

# Networking model checking in Spin

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# Datatypes

We modelled the data as follows:

- ▶ Number of users: 4
- ▶ Number of channels: 1 to be filled with F\_message
- ▶ F\_message - a typedef containing a sender and constraint.

# Datatypes

```
#define MAXUSERS 4  
#define MAXMESSAGES 1
```

```
F_message recMsg;
```

```
typedef F_message  
{  
    byte sender;  
    bit constraint;  
}
```

```
chan recipients[MAXUSERS] = [MAXMESSAGES] of {  
    F_message };
```

# Introduction

- ▶ `receiveMsg( byte receiver );` This process type takes messages from a given channel (`recipients[receiver]`)
- ▶ `sendMsg( byte sender, bit constraint );` Creates messages, and sends them only to the followers (or followers' followers) according to the constraint.
- ▶ `init();` Start to receive and send messages. Only 3 for every user.  
 $Numberofprocesses = 3 * MAXUSER$

## Receive messages

```
proctype receiveMsg( byte receiver )
{
  F_message recMsg;
do
  :: recipients[ receiver ] ? recMsg;
  byte fol1 , fol2 ;
  fol1 = (recMsg.sender+1)%MAXUSERS;
  fol2 = (recMsg.sender+2)%MAXUSERS;
  assert ((fol1 == receiver) || (fol2 == receiver)
  || ((fol1+2)%MAXUSERS == receiver) || ((fol1+2)%
      MAXUSERS == receiver)
  || ((fol2+2)%MAXUSERS == receiver) || ((fol2+2)%
      MAXUSERS == receiver)
  );
od;
}
```

## Sending messages

```
proctype sendMsg ( byte sender; byte constraint )
{
  F_message newMsg;
  byte fol1 , fol2 , fol11 , fol12 , fol21 , fol22;

  /*
   * Locally save followers id's
   */

  d_step{
    newMsg.sender = sender;
    newMsg.constraint = constraint;

    fol1 = ((sender+1)%MAXUSERS);
    fol2 = ((sender+2)%MAXUSERS);
    fol11 = ((fol1+1)%MAXUSERS);
    fol12 = ((fol1+2)%MAXUSERS);
    fol21 = ((fol2+1)%MAXUSERS);
    fol22 = ((fol2+2)%MAXUSERS);
```

## Sending messages (2)

```
}  
do  
  :: if  
  :: (sender > 0) -> atomic{ /*Send to Fo*/  
  if  
  :: fol1 != sender -> recipients[fol1] ! newMsg  
  :: else skip  
  fi;  
  if  
  :: fol2 != sender -> recipients[fol2] ! newMsg  
  :: else skip  
  fi;  
}  
  if
```



## Sending messages (3)

```
:: (constraint == 1) -> atomic{ /*Send to Fo2*/  
  if  
  :: sender != fol11 -> recipients[fol11] ! newMsg  
  :: else skip  
  fi;  
  if  
  :: sender != fol12 -> recipients[fol12] ! newMsg  
  :: else skip  
  fi;  
if  
:: sender != fol21 -> recipients[fol21] ! newMsg  
:: else skip  
fi;
```

## Sending messages (4)

```
if
:: sender != fol22 -> recipients[fol22] ! newMsg
:: else skip
fi
}
:: else skip
fi
:: else skip
fi
od;
}
```

# Initializing

```
init
{
/*
 * Start all receivers and senders, for every
 *   channel (maxusers)
 */

byte i = 0;

atomic{
do
:: (i < MAXUSERS) -> run receiveMsg(i); run sendMsg
    (i,0); run sendMsg (i,1); i++;
:: else -> break;
od;
}
}
```

# Asserts

- ▶ Privacy constraint is never broken
- ▶ No deadlocks
- ▶ Message eventually reaches the receiver

# Work in progress

- ▶ How to check whether all messages eventually reach the right user.
  - ▶ Implicitly done, it would deadlock otherwise, but is this enough?
  - ▶ Checking on the right follower relation in every message.
- ▶ How to express assertions in LTL.
  - ▶ Using never claim