

Sets

Introduction

Sets are **primitive** objects when doing classical, old-school, pen-and-paper mathematics:

- no *definition*;
- only *rules* about how these objects work (unions, intersections, etc.).

That's all you need: do you ever look at $S \rightarrow T$ as $S \subseteq T$?

Objects normally represented by a set are formalised in Lean as *types* with some extra-structure.

So, for Lean, sets are **no longer primitive objects**; yet

- sometimes we still want to speak about *sets* as collections of elements
- we want then to play the usual games.

Definitions

+++ Every set lives in a **given** type, it is a set of elements (*terms*) in it:

```
variable (α : Type) (S : Set α)
```

expresses that α is a type and S is a set of elements/terms of the type α . On the other hand,

```
variable (S : Set)
```

does not mean "let S be a set": it means nothing and it is an error.

+++

+++ A set coincides with the test-function defining it.

Given a type α , a set S (of elements/terms of α) is a *function*

```
S : α → Prop
```

so $(\text{Set } \alpha) = (\alpha \rightarrow \text{Prop})$.

- This function is the "characteristic function" of the set S ;
- the $a \in S$ symbol means that the value of S is **True** when evaluated at the element a ;
- So, the positive integers are a *function*!



Yet, given a function $P : \alpha \rightarrow \text{Prop}$ we prefer to write $\text{setOf } P : \text{Set } \alpha$ rather than $P : \text{Set } \alpha$ to avoid *abusing definitional equality*.

Some examples:

1. How to prove that something belongs to a set?
2. Positive naturals;
3. Even numbers;
4. An abstract set of α given by some P .



+++

+++ Sub(sub-sub-sub)sets are not treated as sets-inside-sets.

Given a (old-style) set S , what is a subset T of S ? At least two answers:

1. Another set such that $x \in T \rightarrow x \in S$.
2. A collection of elements of S .

Now,

1. stresses that T is a honest set satisfying some property;
2. stresses that it is a set whose elements "come from" S .

We take the **first approach**: being a subset is *an implication*

```
def (T ⊆ S : Prop) := ∀ a, a ∈ T → a ∈ S
```



One can also *upgrade* sets to types: $T : \text{Set } S$ for $S : \text{Set } \alpha$ means $T : \text{Set } \uparrow S = \text{Set } (S : \text{Type}^*)$.

Some examples:

1. Double inclusions;
2. Subsets as sets;
3. This upgrade (*coercion*) from $\text{Set } \alpha$ to Type^* .

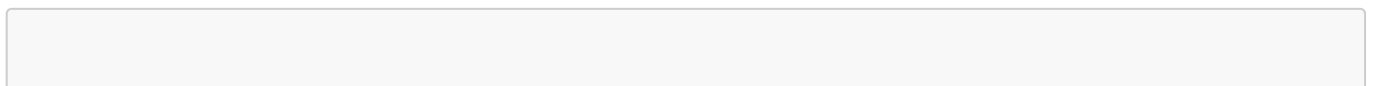


+++

Operations on Sets

+++ **Intersection**

Given sets $S \ T : \text{Set } \alpha$ have the



```
def (S ∩ T : Set α) := fun a ↦ a ∈ S ∧ a ∈ T
```

- Often need **extensionality**: equality of sets can be tested on elements;
- related to *functional extensionality*: two functions are equal if and only if they take the same values on same arguments;
- not strange: sets *are* functions.

⌘

+++

+++ Union

Given sets $S, T : \text{Set } \alpha$ we have the

```
def (S ∪ T : Set α) := fun a ↦ a ∈ S ∨ a ∈ T
```

And if $S : \text{Set } \alpha$ but $T : \text{Set } \beta$? **ERROR!**

⌘

+++

+++ Universal set & Empty set

- The first (containing all terms of α) is the constant function $\text{True} : \text{Prop}$

```
def (univ : Set α) := fun a ↦ True
```

- The second is the constant function $\text{False} : \text{Prop}$

```
def (∅ : Set α) := fun a ↦ False
```

Bonus: There are infinitely many empty sets!

+++

+++ Complement and Difference

- The complement is defined by the negation of the defining property, denoted S^c .

```
Sc = {a : α | ¬a ∈ S}
```

The superscript ^c can be typed as `\^c`.

- The difference $S \setminus T : \text{Set } \alpha$, corresponds to the property

```
def (S \ T : Set α) = fun a ↦ a ∈ S ∧ a ∉ T
```

⌘

+++

+++ Indexed Intersections & Indexed Unions

- Can allow for fancier indexing sets (that will actually be *types*, ça va sans dire): given an index type I and a collection $A : I \rightarrow \text{Set } \alpha$, the union $(\bigcup i, A i) : \text{Set } \alpha$ consists of the union of all the sets $A i$ for $i : I$.
- Similarly, $(\bigcap i, A i) : \text{Set } \alpha$ is the intersection of all the sets $A i$ for $i : I$.
- These symbols can be typed as $\backslash U = \bigcup$ and $\backslash I = \bigcap$.

⌘

+++

Functions

Introduction

Functions among types are *primitive notions*: given two types α and β , the type $\alpha \rightarrow \beta$ exists as a type *axiomatically*. And there is a rule saying

```
a : α, f : α → β ⊢ f a : β
```

that can be read

1. If the type of a is α and the type of f is $\alpha \rightarrow \beta$, then the expression $f a$ has a meaning, and it is a term of type β .
2. f behaves like a function from α to β , sending a to $f a$.

Sometimes we want to speak about functions among *sets* : **functions among sets are different gadgets than functions among types**.

Let's inspect the following code:

```
example (α β : Type) (S : Set α) (T : Set β) (f g : S → T) :
  f = g ↔ ∀ a : α, a ∈ S → f a = g a :=
```

It *seems* to say that $f = g$ if and only if they coincide on every element of the domain, yet... ⌘

+++ Take-home message

To apply $f : \alpha \rightarrow \beta$ to some $s \in S : \text{Set } \alpha$, *restrict* it to the *subtype* $\uparrow S$ attached to S .

+++

Operations

Given a function $f : \alpha \rightarrow \beta$ and sets $(S : \text{Set } \alpha)$, $(T : \text{Set } \beta)$, there are some constructions and properties that we are going to study:

+++ The **image** of S through f , noted $f '' S$.

This is the set $f '' S : \text{Set } \beta$ whose defining property is

$$f '' S := \text{fun } b \mapsto \exists x, x \in S \wedge f x = b$$

Unfortunately it comes with a lot of accents (but we're in France...): and with a space between f and $''$: it is not $f'' S$, it is $f '' S$.

⌘

+++

+++ The **range** of f , equivalent to $f '' \text{univ}$.

I write *equivalent* because the defining property is

$$\text{range } f := (\text{fun } b \mapsto \exists x, f x = b) : \beta \rightarrow \text{Prop} = (\text{Set } \beta)$$

This is not the verbatim definition of $f '' \text{univ}$: there will be an exercise about this.

+++

+++ The **preimage** of T through f , denoted $f^{-1} T$.

This is the set

$$f^{-1} T : \text{Set } \alpha := \text{fun } a \mapsto f a \in T$$

This also comes with one accent and *two* spaces; the symbol $^{-1}$ can be typed as `\^-1`.

⌘

+++

+++ The function f is **injective on** S , denoted by $\text{InjOn } f S$ if it is injective (a notion defined for functions **between two types**) when restricted to S :

$$\text{def } : \text{InjOn } f S := \forall x_1 \in S, \forall x_2 \in S, f x_1 = f x_2 \rightarrow x_1 = x_2$$

In particular, the following equivalence is not a tautology:

```
example : Injective f ↔ InjOn f univ
```

rather, it will be an exercise for you.

✂

+++