

# Time Complexity Practice Questions



## 1. Summing Elements in an Array

```
int sum = 0;
for (int i = 0; i < n; i++) {
    sum = sum + i;
}
```

Time Complexity:  $O(n)$

## 2. Matrix Addition

```
for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
        c[i][j] = a[i][j] + b[i][j];
    }
}
```

Time Complexity:  $O(n^2)$

## 3. Normal Loops

### a. Simple Loop

```
for (int i = 0; i < n; i++) {
    stmt();
}
```

### b. Decrementing Loop

```
for (int i = n; i > 0; i--) {
    stmt();
}
```

Time Complexity:  $O(n)$

#### 4. Increment by Two

```
for (int i = 0; i < n; i += 2) {  
    stmt();  
}
```

Time Complexity:  $O(n)$

#### 5. Nested For Loops

```
for (int i = 0; i < n; i++) {  
    for (int j = 0; j < n; j++) {  
        stmt();  
    }  
}
```

Time Complexity:  $O(n^2)$

#### 6. Dependent For Loops

What happens if the inner loop is dependent on the outer loop?

```
for (int i = 1; i < n; i*=2) {  
    for (int j = 0; j < i; j++) {  
        for(int k=0 ; k<j ; k++){  
            }  
        stmt();  
    }  
}
```

#### 7. Non-Standard Outer Loop Execution

```
int p = 0;  
for (int i = 1; p <= n; i++) {  
    p = p + i;  
    stmt();  
}
```

Time Complexity:  $O(n^{1/2})$

### **8. Multiply i Value**

```
for (int i = 1; i < n; i = i * 2) {  
    stmt();  
}
```

Time Complexity:  $O(\log n)$

### **9. Divide i Value**

```
for (int i = n; i > 0; i = i / 2) {  
    stmt();  
}
```

Time Complexity:  $O(\log n)$

# Practice Problems

1)

```
bool List::Equalize_Occurrences(char d, int maxcount)
{
    Node* ptr = first;
    bool chk=false;
    Node* temp;
    while (ptr != NULL)
    {
        if (ptr->data == d)
        {
            chk = true;
            int count = 0;
            temp = ptr;
            while (ptr != NULL)
            {
                if (ptr->data == d)
                {
                    count++;
                    ptr = ptr->next;
                }
                else
                    break;
            }
            if (count > maxcount)
            {
                while (count > maxcount)
                {
                    del_after(temp);
                    count--;
                }
            }
            else if (maxcount > count)
            {
                while (count < maxcount)
                {
                    ins_after(temp,d);
                    count++;
                }
            } // else
        } // outer if
        ptr = ptr->next;
    } // while
    return chk;
}
```

**2)**

```
for(int i=2 ; i<n ; i=i*i){  
    ;  
}
```

**3)**

```
int m = (int)((15 + Math.round(3.2 / 2)) * (Math.floor(10 / 5.5) / 2.5) *  
Math.pow(2,5));  
for (int i = 0; i < m; i++) {  
    cout<<"hello";  
}
```

**4)**

```
for (int i = 1; i <= N * N; i *= 2)  
{  
    for (int j = 0; j < i; j++) {  
        cout<<"hello";  
    }  
}
```

**5)**

```
for(int i=1 ; i<n ; i*=2){  
    for(int j=0 ; j<i ; j++){  
        x = 0;  
    }  
}
```

**6)**

```
for(int i=1;i<=n;i++){  
    for(int j=2;j<=n;j=j*j){  
        cout<<i<<j<<endl;  
    }  
}
```

**7)**

```
for (i=n/2; i<=n; i++)
    for (j=1; j+n/2<=n; j++)
        for (k=1; k<=n; k = k * 2){
            cout<<"hello ";
            c++;
        }
```

**8)**

```
s=1;
While(s<=n)
{
    for(int i=1; i<=s; i++)
        cout<<" hello";
    s*=2;
}
```

**9)**

```
for (j=1; j<=n; j++)
    for (k=1; k<=j*3; k++)
    {
        cout<<"hello ";
    }
```

**10)**

```
for(int i=n/2;i<=n;i++){
    for(int j=2;j<=n;j=j*j){
        cout<<i<<j<<endl;
    }
}
```

**11)**

```
for (j=1; j<=n; j*=2)
    for (k=n; k>=1; k--) {
        for (i=1; i<=n; i*=3)
            cout<<"hello ";
    }
```

**12)**

```
For ( i = 1 to n)
    For ( j = 1 to i * i)
        If (j mod i == 1)
            Cout << i
```

**13)**

### Algorithm

Initialize count as 0  
Sort all numbers in increasing order using quicksort  
Remove duplicates from the array.  
Let D be the new array  
Do the following for each element A[i], where i varies from 1 to D  
    -> Binary search for A[i] + K in subarray from i+1 to D  
    -> if A[i] + K found, increment count  
Return count

**14)**

```
Int key, j
For (int i = 1; i < size; i++)
    Key = array[i]
    J = i
    While (j > 0 && array[j-1] > key)
        Array[j] = array[j-1]
        J - -
    Array[j] = key
```

**15)**

```
For (int i = 0; i < n; i++)
    For (int j = 1; j <= n*n; j++)
        If (j % 2 == 0)
            For (int k = 0; k < n; k++)
                Cout << "**";
```

**16)**

```
Void fun(int n, int k)
    For (int i = 1; i <=n; i++)
        Int p = pow(i, k)
        For (int j = 1; j <= p; j++)
            Stmt
```

**17)**

```
For (int i = 1; i<= n; i++)
    For (int j = 1; j < i*i; j*=2)
        Stmt
```

**18)**

```
int p = 0;
for (int i = 1; p <= n; i++) {
    p = p + i;
    stmt();
}
```



## Solution:

1)  $O(N)$

2)  $O(\log(\log n))$

3)  $O(1)$

4)

Outer loop executes  $\log_2 n^2 = 2 \log_2 n$  times.

Inner loop then executes

$1+2+4+8+16+\dots+(2^{\log_2 n})^2 = n^2$  times.

Time complexity =  $O(n^2)$

cout runtime =  $n^2$

5)

Outer loop executes  $\log_2 n$  times.

Inner loop executes

$1+2+4+8+16+\dots+2^{\log_2 n} = n$  times.

Time complexity =  $O(n)$

cout runtime =  $n$  times

6)

Outer loop executes  $n$  times.

Inner loop then executes

$2+2^3+2^{3^3}+\dots+2^{3^k}$  times.

$2^{3^k} = n \rightarrow 3^k = \log_2 n \rightarrow k = \log_3 \log_2 n$

Time complexity =  $O(n \log \log n)$

cout runtime =  $n \log \log n$ .

7)

First loop executes  $n/2$  times.

Second loop executes  $n^2/4$  times.

Third loop executes  $n^2 \log_2 n / 4$  times.

Time Complexity =  $O(n^2 \log n)$

cout runtime =  $n^2 \log n$

8)

Outer loop executes  $\log_2 n$  times.  
Inner loop executes  
 $1+2+4+8+16+\dots+2^{\log_2 n} = n$  times.  
Time complexity =  $O(n)$   
cout runtime =  $n$  times

9)

Outer loop executes  $n$  times.  
Inner loop executes  $3+6+9+\dots+3n$   
times =  $3n(n+1)/2$  times.  
Time complexity =  $O(n^2)$   
cout runtime =  $3n(n+1)/2$

10)

Outer loop executes  $n/2$  times.  
Inner loop then executes  
 $2+2^2+2^4+2^8+\dots+2^{2^k}$  times.  
 $2^{2^k} = n \rightarrow 2^k = \log_2 n \rightarrow k = \log_2 \log_2 n$   
Time complexity =  $O(n \log \log n)$   
Cout runtime =  $n \log \log n$ .

11)

First loop executes  $\log_2 n$  times.  
Second loop executes  $n \log_2 n$  times.  
Third loop executes  $n \log_2 n * \log_3 n$  times.  
Time Complexity =  $O(n (\log n)^2)$   
Cout runtime =  $n \log_2 n * \log_3 n$

**12)**  $O(n^3)$

**13)**  $O(n \log n + n + n \log n) = O(n \log n)$

**14)**  $O(n^2)$

**15)**  $O(n^4)$

**16)**  $O(n^{k+1})$

**17)**

Outer =  $O(N)$

Inner =  $O(\lg N)$

Total =  $O(N \lg N)$

**18)**  $O(n^{1/2})$