

Chapter 3 Operating System Structures

3.1 Resource Management Process Management

- The creation, deletion, suspension, and resumption of processes.
- Provide mechanism for
 - process synchronization
 - process communication
 - deadlock handling

Process Management

- Program: A sequence of instructions。被動的實體 (Passive entity)。
- Process: 正在執行的程式。活動的實體 (Active entity)。
- Job: The collection of activities needed to do the work required.

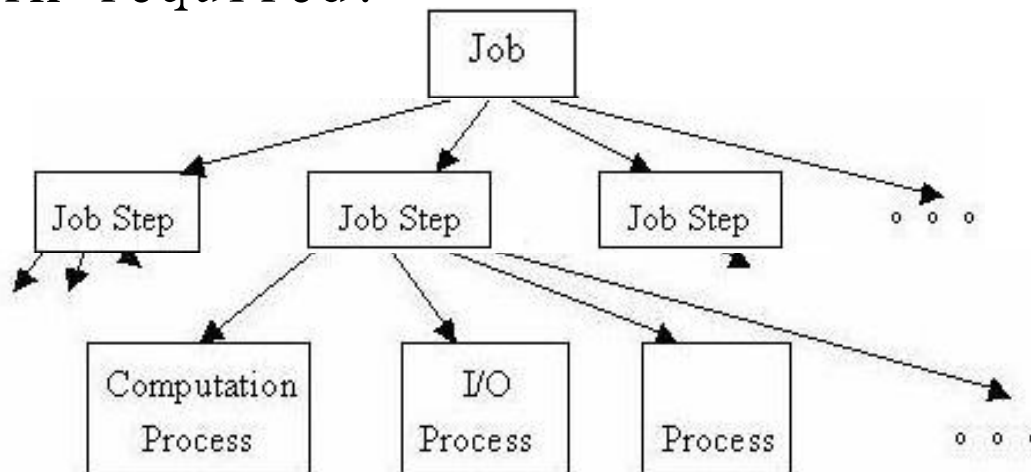


圖 3.2 工作、程式、及處理元之關係

Memory Management

- Keep track of the memory. What parts are used and by whom?
- Decide which processes get memory, when it gets it and how much.
- Allocate and de-allocate memory.

Device Management

- Keep track of the devices, control units.
- Decide what is an efficient way to allocate the device.
- Allocate and de-allocate devices.
 - Buffer caching system
 - Device driver interface
 - Drivers for specific hardware devices (device driver)

Information Management

- Creation and deletion of file systems.
- Keep track of the information (location, used, status).
- The support of primitives for manipulating file systems.
- The mapping of files into secondary storage.

System Call(1)

3.2 System Calls

- When system call, there are three general methods are used to pass parameters to the OS:
 - Registers
 - Pass block address
 - Stack

System Call(2)

- System call can be grouped into five categories
 - Process control
 - File manipulation
 - Device manipulation
 - Information maintenance
 - Communications

System Call(3)

- Process Control
 - end, abort
 - load, execute
 - create process, terminate process
 - get process attributes, set process attributes
 - wait for time
 - wait event, signal event
 - allocate and free memory

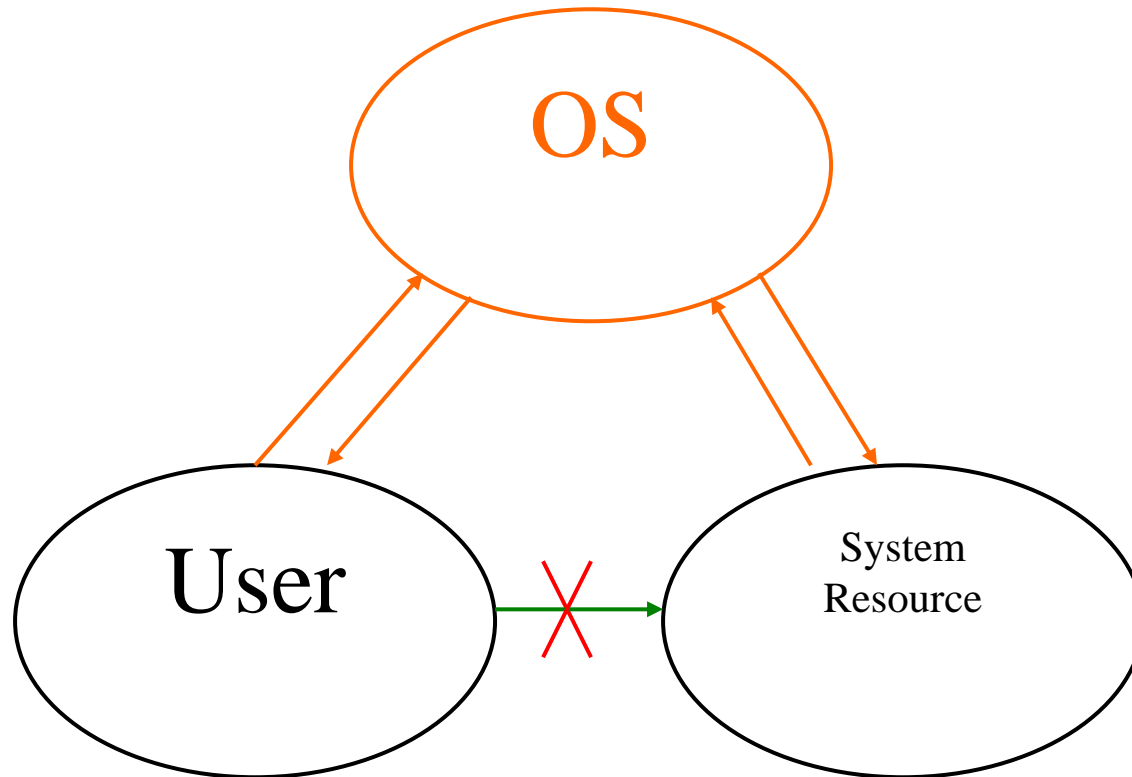
System Call(4)

- File Manipulation
 - create file, delete file
 - open, close
 - read, write, reposition
 - get file attributes, set file attributes
- Device Manipulation
 - request device, release device
 - read, write, reposition
 - get device attributes, set device attributes
 - logically attach or detach devices

System Call(5)

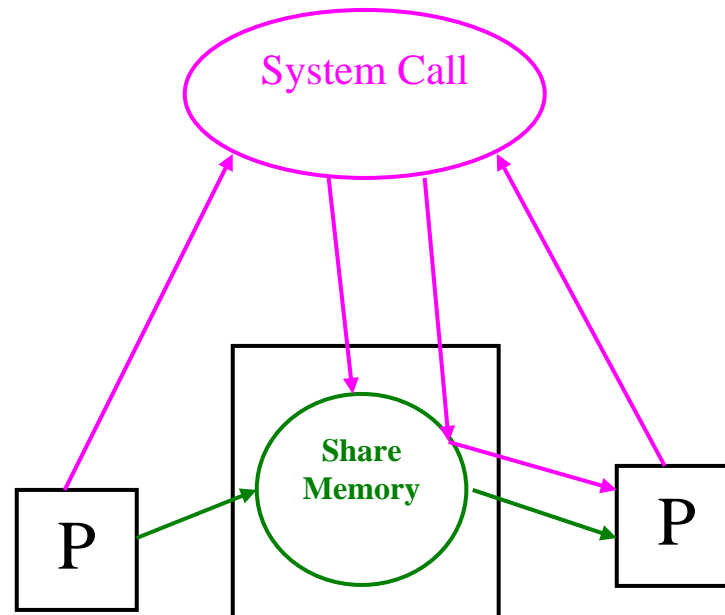
- Information Maintenance
 - get time or date, set time or date
 - get system data, set system data
 - get process, file, or device attributes
 - set process, file, or device attributes
- Communication
 - creat, delete communication connection
 - send, receive messages
 - transfer status information
 - attach or detach remote devices

System Call(6)



Communication

- Message passing model
 - By using **Inter Process Communication** (IPC)
- Shared memory model
 - Processes use **map memory** system calls to gain access to regions of memory owned by other processes.



分層結構(Layer Approach)(1)

3.3 System Structure

- 原則上作業系統之各個管理及功能都是以分層結構方式組合。
- 作業系統也是由軟體及硬體所撰寫完成，而撰寫軟體常強調系統化及模組化。

分層結構(Layer Approach)(2)

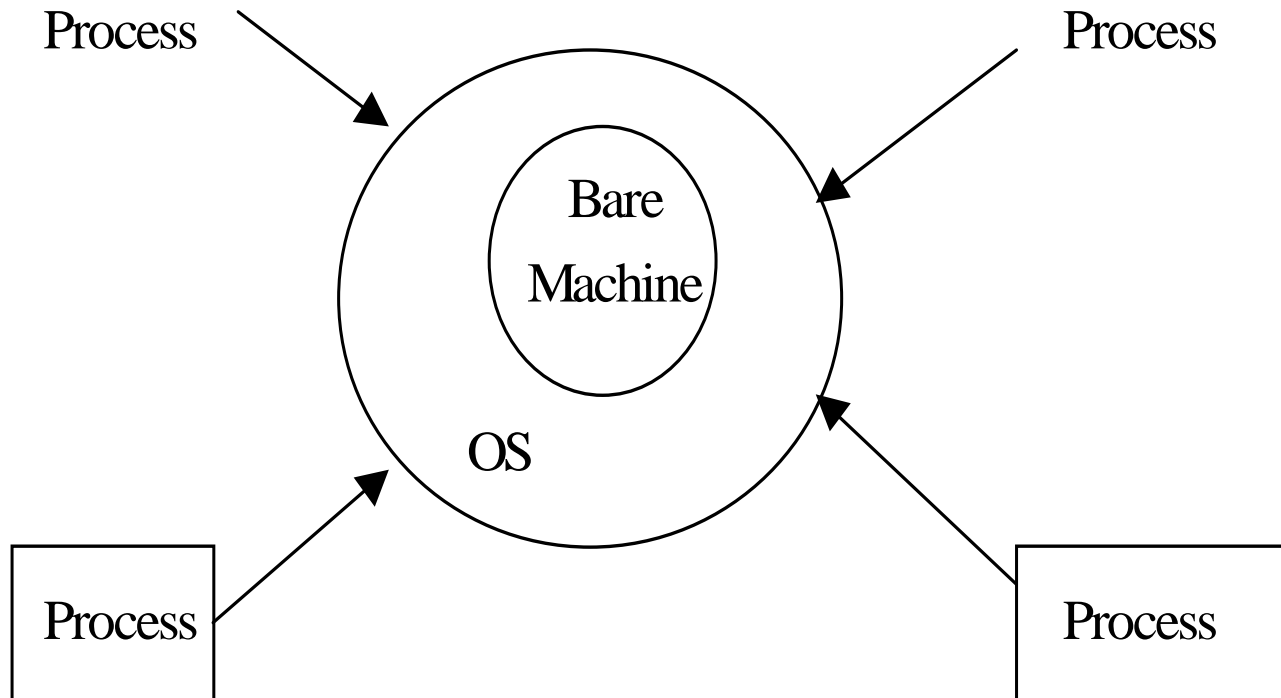


圖 1.22a 非分層結構之作業系統

分層結構(Layer Approach)(3)

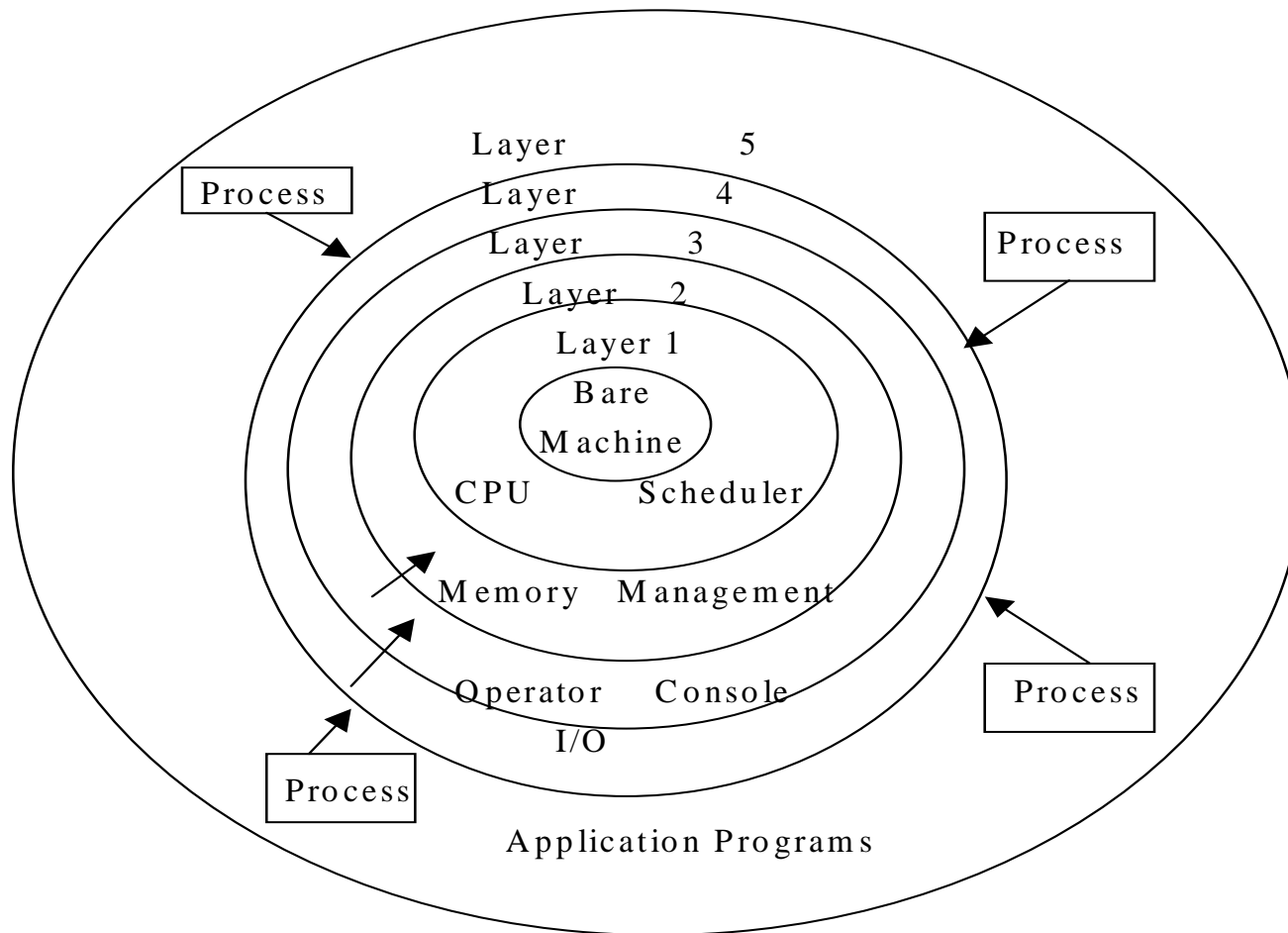


圖 1.22b 分層結構之作業系統

分層結構(Layer Approach)(4)

Users		
Shell and System Library		
System Call Interface to the Kernel		
I/O	Device Driver	File System CPU Scheduling Virtual Memory Demand Paging (Operating System Kernel)
Kernel Interface to the Kernel		
Terminal Controller Terminals	Device Controllers Disks and Tapes	Memory Controllers Physical Memory

圖 1.12 UNIX 作業系統之分層結構

分層結構(Layer Approach)(5)

MS / DOS之分層結構，應用程式可以直接呼叫基本輸出 / 輸入函數，而沒有透過作業系統之系統呼叫後，再經由作業系統核心出面安排處理工作；也因為MS / DOS作業系統的不明顯分層結構，而可以由應用程式直接控制系統資源，以致造成容易當機的問題發生。

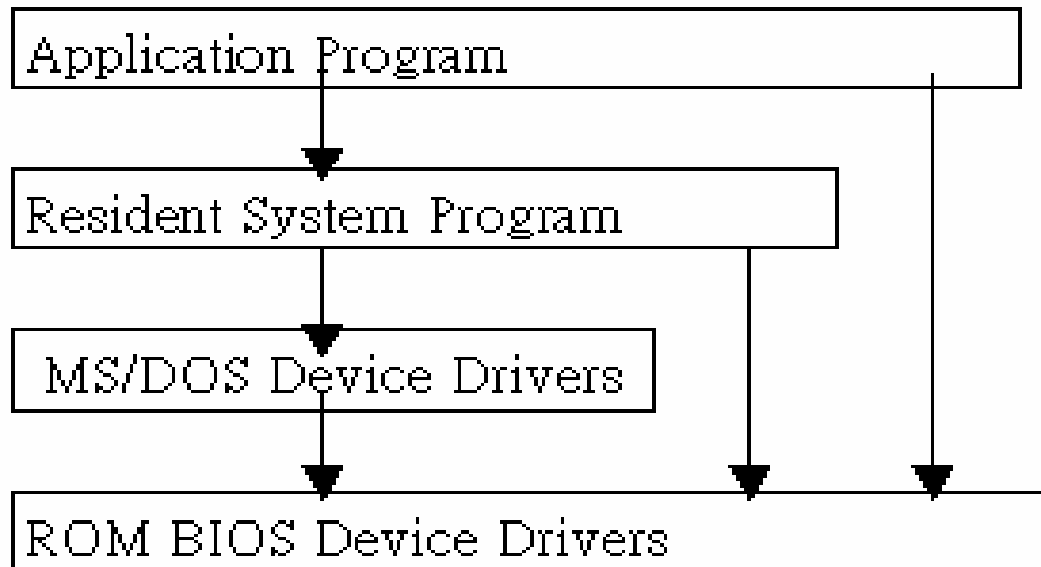


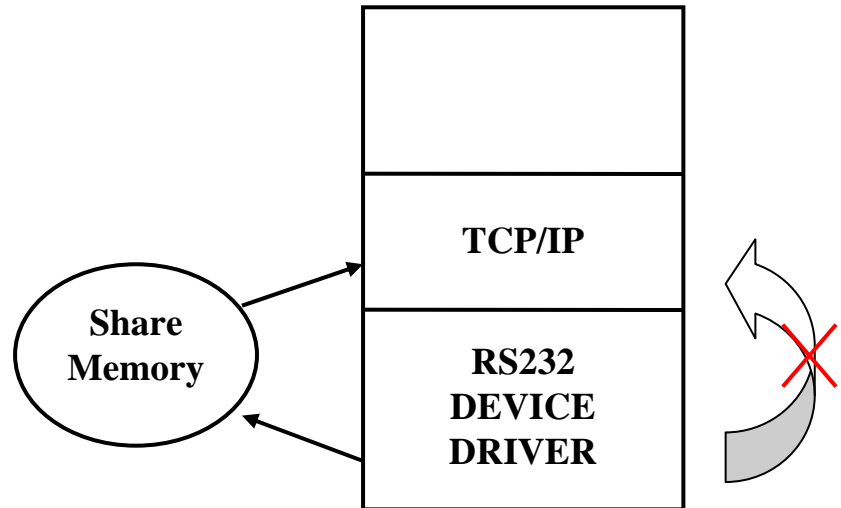
圖 1.13 MS / DOS 之分層結構

分層結構(Layer Approach)(6)

- Direct user access to low level facilities is not allowed.

- 內層不可以呼叫外層。

- 外層只能一層一層呼叫。



- 可以在不更改該介面的條件下，將某一層抽離換新，而不會影響其他幾層。

虛擬機器(Virtual Machine)(1)

3.4 Virtual Machine

- 從使用者的角度來看，每個人都認為電腦是自己所擁有的，電腦資源都是自己獨享的，這就是虛擬機器的概念。
- 每位使用者都擁有自己的虛擬機器，而實際上作業系統會將虛擬機器上的虛擬資源（Virtual Resources），轉換為實際資源（Physical Resources），並去操控它們。

虛擬機器(Virtual Machine)(2)

For each user program execution on **its own processor**(CPU scheduling is managed by OS) with its **own memory virtually**, and the system support **spooling** facility.

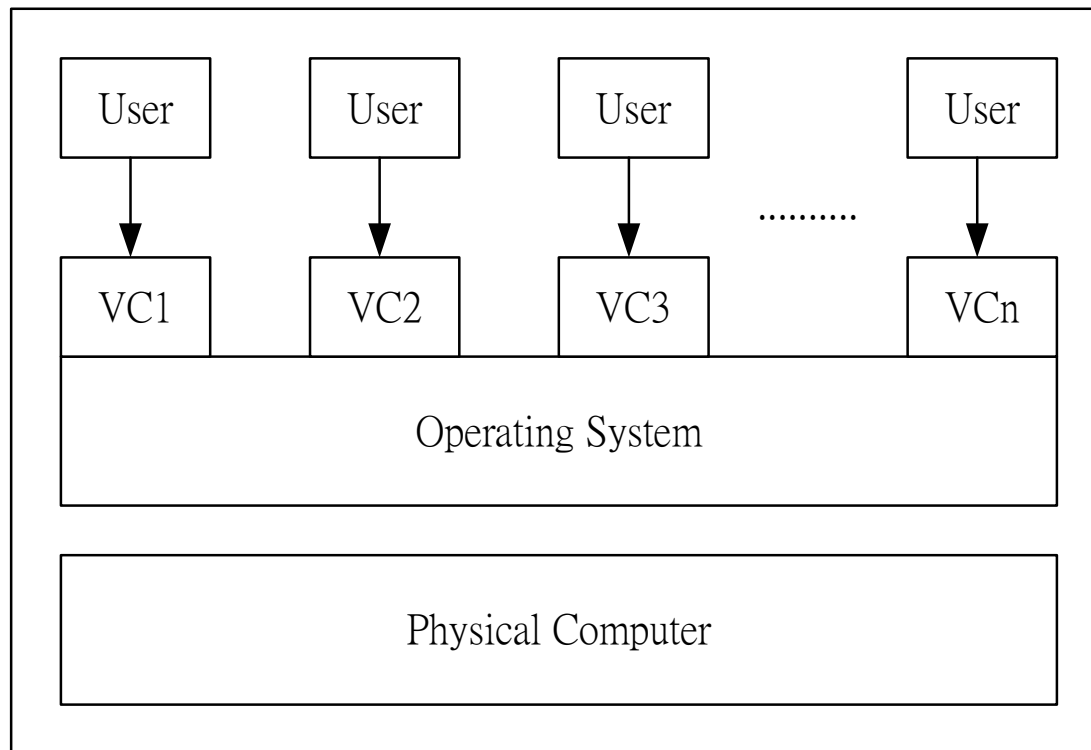


圖2.12 虛擬機器

虛擬機器(Virtual Machine)(3)

- 作業系統提供時間多工 (Time Multiplex) ，使得許多程式可以分時使用中央處理器。
- 提供空間多工 (Space Multiplex) ，使得許多程式可以存在主記憶體內不同位置。

虛擬機器(Virtual Machine)(4)

- 虛擬機器最簡單的講法就是說，一部電腦的作業系統是以分層結構方式完成，使用者純粹以邏輯結構來看電腦及命令電腦工作，而用不著去了解細微的機器控制動作，因為這些細微控制及同步機制均由作業系統內各分層協調完成。
- Layer approach is taken to its logical conclusion in the concept of a virtual machine.

Java Virtual Machine(1)

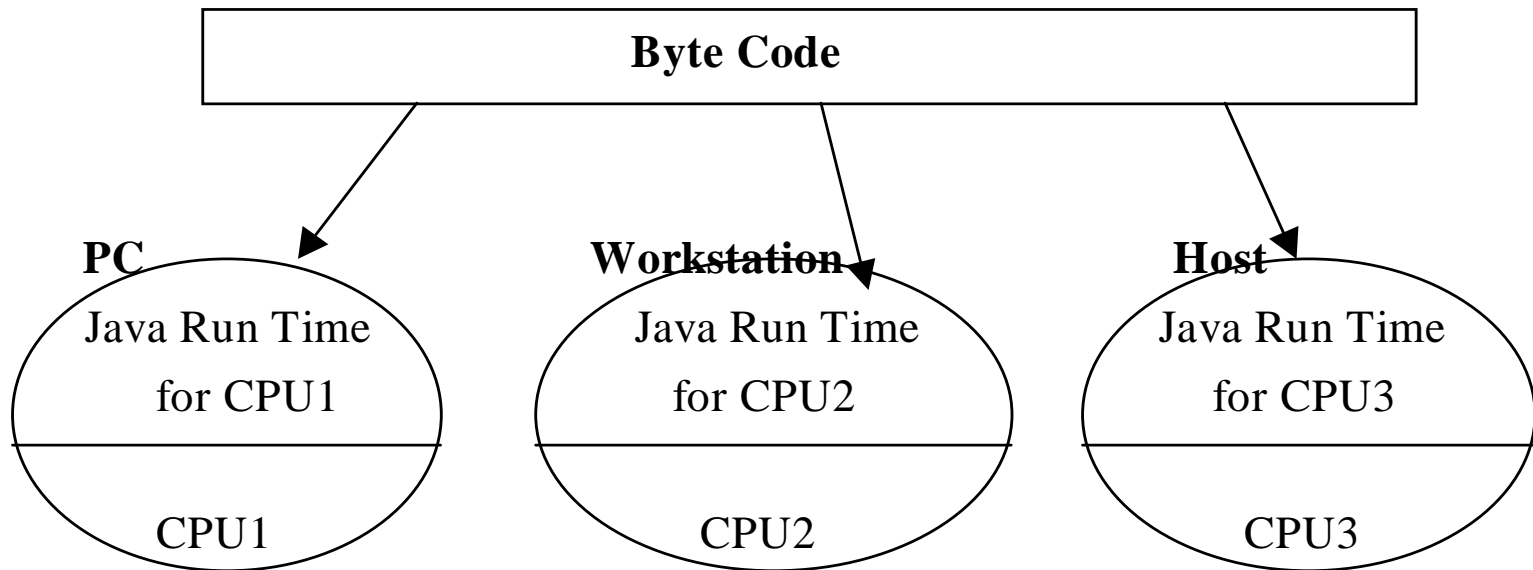


圖 2.13 Java 虛擬機器

Java Virtual Machine(2)

- Java挾其可攜帶性之優點，橫掃跨平台作業系統，它便是一種虛擬機器的代表。
- 從外層的角度來看，所有命令都與機器無關，它們是位元組碼。
- 從內層來看，不同機器的Java執行模組均不相同，因為Java執行模組必須依據所使用的中央處理器，撰寫不同的機器碼，來命令電腦工作。

System Design and Implementation

3.5 System Design and Implementation

- 在設計分層結構時，通常必須將機制（Mechanism）與政策（Policy）分開。
- 所謂機制是一組供使用之基本功能。
 - A set of basic facilities that can be used in many different ways. Determine how to do something. 如何作？
- 所謂政策是使用機制功能作特定工作。
 - The use of a mechanism for a particular purpose. Decide what will be done. 作什麼？

機制（Mechanism）與政策（Policy）

- Mechanism
 - CPU protection: timer, interrupt
 - Protect file: Read, Write, Execute
- Policies
 - CPU protection: how long, what to do
 - Protect file: which files can be R, W, E
- 如果我們確定機制之後，即使政策改變，只要重新定義系統中相關的參數便可以（例如時間片斷減小）；因此將機制與政策分開，可以增加系統的彈性。

System Generation

3.6 System Generation(Installation)

- The system be configured for each specific computer site. e.g. memory size, device, options,...etc.