# PETRINET SIMULATOR

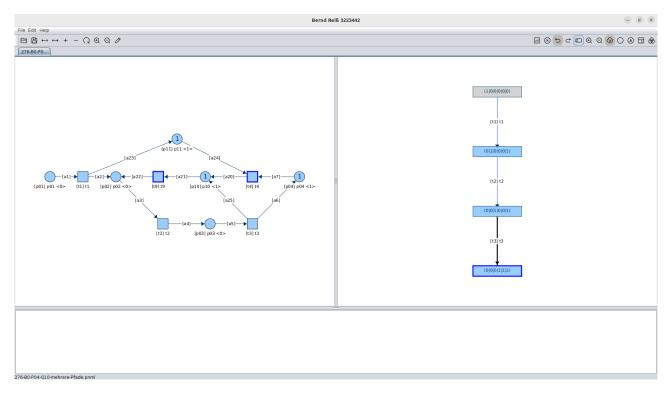
**Program Documentation** 

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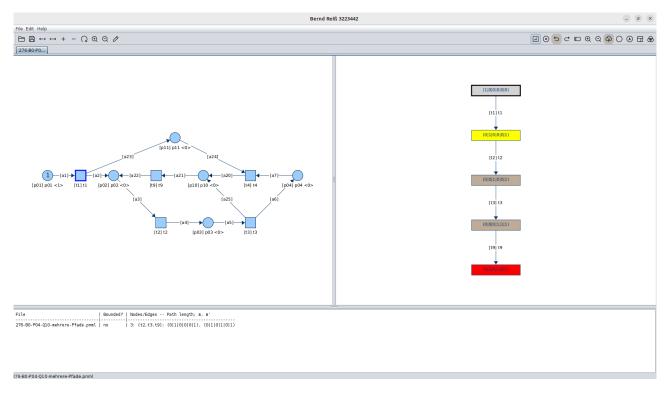
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### 1 Manual

The Petrinet Simulator is able to display a petrinet and gradually build a reachability graph by firing transitions using the GraphStream library<sup>1</sup>. The following picture shows the programs surface:



In the petrinet on the left activated transitions are marked bold. The current state is marked bold in the reachabiltiy graph on the right. Additionally the initial state is marked grey. After performing the boundedness analysis the graphs look like this:



The reachability graph shows whether it encountered a state that is unbounded and mark the beginning

<sup>1</sup>https://graphstream-project.org/

(yellow) and ending (red) nodes on the path signifying the unboundedness of the petrinet. Any states on the paths between these nodes and between the yellow and intial node are marked also (brown).

The following sections give an overview of additional features implemented in the program as well as of the menus and toolbars.

### 1.1 Additional features

**Automatic Boundedness Check** By default building the reachability check automatically checks whether the petrinet is bounded or not (see also 2.2). This can be turned off in the menu.

**Tabs** A new petrinet can be opened in a new tab. This can be done via the menu.

Undo/Redo The reachability graph provides undo/redo functionality.

Custom Layouts The Petrinet Simulator provides two custom layouts: tree layout and circle layout (see 2.3). Users can switch between these two and the automatic layout provided by GraphStream.

Editor Mode Petrinets can be edited in an editor mode (see 1.3.2).

### 1.2 The Menu

### 1.2.1 File

The File menu consists of the following entries:

- New: opens a new empty file for editing  $\rightarrow$  automatically opens the Editor view (see 1.3.2).
- Open: opens saved file from directory.
- Open in new Tab: same as open but in a new tab.
- **Reload:** reload the currently opened file.
- Save: save changes.
- Save as: save changes and choose directory/name.
- Analyse++: analyse many petrinets at once. Results are printed in the text area.
- Close: close currently opened file.
- Exit: close the program.

### 1.2.2 Edit

The *Edit* menu consists of the following entries:

- Open Editor: opens the Editor.
- Close Editor: closes the Editor and reverts back to Viewer.
- Change Look and Feel: change between Nimbus and Metal look and feel.
- Enable Automatic Boundedness Check: automatically check for boundedness when building reachabilty graph.
- **Disable Automatic Boundedness Check:** do noe automatically check for boundedness when building reachabilty graph.

### 1.2.3 Help

The *Help* menu only consists of the element *Info*, which shows information about the Java version used and the user directory.

### 1.3 Toolbar Modes

The program has two distinct toolbar modes: viewer mode and editor mode. This section discusses the two modes and toolbar buttons associated with it. In general the toolbar is fully de- and reattachable (except for the bottom of the program being reserved for the status bar). All icons for the toolbar buttons are taken from https://iconoir.com/.

#### 1.3.1 Viewer Mode

The viewer mode provides functionality for viewing petrinets and building the reachability graph. The toolbar is split into two parts: the petrinet toolbar to the left and the reachability graph toolbar to the right. Below the specifics are discussed.

### Petrinet Toolbar $\square \hookrightarrow + - \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc$

From left to right the buttons provide the following functionality:

- $\square$  Load: see Menu File 1.2.1.
- Bave: see Menu File 1.2.1.
- $\leftarrow$  **Previous:** open the previous file in the current directory  $\rightarrow$  do nothing if current file is the first.
- $\rightarrow$  Next: open the next file in the current directory  $\rightarrow$  do nothing if current file is the last.
- + Increment Place: increment a marked place by one token.
- - Decrement Place: decrement a marked place by one token.
- Q Reset: reset the petrinet graph to initial state → if place was in-/decremented the initial state is being updated.
- ② **Zoom in:** zoom in into petrinet graph.
- Q **Zoom out:** zoom out of patrinet graph.
- **Open Editor:** open the Editor (see 1.3.2).

# Reachbility Graph Toolbar

From left to right the buttons provide the following functionality:

- ■ Analyse Petrinet: analyse the currently opened petrinet → the reachability graph is built and a pop up window shows whether the petrinet is bounded or unbounded. If the petrinet is unbounded the starting and ending nodes of the path signifying unboundedness are marked yellow and red in the reachability graph.
- Seset: delete the reachability graph except for the initial state (marked gray). Also resets the petrinet.
- $\hookrightarrow$  **Undo:** undo the last step in the reachability graph. If last step way analysis it is treated as one single step.
- $\supseteq$  **Redo:** redo the last step in the reachability graph. If last step way analysis it is treated as one single step.

- Clear text area: clears the text area.
- ② Zoom in: zoom in into reachability graph.
- Q **Zoom out:** zoom out of reachability graph.
- • Tree layout: organizes the reachability graph in a tree layout (see 2.3.1).
- O Circle layout: organizes the reachability graph in a circle layout (see 2.3.2).
- ② Automatic Layout: uses the auto layout by GraphStream to organize the reachability graph.
- 🗆 Reset Split Panes: resets the split panes to the default ratio.
- © Change Look and Feel: changes between the different look and feels (*Metal* or *Nimbus*).

### 1.3.2 Editor Mode

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From left to right the buttons provide the following functionality:

- $\boxminus$  Load: see Menu File 1.2.1.
- 🖺 Save: see Menu File 1.2.1.
- • Previous: see Petrinet Toolbar in Viewer Mode (1.3.1).
- • Next: see Petrinet Toolbar in Viewer Mode (1.3.1).
- + Increment Place: see Petrinet Toolbar in Viewer Mode (1.3.1)...
- - Decrement Place: see Petrinet Toolbar in Viewer Mode (1.3.1)...
- ② **Zoom in:** see Petrinet Toolbar in Viewer Mode (1.3.1)...
- $\bigcirc$  **Zoom out:** see Petrinet Toolbar in Viewer Mode (1.3.1)...
- ① Add Place: adds a place above the left top most element.
- 🗷 Add Transition: adds a transition above the left top most element.
- **Remove Element:** removes the marked element (place or transition).
- **Add Edge:** choose beginning and ending node of an edge to add it → beginning and ending nodes have to be of different type (place/transition).
- $\cap$  Remove Edge: choose beginning and ending node of an edge to remove it  $\rightarrow$  the edge has to exist.
- $\square$  Add Label: add label to marked place or transition.
- © Close Editor: closes the Editor and opens the Viewer (see 1.3.1).

### Reachability Graph Toolbar

From left to right the buttons provide the following functionality:

- Example Reset Split Panes: see Reachability Graph Toolbar in Viewer Mode (1.3.1).
- © Change Look and Feel: see Reachability Graph Toolbar in Viewer Mode (1.3.1).

### 2 Design of the Program

### 2.1 General Structure

Figure 1 shows the general control flow of the program. Double rules signify a GUI component, dashed lines signify interfaces / listeners. The figure only demonstrates the basic design of the program and how the essential data models, controllers and GUI components communicate with each other. Other parts of the program are included in the text as well. The following sections give a detailed description of how the data models are structured and how controllers and listeners work. On a general level it can be noted controllers control the input provided via the GUI elements and listeners transfer changes in the data model to the GUI components.

### 2.1.1 Data Models

**Petrinet Model** It handles everything belonging only to the petrinet. It has functions for modifying the petrinet (such as adding or removing elements) and checks whether an action is valid. If an action is not valid it throws an according *PetrinetException*. The petrinet model holds helper data structures *Place* and *Transition*. There is no data structure for arcs, as they are fully defined via the places and transitions connected to them<sup>2</sup>. Rather arcs are defined by lists of elements preceding and succeeding places and transitions. Places contain information about the number of tokens stored in them. Transitions have a status of being activated or not activated.

Reachability Graph Model It provides methods for adding/removing states to/from the reachability graph and giving information about the graph. It contains a list of *PetrinetStates*. Such a state contains predecessor and successor states, information about the tokens on the places in the petrinet as well as a state called m, being the first state on a path that marks the given state as unbounded. When a new state is added the model checks whether the new state completes a path signifying unboudedness (see 2.2. If so, the state at the start of the path is remembered by the state and the reachability graph model remembers the last state on the path. The states also store transitions leading up to them and from them. These are stored in lists in a map.

Reachability Graph Undo Queue This queue records all the steps taken in the reachability graph and keeps track of the actions being taken (i.e. whether something has been added). To this end it has its own data structure called ReachabilityGraphUndoQueueState. The queue can go back or forward or can be rewinded/reset. If it goes backward, alls steps are undone. Going forward means redoing all steps. The reachability model has methods to stop pushing steps onto the queue. Furthermore steps can be recorded as "skippable". A step being skippable means on un-/redo it does not stop but continues going back-/forwards. This enables a multi step process like adding steps when analysing the petrinet to be treated as one big step.

### 2.1.2 Controllers

There are three controllers in the program: the main controller, the petrinet viewer controller and the petrinet editor controller. Additionally the *PetrinetPanelInterface* serves as an intermediary between the main controller and the other controllers.

Main Controller On a general level there is a main frame and a number of petrinets in tabs. The task of the main controller is to pass user interactions with elements belonging to the main frame (menus and toolbars) to the petrinet panel that is currently active. If tabs are switched it is also the task of the main controller to restore the state of the toolbar with regards to highlighted buttons for the active petrinet or switching to editor/viewer mode (see 1.3).

<sup>&</sup>lt;sup>2</sup>For storing ids belonging to arcs a map is used.

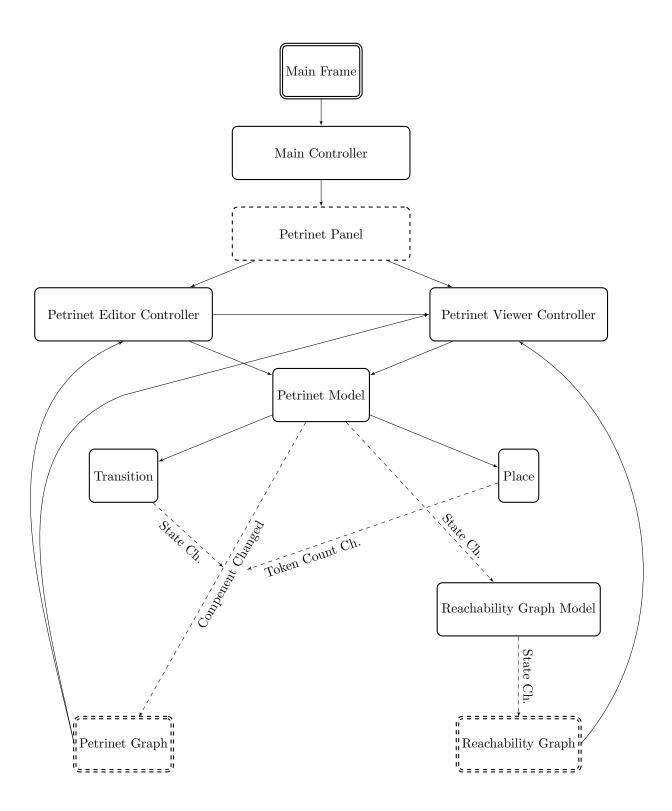


Figure 1: General Design of the Program.

**Petrinet Panel Interface** This interface forwards messages from the main controller. It also handles clicks and forwards them to either the viewer controller or the editor controller. Additionally it keeps track of the toolbar mode and handles GUI related actions for an instance of a petrinet.

**Petrinet Viewer Controller** This controller manages all clicks on the petrinet when in viewer mode and clicks on the reachability graph. It provides methods for interactions with the petrinet and reachability graph except those pertaining to editing the petrinet.

**Petrinet Editor Controller** This controller manages all clicks on the petrinet when in editor mode. It is tasked with providing functionality when editing the given petrinet. This controller could have been part of the petrinet viewer controller. But because the size of both controllers being combined as one would have been quite big and clicks on transitions are handled differently in viewer and editor mode this solution appeared more manageable.

### 2.1.3 Listeners

Petrinet State Changed Listener This listener serves as an intermediary between the petrinet model and the reachability graph model. When the petrinet changes states the reachability graph is informed. It distinguishes between states changing because a transition has been fired and changing states because the petrinet has changed  $\rightarrow$  in the latter case the reachability graph and its initial state have to be reset.

**Petrinet Component Changed Listener** This listener communicates changes in components of the petrinet to the petrinet graph. Whenever elements are added/removed or their attributes are changed it forwards this information. Because it listens for changes in elements it also listens to the *NumberOfTokensChangedListener* and the *TransitionStateListener* and forwards their events.

Number of Tokens Changed Listener Part of the data structure *Place* it listens to changes in the number of tokens on the place. If the number changes it informs the petrinet component changed listener.

**Transition State Listener** Part of the data structure *Transition* it listens to changes in the activation state of the transition. If the state changes it informs the petrinet component changed listener.

Reachability State Changed Listener Listens for changes in the reachability graph model and communicates it to the reachability graph. It contains methods for adding/removing states and also methods for modifying edges.

**Toolbar Changed Listener** Listens for events belonging to the un-/redo buttons and the add/remove edge buttons. If events happen the buttons are highlighted accordingly.

Adjust Arrow Heads Listener This is a special listener for GraphStream. As the arrow heads don't adjust appropriately when adding/removing elements in the graph they need to be adjusted manually. This listener watches for events where arrow heads need to be adjusted.

### 2.2 Petrinet Analyser and the Algorithm for Checking Boundedness

The *PetrinetAnalyser* simulates clicking through the transitions of the petrinet and records the results. For this purpose it starts with the initial state and gets all activated transitions. It fires each transition one by one, gets the resulting new state and recursively calls the analyse method. At each firing of a transition the boundedness of the resulting state is checked. If it signifies the petrinet being unbouded the analyser terminates. Otherwise

### Algorithm 1 Check if Current State is Bounded in a Petrinet

```
function CheckIfCurrentStateIsBounded
   for each predecessor state s of currentState do
      state \leftarrow \text{CheckIfStateIsBounded}(s, \text{ new list, currentState})
      if state \neq null then
          return false
      end if
   end for
   return true
end function
function CheckIfStateIsBounded(state, visitedStates, originalState)
   if state is in visitedStates then
       return null
   end if
   Add state to visitedStates
   if originalState.isBiggerThan(state) then
      return state
   end if
   for each predecessor state s of state do
      newState \leftarrow CheckIfStateIsBounded(s, visitedStates, originalState)
      if newState \neq null then
          return newState
      end if
   end for
   return null
end function
```

it terminates if it does not encounter new states. It therefore uses a depth first search approach to build the reachability graph step by step.

The actual algorithm for checking the boundedness of a given state is implemented in the reachability graph model itself and is represented by algorithm 1. When a state is added the reachability graph model automatically checks for unboundedness. The state is the *currentState* in the algorithm as any added state becomes the current state in the reachability graph model. As each instance of the petrinet state model knows its predecessors it can search for states that mark the current state as unbounded  $\rightarrow$  the algorithm loops through them for the current state and calls the *CHECKIFSTATEISBOUNDED* method for every one of them. This medhod in turn loops through every predecessor of the state passed to it and calls itself recursively. The base case terminating the recursive loop is if the state has already been visitied or if it signifies that the given petrinet is unbounded. To this end the petrinet state model implements a method called *isBiggerThan* that takes a state as a parameter and returns whether the number of tokens on any place is bigger while the tokens on all the other places are the same or bigger also. The algorithm checks the original state (the current in the reachability graph model) and checks it against the state currently passed to the method. This means by performing a depth first search all states in the reachability graph model are checked iteratively until an m is found or all states have been checked.

Algorithm 2 shows how the PetrinetAnalyser works. All the analyser has to do is start from the initial state and simulate all possible configurations of the petrinet at hand. It iteratively performs a depth first approach of firing all transitions  $\rightarrow$  for each state it loops over all activated transitions, fires them and calls the analyse function for the resulting new state in turn firing all activated transitions in this new state. It does so until all configurations have been exhausted, the reachability model informs the analyser that the petrinet is unbounded or the petrinet has been found unbounded in another configuration (this what the variable bounded keeps track of). If the petrinet is unbounded the analyser updates the reachability graph by resetting it to the initial state and only firing the transitions leading up to and including the path marking the petrinet unbouded. This way only the relevant information is shown in the reachability graph.

### Algorithm 2 Analyse State in a Petrinet

```
function AnalyseState(state, visited)
   if not bounded then
      return
   end if
   if state is in visited then
      return
   end if
   numberOfNodes \leftarrow numberOfNodes + 1
   Add state to visited
   petrinet.setState(state)
   for each activated transition t in petrinet do
      petrinet.fireTransition(t)
      stateBounded \leftarrow reachabilityGraphModel.checkIfCurrentStateIsBounded()
      if not stateBounded then
          bounded \leftarrow false
          return
      end if
      numberOfEdges \leftarrow numberOfEdges + 1
      AnalyseState(reachabilityGraphModel.getCurrentPetrinetState(), visited)
      reachabilityGraphModel.setCurrentState(state)
      petrinet.setState(state)
   end for
end function
```

### 2.3 Layout

The program supports three layouts: two custom layouts and the automatic layout<sup>3</sup> provided by GraphStream. Only the two custom layouts are discussed in this section.

### 2.3.1 Tree Layout

The tree layout organizes nodes in the reachability graph as a tree. The level in which a node goes into is defined by the length of the path from the initial node and is defined for each state in the class PetrinetState and therefore provided to the graph whenever a state is added. On each level the nodes are ordered by their proximity to the parent in the level above  $\rightarrow$  nodes should be near below the parent node so edges don't intersect too much with other nodes. Unfortunately this does not prevent intersecting edges altogether. For example if a node has a child directly below and the child has another child directly below and the child of the child has an edge to the grandparent the edge between the grandchild and the grandparent will intersect the child. For this reason adding a new node to the tree performs two additional checks: checkEdgeIntersection and spreadSprites. The former checks whether any edges intersect with any nodes. If so the nodes are pushed to the left or right according to their position on the level. This is repeated until no edges intersect nodes anymore or until the function has looped 20 times. The latter checks three cases: parallel edges and crossing edges. On parallel edges the sprites need to be moved because they are almost guaranteed to cross each other. This is easy since the default position of sprites is in the middle and can be shifted accordingly. Crossing edges are more tricky. An additional check has to be performed to check whether the sprites actually intersect. To this end the reachability Graph Layout package provides data structures and functions to work with rectangles. The layout may not solve all issues but in most cases it provides a pleasant experience.

### 2.3.2 Circle Layout

This layout organizes nodes around a circle – or to be more precise an ellipse. Initially the nodes are ordered by level. Just putting them around a circle would lead to level clusters and most importantly edges clustering at the circles scope. To avoid clusters the initial ordering of nodes is spread evenly around the circle by a certain

 $<sup>^3</sup>$ see Automatic layout at https://graphstream-project.org/doc/Tutorials/Graph-Visualisation/1.0/



<sup>&</sup>lt;sup>4</sup>For a more in depth explanation of how the algorithm works please refer to the comments in the source code.