

assignment2

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1 problem1

(a) if $R1 \cap R2 = \text{empty}$: that is exactly an equivalence relation

else : R : assume $a \in S$

then for all possible a $(a, a) \in R_1$ (reflexive)

and $(a, a) \in R_2$ (reflexive)

thus $(a, a) \in (R1 \cap R2)$

thus $R1 \cap R2$ is reflexive

S : assume $x, y \in S$ that $(x, y) \in (R1 \cap R2)$

it means $(x, y) \in R1$ and $(x, y) \in R2$

if $(x, y) \in R1$ then $(y, x) \in R1$ (symmetric)

if $(x, y) \in R2$ then $(y, x) \in R2$ (symmetric)

thus $(y, x) \in R1$ and $(y, x) \in R2 \rightarrow (y, x) \in (R1 \cap R2)$

thus $R1 \cap R2$ is symmetric

T : assume x, y, z that (x, y) and $(y, z) \in (R1 \cap R2)$

it means there exists $y \in S$ that makes

$\rightarrow (x, y)$ and $(y, z) \in R1$

$\rightarrow (x, y)$ and $(y, z) \in R2$

if (x, y) and $(y, z) \in R1$ then $(x, z) \in R1$ (transitive)

if (x, y) and $(y, z) \in R2$ then $(x, z) \in R2$ (transitive)

thus there exists $y \in S$ that makes

$(x, z) \in R1$, and $(x, z) \in R2 \rightarrow (x, z) \in R1 \cap R2$

thus $R1 \cap R2$ is transitive

thus $R1 \cap R2$ is $R, S, T \rightarrow$ it is a Equivalence relation

(b) assume we get $(x, y) \in R1 \cap R2$

$\rightarrow (x, y) \in R1$ and $(x, y) \in R2$ (definition)

by the definition of equivalence class

$\rightarrow (x, y) \in [x]_1$ and $(x, y) \in [x]_2$

that means $(x, y) \in [x]_1 \cap [x]_2$ (definition)

thus $[x]_1 \cap [x]_2$ is a equivalence class of x under $R1 \cap R2$

$[x] = [x]_1 \cap [x]_2$

c) it is false, for example assume $S = \{1, 2, 3, 4, 5\}$

R_1 as the relation $= (\text{mod}3)$; R_2 as the relation $= (\text{mod}4)$

thus $R_1 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 4), (4, 1), (2, 5), (5, 2)\}$

$R_2 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 5), (5, 1)\}$

elements in $R1 \cup R2$ is
 $\{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 4), (4, 1), (2, 5), (5, 2), (1, 5), (5, 1)\}$
 if it satisfy Transitivity $\rightarrow (2, 5), (5, 1) \in R1 \cup R2 \rightarrow (2, 1) \in R1 \cup R2$
 which is not true
 thus $R1 \cup R2$ is not an equivalence relation

2 problem2

(a) assume $R_1; R_2$ from the definition we know
 $R_1; R_2 = \{(a, c) : \text{there is } a, b \in S \text{ such that } (a, b) \in R1 \text{ and } (b, c) \in R2\}$ (1)
 if $R1 \subseteq S \times S$, $R1$ is reflexive we can ensure that :
 $(a, a) \in R1$ thus $a = b$ is acceptable in (1)
 because $R2$ is reflexive, $R2 \subseteq S \times S$
 $(a, a) \in R2$
 when $a = b$, there can be $a = c$ that satisfy $(b, c) = (a, c) = (a, a) \in R2$
 thus $(a, a) \in R1; R2$
 (b) it is false, for example assume $S = \{1, 2, 3, 4, 5\}$
 R_1 as the relation $= (\text{mod}3)$; R_2 as the relation $= (\text{mod}4)$
 thus $R_1 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 4), (4, 1), (2, 5), (5, 2)\}$
 $R_2 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 5), (5, 1)\}$ $R1, R2$ are symmetric
 by definition we know $(4, 5) \in R1; R2$ since $(4, 1) \in R1, (1, 5) \in R2$
 but $(5, 4) \notin R1; R2$
 (c) it is false, for example assume $S = \{1, 2, 3, 4, 5\}$
 R_1 as the relation $= (\text{mod}3)$; R_2 as the relation $= (\text{mod}4)$
 thus $R_1 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 4), (4, 1), (2, 5), (5, 2)\}$
 $R_2 : \{(1, 1), (2, 2), (3, 3), (4, 4), (5, 5), (1, 5), (5, 1)\}$ $R1, R2$ are transitive
 by definition we know $(2, 1) \in R1; R2$ since $(2, 5) \in R1, (5, 1) \in R2$
 $(1, 4) \in R1; R2$ since $(1, 4) \in R1, (4, 4) \in R2$
 but $(2, 4) \notin R1; R2$

3 problem3

(a) Let $P(x)$ be the proposition that $R^i = R^j$
 we need to show
 1) basic case ($j = i$) : $R^i = R^i$ obviously it is true
 so $P(i)$ holds
 2) then we assume $P(j)$ holds
 we need to show $P(j + 1)$ holds
 $P(j)$ holds \rightarrow there is i that $R^i = R^j$
 $R^{j+1} = R^j \cup (R; R^j)$ (definition)
 $= R^i \cup (R; R^i)$ ($R^i = R^j$)

$= R^{i+1}$ (definition)
 $= R^i$ ($R^i = R^{i+1}$)
 thus $P(j+1)$ holds
 thus $P(j)$ implies $P(j+1)$
 Therefore, by induction, $P(j)$ holds for all $j \geq i$

(b)(1) when $j \geq i$ from (a) we know $R^j = R^i \subseteq R^i$
 (2) when $j < i$
 Let $P(x)$ be the proposition that $R^j \subseteq R^{j+x}$ $x \in N$
 1) basic case ($k=0$) :
 $R^j \subseteq R^j$ so $P(0)$ holds
 2) we assume $P(m)$ holds then we try to proof $P(m+1)$ holds
 when $P(m)$ holds $R^j \subseteq R^{j+m}$
 because $R^{j+m+1} = R^{j+m} \cup (R; R^{j+m}) \rightarrow R^{j+m} \subseteq R^{j+m+1}$
 thus $R^j \subseteq R^{j+m} \subseteq R^{j+m+1}$ thus $P(m+1)$ holds
 we can replace $j+x$ by i since there must be a x that let $j+x=i$
 thus $R^j \subseteq R^{j+x} = R^i$ always holds when $j < i$
 thus it is true

(c) we need to show $R^{k^2} = R^{k^2+1}$
 from (b) we already know that $R^j \subseteq R^{j+x}$ $x \in N$ is true
 let $x=1, j=k^2$ $R^{k^2} \subseteq R^{k^2+1}$
 then we need to show $R^{k^2} \supseteq R^{k^2+1}$
 assume there is $(a, c) \in R^{k^2+1}$, and $(a, c) \notin R^{k^2}$
 because $(a, c) \in R^{k^2+1} = R^{k^2} \cup (R; R^{k^2})$
 $= R^{k^2} \cup \{(a, c) : \text{there is } a, b \in S \text{ such that } (a, b) \in R \text{ and } (b, c) \in R^{k^2}\}$
 thus (a, c) can only $\in \{(a, c) : \text{there is } a, b \in S \text{ such that } (a, b) \in R \text{ and } (b, c) \in R^{k^2}\}$
 thus there exist $(a, b_{k^2}) \in R$ and $(b_{k^2}, c) \in R^{k^2}$
 $(b_{k^2}, c) \in R^{k^2} = R^{k^2-1} \cup (R; R^{k^2-1})$
 (b_{k^2}, c) can not belong to R^{k^2-1} because
 if $(b_{k^2}, c) \in R^{k^2-1}$, $(a, b_{k^2}) \in R$ then $(a, c) \in R^{k^2}$ which violate our assumption
 thus $(b_{k^2}, c) \in (R; R^{k^2-1})$
 \rightarrow thus there exist $(b_{k^2}, b_{k^2-1}) \in R$ and $(b_{k^2-1}, c) \in R^{k^2-1}$
 (b_{k^2-1}, c) can not belong to R^{k^2-1} because
 if $(b_{k^2-1}, c) \in R^{k^2-1}$, $(b_{k^2}, b_{k^2-1}) \in R$ then $(b_{k^2}, c) \in R^{k^2}$ which violate our presupposition
 similarly, we can extend this \rightarrow for $0 \leq n \leq k^2+1$, R^n would get one new element than R^{n-1}
 however whole size of it could only be $|S||S| = k^2$ while need k^2+1 different elements
 there must be some element violate the rule which make $(a, b) \in R^{k^2}$
 or $k=1$ ($R^1 = R^2$ is true with only one element)
 Therefore, $(a, b) \in R^{k^2}$, which means for every element in R^{k^2+1} , it is also in R^{k^2} ,
 so $R^{k^2+1} = R^{k^2}$
 (d) basic case : $R^0; R^m = R^{0+m}$
 $R^{0+m} = R^m$
 $= I; R_m$ (using ass1 - p8 - (b))

$$= R^0; R_m \quad (R^0 = I)$$

so $P(0)$ holds

we assume $P(n)$ holds then need to show $P(n+1)$ holds

if $P(n)$ holds then $R^{n+m} = R^n; R^m$

$$R^{(n+1)+m} = R^{n+m+1} = R^{n+m} \cup (R; R^{m+n}) \quad (\text{definition})$$

$$= (R^n; R^m) \cup (R; (R^n; R^m)) \quad (\text{it holds})$$

$$= (R^n; R^m) \cup ((R; R^n); R^m) \quad \text{using ass1 - p8 - (a)}$$

$$= (R^n \cup (R; R^n)); R^m \quad \text{using ass1 - p8 - (c)}$$

$$= R^{n+1}; R^m \quad (\text{definition})$$

thus $P(n+1)$ holds

Therefore $P(n)$ holds for all $n \in N$

(e) assume there is $(a, b) \in R^{k^2}, (b, c) \in R^{k^2}$

$$\rightarrow \text{from (c)} : R^{k^2} = R^{k^2+1}$$

$$\rightarrow \text{from (d)} R^{k^2}; R^{k^2} = R^{2k^2} = \{(a, c) : \text{there is } a, b \in S \text{ such that } (a, b) \in R^{k^2} \text{ and } (b, c) \in R^{k^2}\}$$

$$\text{thus } (a, c) \in R^{2k^2}$$

because $R^{k^2} = R^{k^2+1}$ so K^2 satisfy our definition of i

from (a) we know $R^{2k^2} = R^{k^2}$ since $2k^2 > k^2$

$$\text{thus } (a, c) \in R^{k^2}$$

thus R^{k^2} is transitive

(f) we need to make $(R \cup R^{\leftarrow})^{k^2}$ reflexive, symmetric and transitive

R : from (b) and definition we know for all $x \in S$

$$(x, x) \in (R \cup R^{\leftarrow})^0 = I \subseteq (R \cup R^{\leftarrow})^{k^2} \text{ thus it is reflexive}$$

T : we can consider $(R \cup R^{\leftarrow})$ as a big R from (e) we know it is transitive

S : For all $x, y \in S$: If $(x, y) \in R$ then $(y, x) \in R$

(1) case1: $(R \cup R^{\leftarrow})^0$ is true since $(R \cup R^{\leftarrow})^0 = I$ and I is symmetric

(2) case 2 $(R \cup R^{\leftarrow})^1$ is true since $(a, b) \in R, (b, a) \in R^{\leftarrow}$ and vice versa

(3) assume there exist $(a, b) \in (R \cup R^{\leftarrow})^{k^2}, (b, a) \notin (R \cup R^{\leftarrow})^{k^2} a \neq b$

$$\text{from (b) we know } (R \cup R^{\leftarrow})^1 \subseteq (R \cup R^{\leftarrow})^{k^2}$$

thus $(a, b) \notin (R \cup R^{\leftarrow})^{k^2}$ which is contradictory to our precondition

thus there do not exist such (a, b)

$(a, b) \in (R \cup R^{\leftarrow})^{k^2}, (b, a) \in (R \cup R^{\leftarrow})^{k^2}$ would always be synchronous

combine(1), (2) it is symmetric

because $(R \cup R^{\leftarrow})^{k^2}$ is $R, T, S \rightarrow$ it is an equivalence relation

(g) we know $R \subseteq (R \cup R^{\leftarrow})^1 \subseteq (R \cup R^{\leftarrow})^{k^2}$ from (c) induction

thus for all $a \in R, a \in (R \cup R^{\leftarrow})^{k^2}$

for (f) we know $(R \cup R^{\leftarrow})^{k^2}$ is an equivalence relation

the definition of minimum: $x \prec y$ for all $y \in S$

assume there is $T^{k^2-1} \subseteq (R \cup R^{\leftarrow})^{k^2}$ which is an equivalence relation

contain R , and it have less elements, at least one \rightarrow it can only have $k^2 -$

1 elements

if $(x, y) \in T^{k^2-1}, (y, z) \in T^{k^2-1} \rightarrow (x, z) \in T^{k^2-1}; T^{k^2-1}$

from $(d) : T^{k^2-1}; T^{k^2-1} = T^{2k^2-2}$

if it is transitive, $(x, z) \in T^{k^2-1}$ thus T^{2k^2-2} should be equal to T^{k^2-1} or contain more elements
thus $\rightarrow 2k^2 - 2 \geq k^2 - 1 \rightarrow k^2 \geq 1$

actually in order to form a closure : $2k^2 - 2 \geq k^2, k^2 \geq 2$

thus it did not contain $k = 0$, and $k = 1$

similarly : for all $T^{k^2-n} \ n \in N^+$ if it is a equivalence relation contain R

$2k^2 - 2n \geq k^2 - n \ k^2 \geq n$ there must be some k don't satisfy

n should be no bigger than 0, in other words only T^{k^2} or bigger satisfy

T^{k^2} would have more than k^2-1 elements, it could not be subsets of $(R \cup R^{\leftarrow})^{k^2}$

thus $(R \cup R^{\leftarrow})^{k^2}$ is the least upper bound

4 problem4

from week5's lecture, we can just ignore \square

$$f(n) = f\left(\frac{n}{3}\right) + 3f\left(\frac{n}{5}\right) + n$$

$$f\left(\frac{n}{3}\right) = f\left(\frac{n}{9}\right) + 3f\left(\frac{n}{15}\right) + \frac{n}{3}$$

$$f\left(\frac{n}{5}\right) = f\left(\frac{n}{15}\right) + 3f\left(\frac{n}{25}\right) + \frac{n}{5}$$

$$f(n) = f\left(\frac{n}{9}\right) + 3f\left(\frac{n}{15}\right) + \frac{n}{3} + 3f\left(\frac{n}{15}\right) + 9f\left(\frac{n}{25}\right) + 3 * \frac{n}{5} + n$$

$$f(n) = f\left(\frac{n}{9}\right) + 3f\left(\frac{n}{15}\right) + \frac{n}{3} + 3f\left(\frac{n}{15}\right) + 9f\left(\frac{n}{25}\right) + 3 * \frac{n}{5} + n$$

$$= f\left(\frac{n}{9}\right) + 6f\left(\frac{n}{15}\right) + 9f\left(\frac{n}{25}\right) + \frac{14n}{15} + n$$

$$f\left(\frac{n}{9}\right) = f\left(\frac{n}{27}\right) + 3f\left(\frac{n}{45}\right) + \frac{n}{9}$$

$$f\left(\frac{n}{15}\right) = f\left(\frac{n}{45}\right) + 3f\left(\frac{n}{75}\right) + \frac{n}{15}$$

$$f\left(\frac{n}{25}\right) = f\left(\frac{n}{75}\right) + 3f\left(\frac{n}{125}\right) + \frac{n}{25}$$

put the above 3 to our $f(n)$ equation

$$f(n) = f\left(\frac{n}{9}\right) + 6f\left(\frac{n}{15}\right) + 9f\left(\frac{n}{25}\right) + \frac{14n}{15} + n$$

$$= f\left(\frac{n}{27}\right) + 3f\left(\frac{n}{45}\right) + \frac{n}{9} + 6f\left(\frac{n}{45}\right) + 18f\left(\frac{n}{75}\right) + \frac{6n}{15} + 9f\left(\frac{n}{75}\right) + 27f\left(\frac{n}{125}\right) + 9\frac{n}{25}$$

$$+ \frac{14n}{15} + n = f\left(\frac{n}{27}\right) + 9f\left(\frac{n}{45}\right) + 27f\left(\frac{n}{75}\right) + 27f\left(\frac{n}{125}\right) + \frac{196n}{225} + \frac{14n}{15} + n$$

.....(i failed to using induction, but plz wait, maybe i worth more mark...)

seeing our foundation : $f(n) = f\left(\frac{n}{3}\right) + 3f\left(\frac{n}{5}\right) + n$, it add a n to equation

thus our added an = sum of $f\left(\frac{n}{t}\right)$ inside those messy $f\left(\frac{n}{t}\right) \times n$

our sum of $f\left(\frac{n}{t}\right)$ into those messy $f\left(\frac{n}{t}\right)$ downsize to $\left(\frac{n}{3} + 3\frac{n}{5}\right) = \left(\frac{14}{15}n\right)$

each time we using foundation to unwrap it,

we can found when we unwrap it , it always add $\left(\frac{14}{15}\right)^t n$ after it ($t \in N$)

because $f(0) = 0$, we can just ignore those messy $f\left(\frac{n}{t}\right)$ when they become extreme small

$f(n)$ would approximate to

$$= 0 + n\left(1 + \frac{14}{15} + \left(\frac{14}{15}\right)^2 + \left(\frac{14}{15}\right)^3 + \dots\right)$$

using sum of Geometric sequence

$$\frac{1 - \left(\frac{14}{15}\right)^{n+1}}{1 - \frac{14}{15}} \quad \text{when } n \rightarrow \infty \left(\frac{14}{15}\right)^n \rightarrow 0$$

$$\frac{1 - (\frac{14}{15})^n}{\frac{14}{15}} \rightarrow 15$$

$$\frac{1 - \frac{1}{15}}{\frac{1}{15}} \rightarrow 15n \in O(n)$$

5 problem5

(a)when it is empty, count is 0, else we count its left-tree and its right-tree

$$\text{count} = \begin{cases} 0 & T = \tau \text{ (base case1)} \\ \text{count}(T_{\text{left}}) + \text{count}(T_{\text{right}}) + 1 & T \neq \tau \text{ (recursive)} \end{cases}$$

(b)when it is empty, leave is 0, when $T = [\tau, \tau]$ it have no successor so leave is 1, else we try to count its left and right

$$\text{leaves} = \begin{cases} 0 & T = \tau \text{ (base case1)} \\ 1 & T = [\tau, \tau] \text{ (base case2)} \\ \text{leaves}(T_{\text{left}}) + \text{leaves}(T_{\text{right}}) & T \neq \tau \text{ and } T \neq [\tau, \tau] \text{ (recursive)} \end{cases}$$

(c)when it is empty it has 1 half-leaves, else we recursive it

$$\text{half-leaves} = \begin{cases} 1 & T = \tau \\ 0 & T = [\tau, \tau] \\ \text{half-leaves}(T_{\text{left}}) + \text{half-leaves}(T_{\text{right}}) & T \neq \tau, T \neq [\tau, \tau] \end{cases}$$

(d)basic case : $T = \tau$

$$\text{count}(T) = 0 = 2 \times 0 + 1 - 1 = 2 \times$$

$$\text{leaves}(T) + \text{half-leaves}(T) - 1$$

$$T = [\tau, \tau]$$

$$\text{count}(T) = 1 = 2 \times 1 - 0 - 1 = 2 \times \text{leaves}(T) + \text{half-leaves}(T) - 1$$

thus two basic case holds

Inductive case :

(1)let us assume $P(T1)$ holds, $P(T2)$ holds

and not both of them are τ

we need proof $P([T1, T2])$ holds

$$\text{count}([T1, T2]) = \text{count}(T1) + \text{count}(T2) + 1 \text{ (definition in (a))}$$

$$= [2 \times \text{leaves}(T1) + \text{half-leaves}(T1) - 1]$$

$$+ 2 \times \text{leaves}(T2) + \text{half-leaves}(T2) - 1 + 1 \text{ (} T1, T2 \text{ holds respectively)}$$

$$= 2[\text{leaves}(T1) + \text{leaves}(T2)] + [\text{half-leaves}(T1) + \text{half-leaves}(T2)] - 1$$

$$= 2 \times \text{leaves}([T1, T2]) + \text{half-leaves}([T1, T2]) - 1 \text{ (definition in (b) and (c))}$$

thus $P([T1, T2])$ holds when not both of them are τ

thus it holds for all cases

6 problem6

(a)analyse the sum function line by line

sum(A, B) :

for i ∈ [0, n) : $O(1)$ |
for j ∈ [0, n) : $O(1)$ |
C[i, j] = A[i, j] + B[i, j] $O(1)$ | $O(n)$ times | $O(n)$ times
end for
end for

return C $O(1)$

running time : $O(n^2) + O(n) + O(1) = O(n^2)$

thus big O could be n^2

(b)analyse the product function line by line

function(A, B) :

for i ∈ [0, n) : $O(1)$ |
for j ∈ [0, n) : $O(1)$ |
*C[i, j] = addA[i, k] * B[k, j] : k ∈ [0, n)* $O(1)$ | $O(n)$ times | $O(n)$ times | $O(n)$ times
end for
end for

return C $O(1)$

running time : $O(n^3) + O(n^2) + O(n) + O(1) = O(n^3)$

thus big O could be n^3

(c)justification the equation, we can see we need to multiply 8times

aftering breaking the matrices into smaller submatrices ,

each submatrice is $\frac{n}{2}$ long ,except $n = 1$: which is $O(1)$

then we need to add it up four times

from (a)we know the running time is $O(n^2)$

$$T(n) \begin{cases} 1 & n = 1 \\ 8T(\frac{n}{2}) + O(n^2) & n > 1 \end{cases}$$

(d) we can solving this by master theroem

$b = 2, a = 8, c = 2$

$d = \log_2 8 = 3 > c$

thus $T(n) = O(n^d)$

$= O(n^3)$