Practical Process Injection techniques

This article is about explaining multiple process injection techniques in detail. Malware had numerous capabilities to achieve its goal and it is improving day by day .By clear understanding of basic techniques will provide strong base to understand day to day improvised techniques. Motto of this article is make things more simpler way . So lets start deep dive into

What is process injection?

By book ,In the Windows operating system, processes are allowed to allocate, read, and write in another process's virtual memory, as well as create new threads, suspend threads, and change these threads' registers, including the instruction pointer (EIP/RIP). Process injection is a technique that's implemented by malware authors so that they can inject code.In simple words ,**writing malicious code from one live userspace process (malware) to another live userspaceprocess (target).**

Why Process Injection ?

To bypass trivial firewalls that block internet connections from all applications except browsers or other signed allowed apps,so malware can speak with c2 server. Evade debuggers and other dynamic analysis or monitoring tools by running the malicious code inside another unmonitored. Maintain persistence for fileless malware. By injecting its code into a background process, malware can maintain persistence on a server that rarely gets rebooted

Is this techniques are still relevant today ?

Many of techniques are comprised or detected by next gen AV's and specially Microsoft win 10 x64 is became popular because of its security features like CFG(prevent indirect calls to non-approved addresses,CIG (only allow modules signed by Microsoft/Microsoft Store/WHQL to be loaded into the process memory), Dynamic Code prevention and few others .But if you want to detect new techniques ,you need to understand basics as well for new techniques they only change few

things like different api core principle of technique is same .In other words we cant paragraphs without learning the alphabets

Process injection building blocks

- Memory allocation
- Memory writing
- Execution

CLASSIC DLL INJECTION:

The malware writes the path to its malicious dynamic-link library (DLL) in the virtual address space of another process, and ensures the remote process loads it by creating a remote thread in the target process .

Before going how it working practical we need to understand below api's and its purpose

CreateToolhelp32Snapshot -Takes a snapshot of the specified processes, as well as the heaps, modules, and threads used by these processes

Process32First -Retrieves information about the first process encountered in a system snapshot.
Process32Next -Retrieves information about the next process recorded in a system snapshot

OpenProcess - The OpenProcess function returns a handle of an existing process object.
VirtualAllocEX - The VirtualAllocEx function is used to allocate the memory and grant the access permissions to the memory address.

WriteProcessMemory - The WriteProcessMemory function writes data to an area of memory in a specified process.

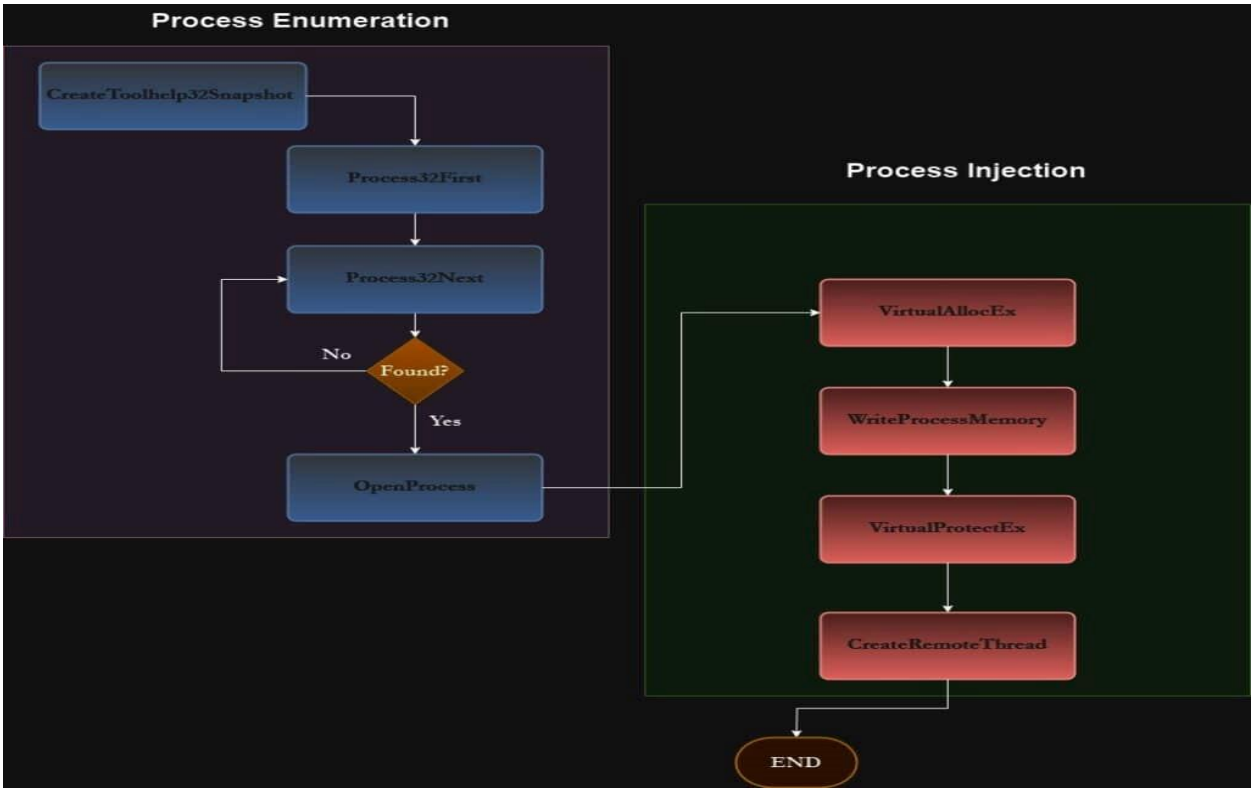
CreateRemoteThread - The CreateRemoteThread function creates a thread that runs in the virtual address space of another process.

First malware need to find the target which is usually legitimate one .for that malware use these api: CreateToolhelp32Snapshot,Process32First,Process32Next**.** In a kitchen you need find a chocolate in out of 10 boxes ,first you need to take 10 boxes before you and open and check the boxes to find chocolate like wise : CreateToolhelp32Snapshot—helps to take

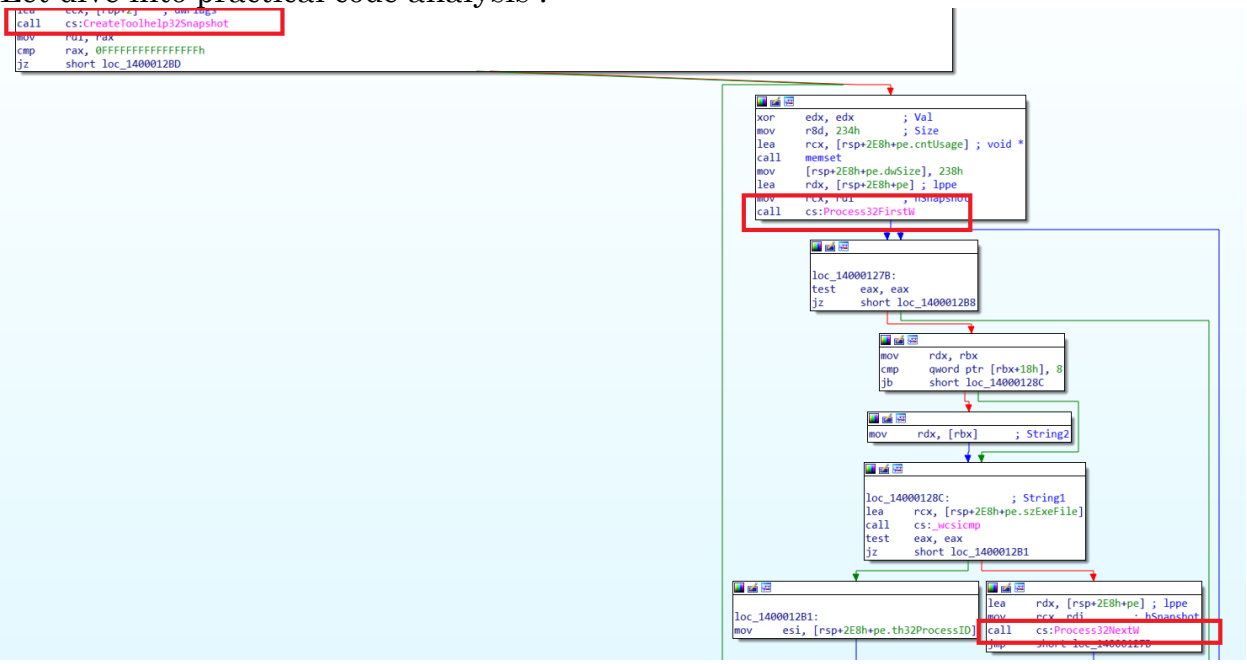
snapshot of all process , Process32First - retrieves information about the first process in the snapshot and Process32Next- is used in a loop to iterate through them.

Then malware calls VirtualAllocEx to have a space to write the path to its DLL. The malware then calls WriteProcessMemory to write the path in the allocated memory. Finally, to have the code executed in another process, the malware calls APIs such as CreateRemoteThread, NtCreateThreadEx, or RtlCreateUserThread.

Usual flow of process injection :

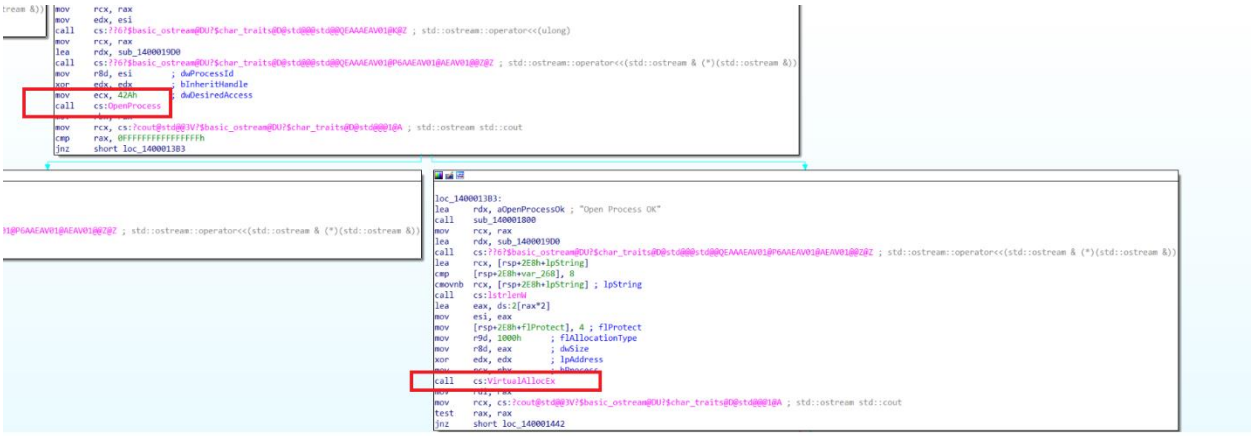


Let dive into practical code analysis :

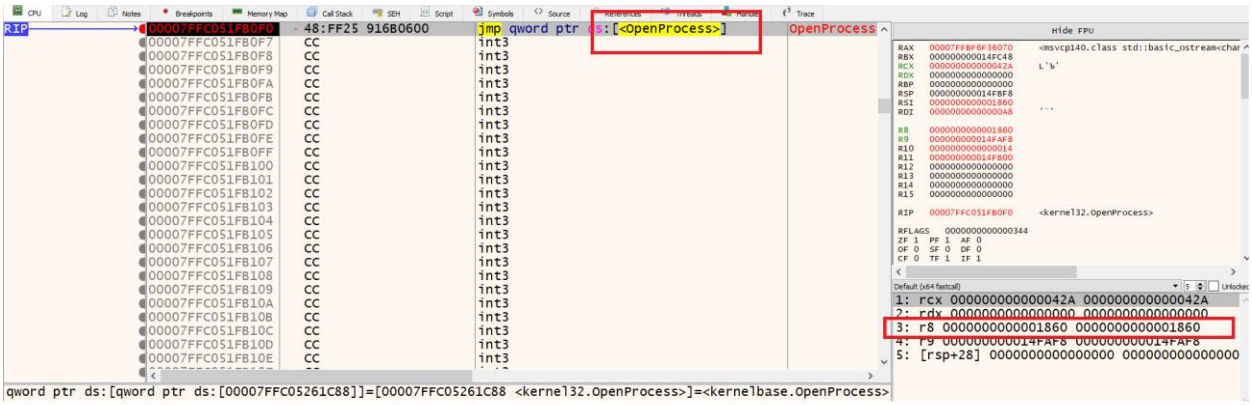


From the above ,we can see taking snapshot of all process and looking into the process and iterating it for targeted process.

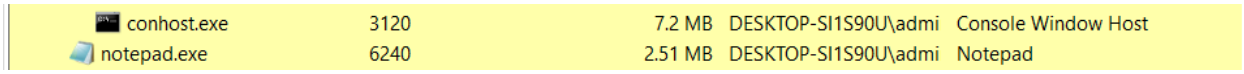
For this sample ,it is target one legitimate file ,we can find out next



And code we can follow if find the target process ,it can open the process allocate memory .

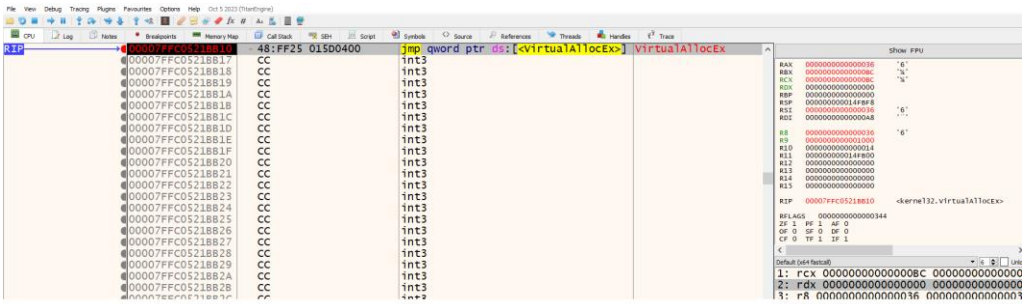


From the image ,debugger stop at open process ,you can see red box marked below you can find the 3rd parameter it shows value 1860.it is target process we are looking and that need to convert to decimal we can see what it is hex 1860 -> dec 6240 PID

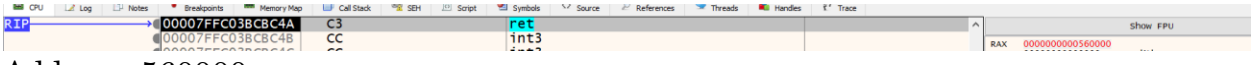


So it is notepad.exe

Now we got the process and next to find in this process where it write



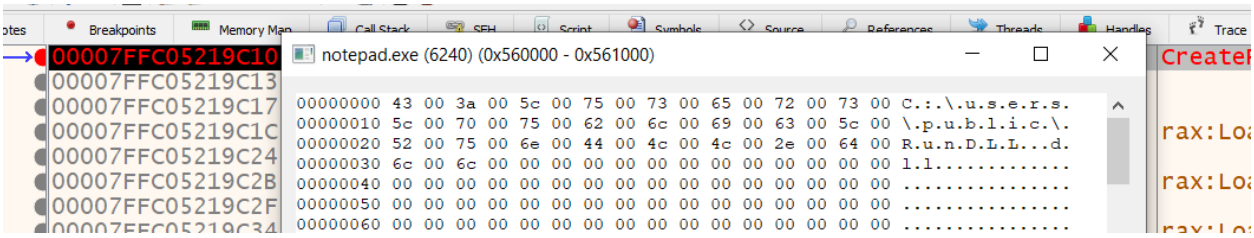
The next stop /bp will be virtualallocex ,after we get this we click on execute till return and eax had address of where it can write .



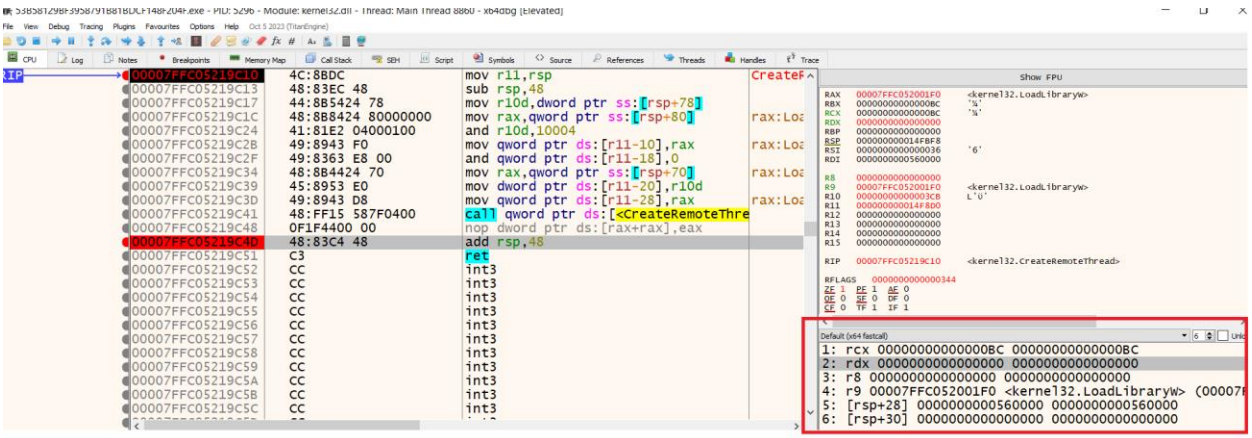
Address :560000

Same we can confirm with next stop writeprocessmemory
Which is 2nd parameter

Now can follow this in process hacker after writing



You can see the path of dll above
After that, the general idea is to pass the address of LoadLibrary to one of these APIs so that a remote process has to execute the DLL on behalf of the malware



Next stop is createremotethread() which can start dll ,do have look handle and parameter (1st one and 5 th one)

In way this we find the
In way this we can analyse and take the details of classic dll injection

1.Pre requisites to use this technique the DLL is on disk; write technique used to write the DLL path to the target process
2.CFG no impact
3.CIG blocks this technique

You might got doubt why this guy this for simple dll injection explained this much because from the next technique onwards I take the liberality to go little fast

PE injection:

For this technique we learn theory from red teaming perspective and this basic helps to understand many in memory techniques

So the idea of this technique is malware can copy its malicious pe into an existing open process and execute by calling CreateRemoteThread or others. One advantage of PE injection over the LoadLibrary technique is that the malware does not have to drop a malicious DLL on the disk. Previously we add path in memory of remote process so there is no issue but we need add to pe file (exe or dll),When running in memory most, but not all, portable executables make use of 2 structures we need to know about: IAT (Import Address Table), and Reloc (Base Relocation Table)

What IAT table do simply allows for the addresses of DLL functions to be set by the PE loader, without having to modify the code of the application.

What Base Relocation Table do ,it is also possible that the application itself is not loaded at the same address every time. For the most part this isn't a problem because the application uses relative addressing, however because absolute addresses will need to be changed if the process base address changes, whenever an absolute address is used, it must be easily located. The Base Relocation Table is a table of pointers to every absolute address used in the code. During process initialization, if the process is not being loaded at its base address, the PE loader will modify all the absolute addresses to work with the new base address.

Before going to technical part ,go with this story -you and your cousin loved the same toy and you want gift that one to him .you bought two same toys and gave one to cousin .technical details really important because it helps to understand in the other tech like process hallowing

Steps :

1.Pe injections usually starts with Get the image base address **imageBase** Parse the PE headers and get its sizeOfImage

2. Allocate a block of memory -size of PE image retrieved(**VirtualAlloc**),), call it localImage,Copy the image of the current process into the newly allocated local memory localImage**(memcopy)**

3.Allocate a new memory block size of PE image retrieved step 1 in a remote process - the target process we want to inject the currently running PE into call it targetImage

4. When a malware injects its PE into another process it will have a new base address which is unpredictable, requiring it to dynamically recompute the fixed addresses of its PE. To overcome this, the malware needs to find its relocation table address in the host process, and resolve the absolute addresses of the copied image by looping through its relocation descriptors.

5.by this way it will do , Calculate the remote address of the function to be executed in the remote process by subtracting the address of the function in the current process by the base address of the current process, then adding it to the address of the allocated memory in the target process.

6.Create a new thread with the start address set to the remote address of the function (CreateRemoteThread).

7. In some cases once the image is executed in the remote process, it may have to fix its own IAT so that it can call functions imported from DLLs, however; DLLs are usually at the same address in all processes, so this wouldn't be necessary.

In practicle we do with easy sample ,injection of pe and execute :

Code flow :

```
push    edi
call    ds:GetCurrentProcessId
mov     ebx, eax
push    ebx
push    ebx
push    offset Format   ; "[+] PID is: %d,0x%x\n"
mov     [ebp+var_C], ebx
call    _printf
add     esp, 0Ch
push    ebx             ; dwProcessId
push    0               ; bInheritHandle
push    1FFFFFh         ; dwDesiredAccess
call    ds:OpenProcess
mov     ebx, eax
push    ebx
push    offset aProcessHandle0 ; "[+] Process handle: 0x%x\n"
call    _printf
add     esp, 8
push    offset ModuleName ; "kernel32.dll"
call    ds:GetModuleHandleA
push    offset ProcName ; "LoadLibraryW"
push    eax             ; hModule
call    ds:GetProcAddress
mov     edi, eax
push    edi
push    offset aLoadlibraryBas ; "[+] LoadLibrary base address is: 0x%x\n"
call    _printf
add     esp, 8
push    4               ; flProtect
push    3000h           ; flAllocationType
push    1000h           ; dwSize
push    0               ; lpAddress
push    ebx             ; hProcess
call    ds:VirtualAllocEx
mov     esi, eax
push    esi
push    offset aAllocatedMemor ; "[+] Allocated memory address in target "...
call    _printf
add     esp, 8
push    0               ; lpNumberOfBytesWritten
push    28000h          ; nSize
```

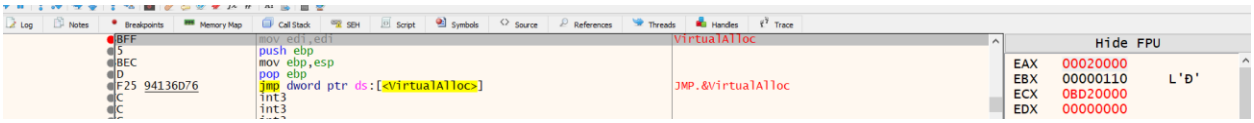
Here it is getting the id ,opening the process ,and allocating the memory

```
push    esi             ; lpBaseAddress
push    ebx             ; hProcess
call    ds:WriteProcessMemory
push    offset aDllNameIsWritt ; "[+] DLL name is written to memory of ta"...
call    _printf
add     esp, 4
mov     [ebp+ThreadId], 0
lea     eax, [ebp+ThreadId]
push    eax             ; lpThreadId
push    0               ; dwCreationFlags
push    esi             ; lpParameter
push    edi             ; lpStartAddress
push    0               ; dwStackSize
push    0               ; lpThreadAttributes
push    ebx             ; hProcess
call    ds:CreateRemoteThread
push    offset aSuccessfullySt ; "[+] Successfully started DLL in target "...
call    _printf
mov     eax, [ebp+ThreadId]
add     esp, 4
pop     edi
pop     esi
pop     ebx
test    eax, eax
jz      short loc_40137A
```

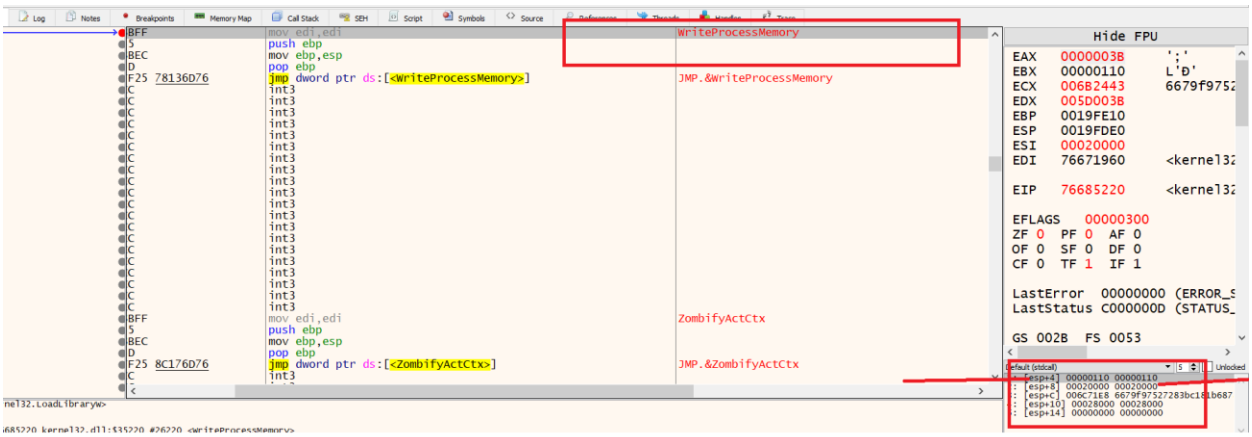
```
push    [ebp+var_C]
push    eax
push    offset aInjectedThread ; "[+] Injected thread id: %u for pid: %u"...
call    _printf
add     esp, 0Ch
```

Now with debugger :

After find the process and open the process ,need to allocate the memory



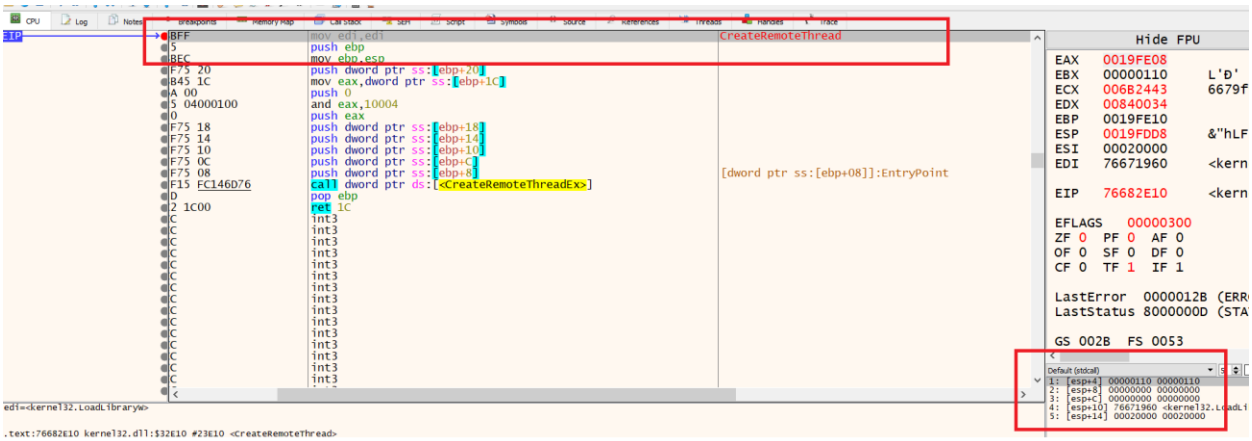
It running the execute till return in virtual alloc call ,we got the address of where is locating .here it is 20000,you can see in the EAX

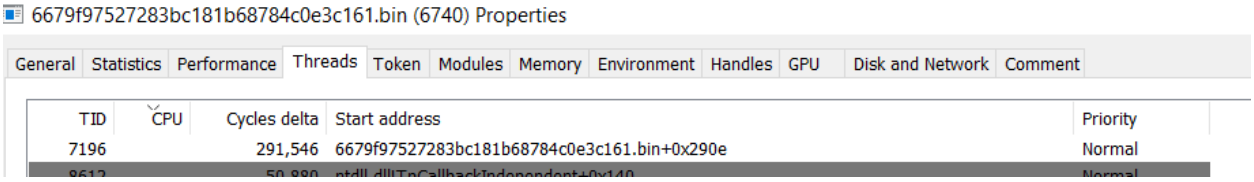


Above you can two lines 1st is 110 it is handle and 2nd is the address where it is writing 20000

6679f97527283bc181b68784c0e3c161.bin (6740) (0x20000 - 0x21000)

```
00000000 4d 5a 90 00 03 00 00 00 04 00 00 00 ff ff 00 00 MZ..............
00000010 b8 00 00 00 00 00 00 00 40 00 00 00 00 00 00 00 ........@.......
00000020 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000030 00 00 00 00 00 00 00 00 00 00 00 00 f8 00 00 00 ................
00000040 0e 1f ba 0e 00 b4 09 cd 21 b8 01 4c cd 21 54 68 ........!..L.!Th
00000050 69 73 20 70 72 6f 67 72 61 6d 20 63 61 6e 6e 6f is program canno
00000060 74 20 62 65 20 72 75 6e 20 69 6e 20 44 4f 53 20 t be run in DOS
00000070 6d 6f 64 65 2e 0d 0d 0a 24 00 00 00 00 00 00 00 mode....$.......
00000080 38 a3 20 2f 7c c2 4e 7c 7c c2 4e 7c 7c c2 4e 7c 8. /|.N||.N||.N|
00000090 7c c2 4e 7c 7d c2 4e 7c c1 8d d8 7c 7e c2 4e 7c |.N|}.N|...|~.N|
000000a0 bf cd 2e 7c 7e c2 4e 7c bf cd 41 7c 7f c2 4e 7c ...|~.N|..A|..N|
000000b0 bf cd 13 7c 7e c2 4e 7c bf cd 11 7c 75 c2 4e 7c ...|~.N|...|u.N|
000000c0 5b 04 20 7c 7d c2 4e 7c 5b 04 35 7c 65 c2 4e 7c [. |}.N|[.5|e.N|
000000d0 7c c2 4f 7c 98 c2 4e 7c 62 90 c4 7c 60 c2 4e 7c |.O|..N|b..|`.N|
000000e0 62 90 df 7c 7d c2 4e 7c 52 69 63 68 7c c2 4e 7c b..|}.N|Rich|.N|
000000f0 00 00 00 00 00 00 00 00 50 45 00 00 4c 01 05 00 ........PE..L...
00000100 b5 ce 20 5a 00 00 00 00 00 00 00 00 e0 00 02 21 .. Z...........!
00000110 0b 01 09 00 00 e0 00 00 00 00 52 00 00 00 00 00 ..........R.....
00000120 c0 fa 00 00 00 10 00 00 00 10 00 00 00 00 00 10 ................
00000130 00 10 00 00 00 02 00 00 05 00 00 00 00 00 00 00 ................
00000140 05 00 00 00 00 00 00 00 00 00 00 00 70 01 00 00 ............p...
00000150 00 00 00 00 02 00 40 01 00 00 10 00 00 10 00 00 ......@.........
00000160 00 00 10 00 00 10 00 00 00 00 00 00 10 00 00 00 ................
00000170 00 00 00 00 00 00 00 00 a8 21 01 00 18 01 00 00 .........!......
00000180 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000190 00 00 00 00 00 00 00 00 00 00 00 00 50 01 00 b0 ............P.....
000001a0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000001b0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000001c0 90 16 01 00 40 00 00 00 00 00 00 00 00 00 00 00 ....@...........
000001d0 00 00 01 00 80 03 00 00 00 00 00 00 00 00 00 00 ................
000001e0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000001f0 2e 66 6c 61 74 00 00 00 9b 02 00 00 00 10 00 00 .flat...........
00000200 00 04 00 00 00 04 00 00 00 00 00 00 00 00 00 00 ................
00000210 00 00 00 00 60 00 00 60 2e 74 65 78 74 00 00 00 ....`..`.text...
00000220 cd de 00 00 00 20 00 00 00 e0 00 00 00 08 00 00 ..... ..........
00000230 00 00 00 00 00 00 00 00 00 00 00 00 20 00 00 60 ............ ..`
00000240 2e 72 64 61 74 61 00 00 dc 34 00 00 00 00 01 00 .rdata...4......
00000250 00 36 00 00 00 e8 00 00 00 00 00 00 00 00 00 00 .6..............
00000260 00 00 00 00 40 00 00 40 2e 64 61 74 61 00 00 00 ....@..@.data...
00000270 b4 08 00 00 00 40 01 00 00 0a 00 00 00 1e 01 00 .....@..........
00000280 00 00 00 00 00 00 00 00 00 00 00 00 40 00 00 c0 ............@...
00000290 2e 72 65 6c 6f 63 00 00 22 11 00 00 00 50 01 00 .reloc.."....P..
000002a0 00 12 00 00 00 28 01 00 00 00 00 00 00 00 00 00 .....(..........
000002b0 00 00 00 00 40 00 00 42 00 00 00 00 00 00 00 00 ....@..B........
000002c0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000002d0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000002e0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
000002f0 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000300 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000310 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000320 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000330 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000340 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000350 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
00000360 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ................
```





Thread created

CFG/CIG-readiness: CIG prevents loading on non-Microsoft signed DLL. An attempt to do

Process Hollowing:

Basically Process hollowing occurs when a malware hollows out the legitimate code from memory of the victim process which is clean , and overwrites the memory space of the victim process (e.g svchost.exe) with a malicious executable or code .

Steps :
1. Create the victim process in a suspended state .
2. Information Gathering of the newly created process.
3. Hollowing the memory of the victim process .
4. Allocate and inject the malicious code into the victim process.
5. Adjust the base address .
6. Set the entrypoint .

In practicle we do discuss the api which are not involved in the steps ,before going i will add few to understand more about this one

NT Headers:

PE structure containing information about the PE signature and the Headers : IMAGE_FILE_HEADER (COFF) and IMAGE_OPTIONAL_HEADER (OptionalHeader), both containing basic information about the PE, such as the ImageBase .

```
typedef struct _IMAGE_NT_HEADERS64 {
    DWORD Signature;
    IMAGE_FILE_HEADER FileHeader;
    IMAGE_OPTIONAL_HEADER64 OptionalHeader;
} IMAGE_NT_HEADERS64, *PIMAGE_NT_HEADERS64;
```

PEB:
It is the representation of a process in user mode, containing valuable information about the running process.

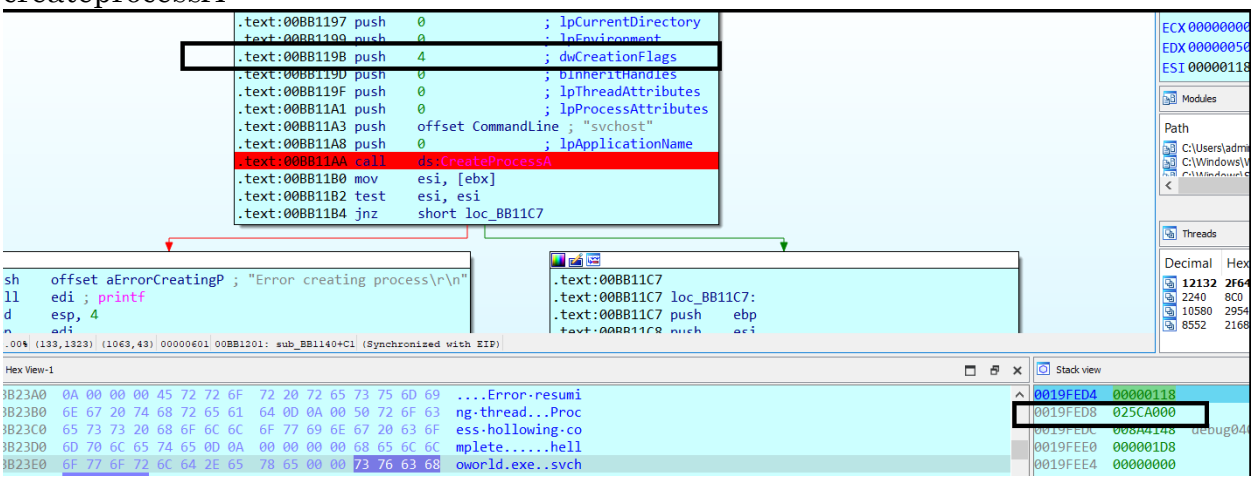
```
0:007> dt ntdll!_PEB
   +0x000 InheritedAddressSpace : UChar
   +0x001 ReadImageFileExecOptions : UChar
   +0x002 BeingDebugged    : UChar
   +0x003 BitField         : UChar
   +0x003 ImageUsesLargePages : Pos 0, 1 Bit
   +0x003 IsProtectedProcess : Pos 1, 1 Bit
   +0x003 IsImageDynamicallyRelocated : Pos 2, 1 Bit
   +0x003 SkipPatchingUser32Forwarders : Pos 3, 1 Bit
   +0x003 IsPackagedProcess : Pos 4, 1 Bit
   +0x003 IsAppContainer   : Pos 5, 1 Bit
   +0x003 IsProtectedProcessLight : Pos 6, 1 Bit
   +0x003 IsLongPathAwareProcess : Pos 7, 1 Bit
   +0x004 Padding0         : [4] UChar
   +0x008 Mutant           : Ptr64 Void
   +0x010 ImageBaseAddress : Ptr64 Void
   +0x018 Ldr              : Ptr64 _PEB_LDR_DATA
   +0x020 ProcessParameters : Ptr64 _RTL_USER_PROCESS_PARAMETERS
   +0x028 SubSystemData    : Ptr64 Void
   +0x030 ProcessHeap      : Ptr64 Void
   +0x038 FastPebLock      : Ptr64 _RTL_CRITICAL_SECTION
   +0x040 AtlThunkSListPtr : Ptr64 _SLIST_HEADER
   +0x048 IFEOKey          : Ptr64 Void
   +0x050 CrossProcessFlags : Uint4B
   +0x050 ProcessInJob     : Pos 0, 1 Bit
   +0x050 ProcessInitializing : Pos 1, 1 Bit
   +0x050 ProcessUsingVEH  : Pos 2, 1 Bit
   +0x050 ProcessUsingVCH  : Pos 3, 1 Bit
   +0x050 ProcessUsingFTH  : Pos 4, 1 Bit
   +0x050 ProcessPreviouslyThrottled : Pos 5, 1 Bit
   +0x050 ProcessCurrentlyThrottled : Pos 6, 1 Bit
   +0x050 ProcessImagesHotPatched : Pos 7, 1 Bit
   +0x050 ReservedBits0    : Pos 8, 24 Bits
   +0x054 Padding1         : [4] UChar
   +0x058 KernelCallbackTable : Ptr64 Void
   +0x058 UserSharedInfoPtr : Ptr64 Void
   +0x060 SystemReserved   : Uint4B
   +0x064 AtlThunkSListPtr32 : Uint4B
```

From the above data of process like 0x002, 0x0010, 0x0018, 0x0020 plays critical role for loader

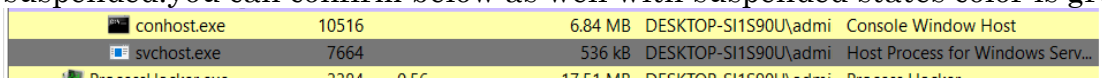
Thread Context:
For each Thread in a process, that Thread needs **to** maintain its **“state”** (after all, they are not always running). So, when the processor tells a Thread that it is not its turn to run, the Thread takes a **snapshot of the CPU registers at the moment it was stopped** and when it is its turn to run again, it can **“pick up where it left off”**
This “state” is called **context** . And how do we change this context? Simple.Through the **GetThreadContext()** and **SetThreadContext()** functions .

Code flow and debug:

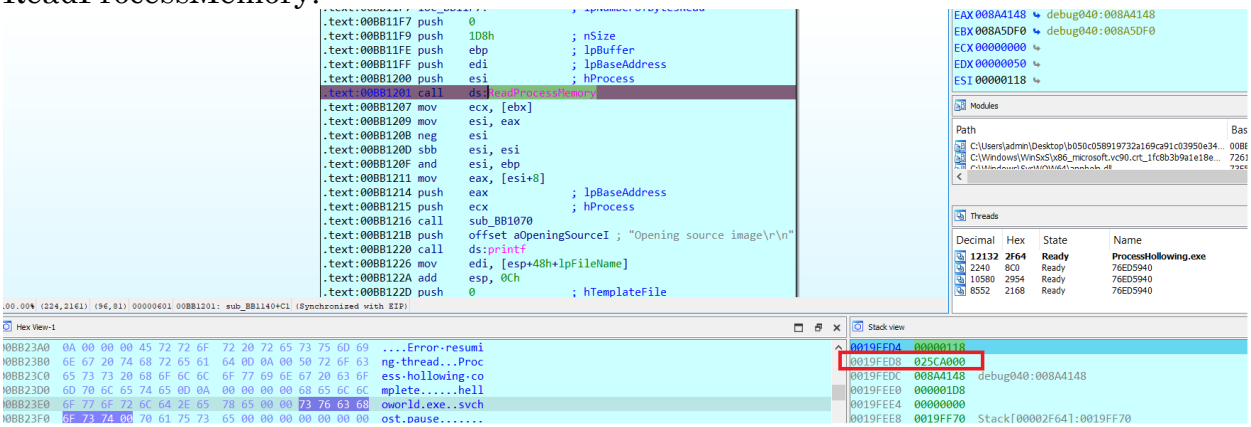
createprocessA



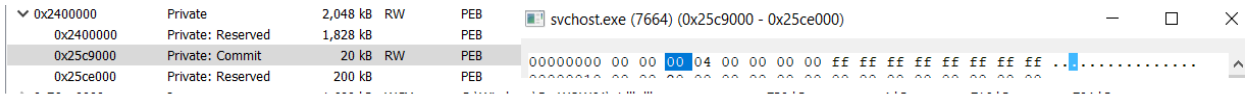
From the above you can see it is creating svc host and if you follow stack 2nd parameter in hexdump you can ,what process it is creating,4 stands for suspended.you can confirm below as well with suspended states color is grey



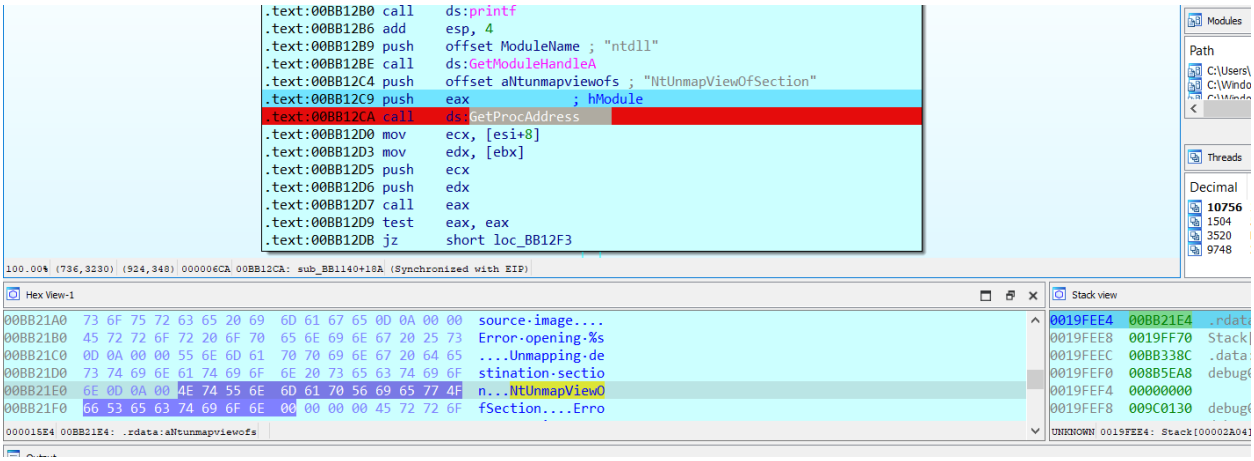
ReadProcessMemory:



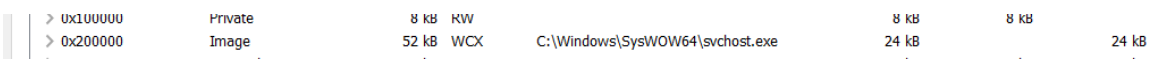
2nd parameter is what is malware want to read a section in victim svchost ,1st parameter is handle .what is in that address?



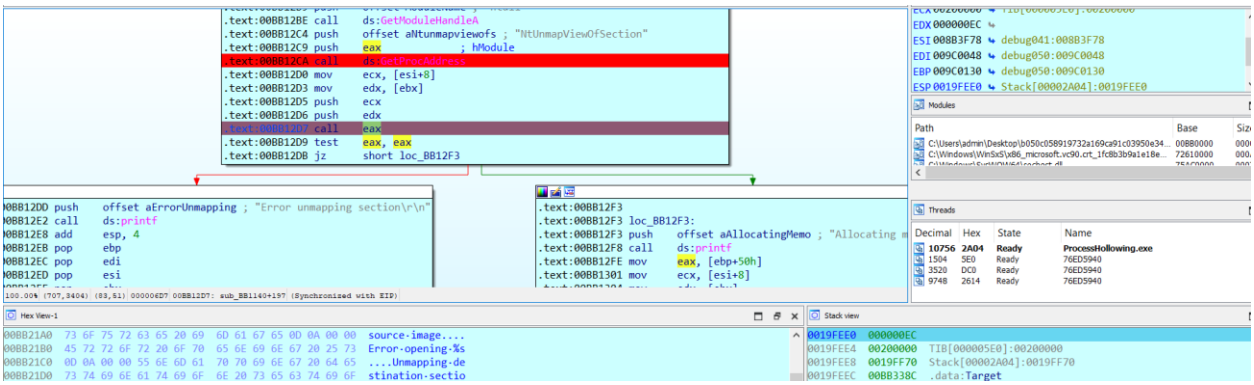
It contains peb structure and address where the executable loaded in the memory (it is little endian format)



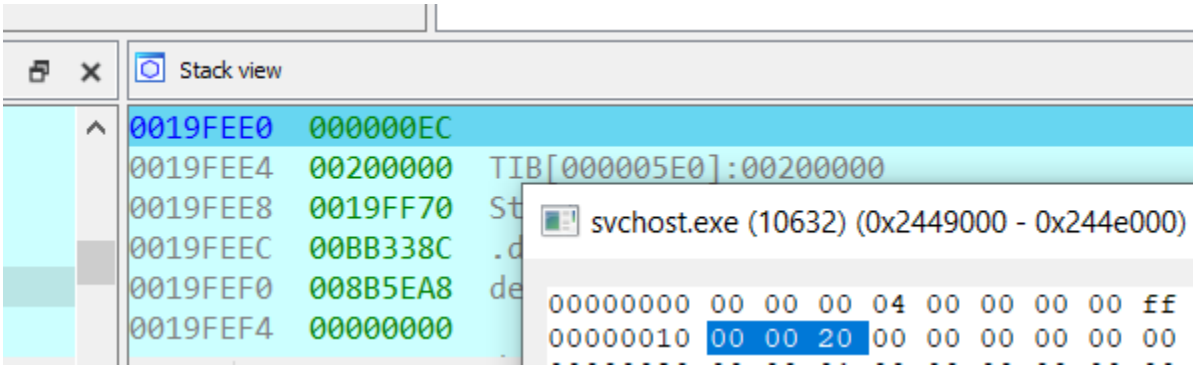
Now it is resolve the api **NTmapViewOfSection function**

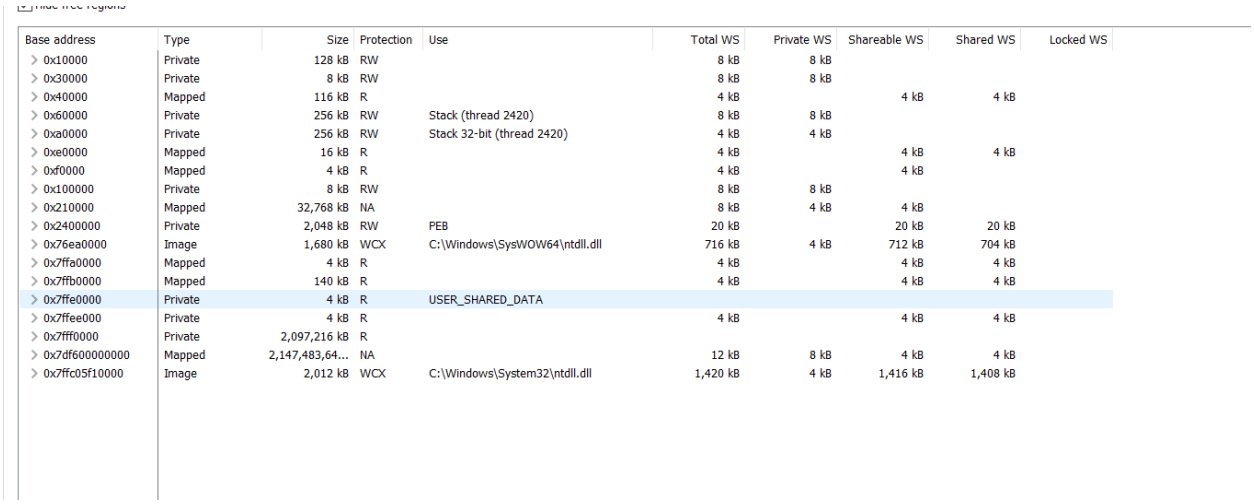


Svchost loaded address



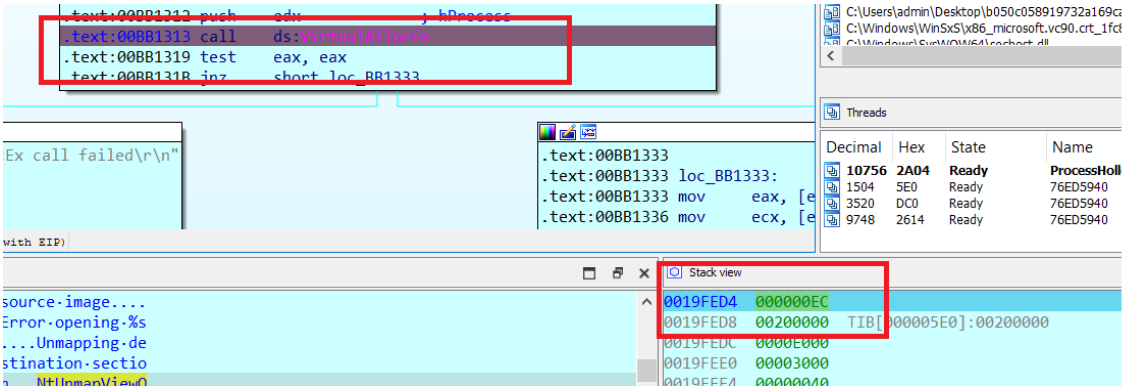
NTmapViewOfSection function called in the eax parameters you can observe





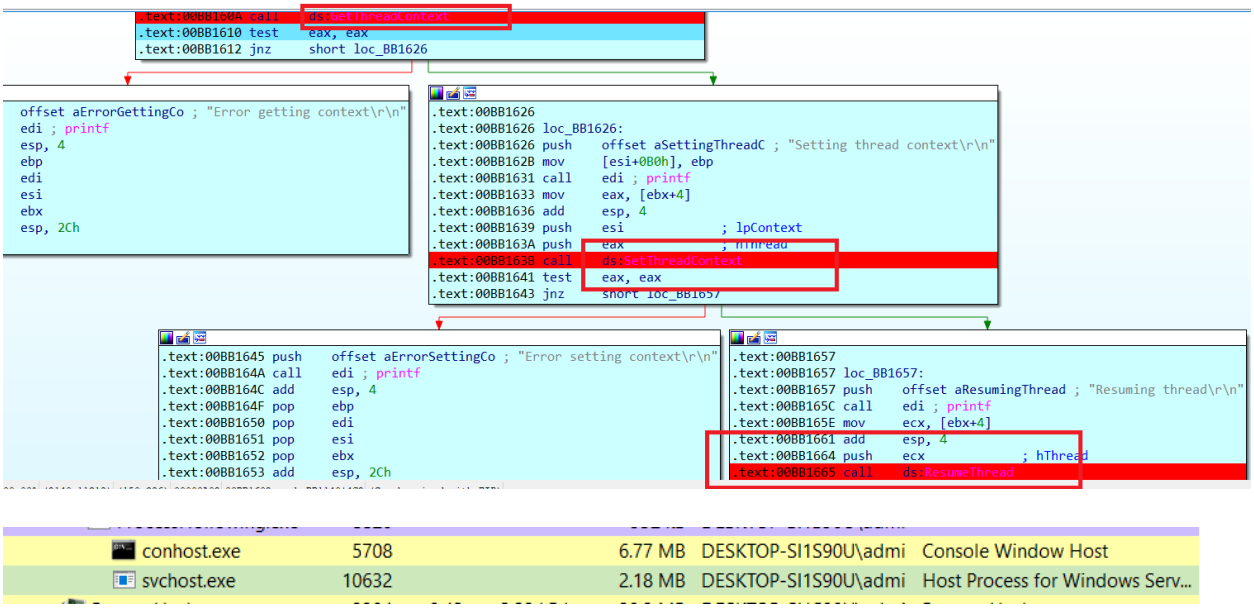
Base address	Type	Size	Protection	Use	Total WS	Private WS	Shareable WS	Shared WS	Locked WS
0x10000	Private	128 kB	RW		8 kB	8 kB			
0x30000	Private	8 kB	RW		8 kB	8 kB			
0x40000	Mapped	116 kB	R		4 kB		4 kB	4 kB	
0x60000	Private	256 kB	RW	Stack (thread 2420)	8 kB	8 kB			
0xa0000	Private	256 kB	RW	Stack 32-bit (thread 2420)	4 kB	4 kB			
0xe0000	Mapped	16 kB	R		4 kB		4 kB	4 kB	
0xf0000	Mapped	4 kB	R		4 kB		4 kB		
0x100000	Private	8 kB	RW		8 kB	8 kB			
0x210000	Mapped	32,768 kB	NA		8 kB	4 kB	4 kB		
0x2400000	Private	2,048 kB	RW	PEB	20 kB		20 kB	20 kB	
0x76ea0000	Image	1,680 kB	WCX	C:\Windows\SysWOW64\ntdll.dll	716 kB	4 kB	712 kB	704 kB	
0x7ffa0000	Mapped	4 kB	R		4 kB		4 kB	4 kB	
0x7ffb0000	Mapped	140 kB	R		4 kB		4 kB	4 kB	
0x7ffe0000	Private	4 kB	R	USER_SHARED_DATA					
0x7ffee000	Private	4 kB	R		4 kB		4 kB	4 kB	
0x7fff0000	Private	2,097,216 kB	R						
0x7df600000000	Mapped	2,147,483,64...	NA		12 kB	8 kB	4 kB	4 kB	
0x7ffc05f10000	Image	2,012 kB	WCX	C:\Windows\System32\ntdll.dll	1,420 kB	4 kB	1,416 kB	1,408 kB	

Now you can see memory address is unmapped



Allocating memory

After that it will do writing headers ,sections ,relocating the base



Then getthreadcontext ,setthreadcontext and resume thread api are used and Setthreadcontext will change the address of entrypoint and then resume the thread

So this the way process hollowing ,I added more details for this technique because one of important techninque and debugged in ida because it had more clarity to show ,dynamic code prevention code will stop this technique ,many I want to use in

demo itself are not worked virtualalloc and to detect need to do is compare the main module of the victim process to its module path. If they're almost the same, the process is not hollowed. If they are significantly different. We do discuss in the second chapter of this topic

THREAD EXECUTION HIJACKING:

This injection technique injects malicious code into the existing thread of a process (thereby avoiding the overhead of creating a new process and thread) and then uses the target to start a thread in itself.

Steps in detail:
OpenProcess: Obtaining a handle to the target process to manipulate it.

VirtualAllocEx: Allocating memory within the target process to store the malicious payload.

WriteProcessMemory: Copying the malicious shellcode into the allocated memory space.

SuspendThread: Temporarily suspending the victim thread to obtain its context.

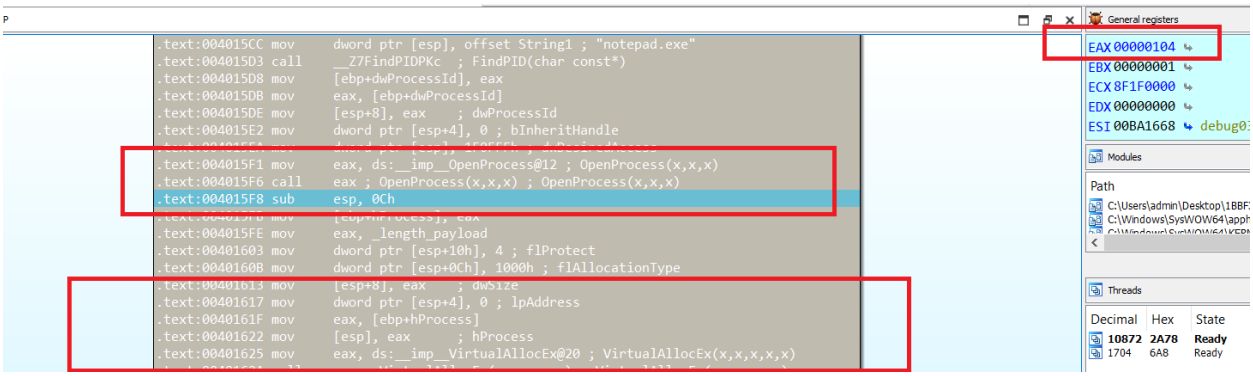
GetThreadContext: Retrieving the context of the suspended thread, including the instruction pointer (RIP).

Modify RIP Address: Changing the RIP address to point to the memory location of the injected shellcode.

SetThreadContext: Setting the modified context back to the thread.

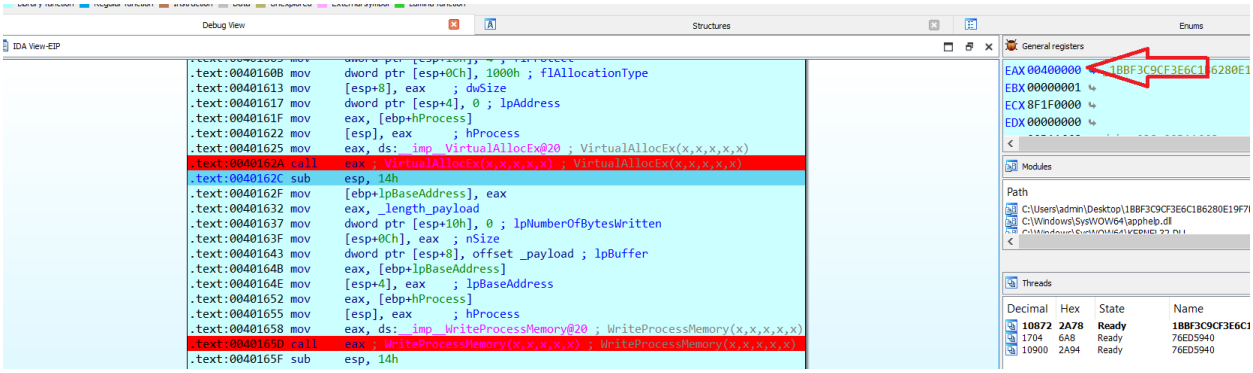
ResumeThread: Resuming the thread's execution, leading to the execution of the injected payload.

Code flow and debugging:



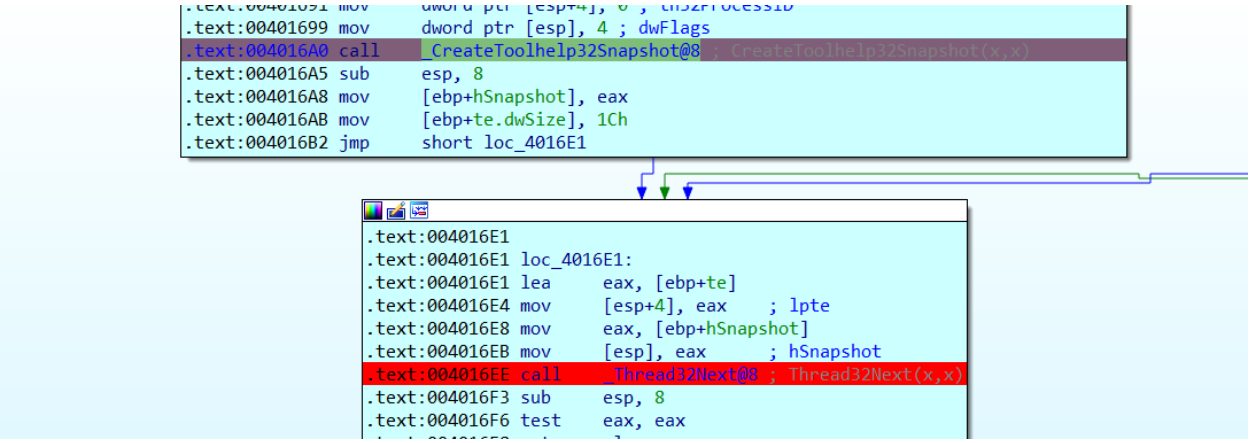
Start with openprocess, handle is 104 .what target it is below

Key	HKLM\SYSTEM\ControlSet001\Control\...	0x110
Mutant	\Sessions\2\BaseNamedObjects\SM0:6...	0xac
Process	notepad.exe (8824)	0x104



Allocated memory and it will write th pe file .

```
003FFFF0  ?? ?? ?? ?? ?? ?? ?? ??  ?? ?? ?? ?? ?? ?? ?? ??  ????????????????
00400000  4D 5A 90 00 03 00 00 00  04 00 00 00 FF FF 00 00  MZ..............
00400010  B8 00 00 00 00 00 00 00  40 00 00 00 00 00 00 00  ........@.......
00400020  00 00 00 00 00 00 00 00  00 00 00 00 00 00 00 00  ................
00400030  00 00 00 00 00 00 00 00  00 00 00 00 80 00 00 00  ................
00400040  0E 1F BA 0E 00 B4 09 CD  21 B8 01 4C CD 21 54 68  ...........L..Th
```



It will change the rights by using virtualprotect ,then call createtoolhelp32 snaphot and thread32next ,to choose the thread .

Then it is similar to last technique ,resume thread once it got context and set as address of entry point

HOOK INJECTION:

This technique similar all previous ones but here we are using this api to achive injecting by SetWindowHookEx–.so what is this api, Windows allow programs to install hooks to monitor various system events such as mouse clicks and keyboard key presses by using SetWindowHookEx. Hooking is a technique used to intercept function calls. Malware can leverage hooking functionality to have their malicious DLL loaded upon an event getting triggered in a specific thread. The SetWindowsHookEx installs a hook routine into the hook chain, which is then invoked whenever certain events are triggered.it has 4 parameters

1.Type of hook

- WH_CALLWNDPROC
- WH_CALLWNDPROCRET
- WH_CBT
- WH_DEBUG
- WH_FOREGROUNDIDLE
- WH_GETMESSAGE
- WH_JOURNALPLAYBACK
- WH_JOURNALRECORD
- WH_KEYBOARD
- WH_KEYBOARD_LL
- WH_MOUSE
- WH_MOUSE_LL
- WH_MSGFILTER
- WH_SHELL
- WH_SYSMSGFILTER

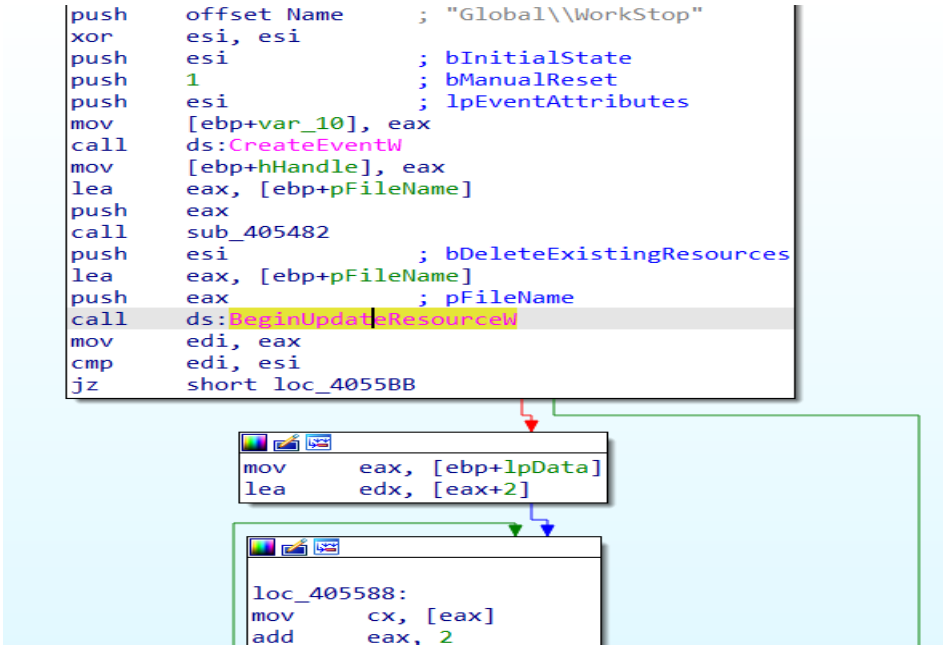
2. Pointer to the function the malware wants to invoke upon the event execution

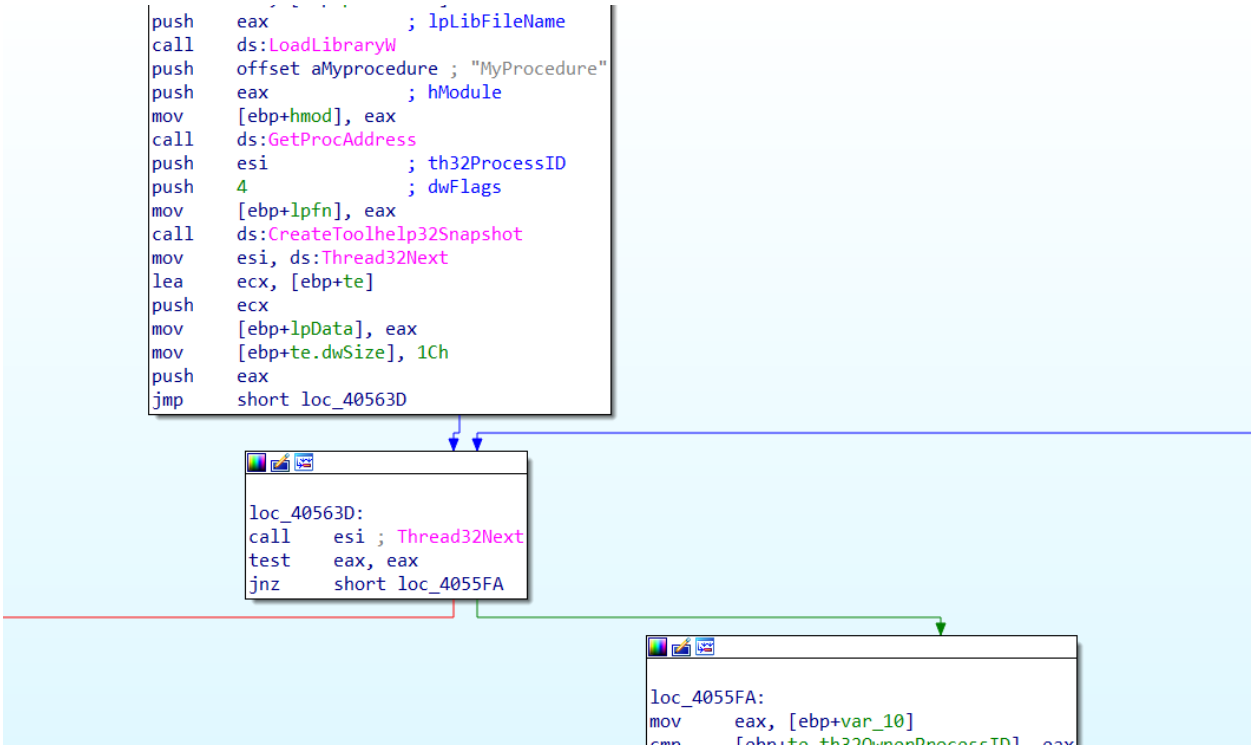
3. A handle to the DLL that contains the hook function (usually LoadLibrary and GetProcAddress are api used before call sethookwindows api)

4. The identifier of the thread, which calls the hook function, f the parameter is 0, the hook will be called by all threads, so we don't have to restrict it to particular thread ID.To avoid that we call

Control flow:

This function starts with create event .





AS we discuced earlier by using loadlibrary and getprocadress ,it loads the function
Then createtoolhelp32snapshot and tread
Then createtoolhelp32snapshot and tread32next findout the thread .

Later that passed the arguments earlier mentioned .

This technique usually block by CIG which prevents indirect call and CFG only mstf singed dll are allowed and dll should on disk for this technique and target must load user32.dll where setwindowshook function exist .

Reference :

I read multiples blogs before writing started with search process injection ,I didn't mentioned links because i might miss some guys and I am always grateful to them

Conclusion :
We learned 5 basic injection techniques here ,in chapter 2 we do discuss the advanced techniques and detection as well .Take care