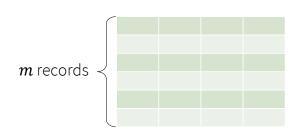
Workshop 6

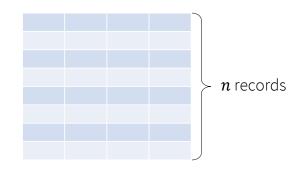
COMP20008 Elements of Data Processing

Learning outcomes

By the end of this class, you should be able to:

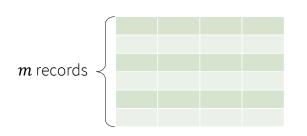
- explain how record comparisons are used to perform data linkage and duplicate detection
- explain the impact of blocking on accuracy and efficiency of data linkage
- produce advanced visualisations in Python: scatter matrix plots and parallel coordinates plots

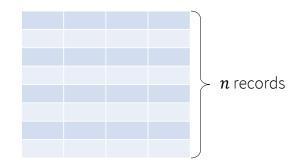




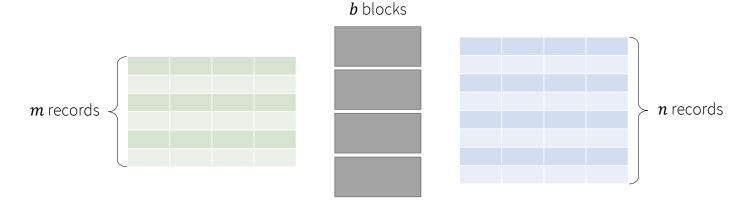
What is the maximum number of matching record pairs?

How many non-matching record pairs are there in this case?

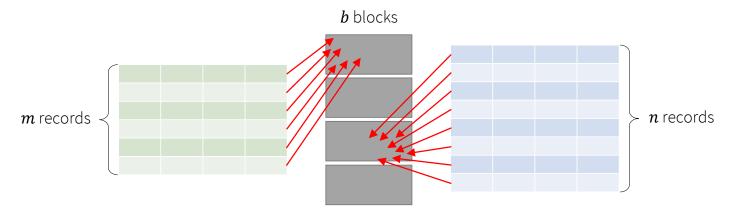




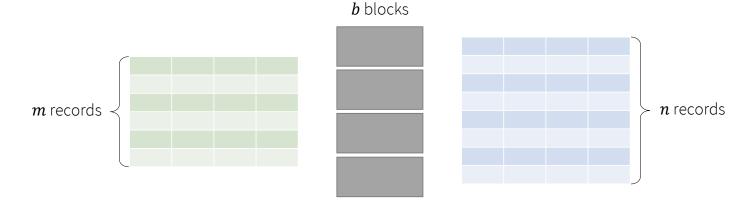
What is the maximum number of matching record pairs? mHow many non-matching record pairs are there in this case? m(n-1)



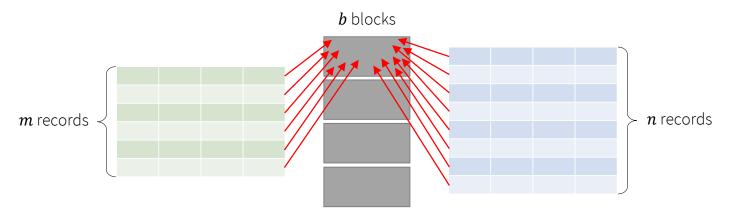
What's the *smallest* number of comparisons we could expect to make using blocking?



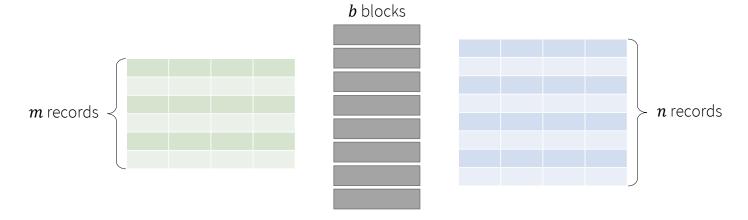
What's the *smallest* number of comparisons we could expect to make using blocking? None



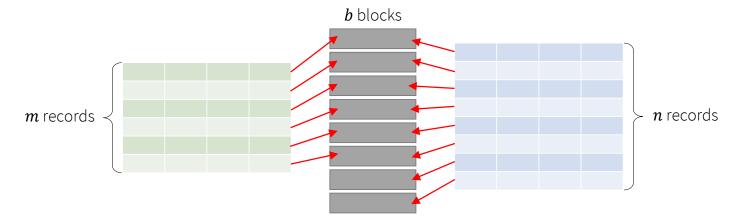
What's the *largest* number of comparisons we could expect to make using blocking?



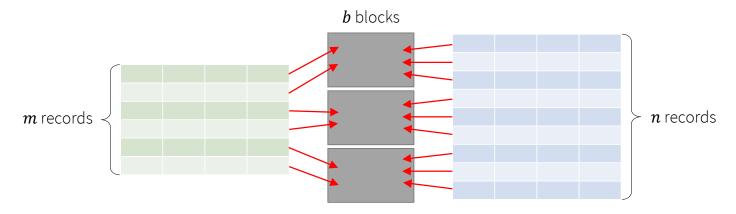
What's the *largest* number of comparisons we could expect to make using blocking? $n \times m$



Suppose there are m record matches. What's the *smallest* number of comparisons needed to return all matches, and how many blocks would there be?



Suppose there are m record matches. What's the *smallest* number of comparisons needed to return all matches, and how many blocks would there be? m comparisons; m+1 to n blocks



How many comparisons are needed if the records are allocated evenly to the blocks? $\frac{m}{b} \times \frac{1}{b} \times b$.

How many comparisons are needed if the records are allocated evenly to the blocks?

- Each block has $\frac{m}{b}$ records from the green data set and $\frac{n}{b}$ records from the blue data set
- Number of comparisons within one block: $\frac{m}{h} \times \frac{n}{h} = \frac{mn}{h^2}$
- Total across all blocks: $\frac{mn}{b^2} \times b = \frac{mn}{b}$

What's the advantage of using a large number of blocks?

efficiently reduce useless comparision

What about a small number of blocks?

precision.

How would your answers change if records are allocated to *multiple* blocks and the block allocation is *uneven*?

block size matter, extreme case block m
with all in one block.

What's the advantage of using a large number of blocks?

Fewer comparisons ⇒ more efficient

What about a small number of blocks?

More comparisons \Rightarrow typically more accurate \checkmark

How would your answers change if records are allocated to *multiple* blocks and the block allocation is *uneven*?

1/.

Conclusions don't change, except now the block sizes determine the number of comparisons

What's the advantage of allowing records to be allocated to multiple blocks?

precision

What's the advantage of allowing records to be allocated to multiple blocks?

- Restricting to a single block is often too inflexible, end up missing matches.
- Consider the movie records example from lectures.

 The release date of a movie may vary across markets. This means we will fail to detect some matches from different markets if we block on release_year. To fix this, we can allocate records to multiple blocks:

 e.g. release_year and release_year + 1.

Q2: Duplicate detection vs. linkage

Duplicate detection:

- Identify records within a single data set that refer to the same entity
- Can cast as linking a data set with itself

Bill	Joy	US	1954	1
Sophie	Wilson	UK	1957	
Stephen	Cook	US	1939	
Radia	Perlman	US	1951	
Barbara	Liskov	US	1939	
William	Joy	US	1954	

Q2: Duplicate detection vs. linkage

Data linkage:

- Identify records across a pair of data sets that refer to the same entity
- Often enforce a 1-to-1 constraint

ill	Joy	US	1954				
				Grace	Hopper	USA	
Sophie	Wilson	UK	1957	Alan		LIK	
Stephen	Cook	US	1939	Alall	Turing	UK]
'				Frances	Allen	USA	1
Radia	Perlman	US	1951	Peter	Naur	Denmark	1
Barbara	Liskov	US	1939				
Donald	Knuth	US	1938	William	Joy	USA	1
Donald	MIIUUI	03	1930				

Why might two records be incorrectly linked by a data linkage system?

Why might two records be incorrectly linked by a data linkage system?

Two reasons:

- 1. If there are not enough informative attributes, some records may match exactly when they actually refer to different entities. Linkage is hopeless in this case.
- 2. System may not be strict enough when considering fuzzy matches. For example, linking a father and son that share the same name and address, but different date of birth.

Evaluation: assess performance by comparing the predicted matches (links) from the system to human-verified truth

Category	Abbr	Description
False positive	FP	Pair of records linked that don't match
False negative	FN	Pair of records not linked that do match
True positive	TP	Pair of records linked that do match
True negative	TN	Pair of records not linked that don't match

What are the relative sizes of these categories?

What are the relative sizes of these categories?

- They should add up to the total number of pairs $m \times n$
- Often represent in a confusion matrix

		Truth		
		Match	Non-match	
Predicted	Link	TP (small)	FP (small)	
	Non-link	FN (small)	TN (large)	

When would minimizing FP be more important?

When would minimizing FN be more important?

When would minimizing FP be more important?

- Cautious about linking: want predicted matches to be correct, even if some matches are missed
- E.g. sending debt notices by linking social security and tax data.

When would minimizing FN be more important?

- Lenient about linking: want to find all matches, even if some predicted matches are incorrect
- E.g. linking child protection data across states. Want to find as many matches as possible. Can use manual review to deal with FP.

Q4: Advanced visualisation

Let's move to JupyterLab