

Design and Analysis of Algorithms I

# QuickSort

# Proof of Correctness

### Induction Review

Let P(n) = assertion parameterized by positive integers n.

For us : P(n) is "Quick Sort correctly sorts every input array of length n"

#### How to prove P(n) for all $n \ge 1$ by induction :

- 1. [base case] directly prove that P(1) holds.
- 2. [inductive step] for every n>=2, prove that:

  If P(k) holds for all k<n, then P(n) holds as well.



## Correctness of QuickSort

P(n) = "QuickSort correctly sorts every input array of length n "

<u>Claim</u>: P(n) holds for every  $n \ge 1$  [no matter how pivot is chosen]

#### **Proof by induction:**

- [base case] every input array of length 1 is already sorted.
   Quick Sort returns the input array which is correct (so P(1) holds)
- 2. [inductive step] Fix  $n \ge 2$ . Fix some input array of length n.

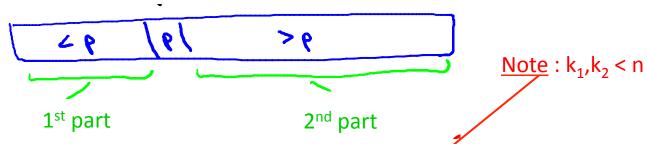
<u>Need to show</u>: if P(k) holds for all k < n, then P(n) holds as well.

INDUCTIVE STEP

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## Correctness of QuickSort (con'd)

Recall: QuickSort first partitions A around some pivot p.



Note: pivot winds up in the correct position.

Let  $k_1, k_2$  = lengths of 1<sup>st</sup>, 2<sup>nd</sup> parts of partitioned array.

Using  $P(k_1)$ ,  $P(k_2)$ 

By inductive hypothesis: 1<sup>st</sup>, 2<sup>nd</sup> parts get sorted correctly by recursive calls. So after recursive calls, entire array correctly sorted.