# 16-350 Planning Techniques for Robotics

## Interleaving Planning and Execution: Real-time Heuristic Search

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#### Planning during Execution

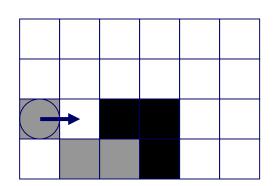
- Planning is a <u>repeated</u> process!
  - partially-known environments
  - dynamic environments
  - imperfect execution of plans
  - imprecise localization

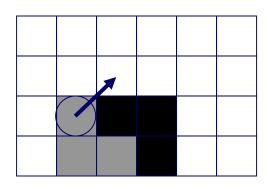
- Need to be able to re-plan fast!
- Several methodologies to achieve this:
  - anytime heuristic search: return the best plan possible within T msecs
  - incremental heuristic search: speed up search by reusing previous efforts
  - real-time heuristic search: plan few steps towards the goal and re-plan later

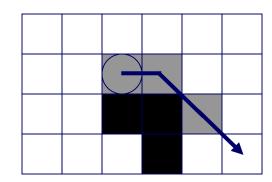
Enforce <u>a strict limit</u> on the amount of computations (no requirement on planning all the way to the goal)

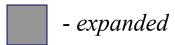
- 1. Compute a partial path by expanding at most N states around the robot
- 2. Move once, incorporate sensor information, and goto step 1

Example in a fully-known terrain:



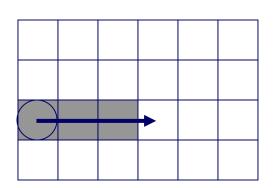


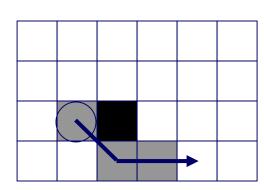


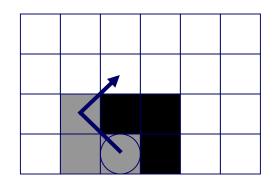


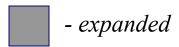
- 1. Compute a partial path by expanding at most N states around the robot
- 2. Move once, incorporate sensor information, and goto step 1

Example in an unknown terrain (planning with Freespace Assumption):









- 1. Compute a partial path by expanding at most N states around the robot
- 2. Move once, incorporate sensor information, and goto step 1

#### Research issues:

- how to compute partial path
- how to guarantee complete behavior (guarantee to reach the goal)
- provide bounds on the number of steps before reaching the goal

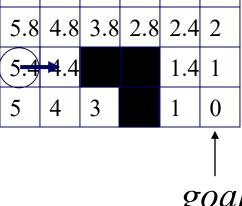
## What should the planner decide for the robot's next move?

$$h(x,y) = \max(abs(x-x_{goal}), abs(y-y_{goal})) + 0.4*\min(abs(x-x_{goal}), abs(y-y_{goal}))$$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	3		1	0
				σ	† oai

• Repeatedly move the robot to the most promising adjacent state, using heuristics

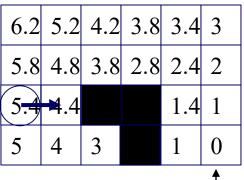
1. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$ 

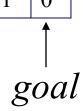


• Repeatedly move the robot to the most promising adjacent state, using heuristics

1. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$ 

 $h(x,y) = \max(abs(x-x_{goal}), \ abs(y-y_{goal})) + 0.4*\min(abs(x-x_{goal}), \ abs(y-y_{goal}))$ 





Any problems?

• Repeatedly move the robot to the most promising adjacent state, using heuristics

1. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$ 

 $h(x,y) = \max(abs(x-x_{goal}), \ abs(y-y_{goal})) + 0.4*\min(abs(x-x_{goal}), \ abs(y-y_{goal}))$ 

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	<b>-</b> 4.4			1.4	1
5	4	3		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	44			1.4	1
5	4	3		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	$\mathbb{C}$		1	0

Any problems?

• Repeatedly move the robot to the most promising adjacent state, using heuristics

1. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$ 

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6.2	5.2	4.2	3.8	3.4	3
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5.4	4.4			1.4	1
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6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	*		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	$\left( \mathbf{T} \right)$		1	0

Local minima problem (myopic or incomplete behavior)

Any solutions?

• Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics

\*\*makes h-values more informed\*

- 1.  $update h(s_{start}) = min_{s \in succ(sstart)} c(s_{start}, s) + h(s)$
- 2. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	<b>-</b> 4.4			1.4	1
5	4	3		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	3		1	0

6.2	5.2	4.2	$\beta$ .8	3.4	3
5.8	4.8	3.8	$\frac{1}{2.8}$	2.4	2
5.4	4.4			1.4	1
5	4	4		1	0

• Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics

- 1.  $update h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
- 2. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4,4			1.4	1
5	5.4	5		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	5.2			1.4	1
5	5.4	5		1	0

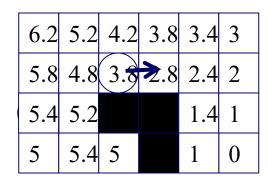
6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	<b>≥</b> .8	2.4	2
5.4	5.2			1.4	1
5	5.4	5		1	0

• Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics

- 1.  $update h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
- 2. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	44			1.4	1
5	5.4	5		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	5.2			1.4	1
5	5.4	5		1	0



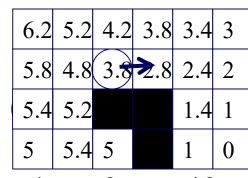
*h*-values guaranteed to remain admissible and consistent

• Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics

- 1.  $update h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
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5.8	4.8	3.8	2.8	2.4	2
5.4	4,4			1.4	1
5	5.4	5		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	5.2			1.4	1
5	5.4	5		1	0



robot is guaranteed to reach goal in finite number of steps if:

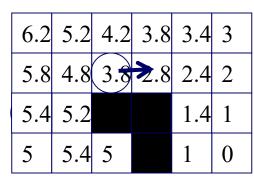
- all costs are bounded from below with  $\Delta > 0$
- graph is of finite size and there exists a finite-cost path to the goal
- all actions are reversible

• Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics

- 1.  $update h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
- 2. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	44			1.4	1
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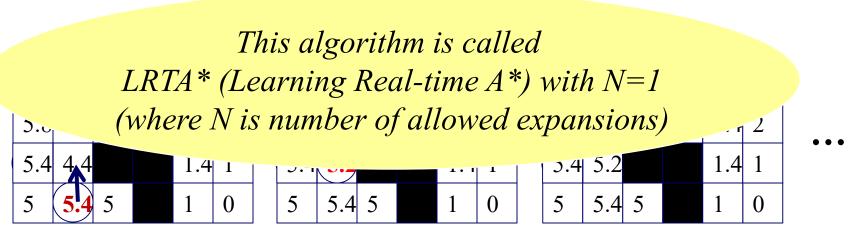
6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	5.2			1.4	1
5	5.4	5		1	0



robot is guaranteed to reach goal in finite number of steps if

- all costs are bounded from below with  $\Delta > 0$ 
  - Why conditions?
- graph is of finite size and there exists a finite-cost path to the goal
- all actions are reversible

- Repeatedly move the robot to the most promising adjacent state, using **and updating** heuristics
  - 1.  $update\ h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
  - 2. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)}c(s_{start}, s) + h(s)$



robot is guaranteed to reach goal in finite number of steps if:

- all costs are bounded from below with  $\Delta > 0$
- graph is of finite size and there exists a finite-cost path to the goal
- all actions are reversible

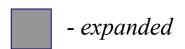
### Learning Real-Time A\* (LRTA\*) with N=1

• expand N = 1 state, make a move towards a state s in OPEN with smallest g(s) + h(s):

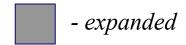
- 1.  $expand s_{start}$
- 2.  $update h(s_{start}) = min_{s \in succ(sstart)}c(s_{start}, s) + h(s)$
- 3. always move as follows:  $s_{start} = argmin_{s \in succ(sstart)} c(s_{start}, s) + h(s)$ =  $argmin_{s \in succ(sstart)} g(s) + h(s)$

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	<b>≠</b> .4			1.4	1
5	4	3		1	0

6.2	5.2	4.2	3.8	3.4	3
5.8	4.8	3.8	2.8	2.4	2
5.4	4.4			1.4	1
5	4	3		1	0



- LRTA\* with  $N \ge 1$  expands
  - necessary for the guarantee to reach the goal
  - 2. update h-values of expanded states by Dynamic Programming (DP)
  - 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

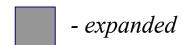


#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

#### state s:

- the state that minimizes cost to it plus heuristic estimate of the remaining distance
- the state that looks most promising in terms of the whole path from current robot state to goal



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = argmin_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
5	4		2	1
4	3	<b>2</b>		0

4-connected grid (robot moves in 4 directions)

example borrowed from ICAPS'06 planning summer school lecture (Koenig & Likhachev)



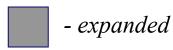
- expanded

#### • LRTA\* with $N \ge 1$ expands

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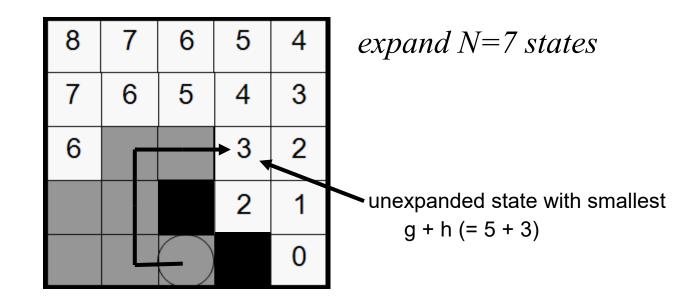
8	7	6	5	4
7	6	5	4	3
6			3	2
			2	1
				0

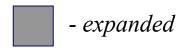
expand N=7 states



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = argmin_{s' \in OPEN} g(s') + h(s')$





• LRTA\* with  $N \ge 1$  expands

How path is found?

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

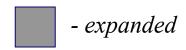
					•
8	7	6	5	4	expand N=7 states
7	6	5	4	3	
6			<b>3</b>	2	
			2	1	unexpanded state with smallest g + h (= 5 + 3)
				0	9 ( 0 . 0)

- expanded

#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

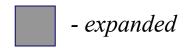
8	7	6	5	4
7	6	5	4	3
6	$\infty$	$\infty$	3	2
$\infty$	$\infty$		2	1
$\infty$	$\infty$	$\infty$		0



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = argmin_{s' \in OPEN} g(s') + h(s')$

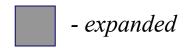
8	7	6	5	4
7	6	5	4	3
6	$\infty$	4	3	2
$\infty$	$\infty$		2	1
$\infty$	$\infty$	$\infty$		0



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
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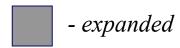
8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
$\infty$	$\infty$		2	1
$\infty$	$\infty$	$\infty$		0



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
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- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
$\infty$	6		2	1
$\infty$	$\infty$	$\infty$		0



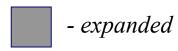
#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
7	6		2	1
$\infty$	$\infty$	$\infty$		0

update h-values of expanded states via DP: set h-values of expanded states to infinity compute  $h(s) = \min_{s' \in succ(s)} (c(s,s') + h(s'))$ until convergence

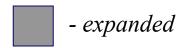
Does it matter in what order?



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

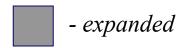
8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
7	6		2	1
$\infty$	7	$\infty$		0



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

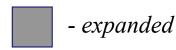
8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
7	6		2	1
8	7	$\infty$		0



#### • LRTA\* with $N \ge 1$ expands

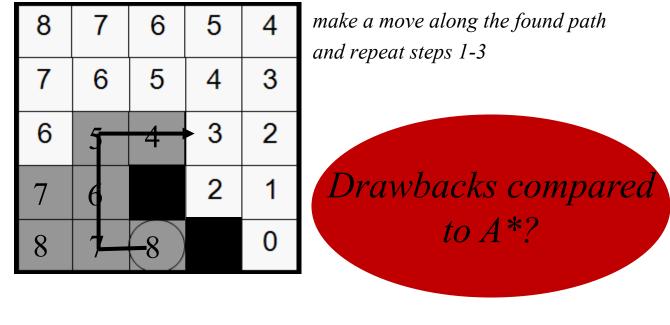
- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
7	6		2	1
8	7	8		0



#### • LRTA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states by Dynamic Programming (DP)
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$



• RTAA\* with  $N \ge 1$  expands: LRTA\*

one linear pass, and even that can be lazy(postponed)

- 1. expand N states
- 2. update h-values of expanded states u by h(u) = f(s) g(u), where  $s = argmin_{s' \in OPEN} g(s') + h(s')$
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4	expand N=7 states
7	6	5	4	3	
6			<b>3</b>	2	
			2	1	unexpanded state $s$ with smallest $g + h (= 5 + 3)$
				0	g · // (= 3 · 3)

- expanded

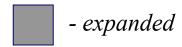
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- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	g=3	g=4	3	2
g=3	g=2		2	1
g=2	g=1	(g=0)		0

 $update \ all \ expanded \ states \ u:$  h(u) = f(s) - g(u)

unexpanded state s with smallest f(s) = 8



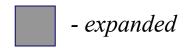
#### • RTAA\* with $N \ge 1$ expands

- 1. expand N states
- 2. update h-values of expanded states u by h(u) = f(s) g(u), where  $s = argmin_{s' \in OPEN} g(s') + h(s')$
- 3. move on the path to state  $s = \operatorname{argmin}_{s' \in OPEN} g(s') + h(s')$

8	7	6	5	4
7	6	5	4	3
6	8-3	8-4	3	2
8-3	8-2		2	1
8-2	8-1	8-0		0

update all expanded states u: h(u) = f(s) - g(u)

unexpanded state s with smallest f(s) = 8



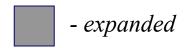
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8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
5	6		2	1
6	7	8		0

 $update \ all \ expanded \ states \ u:$  h(u) = f(s) - g(u)

unexpanded state s with smallest f(s) = 8

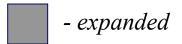


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proof of admissibility: h*() - true cost-to-goal	8	
$g(u) + h^*(u) \ge h^*(s_{start})$	7	
$h^*(u) \ge h^*(s_{start}) - g(u)$ because $f(s) \le h^*(s_{start})$	6	
$h^*(u) \ge f(s) - g(u) \checkmark$	5	
$h^*(u) \ge h_{updated}(u)$	6	,

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
5	6		2	1
6	7	8		0



#### LRTA\* vs. RTAA\*

T	$\mathbf{R}$	$\Gamma \Delta$	<b>*</b>
	/ <b>                                    </b>		<b>\</b>

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
7	6		2	1
8	7	8		0

#### RTAA\*

8	7	6	5	4
7	6	5	4	3
6	5	4	3	2
5	6		2	1
6	7	$\left(8\right)$		0

- Update of *h*-values in RTAA\* is much faster but not as informed
- Both guarantee adimssibility and consistency of heuristics
- For both, heuristics are monotonically increasing
- Both guarantee to reach the goal in a finite number of steps (given the conditions listed previously)

## Summary

- Real-time Heuristic Search puts a hard constraint on planning time (usually, a smaller planning time than what is required to plan a path all the way to the goal)
- Computing a partial path to the goal may result in highly sub-optimal behavior
- It is important to think how to avoid infinite oscillations
  - Updating heuristics is a popular way for doing it
  - Mostly applicable to low-dimensional planning
  - How to extend it to high-dimensional planning is a research question