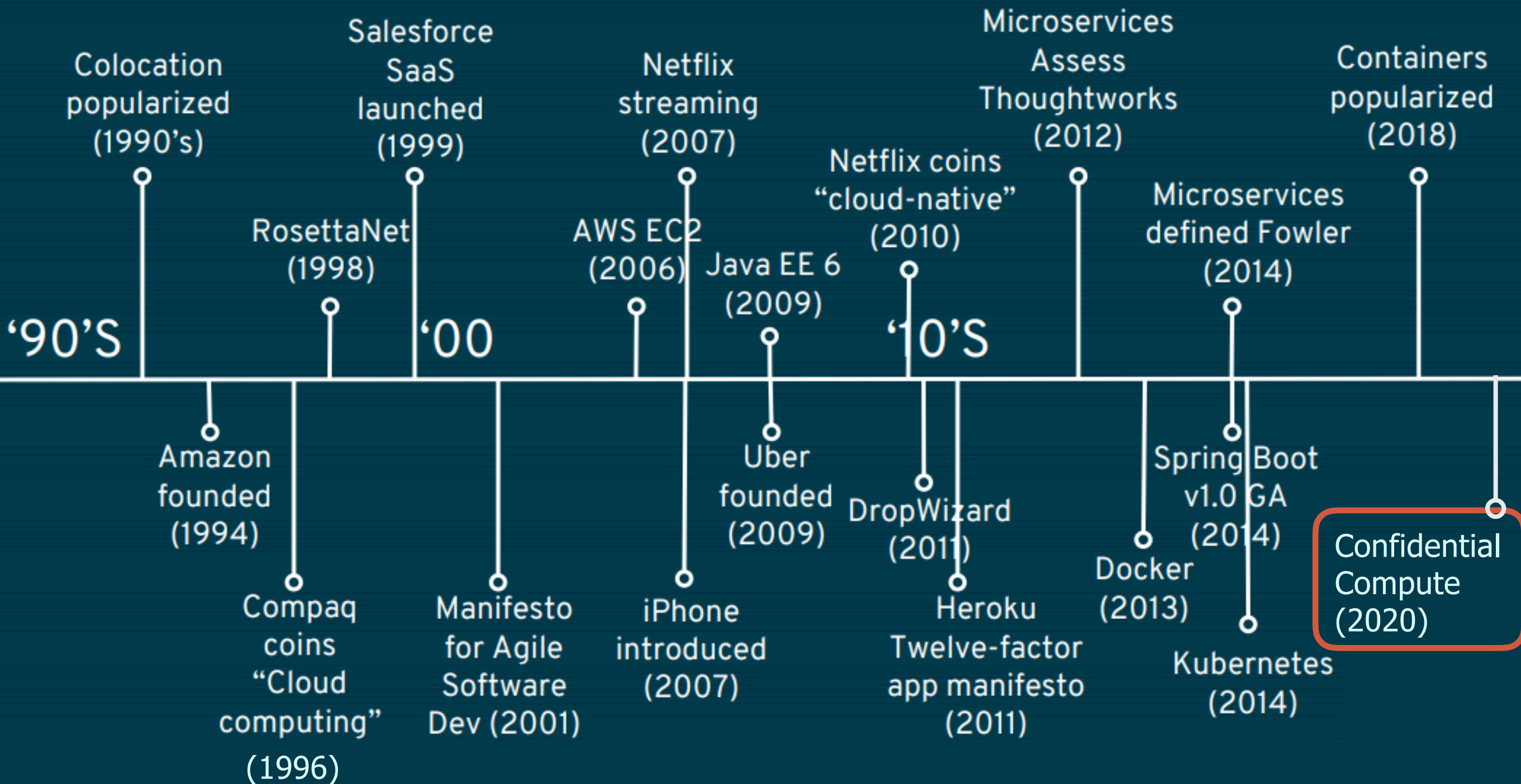


# BUILDING CONFIDENTIAL CLOUD-NATIVE APPLICATIONS WITH THE **SCONE PLATFORM**

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*[christof.fetzer@scontain.com](mailto:christof.fetzer@scontain.com)*  
*<https://scontain.com>*

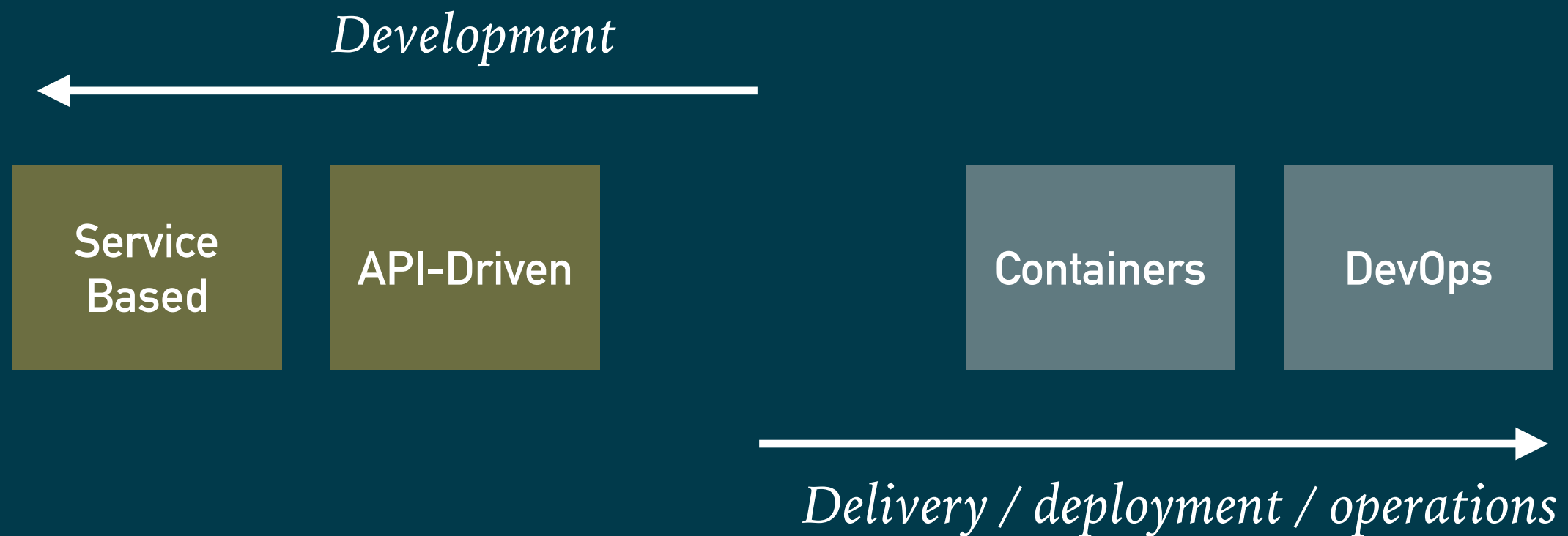
# 2020: CONFIDENTIAL COMPUTING



# CLOUD-NATIVE APPLICATIONS VS TRADITIONAL APPLICATIONS

	Cloud-Native Application Development	Traditional Development
Focus	Speed to market	Longevity & stability
Development Methodology	Agile development, DevOps	Waterfall, semi-agile development
Teams	Collaborative DevOps team	Isolated dev, operations, QA, and security teams
Delivery Cycle	Short and continuous	Long
Application Architecture	Loosely coupled, service-based, API-based	Tightly-coupled, monolithic
Infrastructure	Container-centric, portable, scales horizontally, on-demand capacity, on premise & cloud	Server-centric, infrastructure dependent, scales vertically, provisioned for peak capacity, on premise

# CLOUD-NATIVE APPLICATIONS



## *Cloud-Native Application*

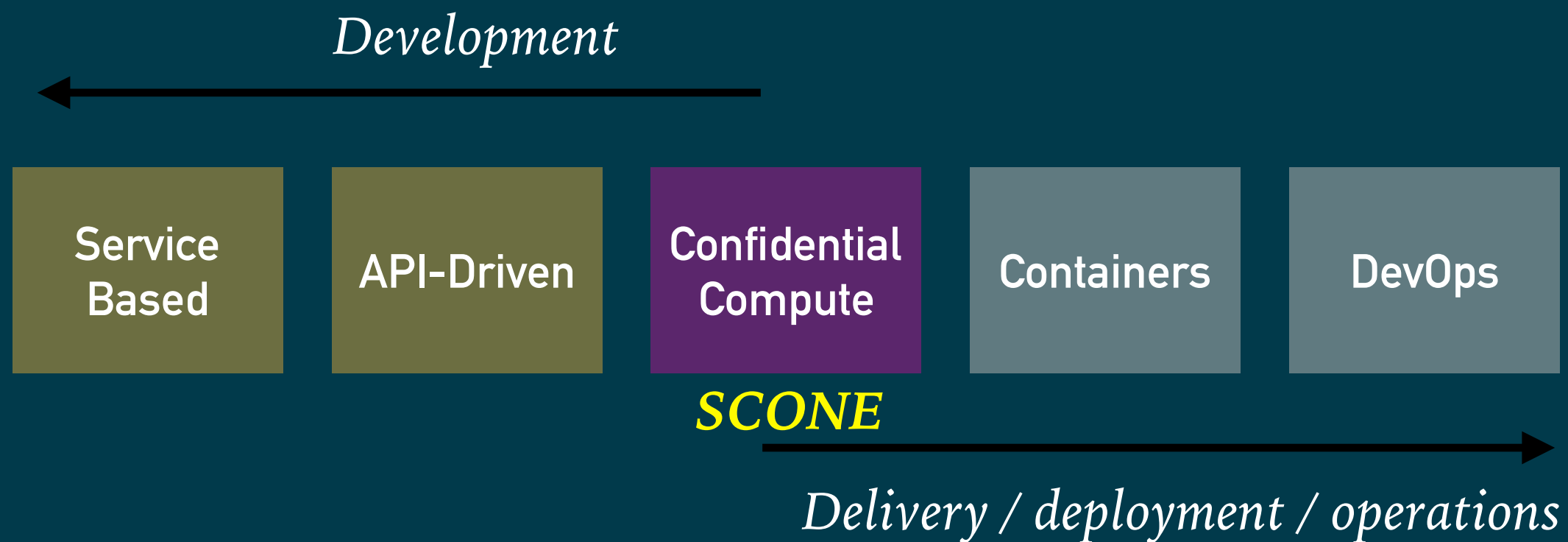
*- an application developed and operated using the cloud-native development/operation model*

# CONFIDENTIAL CLOUD-NATIVE APPLICATIONS

Confidential Cloud-Native Application Development	
Security	<b>data, code, and keys are always encrypted</b> <i>NEW</i> - at rest, in transit, in main memory -
Focus	Speed to market
Development Methodology	Agile development, DevOps
Teams	Collaborative DevOps team
Delivery Cycle	Short and continuous
Application Architecture	Loosely coupled, service-based, API-based communication
Infrastructure	Container-centric, portable, scales horizontally, on-demand capacity

# CONFIDENTIAL CLOUD-NATIVE APPLICATIONS

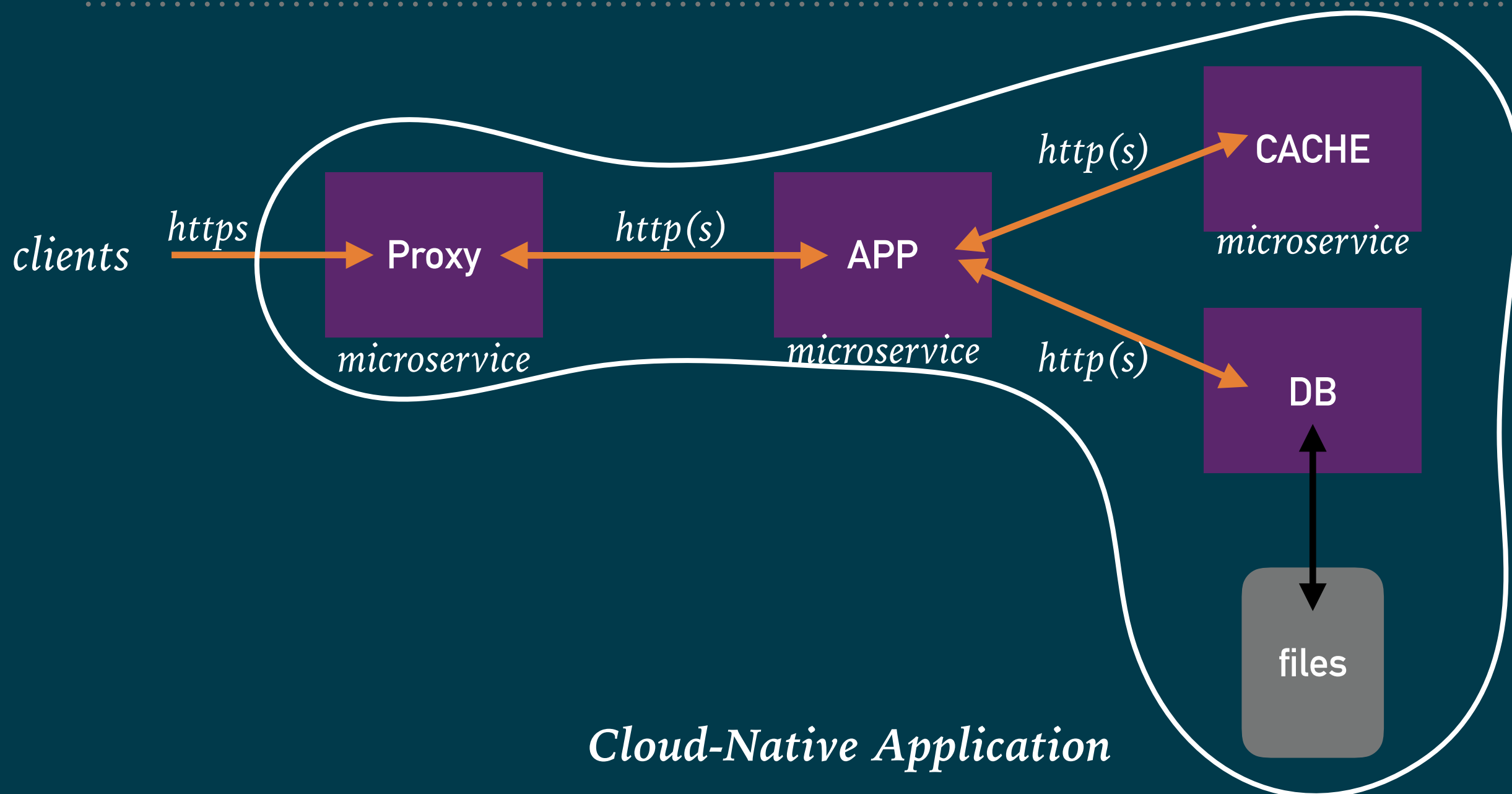
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## *Confidential Cloud-Native Application*

- cloud-native application*
- protect code, data and keys of application*

# CLOUD-NATIVE APPLICATION



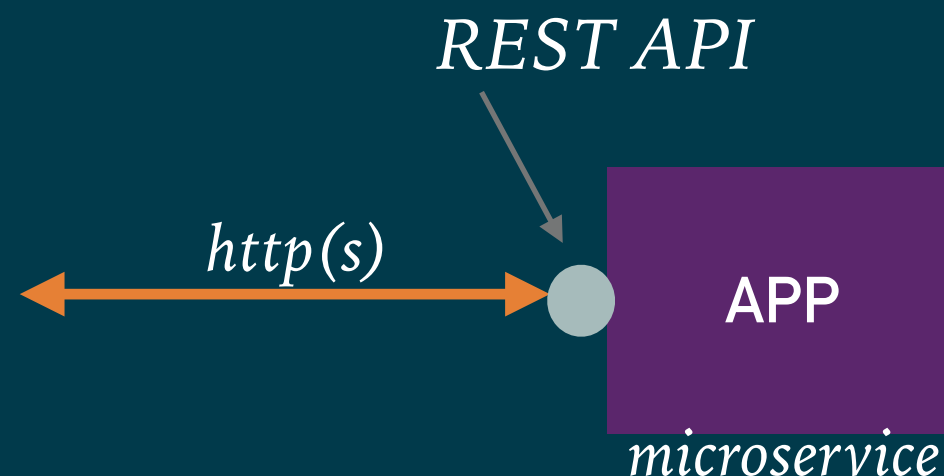
## Cloud-Native Application

- an application developed and operated using the cloud-native development/operation model



# MICROSERVICE

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## *microservice*

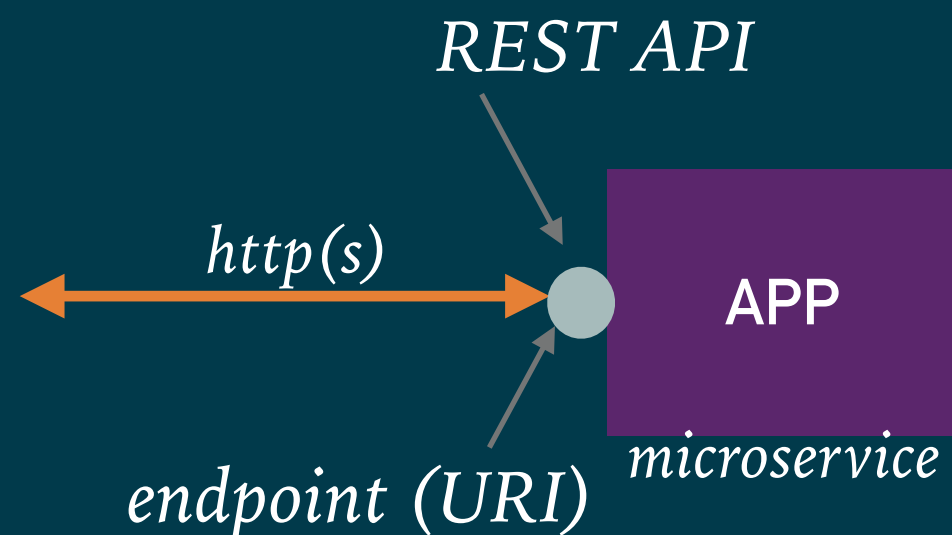
- *focus on a single aspect*
- *microservices are small, autonomous services that work together*

*Cloud-Native Application*



# REST API

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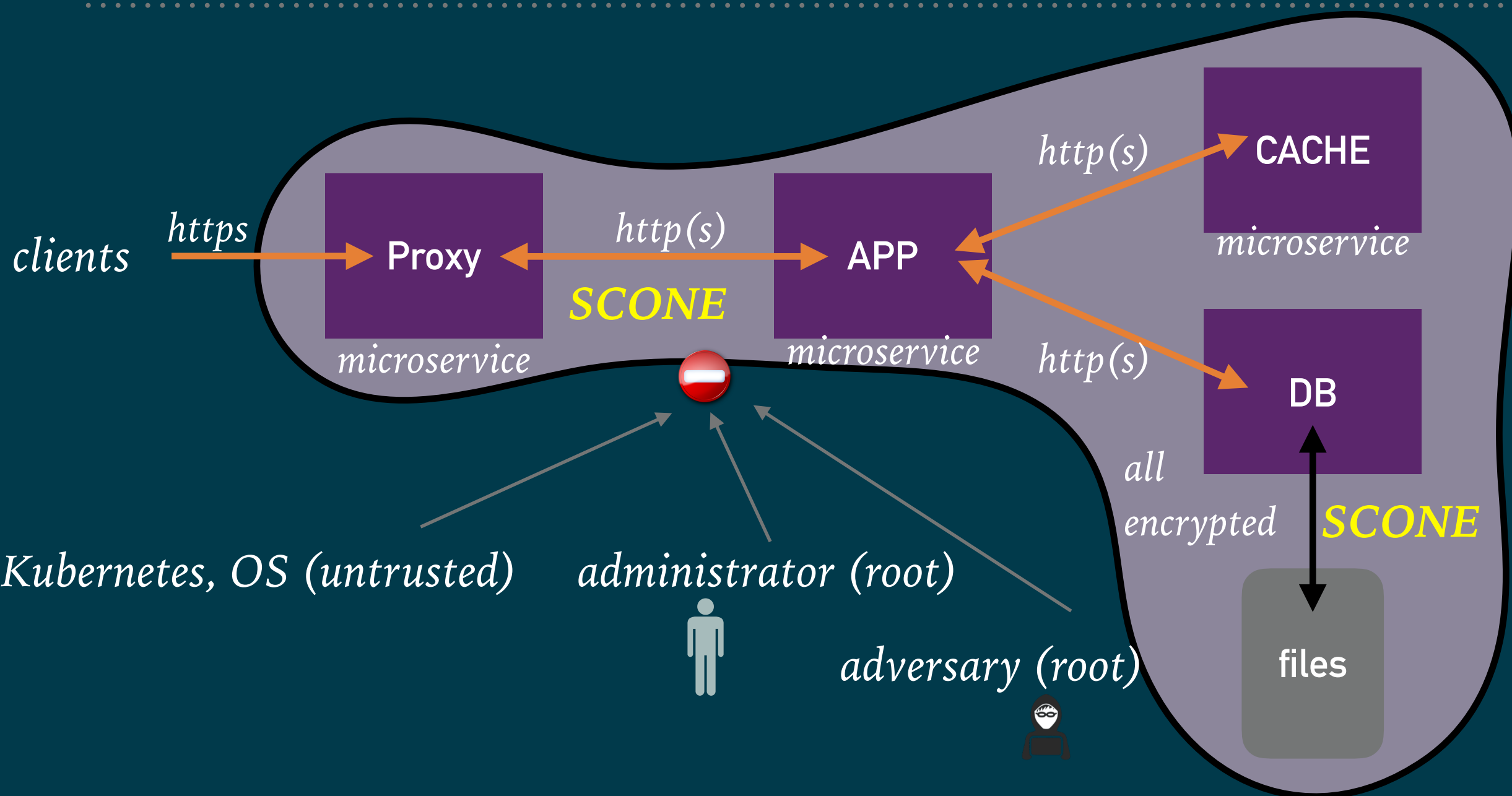
*Cloud-Native Application*

*REST = Representational state transfer*

## *REST API*

- *identify resource in request*
- *Fixed methods (from http):*
  - *Create (POST)*
  - *Retrieve (GET)*
  - *Update (PUT)*
  - *Delete (DELETE)*

# CONFIDENTIAL CLOUD-NATIVE APPLICATION



## Confidential Cloud-Native Application

- cloud-native application
- protect code, data and keys of application

# PROTECTION GOALS OF CONFIDENTIAL COMPUTE

## ➤ Protection of

- **Confidentiality:** information is not made available or disclosed to unauthorized individuals, entities, or processes
- **Integrity:** information cannot be modified by unauthorized individuals, entities, or processes
- **Freshness:** information cannot be replaced by old information by unauthorized individuals, entities, or processes

## ➤ Additional Protection goals:

- **Availability:** probability that information is available when it is needed
- **Durability:** probability that information will survive for one year

SCONE

Kubernetes, Ceph

# WHAT INFORMATION TO PROTECT?

---

- Protection of
  - Code, e.g., modern AI programs written in Python
  - Data, e.g., training data to create AI models
  - Keys, e.g., keys used to encrypt databases

# WHAT INFORMATION TO PROTECT?

---

- Protection of
  - Code, e.g., modern AI programs written in Python
  - Data, e.g., training data to create AI models
  - Keys, e.g., key used to encrypt database
- Example:
  - *MariaDB* supports encryption of database
  - encryption key is stored in configuration file
  - configuration file protected via access control:
    - i.e., can be read and written by *MariaDB* (user) as well as any root (=privileged) user

# IS ACCESS CONTROL SUFFICIENT?

---

- Sufficient if we define that
  - only **authorized user** can become root users

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- Sufficient if we define that
  - only **authorized user** can become root users
- What about
  - **adversary** gaining root access (e.g., stealing credentials)?
  - **authorized user** laid off -> could become an **adversary**?



# IS ACCESS CONTROL SUFFICIENT?

---

- Sufficient if we define that
  - only **authorized user** can become root users
- What about
  - **adversary** gaining root access: **must be addressed**
  - **insider attacks**: **must be addressed**
- **Solution:**
  - Use a **threat model** that gives adversary more power!

# THREAT MODEL

ADVERSARY HAS ROOT & HW ACCESS!

# WHY ASSUME ADVERSARY HAS ROOT ACCESS?

---

- Reasons:

- Legal:

- cloud provider might be legally required to provide access to the data

- Liability:

- too expensive to err on the threat model

- Limits of access control:

- How do we know that we can trust individual user?

- **Software complexity:** cannot assume that software is correct (see Defender's dilemma)

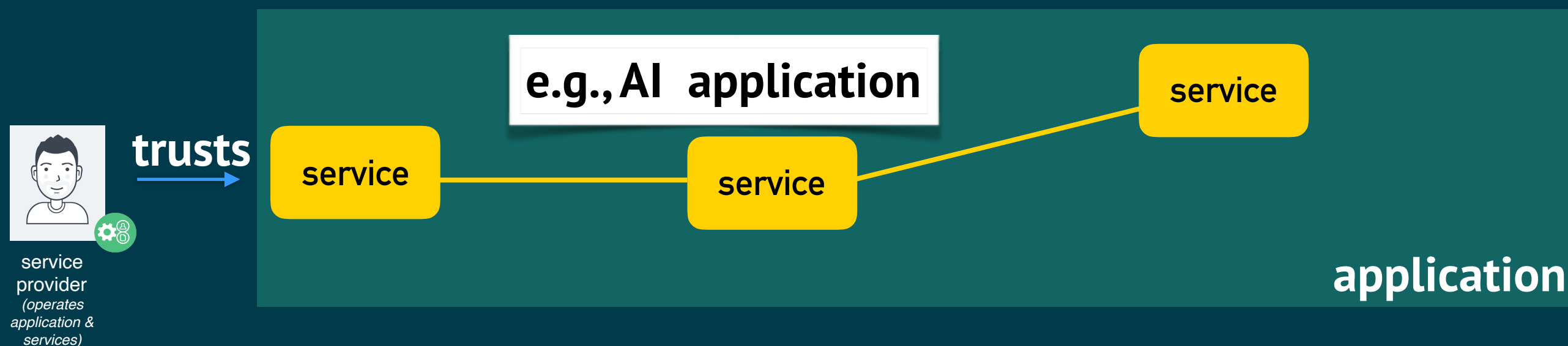
- **Hardware complexity:** cannot assume that hardware is correct (see BMC, firmware, ...)

- ...

# LEGAL REASON

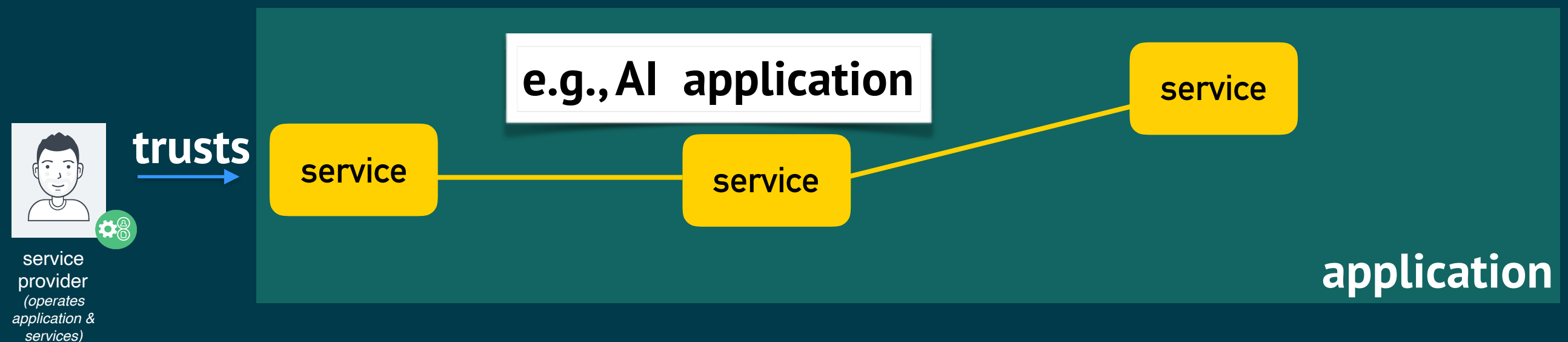
## Definition: TRUST

*we can justifiably assume the security, reliability, or ...  
of someone or something*



**service provider trusts that the services do what they supposed to do**

# APPLICATION-ORIENTED VIEW

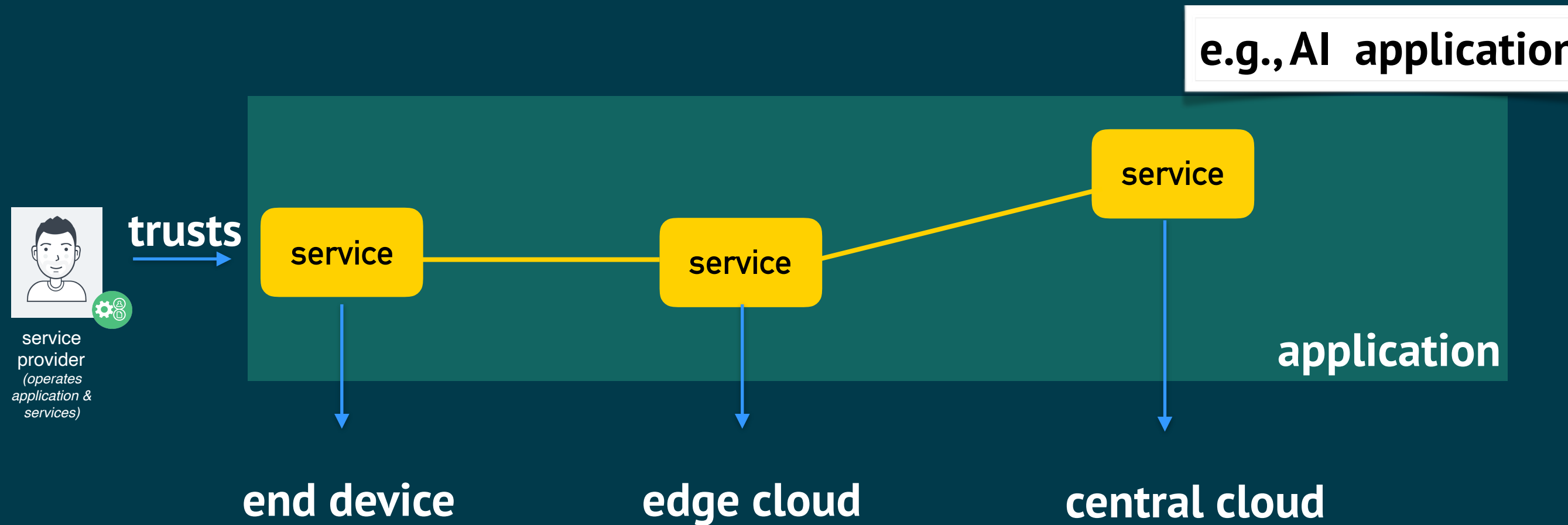


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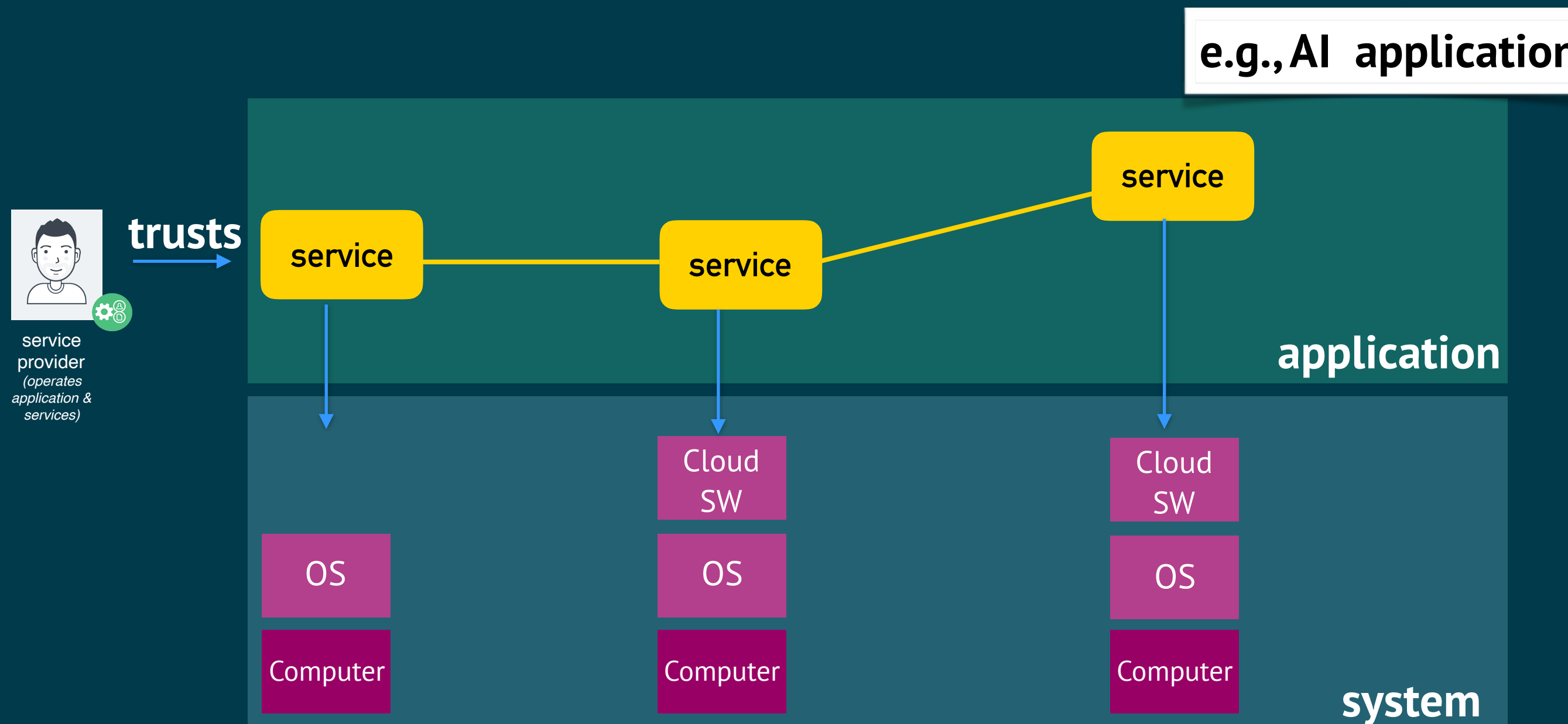
**Why?**

- **service provider controls the source code used in the services**

# APPLICATION-ORIENTED VIEW

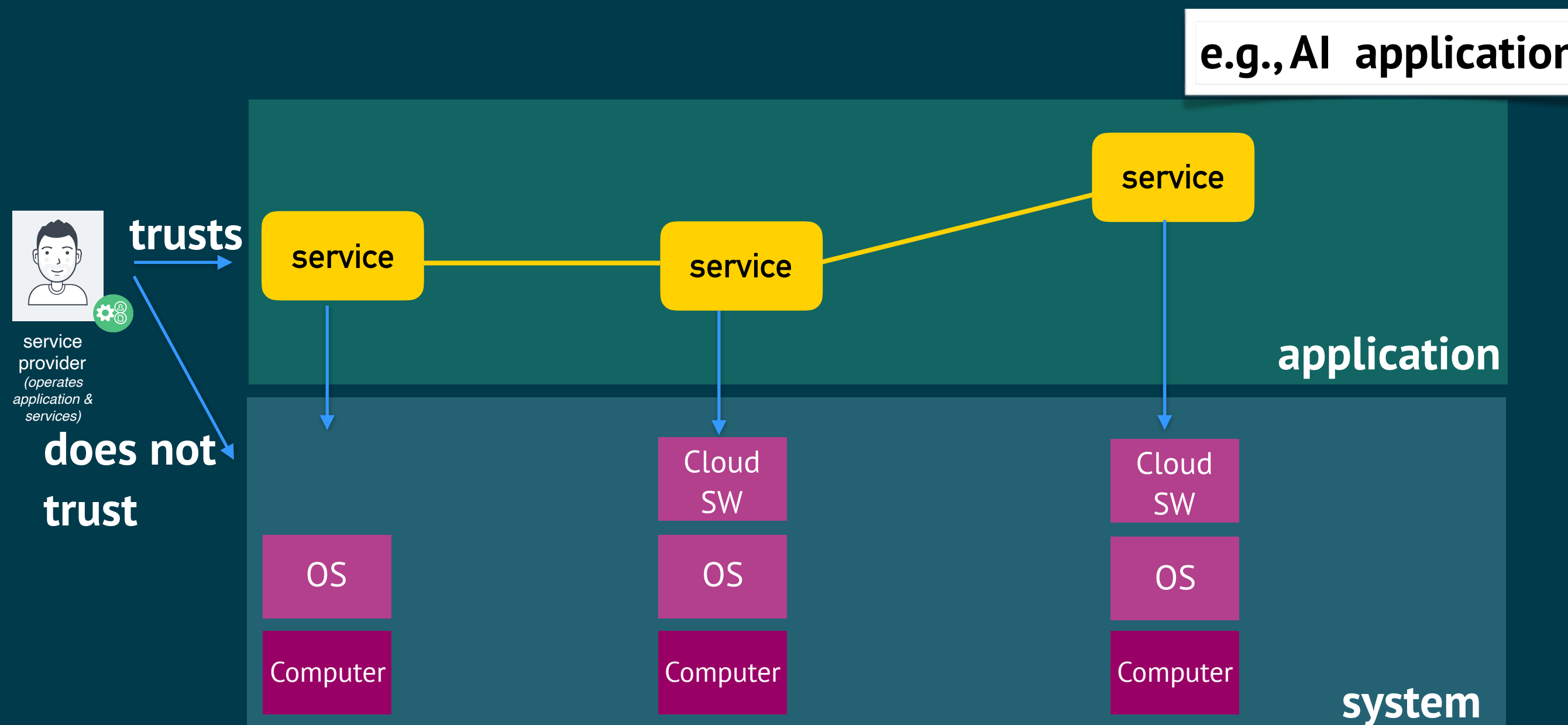


# APPLICATION-ORIENTED VIEW





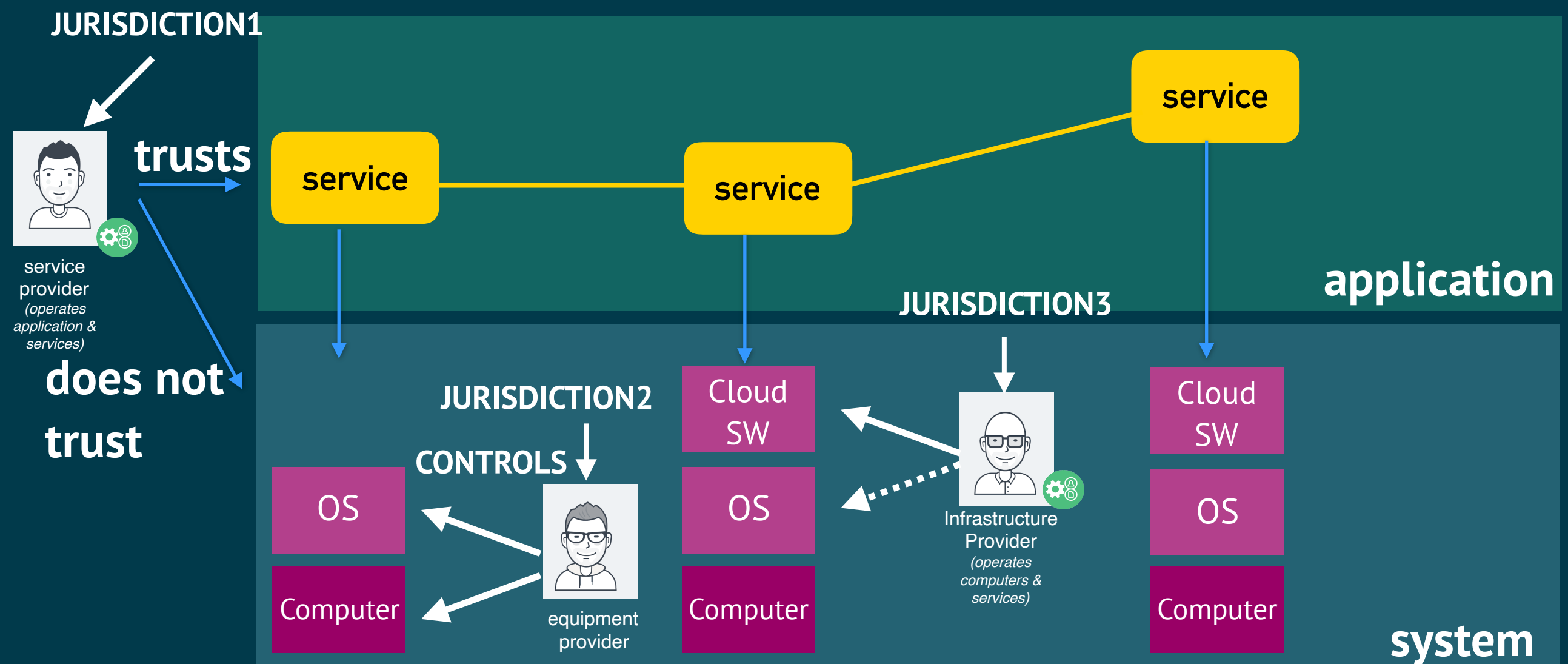
# PROBLEM: TRUST



# PROBLEM: TRUST

Why?

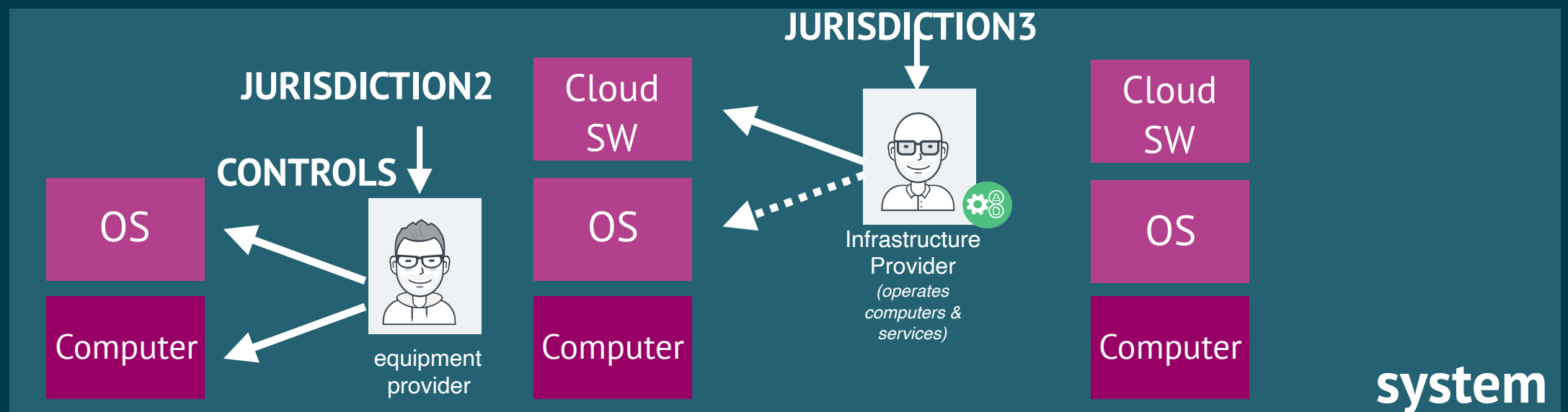
- service provider not in control of the system! Different jurisdictions.



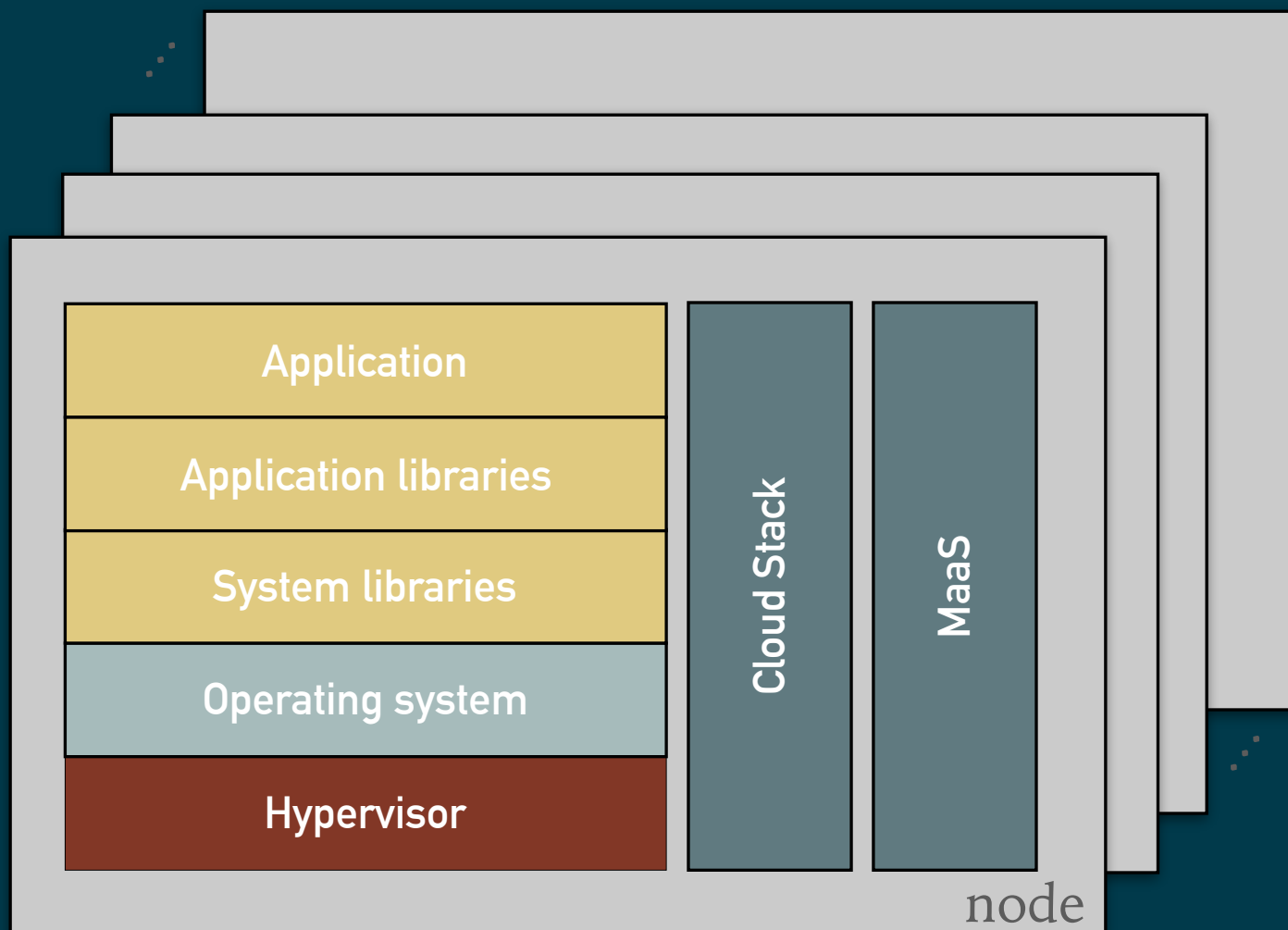
# TRADITIONAL APPROACH: SECURITY CERTIFICATION

## Why not depend on security certification?

- system too complex to understand the security!
- security certification typically shallow & historic software/firmware version



# SOFTWARE COMPLEXITY



cloud software stack

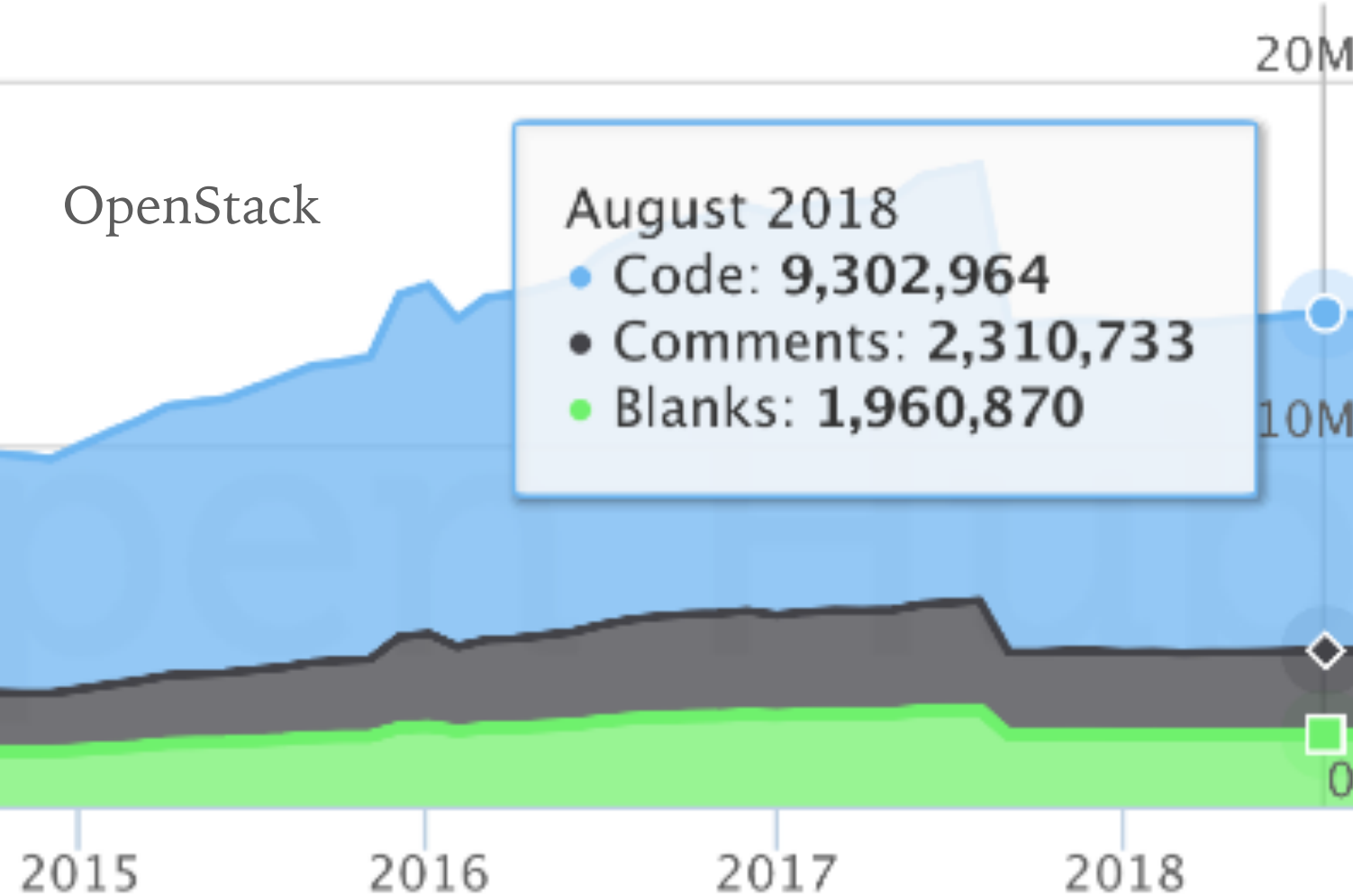
## DEFENDER'S DILEMMA

### ➤ Attackers:

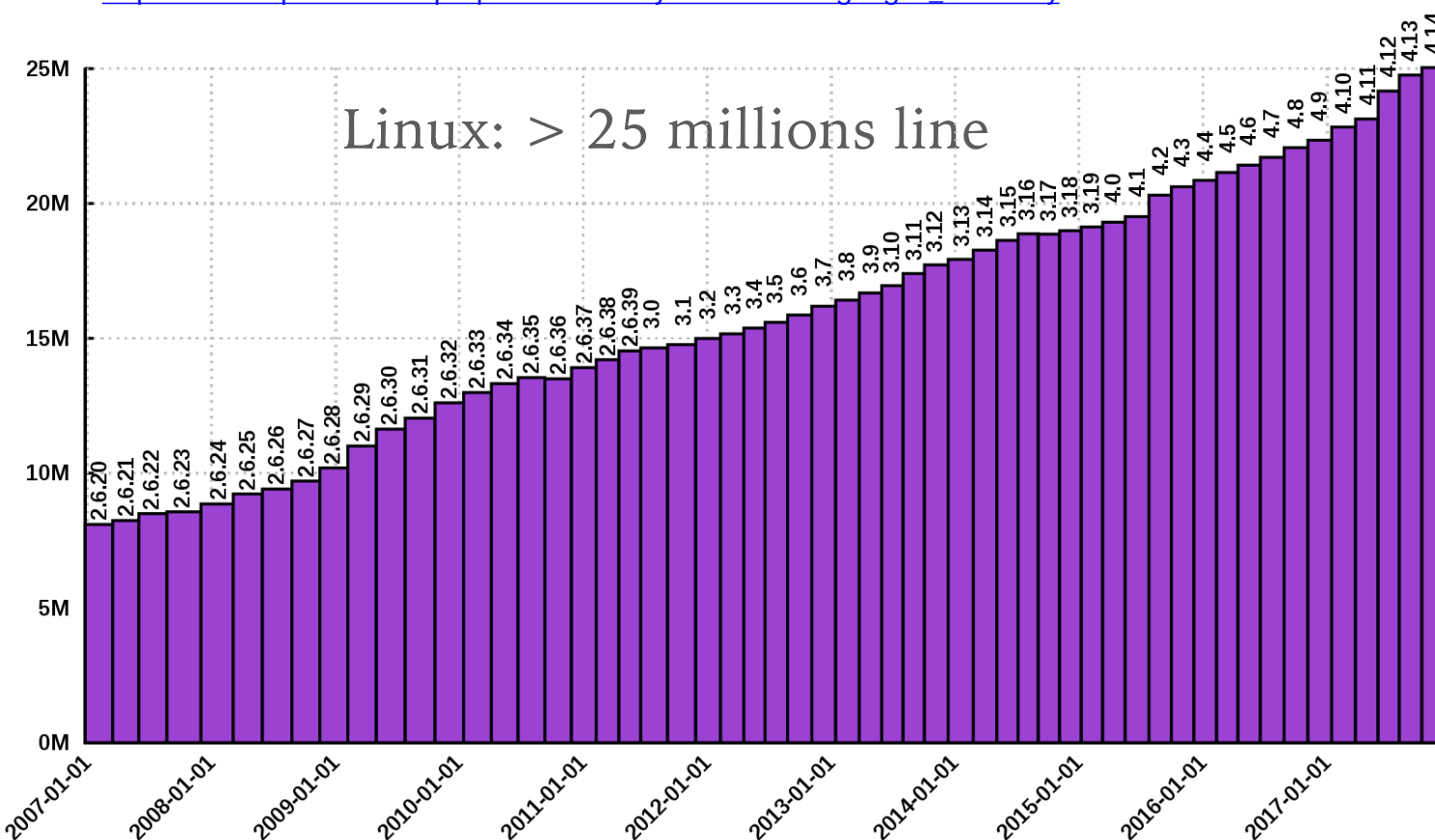
- success by exploiting a single vulnerability

### ➤ Defender:

- must protect against every vulnerability
  - not only in application
- millions of lines of source code
- not all known by service provider



[https://www.openhub.net/p/openstack/analyses/latest/languages\\_summary](https://www.openhub.net/p/openstack/analyses/latest/languages_summary)



# CLOUD SOFTWARE STACK

Applications run on top of software stack

➤ millions of lines of code

Cloud stack consists of

- VM/container engine
- operating system
- hypervisor
- node management service

# VULNERABILITIES

---

- **Coverity reports:**

- 1 defect per approx. 1700 lines of code

- **Kernel self protection project:**

- 500 security bugs fixed in Linux during the last 5 years
- each bug stayed about 5 years inside kernel

- **Xen hypervisor**

- 184 vulnerabilities (2012-2016)

[[http://www.cvedetails.com/product/23463/XEN-XEN.html?vendor\\_id=6276](http://www.cvedetails.com/product/23463/XEN-XEN.html?vendor_id=6276)]

- **Coverity:**

- quality of commercial software is not better than open source software

[KSPP] Kees Cook, The State of Kernel Self Protection Project, Linux Security Summit (LSS), 2016

[Coverity] Open Source Report 2014 - Coverity, [go.coverity.com/rs/157-LQW.../2014-Coverity-Scan-Report.pdf](https://go.coverity.com/rs/157-LQW.../2014-Coverity-Scan-Report.pdf)

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Bugs are especially bad in privileged software as it may result in unrestricted access to the system



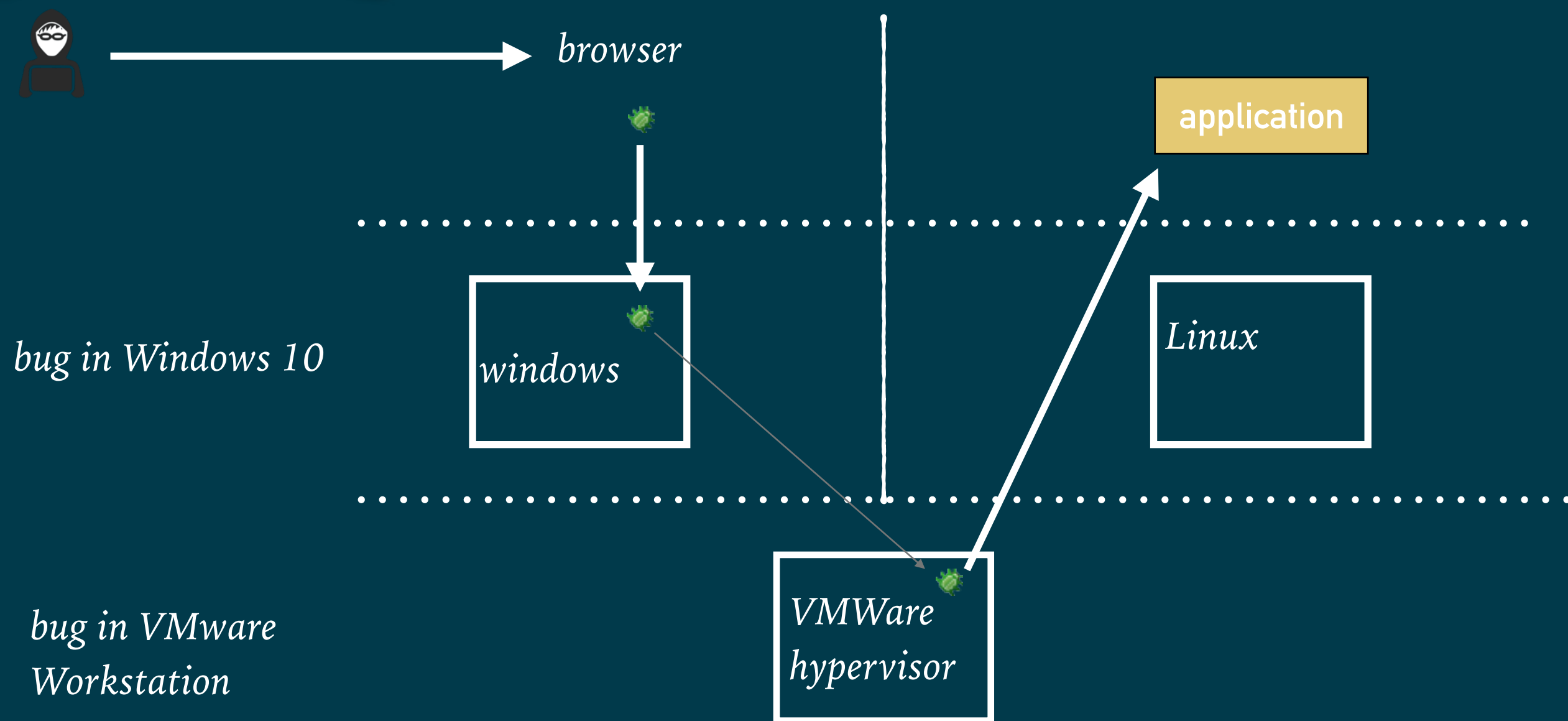
# HARDWARE PROTECTED MODE NOT SUFFICIENT

---

- Protected mode (rings) protects OS from applications, and applications from one another...
  - ... until a malicious applications exploits a flaw to gain full privileges and then tampers with the OS or other applications
- Applications not protected from privileged code attacks
- The attack surface is the whole software stack
  - Applications, OS, VMM, drivers, BIOS...

# ATTACKING VIRTUAL MACHINES

2017 hacking contest



<https://arstechnica.com/security/2017/03/hack-that-escapes-vm-by-exploiting-edge-browser-fetches-105000-at-pwn2own/>

# EXAMPLE: HEARTBLEED BUG



- Serious vulnerability in the popular OpenSSL cryptographic software library
  - Very widely used: apache/nginx (66% of Web servers), email servers, chat servers, VPN, etc.
- Buffer overrun when replying to a heartbeat message
- Allows anyone on the Internet to read the memory of the systems protected by the vulnerable versions of the OpenSSL software
  - The attacker can obtain sensitive data from server's memory: passwords, private keys, ...



# HARDWARE, FIRMWARE, MANAGEMENT FEATURES, ...

*Buggy on multiple levels...*

# CONFIDENTIALITY

---

*information is not made available or disclosed to unauthorized individuals, entities, or processes*

service

keys, data, code

**SECRETS IN  
MAIN MEMORY**

# CONFIDENTIALITY

*information is not made available or disclosed to unauthorized individuals, entities, or processes*

service  
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**SECRETS IN  
MAIN MEMORY**

RUNS ON TOP

**ANY ROOT USER**

**OS/HYPERVISOR HAS  
LIMITED CONTROL  
OVER BMC, IOMMU ...**

READ MAIN MEMORY

BMC, DMA  
READ MAIN MEMORY

OS

Computer



**current  
system**

**BMC, INTEL ME, ... : MANAGEMENT INTERFACE TO ACCESS  
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Computer

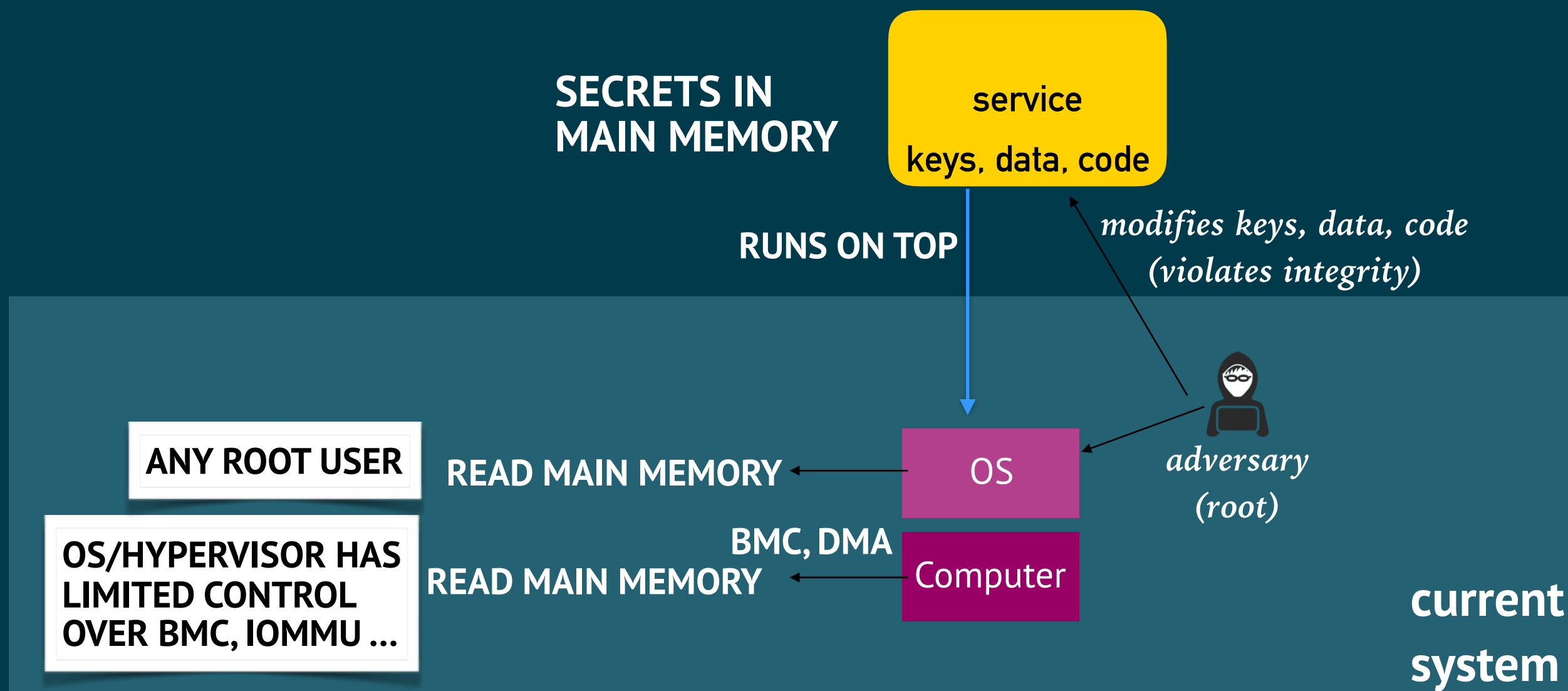
**current  
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**BMC, INTEL ME, ... : MANAGEMENT INTERFACE TO ACCESS  
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**AN APPLICATION CANNOT PROTECT CONFIDENTIALITY OF ITS SECRETS  
AGAINST ADVERSARIES WITH ROOT ACCESS (TRADITIONAL SYSTEMS)**

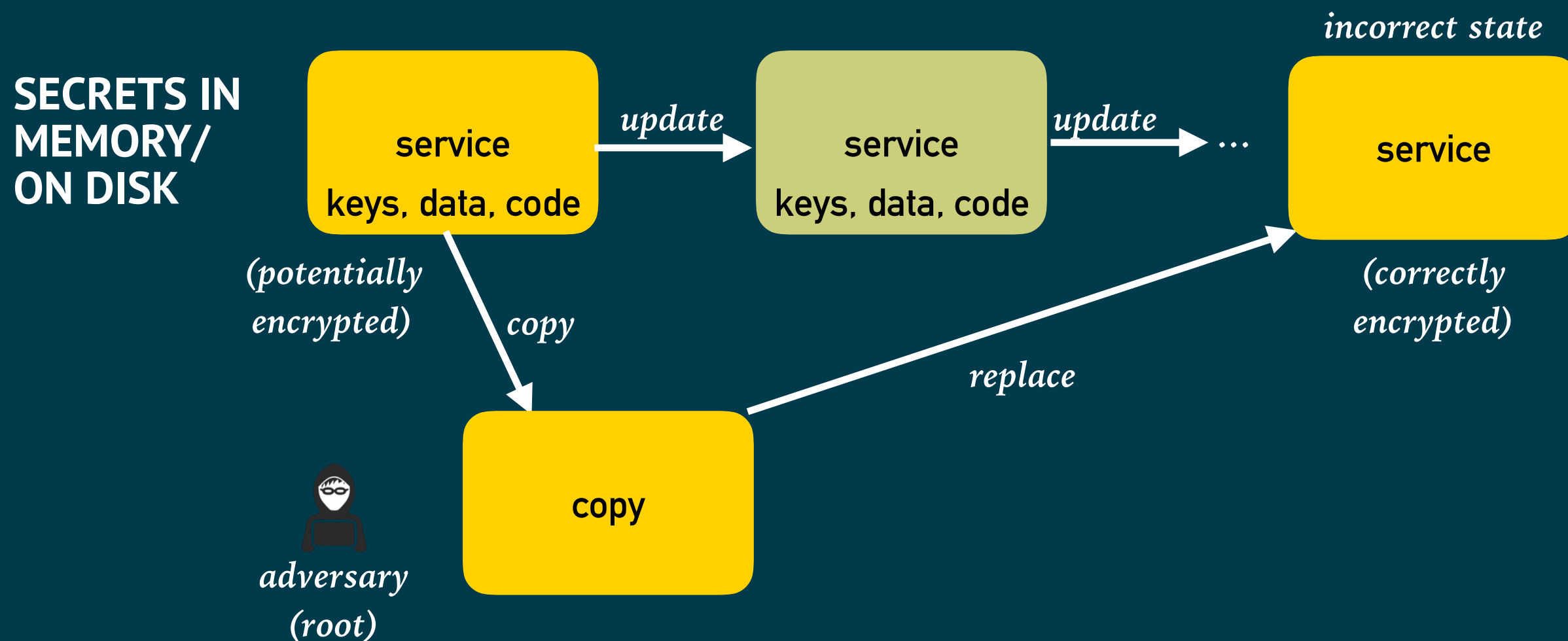


# INTEGRITY



**AN APPLICATION CANNOT PROTECT INTEGRITY OF ITS SECRETS AGAINST ADVERSARIES WITH ROOT ACCESS (TRADITIONAL SYSTEMS)**

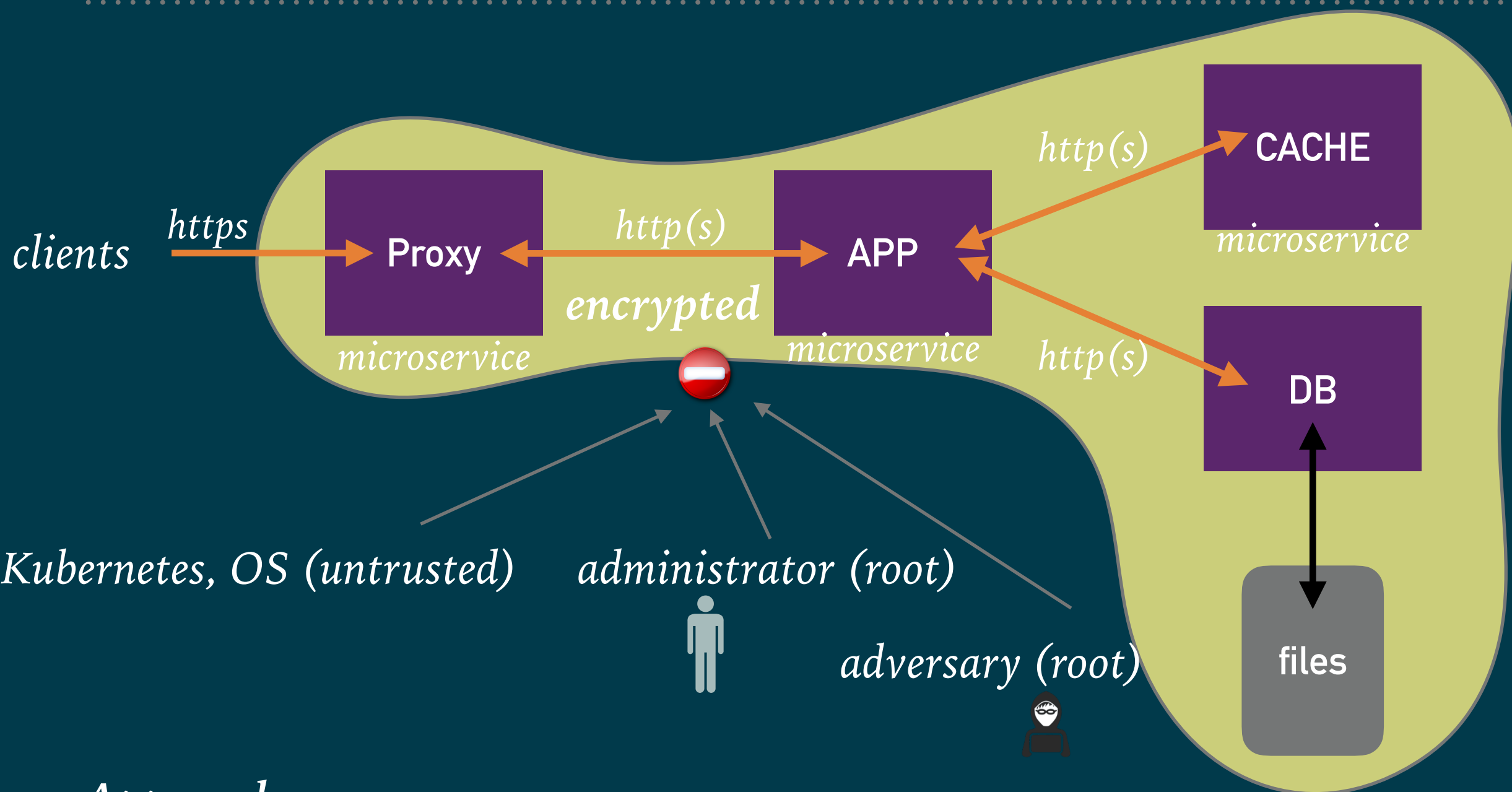
# FRESHNESS



*freshness: information cannot be rolled-back to an old state that has already been replaced*

**PROBLEM: HOW TO ENSURE  
CONFIDENTIALITY, INTEGRITY AND  
FRESHNESS IF ADVERSARY HAS  
ROOT AND HARDWARE ACCESS?**

# CONFIDENTIAL CLOUD-NATIVE APPLICATION



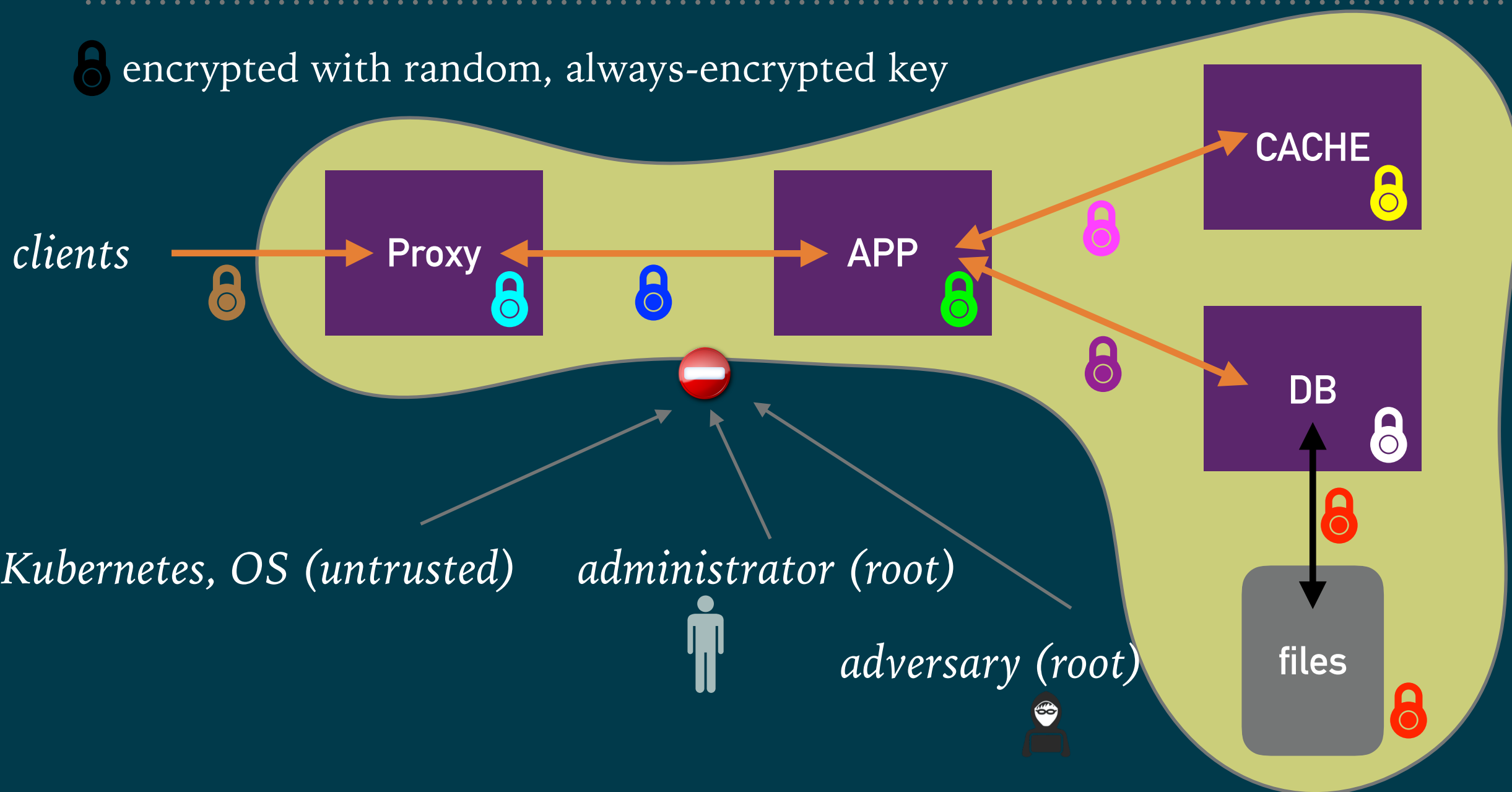
## Approach

- Encrypt everything (code, data, keys), and
- is never unencrypted: always encrypted at least

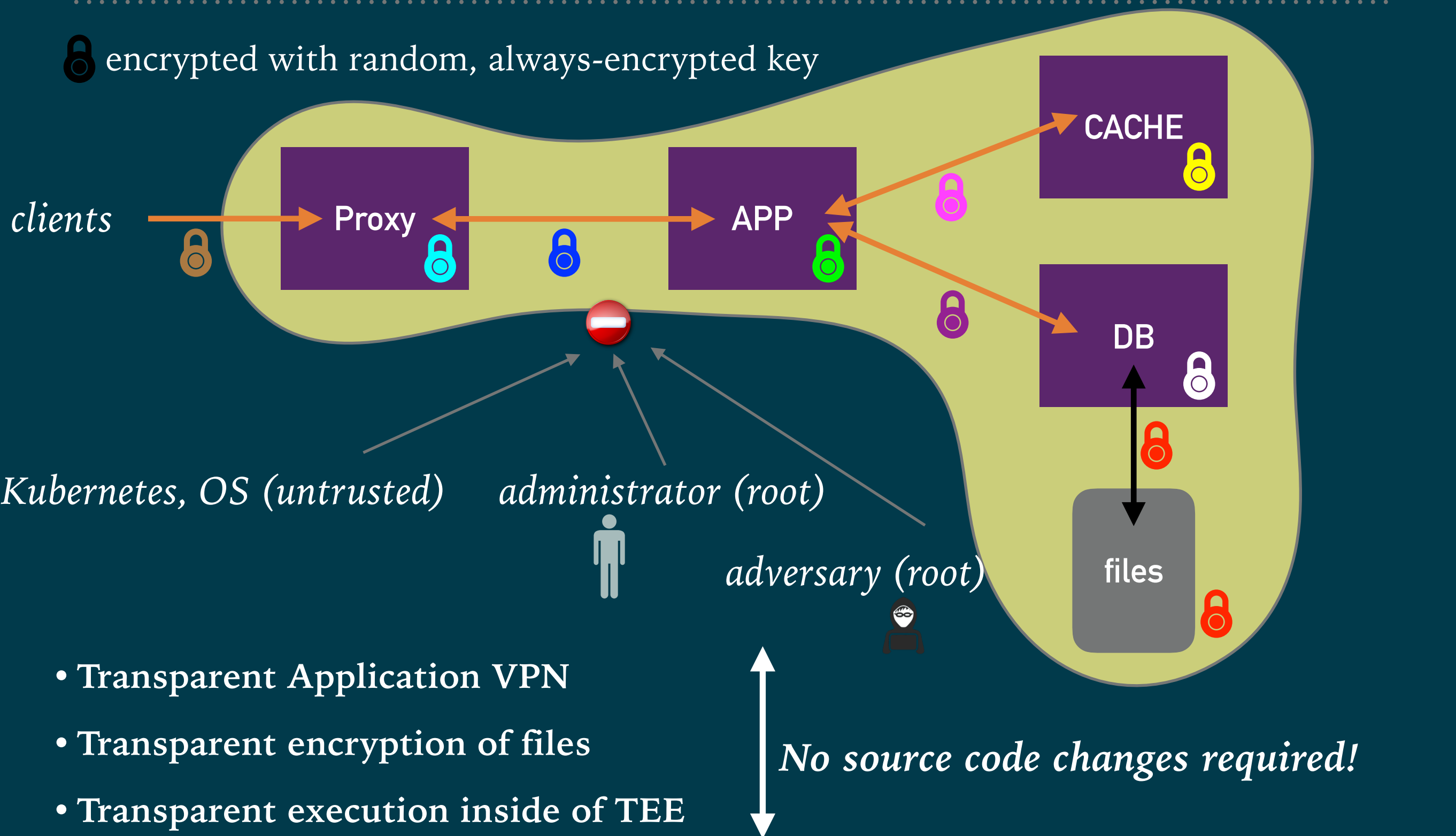
once

# SCONE ENCRYPTS DATA IN TRANSIT, AT REST AND IN MAIN MEMORY

🔒 encrypted with random, always-encrypted key

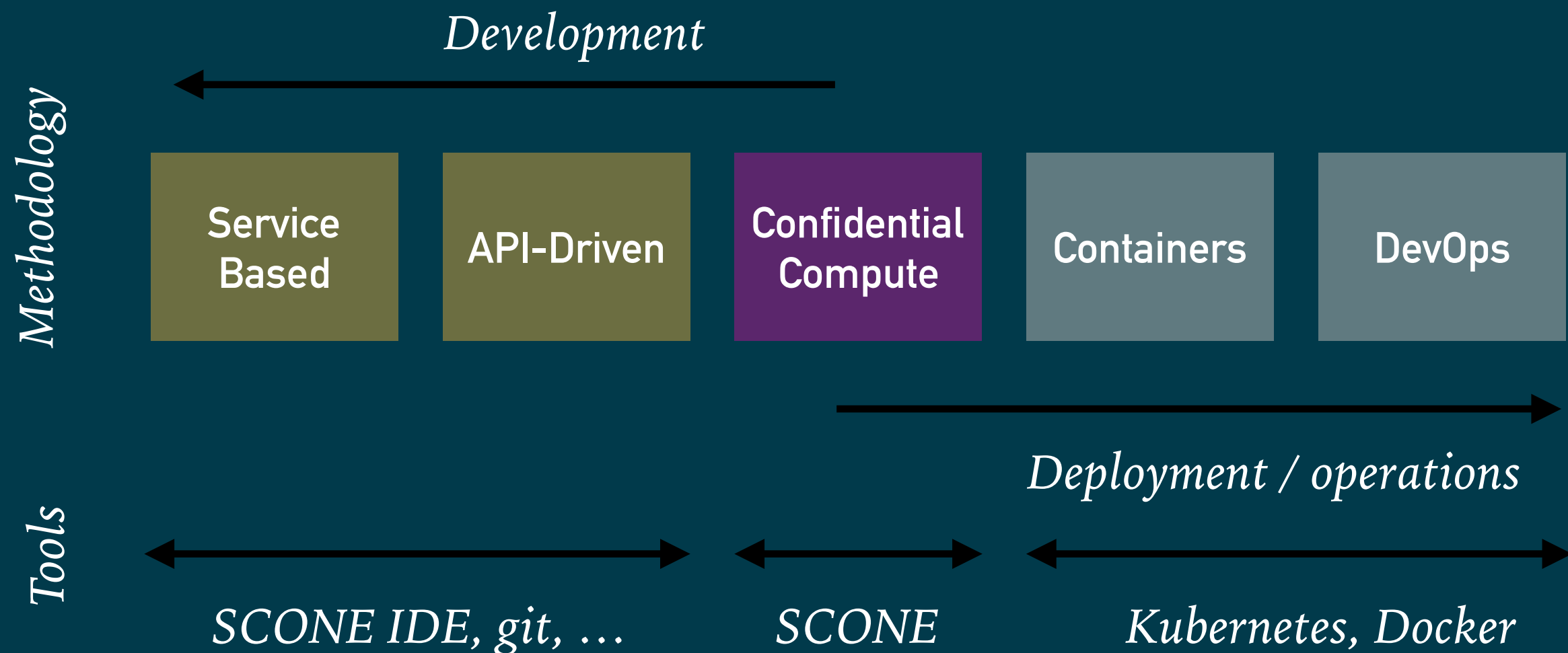


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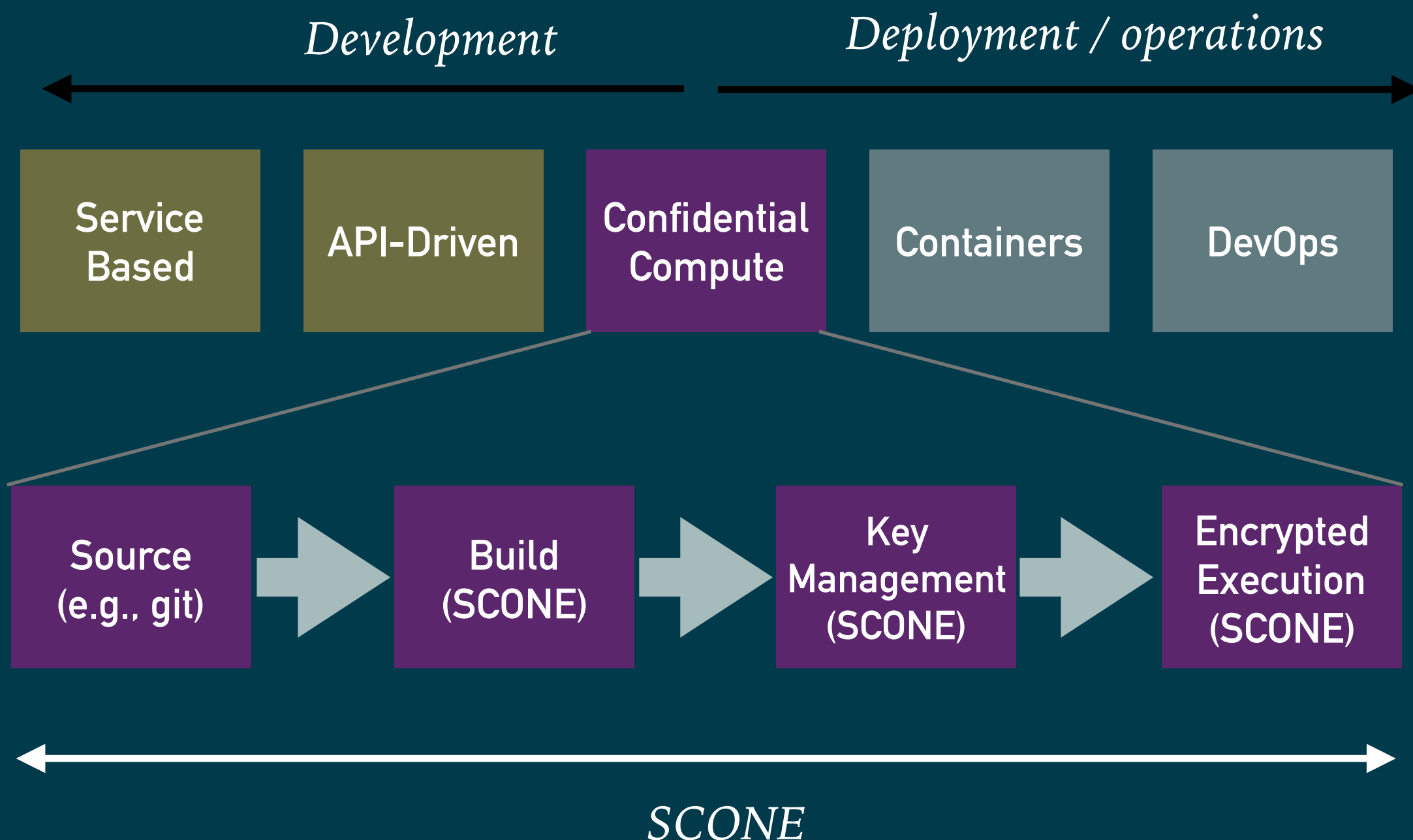
# HOW?

# CONFIDENTIAL CLOUD-NATIVE APPLICATION DEVELOPMENT

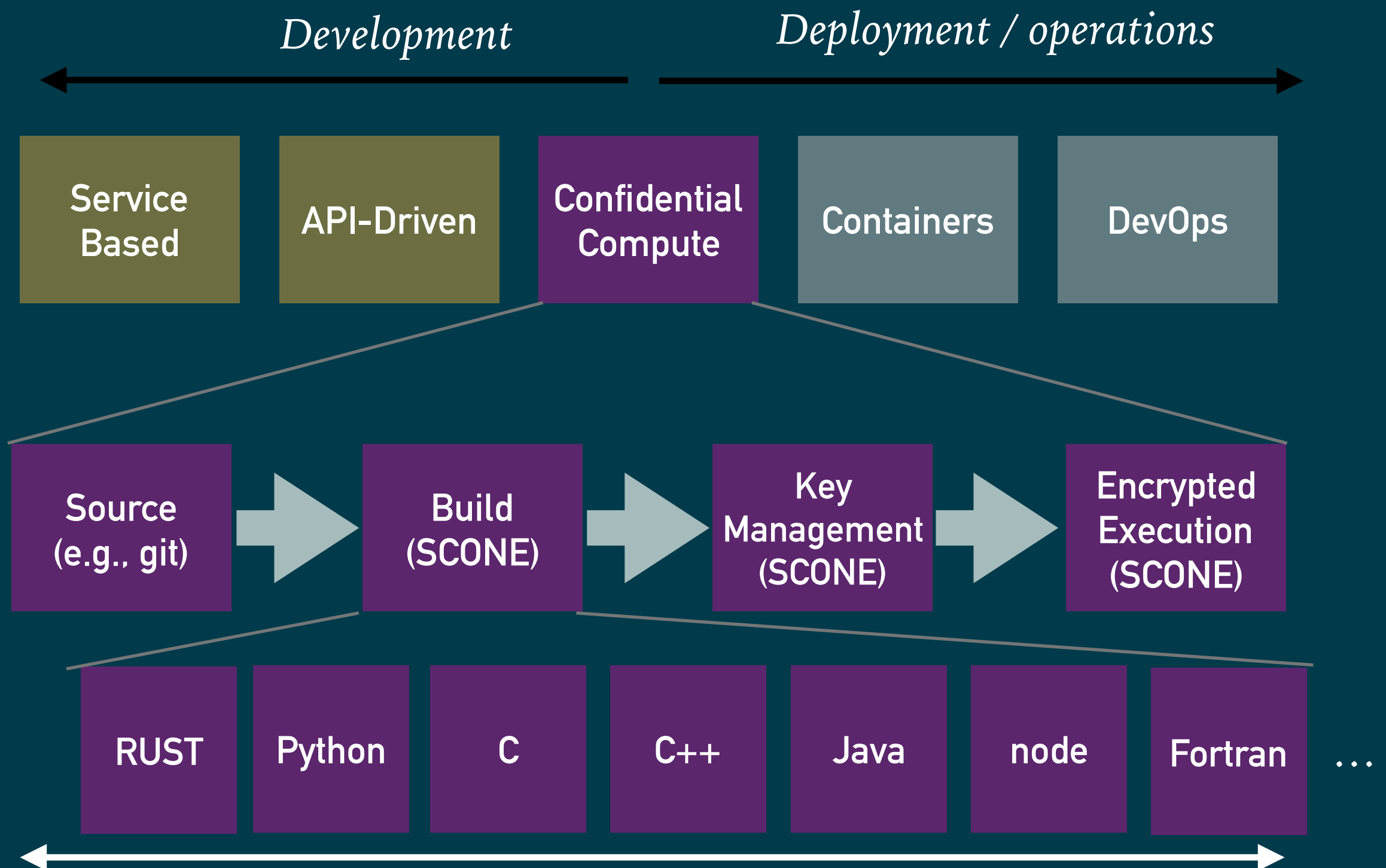




# CONFIDENTIAL CLOUD-NATIVE APPLICATIONS

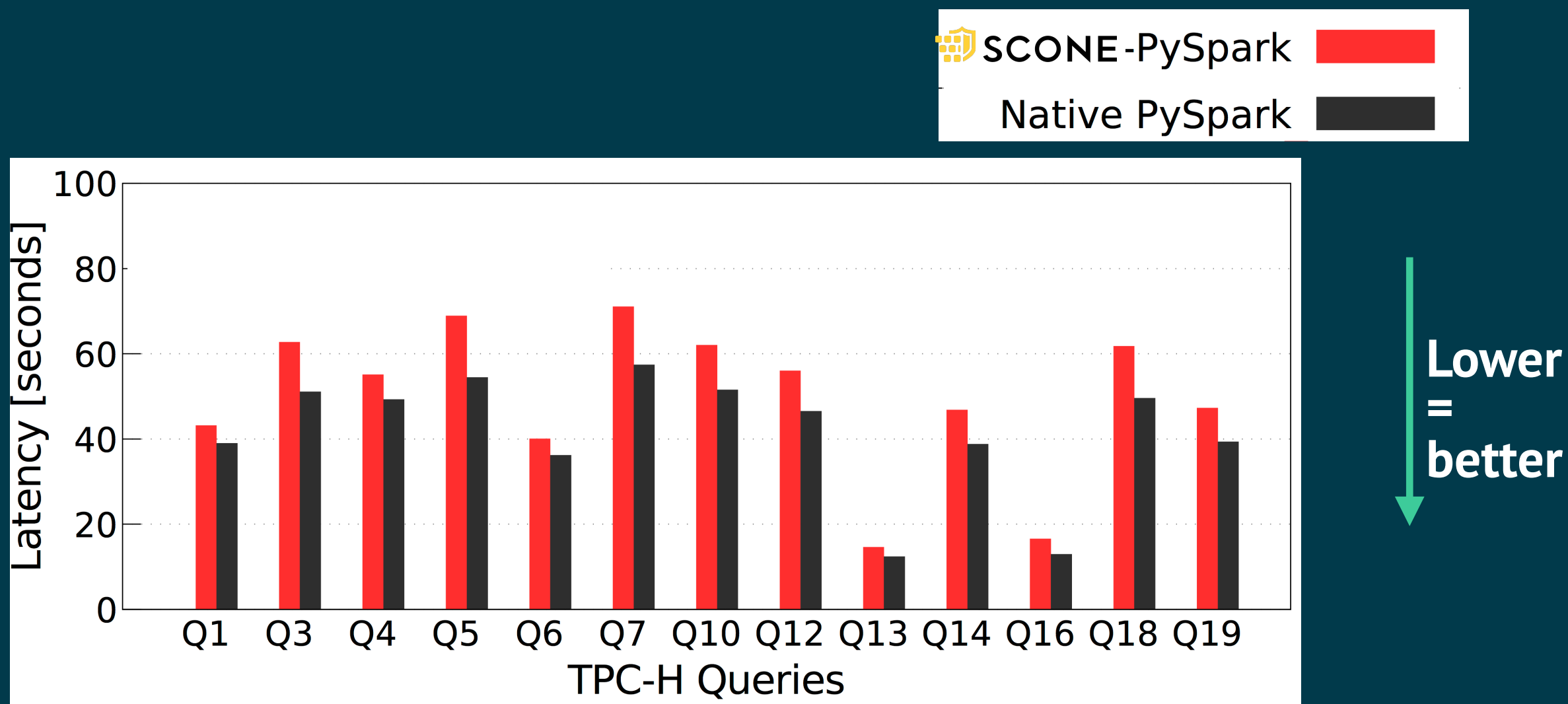


# CONFIDENTIAL CLOUD-NATIVE APPLICATIONS



# PERFORMANCE?

# PERFORMANCE



**< 22 %** overhead compared to native execution

# USE CASES?

# EHEALTH

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- we have built the Electronic Patient Record service
  - a confidential cloud-native application on top of Kubernetes
- highly scalable and efficient:
  - We scaled to 6000 parallel confidential microservices in a single Kubernetes cluster



# AI FRAMEWORKS

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- SCONE supports
  - TensorFlow
  - TensorFlow Lite
  - PyTorch
  - OpenVino
  - Scikit-Learn
  - ...

# CONCLUSION



SCONE supports the development & operation of **confidential cloud-native applications**



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<https://scontain.com>

<https://sconedocs.github.io/>

