

Embedded System Design

Traffic Light

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Abstract

The purpose of this project was to simulate traffic lights at an intersection which is controlled by software and use priorities to tell the operating system how to schedule the different tasks.

Implementation / Design

The first step we took in implementing our project was to break down the functionality of the intersection and figure out the priority of each function. We ended up matching the assignment document exactly in terms of functionality per task and priority of each task. We determined task A to be the lowest priority and each following task to be a higher priority than the previous all the way up to the highest priority task, task F. In regards to the length of each light we did not want to use *actual* intersection timing as it would take for ever to test. We therefore used what we felt were accurate *relative* timings. For instance the primary street green is twice as long as the secondary street green. It really does not matter what the timings are as we are more concerned with sequencing and event control.

Our next design decision was how to display the status of the intersection and how we would control events such as a car being in the turning lane or the walk button being pressed. After evaluating all the IO on the DE1 we decided that the cleanest and clearest way to convey all this information would be to utilize the JTAG UART functionality of the Nios 2 and actually draw the intersection in the terminal using ANSI control sequences to set color and position of different things such as the lights and cars in the turn lane. This allows us to see the whole state of our intersection without trying to have to decode LED sequences or 7-segment characters on the DE1 board. Additionally, we can use the receiving side of the JTAG UART to read characters from the keyboard on the PC. This allows us to logically map keys such as 'w' to the walk button. We found this to be a much cleaner solution than trying to remember which switch or button on the DE1 corresponded to which functionality.

Once our IO functionality was decided, the first step was to write thread-safe drawing functions to draw things to the terminal such as the lanes of the street and the lights at the proper position and color. Thread-safety was achieved by a 'draw_lock' mutex which each drawing function waits to obtain before writing anything to the UART. This ensures that two threads won't interrupt each other when trying to update the screen. To read bytes off the UART a separate task is used which actually has the highest priority in our whole program. This task continuously calls `getc()` to try and read the next byte in the UART buffer and if a byte is read it uses its value to set a flag in a global 'key_down' array. For instance, if the key 'a' is pressed, which has an ASCII value of 97, then `key_down[97]` is set to 1. A separate function called 'key_pressed(char key)' is used to check if a certain key has been pressed. For instance, calling 'key_pressed('w')' will check the `key_down[]` array to see if `key_down['w']` is set to a 1. If it is set it will clear it to a 0, since it's now been handled, and return true. The `key_down` array is protected by a 'input_lock' mutex to avoid reading and writing at the same time.

After establishing all the IO routines we were able to quickly implement each task. In order to keep a modular design each task is independent in the sense that when that task is running it does not depend on any other task to control the lights. A global mutex called 'light_lock' is used to control which task is controlling the lights. A task must acquire this mutex before attempting to modify the light values. This works well because it ensures that higher priority tasks will always gain control of the lights when they need it. Each task, except for manual mode task F, keeps track of what it is currently doing by using a state machine. This is an ideal control scheme for this project because the nature of a traffic light can be viewed as a state machine. For instance, in task A, a switch case is used to go from Primary Green, to Primary Yellow, to Primary Red, etc... In the other tasks there is an initial state which just waits for input such as the walk button being pressed. Once the walk button is pressed it goes to the next state which waits for all lights to turn red. It then goes to the next state which locks the lights and makes the walk light green. It then goes to the next state which makes it yellow, etc.. It's easy to see how using state machines allows the required sequential functionality to be directly implemented.

Program Output

The terminal output showing an 'ASCII art' representation of the intersection is shown below. Using Unicode characters allowed us to draw nice continuous lines, corners, and arrows. ANSI control sequences allowed us to achieve proper colors for the lanes and lights. The top right section shows the key mapping to control the functionality of the intersection and the bottom right section shows status and debug output describing the current state of the intersection. The blue squares in the primary street turn lanes indicate a car is waiting.

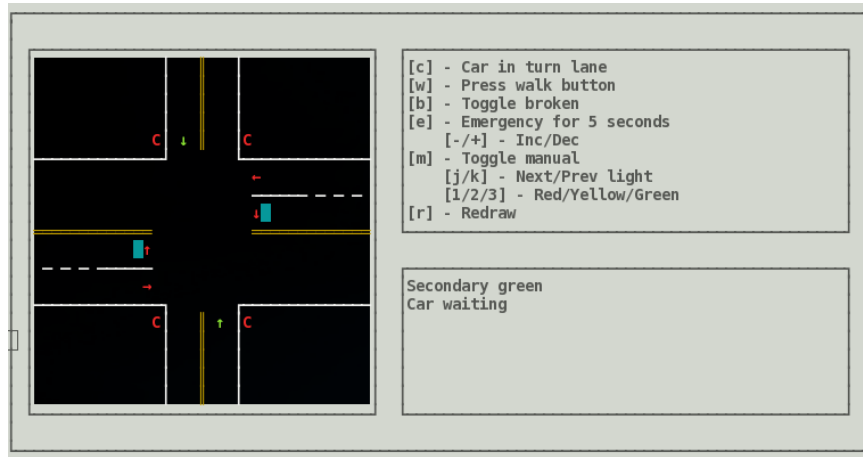


Figure 1: Terminal Output

Testing Strategy

Given the modular approach of our design testing the functionality of each task was extremely easy. We were able to develop each task in parallel and test their individual functionality independently. The only exception to this is task A, which a few other tasks depend on. For instance, task b waits until the primary street is about to turn green and task c waits until all lights are red. Overall the testing and integrating of each task went flawlessly which we attribute mostly to our choice in IO. We were, at a glance, able to see the entire state of our system and quickly test out the different inputs without having to remember if switch 1 was for the walk button or emergency. Had we chosen to use LEDs and switches we might not have had this advantage.

Implementation Resource Usage

Using the 'nios2-elf-size' command we can get a breakdown of our programs resource utilization

```
$ nios2-elf-size traffic.elf
textwidth  data    bss    Dec    hexfilename
105580    7044   71332 183956  2ce94traffic.elf
```

The Quartus build summary provides a summary of the FPGA utilization.

```
Fitter Status : Successful - Mon Nov 26 20:51:57 2012
Quartus II 32-bit Version : 12.0 Build 178 05/31/2012 SJ Web Edition
Revision Name : top
Top-level Entity Name : top
Family : Cyclone II
Device : EP2C20F484C7
Timing Models : Final
Total logic elements : 2,832 / 18,752 ( 15 % )
    Total combinational functions : 2,366 / 18,752 ( 13 % )
    Dedicated logic registers : 1,861 / 18,752 ( 10 % )
Total registers : 1929
Total pins : 80 / 315 ( 25 % )
Total virtual pins : 0
Total memory bits : 44,032 / 239,616 ( 18 % )
Embedded Multiplier 9-bit elements : 0 / 52 ( 0 % )
Total PLLs : 1 / 4 ( 25 % )
```

Implementation Speed and Power

The max clock speeds as indicated by the Timing Analyzer is shown below

```
+-----+
; Slow Model Fmax Summary                                     ;
+-----+-----+-----+-----+
; Fmax           ; Restricted Fmax ; Clock Name           ; Note ;
+-----+-----+-----+-----+
; 80.33 MHz      ; 80.33 MHz        ; u0|altpll_0|sd1|pll|clk[0] ;       ;
; 252.91 MHz     ; 252.91 MHz       ; u0|altpll_0|sd1|pll|clk[2] ;       ;
+-----+-----+-----+-----+
```

The thermal power dissipation as indicated by Power Play is shown below

```
PowerPlay Power Analyzer Status : Successful - Sun Dec  9 21:50:01 2012
Quartus II 32-bit Version : 12.0 Build 178 05/31/2012 SJ Web Edition
Revision Name : top
Top-level Entity Name : top
Family : Cyclone II
Device : EP2C20F484C7
Power Models : Final
Total Thermal Power Dissipation : 149.99 mW
Core Dynamic Thermal Power Dissipation : 48.74 mW
Core Static Thermal Power Dissipation : 47.50 mW
I/O Thermal Power Dissipation : 53.75 mW
Power Estimation Confidence : Low: user provided insufficient toggle rate data
```

Our critical path seems to be the interface with the SDRAM as shown by the Timing Analyzer

```
+-----+
; Fast Model Setup: 'u0|altpll_0|sd1|pll|clk[0] '
+-----+-----+-----+-----+
; Slack      ; From Node
+-----+-----+-----+-----+
; 15.098 ; nios_system:u0|nios_system_sdram_0:sdram_0|nios_system_sdram_0_input_efifo_module
; 15.127 ; nios_system:u0|nios_system_sdram_0:sdram_0|nios_system_sdram_0_input_efifo_module
; 15.167 ; nios_system:u0|nios_system_sdram_0:sdram_0|nios_system_sdram_0_input_efifo_module
; 15.182 ; nios_system:u0|nios_system_sdram_0:sdram_0|nios_system_sdram_0_input_efifo_module
...
```

Team Member Contributions

- **Max Thrun** - Print/drawing functions, UART input task, task F
- **Ian Cathey** - Task A, Task C, Task D
- **Mark Labbato** - Task B, Task E, Task D