

① First & Second Welfare theorem (General ver.) ← Today's topic

Distinct concepts

* Competitive equilibrium: optimization given prices
(market clearing)

* Pareto optimality: welfare

Welfare theorems

- give a connection
- imposing strong conditions

Basic Building Block: Equilibrium with transfers

Economy: $E = ((Y_j)_{j=1}^J; (X_i, \bar{z}_i, W_i, \theta_{i1}, \dots, \theta_{iJ})_{i=1}^I)$

Def An allocation (x^*, y^*) and a price vector p^* form a competitive equilibrium

if ① for each $j=1, \dots, J$, $p^* \cdot y_j^* \geq p^* \cdot y_j$ for each $y_j \in Y_j$ [profit max]

② for each $i=1, \dots, I$

2a) $x_i^* \in B_i^*(p^*) = \{x_i \in X_i : p^* \cdot x_i \leq p^* \cdot \bar{w}_i + \sum_{j=1}^J \theta_{ij} (p^* \cdot y_j^*)\}$ [util. max]

2b) $x_i^* \succeq_i x_i$ for each $x_i \in B_i^*(p^*)$

③ $\sum_{i=1}^I x_i^* = \bar{w} + \sum_{j=1}^J y_j^*$ [Market clearing]

Def An allocation (x^*, y^*) and a price vector p^* form a price equilibrium with transfer

if \exists an assignment of wealth $W = (W_1, \dots, W_I)$ s.t.

① for each $j=1, \dots, J$, $p^* \cdot y_j^* \geq p^* \cdot y_j$ for each $y_j \in Y_j$ [profit max]

② for each $i=1, \dots, I$ 2a) $x_i^* \in B_i(p^*, W_i) = \{x_i \in X_i : p^* \cdot x_i \leq W_i\}$ [util. max]

2b) $x_i^* \succeq_i x_i$ for all $x_i \in B_i(p^*, W_i)$ ↑ Wealth for i

$$\begin{aligned} \textcircled{3} \quad \sum_{i=1}^I x_i^* &= \bar{w} + \sum_{j=1}^J y_j^* \\ \textcircled{4} \quad \sum_{i=1}^I W_i &= p^* \cdot \bar{w} + \sum_{j=1}^J p^* \cdot y_j^* \end{aligned} \quad \left. \vphantom{\sum_{i=1}^I x_i^*} \right\} \text{[Market clearing]}$$

Remark A competitive equilibrium induces a price equilibrium with transfers

$$W_i = p^* \cdot w_i + \sum_{j=1}^J \theta_{ij} p^* \cdot y_j^*$$

\Downarrow

Check $\textcircled{4}$: $\sum_i W_i = p^* \cdot \bar{w} + \sum_{j=1}^J p^* \cdot y_j^*$ (Summation of all i's wealth)

First Welfare Theorem \star C.E. \Rightarrow P.O. Under \sum_i LNS

Idea: Any competitive equilibrium induces a Pareto allocation.

Def A preference relation \succeq_i on $X_i \subseteq \mathbb{R}_+^L$ satisfies Locally NON-satiation if for each $x_i \in X_i$ and each $\varepsilon > 0 \exists$ some $x_i' \in X_i$ s.t.

$$(1) \quad \|x_i - x_i'\| < \varepsilon \quad (2) \quad x_i' \succ x_i$$

c.f. "Rules out":
of LNS $\left\{ \begin{array}{l} \text{fat indifference curve} \\ \text{bliss point preference} \end{array} \right.$

Theorem: First Welfare Theorem Under \sum_i LNS, $\left\{ \begin{array}{l} \text{Price Eq w/ transfers} \\ \text{Competitive eq} \end{array} \right. \Rightarrow \text{P.O.}$

Fix an economy with locally non-satiated preferences.

If $((x^*, y^*), p^*)$ form a price equilibrium with transfers.

then (x^*, y^*) is Pareto Optimal allocation.

\hookrightarrow Corollary: This holds for a competitive equilibrium allocation.

Lemma 1 Suppose \succeq_i is locally non-satiated and let $x_i^* \in B_i(p, w_i)$

that \succeq_i - maximal, i.e., $x_i^* \succeq_i x_i$ for each $x_i \in B_i(p, w_i)$.

Then $x_i \succeq_i x_i^*$ implies $p \cdot x_i \geq w_i$

proof) * If $x_i \succ_i x_i^*$, $\Rightarrow x_i \notin B_i(p, w_i) \Rightarrow p \cdot x_i > p \cdot x_i^* = w_i$

* If $x_i \sim x_i^*$,

LNS: for each $n \geq 1$, there exists some $x_i^n \in X_i$ with $\|x_i^n - x_i\| < \frac{1}{n}$

and $x_i^n \succ_i x_i$

\Rightarrow By transitivity, $x_i^n \succeq_i x_i \sim x_i^*$ i.e., $x_i^n \succeq_i x_i^* \Rightarrow p \cdot x_i^n > w_i$

$\lim_{n \rightarrow \infty} x_i^n = x_i$, so $\lim_{n \rightarrow \infty} p \cdot x_i^n \geq w_i$

\parallel
 $p \cdot x_i$

Thus, $p \cdot x_i \geq w_i$ \parallel

proof of Theorem)

Let (x^*, y^*, p^*) be a price equilibrium with transfers that is supported by (w_1, \dots, w_I) . Suppose, contra hypothesis, that (x^*, y^*) is not Pareto Optimal.

\Rightarrow Then, \exists a feasible allocation (x, y) s.t. (1) $x_i \succeq_i x_i^*$ for each $i=1, \dots, I$
(2) $x_i \succ_i x_i^*$ for some $i=1, \dots, I$.

WTS: (x, y) cannot be feasible.

For each i , $x_i^* \in B_i(p, w_i)$ and x_i^* is \succeq_i -maximal.

- Lemma 1 + (1) : $p^* \cdot x_i \geq w_i$ for each $i=1, \dots, I$

- (2) says there is some i with $p^* \cdot x_i > w_i$

$$\Rightarrow \sum_{i=1}^I p^* \cdot x_i > \sum_{i=1}^I w_i = \underbrace{p^* \cdot \bar{w}}_{\text{By Price Eq. Tr.}} + \sum_{j=1}^J p^* \cdot y_j^* \geq p^* \cdot \bar{w} + \underbrace{\sum_{j=1}^J p^* \cdot y_j}_{\substack{\uparrow \\ \text{By condition (b) of P.E.T.}}}$$

$$\sum_{i=1}^I p^* \cdot x_i > p^* \cdot \bar{w} + \sum_{j=1}^J p^* \cdot y_j$$

\Downarrow

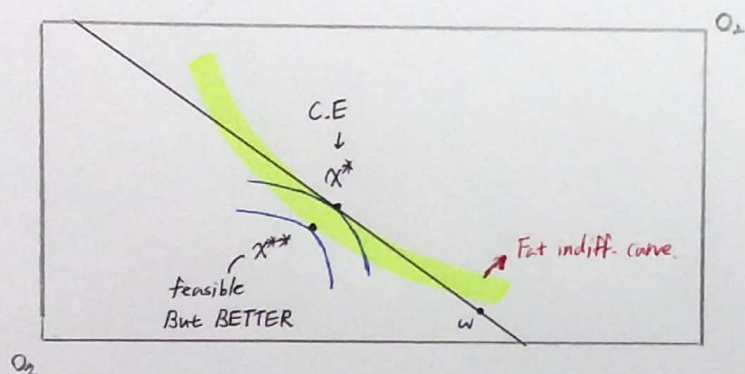
$$\sum_{i=1}^I x_i \neq \bar{w} + \sum_{j=1}^J y_j \quad \leftarrow (x, y) \text{ is NOT feasible} \quad ||$$

(* Market clearing does not hold by the contrahypothesis)

Remark The role of local non-satiation.

To give us: " $x_i \succ x_i^*$ and $x_i^* \in B_i(p^*, w_i)$ is x_i -maximal" $\Rightarrow p \cdot x_i \geq w_i$

Example where LNS violated: Fat indifference curves.



Second Welfare Theorem

P.O \Rightarrow P.E.T / C.E Under "stronger" assumptions

Idea: If we have a Pareto optimal allocation,

then we can sustain it as an equilibrium - potentially after making transfers

\Rightarrow This is why we are in the world of a price equilibrium with transfers

* Will NEED STRONGER Assumptions on preference.

proof strategy: define a "silly equilibrium"

- Show that under [Assumption 1], any Pareto Optimal allocation induces ^{a silly equilibrium}
- Show that under [Assumption 2], any silly eq. is a price eq. with transfers
- Conclude: Second Welfare Theorem [Assumption 1, 2] ^{"NOT Silly"}

NOT all assumptions are created equal