## **NETAJI SUBHAS UNIVERSITY OF TECHNOLOGY**



# Database Management Systems ICCSC09

# Hospital Management System (HMS)

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## **Overview**

The Hospital Management System (HMS) is a comprehensive solution designed to streamline and digitalize various processes within a hospital environment. This project aims to enhance the efficiency, accuracy, and accessibility of information across different departments.

## **Problem Statement**

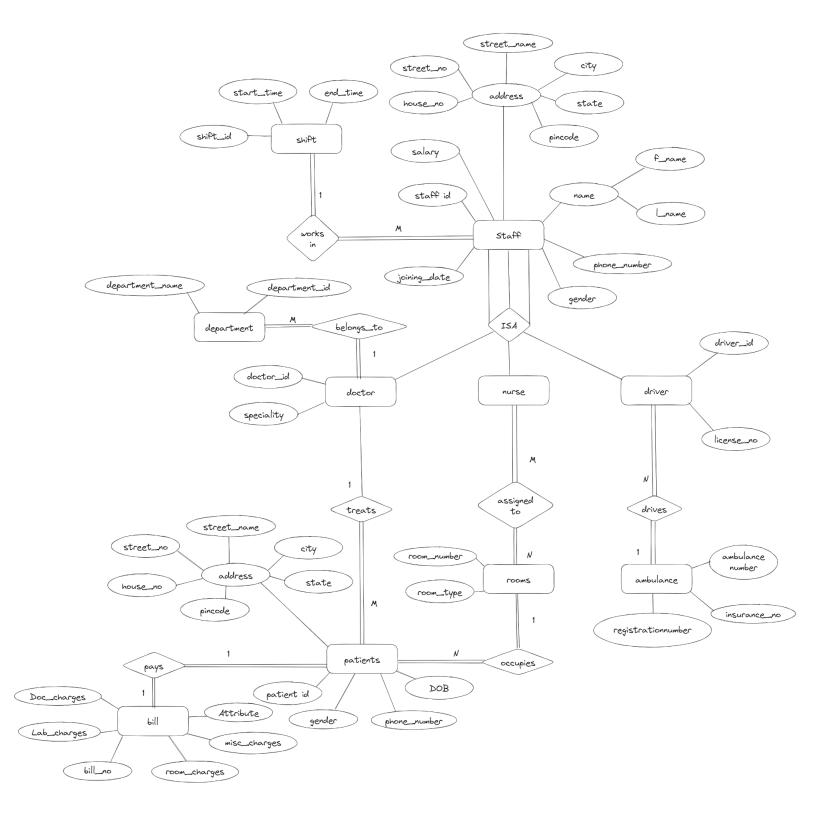
Healthcare institutions often face challenges in managing patient information, appointments, billing, and other administrative tasks. The Hospital Management System is developed to address these challenges by introducing a centralized digital platform.

The current hospital management system is plagued by inefficiencies and limitations that hinder the delivery of optimal healthcare services. These issues range from administrative bottlenecks to patient care challenges, and they impede the overall effectiveness of the healthcare institution. As such, there is an urgent need to develop and implement an advanced hospital management system.

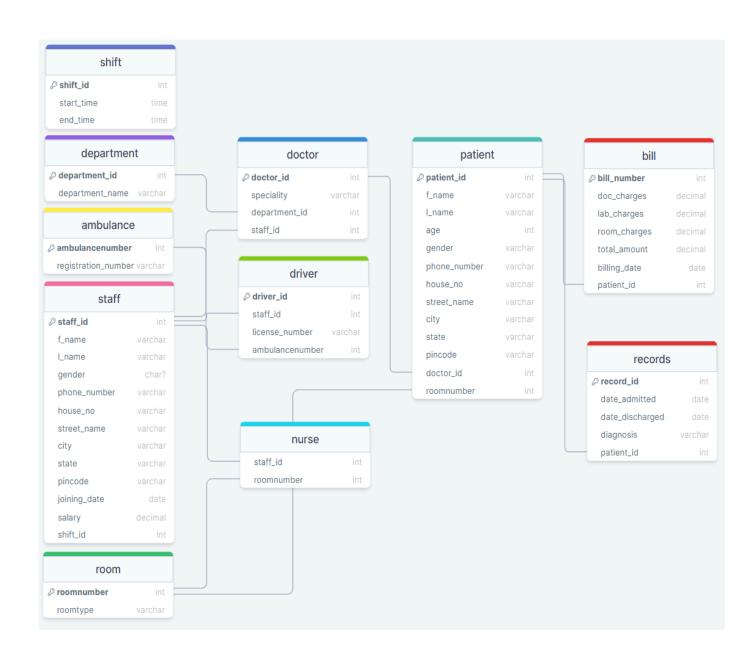
## **Purpose of Project**

- Automate hospital processes for improved efficiency.
- Enhance patient care by facilitating quick access to medical records.
- Improve resource management, including staff scheduling and inventory control.
- Ensure accurate billing and financial tracking.

## **ER Diagram**

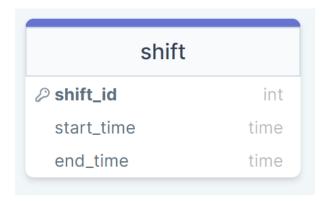


## **Relational Schema**



## **Function Dependencies and Third Normal Form (3NF)**

#### 1. Shift

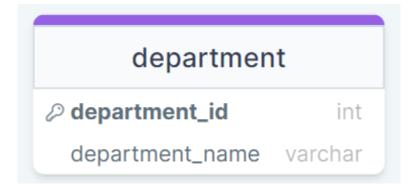


## • Functional Dependencies

- shift\_id → start\_time
- shift\_id → end\_time

- 1st Normal Form (1NF): The table already meets 1NF since there are no repeating groups, and each cell contains atomic values.
- 2nd Normal Form (2NF): Since there's only one candidate key (shift\_id), and both start\_time and end\_time are fully functionally dependent on the entire candidate key, we don't have partial dependencies. Thus, the table is already in 2NF.
- 3rd Normal Form (3NF): In this step, we check for transitive dependencies.
   Since there's only one non-prime attribute, start\_time and end\_time are both attributes that directly depend on the primary key. There are no transitive dependencies, so the table is already in 3NF.

## 2. Department

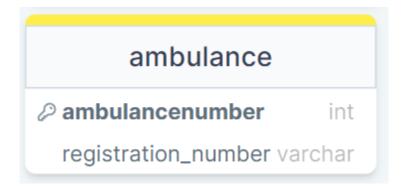


## • Functional Dependencies

department id → department name

- 1st Normal Form (1NF): The table already meets 1NF since there are no repeating groups, and each cell contains atomic values.
- 2nd Normal Form (2NF): Since there's only one candidate key
   (department\_id), and department\_name is fully functionally dependent on
   the entire candidate key, we don't have partial dependencies. Thus, the
   table is already in 2NF.
- 3rd Normal Form (3NF): In this step, we need to ensure that there are no transitive dependencies. Since there's only one non-prime attribute (department\_name), and it is directly dependent on the primary key, there are no transitive dependencies.

#### 3. Ambulance



## • Functional Dependencies

1. ambulance\_number → registration\_number

- First Normal Form (1NF): The table seems to be in 1NF because each attribute contains atomic values, and there are no repeating groups.
- Second Normal Form (2NF): Since there is only one non-prime attribute (registration\_number), there are no partial dependencies. Thus, the table already satisfies the requirements of 2NF.
- Third Normal Form (3NF): To achieve 3NF, we need to ensure that there
  are no transitive dependencies. In this schema, there is only one
  non-prime attribute, so there are no transitive dependencies to address.

## 4. Staff



## • Functional Dependencies

- staff id → speciality: Each doctor's staff id uniquely determines their specialty.
- staff\_id → department\_id:Each doctor's staff\_id uniquely determines department.
- staff\_id → license\_number: Each driver's staff\_id uniquely determines their license number.
- staff\_id → ambulance\_number: Similarly, each driver's staff\_id uniquely determines the ambulance they are assigned to.
- staff\_id → room\_number: Each nurse's staff\_id uniquely determines the room number they are assigned to.

## Normal Form

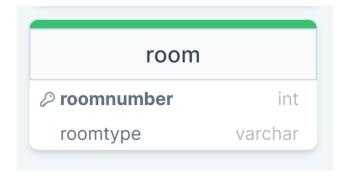
- First Normal Form (1NF): The table already seems to be in 1NF because each attribute contains atomic values, and there are no repeating groups.
- Second Normal Form (2NF): The table is already in 2NF because there are no partial dependencies. Each non-prime attribute depends on the entire primary key.
- Third Normal Form (3NF):To achieve 3NF, we need to ensure there are no transitive dependencies. There is a transitive dependency between staff and shift\_id. Since the shift details might not be directly related to the staff member, we should move this attribute to another table.

#### Normalized Schema:

```
staff_info:
staff(staff_id (primary key), f_name, l_name, gender, phone_number, house_no,
street_name, city, state, pincode, joining_date, salary)
shift_assignment:
shift_assignment(staff_id (foreign key), shift_id (foreign key), primary key(staff_id,
shift_id))
```

- 1. doctor(doctor\_id(primary key), speciality, department\_id, staff\_id(foreign key)):
  - No changes required as it's a separate entity.
- 2. driver(driver\_id (primary key), staff\_id(foreign key), license\_number, ambulance\_number):
  - No changes required as it's a separate entity.
- 3. nurse(staff id (foreign key), room number):
  - No changes required as it's a separate entity.

## 5. Room



## • Functional Dependencies

- roomnumber → roomtype
- roomnumber → staff id
- roomnumber → patient\_id

- First Normal Form (1NF): The table appears to be in 1NF because each attribute contains atomic values, and there are no repeating groups.
- Second Normal Form (2NF): The table is already in 2NF because there are no partial dependencies. Each non-prime attribute depends on the entire primary key.
- Third Normal Form (3NF): To achieve 3NF, we need to ensure there are no transitive dependencies. Since roomnumber is a foreign key in both the "nurse" and "patient" tables, and it directly relates to the room type, there are no transitive dependencies to address.

## 6. Doctor

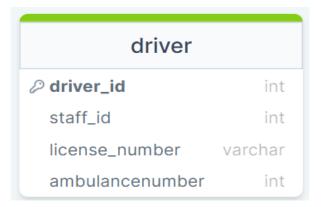


## • Functional Dependencies

- doctor\_id → speciality
- doctor\_id → staff\_id

- First Normal Form (1NF): The table appears to be in 1NF because each attribute contains atomic values, and there are no repeating groups.
- Second Normal Form (2NF): The table is already in 2NF because there are no partial dependencies. Each non-prime attribute depends on the entire primary key.
- Third Normal Form (3NF): To achieve 3NF, we need to ensure there are no transitive dependencies. Since doctor\_id is a foreign key in the "patient" table and it directly relates to the doctor's specialty and staff ID, there are no transitive dependencies to address.

## 7. Driver

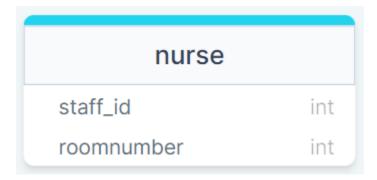


## • Functional Dependencies

- driver\_id → staff\_id
- driver\_id → license\_number
- $\bullet \quad \text{driver\_id} \rightarrow \text{ambulancenumber}$

- First Normal Form (1NF): The table already appears to be in 1NF because each attribute contains atomic values, and there are no repeating groups.
- Second Normal Form (2NF): The table is already in 2NF because there are no partial dependencies. Each non-prime attribute depends on the entire primary key.
- Third Normal Form (3NF): To achieve 3NF, we need to ensure there are no transitive dependencies. There are no transitive dependencies in the original schema.

## 8. Driver



## • Functional Dependencies

• (staff\_id, roomnumber)  $\rightarrow$  None

- First Normal Form (1NF): The table is already in 1NF as there are no repeating groups, and all attributes are atomic.
- Second Normal Form (2NF): Since there's no composite primary key, and both staff\_id and roomnumber are attributes directly dependent on the entire primary key, the table is already in 2NF.
- Third Normal Form (3NF): There are no attributes that depend on any non-primary key attributes, so there are no transitive dependencies. Therefore, the table is already in 3NF.

#### 9. Patient

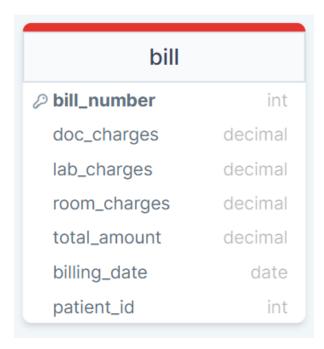


## • Functional Dependencies

 patient\_id → f\_name, l\_name, age, gender, phone\_number, house\_no, street\_name, city, state, pincode, doctor\_id, roomnumber

- First Normal Form (1NF): The table is already in 1NF as there are no repeating groups, and all attributes are atomic.
- Second Normal Form (2NF): Since there's no composite primary key, and all non-prime attributes (f\_name, l\_name, age, gender, phone\_number, house\_no, street\_name, city, state, pincode, doctor\_id, roomnumber) are fully functionally dependent on the entire primary key (patient\_id), the table is already in 2NF.
- Third Normal Form (3NF): There are no attributes that depend on any non-primary key attributes, so there are no transitive dependencies. Therefore, the table is already in 3NF.

#### 10. Bill

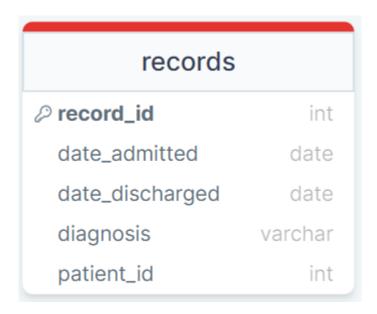


## • Functional Dependencies

 bill\_number → doc\_charges, lab\_charges, room\_charges, total\_amount, billing\_date, patient\_id

- First Normal Form (1NF): Already satisfied.
- Second Normal Form (2NF): Since there's only one candidate key
   (bill\_number), and all other attributes (doc\_charges, lab\_charges,
   room\_charges, total\_amount, billing\_date, patient\_id) are fully
   functionally dependent on the entire candidate key, we don't have partial
   dependencies. Thus, the table is already in 2NF.
- Third Normal Form (3NF): In this step, we need to ensure that there are no transitive dependencies. Since there are no attributes that depend on any non-primary key attributes, there are no transitive dependencies.

## 11. Records



## • Functional Dependencies

- record id → date admitted, date discharged, diagnosis, patient id
- patient id → None

- First Normal Form (1NF): Already satisfied.
- Second Normal Form (2NF): Since there's only one candidate key
   (record\_id), and all other attributes (date\_admitted, date\_discharged,
   diagnosis, patient\_id) are fully functionally dependent on the entire
   candidate key, we don't have partial dependencies. Thus, the table is
   already in 2NF.
- Third Normal Form (3NF): In this step, we need to ensure that there are no transitive dependencies. Since there are no attributes that depend on any non-primary key attributes, there are no transitive dependencies.