# **Hashtable Analysis**

In writing a hashtable with doubling, I wanted to test that my understanding of amortized complexity was correct, and that the table was functioning as intended. To do this I captured three datasets into a single file. I measured how fast the hashtable could perform insertions, retrievals and deletions of n elements, where n doubled for each measurement. The resolution of this measurement was in milliseconds, so I started with  $2^{14}$  elements and finished with  $2^{28}$ . Any fewer and the whole batch of operations took less than one millisecond – already a good sign for the performance of this library!

```
In [40]:
         import re
         from numpy import array
         def next number(file):
             return int(re.search("\d+", next(f)).group())
         data = []
         with open("data.txt") as f:
             while f:
                 try:
                      number_elements = next number(f)
                      insert ms = next number(f)
                      retrieve ms = next number(f)
                      delete ms = next number(f)
                      data.append((number elements, insert ms, retrieve ms, d
         elete ms))
                 except StopIteration:
                      break
         sizes, insert times, retrieve times, delete times = map(array, zip(
         *data))
```

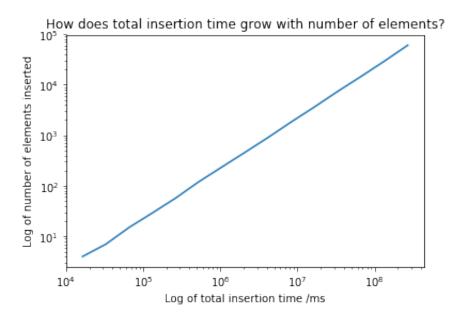
## **Linear Relationships**

In the following three graphs I demonstrate that the time taken to insert, retrieve and delete all the elements in the hashtable is a *linear* function of the number of elements in the hashtable. For example, inserting a hundred elements will take one hundred times as long as inserting only one element.

```
In [98]: from matplotlib import pyplot as plt
from numpy import log, polyfit

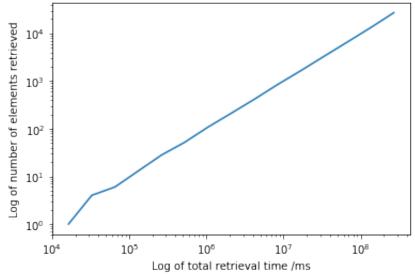
chart, *_ = plt.loglog(sizes, insert_times)
slope, _ = polyfit(log(sizes), log(insert_times), 1)
print(f"The slope of this graph is {slope:.4f} (~1), indicating a l
inear relationship.")
plt.title("How does total insertion time grow with number of elemen
ts?")
plt.xlabel("Log of total insertion time /ms")
plt.ylabel("Log of number of elements inserted")
plt.show()
```

The slope of this graph is 0.9965 ( $\sim$ 1), indicating a linear relationship.



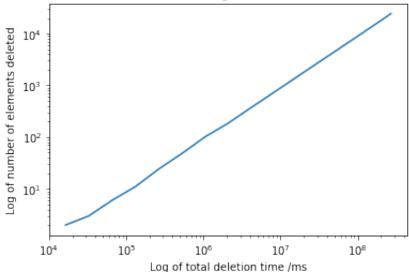
The slope of this graph is  $1.0121 \ (\sim 1)$ , indicating a linear relationship.





The slope of this graph is 0.9884 (~1), indicating a linear relationship.





## **Amortized Complexities**

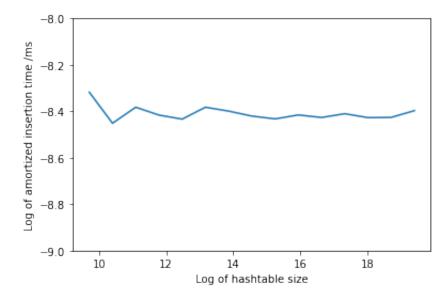
Here I demonstrate that the amortized complexity of insertion, retrieval and deletion is O(1) (or in other words, does not grow meaningfully with the number of elements in the hashtable).

```
from numpy import exp
In [99]:
         ylower, yupper = 9, 8
         amortized time = insert times / sizes
         slope, = polyfit(log(sizes), log(amortized time), 1)
         print(f"The slope of this graph is {slope:.4f} (~0), indicating a c
         onstant relationship.")
         print(f"The Y axis range here is effectively from {exp(-ylower) * 1
         0**6:.0fns to {exp(-yupper) * 10**6:.0f}ns.\n",
               "(Faster == lower on the Y axis)")
         plt.plot(log(sizes), log(amortized_time))
         plt.ylim(-ylower, -yupper)
         plt.ylabel("Log of amortized insertion time /ms")
         plt.xlabel("Log of hashtable size")
         plt.title("Log-Log graph of amortized insertion time against hashta
         ble size\n")
         plt.show()
```

The slope of this graph is -0.0035 ( $\sim 0$ ), indicating a constant relationship.

The Y axis range here is effectively from 123ns to 335ns. (Faster == lower on the Y axis)

Log-Log graph of amortized insertion time against hashtable size

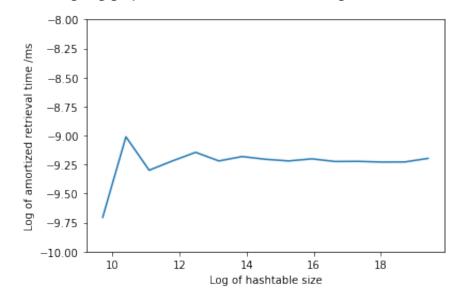


#### In [100]: from numpy import exp ylower, yupper = 10, 8 amortized time = retrieve times / sizes slope, = polyfit(log(sizes), log(amortized time), 1) print(f"The slope of this graph is {slope:.4f} (~0), indicating a c onstant relationship.") print(f"The Y axis range here is effectively from {exp(-ylower) \* 1 0\*\*6:.0fns to {exp(-yupper) \* 10\*\*6:.0f}ns.\n", "(Faster == lower on the Y axis)") plt.plot(log(sizes), log(amortized\_time)) plt.ylim(-ylower, -yupper) plt.ylabel("Log of amortized retrieval time /ms") plt.xlabel("Log of hashtable size") plt.title("Log-Log graph of amortized retrieval time against hashta ble size\n") plt.show()

The slope of this graph is 0.0121 ( $\sim$ 0), indicating a constant relationship.

The Y axis range here is effectively from 45ns to 335ns. (Faster == lower on the Y axis)

Log-Log graph of amortized retrieval time against hashtable size



#### 

plt.xlabel("Log of hashtable size")

plt.ylim(-ylower, -yupper)

le size\n")
plt.show()

The slope of this graph is -0.0116 (-0), indicating a constant relationship. The Y axis range here is effectively from 45ns to 335ns.

plt.title("Log-Log graph of amortized deletion time against hashtab

(Faster == lower on the Y axis)

Log-Log graph of amortized deletion time against hashtable size

plt.ylabel("Log of amortized deletion time /ms")

