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# The Beauty and Joy of Computing

## Lecture #24 Limits of Computing



## SEPARATION OF POWERS



Trouble brews as the CIA spies on Congressional committee computers as they investigate CIA interrogation techniques. Unfortunately, hacking is sometimes used as a powerful weapon between countries as well as within a country itself.

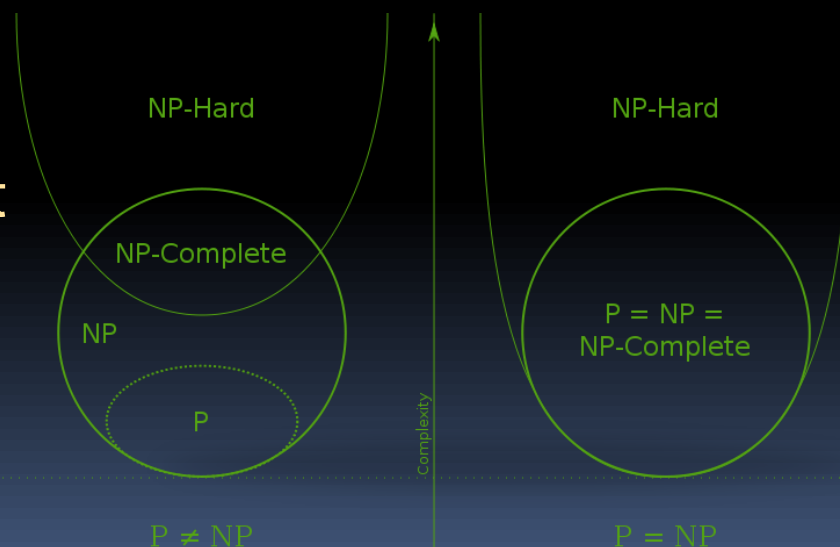
# Computer Science ... A UCB view

- CS research areas:
  - Artificial Intelligence
  - Biosystems & Computational Biology
  - Database Management Systems
  - Graphics
  - Human-Computer Interaction
  - Networking
  - Programming Systems
  - Scientific Computing
  - Security
  - Systems
  - Theory
    - Complexity theory
  - ...



# Let's revisit algorithm complexity

- Problems that...
  - are tractable with efficient solutions in reasonable time
  - are intractable
  - are solvable approximately, not optimally
  - have no known efficient solution
  - are not solvable



<http://en.wikipedia.org/wiki/NP-complete>

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# Tractable with efficient sols in reas time

- Recall our algorithm complexity lecture, we've got several common orders of growth

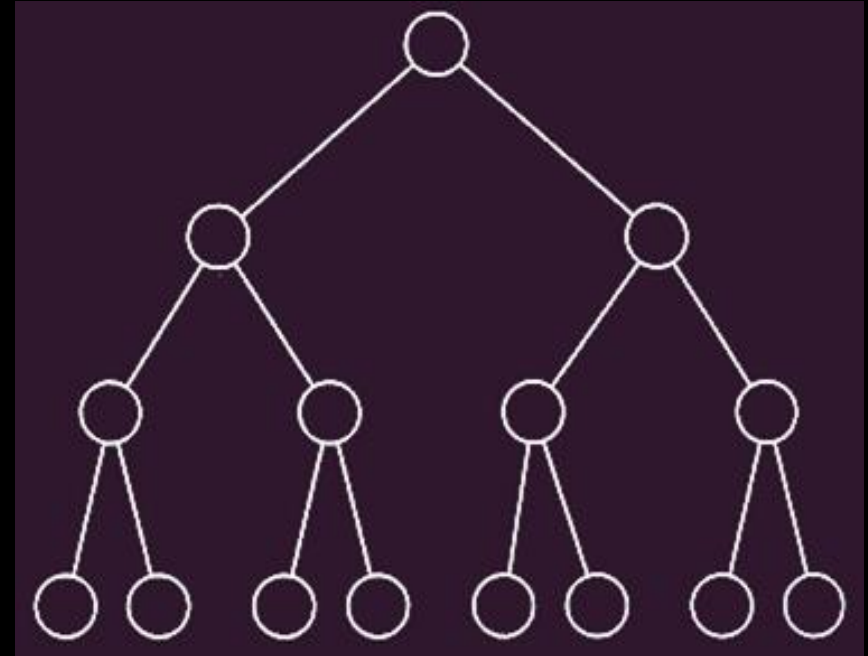
- Constant
- Logarithmic
- Linear
- Quadratic
- Cubic
- Exponential

- Order of growth is polynomial in the size of the problem
- E.g.,
  - Searching for an item in a collection
  - Sorting a collection
  - Finding if two numbers in a collection are same
- These problems are called being "in P" (for polynomial)



# Intractable problems

- Problems that can be solved, but not solved fast enough
- This includes exponential problems
  - E.g.,  $f(n) = 2^n$ 
    - as in the image to the right
- ~~This also includes poly-time algorithm with a huge exponent~~
  - ~~E.g.,  $f(n) = n^{10}$~~
- ~~Only solve for small n~~



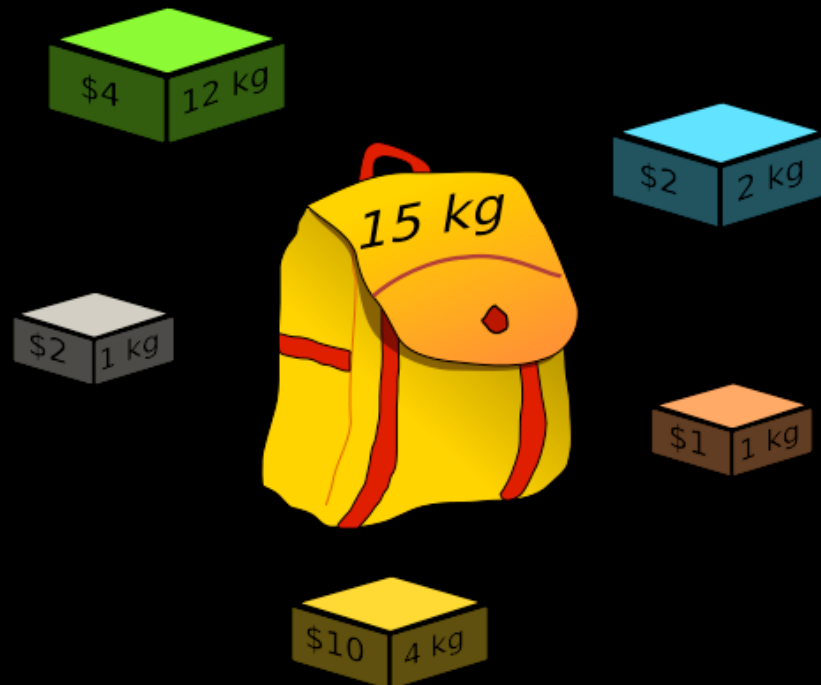
Imagine a program that calculated something important at each of the bottom circles. This tree has height  $n$ , but there are  $2^n$  bottom circles!







What's the most you can put in your knapsack?



- a) \$10
- b) \$15
- c) \$33
- d) \$36
- e) \$40

## Knapsack Problem

You have a backpack with a weight limit (here **15kg**), which boxes (with weights and values) should be taken to maximize value?

**(any # of each box is available)**

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# Solvable approximately, not optimally in reasonable time

- A problem might have an optimal solution that cannot be solved in reasonable time
- BUT if you don't need to know the perfect solution, **there might exist algorithms which could give pretty good answers in reasonable time**



## Knapsack Problem

You have a backpack with a weight limit (here 15kg), which boxes (with weights and values) should be taken to maximize value?



# Have no known efficient solution

- Solving one of them would solve an entire class of them!
  - We can transform one to another, i.e., reduce
  - A problem P is “hard” for a class C if every element of C can be “reduced” to P
- If you’re in both “in NP” and “NP-hard”, then you’re “NP-complete”

-2      -3      15  
 14      7      -10

Subset Sum Problem

Are there a handful of these numbers (at least 1) that add together to get 0?

- If you guess an answer, can I verify it in polynomial time?
    - Called being “in NP”
    - Non-deterministic (the “guess” part)
- Polynomial

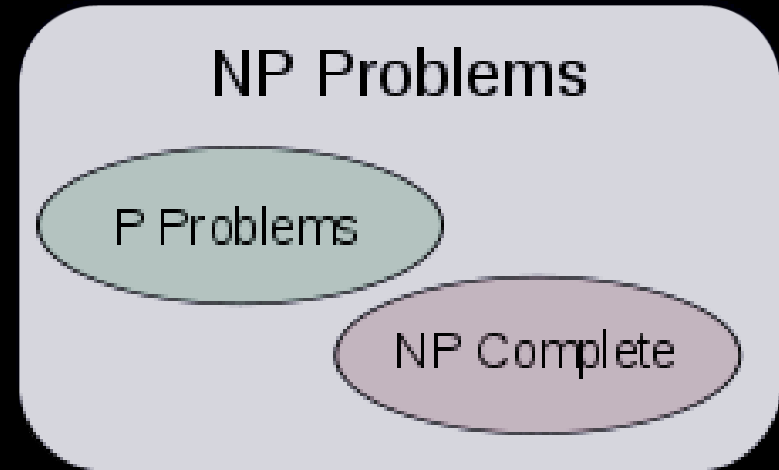




# The fundamental question. Is $P = NP$ ?

- This is THE major unsolved problem in Computer Science!
  - One of 7 “millennium prizes” w/a \$1M reward

If  $P \neq NP$ , then



- All it would take is solving ONE problem in the NP-complete set in polynomial time!!
  - Huge ramifications for cryptography, others
- Other NP-Complete
  - Traveling salesman who needs efficient route below a certain cost to visit all cities and return home



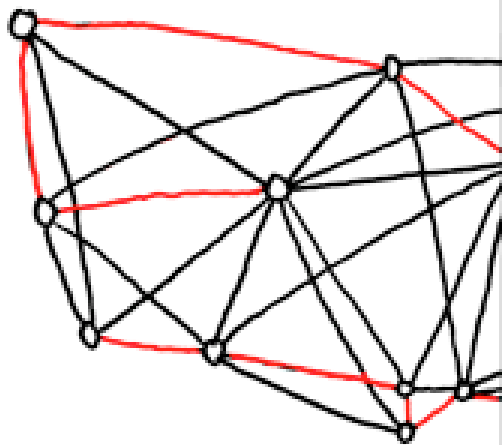
# MY HOBBY: EMBEDDING NP-COMPLETE PROBLEMS IN RESTAURANT ORDERS

CHOTCHKIES RESTAURANT	
~ APPETIZERS ~	
MIXED FRUIT	2.15
FRENCH FRIES	2.75
SIDE SALAD	3.35
HOT WINGS	3.55
MOZZARELLA STICKS	4.20
SAMPLER PLATE	5.80
~ SANDWICHES ~	
BARBECUE	6.55



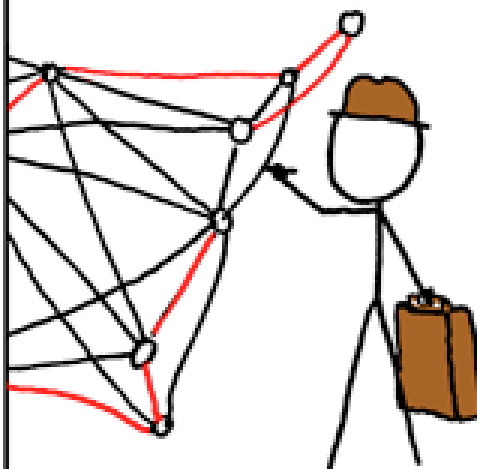
BRUTE-FORCE  
SOLUTION:

$$O(n!)$$



DYNAMIC  
PROGRAMMING  
ALGORITHMS:

$$O(n^2 2^n)$$



SELLING ON EBAY:  
 $O(1)$

STILL WORKING  
ON YOUR ROUTE?

SHUT THE  
HELL UP.

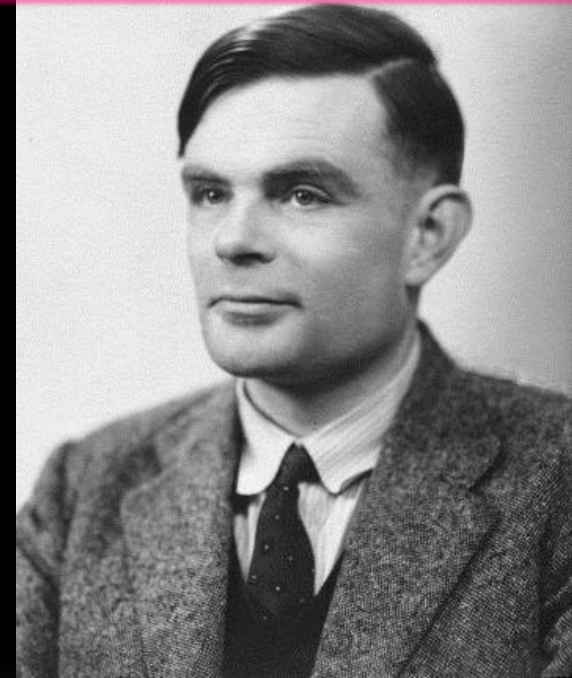




# Problems NOT solvable

- Decision problems answer YES or NO for an infinite # of inputs
  - E.g., is  $N$  prime?
  - E.g., is sentence  $S$  grammatically correct?
- An algorithm is a solution if it correctly answers YES/NO in a finite amount of time
- A problem is decidable if it has a solution

June 23, 2012 was his 100<sup>th</sup> birthday celebration!!



Alan Turing

He asked:

“Are all problems decidable?”  
(people used to believe this was true)

**Turing proved it wasn't for CS!**

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# Review: Proof by Contradiction

- Infinitely Many Primes?
- Assume the contrary, then prove that it's impossible
  - Only a finite set of primes, numbered  $p_1, p_2, \dots, p_n$
  - Consider  $q = (p_1 \cdot p_2 \cdot \dots \cdot p_n) + 1$
  - Dividing  $q$  by  $p_i$  has remainder 1
  - $q$  either prime or composite
    - If prime,  $q$  is not in the set
    - If composite, since no  $p_i$  divides  $q$ , there must be another  $p$  that does that is not in the set.
  - So there's infinitely many primes

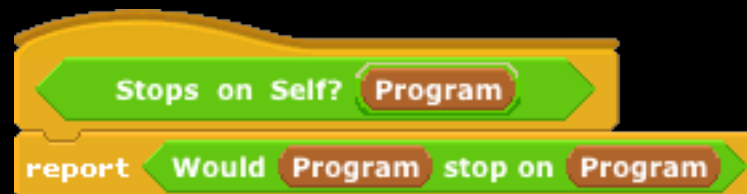


Euclid

[www.hisschemoller.com/wp-content/uploads/2011/01/euclides.jpg](http://www.hisschemoller.com/wp-content/uploads/2011/01/euclides.jpg)

# Turing's proof : The Halting Problem

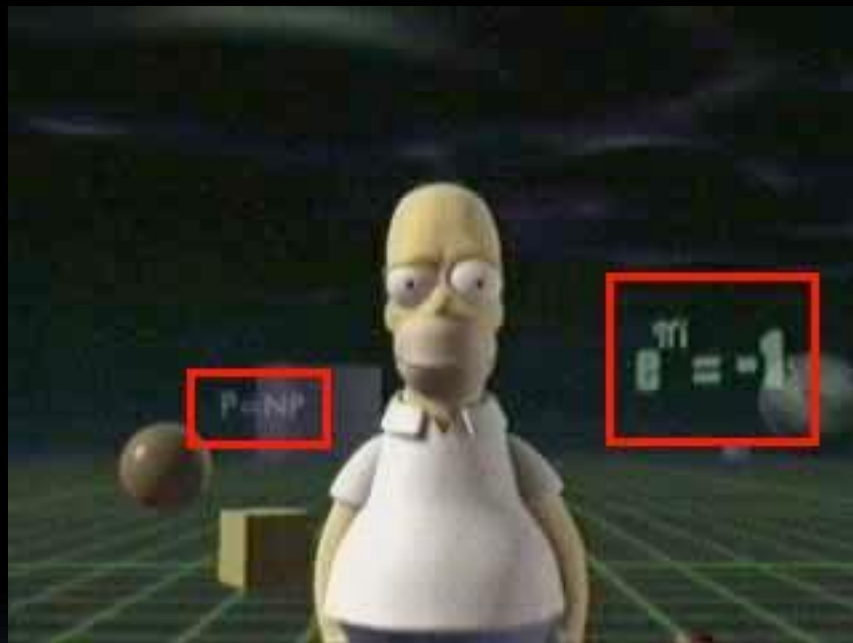
- Given a program and some input, will that program eventually stop? (or will it loop)
- Assume we could write it, then let's prove a contradiction
  - 1. write Stops on Self?
  - 2. Write Weird
  - 3. Call Weird on itself





# Conclusion

- Complexity theory  
**important part of CS**
- If given a hard problem, rather than try to solve it yourself,  
**see if others have tried similar problems**
- If you don't need an exact solution, many  
**approximation algorithms help**
- Some not solvable!



P=NP question even made its way into popular culture, here shown in the Simpsons 3D episode!