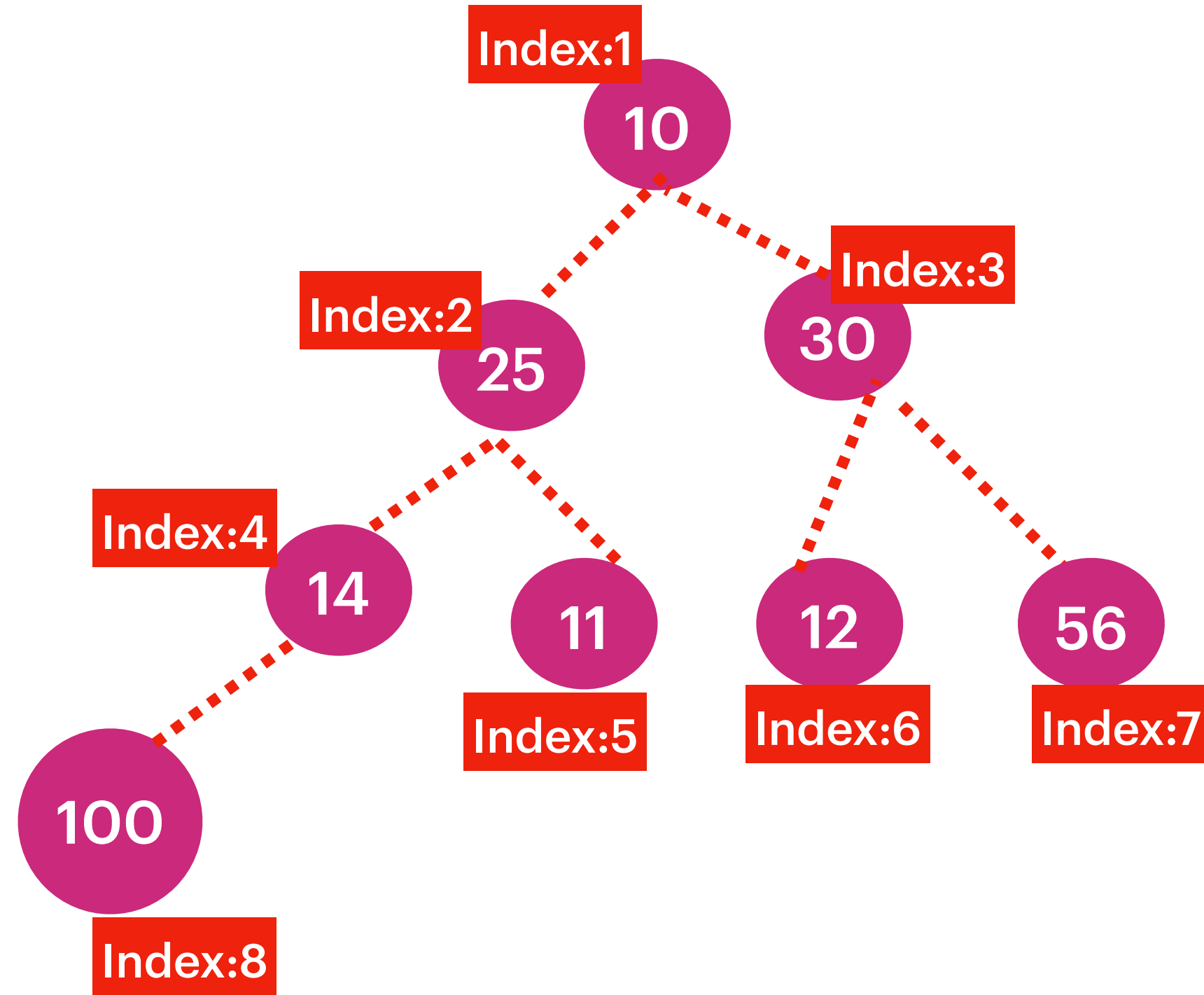


Complete Binary Tree



Parent = $\text{index}/2$;

Ex: Parent Of 25 = $2/2 = \text{index:1} = 10$

Transform Complete Binary Tree to the Hash/Array

How to you find parent ?

How do you find left & right child ?

How do you know that current element is a leaf Node ?

Left Child = $\text{index} * 2$

Right Child = $\text{index} * 2 + 1$

Ex: Left Child of Element 30 = $2 * \text{index:3} = \text{index:6} = 12$

Right Child of Element 30 = $2 * \text{index:3} + 1 = \text{index:7} = 56$

How do you know that current element is a leaf Node ?

LeafNode: $\text{index} > \text{size}/2$

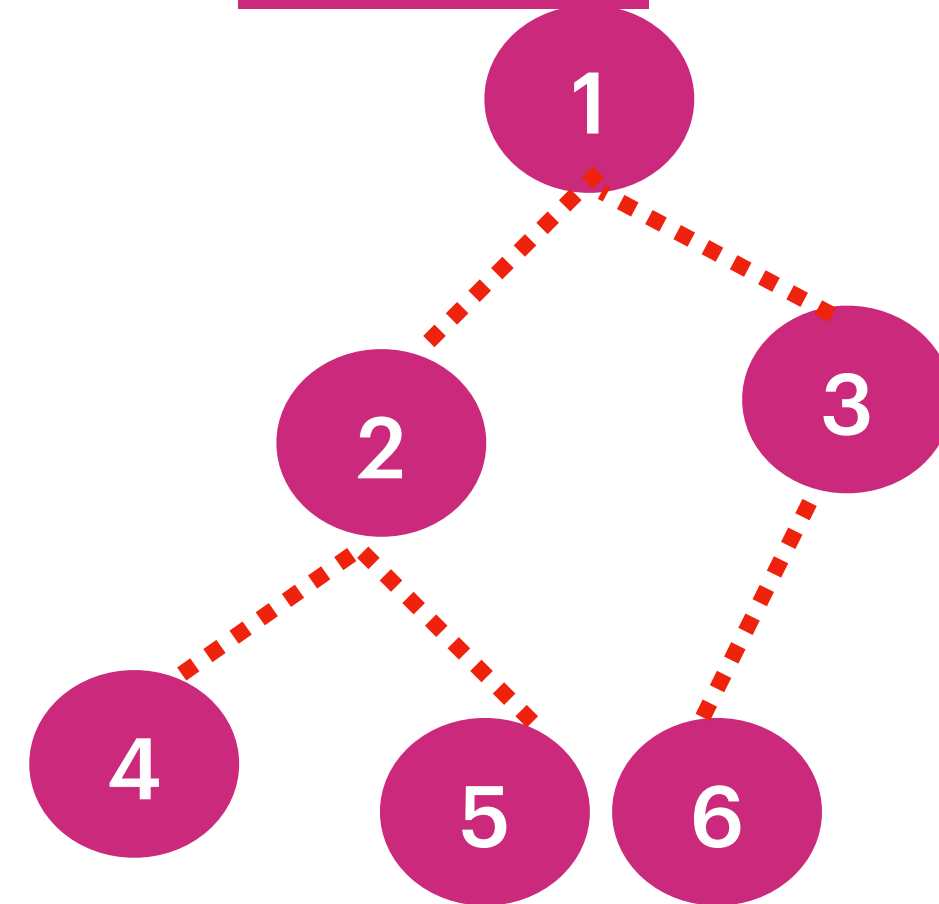


Min Heap

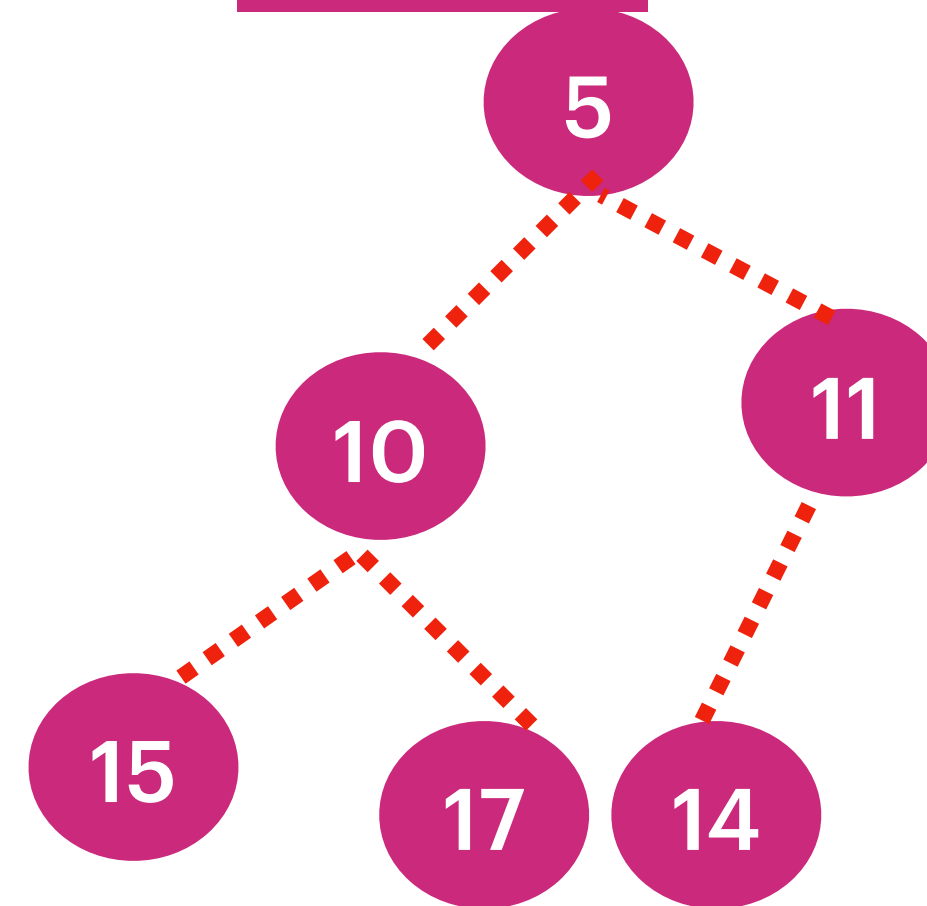
**** The Data Structure
should Complete Binary Tree**

**** Each Parent Node should be \leq it child node's**

Min Heap



Min Heap



Insert Operation On MinHeap

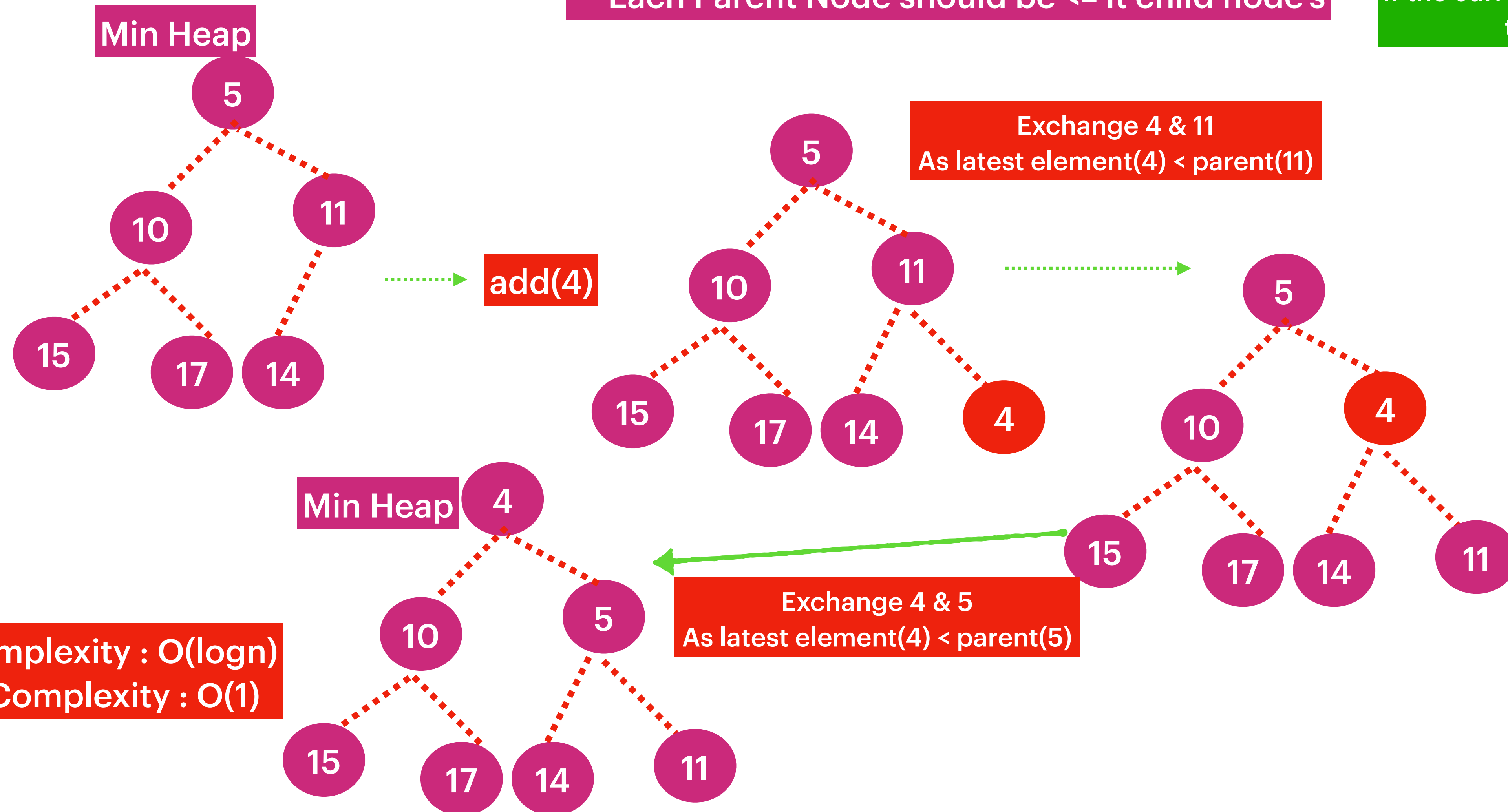
Min Heap

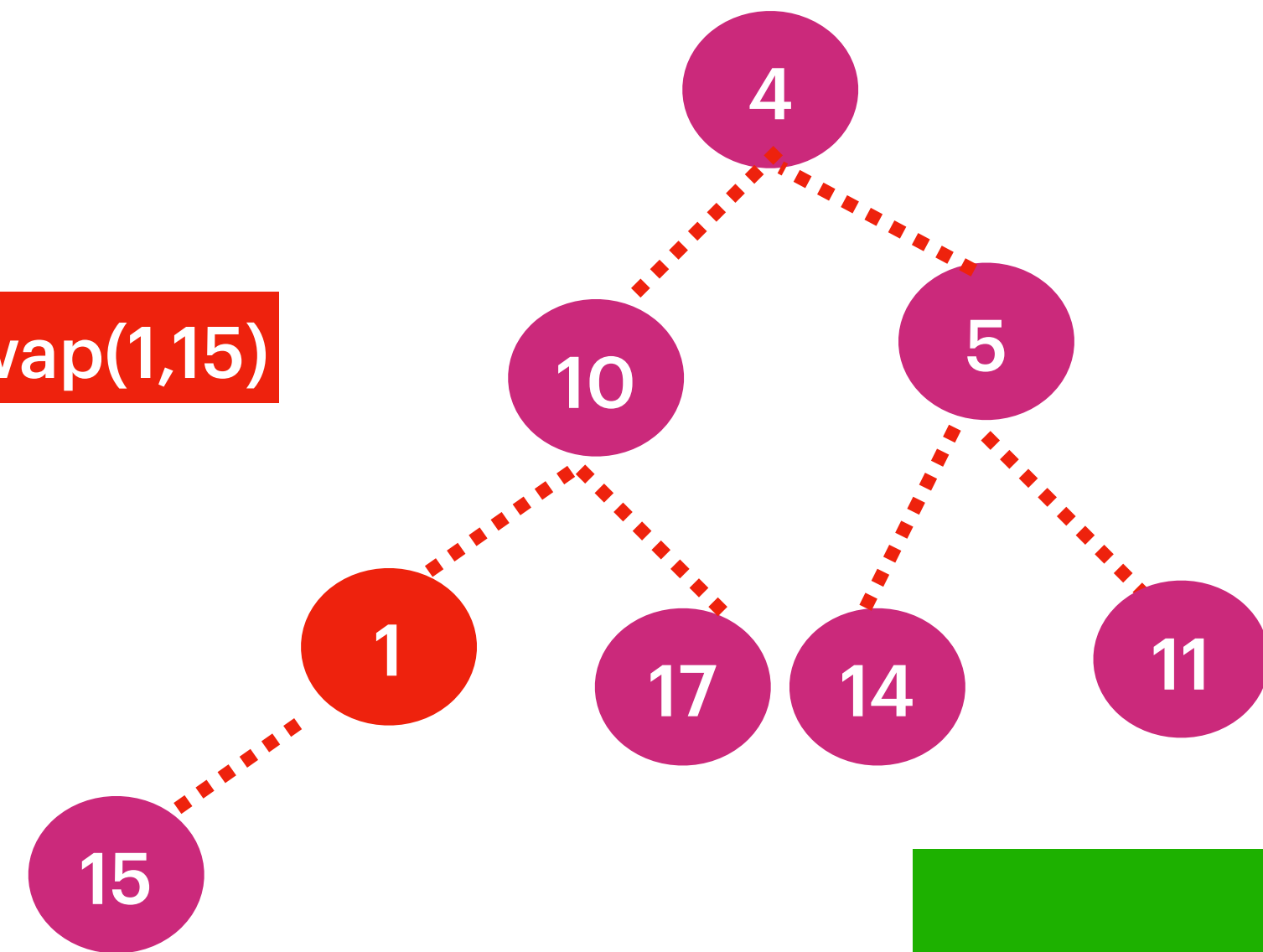
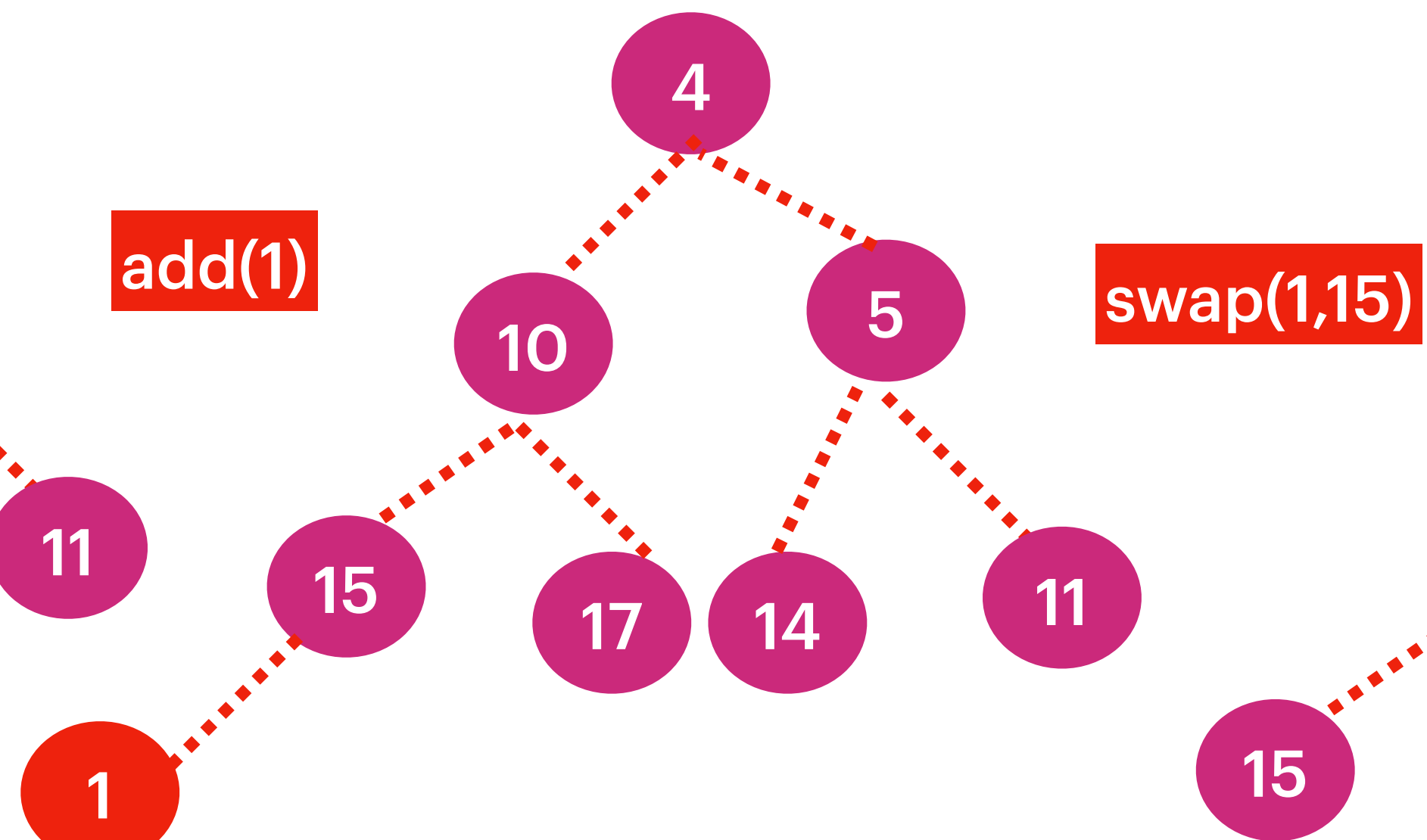
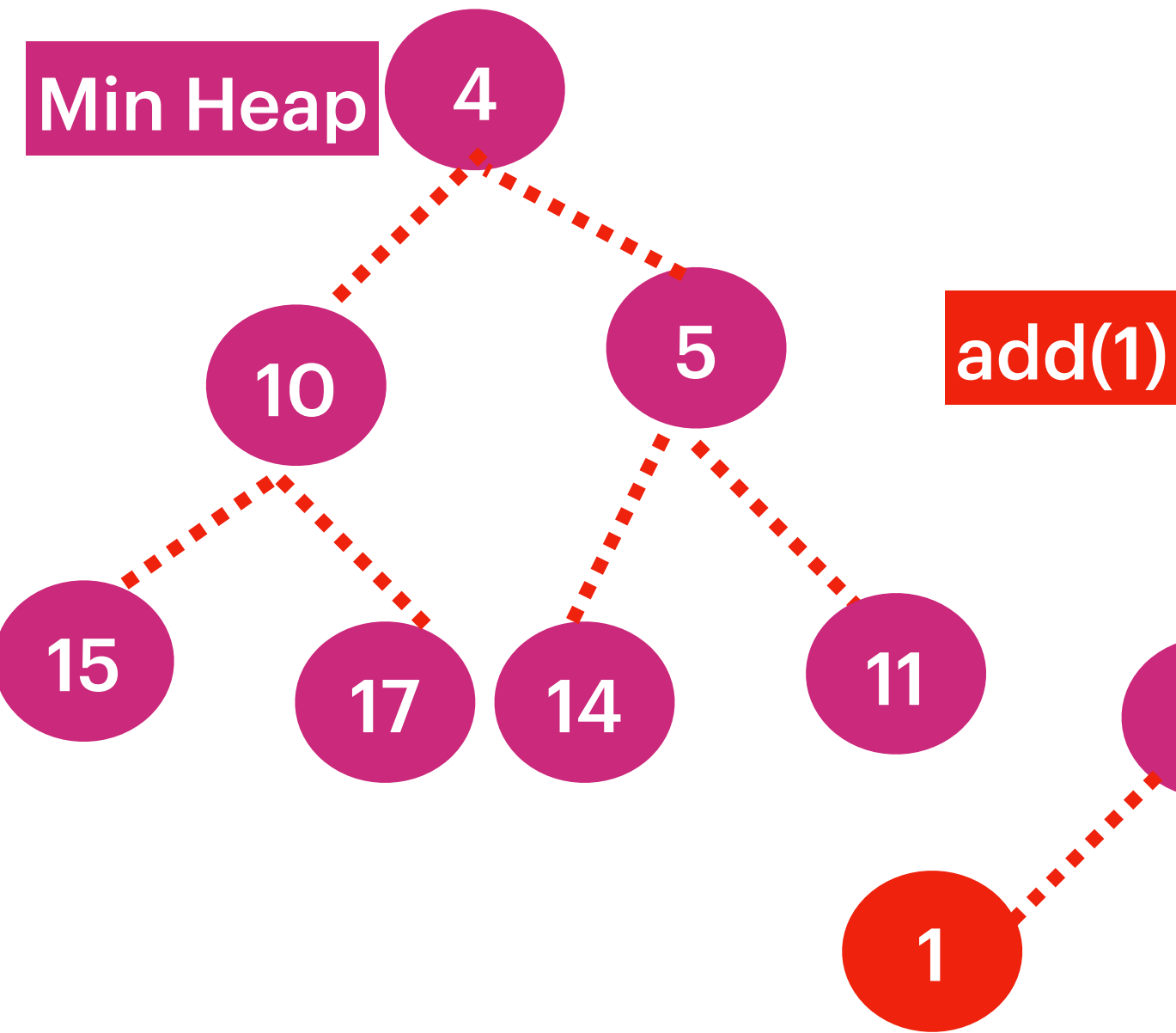
**** The Data Structure
should Complete Binary Tree**

**** Each Parent Node should be \leq it child node's**

Algorithm :

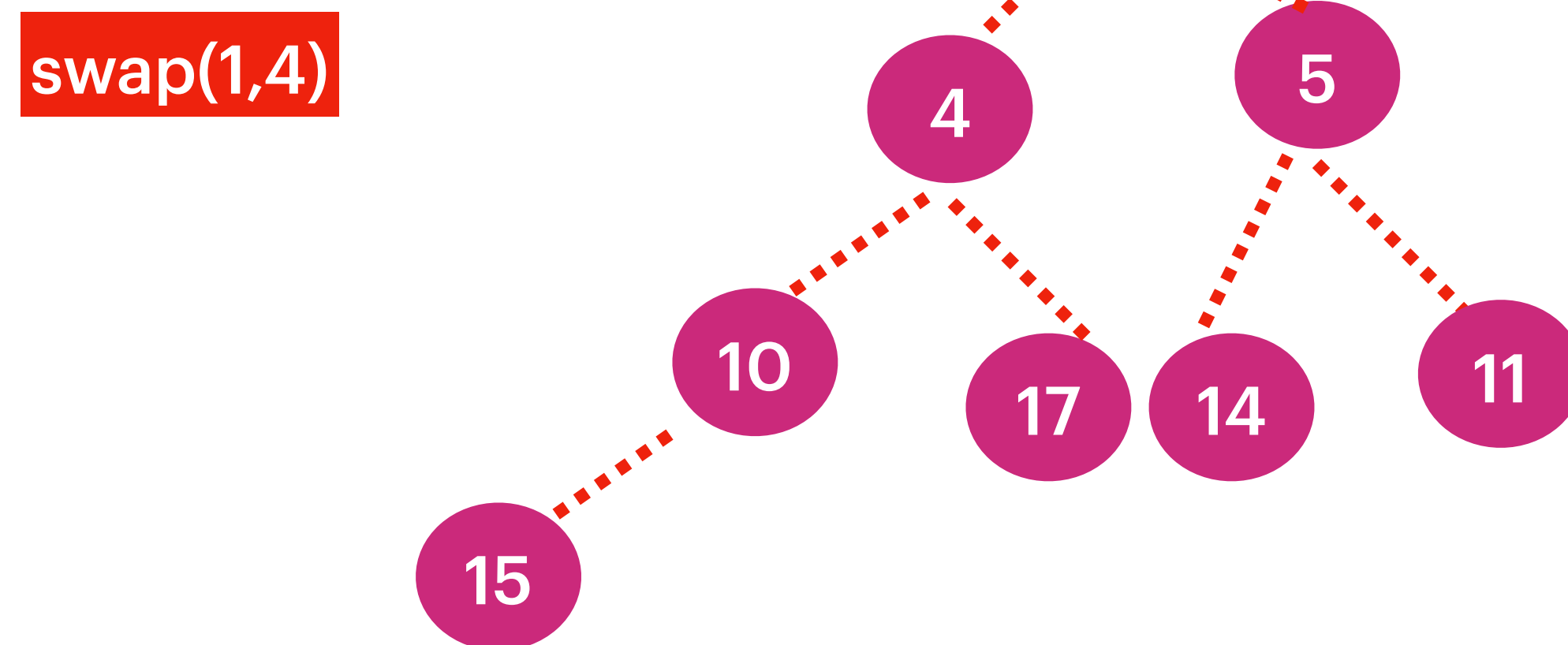
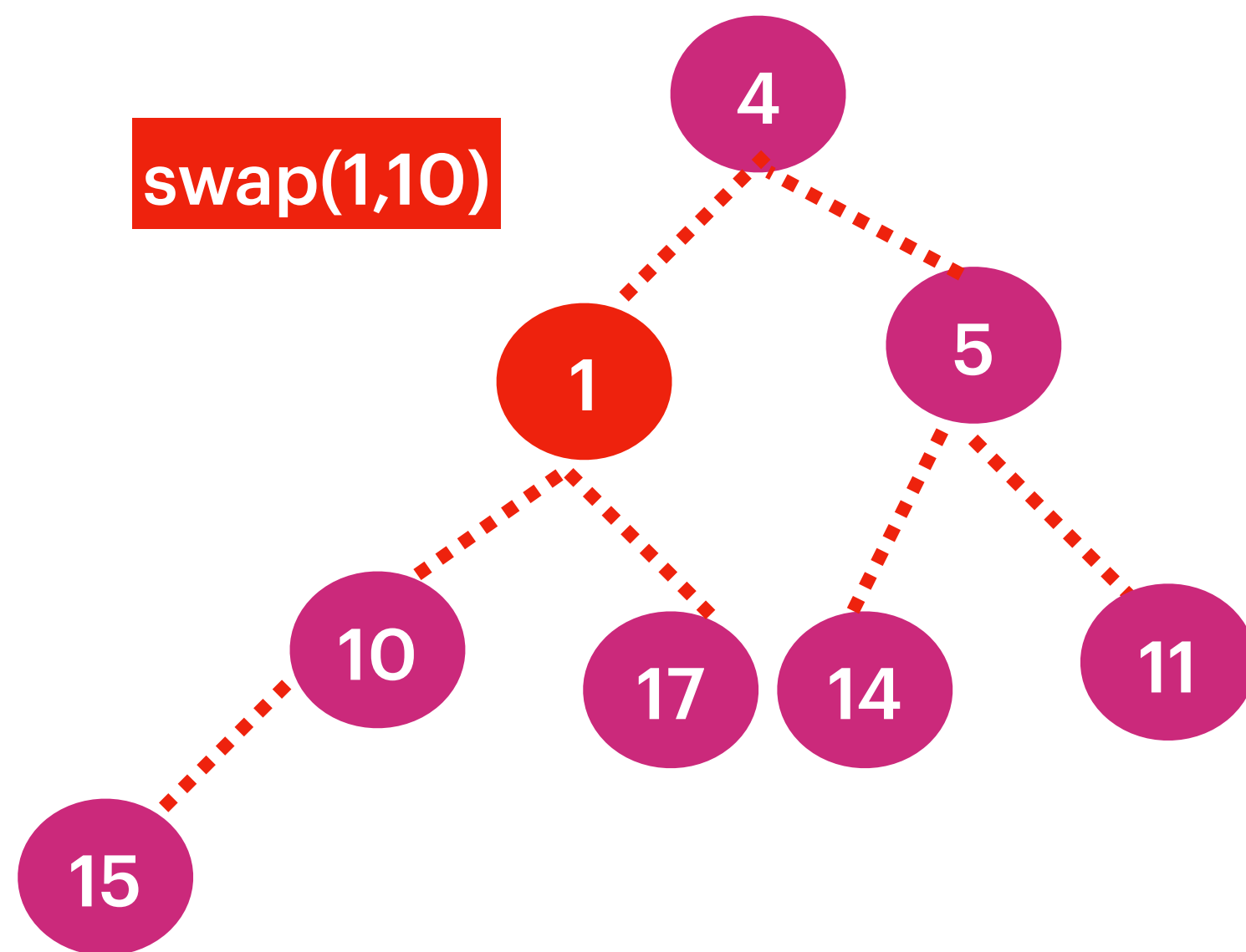
Add to the RightMost position.
If the current element $>$ parent
then swap.



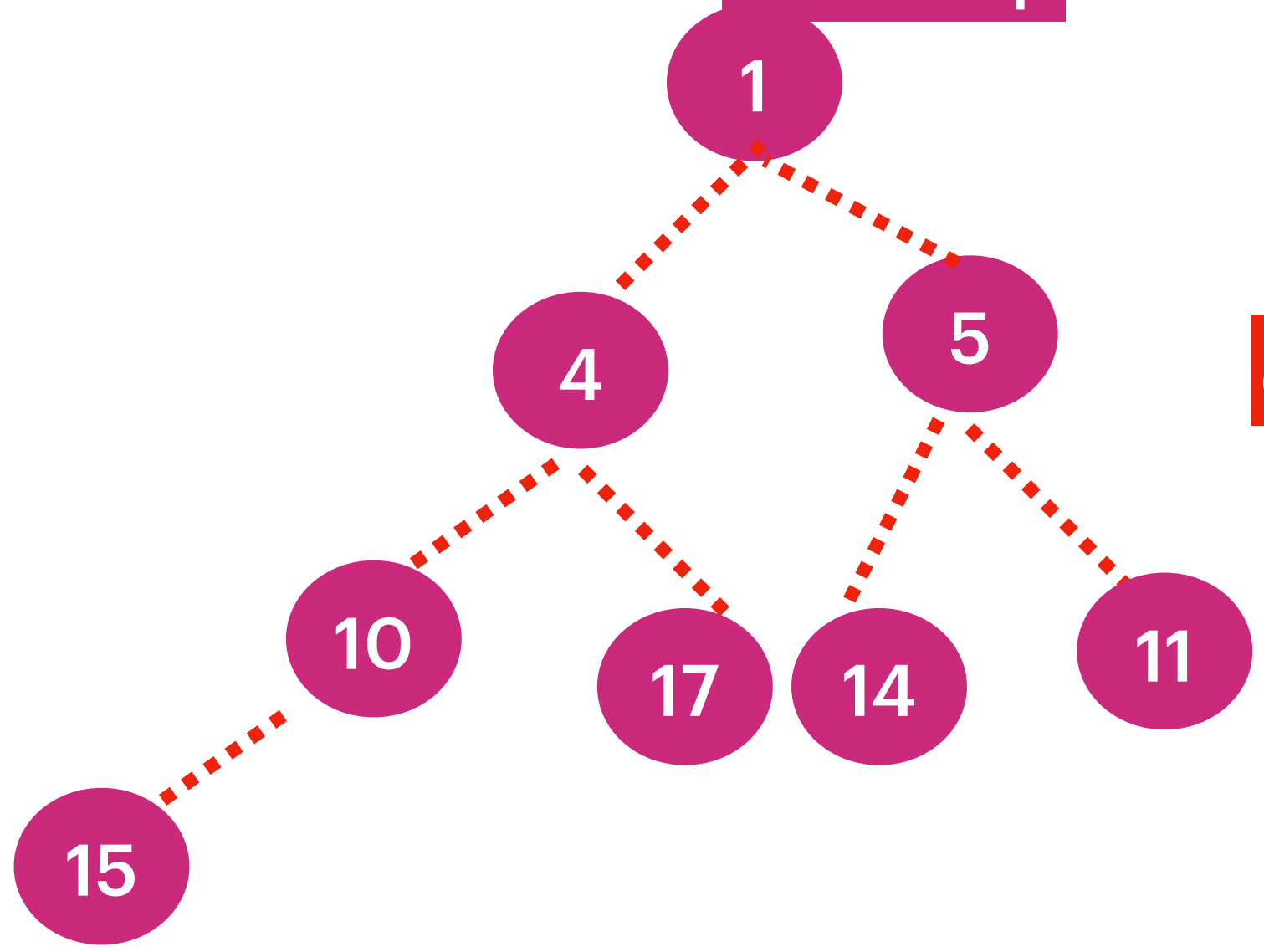


Algorithm :

Add to the RightMost position.
If the current element > parent
then swap.



Min Heap

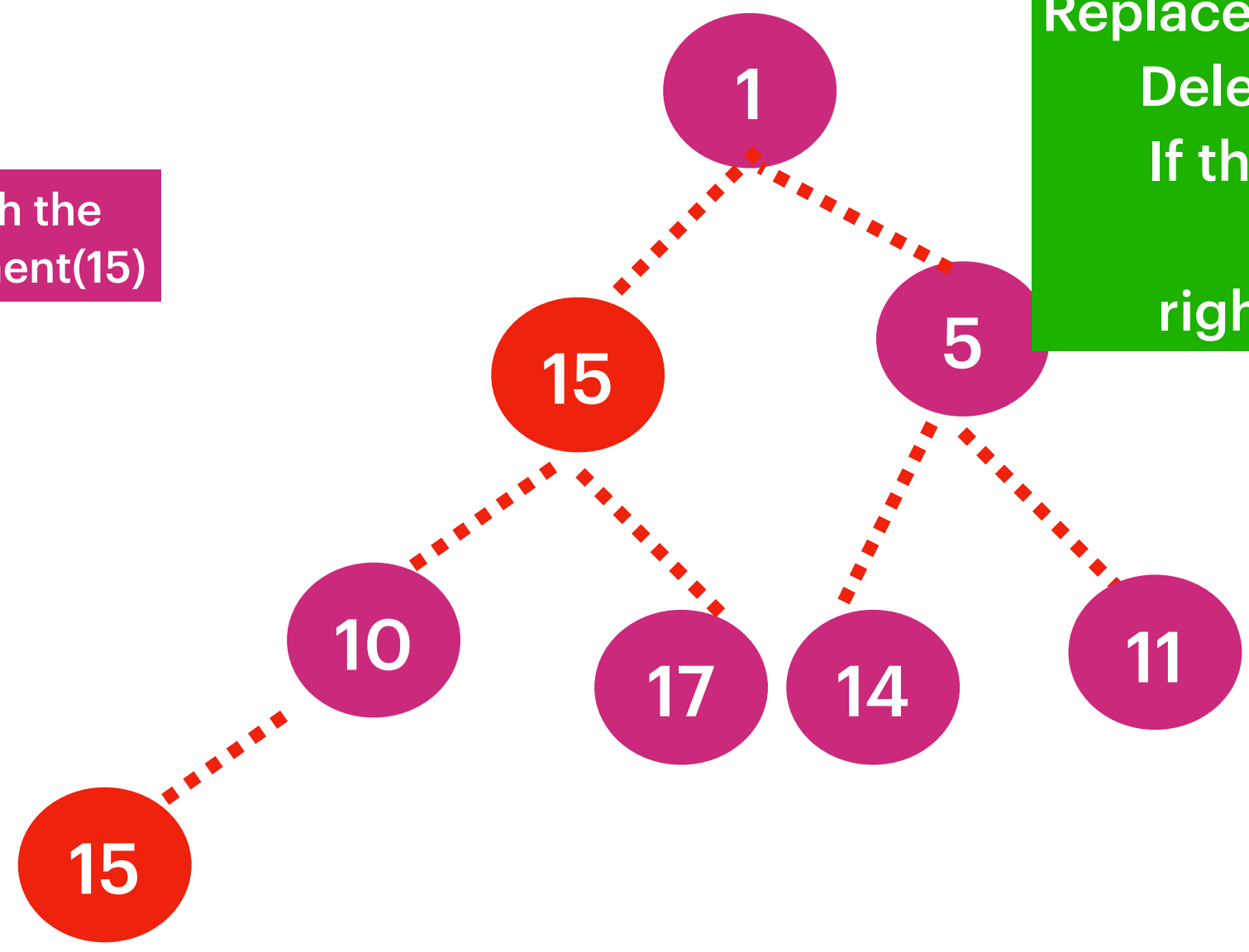


Delete Operation On MinHeap

delete(4)

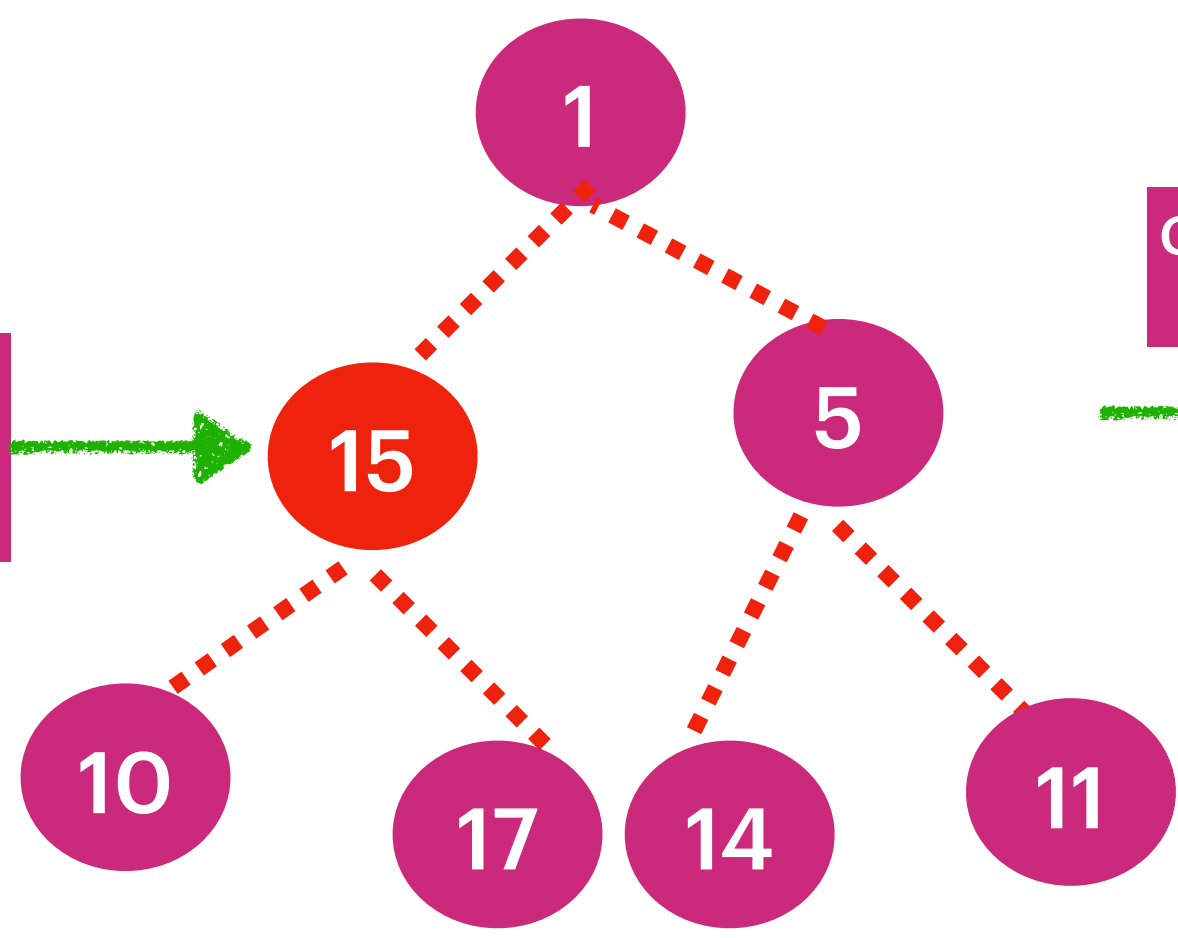


replace(4) with the
Right Most Element(15)

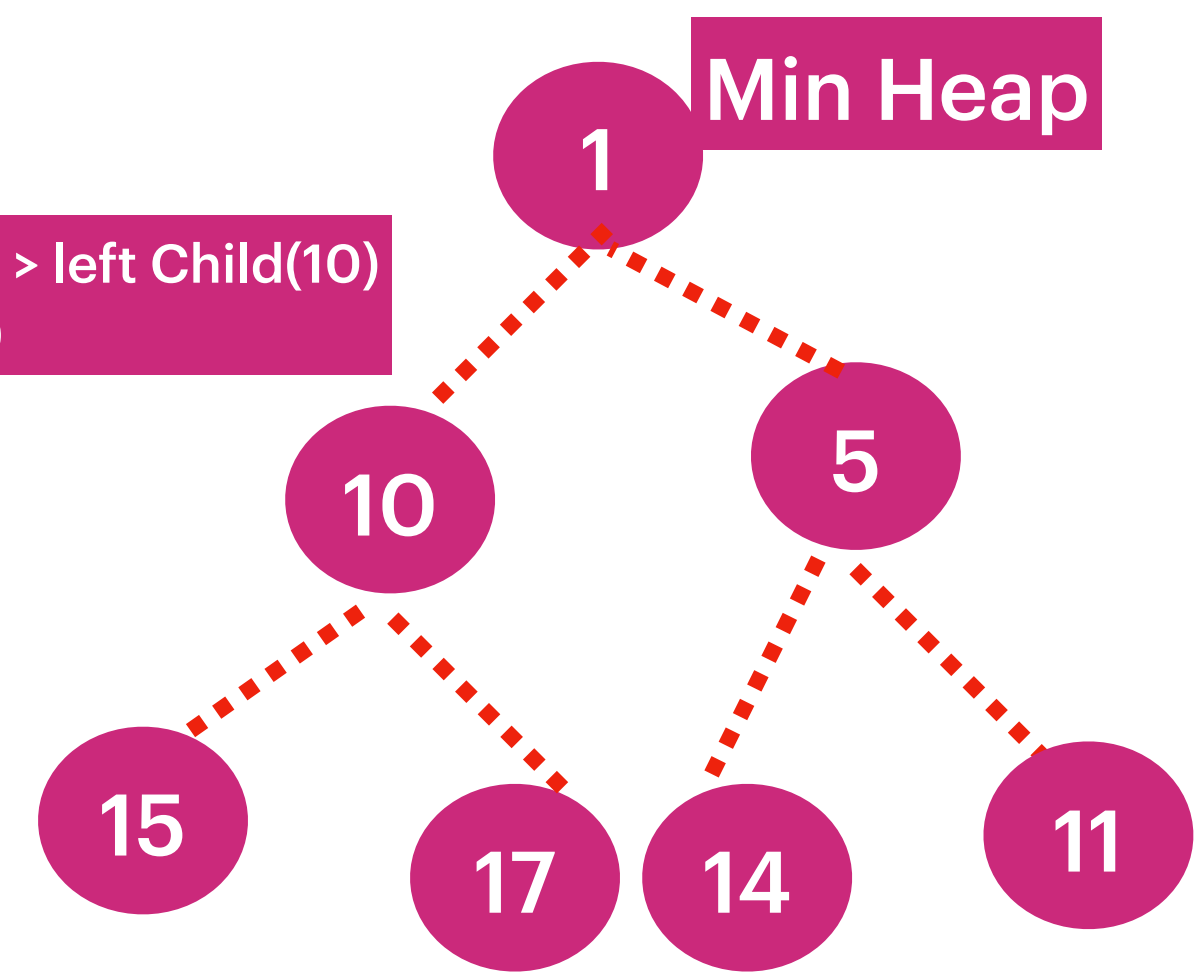


Algorithm :
Replace with the right most element ,
Delete the right most element.
If the current Element either >
(leftNode ||
rightNode) swap accordingly

Delete the
Right Most Element(15) in the
max level of MinHeap

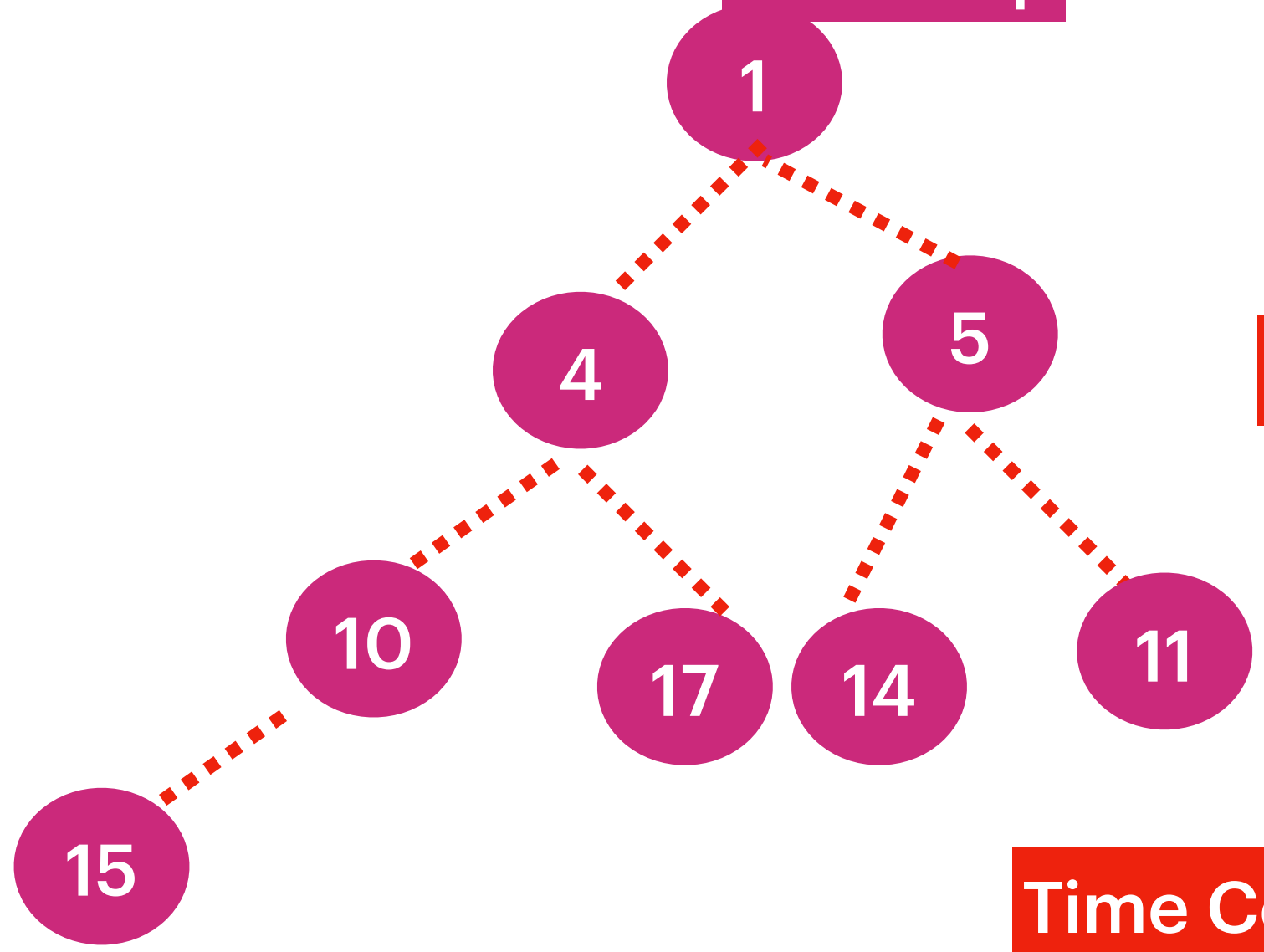


Current Parent(15) which is > left Child(10)
Swap (10,15)



Min Heap

Min Heap

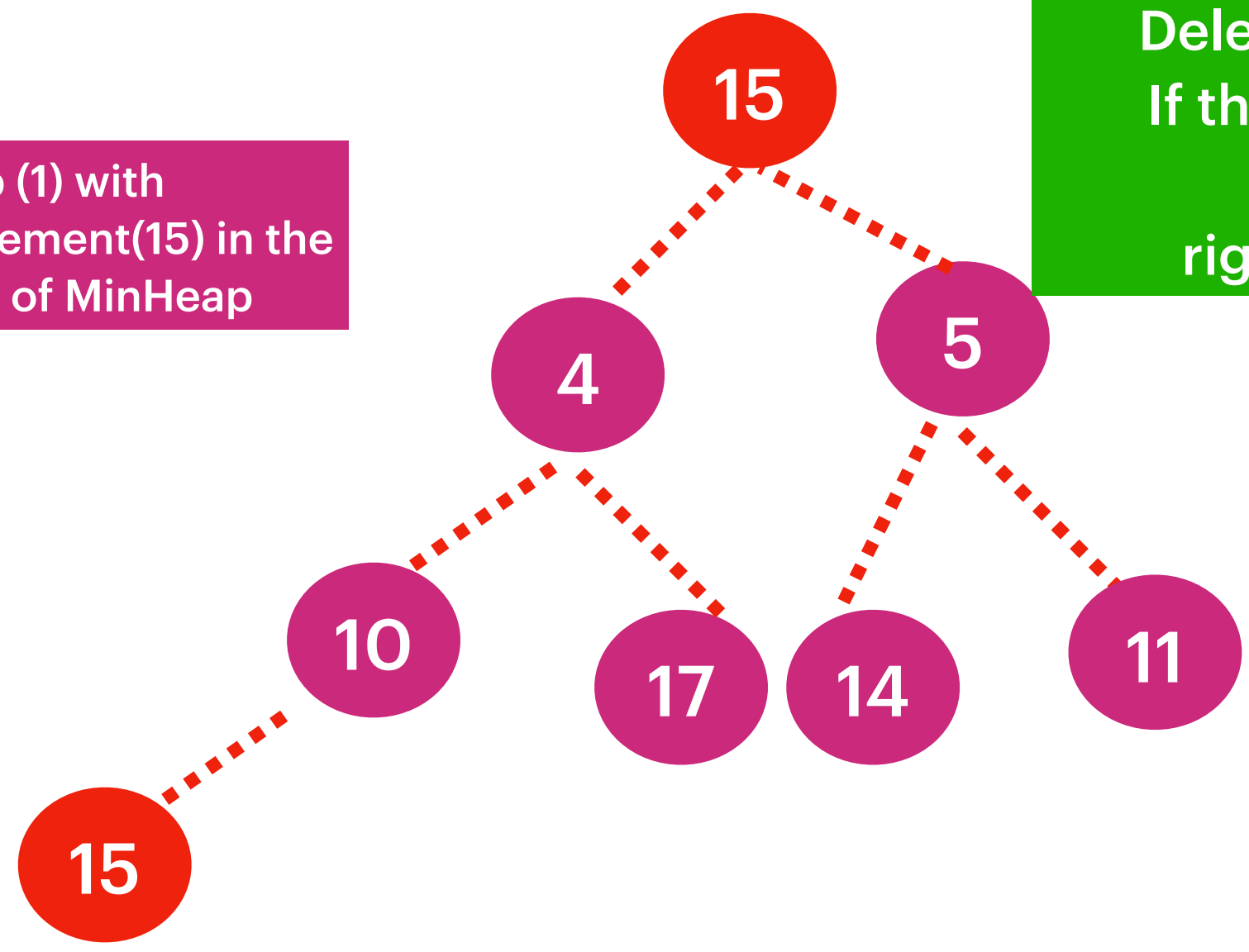


Delete Operation On MinHeap

delete(1)



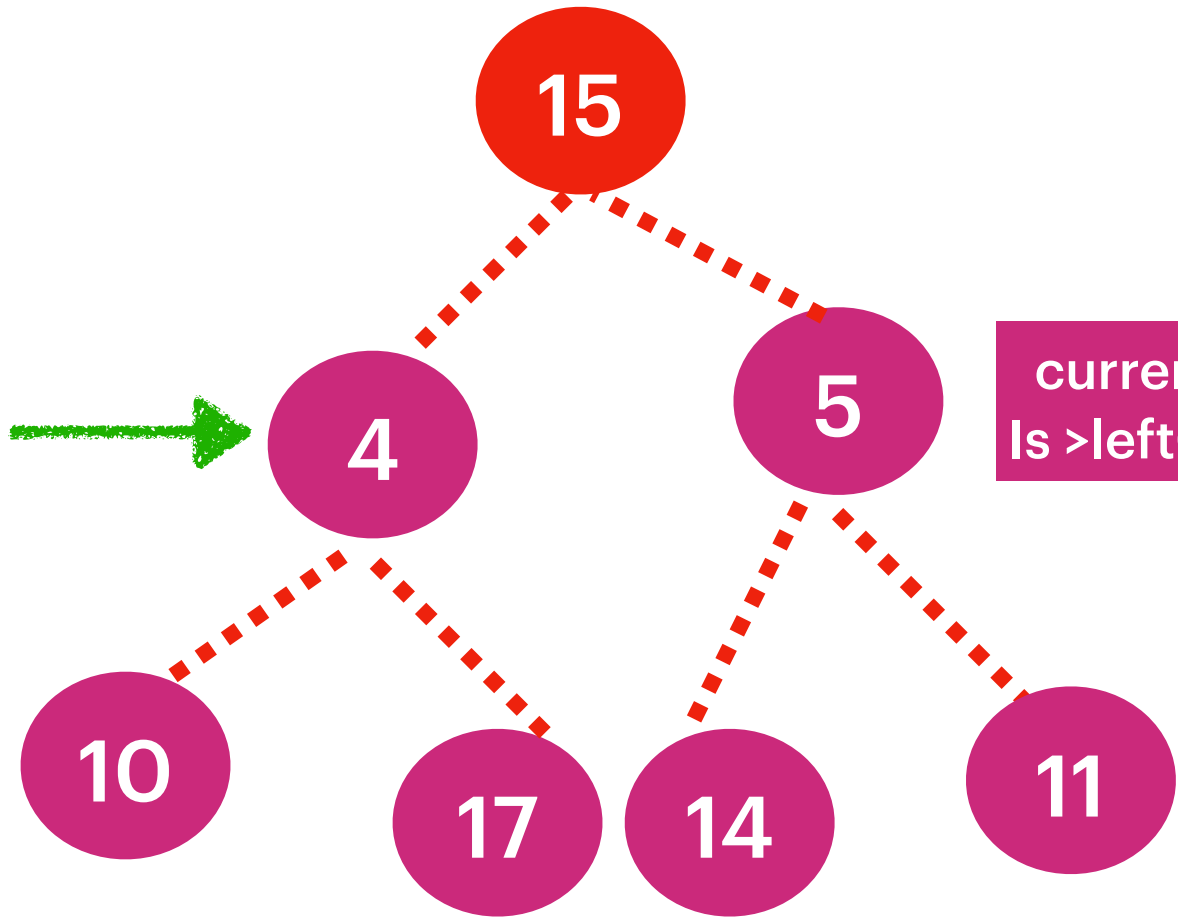
Swap (1) with
Right Most Element(15) in the
max level of MinHeap



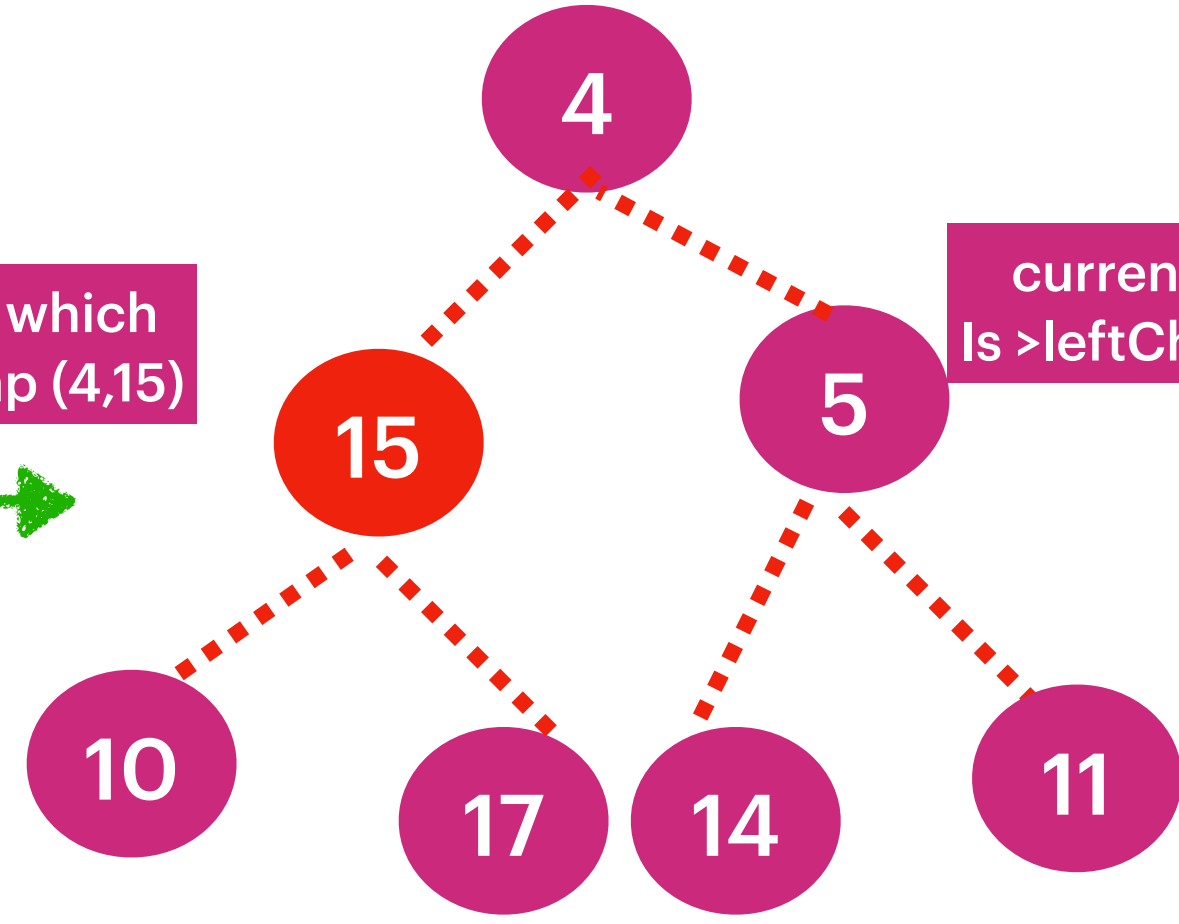
Algorithm :
Replace with the right most element ,
Delete the right most element.
If the current Element either >
(leftNode ||
rightNode) swap Iteratively

Time Complexity : $O(\log n)$
Space Complexity : $O(1)$

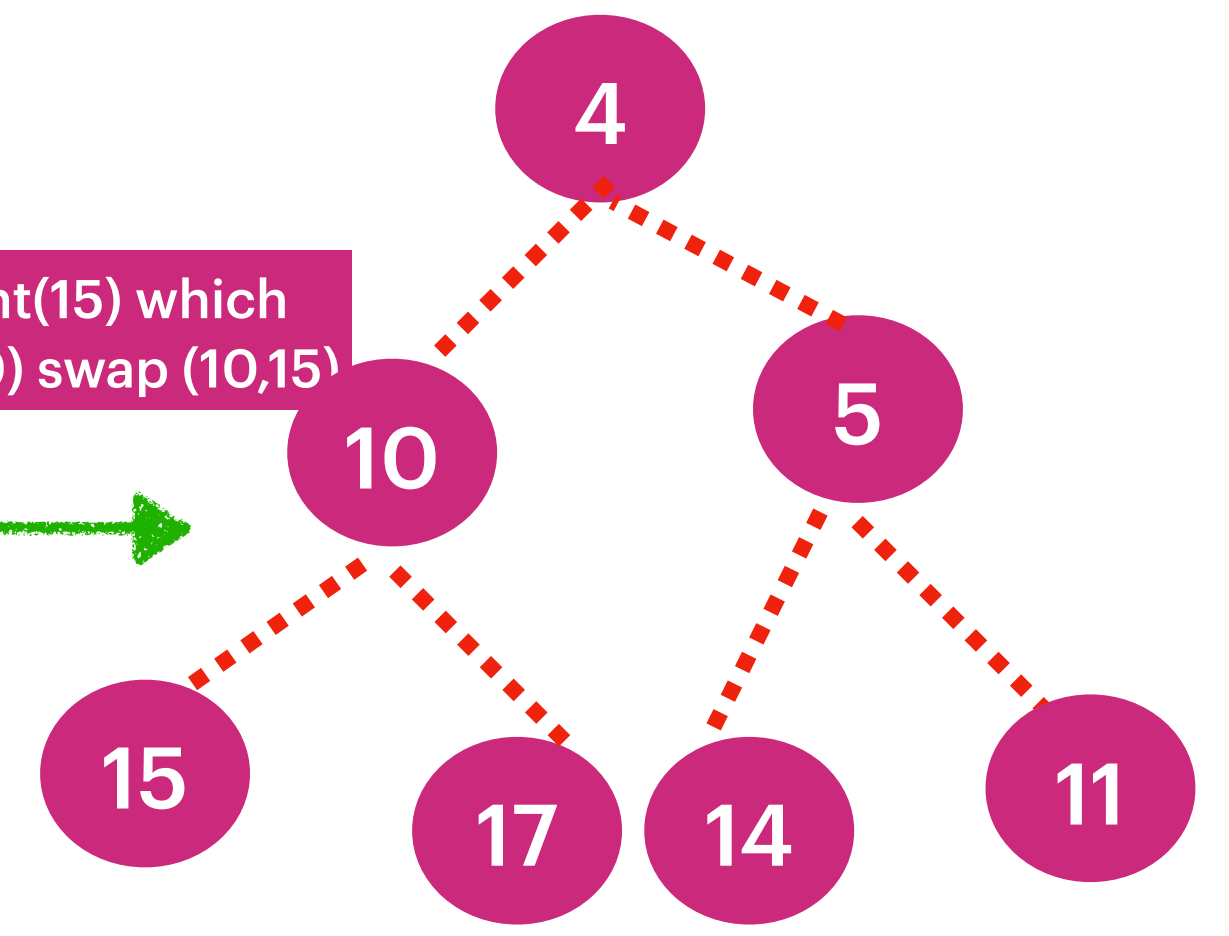
Delete the
Right Most Element(1) in the
max level of MinHeap



currentParent(15) which
Is > leftChild(4) swap (4,15)



currentParent(15) which
Is > leftChild(10) swap (10,15)



HeapSize : 5 [4,5,3,1,2]

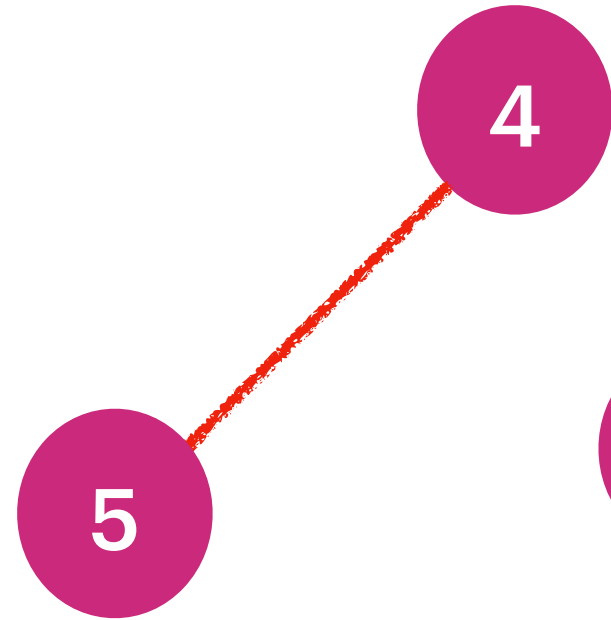
add(4)
add(5)
add(3)
add(1)
add(2)

Parent = index/2 ;
LeftChild = 2 * index
RightChild = 2 * index + 1
leafNode > realSize/2

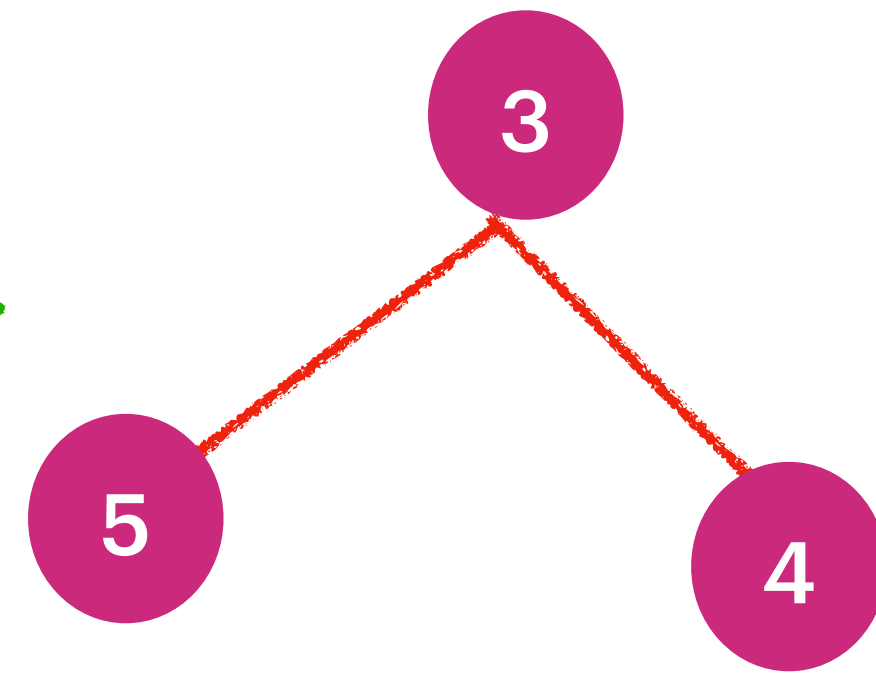
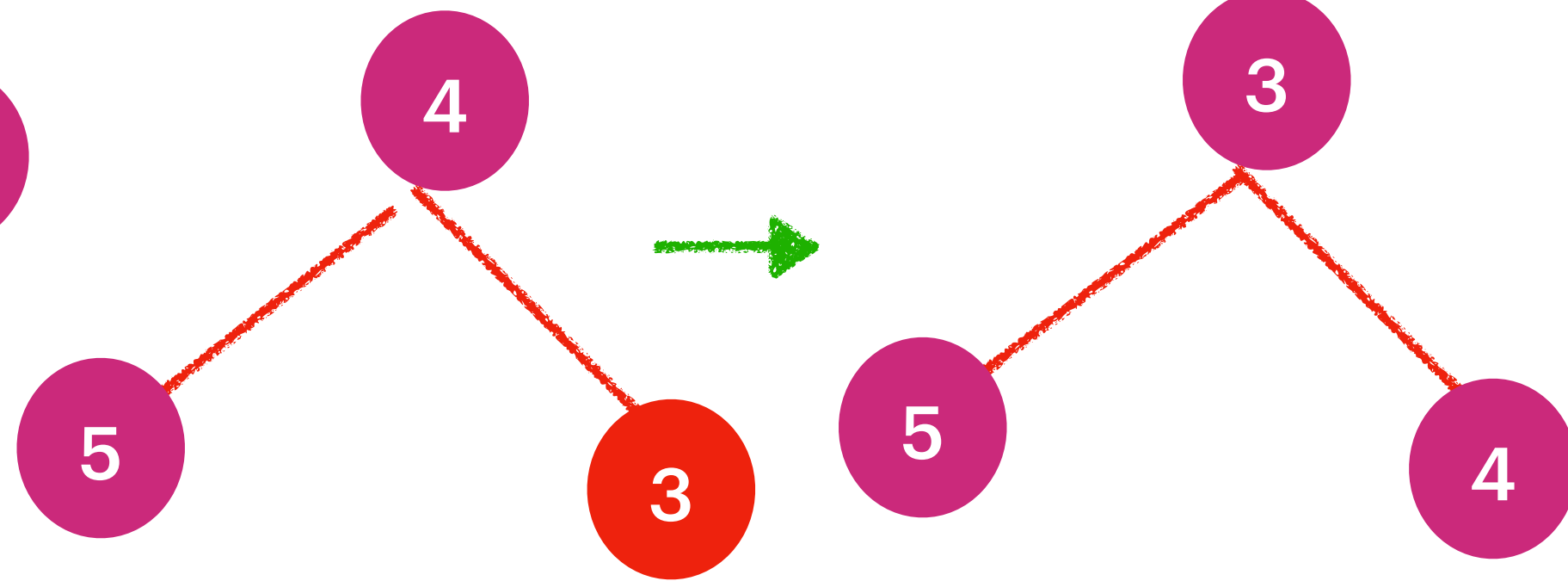
add(4)



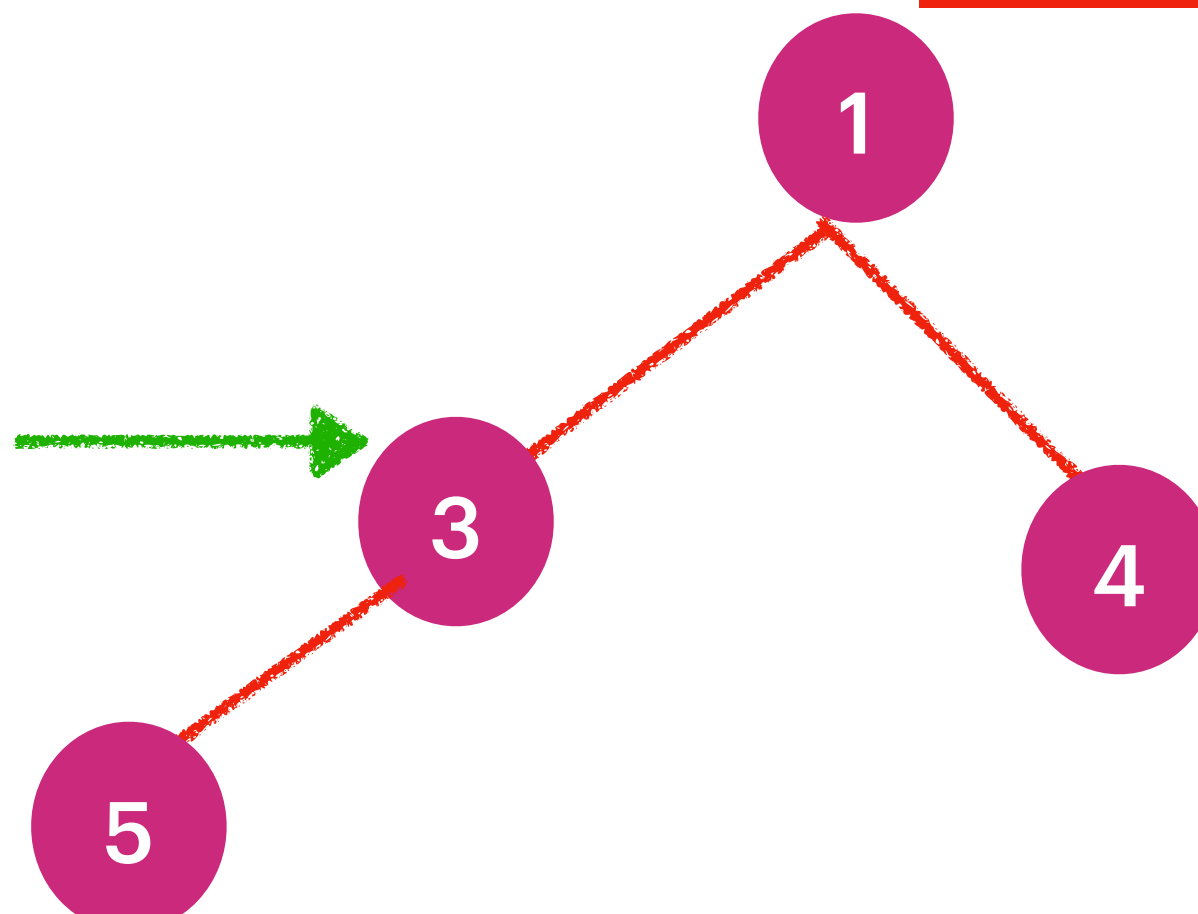
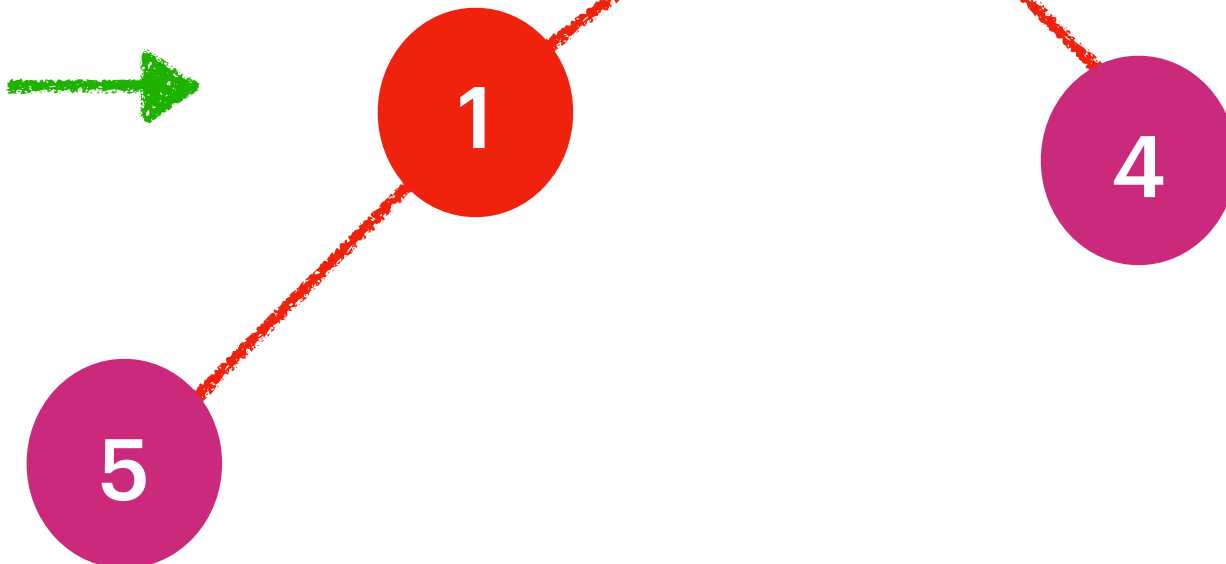
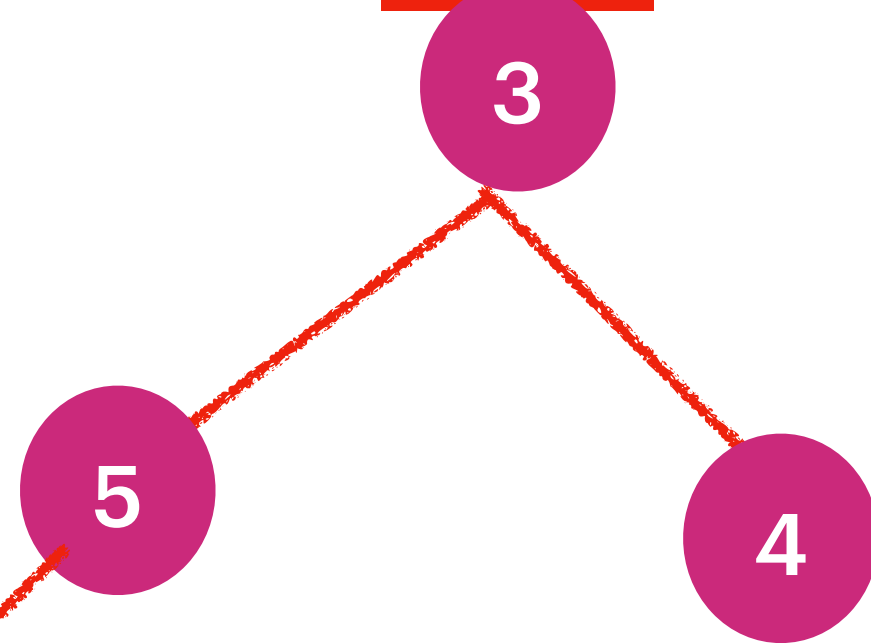
add(5)



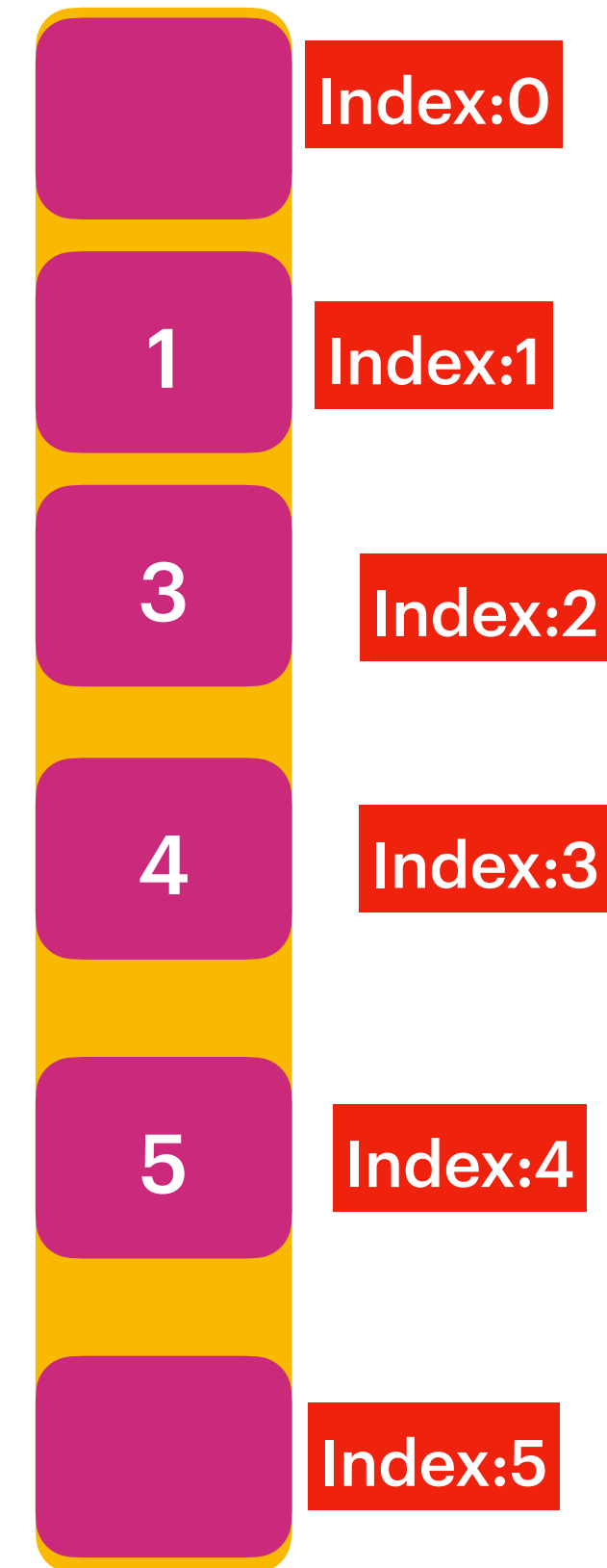
add(3)



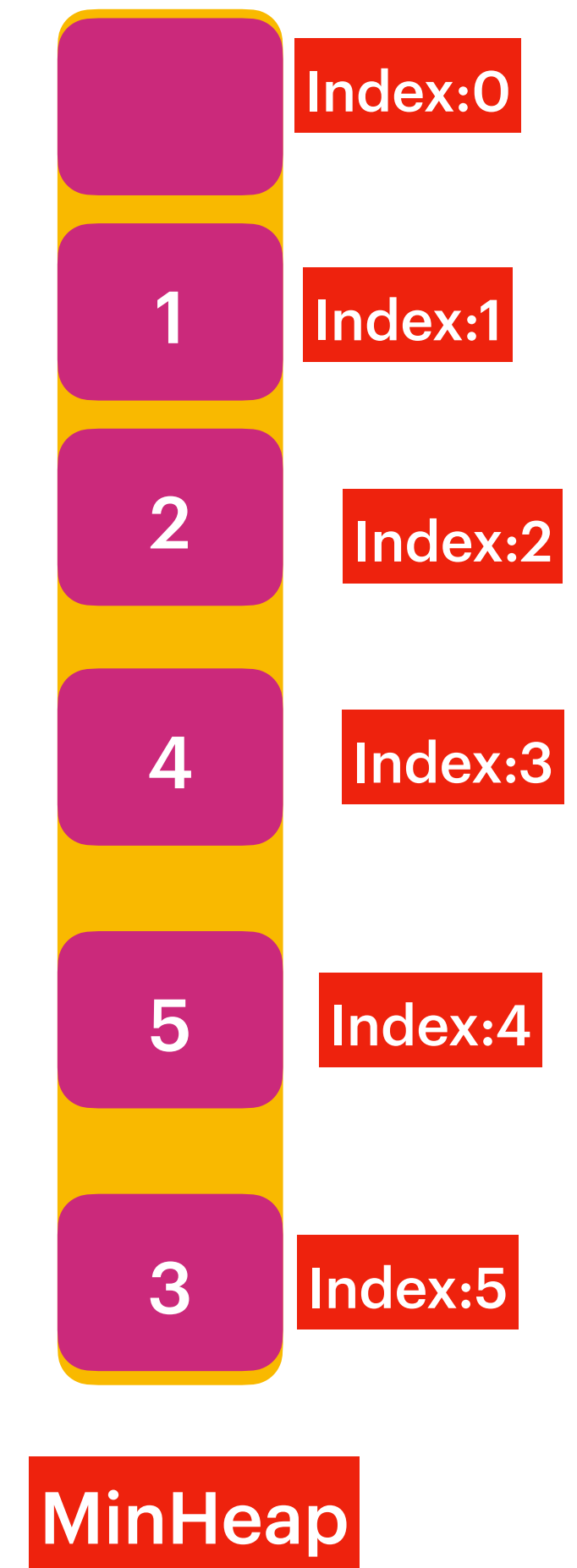
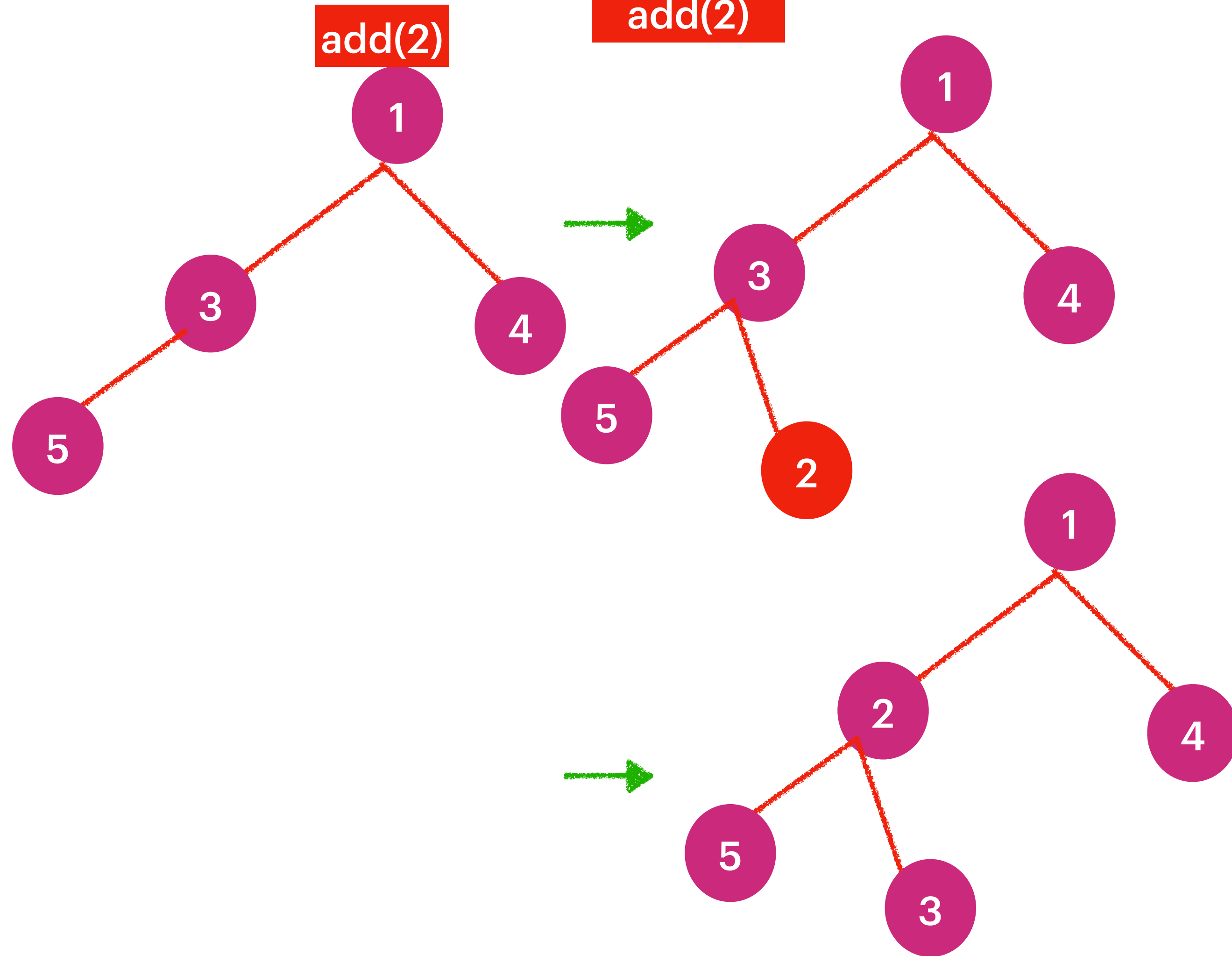
add(1)



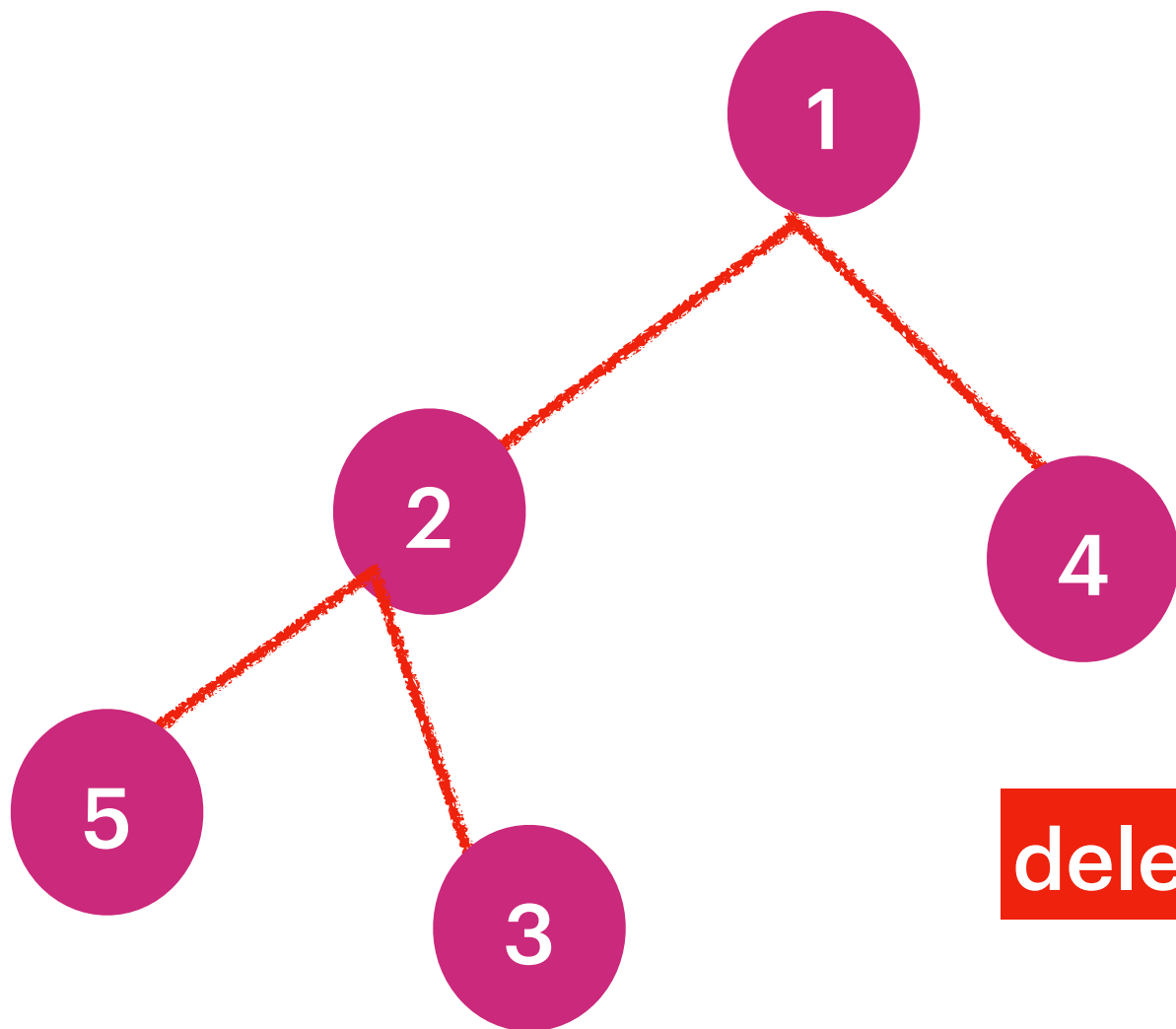
MinHeap



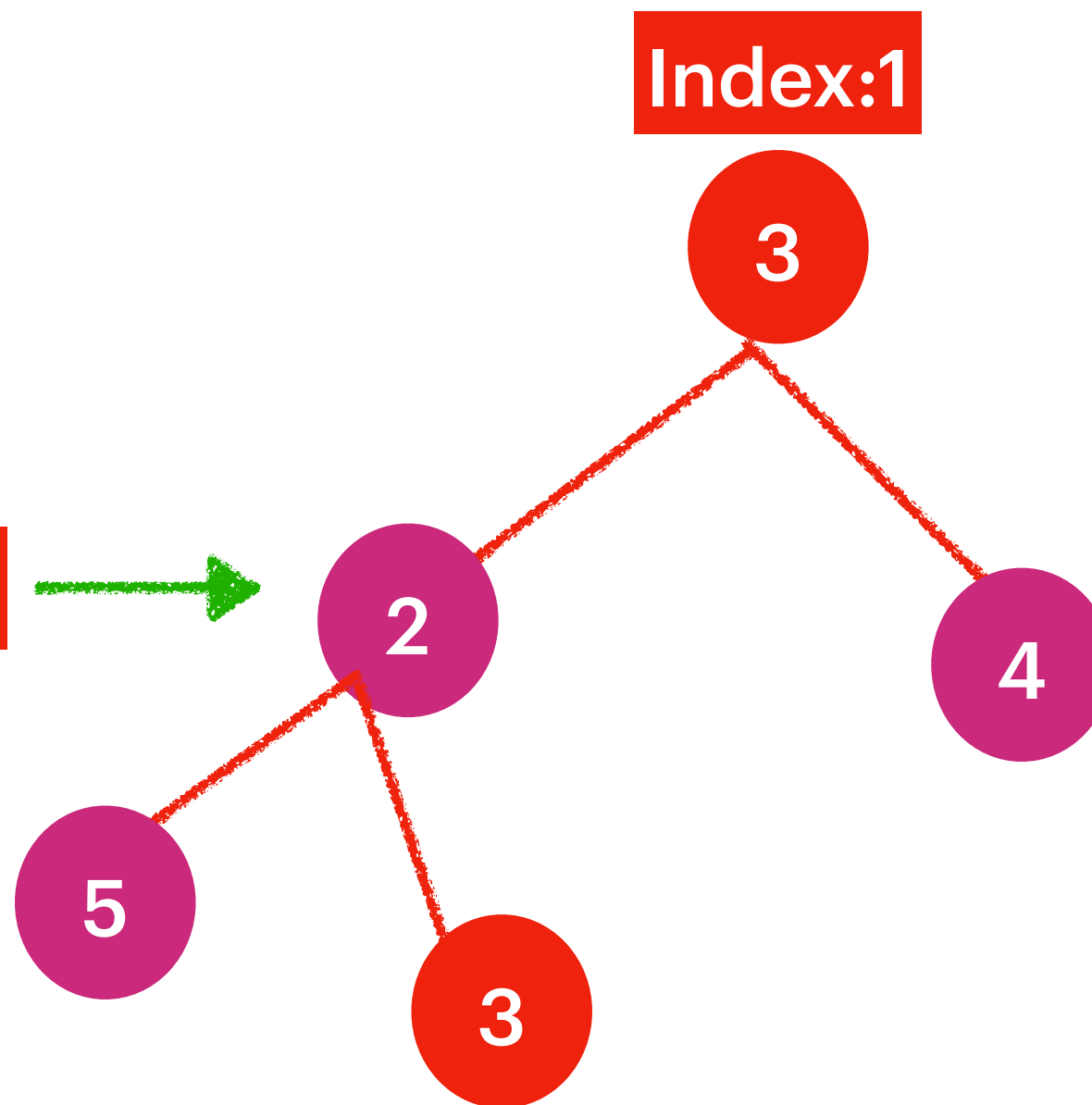
HeapSize : 5 [4,5,3,1,2]
add(4)
add(5)
add(3)
add(1)
add(2)



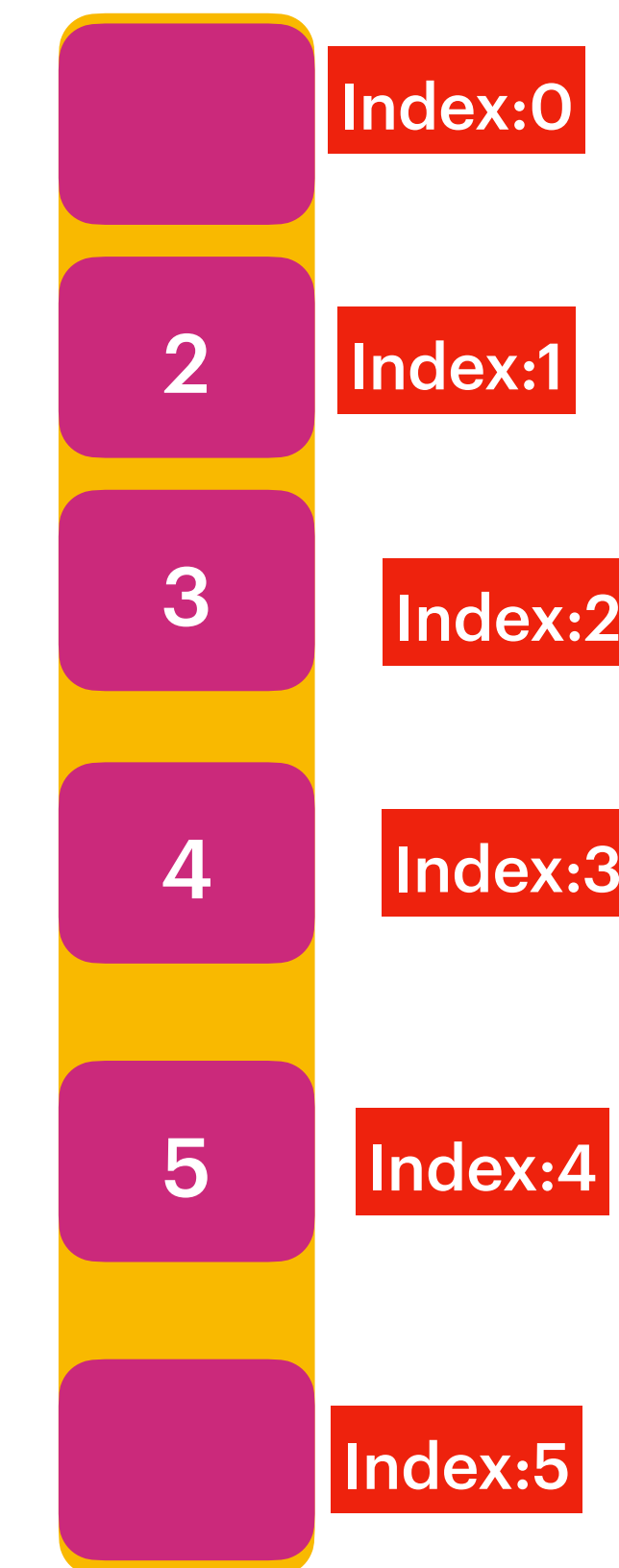
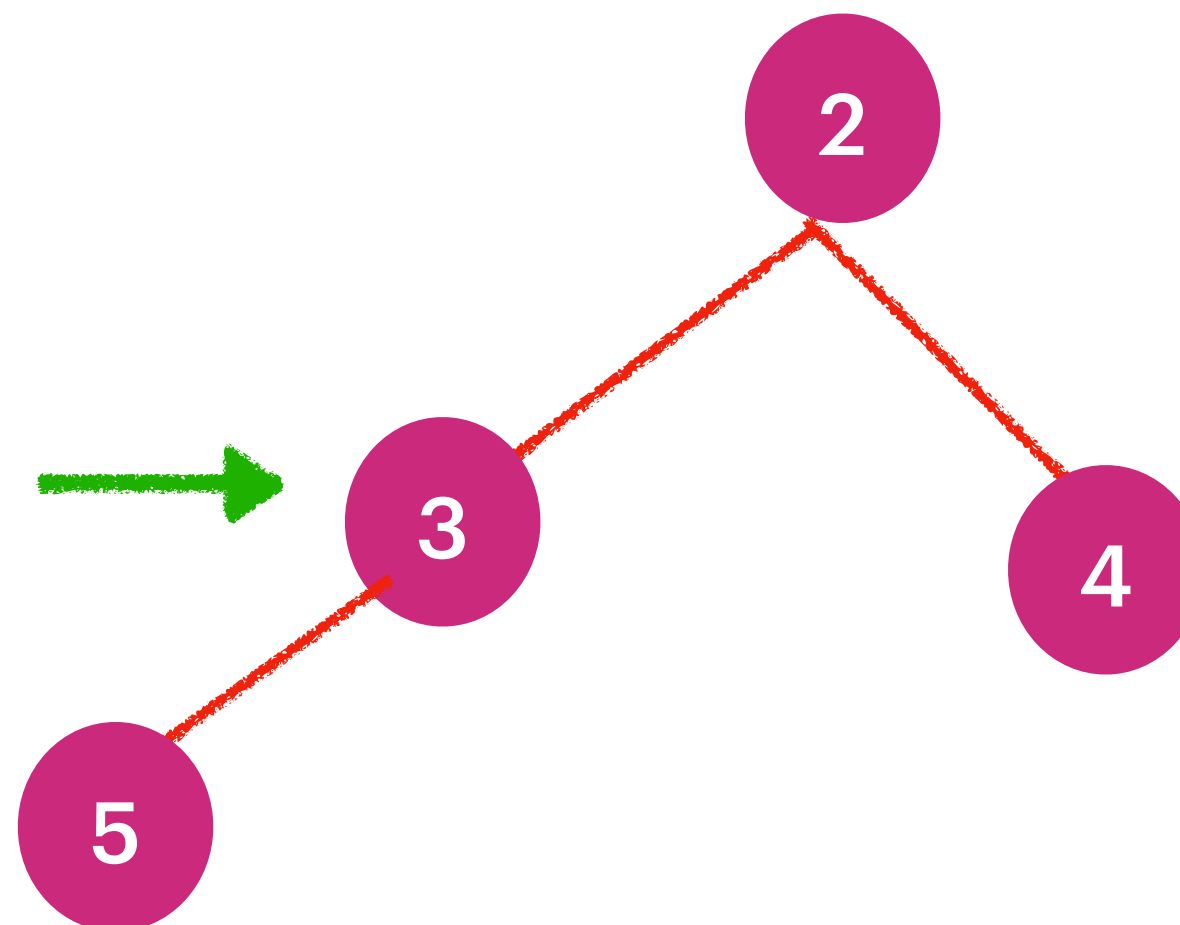
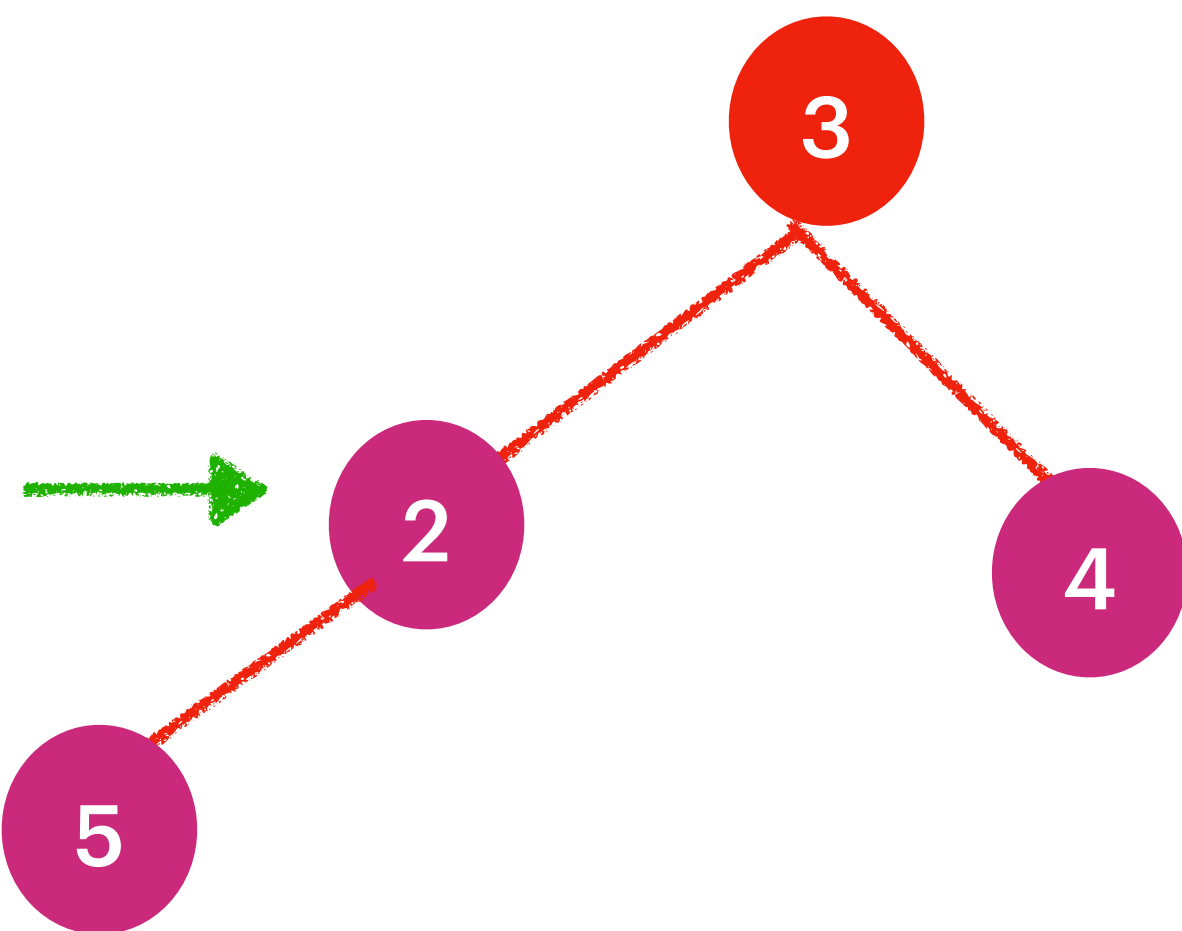
Identify the delete Element Index :
 $O(n) + \text{minHeap } \log(n) = O(n)$



delete(1)



Index:1



MinHeap

Identify the delete Element Index :
 $O(n) + \text{minHeap } \log(n) = O(n)$

