# Compiler Construction

#### Instructor: เกียรติกูล เจียรนัยชนะกิจ

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#### Grading criteria:

Assignments 50 %

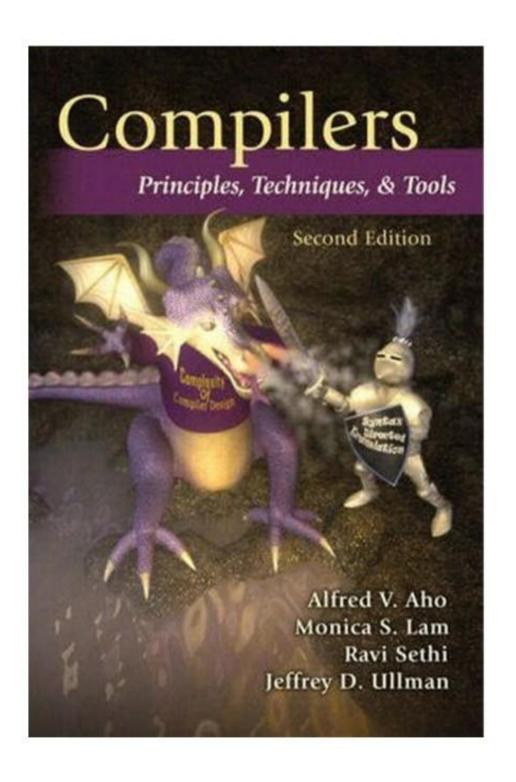
- Lexical analyzer 20 %

- Parser 20 %

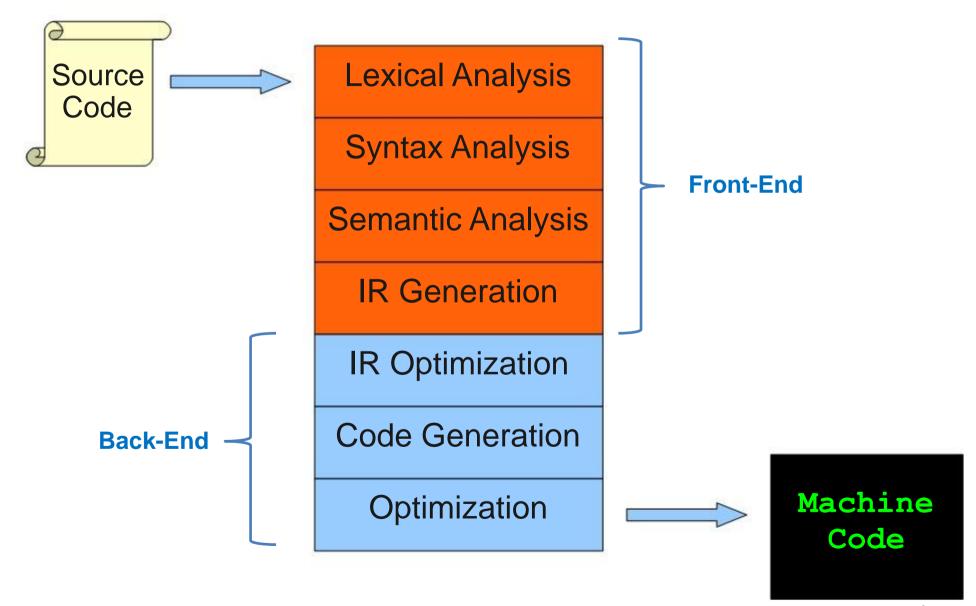
- Report 10 %

Final exam 50%

Slides: Adapted from Compiler course @ Stanford University



#### The Structure of a Modern Compiler



"Identify tokens in the source code."

```
while (y < z) {
    int x = a + b;
    y += x;
}</pre>
```

```
T While
T LeftParen
T Identifier v
T Less
T Identifier z
T RightParen
T OpenBrace
T Int.
T Identifier x
T Assign
T Identifier a
T Plus
T Identifier b
T Semicolon
T Identifier y
T PlusAssign
T Identifier x
T Semicolon
T CloseBrace
```

A list of "Tokens"

**Lexical Analysis** Syntax Analysis Semantic Analysis IR Generation IR Optimization **Code Generation Optimization** 

while (y < z) { int x = a + b;y += x;While Sequence ID ID + + ID ID

"Identify how those tokens relate to each other."

**Lexical Analysis** 

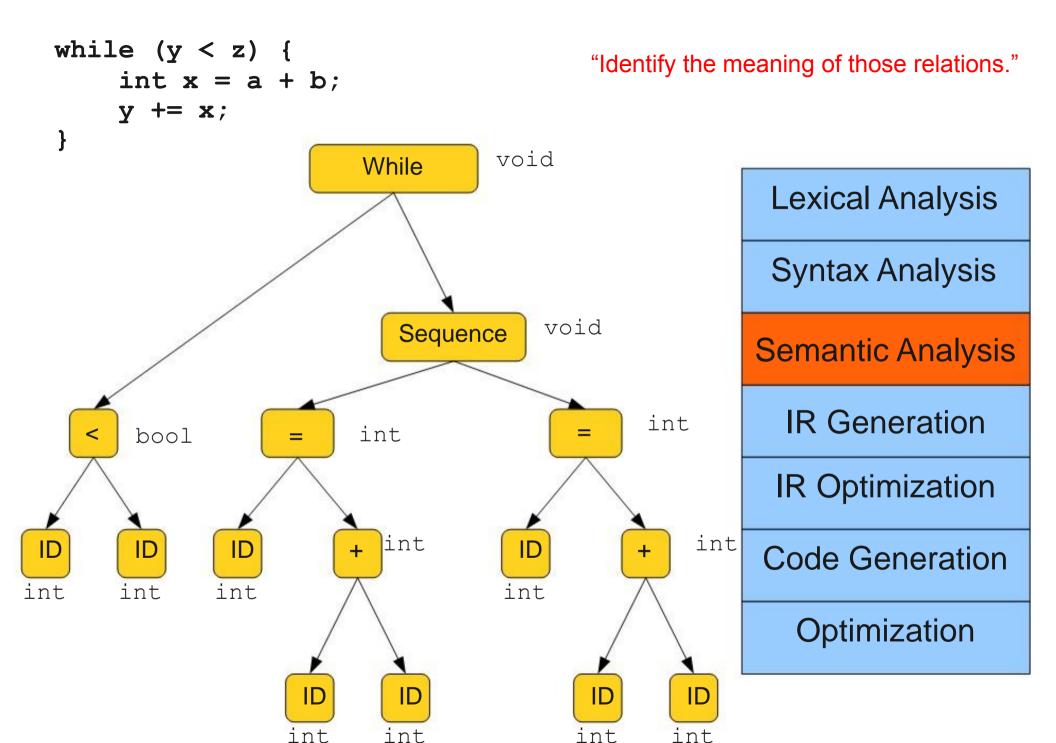
Syntax Analysis

Semantic Analysis

IR Generation

**IR Optimization** 

**Code Generation** 



```
while (y < z) {
    int x = a + b;
    y += x;
}</pre>
```

Exit:

IR: Intermediate Representation

**Lexical Analysis** 

Syntax Analysis

Semantic Analysis

IR Generation

**IR Optimization** 

**Code Generation** 

"Optimize the assembly-like code."

```
while (y < z) {
    int x = a + b;
    y += x;
}</pre>
```

x = a + b

Loop: \_t1 = y < z
if not \_t1 goto Exit

y = x + y

goto Loop

Exit:

**Lexical Analysis** 

Syntax Analysis

Semantic Analysis

**IR** Generation

**IR** Optimization

**Code Generation** 

"Generate the machine code."

```
while (y < z) {
    int x = a + b;
    y += x;
}</pre>
```

Exit:

```
add $1, $2, $3
Loop: slt $6, $4, $5
beq $6, Exit
add $4, $4, $1
b Loop
```

**Lexical Analysis** 

Syntax Analysis

Semantic Analysis

IR Generation

**IR Optimization** 

**Code Generation** 

"Optimize the machine code."

```
while (y < z) {
    int x = a + b;
    y += x;
}</pre>
```

```
add $1, $2, $3

Loop: blt $4, $5, Exit

add $4, $4, $1

b Loop
```

Exit:

**Lexical Analysis** 

Syntax Analysis

Semantic Analysis

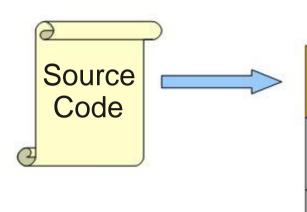
**IR** Generation

**IR Optimization** 

**Code Generation** 

# Lexical Analysis

#### Where We Are



**Lexical Analysis** 

Syntax Analysis

Semantic Analysis

IR Generation

IR Optimization

**Code Generation** 

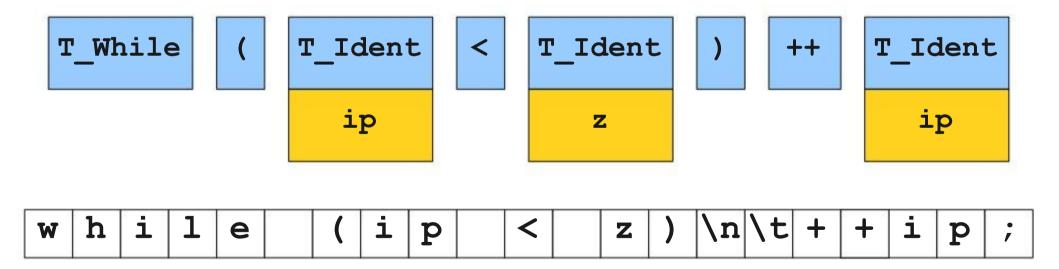
Optimization

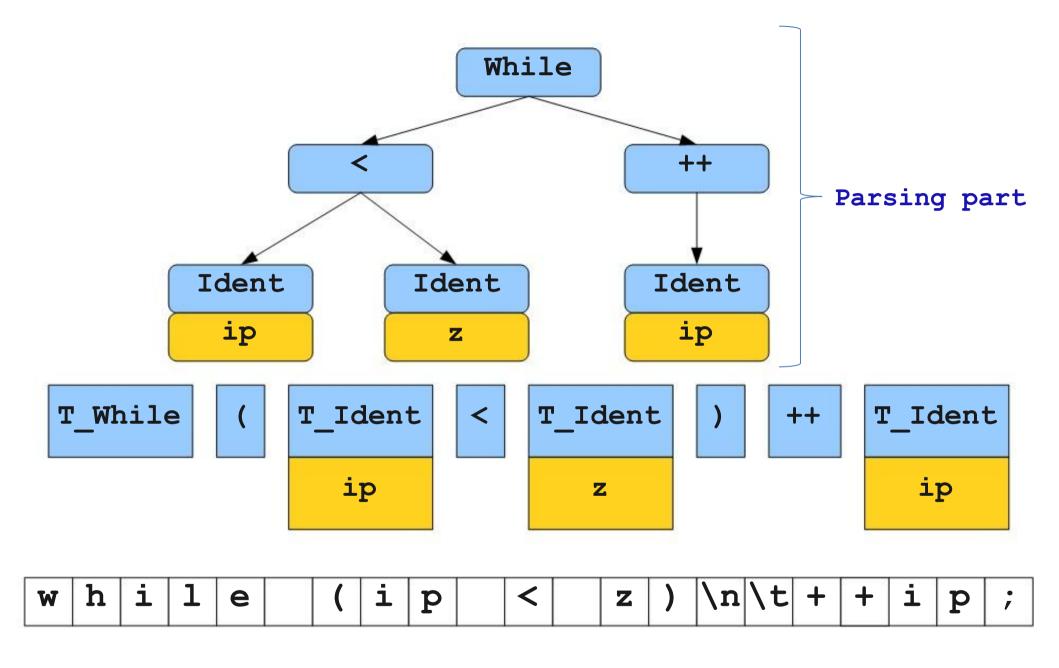


Machine Code

#### Why do Lexical Analysis?

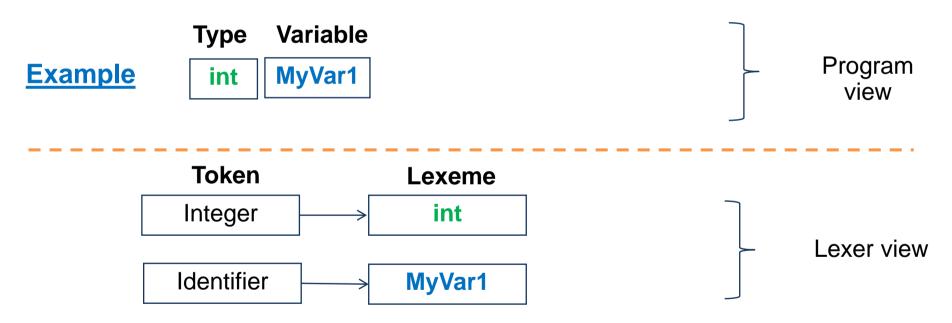
- Dramatically simplify parsing.
  - Eliminate whitespace, comments.
  - Provide input for the parsing part.





#### Goals of Lexical Analysis

 Convert from physical description of a program into sequence of tokens.



 The token stream will be used in the parser to create the syntax tree.

## Challenges in Lexical Analysis

How to partition the program into lexemes?

```
DO 5 I = 1,25

Equivalent to

DO

// repeat loop for I = 1 to 25

5:
```

```
DO 5 I = 1.25

Equivalent to

DO5I = 1.25

Whitespace is irrelevant
```

 FORTRAN requires look ahead symbol to partition the program into lexemes.

## Not-so-Great Moments in Scanning

PL/I: Keywords can be used as identifiers.

```
IF THEN THEN THEN = ELSE; ELSE ELSE = IF
```

 Can be difficult to determine how to label lexemes.

PL/I (Programming Language One) Originally developed by IBM

#### **Choosing Good Tokens**

- Very much dependent on the language.
- Typical kinds of token:
  - Reserve words (If, Then, Else, For, While, Do, etc.).
  - Punctuation symbols (; " '<>+-\*/ % ! [ ] { } && || , etc.)
  - Identifier
  - Number (Integer, floating point, etc.)
  - Discard irrelevant information (whitespace, comments)
- Token should be recognized in maximal munch manner <u>Example</u>:

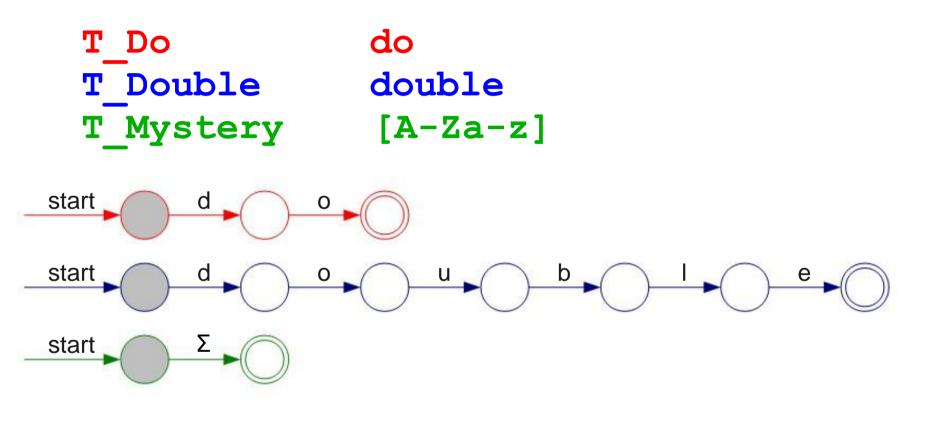
dot should be recognized as an Identifier "dot" (Not a reserved word DO and an identifier "t")

#### Implementing Maximal Munch Lexer

Represent all tokens as regular expressions.

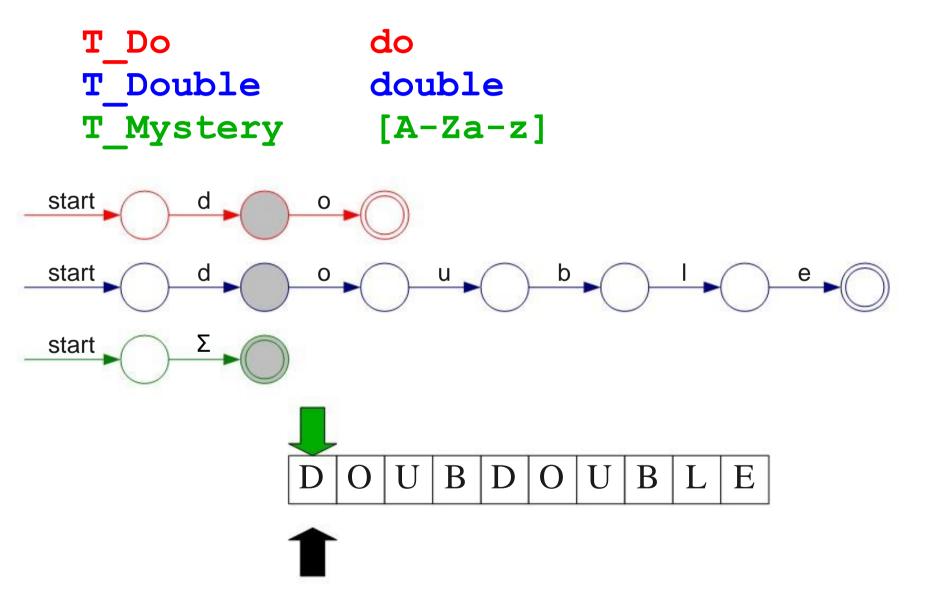
<b>Example</b>	Token	Regular expression
	T_Do	do
	T_IDEN	[a-z][a-zA-Z0-9]*

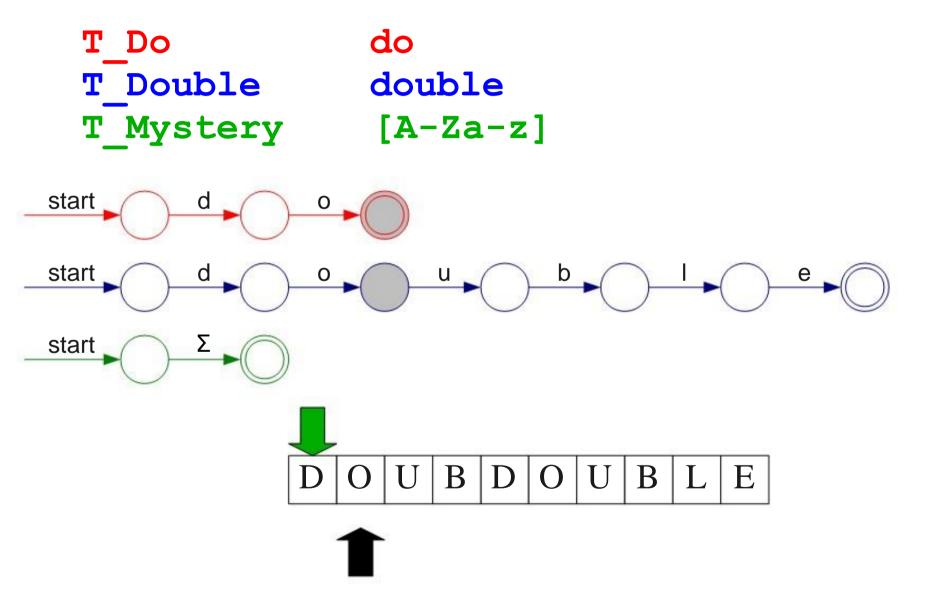
- Convert all regular expressions to NFAs.
- Scan input from Left to Right
- Run all NFAs in parallel, keeping track of the last match.
- When all automata get stuck, report the longest match and restart the search at that point.

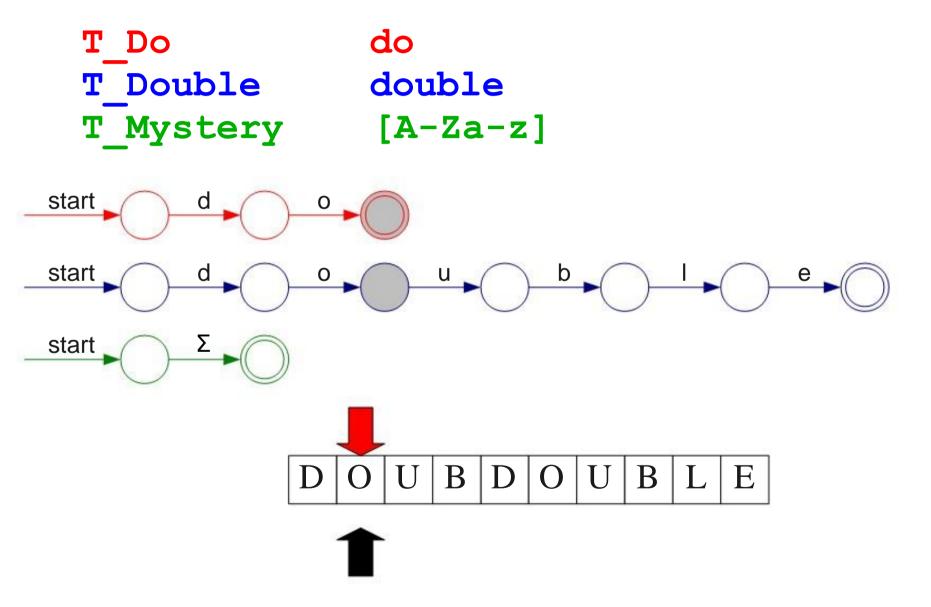


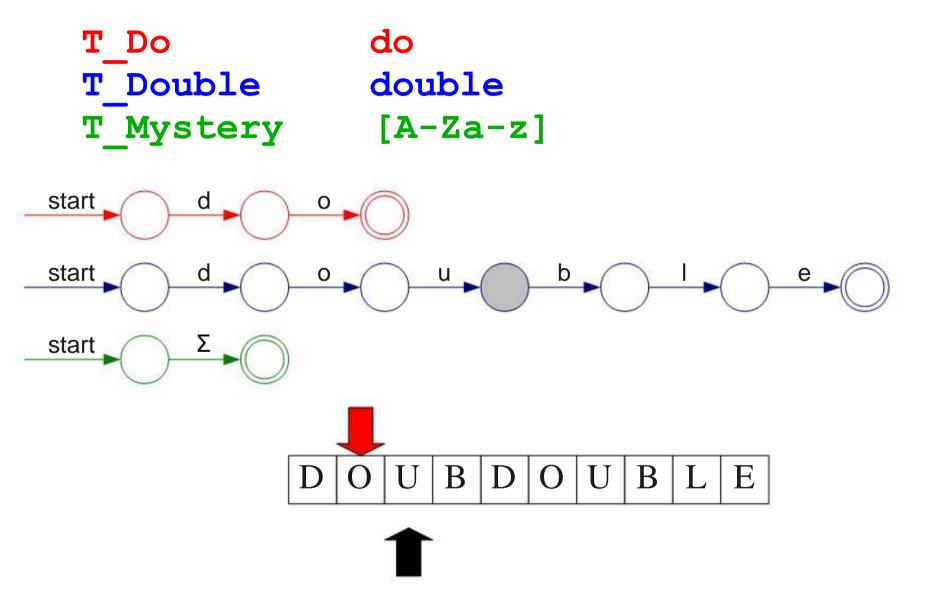


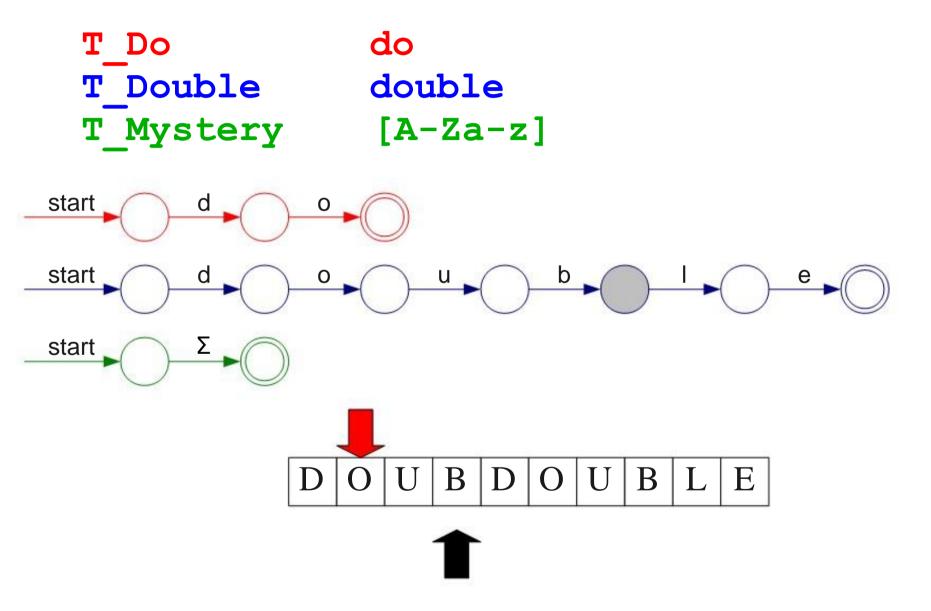


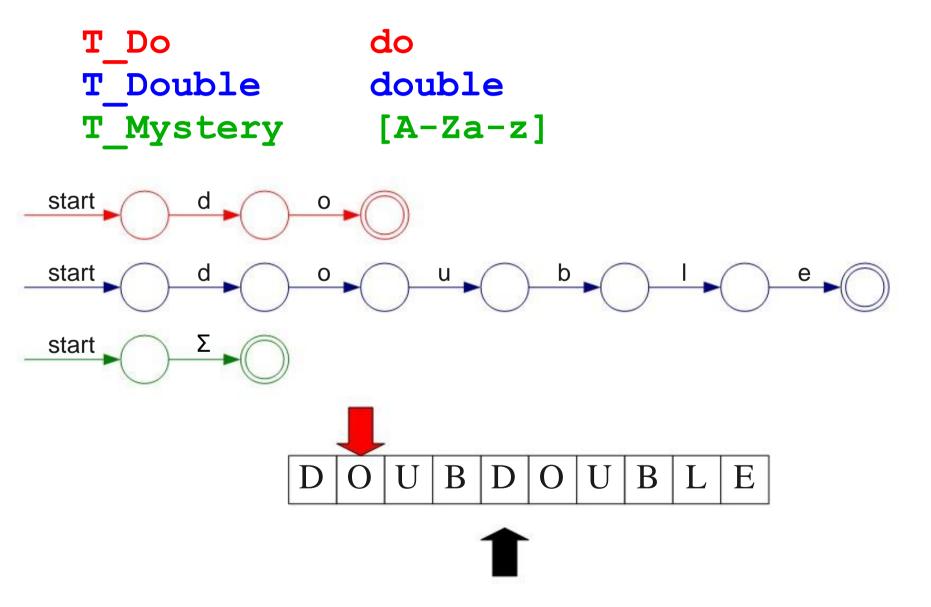


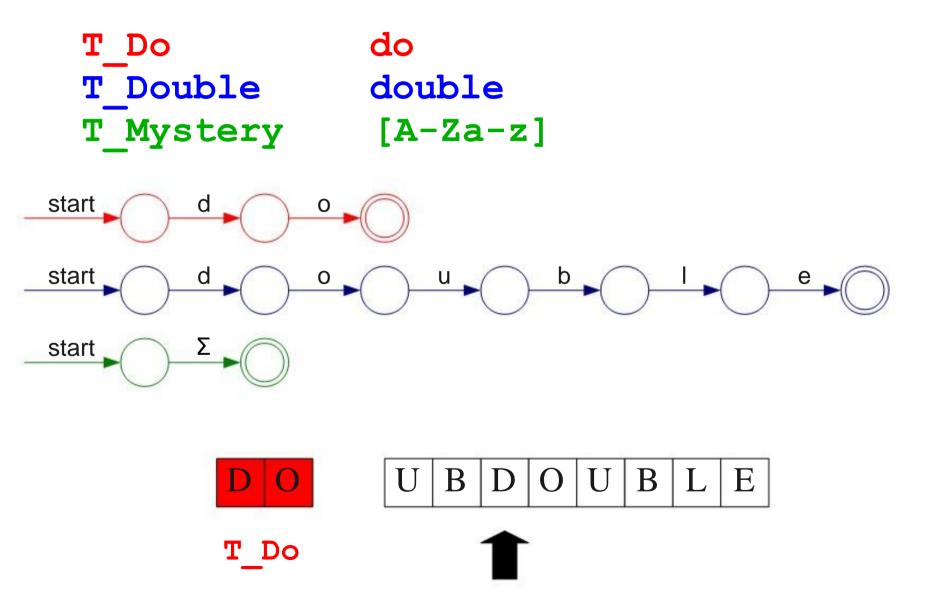


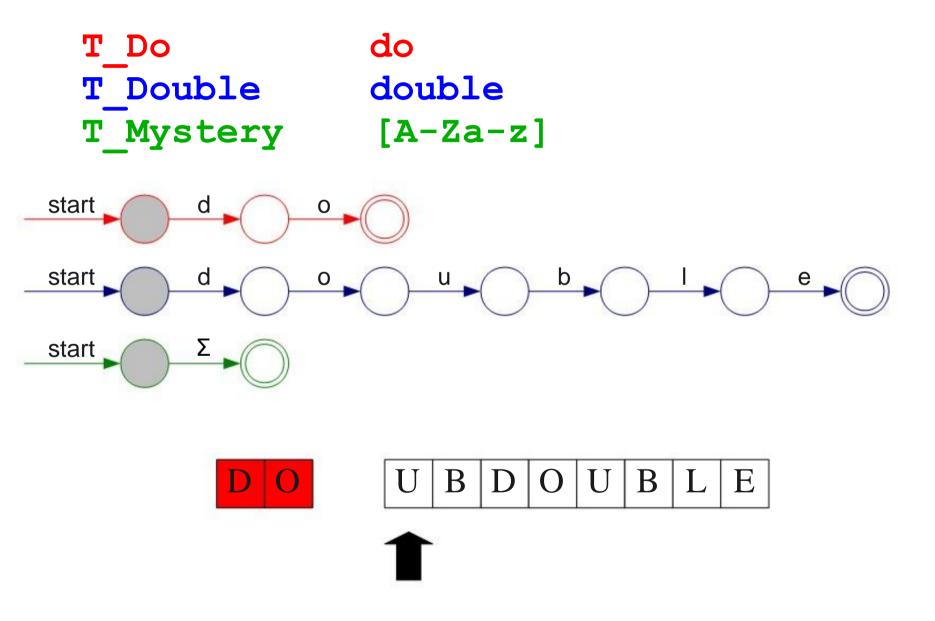


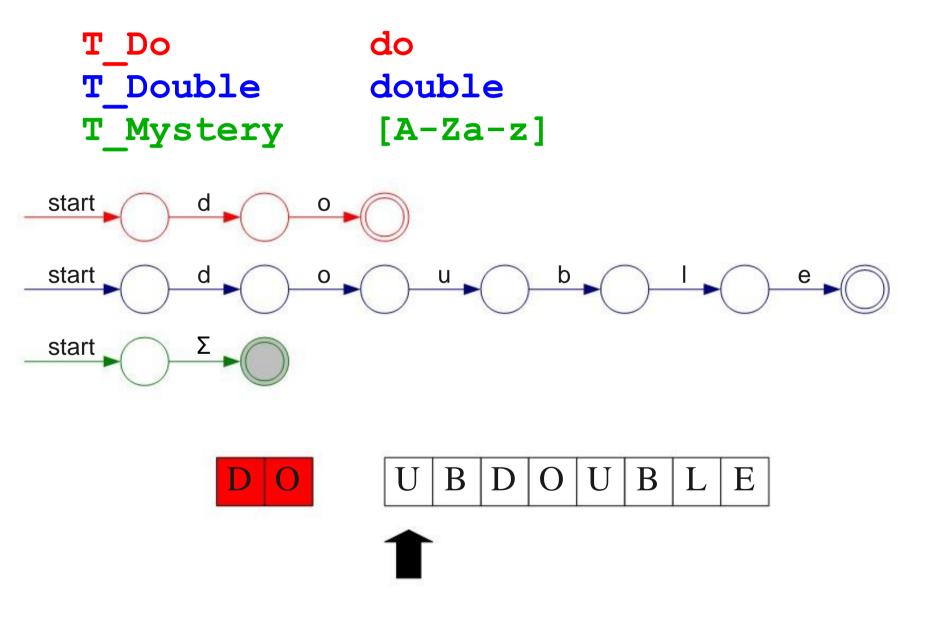


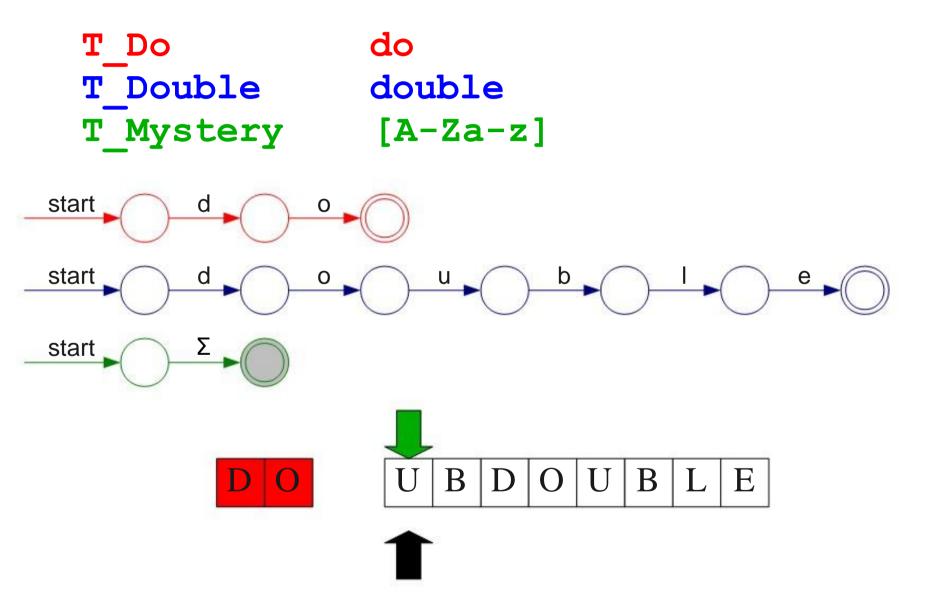


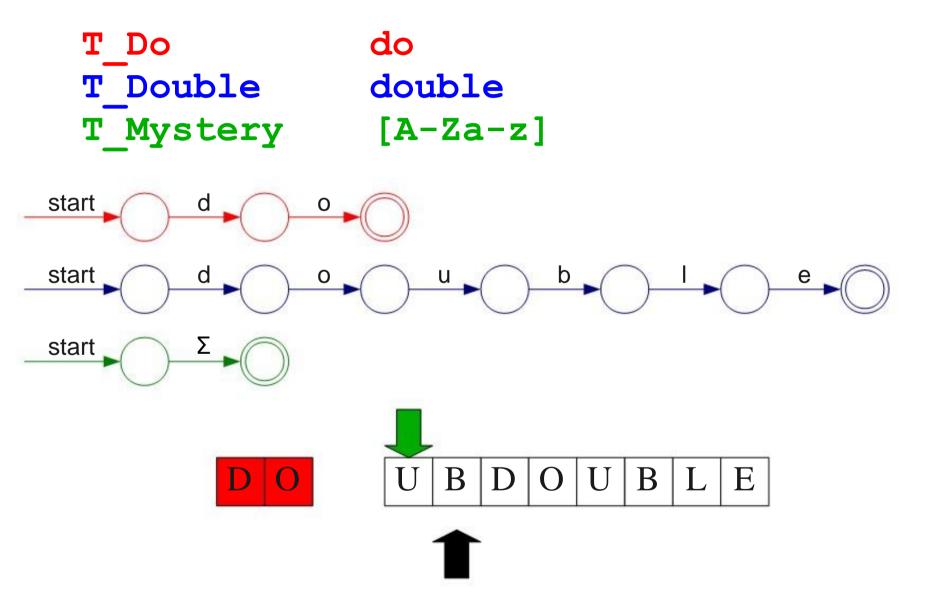


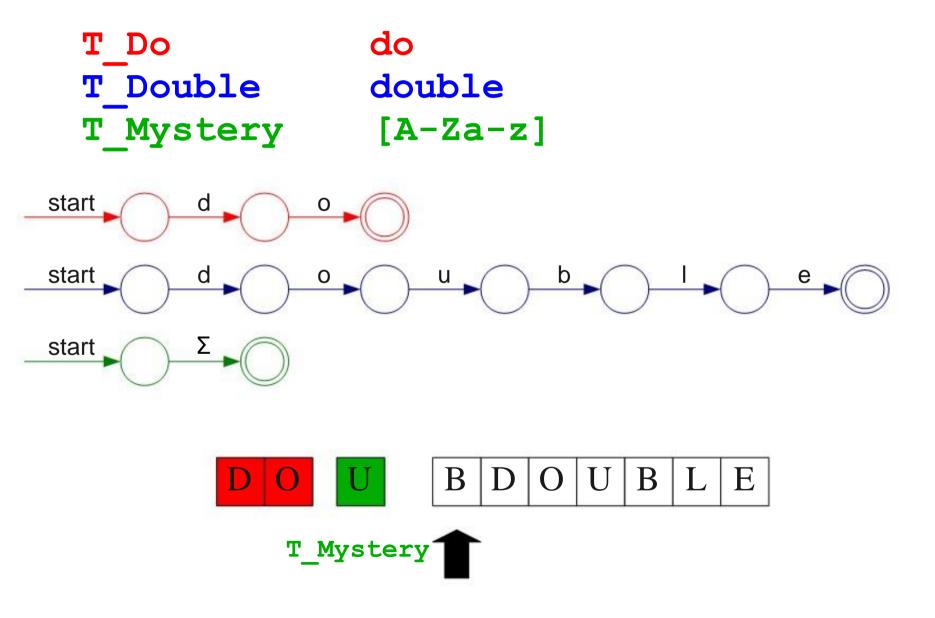


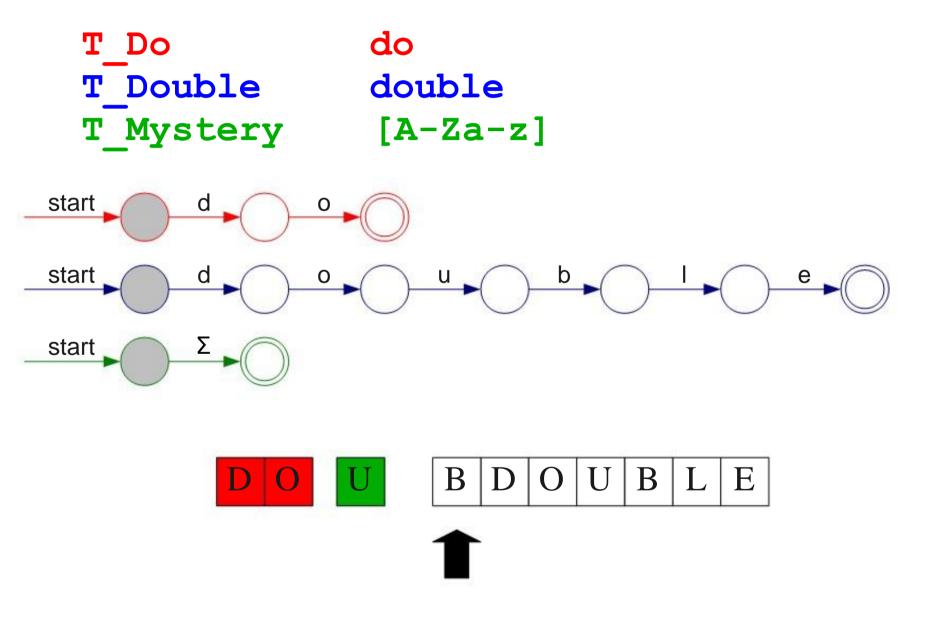


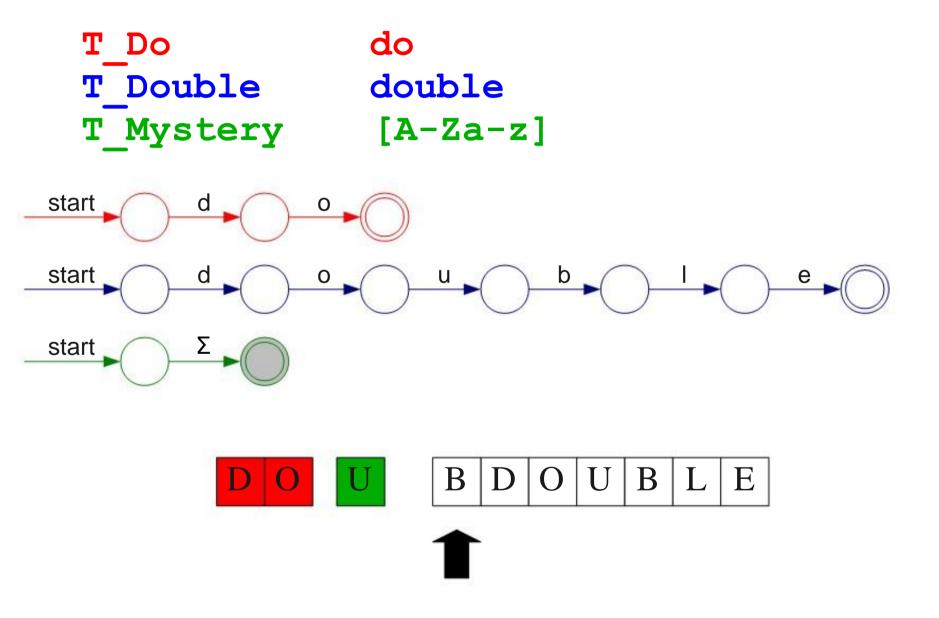


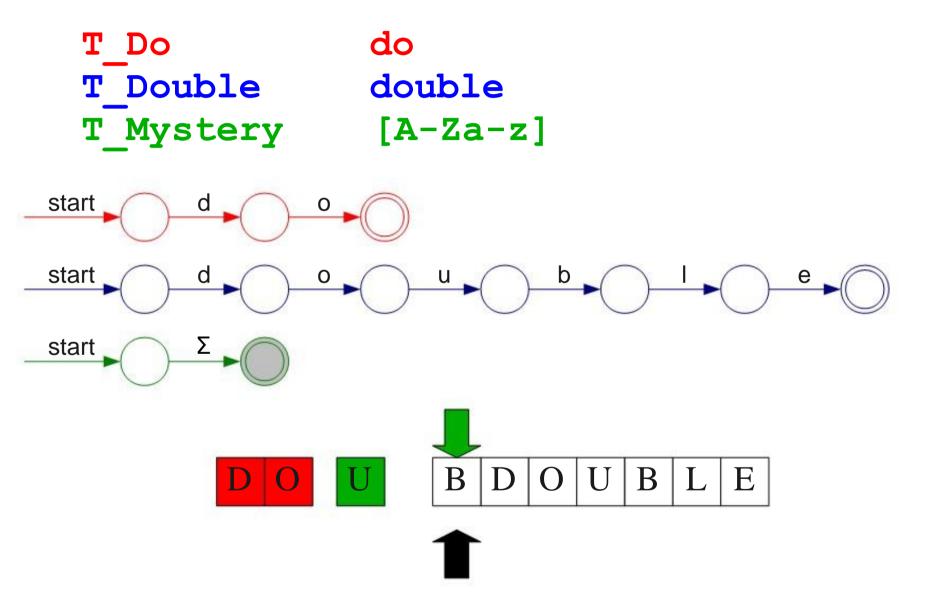


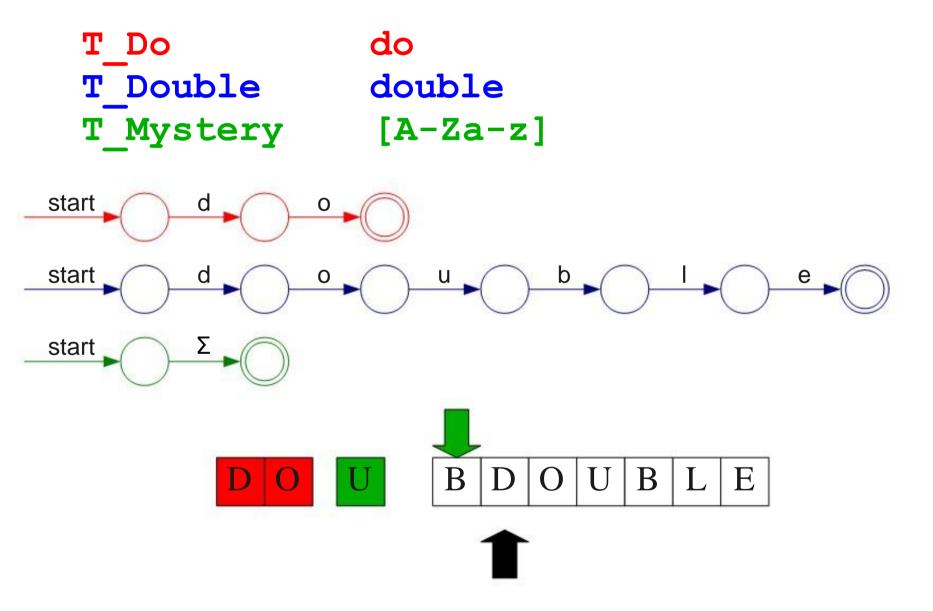


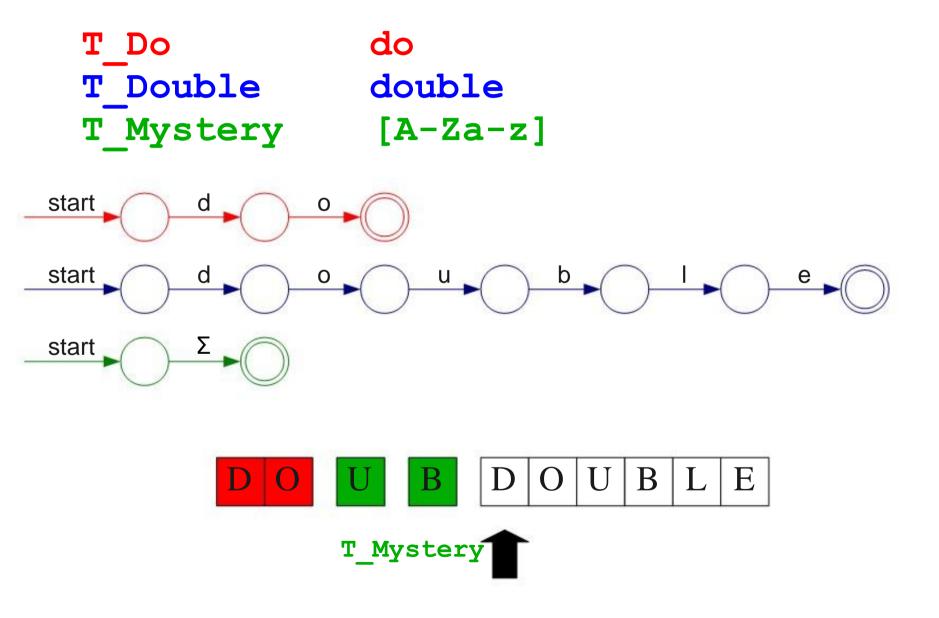


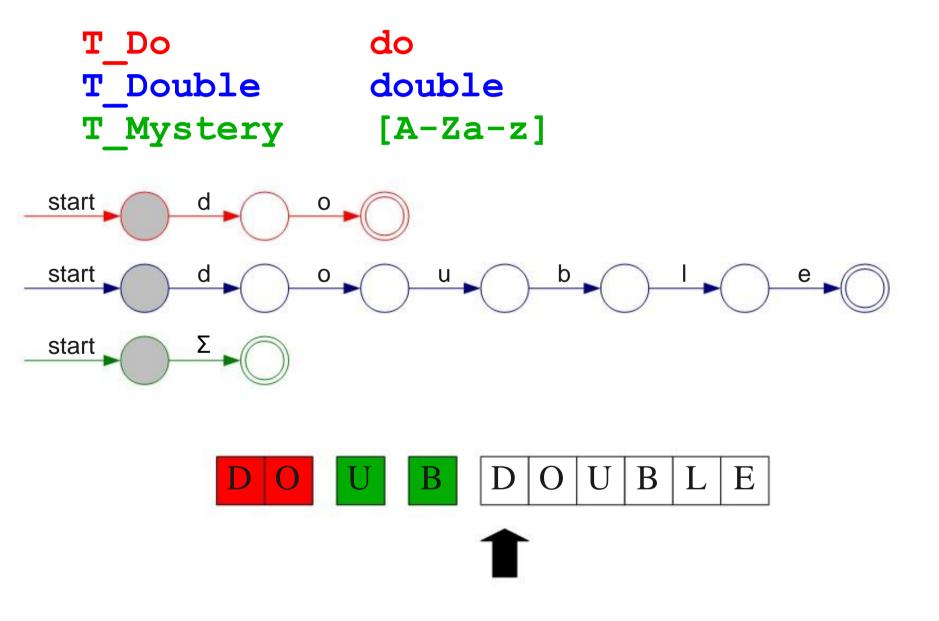


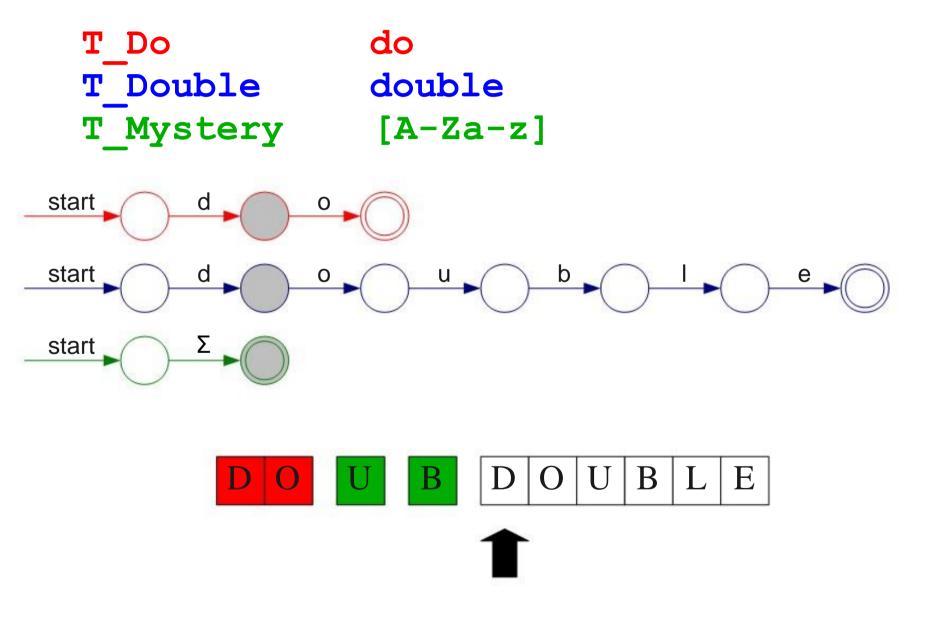


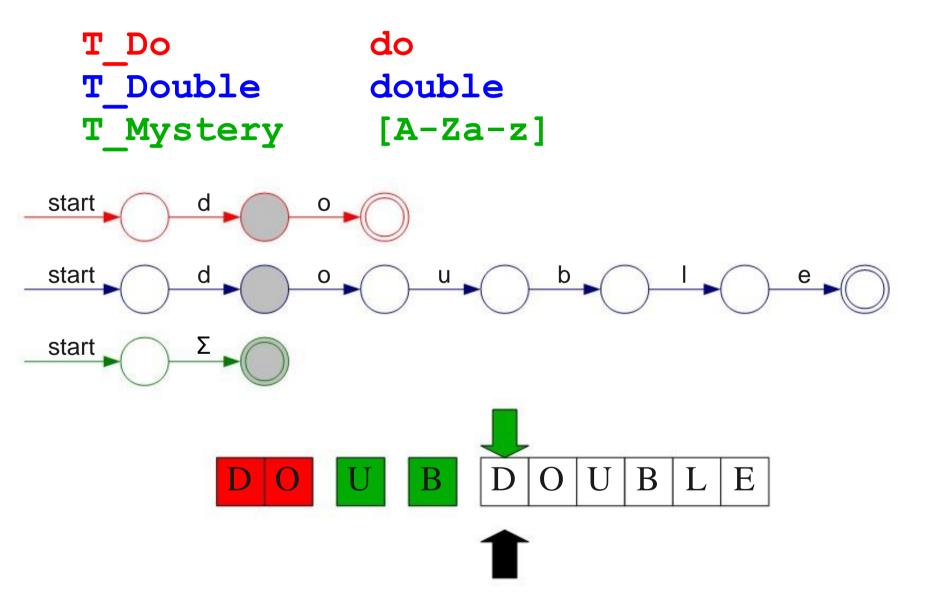


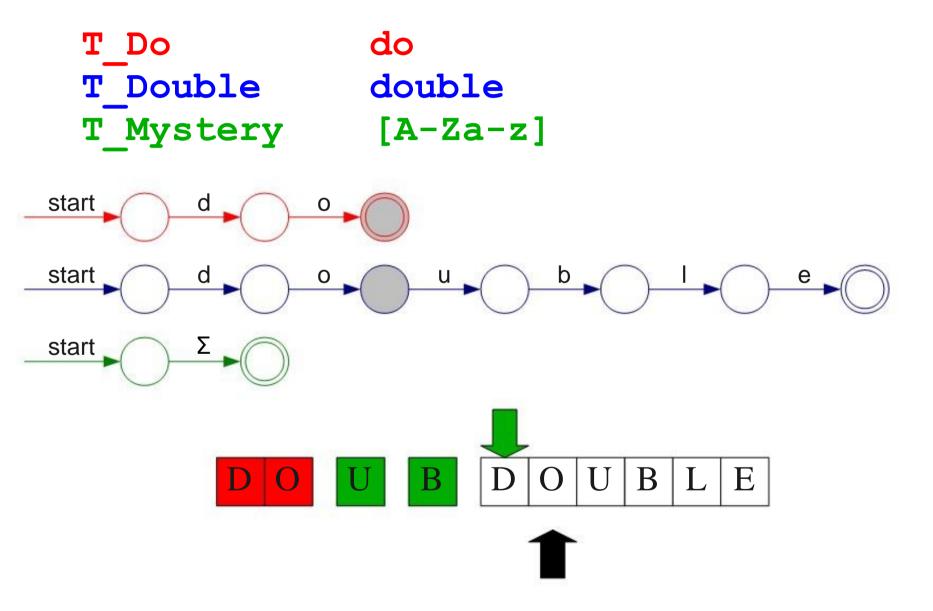


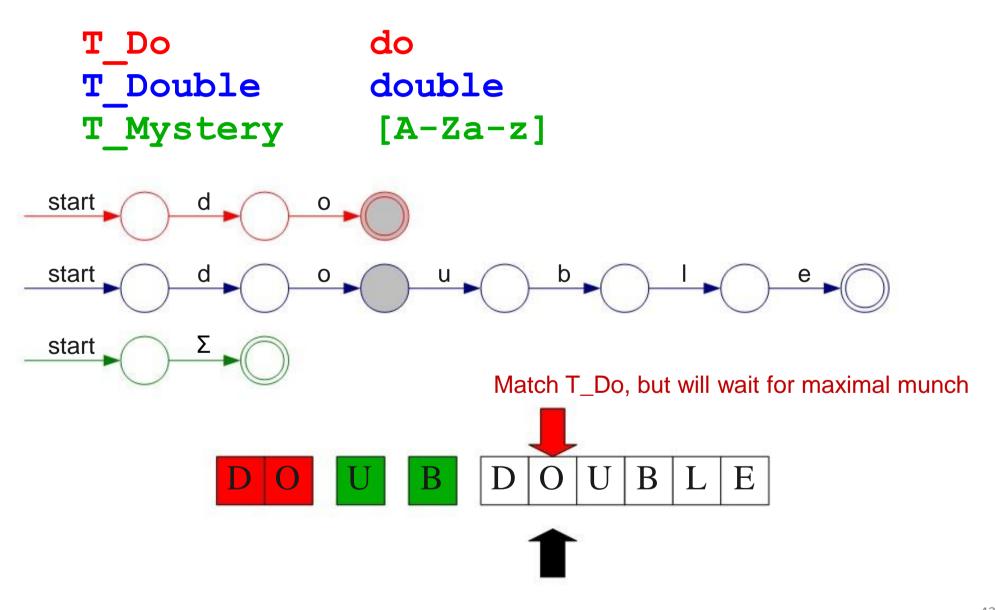


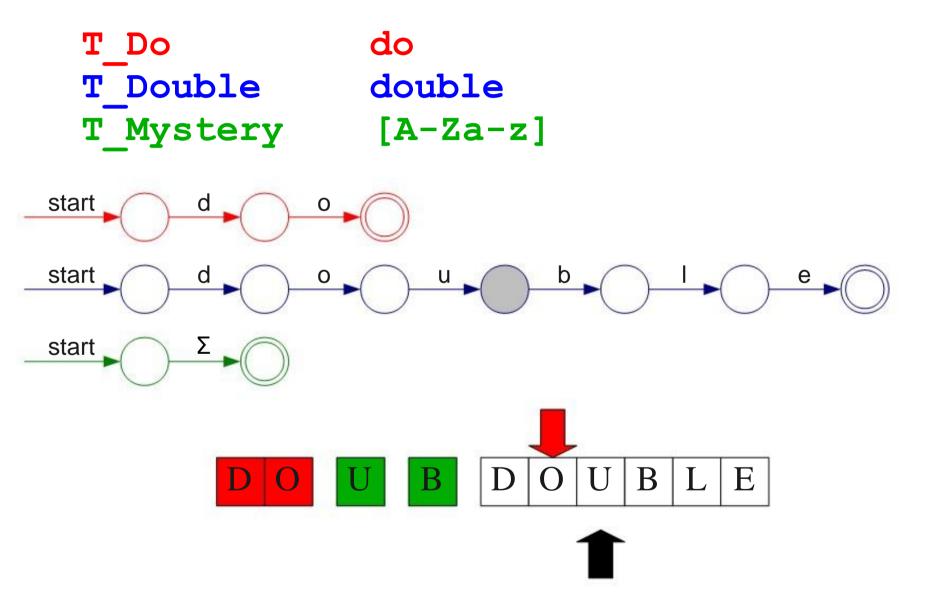


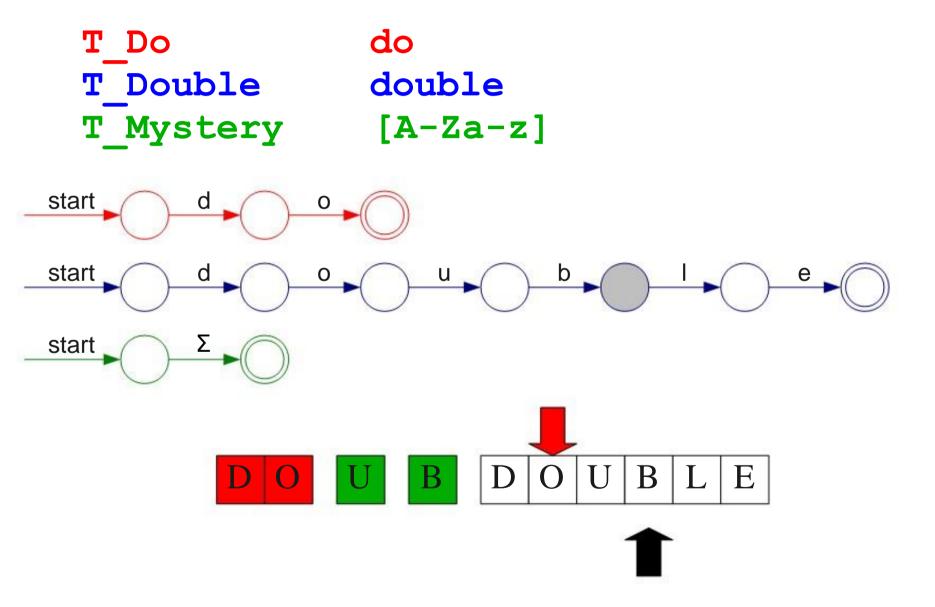


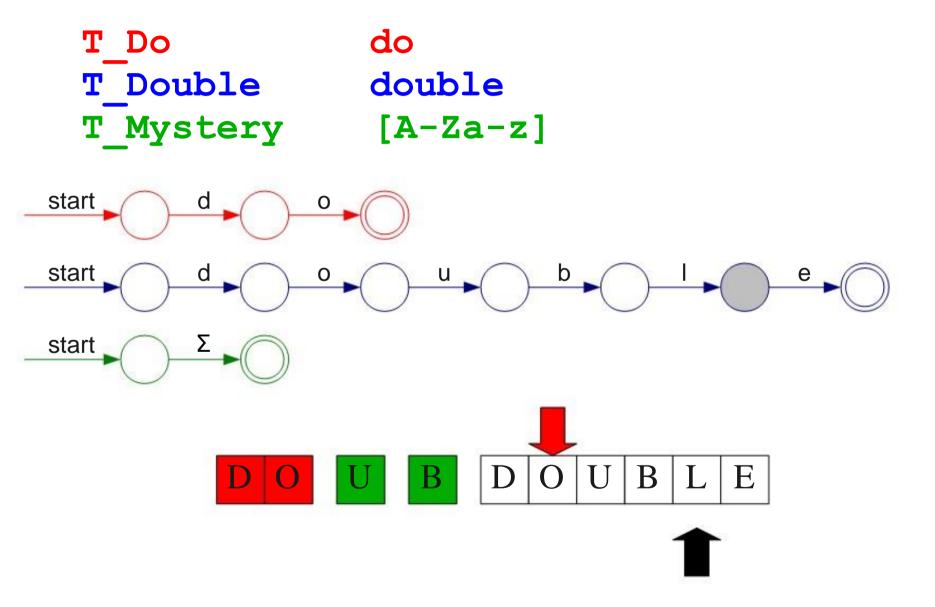


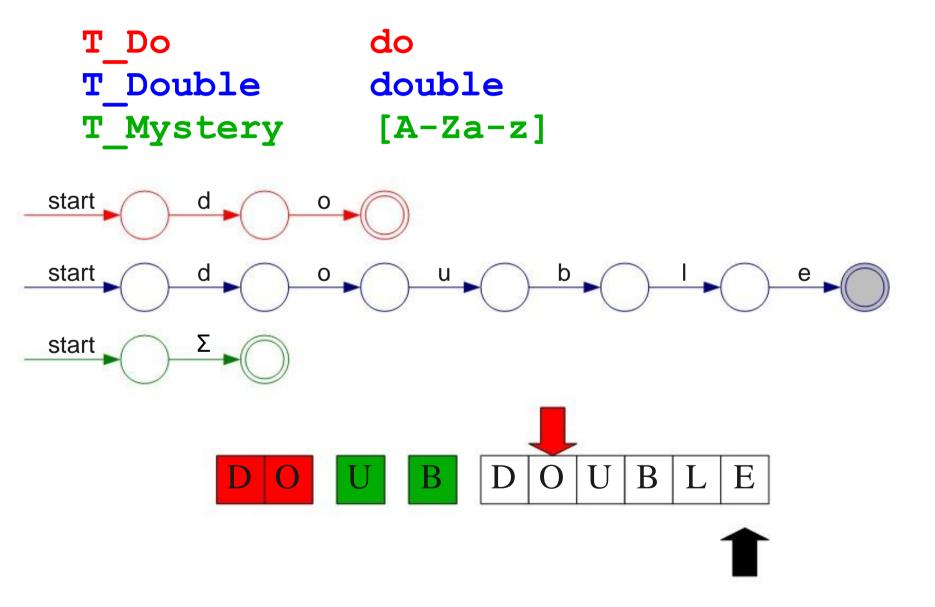


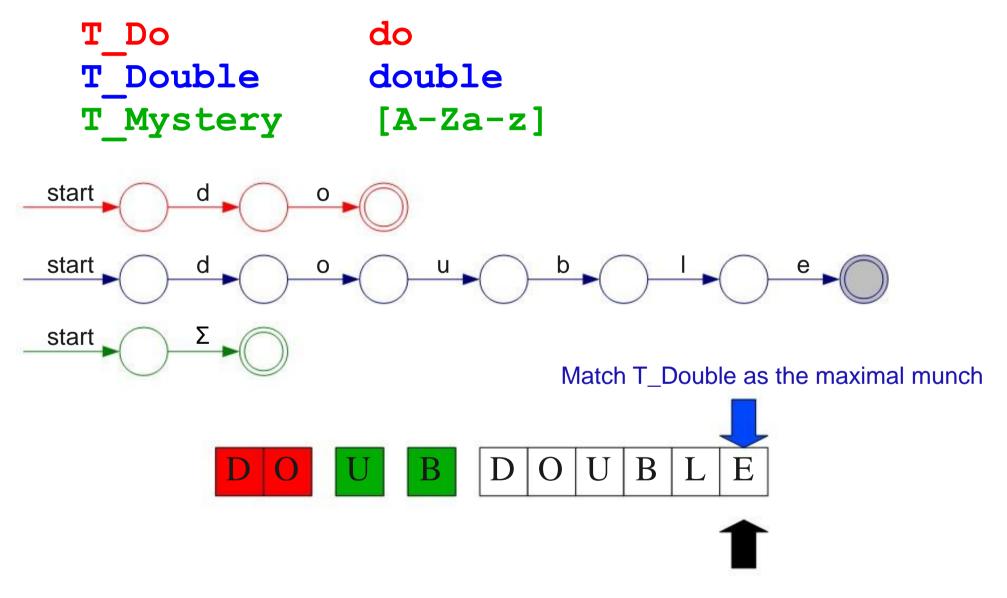


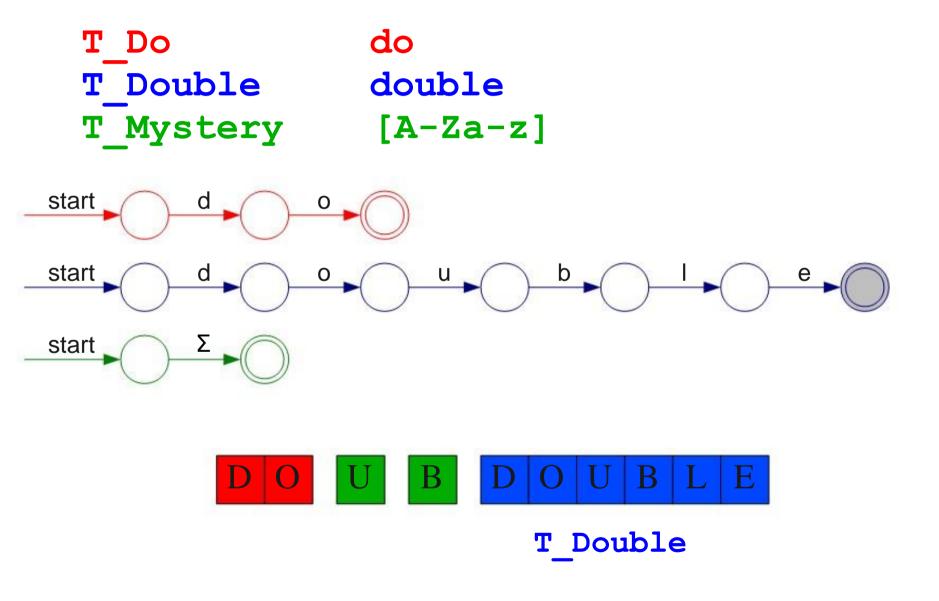




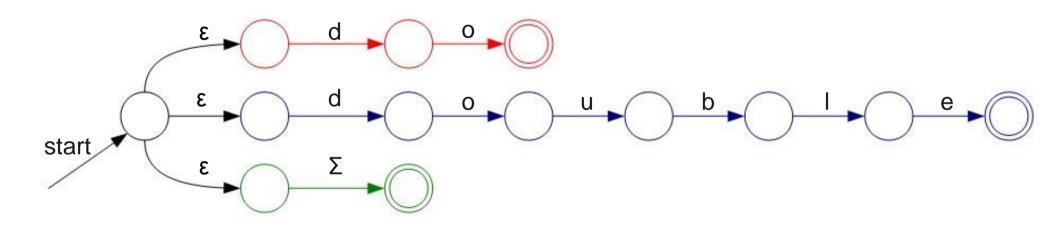








# A Minor Simplification



#### Conflict

```
T_Do do
T_Double double
T_Identifier [A-Za-z_][A-Za-z0-9_]*
```

Multiple rules can be executed for the same lexeme. As a result, one lexeme can be recognized as more than one token.

#### Conflict

```
T_Do do
T_Double double
T_Identifier [A-Za-z_][A-Za-z0-9_]*

Example d o u b l e
```

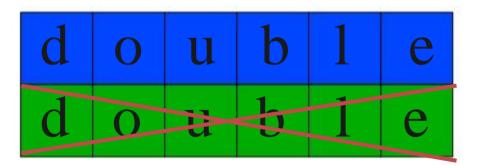
d	0	u	b	1	e	T_Double
d	0	u	b	1	e	T_Identifier

### More Tiebreaking

- When two regular expressions apply, choose the one with the greater "priority."
- Simple priority system: which rule was defined first?

#### Conflict

```
T_Do
    do
T_Double    double
T_Identifier [A-Za-z_][A-Za-z0-9_]*
```



#### One Last Detail...

- What if nothing matches?
- Trick: Add a "catch-all" rule that matches any character and reports an error. Don't forget to put it at the last line (lowest priority).

### Summary of Conflict Resolution

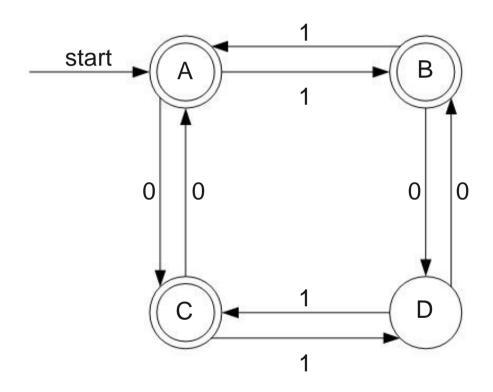
- Construct an automaton for each regular expression.
- Merge them into one automaton by adding a new start state.
- Scan the input, keeping track of the last known match.
- Break ties by choosing higher-precedence matches.
- Have a catch-all rule to handle errors.

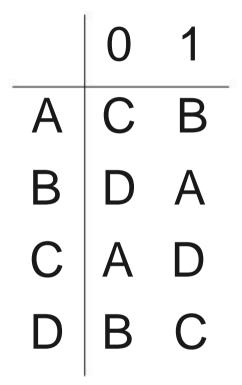
# Speeding up the Scanner

### DFAs (Review)

- The automata we've seen so far have all been NFAs.
- A DFA is like an NFA, but with tighter restrictions:
  - Every state must have exactly one transition defined for every letter.
  - ε-moves are not allowed.

# A Sample DFA



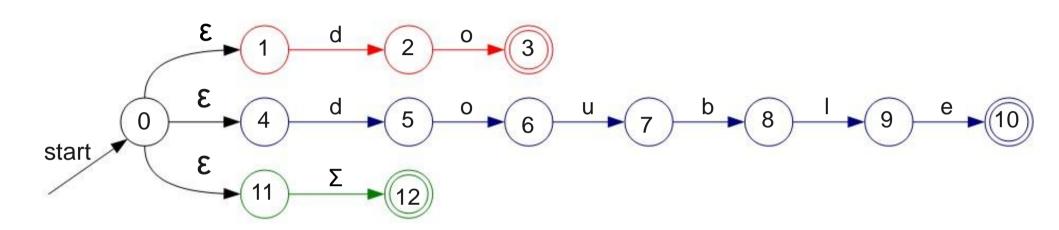


# Speeding up Matching

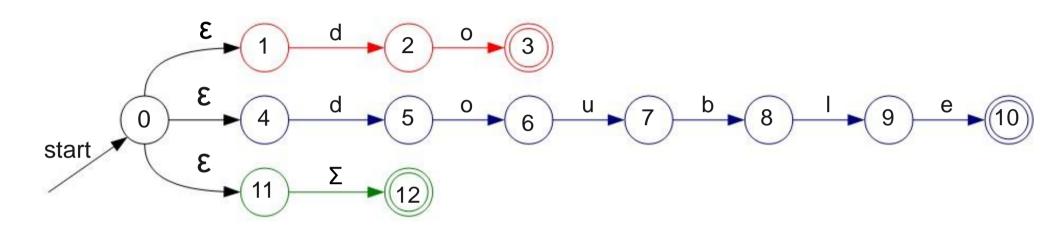
 DFAs usually take shorter execution time than NFAs.

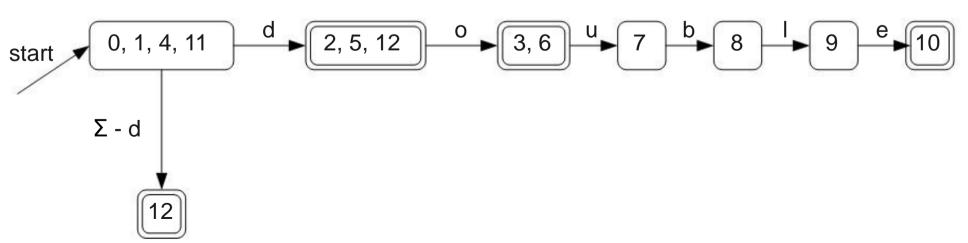


#### From NFA to DFA

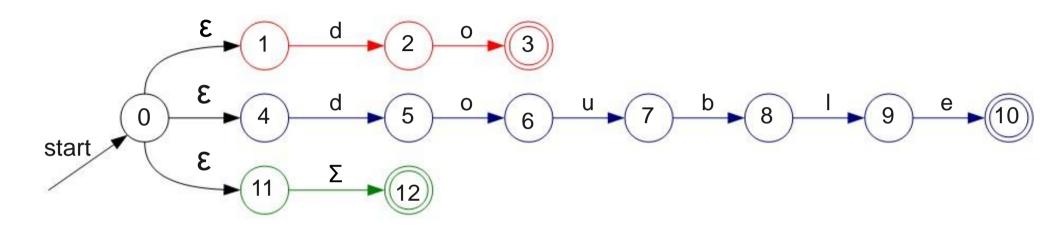


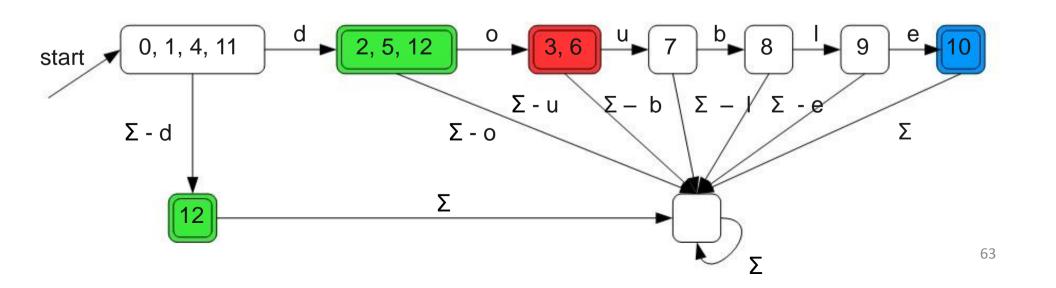
#### From NFA to DFA





#### From NFA to DFA





# Introduction to flex

(Lexical Analysis Tool)

#### What is flex?

- Automated tool for generating scanner (lexer).
- Uses maximal-munch/precedence system.
- Internally, builds a DFA from regular expressions.
- Plus several more features...

### A Simple flex File

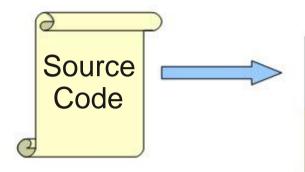
```
%%
[A-Za-z]* printf("Word\n");
[0-9]* printf("Number\n");
[ \t\n] ;
    printf("Undefine Symbol\n");
```

Further reading: http://flex.sourceforge.net/manual/

### Summary

- Lexical analysis splits input text into tokens holding a lexeme and an attribute.
- Lexemes are sets of strings often defined with regular expressions.
- Regular expressions can be converted to NFAs and from there to DFAs.
- Maximal-munch using an automaton allows for fast scanning.
- Not all tokens come directly from the source code.

### **Next Time**



Lexical Analysis

Syntax Analysis

Semantic Analysis

IR Generation

**IR Optimization** 

**Code Generation** 

Optimization



Machine Code