# CMPT 383 Comparative Programming Languages

## Programming Assignment 3

This assignment is due by 11:59pm PT on Tuesday Mar 25, 2025. Please submit it to Canvas.

## Requirements:

- This assignment must be your own work. No collaboration is permitted.
- You can learn the code on slides and start from it.
- You can only use library functions from the following modules: Prelude, System.IO, System.Environment, and Data.Map.Strict. Detailed information can be found on https://hoogle.haskell.org

## Late policy:

Suppose you can get n (out of 100) points based on your code and report

- If you submit before the deadline, you can get all n points.
- If you submit between 11:59pm PT Mar 25 and 11:59pm PT Mar 26, you get n-10 points.
- If you submit between 11:59pm PT Mar 26 and 11:59pm PT Mar 27, you get n-20 points.
- If you submit after 11:59pm PT Mar 27, you get 0 points.

(100 points) Consider the following fragment of the FUN language

$$e ::= c | b | x | '(' e ')'$$
 $| e '+' e | e '-' e | e '==' e$ 
 $| 'if' e 'then' e 'else' e$ 
 $| 'lambda' x ':' t '.' e$ 
 $| 'app' e e$ 
 $| 'let' x ':' t '=' e 'in' e$ 
 $t ::= 'Int' | 'Bool' | '(' t ')'$ 
 $| t '->' t$ 
 $c \in Int b \in Bool x \in Ident$ 

Here, e is the start symbol. c stands for an integer constant, b stands for a boolean constant, and x stands for an identifier (variable). The  $\rightarrow$  operator is **right-associative**.

The typing rules of this language are defined as follows:

$$\frac{\operatorname{Int}\,c}{\Gamma\vdash c:\operatorname{Int}}\,(\operatorname{T-Int}) \qquad \frac{\operatorname{Bool}\,b}{\Gamma\vdash b:\operatorname{Bool}}\,(\operatorname{T-Bool}) \qquad \frac{\Gamma\vdash e_1:\operatorname{Int}}{\Gamma\vdash e_2:\operatorname{Int}}\,(\operatorname{T-Plus}) \qquad \frac{\Gamma\vdash e_1:\operatorname{Int}}{\Gamma\vdash e_2:\operatorname{Int}}\,(\operatorname{T-Minus})$$

$$\frac{T\in\{\operatorname{Int},\operatorname{Bool}\}}{\Gamma\vdash e_1:T\quad\Gamma\vdash e_2:T} \qquad \Gamma\vdash e_2:T\quad\Gamma\vdash e_3:T} \qquad \operatorname{Ident}\,x \qquad \Gamma\vdash e_1:T\quad\Gamma\vdash e_2:T\quad\Gamma\vdash e_3:T}{\Gamma\vdash e_1==e_2:\operatorname{Bool}}\,(\operatorname{T-Eq}) \qquad \frac{\Gamma\vdash e_1:\operatorname{Tohender}}{\Gamma\vdash \operatorname{if}\,e_1\operatorname{then}\,e_2\operatorname{else}\,e_3:T}\,(\operatorname{T-ITE}) \qquad \frac{\Gamma(x)=T}{\Gamma\vdash x:T}\,(\operatorname{T-Ident})$$

$$\frac{\Gamma[x\lhd T_1]\vdash e:T_2}{\Gamma\vdash \operatorname{lambda}\,x:T_1.e:T_1\to T_2}\,(\operatorname{T-Abs}) \qquad \frac{\Gamma\vdash e_1:T_1\to T_2}{\Gamma\vdash e_2:T_1}\,(\operatorname{T-App}) \qquad \frac{\Gamma[x\lhd T_1]\vdash e_1:T_1}{\Gamma\vdash \operatorname{lat}\,x:T_1\vdash e_2:T_2}\,(\operatorname{T-Let})$$

In this assignment, you need to write a type checker in Haskell to type check expressions in the FUN language. Specifically, given an expression e, the type checker computes the type of e if it is well-typed. If e is ill-typed, the type checker should output Type Error.

To avoid the complication of parsing, you can assume the type is represented as a value of the following Type data type in Haskell

where TInt represents the Int type, TBool represents the Bool type, and TArr represents the function type. For example, type Int -> Bool -> Int in FUN is represented as TArr TInt (TArr TBool TInt).

Variable IDs are assumed to be strings:

```
type VarId = String
```

An expression of the FUN language is represented as a value of the following Expr type

As indicated by the names of data constructors, CInt denotes an integer constant, CBool denotes a boolean constant, Var denotes a variable (identifier), Plus denotes the + operator, Minus denotes the - operator, Equal denotes the == operator, ITE denotes the if-then-else expression, Abs denotes the function abstraction (lambda), App denotes the function application, and LetIn denotes the let-in expression. There is no data constructor for parenthesized expressions. For example,

- x1 is represented as Var "x1"
- 1 + 2 is represented as Plus (CInt 1) (CInt 2)
- if True then 1 else 2 is represented as ITE (CBool True) (CInt 1) (CInt 2)
- lambda x:Int.x is represented as Abs "x" TInt (Var "x")
- app (lambda x:Int.x) 1 is represented as App (Abs "x" TInt (Var "x")) (CInt 1)
- let x:Int = 1 in x is represented as LetIn "x" TInt (CInt 1) (Var "x")

Note that you need to use exactly the same definition of Type, VarId, Expr, and their deriving clauses as written in this document. Otherwise, you will lose points because potential grading scripts may not work as expected.

## **Detailed Steps**

1. Use Map from Data.Map.Strict to define the Env type for the typing environment, i.e., finish the following declaration

```
type Env = ...
```

2. Write an auxiliary function typingArith :: Maybe Type -> Maybe Type -> Maybe Type that

- returns Just TInt if both arguments are Just TInt
- returns Nothing otherwise
- 3. Write an auxiliary function typingEq :: Maybe Type -> Maybe Type -> Maybe Type that
  - returns Just TBool if both arguments are Just TInt
  - returns Just TBool if both argumnets are Just TBool
  - returns Nothing otherwise
- 4. Write a function typing :: Env -> Expr -> Maybe Type that takes a typing environment and a FUN expression as input and produces as output the type of that expression based on the typing rules. If there is a type error, it returns Nothing. Note that you can use the auxiliary function typingArith and typingEq in this function.
- 5. Write a **simple** function **readExpr** :: String -> Expr that can read a value of Expr type from the corresponding string, such as "CInt 1".
- 6. Write a function typeCheck :: Expr -> String that takes an expression as input and produces a string as output representing the type checking result. Specifically,
  - $\bullet$  If the expression is well-typed and a value v of type Type is obtained, generate the output using show v.
  - If the expression is ill-typed, output string "Type Error".
- 7. Write a main to handle IO and put everything together.

The program must be in a form that GHC can compile. It needs to take one command-line argument denoting the path to the expression file. Each line of the file contains a string representing a FUN expression, and the program needs to print the result of typeCheck on each expression to the console.

### Sample Input and Output

Suppose we have a file called exprs.txt that contains the following six lines:

```
Var "x1"
Plus (CInt 1) (CInt 2)
ITE (CBool True) (CInt 1) (CInt 2)
Abs "x" TInt (Var "x")
App (Abs "x" TInt (Var "x")) (CInt 1)
LetIn "x" TInt (CInt 1) (Var "x")
```

After compiling, we can run the executable and get

```
$ ./P3_SFUID exprs.txt
Type Error
TInt
TInt
TArr TInt TInt
TInt
TInt
```

#### Deliverable

A zip file called P3\_SFUID.zip that contains at least the followings:

- A file called P3\_SFUID.hs that contains the source code of your Haskell program. You can have multiple source files if you want, but you need to make sure ghc P3\_SFUID.hs can compile.
- A report called P3\_SFUID.pdf that explains the design choices, features, issues (if any), and anything else that you want to explain about your program.