Fast Sketching with Fast Walsh-Hadamard Transform

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Background Information

Matrix sketching preserves some properties of a very large matrix into a much smaller matrix with reduced dimension.

$$\mathbf{A} \in \mathbb{R}^n o \mathbb{R}^d, d < n$$

By having a low rank approximation of a big matrix, it effectively reduces the size of the problem and makes it significantly easier and faster to compute.

We ensure a high probability sketch through **random sampling**. Randomization reduces the complexity of algorithms and increases its speed.

Fast Walsh-Hadamard Transform

When performing matrix sketching by SRHT, we use the **Fast-Walsh Hadamard transform (FWHT)** over the regular Hadamard transform.

Both transformations make use of the Hadamard matrix, which is a square matrix that is defined recursively and has dimensions of a power of 2.

$$H_n=H_2\otimes H_{rac{n}{2}}$$
 $H_2=egin{bmatrix}1&1\1&-1\end{bmatrix}$

where ⊗ is the Kronecker product

Rather than multiplying with the Hadamard matrix, FWHT is computed algorithmically without needing to create the full Hadamard matrix.

$$egin{aligned} x = [y_1 & y_2]
ightarrow x H_4 = [y_1 & y_2] egin{bmatrix} H_2 & H_2 \ H_2 & -H_2 \end{bmatrix} \ x H_4 = egin{bmatrix} y_1 H_2 + y_2 H_2 & y_1 H_2 - y_2 H_2 \end{bmatrix} \end{aligned}$$

Sketching Methods

The current sketching method consists of random sampling.

$$Y=\mathbf{A}\Omega$$

 $\mathbf{A} \subseteq \mathbb{R}^{m,n}$: an original $m \times n$ matrix $\mathbf{\Omega} \subseteq \mathbb{R}^{n,d}$: a Gaussian random matrix \mathbf{Cost} : O(mnd)

To optimize, we produce and modify the Gaussian random matrix using the **Subsampled randomized Hadamard transform (SRHT)**.

$$\Omega = DHS$$

 $\mathbf{D} \in \mathbb{R}^{n,n}$: a diagonal matrix whose diagonal entries are ± 1 $\mathbf{H} \in \mathbb{R}^{n,n}$: Walsh-Hadamard transform matrix $\mathbf{S} \in \mathbb{R}^{n,d}$: a random subset of d columns from an $n \times n$ identity matrix \mathbf{Cost} : $O(mn \log(n))$

Bit-Reversal

Bit-reversal is a permutation where the indices of elements within a sequence are switched according to their binary representation. It simplifies the process of accessing elements in an array, thus making the algorithm run faster.

Bit-reversal can be implemented in FWHT for potential improved fast memory access. When performing the calculations, bit-reversal allows immediate addition and subtraction on the same numbers.

Omitting the permutation results in calculating all additions first on different values, followed by all subtractions in each iteration.

$$x_{0} + x_{1,} = \hat{x}_{0} \rightarrow \hat{x}_{0} + \hat{x}_{2}$$

$$x_{0} - x_{1} = \hat{x}_{1} \rightarrow \hat{x}_{1} + \hat{x}_{3}$$

$$x_{2} + x_{3} = \hat{x}_{2} \rightarrow \hat{x}_{0} - \hat{x}_{2}$$

$$x_{2} - x_{3} = \hat{x}_{3} \rightarrow \hat{x}_{1} - \hat{x}_{3}$$

FWHT without bit-reversal

$$x_{0} + x_{2} = \hat{x}_{0} \rightarrow \hat{x}_{0} + \hat{x}_{2}$$

$$x_{0} - x_{2} = \hat{x}_{1} \rightarrow \hat{x}_{0} + \hat{x}_{2}$$

$$x_{1} + x_{3} = \hat{x}_{2} \rightarrow \hat{x}_{1} + \hat{x}_{3}$$

$$x_{1} - x_{3} = \hat{x}_{3} \rightarrow \hat{x}_{1} + \hat{x}_{3}$$

FWHT with bit-reversal

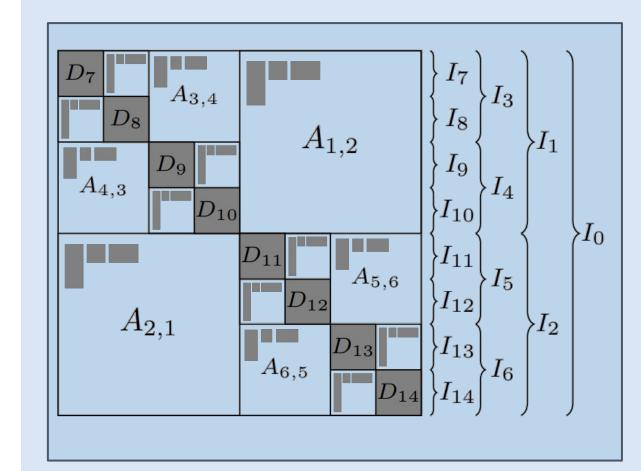
Application

SRHT is one of the many sketching algorithms that can be used in a **Hierarchically Semi-Separable Matrix (HSS Matrix)**.

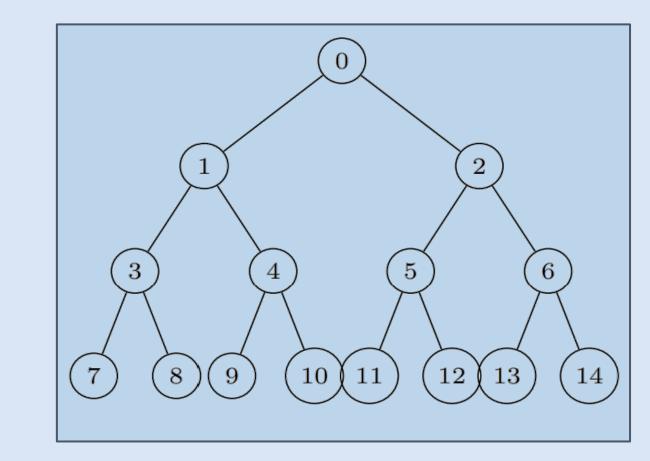
Other sketching algorithms that can be used are:

- Sparse Johnson-Lindenstrauss transform
- Subsampled randomized fast Fourier transform
- Subsampled randomized discrete cosine transform

The HSS matrix is a dense matrix where the off-diagonal blocks are recursively partitioned into smaller rank-deficient blocks.



An HSS matrix with 4 levels. Grey blocks denote the basis matrices, *U*, *B*, *V*, and *D*



The HSS matrix represented as a top-down level-by-level tree.

By compressing a matrix into its HSS form, we are able to perform computations significantly faster with a complexity of O(rn) as opposed to $O(n^2)$, where r is the maximum rank and n is the size of the dense matrix.

References

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