

$$T = \frac{Rf}{1}$$

Look ohm \times 100 μF

$$10000 \times 0.0001$$

$$10000 \times 0.0001$$

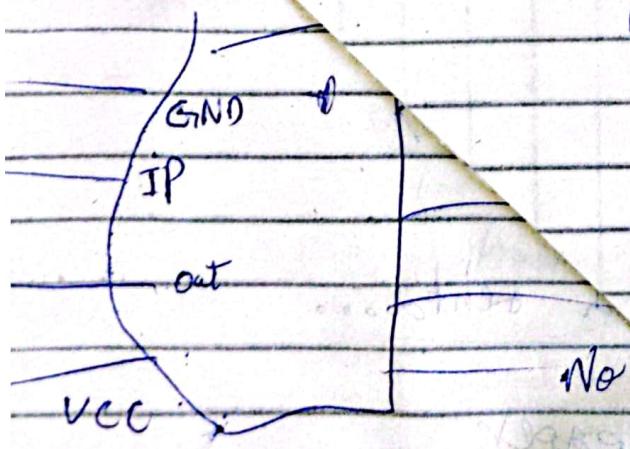
10

$$1 \times 10^5$$

Notes stands for
rem \rightarrow remember
IMP \rightarrow important

OPERATOR

OVERLOADING



Lecture Contents Checklist

(You should be Able to :)

- Overloading Boundary Operators
- Overload I/O "<<",">>" operators
- Overload postfix/prefix increment op.
- Overload Array subscript operator
- Overloading assignment operator

→ Similar to function overload

Operator Overloading

Using function-overloading to overload operators / redefine the functionality of operators is —

* Even when we do

```
int x=2;  
int y=494;  
cout << y-x << "\n";
```

The computer translates

it to operator +(x,y)

The function call.

Because

these are declared and defined

as member functions

inside the class.

under the library file.

* For us to add two objects of type
Complex w/o using operator overloading.

we would need to do that simple
task by using a function like this

```
class complex
```

```
{
```

```
    double real, imag;
```

```
public:
```

```
    double getReal() { return real; }
```

```
    double getImag() { return imag; }
```

```
}
```

```
complex add(const complex &c1, const complex &c2)
```

```
{
```

```
    complex temp;
```

```
    temp.real = c1.real + c2.real;
```

```
    temp.imag = c1.imag + c2.imag;
```

```
    return temp;
```

```
}
```

~~Three~~

THREE WAYS OF IMPLEMENTING OPERATOR

OVERLOADING :

① ~~overloaded~~ operator function implemented as member function of class.

② overloaded operator function implemented as a non-member function.

→ Notes: In this case, the operator function must be declared as a friend function if it requires direct access

to private or protected data members

* otherwise, getter/setter functions

* In the case of overloading, it would be used the binary '+' operator as a member function, when it would be called like in $c3 = c1 + c2;$

then

computer would translate that call to

$c3 = c1.operator +(c2);$

* While in the case of overloading the binary '+' operator as a non-member friend function, then, when it would be called like

$c3 = c1 + c2;$

→ would be translated as

→ $c3 = operator +(c1, c2);$

* In this case, both objects need to be passed to that non-member function.

(Defined as a ~~non-member~~
function)

e.g. code - 1

class complex

{

 double real, imag ;

public:

 complex & operator + (const & complex & passed)

{

 double r = real + (passed.real);

 double i = imag + (passed.imag) ;

 return (complex(r, i));

}

 complex(double r, double i) : real(r), imag(i) {}

};

int main()

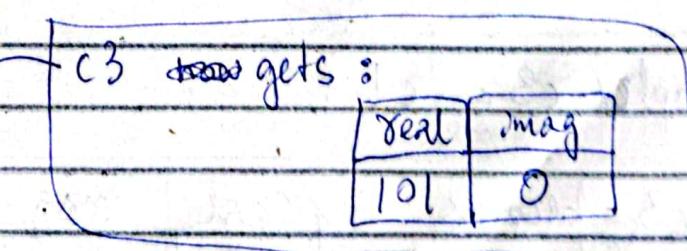
{

 complex c1(2, 4);

 complex c2(99, -4);

 complex c3 = c1 + c2;

}



E.g. Code 2. Defined as a non-member function and not a friend

class complex

{

double real, imag;

public:

complex(double r, double i) : real(r), i

double getReal() { return real; }

double getImaginary() { return imag; }

};

imag(i) { }

complex& operator + (const complex& c1, const complex& c2).

{

double r = (c1.getReal()) + (c2.getReal());

double i = (c2.getImaginary()) + (c1.getImaginary());

return obj (complex(r, i));

};

int main()

{

complex c1(34, -2);

complex c2(64, 4);

complex c3 = c1 + c2;

// c3 becomes

real	mag
98	2

};

e.g code 3

Defined as a friend function

// now class's object must be passed but no need of
getter/setter functions due to friend function.

class complex

{

 double real, imag;

public:

 complex(double r, double i) : real(r), imag(i) {}

// providing declaration of friend function;

friend complex& (const complex& c1, const complex& c2)

{

 double r = c1.real + c2.real;

 double i = c1.imag + c2.imag;

 return (complex obj(r, i));

}

* Note: Being a friend function to a class
only means that the function made as
friend can directly access its
private/protected data member given
that ~~is~~ object of that class is
passed to that function.

i.e. the object of class in main() can not
call it by obj.functionName();

just like the object of class could
directly call its member function for
operator overloading in e.g code-1.

Notes Implementing operator overloading using
non-member function and non-member
FRIEND function is very same.

Except that in the case of non-member
functions, we would've to use

Getters/setters to access private/protected

AND

data

Incase of non-member friend function,
we would ~~not~~ directly access
by using '.' operator with object name.

The
private data

IMP

Given that in both cases, both
objects are passed into
the functions

* However, the previous 3 e.g. codes do not cater such scenarios (scenarios A and B)

(A)

complex c1;

$$c1 = c1 + 2 \cdot 764;$$

(B)

complex b9;

$$b9 = 6 \cdot 781 + b9;$$

→ For Scenario-A

* we can create a member function - overloaded operator function. like:

complex & operator+(double d)

$$\text{double } r = \text{real} + d;$$

$$\text{double } i = \text{imag};$$

return(complex obj(r, i));

For Scenario-B

* AS THE L-OPERAND of the operator is of type double and not the user-defined complex datatype.

* AND we know that L-OPERAND calls the operator (in case if we declare the operator overloaded function as a member function.)

So As L-OPERAND is not of

* Complex datatype, we cannot go

* for having a member function

* to implement this "b9 = 6 · 781 + b9;"

for our case

friend

Scenario-B contd.

∴ we must go for a non-member function to implement this feature.

\hookrightarrow See code.go for this

class complex

5

double seal, mag;

public:

complex<double r, double i> real(s), imag(i)
{ { double d;

{ }

friend complex operator + (double a, const complex<c2>)

3;

complex& operator f. + (double d , const complex& c2) -

3

double r = d + c2.real;

double r = c2.cmag;

```
return (complex obj(r, i));
```

5

~~Summarizing previous stuff~~

class complex

{

double real, imag;

public:

complex(double r, double i) : real(r), imag(i) {}

complex& operator + (complex c2);

complex& operator + (double d);

friend complex& operator + (complex c2);

};

friend complex& operator + (double d, const complex& c2);

friend complex& operator + (double d);

c1 + c2

c1 + double

c1 + double

c1 + double

double + c1

operator + (d, c1)

Test this On Computer

Chaining

complex $c = c1 + c2 + c3;$

→ would automatically chain these multiple calls to the binary operator in sequence as:

~~operator~~

complex $c = (c1 \cdot \text{operator} + (c2)) \cdot \text{operator} + (c3)$

VM

complex $c = c1 + 4.96 + c2;$

would run fine as long as we have implemented

$c1 \cdot \text{operator} + (c2)$

AND

$c1 \cdot \text{operator} + (\text{double } d)$

→ BUT, this would be converted into (by compiler)

$c = (c1 \cdot \text{operator} + (4.96)) \cdot \text{operator} + (c2);$

See
prev.
pages

AND this would return

an object of type complex

(return
type is
Complex)

which could be further act as L-operand for this

Restrictions on Operator Overloading

* The overloaded operator must be an existing or valid operator.
You cannot create your own operator such as `xD` or `:`)

* Certain C++ operators can't be overloaded, such as

- sizeof
- dot operator `.`
- scope resolution operator `::`
- Conditionals like `"?"`, :" (used in ternary operator)

** The overloaded operator must have at least one operand of a user-defined data type.

(You can't overload operators working only on fundamental aka built-in data types)

e.g. `operator + (int a, int b)`

won't work

* Associativity, precedence & no. of arguments of an operator can't be changed.

e.g. `+` would remain left associative
`=` would remain right associative

Streamline insertion

Overloading "operator <<"

Premember aka ~~remember~~. This is a Binary Operator

(takes two operands)

E.g.

complex c1;
cout << c1;

L.H.S is
a C++ object
of the ostream
class which
is in the
iostream library

insertion
operator

object of
type complex

∴ To overload/redefine
the "<<" operator,
we must pass an
object of ostream class
(by reference)

AND

an object of the complex
class (by reference)

Reasons

L-OPERAND
was not of
type class

complex,
in fact we want
to print details

of the object of type complex so it would always
be like `cout << obj;`

To the operator overloaded
non-member friend function.

∴ We couldn't go for member function
to implement this.

(Test on computer, what of cout vs not replaced)

Note, cout is a C++ object of the ostream class.

So similarly,

we can give our own object name rather than "cout" (as its a reserved keyword)

To implement overloading the "<<" operator.

e.g code-1

```
class complex  
{ private:
```

```
    double real, imag;  
public:
```

```
    complex(double r, double i) : real(r), imag(i){};
```

```
    friend ostream& operator <<(ostream&, complex);
```

```
};
```

```
ostream& operator <<(ostream& out, complex c1)
```

```
{
```

```
    out << "InInt#Complex number is: "
```

```
    << c1.real << " + " << c1.imag << " i\n\n";
```

```
    return out;
```

```
}
```

This is done to enable cascading

See this page for revision later 2A

Short trick to rem. overloading insertion
& extraction operators

* Rem: cout & cin are objects of

↓ ↓
object of the object of the
of the istream class
ostream class in the ostream libr
in the ostream library

* AND for overloading

* AND that these are Binary operations

* So, we have to pass object of
ostream / istream & then the object of class.

* AND for cascading, we return
(the object)

therefore, return type is set

as either ostream &

OR

istream &

c.g code 2 for overloading stream extraction operator " >> "



class complex

{
double real, imag;

public:

complex(double r, double i) : real(r), img(i) {}

friend istream& operator >>(istream& in, complex&c)

};

friend istream& operator >>(istream& in, complex&c);
{

in >> c.real;

in >> c.imag;

return in;

}

int main()

{
complex c(4.96728, 13); and imaginary part

cout << "please enter the real part: ";

cin >> c;

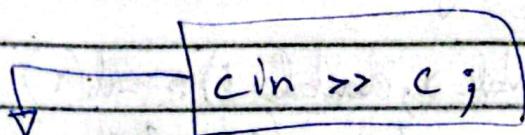
}

* Now c would have the user-entered values.

* Notes Within main() we would use cin & cout, ("in" and "out" were just used to represent objects of istream & ostream classes).

Infact, in the previous code,

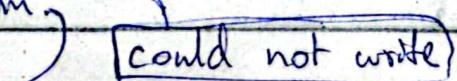
* within main()
the statements:

 cin >> c ;

basically is converted as: (By Compiler)

operator >> (cin, c);

* we could not write "cin" to represent
an object of ostream and
similarly "cout" to represent an object
of ostream,

 could not write

↳ within the function parameters
and the function body for the
operator overloaded functions of
">>" and "<<"

Because, "cout" & "cin"
are reserved keywords

Overloading Unary Operators

- $++$
- $--$
- $=$ (Negation operator)
- $!$ (Logical NOT operator)

Operating
on one

operand only

Overloading Logical NOT operator

① As a member function

```
class complex
```

```
{
```

```
    double real, imag;
```

```
public:
```

```
    complex(double r, double i) : real(r), imag(i) {}
```

```
    bool operator !() const
```

```
{
```

```
    return (real == 0 && imag == 0);
```

```
// TRUE is returned only if real & imaginary
```

```
// part of complex number are zero.
```

```
}
```

```
int main() {
```

```
    complex c1(0, 0);
```

```
    if (!c1) {
```

```
        cout << "In\nInt Empty\nIn";
```

```
}
```

```
} In this case yehi AAYEGA
```

Overloading Logical NOT

② As a non-member friend function

e.g.

```
class complex
```

```
{
```

```
    double real, imag;
```

```
public:
```

```
    complex(double real, double imag) : real(real), imag(imag)
```

```
{ }
```

```
friend bool operator !(complex c);
```

```
{};
```

```
friend bool operator !(complex c)
```

```
{
```

```
    return (c.real == 0 && c.imag == 0);
```

```
// returns TRUE only if condition fulfills.
```

```
{}
```

```
int main()
```

```
{
```

```
    complex c(0, 0);
```

```
    if (!c)
```

```
{
```

```
        cout << "Emptyyyy!!";
```

```
}
```

```
{}
```

Ques 8 Dummy argument (int) for postfix

Overloading unary / prefix operator

e.g.
class complex

{
double real, imag;

complex(double r, double i) : real(r), imag(i) {}

{
complex& operator ++(int) //postfix
{

real++;

imag++;

} return *this; // to enable cascaded
}

complex& operator --(int) //postfix
{

real--;

imag--;

return *this;

{
complex& operator +() {} //prefix

++real;

++imag;

return *this; }

{
complex& operator -() {} //prefix

--real; --imag;

return *this;

}

?:

int main()

{

complex c(4, 6);

$\text{++c}; \quad // (5, 7) \text{ now}$

$\text{c--}; \quad // (4, 7)$

$\text{++}((\text{++c})\text{++}); \quad // (7, 9) \text{ now}$

↑

Cascading was enabled through
return *this;

}

((See professor Hank Stalica's video for revision))

Overloading Array subscript [] operator

- * If a class has an array inside of it as a data member

then overloading [] operator lets us access (Valid()) indexes of that array by using subscript operator with the object inside main()!

① for this we have to pass index as parameter to the operator overloaded function for "[]" and

② It must return a reference to the memory address occupied by that specific index of the array (which we've passed)

So that what's returned could be used as a

L-OPERAND for the assignment operator like to enable

complex c;

c[4] = 496;

* Note: The overloaded "[]" operator should ensure valid index and deal with invalid index passed correctly to prevent undefined behaviour.

~~#include <iostream>~~

e.g. ~~#include <cstdlib>~~

class database

{

int employeeId[5];

public:

database() // CSTR

{

for (int i=0; i<5; i++)

{

employeeId[i] = 0;

}

,

int& operator [](int index)

{

if (!
{
if (index >= 0 && index <= 4))

cout << "Index out of bound!";

exit(EXIT_FAILURE);

} // ↳ this terminates the program.

else

{

return employeeId[index];

}

}

};

```
int main()
{
    database storage1;

    storage1[0] = 1942;
    storage1[1] = 2491;
    storage1[2] = 4219;
    storage1[3] = 2941;
    storage1[4] = 1249;

    cout << "Displaying vals: \n\n"
    for (int i=0; i<5; i++)
    {
        cout "ID-" << i+1 << ":" << storage1[i];
    }
}
```

→ Doing this in next example with
using overloaded stream insertion
operator

✗ Dynamic Array instead of statically allocated array.

```
#include <iostream>
#include <cstdlib> // for use of exit(EXIT_FAILURE)
// to terminate program if
class database // wrong/invalid index is given.
{
    int size;
    int* ptr;
public: int getSize() { return size; }
        database(int s) : size(s), ptr(new int[s]) {}
```

```
int& operator [](int index)
{
```

```
if (! (index >= 0 && index <= size))
{
```

```
cout << "Index out of bound";
exit(EXIT_FAILURE);
```

```
}
```

```
else {
```

```
return ptr[index];
```

```
}
```

```
friend ostream& operator << (ostream& out,
    database storage);
```

```
};
```

```
friend ostream& operator << (ostream& out, database storage)
```

```
out << "\n\nArray-contents as follows:\n\n";
```

```
for (int i=0; i < (storage.size()); i++)
{
```

```
    out << " " << ptr[i];
}
```

```
out << "\n\n";
```

```
}
```

(object name)

```
int main()
{
    storage
    database(3);
    //taking user inputs
    for (int i=0; i < storage.getSize(); i++)
    {
        cout << "Please enter input for index-"
        << i+1 << ". of array: ";
    }
}
```

storage[0] = 4278;

storage[3] = 9642;

storage[1] = 7542;

```
cout << storage; // prints the whole  
}                                Array pointed by  
                                         ptr.
```

Overloading Assignment Operator

- * Assigns values to a pre-existing (already constructed) object.
- * If a user defined Assignment operator is not provided then the default Assignment operator is provided by C++.

```
Complex c; // default CTOR  
Complex c2(4,9); // param. CTOR  
Complex c3(c2); // copy CTOR  
Complex c4 = c1; // copy CTOR
```

~~Complex~~ c4 = c3; // Assignment operator.

- * which always does shallow copy

Difference

Copy CTOR

- * making a new object and copying into it

(in these cases ;)

Assignment operator

Copying the contents into pre-existing object (to the left side of the operator.)

(for revision, see slide 50 of 54)

* If class contains pointers, then of *
still the default assignment operator is
used,

then

if one object deletes that
ptr / deallocates memory,
other object's ptr would
be affected too

As shallow copy was done

(SAME PROBLEM AS WAS WITH
THE DEFAULT COPY CTOR.)

eg. code on next pg

```
#include <string>
#include <iostream>
```

length
of name
including
NULL character

```
Class name {
    int length; char* ptr;
```

public:

```
name(const char* ptr2 = "")
```

{

```
    length = strlen(ptr2)
```

```
    ptr = new char[length];
```

```
    strcpy(ptr, ptr2);
```

}

//IMP //sem

~~// strcpy(destination, source);~~

// Default Assignment operator looks like

fnss

```
void operator = (const name & n2)
```

{

```
    length = n2.length; // (Shallow - Copy)
    ptr = n2.ptr;
```

}

*

Deep Copy ON NEXT PAGE.

Assignment operator for Deep

V. IMP

Note: We are going for deep copy in an Assignment because we have any pointer as data member in our class right?

→ So, if a pointer is there let's say of type characters

* And we assume that there are two ~~two classes A, B~~ objects of class Name.

~~then we are doing~~
* And, if in main() we are doing:

Name name1 ("Atmar");

Name name2 ("Bethumble");

then:

name1 = name2;

* * * DID YOU NOTICE? : the length

{ of the two arrays being pointed to by the ptr in the case of name1 & name2 would be different.

→ THIS IS WHY in implementing deep copy in the Assignment operator, we first delete any dynamic memory allocated to ~~the~~ the source ptr (to which we have to copy.)

The possibility of

- * Bcz of varying length of the arrays being pointed to by the ptr of both objects
i.e. \downarrow
 $\text{object 1} = \text{object 2}$;

* So first we deallocate any memory allocated to object-1's ptr.

* Then we pass the length of Array being pointed to by the ptr of object 2

\hookrightarrow and assign it to the variable for storing length on object 1

* Then we dynamically allocate an array of ~~as~~ the new length that has been recently stored in the (length) variable for object 1 by doing:

\downarrow $\left[\text{ptr} = \text{new char [length];} \right]$

* Now, that we have same length of Arrays, we can easily use strcpy to copy object 2's array into object 1's array:

Definition of Assignment Operator At Last of Page

by doing:

`strcpy(ptr, n2.ptr);`

i.e.:

`strcpy(ptr, object2.ptr);`

Definition of Assignment operation

* * also we have to pass right-operand
i.e. object 2 by reference.

* * overloading subscript operator "[]"
Reimpl. Subscript operator would only be implemented via a member function.

Similarly, overload assignment operator
also as a member function

bcz L-operand of operator is the

object of our user-defined class

→ unlike in the case

of stream operators, that's why we went
for non-member friend functions for them.

Definition:

void operator = (const Name& n2)

{

 delete ptr;

 length = n2.length;

 ptr = new char[length];

 strcpy(ptr, n2.ptr);

}

~~OOD~~
~~REVERSION~~

(PolyMorphism)

Polymorphism & Virtual functions

Lecture contents:

- Static Binding (Compile-time Polymorphism)
 - ⇒ Dynamic Binding (Run-time Polymorphism)
- Pointers and derived classes
- Derived class members invisible to base class pointers.
- Base class members visible to base class pointers
- Polymorphism
- Virtual functions
- Virtual destructors.

With the help of Apna college

~~Learned~~

Polyorphism

Compile-time Polymorphism
A.K.A. Static Binding

Run-time Polymorphism
A.K.A. Dynamic Binding

Function Overloading

Operator Overloading

With the help of virtual functions and function overriding within Inheritance

* This would also be static binding:

display() of base class would be called

Re: If base class's are not made virtual,

then inside main()
if its like: {

derived obj

base *ptr = &obj;

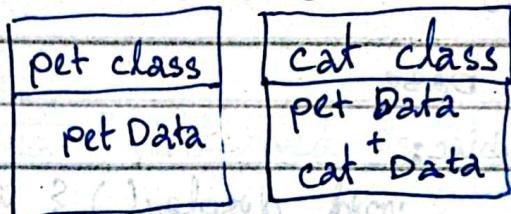
{ (*ptr).display(); }

Visiting Stack Overflowunderstanding

- * Why C++ does not allow pointers of type ~~Base~~ class to point to Derived objects of type Base class

without Casting which is not in our syllabus

Let's picture a scenario:



- There's a base class 'pet'.
- There's a derived class from 'pet' as 'cat'.

- * Now whatever or however is pet defined as by the programmer

* The cat would at least be a pet, and maybe more.

i.e. the cat would have all the characteristics defined of a pet in pet class and the cat could have its own distinct characteristics as well.

Notes from Slides

Static Binding

↳ Binding that takes place during Compilation

↳ Connecting a function call to a function body.

e.g #1

class base

{ public:

void display() {

cout << "In Base class's display called in";
}

}

class derived : public base {

public:

void display() {

cout << "In Derived class's display called in";
}

}

void main() {

base obj1;

obj1.display(); // calls base::display()

derived obj2;

obj2.display(); // calls derived::display()

}

IMP

Static Binding

Eg. 2

```
class base {
```

```
public:
```

```
    void display() {
```

```
        cout << "InBase vs Legend\n";
```

```
    }
```

```
}
```

```
class derived {
```

```
public:
```

```
    void display() {
```

```
        cout << derived "In derived vs Legend\n";
```

```
    }
```

```
}
```

```
int main() {
```

```
    derived obj;
```

```
    base *ptr = &obj;
```

```
(*ptr).display(); // use of static binding  
// function call would depend  
on type of pointer  
// base::display called.
```

```
obj.display(); // use of static binding
```

```
// derived::display called
```

```
base legend;
```

```
legend.display(); // use of static binding
```

```
// base::display called.
```

```
}
```

Remember

Dynamic Binding / runtime polymorphism would only occur once Base class's member functions (overridden in derived classes) are declared as virtual in the base class.

IMP

[Use of reference rather than pointer]

E.g #3 of Static Binding → [On next pg]

Rem: If virtual is not used with the ^{member} functions of base class,

- then at compile-time, ~~the computer~~ if a pointer is dereferenced to call a member function like in E.g. #2

OR

- if an ~~ref~~ object is passed by a reference, and then that reference into some other function

is used to call a member function

like in E.g #3

Then

In Case-1

- The computer would statically bind (at compile time) the function call with the function body
- * BASED ON THE TYPE OF POINTER.

In Case-2

- The computer would statically bind (at compile time) the function call with the function body
- ** BASED ON THE TYPE OF REFERENCE.

~~TMQ~~ * E.g #3 (Use of reference
& STATIC BINDING)

class base { public:

void display() { cout << "In Base Legend\n"; }

}

:public base // inheritance

class derived { public:

void display() { cout << "In Derived Legend\n"; }

}

void fun(* base& obj)

{

 obj.display(); // calls base::display

// no matter of
object-type is
base or derived,
still base::display()
would be called

→ // rem: function call bounded
with function body based
on TYPE OF REFERENCE
in case of static binding
just like TYPE OF POINTER
in case of POINTERS.

bcz

}

int main()

 derived obj;

 fun(obj); // calls base::display()

 // due to static Binding rules...

 // case-2 applied here (see prv. pg.)

 base obj2;

 fun(obj2); // base::display() called.

}

DYNAMIC BINDING

* Delaying the binding-process to the runtime is called —

Occurs

Once the Base class's member functions are declared virtual.

Side-Note: 'virtual' keyword ~~can~~ always comes before or after the return type of function but, ~~as~~ in virtual inheritance it's like:

< class derived : public virtual base { } >
written after public

Polyorphism

The ability to decide at runtime which among the several overridden member functions

is to be called by checking the actual type (the type of the object passed) and then calling that object's definition for that overridden function.

i.e. not based this decision is now not based upon the type of pointer or the type of reference.

But rather the type of object.

DYNAMIC BINDING

Reimplementing the previous

e.g. #2,
e.g. #3

using

```
class base {
```

```
public:
```

```
    virtual void display() {
```

```
        cout << "base is legend\n";
```

```
    };
```

```
class derived : public base
```

```
{
```

```
public:
```

```
    void display() {
```

```
        cout << "derived is legend\n";
```

```
}
```

```
void fun( base& obj )
```

```
{
```

```
    obj.display(); // function binding would not now
```

```
}
```

```
                                // be dependant upon the type
```

```
int main() {
```

```
    // of object, that is passed.
```

```
    base legend1;
```

```
    derived legend2;
```

```
    base *ptr = & legend2
```

```
(*ptr).display(); // calls derived::display()
```

```
fun(legend2); // calls derived::display()
```

```
base *ptr2 = & legend1;
```

```
(*ptr2).display(); // calls base::display()
```

```
fun(legend1); // calls base::display()
```

```
}
```

NOTE: DO NOT CONFUSE/MIX HOW
PROGRAM BEHAVES IN THE
CONCEPT OF ~~SINGLE SIMPLE~~ OVERRIDING & IN THE CASE
OF STATIC BINDING.

SEE THIS:

```
#include <iostream> using namespace std;  
class base  
{ public:  
    void display() {  
        cout << "base legendary\n";  
    }  
};  
class derived  
{ public:  
    void display() {  
        cout << "derived legendary\n";  
    }  
};  
int main() {  
    derived obj;  
    obj.display(); // calls derived::display()  
    // use of overriding  
    base *ptr = &obj;  
    ptr->display(); // calls base::display()  
    // due to static Binding  
}
```

Rem:

- * Derived class pointer can't point to Base class object
[Compiler error is thrown]
[See pg. 3]

- * Base class ptr pointing to object of derived class

Base class's

member function
declared
virtual

Base
Derived class's

member function
NOT declared
virtual

Then

Dynamic Binding

or

Runtime PolyMorph.

Then

Static Binding

or

Compile Time PolyMorph

↓

* binding function
call to function
definition would
be based on
the type of
object passed
or the type of
object being
pointed to.

Now, binding function
call to function
definition would be
based on the
type of reference
stated or the
type of pointer

| See previous code examples |

VIRTUAL DESTRUCTORS

- * Virtual functions and virtual destructor go hand in hand.
- * Making the destructor virtual will ensure that the correct destructor will be called for an object when pointers to the base classes are used.
- * If destructor of base class is not declared as virtual, then the destructor is statically bound based on the type of pointer and the destructor of derived class won't be called when the object is deallocated via the base class pointer.

So the object is partially destroyed.

For Visualization

see in PC:

OOP/Revising-for-final/revisingOOP/more.cpp

* A simple rule: To have a virtual destructor whenever your class has at least one virtual function.

* Note:

(IMP)

If the base class destructor is virtual, the derived classes' destructors will be virtual automatically.

* If Virtual Destructor is used, then:

* Don't forget the concept of

CSTRs & DSTRs

* Note: there's inheritance involved in implementing polymorphism.

So

doing:

baseClass *ptr = new derived;

would call

①. baseClass's CSTR

then → ②. derivedClass's CSTR

AND

doing:

delete ptr;

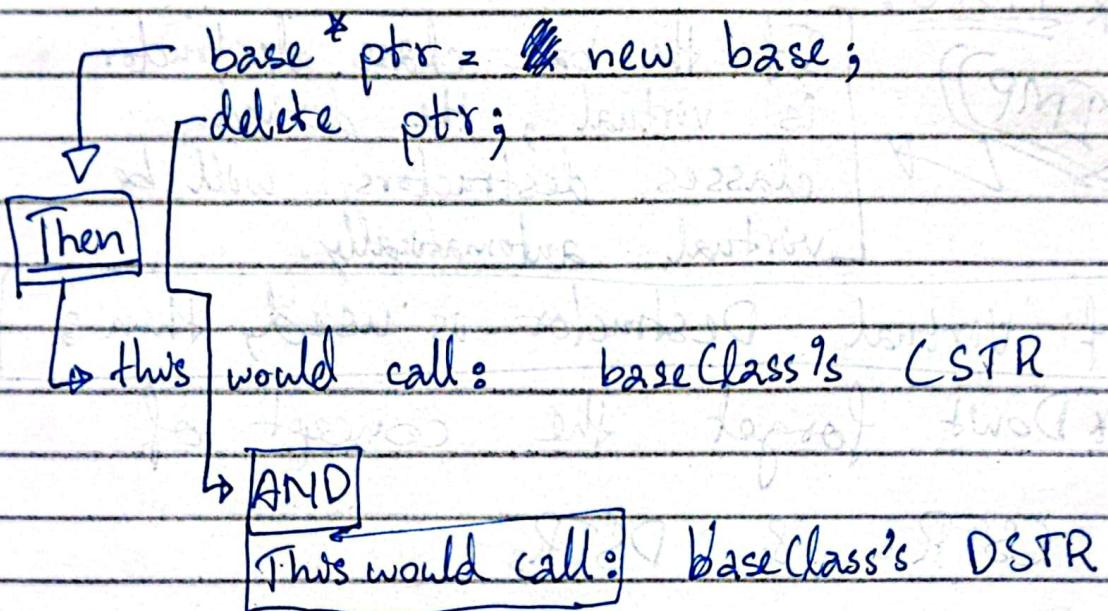
would call

①. derivedClass's DSTR

then → ②. baseClass's DSTR

[See next page]

Similarly, if we already released
the destructor from a base class
of its just:



* For visualization of above

this is $\times D$, see ↗

OOP/Revising-for-final/revisingOOP/nonVirtualDestructor.cpp

→ Purpose: to define an interface/blueprint

Abstract classes: classes from which
we're intended to instantiate

which are too generic

real objects

Abstract

works for the
derived
classes

&

Concrete
classes

instantiate objects.
for every

instance.

becomes an abstract
class by declaring any of its
virtual functions.

=

Abstract classes are not required
to implement their pure virtual functions.

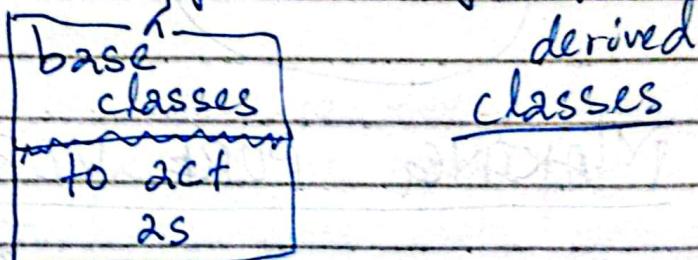
Derived classes MUST define all
pure virtual functions, failure to
define even 1 let's say out of 10 would make

→ Purpose: to define an interface/blueprint

* Abstract classes: classes from which it's never intended to instantiate any objects

→ i.e. classes which are too generic to define real objects

→ Normally used as frameworks for the



* Concrete classes

→ Classes which are used to instantiate objects.

* Must provide implementation for every member function they define.

SIMP: A Base class becomes an abstract base class by declaring any of its a pure virtual function.

* Abstract classes are not required to implement their pure virtual functions.

* Derived classes MUST define all pure virtual functions, failure to define even 1 let's say out of 10 would make

* that derived class an abstract class too.

So that derived class would then become un-instantiable in main().

Disadvantage

↳ Bcz Abstract classes

are

MAKING PURE VIRTUAL

class A

{ public:

 virtual void fun() const = 0;

"Pure Specifier"

Tells compiler that what makes it there's no implementation of it inside this class. PURE

Notes

In multi-level inheritance, it can happen that other than the Abstract base class, some derived class become abstract too (for not providing definitions for all ^{pure} virtual functions)

**** while some derived classes become concrete classes

① ** (maybe due to they having defined all the pure virtual functions)

- OR -

** (maybe due to they defining only some of the total pure virtual functions while definitions for the other pure virtual functions could have been inherited to them by the derived classes above them in the chain of inheritance.)

- OR -

③ → they defining none of the pure virtual functions; but due to inheritance and being too down in the chain of inheritance, they get as inheritance all the definitions for all the pure virtual functions.

→ These 3 scenarios are cases for derived classes to become concrete classes rather than abstract classes.

Replies The rules of polymorphism are used in implementing Abstract classes.

i.e. Notice: In the later examples that once functions are made virtual, member polymorphism automatically comes into play.

i.e. Dynamic Binding would occur with regards to functions when calling functions through Base class pointer.

of derived classes

* Also note that Although abstract classes are un-instantiable, pointers of

type Abstract classes could definitely be created in order to use runtime Polymorphism

↳ bcz member functions of base class would be pure virtual.

i.e. that they can't instantiate any objects.

(Defining pure virtual functions for the Abstract Base Class as well,

- * Also note that pure virtual functions, although meant to be left undefined in the Abstract base class ; can still be defined for that Abstract class to provide a common piece of code for the derived classes in order to facilitate reuse.

→ But they can not be defined within the class's curly brackets → Abstract Base

- * i.e. they cannot be defined as inline member functions
- * However, we can define them outside the class using scope resolution operator:

For this See

Code e.g. 2 coming

Later

e.g. 1 of Abstract Classes

```
class A {  
public:  
    virtual void x() = 0;  
    virtual void y() const = 0;  
};
```

```
class B : public A  
{ public:  
    virtual void x(); // provides some  
}; // definition as well
```

```
class C : public B  
{ public:  
    virtual void y(); // provides some  
}; // definition as well
```

```
int main(){  
    // A obj; // totally invalid  
    // Abstract classes can't instantiate objects
```

```
    // B obj; // Abstract class can't instantiate  
    // guess why B is abstract
```

```
# C obj; // totally legit!!! (Concrete class)
```

```
A * ptr; // legit  
B * ptr; // legit
```

*** pointers of type Abstract class can
be made to enable polymorphism

Refinement

* Pure virtual functions

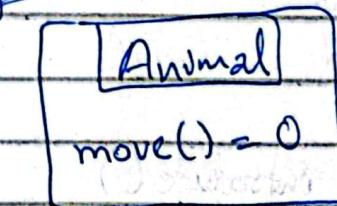
Let us specify
interfaces appropriately
& retain implementation
decisions till later
derived/sub-classes.

* As the hierarchy of

classes extends,
the pure virtual
functions of the
Abstract Base class
are replaced by
virtual functions

that fill in and
may override the
way of implementation.

e.g.



Fish
`move()`
`swim()`

Mammal
`move()`
`run()`

Bird
`move()`

Sparrow
`walk()`
`fly()`

Penguin
`waddle()`
`swim()`

code e.g. 2 of Abstract classes

```
class Animal // Abstract base class.
```

```
{ public:
```

```
    virtual void introduce() const = 0;
```

```
}
```

```
void Animal::introduce()
```

```
{
```

```
    cout << "\nI am an animal! Yeeheeee!\n";
```

```
}
```

```
Class cat : public Animal
```

```
{
```

```
public:
```

```
    virtual void introduce()
```

```
{
```

```
    // code-reuse with the help of  
    // calling Animal::introduce here;
```

```
    Animal::introduce();
```

```
    cout << "\nI am a cat,"
```

```
    << " how may I not"
```

```
    << " help you\n";
```

```
}
```

```
}
```

```
int main() {
```

```
    cat persian;
```

```
    persian.introduce();
```

```
}
```

* In the prev. exq. code,
run-time polymorphism could be added
by doing this - inside main().
}

cat persian;

Animal *ptr = &persian;

(*ptr).introduce();

{ } of base class
* As member functions are virtual
* So Dynamic binding would occur
and cat::introduce() would be
called.

((Payroll System

Implement Karna zehda...

Do it b4 paper

(E.A.))