$$\left. \begin{array}{l} \text{for } j = 0, \dots, n-1 \\ \quad \text{for } i = 0, \dots, m-1 \\ \quad \quad \text{for } p = 0, \dots, k-1 \\ \quad \quad \quad \gamma_{ij} := \alpha_{ip} \beta_{pj} + \gamma_{ij} \\ \quad \text{endfor} \\ \quad \text{endfor} \\ \text{endfor} \end{array} \right\} c_j := Ab_j + c_j \quad \text{or} \quad \left. \begin{array}{l} \text{for } j = 0, \dots, n-1 \\ \quad \text{for } p = 0, \dots, k-1 \\ \quad \quad \text{for } i = 0, \dots, m-1 \\ \quad \quad \quad \gamma_{ij} := \alpha_{ip} \beta_{pj} + \gamma_{ij} \\ \quad \quad \text{endfor} \\ \quad \text{endfor} \\ \text{endfor} \end{array} \right\} c_j := Ab_j + c_j$$

Exchanging the order of the two inner-most loops merely means we are using a different algorithm (dot product vs. AXPY) for the matrix-vector multiplication $c_j := Ab_j + c_j$.

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Columns

Algorithm: $C := \text{GEMM_UNB_VAR1}(A, B, C)$

Partition $B \rightarrow (B_L \mid B_R), C \rightarrow (C_L \mid C_R)$
where B_L has 0 columns, C_L has 0 columns

while $n(B_L) < n(B)$ do

Repartition

$$(B_L \mid B_R) \rightarrow (B_0 \mid b_1 \mid B_2), (C_L \mid C_R) \rightarrow (C_0 \mid c_1 \mid C_2)$$

where b_1 has 1 column, c_1 has 1 column

$$c_1 := Ab_1 + c_1$$

Continue with

$$(B_L \mid B_R) \leftarrow (B_0 \mid b_1 \mid B_2), (C_L \mid C_R) \leftarrow (C_0 \mid c_1 \mid C_2)$$

endwhile

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Rows

In Theorem 5.1 partition C and A by rows and do not partition B . In other words, let $M = m, m_i = 1, i = 0, \dots, m-1; N = 1, n_0 = n$; and $K = 1, k_0 = k$. Then

$$C = \begin{pmatrix} c_0^T \\ c_1^T \\ \vdots \\ c_{m-1}^T \end{pmatrix} \quad \text{and} \quad A = \begin{pmatrix} a_0^T \\ a_1^T \\ \vdots \\ a_{m-1}^T \end{pmatrix}$$

so that

$$\begin{pmatrix} c_0^T \\ c_1^T \\ \vdots \\ c_{m-1}^T \end{pmatrix} = C = AB = \begin{pmatrix} a_0^T \\ a_1^T \\ \vdots \\ a_{m-1}^T \end{pmatrix} B = \begin{pmatrix} a_0^T B \\ a_1^T B \\ \vdots \\ a_{m-1}^T B \end{pmatrix}.$$



Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Rows

Homework 5.3.3.2

$$\begin{aligned} & \cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \\ 2 & 1 & -1 \\ 1 & -1 & 2 \end{pmatrix} = \\ & \cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \\ 2 & 1 & -1 \\ 1 & -1 & 2 \end{pmatrix} = \\ & \cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \\ 2 & 1 & -1 \\ 1 & -1 & 2 \end{pmatrix} = \end{aligned}$$

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Rows

In the algorithm for computing $C = AB + C$ the loop indexed by i can be moved to the outside so that

```

for i = 0, ..., m-1
  for j = 0, ..., n-1
    for p = 0, ..., k-1
       $\gamma_{i,j} := \alpha_{i,p} \beta_{p,j} + \gamma_{i,j}$ 
    endfor
  endfor
endfor

```

or

```

for i = 0, ..., m-1
  for p = 0, ..., k-1
    for j = 0, ..., n-1
       $\gamma_{i,j} := \alpha_{i,p} \beta_{p,j} + \gamma_{i,j}$ 
    endfor
  endfor
endfor

```

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication by Rows

Algorithm: $C := \text{GEMM_UNB_VAR2}(A, B, C)$

Partition $A \rightarrow \begin{pmatrix} A_T \\ A_B \end{pmatrix}, C \rightarrow \begin{pmatrix} C_T \\ C_B \end{pmatrix}$
 where A_T has 0 rows, C_T has 0 rows
 while $m(A_T) < m(A)$ do

Repartition

$$\begin{pmatrix} A_T \\ A_B \end{pmatrix} \rightarrow \begin{pmatrix} A_0 \\ a_1^T \\ A_2 \end{pmatrix}, \begin{pmatrix} C_T \\ C_B \end{pmatrix} \rightarrow \begin{pmatrix} C_0 \\ c_1^T \\ C_2 \end{pmatrix}$$

where a_1 has 1 row, c_1 has 1 row

$$c_1^T := a_1^T B + c_1^T$$

Continue with

$$\begin{pmatrix} A_T \\ A_B \end{pmatrix} \leftarrow \begin{pmatrix} A_0 \\ a_1^T \\ A_2 \end{pmatrix}, \begin{pmatrix} C_T \\ C_B \end{pmatrix} \leftarrow \begin{pmatrix} C_0 \\ c_1^T \\ C_2 \end{pmatrix}$$

endwhile

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication with Rank-1 Updates

In Theorem 5.1 partition A and B by columns and rows, respectively, and do not partition C . In other words, let $M = 1$, $m_0 = m$, $N = 1$, $n_0 = n$; and $K = k$, $k_p = 1$, $p = 0, \dots, k-1$. Then

$$A = \begin{pmatrix} a_0 & a_1 & \dots & a_{k-1} \end{pmatrix} \quad \text{and} \quad B = \begin{pmatrix} b_0^T \\ b_1^T \\ \vdots \\ b_{k-1}^T \end{pmatrix}$$

so that

$$C = AB = \begin{pmatrix} a_0 & a_1 & \dots & a_{k-1} \end{pmatrix} \begin{pmatrix} b_0^T \\ b_1^T \\ \vdots \\ b_{k-1}^T \end{pmatrix} = a_0 b_0^T + a_1 b_1^T + \dots + a_{k-1} b_{k-1}^T.$$

Notice that each term $a_p b_p^T$ is an outer product of a_p and b_p . Thus, if we start with $C := 0$, the zero matrix, then we can compute $C := AB + C$ as

$$C := a_{k-1} b_{k-1}^T + (\dots + (a_p b_p^T + (\dots + (a_1 b_1^T + (a_0 b_0^T + C)) \dots)) \dots),$$

which illustrates that $C := AB$ can be computed by first setting C to zero, and then repeatedly updating it with rank-1 updates.



Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication with Rank-1 Updates

Homework 5.3.4.1

$$\begin{aligned} & \cdot \begin{pmatrix} 1 \\ -1 \\ 0 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \end{pmatrix} = \\ & \cdot \begin{pmatrix} -2 \\ 2 \\ 1 \end{pmatrix} \begin{pmatrix} 2 & 1 & -1 \end{pmatrix} = \\ & \cdot \begin{pmatrix} 2 \\ 1 \\ 2 \end{pmatrix} \begin{pmatrix} 1 & -1 & 2 \end{pmatrix} = \\ & \cdot \begin{pmatrix} 1 & -2 & 2 \\ -1 & 2 & 1 \\ 0 & 1 & 2 \end{pmatrix} \begin{pmatrix} -1 & 0 & 1 \\ 2 & 1 & -1 \\ 1 & -1 & 2 \end{pmatrix} = \end{aligned}$$

Algorithms for Computing Matrix-Matrix Multiplication

• Matrix-Matrix Multiplication with Rank-1 Updates

Algorithm: $C := \text{GEMM_UNB_VAR3}(A, B, C)$

Partition $A \rightarrow \begin{pmatrix} A_L & A_R \end{pmatrix}$, $B \rightarrow \begin{pmatrix} B_T \\ B_B \end{pmatrix}$
 where A_L has 0 columns, B_T has 0 rows
 while $n(A_L) < n(A)$ do

Repartition

$$\begin{pmatrix} A_L & A_R \end{pmatrix} \rightarrow \begin{pmatrix} A_0 & a_1 & A_2 \end{pmatrix} \cdot \begin{pmatrix} B_T \\ B_B \end{pmatrix} \rightarrow \begin{pmatrix} B_0 \\ b_1^T \\ B_2 \end{pmatrix}$$

where a_1 has 1 column, b_1 has 1 row

$$C := a_1 b_1^T + C$$

Continue with

$$\begin{pmatrix} A_L & A_R \end{pmatrix} \leftarrow \begin{pmatrix} A_0 & a_1 & A_2 \end{pmatrix} \cdot \begin{pmatrix} B_T \\ B_B \end{pmatrix} \leftarrow \begin{pmatrix} B_0 \\ b_1^T \\ B_2 \end{pmatrix}$$

endwhile

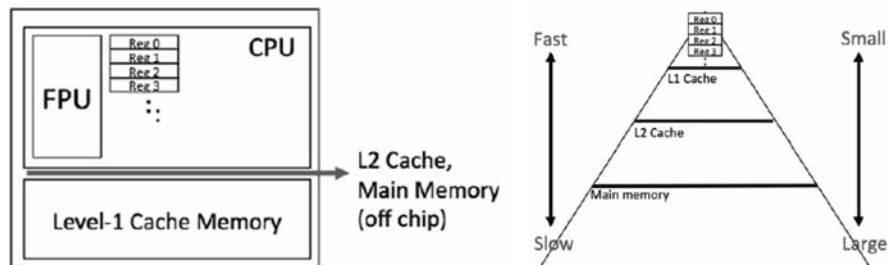
Spark

Enrichment

• Slicing and Dicing for Performance

— CPU : Central Processing Unit

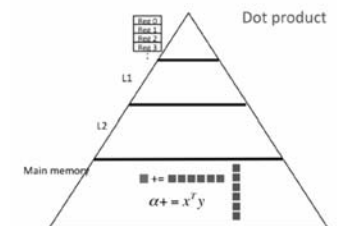
- FPU : Floating Point Unit
- Registers



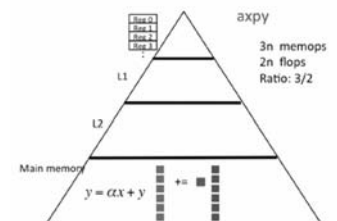
Enrichment

• Slicing and Dicing for Performance

— Vector-Vector Computations

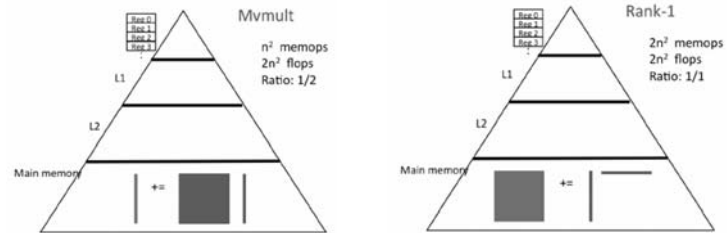


— Matrix-Vector Computations



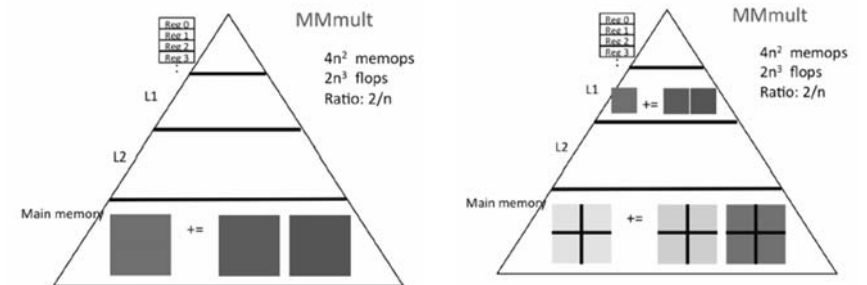
Enrichment

- Slicing and Dicing for Performance
 - Matrix-Vector Computations



Enrichment

- Slicing and Dicing for Performance
 - Matrix-Matrix Computations



Questions and Answers

