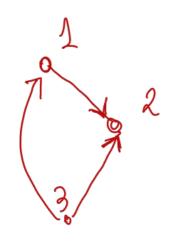
# Busca em Profundidade em Grafos Direcionados

#### Busca em Largura

```
BuscaLargura(Grafo G, vértice S)
 s.visitado = 1;
 Cria fila vazia F;
 ENFILEIRA (F,s);
 Enquanto F.tamanho > 0 faça
   u = DESENFILEIRA(F);
  Para todo vértice v € N+ (u) faça
    Se v.visitado==0 então
      v.visitado = 1;
      v.predecessor = u;
      ENFILEIRA (F,s);
```

#### Busca em Profundidade

Dado um grafo GDesmarcar os vértices
Definir uma pilha QDefinir uma raiz  $s \in V$  P(G,s)



```
P(Grafo G, vértice V)

marcar V

colocar V na pilha Q

para W \in \mathbb{N}^+(V)

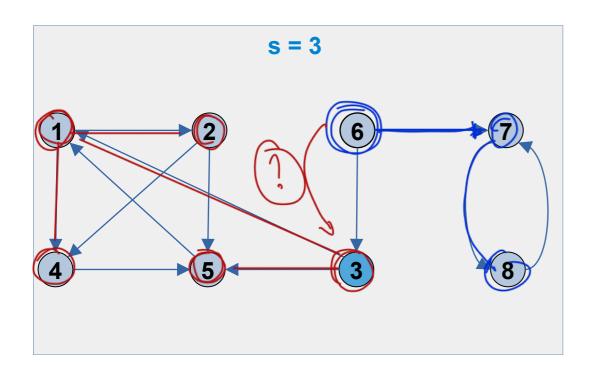
visitar (V, W)

se W não é marcado, então P(W)

retirar V de Q
```

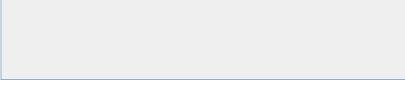
$$1/2$$
  $2 \in N^{+}(1)$   $N^{+}(3) = \{2, 1\}$   
 $N^{+}(2) = \emptyset$ 

# Percurso em Largura

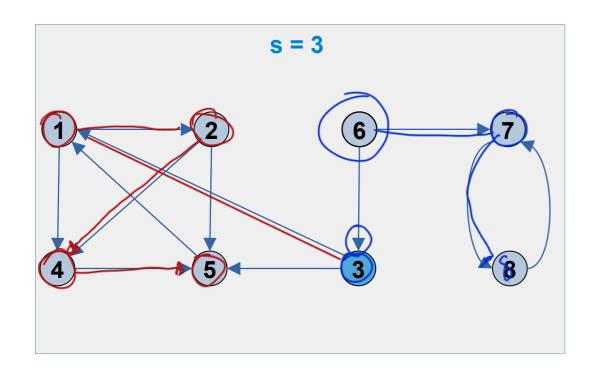


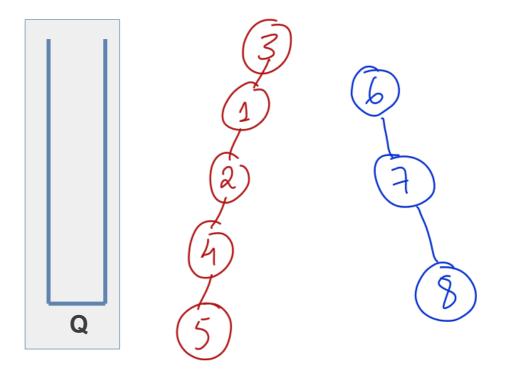
<u>(1</u>	(3)	
2 (	4	6
		8

Vértices	1	2	3	4	5	6	7	8
Visitado	1	1	1	1	1	7	1	1
Predecessor	3	1	_	1	3	_	6	7

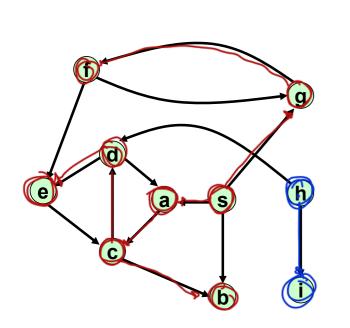


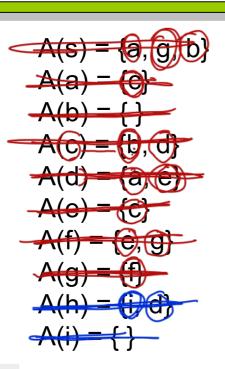
#### Percurso em Profundidade

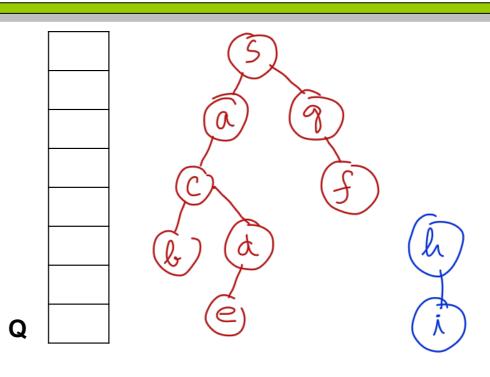




Vértices	1	2	3	4	5	6	7	8
Visitado	Į	1	1	1	7	1	1	1
Predecessor	3	1	-	2	4	_	6	7







```
P(Grafo G, vértice v)

marcar v

colocar v na pilha Q

para w \in N^+ (v)

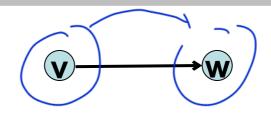
visitar (v, w)

se w não é marcado, então P(w)

retirar v de Q
```

Vértices	S	a	b	С	d	е	f	g	h	i
Visitado	1	1	1	1	1	1	1	1	1	1
Predecessor	_	5	C	$\alpha$	C	d	8	5	_	h
P(E)	1	2	4	3	5	6	8	<u></u>	9	10
P(S)	8	5	1	4	3	2	6	7	10	9

Considere a visita a aresta (v,w)



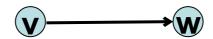
Caso 1. v é alcançado antes de w na busca

1.1 − Se w estava desmarcado antes da visita, então (v,w) é aresta da árvore (floresta) de profundidade (floresta) de profundidade

1.2 - Se w estava marcado antes da visita, então (v,w) é aresta de avanço

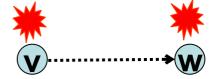


Considere a visita a aresta (v,w)



Caso 2. w é alcançado antes de v na busca

2.1 − Se  $w \in Q$  no momento da visita, então (v,w) denomina-se aresta de retorno



2.2 - Se  $w \notin Q$  no momento da visita, então (v,w) é denominada aresta de cruzamento



Dado um grafo G

Desmarcar os vértices

Definir uma pilha Q

Definir uma raiz  $s \in V$ P(G,s)

```
P(Grafo G, vértice v)

marcar v

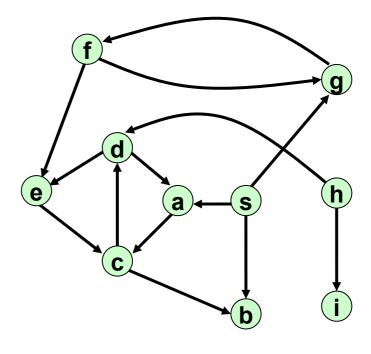
colocar v na pilha Q

para w \in N^+(v)

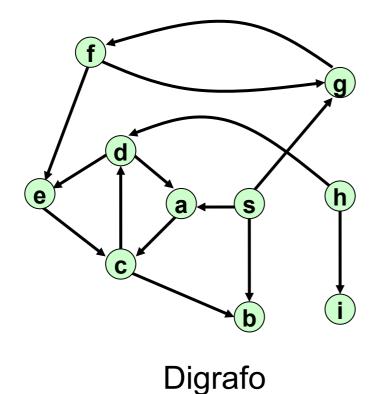
visitar (v, w)

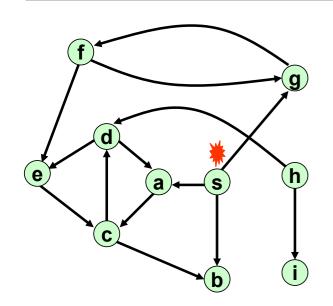
se w não é marcado, então P(w)

retirar v de Q
```



Digrafo





```
S
```

```
S
```

```
P(Grafo G, vértice v)

marcar v

colocar v na pilha Q

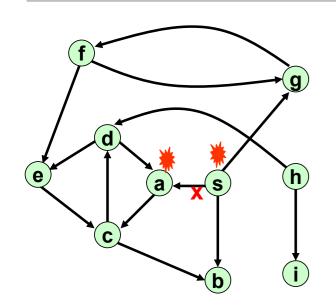
para w \in N^+ (v)

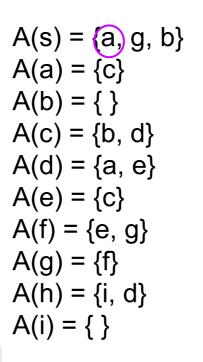
visitar (v, w)

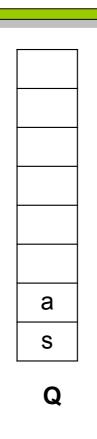
se w não é marcado, então P(w)

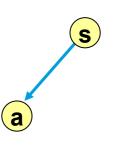
retirar v de Q
```

Q









```
P(Grafo G, vértice v)

marcar v

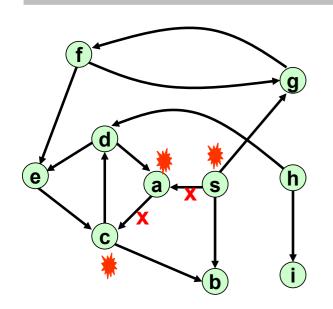
colocar v na pilha Q

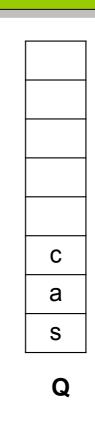
para w \in N^+(v)

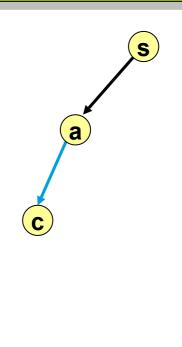
visitar (v, w)

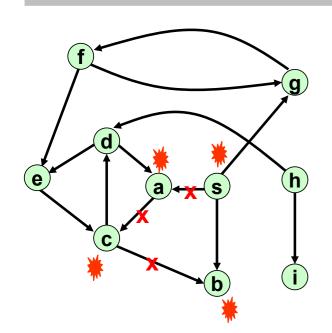
se w não é marcado, então P(w)

retirar v de Q
```

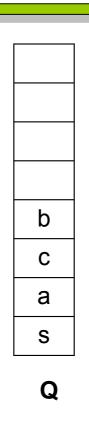


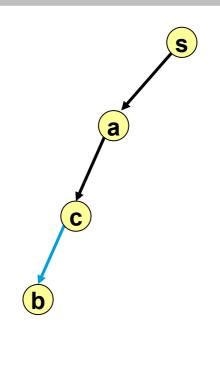




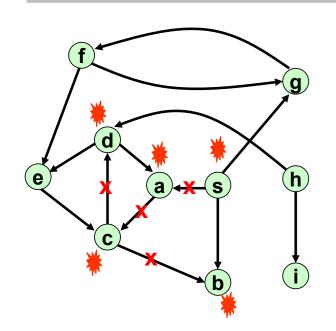


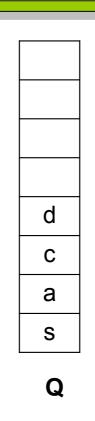
$$A(s) = \{a, g, b\}$$
 $A(a) = \{c\}$ 
 $A(b) = \{c\}$ 
 $A(c) = \{b, d\}$ 
 $A(d) = \{a, e\}$ 
 $A(e) = \{c\}$ 
 $A(f) = \{e, g\}$ 
 $A(g) = \{f\}$ 
 $A(h) = \{i, d\}$ 
 $A(i) = \{ \}$ 

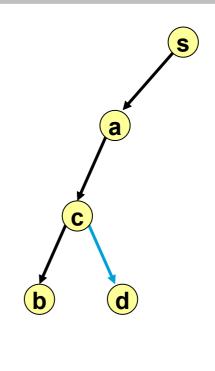


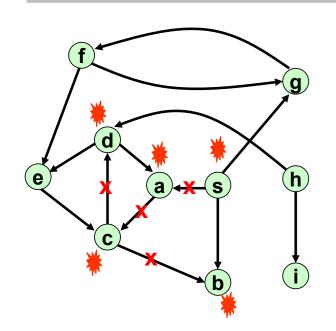


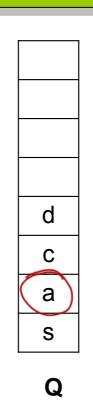
b é removido da pilha

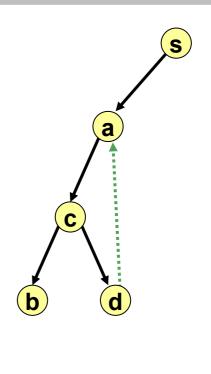


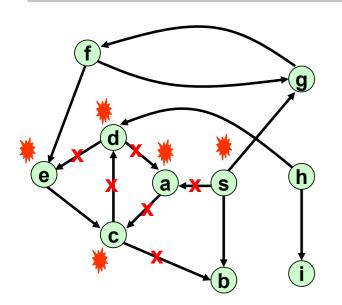


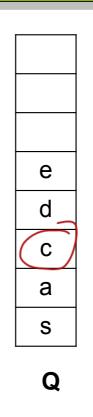


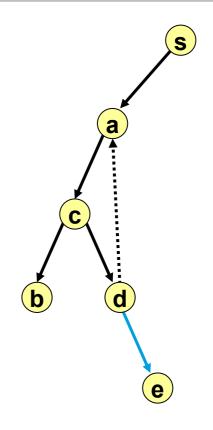


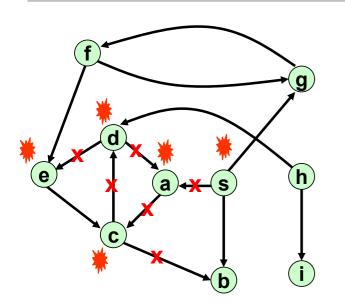


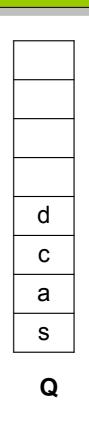


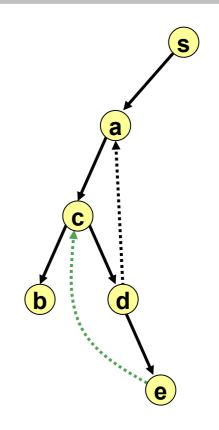






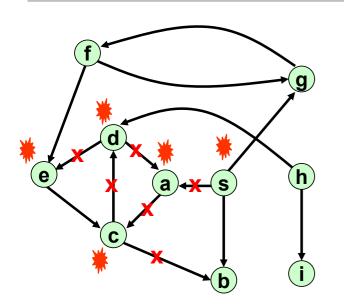


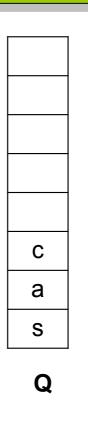


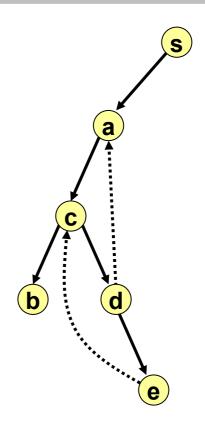


e é removido da pilha

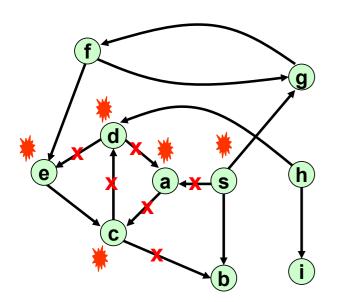
Considere a aresta (v,w). Se  $w \in Q$  no momento da visita, então (v,w) denomina-se **aresta de retorno** 

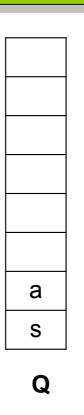


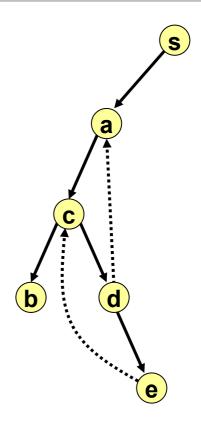




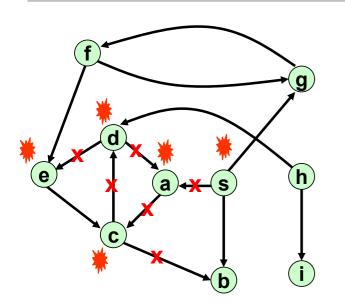
d é removido da pilha

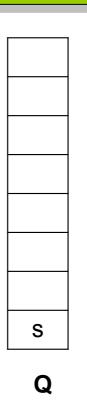


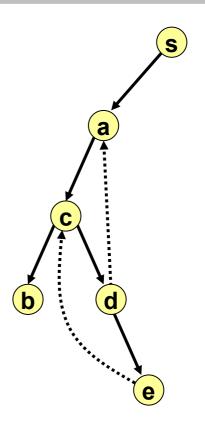




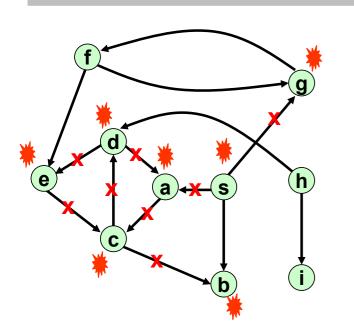
c é removido da pilha

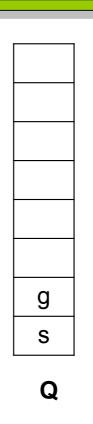


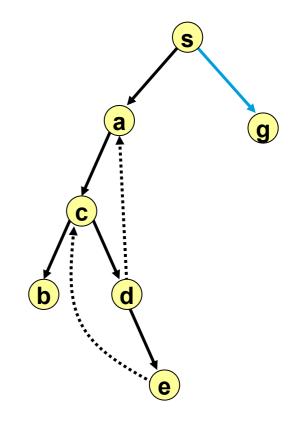


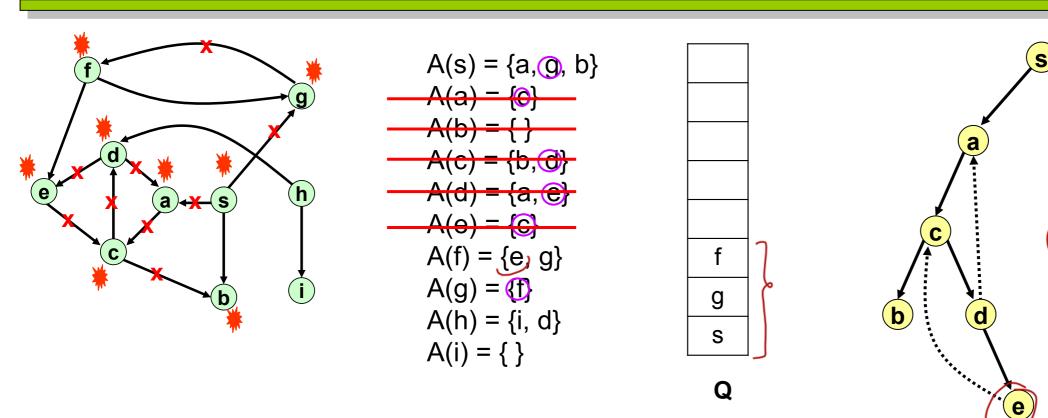


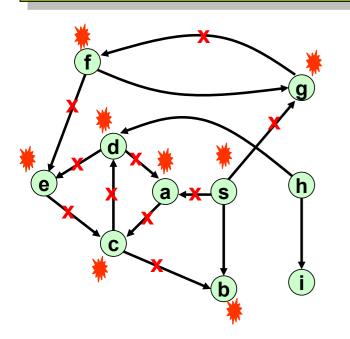
a é removido da pilha

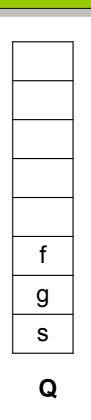


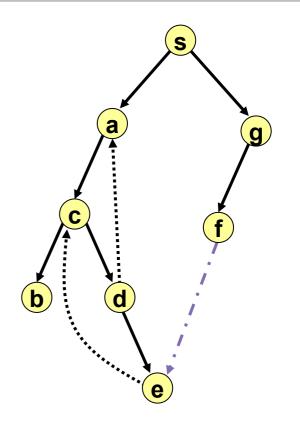




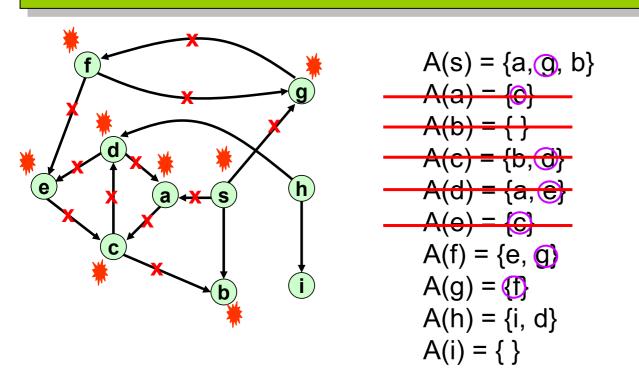


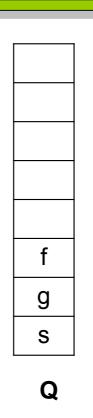


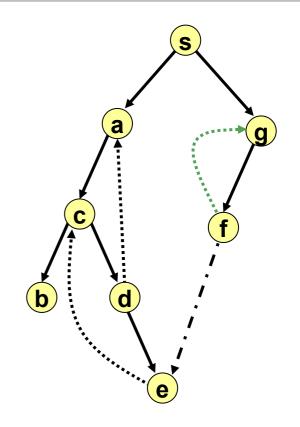


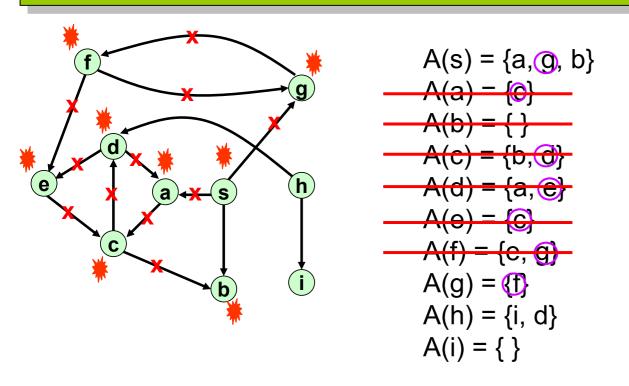


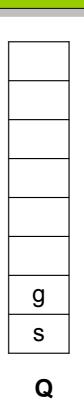
Considere a aresta (v,w). Se  $w \notin Q$  no momento da visita, então (v,w) é denominada aresta de cruzamento

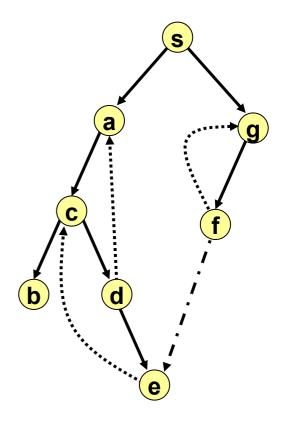




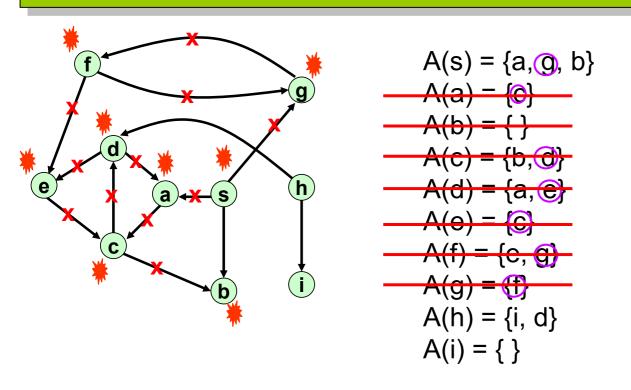


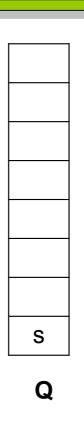


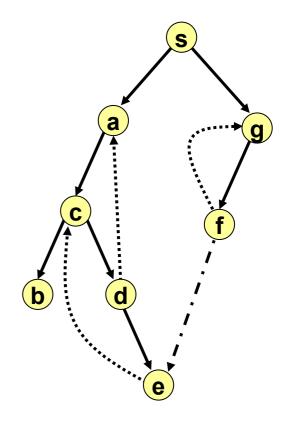




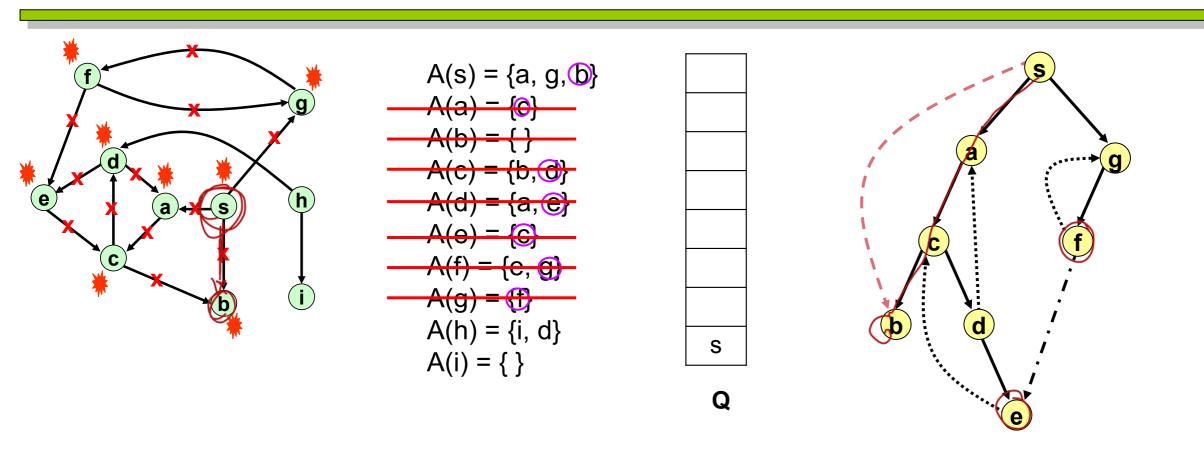
f é removido da pilha



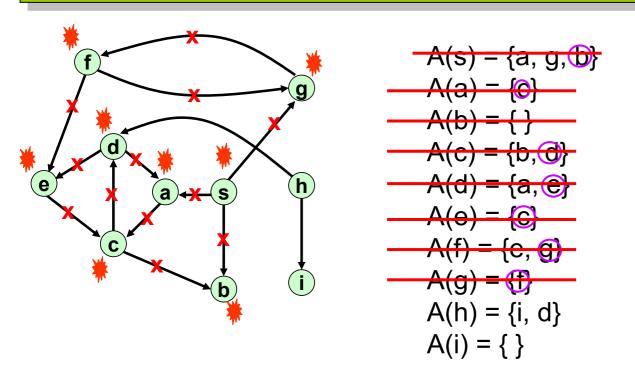


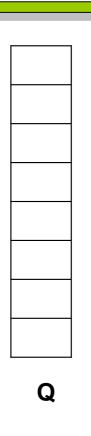


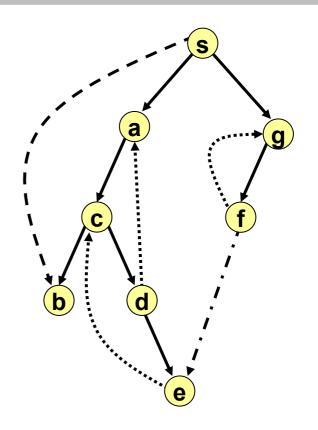
g é removido da pilha



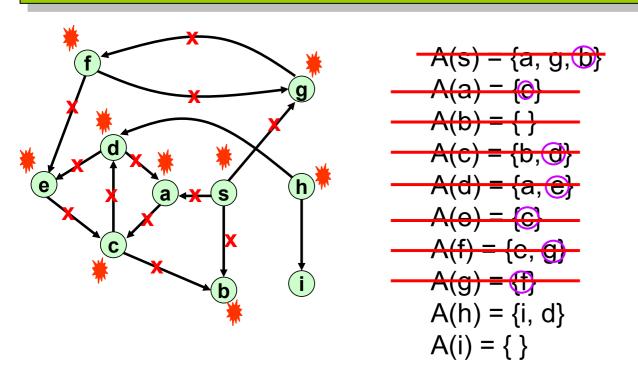
Considere a aresta (*v*,*w*). Se *w* estava marcado antes da visita, então (*v*,*w*) é **aresta de avanço** (*s* é ancestral de *b*)

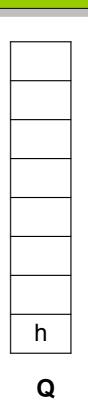


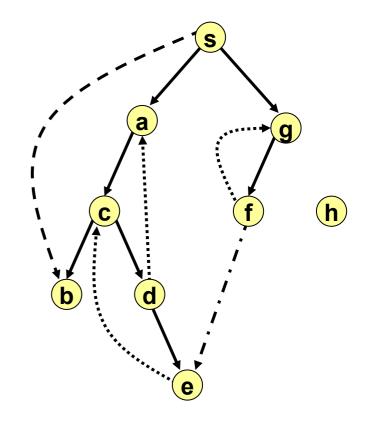


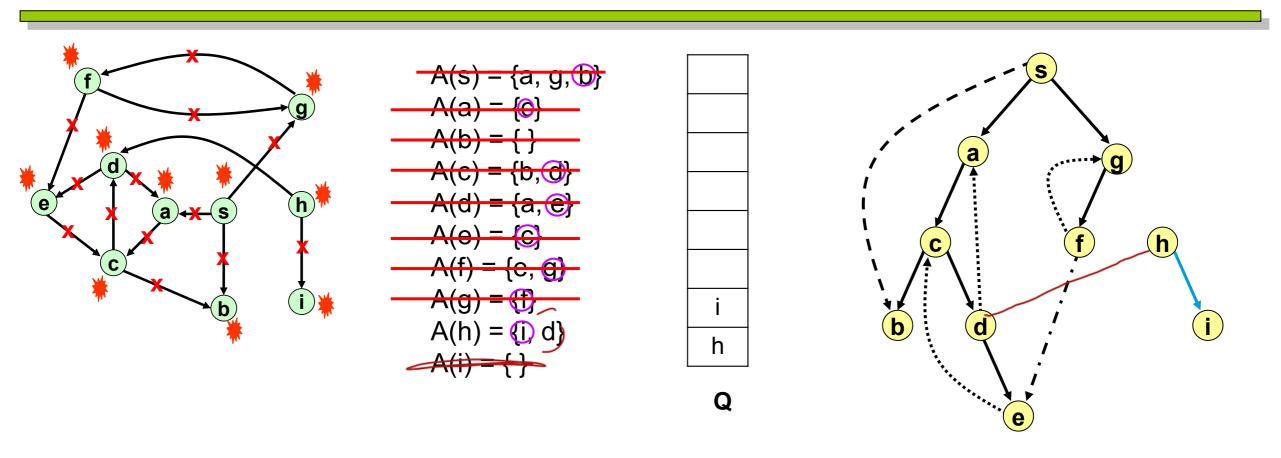


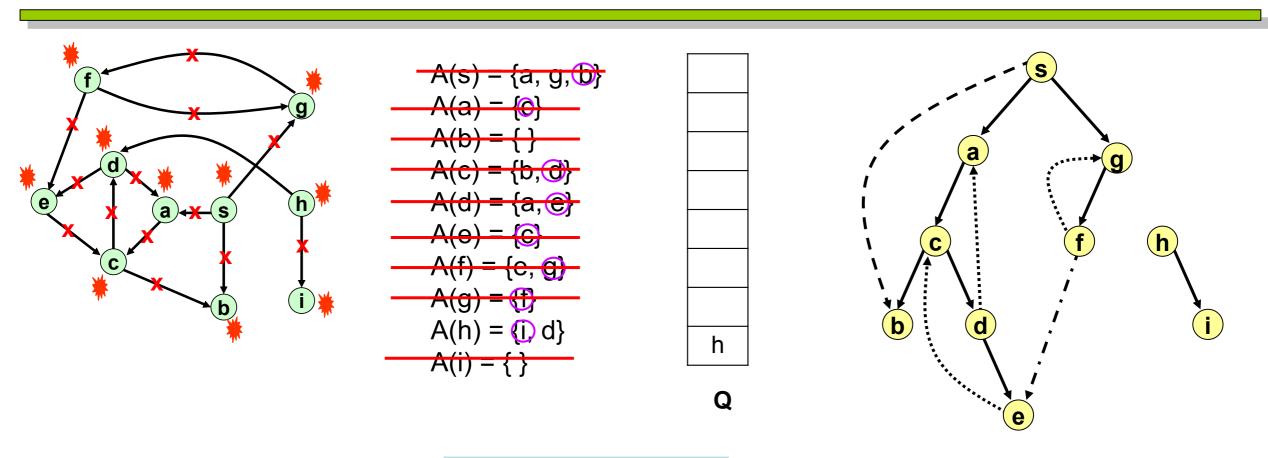
s é removido da pilha



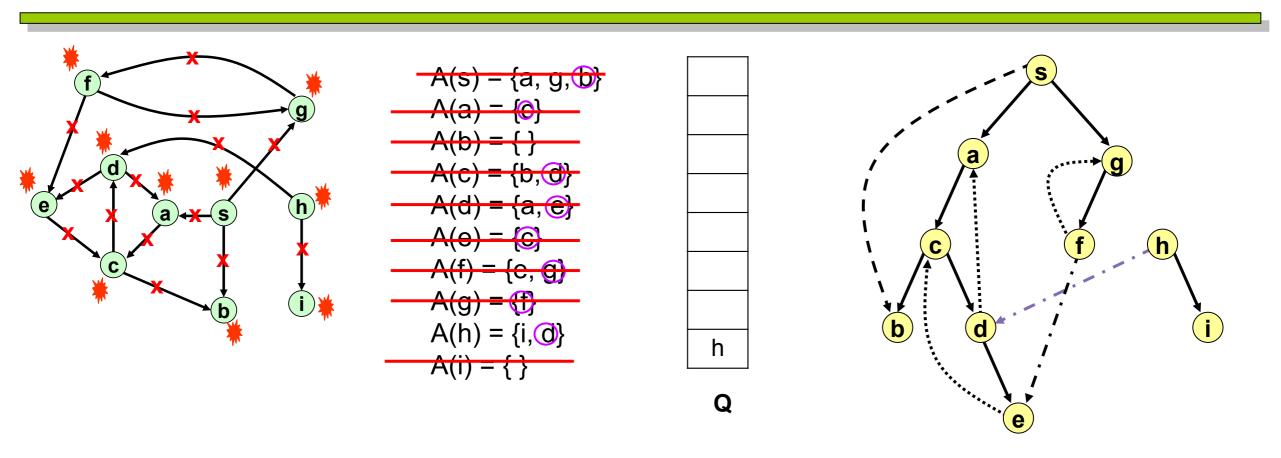




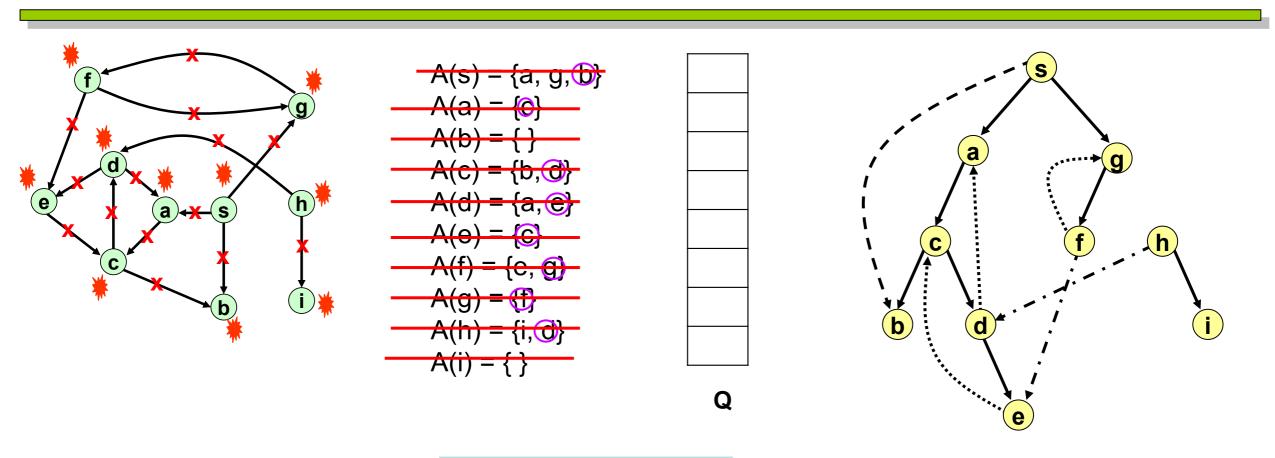




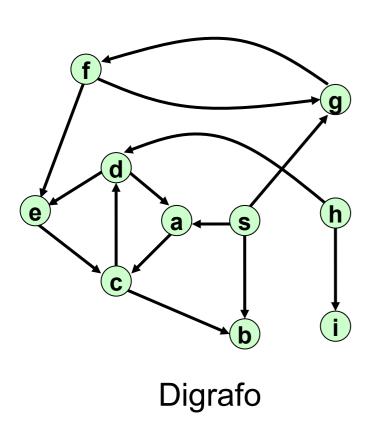
i é removido da pilha



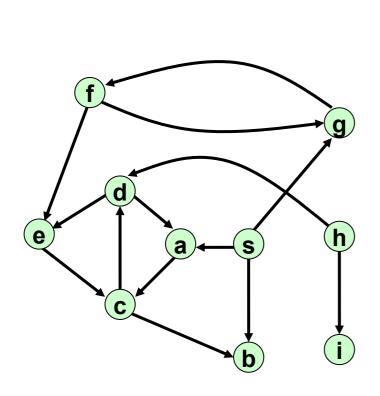
Considere a aresta (v,w). Se  $w \notin Q$  no momento da visita, então (v,w) é denominada **aresta de cruzamento** 



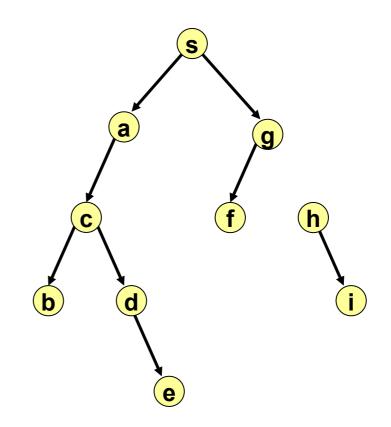
h é removido da pilha



Aresta da árvore ---→ Aresta de avanço Aresta de retorno Aresta de cruzamento

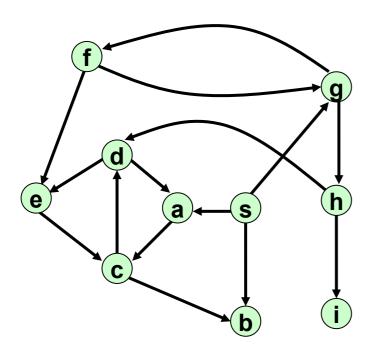


Digrafo

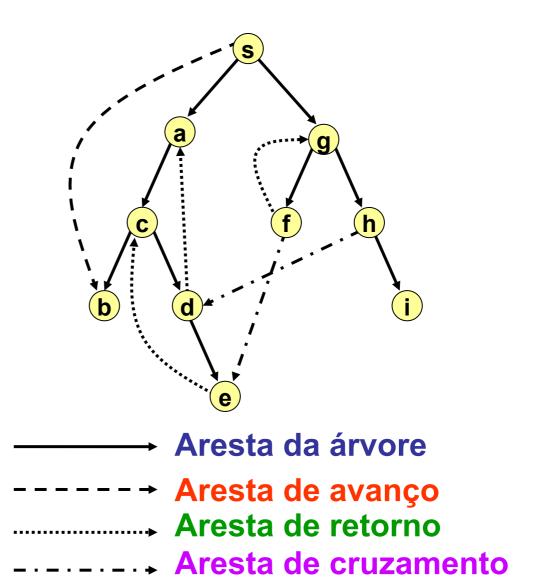


Floresta de Profundidade

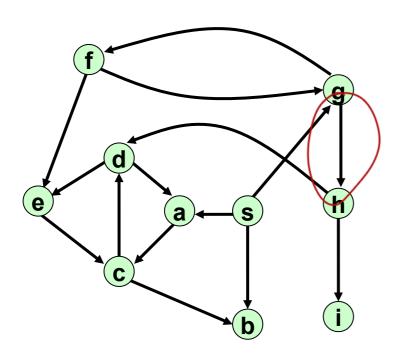
#### **Outro Exemplo:**



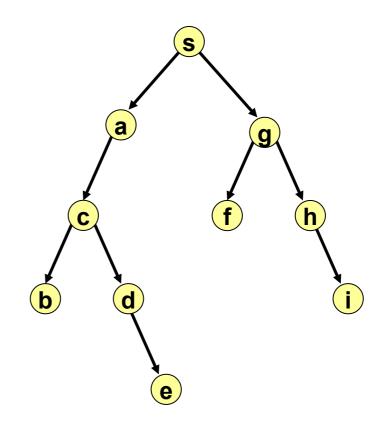
Digrafo



#### **Outro Exemplo:**



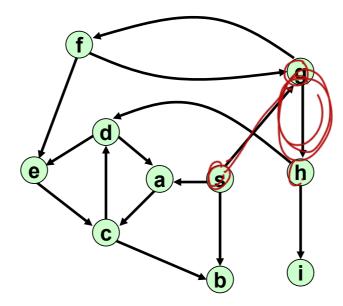
Digrafo

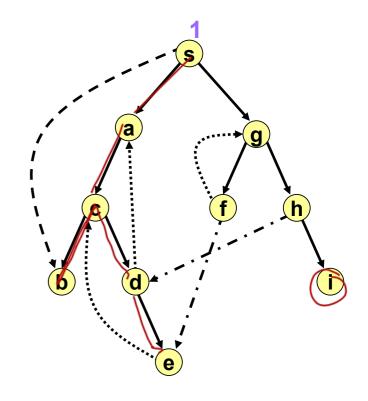


**Árvore de Profundidade** 

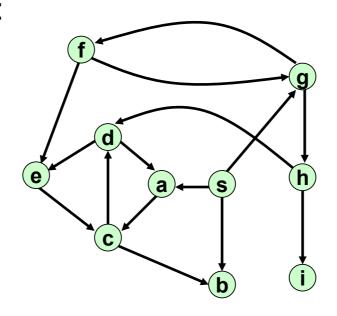
Assim como para grafos simples, a ordem de entrada e saída dos vértices na pilha Q são importantes para algumas aplicações.

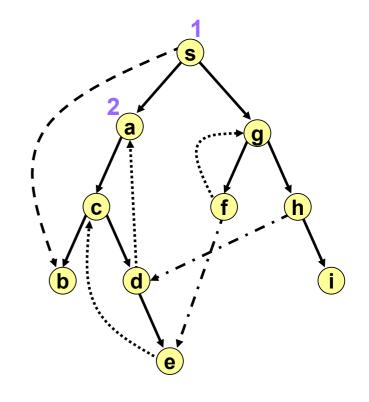
Assim sendo, para cada vértice v define-se profundidade de entrada de v, PE(v), e profundidade de saída de v, PS(v), respectivamente, como sendo o número de ordem que v foi incluído e excluído da pilha Q.



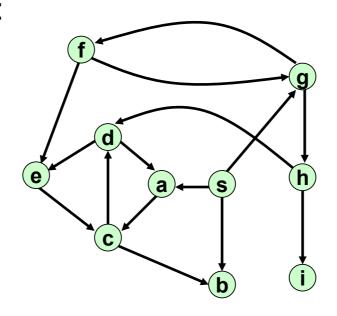


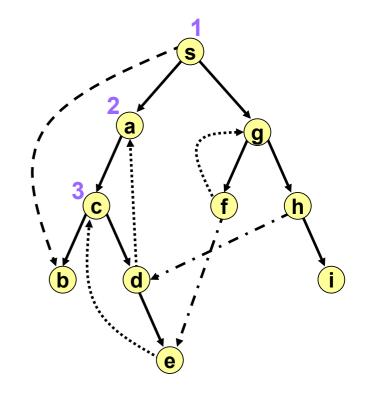
V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3	5	6	8	7	9	10
PS(v)	10	5	1	4	3	2	6	9	8	7



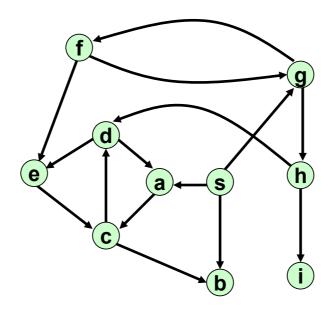


V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2								
PS(v)										

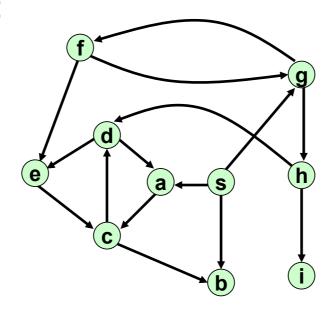


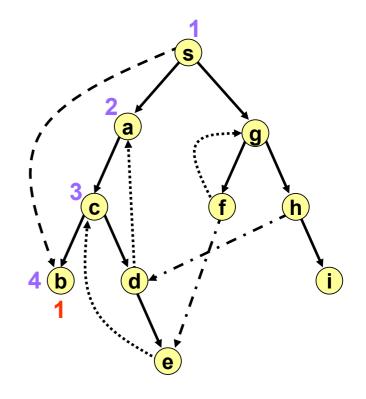


V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2		3						
PS(v)										

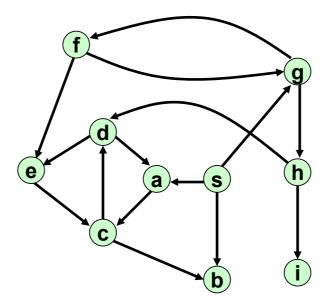


V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3						
PS(v)										



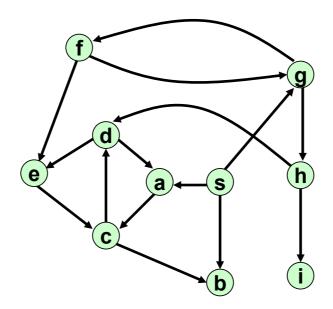


V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3						
PS(v)			1							



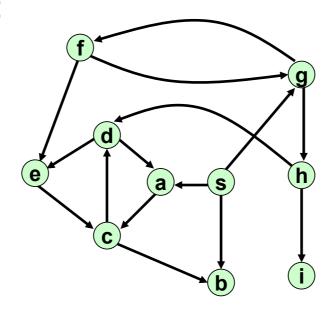
3 c f h
·e

V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3	5					
PS(v)			1							



2 a	<b>.</b> g
3 c 4 b 5 d	f h
1 6 e	i

V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3	5	6				
PS(v)			1							

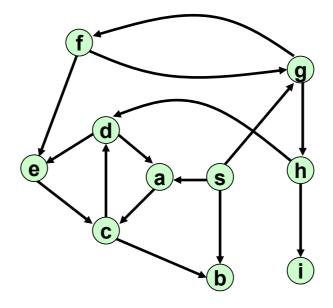


3 c f h i
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V	S	а	b	С	d	е	f	g	h	ij
PE(v)	1	2	4	3	5	6				
PS(v)			1			2				

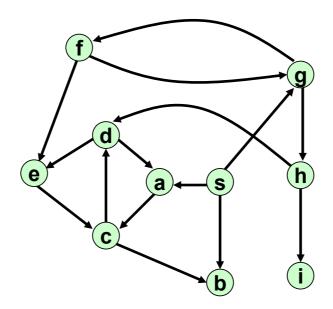
Para o exemplo anterior, temos:

41,1343



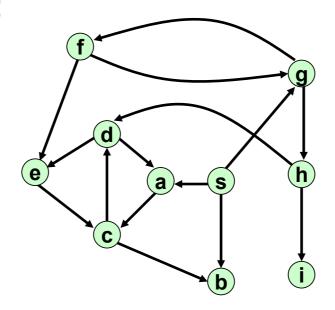
3 c f h	)
6 e 2	

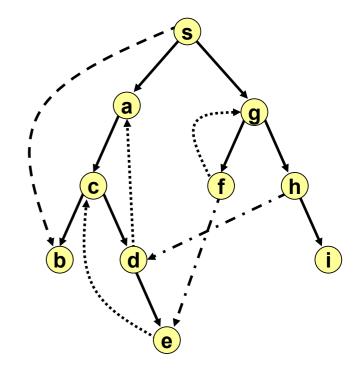
V	S	а	b	С	d	е	f	g	h	.—
PE(v)	1	2	4	3	5	6				
PS(v)			1		3	2				



3 c 4 f h
1

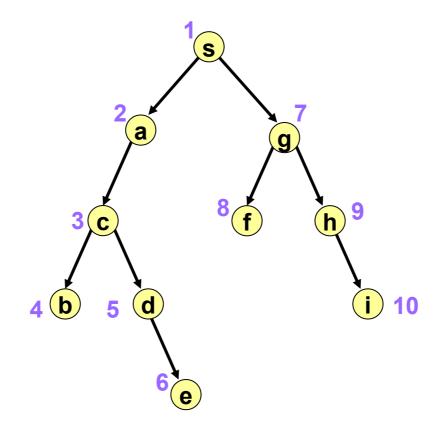
V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3	5	6				
PS(v)			1	4	3	2				





V	S	а	b	С	d	е	f	g	h	i
PE(v)	1	2	4	3	5	6	8	7	9	10
PS(v)	10	5	1	4	3	2	6	9	8	7

Observa-se que a sequência dos vértices v de um digrafo D em ordem crescente de PE(v) corresponde a um caminhamento pré-ordem na árvore de profundidade produzida pela busca correspondente.



Os valores de entrada e saída podem ser utilizados para classificar as arestas de um digrafo através de uma busca em profundidade com o seguinte lema:

Lema: Considere D(V,E) um digrafo de raiz s, (v,w) uma aresta de D, e B um percurso em profundidade com s a raiz da busca, então:

- 1. (v,w) é aresta da árvore ou avanço se e somente se PE(v) < PE(w)
- 2. (v,w) é aresta de retorno se e somente se PE(v) > PE(w) e PS(v) < PS(w)
- 3. (v,w) é aresta de cruzamento se e somente se PE(v) > PE(w) e PS(v) > PS(w)